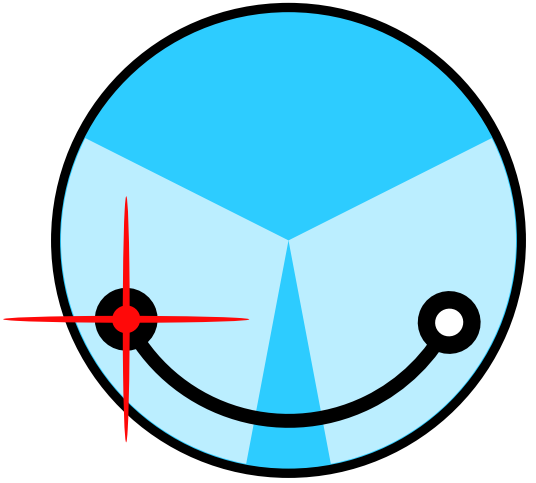
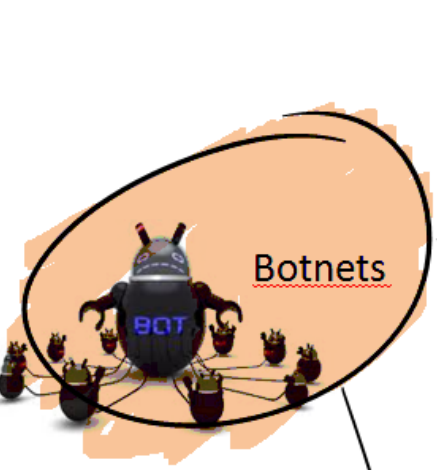


Minemu: Protecting buggy binaries from memory corruption attacks

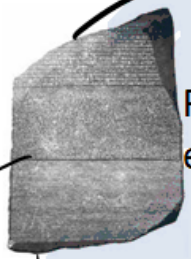
Erik Bosman <erik@minemu.org>

VU VRIJE UNIVERSITEIT AMSTERDAM

Some things we do



Botnets



Reverse engineering

Obfuscation + deobfuscation

Taint tracking

minemu

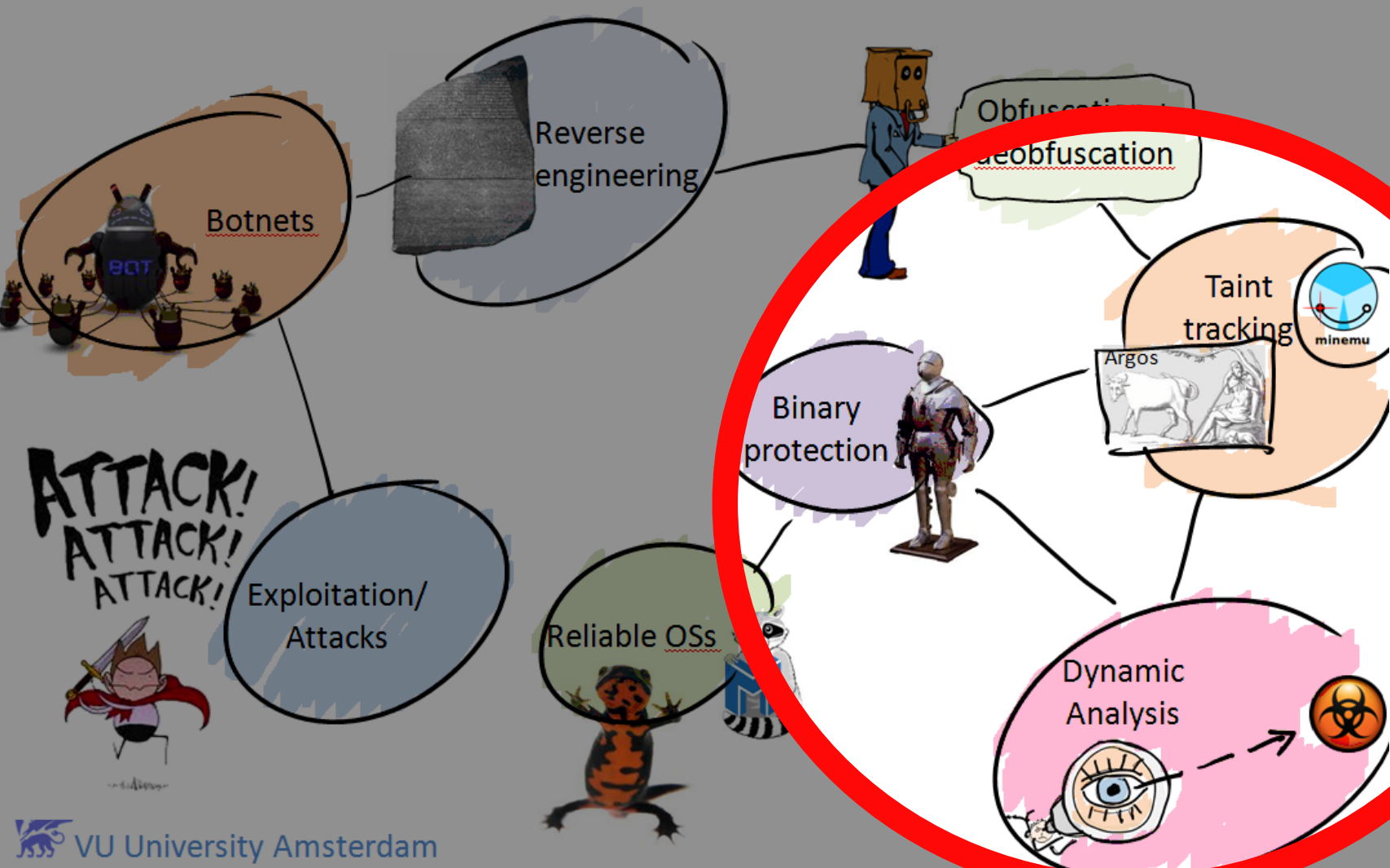
Binary protection

Exploitation/ Attacks

Reliable OSs

Dynamic Analysis

Some things we do



WARNING

THIS PRESENTATION
MAY CONTAIN POINTERS

Programming Languages

type-safe vs. not type-safe

Programming Languages

type-safe vs. not type-safe

Java
Python
Ruby
Javascript

Programming Languages

type-safe	vs.	not type-safe
Java		C
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MEMORY CORRUPTIONS!

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type-safe vs. not type-safe

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MEMORY CORRUPTIONS!

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Programming Languages

type-safe vs. not type-safe

Java
Python
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C++

MEMORY CORRUPTIONS! but not your fault

MEMORY CORRUPTIONS!

The Stack

```
[code]

run(char *name)
{
  char buf[16];

  print("hello ");
  print("world\n")
}
```

The Stack

[code]

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[stack]

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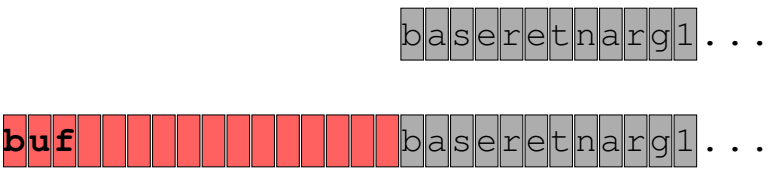

The Stack

```
[address] [code]

 8048751: run(char *name)
          {
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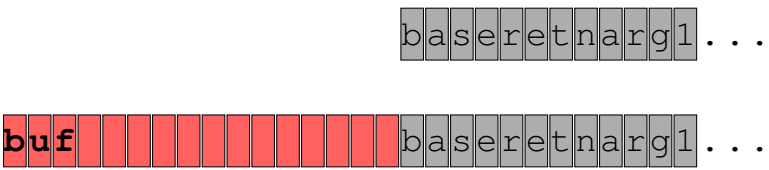
 8048770:   print("hello ");
 8048798:   print("world\n")
          }
```



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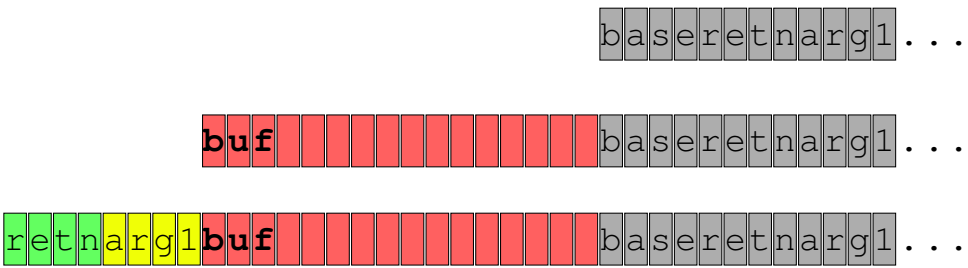

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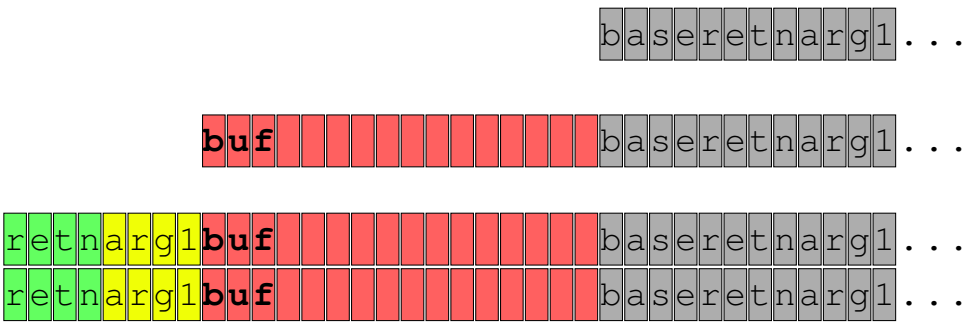
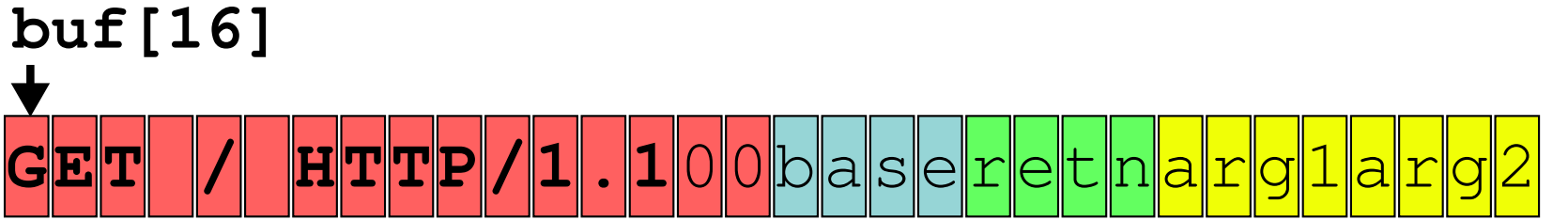


Traditional Stack Smashing



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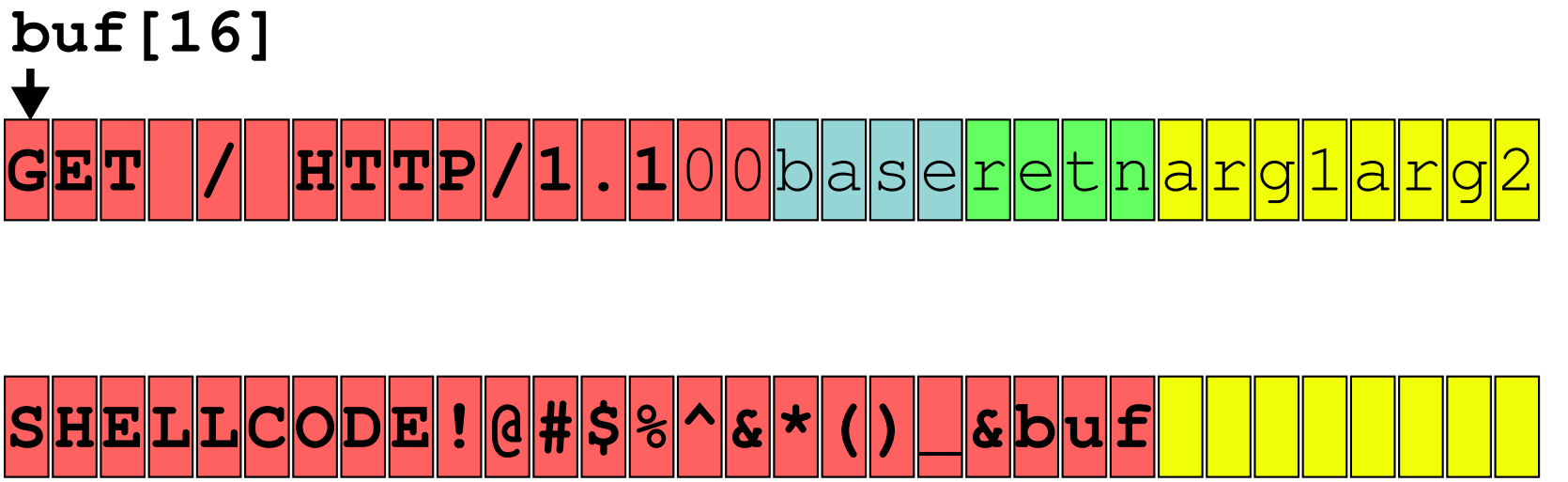
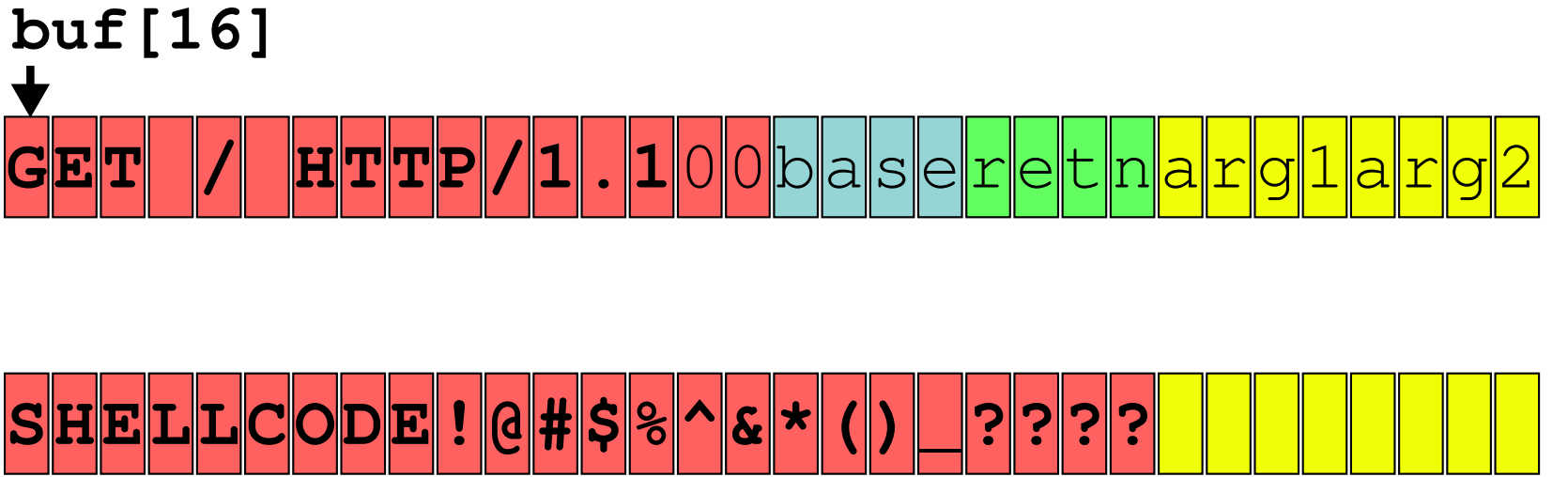
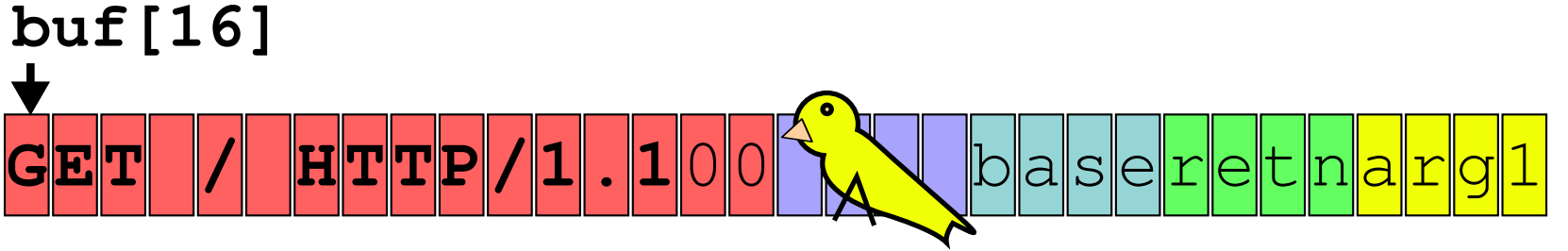


