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```

### Basic

### 1.1 vimrc

```
"This file should be placed at ~/.vimrc"
se nu ai hls et ru ic is sc cul
se re=1 ts=4 sts=4 sw=4 ls=2 mouse=a
syntax on
hi cursorline cterm=none ctermbg=89
set ba=dark
inoremap {<CR> {<CR>}<Esc>ko<tab>
```

### 1.2 readchar

```
inline char readchar() {
  static const size_t bufsize = 65536;
  static char buf[bufsize];
  static char *p = buf, *end = buf;
  if (p == end) end = buf +
       fread unlocked(buf, 1, bufsize, stdin), p = buf;
  return *p++:
}
```

### 1.3 Black Magic

```
#include <ext/pb_ds/priority_queue.hpp>
#include <ext/pb_ds/assoc_container.hpp> // rb_tree
#include <ext/rope> // rope
using namespace __gnu_pbds;
using namespace __gnu_cxx; // rope
typedef __gnu_pbds::priority_queue<int> heap;
int main() {
  heap h1, h2; // max heap
   h1.push(1), h1.push(3), h2.push(2), h2.push(4);
   h1.join(h2); // h1 = {1, 2, 3, 4}, h2 = {};
   tree<ll, null_type, less<ll>, rb_tree_tag
           tree_order_statistics_node_update> st;
   tree<ll, ll, less<ll>, rb_tree_tag
         , tree_order_statistics_node_update > mp;
   for (int x : {0, 2, 3, 4}) st.insert(x);
cout << *st.find_by_order</pre>
   (2) << st.order_of_key(1) << endl; //31
rope<char> *root[10]; // nsqrt(n)
root[0] = new rope<char>();
   root[1] = new rope < char > (*root[0]);
   // root[1]->insert(pos,
   // root[1]->at(pos); 0-base
   // root[1]->erase(pos, size);
// __int128_t,__float128_t
// for (int i = bs._Find_first
      (); i < bs.size(); i = bs._Find_next(i));
```

# 2 Graph

### 2.1 BCC Vertex\*

```
vector < int > G[N]; // 1-base
vector < int > nG[N * 2], bcc[N];
int low[N], dfn[N], Time;
int bcc_id[N], bcc_cnt; // 1-base
bool is_cut[N]; // whether is av
bool cir[N * 2];
int st[N], top;
void dfs(int u, int pa = -1) {
  int child = 0;
  low[u] = dfn[u] = ++Time;
  st[top++] = u;
for (int v : G[u])
     if (!dfn[v]) {
       dfs(v, u), ++child;
low[u] = min(low[u], low[v]);
       if (dfn[u] <= low[v]) {
  is_cut[u] = 1;</pre>
          bcc[++bcc_cnt].clear();
          int t:
          do {
            bcc_id[t = st[--top]] = bcc_cnt;
            bcc[bcc_cnt].push_back(t);
          } while (t != v);
          bcc_id[u] = bcc_cnt;
          bcc[bcc_cnt].pb(u);
     } else if (dfn[v] < dfn[u] && v != pa)</pre>
       low[u] = min(low[u], dfn[v]);
  if (pa == -1 && child < 2) is_cut[u] = 0;</pre>
}
void bcc_init(int n) { // TODO: init {nG, cir}[1..2n]
  Time = bcc_cnt = top = 0;
  for (int i = 1; i <= n; ++i)
  G[i].clear(), dfn[i] = bcc_id[i] = is_cut[i] = 0;</pre>
}
void bcc_solve(int n) {
  for (int i = 1; i <= n; ++i)</pre>
     if (!dfn[i]) dfs(i);
   // block-cut tree
  for (int i = 1; i <= n; ++i)</pre>
     if (is_cut[i])
       bcc_id[i] = ++bcc_cnt, cir[bcc_cnt] = 1;
  for (int i = 1; i <= bcc_cnt && !cir[i]; ++i)
for (int j : bcc[i])</pre>
       if (is_cut[j])
          nG[i].pb(bcc_id[j]), nG[bcc_id[j]].pb(i);
}
```

### 2.2 Bridge\*

```
int low[N], dfn[N], Time; // 1-base
vector<pii> G[N], edge;
vector<bool> is_bridge;
void init(int n) {
 Time = 0;
  for (int i = 1; i <= n; ++i)</pre>
    G[i].clear(), low[i] = dfn[i] = 0;
void add_edge(int a, int b) {
 G[a].pb(pii(b, SZ(edge))), G[b].pb(pii(a, SZ(edge)));
  edge.pb(pii(a, b));
void dfs(int u, int f) {
 dfn[u] = low[u] = ++Time;
  for (auto i : G[u])
    if (!dfn[i.X])
      dfs(i.X, i.Y), low[u] = min(low[u], low[i.X]);
    else if (i.Y != f) low[u] = min(low[u], dfn[i.X]);
  if (low[u] == dfn[u] && f != -1) is_bridge[f] = 1;
void solve(int n) {
 is_bridge.resize(SZ(edge));
  for (int i = 1; i <= n; ++i)</pre>
    if (!dfn[i]) dfs(i, -1);
```

### 2.3 2SAT (SCC)\*

```
struct SAT { // 0-base
  int low[N], dfn[N], bln[N], n, Time, nScc;
  bool instack[N], istrue[N];
  stack<int> st:
  vector<int> G[N], SCC[N];
  void init(int _n) {
    n = _n; // assert(n * 2 <= N);</pre>
    for (int i = 0; i < n + n; ++i) G[i].clear();</pre>
  void add_edge(int a, int b) { G[a].pb(b); }
  int rv(int a) {
    if (a >= n) return a - n;
    return a + n;
  void add_clause(int a, int b) {
    add_edge(rv(a), b), add_edge(rv(b), a);
  void dfs(int u) {
   dfn[u] = low[u] = ++Time;
     instack[u] = 1, st.push(u);
     for (int i : G[u])
       if (!dfn[i])
         dfs(i), low[u] = min(low[i], low[u]);
       else if (instack[i] && dfn[i] < dfn[u])</pre>
         low[u] = min(low[u], dfn[i]);
     if (low[u] == dfn[u]) {
       int tmp;
      do {
         tmp = st.top(), st.pop();
         instack[tmp] = 0, bln[tmp] = nScc;
       } while (tmp != u);
       ++nScc:
    }
  bool solve() {
    Time = nScc = 0;
    for (int i = 0; i < n + n; ++i)</pre>
    SCC[i].clear(), low[i] = dfn[i] = bln[i] = 0;
for (int i = 0; i < n + n; ++i)
      if (!dfn[i]) dfs(i);
     for (int i = 0; i < n + n; ++i) SCC[bln[i]].pb(i);</pre>
     for (int i = 0; i < n; ++i) {</pre>
      if (bln[i] == bln[i + n]) return false;
       istrue[i] = bln[i] < bln[i + n];</pre>
       istrue[i + n] = !istrue[i];
     return true;
  }
};
```

# 2.4 MinimumMeanCycle\*

```
ll road[N][N]; // input here
struct MinimumMeanCycle {
  ll dp[N + 5][N], n;
  pll solve() {
     ll a = -1, b = -1, L = n + 1;
     for (int i = 2; i <= L; ++i)</pre>
       for (int k = 0; k < n; ++k)
          for (int j = 0; j < n; ++j)</pre>
            dp[i][j] =
               min(dp[i - 1][k] + road[k][j], dp[i][j]);
     for (int i = 0; i < n; ++i) {
  if (dp[L][i] >= INF) continue;
       ll ta = 0, tb = 1;
for (int j = 1; j < n; ++j)
          if (dp[j][i] < INF &&</pre>
            ta * (L - j) < (dp[L][i] - dp[j][i]) * tb)
ta = dp[L][i] - dp[j][i], tb = L - j;</pre>
       if (ta == 0) continue;
       if (a == -1 || a * tb > ta * b) a = ta, b = tb;
     if (a != -1) {
                  _gcd(a, b);
       ll q =
       return pll(a / g, b / g);
     return pll(-1LL, -1LL);
  void init(int _n) {
     n = _n;
for (int i = 0; i < n; ++i)</pre>
       for (int j = 0; j < n; ++j) dp[i + 2][j] = INF;</pre>
};
```

### 2.5 Virtual Tree\*

```
vector<int> vG[N];
int top, st[N];
void insert(int u) {
  if (top == -1) return st[++top] = u, void();
  int p = LCA(st[top], u);
  if (p == st[top]) return st[++top] = u, void();
  while (top >= 1 && dep[st[top - 1]] >= dep[p])
    vG[st[top - 1]].pb(st[top]), --top;
  if (st[top] != p)
    vG[p].pb(st[top]), --top, st[++top] = p;
  st[++top] = u;
}
void reset(int u) {
  for (int i : vG[u]) reset(i);
  vG[u].clear();
void solve(vector<int> &v) {
  top = -1;
  sort(ALL(v),
   [&](int a, int b) { return dfn[a] < dfn[b]; });</pre>
  for (int i : v) insert(i);
  while (top > 0) vG[st[top - 1]].pb(st[top]), --top;
  // do something
  reset(v[0]);
```

### 2.6 Maximum Clique Dyn\*

```
struct MaxClique { // fast when N <= 100</pre>
  bitset<N> G[N], cs[N];
  int ans, sol[N], q, cur[N], d[N], n;
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i) G[i].reset();</pre>
  void add_edge(int u, int v) {
    G[u][v] = G[v][u] = 1;
  void pre_dfs(vector<int> &r, int l, bitset<N> mask) {
    if (1 < 4) {
      for (int i : r) d[i] = (G[i] & mask).count();
      sort(ALL(r)
          , [&](int x, int y) { return d[x] > d[y]; });
    vector<int> c(SZ(r));
    int lft = max(ans - q + 1, 1), rgt = 1, tp = 0;
    cs[1].reset(), cs[2].reset();
    for (int p : r) {
```

```
int k = 1:
      while ((cs[k] & G[p]).any()) ++k;
      if (k > rgt) cs[++rgt + 1].reset();
      cs[k][p] = 1;
      if (k < lft) r[tp++] = p;</pre>
    for (int k = lft; k <= rgt; ++k)</pre>
      for (int p = cs[k]._Find_first
           (); p < N; p = cs[k]._Find_next(p))
        r[tp] = p, c[tp] = k, ++tp;
    dfs(r, c, l + 1, mask);
  void dfs(vector<</pre>
       int> &r, vector<int> &c, int l, bitset<N> mask) {
    while (!r.empty()) {
      int p = r.back();
      r.pop_back(), mask[p] = 0;
      if (q + c.back() <= ans) return;</pre>
      cur[q++] = p;
      vector<int> nr;
      for (int i : r) if (G[p][i]) nr.pb(i);
      if (!nr.empty()) pre_dfs(nr, l, mask & G[p]);
      else if (q > ans) ans = q, copy_n(cur, q, sol);
      c.pop_back(), --q;
    }
  int solve() {
    vector<int> r(n);
    ans = q = 0, iota(ALL(r), \theta);
    pre_dfs(r, 0, bitset<N>(string(n, '1')));
    return ans:
};
```

### 2.7 Minimum Steiner Tree\*

struct SteinerTree { // 0-base

```
int n, dst[N][N], dp[1 << T][N], tdst[N];</pre>
   int vcst[N]; // the cost of vertexs
   void init(int _n) {
     n = _n;
for (int i = 0; i < n; ++i) {</pre>
       fill_n(dst[i], n, INF);
       dst[i][i] = vcst[i] = 0;
     }
   void chmin(int &x, int val) {
    x = min(x, val);
   void add_edge(int ui, int vi, int wi) {
     chmin(dst[ui][vi], wi);
   void shortest_path() {
     for (int k = 0; k < n; ++k)
  for (int i = 0; i < n; ++i)</pre>
          for (int j = 0; j < n; ++j)</pre>
            chmin(dst[i][j], dst[i][k] + dst[k][j]);
   int solve(const vector<int>& ter) {
     shortest_path();
     int t = SZ(ter), full = (1 << t) - 1;</pre>
     for (int i = 0; i <= full; ++i)
  fill_n(dp[i], n, INF);</pre>
     copy_n(vcst, n, dp[0]);
     for (int msk = 1; msk <= full; ++msk) {</pre>
        if (!(msk & (msk - 1))) {
          int who = __lg(msk);
for (int i = 0; i < n; ++i)</pre>
            dp[msk
                 ][i] = vcst[ter[who]] + dst[ter[who]][i];
        for (int i = 0; i < n; ++i)</pre>
          for (int sub = (
              msk - 1) & msk; sub; sub = (sub - 1) & msk)
            chmin(dp[msk][i],
                dp[sub][i] + dp[msk ^ sub][i] - vcst[i]);
        for (int i = 0; i < n; ++i) {</pre>
          tdst[i] = INF;
          for (int j = 0; j < n; ++j)</pre>
            chmin(tdst[i], dp[msk][j] + dst[j][i]);
       copy_n(tdst, n, dp[msk]);
     return *min_element(dp[full], dp[full] + n);
}; // O(V 3^T + V^2 2^T)
```

### 2.8 Dominator Tree\*

```
struct dominator_tree { // 1-base
   vector<int> G[N], rG[N];
   int n, pa[N], dfn[N], id[N], Time;
   int semi[N], idom[N], best[N];
vector < int > tree[N]; // dominator_tree
   void init(int _n) {
     n = _n;
for (int i = 1; i <= n; ++i)
       G[i].clear(), rG[i].clear();
   void add_edge(int u, int v) {
     G[u].pb(v), rG[v].pb(u);
   void dfs(int u) {
     id[dfn[u] = ++Time] = u;
     for (auto v : G[u])
       if (!dfn[v]) dfs(v), pa[dfn[v]] = dfn[u];
   int find(int y, int x) {
     if (y <= x) return y;</pre>
     int tmp = find(pa[y], x);
     if (semi[best[y]] > semi[best[pa[y]]])
       best[y] = best[pa[y]];
     return pa[y] = tmp;
   void tarjan(int root) {
     Time = 0;
     for (int i = 1; i <= n; ++i) {</pre>
       dfn[i] = idom[i] = 0;
       tree[i].clear();
       best[i] = semi[i] = i;
     dfs(root);
     for (int i = Time; i > 1; --i) {
       int u = id[i];
       for (auto v : rG[u])
         if (v = dfn[v]) {
           find(v, i);
            semi[i] = min(semi[i], semi[best[v]]);
       tree[semi[i]].pb(i);
       for (auto v : tree[pa[i]]) {
         find(v, pa[i]);
         idom[v] =
            semi[best[v]] == pa[i] ? pa[i] : best[v];
       tree[pa[i]].clear();
     for (int i = 2; i <= Time; ++i) {
   if (idom[i] != semi[i]) idom[i] = idom[idom[i]];</pre>
       tree[id[idom[i]]].pb(id[i]);
  }
};
```

### 2.9 Minimum Arborescence\*

```
/* TODO
DSU: disjoint set
- DSU(n), .boss(x), .Union(x, y)
min_heap <
    T, Info>: min heap for type {T, Info} with lazy tag
 .push({w, i}),
    .top(), .join(heap), .pop(), .empty(), .add_lazy(v)
struct E { int s, t; ll w; }; // 0-base
vector<int> dmst(const vector<E> &e, int n, int root) {
  vector<min_heap<ll, int>> h(n * 2);
for (int i = 0; i < SZ(e); ++i)</pre>
    h[e[i].t].push({e[i].w, i});
  DSU dsu(n * 2);
  vector<int> v(n * 2, -1), pa(n * 2, -1), r(n * 2);
  v[root] = n + 1;
  int pc = n;
  for (int i = 0; i < n; ++i) if (v[i] == -1) {</pre>
    for (int p = i; v[p]
         == -1 || v[p] == i; p = dsu.boss(e[r[p]].s)) {
      if (v[p] == i) {
         int q = p; p = pc++;
         do {
          h[q].add_lazy(-h[q].top().X);
           pa[q] = p, dsu.Union(p, q), h[p].join(h[q]);
        } while ((q = dsu.boss(e[r[q]].s)) != p);
```

```
v[p] = i:
      while (!h[p].
          empty() \&\& dsu.boss(e[h[p].top().Y].s) == p)
        h[p].pop();
      if (h[p].empty()) return {}; // no solution
      r[p] = h[p].top().Y;
   }
  }
  vector<int> ans;
  for (int i = pc
       - 1; i >= 0; i--) if (i != root && v[i] != n) {
    for (int f = e[r[i]].t; \sim f \&\& v[f] != n; f = pa[f])
      v[f] = n;
    ans.pb(r[i]);
  return ans; // default minimize, returns edgeid array
\} // O(Ef(E)), f(E) from min_heap
```

