"This file should be placed at ~/.vimrc"

se nu ai hls et ru ic is sc cul

#### Contents 14 6.13 PiCount\* se re=1 ts=4 sts=4 sw=4 ls=2 mouse=a 6.14 Discrete Log\* . . . . . 14 svntax on 6.15 Berlekamp Massey . . 14 hi cursorline cterm=none ctermbg=89 1 Basic 6.16 Primes . . . . . . . . . 15 set bg=dark 1.1 vimrc . . . . . . . . . 6.17 Theorem . . . . . . . inoremap {<CR> {<CR>}<Esc>ko<tab> 1.2 readchar . . . . . . 6.18 Estimation . . . . . . 1.3 Black Magic . . . . . . 6.19 Euclidean Algorithms 1.2 readchar 15 6.20 General Purpose 2 Graph Numbers 15 inline char readchar() { 2.1 BCC Vertex\* . . . . . 6.21 Tips for Generating static const size\_t bufsize = 65536; Functions . . . . . . 16 static char buf[bufsize]; static char \*p = buf, \*end = buf; 2.4 MinimumMeanCycle\* Polvnomial 16 if (p == end) end = buf + 2.5 Virtual Tree\* . . . . . 7.1 Fast Fourier Transform 16 fread\_unlocked(buf, 1, bufsize, stdin), p = buf; 2.6 Maximum Clique Dyn\* Theory 7.2 Number return \*p++: 2.7 Minimum Steiner Tree\* Transform\* . . . . . . 16 } Dominator Tree\* . . . 2.8 3 Fast Walsh Transform\* 16 Minimum Arbores-7.4 Polynomial Operation 16 1.3 Black Magic cence\* . . . . . . . . . 7.5 Value Polynomial . . . 17 2.10 Vizing's theorem\* . . 7.6 Newton's Method . . 17 #include <ext/pb\_ds/priority\_queue.hpp> 2.11 Minimum Clique Cover\* 4 #include <ext/pb\_ds/assoc\_container.hpp> // rb\_tree 2.12 NumberofMaximalClique\* 4g Geometry 17 #include <ext/rope> // rope using namespace \_\_gnu\_pbds; using namespace \_\_gnu\_cxx; // rope **Data Structure** 18 3.1 Discrete Trick . . . . 18 typedef \_ \_gnu\_pbds::priority\_queue<**int**> heap; 3.2 BIT kth\* . . . . . . . 8.4 point in circle . . . . . 18 int main() { 3.3 Interval Container\* .3.4 Leftist Tree . . . . . 8.5 Convex hull\* . . . . . 18 heap h1, h2; // max heap PointInConvex\* . . . h1.push(1), h1.push(3), h2.push(2), h2.push(4); Heavy light Decom-3.5 8.7 TangentPointToHull\* h1.join(h2); // h1 = {1, 2, 3, 4}, h2 = {}; tree<ll, null\_type, less<ll>, rb\_tree\_tag 5 position . . . . . . . 8.8 Intersection of line Centroid Decomposiand convex . . . . . . 18 tree\_order\_statistics\_node\_update > st; tion\* . . . . . . . . . . . . 8.9 minMaxEnclosingRectangle\* LiChaoST\* . . . . . . tree<ll, ll, less<ll>, rb\_tree\_tag 8.10 VectorInPoly\* . . . . . , tree\_order\_statistics\_node\_update > mp; Link cut tree\* . . . . . 8.11 PolyUnion\* . . . . . . . . 8.12 PolyCut . . . . . . . 3.8 19 for (int x : {0, 3, 20, 50}) st.insert(x); 3.9 KDTree . . . . . . . . 19 assert(st. 8.13 Trapezoidalization . . 19 order\_of\_key(3) == 1 && st.order\_of\_key(4) == 2); Flow/Matching 8.14 Polar Angle Sort\* . . . 20 assert(\*st.find\_by\_order 4.1 Kuhn Munkres\* . . . . 8.15 Half plane intersection\* 20 (2) == 20 && \*st.lower bound(4) == 20);4.2 MincostMaxflow\* . . 8.16 RotatingSweepLine . rope<char> \*root[10]; // nsqrt(n) 4.3 Maximum Simple 8.17 Minimum Enclosing root[0] = new rope<char>(); Graph Matching\* . . . 8 Circle\* . . . . . . . . . 20 root[1] = new rope < char > (\*root[0]); Maximum Weight 8.18 Intersection of two Matching\* . . . . . . // root[1]->insert(pos, circles\* . . . . . . . 20 // root[1]->at(pos); 0-base 45 SW-mincut . . . . . 9 8.19 Intersection of poly-// root[1]->erase(pos, size); BoundedFlow\*(Dinic\*) q 4.6 gon and circle\* . . . . 21 4.7 Gomory Hu tree\* . . . 10 8.20 Intersection of line // \_\_int128\_t,\_\_float128\_t // for (int i = bs.\_Find\_first Minimum Cost Circuand circle\* . . . . . . 21 lation\* . . . . . . . . . 8.21 Tangent line of two (); i < bs.size(); i = bs.\_Find\_next(i)); Flow Models . . . . . 10 4.9 circles . . . . . . . . . 21 8.22 CircleCover\* . . . . . 21 2 10 Graph String 8.23 3Dpoint\* . . . . . . . 5.1 KMP . . . . . . . . . 10 8.24 Convexhull3D\* . . . . 2.1 BCC Vertex\* 5.2 Z-value\* . . . . . . . 8.25 DelaunayTriangulation\* 22 Manacher\* . . . . . . 5.3 11 8.26 Triangulation Vonoroi\* 23 vector < int > G[N]; // 1-base vector < int > nG[N \* 2], bcc[N]; int low[N], dfn[N], Time; SAIS\* . . . . . . . . . 5.4 11 8.27 Minkowski Sum\* . . . 23 5.5 Fise int bcc\_id[N], bcc\_cnt; // 1-base Smallest Rotation . . 11 9.1 Cyclic Ternary Search\* 23 De Bruijn sequence\* . bool is\_cut[N]; // whether is av 5.7 11 Mo's Alogrithm(With Extended SAM\* . . . bool cir[N \* 2]; 5.8 modification) . . . . . int st[N], top; 5.9 PalTree\* . . . . . . . Mo's Alogrithm On Tree 23 Additional Mo's Algovoid $dfs(int u, int pa = -1) {$ Math rithm Trick . . . . . . int child = 0; 6.1 ax+by=gcd(only Hilbert Curve . . . exgcd \*) . . . . . . . . low[u] = dfn[u] = ++Time; 9.6 DynamicConvexTrick\* 24 st[top++] = u;Floor and Ceil . . . . . 12 24 for (int v : G[u]) Floor Enumeration . . 12 24 **if** (!dfn[v]) { 6.4 Mod Min . . . . . . . 9.9 Matroid Intersection . 25 Gaussian integer gcd . 13 Miller Rabin\* . . . . . 13 dfs(v, u), ++child; 6.5 9.10 AdaptiveSimpson\* . . low[u] = min(low[u], low[v]); 6.6 9.11 Simulated Annealing . 25 if (dfn[u] <= low[v]) { is\_cut[u] = 1;</pre> Simultaneous Equa-9.12 Tree Hash\* . . . . . . 25 9.13 Binary Search On bcc[++bcc\_cnt].clear(); 13 Fraction . . . . . . . Simplex Algorithm . . 13 9.14 Min Plus Convolution\* 25 **do** { 6.9.1 Construction . 13 9.15 Bitset LCS . . . . . . . 25 6.10 chineseRemainder . . 14 bcc\_id[t = st[--top]] = bcc\_cnt; 9.16 N Queens Problem . . 25 bcc[bcc\_cnt].push\_back(t); 6.11 Factorial without prime factor\* . . . . . 14 6.12 QuadraticResidue\* . . 14 } while (t != v); 10 Python 26 bcc\_id[u] = bcc\_cnt; 10.1 Misc . . . . . . . . . . bcc[bcc\_cnt].pb(u); Basic } else if (dfn[v] < dfn[u] && v != pa)</pre> low[u] = min(low[u], dfn[v]); 1.1 vimrc if (pa == -1 && child < 2) is\_cut[u] = 0;</pre>

}

void bcc\_init(int n) { // TODO: init {nG, cir}[1..2n]

```
Time = bcc_cnt = top = 0;
    for (int i = 1; i <= n; ++i)
        G[i].clear(), dfn[i] = bcc_id[i] = is_cut[i] = 0;
}

void bcc_solve(int n) {
    for (int i = 1; i <= n; ++i)
        if (!dfn[i]) dfs(i);

    // block-cut tree
    for (int i = 1; i <= n; ++i)
        if (is_cut[i])
        bcc_id[i] = ++bcc_cnt, cir[bcc_cnt] = 1;
    for (int i = 1; i <= bcc_cnt && !cir[i]; ++i)
        for (int j : bcc[i])
        if (is_cut[j])
        nG[i].pb(bcc_id[j]), nG[bcc_id[j]].pb(i);
}</pre>
```

# 2.2 Bridge\*

```
int low[N], dfn[N], Time; // 1-base
vector<pii> G[N], edge;
vector<bool> is_bridge;
void init(int n) {
  for (int i = 1; i <= n; ++i)</pre>
    G[i].clear(), low[i] = dfn[i] = 0;
void add_edge(int a, int b) {
   G[a].pb(pii(b, SZ(edge))), G[b].pb(pii(a, SZ(edge)));
  edge.pb(pii(a, b));
void dfs(int u, int f) {
  dfn[u] = low[u] = ++Time;
  for (auto i : G[u])
    if (!dfn[i.X])
      dfs(i.X, i.Y), low[u] = min(low[u], low[i.X]);
    else if (i.Y != f) low[u] = min(low[u], dfn[i.X]);
  if (low[u] == dfn[u] && f != -1) is_bridge[f] = 1;
void solve(int n) {
  is_bridge.resize(SZ(edge));
  for (int i = 1; i <= n; ++i)</pre>
    if (!dfn[i]) dfs(i, -1);
```

# 2.3 2SAT (SCC)\*

```
struct SAT { // 0-base
  int low[N], dfn[N], bln[N], n, Time, nScc;
  bool instack[N], istrue[N];
  stack<int> st;
  vector<int> G[N], SCC[N];
  void init(int _n) {
   n = _n; // assert(n * 2 <= N);
    for (int i = 0; i < n + n; ++i) G[i].clear();</pre>
  void add_edge(int a, int b) { G[a].pb(b); }
  int rv(int a) {
    if (a >= n) return a - n;
    return a + n:
  void add_clause(int a, int b) {
    add_edge(rv(a), b), add_edge(rv(b), a);
  void dfs(int u) {
    dfn[u] = low[u] = ++Time;
    instack[u] = 1, st.push(u);
    for (int i : G[u])
      if (!dfn[i])
      dfs(i), low[u] = min(low[i], low[u]);
else if (instack[i] && dfn[i] < dfn[u])</pre>
        low[u] = min(low[u], dfn[i]);
    if (low[u] == dfn[u]) {
      int tmp;
      do {
        tmp = st.top(), st.pop();
        instack[tmp] = 0, bln[tmp] = nScc;
      } while (tmp != u);
      ++nScc;
    }
 }
```

```
bool solve() {
    Time = nScc = 0;
    for (int i = 0; i < n + n; ++i)
        SCC[i].clear(), low[i] = dfn[i] = bln[i] = 0;
    for (int i = 0; i < n + n; ++i)
        if (!dfn[i]) dfs(i);
    for (int i = 0; i < n + n; ++i) SCC[bln[i]].pb(i);
    for (int i = 0; i < n; ++i) {
        if (bln[i] == bln[i + n]) return false;
        istrue[i] = bln[i] < bln[i + n];
        istrue[i + n] = !istrue[i];
    }
    return true;
}
</pre>
```

# 2.4 MinimumMeanCycle\*

```
ll road[N][N]; // input here
struct MinimumMeanCvcle {
  ll dp[N + 5][N], n;
  pll solve() {
     ll a = -1, b = -1, L = n + 1;
     for (int i = 2; i <= L; ++i)
  for (int k = 0; k < n; ++k)</pre>
          for (int j = 0; j < n; ++j)</pre>
            dp[i][j] =
              min(dp[i - 1][k] + road[k][j], dp[i][j]);
     for (int i = 0; i < n; ++i) {
  if (dp[L][i] >= INF) continue;
       ll ta = 0, tb = 1;
       for (int j = 1; j < n; ++j)</pre>
          if (dp[j][i] < INF &&</pre>
            ta * (L - j) < (dp[L][i] - dp[j][i]) * tb)
ta = dp[L][i] - dp[j][i], tb = L - j;
       if (ta == 0) continue;
       if (a == -1 || a * tb > ta * b) a = ta, b = tb;
     if (a != -1) {
       ll g = __gcd(a, b);
       return pll(a / g, b / g);
     return pll(-1LL, -1LL):
  void init(int _n) {
     n = _n;
for (int i = 0; i < n; ++i)</pre>
       for (int j = 0; j < n; ++j) dp[i + 2][j] = INF;</pre>
};
```

## 2.5 Virtual Tree\*

```
vector<int> vG[N];
int top. st[N]:
void insert(int u) {
  if (top == -1) return st[++top] = u, void();
  int p = LCA(st[top], u);
  if (p == st[top]) return st[++top] = u, void();
  while (top >= 1 && dep[st[top - 1]] >= dep[p])
    vG[st[top - 1]].pb(st[top]), --top;
  if (st[top] != p)
    vG[p].pb(st[top]), --top, st[++top] = p;
  st[++top] = u;
void reset(int u) {
 for (int i : vG[u]) reset(i);
  vG[u].clear();
void solve(vector<int> &v) {
 top = -1;
  sort(ALL(v),
    [&](int a, int b) { return dfn[a] < dfn[b]; });
  for (int i : v) insert(i);
  while (top > 0) vG[st[top - 1]].pb(st[top]), --top;
  // do something
  reset(v[0]);
```

#### 2.6 Maximum Clique Dyn\*

```
struct MaxClique { // fast when N <= 100
bitset<N> G[N], cs[N];
```

```
int ans, sol[N], q, cur[N], d[N], n;
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i) G[i].reset();</pre>
  void add_edge(int u, int v) {
    G[u][v] = G[v][u] = 1;
  void pre_dfs(vector<int> &r, int l, bitset<N> mask) {
    if (1 < 4) {
      for (int i : r) d[i] = (G[i] & mask).count();
      sort(ALL(r)
           , [&](int x, int y) { return d[x] > d[y]; });
    vector<int> c(SZ(r));
    int lft = max(ans - q + 1, 1), rgt = 1, tp = 0;
    cs[1].reset(), cs[2].reset();
    for (int p : r) {
      int k = 1;
      while ((cs[k] & G[p]).any()) ++k;
      if (k > rgt) cs[++rgt + 1].reset();
      cs[k][p] = 1;
      if (k < lft) r[tp++] = p;</pre>
    for (int k = lft; k <= rgt; ++k)</pre>
      for (int p = cs[k]._Find_first
           (); p < N; p = cs[k]._Find_next(p))
        r[tp] = p, c[tp] = k, ++tp;
    dfs(r, c, l + 1, mask);
  void dfs(vector<</pre>
      int> &r, vector<int> &c, int l, bitset<N> mask) {
    while (!r.empty()) {
      int p = r.back();
      r.pop_back(), mask[p] = 0;
      if (q + c.back() <= ans) return;</pre>
      cur[q++] = p;
      vector<int> nr;
      for (int i : r) if (G[p][i]) nr.pb(i);
      if (!nr.empty()) pre_dfs(nr, l, mask & G[p]);
      else if (q > ans) ans = q, copy_n(cur, q, sol);
      c.pop_back(), --q;
    }
  int solve() {
    vector<int> r(n);
    ans = q = 0, iota(ALL(r), 0);
    pre_dfs(r, 0, bitset<N>(string(n, '1')));
    return ans;
 }
};
```

#### 2.7 Minimum Steiner Tree\*

```
struct SteinerTree { // 0-base
  int n, dst[N][N], dp[1 << T][N], tdst[N];</pre>
  int vcst[N]; // the cost of vertexs
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i) {</pre>
      fill_n(dst[i], n, INF);
      dst[i][i] = vcst[i] = 0;
    }
  void chmin(int &x, int val) {
    x = min(x, val);
  void add_edge(int ui, int vi, int wi) {
    chmin(dst[ui][vi], wi);
  void shortest_path() {
    for (int k = 0; k < n; ++k)
      for (int i = 0; i < n; ++i)</pre>
        for (int j = 0; j < n; ++j)
           chmin(dst[i][j], dst[i][k] + dst[k][j]);
  int solve(const vector<int>& ter) {
    shortest_path();
    int t = SZ(ter), full = (1 << t) - 1;</pre>
    for (int i = 0; i <= full; ++i)
  fill_n(dp[i], n, INF);</pre>
    copy_n(vcst, n, dp[0]);
        (int msk = 1; msk <= full; ++msk) {
      if (!(msk & (msk - 1))) {
        int who = __lg(msk);
        for (int i = 0; i < n; ++i)</pre>
```

#### 2.8 Dominator Tree\*

```
struct dominator_tree { // 1-base
  vector < int > G[N], rG[N];
  int n, pa[N], dfn[N], id[N], T
int semi[N], idom[N], best[N];
                                 Time;
  vector<int> tree[N]; // dominator_tree
  void init(int _n) {
    n = n;
    for (int i = 1; i <= n; ++i)
      G[i].clear(), rG[i].clear();
  void add_edge(int u, int v) {
    G[u].pb(v), rG[v].pb(u);
  void dfs(int u) {
    id[dfn[u] = ++Time] = u;
    for (auto v : G[u])
      if (!dfn[v]) dfs(v), pa[dfn[v]] = dfn[u];
  int find(int y, int x) {
    if (y <= x) return y;</pre>
    int tmp = find(pa[y], x);
    if (semi[best[y]] > semi[best[pa[y]]])
      best[y] = best[pa[y]];
    return pa[y] = tmp;
  void tarjan(int root) {
    Time = 0;
    for (int i = 1; i <= n; ++i) {</pre>
      dfn[i] = idom[i] = 0;
      tree[i].clear();
      best[i] = semi[i] = i;
    dfs(root);
    for (int i = Time; i > 1; --i) {
      int u = id[i];
      for (auto v : rG[u])
        if (v = dfn[v]) {
           find(v, i);
           semi[i] = min(semi[i], semi[best[v]]);
      tree[semi[i]].pb(i);
      for (auto v : tree[pa[i]]) {
         find(v, pa[i]);
         idom[v] =
           semi[best[v]] == pa[i] ? pa[i] : best[v];
      tree[pa[i]].clear();
    for (int i = 2; i <= Time; ++i) {</pre>
      if (idom[i] != semi[i]) idom[i] = idom[idom[i]];
      tree[id[idom[i]]].pb(id[i]);
  }
};
```