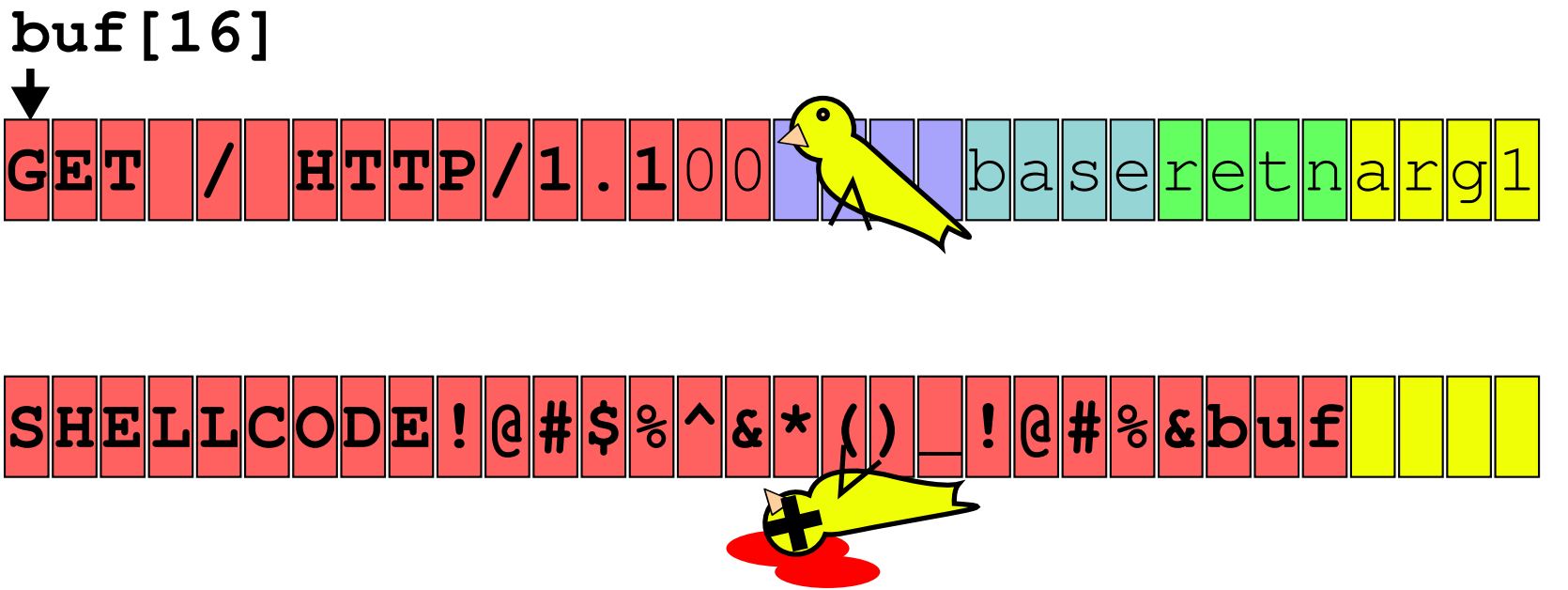
Address Space Layout Randomisation (ASLR)



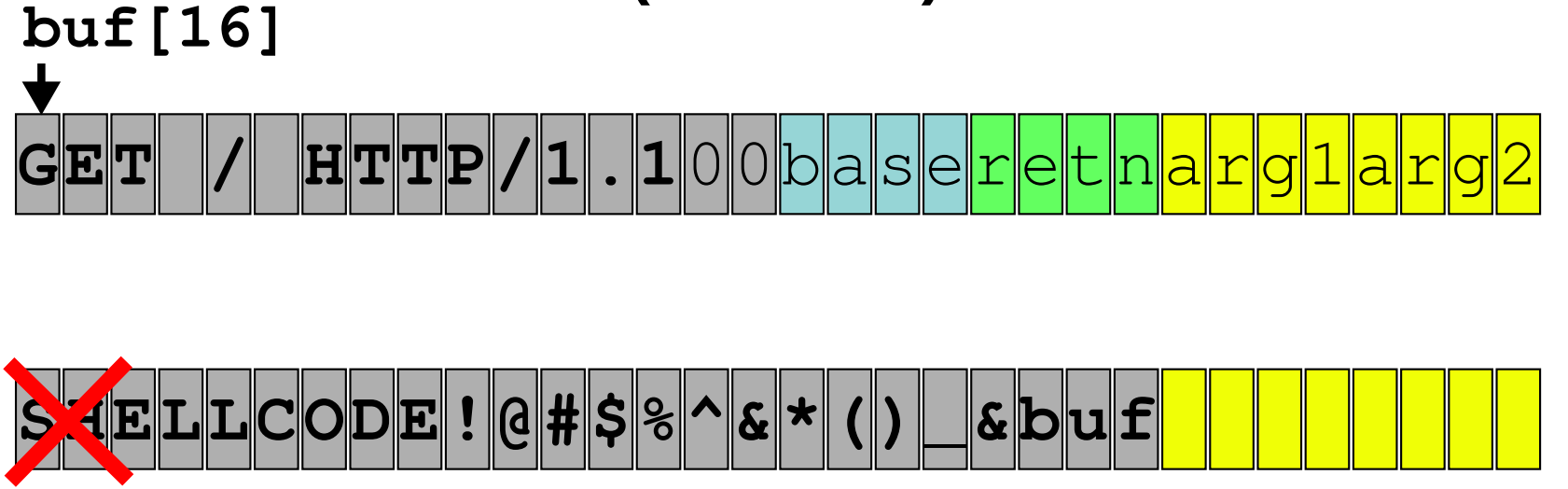
Stack Canaries



Stack Canaries

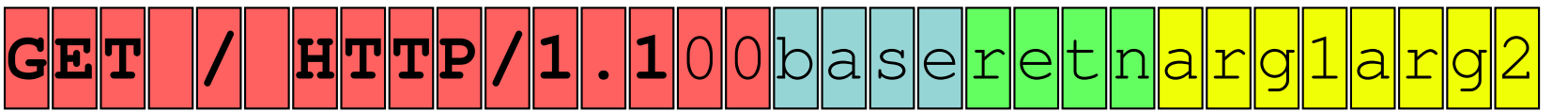


Non-executable data (DEP / NX)



Fortify Source

```
char buf[16];
memcpy(buf, r->buf, r->len);
```



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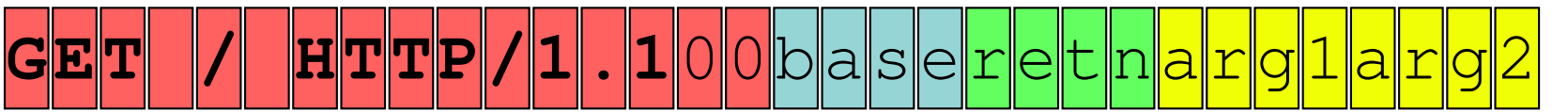
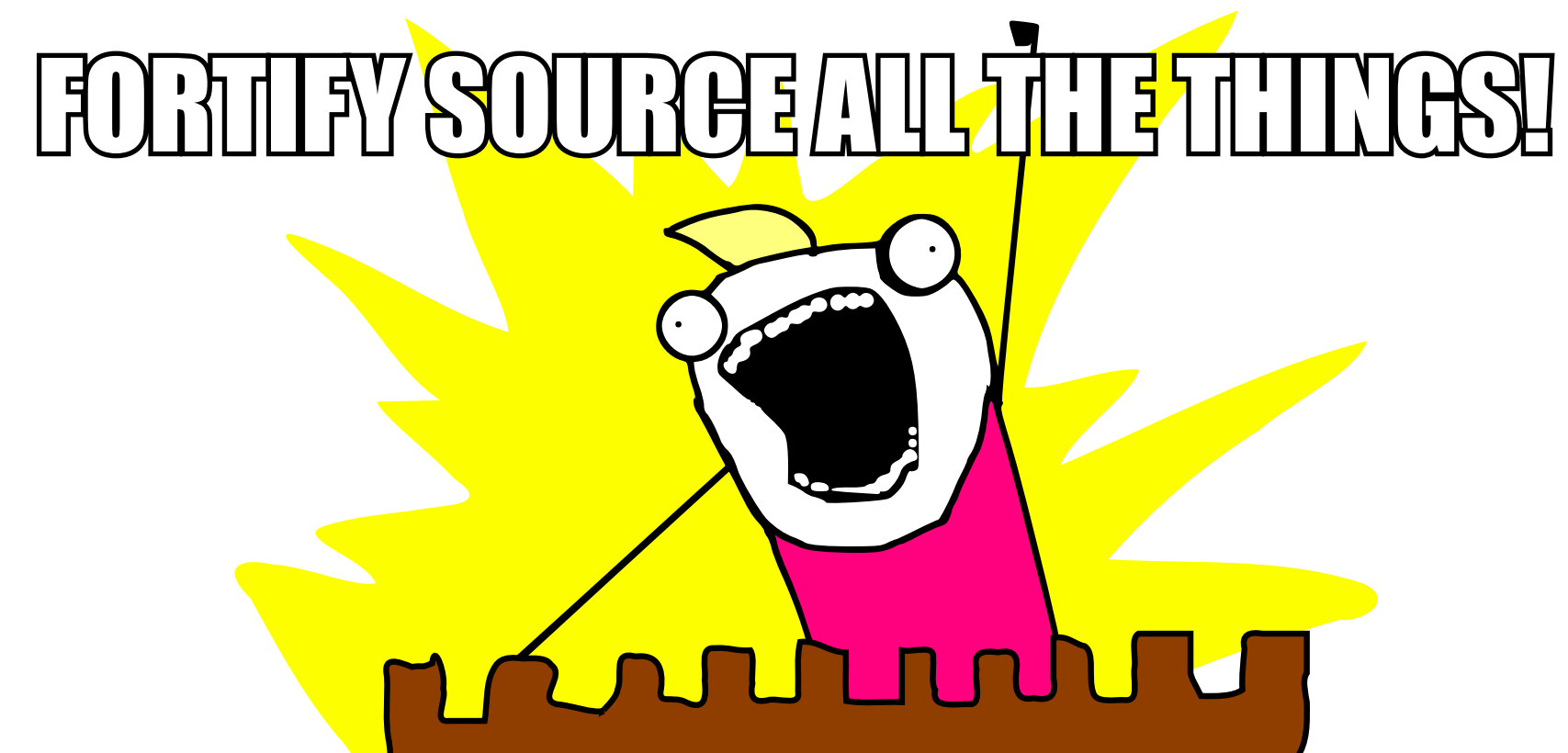


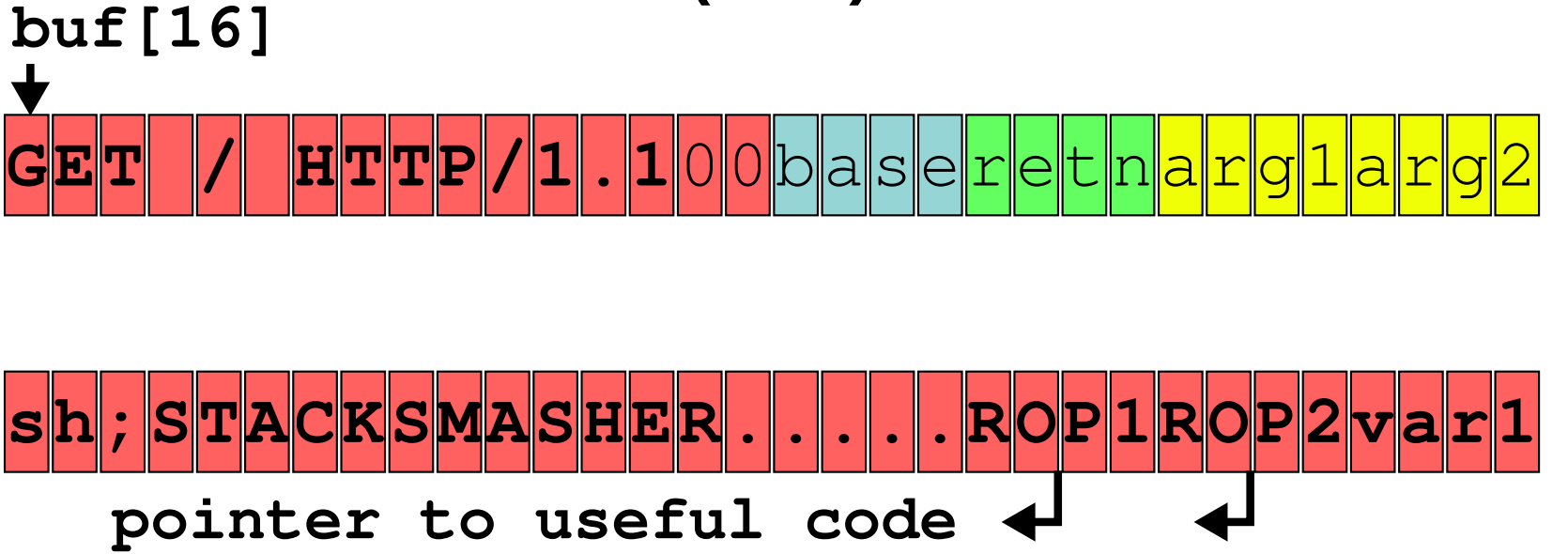

```
char buf[16];
memcpy_chk(buf, r->buf, r->len, 16);
```

```
*** buffer overflow detected ***: /my/fortified/binary terminated
======= Backtrace: =========
/lib/i386-linux-gnu/i686/cmov/libc.so.6(__fortify_fail+0x50)[0xb774a4d0]
/lib/i386-linux-gnu/i686/cmov/libc.so.6(+0xe040a)[0xb774940a]
/my/fortified/binary[0x8048458]
/lib/i386-linux-gnu/i686/cmov/libc.so.6(__libc_start_main+0xe6)[0xb767fe46]
/my/fortified/binary[0x8048371]
======= Memory map: ========
08048000-08049000 r-xp 00000000 fe:00 282465     /my/fortified/binary
08049000-0804a000 rw-p 00000000 fe:00 282465     /my/fortified/binary
08600000-08621000 rw-p 00000000 00:00 0          [heap]
b764b000-b7667000 r-xp 00000000 fe:00 131602     /lib/i386-linux-gnu/libgcc_s.so.1
b7667000-b7668000 rw-p 0001b000 fe:00 131602     /lib/i386-linux-gnu/libgcc_s.so.1
b7668000-b7669000 rw-p 00000000 00:00 0
...

Aborted
```


FORTIFY SOURCE ALL THE THINGS!

