



Figure 1: Example of faces

With Figure 1 as an example of faces.

For a face f of a planar drawing \tilde{G} of a planar graph G :

Definition List

- The set of vertices and edges on the perimeter of f form the *boundary* of f (also said to be *incident* to f)
- Faces f and f' are said to be *adjacent* if they are incident to a common edge
- with a face f and a vertex v_0 incident to f , a *boundary walk* $w(f)$ of f is a sequence $v_0e_1v_1e_2v_2\dots e_kv_k$ (where $v_k = v_0$) of all the vertices and edges of the boundary of f . These are listed in the order that they are encountered as one moves around the whole perimeter of f .

Note: Imagine this like you're trailing your hand along one side of the fence. You can't just stop when you've re-encountered a vertex, but you have to be in the position that you'd be on the "starting side" to the fence.

- the degree of f is the length of $w(f)$. (not doubling up on the last vertex)

Note that edges e in the boundary of f contribute 1 to $\deg(f)$ if e is not a bridge and 2 if e is a bridge. This is because if C is a cycle then in any planar drawing, each edge of C is incident to 2 faces.

If G is a tree, then any planar drawing of G has 1 face (the outer face), which has degree $2(p - 1)$ where $p = |V(G)|$ since you'll be counting every edge twice.

Lemma Let \tilde{G} be a planar drawing of a connected planar graph. Let $F(\tilde{G})$ denote the set of faces of \tilde{G} then

$$\sum_{f \in F(\tilde{G})} \deg(f) = 2|E(G)|$$

(since every edge is counted twice - either separating two faces from each other, or a bridge)

Proof Each edge e of G contributes 1 to the 2 distinct faces in the sum to the Left Hand Side. If e is not a bridge, and 2 to the same face if e is a bridge. So we get $2|E(G)|$.

(sometimes called "Face Shake Lemma")

Theorem (Euler's Formula) Let \tilde{G} be the planar drawing of a connected planar graph G . Let $p = |V(G)|$, $q = |E(G)|$, and $s = |F(\tilde{G})|$. Then

$$p - q + s = 2$$

(basically an invariant of the plane. Special number that goes along with \mathbb{R}^2)

Proof We fix p and use induction on q

Base Case $q = p - 1$ smallest it can be is minimally connected, and therefore a tree by our earlier lemma. Then $s = 1$. So $p - q + s = p - (p - 1) + 1 = 2$

Inductive Hypothesis Assume that $q \geq p$ and that for all $q' - q$, and all planar drawings of connected planar graphs G' with p vertices and q' edges and s' faces, $p - q' + s' = 2$.

Let \tilde{G} have p vertices, q edges and s faces.

Then G contains a cycle, hence we can choose an edge e of G that is not a bridge of G .

Then e is incident to 2 distinct faces of \tilde{G}

Then $G - e$ is connected (because e is not a bridge), and the drawing \tilde{G} has p vertices, $q - 1$ edges, and $s - 1$ faces.

Hence by the inductive hypothesis, $\tilde{G} - e$ has $p - (q - 1) + (s - 1) = 2$ and therefore $p - q + s = 2$ since when e was removed we removed one face.

Therefore by induction the theorem holds \square