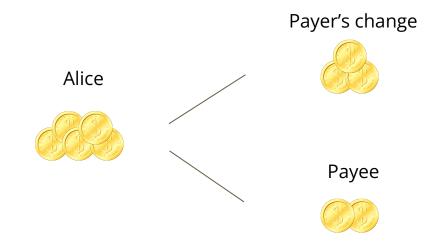
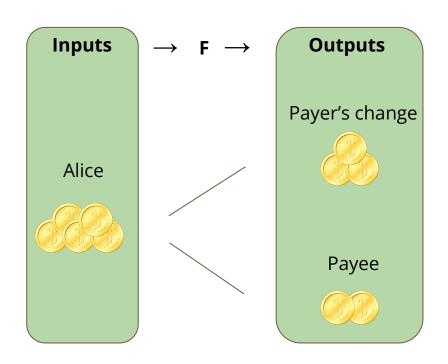
Utreexo

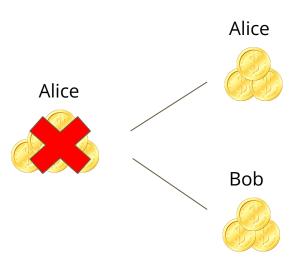


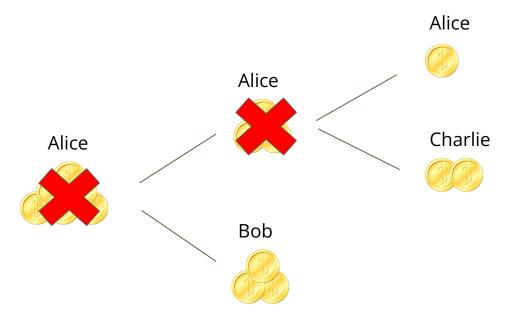


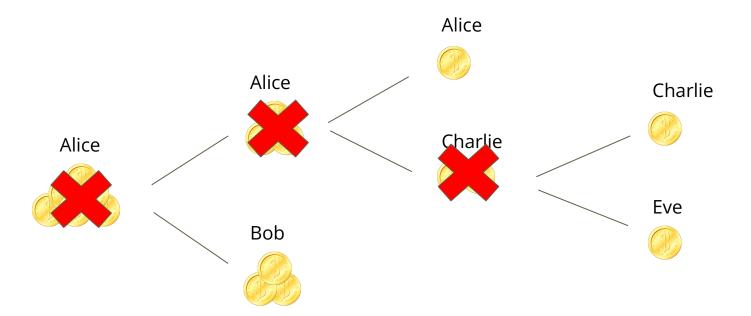


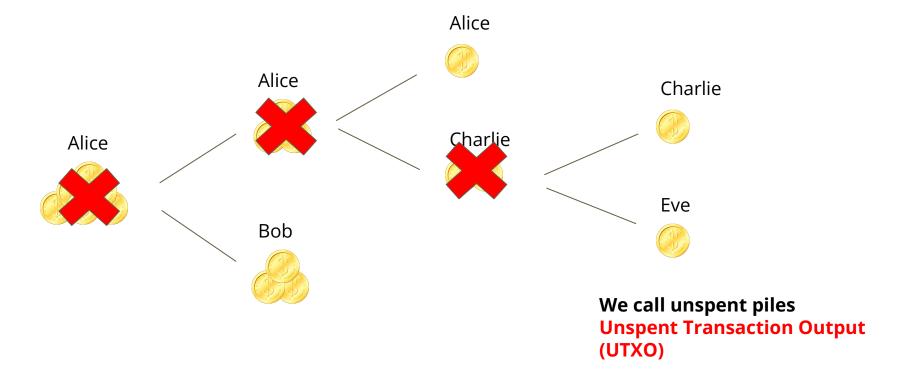
Alice











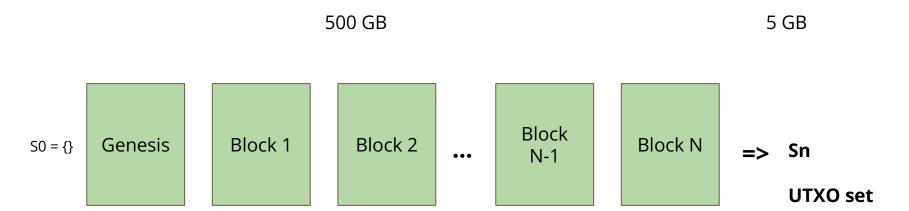
UTXO set

A set of unspent coins



Bitcoin UTXO set has over **130 MILLION** piles of coins (outputs) which is ~5GB of data.

Bitcoin sync



We download all the blocks and execute them. When we've executed all of the blocks, we arrive at the current state of the chain i.e. the current UTXO set.

Bitcoin sync with pruning

5 GB 500 GB $S0 = \{\}$ **UTXO** set We download all the blocks and execute them. When we've executed all We only keep the UTXO set!

of the blocks, we arrive at the current state of the chain i.e. the current UTXO set.

Non-dormant state

UTXO set size: ~130M

Block inputs: ~5000

A very small set of outputs is active in a block. Do we really need to know about every output on each block?

What if instead of caring about every UTXO, we only cared about those that are actually being touched in the update/block?

Cryptographic Accumulators

Given a set of elements, we can:

- Add element to the set
- Remove element from the set
- Prove element is in the set
- Prove element is not in the set

without revealing other elements in the set.

Based on the operations the accumulator supports, we mark it as static, dynamic or universal accumulator.

Utreexo (dynamic accumulator) •

- 1. Supports element addition, removal and inclusion proofs
- 2. Accumulator size is ~1kb
- 3. Optimized for Bitcoin



How does it work?

Utreexo (dynamic accumulator)

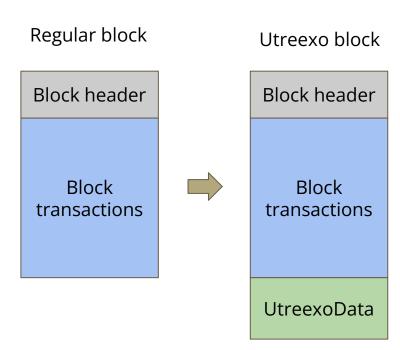
- 1. Supports element addition, removal and inclusion proofs
- 2. Accumulator size is ~1kb
- 3. Optimized for Bitcoin



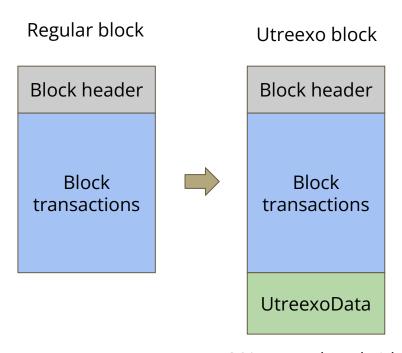
How does it work?



