

LC-3 Assembly language

Lecture Topics

LC-3 assembly language

The assembly process

Example

Lecture materials

Textbook Ch. 7

Homework

Machine problem

MP2 due March 2, 2011 at 5pm submitted electronically.

Announcements

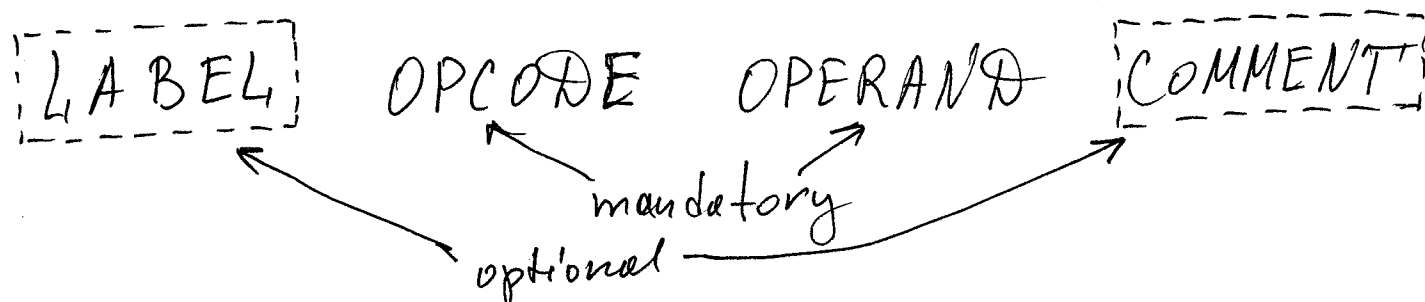
LC-3 assembly language

- LC-3 assembly language is a low-level language specifically invented for LC-3 computer
 - It is machine-specific; different processor have different assembly languages
 - Each assembly language instruction corresponds to a single ISA instruction
- Assembly language makes it much easier to program by using human-readable language instead of binary words, while still providing the programmer with the fine-grain control over the instructions and data
- Assembly language allows to refer to *instructions*, *memory locations*, and *values* **symbolically**.

Program from last lecture, but now in LC-3 Assembly language

```
; Program to count occurrences of a character in a file
;
; Initialization
;
    .ORIG        x3000
    AND R2, R2, #0 ; R2 is counter, initially 0
    LD  R3, PTR    ; R3 is pointer to characters
    GETC          ; R0 gets character input
    LDR R1, R3, #0 ; R1 gets first character
;
; Test character for end of file
;
TEST    ADD     R4, R1, #-4 ; Test for EOT (ASCII x04)
        BRz OUTPUT        ; If done, prepare the output
;
; Test character for match. If a match, increment count.
;
    NOT R1, R1
    ADD R1, R1, R0 ; If match, R1 = xFFFF
    NOT R1, R1     ; If match, R1 = x0000
    BRnp GETCHAR   ; If no match, do not increment
    ADD R2, R2, #1
;
; Get next character from file.
;
GETCHAR ADD     R3, R3, #1 ; Point to next character.
        LDR R1, R3, #0    ; R1 gets next char to test
        BRnzp TEST
;
; Output the count.
;
OUTPUT LD R0, ASCII ; Load the ASCII template
        ADD R0, R0, R2 ; Covert binary count to ASCII
        OUT          ; ASCII code in R0 is displayed.
        HALT         ; Halt machine
;
; Storage for pointer and ASCII template
;
ASCII .FILL x0030
PTR   .FILL x3016
.END
```

- Few observations about the program
 - semicolon (;) is used to indicate a comment
 - some lines contain ~~a~~ familiar instructions, e.g., LD, ADD. These are assembly language instructions that will be executed by the computer
 - some lines contain yet unfamiliar instructions, e.g., .FILL, .END. These are pseudo-ops, messages to the translating program to help with the translation to ISA instructions.
- An instruction in LC-3's assembly language consists from 4 parts



→ OPCODEs

- are reserved symbols that correspond to LC-3 instructions (it is easier to remember a name than a sequence of 0's and 1's!)
- 15 opcodes in the LC-3's ISA.

