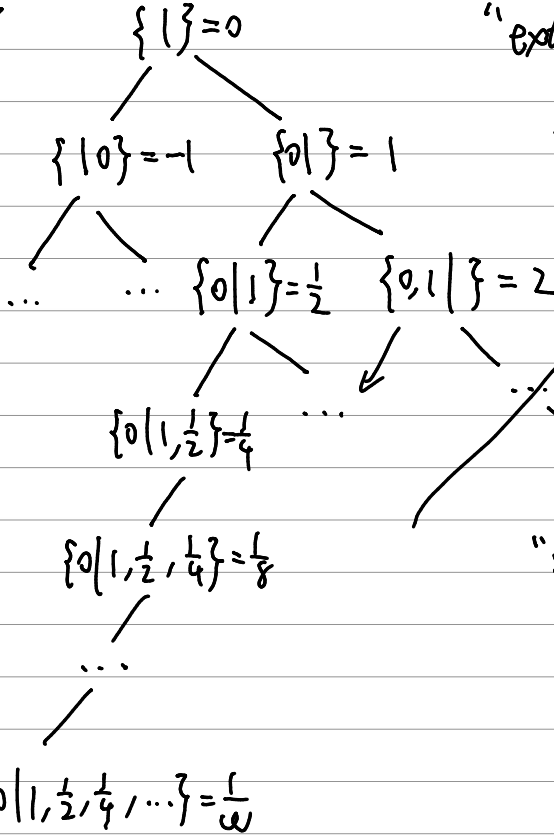


$\frac{f}{g}$



"extension of Dedekind cut"

finite, algorithmic

infinite, non-algorithmic  
countable ordinals

$$\{0, 1, 2, \dots\} = \omega$$

"smallest number greater than all positive integers"

$$\{0, 1, 2, \dots, \omega\} = \omega + 1$$

inverse multiplication of  $w$

(closer to 0 than any other positive rational numbers)

restricting to one unique approximation (infinite countable sequence) to each irrational number  $\Rightarrow$  wrong

$$G = \{G^L | G^R\}$$

$$\Rightarrow G = \{G^L | G^M | G^R\}$$

two independent game states

$$\Rightarrow G = \{G^0 | G^1 | G^2 | \dots\}$$

$$\begin{array}{ccccccc}
 & & 0 & & & & \\
 & -1 & & 1 & & & \\
 -2 & -\frac{1}{2} & & \frac{1}{2} & 2 & & \\
 & -\frac{1}{4} & & \frac{1}{4} & \frac{3}{4} & & \\
 & & -\frac{1}{8} & & & & \\
 & & \vdots & & & & 
 \end{array}$$

$$\frac{2^0}{2^1} \quad \frac{2^1+2^0}{2^2}$$

$$\frac{2}{3} \in \left(\frac{1}{2}, \frac{3}{4}\right) \subset (0, 1)$$

$$\frac{2}{3} = \frac{b_{n-1}2^{n-1} + \dots + b_0 2^0}{2^n}$$

binary search

$$\left\{0, \frac{1}{2}, \frac{5}{8}, \dots \mid 1, \frac{3}{4}, \frac{11}{16}, \dots\right\} = \frac{2}{3}$$

$$\{0, 1, 2, \dots \mid \omega, \omega-1, \omega-2, \dots\} = \frac{\omega}{2}$$

$$\{0, 1, 2, \dots \mid \omega, \frac{\omega}{2}, \frac{\omega}{2^2}, \dots\} = \omega^{\frac{1}{2}}$$

defined by transfinite induction: — if a property holds  $G$ ,  
provided it holds for  $G^L$  and  $G^R$ ,  
then it holds for  $G$ .

$\{ \cdot \mid \cdot \}$  an ordered pair of two sets

$$0 = \{\emptyset \mid \emptyset\} \quad (?)$$

(set, holds for all elements, set)

$$G = \{G^L \mid G^R\}$$

$$G + H =$$

$$H = \{H^L \mid H^R\}$$

$$\{G^L + H, G + H^L \mid G^R + H, G + H^R\}$$

$$\begin{array}{ccc}
 \{-1 \mid 0\} & + & \{0 \mid 1\} \\
 -\frac{1}{2} & & \frac{1}{2}
 \end{array}
 = \left\{ -1 + \frac{1}{2}, -\frac{1}{2} + 0 \mid 0 + \frac{1}{2}, -\frac{1}{2} + 1 \right\}$$

$$\begin{matrix} \{ -1 | 0 \} & + & \{ 0 | 1 \} & = & \{ -1 + \frac{1}{2}, -\frac{1}{2} + 0 & | & 0 + \frac{1}{2}, -\frac{1}{2} + 1 \} \\ -\frac{1}{2} & & \frac{1}{2} & & \textcircled{1} & \textcircled{2} & \textcircled{3} & \textcircled{4} \end{matrix}$$