Utreexo block

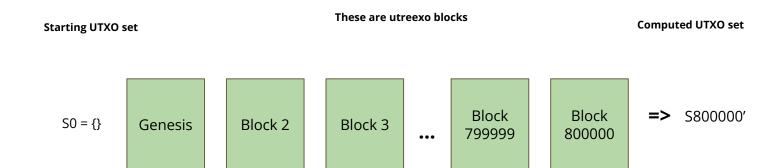


Utreexo block



~30% more bandwidth

Utreexo sync



AssumeUtxo

Starting UTXO set Computed UTXO set

S0 = {}

Genesis

Block 2

Block 3

Block 399999

Block 400000 **=>** \$400000'

S400000

Block 400001 Block 400002

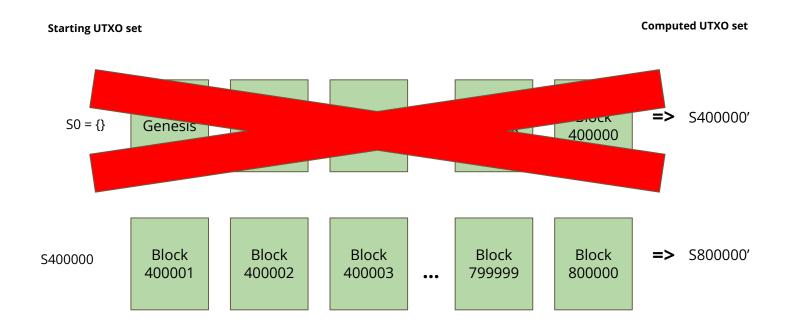
Block 400003

•••

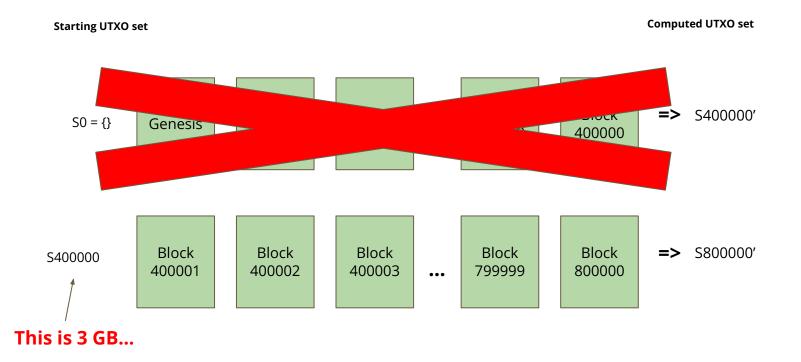
Block 799999 Block 800000

=> S800000'

AssumeUtxo



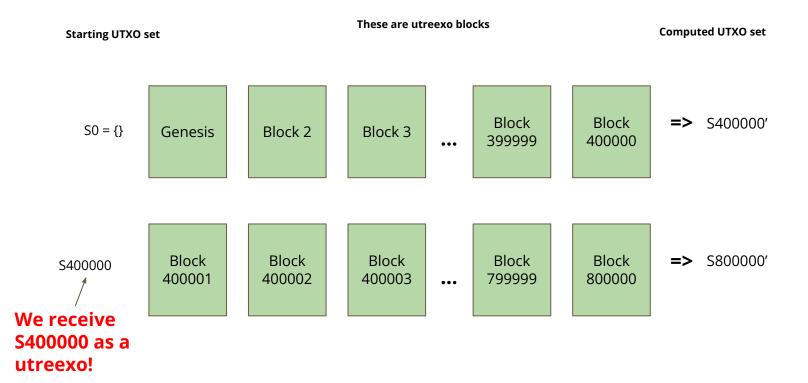
AssumeUtxo



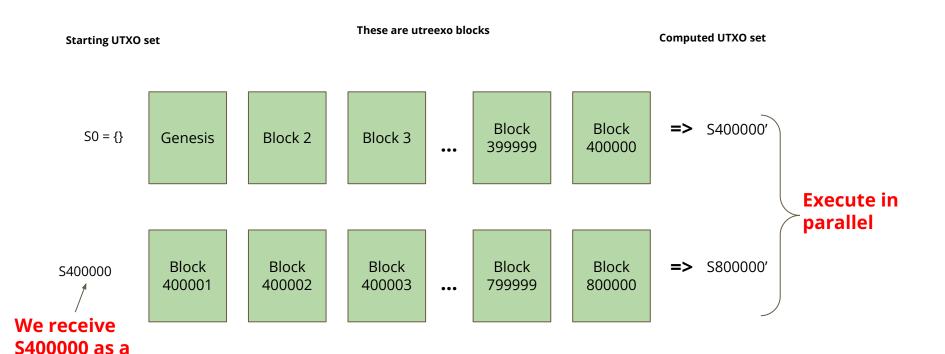
Wait a minute... \$400000 can be 1kb!

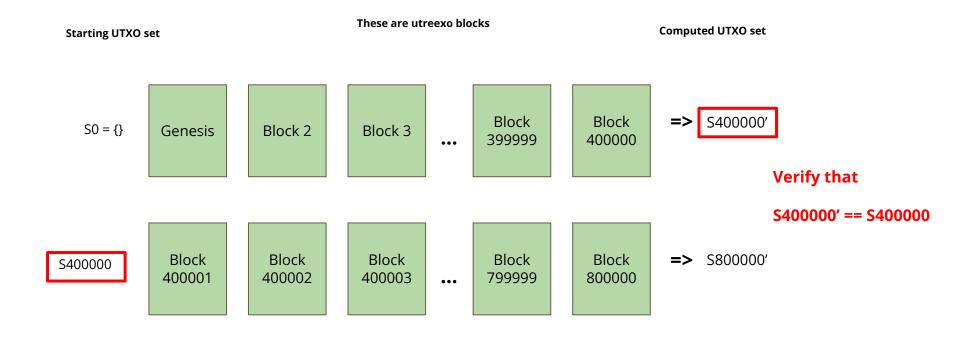
We can now make the **Initial Block Download** run in **parallel**!