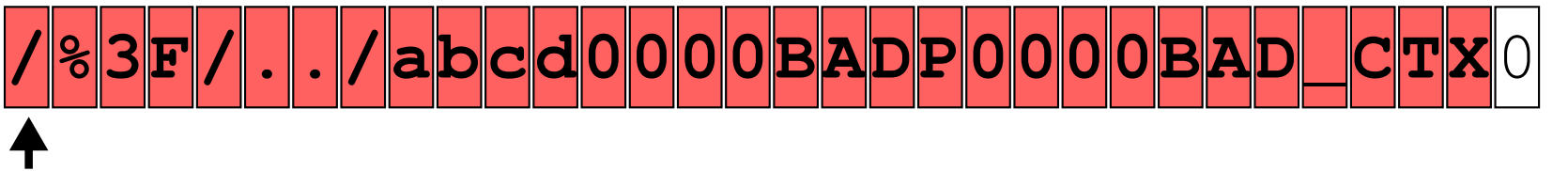
Return Oriented Programming (ROP)



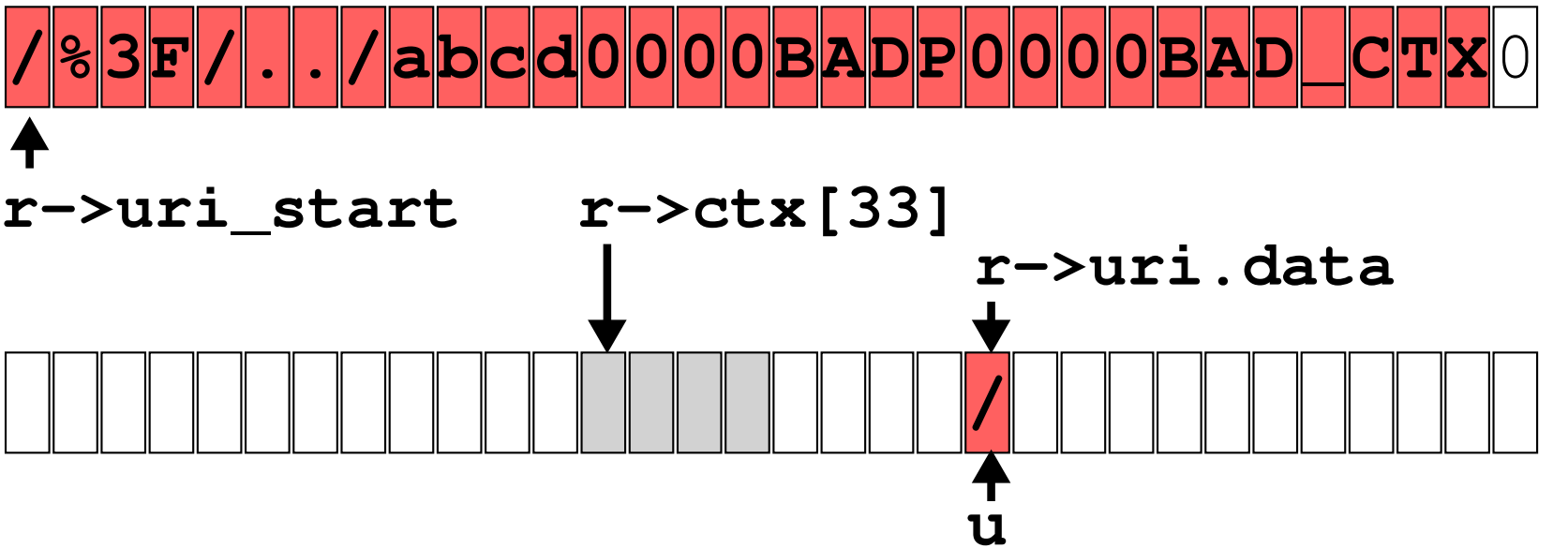
Some exploits still work with all these defense measures.

Example: nginx buffer underrun (CVE-2009-2629)

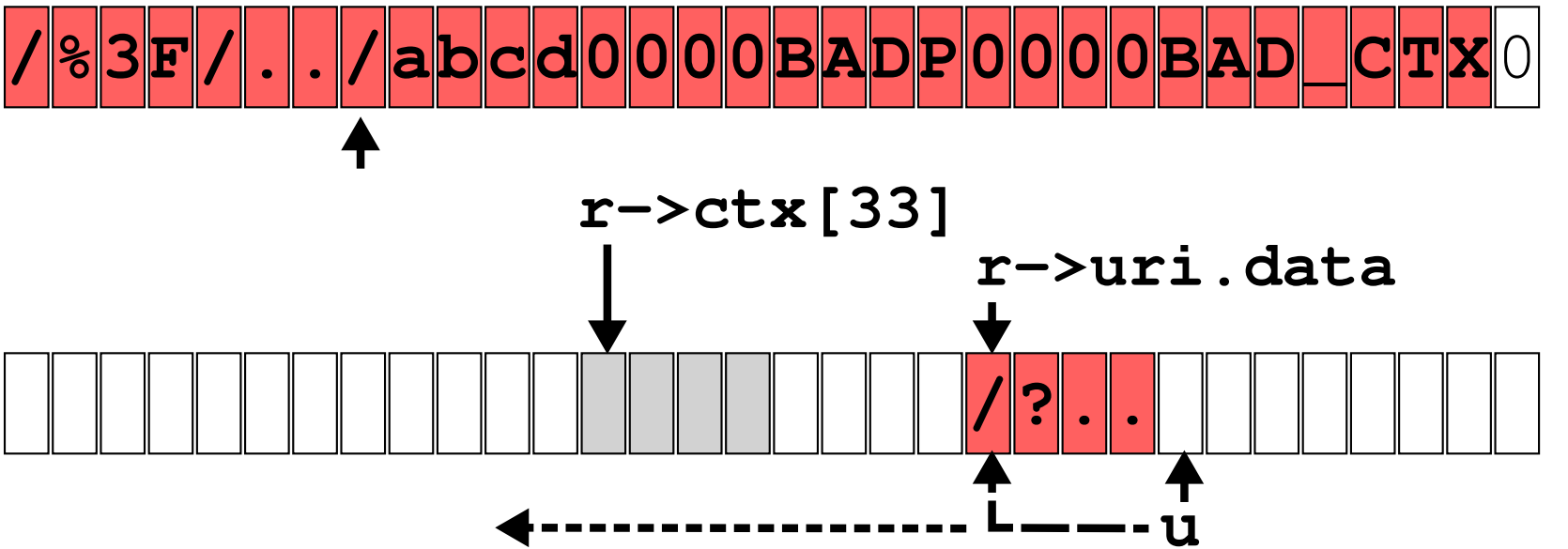
CVE-2009-2629



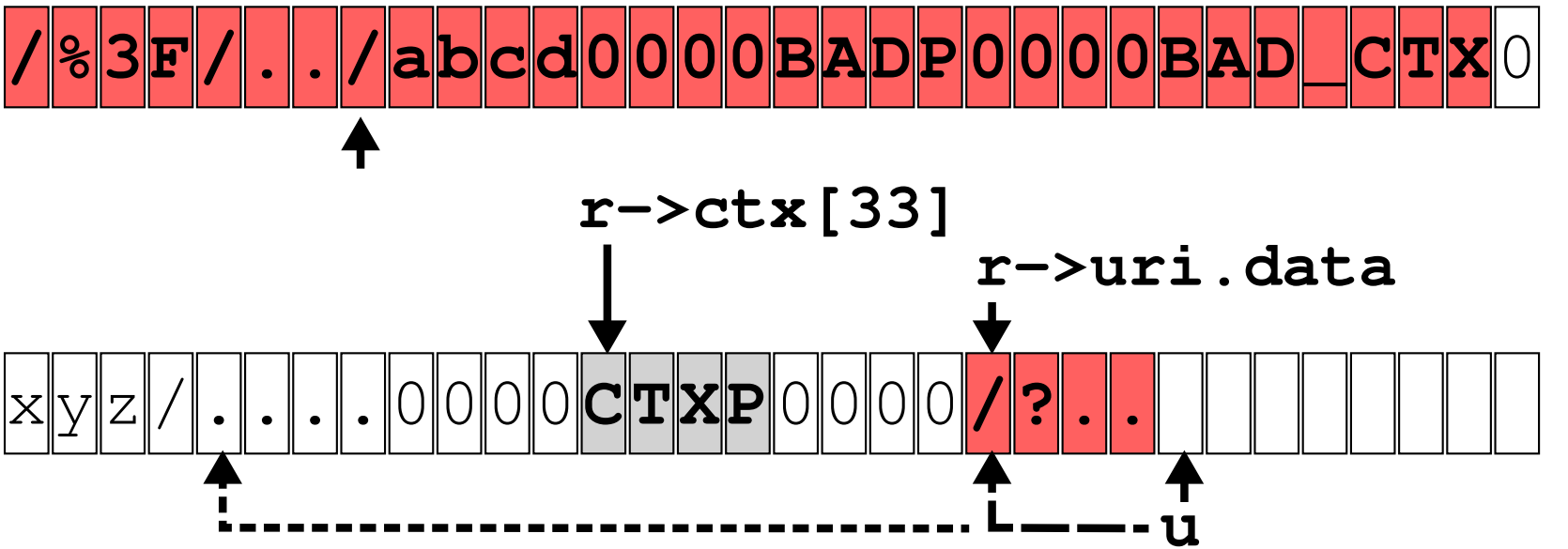
CVE-2009-2629



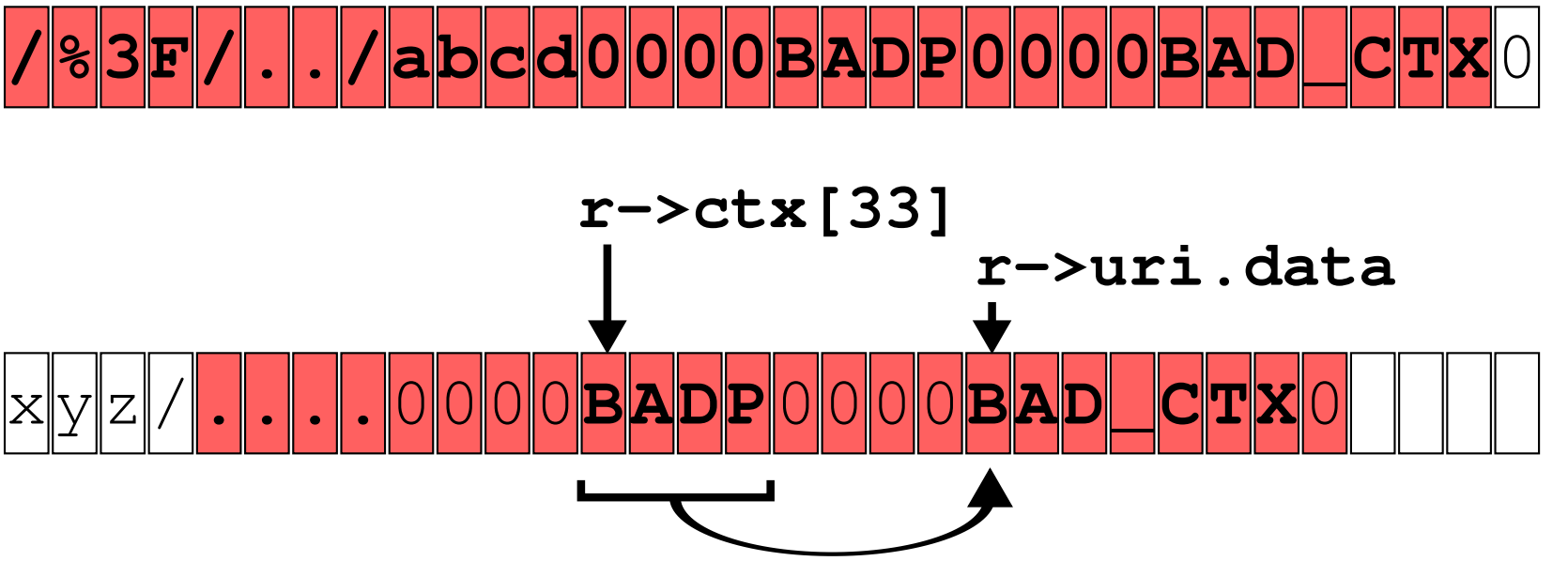
CVE-2009-2629



CVE-2009-2629



CVE-2009-2629



```
typedef struct {
    ngx_buf_t                    *buf;
    ngx_chain_t                  *in;
    ngx_chain_t                  *free;
    ngx_chain_t                  *busy;

    unsigned                      sendfile;
    unsigned                      need_in_memory;
    unsigned                      need_in_temp;

    ngx_pool_t                   *pool;
    ngx_int_t                     allocated;
    ngx_bufs_t                    bufs;
    ngx_buf_tag_t                 tag;

    ngx_output_chain_filter_pt    output_filter;
    void                         *filter_ctx;
} ngx_output_chain_ctx_t;
```

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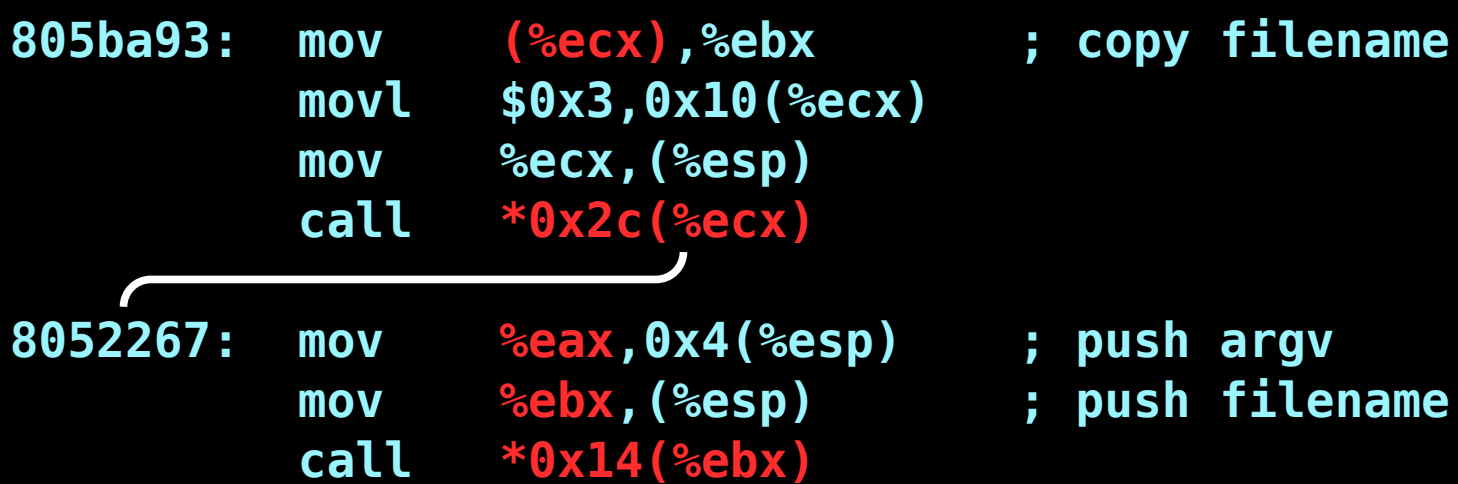
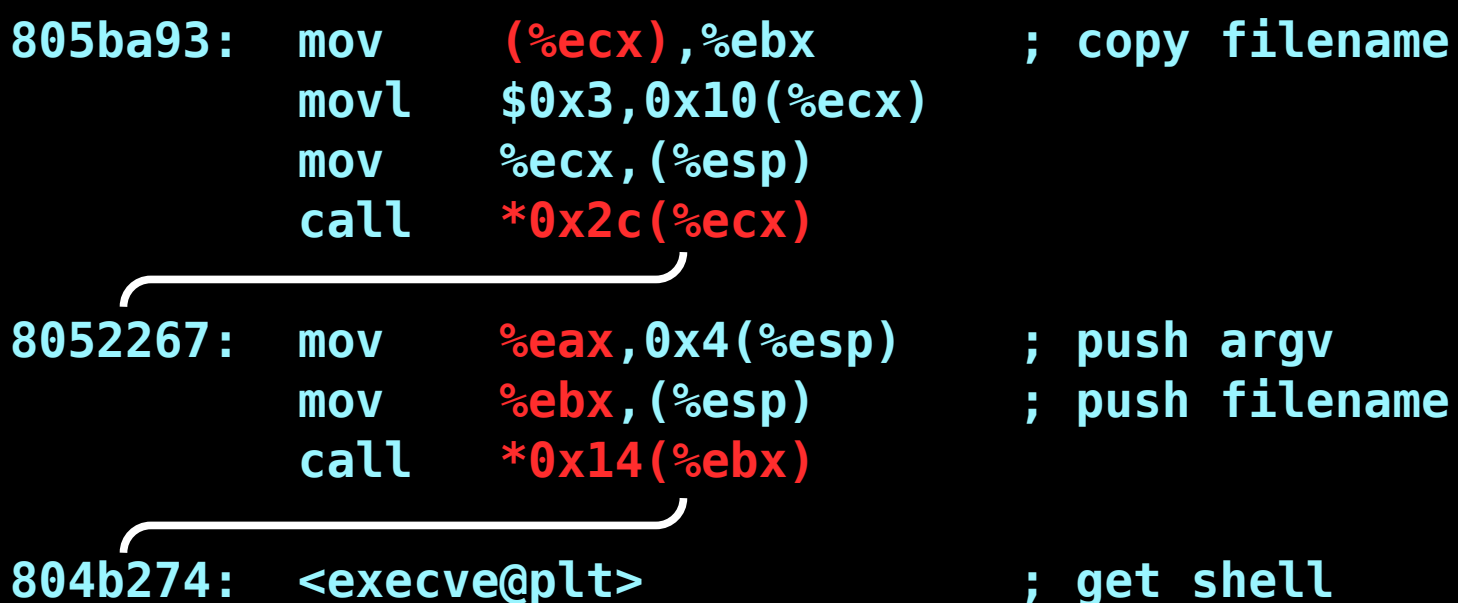
```

function pointer