# 2.9 Minimum Arborescence\*

```
/* TODO
DSU: disjoint set
- DSU(n), .boss(x), .Union(x, y)
min_heap<
    T, Info>: min heap for type {T, Info} with lazy tag
- .push({w, i}),
    .top(), .join(heap), .pop(), .empty(), .add_lazy(v)
*/
```

```
struct E { int s, t; ll w; }; // O-base
vector<int> dmst(const vector<E> &e, int n, int root) {
  vector<min_heap<ll, int>> h(n * 2);
for (int i = 0; i < SZ(e); ++i)</pre>
    h[e[i].t].push({e[i].w, i});
  DSU dsu(n * 2);
  vector<int> v(n * 2, -1), pa(n * 2, -1), r(n * 2);
  v[root] = n + 1;
  int pc = n;
  for (int i = 0; i < n; ++i) if (v[i] == -1) {</pre>
    for (int p = i; v[p]
          == -1 || v[p] == i; p = dsu.boss(e[r[p]].s)) {
       if (v[p] == i) {
         int q = p; p = pc++;
           h[q].add_lazy(-h[q].top().X);
           pa[q] = p, dsu.Union(p, q), h[p].join(h[q]);
        } while ((q = dsu.boss(e[r[q]].s)) != p);
      v[p] = i;
      while (!h[p].
           empty() && dsu.boss(e[h[p].top().Y].s) == p)
        h[p].pop();
       if (h[p].empty()) return {}; // no solution
      r[p] = h[p].top().Y;
    }
  }
  vector<int> ans;
  for (int i = pc
        - 1; i >= 0; i--) if (i != root && v[i] != n) {
    for (int f = e[r[i]].t; ~f && v[f] != n; f = pa[f])
      v[f] = n;
    ans.pb(r[i]);
  return ans; // default minimize, returns edgeid array
\} // O(Ef(E)), f(E) from min_heap
```

# 2.10 Vizing's theorem\*

```
namespace vizing { // returns
  edge coloring in adjacent matrix G. 1 - based
const int N = 105;
int C[N][N], G[N][N], X[N], vst[N], n;
void init(int _n) { n = _n;
for (int i = 0; i <= n; ++i)
    for (int j = 0; j <= n; ++j)</pre>
       C[i][j] = G[i][j] = 0;
void solve(vector<pii> &E) {
  auto update = [&](int u)
{ for (X[u] = 1; C[u][X[u]]; ++X[u]); };
  auto color = [&](int u, int v, int c) {
     int p = G[u][v];
    G[u][v] = G[v][u] = c;
C[u][c] = v, C[v][c] = u;
    C[u][p] = C[v][p] = 0;
     if (p) X[u] = X[v] = p;
     else update(u), update(v);
     return p;
  }:
  auto flip = [&](int u, int c1, int c2) {
     int p = C[u][c1];
     swap(C[u][c1], C[u][c2]);
     if (p) G[u][p] = G[p][u] = c2;
     if (!C[u][c1]) X[u] = c1;
     if (!C[u][c2]) X[u] = c2;
     return p;
  fill_n(X + 1, n, 1);
for (int t = 0; t < SZ(E); ++t) {
     int u = E[t
          ].X, v0 = E[t].Y, v = v0, c0 = X[u], c = c0, d;
     vector<pii> L;
     fill_n(vst + 1, n, 0);
     while (!G[u][v0]) {
       L.emplace_back(v, d = X[v]);
       if (!C[v][c]) for (int a = SZ(
                  1; a >= 0; --a) c = color(u, L[a].X, c);
            L) -
       else if (!C[u][d]) for (int a = SZ(L
) - 1; a >= 0; --a) color(u, L[a].X, L[a].Y);
else if (vst[d]) break;
       else vst[d] = 1, v = C[u][d];
     if (!G[u][v0]) {
       for (; v; v = flip(v, c, d), swap(c, d));
       if (int a; C[u][c0]) {
```

# 2.11 Minimum Clique Cover\*

```
struct Clique_Cover { // 0-base, 0(n2^n)
   int co[1 << N], n, E[N];</pre>
   int dp[1 << N];</pre>
   void init(int _n) {
    n = _n, fill_n(dp, 1 << n, 0);
</pre>
      fill_n(E, n, 0), fill_n(co, 1 << n, 0);
   void add_edge(int u, int v) {
     E[u] \mid = 1 << v, E[v] \mid = 1 << u;
   int solve() {
     for (int i = 0; i < n; ++i)</pre>
       co[1 << i] = E[i] | (1 << i);
     co[0] = (1 << n) - 1;

dp[0] = (n & 1) * 2 - 1;
      for (int i = 1; i < (1 << n); ++i) {</pre>
        int t = i & -i;
        dp[i] = -dp[i ^ t];
        co[i] = co[i ^ t] & co[t];
      for (int i = 0; i < (1 << n); ++i)</pre>
        co[i] = (co[i] & i) == i;
     fwt(co, 1 << n, 1);
for (int ans = 1; ans < n; ++ans) {
  int sum = 0; // probabilistic</pre>
        for (int i = 0; i < (1 << n); ++i)</pre>
          sum += (dp[i] *= co[i]);
        if (sum) return ans;
     }
     return n;
   }
};
```

# 2.12 NumberofMaximalClique\*

```
struct BronKerbosch { // 1-base
   int n, a[N], g[N][N];
  int S, all[N][N], some[N][N], none[N][N];
  void init(int _n) {
     for (int i = 1; i <= n; ++i)</pre>
       for (int j = 1; j <= n; ++j) g[i][j] = 0;</pre>
  void add_edge(int u, int v) {
    g[u][v] = g[v][u] = 1;
  void dfs(int d, int an, int sn, int nn) {
     if (S > 1000) return; // pruning
     if (sn == 0 && nn == 0) ++S;
     int u = some[d][0];
     for (int i = 0; i < sn; ++i) {</pre>
       int v = some[d][i];
       if (g[u][v]) continue;
       int tsn = 0, tnn = 0;
       copy_n(all[d], an, all[d + 1]);
       all[d + 1][an] = v;
       for (int j = 0; j < sn; ++j)
         if (g[v][some[d][j]])
           some[d + 1][tsn++] = some[d][j];
       for (int j = 0; j < nn; ++j)</pre>
         if (g[v][none[d][j]])
           none[d + 1][tnn++] = none[d][j];
       dfs(d + 1, an + 1, tsn, tnn);
       some[d][i] = 0, none[d][nn++] = v;
  int solve() {
    iota(some[0], some[0] + n, 1);
     S = 0, dfs(0, 0, n, 0);
     return S;
};
```

## 3 Data Structure

#### 3.1 Discrete Trick

```
vector < int > val;
// build
sort(ALL
          (val)), val.resize(unique(ALL(val)) - val.begin());
// index of x
upper_bound(ALL(val), x) - val.begin();
// max idx <= x
upper_bound(ALL(val), x) - val.begin();
// max idx < x
lower_bound(ALL(val), x) - val.begin();</pre>
```

#### 3.2 BIT kth\*

#### 3.3 Interval Container\*

```
/* Add and
     remove intervals from a set of disjoint intervals.
 * Will merge the added interval with
      any overlapping intervals in the set when adding.
 * Intervals are [inclusive, exclusive). */
set<pii>::
    iterator addInterval(set<pii>% is, int L, int R) {
  if (L == R) return is.end();
  auto it = is.lower_bound({L, R}), before = it;
  while (it != is.end() && it->X <= R) {</pre>
   R = max(R, it->Y);
    before = it = is.erase(it);
  if (it != is.begin() && (--it)->Y >= L) {
   L = min(L, it->X);
R = max(R, it->Y);
    is.erase(it);
  return is.insert(before, pii(L, R));
void removeInterval(set<pii>& is, int L, int R) {
  if (L == R) return;
  auto it = addInterval(is, L, R);
  auto r2 = it->Y;
  if (it->X == L) is.erase(it);
  else (int&)it->Y = L;
  if (R != r2) is.emplace(R, r2);
```

#### 3.4 Leftist Tree

```
struct node {
   ll v, data, sz, sum;
node *l, *r;
   node(ll k)
      : v(0), data(k), sz(1), l(0), r(0), sum(k) {}
ll sz(node *p) { return p ? p->sz : 0; }
ll V(node *p) { return p ? p->v : -1; }
ll sum(node *p) { return p ? p->sum : 0; }
node *merge(node *a, node *b) {
   if (!a || !b) return a ? a : b;
   if (a->data < b->data) swap(a, b);
   a - r = merge(a - r, b);
   if (V(a->r) > V(a->l)) swap(a->r, a->l);
   a \rightarrow v = V(a \rightarrow r) + 1, a \rightarrow sz = sz(a \rightarrow l) + sz(a \rightarrow r) + 1; a \rightarrow sum = sum(a \rightarrow l) + sum(a \rightarrow r) + a \rightarrow data;
   return a;
void pop(node *&o) {
   node *tmp = o;
   o = merqe(o->l, o->r);
   delete tmp;
```

# 3.5 Heavy light Decomposition

```
struct Heavy_light_Decomposition { // 1-base
  int n, ulink[N], deep[N], mxson[N], w[N], pa[N];
  int t, pl[N], data[N], dt[N], bln[N], edge[N], et;
  vector<pii> G[N];
  void init(int _n) {
    n = _n, t = 0, et = 1;
for (int i = 1; i <= n; ++i)
  G[i].clear(), mxson[i] = 0;</pre>
  void add_edge(int a, int b, int w) {
    G[a].pb(pii(b, et));
     G[b].pb(pii(a, et));
     edge[et++] = w;
  void dfs(int u, int f, int d) {
    w[u] = 1, pa[u] = f, deep[u] = d++;
     for (auto &i : G[u])
       if (i.X != f) {
         dfs(i.X, u, d), w[u] += w[i.X];
if (w[mxson[u]] < w[i.X]) mxson[u] = i.X;</pre>
       } else bln[i.Y] = u, dt[u] = edge[i.Y];
  void cut(int u, int link) {
  data[pl[u] = t++] = dt[u], ulink[u] = link;
     if (!mxson[u]) return;
     cut(mxson[u], link);
     for (auto i : G[u])
       if (i.X != pa[u] && i.X != mxson[u])
         cut(i.X, i.X);
  void build() { dfs(1, 1, 1), cut(1, 1), /*build*/; }
int query(int a, int b) {
     int ta = ulink[a], tb = ulink[b], re = 0;
     while (ta != tb)
       if (deep[ta] < deep[tb])</pre>
         /*query*/, tb = ulink[b = pa[tb]];
       else /*query*/, ta = ulink[a = pa[ta]];
     if (a == b) return re;
     if (pl[a] > pl[b]) swap(a, b);
     /*query*,
     return re:
};
```

#### 3.6 Centroid Decomposition\*

```
struct Cent_Dec { // 1-base
  vector<pli> G[N];
  pll info[N]; // store info. of itself
pll upinfo[N]; // store info. of climbing up
int n, pa[N], layer[N], sz[N], done[N];
  ll dis[__lg(N) + 1][N];
  void init(int _n) {
    n = _n, layer[0] = -1;
    fill_n(pa + 1, n, 0), fill_n(done + 1, n, 0);
    for (int i = 1; i <= n; ++i) G[i].clear();</pre>
  void add_edge(int a, int b, int w) {
    G[a].pb(pll(b, w)), G[b].pb(pll(a, w));
  void get_cent(
    int u, int f, int &mx, int &c, int num) {
    int mxsz = 0:
    sz[u] = 1;
    for (pll e : G[u])
       if (!done[e.X] && e.X != f) {
         get_cent(e.X, u, mx, c, num);
         sz[u] += sz[e.X], mxsz = max(mxsz, sz[e.X]);
    if (mx > max(mxsz, num - sz[u]))
       mx = max(mxsz, num - sz[u]), c = u;
  void dfs(int u, int f, ll d, int org) {
   // if required, add self info or climbing info
    dis[layer[org]][u] = d;
     for (pll e : G[u])
       if (!done[e.X] && e.X != f)
         dfs(e.X, u, d + e.Y, org);
  int cut(int u, int f, int num) {
    int mx = 1e9, c = 0, lc;
    get_cent(u, f, mx, c, num);
    done[c] = 1, pa[c] = f, layer[c] = layer[f] + 1;
    for (pll e : G[c])
```

```
if (!done[e.X]) {
        if (sz[e.X] > sz[c])
         lc = cut(e.X, c, num - sz[c]);
else lc = cut(e.X, c, sz[e.X]);
        upinfo[lc] = pll(), dfs(e.X, c, e.Y, c);
    return done[c] = 0, c;
  void build() { cut(1, 0, n); }
  void modify(int u) {
    for (int a = u, ly = layer[a]; a;
         a = pa[a], --ly) {
      info[a].X += dis[ly][u], ++info[a].Y;
      if (pa[a])
        upinfo[a].X += dis[ly - 1][u], ++upinfo[a].Y;
    }
  ll query(int u) {
    ll rt = 0;
    for (int a = u, ly = layer[a]; a;
         a = pa[a], --ly) {
      rt += info[a].X + info[a].Y * dis[ly][u];
      if (pa[a])
          upinfo[a].X + upinfo[a].Y * dis[ly - 1][u];
    return rt;
};
```

# 3.7 LiChaoST\*

```
struct L {
  ll m, k, id;
  L() : id(-1) {}
  L(ll a, ll b, ll c): m(a), k(b), id(c) {} ll at(ll x) { return m * x + k; }
class LiChao { // maintain max
private:
  int n; vector<L> nodes;
  void insert(int l, int r, int rt, L ln) {
    int m = (l + r) \gg 1;
    if (nodes[rt].id == -1)
      return nodes[rt] = ln, void();
    bool atLeft = nodes[rt].at(l) < ln.at(l);</pre>
    if (nodes[rt].at(m) < ln.at(m))</pre>
      atLeft ^= 1, swap(nodes[rt], ln);
    if (r - l == 1) return;
    if (atLeft) insert(l, m, rt << 1, ln);</pre>
    else insert(m, r, rt << 1 | 1, ln);</pre>
  ll query(int l, int r, int rt, ll x) {
    int m = (l + r) >> 1; ll ret = -INF;
    if (nodes[rt].id != -1) ret = nodes[rt].at(x);
    if (r - l == 1) return ret;
    if (x
         < m) return max(ret, query(l, m, rt << 1, x));</pre>
    return max(ret, query(m, r, rt << 1 | 1, x));</pre>
 ll query(ll x) { return query(0, n, 1, x); }
```

#### 3.8 Link cut tree\*

```
struct Splay { // xor-sum
 static Splay nil;
  Splay *ch[2], *f;
  int val, sum, rev, size;
  Splay (int
      _val = 0) : val(_val), sum(_val), rev(0), size(1)
  {f = ch[0] = ch[1] = &nil;}
 bool isr()
  { return f->ch[0] != this && f->ch[1] != this; }
  int dir()
  { return f->ch[0] == this ? 0 : 1; }
  void setCh(Splay *c, int d) {
    ch[d] = c;
    if (c != &nil) c->f = this;
    pull();
  void give_tag(int r) {
    if (r) swap(ch[0], ch[1]), rev ^= 1;
```

```
void push() {
    if (ch[0] != &nil) ch[0]->give_tag(rev);
if (ch[1] != &nil) ch[1]->give_tag(rev);
  void pull() {
    // take care of the nil!
    size = ch[0]->size + ch[1]->size + 1;
    sum = ch[\theta]->sum ^ ch[1]->sum ^ val;
    if (ch[0] != &nil) ch[0] ->f = this;
    if (ch[1] != &nil) ch[1]->f = this;
} Splay::nil;
Splay *nil = &Splay::nil;
void rotate(Splay *x) {
  Splay *p = x - > f;
  int d = x->dir();
  if (!p->isr()) p->f->setCh(x, p->dir());
  else x->f = p->f;
  p->setCh(x->ch[!d], d);
  x->setCh(p, !d);
  p->pull(), x->pull();
void splay(Splay *x) {
  vector < Splay*> splayVec;
  for (Splay *q = x;; q = q->f) {
    splayVec.pb(q);
    if (q->isr()) break;
  reverse(ALL(splayVec));
  for (auto it : splayVec) it->push();
  while (!x->isr()) {
    if (x->f->isr()) rotate(x);
    else if (x->dir() == x->f->dir())
    rotate(x->f), rotate(x);
else rotate(x), rotate(x);
  }
Splay* access(Splay *x) {
  Splay *q = nil;
  for (; x != nil; x = x->f)
    splay(x), x -> setCh(q, 1), q = x;
  return q;
void root_path(Splay *x) { access(x), splay(x); }
void chroot(Splay *x){
  root_path(x), x->give_tag(1);
  x->push(), x->pull();
void split(Splay *x, Splay *y) {
  chroot(x), root_path(y);
void link(Splay *x, Splay *y) {
  root_path(x), chroot(y);
  x->setCh(y, 1);
void cut(Splay *x, Splay *y) {
  split(x, y);
  if (v->size != 5) return;
  y->push();
  y - ch[0] = y - ch[0] - f = nil;
Splay* get_root(Splay *x) {
  for (root_path(x); x->ch[0] != nil; x = x->ch[0])
   x->push();
  splay(x);
  return x;
bool conn(Splay *x, Splay *y) {
  return get_root(x) == get_root(y);
Splay* lca(Splay *x, Splay *y) {
  access(x), root_path(y);
  if (y->f == nil) return y;
  return y->f;
void change(Splay *x, int val) {
  splay(x), x->val = val, x->pull();
int query(Splay *x, Splay *y) {
 split(x, y);
  return y->sum;
```

#### 3.9 KDTree

```
namespace kdt {
int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn],
  yl[maxn], yr[maxn];
point p[maxn];
int build(int l, int r, int dep = 0) {
  if (l == r) return -1;
function < bool(const point &, const point &) > f =
     [dep](const point &a, const point &b) {
       if (dep & 1) return a.x < b.x;</pre>
       else return a.y < b.y;</pre>
     };
  int m = (l + r) >> 1;
  nth_element(p + l, p + m, p + r, f);
   xl[m] = xr[m] = p[m].x;
   yl[m] = yr[m] = p[m].y;
   lc[m] = build(l, m, dep + 1);
   if (~lc[m]) {
    xl[m] = min(xl[m], xl[lc[m]]);
    xr[m] = max(xr[m], xr[lc[m]]);
yl[m] = min(yl[m], yl[lc[m]]);
     yr[m] = max(yr[m], yr[lc[m]]);
   rc[m] = build(m + 1, r, dep + 1);
   if (~rc[m]) {
    xl[m] = min(xl[m], xl[rc[m]]);
    xr[m] = max(xr[m], xr[rc[m]]);
yl[m] = min(yl[m], yl[rc[m]]);
    yr[m] = max(yr[m], yr[rc[m]]);
   return m;
bool bound(const point &q, int o, long long d) {
   double ds = sqrt(d + 1.0);
   if (q.x < xl[o] - ds || q.x > xr[o] + ds ||
     q.y < yl[o] - ds || q.y > yr[o] + ds)
     return false:
   return true:
long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 1ll * (a.x - b.x) +
  (a.y - b.y) * 1ll * (a.y - b.y);
void dfs(
   const point &q, long long &d, int o, int dep = 0) {
  if (!bound(q, o, d)) return;
long long cd = dist(p[o], q);
   if (cd != 0) d = min(d, cd);
   if ((dep & 1) && q.x < p[o].x ||
     !(dep & 1) && q.y < p[o].y) {
     if (~lc[o]) dfs(q, d, lc[o], dep + 1);
     if (~rc[o]) dfs(q, d, rc[o], dep + 1);
  } else {
     if (~rc[o]) dfs(q, d, rc[o], dep + 1);
if (~lc[o]) dfs(q, d, lc[o], dep + 1);
  }
void init(const vector<point> &v) {
  for (int i = 0; i < v.size(); ++i) p[i] = v[i];</pre>
  root = build(0, v.size());
long long nearest(const point &q) {
  long long res = 1e18;
   dfs(q, res, root);
   return res;
} // namespace kdt
```

