

Category theory in context: 4.4 — Calculus of Adjunctions

Siddharth Bhat

Monsoon, second year of the plague

1 PROPOSITION 4.4.1

If F, F' are both left adjoint to G , then $F \simeq F'$. Moreover, there is a unique iso $\theta : F \simeq F'$ commuting with units and counits of adjunctions:

$$\begin{array}{ccc} 1_C & \xrightarrow{\eta} & GF \\ & \searrow \eta' & \downarrow G\theta \\ & & GF' \end{array} \quad \begin{array}{ccc} FG & \xrightarrow{\epsilon} & 1_D \\ \theta G \downarrow & & \nearrow \epsilon' \\ F'G & & \end{array}$$

1.1 Proof by unit/counit

Let's consider the data we need to define for an iso $\theta : F \Rightarrow F'$. Drawing out the naturality square, we need the arrows:

$$\begin{array}{ccc} Fc & \xrightarrow{\theta_c} & F'c \\ \downarrow Ff & & \downarrow F'f \\ Fc' & \xrightarrow{\theta_{c'}} & F'c' \end{array}$$

By adjunction, defining a commutative diagram with $Fc \rightarrow d$ is the same as defining a commutative diagram with $c \rightarrow Gd$:

$$\begin{array}{ccc} c & \xrightarrow{\theta_c^\#} & GF'c \\ f \downarrow & & \downarrow GF'f \\ c' & \xrightarrow{\theta_{c'}^\#} & GF'c' \end{array}$$

We define $\theta^\# \equiv \eta' : 1 \rightarrow GF'$, since the types match. Using this, we compute a formula for θ as the transpose of $\theta^\#$. [TODO: how did we compute this in the first place?]

$$\theta \equiv F \xRightarrow{F\eta'} FGF' \xRightarrow{\epsilon F'} F'$$

Exchanging the roles of F with F' , η with η' , and ϵ with ϵ' , this also computes a formula for θ' given by:

$$\theta' \equiv F' \xRightarrow{F'\eta} F'GF \xRightarrow{\epsilon' F} F$$

The hope is that θ and θ' are inverse natural transforms. We need to check that $\theta' \circ \theta = 1_F$. We claim that it suffices to check that $GF(\theta' \circ \theta) \circ \eta = \eta$. [TODO: why does this suffice?]

Writing out $G(\theta' \circ \theta) \circ \eta$, which is equal to $G\theta' \circ G\theta \circ \eta$:

$$\begin{aligned} 1 &\xRightarrow{\eta} GF \xrightarrow{G\theta} GF' \xrightarrow{G\theta'} GF \\ 1 &\xRightarrow{\eta} GF \xrightarrow{GF\eta'} GFGF' \xrightarrow{G\epsilon F'} GF' \xrightarrow{GF'\eta} GF'GF \xrightarrow{G\epsilon' F} GF \end{aligned}$$

We wish to swap η with $GF\eta'$ (at the first two terms) to bring the η and ϵ close together (at the first three terms) so we can use the triangle identities. To do this, we consider the commutative square, where we transport the morphism $c \xrightarrow{\eta'_c} GF'c$ along $\eta : 1_C x \rightarrow GFx$ to give:

$$\begin{array}{ccccc} & & 1_C(x) & \xrightarrow{\eta_x} & GF(x) \\ & & & & \\ c & & 1_C(c) & \xrightarrow{\eta_c} & GF(c) \\ \downarrow \eta'_c & & \downarrow 1_C\eta'_c & \eta \text{ natural} & \downarrow GF\eta'_c \\ GF'c & & 1_C(GF'c) & \xrightarrow{\eta_{GF'c}} & GF(GF'c) \end{array}$$

- See that this square contains $1 \xRightarrow{\eta} GF \xrightarrow{GF\eta'} GFGF'$, by following right and top. The commutativity
- of the square witnesses that this is equal to $1 \xRightarrow{\eta'} GF' \xrightarrow{\eta_{GF'}} GFGF'$.
- See that $\eta_{GF'}$ equals $\eta GF'$ since $\eta GF'(x) \equiv \eta_{GF'} GF'x$, which is the same as $\eta_{GF'}(GF'x)$.
- So, in total, the commutativity of this naturality square allows us to rewrite the segment $1 \xRightarrow{\eta} GF \xrightarrow{GF\eta'} GFGF'$ with $1 \xRightarrow{\eta'} GF' \xrightarrow{\eta_{GF'}} GFGF'$.

This gives us the diagram:

$$\begin{aligned} 1 &\xRightarrow{\eta} GF \xrightarrow{GF\eta'} GFGF' \xrightarrow{G\epsilon F'} GF' \xrightarrow{GF'\eta} GF'GF \xrightarrow{G\epsilon' F} GF \\ 1 &\xRightarrow{\eta'} GF' \xrightarrow{\eta_{GF'}} GFGF' \xrightarrow{G\epsilon F'} GF' \xrightarrow{GF'\eta} GF'GF \xrightarrow{G\epsilon' F} GF \end{aligned}$$