```
805ba93:  mov    (%ecx),%ebx          ; copy filename
          movl   $0x3,0x10(%ecx)
          mov    %ecx,(%esp)
          call   *0x2c(%ecx)
```


805ba93: mov (%ecx),%ebx ; copy filename
movl $0x3,0x10(%ecx)
mov %ecx,(%esp)
call *0x2c(%ecx)
8052267: mov %eax,0x4(%esp) ; push argv
mov %ebx,(%esp) ; push filename
call *0x14(%ebx)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movl $0x3,0x10(%ecx)
mov %ecx,(%esp)
call *0x2c(%ecx)
8052267: mov %eax,0x4(%esp) ; push argv
mov %ebx,(%esp) ; push filename
call *0x14(%ebx)
804b274: <execve@plt> ; get shell

- defeats address randomisation (through info leak)

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- defeats non-executable data protection

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- no standard copy function (no fortify src protections)

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- no standard copy function (no fortify src protections)

- not return oriented, so stack smash protection does not matter

But the situation is even worse

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- needs to be enabled at compile time, and
  there is a lot of old code out there

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- many packages do not apply these defence mechanisms even today

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- needs to be enabled at compile time, and there is a lot of old code out there

- many packages do not apply these defence mechanisms even today

- implementation flaws

Can we do more?

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>> Non-executable data prevents untrusted data from being run as code

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<< Return oriented programming replaces untrusted code with pointers to original code.

Can we do more?

>> Non-executable data prevents untrusted data from being run as code

<< Return oriented programming replaces untrusted code with pointers to original code.

>> Can we prevent untrusted pointers from being used as jump addresses?

Taint analysis