# 4 Flow/Matching

# 4.1 Kuhn Munkres\*

```
struct KM { // 0-base
    ll w[N][N], hl[N], hr[N], slk[N];
    int fl[N], fr[N], pre[N], qu[N], ql, qr, n;
    bool vl[N], vr[N];
    void init(int _n) {
        n = _n;
        for (int i = 0; i < n; ++i)
            fill_n(w[i], n, -INF);
    }
    void add_edge(int a, int b, ll wei) {
        w[a][b] = wei;
    }
    bool Check(int x) {
        if (vl[x] = 1, ~fl[x])
            return vr[qu[qr++] = fl[x]] = 1;</pre>
```

```
while (\sim x) swap(x, fr[fl[x] = pre[x]]);
     return 0:
   void bfs(int s) {
     fill_n(slk
     , n, INF), fill_n(vl, n, 0), fill_n(vr, n, 0); ql = qr = 0, qu[qr++] = s, vr[s] = 1;
     for (ll d;;) {
       while (ql < qr)</pre>
         for (int x = 0, y = qu[ql++]; x < n; ++x)
           if (!vl[x] && slk
                [x] >= (d = hl[x] + hr[y] - w[x][y])) {
              if (pre[x] = y, d) slk[x] = d;
              else if (!Check(x)) return;
       d = INF;
       for (int x = 0; x < n; ++x)
         if (!vl[x] && d > slk[x]) d = slk[x];
       for (int x = 0; x < n; ++x) {</pre>
         if (vl[x]) hl[x] += d;
         else slk[x] -= d;
         if (vr[x]) hr[x] -= d;
       for (int x = 0; x < n; ++x)
         if (!vl[x] && !slk[x] && !Check(x)) return;
    }
  ll solve() {
     fill_n(fl
           n, -1), fill_n(fr, n, -1), fill_n(hr, n, 0);
     for (int i = 0; i < n; ++i)</pre>
      hl[i] = *max_element(w[i], w[i] + n);
     for (int i = 0; i < n; ++i) bfs(i);</pre>
     ll res = 0;
     for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
     return res:
};
```

#### 4.2 MincostMaxflow\*

```
struct MinCostMaxFlow { // 0-base
  struct Edge {
    ll from, to, cap, flow, cost, rev;
  } *past[N];
  vector < Edge > G[N];
  int inq[N], n, s, t;
  ll dis[N], up[N], pot[N];
  bool BellmanFord() {
    fill_n(dis, n, INF), fill_n(inq, n, 0);
    queue < int > q;
    auto relax = [&](int u, ll d, ll cap, Edge *e) {
  if (cap > 0 && dis[u] > d) {
         dis[u] = d, up[u] = cap, past[u] = e;
         if (!inq[u]) inq[u] = 1, q.push(u);
      }
    };
    relax(s, 0, INF, 0);
    while (!q.empty()) {
      int u = q.front();
       q.pop(), inq[u] = 0;
       for (auto &e : G[u]) {
         ll d2 = dis[u] + e.cost + pot[u] - pot[e.to];
         relax
              (e.to, d2, min(up[u], e.cap - e.flow), &e);
      }
    }
    return dis[t] != INF;
  void solve(int _s
    , int _t, ll &flow, ll &cost, bool neg = true) {
    s = _s, t = _t, flow = 0, cost = 0;
    if (neg) BellmanFord(), copy_n(dis, n, pot);
    for (; BellmanFord(); copy_n(dis, n, pot)) {
       for (int
       \dot{i} = 0; i < n; ++i) dis[i] += pot[i] - pot[s]; flow += up[t], cost += up[t] * dis[t];
       for (int i = t; past[i]; i = past[i]->from) {
         auto &e = *past[i];
         e.flow += up[t], G[e.to][e.rev].flow -= up[t];
      }
    }
  void init(int n) {
    n = _n, fill_n(pot, n, \theta);
    for (int i = 0; i < n; ++i) G[i].clear();</pre>
```

```
}
void add_edge(ll a, ll b, ll cap, ll cost) {
    G[a].pb(Edge{a, b, cap, 0, cost, SZ(G[b])});
    G[b].pb(Edge{b, a, 0, 0, -cost, SZ(G[a]) - 1});
}
```

# 4.3 Maximum Simple Graph Matching\*

```
struct Matching { // 0-base
  queue<int> q; int n;
  vector<int> fa, s, vis, pre, match;
  vector<vector<int>> G;
  int Find(int u)
  { return u == fa[u] ? u : fa[u] = Find(fa[u]); }
  int LCA(int x, int y) {
     static int tk = 0; tk++; x = Find(x); y = Find(y);
     for (;; swap(x, y)) if (x != n) {
       if (vis[x] == tk) return x;
       vis[x] = tk;
       x = Find(pre[match[x]]);
  void Blossom(int x, int y, int l) {
   for (; Find(x) != l; x = pre[y]) {
       pre[x] = y, y = match[x];
       if (s[y] == 1) q.push(y), s[y] = 0;
       for (int z: {x, y}) if (fa[z] == z) fa[z] = l;
    }
  bool Bfs(int r) {
     iota(ALL(fa), 0); fill(ALL(s), -1);
     q = queue < int > (); q.push(r); s[r] = 0;
     for (; !q.empty(); q.pop()) {
       for (int x = q.front(); int u : G[x])
         if (s[u] == -1) {
            if (pre[u] = x, s[u] = 1, match[u] == n) {
   for (int a = u, b = x, last;
                   b != n; a = last, b = pre[a])
                     match[b], match[b] = a, match[a] = b;
              return true;
            q.push(match[u]); s[match[u]] = 0;
         } else if (!s[u] && Find(u) != Find(x)) {
            int l = LCA(u, x);
            Blossom(x, u, l); Blossom(u, x, l);
     return false;
  \mathsf{Matching}(\mathbf{int} \ \_\mathsf{n}) \ : \ \mathsf{n}(\_\mathsf{n}), \ \mathsf{fa}(\mathsf{n} + 1), \ \mathsf{s}(\mathsf{n} + 1), \ \mathsf{vis}
  (n + 1), pre(n + 1, n), match(n + 1, n), G(n) \{\} void add_edge(int u, int v)
   { G[u].pb(v), G[v].pb(u); }
  int solve() {
     int ans = 0;
     for (int x = 0; x < n; ++x)</pre>
       if (match[x] == n) ans += Bfs(x);
     return ans;
  } // match[x] == n means not matched
}:
```

#### 4.4 Maximum Weight Matching\*

```
#define rep(i, l, r) for (int i=(l); i<=(r); ++i)</pre>
struct WeightGraph { // 1-based
  struct edge { int u, v, w; }; int n, nx;
  vector<int> lab; vector<vector<edge>> g;
  vector < int > slack, match, st, pa, S, vis;
vector < vector < int >> flo, flo_from; queue < int > q;
WeightGraph(int n_) : n(n_), nx(n * 2), lab(nx + 1),
     g(nx + 1, vector < edge > (nx + 1)), slack(nx + 1),
     flo(nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
     match = st = pa = S = vis = slack;
     rep(u, 1, n) rep(v, 1, n) g[u][v] = \{u, v, 0\};
  int ED(edge e)
  { return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2; }
  void update_slack(int u, int x, int &s)
{ if (!s || ED(g[u][x]) < ED(g[s][x])) s = u; }</pre>
  void set_slack(int x) {
     slack[x] = 0;
     for (int u = 1; u <= n; ++u)</pre>
       if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
          update_slack(u, x, slack[x]);
```

```
void q_push(int x) {
  if (x <= n) q.push(x);</pre>
  else for (int y : flo[x]) q_push(y);
void set_st(int x, int b) {
  st[x] = b;
  if (x > n) for (int y : flo[x]) set_st(y, b);
vector<int> split_flo(auto &f, int xr) {
  auto it = find(ALL(f), xr);
  if (auto pr = it - f.begin(); pr % 2 == 1)
    reverse(1 + ALL(f)), it = f.end() - pr;
  auto res = vector(f.begin(), it);
  return f.erase(f.begin(), it), res;
void set_match(int u, int v) {
  match[u] = g[u][v].v;
  if (u <= n) return;</pre>
  int xr = flo_from[u][g[u][v].u];
  auto &f = flo[u], z = split_flo(f, xr);
  rep(i, 0, SZ(z)-1) set_match(z[i], z[i ^ 1]);
  set_match(xr, v); f.insert(f.end(), ALL(z));
void augment(int u, int v) {
  for (;;) {
    int xnv = st[match[u]]; set_match(u, v);
    if (!xnv) return;
    set_match(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
  }
int lca(int u, int v) {
  static int t = 0; ++t;
  for (++t; u || v; swap(u, v)) if (u) {
    if (vis[u] == t) return u;
    vis[u] = t; u = st[match[u]];
    if (u) u = st[pa[u]];
  return 0:
void add_blossom(int u, int o, int v) {
  int b = find(n + 1 + ALL(st), \theta) - begin(st);
  lab[b] = 0, S[b] = 0; match[b] = match[o];
  vector < int > f = {o};
  for (int x = u, y; x != o; x = st[pa[y]])
    f.pb(x), f.pb(y = st[match[x]]), q_push(y);
  reverse(1 + ALL(f));
  for (int x = v, y; x != o; x = st[pa[y]])
  f.pb(x), f.pb(y = st[match[x]]), q_push(y);
  flo[b] = f; set_st(b, b);
  for (int x = 1; x <= nx; ++x)
    g[b][x].w = g[x][b].w = 0;
  fill(ALL(flo_from[b]), 0);
  for (int xs : flo[b]) {
    for (int x = 1; x <= nx; ++x)</pre>
      if (g[
           b][x].w == 0 \mid\mid ED(g[xs][x]) < ED(g[b][x]))
    g[b][x] = g[xs][x], g[x][b] = g[x][xs];
for (int x = 1; x <= n; ++x)
      if (flo_from[xs][x]) flo_from[b][x] = xs;
  set_slack(b);
void expand_blossom(int b) {
  for (int x : flo[b]) set_st(x, x);
  int xr = flo_from[b][g[b][pa[b]].u], xs = -1;
  for (int x : split_flo(flo[b], xr)) {
    if (xs == -1) { xs = x; continue; }
pa[xs] = g[x][xs].u; S[xs] = 1, S[x] = 0;
    slack[xs] = 0; set_slack(x); q_push(x); xs = -1;
  for (int x : flo[b])
    if (x == xr) S[x] = 1, pa[x] = pa[b];
else S[x] = -1, set_slack(x);
  st[b] = 0;
bool on_found_edge(const edge &e) {
  if (int u = st[e.u], v = st[e.v]; S[v] == -1) {
  int nu = st[match[v]]; pa[v] = e.u; S[v] = 1;
    slack[v] = slack[nu] = 0; S[nu] = 0; q_push(nu);
  } else if (S[v] == 0) {
    if (int o = lca(u, v)) add_blossom(u, o, v);
    else return augment(u, v), augment(v, u), true;
  return false:
```

```
bool matching() {
  fill(ALL(S), -1), fill(ALL(slack), 0);
  q = queue < int >();
  for (int x = 1; x <= nx; ++x)</pre>
    if (st[x] == x && !match[x])
  pa[x] = 0, S[x] = 0, q_push(x);
  if (q.empty()) return false;
  for (;;) {
     while (SZ(q)) {
       int u = q.front(); q.pop();
if (S[st[u]] == 1) continue;
       for (int v = 1; v <= n; ++v)</pre>
          if (g[u][v].w > 0 && st[u] != st[v]) {
            if (ED(g[u][v]) != 0)
               update_slack(u, st[v], slack[st[v]]);
            else if
                   (on_found_edge(g[u][v])) return true;
     int d = INF;
    for (int b = n + 1; b <= nx; ++b)</pre>
       if (st[b] == b && S[b] == 1)
         d = min(d, lab[b] / 2);
     for (int x = 1; x <= nx; ++x)</pre>
       if (int
            s = slack[x]; st[x] == x && s && s[x] <= 0)
          d = min(d, ED(g[s][x]) / (S[x] + 2));
     for (int u = 1; u <= n; ++u)
  if (S[st[u]] == 1) lab[u] += d;</pre>
       else if (S[st[u]] == 0) {
  if (lab[u] <= d) return false;</pre>
          lab[u] -= d;
    rep(b, n + 1, nx) if (st[b] == b && S[b] >= 0)
lab[b] += d * (2 - 4 * S[b]);
for (int x = 1; x <= nx; ++x)</pre>
       if (int s = slack[x]; st[x] == x &&
            s \&\& st[s] != x \&\& ED(g[s][x]) == 0)
          if (on_found_edge(g[s][x])) return true;
    for (int b = n + 1; b <= nx; ++b)
  if (st[b] == b && S[b] == 1 && lab[b] == 0)</pre>
          expand_blossom(b);
  return false:
pair<ll, int> solve() {
  fill(ALL(match), 0);
  rep(u, 0, n) st[u] = u, flo[u].clear();
  int w_max = 0;
  rep(u, 1, n) rep(v, 1, n) {
  flo_from[u][v] = (u == v ? u : 0);
     w_{max} = max(w_{max}, g[u][v].w);
  fill(ALL(lab), w_max);
int n_matches = 0; ll tot_weight = 0;
  while (matching()) ++n_matches;
  rep(u, 1, n) if (match[u] && match[u] < u)</pre>
     tot_weight += g[u][match[u]].w;
  return make_pair(tot_weight, n_matches);
void add_edge(int u, int v, int w)
\{ g[u][v].w = g[v][u].w = w; \}
```

#### 4.5 SW-mincut

```
struct SW{ // global min cut, O(V^3)
 #define REP for (int i = 0; i < n; ++i)
static const int MXN = 514, INF = 2147483647;
  int vst[MXN], edge[MXN][MXN], wei[MXN];
  void init(int n) {
    REP fill_n(edge[i], n, 0);
  void addEdge(int u, int v, int w){
    edge[u][v] += w; edge[v][u] += w;
  int search(int &s, int &t, int n){
   fill_n(vst, n, 0), fill_n(wei, n, 0);
    s = t = -1:
    int mx, cur;
    for (int j = 0; j < n; ++j) {
      mx = -1, cur = 0;
      REP if (wei[i] > mx) cur = i, mx = wei[i];
      vst[cur] = 1, wei[cur] = -1;
      s = t; t = cur;
```

```
REP if (!vst[i]) wei[i] += edge[cur][i];
}
return mx;
}
int solve(int n) {
  int res = INF;
  for (int x, y; n > 1; n--){
    res = min(res, search(x, y, n));
    REP edge[i][x] = (edge[x][i] += edge[y][i]);
    REP {
      edge[y][i] = edge[n - 1][i];
      edge[i][y] = edge[i][n - 1];
    } // edge[y][y] = 0;
}
return res;
}
sw;
```

# 4.6 BoundedFlow\*(Dinic\*)

```
struct BoundedFlow { // 0-base
  struct edge {
    int to, cap, flow, rev;
  vector<edge> G[N];
  int n, s, t, dis[N], cur[N], cnt[N];
  void init(int _n) {
    n = _n;
for (int i = 0; i < n + 2; ++i)</pre>
      G[i].clear(), cnt[i] = 0;
  void add_edge(int u, int v, int lcap, int rcap) {
    cnt[u] -= lcap, cnt[v] += lcap;
G[u].pb(edge{v, rcap, lcap, SZ(G[v])});
    G[v].pb(edge{u, 0, 0, SZ(G[u]) - 1});
  void add_edge(int u, int v, int cap) {
    G[u].pb(edge{v, cap, 0, SZ(G[v])});
G[v].pb(edge{u, 0, 0, SZ(G[u]) - 1});
  int dfs(int u, int cap) {
    if (u == t || !cap) return cap;
    for (int &i = cur[u]; i < SZ(G[u]); ++i) {</pre>
       edge &e = G[u][i];
       if (dis[e.to] == dis[u] + 1 && e.cap != e.flow) {
         int df = dfs(e.to, min(e.cap - e.flow, cap));
         if (df) {
           e.flow += df, G[e.to][e.rev].flow -= df;
           return df;
         }
      }
    dis[u] = -1;
    return 0;
  bool bfs() {
    fill_n(dis, n + 3, -1);
    queue < int > q;
    q.push(s), dis[s] = 0;
    while (!q.empty()) {
  int u = q.front();
       q.pop();
       for (edge &e : G[u])
         if (!~dis[e.to] && e.flow != e.cap)
           q.push(e.to), dis[e.to] = dis[u] + 1;
    return dis[t] != -1;
  int maxflow(int _s, int _t) {
    s = _s, t = _t;
int flow = 0, df;
    while (bfs()) {
      fill_n(cur, n + 3, 0);
while ((df = dfs(s, INF))) flow += df;
    return flow;
  bool solve() {
    int sum = 0;
    for (int i = 0; i < n; ++i)</pre>
      if (cnt[i] > 0)
         add_edge(n + 1, i, cnt[i]), sum += cnt[i];
       else if (cnt[i] < 0) add_edge(i, n + 2, -cnt[i]);</pre>
    if (sum != maxflow(n + 1, n + 2)) sum = -1;
    for (int i = 0; i < n; ++i)</pre>
      if (cnt[i] > 0)
```

```
G[n + 1].pop_back(), G[i].pop_back();
else if (cnt[i] < 0)
    G[i].pop_back(), G[n + 2].pop_back();
return sum != -1;
}
int solve(int _s, int _t) {
   add_edge(_t, _s, INF);
   if (!solve()) return -1; // invalid flow
   int x = G[_t].back().flow;
   return G[_t].pop_back(), G[_s].pop_back(), x;
}
};</pre>
```

# 4.7 Gomory Hu tree\*

```
MaxFlow Dinic;
int g[MAXN];
void GomoryHu(int n) { // 0-base
  fill_n(g, n, 0);
  for (int i = 1; i < n; ++i) {
    Dinic.reset();
    add_edge(i, g[i], Dinic.maxflow(i, g[i]));
    for (int j = i + 1; j <= n; ++j)
        if (g[j] == g[i] && ~Dinic.dis[j])
        g[j] = i;
  }
}</pre>
```