This is regrouped using $G\epsilon \circ \eta G = 1_G$ into:

$$\begin{aligned}
1 &\xRightarrow{\eta'} GF' \xRightarrow{\eta GF'} GFGF' \xRightarrow{G\epsilon F'} GF' \xRightarrow{GF'\eta} GF'GF \xRightarrow{G\epsilon' F} GF \\
1 &\xRightarrow{\eta'} GF' \xRightarrow{\eta GF'} GFGF' \xRightarrow{G\epsilon F'} GF' \xRightarrow{GF'\eta} GF'GF \xRightarrow{G\epsilon' F} GF \\
1 &\xRightarrow{\eta'} GF' \xRightarrow{(\eta G; G\epsilon)F'} GF' \xRightarrow{GF'\eta} GF'GF \xRightarrow{G\epsilon' F} GF \\
1 &\xRightarrow{\eta'} GF' \xRightarrow{GF'\eta} GF'GF \xRightarrow{G\epsilon' F} GF
\end{aligned}$$

Next, we use the naturality of η to swap η' with $GF'\eta$:

$$\begin{aligned}
1 &\xRightarrow{\eta'} GF' \xRightarrow{GF'\eta} GF'GF \xRightarrow{G\epsilon' F} GF \\
1 &\xRightarrow{\eta'} GF' \xRightarrow{GF'\eta} GF'GF \xRightarrow{G\epsilon' F} GF \\
1 &\xRightarrow{\eta} GF \xRightarrow{\eta' GF} GF'GF \xRightarrow{G\epsilon' F} GF
\end{aligned}$$

Finally, we use the identity $G\epsilon' \circ \eta' G = 1_G$ to reduce the equation:

$$\begin{aligned}
1 &\xRightarrow{\eta} GF \xRightarrow{\eta' GF} GF'GF \xRightarrow{G\epsilon' F} GF \\
1 &\xRightarrow{\eta} GF \xRightarrow{\eta' GF} GF'GF \xRightarrow{G\epsilon' F} GF \\
1 &\xRightarrow{\eta} GF \xRightarrow{(\eta' G; G\epsilon)F} GF \\
1 &\xRightarrow{\eta} GF
\end{aligned}$$

1.2 Proof by Yoneda

- Since $F \vdash G$, we have that $D(Fc, d) \simeq C(c, Gd)$.
- Similarly, since $F' \vdash G$, we have $C(c, Gd) \simeq D(F'c, d)$.
- Together, this gives $D(Fc, d) \simeq D(F'c, d)$, natural in both c and d .
- This implies that $D(Fc, -) \simeq D(F'c, -)$, natural in c , or by Yoneda, that $Fc \simeq F'c$, natural in c .
- The naturality in c allows us to deduce that $F \simeq F'$.
- We can identify the morphism which sends Fc to $F'c$ by choosing $d = Fc$. This will start at $D(Fc, d = Fc)$ and ends at $D(F'c, d = Fc)$.

We compute θ_c by contemplating the diagram below, and setting $d = Fc$ to arrive at a morphism from $1_{Fc} \in D(Fc, d = Fc)$ to $\theta'_c \in D(F'c, d = Fc)$:
[TODO: fill in the ?]

$$D(Fc, d) \longrightarrow C(c, Gd) \longrightarrow D(F'c, d)$$

$$f : Fc \rightarrow d \longmapsto c \xrightarrow{\eta_c} GFc \xrightarrow{Gf} Gd$$

$$g : c \rightarrow Gd \longmapsto F'c \xrightarrow{F'g} F'Gd \xrightarrow{\epsilon'_d} d$$

$$1_{Fc} \in D(Fc, Fc) \longrightarrow ?$$

2 PROPOSITION 4.4.4

Given adjunctions $F \vdash G$ and $F' \vdash G'$, their composite FF' is left adjoint to the composite GG' :

$$C \begin{array}{c} \xrightarrow{F} \\ \perp \\ \xleftarrow{G} \end{array} D \begin{array}{c} \xrightarrow{F'} \\ \perp \\ \xleftarrow{G'} \end{array} E \quad \rightsquigarrow \quad C \begin{array}{c} \xrightarrow{F'F} \\ \perp \\ \xleftarrow{GG'} \end{array} E$$

2.1 Proof by unit/counit

- The only “reasonable” definition of $\bar{\eta} : 1_C \Rightarrow GG'F'F$ is given by:

$$\bar{\eta} \equiv 1_c \xRightarrow{\eta} GF \xRightarrow{G\eta'F} GG'FF'$$

- A point to note: morally, the reason we build $G\eta'F$ is for the types to work; $\eta' : 1_D \rightarrow G'F'$. To mutate GF , it is the only type valid choice among $(\eta'GF, G\eta'F, \text{ and } GF\eta')$.
- Similarly, the only reasonable definition of $\bar{\epsilon} : F'FGG' \Rightarrow 1_E$ is given by the other expression as in the text.
- I dare not perform the “entertaining” diagram chase.

2.2 Proof by Yoneda

The pleasant proof by yoneda:

$$\begin{aligned} E(F'Fc, e) &\simeq D(F'c, Gd) & (F \vdash G) \\ D(F'c, d) &\simeq C(c, G'Gd) & (F' \vdash G') \end{aligned}$$

which establishes a natural bijection $E(FF'c, e) \simeq C(c, G'Gd)$, which means $FF' \vdash G'G$ by the Hom-set definition of Yoneda.

3 4.4.3

4 EXERCISES