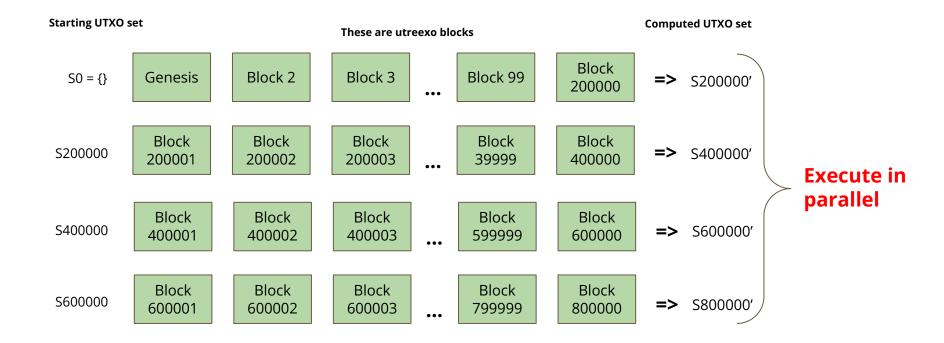




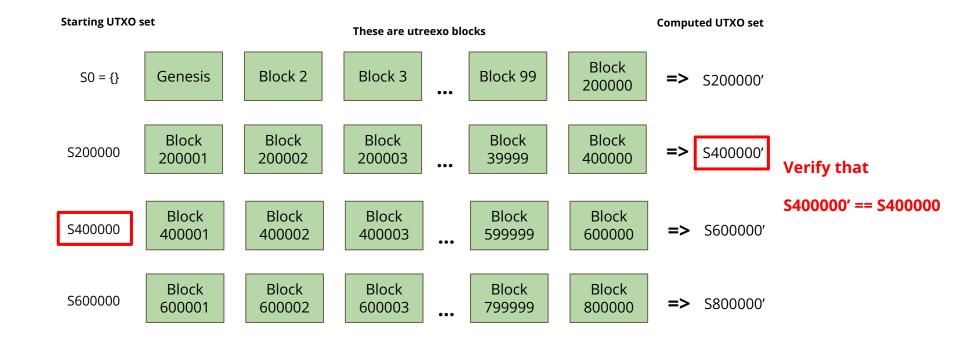
utreexo!







| Starting UTXO set | | | These are utreexo blocks | | | | Computed UTXO set | | |
|-------------------|-----------------|-----------------|--------------------------|-----|-----------------|-----------------|-------------------|----------|---------------------|
| S0 = {} | Genesis | Block 2 | Block 3 | ••• | Block 99 | Block 200000 | => | S200000' | |
| S200000 | Block 200001 | Block 200002 | Block 200003 | ••• | Block 39999 | Block 400000 | => | S400000' | Verify that |
| S400000 | Block 400001 | Block 400002 | Block 400003 | ••• | Block 599999 | Block 600000 | => | S600000' | S200000' == S200000 |
| S600000 | Block 600001 | Block 600002 | Block 600003 | ••• | Block 799999 | Block 800000 | => | S800000′ | |



| Starting UTXO set | | | These are utreexo blocks | | | | Computed UTXO set | | |
|-------------------|-----------------|-----------------|--------------------------|-----|-----------------|-----------------|-------------------|---------------------|--|
| S0 = {} | Genesis | Block 2 | Block 3 | ••• | Block 99 | Block 200000 | => S200000' | | |
| S200000 | Block 200001 | Block 200002 | Block 200003 | ••• | Block 39999 | Block 400000 | => S400000' | Verify that | |
| S400000 | Block 400001 | Block 400002 | Block 400003 | ••• | Block 599999 | Block 600000 | => S600000' | S600000' == S600000 | |
| S600000 | Block 600001 | Block 600002 | Block 600003 | ••• | Block 799999 | Block 800000 | => S800000' | | |

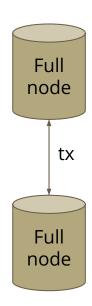
| Starting UTXO set | | | These are utreexo blocks | | | | Computed UTXO set | | |
|-------------------|-----------------|-----------------|--------------------------|-----|-----------------|-----------------|------------------------|------------------------------|--|
| S0 = {} | Genesis | Block 2 | Block 3 | ••• | Block 99 | Block 200000 | => S200000′ | | |
| S200000 | Block 200001 | Block 200002 | Block 200003 | ••• | Block 39999 | Block 400000 | => \$400000′ | S800000' is the current UTXO | |
| S400000 | Block 400001 | Block 400002 | Block 400003 | ••• | Block 599999 | Block 600000 | => \$600000′ | set | |
| S600000 | Block 600001 | Block 600002 | Block 600003 | ••• | Block 799999 | Block 800000 | => S800000' | | |

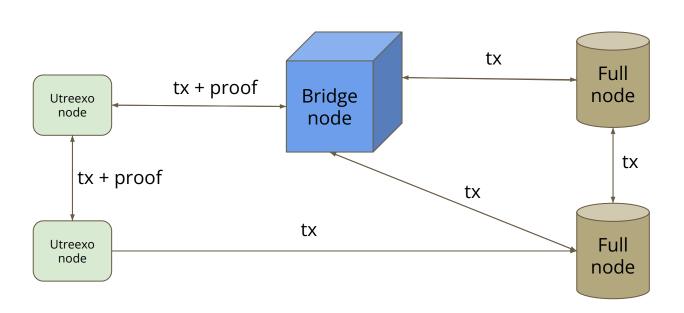


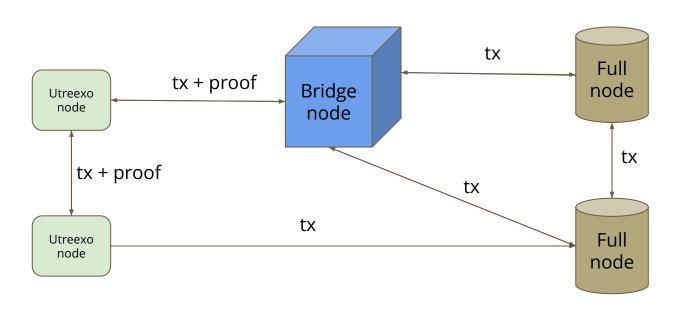
Utreexo node



Where do I get blocks with proofs from? No way to bootstrap.