#### 4.8 Minimum Cost Circulation\*

```
struct MinCostCirculation { // 0-base
  struct Edge {
     ll from, to, cap, fcap, flow, cost, rev;
  } *past[N]:
  vector < Edge > G[N];
  ll dis[N], inq[N], n;
  void BellmanFord(int s) {
    fill_n(dis, n, INF), fill_n(inq, n, 0);
    queue<int> q;
     auto relax = [&](int u, ll d, Edge *e) {
       if (dis[u] > d) {
         dis[u] = d, past[u] = e;
         if (!inq[u]) inq[u] = 1, q.push(u);
      }
    };
     relax(s, 0, 0);
    while (!q.empty()) {
      int u = q.front();
       q.pop(), inq[u] = 0;
       for (auto &e : G[u])
         if (e.cap > e.flow)
           relax(e.to, dis[u] + e.cost, &e);
  }
  void try_edge(Edge &cur) {
    if (cur.cap > cur.flow) return ++cur.cap, void();
     BellmanFord(cur.to);
    if (dis[cur.from] + cur.cost < 0) {</pre>
       ++cur.flow, --G[cur.to][cur.rev].flow;
       for (int
            i = cur.from; past[i]; i = past[i]->from) {
         auto &e = *past[i];
         ++e.flow, --G[e.to][e.rev].flow;
      }
    }
     ++cur.cap;
  void solve(int mxlg) {
  for (int b = mxlg; b >= 0; --b) {
       for (int i = 0; i < n; ++i)</pre>
         for (auto &e : G[i])
       e.cap *= 2, e.flow *= 2;

for (int i = 0; i < n; ++i)
         for (auto &e : G[i])
           if (e.fcap >> b & 1)
             try_edge(e);
    }
  void init(int _n) { n = _n;
    for (int i = 0; i < n; ++i) G[i].clear();</pre>
  void add_edge(ll a, ll b, ll cap, ll cost) {
    G[a].pb(Edge
         {a, b, 0, cap, 0, cost, SZ(G[b]) + (a == b)});
    G[b].pb(Edge{b, a, 0, 0, 0, -cost, SZ(G[a]) - 1});
} mcmf; // O(VE * ElogC)
```

#### 4.9 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source S and sink T.
  - 2. For each edge (x,y,l,u), connect  $x \rightarrow y$  with capacity u-l.
  - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If in(v)>0, connect  $S\to v$  with capacity in(v), otherwise, connect  $v\to T$  with capacity -in(v).
    - To maximize, connect  $t \to s$  with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from s to t is the answer.
    - To minimize, let f be the maximum flow from S to T. Connect  $t \to s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$ , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e+f_e$ , where  $f_e$  corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
  - 1. Redirect every edge:  $y \rightarrow x$  if  $(x,y) \in M$ ,  $x \rightarrow y$  otherwise.
  - 2. DFS from unmatched vertices in X.
  - 3.  $x \in X$  is chosen iff x is unvisited.
  - 4.  $y \in Y$  is chosen iff y is visited.
- · Minimum cost cyclic flow
  - 1. Consruct super source  ${\cal S}$  and sink  ${\cal T}$
  - 2. For each edge (x,y,c), connect  $x \to y$  with (cost,cap)=(c,1) if c>0, otherwise connect  $y \to x$  with (cost,cap)=(-c,1)
  - 3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
  - 4. For each vertex v with d(v)>0, connect  $S\to v$  with (cost,cap)=(0,d(v))
  - 5. For each vertex v with d(v) < 0 , connect  $v \rightarrow T$  with (cost, cap) = (0, -d(v))
  - 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer  ${\cal T}$
  - 2. Construct a max flow model, let K be the sum of all weights
  - 3. Connect source  $s \rightarrow v$ ,  $v \in G$  with capacity K
  - 4. For each edge (u,v,w) in G, connect  $u \to v$  and  $v \to u$  with capacity w
  - 5. For  $v\in G$ , connect it with sink  $v\to t$  with capacity  $K+2T-(\sum_{e\in E(v)}w(e))-2w(v)$
  - 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v' , and connect  $u' \to v'$  with weight w(u,v) .
  - 2. Connect  $v \to v'$  with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to v.
- 3. Find the minimum weight perfect matching on G'.
- Project selection problem
  - 1. If  $p_v>0$ , create edge (s,v) with capacity  $p_v$ ; otherwise, create edge (v,t) with capacity  $-p_v$ .
  - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
  - 3. The mincut is equivalent to the maximum profit of a subset of projects.
- · Dual of minimum cost maximum flow
  - 1. Capacity  $c_{uv}$  , Flow  $f_{uv}$  , Cost  $w_{uv}$  , Required Flow difference for vertex  $b_u$  .
  - 2. If all  $w_{uv}$  are integers, then optimal solution can happen when all  $p_u$  are integers.

$$\begin{aligned} & \min \sum_{uv} w_{uv} f_{uv} \\ & -f_{uv} \geq -c_{uv} \Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv}) \\ & \sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_{u} \end{aligned}$$

# 5 String

#### 5.1 KMP

```
int F[MAXN];
vector<int> match(string A, string B) {
  vector<int> ans;
  F[0] = -1, F[1] = 0;
  for (int i = 1, j = 0; i < SZ(B); F[++i] = ++j) {
    if (B[i] == B[j]) F[i] = F[j]; // optimize
    while (j != -1 && B[i] != B[j]) j = F[j];
  }
  for (int i = 0, j = 0; i < SZ(A); ++i) {
    while (j != -1 && A[i] != B[j]) j = F[j];
    if (++j == SZ(B)) ans.pb(i + 1 - j), j = F[j];
  }
  return ans;
}</pre>
```

#### 5.2 Z-value\*

```
int z[MAXn];
void make_z(const string &s) {
  int l = 0, r = 0;
  for (int i = 1; i < SZ(s); ++i) {
    for (z[i] = max(0, min(r - i + 1, z[i - l]));
        i + z[i] < SZ(s) && s[i + z[i]] == s[z[i]];
        ++z[i])
    ;
  if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
  }
}
```

#### 5.3 Manacher\*

#### 5.4 **SAIS\***

```
\begin{tabular}{ll} \textbf{namespace} & \textbf{sfx} & \{ \end{tabular}
bool _t[N * 2];
int SA[N * 2], H[N], RA[N];
int _s[N * 2], _c[N * 2], x[N], _p[N], _q[N * 2];
// zero based, string content MUST > 0
// SA[i]: SA[i]-th
     suffix is the i-th lexigraphically smallest suffix.
// H[i]: longest
     common prefix of suffix SA[i] and suffix SA[i - 1].
void pre(int *sa, int *c, int n, int z)
{ fill_n(sa, n, \theta), copy_n(c, z, x); }
void induce
     (int *sa, int *c, int *s, bool *t, int n, int z) {
  copy_n(c, z - 1, x + 1);
  for (int i = 0; i < n; ++i)
if (sa[i] && !t[sa[i] - 1])
       sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
  copy_n(c, z, x);
for (int i = n - 1; i >= 0; --i)
     if (sa[i] && t[sa[i] - 1])
       sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa
  , int *p, int *q, bool *t, int *c, int n, int z) { bool uniq = t[n - 1] = true;
  int nn = 0.
       nmxz = -1, *nsa = sa + n, *ns = s + n, last = -1;
  fill_n(c, z, \theta);
  for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;</pre>
  partial_sum(c, c + z, c);
  if (uniq) {
     for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;</pre>
     return;
  for (int i = n - 2; i >= 0; --i)
     t[i] = (
         s[i] == s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
  pre(sa, c, n, z);
for (int i = 1; i <= n - 1; ++i)</pre>
     if (t[i] && !t[i - 1])
       sa[--x[s[i]]] = p[q[i] = nn++] = i;
  induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i)</pre>
     if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
       bool neq = last < 0 || !equal</pre>
            (s + sa[i], s + p[q[sa[i]] + 1], s + last);
       ns[q[last = sa[i]]] = nmxz += neq;
    }
  sais(ns,
        nsa, p + nn, q + n, t + n, c + z, nn, nmxz + 1);
  pre(sa, c, n, z);
  for (int i = nn - 1; i >= 0; --i)
```

```
sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
induce(sa, c, s, t, n, z);
}
void mkhei(int n) {
  for (int i = 0, j = 0; i < n; ++i) {
    if (RA[i])
      for (; _s[i + j] == _s[SA[RA[i] - 1] + j]; ++j);
    H[RA[i]] = j, j = max(0, j - 1);
  }
}
void build(int *s, int n) {
  copy_n(s, n, _s), _s[n] = 0;
  sais(_s, SA, _p, _q, _t, _c, n + 1, 256);
  copy_n(SA + 1, n, SA);
  for (int i = 0; i < n; ++i) RA[SA[i]] = i;
  mkhei(n);
}
</pre>
```

## 5.5 Aho-Corasick Automatan

```
const int len = 400000, sigma = 26;
struct AC_Automatan {
  int nx[len][sigma], fl[len], cnt[len], pri[len], top;
  int newnode() {
    fill(nx[top], nx[top] + sigma, -1);
    return top++;
  void init() { top = 1, newnode(); }
  int input(
    string &s) { // return the end_node of string
    int X = 1;
    for (char c : s) {
      if (!~nx[X][c - 'a']) nx[X][c - 'a'] = newnode();
      X = nx[X][c - 'a'];
    }
    return X;
  void make fl() {
    queue < int > q;
q.push(1), fl[1] = 0;
    for (int t = 0; !q.empty();) {
      int R = q.front();
      q.pop(), pri[t++] = R;
for (int i = 0; i < sigma; ++i)</pre>
        if (~nx[R][i]) {
          int X = nx[R][i], Z = fl[R];
           for (; Z && !~nx[Z][i];) Z = fl[Z];
           fl[X] = Z ? nx[Z][i] : 1, q.push(X);
    }
  void get_v(string &s) {
    int X = 1;
    fill(cnt, cnt + top, 0);
    for (char c : s) {
      while (X \&\& !~nx[X][c - 'a']) X = fl[X];
      X = X ? nx[X][c -
                          'a'] : 1, ++cnt[X];
    for (int i = top - 2; i > 0; --i)
      cnt[fl[pri[i]]] += cnt[pri[i]];
```

#### 5.6 Smallest Rotation

```
string mcp(string s) {
  int n = SZ(s), i = 0, j = 1;
  s += s;
  while (i < n && j < n) {
    int k = 0;
    while (k < n && s[i + k] == s[j + k]) ++k;
    if (s[i + k] <= s[j + k]) j += k + 1;
    else i += k + 1;
    if (i == j) ++j;
}
int ans = i < n ? i : j;
  return s.substr(ans, n);
}</pre>
```

## 5.7 De Bruijn sequence\*

```
constexpr int MAXC = 10, MAXN = 1e5 + 10;
struct DBSeq {
  int C, N, K, L, buf[MAXC * MAXN]; // K <= C^N
  void dfs(int *out, int t, int p, int &ptr) {
    if (ptr >= L) return;
```

```
if (t > N) {
        if (N % p) return;
        for (int i = 1; i <= p && ptr < L; ++i)</pre>
          out[ptr++] = buf[i];
         buf[t] = buf[t - p], \ dfs(out, t + 1, p, ptr); \\  for (int j = buf[t - p] + 1; \ j < C; ++j) 
          buf[t] = j, dfs(out, t + 1, t, ptr);
   void solve(int _c, int _n, int _k, int *out) {
     int p = 0;
     C = _{c}, N = _{n}, K = _{k}, L = N + K - 1;
dfs(out, 1, 1, p);
     if (p < L) fill(out + p, out + L, 0);</pre>
} dbs;
```

## 5.8 Extended SAM\*

```
struct exSAM {
  int len[N * 2], link[N * 2]; // maxlength, suflink
  int next[N * 2][CNUM], tot; // [0, tot), root = 0
  int lenSorted[N * 2]; // topo. order
  int cnt[N * 2]; // occurence
  int newnode() {
    fill_n(next[tot], CNUM, 0);
    len[tot] = cnt[tot] = link[tot] = 0;
    return tot++;
  void init() { tot = 0, newnode(), link[0] = -1; }
  int insertSAM(int last, int c) {
    int cur = next[last][c];
    len[cur] = len[last] + 1;
    int p = link[last];
    while (p != -1 && !next[p][c])
  next[p][c] = cur, p = link[p];
if (p == -1) return link[cur] = 0, cur;
    int q = next[p][c];
    if (len
        [p] + 1 == len[q]) return link[cur] = q, cur;
    int clone = newnode();
    for (int i = 0; i < CNUM; ++i)</pre>
      next[
           clone][i] = len[next[q][i]] ? next[q][i] : 0;
    len[clone] = len[p] + 1;
    while (p != -1 && next[p][c] == q)
      next[p][c] = clone, p = link[p];
    link[link[cur] = clone] = link[q];
    link[q] = clone;
    return cur;
  void insert(const string &s) {
    int cur = 0;
    for (auto ch : s) {
      int &nxt = next[cur][int(ch - 'a')];
      if (!nxt) nxt = newnode();
      cnt[cur = nxt] += 1;
    }
  void build() {
    queue < int > q;
    q.push(0);
    while (!q.empty()) {
      int cur = q.front();
      q.pop();
       for (int i = 0; i < CNUM; ++i)</pre>
        if (next[cur][i])
           q.push(insertSAM(cur, i));
    vector<int> lc(tot);
    for (int i = 1; i < tot; ++i) ++lc[len[i]];</pre>
    partial_sum(ALL(lc), lc.begin());
    for (int i
         = 1; i < tot; ++i) lenSorted[--lc[len[i]]] = i;
  void solve() {
    for (int i = tot - 2; i >= 0; --i)
      cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
};
```

## 5.9 PalTree\*

```
struct palindromic_tree {
 struct node {
```

```
int next[26], fail, len;
    int cnt, num; // cnt: appear times, num: number of
    // pal. suf.
node(int l = 0) : fail(0), len(l), cnt(0), num(0) {
      for (int i = 0; i < 26; ++i) next[i] = 0;</pre>
    }
  };
  vector<node> St;
  vector < char > s;
  int last, n;
  palindromic_tree() : St(2), last(1), n(0) {
    St[0].fail = 1, St[1].len = -1, s.pb(-1);
  inline void clear() {
    St.clear(), s.clear(), last = 1, n = 0;
    St.pb(0), St.pb(-1);
St[0].fail = 1, s.pb(-1);
  inline int get_fail(int x) {
    while (s[n - St[x].len - 1] != s[n])
      x = St[x].fail;
    return x;
  inline void add(int c) {
  s.push_back(c -= 'a'), ++n;
    int cur = get_fail(last);
    if (!St[cur].next[c]) {
      int now = SZ(St);
      St.pb(St[cur].len + 2);
      St[now].fail =
         St[get_fail(St[cur].fail)].next[c];
      St[cur].next[c] = now;
      St[now].num = St[St[now].fail].num + 1;
    last = St[cur].next[c], ++St[last].cnt;
  inline void count() { // counting cnt
    auto i = St.rbegin();
    for (; i != St.rend(); ++i) {
      St[i->fail].cnt += i->cnt;
    }
  inline int size() { // The number of diff. pal.
    return SZ(St) - 2;
};
```

#### Math 6

# 6.1 ax+by=gcd(only exgcd \*)

```
pll exgcd(ll a, ll b) {
  if (b == 0) return pll(1, 0);
  ll p = a / b;
   pll q = exgcd(b, a % b);
   return pll(q.Y, q.X - q.Y * p);
}
/* ax+by=res, let x be minimum non-negative
g, p = gcd(a, b), exgcd(a, b) * res / g

if p.X < 0: t = (abs(p.X) + b / g - 1) / (b / g)
else: t = -(p.X / (b / g))

p += (b / g, -a / g) * t */
```

# 6.2 Floor and Ceil

```
int floor(int a. int b)
{ return a / b - (a % b && (a < 0) ^ (b < 0)); } int ceil(int a, int b)
{ return a / b + (a % b && (a < 0) ^ (b > 0)); }
```

#### 6.3 Floor Enumeration

```
'/ enumerating x = floor(n / i), [l, r]
 for (int l = 1, r; l <= n; l = r + 1) {</pre>
  int x = n / l;
  r = n / x;
}
```

#### 6.4 Mod Min

```
// min\{k \mid l \ll ((ak) \mod m) \ll r\}, no solution -> -1
ll mod_min(ll a, ll m, ll l, ll r) {
  if (a == 0) return l ? -1 : 0;
  if (ll k = (l + a - 1) / a; k * a <= r)
    return k;
  ll b = m / a, c = m % a;
```

```
if (ll y = mod_min(c, a, a - r % a, a - l % a))
    return (l + y * c + a - 1) / a + y * b;
    return -1;
}
```

# 6.5 Gaussian integer gcd

```
cpx gaussian_gcd(cpx a, cpx b) {
#define rnd
    (a, b) ((a >= 0 ? a * 2 + b : a * 2 - b) / (b * 2))
    ll c = a.real() * b.real() + a.imag() * b.imag();
    ll d = a.imag() * b.real() - a.real() * b.imag();
    ll r = b.real() * b.real() + b.imag() * b.imag();
    if (c % r == 0 && d % r == 0) return b;
    return gaussian_gcd
        (b, a - cpx(rnd(c, r), rnd(d, r)) * b);
}
```

#### 6.6 Miller Rabin\*

# 6.7 Simultaneous Equations

```
struct matrix { //m variables, n equations
  fraction M[MAXN][MAXN + 1], sol[MAXN];
  int solve() { //-1: inconsistent, >= 0: rank
    for (int i = 0; i < n; ++i) {</pre>
      int piv = 0;
       while (piv < m && !M[i][piv].n) ++piv;</pre>
       if (piv == m) continue;
      for (int j = 0; j < n; ++j) {
         if (i == j) continue;
         fraction tmp = -M[j][piv] / M[i][piv];
for (int k = 0; k <=</pre>
               m; ++k) M[j][k] = tmp * M[i][k] + M[j][k];
      }
    }
    int rank = 0;
    for (int i = 0; i < n; ++i) {</pre>
      int piv = 0;
      while (piv < m && !M[i][piv].n) ++piv;</pre>
      if (piv == m && M[i][m].n) return -1;
       else if (piv
            < m) ++rank, sol[piv] = M[i][m] / M[i][piv];</pre>
    return rank;
  }
};
```

#### 6.8 Pollard Rho\*

```
map<ll, int> cnt;
void PollardRho(ll n) {
  if (n == 1) return;
  if (prime(n)) return ++cnt[n], void();
  if (n % 2
       == 0) return PollardRho(n / 2), ++cnt[2], void();
  ll x = 2, y = 2, d = 1, p = 1;
  #define f(x, n, p) ((mul(x, x, n) + p) % n)
  while (true) {
    if (d != n && d != 1) {
      PollardRho(n / d);
      PollardRho(d);
      return:
    if (d == n) ++p;
    x = f(x, n, p), y = f(f(y, n, p), n, p);
     d = gcd(abs(x - y), n);
}
```