### 2.10 Vizing's theorem\*

```
namespace vizing { // returns
  edge coloring in adjacent matrix G. 1 - based
const int N = 105;
int C[N][N], G[N][N], X[N], vst[N], n;
void init(int _n) { n = _n;
  for (int i = 0; i <= n; ++i)</pre>
    for (int j = 0; j <= n; ++j)</pre>
      C[i][j] = G[i][j] = 0;
void solve(vector<pii> &E) {
  auto update = [&](int u)
  { for (X[u] = 1; C[u][X[u]]; ++X[u]); };
  auto color = [&](int u, int v, int c) {
    int p = G[u][v];
    G[u][v] = G[v][u] = c;
    C[u][c] = v, C[v][c] = u;
    C[u][p] = C[v][p] = 0;
    if (p) X[u] = X[v] = p;
    else update(u), update(v);
    return p;
  auto flip = [&](int u, int c1, int c2) {
    int p = C[u][c1];
    swap(C[u][c1], C[u][c2]);
    if (p) G[u][p] = G[p][u] = c2;
    if (!C[u][c1]) X[u] = c1;
    if (!C[u][c2]) X[u] = c2;
    return p;
  fill_n(X + 1, n, 1);
  for (int t = 0; t < SZ(E); ++t) {</pre>
    int u = E[t
         ].X, v0 = E[t].Y, v = v0, c0 = X[u], c = c0, d;
    vector<pii> L;
    fill_n(vst + 1, n,
    while (!G[u][v0]) {
      L.emplace_back(v, d = X[v]);
if (!C[v][c]) for (int a = SZ(
           L) - 1; a >= 0; --a) c = color(u, L[a].X, c);
       else if (!C[u][d]) for (int a = SZ(L
           ) - 1; a >= 0; --a) color(u, L[a].X, L[a].Y);
      else if (vst[d]) break;
       else vst[d] = 1, v = C[u][d];
    if (!G[u][v0]) {
      for (; v; v = flip(v, c, d), swap(c, d));
if (int a; C[u][c0]) {
         for (
             a = SZ(L) - 2; a >= 0 && L[a].Y != c; --a);
         for (; a >= 0; --a) color(u, L[a].X, L[a].Y);
      else --t;
    }
  }
} // namespace vizing
```

### 2.11 Minimum Clique Cover\*

```
struct Clique_Cover { // 0-base, 0(n2^n)
  int co[1 << N], n, E[N];
  int dp[1 << N];
  void init(int _n) {
    n = _n, fill_n(dp, 1 << n, 0);
    fill_n(E, n, 0), fill_n(co, 1 << n, 0);
}</pre>
```

```
void add_edge(int u, int v) {
     E[u] |= 1 << v, E[v] |= 1 << u;
   int solve() {
     for (int i = 0; i < n; ++i)</pre>
       co[1 << i] = E[i] | (1 << i);
     co[0] = (1 << n) - 1;

dp[0] = (n & 1) * 2 - 1;
     for (int i = 1; i < (1 << n); ++i) {</pre>
       int t = i & -i;
       dp[i] = -dp[i ^ t];
       co[i] = co[i ^ t] & co[t];
     for (int i = 0; i < (1 << n); ++i)</pre>
       co[i] = (co[i] & i) == i;
     fwt(co, 1 << n, 1);
for (int ans = 1; ans < n; ++ans) {</pre>
       int sum = 0; // probabilistic
       for (int i = 0; i < (1 << n); ++i)</pre>
         sum += (dp[i] *= co[i]);
       if (sum) return ans;
     return n;
};
```

### 2.12 NumberofMaximalClique\*

```
struct BronKerbosch { // 1-base
   int n, a[N], g[N][N];
int S, all[N][N], some[N][N], none[N][N];
   void init(int _n) {
     n = _n;
for (int i = 1; i <= n; ++i)</pre>
       for (int j = 1; j <= n; ++j) g[i][j] = 0;</pre>
   void add_edge(int u, int v) {
     g[u][v] = g[v][u] = 1;
   void dfs(int d, int an, int sn, int nn) {
  if (S > 1000) return; // pruning
     if (sn == 0 && nn == 0) ++S;
     int u = some[d][0];
     for (int i = 0; i < sn; ++i) {</pre>
       int v = some[d][i];
       if (g[u][v]) continue;
       int tsn = 0, tnn = 0;
       copy_n(all[d], an, all[d + 1]);
all[d + 1][an] = v;
       for (int j = 0; j < sn; ++j)
          if (g[v][some[d][j]])
            some[d + 1][tsn++] = some[d][j];
        for (int j = 0; j < nn; ++j)</pre>
          if (g[v][none[d][j]])
            none[d + 1][tnn++] = none[d][j];
       dfs(d + 1, an + 1, tsn, tnn);
        some[d][i] = 0, none[d][nn++] = v;
   int solve() {
     iota(some[0], some[0] + n, 1);
     S = 0, dfs(0, 0, n, 0);
     return S;
  }
};
```

### 3 Data Structure

### 3.1 Discrete Trick

```
vector < int > val;
// build
sort(ALL
     (val)), val.resize(unique(ALL(val)) - val.begin());
// index of x
upper_bound(ALL(val), x) - val.begin();
// max idx <= x
upper_bound(ALL(val), x) - val.begin();
// max idx < x
lower_bound(ALL(val), x) - val.begin();</pre>
```

### 3.2 Interval Container\*

```
/* Add and
    remove intervals from a set of disjoint intervals.
```

```
* Will merae the added interval with
      any overlapping intervals in the set when adding.
 * Intervals are [inclusive, exclusive). */
set<pii>::
    iterator addInterval(set<pii>& is, int L, int R) {
  if (L == R) return is.end();
  auto it = is.lower_bound({L, R}), before = it;
  while (it != is.end() && it->X <= R) {</pre>
   R = max(R, it->Y);
    before = it = is.erase(it);
  if (it != is.begin() && (--it)->Y >= L) {
   L = min(L, it->X);
R = max(R, it->Y);
    is.erase(it);
 }
  return is.insert(before, pii(L, R));
void removeInterval(set<pii>& is, int L, int R) {
  if (L == R) return;
  auto it = addInterval(is, L, R);
  auto r2 = it->Y;
  if (it->X == L) is.erase(it);
  else (int&)it->Y = L;
  if (R != r2) is.emplace(R, r2);
```

### 3.3 Leftist Tree

```
struct node {
  ll v, data, sz, sum;
node *l, *r;
  node(ll k)
    : v(0), data(k), sz(1), l(0), r(0), sum(k) {}
ll sz(node *p) { return p ? p->sz : 0; }
ll V(node *p) { return p ? p->v : -1; }
ll sum(node *p) { return p ? p->sum : 0; }
node *merge(node *a, node *b) {
  if (!a || !b) return a ? a : b;
  if (a->data < b->data) swap(a, b);
  a - r = merge(a - r, b);
  if (V(a->r) > V(a->l)) swap(a->r, a->l);
a->v = V(a->r) + 1, a->sz = sz(a->l) + sz(a->r) + 1;
  a->sum = sum(a->l) + sum(a->r) + a->data;
  return a;
void pop(node *&o) {
  node *tmp = o;
  o = merge(o->l, o->r);
  delete tmp;
```

### 3.4 Heavy light Decomposition

```
struct Heavy_light_Decomposition { // 1-base
  int n, ulink[N], deep[N], mxson[N], w[N], pa[N];
  int t, pl[N], data[N], dt[N], bln[N], edge[N], et;
  vector<pii> G[N];
  void init(int _n) {
  n = _n, t = 0, et = 1;
  for (int i = 1; i <= n; ++i)</pre>
       G[i].clear(), mxson[i] = 0;
  void add_edge(int a, int b, int w) {
    G[a].pb(pii(b, et));
    G[b].pb(pii(a, et));
    edge[et++] = w;
  void dfs(int u, int f, int d) {
  w[u] = 1, pa[u] = f, deep[u] = d++;
     for (auto &i : G[u])
       if (i.X != f) {
         dfs(i.X, u, d), w[u] += w[i.X];
if (w[mxson[u]] < w[i.X]) mxson[u] = i.X;</pre>
       } else bln[i.Y] = u, dt[u] = edge[i.Y];
  void cut(int u, int link) {
  data[pl[u] = t++] = dt[u], ulink[u] = link;
    if (!mxson[u]) return;
     cut(mxson[u], link);
     for (auto i : G[u])
       if (i.X != pa[u] && i.X != mxson[u])
         cut(i.X, i.X);
  void build() { dfs(1, 1, 1), cut(1, 1), /*build*/; }
```

```
int query(int a, int b) {
   int ta = ulink[a], tb = ulink[b], re = 0;
   while (ta != tb)
      if (deep[ta] < deep[tb])
        /*query*/, tb = ulink[b = pa[tb]];
      else /*query*/, ta = ulink[a = pa[ta]];
   if (a == b) return re;
   if (pl[a] > pl[b]) swap(a, b);
      /*query*/
   return re;
}
```

### 3.5 Centroid Decomposition\*

```
struct Cent_Dec { // 1-base
  vector<pll> G[N];
  pll info[N]; // store info. of itself
  pll upinfo[N]; // store info. of climbing up
int n, pa[N], layer[N], sz[N], done[N];
  ll dis[__lg(N) + 1][N];
  void init(int _n) {
    n = _n, layer[0] = -1;
    fill_n(pa + 1, n, 0), fill_n(done + 1, n, 0);
    for (int i = 1; i <= n; ++i) G[i].clear();</pre>
  void add_edge(int a, int b, int w) {
    G[a].pb(pll(b, w)), G[b].pb(pll(a, w));
  void get_cent(
    int u, int f, int &mx, int &c, int num) {
    int mxsz = 0;
    sz[u] = 1;
     for (pll e : G[u])
       if (!done[e.X] && e.X != f) {
         get_cent(e.X, u, mx, c, num);
         sz[u] += sz[e.X], mxsz = max(mxsz, sz[e.X]);
    if (mx > max(mxsz, num - sz[u]))
      mx = max(mxsz, num - sz[u]), c = u;
  void dfs(int u, int f, ll d, int org) {
     // if required, add self info or climbing info
     dis[layer[org]][u] = d;
     for (pll e : G[u])
      if (!done[e.X] && e.X != f)
         dfs(e.X, u, d + e.Y, org);
  int cut(int u, int f, int num) {
    int mx = 1e9, c = 0, lc;
    get_cent(u, f, mx, c, num);
    done[c] = 1, pa[c] = f, layer[c] = layer[f] + 1;
for (pll e : G[c])
       if (!done[e.X]) {
         if (sz[e.X] > sz[c])
         lc = cut(e.X, c, num - sz[c]);
else lc = cut(e.X, c, sz[e.X]);
         upinfo[lc] = pll(), dfs(e.X, c, e.Y, c);
    return done[c] = 0, c;
  void build() { cut(1, 0, n); }
  void modify(int u) {
    for (int a = u, ly = layer[a]; a;
    a = pa[a], --ly) {
       info[a].X += dis[ly][u], ++info[a].Y;
       if (pa[a])
         upinfo[a].X += dis[ly - 1][u], ++upinfo[a].Y;
  ll query(int u) {
    ll rt = 0;
    for (int a = u, ly = layer[a]; a;
          a = pa[a], --ly) {
       rt += info[a].X + info[a].Y * dis[ly][u];
       if (pa[a])
         rt -=
           upinfo[a].X + upinfo[a].Y * dis[ly - 1][u];
     return rt;
};
```

### 3.6 Link cut tree\*

```
struct Splay { // xor-sum
```

```
static Splay nil;
  Splay *ch[2], *f;
  int val, sum, rev,
                     size;
  Splay(int _val = 0)
    : val(\_val), sum(\_val), rev(0), size(1) {
    f = ch[0] = ch[1] = &nil;
  bool isr() {
    return f->ch[0] != this && f->ch[1] != this;
  int dir() { return f->ch[0] == this ? 0 : 1; }
  void setCh(Splay *c, int d) {
    ch[d] = c;
    if (c != &nil) c->f = this;
    pull();
  void give_tag(int r)
  { if (r) swap(ch[0], ch[1]), rev ^= 1; }
  void push() {
    if (ch[0] != &nil) ch[0]->give_tag(rev);
    if (ch[1] != &nil) ch[1]->give_tag(rev);
    rev = 0;
  void pull() {
   // take care of the nil!
    size = ch[0]->size + ch[1]->size + 1;
    sum = ch[0] -> sum ^ ch[1] -> sum ^ val;
    if (ch[0] != &nil) ch[0]->f = this;
    if (ch[1] != &nil) ch[1]->f = this;
} Splay::nil;
Splay *nil = &Splay::nil;
void rotate(Splay *x) {
  Splay *p = x - > f;
  int d = x->dir();
 if (!p->isr()) p->f->setCh(x, p->dir());
 else x->f = p->f;
 p->setCh(x->ch[!d], d);
 x->setCh(p, !d);
 p->pull(), x->pull();
void splay(Splay *x) {
 vector < Splay *> splayVec;
  for (Splay *q = x;; q = q->f) {
    splayVec.pb(q);
    if (q->isr()) break;
  reverse(ALL(splayVec));
  for (auto it : splayVec) it->push();
 while (!x->isr()) {
    if (x->f->isr()) rotate(x);
    else if (x->dir() == x->f->dir())
    rotate(x->f), rotate(x);
else rotate(x), rotate(x);
 }
Splay *access(Splay *x) {
  Splay *q = nil;
  for (; x != nil; x = x->f)
   splay(x), x - setCh(q, 1), q = x;
  return q;
void root_path(Splay *x) { access(x), splay(x); }
void chroot(Splay *x) {
 root_path(x), x->rev ^= 1:
 x->push(), x->pull();
void split(Splay *x, Splay *y) {
 chroot(x), root_path(y);
void link(Splay *x, Splay *y) {
 root_path(x), chroot(y);
 x->setCh(y, 1);
void cut(Splay *x, Splay *y) {
  split(x, y);
  if (y->size != 5) return;
 y->push();
 y - ch[0] = y - ch[0] - f = nil;
Splay *get_root(Splay *x) {
 for (root_path(x); x->ch[0] != nil; x = x->ch[0])
   x->push();
  splay(x);
 return x:
```

```
bool conn(Splay *x, Splay *y) {
   return get_root(x) == get_root(y);
}
Splay *lca(Splay *x, Splay *y) {
   access(x), root_path(y);
   if (y->f == nil) return y;
   return y->f;
}
void change(Splay *x, int val) {
   splay(x), x->val = val, x->pull();
}
int query(Splay *x, Splay *y) {
   split(x, y);
   return y->sum;
}
```

### 3.7 KDTree

```
namespace kdt {
int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn],
  yl[maxn], yr[maxn];
point p[maxn];
int build(int l, int r, int dep = 0) {
  if (l == r) return -1;
  function < bool(const point &, const point &) > f =
    [dep](const point &a, const point &b) {
      if (dep & 1) return a.x < b.x;</pre>
      else return a.y < b.y;</pre>
    };
  int m = (l + r) >> 1;
nth_element(p + l, p + m, p + r, f);
  xl[m] = xr[m] = p[m].x;
  yl[m] = yr[m] = p[m].y;
  lc[m] = build(l, m, dep + 1);
  if (~lc[m]) {
    xl[m] = min(xl[m], xl[lc[m]]);
    xr[m] = max(xr[m], xr[lc[m]]);
    yl[m] = min(yl[m], yl[lc[m]]);
    yr[m] = max(yr[m], yr[lc[m]]);
  rc[m] = build(m + 1, r, dep + 1);
  if (~rc[m]) {
    xl[m] = min(xl[m], xl[rc[m]]);
    xr[m] = max(xr[m], xr[rc[m]]);
    yl[m] = min(yl[m], yl[rc[m]]);
    yr[m] = max(yr[m], yr[rc[m]]);
  return m:
bool bound(const point &q, int o, long long d) {
  double ds = sqrt(d + 1.0);
  if (q.x < xl[o] - ds || q.x > xr[o] + ds ||
    q.y < yl[o] - ds || q.y > yr[o] + ds)
    return false;
  return true;
long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 1ll * (a.x - b.x) +
    (a.y - b.y) * 1ll * (a.y - b.y);
void dfs(
  const point &q, long long &d, int o, int dep = \theta) {
  if (!bound(q, o, d)) return;
  long long cd = dist(p[o], q);
  if (cd != 0) d = min(d, cd);
  if ((dep & 1) && q.x < p[o].x ||</pre>
    !(dep & 1) && q.y < p[o].y) {
    if (~lc[o]) dfs(q, d, lc[o], dep + 1);
    if (~rc[o]) dfs(q, d, rc[o], dep + 1);
  } else {
    if (~rc[o]) dfs(q, d, rc[o], dep + 1);
    if (~lc[o]) dfs(q, d, lc[o], dep + 1);
  }
void init(const vector<point> &v) {
  for (int i = 0; i < v.size(); ++i) p[i] = v[i];</pre>
  root = build(0, v.size());
long long nearest(const point &q) {
  long long res = 1e18;
  dfs(q, res, root);
  return res;
} // namespace kdt
```

# 4 Flow/Matching

### 4.1 Kuhn Munkres\*

```
struct KM { // 0-base
  ll w[N][N], hl[N], hr[N], slk[N];
int fl[N], fr[N], pre[N], qu[N], ql, qr, n;
  bool vl[N], vr[N];
  void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i)
  fill_n(w[i], n, -INF);</pre>
  void add_edge(int a, int b, ll wei) {
    w[a][b] = wei;
  bool Check(int x) {
    if (vl[x] = 1, \sim fl[x])
       return vr[qu[qr++] = fl[x]] = 1;
     while (~x) swap(x, fr[fl[x] = pre[x]]);
     return 0:
  void bfs(int s) {
     fill_n(slk
          , n, INF), fill_n(vl, n, \theta), fill_n(vr, n, \theta);
     ql = qr = 0, qu[qr++] = s, vr[s] = 1;
     for (ll d;;) {
       while (ql < qr)</pre>
         for (int x = 0, y = qu[ql++]; x < n; ++x)
  if (!vl[x] && slk</pre>
                 [x] \rightarrow = (d = hl[x] + hr[y] - w[x][y])) {
              if (pre[x] = y, d) slk[x] = d;
              else if (!Check(x)) return;
         }
       d = INF;
       for (int x = 0; x < n; ++x)</pre>
         if (!vl[x] && d > slk[x]) d = slk[x];
       for (int x = 0; x < n; ++x) {</pre>
         if (vl[x]) hl[x] += d;
         else slk[x] -= d;
         if (vr[x]) hr[x] -= d;
       for (int x = 0; x < n; ++x)
         if (!vl[x] && !slk[x] && !Check(x)) return;
    }
  ll solve() {
    fill_n(fl
          , n, -1), fill_n(fr, n, -1), fill_n(hr, n, 0);
     for (int i = 0; i < n; ++i)</pre>
     hl[i] = *max_element(w[i], w[i] + n);
for (int i = 0; i < n; ++i) bfs(i);
     ll res = 0;
     for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
     return res:
  }
};
```