```
0805be60   00 00 00 00 00 00 00 00  00 00 00 00 00 00 00 00   |................|
0805be70   00 00 00 00 00 00 00 00  00 00 00 00 00 00 00 00   |................|
0805be80   00 00 00 00 02 00 00 00  d8 4b 06 08 a0 2e 05 08   |.........K......|
0805be90   94 be 05 08 78 a0 04 08  ef be ad de a4 be 05 08   |....x...........|
0805bea0   ac be 05 08 2f 62 69 6e  2f 73 68 00 a4 be 05 08   |..../bin/sh.....|
0805beb0   00 00 00 00 53 41 4d 45  54 48 49 4e 47 57 45 44   |....SAMETHINGWED|
0805bec0   4f 45 56 45 52 59 4e 49  47 48 54 50 49 4e 4b 59   |OEVERYNIGHTPINKY|
0805bed0   00 00 00 00 4e 41 52 46  90 be 05 08 ef 1f 05 08   |....NARF........|
0805bee0   ff fa 26 08 ff f0 00 00  00 00 00 00 00 00 00 00   |..&.............|
0805bef0   00 00 00 00 00 00 00 00  00 00 00 00 00 00 00 00   |................|
0805bf00   00 00 00 00 00 00 00 00  00 00 00 00 00 00 00 00   |................|
```

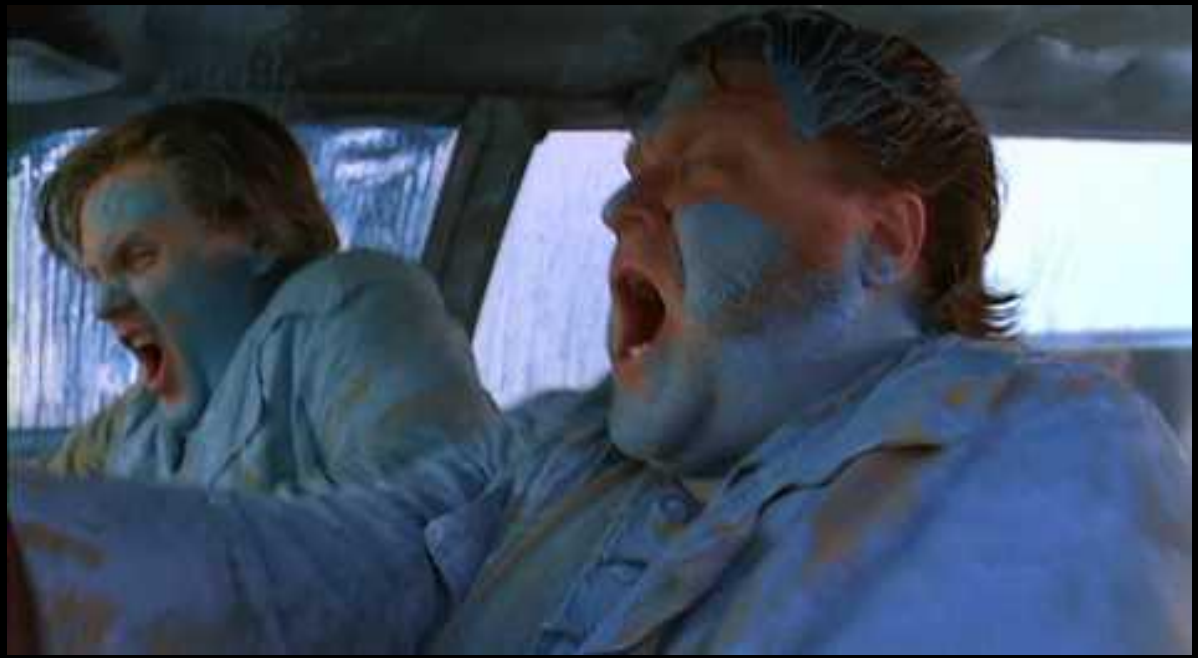
Taint tracking (1/2):

- remember whether data is trusted or not
- untrusted data is 'tainted'
- when data is copied, its taint is copied along
- taint is ORed for arithmetic operations

Taint tracking (2/2):

When the code jumps to an address in memory,
the source of this address is checked for taint.

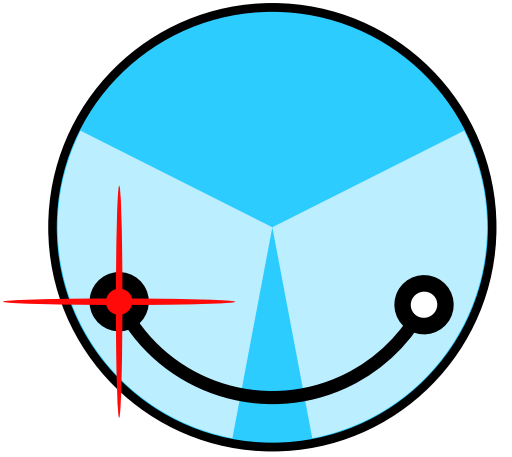
eg.:
- RET
- CALL *%eax
- JMP *0x1c(%ebx)



Taint tracking

useful, but slow as hell

Is this slowness fundamental?



minemu

fast emulator
memory layout
