

# Category theory in context

Siddharth Bhat

Monsoon of the second Year of the Plague



# Contents

<b>1</b>	<b>Categories, Functors, Natural transformations</b>	<b>5</b>
1.1	Abstract and concrete categories . . . . .	5
1.2	Duality . . . . .	5
1.2.1	Musing . . . . .	5
1.2.2	Solutions . . . . .	5



## Chapter 1

# Categories, Functors, Natural transformations

### 1.1 Abstract and concrete categories

### 1.2 Duality

#### 1.2.1 Musing

How does one remember  $gk = gl \implies k = l$  and vice versa?

#### 1.2.2 Solutions

**Lemma 1.2.3**  $f : x \rightarrow y$  is an isomorphism iff it defines a bijection  $f_* : C(c, x) \rightarrow C(c, y)$ .

*Proof (f is iso  $\implies$  post composition with f induces bijection):* Let  $f : x \rightarrow y$  be an isomorphism. Thus we have an inverse arrow  $g : y \rightarrow x$  such that  $fg = \text{id}_y$ ,  $gf = \text{id}_x$ . The map:

$$C(c, x) \xrightarrow{f_*} C(c, y) : (\alpha : c \rightarrow x) \mapsto (f\alpha : c \rightarrow y)$$

has a two sided inverse:

$$C(c, y) \xrightarrow{g_*} C(c, x) : (\beta : c \rightarrow y) \mapsto (g\beta : c \rightarrow x)$$

which can be checked as  $g_*(f_*(\alpha)) = g_*(f\alpha) = gf\alpha = \text{id}_x\alpha = \alpha$ , and similarly for  $f_*(g_*(\beta))$ . Hence we are done, as the iso induces a bijection of hom-sets.  $\square$

*Proof (post-composition with f is bijection implies f is iso):* We are given that the post composition by  $f$ ,  $f_* : C(c, x) \rightarrow C(c, y)$  is a bijection. We need to show that  $f$  is an isomorphism, which means that there exists a function  $g$  such that  $fg = \text{id}_y$  and  $gf = \text{id}_x$ . Since post-composition is a bijection for all  $c$ , pick  $c = y$ . This tells us that the post-composition  $f_* : C(y, x) \rightarrow C(y, y)$  is a bijection. Since  $\text{id}_y \in C(y, y)$ ,  $\text{id}_y$  an inverse image  $g \equiv f_*^{-1}(\text{id}_y)$ . [We choose to call this map  $g$ ]. By definition of  $f_*^{-1}$ , we have that  $f_*(f_*^{-1}(\text{id}_y)) = \text{id}_y$ , which means that  $fg = \text{id}_y$ . We also need to show that  $gf = \text{id}_x$ . To show this, consider  $f_*(gf) = fgf = (fg)f = (\text{id}_y)f = f$ . We also have that  $f_*(\text{id}_x) = f\text{id}_x = f$ . Since  $f_*$  is a bijection, we have that  $\text{id}_x = gf$  and we are done.  $\square$

$$\begin{array}{ccc}
 C(y, x) & \xrightarrow{f_*} & C(y, y) \\
 \downarrow \psi & & \downarrow \text{id}_y \\
 g = f_*^{-1}(\text{id}_y) & \xleftarrow{f_*^{-1}} & \text{id}_y \\
 \uparrow f_*^{-1} & & \uparrow \\
 & & f_* \text{ is bijective.}
 \end{array}$$

by defn:

$$\textcircled{a} \quad f_* (f_*^{-1}(\text{id}_y)) = \text{id}_y \Rightarrow f_* g = \text{id}_y$$

$$\begin{aligned}
 \textcircled{b} \quad f_* (g b) &= f_* g b = (f_* g) b = \text{id}_y b = b = f_* \text{id}_x = f_* (\text{id}_x) \\
 f_* (g b) &= f_* (\text{id}_x) \Rightarrow g b = \text{id}_x \\
 &\quad \underbrace{f_* \text{ is injective}}
 \end{aligned}$$

Iso is bijection of hom-sets

**Q 1.2.ii:** Show that  $f : x \rightarrow y$  is split epi iff for all  $c \in C$ , post composition  $f \circ - : C(c, x) \rightarrow C(c, y)$  is a surjection.