No soft fork needed!



Tradeoffs?

Pros:

- Compressed state from 5GB to ~1kb
- No disk IO
- No database! (levelDB)
- Parallel IBD
- No soft forks

Cons:

- More bandwidth
- Proof construction/updating

Quick summary

- 1. We have a large set of UTXOs (130 million)
- 2. We can already prune blocks, but we're left with 5GB state (utxo set)
- 3. We can create a compact representation of that set (~1kb)
- 4. We can prove elements are in the compact set representation as well as add/delete them

We now have a node that takes very little storage, but requires a bit more data to download.

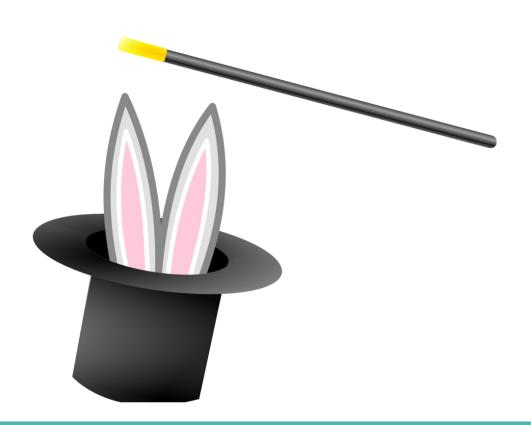
What does this enable?

Compact full nodes.

Any device or process connected to the internet could run a compact full node without using much storage.

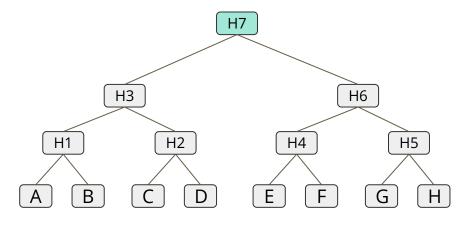
- Web wallet Chrome extension node (every time we open a web browser, an extension syncs the utreexo node)
- Router node Router could run a node in the background

Ok, but what is this magic?



Merkle trees

A merkle tree commits to a sequence of elements.



H1 commits to A and B.

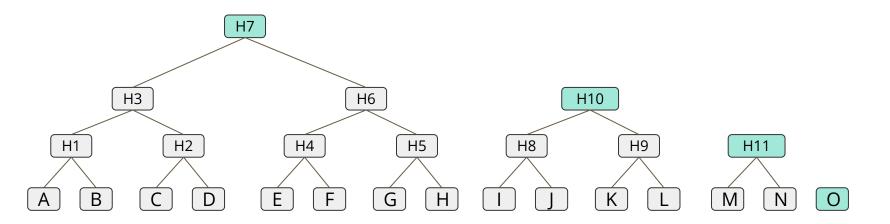
H2 commits to **C** and **D**.

H3 commits to H1 and H2 => A, B, C and D.

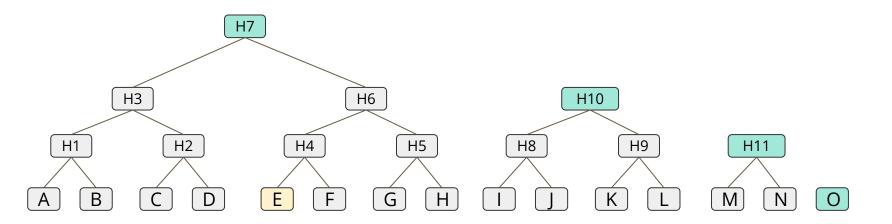
The root of this merkle tree (H7) commits to A, B, C, D, E, F, G and H which are all the 8 elements in the tree.

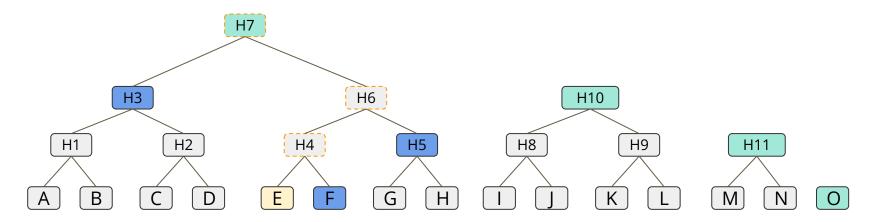
Utreexo - What does it look like?

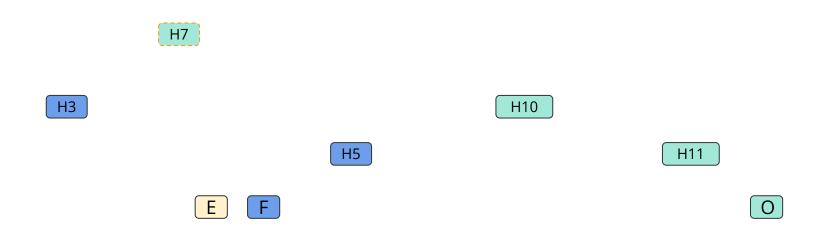
It's a forest of perfect merkle trees. A utreexo node only keeps the roots of the trees.

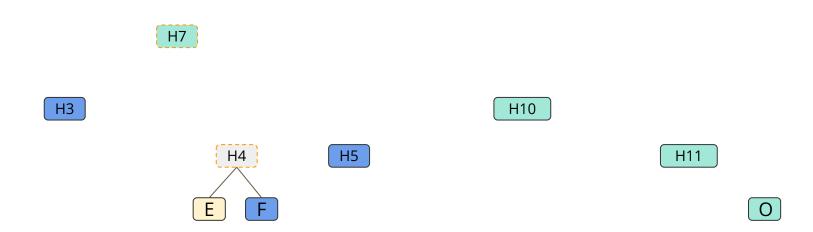


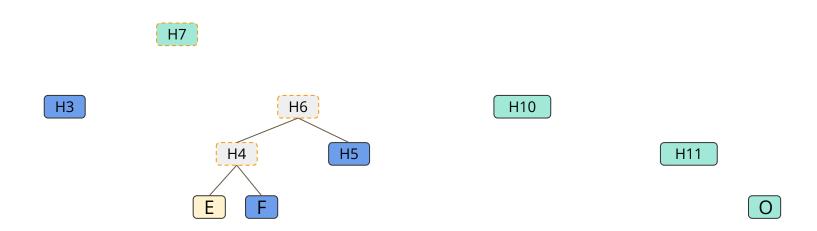
How do we prove E is in the set?

