

# Heterogeneous LDPC Decoder Algorithm on ARM and GPU of Mobile Devices

Roohollah Amiri  
Department of Electrical and  
Computer Engineering  
Boise State University

Email: roohollahamiri@boisestate.edu

Hani Mehrpouyan  
Department of Electrical and  
Computer Engineering  
Boise State University

Email: hanimehrpouyan@boisestate.edu

**Abstract**—Low density parity check (LDPC) codes have been extensively applied in mobile communication systems due to their excellent error correcting capabilities. However, their wide adoption has been hindered by the high complexity of the LDPC decoder. Although to date, dedicated hardware has been used to implement low latency LDPC decoders, recent advancements in the architecture of mobile processors has made it possible to develop software solutions. Here, unlike prior solutions that are based on either graphic processing units (GPUs) or advanced RISC machine (ARM) architectures, we propose a heterogeneous LDPC decoder that uses both the ARM and GPU processors of a mobile device to achieve efficient real-time decoding. The proposed solution is implemented on an NVIDIA development kit, where our results indicate that we can reduce the load on either the GPU or the ARM processor through the proposed heterogeneous structure, which in turn allows these resources to support other applications, simultaneously.

## I. INTRODUCTION

Originally proposed by Robert Gallager in 1962 [1] and rediscovered by MacKay and Neal in 1996 [2] Low Density Parity Check (LDPC) codes have been adopted by a wide range of applications including many communication standards such as WiFi (IEEE 802.11n), 10 Gigabit Ethernet (IEEE 802.3an), Long Term Evolution (LTE), and DVB-S2. Recently, Chung and Richardson [3] showed that a class of LDPC codes can approach the Shannon limit to within 0.0045 dB. However, the error correcting strength of these codes comes at the cost of very high decoding complexity [4]. Moreover, to date, there are no closed-form solutions to determine the performance of LDPC codes in various wireless channels and systems. Thus, performance evaluation is typically carried out via simulations on computers or dedicated hardwares [5].

Since LDPC decoding algorithms are computationally-intensive and need powerful computer architecture to result in low latency and high throughput, to date, most LDPC decoders are implemented using application-specific integrated circuits(ASIC) or field-programmable gate array (FPGA) circuits [6]. However, their high speed often comes at a price of high development cost and low programming flexibility [7] and it is very challenging to design decoder hardware that supports various standards and multiple data rates [8]. On the other hand, iterative LDPC decoding schemes based on the sum-product algorithm (SPA) can fully be parallelized,

leading to high-speed decoding [3]. For these reasons, designers have recently focused on software implementations of LDPC decoders on multi/many-core devices [9] to meet the performance requirements of current communication systems through Software Defined Radio (SDR). As in terms of multicore architectures, researchers have used CPUs [10], [11], GPUs [5], [9], [12] and advanced RISC machine (ARM) [11], [13] architectures to develop high throughput, low latency SDRs.

In microarchitectures, increasing clock frequencies to obtain faster processing performance has reached the limits of silicon based architectures. Hence, to achieve gains in processing performance, other techniques based on parallel processing is being investigated [4]. Today's multicore architectures support Single Instruction Multiple Data (SIMD), Single Program Multiple Data (SPMD) and Single Instruction Multiple Threads (SMT). The general purpose multicore processors replicate a single core in a homogeneous way, typically with a x86 instruction set, and provide shared memory hardware mechanisms [9]. Such multi-core structures can be programmed at a high level by using different software technologies [14]. OpenMP [15] provides an effective and relatively straightforward approach for programming general-purpose multicores. On the other hand newer microarchitectures are trying to provide larger SIMD units for vector processing like Streaming SIMD Extensions (SSE), Advanced Vector Extensions (AVX) and AVX2 [16] on Intel Architectures. In [4], the authors have used Intel SSE/AVX2 SIMD Units to efficiently implement a high throughput LDPC decoder. In [8], OpenMP is used to generate address patterns with parity check H-matrix.

Mainly due to the demands for visualization technology in the gaming industry, the performance of graphics processing units (GPUs) has significantly improved over the last decade. With many cores driven by a considerable memory bandwidth, recent GPUs are targeted for solving computationally intensive algorithms in a multithreaded and highly parallel fashion. Hence, researchers in the high-performance computing field are applying GPUs to general-purpose applications (GPGPU) [5], [8], [12], [17]–[19]. Pertaining to the field of communication, researchers have used Compute Unified Device Architecture (CUDA) from NVIDIA [20] and Open Computing

Language (OpenCL) platforms to develop LDPC decoders on GPUs.