use SSE registers to hold taint

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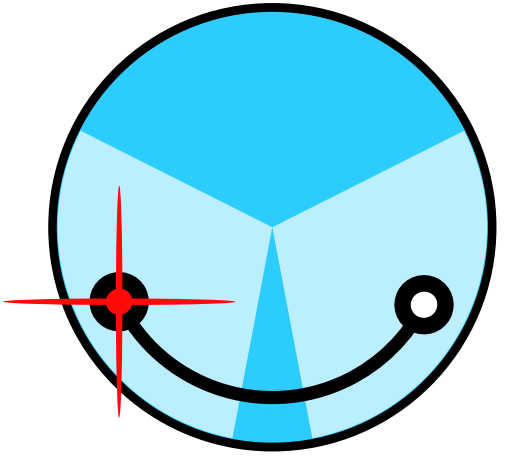

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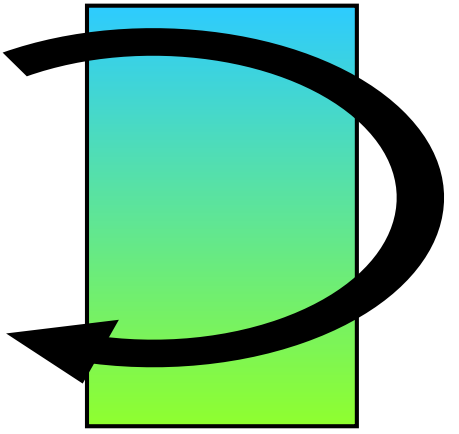
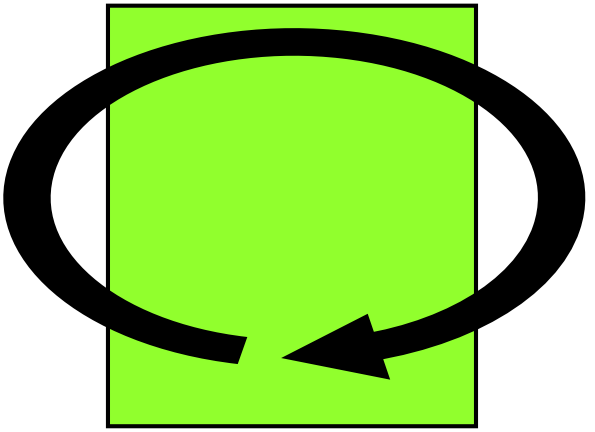
Emulator

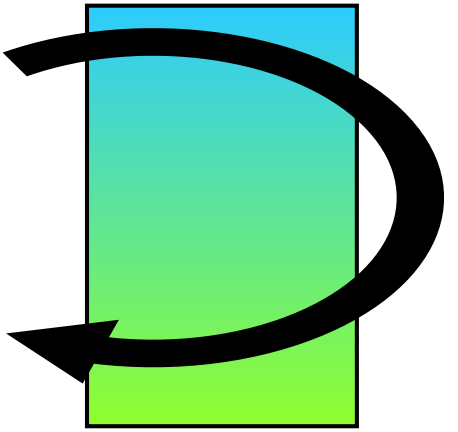


compile jit code

Emulator

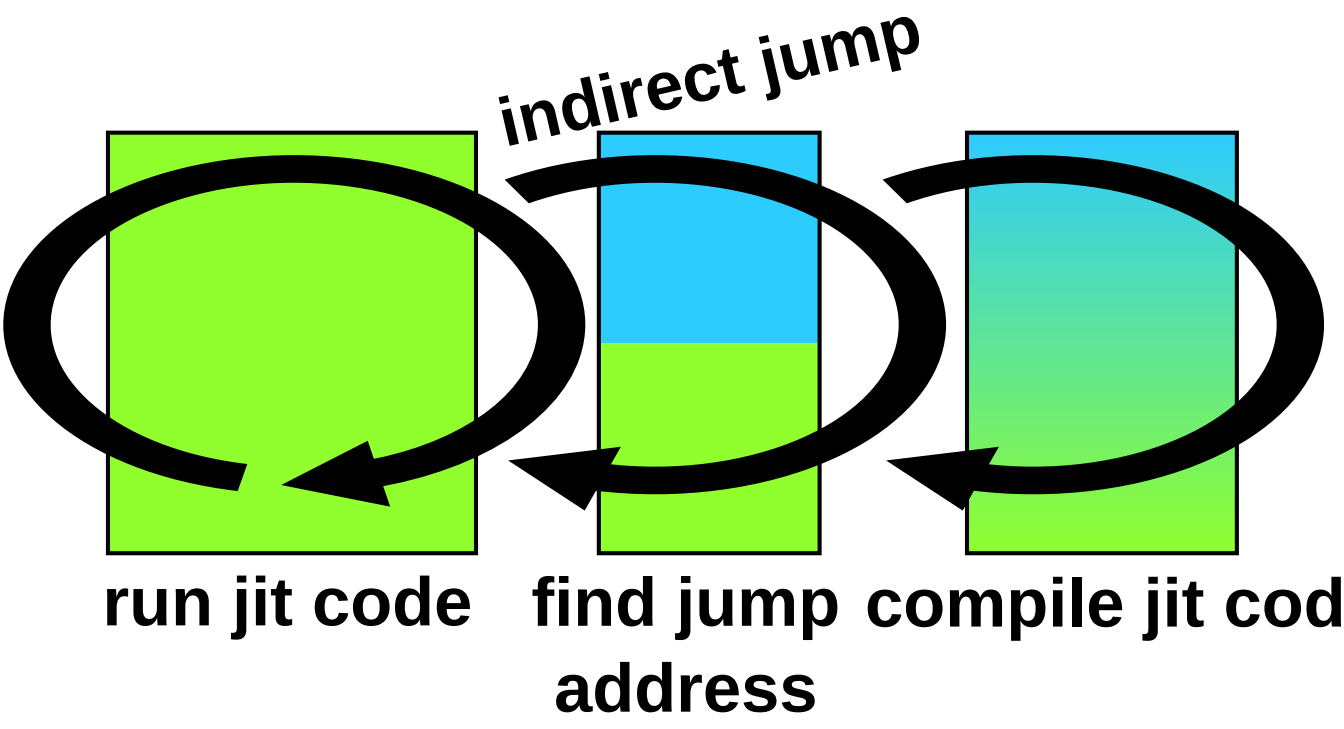


run jit code

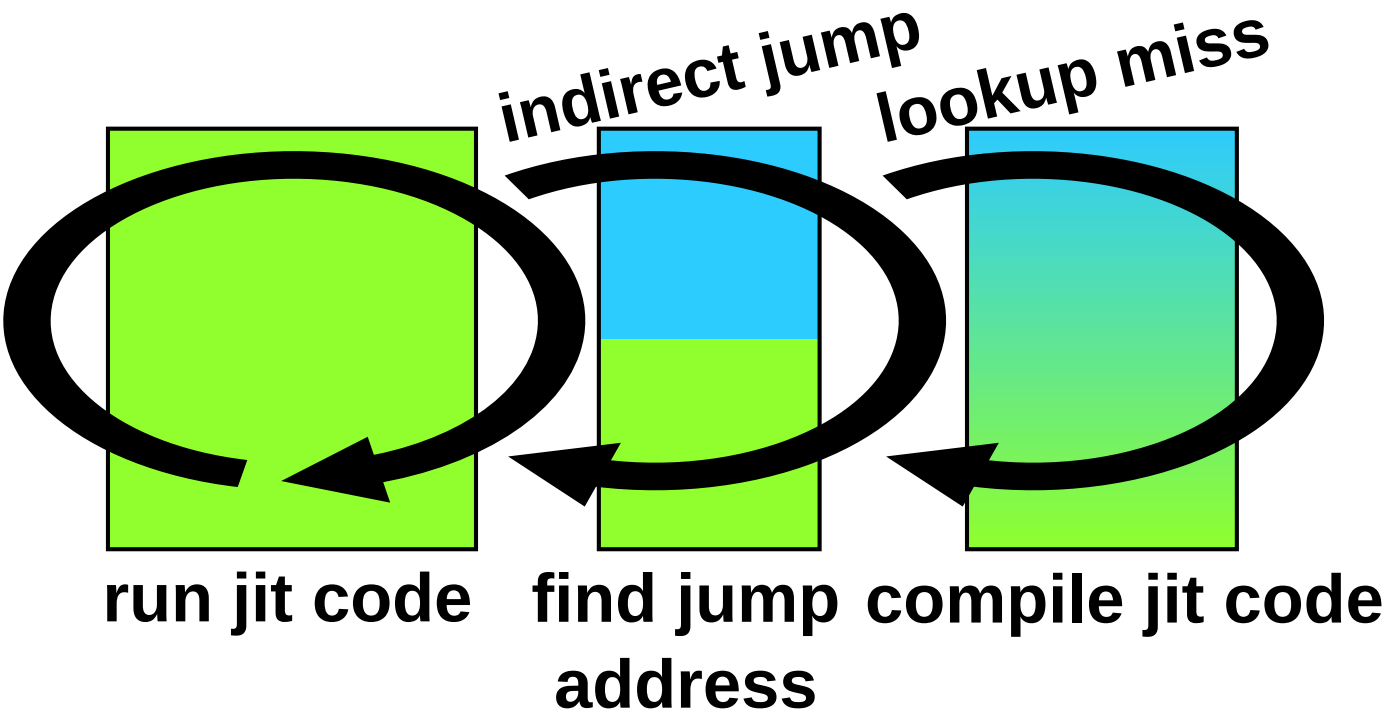


compile jit code

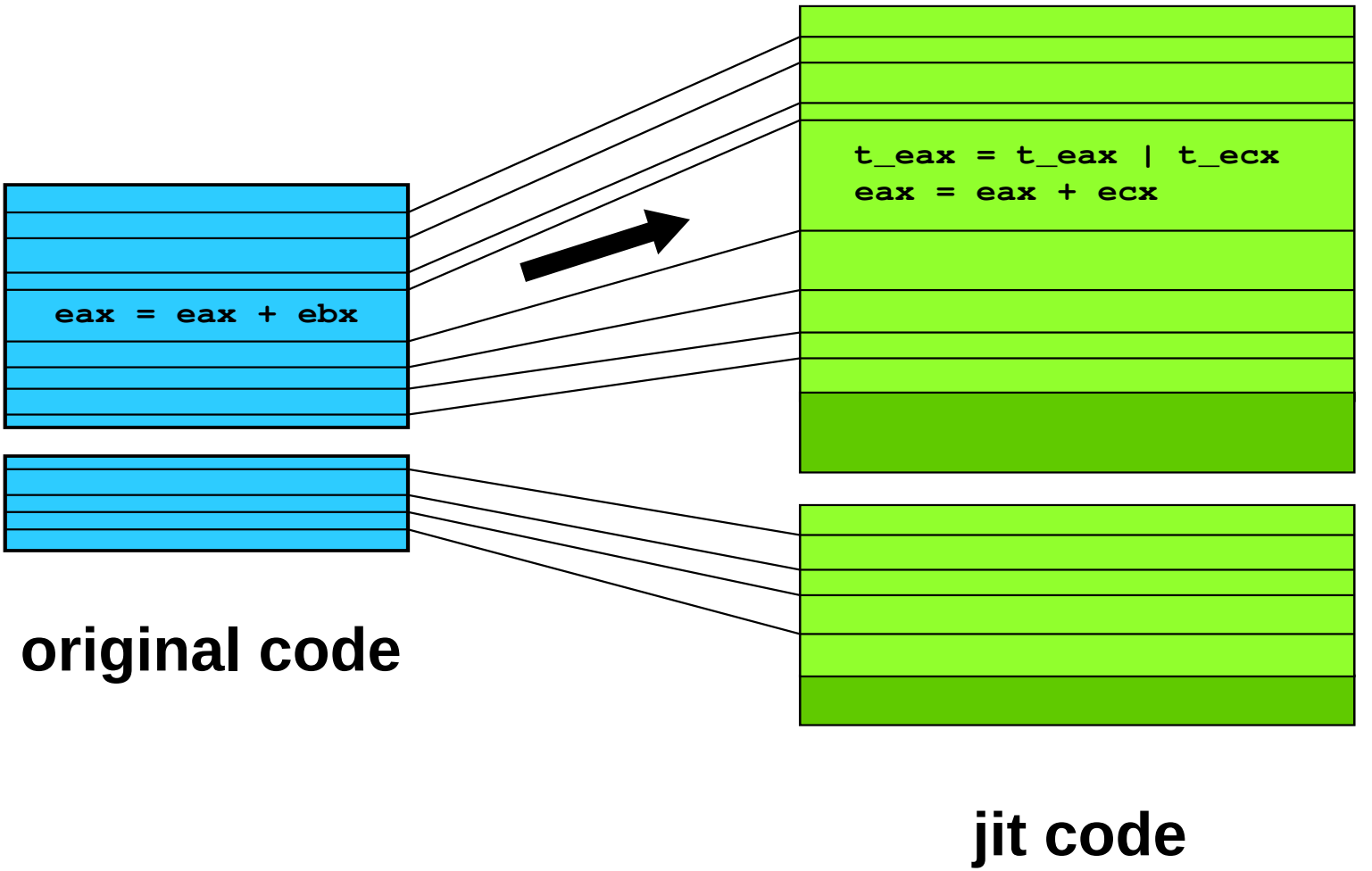
Emulator


indirect jump
run jit code
find jump address
compile jit code

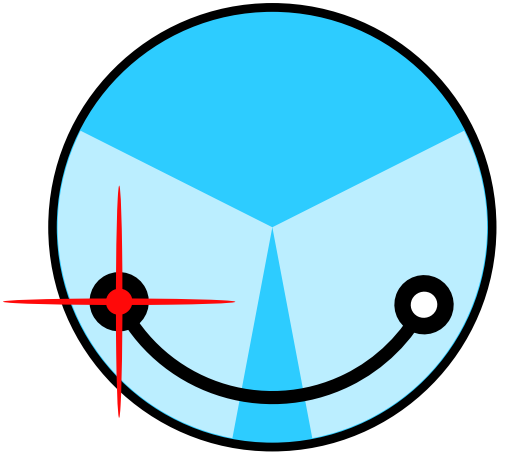
Emulator


indirect jump
lookup miss
run jit code
find jump address
compile jit code

Dynamic instrumentation

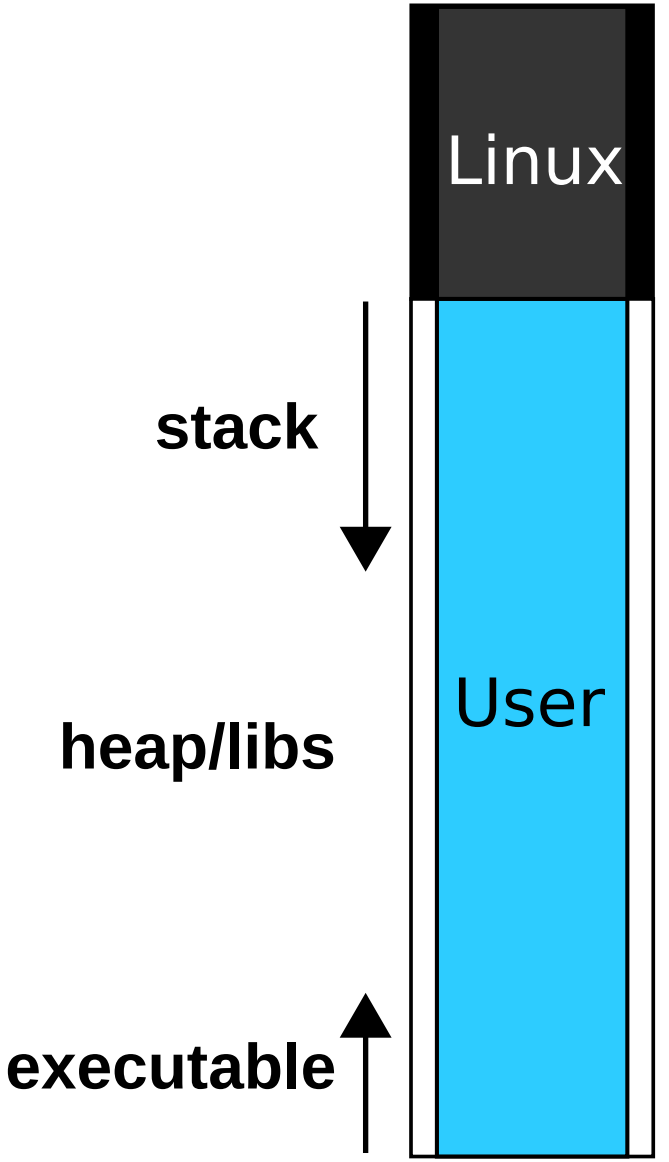


Is this slowness fundamental?

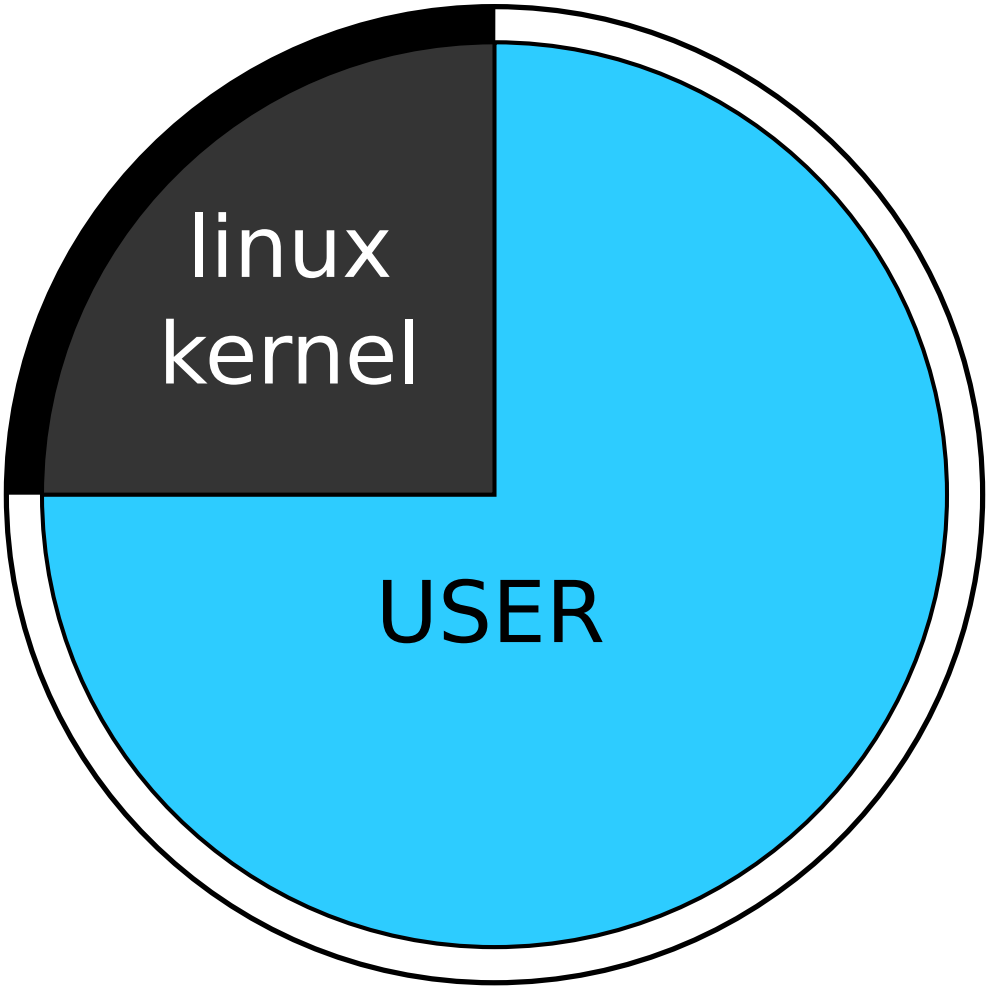


minemu

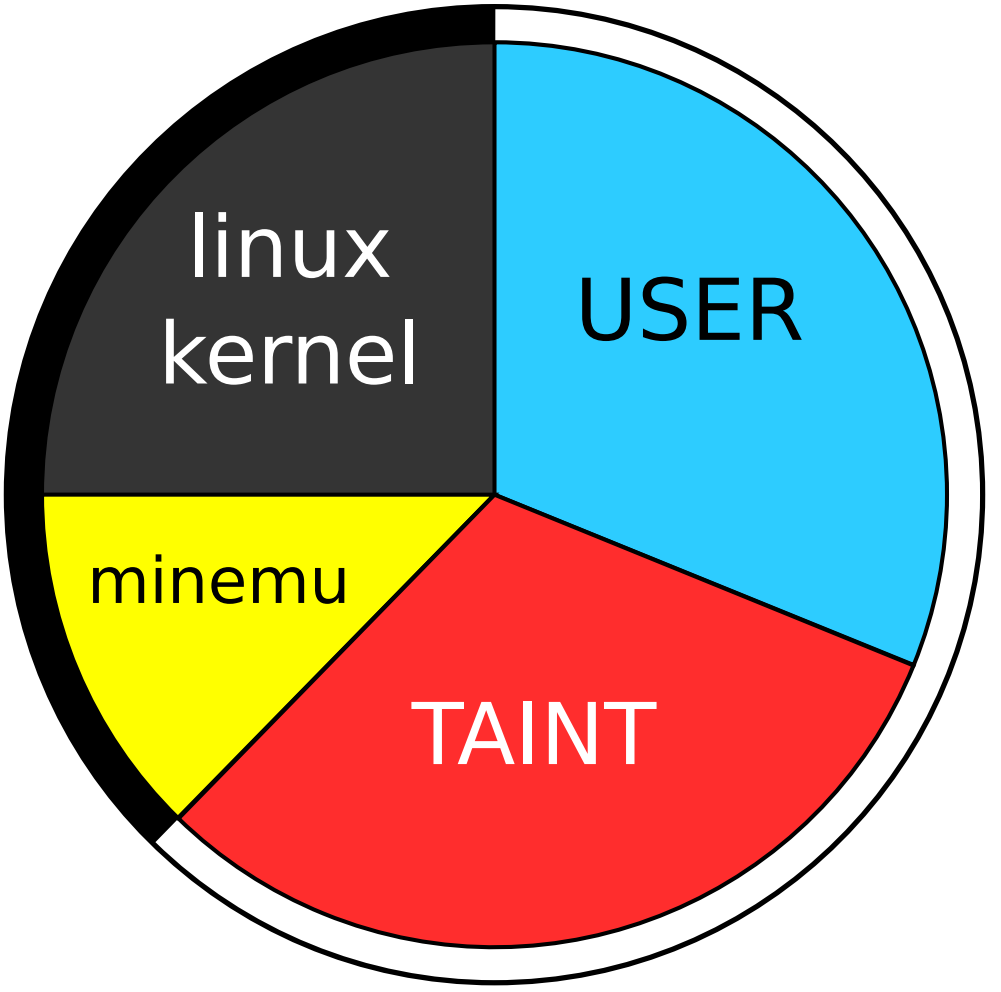
- fast emulator
- memory layout
- use SSE registers to hold taint


Linux
stack
User
heap/libs
executable

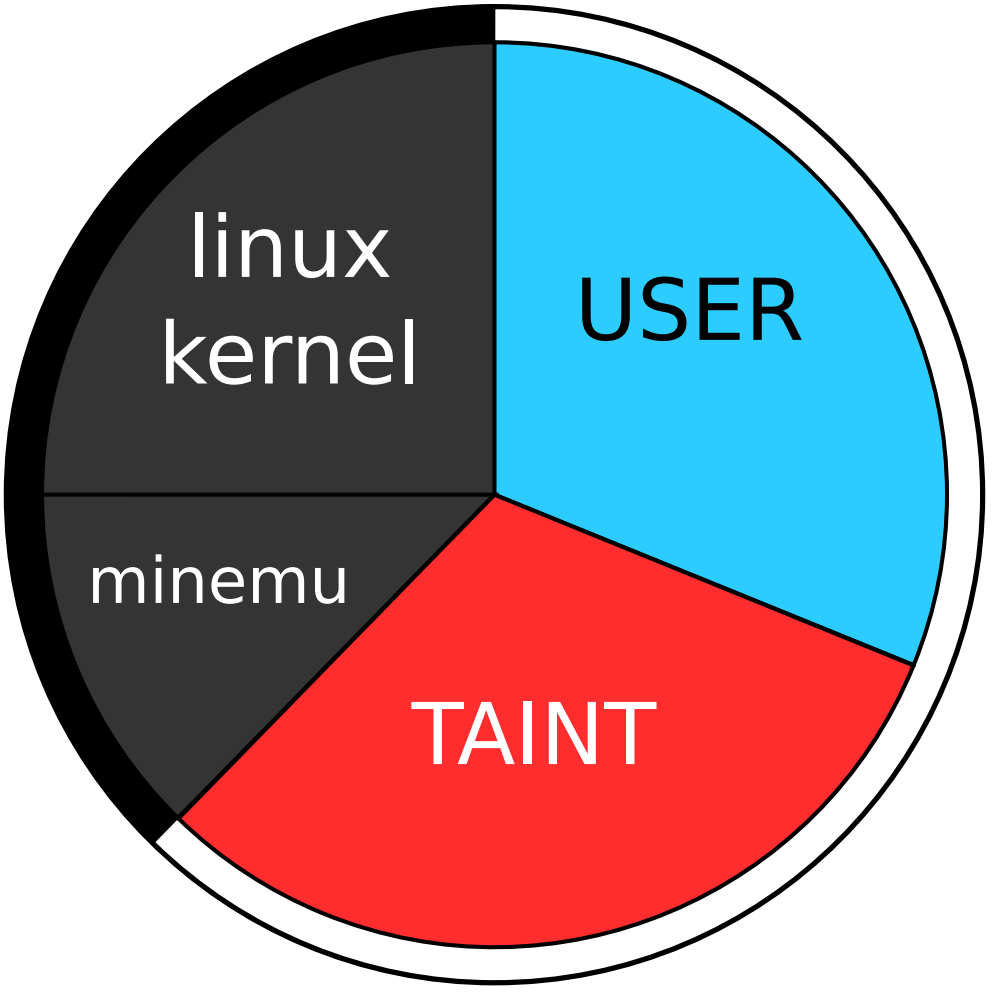
Memory layout (linux)



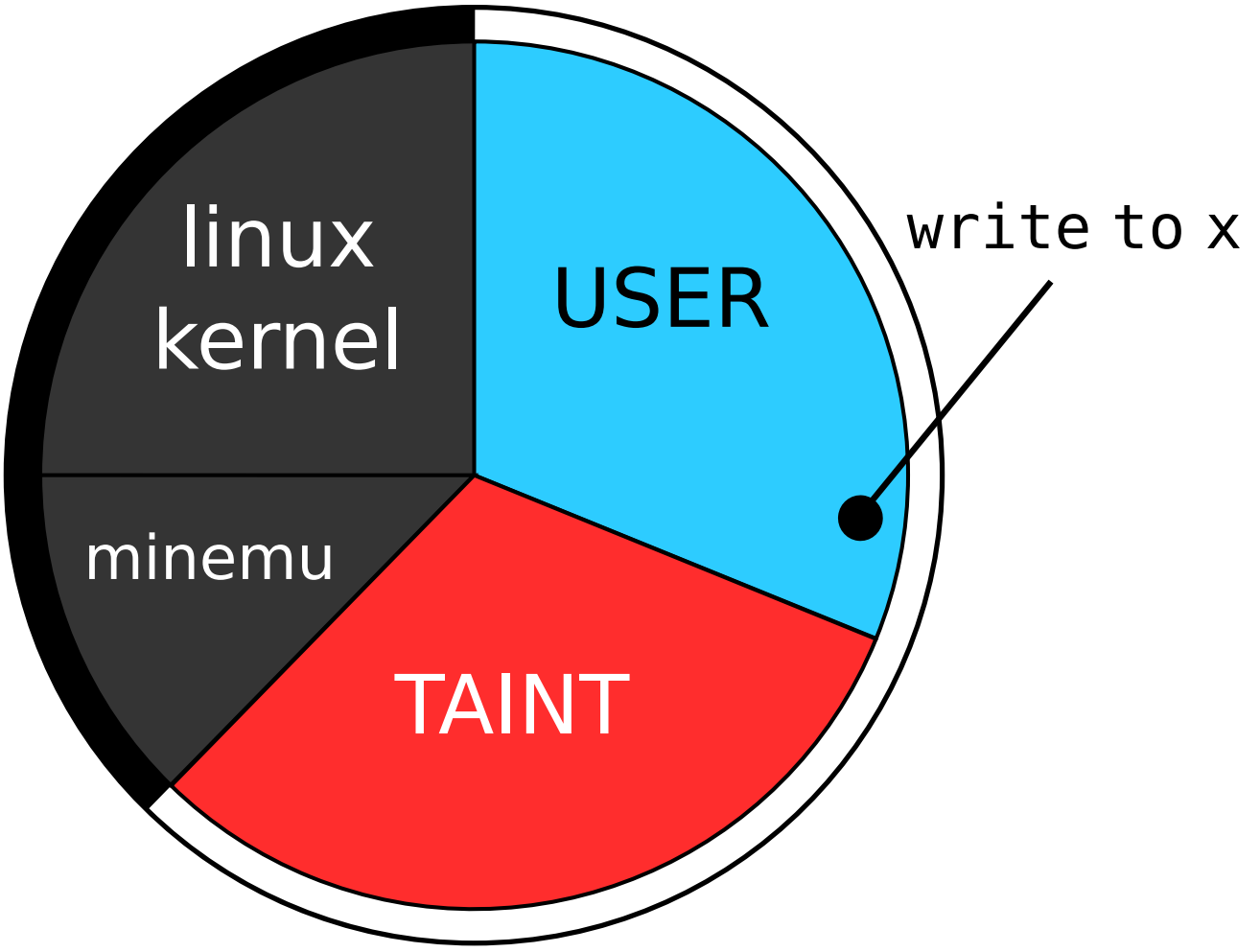
Memory layout (minemu)



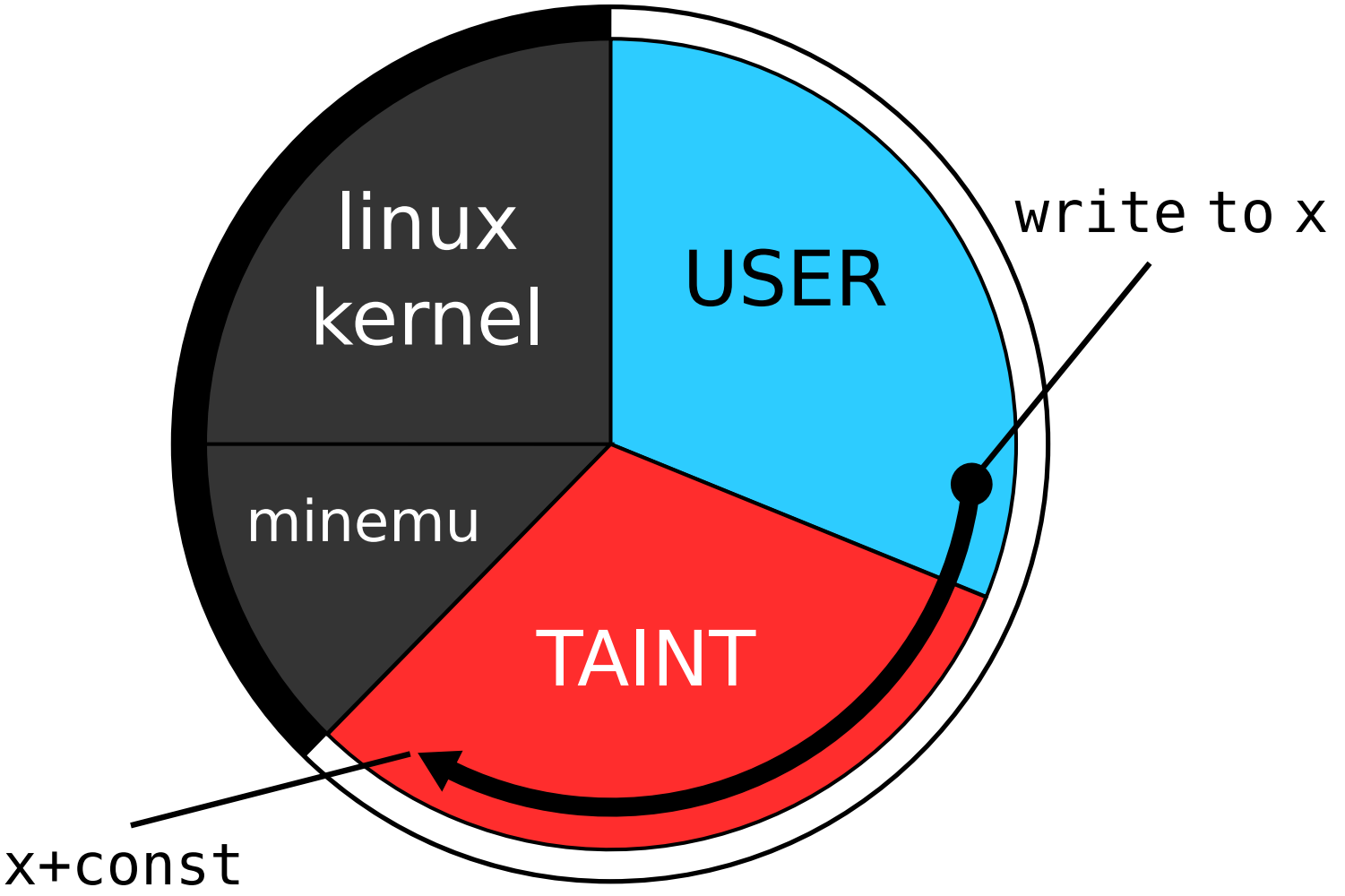
Memory layout (minemu)



Memory layout (minemu)



Memory layout (minemu)



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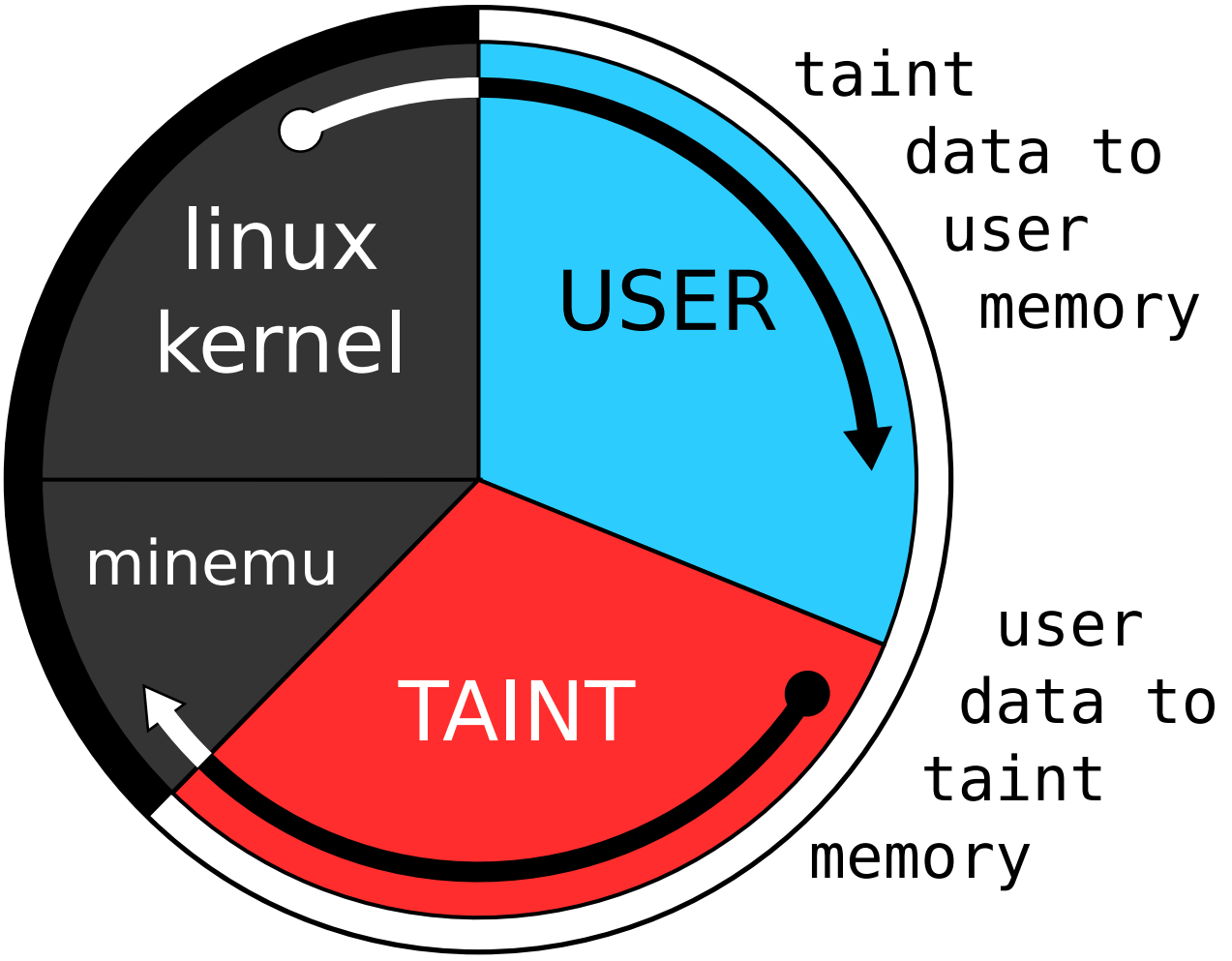

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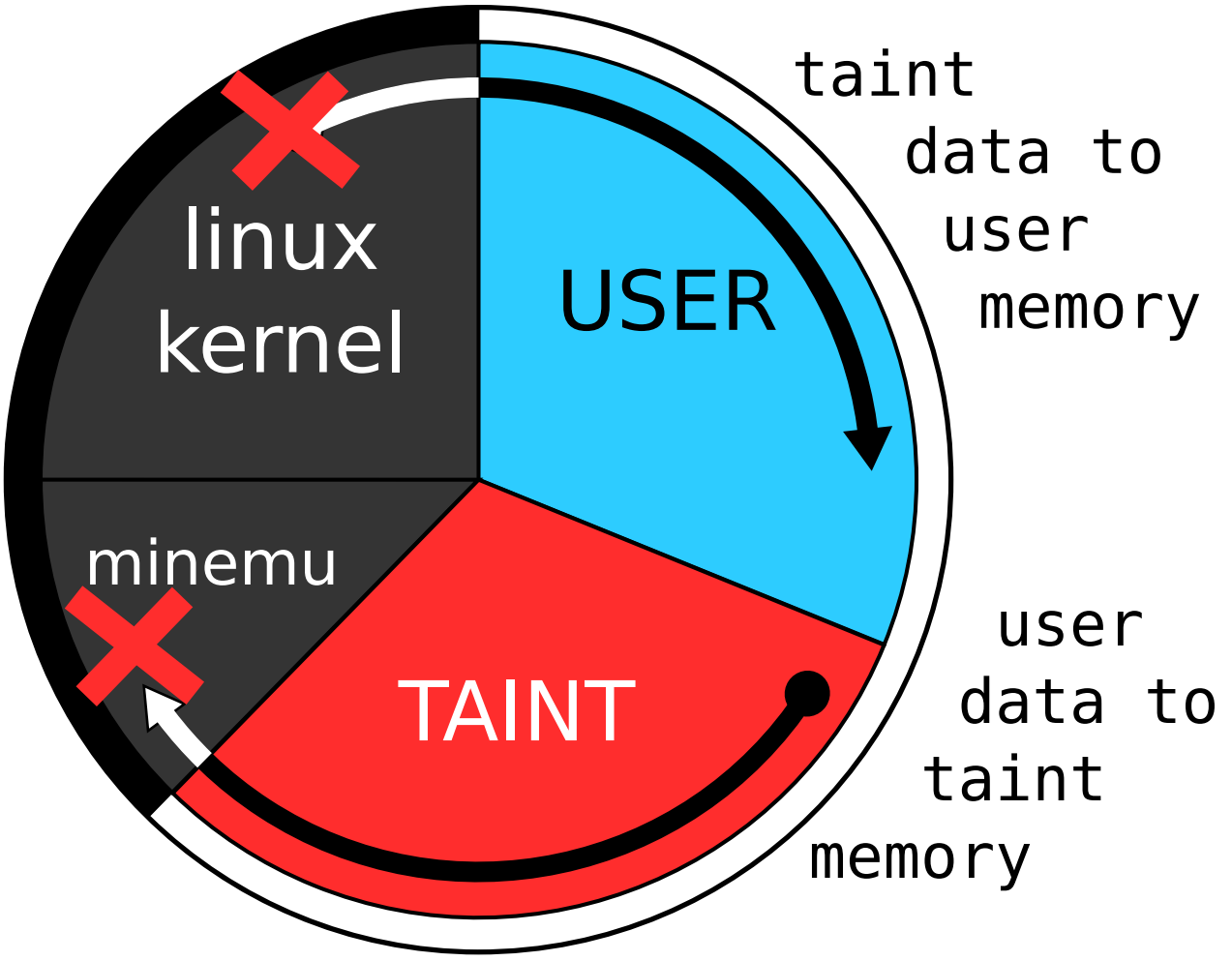

Addressing shadow memory