→ OPERANDs

- each instruction has a number of instruction-specific operands, in specific order
- operands are one of the following:
 - registers: R_n where n is the register number
 - numbers: # (decimal), x (hex)
 - labels: symbolic names of memory locations
- operands are separated by comma (,)

→ LABELS

- symbolic names used to identify memory locations that are explicitly referred to in the program.
- consist of one-to-20 alphanumeric characters
- used for 2 reasons
 - the memory location contains the target of a branch instruction
 - the memory location contains a value that is loaded or stored by LD-ST family of instructions.
 - value stored in the memory location explicitly referenced ~~as~~ with the label will be used in the LD/STx instruction.

→ COMMENTS

- intended only for human consumption
- their purpose is to make the program more readable by a human
- ignored by the assembler - more on this later
- should provide more insight, not just restate the obvious.
- text after a semicolon (;) is considered to be a comment.

→ Pseudo ops

- they do not refer to operations that will be actually executed by the computer
- they really are assembler directives that are used during the assembler language translation into the binary form - more to follow
- they do not appear in the final binary program.
- start with dot (.)

→ Pseudo-ops

→ LC-3 assembly language contains ⁵ pseudo-ops

→ **.ORIG** specifies where in memory the program should be placed

→ format .ORIG address

→ **.END** indicates the end of program. It has nothing to do with HALT TRAP, it just says that any characters after it should be ignored

→ format .END

→ **.FILL** indicates that a location in memory will be filled with a value specified in the operand

→ format .FILL n

→ **.BLKW** indicates that some number of memory locations will be reserved (block of words)

→ format .BLKW n

→ **.STRINGZ** indicates that (n+1) memory locations will be reserved and ~~allocated~~ initialized with the characters provided in the pseudo-instruction.

→ format .STRINGZ "text"

→ example: .ORIG x3010
.STRINGZ "test"

⇓

x3010 't'
x3011 'e'
x3012 's'
x3013 't'
x3014 0

← null-terminated string.

→ Trap Code pseudo-instructions

{

 → HALT (TRAP x25)

 → IN (TRAP x23)

 → OUT (TRAP x21)

 }

 → GETC (TRAP x20) - read one char from Keyboard

 → PUTS (TRAP x22) - write null-terminated string to console

you should know these by now

→ Example #1

Count occurrence of a character in a file.
 Character to be input from the Keyboard.
 Result to be displayed on the monitor.
 Same example as in last lecture.

→ Example #2.

Given a sequence of 100 values, stored in 2's complement format starting at x4000, find maximum value. R2 to contain the max value, R3 to contain the address of the value.

→ Assembly process

→ > lc3as countchar.asm

STARTING PASS 1

0 errors found in first pass.

STARTING PASS 2

0 errors found in second pass

> ls

countchar.asm countchar.obj

→ LC3as translated ~~countchar.asm~~ program stored in countchar.asm into a binary code stored in countchar.obj which now can be executed by LC-3 simulator / computer.

→ LC3as is ~~call~~ an LC-3 assembler, a program that translates an assembly language program into a machine language program.

→ 2-pass assembly process

→ 1: creating the symbol table

→ symbol table is a correspondence of symbolic names and their 16-bit memory addresses