### 4.2 MincostMaxflow\*

```
struct MinCostMaxFlow { // 0-base
  struct Edge {
    ll from, to, cap, flow, cost, rev;
  } *past[N];
  vector < Edge > G[N];
  int inq[N], n, s, t;
  ll dis[N], up[N], pot[N];
  bool BellmanFord() {
    fill_n(dis, n, INF), fill_n(inq, n, 0);
    queue<int> q;
    auto relax = [&](int u, ll d, ll cap, Edge *e) {
      if (cap > 0 && dis[u] > d) {
        dis[u] = d, up[u] = cap, past[u] = e;
if (!inq[u]) inq[u] = 1, q.push(u);
      }
    }:
    relax(s, 0, INF, 0);
    while (!q.empty()) {
      int u = q.front();
      q.pop(), inq[u] = 0;
      for (auto &e : G[u]) {
        ll d2 = dis[u] + e.cost + pot[u] - pot[e.to];
             (e.to, d2, min(up[u], e.cap - e.flow), &e);
      }
    }
```

```
return dis[t] != INF;
   void solve(int
     , int _t, ll &flow, ll &cost, bool neg = true) {
s = _s, t = _t, flow = 0, cost = 0;
     if (neg) BellmanFord(), copy_n(dis, n, pot);
     for (; BellmanFord(); copy_n(dis, n, pot)) {
       for (int
           i = 0; i < n; ++i) dis[i] += pot[i] - pot[s];
       flow += up[t], cost += up[t] * dis[t];
       for (int i = t; past[i]; i = past[i]->from) {
         auto &e = *past[i];
         e.flow += up[t], G[e.to][e.rev].flow -= up[t];
    }
  }
   void init(int _n) {
     n = _n, fill_n(pot, n, 0);
     for (int i = 0; i < n; ++i) G[i].clear();</pre>
   void add_edge(ll a, ll b, ll cap, ll cost) {
     G[a].pb(Edge{a, b, cap, 0, cost, SZ(G[b])});
     G[b].pb(Edge{b, a, 0, 0, -cost, SZ(G[a]) - 1});
};
```

# 4.3 Maximum Simple Graph Matching\*

```
struct GenMatch { // 1-base
  int V, pr[N];
  bool el[N][N], inq[N], inp[N], inb[N];
  int st, ed, nb, bk[N], djs[N], ans;
  void init(int _V) {
    for (int i = 0; i <= V; ++i) {
  for (int j = 0; j <= V; ++j) el[i][j] = 0;</pre>
      pr[i] = bk[i] = djs[i] = 0;
      inq[i] = inp[i] = inb[i] = 0;
  }
  void add_edge(int u, int v) {
    el[u][v] = el[v][u] = 1;
  int lca(int u, int v) {
    fill_n(inp, V + 1, 0);
    while (1)
      if (u = djs[u], inp[u] = true, u == st) break;
      else u = bk[pr[u]];
    while (1)
      if (v = djs[v], inp[v]) return v;
      else v = bk[pr[v]];
    return v;
  void upd(int u) {
    for (int v; djs[u] != nb;) {
     v = pr[u], inb[djs[u]] = inb[djs[v]] = true;
      u = bk[v];
      if (djs[u] != nb) bk[u] = v;
  void blo(int u, int v, queue < int > &qe) {
    nb = lca(u, v), fill_n(inb, V + 1, 0);
    upd(u), upd(v);
    if (djs[u] != nb) bk[u] = v;
    if (djs[v] != nb) bk[v] = u;
    for (int tu = 1; tu <= V; ++tu)</pre>
      if (inb[djs[tu]])
        if (djs[tu] = nb, !inq[tu])
          qe.push(tu), inq[tu] = 1;
  void flow() {
    fill_n(inq + 1, V, 0), fill_n(bk + 1, V, 0);
    iota(djs + 1, djs + V + 1, 1);
    queue<int> qe;
    qe.push(st), inq[st] = 1, ed = 0;
    while (!qe.empty()) {
      int u = qe.front();
      qe.pop();
      for (int v = 1; v <= V; ++v)</pre>
        if (el[u][v] && djs[u] != djs[v] &&
          pr[u] != v) {
          if ((v == st) ||
            (pr[v] > 0 \&\& bk[pr[v]] > 0)) {
            blo(u, v, qe);
          } else if (!bk[v]) {
            if (bk[v] = u, pr[v] > 0) {
```

```
if (!inq[pr[v]]) qe.push(pr[v]);
            } else {
              return ed = v, void();
          }
        }
    }
  }
  void aug() {
    for (int u = ed, v, w; u > 0;)
      v = bk[u], w = pr[v], pr[v] = u, pr[u] = v,
  int solve() {
    fill_n(pr, V + 1, 0), ans = 0;
    for (int u = 1; u <= V; ++u)</pre>
      if (!pr[u])
        if (st = u, flow(), ed > \theta) aug(), ++ans;
    return ans;
  }
}:
```

# 4.4 Minimum Weight Matching (Clique version)\*

```
struct Graph { // O-base (Perfect Match), n is even
  int n, match[N], onstk[N], stk[N], tp;
  ll edge[N][N], dis[N];
  void init(int _n) {
    n = _n, tp = 0;
     for (int i = 0; i < n; ++i) fill_n(edge[i], n, 0);</pre>
  void add edge(int u. int v. ll w) {
     edge[u][v] = edge[v][u] = w;
  bool SPFA(int u) {
     stk[tp++] = u, onstk[u] = 1;
     for (int v = 0; v < n; ++v)
  if (!onstk[v] && match[u] != v) {</pre>
         int m = match[v];
         if (dis[m] >
           dis[u] - edge[v][m] + edge[u][v]) {
           dis[m] = dis[u] - edge[v][m] + edge[u][v];
           onstk[v] = 1, stk[tp++] = v;
           if (onstk[m] || SPFA(m)) return 1;
            --tp, onstk[v] = 0;
        }
       }
     onstk[u] = 0, --tp;
     return 0;
  ll solve() { // find a match
     for (int i = 0; i < n; ++i) match[i] = i ^ 1;</pre>
     while (1) {
       int found = 0;
       fill_n(dis, n, 0);
      fill_n(onstk, n, 0);
for (int i = 0; i < n; ++i)</pre>
         if (tp = 0, !onstk[i] && SPFA(i))
           for (found = 1; tp >= 2;) {
              int u = stk[--tp];
             int v = stk[--tp];
             match[u] = v, match[v] = u;
       if (!found) break;
     ll ret = 0:
     for (int i = 0; i < n; ++i)</pre>
      ret += edge[i][match[i]];
     return ret >> 1;
};
```

### 4.5 SW-mincut

```
struct SW{ // global min cut, O(V^3)
  #define REP for (int i = 0; i < n; ++i)
  static const int MXN = 514, INF = 2147483647;
  int vst[MXN], edge[MXN][MXN], wei[MXN];
  void init(int n) {
    REP fill_n(edge[i], n, 0);
  }
  void addEdge(int u, int v, int w){
    edge[u][v] += w; edge[v][u] += w;
  }
  int search(int &s, int &t, int n){</pre>
```

```
fill_n(vst, n, 0), fill_n(wei, n, 0);
     s = t = -1:
     int mx, cur;
     for (int j = 0; j < n; ++j) {</pre>
      mx = -1, cur = 0;
       REP if (wei[i] > mx) cur = i, mx = wei[i];
       vst[cur] = 1, wei[cur] = -1;
       s = t; t = cur;
       REP if (!vst[i]) wei[i] += edge[cur][i];
     return mx;
  int solve(int n) {
     int res = INF;
     for (int x, y; n > 1; n--){
       res = min(res, search(x, y, n));
       REP edge[i][x] = (edge[x][i] += edge[y][i]);
       REP {
         edge[y][i] = edge[n - 1][i];
         edge[i][y] = edge[i][n - 1];
       } // edge[y][y] = 0;
    }
     return res;
} sw;
```

### 4.6 BoundedFlow\*(Dinic\*)

```
struct BoundedFlow { // 0-base
  struct edge {
    int to, cap, flow, rev;
  vector<edge> G[N];
  int n, s, t, dis[N], cur[N], cnt[N];
  void init(int _n) {
    G[i].clear(), cnt[i] = 0;
  void add_edge(int u, int v, int lcap, int rcap) {
  cnt[u] -= lcap, cnt[v] += lcap;
    G[u].pb(edge{v, rcap, lcap, SZ(G[v])});
G[v].pb(edge{u, 0, 0, SZ(G[u]) - 1});
  void add_edge(int u, int v, int cap) {
    G[u].pb(edge\{v, cap, 0, SZ(G[v])\});
    G[v].pb(edge{u, 0, 0, SZ(G[u]) - 1});
  int dfs(int u, int cap) {
    if (u == t || !cap) return cap;
    for (int &i = cur[u]; i < SZ(G[u]); ++i) {</pre>
      edge &e = G[u][i];
      if (dis[e.to] == dis[u] + 1 && e.cap != e.flow) {
        int df = dfs(e.to, min(e.cap - e.flow, cap));
        if (df) {
          e.flow += df, G[e.to][e.rev].flow -= df;
          return df;
      }
    dis[u] = -1;
    return 0;
  bool bfs() {
    fill_n(dis, n + 3, -1);
    queue < int > q;
    q.push(s), dis[s] = 0;
    while (!q.empty()) {
      int u = q.front();
      q.pop():
      for (edge &e : G[u])
        if (!~dis[e.to] && e.flow != e.cap)
          q.push(e.to), dis[e.to] = dis[u] + 1;
    return dis[t] != -1;
  int maxflow(int _s, int _t) {
         _s, t = _t;
    int flow = 0, df;
    while (bfs()) {
      fill_n(cur, n + 3, 0);
      while ((df = dfs(s, INF))) flow += df;
    return flow;
  bool solve() {
```

```
int sum = 0;
     for (int i = 0; i < n; ++i)</pre>
       if (cnt[i] > 0)
          add_edge(n + 1, i, cnt[i]), sum += cnt[i];
       else if (cnt[i] < 0) add_edge(i, n + 2, -cnt[i]);</pre>
     if (sum != maxflow(n + 1, n + 2)) sum = -1;
     for (int i = 0; i < n; ++i)</pre>
       if (cnt[i] > 0)
       G[n + 1].pop_back(), G[i].pop_back();
else if (cnt[i] < 0)</pre>
          G[i].pop_back(), G[n + 2].pop_back();
     return sum != -1;
  int solve(int _s, int _t) {
     add_edge(_t, _s, INF);
if (!solve()) return -1; // invalid flow
int x = G[_t].back().flow;
     return G[_t].pop_back(), G[_s].pop_back(), x;
};
```

#### 4.7 Gomory Hu tree\*

```
MaxFlow Dinic:
int g[MAXN];
void GomoryHu(int n) { // 0-base
  fill_n(g, n, 0);
  for (int i = 1; i < n; ++i) {</pre>
    Dinic.reset();
    add_edge(i, g[i], Dinic.maxflow(i, g[i]));
    for (int j = i + 1; j <= n; ++j)</pre>
      if (g[j] == g[i] && ~Dinic.dis[j])
        g[j] = i;
  }
}
```

#### 4.8 Minimum Cost Circulation\*

```
struct MinCostCirculation { // 0-base
  struct Edge {
    ll from, to, cap, fcap, flow, cost, rev;
  } *past[N];
  vector < Edge > G[N];
  ll dis[N], inq[N], n;
  void BellmanFord(int s) {
    fill_n(dis, n, INF), fill_n(inq, n, 0);
    queue<int> q;
    auto relax = [&](int u, ll d, Edge *e) {
      if (dis[u] > d) {
         dis[u] = d, past[u] = e;
if (!inq[u]) inq[u] = 1, q.push(u);
      }
    };
    relax(s, 0, 0);
    while (!q.empty()) {
  int u = q.front();
      q.pop(), inq[u] = 0;
       for (auto &e : G[u])
         if (e.cap > e.flow)
           relax(e.to, dis[u] + e.cost, &e);
    }
  }
  void try_edge(Edge &cur) {
    if (cur.cap > cur.flow) return ++cur.cap, void();
    BellmanFord(cur.to);
if (dis[cur.from] + cur.cost < 0) {</pre>
       ++cur.flow, --G[cur.to][cur.rev].flow;
       for (int
            i = cur.from; past[i]; i = past[i]->from) {
         auto &e = *past[i];
++e.flow, --G[e.to][e.rev].flow;
      }
    }
    ++cur.cap;
  }
  void solve(int mxlg) {
    for (int b = mxlg; b >= 0; --b) {
      for (int i = 0; i < n; ++i)</pre>
         for (auto &e : G[i])
      e.cap *= 2, e.flow *= 2;
for (int i = 0; i < n; ++i)
         for (auto &e : G[i])
           if (e.fcap >> b & 1)
             try_edge(e);
```

```
void init(int _n) { n = _n;
    for (int i = 0; i < n; ++i) G[i].clear();</pre>
   void add_edge(ll a, ll b, ll cap, ll cost) {
      G[a].pb(Edge
      {a, b, 0, cap, 0, cost, SZ(G[b]) + (a == b)});
G[b].pb(Edge{b, a, 0, 0, 0, -cost, SZ(G[a]) - 1});
} mcmf; // O(VE * ElogC)
```

### 4.9 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source S and sink T.
  - 2. For each edge (x,y,l,u), connect  $x \to y$  with capacity u-l.
  - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If in(v) > 0, connect  $S \to v$  with capacity in(v), otherwise, connect  $v \rightarrow T$  with capacity -in(v).
    - To maximize, connect  $t\to s$  with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T.If  $f 
      eq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from s to t is the answer.
    - To minimize, let f be the maximum flow from S to T. Connect  $t\, o\,s$  with capacity  $\infty$  and let the flow from S to T be f' . If  $f+f' 
      eq \sum_{v \in V, in(v)>0} in(v)$ , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e+f_e$ , where  $f_e$  corresponds to the flow of edge  $\boldsymbol{e}$  on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph(X,Y)
  - 1. Redirect every edge:  $y \rightarrow x$  if  $(x,y) \in M$ ,  $x \rightarrow y$  otherwise.
  - 2. DFS from unmatched vertices in X.
- 3.  $x \in X$  is chosen iff x is unvisited.
- 4.  $y \in Y$  is chosen iff y is visited.
- · Minimum cost cyclic flow
  - 1. Consruct super source  ${\cal S}$  and sink  ${\cal T}$
  - 2. For each edge (x,y,c), connect  $x \to y$  with (cost,cap) = (c,1) if c > 0, otherwise connect  $y \rightarrow x$  with (cost, cap) = (-c, 1)
  - 3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
  - 4. For each vertex v with d(v) > 0, connect  $S \rightarrow$ (cost, cap) = (0, d(v))
  - 5. For each vertex v with d(v) < 0, connect  $v \rightarrow T$  with (cost, cap) = (0, -d(v))
- 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer T
  - 2. Construct a max flow model, let K be the sum of all weights
  - 3. Connect source  $s \rightarrow v$  ,  $v \in G$  with capacity K
  - 4. For each edge (u,v,w) in G, connect  $u \to v$  and  $v \to u$  with capacity w
  - 5. For  $v \in G$ , connect it with sink  $v \to t$  with capacity  $K+2T-(\sum_{e \in E(v)} w(e))-2w(v)$
  - 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v', and connect  $u' \to v'$  with weight w(u,v).
  - 2. Connect v 
    ightarrow v' with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to v.
  - 3. Find the minimum weight perfect matching on G'.
- · Project selection problem
  - 1. If  $p_v > 0$ , create edge (s, v) with capacity  $p_v$ ; otherwise, create edge (v,t) with capacity  $-p_v$
  - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
  - 3. The mincut is equivalent to the maximum profit of a subset of projects.
- · Dual of minimum cost maximum flow
  - 1. Capacity  $c_{uv}$ , Flow  $f_{uv}$ , Cost  $w_{uv}$ , Required Flow difference for vertex
  - 2. If all  $w_{uv}$  are integers, then optimal solution can happen when all  $p_u$ are integers.

are integers. 
$$\begin{aligned} \min & \sum_{uv} w_{uv} f_{uv} \\ & -f_{uv} \geq -c_{uv} \Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv}) \\ & \sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_{u} \end{aligned}$$

# 5 String 5.1 KMP

```
int F[MAXN];
vector < int > match(string A, string B) {
  vector < int > ans;
  F[0] = -1, F[1] = 0;
  for (int i = 1, j = 0; i < SZ(B); F[++i] = ++j) {
    if (B[i] == B[j]) F[i] = F[j]; // optimize
    while (j != -1 && B[i] != B[j]) j = F[j];
}
for (int i = 0, j = 0; i < SZ(A); ++i) {
    while (j != -1 && A[i] != B[j]) j = F[j];
    if (++j == SZ(B)) ans.pb(i + 1 - j), j = F[j];
}
return ans;
}</pre>
```

### 5.2 Z-value\*

```
int z[MAXn];
void make_z(const string &s) {
  int l = 0, r = 0;
  for (int i = 1; i < SZ(s); ++i) {
    for (z[i] = max(0, min(r - i + 1, z[i - l]));
        i + z[i] < SZ(s) && s[i + z[i]] == s[z[i]];
        ++z[i])
    ;
  if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
  }
}
```