$$\textcircled{1} \begin{matrix} \{ 1 | 0 \} & + & \{ 0 | 1 \} & = & \{ \text{" " } + \frac{1}{2}, -1 + 0 & | & 0 + \frac{1}{2}, -1 + 1 \} \\ -1 & & \frac{1}{2} & & \end{matrix}$$

$$\textcircled{2} \begin{matrix} \{ -1 | 0 \} & + & \{ 1 | \} & = & \{ -1 + 0, -\frac{1}{2} + \text{" " } & | & 0 + 0, -\frac{1}{2} + \text{" " } \} \\ -\frac{1}{2} & & 0 & & \end{matrix}$$

$$\textcircled{3} \begin{matrix} \{ 1 | \} & + & \{ 0 | 1 \} & = & \{ \text{" " } + \frac{1}{2}, 0 + 0 & | & \text{" " } + \frac{1}{2}, 0 + 1 \} \\ 0 & & \frac{1}{2} & & = \{ 0 | 1 \} \end{matrix}$$

$$\textcircled{4} \begin{matrix} \{ -1 | 0 \} & + & \{ 0 | \} & = & \{ -1 + 1, -\frac{1}{2} + 0 & | & 0 + 1, -\frac{1}{2} + \text{" " } \} \\ -\frac{1}{2} & & 1 & & = \{ 0, -\frac{1}{2} & | & 1 \} \end{matrix}$$

$$\text{" " } + x = \text{" " }$$

reduction?

Figure

1>

$$2> \{ \text{" "}, x | y \} = \{ x | y \}$$

$$3> \{ x | \text{" "}, y \} = \{ x | y \}$$

$$4> \{ -1 | 1 \} = \{ 1 | \}$$

$$\{ -2 | 2 \} = \{ 1 | \}$$

⋮

$$\textcircled{2}.\textcircled{3} \begin{matrix} \{ 1 | \} & + & \{ 1 | \} & = & \{ \text{" " } + 0, 0 + \text{" " } & | & \text{" " } + 0, 0 + \text{" " } \} \\ 0 & & 0 & & = \{ 1 | \} \end{matrix}$$

$$\textcircled{2}.\textcircled{1} \begin{matrix} \{ 1 | 0 \} & + & \{ 1 | \} & = & \{ \text{" " } + 0, -1 + \text{" " } & | & 0 + 0, -1 + \text{" " } \} \\ -1 & & 0 & & = \{ 1 | 0 \} \end{matrix}$$

$$\textcircled{2} \begin{array}{c} \{ -1 | 0 \} \\ -\frac{1}{2} \end{array} + \begin{array}{c} \{ 1 \} \\ 0 \end{array} = \begin{array}{c} \{ -1 + 0, -\frac{1}{2} + " " | 0 + 0, -\frac{1}{2} + " " \} \\ -\frac{1}{2} \end{array} = \begin{array}{c} \{ -1 | 0 \} \\ -\frac{1}{2} \end{array}$$

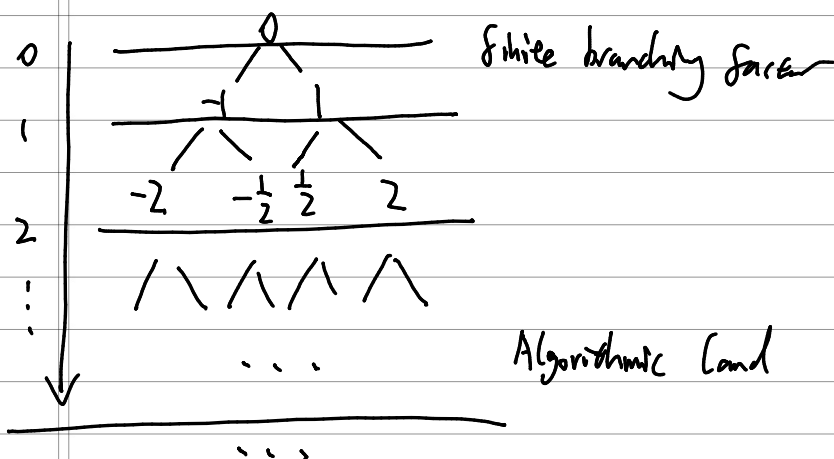
$$\textcircled{1} \cdot \textcircled{4} \begin{array}{c} \{ | 0 \} \\ -1 \end{array} + \begin{array}{c} \{ 0 | \} \\ 1 \end{array} = \begin{array}{c} \{ " " + 1, -1 + 0 | 0 + 1, -1 + " " \} \\ -1 \end{array} = \begin{array}{c} \{ -1 | 1 \} \\ 1 \end{array} \stackrel{?}{=} \begin{array}{c} \{ 1 \} \\ 0 \end{array} = 0$$