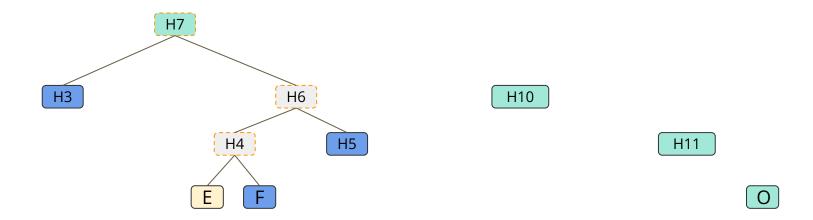


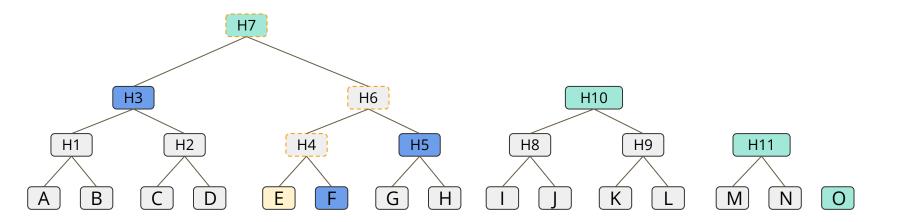


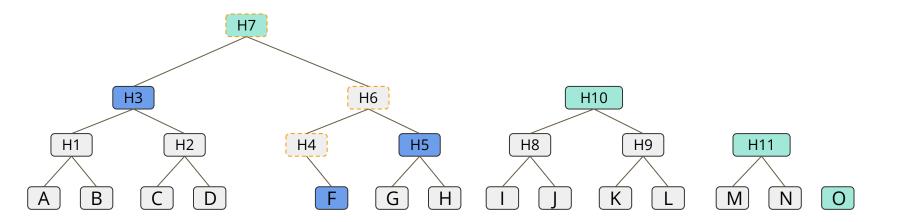


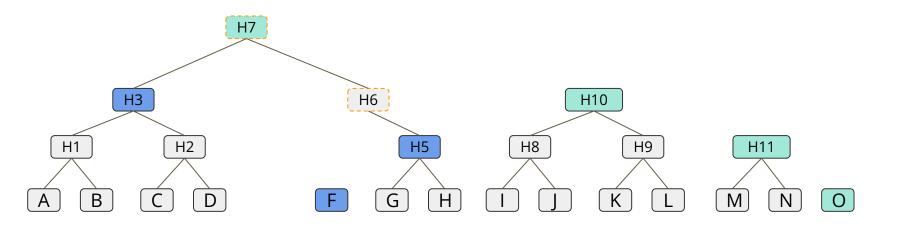


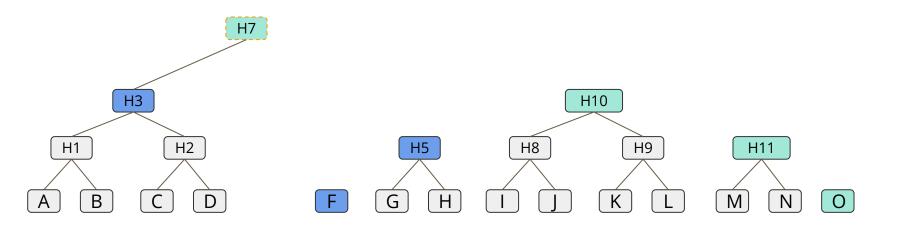


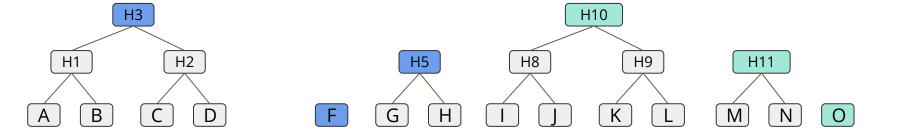


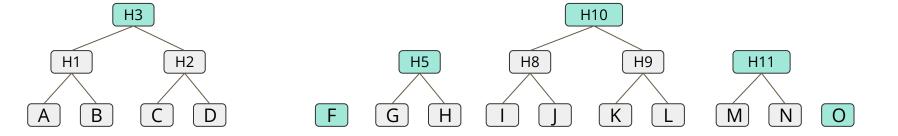


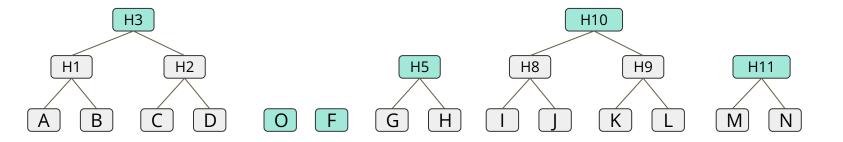


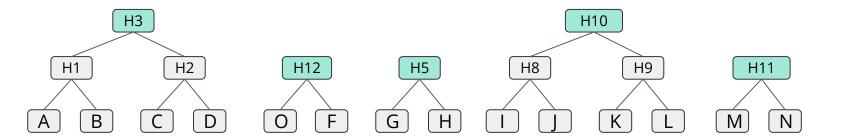


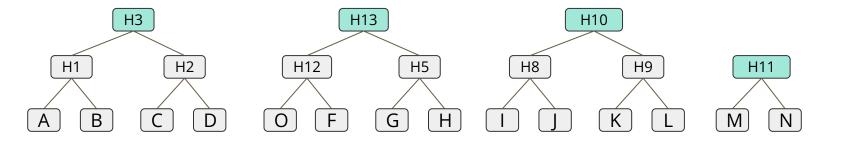


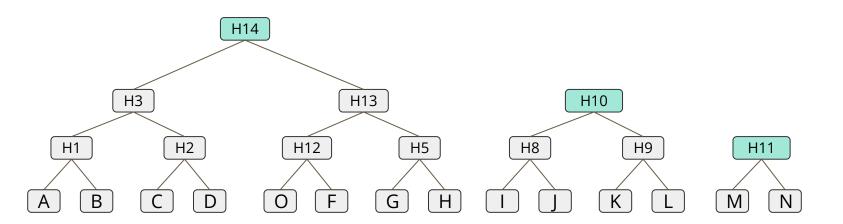


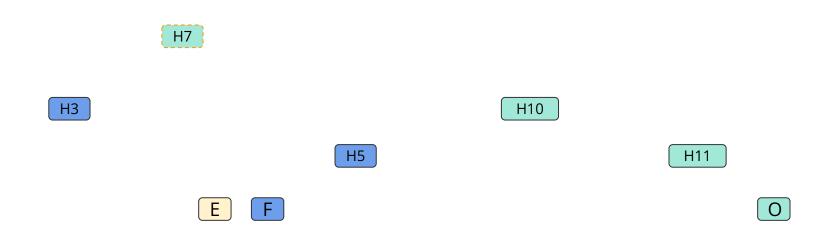


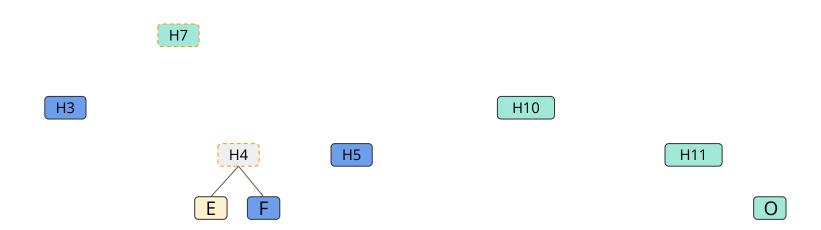


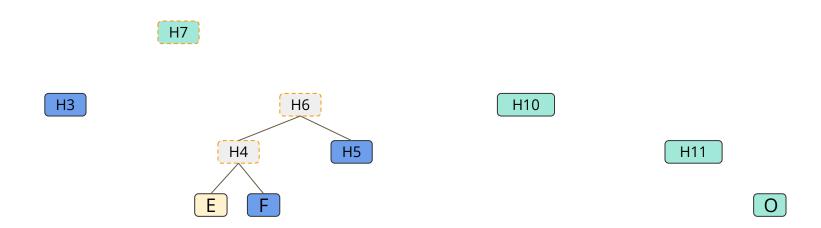


