CSCI432 HW5

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1 Prove that the Frechet distance is a distance metric

Prove that the Frechet distance is a distance metric.

To be a distance metric, you must satisfy 4 requirements:
Let X = discrete space
A metric is a function d: X x X $\rightarrow \mathbb{R}$ such that:

1. d(x,y) = d(y,x)

2. d(x,y) = 0 $\Leftrightarrow$ x = y

3. d(x,y) + d(y,z) $\geq$ d(x,z)

4. d(x,y) $\geq$ 0

So,

2 Recurrence Relations

2.a T(n) = 2T(n/4) + n^2

Master's Theorem:
T(n) = aT($\frac{n}{b}$) + f(n)
Where, a $\geq$ 1, b > 1, f(n) is asymptotically positive

T(n) = 2T(n/4) + n^2
a = 2, b = 4, f(n) = n^2
n$^{log_b a}$ $\Rightarrow$ n$^{log_4 2}$ $\Rightarrow$ n$^{1/2}$
Case 3:
if f(n) is $\Omega(n^{log_b a + \epsilon})$ for some $\epsilon > 0$ and if f(n/b) $\leq$ f(n), then T(n) = $\theta(f(n))$
Therefore, T(n) = $\theta(n^2)$

2.b   T(n) = 4T(n/2) + n

Master's Theorem:
T(n) = aT($\frac{n}{b}$) + f(n)
Where, a $\geq$ 1, b > 1, f(n) is asymptotically positive

T(n) = 4T(n/2) + n
a = 4, b = 2, f(n) = n
$n^{log_b a} \Rightarrow n^{log_2 4} \Rightarrow n^2$
Case 1:
if f(n) = O($n^{log_b a - \epsilon}$) for some $\epsilon > 0$ then T(n) = $\Theta(n^2)$.
Therefore, T(n) = $\Theta(n^2)$

2.c   T(n) = 3T(2n/3) + 4n

Master's Theorem:
T(n) = aT($\frac{n}{b}$) + f(n)
Where, a $\geq$ 1, b > 1, f(n) is asymptotically positive
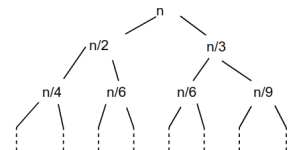
T(n) = 3T(2n/3) + 4n
a = 3, b = 2/3, f(n) = 4n
$n^{log_b a} \Rightarrow n^{log_{3/2} 3} \Rightarrow n^{2.71}$
Case 1:
if f(n) = O($n^{log_b a - \epsilon}$) for some $\epsilon > 0$ then T(n) = $\Theta(n^2)$.
Therefore, T(n) = $\Theta(n^2)$

2.d   T(n) = T(n/2) + T(n/3)

Our longest path in this tree is the rightmost path, following a sequence: $log_3 n$.
So our initial guess is for this recurrence is $O(n log n)$.

2.e   2T(n/2) + O(log n)

3   Climbing Stairs Problem