### 5.3 Manacher\*

### 5.4 **SAIS\***

```
namespace sfx {
      _t[N * 2];
int SA[N * 2], H[N], RA[N];
int _s[N * 2], _c[N * 2], x[N], _p[N], _q[N * 2];
// zero based, string content MUST > 0
// SA[i]: SA[i]-th
    suffix is the i-th lexigraphically smallest suffix.
// H[i]: longest
    common prefix of suffix SA[i] and suffix SA[i - 1].
void pre(int *sa, int *c, int n, int z)
{ fill_n(sa, n, 0), copy_n(c, z, x); }
void induce
    (int *sa, int *c, int *s, bool *t, int n, int z) {
  copy_n(c, z - 1, x + 1);
  for (int i = 0; i < n; ++i)</pre>
    if (sa[i] && !t[sa[i] - 1])
      sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
  copy_n(c, z, x);
for (int i = n - 1; i >= 0; --i)
    if (sa[i] && t[sa[i] - 1])
      sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa
      int *p, int *q, bool *t, int *c, int n, int z) {
  bool uniq = t[n - 1] = true;
  int nn = 0,
       nmxz = -1, *nsa = sa + n, *ns = s + n, last = -1;
  fill_n(c, z, 0);
  for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;</pre>
  partial_sum(c, c + z, c);
  if (uniq) {
    for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
    return;
```

```
for (int i = n - 2; i >= 0; --i)
     t[i] = (
         s[i] == s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
  pre(sa, c, n, z);
   for (int i = 1; i <= n - 1; ++i)</pre>
     if (t[i] && !t[i - 1])
       sa[--x[s[i]]] = p[q[i] = nn++] = i;
  induce(sa, c, s, t, n, z);
  for (int i = 0; i < n; ++i)</pre>
     if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
       bool neq = last < 0 || !equal
            (s + sa[i], s + p[q[sa[i]] + 1], s + last);
       ns[q[last = sa[i]]] = nmxz += neq;
  sais(ns,
        nsa, p + nn, q + n, t + n, c + z, nn, nmxz + 1);
  pre(sa, c, n, z);
for (int i = nn - 1; i >= 0; --i)
     sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
  induce(sa, c, s, t, n, z);
void mkhei(int n) {
  for (int i = 0, j = 0; i < n; ++i) {</pre>
     if (RA[i])
    for (; _s[i + j] == _s[SA[RA[i] - 1] + j]; ++j);
H[RA[i]] = j, j = max(0, j - 1);
  }
void build(int *s, int n) {
  copy_n(s, n, _s), _s[n] = 0;
sais(_s, SA, _p, _q, _t, _c, n + 1, 256);
  copy_n(SA + 1, n, SA);
  for (int i = 0; i < n; ++i) RA[SA[i]] = i;</pre>
  mkhei(n);
}}
```

### 5.5 Aho-Corasick Automatan

```
const int len = 400000, sigma = 26;
struct AC_Automatan {
  int nx[len][sigma], fl[len], cnt[len], pri[len], top;
  int newnode() {
    fill(nx[top], nx[top] + sigma, -1);
    return top++;
  void init() { top = 1, newnode(); }
  int input(
    string &s) { // return the end_node of string
    int X = 1;
    for (char c : s) {
      if (!~nx[X][c - 'a']) nx[X][c - 'a'] = newnode();
X = nx[X][c - 'a'];
    return X;
  void make_fl() {
    queue < int > q;
    q.push(1), fl[1] = 0;
    for (int t = 0; !q.empty();) {
      int R = q.front();
      q.pop(), pri[t++] = R;
for (int i = 0; i < sigma; ++i)</pre>
        if (~nx[R][i]) {
           int X = nx[R][i], Z = fl[R];
           for (; Z && !~nx[Z][i];) Z = fl[Z];
           fl[X] = Z ? nx[Z][i] : 1, q.push(X);
        }
   }
  void get_v(string &s) {
    int X = 1;
    fill(cnt, cnt + top, \theta);
    for (char c : s) {
      while (X \&\& !\sim nx[X][c - 'a']) X = fl[X];
      X = X ? nx[X][c - 'a'] : 1, ++cnt[X];
    for (int i = top - 2; i > 0; --i)
      cnt[fl[pri[i]]] += cnt[pri[i]];
```

### 5.6 Smallest Rotation

```
string mcp(string s) {
  int n = SZ(s), i = 0, j = 1;
```

```
s += s;
while (i < n && j < n) {
   int k = 0;
   while (k < n && s[i + k] == s[j + k]) ++k;
   if (s[i + k] <= s[j + k]) j += k + 1;
   else i += k + 1;
   if (i == j) ++j;
}
int ans = i < n ? i : j;
return s.substr(ans, n);
}</pre>
```

### 5.7 De Bruijn sequence\*

```
constexpr int MAXC = 10, MAXN = 1e5 + 10;
struct DBSeq {
  int C, N, K, L, buf[MAXC * MAXN]; // K <= C^N</pre>
  void dfs(int *out, int t, int p, int &ptr) {
    if (ptr >= L) return;
    if (t > N) {
      if (N % p) return;
      for (int i = 1; i <= p && ptr < L; ++i)</pre>
        out[ptr++] = buf[i];
     else -
      buf[t] = buf[t - p], dfs(out, t + 1, p, ptr);
      for (int j = buf[t - p] + 1; j < C; ++j)</pre>
        buf[t] = j, dfs(out, t + 1, t, ptr);
    }
  void solve(int _c, int _n, int _k, int *out) {
    int p = 0;
    C = c, N = n, K = k, L = N + K - 1; dfs(out, 1, 1, p);
    if (p < L) fill(out + p, out + L, 0);</pre>
} dbs;
```

### 5.8 Extended SAM\*

```
struct exSAM {
  int len[N * 2], link[N * 2]; // maxlength, suflink int next[N * 2][CNUM], tot; // [\theta, tot), root = \theta int lenSorted[N * 2]; // topo. order
  int cnt[N * 2]; // occurence
  int newnode() {
    fill_n(next[tot], CNUM, 0);
    len[tot] = cnt[tot] = link[tot] = 0;
    return tot++;
  void init() { tot = 0, newnode(), link[0] = -1; }
  int insertSAM(int last, int c) {
    int cur = next[last][c];
    len[cur] = len[last] + 1;
    int p = link[last];
    while (p != -1 && !next[p][c])
  next[p][c] = cur, p = link[p];
    if (p == -1) return link[cur] = 0, cur;
    int q = next[p][c];
    if (len
         [p] + 1 == len[q]) return link[cur] = q, cur;
    int clone = newnode();
    for (int i = 0; i < CNUM; ++i)</pre>
      next[
           clone][i] = len[next[q][i]] ? next[q][i] : 0; | };
    len[clone] = len[p] + 1;
while (p != -1 && next[p][c] == q)
      next[p][c] = clone, p = link[p];
    link[link[cur] = clone] = link[q];
    link[q] = clone;
    return cur:
  void insert(const string &s) {
    int cur = 0;
    for (auto ch : s) {
       int &nxt = next[cur][int(ch - 'a')];
       if (!nxt) nxt = newnode();
       cnt[cur = nxt] += 1;
    }
  }
  void build() {
    queue < int > q;
    q.push(0);
    while (!q.empty()) {
      int cur = q.front();
      q.pop();
```

for (int i = 0; i < CNUM; ++i)</pre>

### 5.9 PalTree\*

```
struct palindromic_tree {
  struct node {
    int next[26], fail, len;
    node(int l = 0) : fail(0), len(l), cnt(0), num(0) {
      for (int i = 0; i < 26; ++i) next[i] = 0;</pre>
  };
  vector<node> St;
  vector < char > s;
  int last, n;
  palindromic_tree() : St(2), last(1), n(0) {
    St[0].fail = 1, St[1].len = -1, s.pb(-1);
  inline void clear() {
    St.clear(), s.clear(), last = 1, n = 0;
    St.pb(0), St.pb(-1);
    St[0].fail = 1, s.pb(-1);
  inline int get_fail(int x) {
  while (s[n - St[x].len - 1] != s[n])
      x = St[x].fail;
    return x;
  inline void add(int c) {
  s.push_back(c -= 'a'), ++n;
    int cur = get_fail(last);
    if (!St[cur].next[c]) {
      int now = SZ(St);
      St.pb(St[cur].len + 2);
      St[now].fail =
        St[get_fail(St[cur].fail)].next[c];
      St[cur].next[c] = now;
      St[now].num = St[St[now].fail].num + 1;
    last = St[cur].next[c], ++St[last].cnt;
  inline void count() { // counting cnt
    auto i = St.rbegin();
    for (; i != St.rend(); ++i) {
      St[i->fail].cnt += i->cnt;
  inline int size() { // The number of diff. pal.
    return SZ(St) - 2;
```

### 6 Math

### 6.1 ax+by=gcd(only exgcd \*)

```
pll exgcd(ll a, ll b) {
   if (b == 0) return pll(1, 0);
   ll p = a / b;
   pll q = exgcd(b, a % b);
   return pll(q.Y, q.X - q.Y * p);
}
/* ax+by=res, let x be minimum non-negative
g, p = gcd(a, b), exgcd(a, b) * res / g
if p.X < 0: t = (abs(p.X) + b / g - 1) / (b / g)
else: t = -(p.X / (b / g))
p += (b / g, -a / g) * t */</pre>
```

### 6.2 Floor and Ceil

```
int floor(int a, int b)
{ return a / b - (a % b && (a < 0) ^ (b < 0)); }
int ceil(int a, int b)
{ return a / b + (a % b && (a < 0) ^ (b > 0)); }
```

### 6.3 Floor Enumeration

```
// enumerating x = floor(n / i), [l, r]
for (int l = 1, r; l <= n; l = r + 1) {
  int x = n / l;
  r = n / x;
}</pre>
```

### 6.4 Mod Min

```
// min{k / l <= ((ak) mod m) <= r}, no solution -> -1
ll mod_min(ll a, ll m, ll l, ll r) {
  if (a == 0) return l ? -1 : 0;
  if (ll k = (l + a - 1) / a; k * a <= r)
    return k;
  ll b = m / a, c = m % a;
  if (ll y = mod_min(c, a, a - r % a, a - l % a))
    return (l + y * c + a - 1) / a + y * b;
  return -1;
}</pre>
```

### 6.5 Gaussian integer gcd

```
cpx gaussian_gcd(cpx a, cpx b) {
#define rnd
    (a, b) ((a >= 0 ? a * 2 + b : a * 2 - b) / (b * 2))
    ll c = a.real() * b.real() + a.imag() * b.imag();
    ll d = a.imag() * b.real() - a.real() * b.imag();
    ll r = b.real() * b.real() + b.imag() * b.imag();
    if (c % r == 0 && d % r == 0) return b;
    return gaussian_gcd
        (b, a - cpx(rnd(c, r), rnd(d, r)) * b);
}
```

### 6.6 Miller Rabin\*

### 6.7 Simultaneous Equations

```
struct matrix { //m variables, n equations
  int n, m;
  fraction M[MAXN][MAXN + 1], sol[MAXN];
  int solve() { //-1: inconsistent, >= 0: rank
    for (int i = 0; i < n; ++i) {</pre>
      int piv = 0;
      while (piv < m && !M[i][piv].n) ++piv;</pre>
       if (piv == m) continue;
      for (int j = 0; j < n; ++j) {
         if (i == j) continue;
         fraction tmp = -M[j][piv] / M[i][piv];
         for (int k = 0; k <=</pre>
              m; ++k) M[j][k] = tmp * M[i][k] + M[j][k];
      }
    int rank = 0;
    for (int i = 0; i < n; ++i) {</pre>
      int piv = 0;
      while (piv < m && !M[i][piv].n) ++piv;</pre>
      if (piv == m && M[i][m].n) return -1;
      else if (piv
            < m) ++rank, sol[piv] = M[i][m] / M[i][piv];</pre>
    return rank;
  }
};
```

### 6.8 Pollard Rho\*

```
map<ll, int> cnt;
void PollardRho(ll n) {
  if (n == 1) return;
  if (prime(n)) return ++cnt[n], void();
  if (n % 2
      == 0) return PollardRho(n / 2), ++cnt[2], void();
  ll x = 2, y = 2, d = 1, p = 1;
  #define f(x, n, p) ((mul(x, x, n) + p) % n)
  while (true) {
    if (d != n && d != 1) {
      PollardRho(n / d);
      PollardRho(d);
      return;
    if (d == n) ++p;
    x = f(x, n, p), y = f(f(y, n, p), n, p);
d = gcd(abs(x - y), n);
  }
}
```

### 6.9 Simplex Algorithm

```
const int MAXN = 11000, MAXM = 405;
const double eps = 1E-10;
double a[MAXN][MAXM], b[MAXN], c[MAXM];
double d[MAXN][MAXM], x[MAXM];
int ix[MAXN + MAXM]; // !!! array all indexed from 0
// max{cx} subject to {Ax<=b,x>=0}
// n: constraints, m: vars !!!
// x[] is the optimal solution vector
// usage :
// value = simplex(a, b, c, N, M);
double simplex(int n, int m){
  fill_n(d[n], m + 1, 0);
fill_n(d[n + 1], m + 1, 0);
  iota(ix, ix + n + m, 0);
  int r = n, s = m - 1;
for (int i = 0; i < n; ++i) {
     for (int j = 0; j < m - 1; ++j) d[i][j] = -a[i][j];</pre>
     d[i][m - 1] = 1;
d[i][m] = b[i];
     if (d[r][m] > d[i][m]) r = i;
  copy_n(c, m - 1, d[n]);
  d[n + 1][m - 1] = -1;
for (double dd;; ) {
     if (r < n) {
       swap(ix[s], ix[r + m]);
       d[r][s] = 1.0 / d[r][s];
for (int j = 0; j <= m; ++j)</pre>
         if (j != s) d[r][j] *= -d[r][s];
       for (int i = 0; i <= n + 1; ++i) if (i != r) {
  for (int j = 0; j <= m; ++j) if (j != s)</pre>
            d[i][j] += d[r][j] * d[i][s];
          d[i][s] *= d[r][s];
       }
     }
     \Gamma = S = -1;
     for (int j = 0; j < m; ++j)
       if (s < 0 || ix[s] > ix[j]) {
          if (d[n + 1][j] > eps ||
              (d[n + 1][j] > -eps && d[n][j] > eps))
            s = j;
     if (s < 0) break;</pre>
     for (int i = 0; i < n; ++i) if (d[i][s] < -eps) {</pre>
       if (r < 0 ||
            (dd = d[r][m]
                  / d[r][s] - d[i][m] / d[i][s]) < -eps ||
            (dd < eps && ix[r + m] > ix[i + m]))
          r = i:
     if (r < 0) return -1; // not bounded</pre>
  if (d[n + 1][m] < -eps) return -1; // not executable</pre>
  double ans = 0;
  fill_n(x, m, 0);
  for (int i = m; i <
        n + m; ++i) { // the missing enumerated x[i] = 0
     if (ix[i] < m - 1){
  ans += d[i - m][m] * c[ix[i]];</pre>
       x[ix[i]] = d[i-m][m];
```

```
return ans;
}
```

### 6.9.1 Construction

```
Standard form: maximize \mathbf{c}^T\mathbf{x} subject to A\mathbf{x} \leq \mathbf{b} and \mathbf{x} \geq 0. Dual LP: minimize \mathbf{b}^T\mathbf{y} subject to A^T\mathbf{y} \geq \mathbf{c} and \mathbf{y} \geq 0. \bar{\mathbf{x}} and \bar{\mathbf{y}} are optimal if and only if for all i \in [1,n], either \bar{x}_i = 0 or \sum_{j=1}^m A_{ji}\bar{y}_j = c_i holds and for all i \in [1,m] either \bar{y}_i = 0 or \sum_{j=1}^n A_{ij}\bar{x}_j = b_j holds.

1. In case of minimization, let c_i' = -c_i
2. \sum_{1 \leq i \leq n} A_{ji}x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji}x_i \leq -b_j
3. \sum_{1 \leq i \leq n} A_{ji}x_i = b_j
• \sum_{1 \leq i \leq n} A_{ji}x_i \leq b_j
• \sum_{1 \leq i \leq n} A_{ji}x_i \geq b_j
4. If x_i has no lower bound, replace x_i with x_i - x_i'
```

### 6.10 chineseRemainder

```
ll solve(ll x1, ll m1, ll x2, ll m2) {
    ll g = gcd(m1, m2);
    if ((x2 - x1) % g) return -1; // no sol
    m1 /= g; m2 /= g;
    pll p = exgcd(m1, m2);
    ll lcm = m1 * m2 * g;
    ll res = p.first * (x2 - x1) * m1 + x1;
    // be careful with overflow
    return (res % lcm + lcm) % lcm;
}
```

### 6.11 Factorial without prime factor\*

```
// O(p^k + log^2 n), pk = p^k
ll prod[MAXP];
ll fac_no_p(ll n, ll p, ll pk) {
  prod[0] = 1;
  for (int i = 1; i <= pk; ++i)
    if (i % p) prod[i] = prod[i - 1] * i % pk;
    else prod[i] = prod[i - 1];
ll rt = 1;
  for (; n; n /= p) {
    rt = rt * mpow(prod[pk], n / pk, pk) % pk;
    rt = rt * prod[n % pk] % pk;
}
return rt;
} // (n! without factor p) % p^k</pre>
```