given assumption 4

$$\textcircled{1} \begin{array}{c} \{ | 0 \} \\ -1 \end{array} + \begin{array}{c} \{ 0 | 1 \} \\ \frac{1}{2} \end{array} = \begin{array}{c} \{ " " + \frac{1}{2}, -1 + 0 | 0 + \frac{1}{2}, -1 + 1 \} \\ -1 \end{array} = \begin{array}{c} \{ -1 | \frac{1}{2}, 0 \} \\ -\frac{1}{2} \end{array}$$

$$\underbrace{\begin{array}{c} -\frac{1}{2} \\ \{ -1 | 0 \} \end{array}}$$

enumerate  $\mathbb{R}$  ?



$\frac{1}{\omega}$ ,  $\omega$

↓ a projection into finite basis

human : finite

complexity / information

(measured algorithmically)

continuity : infinite

(no such an algorithm  
that can measure its  
complexity)

all math systems / human knowledge systems

is finite, and have an algorithmic description.

## diagonal proof of uncountability

— algorithmic operations on top of  
an non-algorithmic object.

— give us some false confidence/safety  
because of its algorithmic camouflage.

the same as arithmetic operations on  $\omega$ .

$$\omega = \{0, 1, 2, \dots \mid \}$$

a boundary whose existence cannot be proved  
but "constructed" by imagination.

an unique symbol



not  
clearly  
defined

$$\overline{\omega} + 1 = \{0, 1, 2, \dots, \omega \mid \}$$

$$\frac{1}{\omega} = \{0, \frac{1}{2}, \frac{1}{4}, \dots \mid \} = \{0, \frac{1}{3}, \frac{1}{9}, \dots \mid \} = \dots$$

$$\frac{\omega}{2} = \{0, 1, 2, \dots \mid \overline{\omega}, \overline{\omega-1}, \overline{\omega-2}, \dots \}$$

$$\omega^{\frac{1}{2}} = \{0, 1, 2, \dots \mid \overline{\omega}, \overline{\frac{\omega}{2}}, \overline{\frac{\omega}{4}}, \dots \}$$

$$\sinh(X) = \sum_{n=0}^{\infty} \frac{(-1)^n X^{2n+1}}{(2n+1)!}$$

$$\cosh(X) = \sum_{n=0}^{\infty} \frac{(-1)^n X^{2n}}{(2n)!}$$

1. factor polynomials  
into products of  
linear, quadratic  
factors

2. integration by partial fractions

3. allows solutions of recurrence  
relations, linear differential equations,  
eigen problems

Nature of continuum

Fundamental Theorem of Algebra. (I)

Every polynomial  $p(x) = a_0 + a_1 x + \dots + a_n x^n$

of degree  $n \geq 1$  has a complex zero  $r = a + bi$ ,  
that is  $p(r) = 0$ .

$$a_0, a_1, \dots, a_n \in \mathbb{R}$$

FTA(II)

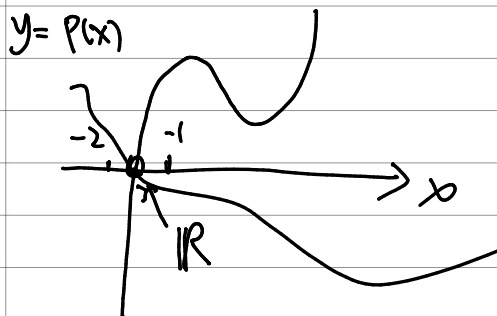
Every polynomial  $p(x) = a_0 + a_1 x + \dots + a_n x^n$ ,  $n \geq 1$ ,

can be factored as a product of linear factors

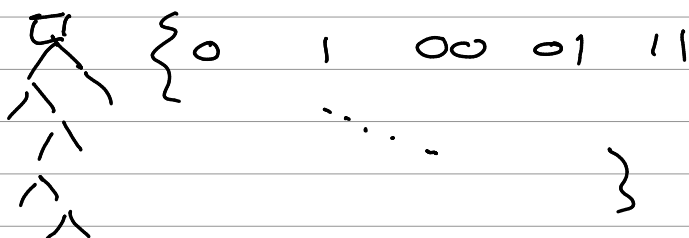
$$p(x) = \lambda (x - r_1)(x - r_2) \dots (x - r_n)$$

for some  $\lambda \in \mathbb{R}$ ,  $r_1, r_2, \dots, r_n \in \mathbb{C}$

$$a_0, a_1, \dots, a_n \in \mathbb{R}$$







"Rationals"

"Countable" ?

~~"Uncountable"~~  
"Countable"