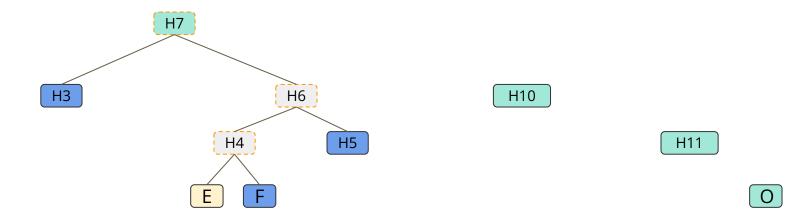


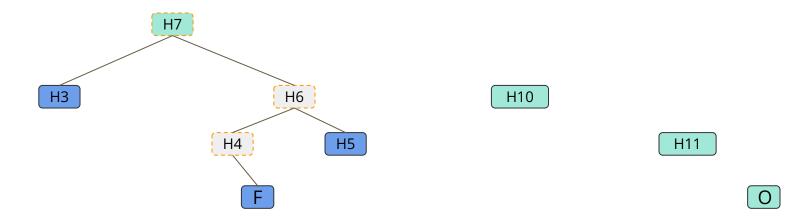


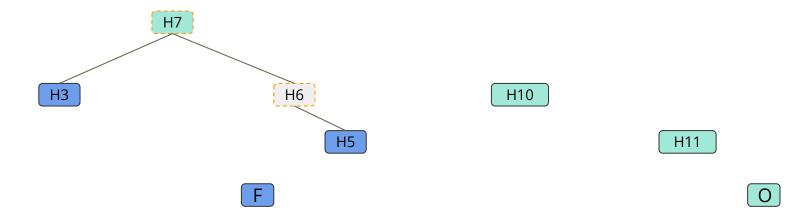


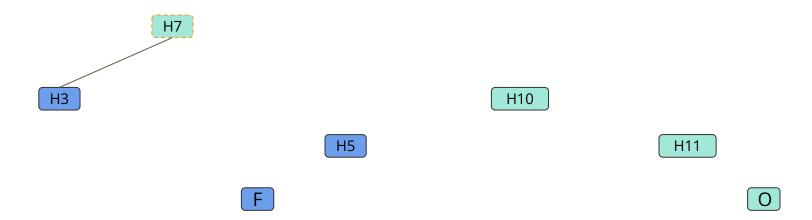












We use inclusion proofs to delete E and find another element to move to the position of E.

H3 H10 H11

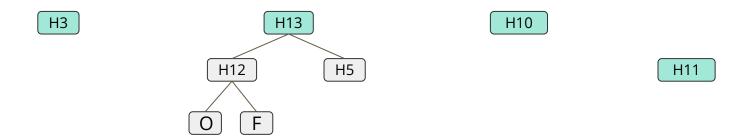
We use inclusion proofs to delete E and find another element to move to the position of E.

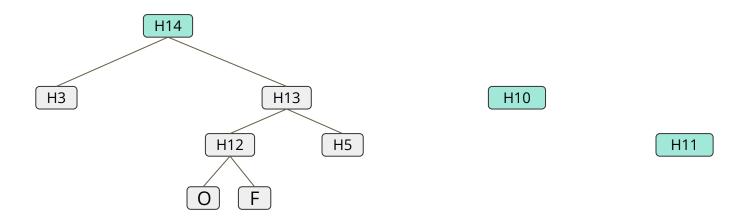
H3 H10 H11 H11

We use inclusion proofs to delete E and find another element to move to the position of E.