# 6.9 Simplex Algorithm

```
const int MAXN = 11000, MAXM = 405;
const double eps = 1E-10;
double a[MAXN][MAXM], b[MAXN], c[MAXM];
double d[MAXN][MAXM], x[MAXM];
int ix[MAXN + MAXM]; // !!! array all indexed from 0
// max{cx} subject to {Ax<=b,x>=0}
// n: constraints, m: vars !!!
// x[] is the optimal solution vector
// usage :
// value = simplex(a, b, c, N, M);
double simplex(int n, int m){
  fill_n(d[n], m + 1, 0);
fill_n(d[n + 1], m + 1, 0);
  iota(ix, ix + n + m, \theta);
   int r = n, s = m - 1;
   for (int i = 0; i < n; ++i) {</pre>
     for (int j = 0; j < m - 1; ++j) d[i][j] = -a[i][j];
d[i][m - 1] = 1;</pre>
     d[i][m] = b[i];
     if (d[r][m] > d[i][m]) r = i;
  copy_n(c, m - 1, d[n]);
d[n + 1][m - 1] = -1;
for (double dd;; ) {
     if (r < n) {
       swap(ix[s], ix[r + m]);
       d[r][s] = 1.0 / d[r][s];
for (int j = 0; j <= m; ++j)
  if (j != s) d[r][j] *= -d[r][s];</pre>
       for (int i = 0; i <= n + 1; ++i) if (i != r) {</pre>
          for (int j = 0; j <= m; ++j) if (j != s)</pre>
            d[i][j] += d[r][j] * d[i][s];
          d[i][s] *= d[r][s];
       }
     }
     r = s = -1;
     for (int j = 0; j < m; ++j)
  if (s < 0 || ix[s] > ix[j]) {
         if (d[n + 1][j] > eps ||
              (d[n + 1][j] > -eps && d[n][j] > eps))
            s = j;
     if (s < 0) break;</pre>
     for (int i = 0; i < n; ++i) if (d[i][s] < -eps) {</pre>
       if (r < 0 ||
            (dd = d[r][m]
                  / d[r][s] - d[i][m] / d[i][s]) < -eps ||
            (dd < eps && ix[r + m] > ix[i + m]))
         r = i:
     if (r < 0) return -1; // not bounded
  if (d[n + 1][m] < -eps) return -1; // not executable
  double ans = 0;
  fill n(x, m, 0);
  for (int i = m; i <
     n + m; ++i) { // the missing enumerated x[i] = 0 if (ix[i] < m - 1){
       ans += d[i - m][m] * c[ix[i]];
       x[ix[i]] = d[i-m][m];
  return ans;
}
```

#### 6.9.1 Construction

Standard form: maximize  $\mathbf{c}^T\mathbf{x}$  subject to  $A\mathbf{x} \leq \mathbf{b}$  and  $\mathbf{x} \geq 0$ . Dual LP: minimize  $\mathbf{b}^T\mathbf{y}$  subject to  $A^T\mathbf{y} \geq \mathbf{c}$  and  $\mathbf{y} \geq 0$ .  $\bar{\mathbf{x}}$  and  $\bar{\mathbf{y}}$  are optimal if and only if for all  $i \in [1,n]$ , either  $\bar{x}_i = 0$  or  $\sum_{j=1}^m A_{ji}\bar{y}_j = c_i$  holds and for all  $i \in [1,m]$  either  $\bar{y}_i = 0$  or  $\sum_{j=1}^n A_{ij}\bar{x}_j = b_j$  holds.

```
1. In case of minimization, let c'_i = -c_i
```

- 2.  $\sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j$
- 3.  $\sum_{1 \le i \le n}^{-} A_{ji} x_i = b_j$ 
  - $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$ •  $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
- 4. If  $x_i$  has no lower bound, replace  $x_i$  with  $x_i x_i'$

#### 6.10 chineseRemainder

```
ll solve(ll x1, ll m1, ll x2, ll m2) {
  ll g = gcd(m1, m2);
  if ((x2 - x1) % g) return -1; // no sol
  m1 /= g; m2 /= g;
 pll p = exgcd(m1, m2);
ll lcm = m1 * m2 * g;
ll res = p.first * (x2 - x1) * m1 + x1;
 // be careful with overflow
  return (res % lcm + lcm) % lcm;
```

# 6.11 Factorial without prime factor\*

```
1/0(p^k + \log^2 n), pk = p^k
ll prod[MAXP];
ll fac_no_p(ll n, ll p, ll pk) {
  prod[0] = 1;
  for (int i = 1; i <= pk; ++i)</pre>
    if (i % p) prod[i] = prod[i - 1] * i % pk;
    else prod[i] = prod[i - 1];
  ll rt = 1;
  for (; n; n /= p) {
   rt = rt * mpow(prod[pk], n / pk, pk) % pk;
    rt = rt * prod[n % pk] % pk;
  return rt;
} // (n! without factor p) % p^k
```

# 6.12 QuadraticResidue\*

```
int Jacobi(int a, int m) {
  int s = 1;
  for (; m > 1; ) {
   a %= m;
    if (a == 0) return 0;
    const int r = __builtin_ctz(a);
    if ((r & 1) & ((m + 2) & 4)) s = -s;
    a >>= r;
    if (a & m & 2) s = -s;
    swap(a, m);
 }
  return s;
}
int QuadraticResidue(int a, int p) {
 if (p == 2) return a & 1;
  const int jc = Jacobi(a, p);
  if (jc == 0) return 0;
  if (jc == -1) return -1;
  int b, d;
  for (; ; ) {
   b = rand() % p;
d = (1LL * b * b + p - a) % p;
    if (Jacobi(d, p) == -1) break;
  int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
  for (int e = (1LL + p) >> 1; e; e >>= 1) {
   if (e & 1) {
     tmp = (1LL *
         g0 * f0 + 1LL * d * (1LL * g1 * f1 % p)) % p;
      g1 = (1LL * g0 * f1 + 1LL * g1 * f0) % p;
      g0 = tmp;
    }
    tmp = (1LL)
        * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p)) % p; 6.15 Berlekamp Massey
    f1 = (2LL * f0 * f1) % p;
    f0 = tmp;
  return q0;
```

#### 6.13 PiCount\*

```
ll PrimeCount(ll n) { // n \sim 10^13 => < 2s
  if (n <= 1) return 0;</pre>
  int v = sqrt(n), s = (v + 1) / 2, pc = 0;
  vector<int> smalls(v + 1), skip(v + 1), roughs(s);
  vector<ll> larges(s);
  for (int i = 2; i <= v; ++i) smalls[i] = (i + 1) / 2;</pre>
  for (int i = 0; i < s; ++i) {</pre>
    roughs[i] = 2 * i + 1;
    larges[i] = (n / (2 * i + 1) + 1) / 2;
  for (int p = 3; p <= v; ++p) {</pre>
```

```
if (smalls[p] > smalls[p - 1]) {
       int q = p * p;
       if (1LL * q * q > n) break;
       skip[p] = 1;
       for (int i = q; i <= v; i += 2 * p) skip[i] = 1;</pre>
       int ns = 0;
       for (int k = 0; k < s; ++k) {
         int i = roughs[k];
         if (skip[i]) continue;
         ll d = 1LL * i * p;
larges[ns] = larges[k] - (d <= v ? larges
              [smalls[d] - pc] : smalls[n / d]) + pc;
          roughs[ns++] = i;
       }
       s = ns;
       for (int j = v / p; j >= p; --j) {
         int c =
               smalls[j] - pc, e = min(j * p + p, v + 1);
          for (int i = j * p; i < e; ++i) smalls[i] -= c;</pre>
       }
    }
   for (int k = 1; k < s; ++k) {</pre>
     const ll m = n / roughs[k];
ll t = larges[k] - (pc + k -
                                     1);
     for (int l = 1; l < k; ++l) {
       int p = roughs[l];
       if (1LL * p * p > m) break;
       t -= smalls[m / p] - (pc + l - 1);
     larges[0] -= t;
  }
   return larges[0];
}
```

# 6.14 Discrete Log\*

```
int DiscreteLog(int s, int x, int y, int m) {
   constexpr int kStep = 32000;
   unordered_map<int, int> p;
   int b = 1;
   for (int i = 0; i < kStep; ++i) {</pre>
     p[y] = i;
y = 1LL * y * x % m;
     b = 1LL * b * x % m;
   for (int i = 0; i < m + 10; i += kStep) {
   s = 1LL * s * b % m;</pre>
     if (p.find(s) != p.end()) return i + kStep - p[s];
   return -1:
int DiscreteLog(int x, int y, int m) {
   if (m == 1) return 0;
   int s = 1;
   for (int i = 0; i < 100; ++i) {</pre>
     if (s == y) return i;
s = 1LL * s * x % m;
   if (s == y) return 100;
   int p = 100 + DiscreteLog(s, x, y, m);
   if (fpow(x, p, m) != y) return -1;
   return p;
```

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
  vector <T> d(SZ(output) + 1), me, he;
for (int f = 0, i = 1; i <= SZ(output); ++i) {</pre>
     for (int j = 0; j < SZ(me); ++j)
  d[i] += output[i - j - 2] * me[j];</pre>
     if ((d[i] -= output[i - 1]) == 0) continue;
if (me.empty()) {
       me.resize(f = i);
       continue;
     vector<T> o(i - f - 1);
T k = -d[i] / d[f]; o.pb(-k);
     for (T x : he) o.pb(x * k);
     o.resize(max(SZ(o), SZ(me)));
     for (int j = 0; j < SZ(me); ++j) o[j] += me[j];</pre>
     if (i - f + SZ(he) >= SZ(me)) he = me, f = i;
     me = o;
```

```
return me:
```

# 6.16 Primes

```
/* 12721 13331 14341 75577 123457 222557
     556679 999983 1097774749 1076767633 100102021
    999997771 1001010013 1000512343 987654361 999991231
     999888733 98789101 987777733 999991921 1010101333
     1010102101 1000000000039 100000000000037
     2305843009213693951 4611686018427387847
     9223372036854775783 18446744073709551557 */
```

#### 6.17 Theorem

· Cramer's rule

$$ax+by=e \Rightarrow x = \frac{ed-bf}{ad-bc}$$

$$cx+dy=f \Rightarrow y = \frac{af-ec}{ad-bc}$$

Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

Kirchhoff's Theorem

Denote L be a  $n \times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii} = d(i)$ ,  $L_{ij} = -c$  where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .
- Tutte's Matrix

Let D be a n imes n matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniformly at random) if i < j and  $(i,j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the maximum matching on G.

- Cayley's Formula
  - Given a degree sequence  $d_1, d_2, ..., d_n$  for each  $\emph{labeled}$  vertices, there are  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.
  - Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex  $1,2,\ldots,k$  belong to different components. Then  $T_{n,k}\!=\!kn^{n-k-1}$ .
- Erdős–Gallai theorem

A sequence of nonnegative integers  $d_1 \geq \cdots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if

$$d_1+\dots+d_n \text{ is even and } \sum_{i=1}^k d_i \leq k(k-1)+\sum_{i=k+1}^n \min(d_i,k) \text{ holds for every } 1\leq k\leq n.$$

Gale Ryser theorem

A pair of sequences of nonnegative integers  $a_1 \ge \cdots \ge a_n$  and  $b_1, \dots, b_n$  is bigraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$  and  $\sum_{i=1}^k a_i \le \sum_{i=1}^n \min(b_i, k)$  holds for every  $1 \le k \le n$ .

Fulkerson–Chen–Anstee theorem A sequence  $(a_1,\ b_1),\ \dots,\ (a_n,\ b_n)$  of nonnegative integer pairs with  $a_1 \geq \cdots \geq a_n$  is digraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$  and

$$\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i,k-1) + \sum_{i=k+1}^n \min(b_i,k) \text{ holds for every } 1 \leq k \leq n.$$

For simple polygon, when points are all integer, we have  $A = \#\{\text{lattice points in the interior}\} + \frac{\#\{\text{lattice points on the boundary}\}}{2} - 1.$ 

- Möbius inversion formula

  - $f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$   $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$
- - A portion of a sphere cut off by a plane.
  - r: sphere radius, a: radius of the base of the cap, h: height of the cap,  $\theta$ : arcsin(a/r).
  - Volume =  $\pi h^2 (3r h)/3 = \pi h (3a^2 + h^2)/6 = \pi r^3 (2 + \cos \theta)(1 \sin \theta)$
  - Area  $= 2\pi rh = \pi(a^2 + h^2) = 2\pi r^2(1 \cos\theta)$ .
- Lagrange multiplier
  - Optimize  $f(x_1,...,x_n)$  when k constraints  $g_i(x_1,...,x_n) = 0$ .
  - Lagrangian function  $\mathcal{L}(x_1,\ldots,x_n,\lambda_1,\ldots,\lambda_k)=f(x_1,\ldots,x_n)=$  $\sum_{i=1}^k \lambda_i g_i(x_1,...,x_n).$
  - The solution corresponding to the original constrained optimization is always a saddle point of the Lagrangian function.
- Nearest points of two skew lines
  - Line 1: ${m v}_1\!=\!{m p}\!+\!t_1{m d}_1$ - Line 2:  $v_2 = p + t_2 d_2$

- $n = d_1 \times d_2$
- $n_1 = d_1 \times n$
- $\boldsymbol{n}_2 = \boldsymbol{d}_2 \times \boldsymbol{n}$
- $c_1 = p_1 + \frac{(p_2 p_1) \cdot n_2}{d_1 \cdot n_2} d_1$   $c_2 = p_2 + \frac{(p_1 p_2) \cdot n_1}{d_2 \cdot n_1} d_2$

#### 6.18 Estimation

- Estimation
  - The number of divisors of n is at most around 100 for n < 5e4, 500 for n < 1e7, 2000 for n < 1e10, 200000 for n < 1e19.
  - The number of ways of writing n as a sum of positive integers, disregarding the order of the summands. 1,1,2,3,5,7,11,15,22,30 for  $n = 0 \sim 9$ , 627 for n = 20,  $\sim 2e5$  for n = 50,  $\sim 2e8$  for n = 100.
  - Total number of partitions of n distinct elements: B(n)1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, 115975, 678570, 4213597,27644437,190899322,....

# 6.19 Euclidean Algorithms

- $m = \lfloor \frac{an+b}{c} \rfloor$
- Time complexity:  $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ + f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^{n} i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ + g(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c - b - 1, a, m - 1) \\ -h(c, c - b - 1, a, m - 1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \operatorname{mod} c, b \operatorname{mod} c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \operatorname{mod} c, b \operatorname{mod} c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \operatorname{mod} c, b \operatorname{mod} c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ - 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

## 6.20 General Purpose Numbers

$$\begin{array}{l} \bullet \ \, \text{Bernoulli numbers} \\ B_0-1, B_1^{\pm}=\pm\frac{1}{2}, B_2=\frac{1}{6}, B_3=0 \\ \sum_{j=0}^m {m+1 \choose j} B_j=0, \text{EGF is } B(x)=\frac{x}{e^x-1}=\sum_{n=0}^{\infty} B_n \frac{x^n}{n!}. \\ S_m(n)=\sum_{k=1}^n k^m=\frac{1}{m+1}\sum_{k=0}^m {m+1 \choose k} B_k^+ n^{m+1-k}. \end{array}$$

• Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1$$
 
$$S(n,k) = \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} {k \choose i} i^n$$
 
$$x^n = \sum_{i=0}^n S(n,i)(x)_i$$
 • Pentagonal number theorem

• Pentagonal number theorem 
$$\prod_{n=1}^{\infty}(1-x^n)=1+\sum_{k=1}^{\infty}(-1)^k\left(x^{k(3k+1)/2}+x^{k(3k-1)/2}\right)$$
 • Catalan numbers 
$$C_n^{(k)}=\frac{1}{(k-1)n+1}\binom{kn}{n}$$

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$

$$C_n^{(k)}(x) = 1 + x[C_n^{(k)}(x)]^k$$

Eulerian numbers

Number of permutations  $\pi \in S_n$  in which exactly k elements are greater than the previous element. k j:s s.t.  $\pi(j) > \pi(j+1)$ , k+1 j:s s.t.  $\pi(j) \ge j$ , k j:s s.t.  $\pi(j) > j$ .

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n,n-1) = 1$$

$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$$

# 6.21 Tips for Generating Functions

```
 \begin{array}{ll} \bullet & \text{Ordinary Generating Function } A(x) = \sum_{i \geq 0} a_i x^i \\ & - A(rx) \Rightarrow r^n a_n \\ & - A(x) + B(x) \Rightarrow a_n + b_n \\ & - A(x) B(x) \Rightarrow \sum_{i=0}^n a_i b_{n-i} \\ & - A(x)^k \Rightarrow \sum_{i_1 + i_2 + \dots + i_k = n} a_{i_1} a_{i_2} \dots a_{i_k} \\ & - x A(x)' \Rightarrow n a_n \\ & - \frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^n a_i \\ \bullet & \text{Exponential Generating Function } A(x) = \sum_{i \geq 0} \frac{a_i}{i!} x_i \\ & - A(x) + B(x) \Rightarrow a_n + b_n \\ & - A(x) B(x) \Rightarrow \sum_{i=0}^n \binom{n}{i} a_i b_{n-i} \\ & - A(x) B(x) \Rightarrow \sum_{i=0}^n \binom{n}{i} a_i b_{n-i} \\ & - A(x) \Rightarrow \sum_{i_1 + i_2 + \dots + i_k = n} \binom{n}{i_1, i_2, \dots, i_k} a_{i_1} a_{i_2} \dots a_{i_k} \\ & - x A(x) \Rightarrow n a_n \\ \bullet & \text{Special Generating Function} \\ & - (1+x)^n = \sum_{i \geq 0} \binom{n}{i} x^i \\ & - \frac{1}{(1-x)^n} = \sum_{i \geq 0} \binom{n}{i-1} x^i \end{array}
```

# 7 Polynomial

## 7.1 Fast Fourier Transform

```
template < int MAXN >
struct FFT {
    using val_t = complex < double >;
    const double PI = acos(-1);
    val_t w[MAXN];
    FFT() {
        for (int i = 0; i < MAXN; ++i) {
            double arg = 2 * PI * i / MAXN;
            w[i] = val_t(cos(arg), sin(arg));
        }
    void bitrev(val_t *a, int n); // see NTT
    void trans
        (val_t *a, int n, bool inv = false); // see NTT;
    // remember to replace LL with val_t
};</pre>
```

#### 7.2 Number Theory Transform\*

```
//(2^16)+1, 65537, 3
//7*17*(2^23)+1, 998244353, 3
//1255*(2^20)+1, 1315962881, 3
//51*(2^25)+1, 1711276033, 29
template < int MAXN, ll P, ll RT > //MAXN must be 2^k
struct NTT {
  ll w[MAXN];
  ll mpow(ll a, ll n);
  ll minv(ll a) { return mpow(a, P - 2); }
  NTT() {
    ll dw = mpow(RT, (P - 1) / MAXN);
    w[0] = 1;
    for (int
         i = 1; i < MAXN; ++i) w[i] = w[i - 1] * dw % P;
  void bitrev(ll *a, int n) {
    int i = 0:
    for (int j = 1; j < n - 1; ++j) {</pre>
      for (int k = n >> 1; (i ^= k) < k; k >>= 1);
       if (j < i) swap(a[i], a[j]);</pre>
    }
  }
  void operator()(
       ll *a, int n, bool inv = false) { //\theta <= a[i] < P
    bitrev(a, n);
    for (int L = 2; L <= n; L <<= 1) {</pre>
      int dx = MAXN / L, dl = L >> 1;
for (int i = 0; i < n; i += L) {
         for (int
               j = i, x = 0; j < i + dl; ++j, x += dx) {
           ll tmp = a[j + dl] * w[x] % P;
           if ((a[j
                 + dl] = a[j] - tmp) < 0) a[j + dl] += P;
           if ((a[j] += tmp) >= P) a[j] -= P;
      }
    if (inv) {
       reverse(a + 1, a + n);
       ll invn = minv(n);
      for (int
            i = 0; i < n; ++i) a[i] = a[i] * invn % P;
    }
```