*Proof (split epi implies post composition is surjective):* Let  $f : e \rightarrow b$  be split epi, and thus possess a section  $s : b \rightarrow e$  such that  $fs = \text{id}_b$ . We wish to show that post composition  $C(c, e) \xrightarrow{f_*} C(c, b)$  is surjective. So pick any  $g \in C(c, b)$ . Define  $sg \in C(c, e)$ . See:

$$f_*(sg) = fsg = (fs)g = \text{id}_b g = g$$

. Hence, for all  $g \in C(c, b)$  there exists a pre-image under  $f_*$ ,  $sg \in C(c, e)$ . Thus,  $f_*$  is surjective since every element of codomain has a pre-image.  $\square$

*Proof (post composition is surjective implies split epi):* Let  $f : e \rightarrow b$  be a morphism such that for all  $c \in C$ , we have  $C(c, e) \xrightarrow{f_*} C(c, b)$  is surjective. We need to show that there exists a morphism  $s : b \rightarrow e$  such that  $fs = \text{id}_b$ . Set  $c = b$ . This gives us a surjection  $C(b, e) \xrightarrow{f_*} C(b, b)$ . Pick an inverse image of  $\text{id}_b \in C(b, b)$ . That is, pick any function  $s \in f_*^{-1}(\text{id}_b)$ . By definition, of  $s$  being in the fiber of  $\text{id}_b$ , we have that  $f_*(s) = fs = \text{id}_b$ . Thus means that we have found a function  $s$  such that  $fs = \text{id}_b$ . Thus we are done.  $\square$

**Q 1.2.iii:** Mono is closed under composition, and if  $gf$  is monic then so is  $f$ .

*Proof (Mono is closed under composition):* Let  $f : x \rightarrow y, g : y \rightarrow z$  be monomorphisms (Recall that  $f$  is a monomorphism iff for any  $\alpha, \beta$ , if  $f\alpha = f\beta$  then  $\alpha = \beta$ ). We are to show that  $gf : x \rightarrow z$  is monic. Consider this diagram which shows that  $gfk = gfl$  for arbitrary  $k, l : a \rightarrow x$ . We wish to show that  $k = l$ .

$$\begin{array}{ccccc}
 a & \xrightarrow{k} & x & \xrightarrow{f} & y & \xrightarrow{g} & z \\
 a & \xrightarrow{l} & x & \xrightarrow{f} & y & \xrightarrow{g} & z
 \end{array}$$

Since  $g$  is mono, we can cancel it from  $gfk = gfl$ , giving us  $fk = fl$ . Since  $f$  is mono, we can once again cancel it, giving us  $k = l$  as desired. Hence, we are done.  $\square$

*Proof (If  $gf$  is monic then so is  $f$ ):* Let us assume that  $fk = fl$  for arbitrary  $l$ . We wish to show that  $k = l$ . We show this by applying  $g$ , giving us  $fk = fl \Rightarrow gfk = gfl$ . As  $gf$  is monic, we can cancel, giving us  $gfk = gfl \Rightarrow k = l$ .  $\square$

**Q 1.2.iv** What are monomorphisms in category of fields?

*Proof* : Claim: All morphisms are monomorphisms in the category of fields. Let  $f : K \rightarrow L$  be an arbitrary field morphism. Consider the kernel of  $f$ . It can either be  $\{0\}$  or  $K$ , since those are the only two ideals of  $K$ . However, the kernel can't be  $K$ , since that would send 1 to 0 which is an illegal ring map. Thus, the map  $f$  has trivial kernel, therefore is an injection, is left-cancellable, is a monomorphism.  $\square$

**Q 1.2.v** Show that the ring map  $i : \mathbb{Z} \rightarrow \mathbb{Q}$  is both monic and epic but not iso.

*Proof*  $i$  is not iso: No ring map  $i : \mathbb{Z} \rightarrow \mathbb{Q}$  can be iso since the rings are different (eg.  $\mathbb{Q}$  is a field).  $\square$