### 6.12 OuadraticResidue\*

```
int Jacobi(int a, int m) {
  int s = 1;
  for (; m > 1; ) {
   a %= m:
   if (a == 0) return 0;
    const int r = __builtin_ctz(a);
    if ((r & 1) & ((m + 2) & 4)) s = -s;
    if (a \& m \& 2) s = -s;
    swap(a, m);
 }
int QuadraticResidue(int a, int p) {
 if (p == 2) return a & 1;
  const int jc = Jacobi(a, p);
  if (jc == 0) return 0;
  if (jc == -1) return -1;
  int b, d;
  for (; ; ) {
   b = rand() % p;
    d = (1LL * b * b + p - a) \% p;
   if (Jacobi(d, p) == -1) break;
  int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
  for (int e = (1LL + p) >> 1; e; e >>= 1) {
    if (e & 1) {
     tmp = (1LL *
          g0 * f0 + 1LL * d * (1LL * g1 * f1 % p)) % p;
     g1 = (1LL * g0 * f1 + 1LL * g1 * f0) % p;
     g0 = tmp;
    tmp = (1LL)
        * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p)) % p;
```

```
f1 = (2LL * f0 * f1) % p;
f0 = tmp;
}
return g0;
}
```

### 6.13 PiCount\*

```
ll PrimeCount(ll n) { // n \sim 10^13 => < 2s
  if (n <= 1) return 0;
  int v = sqrt(n), s = (v + 1) / 2, pc = 0;
  vector < int > smalls(v + 1), skip(v + 1), roughs(s);
  vector<ll> larges(s);
  for (int i = 2; i <= v; ++i) smalls[i] = (i + 1) / 2;</pre>
  for (int i = 0; i < s; ++i) {</pre>
    roughs[i] = 2 * i + 1;
    larges[i] = (n / (2 * i + 1) + 1) / 2;
  for (int p = 3; p <= v; ++p) {</pre>
    if (smalls[p] > smalls[p - 1]) {
      int q = p * p;
      ++pc;
      if (1LL * q * q > n) break;
      skip[p] = 1;
      for (int i = q; i <= v; i += 2 * p) skip[i] = 1;</pre>
      int ns = 0;
      for (int k = 0; k < s; ++k) {</pre>
        int i = roughs[k];
        if (skip[i]) continue;
        ll d = 1LL * i * p;
        larges[ns] = larges[k] - (d <= v ? larges
            [smalls[d] - pc] : smalls[n / d]) + pc;
        roughs[ns++] = i;
      }
      s = ns;
      for (int j = v / p; j >= p; --j) {
              smalls[j] - pc, e = min(j * p + p, v + 1);
        for (int i = j * p; i < e; ++i) smalls[i] -= c;</pre>
      }
   }
  for (int k = 1; k < s; ++k) {</pre>
    const ll m = n / roughs[k];
    ll t = larges[k] - (pc + k -
    for (int l = 1; l < k; ++l) {</pre>
      int p = roughs[l];
      if (1LL * p * p > m) break;
t -= smalls[m / p] - (pc + l - 1);
    larges[0] -= t:
  return larges[0];
```

### 6.14 Discrete Log\*

```
int DiscreteLog(int s, int x, int y, int m) {
  constexpr int kStep = 32000;
  unordered map < int , int > p;
  int b = 1;
  for (int i = 0; i < kStep; ++i) {</pre>
    p[y] = i;
    y = 1LL * y * x % m;
    b = 1LL * b * x % m;
  for (int i = 0; i < m + 10; i += kStep) {</pre>
    s = 1LL * s * b % m;
    if (p.find(s) != p.end()) return i + kStep - p[s];
  return -1:
int DiscreteLog(int x, int y, int m) {
  if (m == 1) return 0;
  int s = 1;
  for (int i = 0; i < 100; ++i) {</pre>
    if (s == y) return i;
    s = 1LL * s * x % m;
  if (s == y) return 100;
  int p = 100 + DiscreteLog(s, x, y, m);
  if (fpow(x, p, m) != y) return -1;
  return p;
```

### 6.15 Berlekamp Massey

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
   vector<T> d(SZ(output) + 1), me, he;
   for (int f = 0, i = 1; i <= SZ(output); ++i) {
  for (int j = 0; j < SZ(me); ++j)
    d[i] += output[i - j - 2] * me[j];
  if ((d[i] -= output[i - 1]) == 0) continue;</pre>
      if (me.empty()) {
         me.resize(f = i);
          continue;
      vector<T> o(i - f - 1);
      T k = -d[i] / d[f]; o.pb(-k);
      for (T x : he) o.pb(x * k);
      o.resize(max(SZ(o), SZ(me)));

for (int j = 0; j < SZ(me); ++j) o[j] += me[j];

if (i - f + SZ(he) >= SZ(me)) he = me, f = i;
   return me;
```

### 6.16 Primes

```
/* 12721 13331 14341 75577 123457 222557
     556679 999983 1097774749 1076767633 100102021
    999997771 1001010013 1000512343 987654361 999991231
     999888733 98789101 987777733 999991921 1010101333
     1010102101 1000000000039 100000000000037
     2305843009213693951 4611686018427387847
     9223372036854775783 18446744073709551557 */
```

### 6.17 Theorem

Cramer's rule

$$\begin{array}{l} ax+by=e \\ cx+dy=f \\ \end{array} \Rightarrow \begin{array}{l} x=\frac{ed-bf}{ad-bc} \\ y=\frac{af-ec}{ad-bc} \end{array}$$

· Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

· Kirchhoff's Theorem

Denote L be a  $n \times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii} = d(i)$ ,  $L_{ij} = -c$  where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .
- Tutte's Matrix

Let D be a n imes n matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniformly at random) if i < j and  $(i,j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the  ${\rm maximum\ matching\ on\ } G.$ 

- Cayley's Formula
  - Given a degree sequence  $d_1, d_2, ..., d_n$  for each labeled vertices, there are  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.
  - Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex 1, 2, ..., k belong to different components. Then  $T_{n,k} = kn^{n-k-1}$ .
- Erdős–Gallai theorem

A sequence of nonnegative integers  $d_1 \ge \cdots \ge d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if

$$d_1+\dots+d_n \text{ is even and } \sum_{i=1}^k d_i \leq k(k-1)+\sum_{i=k+1}^n \min(d_i,k) \text{ holds for every } 1\leq k\leq n.$$

Gale–Ryser theorem

A pair of sequences of nonnegative integers  $a_1 \ge \cdots \ge a_n$  and  $b_1, \dots, b_n$ is bigraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$  and  $\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i,k)$  holds for

every  $1 \leq k \leq n$ . Fulkerson–Chen–Anstee theorem

A sequence  $(a_1,\ b_1),\ ...\ ,\ (a_n,\ b_n)$  of nonnegative integer pairs with  $a_1 \geq \cdots \geq a_n$  is digraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$  and

$$\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i, k-1) + \sum_{i=k+1}^n \min(b_i, k) \text{ holds for every } 1 \leq k \leq n.$$
 Möbius inversion formula

- Möbius inversion formula
  - $f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$
  - $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$

- Spherical cap
  - A portion of a sphere cut off by a plane.
  - r: sphere radius, a: radius of the base of the cap, h: height of the cap,  $\theta$ : arcsin(a/r)
  - Volume =  $\pi h^2 (3r h)/3 = \pi h(3a^2 + h^2)/6 = \pi r^3 (2 + \cos \theta)(1 h^2)$
  - Area  $= 2\pi rh = \pi(a^2 + h^2) = 2\pi r^2(1 \cos\theta)$ .
- Lagrange multiplier
  - Optimize  $f(x_1,...,x_n)$  when k constraints  $g_i(x_1,...,x_n) = 0$ .
  - Lagrangian function  $\mathcal{L}(x_1,\ldots,x_n,\lambda_1,\ldots,\lambda_k)=f(x_1,\ldots,x_n)=$  $\sum_{i=1}^{n} \lambda_i g_i(x_1,...,x_n).$
  - The solution corresponding to the original constrained optimization is always a saddle point of the Lagrangian function.

#### Estimation 6.18

- Estimation
  - The number of divisors of n is at most around 100 for n < 5e4, 500 for n < 1e7, 2000 for n < 1e10, 200000 for n < 1e19.
  - The number of ways of writing n as a sum of positive integers, disregarding the order of the summands. 1,1,2,3,5,7,11,15,22,30 for  $n = 0 \sim 9$ , 627 for n = 20,  $\sim 2e5$  for n = 50,  $\sim 2e8$  for n = 100.
  - Total number of partitions of n distinct elements: B(n) =1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, 115975, 678570, 4213597,27644437,190899322,...

#### 6.19 **Euclidean Algorithms**

- m = | <sup>an+b</sup>/<sub>an+b</sub> |
- Time complexity:  $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ + f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ + g(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1)) \\ - h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c - b - 1, a, m - 1) \\ - 2f(c, c - b - 1, a, m - 1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

#### 6.20 General Purpose Numbers

Bernoulli numbers

Sernouth numbers 
$$B_0-1, B_1^{\pm}=\pm\frac{1}{2}, B_2=\frac{1}{6}, B_3=0$$
 
$$\sum_{j=0}^m {m+1 \choose j} B_j=0, \text{EGF is } B(x)=\frac{x}{e^x-1}=\sum_{n=0}^\infty B_n \frac{x^n}{n!}.$$
 
$$S_m(n)=\sum_{k=1}^n k^m=\frac{1}{m+1}\sum_{k=0}^m {m+1 \choose k} B_k^+ n^{m+1-k}$$

Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1$$
 
$$S(n,k) = \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} {k \choose i} i^n$$
 
$$x^n = \sum_{i=0}^n S(n,i)(x)_i$$
 • Pentagonal number theorem

• Pentagonal number theorem 
$$\prod_{n=1}^{\infty}(1-x^n)=1+\sum_{k=1}^{\infty}(-1)^k\left(x^{k(3k+1)/2}+x^{k(3k-1)/2}\right)$$
 • Catalan numbers 
$$C_n^{(k)}=\frac{1}{(k-1)n+1}\binom{kn}{n}$$
 
$$C_n^{(k)}(x)=1+x[C_n^{(k)}(x)]^k$$

$$C_n^{(k)} = \frac{1}{(k-1)n+1} \binom{kn}{n}$$
 
$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

```
• Eulerian numbers  \begin{array}{ll} \text{Number of permutations } \pi \in S_n \text{ in which exactly } k \text{ elements are} \\ \text{greater than the previous element. } k \text{ $j$:s s.t. } \pi(j) > \pi(j+1), k+1 \text{ $j$:s s.t.} \\ \pi(j) \geq j, k \text{ $j$:s s.t. } \pi(j) > j. \\ E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k) \\ E(n,0) = E(n,n-1) = 1 \\ E(n,k) = \sum_{j=0}^k (-1)^j \binom{n+1}{j} (k+1-j)^n \end{array}
```

### 6.21 Tips for Generating Functions

```
 \begin{array}{ll} \bullet & \text{Ordinary Generating Function } A(x) = \sum_{i \geq 0} a_i x^i \\ & - A(rx) \Rightarrow r^n a_n \\ & - A(x) + B(x) \Rightarrow a_n + b_n \\ & - A(x) B(x) \Rightarrow \sum_{i=0}^n a_i b_{n-i} \\ & - A(x)^k \Rightarrow \sum_{i_1 + i_2 + \dots + i_k = n} a_{i_1} a_{i_2} \dots a_{i_k} \\ & - x A(x)' \Rightarrow n a_n \\ & - \frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^n a_i \\ \bullet & \text{Exponential Generating Function } A(x) = \sum_{i \geq 0} \frac{a_i}{i!} x_i \\ & - A(x) + B(x) \Rightarrow a_n + b_n \\ & - A(x) + B(x) \Rightarrow a_n + b_n \\ & - A(x) B(x) \Rightarrow \sum_{i=0}^n \binom{n}{i} a_i b_{n-i} \\ & - A(x)^k \Rightarrow \sum_{i_1 + i_2 + \dots + i_k = n} \binom{n}{i_1, i_2, \dots, i_k} a_{i_1} a_{i_2} \dots a_{i_k} \\ & - x A(x) \Rightarrow n a_n \\ \bullet & \text{Special Generating Function} \\ & - (1+x)^n = \sum_{i \geq 0} \binom{n}{i} x^i \\ & - \frac{1}{(1-x)^n} = \sum_{i \geq 0} \binom{n}{i-1} x^i \end{array}
```

# 7 Polynomial

### 7.1 Fast Fourier Transform

```
template < int MAXN >
struct FFT {
  using val_t = complex < double >;
  const double PI = acos(-1);
  val_t w[MAXN];
  FFT() {
    for (int i = 0; i < MAXN; ++i) {</pre>
      double arg = 2 * PI * i / MAXN;
      w[i] = val_t(cos(arg), sin(arg));
    }
  }
  void bitrev(val_t *a, int n); // see NTT
  void trans
      (val_t *a, int n, bool inv = false); // see NTT;
    remember to replace LL with val_t
};
```

### 7.2 Number Theory Transform\*

```
//(2^16)+1, 65537,
 /7*17*(2^23)+1, 998244353, 3
//1255*(2^20)+1, 1315962881, 3
//51*(2^25)+1, 1711276033, 29
template <int MAXN, ll P, ll RT> //MAXN must be 2^k
struct NTT {
  ll w[MAXN];
  ll mpow(ll a, ll n);
  ll minv(ll a) { return mpow(a, P - 2); }
  NTT() {
    ll dw = mpow(RT, (P - 1) / MAXN);
    w[0] = 1;
    for (int
        i = 1; i < MAXN; ++i) w[i] = w[i - 1] * dw % P;
  void bitrev(ll *a, int n) {
    int i = 0;
    for (int j = 1; j < n - 1; ++j) {</pre>
      for (int k = n >> 1; (i ^= k) < k; k >>= 1);
      if (j < i) swap(a[i], a[j]);</pre>
    }
  void operator()(
      ll *a, int n, bool inv = false) { //0 <= a[i] < P
    bitrev(a, n);
    for (int L = 2; L <= n; L <<= 1) {</pre>
       int dx = MAXN / L, dl = L >> 1;
      for (int i = 0; i < n; i += L) {</pre>
        for (int
          j = i, x = 0; j < i + dl; ++j, x += dx) { ll tmp = a[j + dl] * w[x] % P;
           if ((a[j
                 + dl] = a[j] - tmp) < 0) a[j + dl] += P;
          if ((a[j] += tmp) >= P) a[j] -= P;
        }
```

### 7.3 Fast Walsh Transform\*

```
/* x: a[j], y: a[j + (L >> 1)]
or: (y += x * op), and: (x += y * op)
xor: (x, y = (x + y) * op, (x - y) * op)
invop: or, and, xor = -1, -1, 1/2 */
void fwt(int *a, int n, int op) { //or
  for (int L = 2; L <= n; L <<= 1)
  for (int i = 0; i < n; i += L)</pre>
        for (int j = i; j < i + (L >> 1); ++j)
           a[j + (L >> 1)] += a[j] * op;
const int N = 21;
int f[
     N][1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N];
void
      subset_convolution(int *a, int *b, int *c, int L) {
   // c_k = \sum_{i=0}^{n} \{i \mid j = k, i \& j = 0\} a_i * b_j
   int n = 1 << L;
   for (int i = 1; i < n; ++i)</pre>
     ct[i] = ct[i & (i - 1)] + 1;
   for (int i = 0; i < n; ++i)</pre>
     f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
   for (int i = 0; i <= L; ++i)</pre>
     fwt(f[i], n, 1), fwt(g[i], n, 1);
   for (int i = 0; i <= L; ++i)</pre>
     for (int j = 0; j <= i; ++j)
  for (int x = 0; x < n; ++x)
    h[i][x] += f[j][x] * g[i - j][x];</pre>
   for (int i = 0; i <= L; ++i)
  fwt(h[i], n, -1);
for (int i = 0; i < n; ++i)</pre>
     c[i] = h[ct[i]][i];
```

### 7.4 Polynomial Operation

```
fi(s, n) for (int i = (int)(s); i < (int)(n); ++i)
template < int MAXN, ll P, ll RT> // MAXN = 2^k
struct Poly : vector<ll> { // coefficients in [0, P)
  using vector<ll>::vector;
  static NTT < MAXN, P, RT > ntt;
  int n() const { return (int)size(); } // n() >= 1
Poly(const Poly &p, int m) : vector<ll>(m) {
    copy_n(p.data(), min(p.n(), m), data());
  Poly& irev()
  { return reverse(data(), data() + n()), *this; }
Poly& isz(int m) { return resize(m), *this; }
  Poly& iadd(const Poly &rhs) { // n() == rhs.n()
    fi(0, n()) if
          (((*this)[i] += rhs[i]) >= P) (*this)[i] -= P;
    return *this;
  Poly& imul(ll k) {
    fi(0, n()) (*this)[i] = (*this)[i] * k % P;
  Poly Mul(const Poly &rhs) const {
    int m = 1;
    while (m < n() + rhs.n() - 1) m <<= 1;</pre>
    Poly X(*this, m), Y(rhs, m);
    ntt(X.data(), m), ntt(Y.data(), m);
    fi(0, m) X[i] = X[i] * Y[i] % P;
    ntt(X.data(), m, true);
    return X.isz(n() + rhs.n() - 1);
  Poly Inv() const { // (*this)[0] != 0, 1e5/95ms
    if (n() == 1) return {ntt.minv((*this)[0])};
    int m = 1;
    while (m < n() * 2) m <<= 1;</pre>
    Poly Xi = Poly(*this, (n() + 1) / 2). Inv().isz(m);
    Poly Y(*this, m);
    ntt(Xi.data(), m), ntt(Y.data(), m);
```