```
|};<sup>*</sup>
7.3     Fast Walsh Transform*
```

}

```
/* x: a[j], y: a[j + (L >> 1)]
or: (y += x * op), and: (x += y * op)

xor: (x, y = (x + y) * op, (x - y) * op)

invop: or, and, xor = -1, -1, 1/2 */
void fwt(int *a, int n, int op) { //or
  for (int L = 2; L <= n; L <<= 1)
     for (int i = 0; i < n; i += L)
  for (int j = i; j < i + (L >> 1); ++j)
          a[j + (L >> 1)] += a[j] * op;
const int N = 21;
int f[
     N][1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N];
void
     subset_convolution(int *a, int *b, int *c, int L) {
   // c_k = \sum_{i | j = k, i & j = 0} a_i * b_j
  int n = 1 << L;
  for (int i = 1; i < n; ++i)</pre>
     ct[i] = ct[i & (i - 1)] + 1;
  for (int i = 0; i < n; ++i)</pre>
     f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
  for (int i = 0; i <= L; ++i)
  fwt(f[i], n, 1), fwt(g[i], n, 1);</pre>
  for (int i = 0; i <= L; ++i)</pre>
     for (int j = 0; j <= i; ++j)</pre>
       for (int x = 0; x < n; ++x)
h[i][x] += f[j][x] * g[i - j][x];
  for (int i = 0; i <= L; ++i)</pre>
     fwt(h[i], n, -1);
  for (int i = 0; i < n; ++i)
     c[i] = h[ct[i]][i];
```

# 7.4 Polynomial Operation

```
#define
fi(s, n) for (int i = (int)(s); i < (int)(n); ++i)
template < int MAXN, ll P, ll RT > // MAXN = 2^k
struct Poly : vector < ll > { // coefficients in [0, P)
  using vector<ll>::vector;
  static NTT<MAXN, P, RT> ntt;
int n() const { return (int)size(); } // n() >= 1
  Poly(const Poly &p, int m) : vector<ll>(m) {
    copy_n(p.data(), min(p.n(), m), data());
  Poly& irev()
        { return reverse(data(), data() + n()), *this; }
  Poly& isz(int m) { return resize(m), *this; }
  Poly& iadd(const Poly &rhs) { // n() == rhs.n()
    fi(0, n()) if
          (((*this)[i] += rhs[i]) >= P) (*this)[i] -= P;
    return *this;
  Poly& imul(ll k) {
    fi(0, n()) (*this)[i] = (*this)[i] * k % P;
    return *this:
  Poly Mul(const Poly &rhs) const {
    int m = 1;
    while (m < n() + rhs.n() - 1) m <<= 1;</pre>
    Poly X(*this, m), Y(rhs, m);
ntt(X.data(), m), ntt(Y.data(), m);
    fi(0, m) X[i] = X[i] * Y[i] % P;
    ntt(X.data(), m, true);
    return X.isz(n() + rhs.n() - 1);
  Poly Inv() const { // (*this)[0] != 0, 1e5/95ms
    if (n() == 1) return {ntt.minv((*this)[0])};
    int m = 1;
    while (m < n() * 2) m <<= 1;
    Poly Xi = Poly(*this, (n() + 1) / 2).Inv().isz(m);
    Poly Y(*this, m);
    ntt(Xi.data(), m), ntt(Y.data(), m);
    fi(0, m) {
      Xi[i] *= (2 - Xi[i] * Y[i]) % P;
      if ((Xi[i] %= P) < 0) Xi[i] += P;</pre>
    ntt(Xi.data(), m, true);
    return Xi.isz(n());
  Poly Sqrt()
        const { // Jacobi((*this)[0], P) = 1, 1e5/235ms
```

```
if (n()
      == 1) return {QuadraticResidue((*this)[0], P)};
      X = Poly(*this, (n() + 1) / 2).Sqrt().isz(n());
       X.iadd(Mul(X.Inv()).isz(n())).imul(P / 2 + 1);
pair<Poly, Poly> DivMod
    (const Poly &rhs) const { // (rhs.)back() != 0
  if (n() < rhs.n()) return {{0}, *this};</pre>
  const int m = n() - rhs.n() + 1;
  Poly X(rhs); X.irev().isz(m);
  Poly Y(*this); Y.irev().isz(m);
  Poly Q = Y.Mul(X.Inv()).isz(m).irev();
  X = rhs.Mul(Q), Y = *this;
  fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;
return {0, Y.isz(max(1, rhs.n() - 1))};
Poly Dx() const {
  Poly ret(n() - 1);
  fi(⊖,
      ret.n()) ret[i] = (i + 1) * (*this)[i + 1] % P;
  return ret.isz(max(1, ret.n()));
Poly Sx() const {
  Poly ret(n() + 1);
  fi(0, n())
       ret[i + 1] = ntt.minv(i + 1) * (*this)[i] % P;
  return ret;
Polv
     _tmul(int nn, const Poly &rhs) const {
  Poly Y = Mul(rhs).isz(n() + nn - 1);
  return Poly(Y.data() + n() - 1, Y.data() + Y.n());
vector<ll> _eval(const
    vector<ll> &x, const vector<Poly> &up) const {
  const int m = (int)x.size();
  if (!m) return {};
  vector<Poly> down(m * 2);
  // down[1] = DivMod(up[1]).second;
  down[1] = Poly(up[1])
       .irev().isz(n()).Inv().irev()._tmul(m, *this);
  fi(2, m * 2) down[i]
       = up[i ^ 1]._tmul(up[i].n() - 1, down[i / 2]);
  vector<ll> y(m);
  fi(0, m) y[i] = down[m + i][0];
  return y;
}
static vector<Poly> _tree1(const vector<ll> &x) {
  const int m = (int)x.size();
  vector<Poly> up(m * 2);
  fi(0, m) up[m + i] = \{(x[i] ? P - x[i] : 0), 1\};
  for (int i = m - 1; i
      > 0; --i) up[i] = up[i * 2].Mul(up[i * 2 + 1]);
  return up;
}
vector
    <ll> Eval(const vector<ll> &x) const { // 1e5, 1s
  auto up = _tree1(x); return _eval(x, up);
static Poly Interpolate(const vector
    <ll> &x, const vector<ll> &y) { // 1e5, 1.4s
  const int m = (int)x.size();
  vector<Poly> up = _tree1(x), down(m * 2);
vector<ll> z = up[1].Dx()._eval(x, up);
  fi(0, m) z[i] = y[i] * ntt.minv(z[i]) % P;
  fi(0, m) down[m + i] = {z[i]};
  for (int i = m
      .
1; i > 0; --i) down[i] = down[i * 2].Mul(up[i
* 2 + 1]).iadd(down[i * 2 + 1].Mul(up[i * 2]));
  return down[1];
Poly Ln() const { // (*this)[0] == 1, 1e5/170ms
  return Dx().Mul(Inv()).Sx().isz(n());
Poly Exp() const { // (*this)[0] == 0, 1e5/360ms
  if (n() == 1) return {1};
Poly X = Poly(*this, (n() + 1) / 2).Exp().isz(n());
  Poly Y = X.Ln(); Y[0] = P - 1;
  fi(0, n())
       if ((Y[i] = (*this)[i] - Y[i]) < 0) Y[i] += P;</pre>
  return X.Mul(Y).isz(n());
// M := P(P - 1). If k >= M, k := k % M + M.
```

```
Poly Pow(ll k) const {
    int nz = 0;
    while (nz < n() && !(*this)[nz]) ++nz;</pre>
    if (nz * min(k, (ll)n()) >= n()) return Poly(n());
    if (!k) return Poly(Poly {1}, n());
    Poly X(data() + nz, data() + nz + n() - nz * k);
    const ll c = ntt.mpow(X[\theta], k % (P - 1));
    return X.Ln().imul
         (k % P).Exp().imul(c).irev().isz(n()).irev();
  static ll
      LinearRecursion(const vector<ll> &a, const vector
       <ll> &coef, ll n) { // a_n = |sum c_j a_n(n-j)|
    const int k = (int)a.size();
    assert((int)coef.size() == k + 1);
    Poly C(k + 1), W(Poly \{1\}, k), M = \{0, 1\}; fi(1, k + 1) C[k - i] = coef[i] ? P - coef[i] : 0;
    C[k] = 1;
    while (n)
      if (n % 2) W = W.Mul(M).DivMod(C).second;
      n /= 2, M = M.Mul(M).DivMod(C).second;
    ll ret = 0;
    fi(0, k) ret = (ret + W[i] * a[i]) % P;
    return ret;
  }
};
#undef fi
using Poly_t = Poly<131072 * 2, 998244353, 3>;
template <> decltype(Poly_t::ntt) Poly_t::ntt = {};
```

# 7.5 Value Polynomial

```
struct Poly {
  mint base; // f(x) = poly[x - base]
  vector<mint> poly;
  Poly(mint b = 0, mint x = 0): base(b), poly(1, x) {}
  mint get_val(const mint &x) {
    if (x >= base && x < base + SZ(poly))
      return poly[x - base];
    mint rt = 0;
    vector<mint> lmul(SZ(poly), 1), rmul(SZ(poly), 1);
    for (int i = 1; i < SZ(poly); ++i)</pre>
      lmul[i] = lmul[i - 1] * (x - (base + i - 1));
    for (int i = SZ(poly) - 2; i >= 0; --i)
      rmul[i] = rmul[i + 1] * (x - (base + i + 1));
    for (int i = 0; i < SZ(poly); ++i)</pre>
      rt += poly[i] * ifac[i] * inegfac

[SZ(poly) - 1 - i] * lmul[i] * rmul[i];
    return rt;
  void raise() { // g(x) = sigma\{base:x\} f(x)
    if (SZ(poly) == 1 && poly[0] == 0)
      return:
    mint nw = get_val(base + SZ(poly));
    poly.pb(nw);
    for (int i = 1; i < SZ(poly); ++i)</pre>
      poly[i] += poly[i - 1];
```

#### 7.6 Newton's Method

Given F(x) where

$$F(x) = \sum_{i=0}^{\infty} \alpha_i (x - \beta)^i$$

for  $\beta$  being some constant. Polynomial P such that F(P)=0 can be found iteratively. Denote by  $Q_k$  the polynomial such that  $F(Q_k)=0$  (mod  $x^{2^k}$ ), then

$$Q_{k+1} = Q_k - \frac{F(Q_k)}{F'(Q_k)} \pmod{x^{2^{k+1}}}$$

# 8 Geometry

# 8.1 Default Code

```
typedef pair < double, double > pdd;
typedef pair < pdd, pdd > Line;
struct Cir{ pdd 0; double R; };
const double eps = 1e - 8;
pdd operator + (pdd a, pdd b)
{ return pdd(a.X + b.X, a.Y + b.Y); }
pdd operator - (pdd a, pdd b)
{ return pdd(a.X - b.X, a.Y - b.Y); }
```

```
pdd operator*(pdd a, double b)
{ return pdd(a.X * b, a.Y * b); }
pdd operator/(pdd a, double b)
double dot(pdd a, pdd b)
{ return a.X * b.X + a.Y * b.Y; }
double cross(pdd a, pdd b)
{ return a.X * b.Y - a.Y * b.X; }
double abs2(pdd a)
{ return dot(a, a); }
double abs(pdd a)
{ return sqrt(dot(a, a)); }
int sign(double a)
{ return fabs(a) < eps ? 0 : a > 0 ? 1 : -1; }
int ori(pdd a, pdd b, pdd c)
{ return sign(cross(b - a, c - a)); }
bool collinearity(pdd p1, pdd p2, pdd p3)
{ return sign(cross(p1 - p3, p2 - p3)) == 0; }
bool btw(pdd p1, pdd p2, pdd p3) {
  if (!collinearity(p1, p2, p3)) return 0;
  return sign(dot(p1 - p3, p2 - p3)) <= 0;</pre>
bool seg_intersect(pdd p1, pdd p2, pdd p3, pdd p4) {
  int a123 = ori(p1, p2, p3);
int a124 = ori(p1, p2, p4);
  int a341 = ori(p3, p4, p1);
  int a342 = ori(p3, p4, p2);
  if (a123 == 0 && a124 == 0)
  return btw(p1, p2, p3) || btw(p1, p2, p4) ||
btw(p3, p4, p1) || btw(p3, p4, p2);
return a123 * a124 <= 0 && a341 * a342 <= 0;</pre>
pdd intersect(pdd p1, pdd p2, pdd p3, pdd p4) {
  double a123 = cross(p2 - p1, p3 - p1);
double a124 = cross(p2 - p1, p4 - p1);
  return (p4
      * a123 - p3 * a124) / (a123 - a124); // C^3 / C^2
pdd perp(pdd p1)
{ return pdd(-p1.Y, p1.X); }
pdd projection(pdd p1, pdd p2, pdd p3)
{ return p1 + (
    p2 - p1) * dot(p3 - p1, p2 - p1) / abs2(p2 - p1); }
pdd reflection(pdd p1, pdd p2, pdd p3)
{ return p3 + perp(p2 - p1
    ) * cross(p3 - p1, p2 - p1) / abs2(p2 - p1) * 2; }
pdd linearTransformation
    (pdd p0, pdd p1, pdd q0, pdd q1, pdd r) {
  pdd dp = p1 - p0
      , dq = q1 - q0, num(cross(dp, dq), dot(dp, dq));
  return q0 + pdd(
       cross(r - p0, num), dot(r - p0, num)) / abs2(dp);
} // from line p0--p1 to q0--q1, apply to r
```

# 8.2 PointSegDist\*

```
double PointSegDist(pdd q0, pdd q1, pdd p) {
  if (sign(abs(q0 - q1)) == 0) return abs(q0 - p);
  if (sign(dot(q1 - q0,
        p - q0)) >= 0 && sign(dot(q0 - q1, p - q1)) >= 0)
    return fabs(cross(q1 - q0, p - q0) / abs(q0 - q1));
  return min(abs(p - q0), abs(p - q1));
}
```

# 8.3 Heart

```
pdd circenter
    (pdd p0, pdd p1, pdd p2) { // radius = abs(center)
  p1 = p1 - p0, p2 = p2 - p0;
  double x1 = p1.X, y1 = p1.Y, x2 = p2.X, y2 = p2.Y;
  double m = 2. * (x1 * y2 - y1 * x2);
  center.X = (x1 * x1)
       * y2 - x2 * x2 * y1 + y1 * y2 * (y1 - y2)) / m;
  center.\dot{Y} = (x1 * x2)
       * (x2 - x1) - y1 * y1 * x2 + x1 * y2 * y2) / m;
  return center + p0;
pdd incenter
    (pdd p1, pdd p2, pdd p3) { // radius = area / s * 2
  double a =
      abs(p2 - p3), b = abs(p1 - p3), c = abs(p1 - p2);
  double s = a + b + c;
  return (a * p1 + b * p2 + c * p3) / s;
pdd masscenter(pdd p1, pdd p2, pdd p3)
{ return (p1 + p2 + p3) / 3; }
```

# 8.4 point in circle

```
// return q'
    s relation with circumcircle of tri(p[0],p[1],p[2])
bool in_cc(const array<pll, 3> &p, pll q) {
    __int128 det = 0;
    for (int i = 0; i < 3; ++i)
        det += __int128(abs2(p[i]) - abs2(q)) *
            cross(p[(i + 1) % 3] - q, p[(i + 2) % 3] - q);
    return det > 0; // in: >0, on: =0, out: <0
}</pre>
```

# 8.5 Convex hull\*

#### 8.6 PointInConvex\*

```
bool PointInConvex
    (const vector<pll> &C, pll p, bool strict = true) {
    int a = 1, b = SZ(C) - 1, r = !strict;
    if (SZ(C) == 0) return false;
    if (SZ(C) < 3) return r && btw(C[0], C.back(), p);
    if (ori(C[0], C[a], C[b]) > 0) swap(a, b);
    if (ori
        (C[0], C[a], p) >= r || ori(C[0], C[b], p) <= -r)
        return false;
    while (abs(a - b) > 1) {
        int c = (a + b) / 2;
        (ori(C[0], C[c], p) > 0 ? b : a) = c;
    }
    return ori(C[a], C[b], p) < r;
}</pre>
```

## 8.7 TangentPointToHull\*

```
/* The point should be strictly out of hull
  return arbitrary point on the tangent line */
pii get_tangent(vector<pll> &C, pll p) {
  auto gao = [&](int s) {
    return cyc_tsearch(SZ(C), [&](int x, int y)
      { return ori(p, C[x], C[y]) == s; });
  };
  return pii(gao(1), gao(-1));
} // return (a, b), ori(p, C[a], C[b]) >= 0
```

# 8.8 Intersection of line and convex

```
int TangentDir(vector<pll> &C, pll dir) {
  return cyc_tsearch(SZ(C), [&](int a, int b) {
    return cross(dir, C[a]) > cross(dir, C[b]);
  });
#define cmpL(i) sign(cross(C[i] - a, b - a))
pii lineHull(pll a, pll b, vector<pll> &C) {
  int A = TangentDir(C, a - b);
  int B = TangentDir(C, b - a);
  int n = SZ(C);
  if (cmpL(A) < 0 \mid \mid cmpL(B) > 0)
    return pii(-1, -1); // no collision
  auto gao = [&](int l, int r) {
  for (int t = l; (l + 1) % n != r; ) {
       int m = ((l + r + (l < r? 0 : n)) / 2) % n;
       (cmpL(m) == cmpL(t) ? l : r) = m;
    return (l + !cmpL(r)) % n;
  pii res = pii(gao(B, A), gao(A, B)); // (i, j)
  if (res.X == res.Y) // touching the corner i
  return pii(res.X, -1);
  if (!
       cmpL(res.X) && !cmpL(res.Y)) // along side i, i+1
```

```
switch ((res.X - res.Y + n + 1) % n) {
   case 0: return pii(res.X, res.X);
   case 2: return pii(res.Y, res.Y);
  }
/* crossing sides (i, i+1) and (j, j+1)
crossing corner i is treated as side (i, i+1)
returned
   in the same order as the line hits the convex */
return res;
} // convex cut: (r, l]
```

# 8.9 minMaxEnclosingRectangle\*

```
const double INF = 1e18, qi = acos(-1) / 2 * 3;
pdd solve(vector<pll> &dots) {
#define diff(u, v) (dots[u] - dots[v])
#define vec(v) (dots[v] - dots[i])
  hull(dots);
  double Max = 0, Min = INF, deg;
  int n = SZ(dots);
  dots.pb(dots[0]);
  for (int i = 0, u = 1, r = 1, l = 1; i < n; ++i) {</pre>
    pll nw = vec(i + 1);
    while (cross(nw, vec(u + 1)) > cross(nw, vec(u)))
      u = (u + 1) \% n;
    while (dot(nw, vec(r + 1)) > dot(nw, vec(r)))
      \Gamma = (\Gamma + 1) \% n;
    if (!i) l = (r + 1) % n;
    while (dot(nw, vec(l + 1)) < dot(nw, vec(l)))
      l = (l + 1) \% n;
    Min = min(Min, (double)(dot(nw, vec(r)) - dot
        (nw, vec(l))) * cross(nw, vec(u)) / abs2(nw));
    deg = acos(dot(diff(r
    , l), vec(u)) / abs(diff(r, l)) / abs(vec(u)));
deg = (qi - deg) / 2;
    Max = max(Max, abs(diff))
        (r, l)) * abs(vec(u)) * sin(deg) * sin(deg));
 }
  return pdd(Min, Max);
```

# 8.10 VectorInPoly\*

#### 8.11 PolyUnion\*