~~symbol address~~
~~PTR x3013~~

for example #1

symbol	address
TEST	x3004
GETCHAR	x300B
OUTPUT	x300E
ASCII#	x3012
PTR	x3013

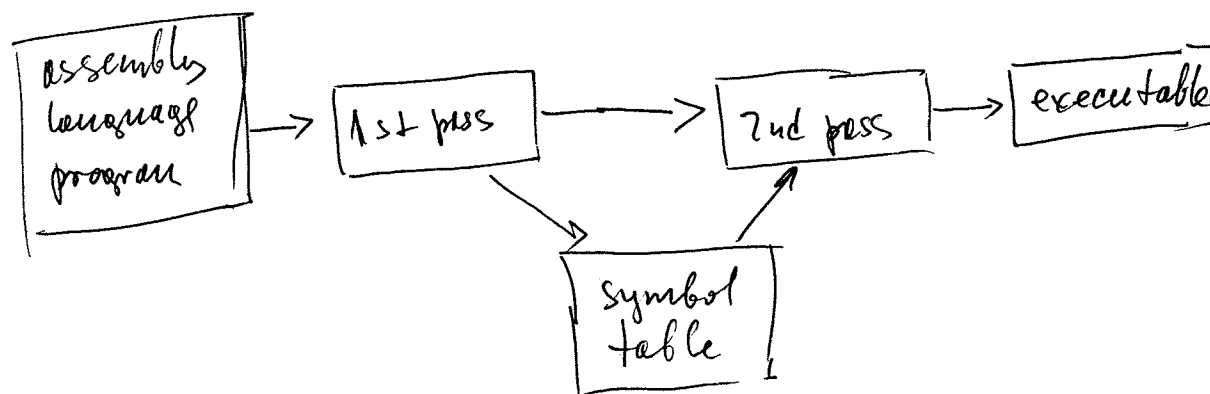
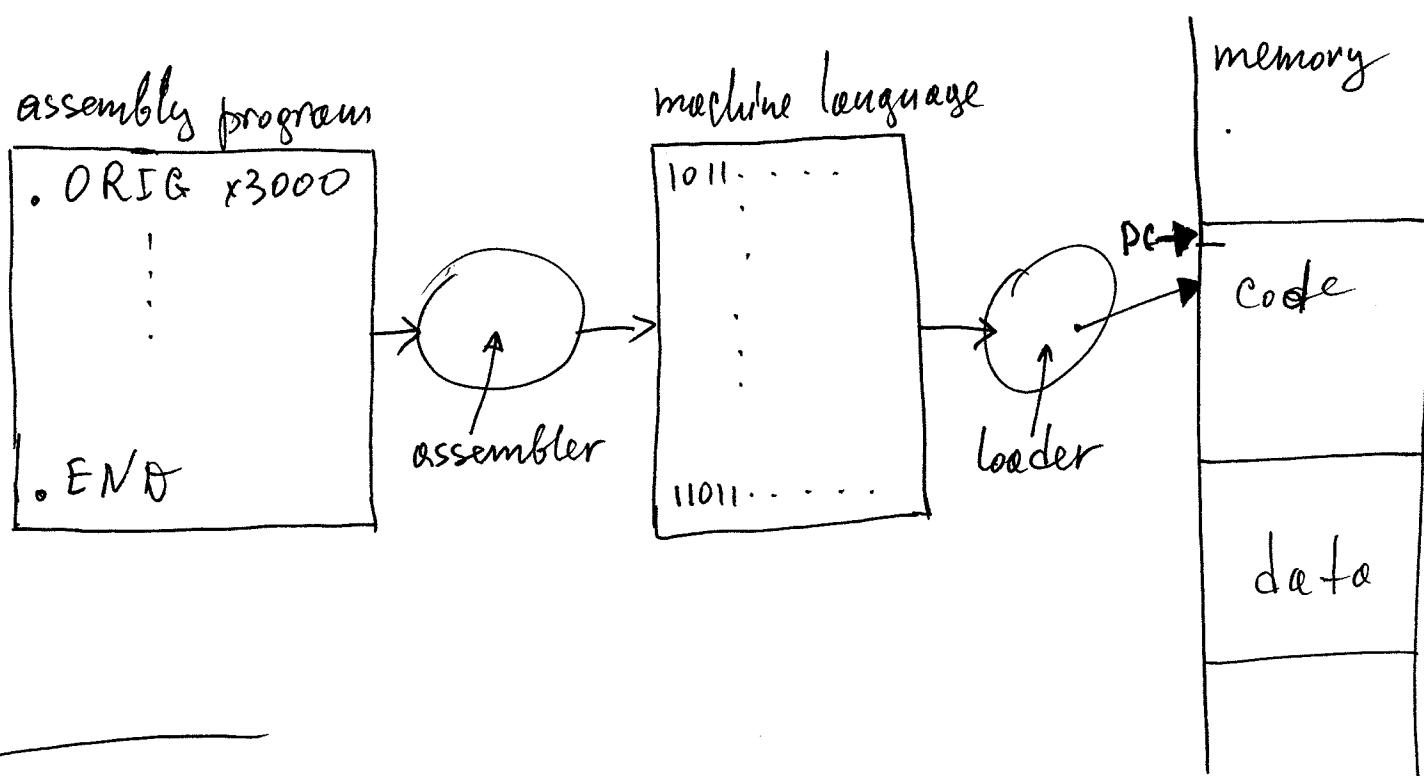
→ each time a new symbol is found, ~~its~~ the PC value for the line with the symbol is noted.