```
if ((Xi[i] %= P) < 0) Xi[i] += P;</pre>
  ntt(Xi.data(), m, true);
  return Xi.isz(n());
Poly Sqrt()
     const { // Jacobi((*this)[0], P) = 1, 1e5/235ms
  if (n()
      == 1) return {QuadraticResidue((*this)[0], P)};
  Poly
      X = Poly(*this, (n() + 1) / 2).Sqrt().isz(n());
  return
       X.iadd(Mul(X.Inv()).isz(n())).imul(P / 2 + 1);
}
pair < Poly , Poly > DivMod
    (const Poly &rhs) const { // (rhs.)back() != 0
  if (n() < rhs.n()) return {{0}, *this};</pre>
  const int m = n() - rhs.n() + 1;
  Poly X(rhs); X.irev().isz(m);
  Poly Y(*this); Y.irev().isz(m);
  Poly Q = Y.Mul(X.Inv()).isz(m).irev();
  X = rhs.Mul(Q), Y = *this;
  fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;
  return {Q, Y.isz(max(1, rhs.n() - 1))};
Poly Dx() const {
  Poly ret(n() - 1);
  fi(0,
      ret.n()) ret[i] = (i + 1) * (*this)[i + 1] % P;
  return ret.isz(max(1, ret.n()));
Poly Sx() const {
  Poly ret(n() + 1);
  fi(0, n())
       ret[i + 1] = ntt.minv(i + 1) * (*this)[i] % P;
Poly _tmul(int nn, const Poly &rhs) const {
  Poly Y = Mul(rhs).isz(n() + nn - 1);
return Poly(Y.data() + n() - 1, Y.data() + Y.n());
vector<ll> _eval(const
     vector<ll> &x, const vector<Poly> &up) const {
  const int m = (int)x.size();
  if (!m) return {};
  vector<Poly> down(m * 2);
  // down[1] = DivMod(up[1]).second;
  // fi(2, m *
       2) down[i] = down[i / 2].DivMod(up[i]).second;
  down[1] = Poly(up[1])
      .irev().isz(n()).Inv().irev()._tmul(m, *this);
  fi(2, m * 2) down[i]
       = up[i ^ 1]._tmul(up[i].n() - 1, down[i / 2]);
  vector<ll> y(m);
  fi(0, m) y[i] = down[m + i][0];
  return y;
static vector<Poly> _tree1(const vector<ll> &x) {
  const int m = (int)x.size();
  vector<Poly> up(m * 2);
  fi(0, m) up[m + i] = \{(x[i] ? P - x[i] : 0), 1\};
  for (int i = m - 1; i
      > 0; --i) up[i] = up[i * 2].Mul(up[i * 2 + 1]);
  return up;
    <ll> Eval(const vector<ll> &x) const { // 1e5, 1s
  auto up = _tree1(x); return _eval(x, up);
static Poly Interpolate(const vector
    <ll> &x, const vector<ll> &y) { // 1e5, 1.4s
  const int m = (int)x.size();
  vector<Poly> up = _tree1(x), down(m * 2);
vector<ll> z = up[1].Dx()._eval(x, up);
  fi(\theta, m) z[i] = y[i] * ntt.minv(z[i]) % P;
  fi(0, m) down[m + i] = {z[i]};
  for (int i = m
       1; i > 0; --i) down[i] = down[i * 2].Mul(up[i
      * 2 + 1]).iadd(down[i * 2 + 1].Mul(up[i * 2]));
  return down[1];
Poly Ln() const { // (*this)[0] == 1, 1e5/170ms
  return Dx().Mul(Inv()).Sx().isz(n());
```

```
Poly Exp() const \{ // (*this)[0] == 0, 1e5/360ms \}
           if (n() == 1) return {1};
           Poly X = Poly(*this, (n() + 1) / 2).Exp().isz(n());
           Poly Y = X.Ln(); Y[0] = P - 1;
           fi(0, n())
                         if ((Y[i] = (*this)[i] - Y[i]) < 0) Y[i] += P;</pre>
            return X.Mul(Y).isz(n());
      // M := P(P - 1). If k >= M, k := k % M + M.
      Poly Pow(ll k) const {
           int nz = 0;
            while (nz < n() && !(*this)[nz]) ++nz;</pre>
            if (nz * min(k, (ll)n()) >= n()) return Poly(n());
            if (!k) return Poly(Poly {1}, n());
            Poly X(data() + nz, data() + nz + n() - nz * k);
            const ll c = ntt.mpow(X[0], k % (P - 1));
            return X.Ln().imul
                       (k % P).Exp().imul(c).irev().isz(n()).irev();
      static ll
                 LinearRecursion(const vector<ll> &a, const vector<ll> &coef, ll n) { // a_n = |sum\ c_j\ 
            const int k = (int)a.size();
            assert((int)coef.size() == k + 1);
            Poly C(k + 1), W(Poly \{1\}, k), M = \{0, 1\};
            fi(1, k + 1) C[k - i] = coef[i] ? P - coef[i] : 0;
           C[k] = 1;
            while (n) {
                 if (n % 2) W = W.Mul(M).DivMod(C).second;
                 n /= 2, M = M.Mul(M).DivMod(C).second;
            ll ret = 0;
           fi(0, k) ret = (ret + W[i] * a[i]) % P;
            return ret;
     }
};
#undef fi
using Poly_t = Poly<131072 * 2, 998244353, 3>;
template <> decltype(Poly_t::ntt) Poly_t::ntt = {};
```

### 7.5 Value Polynomial

```
struct Poly {
   mint base; // f(x) = poly[x - base]
   vector<mint> poly;
   Poly(mint b = 0, mint x = 0): base(b), poly(1, x) {}
   mint get_val(const mint &x) {
     if (x >= base && x < base + SZ(poly))
       return poly[x - base];
     mint rt = 0;
     vector<mint> lmul(SZ(poly), 1), rmul(SZ(poly), 1);
     for (int i = 1; i < SZ(poly); ++i)</pre>
       lmul[i] = lmul[i - 1] * (x - (base + i - 1));
     for (int i = SZ(poly) - 2; i >= 0; --i)
       rmul[i] = rmul[i + 1] * (x - (base + i + 1));
     for (int i = 0; i < SZ(poly); ++i)</pre>
       rt += poly[i] * ifac[i] * inegfac
           [SZ(poly) - 1 - i] * lmul[i] * rmul[i];
     return rt;
   void raise() { // g(x) = sigma\{base:x\} f(x)
     if (SZ(poly) == 1 && poly[0] == 0)
       return;
     mint nw = get_val(base + SZ(poly));
     poly.pb(nw);
     for (int i = 1; i < SZ(poly); ++i)</pre>
       poly[i] += poly[i - 1];
   }
| };
```

### 7.6 Newton's Method

Given F(x) where

$$F(x) = \sum_{i=0}^{\infty} \alpha_i (x - \beta)^i$$

for  $\beta$  being some constant. Polynomial P such that F(P)=0 can be found iteratively. Denote by  $Q_k$  the polynomial such that  $F(Q_k)=0$  (mod  $x^{2^k}$ ), then

$$Q_{k+1}\!=\!Q_k\!-\!\frac{F(Q_k)}{F'(Q_k)}\pmod{x^{2^{k+1}}}$$

# B Geometry B.1 Default Code

```
typedef pair<double, double> pdd;
typedef pair<pdd, pdd> Line;
struct Cir{ pdd 0; double R; };
const double eps = 1e-8;
pdd operator+(pdd a, pdd b)
{ return pdd(a.X + b.X, a.Y + b.Y); }
pdd operator - (pdd a, pdd b)
{ return pdd(a.X - b.X, a.Y - b.Y); }
pdd operator*(pdd a, double b)
{ return pdd(a.X * b, a.Y * b); }
pdd operator/(pdd a, double b)
{ return pdd(a.X / b, a.Y / b); }
double dot(pdd a, pdd b)
 return a.X * b.X + a.Y * b.Y; }
double cross(pdd a, pdd b)
{ return a.X * b.Y - a.Y * b.X; }
double abs2(pdd a)
{ return dot(a, a); }
double abs(pdd a)
{ return sqrt(dot(a, a)); }
int sign(double a)
{ return fabs(a) < eps ? 0 : a > 0 ? 1 : -1; }
int ori(pdd a, pdd b, pdd c)
{ return sign(cross(b - a, c - a)); }
bool collinearity(pdd p1, pdd p2, pdd p3)
{ return sign(cross(p1 - p3, p2 - p3)) == 0; }
bool btw(pdd p1, pdd p2, pdd p3) {
  if (!collinearity(p1, p2, p3)) return 0;
  return sign(dot(p1 - p3, p2 - p3)) <= θ;</pre>
bool seg_intersect(pdd p1, pdd p2, pdd p3, pdd p4) {
  int a123 = ori(p1, p2, p3);
  int a124 = ori(p1, p2, p4);
  int a341 = ori(p3, p4, p1);
  int a342 = ori(p3, p4, p2);
  if (a123 == 0 && a124 == 0)
    return btw(p1, p2, p3) || btw(p1, p2, p4) ||
      btw(p3, p4, p1) || btw(p3, p4, p2);
  return a123 * a124 <= 0 && a341 * a342 <= 0;
pdd intersect(pdd p1, pdd p2, pdd p3, pdd p4) {
  double a123 = cross(p2 - p1, p3 - p1);
  double a124 = cross(p2 - p1, p4 - p1);
  return (p4
      * a123 - p3 * a124) / (a123 - a124); // C^3 / C^2
pdd perp(pdd p1)
{ return pdd(-p1.Y, p1.X); }
pdd projection(pdd p1, pdd p2, pdd p3)
{ return p1 + (
    p2 - p1) * dot(p3 - p1, p2 - p1) / abs2(p2 - p1); }
pdd reflection(pdd p1, pdd p2, pdd p3)
{ return p3 + perp(p2 - p1)
    ) * cross(p3 - p1, p2 - p1) / abs2(p2 - p1) * 2; }
pdd linearTransformation
    (pdd p0, pdd p1, pdd q0, pdd q1, pdd r) \{
  pdd dp = p1 - p0
      , dq = q1 - q0, num(cross(dp, dq), dot(dp, dq));
  return q0 + pdd(
      cross(r - p0, num), dot(r - p0, num)) / abs2(dp);
\} // from line p0--p1 to q0--q1, apply to r
```

### 8.2 PointSeqDist\*

```
double PointSegDist(pdd q0, pdd q1, pdd p) {
   if (sign(abs(q0 - q1)) == 0) return abs(q0 - p);
if (sign(dot(q1 - q0,
        p - q0)) >= 0 && sign(dot(q0 - q1, p - q1)) >= 0)
   return fabs(cross(q1 - q0, p - q0) / abs(q0 - q1));
return min(abs(p - q0), abs(p - q1));
}
```

### 8.3 Heart

```
pdd circenter
    (pdd p0, pdd p1, pdd p2) { // radius = abs(center)
  p1 = p1 - p0, p2 = p2 - p0;
  double x1 = p1.X, y1 = p1.Y, x2 = p2.X, y2 = p2.Y;
  double m = 2. * (x1 * y2 - y1 * x2);
  center.X = (x1 * x1)
       * y2 - x2 * x2 * y1 + y1 * y2 * (y1 - y2)) / m;
  center.Y = (x1 * x2)
       * (x2 - x1) - y1 * y1 * x2 + x1 * y2 * y2) / m;
  return center + p0;
pdd incenter
    (pdd p1, pdd p2, pdd p3) { // radius = area / s * 2 | #define cmpL(i) sign(cross(C[i] - a, b - a))
```

```
double a =
      abs(p2 - p3), b = abs(p1 - p3), c = abs(p1 - p2);
  double s = a + b + c;
  return (a * p1 + b * p2 + c * p3) / s;
pdd masscenter(pdd p1, pdd p2, pdd p3)
{ return (p1 + p2 + p3) / 3; }
pdd orthcenter(pdd p1, pdd p2, pdd p3)
{ return masscenter
    (p1, p2, p3) * 3 - circenter(p1, p2, p3) * 2; }
```

### 8.4 point in circle

```
// return
       p4 is strictly in circumcircle of tri(p1,p2,p3)
ll sqr(ll x) { return x * x; }
bool in_cc(const pll&
   p1, const pll& p2, const pll& p3, const pll& p4) {
ll u11 = p1.X - p4.X; ll u12 = p1.Y - p4.Y;
ll u21 = p2.X - p4.X; ll u22 = p2.Y - p4.Y;
   ll\ u31 = p3.X - p4.X;\ ll\ u32 = p3.Y - p4.Y;
   ll u13
         = sqr(p1.X) - sqr(p4.X) + sqr(p1.Y) - sqr(p4.Y);
   ll u23
         = sqr(p2.X) - sqr(p4.X) + sqr(p2.Y) - sqr(p4.Y);
   ll u33
         = sqr(p3.X) - sqr(p4.X) + sqr(p3.Y) - sqr(p4.Y);
   __int128 det = (__int128)-u13 * u22 * u31
        + (__int128)u12 * u23 * u31 + (__int128)u13 * u21 * u32 - (__int128)u11 * u23 * u32 - (__int128)
        )u12 * u21 * u33 + (__int128)u11 * u22 * u33;
   return det > eps;
}
```

### 8.5 Convex hull\*

```
void hull(vector<pll> &dots) { // n=1 => ans = {}
  sort(dots.begin(), dots.end());
  vector<pll> ans(1, dots[0]);
  for (int ct = 0; ct < 2; ++ct, reverse(ALL(dots)))</pre>
    for (int i = 1,
         t = SZ(ans); i < SZ(dots); ans.pb(dots[i++]))
      while (SZ(ans) > t && ori
          (ans[SZ(ans) - 2], ans.back(), dots[i]) \le 0)
        ans.pop_back();
  ans.pop_back(), ans.swap(dots);
```

### 8.6 PointInConvex\*

```
bool PointInConvex
    (const vector<pll> &C, pll p, bool strict = true) {
  int a = 1, b = SZ(C) - 1, r = !strict;
  if (SZ(C) == 0) return false;
  if (SZ(C) < 3) return r && btw(C[0], C.back(), p);</pre>
  if (ori(C[0], C[a], C[b]) > 0) swap(a, b);
  if (ori
      (C[0], C[a], p) >= r \mid\mid ori(C[0], C[b], p) <= -r)
    return false;
  while (abs(a - b) > 1) {
    int c = (a + b) / 2;
    (ori(C[0], C[c], p) > 0 ? b : a) = c;
  return ori(C[a], C[b], p) < r;</pre>
```

### 8.7 TangentPointToHull\*

```
/* The point should be strictly out of hull
  return arbitrary point on the tangent line */
pii get_tangent(vector<pll> &C, pll p) {
  auto gao = [&](int s) {
    return cyc_tsearch(SZ(C), [&](int x, int y)
    { return ori(p, C[x], C[y]) == s; });
  };
  return pii(gao(1), gao(-1));
} // return (a, b), ori(p, C[a], C[b]) >= 0
```

### 8.8 Intersection of line and convex

```
int TangentDir(vector<pll> &C, pll dir) {
  return cyc_tsearch(SZ(C), [&](int a, int b) {
    return cross(dir, C[a]) > cross(dir, C[b]);
  });
```

```
pii lineHull(pll a, pll b, vector<pll> &C) {
  int A = TangentDir(C, a - b);
  int B = TangentDir(C, b - a);
  int n = SZ(C);
  if (cmpL(A) < 0 \mid | cmpL(B) > 0)
  return pii(-1, -1); // no collision
auto gao = [&](int l, int r) {
     for (int t = l; (l + 1) % n != r; ) {
  int m = ((l + r + (l < r ? 0 : n)) / 2) % n;</pre>
       (cmpL(m) == cmpL(t) ? l : r) = m;
     return (l + !cmpL(r)) % n;
  pii res = pii(gao(B, A), gao(A, B)); // (i, j)
  if (res.X == res.Y) // touching the corner i
     return pii(res.X, -1);
  if (!
     cmpL(res.X) && !cmpL(res.Y)) // along side i, i+1 switch ((res.X - res.Y + n + 1) % n) {
       case 0: return pii(res.X, res.X);
       case 2: return pii(res.Y, res.Y);
  /* crossing sides (i, i+1) and (j, j+1)
  crossing corner i is treated as side (i, i+1)
  returned
        in the same order as the line hits the convex */
  return res;
} // convex cut: (r, l]
```

### 8.9 minMaxEnclosingRectangle\*

```
const double INF = 1e18, qi = acos(-1) / 2 * 3;
pdd solve(vector<pll> &dots) {
#define diff(u, v) (dots[u] - dots[v])
#define vec(v) (dots[v] - dots[i])
 hull(dots);
  double Max = 0, Min = INF, deg;
  int n = SZ(dots);
  dots.pb(dots[0]);
  for (int i = 0, u = 1, r = 1, l = 1; i < n; ++i) {</pre>
   pll nw = vec(i + 1);
    while (cross(nw, vec(u + 1)) > cross(nw, vec(u)))
     u = (u + 1) \% n;
   while (dot(nw, vec(r + 1)) > dot(nw, vec(r)))
     \Gamma = (\Gamma + 1) \% n;
   if (!i) l = (r + 1) % n;
   while (dot(nw, vec(l + 1)) < dot(nw, vec(l)))</pre>
     l = (l + 1) \% n;
   deg = acos(dot(diff(r
        , l), vec(u)) / abs(diff(r, l)) / abs(vec(u)));
    deg = (qi - deg) / 2;
   Max = max(Max, abs(diff))
        (r, l)) * abs(vec(u)) * sin(deg) * sin(deg));
  return pdd(Min, Max);
```

### 8.10 VectorInPoly\*

### 8.11 PolyUnion\*

```
double rat(pll a, pll b) {
  return sign
     (b.X) ? (double)a.X / b.X : (double)a.Y / b.Y;
} // all poly. should be ccw
double polyUnion(vector<vector<pll>>> &poly) {
  double res = 0;
  for (auto &p : poly)
     for (int a = 0; a < SZ(p); ++a) {</pre>
```