Due to large computing capacity of multicore devices, software based LDPC decoders have met the required throughputs of communication standards, although power consumption of x86 and GPU devices is incompatible with most of embedded systems [13]. To solve this issue, ARM-based SDR systems have been proposed in recent years [6], [11], [13] with the goal to develop an SDR based LDPC decoder that provides high throughput and low latency on a low-power embedded system. The authors in [13] have used ARM processor's SIMD and SMT programming models to implement a LDPC decoder that is based on parallel decoding of data frames. This approach allows reaching high throughput while maintaining low-latency. However, the proposed ARM based solution in [13], is based on the assumption that the ARM processor is solely used for LDPC decoding and does not take advantage of the available GPU processing on mobile device. Due to restrictions in an embedded system, using all resources of the system is a crucial task. Recent works in SDR LDPC embedded systems are missing the fact that today's mobile devices have powerful CUDA enabled GPUs. This paper has proposed a new algorithm that exploits ARM NEON SIMD Units and GPU together to reach a high throughput, low latency LDPC decoder. The main specification of the algorithm is that it divides processing task between system's resources.

## II. LDPC CODES AND THEIR DECODING PROCESSES

Many works as in [6], [9], [11], [17] focused on mapping LDPC decoders on multicore architectures. Most of these works are based on the standard Two-Phase Message Passing (TPMP) schedule described in [9]. This algorithm works in two phases. In the first phase, all the variable nodes send messages to their neighboring parity check nodes, and in the second phase the parity check nodes send messages to their neighboring variable nodes. Due to transcendental operations and relying of Sum-Product algorithm to the estimation of noise standard deviation, in practice Min-Sum (MS) variants are preferred by designers [13]. More efficient layered schedules, such as horizontal layered-based decoding algorithm, allow updated information to be utilized more quickly in the algorithm thus speeding up the decoding [19], [21]. In fact, the parity check matrix can be viewed as a layered graph decoded sequentially. The work in [17] has applied a form of layered belief propagation to irregular LDPC codes to reach 2x faster convergence in a given error rate. By using this method they have reduced memory bits usage by 45-50%. The major limitation of layered algorithm is its irregular memory access although it is composed of a single loop kernel compared to two sequential kernels in standard algorithms. To solve the irregular memory access a data interleaving/deinterleaving process is being used before and after the decoding process [13], [17].

In this paper the interleaving/deinterleaving process is done by using ARM Vector processing units and frame decoding is being done in GPU of a mobile device.

## III. PARALLEL FRAME PROCESSING

The proposed LDPC decoder is implemented on Jetson TK1 SoCs which contains 4 Cortex-A15 processors. Each core includes a NEON SIMD unit. To achieve high throughput performance on such low-power embedded processors, the following programming model is exploited in the proposed LDPC decoder.

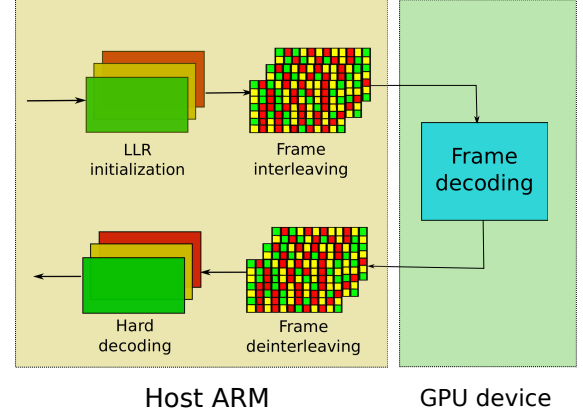


Fig. 1. Proposed Setup for Heterogeneous LDPC decoding

Typically, there are two ways to deliver messages in LDPC decoding. One is to use probabilities, and the other is to use log-likelihood ratios (LLRs). In general, using LLRs is favored since that allows us to replace expensive multiplication operations with inexpensive addition operations [8]. So the host is in charge of Initialization of Check Nodes (CNS), Frame interleaving before decoding and frame deinterleaving after decoding. From decoder point of view, host sends/receives data to/from the GPU device as the decoder. The GPU device is responsible for all CNs to Variable Nodes (VNs) computations that is done in one kernel (see figure 1). At the end of decoding, hard decision decodings are taken and decisions are sent back to the host. SIMD programming model in host enables each processor core to interleave  $F$  frames in parallel with 8-bitxF the width in bits of SIMD unit. So there is  $C$  (number of host cores) set of  $F$  frame streams of data into GPU device. Each processor controls its own stream to GPU. On the GPU there are  $C$  similar kernels running. As long as the memory that is used in GPU is bigger than  $C \times F \times 8$ -bits, there would be no problem in memory allocation.

## IV. EXPERIMENTAL RESULTS

The experiments were carried out by decoding LDPC codes using NVIDIA Tegra K1 SoCs. The programs compiled with GCC-4.8 and CUDA 6.5. The TK1 is composed of 4 Cortex-A15 ARM processors and one NVIDIA Kepler "GK20a" GPU with 192 SM3.2 CUDA cores. The host platform uses a GNU/Linux kernel 3.10.40-gdacc96.

## V. CONCLUSION

The conclusion goes here.

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