```
double rat(pll a, pll b) {
  return sign
      (b.X) ? (double)a.X / b.X : (double)a.Y / b.Y;
} // all poly. should be ccw
double polyUnion(vector<vector<pll>>> &poly) {
  double res = 0;
  for (auto &p : poly)
    for (int a = 0; a < SZ(p); ++a) {</pre>
      pll A = p[a], B = p[(a + 1) \% SZ(p)];
      vector
           <pair<double, int>> segs = {{0, 0}, {1, 0}};
      for (auto &q : poly) {
        if (&p == &q) continue;
        for (int b = 0; b < SZ(q); ++b) {
          pll C = q[b], D = q[(b + 1) \% SZ(q)];
           int sc = ori(A, B, C), sd = ori(A, B, D);
           if (sc != sd && min(sc, sd) < 0) {</pre>
             double sa = cross(D
                  - C, A - C), sb = cross(D - C, B - C);
             segs.emplace back
                 (sa / (sa - sb), sign(sc - sd));
           if (!sc && !sd &&
               &q < &p && sign(dot(B - A, D - C)) > 0) {
             segs.emplace_back(rat(C - A, B - A), 1);
segs.emplace_back(rat(D - A, B - A), -1);
```

```
}
    }
    Sort(ALL(segs));
    for (auto &s : segs) s.X = clamp(s.X, 0.0, 1.0);
    double sum = 0;
    int cnt = segs[0].second;
    for (int j = 1; j < SZ(segs); ++j) {
        if (!cnt) sum += segs[j].X - segs[j - 1].X;
        cnt += segs[j].Y;
    }
    res += cross(A, B) * sum;
    }
    return res / 2;
}</pre>
```

# 8.12 PolyCut

# 8.13 Trapezoidalization

```
template < class T>
struct SweepLine {
  struct cmp {
    cmp(const SweepLine &_swp): swp(_swp) {}
    bool operator()(int a, int b) const {
  if (abs(swp.get_y(a) - swp.get_y(b)) <= swp.eps)</pre>
        return swp.slope_cmp(a, b);
      return swp.get_y(a) + swp.eps < swp.get_y(b);</pre>
    const SweepLine &swp;
   _cmp:
  T curTime, eps, curQ;
  vector<Line> base;
  multiset < int , cmp > sweep;
  multiset<pair<T, int>> event;
  vector<typename multiset<int, cmp>::iterator> its;
  vector
      <typename multiset<pair<T, int>>::iterator> eits;
  bool slope_cmp(int a, int b) const {
    assert(a != -1);
    if (b == -1) return 0;
    return sign(cross(base
        [a].Y - base[a].X, base[b].Y - base[b].X)) < 0;
  T get_y(int idx) const {
    if (idx == -1) return cur0:
    Line l = base[idx];
    if (l.X.X == l.Y.X) return l.Y.Y;
    return ((curTime - l.X.X) * l.Y.Y
        + (l.Y.X - curTime) * l.X.Y) / (l.Y.X - l.X.X);
  void insert(int idx) {
    its[idx] = sweep.insert(idx);
    if (its[idx] != sweep.begin())
      update_event(*prev(its[idx]));
    update_event(idx);
    event.emplace(base[idx].Y.X, idx + 2 * SZ(base));
  void erase(int idx) {
    assert(eits[idx] == event.end());
    auto p = sweep.erase(its[idx]);
    its[idx] = sweep.end();
    if (p != sweep.begin())
      update_event(*prev(p));
  void update_event(int idx) {
    if (eits[idx] != event.end())
      event.erase(eits[idx]);
    eits[idx] = event.end();
    auto nxt = next(its[idx]);
    if (nxt ==
         sweep.end() || !slope_cmp(idx, *nxt)) return;
```

```
auto t = intersect(base[idx].
        X, base[idx].Y, base[*nxt].X, base[*nxt].Y).X;
    if (t + eps < curTime || t</pre>
          >= min(base[idx].Y.X, base[*nxt].Y.X)) return;
    eits[idx] = event.emplace(t, idx + SZ(base));
  void swp(int idx) {
    assert(eits[idx] != event.end());
    eits[idx] = event.end();
    int nxt = *next(its[idx]);
    swap((int&)*its[idx], (int&)*its[nxt]);
    swap(its[idx], its[nxt]);
    if (its[nxt] != sweep.begin())
      update_event(*prev(its[nxt]));
    update_event(idx);
  }
  // only expected to call the functions below
  SweepLine(T t, T e, vector
      <Line> vec): _cmp(*this), curTime(t), eps(e)
        curQ(), base(vec), sweep(_cmp), event(), its(SZ
    (vec), sweep.end()), eits(SZ(vec), event.end()) {
for (int i = 0; i < SZ(base); ++i) {</pre>
      auto &[p, q] = base[i];
      if (p > q) swap(p, q);
      if (p.X <= curTime && curTime <= q.X)</pre>
        insert(i);
      else if (curTime < p.X)</pre>
        event.emplace(p.X, i);
  void setTime(T t, bool ers = false) {
    assert(t >= curTime);
    while (!event.empty() && event.begin()->X <= t) {</pre>
      auto [et, idx] = *event.begin();
      int s = idx / SZ(base);
      idx %= SZ(base);
      if (abs(et - t) <= eps && s == 2 && !ers) break;</pre>
      curTime = et;
      event.erase(event.begin());
      if (s == 2) erase(idx);
      else if (s == 1) swp(idx);
      else insert(idx);
    curTime = t;
  T nextEvent() {
    if (event.empty()) return INF;
    return event.begin()->X;
  int lower_bound(T y) {
    curQ = y;
    auto p = sweep.lower_bound(-1);
    if (p == sweep.end()) return -1;
    return *p:
  }
};
```

#### 8.14 Polar Angle Sort\*

```
int cmp(pll a, pll b, bool same = true) {
#define is_neg(k) (
    sign(k.Y) < 0 || (sign(k.Y) == 0 && sign(k.X) < 0))
int A = is_neg(a), B = is_neg(b);
if (A != B)
    return A < B;
if (sign(cross(a, b)) == 0)
    return same ? abs2(a) < abs2(b) : -1;
return sign(cross(a, b)) > 0;
}
```

# 8.15 Half plane intersection\*

```
sort(ALL(arr), [&](Line a, Line b) -> int {
  if (cmp(a.Y - a.X, b.Y - b.X, 0) != -1)
    return cmp(a.Y - a.X, b.Y - b.X, 0);
  return ori(a.X, a.Y, b.Y) < 0;</pre>
});
deque<Line> dq(1, arr[0]);
for (auto p : arr) {
  if (cmp(
       dq.back().Y - dq.back().X, p.Y - p.X, \theta) == -1)
     continue;
  while (SZ(dq
       ) >= 2 && !isin(p, dq[SZ(dq) - 2], dq.back()))
    dq.pop_back();
  while (SZ(dq) >= 2 && !isin(p, dq[0], dq[1]))
    dq.pop_front();
  dq.pb(p);
while (SZ(dq)
      >= 3 && !isin(dq[0], dq[SZ(dq) - 2], dq.back()))
  dq.pop_back();
while (SZ(dq) >= 3 && !isin(dq.back(), dq[0], dq[1]))
  dq.pop_front();
return vector<Line>(ALL(dq));
```

# 8.16 RotatingSweepLine

```
void rotatingSweepLine(vector<pii> &ps) {
   int n = SZ(ps), m = 0;
   vector<int> id(n), pos(n);
   vector<pii> line(n * (n - 1));
   for (int i = 0; i < n; ++i)</pre>
     for (int j = 0; j < n; ++j)</pre>
   if (i != j) line[m++] = pii(i, j);
sort(ALL(line), [&](pii a, pii b) {
     return cmp(ps[a.Y] - ps[a.X], ps[b.Y] - ps[b.X]);
   }); // cmp(): polar angle compare
   iota(ALL(id), 0);
   sort(ALL(id), [&](int a, int b) {
  if (ps[a].Y != ps[b].Y) return ps[a].Y < ps[b].Y;</pre>
     return ps[a] < ps[b];</pre>
   }); // initial order, since (1, 0) is the smallest
for (int i = 0; i < n; ++i) pos[id[i]] = i;</pre>
   for (int i = 0; i < m; ++i) {</pre>
     auto l = line[i];
     // do something
     tie(pos[l.X], pos[l.Y], id[pos[l.X]], id[pos[l.Y
          ]]) = make_tuple(pos[l.Y], pos[l.X], l.Y, l.X);
}
```

# 8.17 Minimum Enclosing Circle\*

#### 8.18 Intersection of two circles\*

# 8.19 Intersection of polygon and circle\*

```
// Divides into multiple triangle, and sum up
const double PI=acos(-1);
double _area(pdd pa, pdd pb, double r){
  if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
  if(abs(pb)<eps) return 0;</pre>
  double S, h, theta;
  double a=abs(pb),b=abs(pa),c=abs(pb-pa);
  double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
  double cosC = dot(pa,pb) / a / b, C = acos(cosC);
  if(a > r){
    S = (C/2)*r*r;
    h = a*b*sin(C)/c;
    if (h < r && B
        < PI/2) S -= (acos(h/r)*r*r - h*sqrt(r*r-h*h));
 else if(b > r){
    theta = PI - B - asin(sin(B)/r*a);
    S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
  else S = .5*sin(C)*a*b;
 return S;
double area_poly_circle(const
     vector<pdd> poly,const pdd &0,const double r){
  double S=0;
  for(int i=0;i<SZ(poly);++i)</pre>
    S+=\_area(poly[i]-0,poly[(i+1)\%SZ(poly
        )]-0,r)*ori(0,poly[i],poly[(i+1)%SZ(poly)]);
  return fabs(S);
```

#### 8.20 Intersection of line and circle\*

# 8.21 Tangent line of two circles

```
vector<Line
    > go( const Cir& c1 , const Cir& c2 , int sign1 ){
  // sign1 = 1 for outer tang, -1 for inter tang
  vector<Line> ret;
  double d_sq = abs2(c1.0 - c2.0);
  if (sign(d_sq) == 0) return ret;
  double d = sqrt(d_sq);
  pdd v = (c2.0 - c1.0) / d;
  double c = (c1.R - sign1 * c2.R) / d;
  if (c * c > 1) return ret;
  double h = sqrt(max(0.0, 1.0 - c * c));
  for (int sign2 = 1; sign2 >= -1; sign2 -= 2) {
  pdd n = pdd(v.X * c - sign2 * h * v.Y,
      v.Y * c + sign2 * h * v.X);
    pdd p1 = c1.0 + n * c1.R;
    pdd p2 = c2.0 + n * (c2.R * sign1);
    if (sign(p1.X - p2.X) == 0 and
         sign(p1.Y - p2.Y) == 0)
       p2 = p1 + perp(c2.0 - c1.0);
    ret.pb(Line(p1, p2));
  return ret:
}
```

#### 8.22 CircleCover\*

```
const int N = 1021;
struct CircleCover {
  int C;
  Cir c[N];
  bool g[N][N], overlap[N][N];
  // Area[i] : area covered by at least i circles
```

```
double Area[ N ];
  void init(int _C){ C = _C;}
   struct Teve {
     pdd p; double ang; int add;
     Teve() {}
     Teve(pdd _a
     , double _b, int _c):p(_a), ang(_b), add(_c){}
bool operator<(const Teve &a)const</pre>
     {return ang < a.ang;}
  }eve[N * 2];
  // strict: x = 0, otherwise x = -1
bool disjuct(Cir &a, Cir &b, int x)
{return sign(as(a.0 - b.0) - a.R - b.R) > x;}
   bool contain(Cir &a, Cir &b, int x)
  {return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
  bool contain(int i, int j) {
     /* c[j] is non-strictly in c[i]. */
     return (sign
          (c[i].R - c[j].R) > 0 \mid | (sign(c[i].R - c[j].
          R) == 0 \&\& i < j)) \&\& contain(c[i], c[j], -1);
  void solve(){
     fill_n(Area, C + 2, 0);
     for(int i = 0; i < C; ++i)
  for(int j = 0; j < C; ++j)</pre>
         overlap[i][j] = contain(i, j);
     for(int i = 0; i < C; ++i)</pre>
       for(int j = 0; j < C; ++j)</pre>
         g[i][j] = !(overlap[i][j] || overlap[j][i] ||
              disjuct(c[i], c[j], -1));
     for(int i = 0; i < C; ++i){</pre>
       int E = 0, cnt = 1;
       for(int j = 0; j < C; ++j)</pre>
          if(j != i && overlap[j][i])
            ++cnt;
       for(int j = 0; j < C; ++j)</pre>
          if(i != j && g[i][j]) {
            pdd aa, bb;
            CCinter(c[i], c[j], aa, bb);
            double A =
                  atan2(aa.Y - c[i].O.Y, aa.X - c[i].O.X);
            double B =
                  atan2(bb.Y - c[i].0.Y, bb.X - c[i].0.X);
            eve[E++] = Teve
                (bb, B, 1), eve[E++] = Teve(aa, A, -1);
            if(B > A) ++cnt;
       if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
       else{
         sort(eve, eve + E);
          eve[E] = eve[0];
          for(int j = 0; j < E; ++j){</pre>
            cnt += eve[j].add;
            Area[cnt
                 ] += cross(eve[j].p, eve[j + 1].p) * .5;
            double theta = eve[j + 1].ang - eve[j].ang; if (theta < 0) theta += 2. * pi;
            Area[cnt] += (theta
                  - sin(theta)) * c[i].R * c[i].R * .5;
         }
       }
     }
  }
};
```

#### 8.23 3Dpoint\*

```
double dot(Point p1, Point p2)
{ return p1.x * p2.x + p1.y * p2.y + p1.z * p2.z; }
double abs(Point a)
{ return sqrt(dot(a, a)); }
Point cross3(Point a, Point b, Point c)
{ return cross(b - a, c - a); }
double area(Point a, Point b, Point c)
{ return abs(cross3(a, b, c)); }
double volume(Point a, Point b, Point c, Point d)
{ return dot(cross3(a, b, c), d - a); }
//Azimuthal
     angle (longitude) to x-axis in interval [-pi, pi]
double phi(Point p) { return atan2(p.y, p.x); }
     angle (latitude) to the z-axis in interval [0, pi]
double theta(Point p)
    { return atan2(sqrt(p.x * p.x + p.y * p.y), p.z); }
Point masscenter(Point a, Point b, Point c, Point d)
  return (a + b + c + d) / 4; }
pdd proj(Point a, Point b, Point c, Point u) {
// proj. u to the plane of a, b, and c
  Point e1 = b - a;
  Point e2 = c - a;
  e1 = e1 / abs(e1);
  e2 = e2 - e1 * dot(e2, e1);
  e2 = e2 / abs(e2);
  Point p = u - a;
  return pdd(dot(p, e1), dot(p, e2));
Point
     rotate_around(Point p, double angle, Point axis) {
  double s = sin(angle), c = cos(angle);
  Point u = axis / abs(axis);
        * dot(u, p) * (1 - c) + p * c + cross(u, p) * s;
}
```

#### 8.24 Convexhull3D\*

```
struct convex_hull_3D {
struct Face {
  int a, b, c;
  Face(int ta, int tb, int tc): a(ta), b(tb), c(tc) {}
}; // return the faces with pt indexes
vector < Face > res;
vector < Point > P:
convex_hull_3D(const vector<Point> &_P): res(), P(_P) {
  all points coplanar case will WA, O(n^2)
  int n = SZ(P);
  if (n <= 2) return; // be careful about edge case</pre>
  // ensure first 4 points are not coplanar
  swap(P[1], *find_if(ALL(P), [&](
  auto p) { return sign(abs2(P[0] - p)) != 0; }));
swap(P[2], *find_if(ALL(P), [&](auto p) { return
    sign(abs2(cross3(p, P[0], P[1]))) != 0; }));
  swap(P[3], *find_if(ALL(P), [&](auto p) { return
        sign(volume(P[0], P[1], P[2], p)) != 0; }));
  vector<vector<int>> flag(n, vector<int>(n));
  res.emplace_back(0, 1, 2); res.emplace_back(2, 1, 0);
  for (int i = 3; i < n; ++i) {</pre>
    vector < Face > next;
    for (auto f : res) {
      int d
           = sign(volume(P[f.a], P[f.b], P[f.c], P[i]));
      if (d <= 0) next.pb(f);</pre>
      int ff = (d > 0) - (d < 0);
       flag[f.a][
           f.b] = flag[f.b][f.c] = flag[f.c][f.a] = ff;
    for (auto f : res) {
       auto F = [&](int x, int y) {
        if (flag[x][y] > 0 && flag[y][x] \ll 0)
           next.emplace_back(x, y, i);
      F(f.a, f.b); F(f.b, f.c); F(f.c, f.a);
    res = next;
  }
bool same(Face s, Face t) {
  if (sign(volume
       (P[s.a], P[s.b], P[s.c], P[t.a])) != 0) return 0;
  if (sign(volume
       (P[s.a], P[s.b], P[s.c], P[t.b])) != 0) return 0;
  if (sign(volume
       (P[s.a], P[s.b], P[s.c], P[t.c])) != 0) return 0;
```

```
return 1:
int polygon_face_num() {
  for (int i = 0; i < SZ(res); ++i)</pre>
    ans += none_of(res.begin(), res.begin()
         + i, [&](Face g) { return same(res[i], g); });
  return ans:
double get_volume() {
  double ans = 0;
  for (auto f : res)
    ans +=
        volume(Point(0, 0, 0), P[f.a], P[f.b], P[f.c]);
  return fabs(ans / 6);
double get_dis(Point p, Face f) {
  Point p1 = P[f.a], p2 = P[f.b], p3 = P[f.c];
  double a = (p2.y - p1.y)
        * (p3.z - p1.z) - (p2.z - p1.z) * (p3.y - p1.y);
  double b = (p2.z - p1.z)
         (p3.x - p1.x) - (p2.x - p1.x) * (p3.z - p1.z);
  double c = (p2.x - p1.x)
  * (p3.y - p1.y) - (p2.y - p1.y) * (p3.x - p1.x);
double d = 0 - (a * p1.x + b * p1.y + c * p1.z);
  return fabs(a * p.x + b *
      p.y + c * p.z + d) / sqrt(a * a + b * b + c * c);
};
// n^2 delaunay: facets with negative z normal of
// convexhull of (x, y, x^2 + y^2), use a pseudo-point // (0, 0, inf) to avoid degenerate case
```

# 8.25 DelaunayTriangulation\*

