WFQ IS-based Simulation

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WFQ Scheduling Algorithm: Introduction



- WFQ is proposed by Demers, Keshav, and Shenker 1989* as a natural extension of bit-by-bit round robin.
- Given that a scheduler handling K flows is configured with one weight w_k for each flow, the flow k will achieve an average data rate of

$$\frac{w_k}{w_1+w_2+\cdots+w_K}\mu,$$

where μ is the link rate.

^{*} Demers, Alan, Srinivasan Keshav, and Scott Shenker (1989). "Analysis and simulation of a fair queueing algorithm". In: ACM SIGCOMM Computer Communication Review 19.4, pp. 1-12.

WFQ: Bit-by-Bit Round Robin



• Assumptions:

- The scheduler transmits the data bit-by-bit
- For each round, the flow k can transmits w_k bits (suppose w_k is an integer)

• Notations:

- N(t): the number of active flow at time t.
- V(t): the number of rounds made in the round-robin service discipline up to time t. Note that it is set as a continuous function, with the fractional part indicating partially completed rounds.
- Key: Virtual Time V(t)

If N(t) is unchanged during time $[t_1, t_2)$, the increment of the number of rounds is

$$V(t_2) - V(t_1) = \frac{\mu}{\sum_{j \in N(t_1)} w_j} (t_2 - t_1)$$

WFQ Scheduling Algorithm



- Define:
 - $T_{k,n}$: arrival time of packet n of flow k
 - $S_{k,n}$: packet size (in bits) of packet n of flow k
 - $\widetilde{D}_{k,n}$: virtual departure time (in round number) of packet n of flow k
- Relation:

$$\widetilde{D}_{k,n+1} = \max(\widetilde{D}_{k,n}, V(T_{k,n+1})) + \frac{S_{k,n+1}}{w_k}$$

- From Bit-by-Bit Round Robin to Weighted Fair Queue:
 - Select the packet with smallest virtual finishing time.

WFQ Simulation: Problem Setting



- Assumptions:
 - Arrival process: Poisson process with the rate λ_k .
 - Packet size distribution: exponential distribution with mean $1/\mu_k$ bit.
- Target: the p-tile of steady-state response time of class-k jobs

$$Q_k(p) = \inf\{\gamma: P(R_{k,\infty} > \gamma) < 1 - p\}$$

WFQ Simulation: Regenerative Method



- Similarly as the priority queue, we define that the system regenerates whenever a job finds the server idle upon departure.
- Recall that the regenerative method is

$$P(R_{k,\infty} > \gamma) = \frac{E[\sum_{n=1}^{\alpha_k} I(R_{k,n} > \gamma)]}{E[\alpha_k]}$$

• The denominator can be estimated accurately by naïve simulation, but the nominator can be improved by importance sampling.

WFQ Simulation: Importance Sampling



- Switching change of measure:
 - Original densities: f_k^A , f_k^S . IS densities: \tilde{f}_k^A , \tilde{f}_k^S .
 - Stage 1: Starting from time 0, we generate the inter-arrival times and service times according to some IS densities \tilde{f}_k^A , \tilde{f}_k^S and simulate the system dynamics accordingly.
 - Stage 2: the simulator will switch back to the original distribution f_k^A , f_k^S once the event of interest $\{R_{k,n} > \gamma_{\max}\}$ is observed.
 - Define $R'_{k,n}$ as the sum of service times of all jobs who have less virtual departure time than itself $\widetilde{D}_{k,n}$ and the remaining service time of the job in service. And define a stopping time

$$\tau_k = \inf \{ T_{k,n} : R'_{k,n} > \gamma_{\max} \}$$

WFQ Simulation: Importance Sampling



- Notations:
 - $G_l(\cdot)$: the tail probability of inter-arrival times in flow l.
 - $N_l(t)$: the number of job arrivals of class l by time t.

• The likelihood ratio upon time
$$t$$
 can be computed as
$$L(t) = \frac{\prod_{l=1}^{K} \left(\prod_{n=1}^{N_l(t \wedge \tau_k)} \tilde{f}_l^A(A_{l,n}) \tilde{f}_l^S(S_{l,n}) \cdot G_l(H_l)\right)}{\prod_{l=1}^{K} \left(\prod_{n=1}^{N_l(t \wedge \tau_k)} \tilde{f}_l^A(A_{l,n}) \tilde{f}_l^S(S_{l,n}) \cdot \tilde{G}_l(H_l)\right)}$$

where H_l is the age since last arrival of class l at time $t \wedge \tau_k$.

• Importance Sampling for nominator:

•
$$E\left[\sum_{n=1}^{\alpha_k} I(R_{k,n} > \gamma)\right] = \tilde{E}\left[\sum_{n=1}^{\alpha_k} I(R_{k,n} > \gamma)L(E_{k,n})\right]$$
 where $E_{k,n}$ is starting service time.

WFQ Simulation: Cross-Entropy Method



The optimization problem is

$$\max_{f'} E\left[\sum_{n=1}^{\alpha_k} I(R_{k,n} > \gamma) \cdot \log L'(\alpha)\right]$$

• The adaptive cross-entropy method can also be used.

WFQ Simulation: Quantile Estimator



- Sample m_1 cycles under the IS distribution and get response times $\{\tilde{R}_{k,1}, \dots, \tilde{R}_{k,\tilde{\beta}_{m_1}}\}$. Also, sample m_2 cycles under the original distribution.
- The quantile estimator is

$$\widehat{Q}_{m_1,m_2}(p) = \inf\{\gamma : \widehat{P}_{m_1,m_2}(R_{k,\infty} > \gamma) < 1 - p\}$$

$$= \inf\left\{\gamma : \sum_{i=1}^{\widetilde{\beta}_{m_1}} L_{k,i} I(\widetilde{R}_{k,\infty} > \gamma) < \frac{(1-p)m_1}{m_2} \sum_{m=1}^{m_2} \alpha_k^i\right\}$$

• To get $\widehat{Q}(p)$, we can sort $(\widetilde{R}_1, \widetilde{R}_2, \dots, \widetilde{R}_{\widetilde{\beta}_{m_1}})$ in ascending order, and $\widehat{Q}_{m_1, m_2}(p) = \widetilde{R}_{k, (n_p)} \quad and \quad n_p = \min \left\{ n: \sum_{i=n}^{\widetilde{\beta}_{m_1}} L_{k, (i)} < \frac{(1-p)m_1}{m_2} \sum_{m=1}^{m_2} \alpha_k^i \right\}$