```
pll A = p[a], B = p[(a + 1) \% SZ(p)];
      vector
           <pair < double , int >> segs = {{0, 0}, {1, 0}};
      for (auto &q : poly) {
         if (&p == &q) continue;
        for (int b = 0; b < SZ(q); ++b) {
           pll C = q[b], D = q[(b + 1) \% SZ(q)];
           int sc = ori(A, B, C), sd = ori(A, B, D);
           if (sc != sd && min(sc, sd) < 0) {</pre>
             double sa = cross(D
                  - C, A - C), sb = cross(D - C, B - C);
             segs.emplace_back
                 (sa / (sa - sb), sign(sc - sd));
           if (!sc && !sd &&
               &q < &p && sign(dot(B - A, D - C)) > 0) {
             segs.emplace_back(rat(C - A, B - A), 1);
             segs.emplace_back(rat(D - A, B - A), -1);
        }
      }
      sort(ALL(segs));
      for (auto &s : segs) s.X = clamp(s.X, 0.0, 1.0);
       double sum = 0;
      int cnt = segs[0].second;
      for (int j = 1; j < SZ(segs); ++j) {
  if (!cnt) sum += segs[j].X - segs[j - 1].X;</pre>
        cnt += segs[j].Y;
      res += cross(A, B) * sum;
  return res / 2;
}
```

### 8.12 PolyCut

### 8.13 Polar Angle Sort\*

```
int cmp(pll a, pll b, bool same = true) {
#define is_neg(k) (
    sign(k.Y) < 0 || (sign(k.Y) == 0 && sign(k.X) < 0))
int A = is_neg(a), B = is_neg(b);
if (A != B)
    return A < B;
if (sign(cross(a, b)) == 0)
    return same ? abs2(a) < abs2(b) : -1;
return sign(cross(a, b)) > 0;
}
```

### 8.14 Half plane intersection\*

for (auto p : arr) {

```
pll area_pair(Line a, Line b)
{ return pll(cross(a.Y
      - a.X, b.X - a.X), cross(a.Y - a.X, b.Y - a.X)); }
bool isin(Line l0, Line l1, Line l2) {
  // Check inter(l1, l2) strictly in l0
  auto [a02X, a02Y] = area_pair(l0, l2);
  auto [a12X, a12Y] = area_pair(l1, l2);
  if (a12X - a12Y < 0) a12X *= -1, a12Y *= -1;</pre>
  return (__int128)
        a02Y * a12X - (__int128) a02X * a12Y > 0; // C^4
/* Having solution, check size > 2 */
/* --^-- Line.X --^-- Line.Y --^-- */
vector<Line> halfPlaneInter(vector<Line> arr) {
  sort(ALL(arr), [&](Line a, Line b) -> int {
  if (cmp(a.Y - a.X, b.Y - b.X, 0) != -1)
       return cmp(a.Y - a.X, b.Y - b.X, 0);
    return ori(a.X, a.Y, b.Y) < 0;</pre>
  });
  deque<Line> dq(1, arr[0]);
```

### 8.15 RotatingSweepLine

```
void rotatingSweepLine(vector<pii> &ps) {
  int n = SZ(ps), m = 0;
  vector<int> id(n), pos(n);
  vector<pii> line(n * (n - 1));
  for (int i = 0; i < n; ++i)
  for (int j = 0; j < n; ++j)</pre>
       if (i != j) line[m++] = pii(i, j);
  sort(ALL(line), [&](pii a, pii b) {
    return cmp(ps[a.Y] - ps[a.X], ps[b.Y] - ps[b.X]);
  }); // cmp(): polar angle compare
  iota(ALL(id), 0);
  sort(ALL(id), [&](int a, int b) {
  if (ps[a].Y != ps[b].Y) return ps[a].Y < ps[b].Y;</pre>
    return ps[a] < ps[b];</pre>
  }); // initial order, since (1, 0) is the smallest
  for (int i = 0; i < n; ++i) pos[id[i]] = i;</pre>
  for (int i = 0; i < m; ++i) {</pre>
    auto l = line[i];
     // do somethina
     tie(pos[l.X], pos[l.Y], id[pos[l.X]], id[pos[l.Y
         ]]) = make_tuple(pos[l.Y], pos[l.X], l.Y, l.X);
}
```

# 8.16 Minimum Enclosing Circle\*

### 8.17 Intersection of two circles\*

### 8.18 Intersection of polygon and circle\*

```
// Divides into multiple triangle, and sum up
const double PI=acos(-1);
double _area(pdd pa, pdd pb, double r){
  if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
  if(abs(pb)<eps) return 0;</pre>
  double S, h, theta;
  double a=abs(pb),b=abs(pa),c=abs(pb-pa);
  double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
  double cosC = dot(pa,pb) / a / b, C = acos(cosC);
  if(a > r){
    S = (C/2)*r*r;
    h = a*b*sin(C)/c;
    if (h < r && B
         < PI/2) S -= (acos(h/r)*r*r - h*sqrt(r*r-h*h));
  else if(b > r){
  theta = PI - B - asin(sin(B)/r*a);
    S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
  else S = .5*sin(C)*a*b;
  return S;
double area_poly_circle(const
     vector<pdd>> poly,const pdd &0,const double r){
  double S=0;
  for(int i=0;i<SZ(poly);++i)</pre>
    S+=_area(poly[i]-0,poly[(i+1)%SZ(poly
        )]-0,r)*ori(0,poly[i],poly[(i+1)%SZ(poly)]);
  return fabs(S);
```

### 8.19 Intersection of line and circle\*

### 8.20 Tangent line of two circles

```
vector<Line
     > go( const Cir& c1 , const Cir& c2 , int sign1 ){
   // sign1 = 1 for outer tang, -1 for inter tang
  vector < Line > ret;
  double d_sq = abs2(c1.0 - c2.0);
  if (sign(d_sq) == 0) return ret;
  double d = sqrt(d_sq);
  pdd v = (c2.0 - c1.0) / d;

double c = (c1.R - sign1 * c2.R) / d;
  if (c * c > 1) return ret;
  double h = sqrt(max(0.0, 1.0 - c * c));
  for (int sign2 = 1; sign2 >= -1; sign2 -= 2) {
  pdd n = pdd(v.X * c - sign2 * h * v.Y,
       v.Y * c + sign2 * h * v.X);
     pdd p1 = c1.0 + n * c1.R;
     pdd p2 = c2.0 + n * (c2.R * sign1);
     if (sign(p1.X - p2.X) == 0 and
         sign(p1.Y - p2.Y) == 0)
       p2 = p1 + perp(c2.0 - c1.0);
    ret.pb(Line(p1, p2));
  return ret;
}
```

### 8.21 CircleCover\*

```
bool operator < (const Teve &a)const
    {return ang < a.ang;}
  }eve[N * 2];
  // strict: x = 0, otherwise x = -1
  bool disjuct(Cir &a, Cir &b, int x)
  {return sign(abs(a.0 - b.0) - a.R - b.R) > x;}
  bool contain(Cir &a, Cir &b, int x)
  {return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
  bool contain(int i, int j) {
    /* c[j] is non-strictly in c[i]. */
    return (sign
        (c[i].R - c[j].R) > 0 \mid | (sign(c[i].R - c[j].
        R) == 0 && i < j)) && contain(c[i], c[j], -1);
  void solve(){
    fill_n(Area, C + 2, 0);
    for(int i = 0; i < C; ++i)
      for(int j = 0; j < C; ++j)</pre>
        overlap[i][j] = contain(i, j);
    for(int i = 0; i < C; ++i)</pre>
      for(int j = 0; j < C; ++j)</pre>
        g[i][j] = !(overlap[i][j] || overlap[j][i] ||
            disjuct(c[i], c[j], -1));
    for(int i = 0; i < C; ++i){</pre>
      int E = 0, cnt = 1;
      for(int j = 0; j < C; ++j)</pre>
        if(j != i && overlap[j][i])
          ++cnt;
      for(int j = 0; j < C; ++j)</pre>
        if(i != j && g[i][j]) {
          pdd aa, bb;
          CCinter(c[i], c[j], aa, bb);
          double A =
                atan2(aa.Y - c[i].0.Y, aa.X - c[i].0.X);
          double B =
                atan2(bb.Y - c[i].O.Y, bb.X - c[i].O.X);
          eve[E++] = Teve
               (bb, B, 1), eve[E++] = Teve(aa, A, -1);
          if(B > A) ++cnt;
      if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
      else{
        sort(eve, eve + E);
        eve[E] = eve[0];
        for(int j = 0; j < E; ++j){</pre>
          cnt += eve[j].add;
          Area[cnt
               ] += cross(eve[j].p, eve[j + 1].p) * .5;
          double theta = eve[j + 1].ang - eve[j].ang;
          if (theta < 0) theta += 2. * pi;
          Area[cnt] += (theta
                - sin(theta)) * c[i].R * c[i].R * .5;
      }
    }
 }
};
```

### 8.22 3Dpoint\*

```
struct Point {
  double x, y, z;
  Point(double _x = 0, double _y = 0, double _z = 0): x(_x), y(_y), z(_z){}
  Point(pdd p) { x = p.X, y = p.Y, z = abs2(p); }
Point operator - (Point p1, Point p2)
{ return
     Point(p1.x - p2.x, p1.y - p2.y, p1.z - p2.z); }
Point operator+(Point p1, Point p2)
{ return
     Point(p1.x + p2.x, p1.y + p2.y, p1.z + p2.z); }
Point operator*(Point p1, double v)
{ return Point(p1.x * v, p1.y * v, p1.z * v); }
Point operator/(Point p1, double v)
{ return Point(p1.x / v, p1.y / v, p1.z / v); }
Point cross(Point p1, Point p2)
{ return Point(p1.y * p2.z - p1.z * p2.y, p1.z
    * p2.x - p1.x * p2.z, p1.x * p2.y - p1.y * p2.x); }
double dot(Point p1, Point p2)
{ return p1.x * p2.x + p1.y * p2.y + p1.z * p2.z; }
double abs(Point a)
{ return sqrt(dot(a, a));
Point cross3(Point a, Point b, Point c)
{ return cross(b - a, c - a); }
double area(Point a, Point b, Point c)
```

```
{ return abs(cross3(a, b, c)); }
double volume(Point a, Point b, Point c, Point d)
{ return dot(cross3(a, b, c), d - a); }
//Azimuthal
     angle (longitude) to x-axis in interval [-pi, pi]
double phi(Point p) { return atan2(p.y, p.x); }
//Zenith
     angle (latitude) to the z-axis in interval [0, pi]
double theta(Point p)
    { return atan2(sqrt(p.x * p.x + p.y * p.y), p.z); }
Point masscenter(Point a, Point b, Point c, Point d)
{ return (a + b + c + d) / 4; }
pdd proj(Point a, Point b, Point c, Point u) {
// proj. u to the plane of a, b, and c
  Point e1 = b - a;
  Point e2 = c - a;
  e1 = e1 / abs(e1);
  e2 = e2 - e1 * dot(e2, e1);
  e2 = e2 / abs(e2);
  Point p = u - a;
  return pdd(dot(p, e1), dot(p, e2));
Point
     rotate_around(Point p, double angle, Point axis) {
  double s = sin(angle), c = cos(angle);
  Point u = axis / abs(axis);
  return u
       * dot(u, p) * (1 - c) + p * c + cross(u, p) * s;
```

### 8.23 Convexhull3D\*

```
struct convex_hull_3D {
struct Face {
  int a, b, c;
  Face(int ta, int tb, int tc): a(ta), b(tb), c(tc) {}
}; // return the faces with pt indexes
vector < Face > res;
vector<Point> P;
convex_hull_3D(const vector<Point> &_P): res(), P(_P) {
// all points coplanar case will WA, O(n^2)
  int n = SZ(P);
  if (n <= 2) return; // be careful about edge case</pre>
  // ensure first 4 points are not coplanar
  swap(P[1], *find_if(ALL(P), [&](
  auto p) { return sign(abs2(P[0] - p)) != 0; }));
swap(P[2], *find_if(ALL(P), [&](auto p) { return
       sign(abs2(cross3(p, P[0], P[1]))) != 0; }));
  swap(P[3], *find_if(ALL(P), [&](auto p) { return
       sign(volume(P[0], P[1], P[2], p)) != 0; }));
  vector<vector<int>> flag(n, vector<int>(n));
  res.emplace_back(0, 1, 2); res.emplace_back(2, 1, 0); for (int i = 3; i < n; ++i) {
    vector < Face > next;
    for (auto f : res) {
      int d
           = sign(volume(P[f.a], P[f.b], P[f.c], P[i]));
      if (d <= 0) next.pb(f);</pre>
      int ff = (d > 0) - (d < 0);</pre>
      flag[f.a][
           f.b] = flag[f.b][f.c] = flag[f.c][f.a] = ff;
    for (auto f : res) {
      auto F = [&](int x, int y) {
        if (flag[x][y] > 0 && flag[y][x] <= 0)
           next.emplace_back(x, y, i);
      F(f.a, f.b); F(f.b, f.c); F(f.c, f.a);
    }
    res = next;
  }
bool same(Face s, Face t) {
  if (sign(volume
       (P[s.a], P[s.b], P[s.c], P[t.a])) != 0) return 0;
  if (sign(volume
      (P[s.a], P[s.b], P[s.c], P[t.b])) != 0) return 0;
  if (sign(volume
      (P[s.a], P[s.b], P[s.c], P[t.c])) != 0) return 0;
  return 1;
int polygon_face_num() {
  int ans = 0:
  for (int i = 0; i < SZ(res); ++i)</pre>
    ans += none_of(res.begin(), res.begin()
         + i, [&](Face g) { return same(res[i], g); });
```

```
return ans;
double get_volume() {
  double ans = 0;
  for (auto f : res)
     ans +=
          volume(Point(0, 0, 0), P[f.a], P[f.b], P[f.c]);
  return fabs(ans / 6);
double get_dis(Point p, Face f) {
  Point p1 = P[f.a], p2 = P[f.b], p3 = P[f.c];

double a = (p2.y - p1.y)

* (p3.z - p1.z) - (p2.z - p1.z) * (p3.y - p1.y);
  double b = (p2.z - p1.z)
         * (p3.x - p1.x) - (p2.x - p1.x) * (p3.z - p1.z);
  double c = (p2.x - p1.x)
          * (p3.y - p1.y) - (p2.y - p1.y) * (p3.x - p1.x);
  double d = 0 - (a * p1.x + b * p1.y + c * p1.z);
return fabs(a * p.x + b *
       p.y + c * p.z + d) / sqrt(a * a + b * b + c * c);
// n^2 delaunay: facets with negative z normal of
// convexhull of (x, y, x^2 + y^2), use a pseudo-point // (0, 0, inf) to avoid degenerate case
```

### 8.24 DelaunayTriangulation\*

```
/* Delaunay Triangulation:
Given a sets of points on 2D plane, find a
triangulation such that no points will strictly
inside circumcircle of any triangle.
find : return a triangle contain given point
add_point : add a point into triangulation
A Triangle is in triangulation iff. its has_chd is 0.
Region of triangle u: iterate each u.edge[i].tri,
each points are u.p[(i+1)%3], u.p[(i+2)%3]
Voronoi diagram: for each triangle in triangulation,
the bisector of all its edges will split the region.
nearest point will belong to the triangle containing it
     ll inf = MAXC * MAXC * 100; // lower_bound unknown
struct Tri;
struct Edge {
  Tri* tri; int side;
  Edge(): tri(0), side(0){}
Edge(Tri* _tri, int _side): tri(_tri), side(_side){}
struct Tri {
  pll p[3];
  Edge edge[3];
  Tri* chd[3];
  Tri() {}
  Tri(const pll& p0, const pll& p1, const pll& p2) {
  p[0] = p0; p[1] = p1; p[2] = p2;
    chd[0] = chd[1] = chd[2] = 0;
  bool has_chd() const { return chd[0] != 0; }
  int num_chd() const {
  return !!chd[0] + !!chd[1] + !!chd[2];
  bool contains(pll const& q) const {
    for (int i = 0; i < 3; ++i)</pre>
      if (ori(p[i], p[(i + 1) % 3], q) < 0)</pre>
        return 0;
    return 1;
} pool[N * 10], *tris;
void edge(Edge a, Edge b) {
  if(a.tri) a.tri->edge[a.side] = b;
  if(b.tri) b.tri->edge[b.side] = a;
struct Trig { // Triangulation
  Trig() {
         = // Tri should at least contain all points
      new(tris++) Tri(pll(-inf, -inf),
            pll(inf + inf, -inf), pll(-inf, inf + inf));
  Tri* find(pll p) { return find(the_root, p); }
  void add_point(const
       pll &p) { add_point(find(the_root, p), p); }
  Tri* the_root;
  static Tri* find(Tri* root, const pll &p) {
    while (1) {
```