→ 2: generating machine language program

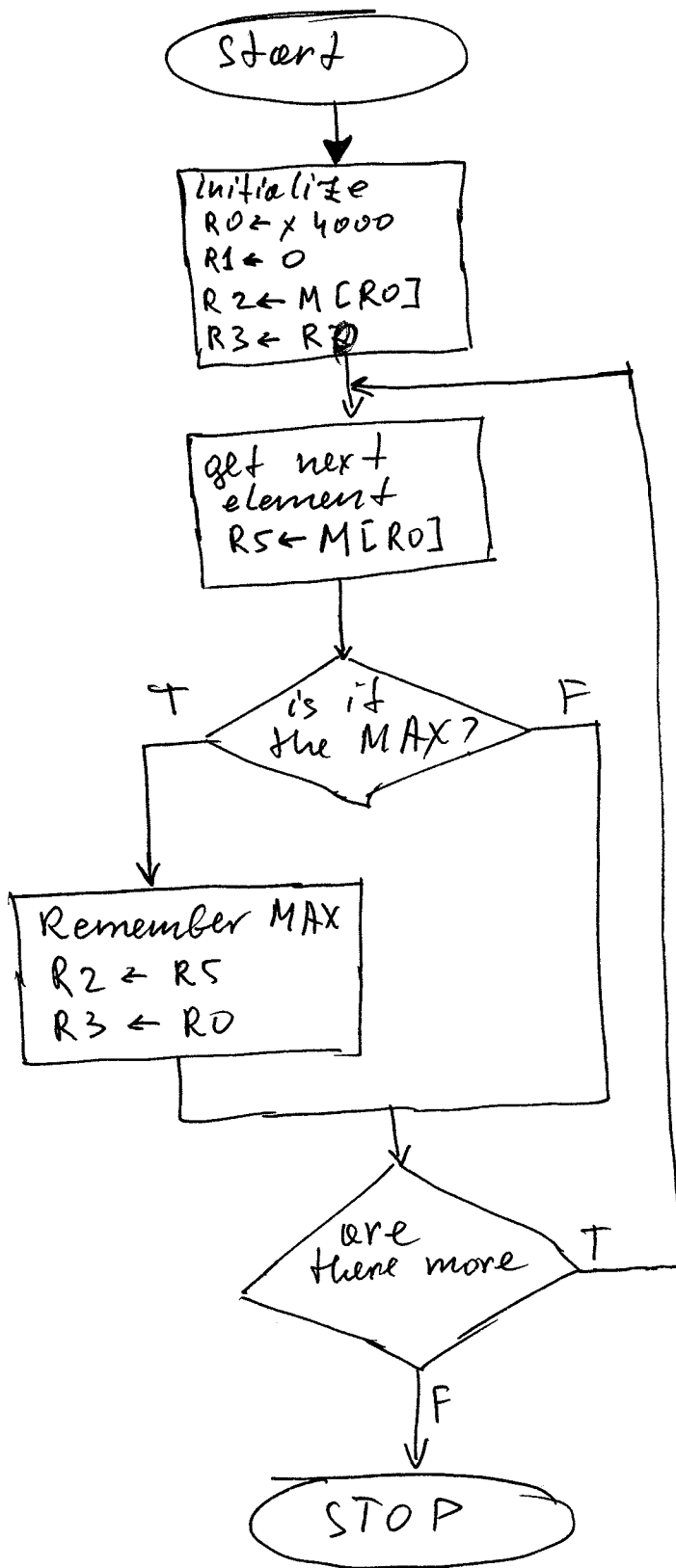
→ line-by-line, each assembler language instruction is translated into machine (binary) language.

→ each time a symbol is found as one of instruction's operands, a value from the symbol table is used to resolve the memory address of PC offset.

→ from assembler to a running program



Example 2



R0	address of next element
R1	count
R2	current max value
R3	address of current MAX
R4	-100
R5	temp
R6	temp

symbol table

Symbol	address
Loop	x 3005
SKIP	x 300C
ARRAY-START	x 3011
ARRAY-END	x 3012

```

. ORG x3000
;
; Initialize
LD R0, ARRAY-START
AND R1, R1, #0
LDR R2, R0, #0
ADD R3, R0, #0
LD R4, ARRAY-END

```

LOOP

```

; Get next element
LDR R5, R0, #0
; is it the max?
NOT R6, R2
ADD R6, R6, #1
ADD R6, R5, R6 ; R6 = R5 - R2
BRn SKIP

```

```

; Remember the Max value
ADD R2, R5, #0
ADD R3, R0, #0

```

SKIP

```

;
ADD R1, R1, #1
ADD R0, R0, #1
;
; Are there more?
ADD R5, R1, R4 ; R5 = R1 - 100
BRn LOOP
HALT

```

```

ARRAY-START .FILL x4000
ARRAY-END .FILL #-100
.END

```