H10
H5
H11



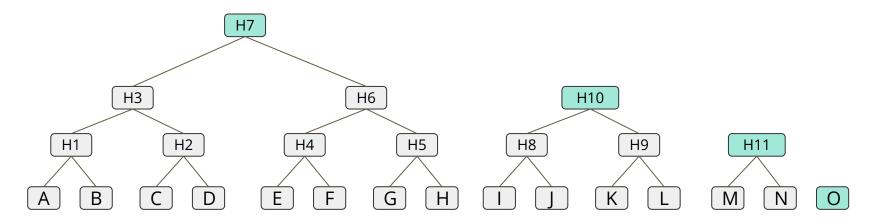




Delete doesn't work exactly like this in the implementation. There are several possible ways of implementing it...

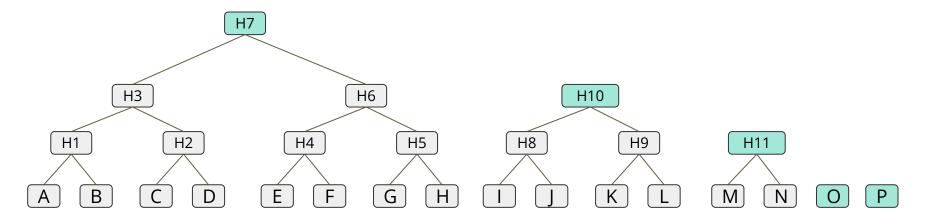
Utreexo - add •

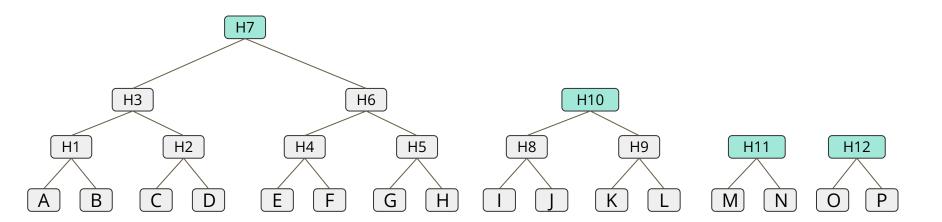
We just add an element to the right and compute the hashes that we can.

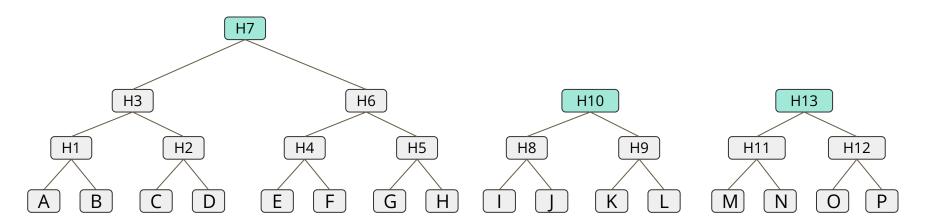


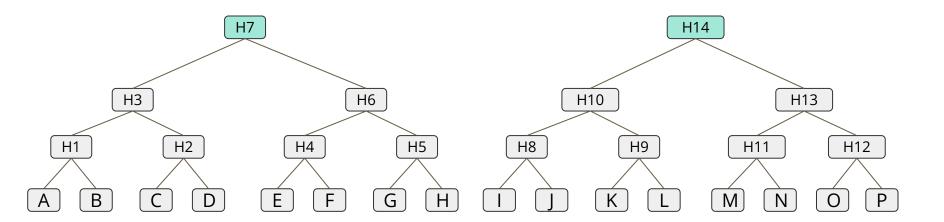
Utreexo - add •

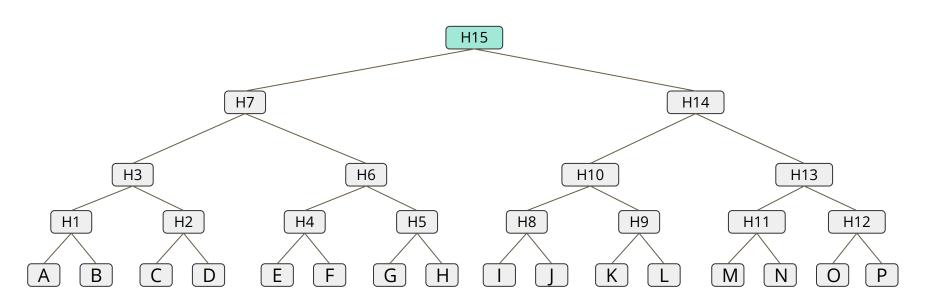
We just add an element to the right and compute the hashes that we can.











We just add an element to the right and compute the hashes that we can.

Н7

H10



We just add an element to the right and compute the hashes that we can.

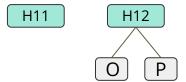
Н7

H10

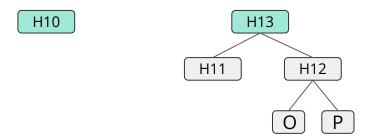


We just add an element to the right and compute the hashes that we can.

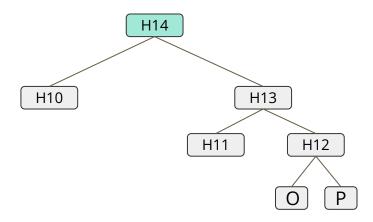
H7

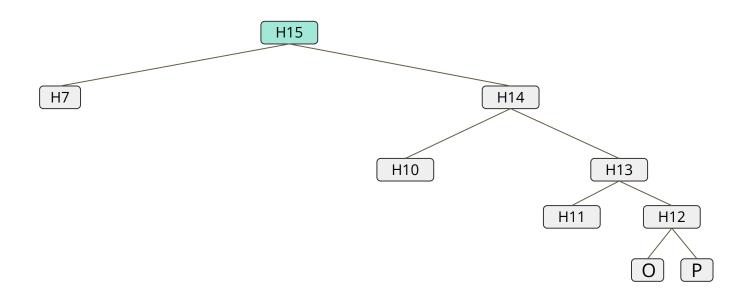


We just add an element to the right and compute the hashes that we can.



We just add an element to the right and compute the hashes that we can.