*Proof*  $i$  is epic: To show that it's epic, we must show that given for arbitrary  $f, g : \mathbb{Q} \rightarrow R$  that  $fi = gi$ :

$$\begin{array}{ccc} \mathbb{Z} & \xrightarrow{i} & \mathbb{Q} & \xrightarrow{f} & R \\ \mathbb{Z} & \xrightarrow{i} & \mathbb{Q} & \xrightarrow{g} & R \end{array}$$

implies that  $f = g$ . Let  $fi : \mathbb{Z} \rightarrow R = gi$ . Then, the functions  $f, g$  are uniquely determined since  $\mathbb{Q}$  is the field of fractions of  $\mathbb{Z}$ , thus a ring map  $\mathbb{Z} \rightarrow R$  extends uniquely to a ring map  $\mathbb{Q} \rightarrow R$ . Let's assume that  $f(i(z)) = g(i(z))$  for all  $z$ , and show that  $f = g$ . Consider arbitrary  $p/q \in \mathbb{Q}$  for  $p, q \in \mathbb{Z}$ . Let's evaluate:

$$f(p/q) = f(p)f(q)^{-1} = f(i(p)) \cdot f(i(q))^{-1} = g(i(p)) \cdot g(i(q))^{-1} = g(p/q)$$

which shows that  $f(p/q) = g(p/q)$  for all  $p, q$ . Thus, we can extend a ring function defined on the integers to rationals uniquely, hence  $fi = gi \implies f = g$  showing that  $i$  is epic.  $\square$

*Proof*  $i$  is monic: given two arbitrary maps  $k, l : R \rightarrow \mathbb{Z}$ , if  $ik = il$  then we must have  $k = l$ . Given  $ik = il$ , since  $i$  is an injection of  $\mathbb{Z}$  into  $\mathbb{Q}$ , we must have  $k = l$ .

**Q 1.2.vi** Mono + split epi iff iso.

*Proof* Iso is mono + split epi: Iso is both left and right cancellable. Hence it's mono and epi. It splits because the inverse of the iso splits it.  $\square$ .

*Proof* mono + split epi is iso: Let  $f : e \rightarrow b$  be mono (for all  $k, l : p \rightarrow e$ ,  $fk = fl \implies k = l$ ) and split epi (there exists  $s : b \rightarrow e$  such that  $fs : b \rightarrow b = id_b$ ). We need to show it's iso. That is, there exists a  $g : b \rightarrow e$  such that  $fg = id_b$  and  $gf = id_e$ . I claim that  $g \equiv s$ . We already know that  $fg = fs = id_b$  from  $f$  being split epi. We need to check that  $gf = sf = id_e$ . Consider:

$$fsf = (fs)f = id_b f = f = fid_e$$

Hence, we have that  $f(sf) = f(id_e)$ . Since  $f$  is mono, we conclude that  $sf = id_e$ . We are done since we have found a map  $s$  such that  $fs = id_b, sf = id_e$ .

**1.2.vii** Regarding a poset a category, define the supremum of a subcollection, such that the dual gives the infimum. *Proof* : We regard an arrow  $a \rightarrow b$  as witnessing that  $a \leq b$ . First define an upper bound of a set  $O$  to be an object  $u$  such that for all  $o \in O$ , we have  $o \leq u$ . Now, the supremum of  $O$  is the least upper bound of  $O$ . That is,  $s$  is a supremum iff  $s$  is an upper bound, and for all other upper bounds  $t$  of  $O$ , we have that  $s \leq t$ . So we draw a diagram showing upper bounds and suprema:

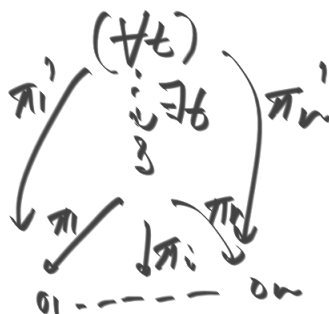
① UPPER BOUND OF  $\Theta$ :



$U$  is an upper bound of  $\Theta$ :

$$\forall o_i \in \Theta, \underline{\underline{o_i \leq U}}$$

② SUPREMUM / LEAST UPPER BOUND  
 $s$  is a LUB:



①  $s$  is an upper bound

② All other upper bounds  $t$  are  $\Rightarrow s$ .

Upper bound and supremum