

# CSCI338 HW4

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## Question 1

Let  $\beta$  be the set of all infinite sequences over  $\{a,b\}$ . Show that  $\beta$  is uncountable, using a proof by diagonalization.

Assume that  $\beta$  is countable, then we can list them as  $\beta_1, \beta_2$

i	f(i)		
1	b <sub>11</sub>	b <sub>12</sub>	b <sub>13</sub>
2	b <sub>21</sub>	b <sub>22</sub>	b <sub>23</sub>
3	b <sub>31</sub>	b <sub>32</sub>	b <sub>33</sub>

Now construct an  $x$  such that the  $i^{th}$  bit of  $x$  is  $\neq$  the  $i^{th}$  bit of  $\beta$

i	f(i)		
1	a	a	a...
2	b	a	b...
3	b	b	b...

Lets define  $x$  as  $x = (b, b, a...)$  thus making  $x \neq \beta$  for any  $i^{th}$  sequence in the  $i^{th}$  bit. Therefore  $\beta$  is uncountable.

## Question 2

Let  $T = \{(i,j,k) \mid i,j,k \in \mathbb{N}\}$ . Show that  $T$  is countable.

## Question 3

Show that  $INFINITE_{PDA}$  is decidable

To decide  $INFINITE_{PDA}$  convert PDA into equivalent CFG, let  $P$  be the pumping length. Construct regular language  $R$  that accepts strings longer than  $P$ . Intersect of CFL and regular languages is a CFL. Test this intersection for emptiness, accept if  $L$  of the intersection is empty, reject otherwise.

## Question 4

## Question 5

Show that  $EQ_{CFG}$  is undecidable.

Assume that  $EQ_{CFG}$  is decidable

$EQ_{CFG} = \{ \langle G, H \rangle \mid G, H \text{ are CFG's and } L(G) = L(H) \}$

Reduce  $ALL_{CFG}$  to  $EQ_{CFG}$  Such That:

$ALL_{CFG} = \{ \langle G \rangle \mid G \text{ is a CFG's and } L(G) = \Sigma^* \}$

Let R be a decider for  $EQ_{CFG}$  and construct TM S to decide  $ALL_{CFG}$ .

Construct CFG T such that  $L(T) = \Sigma^*$

1. Run R on input  $\langle G, T_0 \rangle$
2. If R accepts, accept.
3. If R rejects, reject

R decides if  $L(G) = L(T)$ . S decides  $ALL_{CFG}$  but it is undecidable, therefore  $EQ_{CFG}$  must also be undecidable

## Question 6

Show that  $EQ_{CFG}$  is co-Turing-recognizable.

A language is co-turning recognizable if and only if its complement is a turning-recognizable language.

Convert G and H into Chomsky normal form. Begin iterating through the strings in  $\Sigma^*$ . If both G or H can generate or not generate the string a TM will continue on the iteration of strings, however, if one CFG accepts a string and the other does not, that means that the CFG's are not equivalent and the TM accepts. Therefore  $EQ_{CFG}$  is co-turning recognizable.

## Question 7