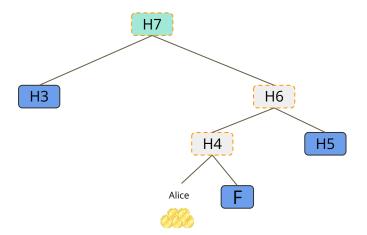




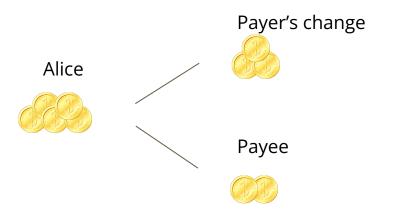
Utreexo transaction

Transaction has additional proofs added to prove the inputs are in the UTXO set.

1) Proofs for transaction inputs

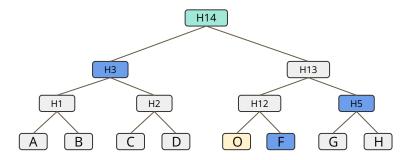


2) Transaction

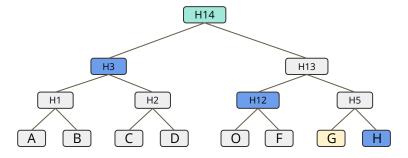


Utreexo optimizations: Combining proofs

Inclusion proof for O

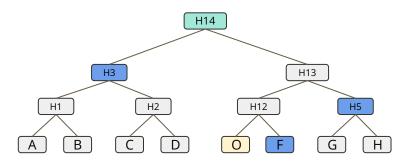


Inclusion proof for G

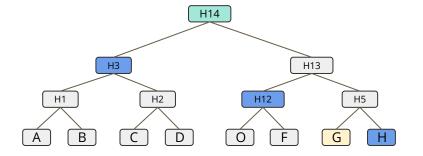


Utreexo optimizations: Combining proofs

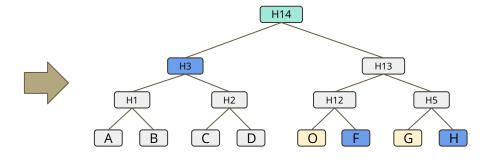
Inclusion proof for O



Inclusion proof for G



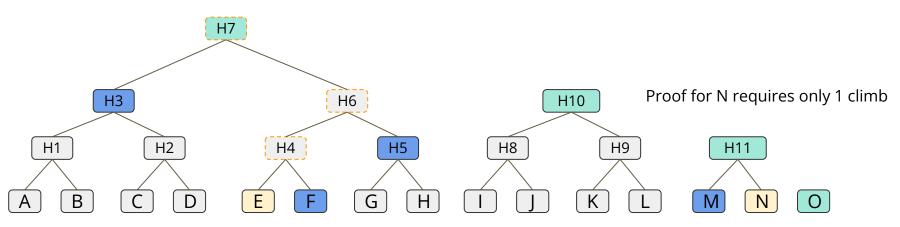
Inclusion proof for O and G



Instead of sending 2*4 pieces, we send only 5 and can still prove both O and G.

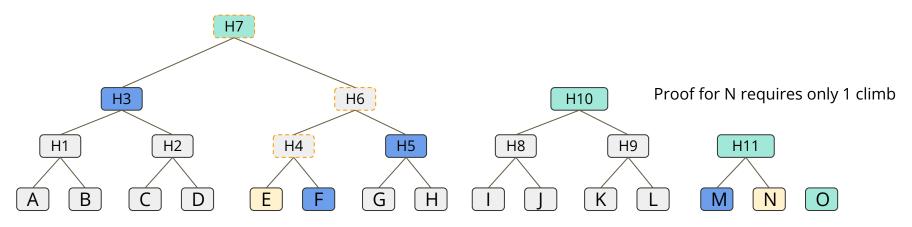
Utreexo optimizations: Proof length reduction

Proof for E requires 3 climbs



Utreexo optimizations: Proof length reduction

Proof for E requires 3 climbs



In Bitcoin, the majority of outputs are spent soon after they're created. **If we make sure new outputs are on the right, they'll have shorter proofs!**

Utreexo optimizations: Look-ahead caching

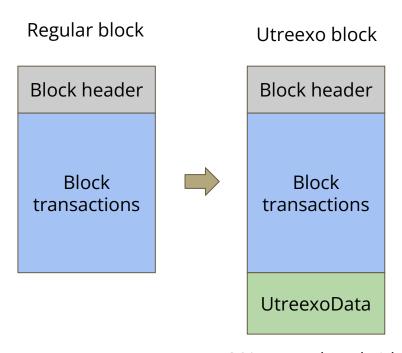
The node knows what happens in the future and can give us hints to save bandwidth and computation.

Ricevuto!

Memory: A, B, C

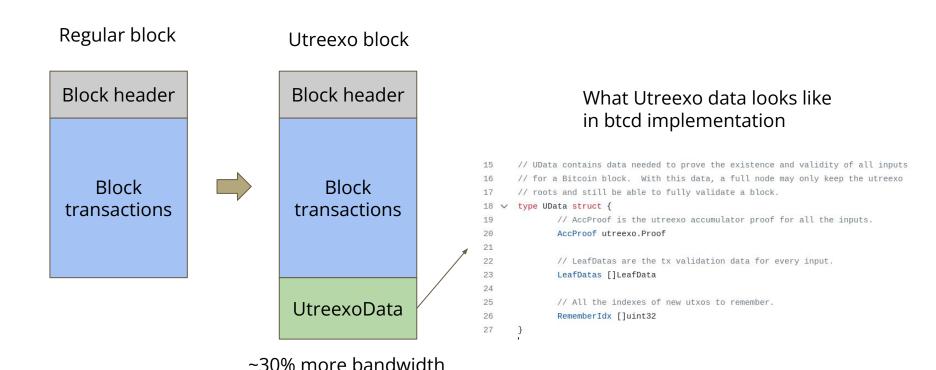
Yo!
Remember UTXOs A, B
and C from this block
for 10 blocks because
they'll get used.

Utreexo block



~30% more bandwidth

Utreexo block



Status

Focus:

- Btcd fork implementation (no work on Core at the moment)
- Very likely to get into some wallets (TBD)

Core things = 3 BIPS + 1 BIP Fix

- BIP for the accumulator
- BIP for how block/tx verification works with the accumulator
- BIP for P2P messaging with utreexo proofs
- BIP that fixes BIP30

Links:

Utreexo full node - https://github.com/utreexo/utreexo/utreexo
Utreexo - https://github.com/utreexo/utreexo

Questions?