```
/* Delaunay Triangulation:
Given a sets of points on 2D plane, find a
triangulation such that no points will strictly
inside circumcircle of any triangle. */
struct Edge {
  int id; // oidx[id]
  list<Edge>::iterator twin;
  Edge(int _id = 0):id(_id) {}
struct Delaunay { // 0-base
  int n, oidx[N];
  list<Edge> head[N]; // result udir. graph
  pll p[N];
  void init(int _n, pll _p[]) {
    n = _n, iota(oidx, oidx + n, 0);
for (int i = 0; i < n; ++i) head[i].clear();</pre>
    sort(oidx, oidx + n, [&](int a, int b)
    { return _p[a] < _p[b]; });
for (int i = 0; i < n; ++i) p[i] = _p[oidx[i]];
    divide(0, n - 1);
  void addEdge(int u, int v) {
    head[u].push_front(Edge(v));
    head[v].push_front(Edge(u));
    head[u].begin()->twin = head[v].begin();
    head[v].begin()->twin = head[u].begin();
  void divide(int l, int r) {
    if (l == r) return;
    if (l + 1 == r) return addEdge(l, l + 1);
    int mid = (l + r) >> 1;
    divide(l, mid), divide(mid + 1, r);
    int nowl = l, nowr = r;
    for (int update = 1; update;) {
      update = 0:
      pll ptL = p[nowl], ptR = p[nowr];
      for (auto it : head[nowl]) {
        pll t = p[it.id];
        int v = ori(ptR, ptL, t);
        if (v > 0 || (v
             == 0 && abs2(ptR - t) < abs2(ptR - ptL))) {
          nowl = it.id, update = 1;
           break;
        }
      if (update) continue;
      for (auto it : head[nowr]) {
        pll t = p[it.id];
        int v = ori(ptL, ptR, t);
        if (v < 0 || (v
             == 0 && abs2(ptL - t) < abs2(ptL - ptR))) {
```

```
nowr = it.id, update = 1;
         break;
       }
     }
    addEdge(nowl, nowr); // add tangent
    while (true) {
      pll ptL = p[nowl], ptR = p[nowr];
      int ch = -1, side = 0;
      for (auto it : head[nowl])
       ch = it.id, side = -1
      for (auto it : head[nowr])
        if (ori(ptR, p[it.id], ptL) > 0 && (ch ==
             -1 || in_cc({ptL, ptR, p[ch]}, p[it.id])))
          ch = it.id, side = 1;
      if (ch == -1) break; // upper common tangent
      if (side == -1) {
        for (auto it = head
            [nowl].begin(); it != head[nowl].end(); )
          if (seg_strict_intersect
              (ptL, p[it->id], ptR, p[ch]))
            head[it->id].erase
               (it->twin), head[nowl].erase(it++);
         else ++it;
       nowl = ch, addEdge(nowl, nowr);
      else {
       for (auto it = head
            [nowr].begin(); it != head[nowr].end(); )
          if (seg_strict_intersect
              (ptR, p[it->id], ptL, p[ch]))
            head[it->id].erase
               (it->twin), head[nowr].erase(it++);
         else ++it;
       nowr = ch, addEdge(nowl, nowr);
   }
 }
} tool;
```

# 8.26 Triangulation Vonoroi\*

#### 8.27 Minkowski Sum\*

```
vector<pll> Minkowski(vector<pll> A, vector<pll> B) {
  hull(A), hull(B);
  vector<pll> C(1, A[0] + B[0]), s1, s2;
  for (int i = 0; i < SZ(A); ++i)
     s1.pb(A[(i + 1) % SZ(A)] - A[i]);
  for (int i = 0; i < SZ(B); i++)
     s2.pb(B[(i + 1) % SZ(B)] - B[i]);
  for (int i = 0, j = 0; i < SZ(A) || j < SZ(B);)
     if (j >= SZ
          (B) || (i < SZ(A) && cross(s1[i], s2[j]) >= 0))
        C.pb(B[j % SZ(B)] + A[i++]);
  else
        C.pb(A[i % SZ(A)] + B[j++]);
  return hull(C), C;
}
```

# 9 Else

## 9.1 Cyclic Ternary Search\*

```
/* bool pred(int a, int b);
f(0) ~ f(n - 1) is a cyclic-shift U-function
return idx s.t. pred(x, idx) is false forall x*/
int cyc_tsearch(int n, auto pred) {
  if (n == 1) return 0;
```

```
int l = 0, r = n; bool rv = pred(1, 0);
while (r - l > 1) {
   int m = (l + r) / 2;
   if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
   else l = m;
}
return pred(l, r % n) ? l : r % n;
}
```

# 9.2 Mo's Alogrithm(With modification)

```
Mo's Algorithm With modification
 Block: N^{2/3}, Complexity: N^{5/3}
 struct Query {
   int L, R, LBid, RBid, T;
   Query(int l, int r, int t):
     L(l), R(r), LBid(l / blk), RBid(r / blk), T(t) {}
   bool operator < (const Query &q) const {</pre>
     if (LBid != q.LBid) return LBid < q.LBid;</pre>
     if (RBid != q.RBid) return RBid < q.RBid;</pre>
     return T < b.T;</pre>
  }
};
 void solve(vector<Query> query) {
   sort(ALL(query));
   int L=0, R=0, T=-1;
   for (auto q : query) {
     while (T < q.T) addTime(L, R, ++T); // TODO
     while (T > q.T) subTime(L, R, T--); // TODO
     while (R < q.R) add(arr[++R]); // TODO</pre>
     while (L > q.L) add(arr[--L]); // TODO
     while (R > q.R) sub(arr[R--]); // TODO
     while (L < q.L) sub(arr[L++]); // TODO</pre>
     // answer query
}
```

# 9.3 Mo's Alogrithm On Tree

```
Mo's Algorithm On Tree
Preprocess:
1) LCA
2) dfs with in[u] = dft++, out[u] = dft++
3) ord[in[u]] = ord[out[u]] = u
4) bitset < MAXN > inset
struct Query {
  int L, R, LBid, lca;
  Query(int u, int v) {
    int c = LCA(u, v);
    if (c == u || c == v)
      q.lca = -1, q.L = out[c ^ u ^ v], q.R = out[c];
    else if (out[u] < in[v])</pre>
      q.lca = c, q.L = out[u], q.R = in[v];
    bool operator < (const Query &q) const {</pre>
    if (LBid != q.LBid) return LBid < q.LBid;</pre>
    return R < q.R;</pre>
  }
};
void flip(int x) {
    if (inset[x]) sub(arr[x]); // TODO
    else add(arr[x]); // TODO
    inset[x] = ~inset[x];
void solve(vector<Query> query) {
  sort(ALL(query));
  int L = 0, R = 0;
  for (auto q : query) {
    while (R < q.R) flip(ord[++R]);
while (L > q.L) flip(ord[--L]);
    while (R > q.R) flip(ord[R--]);
    while (L < q.L) flip(ord[L++]);</pre>
    if (~q.lca) add(arr[q.lca]);
    // answer query
    if (~q.lca) sub(arr[q.lca]);
```

# 9.4 Additional Mo's Algorithm Trick

- · Mo's Algorithm With Addition Only
  - Sort querys same as the normal Mo's algorithm.
  - For each query [l,r]:
  - If l/blk = r/blk, brute-force.
  - If  $l/blk \neq curL/blk$ , initialize  $curL := (l/blk+1) \cdot blk$ , curR := curL-1
  - If r > curR, increase curR
  - decrease curL to fit l, and then undo after answering
- · Mo's Algorithm With Offline Second Time
  - Require: Changing answer  $\equiv$  adding f([l,r],r+1).
  - Require: f([l,r],r+1) = f([1,r],r+1) f([1,l),r+1). Part1: Answer all f([1,r],r+1) first.

  - Part2: Store  $curR \rightarrow R$  for curL (reduce the space to O(N)), and then answer them by the second offline algorithm.
  - Note: You must do the above symmetrically for the left boundaries.

### 9.5 Hilbert Curve

```
ll hilbert(int n, int x, int y) {
  ll res = 0;
  for (int s = n / 2; s; s >>= 1) {
    int rx = (x \& s) > 0;
    int ry = (y \& s) > 0;
    res += s * 1ll * s * ((3 * rx) ^ ry);
    if (ry == 0) {
      if (rx == 1) x = s - 1 - x, y = s - 1 - y;
      swap(x, y);
   }
 }
  return res;
 // n = 2^k
```

# 9.6 DvnamicConvexTrick\*

```
// only works for integer coordinates!! maintain max
struct Line {
  mutable ll a, b, p;
  bool operator
      <(const Line &rhs) const { return a < rhs.a; }
  bool operator<(ll x) const { return p < x; }</pre>
struct DynamicHull : multiset<Line, less<>> {
  static const ll kInf = 1e18;
  ll Div(ll a,
       ll b) { return a / b - ((a ^ b) < 0 && a % b); }
  bool isect(iterator x, iterator y) {
    if (y == end()) { x->p = kInf; return 0; }
    if (x
        ->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
    else x - > p = Div(y - > b - x - > b, x - > a - y - > a);
    return x->p >= y->p;
  void addline(ll a, ll b) {
    auto z = insert({a, b, 0}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin
        () && isect(--x, y)) isect(x, y = erase(y));
    while ((y = x) != begin
        () && (--x)->p >= y->p) isect(x, erase(y));
  ll query(ll x) {
    auto l = *lower_bound(x);
    return l.a * x + l.b;
  }
};
```

# 9.7 All LCS\*

```
void all_lcs(string s, string t) { // 0-base
  vector<int> h(SZ(t));
  iota(ALL(h), 0);
  for (int a = 0; a < SZ(s); ++a) {</pre>
     int v = -1;
     for (int c = 0; c < SZ(t); ++c)</pre>
       if (s[a] == t[c] || h[c] < v)</pre>
     swap(h[c], v);

// LCS(s[0, a], t[b, c]) =

// c - b + 1 - sum([h[i] >= b] / i <= c)
     // h[i] might become -1 !!
}
```

#### 9.8 DLX\*

```
#define TRAV(i, link, start)
     for (int i = link[start]; i != start; i = link[i])
template <
    bool E> // E: Exact, NN: num of 1s, RR: num of rows
struct DLX {
  int lt[NN], rg[NN], up[NN], dn[NN
      ], rw[NN], cl[NN], bt[NN], s[NN], head, sz, ans;
  int rows, columns;
  bool vis[NN];
  bitset<RR> sol, cur; // not sure
  void remove(int c) {
    if (E) lt[rg[c]] = lt[c], rg[lt[c]] = rg[c];
    TRAV(i, dn, c) {
      if (E) {
        TRAV(j, rg, i)
          up[dn[j]]
                = up[j], dn[up[j]] = dn[j], --s[cl[j]];
      } else {
        lt[rg[i]] = lt[i], rg[lt[i]] = rg[i];
      }
    }
  void restore(int c) {
    TRAV(i, up, c) {
      if (E) {
        TRAV(j,
          ++s[cl[j]], up[dn[j]] = j, dn[up[j]] = j;
      } else {
        lt[rg[i]] = rg[lt[i]] = i;
      }
    if (E) lt[rg[c]] = c, rg[lt[c]] = c;
  void init(int c) {
    rows = 0, columns = c;

for (int i = 0; i < c;
                            ++i) {
      up[i] = dn[i] = bt[i] = i;
      lt[i] = i == 0 ? c : i - 1;
rg[i] = i == c - 1 ? c : i + 1;
      s[i] = 0;
    rg[c] = 0, lt[c] = c - 1;
    up[c] = dn[c] = -1;
    head = c, sz = c + 1;
  void insert(const vector<int> &col) {
    if (col.empty()) return;
    int f = sz;
    for (int i = 0; i < (int)col.size(); ++i) {</pre>
      int c = col[i], v = sz++;
      dn[bt[c]] = v;
      up[v] = bt[c], bt[c] = v;
      rg[v] = (i + 1 == (int)col.size() ? f : v + 1);
      rw[v] = rows, cl[v] = c;
      ++s[c];
      if (i > 0) lt[v] = v - 1;
    ++rows, lt[f] = sz - 1;
  int h() {
    int ret = 0;
    fill_n(vis, sz, false);
TRAV(x, rg, head) {
      if (vis[x]) continue;
      vis[x] = true, ++ret;
      TRAV(i, dn, x) TRAV(j, rg, i) vis[cl[j]] = true;
    return ret:
  void dfs(int dep) {
    if (dep + (E ? 0 : h()) >= ans) return;
    if (rg[head
        ] == head) return sol = cur, ans = dep, void();
    if (dn[rg[head]] == rg[head]) return;
    int w = rg[head];
    TRAV(x, rg, head) if (s[x] < s[w]) w = x;
    if (E) remove(w);
    TRAV(i, dn, w) {
      if (!E) remove(i);
      TRAV(j, rg, i) remove(E ? cl[j] : j);
      cur.set(rw[i]), dfs(dep + 1), cur.reset(rw[i]);
      TRAV(j, lt, i) restore(E ? cl[j] : j);
      if (!E) restore(i);
```

```
if (E) restore(w);
}
int solve() {
   for (int i = 0; i < columns; ++i)
      dn[bt[i]] = i, up[i] = bt[i];
   ans = 1e9, sol.reset(), dfs(0);
   return ans;
}
};</pre>
```

## 9.9 Matroid Intersection

Start from  $S=\emptyset$ . In each iteration, let  $Y_1=\{x\not\in S\,|\, S\cup\{x\}\in I_1\}$  •  $Y_2=\{x\not\in S\,|\, S\cup\{x\}\in I_2\}$  If there exists  $x\in Y_1\cap Y_2$ , insert x into S. Otherwise for each  $x\in S, y\not\in S$ , create edges •  $x\to y$  if  $S-\{x\}\cup\{y\}\in I_1$ . •  $y\to x$  if  $S-\{x\}\cup\{y\}\in I_2$ .

Find a shortest path (with BFS) starting from a vertex in  $Y_1$  and ending at a vertex in  $Y_2$  which doesn't pass through any other vertices in  $Y_2$ , and alternate the path. The size of S will be incremented by 1 in each iteration. For the weighted case, assign weight w(x) to vertex x if  $x \in S$  and -w(x) if  $x \notin S$ . Find the path with the minimum number of edges among all minimum length paths and alternate it.

# 9.10 AdaptiveSimpson\*

```
template < typename Func, typename d = double >
struct Simpson {
  using pdd = pair<d, d>;
  Func f;
  pdd mix(pdd l, pdd r, optional<d> fm = {}) {
   d h = (r.X - l.X) / 2, v = fm.value_or(f(l.X + h));
   return {v, h / 3 * (l.Y + 4 * v + r.Y)};
  d eval(pdd l, pdd r, d fm, d eps) {
  pdd m((l.X + r.X) / 2, fm);
     d s = mix(l, r, fm).second;
     auto [flm, sl] = mix(l, m);
     auto [fmr, sr] = mix(m, r);
     d delta = sl + sr - s;
     if (abs(delta
          ) <= 15 * eps) return sl + sr + delta / 15;
     return eval(l, m, flm, eps / 2) +
  eval(m, r, fmr, eps / 2);
  d eval(d l, d r, d eps) {
     return eval
          (\{l, f(l)\}, \{r, f(r)\}, f((l + r) / 2), eps);
  d eval2(d l, d r, d eps, int k = 997) {
  d h = (r - l) / k, s = 0;
     for (int i = 0; i < k; ++i, l += h)</pre>
       s += eval(l, l + h, eps / k);
     return s:
  }
template < typename Func >
Simpson<Func> make simpson(Func f) { return {f}; }
```

#### 9.11 Simulated Annealing

```
double factor = 100000;
const int base = 1e9; // remember to run ~ 10 times
for (int it = 1; it <= 1000000; ++it) {
    // ans:
        answer, nw: current value, rnd(): mt19937 rnd()
    if (exp(-(nw - ans
        ) / factor) >= (double)(rnd() % base) / base)
        ans = nw;
    factor *= 0.99995;
}
```

#### 9.12 Tree Hash\*

```
ull seed;
ull shift(ull x) {
    x ^= x << 13;
    x ^= x >> 7;
    x ^= x << 17;
    return x;
}
ull dfs(int u, int f) {
    ull sum = seed;
    for (int i : G[u])
        if (i != f)</pre>
```

```
sum += shift(dfs(i, u));
return sum;
}
```

# 9.13 Binary Search On Fraction

```
struct Q {
  ll p, q;
  Q go(Q b, ll d) { return {p + b.p*d, q + b.q*d}; }
// returns smallest p/q in [lo, hi] such that // pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(ll N) {
  Q lo{0, 1}, hi{1, 0};
  if (pred(lo)) return lo;
  assert(pred(hi));
  bool dir = 1, L = 1, H = 1;
  for (; L || H; dir = !dir) {
    ll len = 0, step = 1;
    for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
       if (Q mid = hi.go(lo, len + step);
           mid.p > N || mid.q > N || dir ^ pred(mid))
         t++:
      else len += step;
    swap(lo, hi = hi.go(lo, len));
    (dir ? L : H) = !!len;
  return dir ? hi : lo;
}
```

# 9.14 Min Plus Convolution\*

```
// a is convex a[i+1]-a[i] <= a[i+2]-a[i+1]
vector <int> min_plus_convolution
        (vector <int> &a, vector <int> &b) {
    int n = SZ(a), m = SZ(b);
    vector <int> c(n + m - 1, INF);
    auto dc = [&](auto Y, int l, int r, int jl, int jr) {
        if (l > r) return;
        int mid = (l + r) / 2, from = -1, &best = c[mid];
        for (int j = jl; j <= jr; ++j)
         if (int i = mid - j; i >= 0 && i < n)
            if (best > a[i] + b[j])
            best = a[i] + b[j], from = j;
        Y(Y, l,
            mid - 1, jl, from), Y(Y, mid + 1, r, from, jr);
    };
    return dc(dc, 0, n - 1 + m - 1, 0, m - 1), c;
}
```

## 9.15 Bitset LCS

```
cin >> n >> m;
for (int i = 1, x; i <= n; ++i)
   cin >> x, p[x].set(i);
for (int i = 1, x; i <= m; i++) {
   cin >> x, (g = f) |= p[x];
   f.shiftLeftByOne(), f.set(0);
   ((f = g - f) ^= g) &= g;
}
cout << f.count() << '\n';</pre>
```

#### 9.16 N Queens Problem

```
void solve
    (vector < int > &ret, int n) { // no sol when n=2,3
    if (n % 6 == 2) {
        for (int i = 2; i <= n; i += 2) ret.pb(i);
        ret.pb(3); ret.pb(1);
        for (int i = 7; i <= n; i += 2) ret.pb(i);
        ret.pb(5);
    } else if (n % 6 == 3) {
        for (int i = 4; i <= n; i += 2) ret.pb(i);
        ret.pb(2);
        for (int i = 5; i <= n; i += 2) ret.pb(i);
        ret.pb(1); ret.pb(3);
    } else {
        for (int i = 2; i <= n; i += 2) ret.pb(i);
        for (int i = 1; i <= n; i += 2) ret.pb(i);
    }
}</pre>
```

# 10 Python

# 10.1 Misc