```
if (!root->has chd())
        return root;
      for (int i = 0; i < 3 && root->chd[i]; ++i)
        if (root->chd[i]->contains(p)) {
          root = root->chd[i];
          break:
    assert(0); // "point not found"
  void add_point(Tri* root, pll const& p) {
    Tri* t[3];
     /* split it into three triangles */
    for (int i = 0; i < 3; ++i)</pre>
      t[i] = new(tris
           ++) Tri(root->p[i], root->p[(i + 1) % 3], p);
    for (int i = 0; i < 3; ++i)</pre>
      edge(Edge(t[i], 0), Edge(t[(i + 1) % 3], 1));
    for (int i = 0; i < 3; ++i)</pre>
      edge(Edge(t[i], 2), root->edge[(i + 2) % 3]);
    for (int i = 0; i < 3; ++i)</pre>
      root->chd[i] = t[i];
    for (int i = 0; i < 3; ++i)</pre>
      flip(t[i], 2);
  void flip(Tri* tri, int pi) {
    Tri* trj = tri->edge[pi].tri;
    int pj = tri->edge[pi].side;
    if (!trj) return;
    if (!in_cc(tri->p
        [0], tri->p[1], tri->p[2], trj->p[pj])) return;
    /* flip edge between tri,trj */
    Tri* trk = new(tris++) Tri
        (tri->p[(pi + 1) \% 3], trj->p[pj], tri->p[pi]);
    Tri* trl = new(tris++) Tri
        (trj->p[(pj + 1) % 3], tri->p[pi], trj->p[pj]);
    edge(Edge(trk, 0), Edge(trl, 0));
    edge(Edge(trk, 1), tri->edge[(pi + 2) % 3]);
    edge(Edge(trk, 2), trj->edge[(pj + 1) % 3]);
    edge(Edge(trl, 1), trj->edge[(pj + 2) \% 3]);
    edge(Edge(trl, 2), tri->edge[(pi + 1) % 3]);
    tri->chd
        [0] = trk; tri->chd[1] = trl; tri->chd[2] = 0;
    trj->chd
        [0] = trk; trj->chd[1] = trl; trj->chd[2] = 0;
    flip(trk, 1); flip(trk, 2);
flip(trl, 1); flip(trl, 2);
};
vector<Tri*> triang; // vector of all triangle
set<Tri*> vst;
void go(Tri* now) { // store all tri into triang
  if (vst.find(now) != vst.end())
    return:
  vst.insert(now);
  if (!now->has_chd())
    return triang.pb(now);
  for (int i = 0; i < now->num_chd(); ++i)
    go(now->chd[i]);
void build(int n, pll* ps) { // build triangulation
  tris = pool; triang.clear(); vst.clear();
  random_shuffle(ps, ps + n);
  Trig tri; // the triangulation structure
for (int i = 0; i < n; ++i)</pre>
    tri.add_point(ps[i]);
  go(tri.the_root);
```

### 8.25 Triangulation Vonoroi\*

```
vector<Line> ls[N];
pll arr[N];
Line make_line(pdd p, Line l) {
  pdd d = l.Y - l.X; d = perp(d);
  pdd m = (1.X + 1.Y) / 2;
  l = Line(m, m + d);
  if (ori(l.X, l.Y, p) < 0)
    l = Line(m + d, m);
  return l;
double calc_area(int id) {
  // use to calculate
       the area of point "strictly in the convex hull"
  vector<Line> hpi = halfPlaneInter(ls[id]);
  vector<pdd> ps;
```

```
for (int i = 0; i < SZ(hpi); ++i)
  ps.pb(intersect(hpi[i].X, hpi[i].Y, hpi[(i</pre>
           + 1) % SZ(hpi)].X, hpi[(i + 1) % SZ(hpi)].Y));
  double rt = 0;
  for (int i = 0; i < SZ(ps); ++i)</pre>
     rt += cross(ps[i], ps[(i + 1) % SZ(ps)]);
  return fabs(rt) / 2;
void solve(int n, pii *oarr) {
  map<pll, int> mp;
  for (int i = 0; i < n; ++i)</pre>
     arr[i] = pll(oarr[i].X, oarr[i].Y), mp[arr[i]] = i;
  build(n, arr); // Triangulation
  for (auto *t : triang) {
     vector<int> p;
     for (int i = 0; i < 3; ++i)</pre>
       if (mp.find(t->p[i]) != mp.end())
          p.pb(mp[t->p[i]]);
     for (int i = 0; i < SZ(p); ++i)
  for (int j = i + 1; j < SZ(p); ++j) {</pre>
         Line l(oarr[p[i]], oarr[p[j]]);
ls[p[i]].pb(make_line(oarr[p[i]], l));
          ls[p[j]].pb(make_line(oarr[p[j]], l));
  }
}
```

### 8.26 Minkowski Sum\*

```
vector<pll> Minkowski(vector<pll> A, vector<pll> B) {
  hull(A), hull(B);
  vector<pll> C(1, A[\theta] + B[\theta]), s1, s2;
  for (int i = 0; i < SZ(A); ++i)</pre>
    s1.pb(A[(i + 1) % SZ(A)] - A[i]);
  for (int i = 0; i < SZ(B); i++)</pre>
    s2.pb(B[(i + 1) % SZ(B)] - B[i]);
  for (int i = 0, j = 0; i < SZ(A) || j < SZ(B);)</pre>
    if (j >= SZ
        (B) || (i < SZ(A) && cross(s1[i], s2[j]) >= 0))
      C.pb(B[j % SZ(B)] + A[i++]);
    else
      C.pb(A[i % SZ(A)] + B[j++]);
  return hull(C), C;
```

### 9 Else

### 9.1 Cyclic Ternary Search\*

```
/* bool pred(int a, int b);
f(\theta) \sim f(n-1) is a cyclic-shift U-function
return idx s.t. pred(x, idx) is false forall x*/
int cyc_tsearch(int n, auto pred) {
  if (n == 1) return 0;
  int l = 0, r = n; bool rv = pred(1, 0);
while (r - l > 1) {
     int m = (l + r) / 2;
     if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
  return pred(l, r % n) ? l : r % n;
}
```

### Mo's Alogrithm(With modification)

```
Mo's Algorithm With modification
Block: N^{2/3}, Complexity: N^{5/3}
struct Query {
  int L, R, LBid, RBid, T;
  Query(int l, int r, int t):
    L(l), R(r), LBid(l / blk), RBid(r / blk), T(t) {}
  bool operator<(const Query &q) const {</pre>
    if (LBid != q.LBid) return LBid < q.LBid;</pre>
    if (RBid != q.RBid) return RBid < q.RBid;</pre>
    return T < b.T;</pre>
 }
};
void solve(vector<Query> query) {
  sort(ALL(query));
  int L=0, R=0, T=-1;
  for (auto q : query) {
    while (T < q.T) addTime(L, R, ++T); // TODO
    while (T > q.T) subTime(L, R, T--); // TODO
    while (R < q.R) add(arr[++R]); // TODO</pre>
```

```
while (L > q.L) add(arr[--L]); // TODO
while (R > q.R) sub(arr[R--]); // TODO
       while (L < q.L) sub(arr[L++]); // TODO</pre>
       // answer query
}
```

### 9.3 Mo's Alogrithm On Tree

```
Mo's Algorithm On Tree
Preprocess:
1) LCA
2) dfs with in[u] = dft++, out[u] = dft++
   ord[in[u]] = ord[out[u]] = u
4) bitset < MAXN > inset
struct Query {
  int L, R, LBid, lca;
  Query(int u, int v) {
     int c = LCA(u, v);
     if (c == u || c == v)
      q.lca = -1, q.L = out[c ^ u ^ v], q.R = out[c];
     else if (out[u] < in[v])</pre>
      q.lca = c, q.L = out[u], q.R = in[v];
     else
    bool operator < (const Query &q) const {</pre>
     if (LBid != q.LBid) return LBid < q.LBid;</pre>
     return R < q.R;</pre>
  }
void flip(int x) {
     if (inset[x]) sub(arr[x]); // TODO
     else add(arr[x]); // TODO
     inset[x] = ~inset[x];
void solve(vector<Query> query) {
  sort(ALL(query));
   int L = 0, R = 0;
  for (auto q : query) {
     while (R < q.R) flip(ord[++R]);</pre>
     while (L > q.L) flip(ord[--L]);
     while (R > q.R) flip(ord[R--]);
     while (L < q.L) flip(ord[L++]);</pre>
     if (~q.lca) add(arr[q.lca]);
     // answer query
     if (~q.lca) sub(arr[q.lca]);
}
```

### 9.4 Additional Mo's Algorithm Trick

- · Mo's Algorithm With Addition Only
  - Sort querys same as the normal Mo's algorithm.
  - For each query [l,r]:
  - If l/blk = r/blk, brute-force.
  - If  $l/blk \neq curL/blk$ , initialize  $curL := (l/blk+1) \cdot blk$ , curR := curL-1
  - If r > curR, increase curR
  - ullet decrease curL to fit l , and then undo after answering
- Mo's Algorithm With Offline Second Time
  - Require: Changing answer  $\equiv$  adding f([l,r],r+1).
  - Require: f([l,r],r+1) = f([1,r],r+1) f([1,l),r+1). Part1: Answer all f([1,r],r+1) first.

  - Part2: Store  $curR \to R$  for curL (reduce the space to O(N)), and then answer them by the second offline algorithm.
  - Note: You must do the above symmetrically for the left boundaries.

### 9.5 Hilbert Curve

```
ll hilbert(int n, int x, int y) {
   ll res = 0;
   for (int s = n / 2; s; s >>= 1) {
     int rx = (x \& s) > 0;
     int ry = (y & s) > 0;
res += s * 1ll * s * ((3 * rx) ^ ry);
     if (ry == 0) {
       if (rx == 1) x = s - 1 - x, y = s - 1 - y;
       swap(x, y);
     }
  }
   return res;
\} // n = 2^k
```

### 9.6 DynamicConvexTrick\*

```
// only works for integer coordinates!! maintain max
struct Line {
  mutable ll a, b, p;
  bool operator
      <(const Line &rhs) const { return a < rhs.a; }
  bool operator<(ll x) const { return p < x; }</pre>
struct DynamicHull : multiset<Line, less<>>> {
  static const ll kInf = 1e18;
  ll Div(ll a,
       ll b) { return a / b - ((a ^ b) < 0 && a % b); }
  bool isect(iterator x, iterator y) {
    if (y == end()) { x->p = kInf; return 0; }
    if (x
        ->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
    else x - p = Div(y - b - x - b, x - a - y - a);
    return x->p >= y->p;
  void addline(ll a, ll b) {
    auto z = insert({a, b, 0}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin
        () && isect(--x, y)) isect(x, y = erase(y));
    while ((y = x) != begin
        () && (--x)->p >= y->p) isect(x, erase(y));
  ll query(ll x) {
    auto l = *lower_bound(x);
    return l.a * x + l.b;
 }
};
```

### 9.7 All LCS\*

```
void all_lcs(string s, string t) { // 0-base
  vector < int > h(SZ(t));
  iota(ALL(h), 0);
  for (int a = 0; a < SZ(s); ++a) {
    int v = -1;
    for (int c = 0; c < SZ(t); ++c)
        if (s[a] == t[c] || h[c] < v)
            swap(h[c], v);
        // LCS(s[0, a], t[b, c]) =
        // c - b + 1 - sum([h[i] >= b] | i <= c)
        // h[i] might become -1 !!
  }
}</pre>
```

### 9.8 DLX\*

```
#define TRAV(i, link, start)
     for (int i = link[start]; i != start; i = link[i])
template <
    bool E> // E: Exact, NN: num of 1s, RR: num of rows
struct DLX {
  int lt[NN], rg[NN], up[NN], dn[NN
      ], rw[NN], cl[NN], bt[NN], s[NN], head, sz, ans;
  int rows, columns;
 bool vis[NN];
 bitset<RR> sol, cur; // not sure
  void remove(int c) {
    if (E) lt[rg[c]] = lt[c], rg[lt[c]] = rg[c];
    TRAV(i, dn, c) {
   if (E) {
        TRAV(j, rg, i)
          up[dn[j]]
               = up[j], dn[up[j]] = dn[j], --s[cl[j]];
      } else {
        lt[rg[i]] = lt[i], rg[lt[i]] = rg[i];
   }
  void restore(int c) {
    TRAV(i, up, c) {
      if (E) {
        TRAV(i.
                lt. i)
          ++s[cl[j]], up[dn[j]] = j, dn[up[j]] = j;
      } else
        lt[rg[i]] = rg[lt[i]] = i;
      }
    if (E) lt[rg[c]] = c, rg[lt[c]] = c;
  void init(int c) {
```

```
rows = 0, columns = c;
     for (int i = 0; i < c; ++i) {</pre>
       up[i] = dn[i] = bt[i] = i;
       lt[i] = i == 0 ? c : i - 1;
       rg[i] = i == c - 1 ? c : i + 1;
       s[i] = 0;
     rg[c] = 0, lt[c] = c - 1;
     up[c] = dn[c] = -1;
     head = c, sz = c + 1;
   void insert(const vector<int> &col) {
     if (col.empty()) return;
     int f = sz;
     for (int i = 0; i < (int)col.size(); ++i) {</pre>
       int c = col[i], v = sz++;
       dn[bt[c]] = v;
       up[v] = bt[c], bt[c] = v;
       rg[v] = (i + 1 == (int)col.size() ? f : v + 1);
       rw[v] = rows, cl[v] = c;
       ++s[c];
       if (i > 0) lt[v] = v - 1;
     ++rows, lt[f] = sz - 1;
   int h() {
     int ret = 0;
     fill_n(vis, sz, false);
     TRAV(x, rg, head) {
       if (vis[x]) continue;
       vis[x] = true, ++ret;
TRAV(i, dn, x) TRAV(j, rg, i) vis[cl[j]] = true;
     return ret;
   void dfs(int dep) {
     if (dep + (E ? 0 : h()) >= ans) return;
     if (rg[head
         ] == head) return sol = cur, ans = dep, void();
     if (dn[rg[head]] == rg[head]) return;
     int w = rg[head];
     TRAV(x, rg, head) if (s[x] < s[w]) w = x;
     if (E) remove(w);
     TRAV(i, dn, w) {
       if (!E) remove(i);
       TRAV(j, rg, i) remove(E ? cl[j] : j);
       cur.set(rw[i]), dfs(dep + 1), cur.reset(rw[i]);
       TRAV(j, lt, i) restore(E ? cl[j] : j);
       if (!E) restore(i);
     if (E) restore(w);
   int solve() {
     for (int i = 0; i < columns; ++i)</pre>
      dn[bt[i]] = i, up[i] = bt[i];
     ans = 1e9, sol.reset(), dfs(0);
     return ans;
};
```

### 9.9 Matroid Intersection

Start from  $S = \emptyset$ . In each iteration, let

•  $Y_1 = \{x \notin S \mid S \cup \{x\} \in I_1\}$ 

•  $Y_2 = \{x \notin S \mid S \cup \{x\} \in I_2\}$ 

If there exists  $x \in Y_1 \cap Y_2$ , insert x into S. Otherwise for each  $x \in S, y \notin S$ , create edges

•  $x \to y \text{ if } S - \{x\} \cup \{y\} \in I_1.$ 

•  $y \to x$  if  $S - \{x\} \cup \{y\} \in I_2$ .

Find a *shortest* path (with BFS) starting from a vertex in  $Y_1$  and ending at a vertex in  $Y_2$  which doesn't pass through any other vertices in  $Y_2$ , and alternate the path. The size of S will be incremented by 1 in each iteration. For the weighted case, assign weight w(x) to vertex x if  $x \in S$  and -w(x) if  $x \notin S$ . Find the path with the minimum number of edges among all minimum length paths and alternate it.

### 9.10 AdaptiveSimpson

```
template < typename Func, typename d = double >
struct Simpson {
  using pdd = pair < d, d >;
  Func f;
  pdd mix(pdd l, pdd r, optional < d > fm = {}) {
    d h = (r.X - l.X) / 2, v = fm.value_or(f(l.X + h));
    return {v, h / 3 * (l.Y + 4 * v + r.Y)};
  }
  d eval(pdd l, pdd r, d fm, d eps) {
```

```
pdd m((l.X + r.X) / 2, fm);
    d s = mix(l, r, fm).second;
    auto [flm, sl] = mix(l, m);
    auto [fmr, sr] = mix(m, r);
d delta = sl + sr - s;
    if (abs(delta
         ) <= 15 * eps) return sl + sr + delta / 15;
    return eval(l, m, flm, eps / 2) +
      eval(m, r, fmr, eps / 2);
  d eval(d l, d r, d eps) {
    return eval
         ({l, f(l)}, {r, f(r)}, f((l + r) / 2), eps);
  d eval2(d l, d r, d eps, int k = 997) {
    d h = (r - l) / k, s = 0;
for (int i = 0; i < k; ++i, l += h)
      s += eval(l, l + h, eps / k);
 }
};
template < typename Func >
Simpson<Func> make_simpson(Func f) { return {f}; }
```

### 

### 9.11 Simulated Annealing

```
double factor = 100000;
const int base = 1e9; // remember to run ~ 10 times
for (int it = 1; it <= 1000000; ++it) {
    // ans:
        answer, nw: current value, rnd(): mt19937 rnd()
    if (exp(-(nw - ans
        ) / factor) >= (double)(rnd() % base) / base)
        ans = nw;
    factor *= 0.99995;
}
```

### 9.12 Tree Hash\*

```
ull seed;
ull shift(ull x) {
    x ^= x << 13;
    x ^= x >> 7;
    x ^= x << 17;
    return x;
}
ull dfs(int u, int f) {
    ull sum = seed;
    for (int i : G[u])
        if (i != f)
            sum += shift(dfs(i, u));
    return sum;
}</pre>
```

# 9.13 Binary Search On Fraction

```
struct Q {
  ll p, q;
  Q go(Q b, ll d) { return {p + b.p*d, q + b.q*d}; }
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(ll N) {
  Q lo{0, 1}, hi{1, 0};
  if (pred(lo)) return lo;
  assert(pred(hi));
  bool dir = 1, L = 1, H = 1;
  for (; L || H; dir = !dir) {
    ll len = 0, step = 1;
    for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
      if (Q mid = hi.go(lo, len + step);
         mid.p > N || mid.q > N || dir ^ pred(mid))
      else len += step;
    swap(lo, hi = hi.go(lo, len));
    (dir ? L : H) = !!len;
  return dir ? hi : lo;
```

# 10 Python

### 10.1 Misc