```
mov EAX, (EDX)
```

Addressing shadow memory

```
mov EAX, (EDX)

address:

    EDX
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```
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taint:

    EDX+const
```

Addressing shadow memory

```
mov EAX, (EDX+EBX*4)
```

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mov EAX, (EDX+EBX*4)

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```

Addressing shadow memory

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mov EAX, (EDX+EBX*4)

address:

    EDX+EBX*4

taint:

    EDX+EBX*4+const
```

Addressing shadow memory

```
push ESI
```

Addressing shadow memory

```
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address:

    ESP
```

Addressing shadow memory

```
push ESI

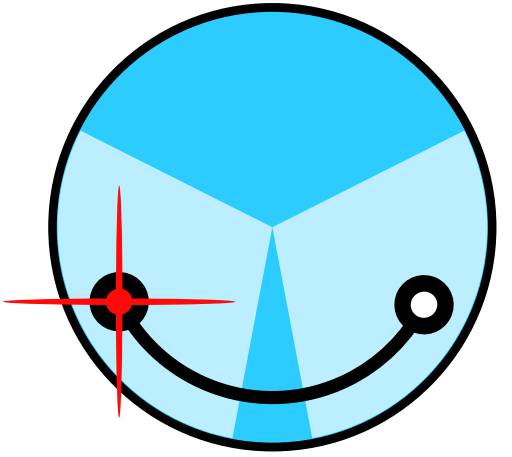
address:

    ESP

taint:

    ESP+const
```

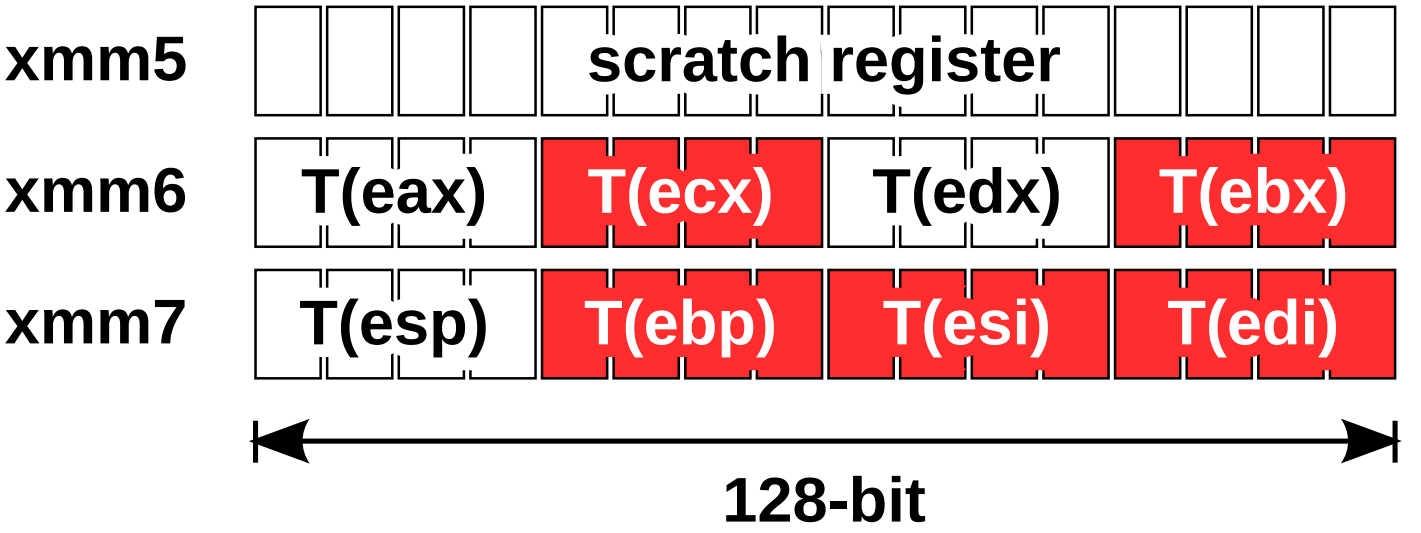
Is this slowness fundamental?



minemu

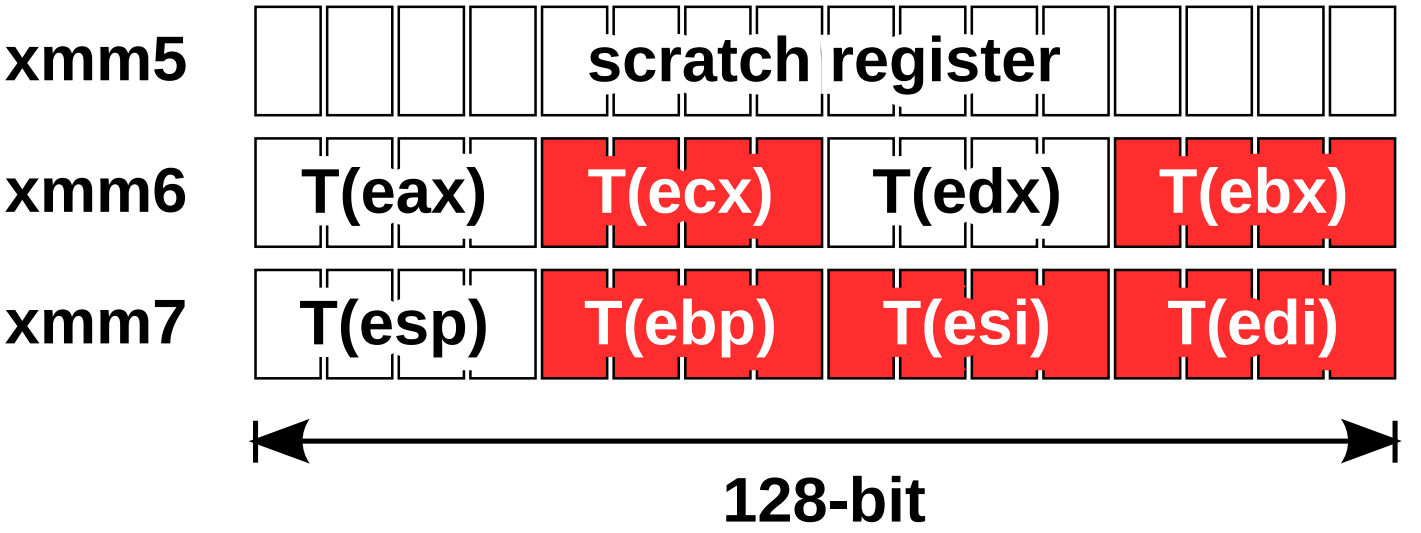
- fast emulator
- memory layout
- use SSE registers to hold taint

Taint propagation in SSE registers



Taint propagation in SSE registers

```
add EDX, x
```



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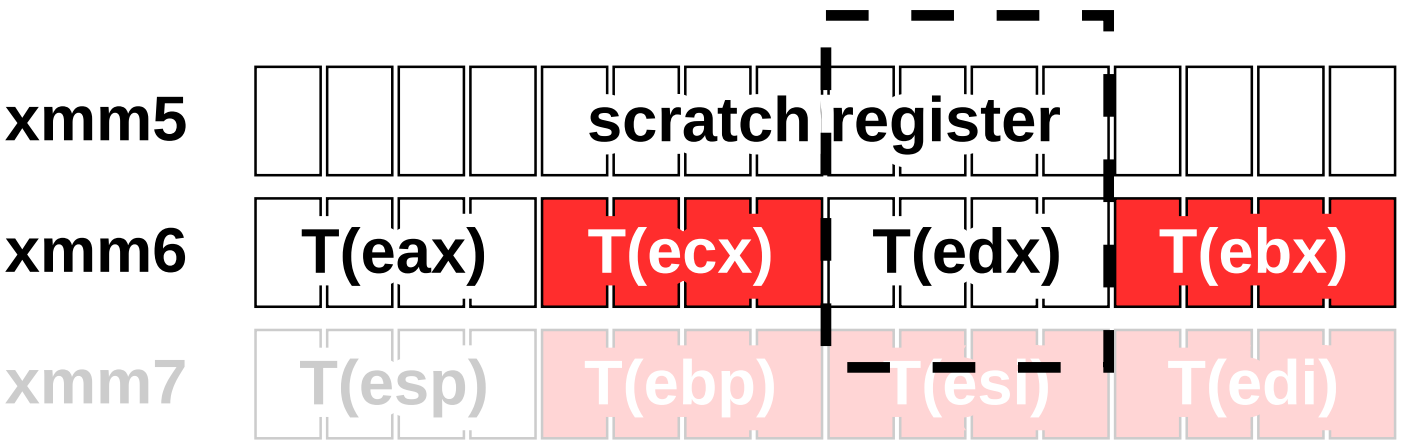
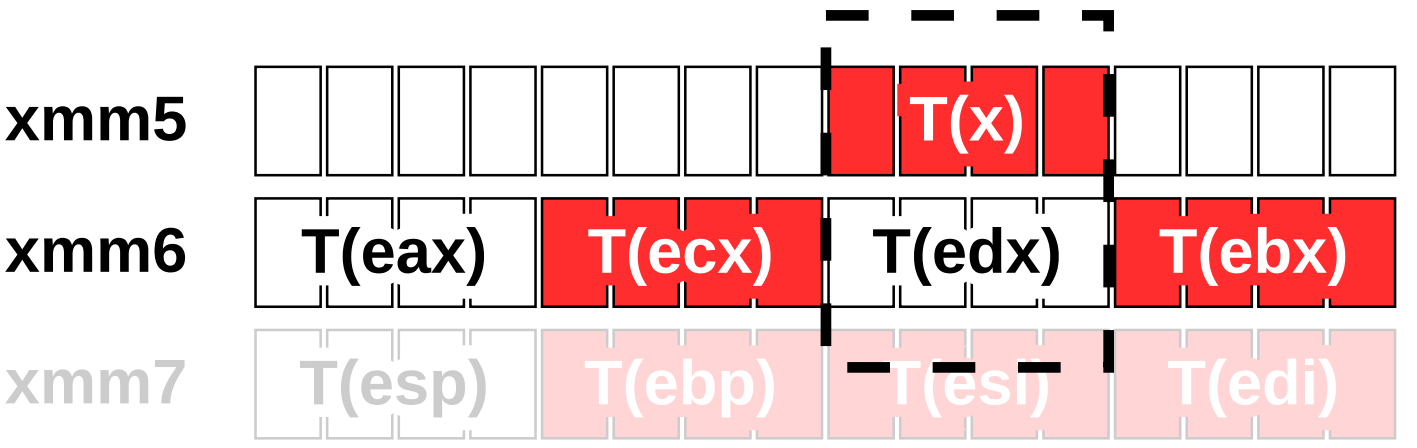


Taint propagation in SSE registers

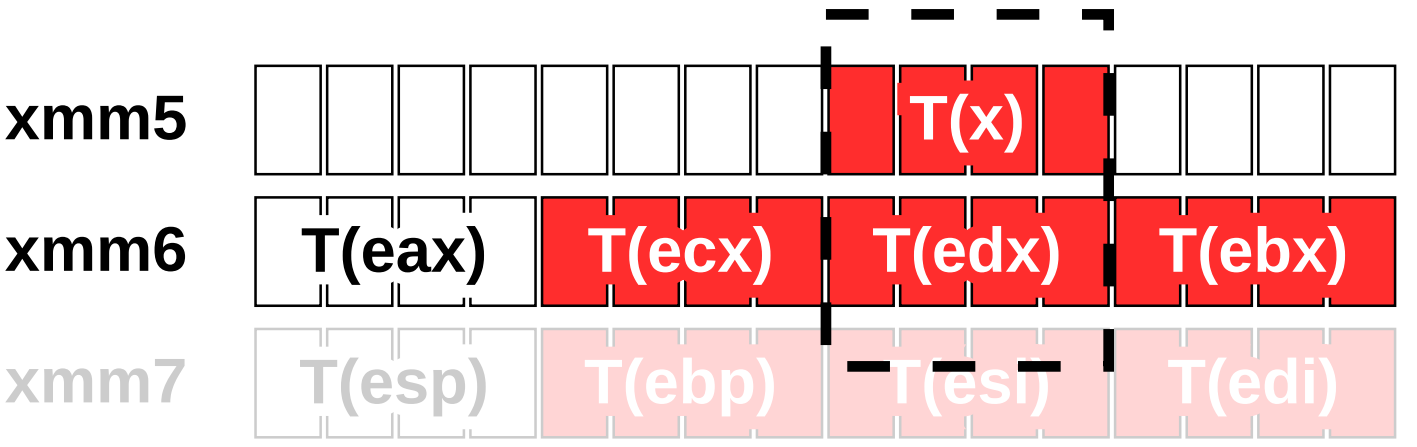
`add EDX, x`



`vector insert`

Taint propagation in SSE registers

`add EDX, x`

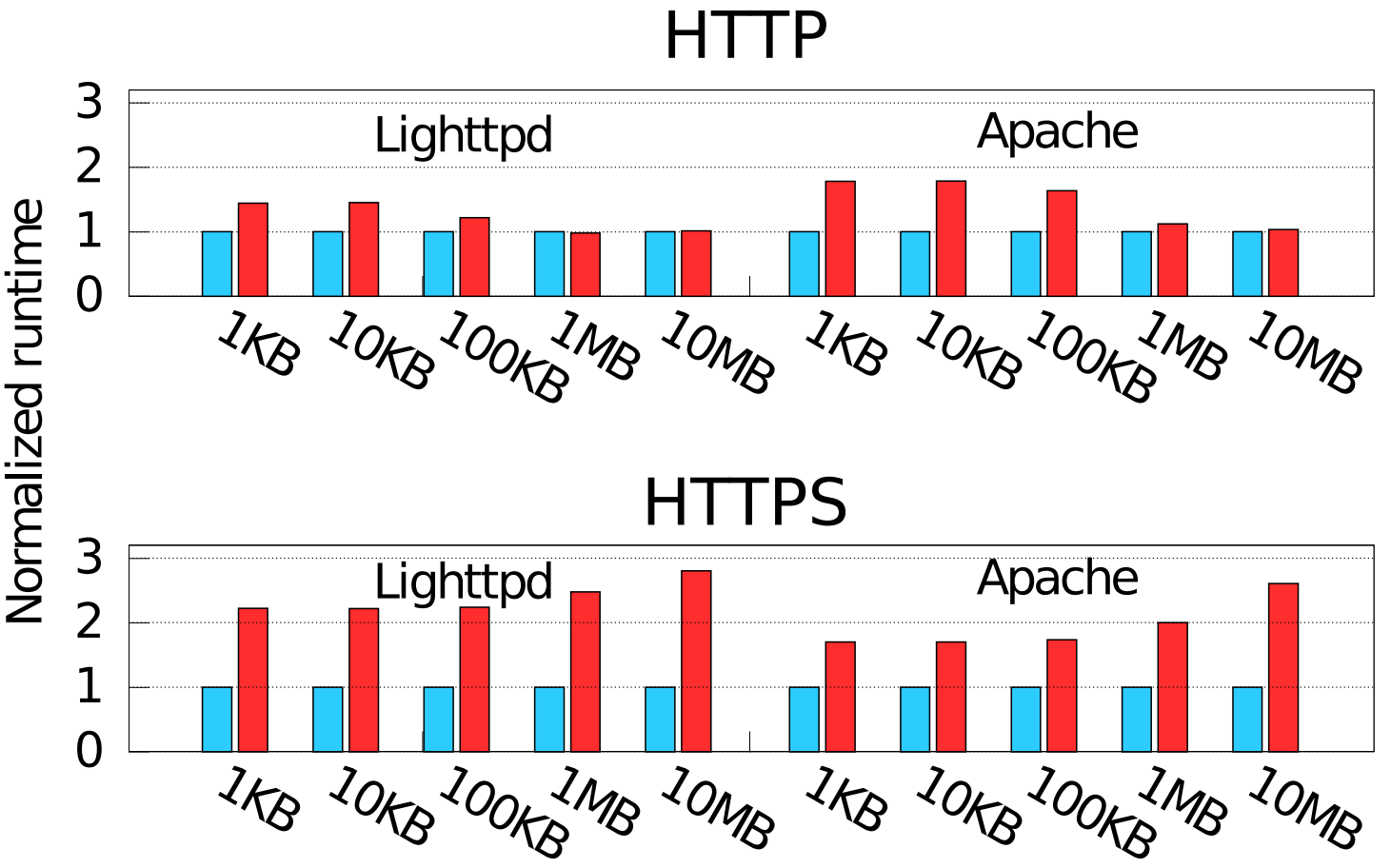


or

Effectiveness

Application	Type of vulnerability	Security advisory
Snort 2.4.0	Stack overflow	CVE-2005-3252
Cyrus imapd 2.3.2	Stack overflow	CVE-2006-2502
Samba 3.0.22	Heap overflow	CVE-2007-2446
Memcached 1.1.12	Heap overflow	CVE-2009-2415
Nginx 0.6.32	Buffer underrun	CVE-2009-2629
Proftpd 1.3.3a	Stack overflow	CVE-2010-4221
Samba 3.2.5	Heap overflow	CVE-2010-2063
Telnetd 1.6	Heap overflow	CVE-2011-4862
Ncompress 4.2.4	Stack overflow	CVE-2001-1413
Iwconfig V.26	Stack overflow	CVE-2003-0947
Aspell 0.50.5	Stack overflow	CVE-2004-0548
Htget 0.93	Stack overflow	CVE-2004-0852
Socat 1.4	Format string	CVE-2004-1484
Aeon 0.2a	Stack overflow	CVE-2005-1019
Exim 4.41	Stack overflow	EDB-ID#796
Htget 0.93	Stack overflow	
Tipxd 1.1.1	Format string	OSVDB-ID#12346

Performance



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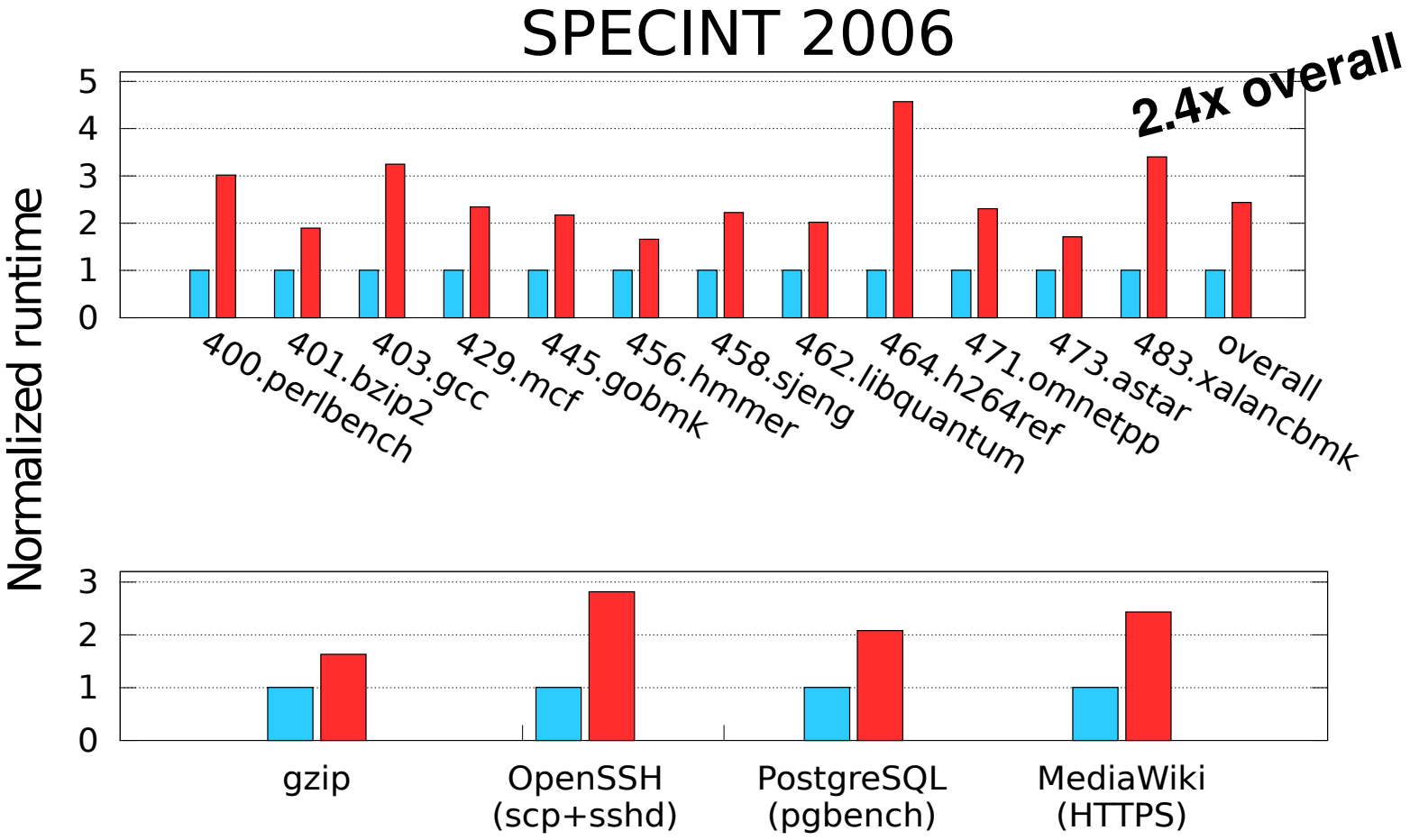

Limitations

Limitations

Doesn't prevent memory corruption, only acts when the untrusted data is used for arbitrary code execution.

Limitations

Tainted pointer dereferences

```
tainted_pointer->some_field = useful_untainted_value;
```

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```
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```

propagation can lead to false positives:

```
dispatch_table[checked_input]();
```

Limitations

Taint whitewashing

```
out = latin1_to_ascii[in];
```

Limitations

Format string attacks:

```
printf("%65534s %123$hn"); // Propagates taint in glibc

printf("FillerFiller...%123$hn"); // Does not :-(
```

Limitations

Does not protect against non-control-flow exploits

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```
void try_system(char *username, char *cmd)
{
    int user_rights = get_credentials(username);
    char buf[16] ; strcpy(buf, username);
    if (user_rights & ALLOW_SYSTEM)
        system(cmd);
    else
        log_error("user %s attempted login", buf);
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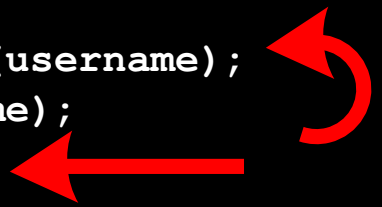

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in some cases we can add validation hooks.

- `mysql_query()` can be hooked to check for taint outside of literals in SQL queries.

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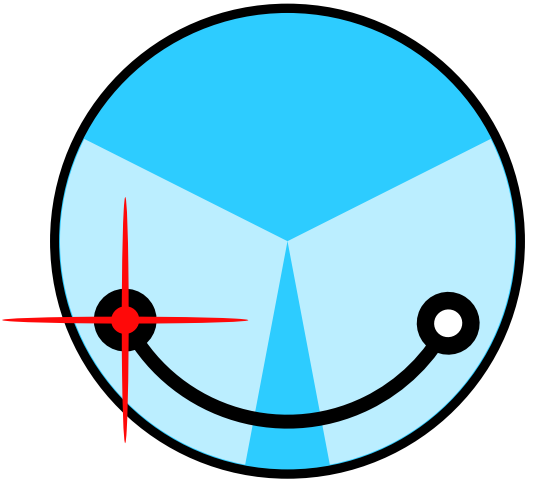
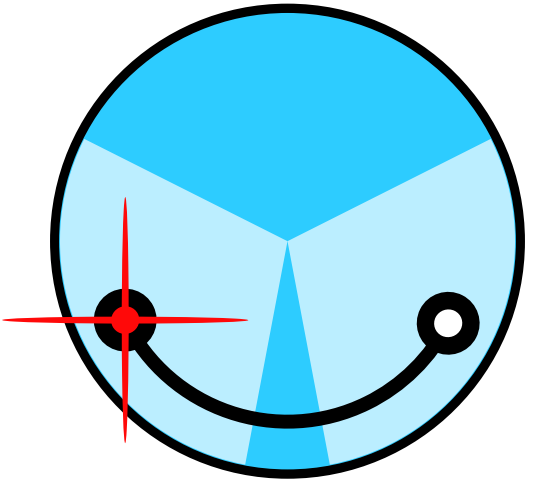
`_IO_vfprintf()` in glibc can be hooked to check format strings for taint.



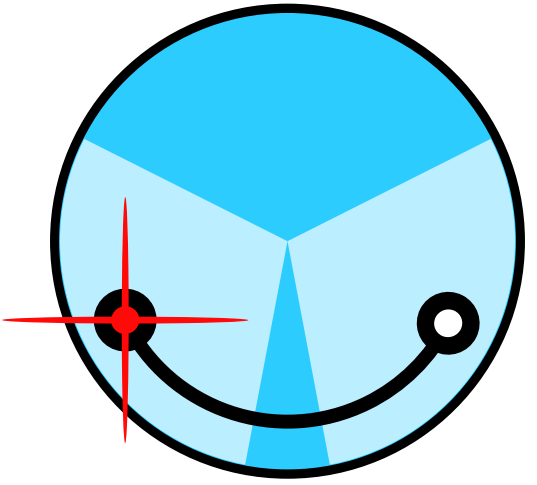
Demo

```
demo@demo:~#  ./minemu bash
```



Minemu

```
git clone https://minemu.org/code/minemu.git
```



Minemu

```
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```

any questions?

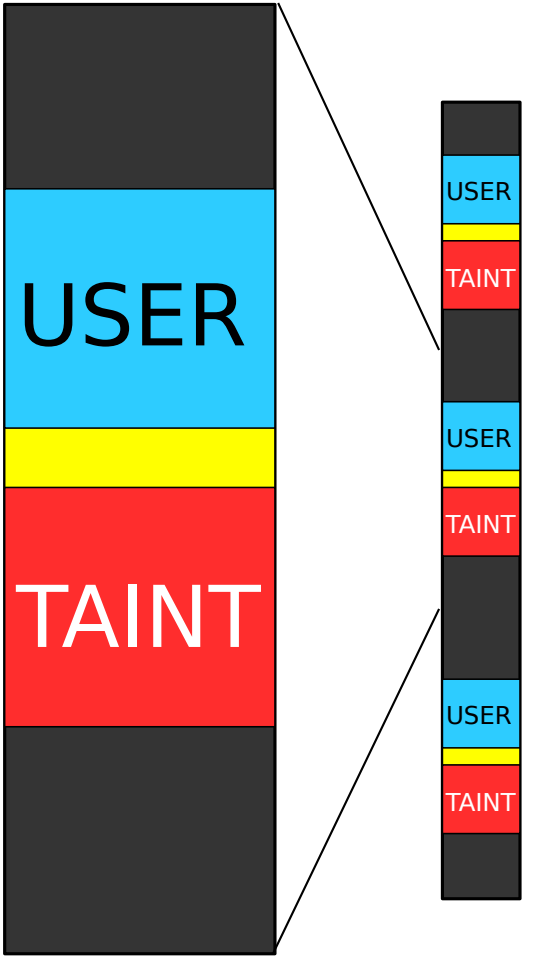
VU VRIJE UNIVERSITEIT AMSTERDAM

```
https://minemu.org/vms/

5f1ee00029e2c68699a7670de7aef02e  minemu-demo.ova
c4ee74155a858676bfb54e1fcfb6db0e  minemu-demo.qcow2
5b8b910c38901f43d406a21fe9767822  minemu-demo.vdi.gz
7ba81ae9d35bfa05a70068a804a331ac  minemu-demo.vmdk.gz
c37acdc455ebac700139f60da621bc38  minemu-demo.xml

minemu needs CPU with SSE 4.1
```

Memory layout (64 bit)



Memory layout (64 bit) alternative