Types, Tableaus and Gödel's God in Isabelle/HOL

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Abstract

A computer-formalisation of the essential parts of Fitting's textbook *Types, Tableaus and Gödel's Go*d in Isabelle/HOL is presented. In particular, Fitting's (and Anderson's) variant of the ontological argument is verified and confirmed. This variant avoids the modal collapse, which has been criticised as an undesirable side-effect of Kurt Gödel's (and Dana Scott's) versions of the ontological argument. Fitting's work is employing an intensional higher-order modal logic, which we shallowly embed here in classical higher-order logic. We then utilize the embedded logic for the formalisation of Fitting's argument.

Contents

1 Introduction

We present a study on Computational Metaphysics: a computer-formalisation and verification of Fitting's variant of the ontological argument (for the existence of God) as presented in his textbook *Types, Tableaus and Gödel's God* [10]. Fitting's argument is an emendation of Kurt Gödel's modern variant [13] (resp. Dana Scott's variant [15]) of the ontological argument.

The motivation is to avoid the *modal collapse* [16, 17], which has been criticised as an undesirable side-effect of the axioms of Gödel resp. Scott. The modal collapse essentially states that there are no contingent truths and that everything is determined. Several authors (e.g. [2, 1, 14, 9]) have proposed emendations of the argument with the aim of maintaining the essential result (the necessary existence of God) while at the same time avoiding the modal collapse. Related work has formalised several of these variants on the computer and verified or falsified them. For example, Gödel's axioms [13] have been shown inconsistent [7, 8] while Scott's version has been verified [4]. Further experiments, contributing amongst others to the clarification of a related debate between Hájek and Anderson, are presented and discussed in [5]. The enabling technique in all of these experiments has been shallow semantical embeddings of (extensional) higher-order modal logics in classical higher-order logic (see [5, 3] and the references therein).

Fitting's emendation also intends to avoid the modal collapse. However, in contrast to the above variants, Fitting's solution is based on the use of an intensional as opposed to an extensional higher-order modal logic. For our work this imposed the additional challenge to provide a shallow embedding of this more advanced logic. The experiments presented below confirm that Fitting's argument as presented in his textbook [10] is valid and that it avoids the modal collapse as intended.

The work presented here originates from the *Computational Metaphysics* lecture course held at FU Berlin in Summer 2016 [6].

2 Embedding of Intensional Higher-Order Modal Logic

The object logic being embedded (IHOML) is a modification of the intentional logic developed by Montague and Gallin (see [12]). IHOML is introduced by Fitting in the second part of the book [10] in order to formalise his emendation of Gödel's ontological argument. We offer here a shallow embedding of this logic in Isabelle/HOL, which has been inspired by previous work on the semantical embedding of multimodal logics with quantification [5]. We expand this approach to allow for actualist quantifiers, intensional types and their related operations.

2.1 Type Declarations

Since IHOML and Isabelle/HOL are both typed languages, we introduce a type-mapping between them by following as closely as possible the syntax given by Fitting (see p. 86).

typedecl i		— type for possible worlds
type-synonym io =	$(i \Rightarrow bool)$	— formulas with world-dependent truth-value
typedecl e	$(\mathbf{0})$	— individuals

Aliases for common unary predicate types:

type-synonym ie =	$(i \Rightarrow \mathbf{0})$	$(\uparrow\mathbf{0})$
type-synonym se =	$(\mathbf{0} \Rightarrow bool)$	$(\langle\mathbf{0}\rangle)$
type-synonym ise =	$(\mathbf{0} \Rightarrow io)$	$(\uparrow\langle\mathbf{0}\rangle)$
type-synonym sie =	$(\uparrow\mathbf{0} \Rightarrow bool)$	$(\langle\uparrow\mathbf{0}\rangle)$
type-synonym $isie$ =	$(\uparrow\mathbf{0} \Rightarrow io)$	$(\uparrow\langle\uparrow\mathbf{0}\rangle)$
type-synonym $sise$ =	$(\uparrow\langle\mathbf{0}\rangle \Rightarrow bool)$	$(\langle\uparrow\langle\mathbf{0}\rangle\rangle)$
type-synonym $isise$ =	$(\uparrow\langle\mathbf{0}\rangle \Rightarrow io)$	$(\uparrow\langle\uparrow\langle\mathbf{0}\rangle\rangle)$
type-synonym $sisise$=	$(\uparrow\langle\uparrow\langle\mathbf{0}\rangle\rangle \Rightarrow bool)$	$(\langle\uparrow\langle\uparrow\langle\mathbf{0}\rangle\rangle\rangle)$
type-synonym $isisise$=	$(\uparrow\langle\uparrow\langle\mathbf{0}\rangle\rangle \Rightarrow io)$	$(\uparrow\langle\uparrow\langle\uparrow\langle\mathbf{0}\rangle\rangle\rangle)$
type-synonym sse =	$\langle\mathbf{0}\rangle \Rightarrow bool$	$(\langle\langle\mathbf{0}\rangle\rangle)$
type-synonym $isse$ =	$\langle\mathbf{0}\rangle \Rightarrow io$	$(\uparrow\langle\langle\mathbf{0}\rangle\rangle)$

Aliases for common binary relation types:

type-synonym see =	$(\mathbf{0} \Rightarrow \mathbf{0} \Rightarrow bool)$	$(\langle\mathbf{0},\mathbf{0}\rangle)$
type-synonym $isee$ =	$(\mathbf{0} \Rightarrow \mathbf{0} \Rightarrow io)$	$(\uparrow\langle\mathbf{0},\mathbf{0}\rangle)$
type-synonym $sieie$ =	$(\uparrow\mathbf{0} \Rightarrow \uparrow\mathbf{0} \Rightarrow bool)$	$(\langle\uparrow\mathbf{0},\uparrow\mathbf{0}\rangle)$
type-synonym $isieie$ =	$(\uparrow\mathbf{0} \Rightarrow \uparrow\mathbf{0} \Rightarrow io)$	$(\uparrow\langle\uparrow\mathbf{0},\uparrow\mathbf{0}\rangle)$
type-synonym $ssese$ =	$(\langle\mathbf{0}\rangle \Rightarrow \langle\mathbf{0}\rangle \Rightarrow bool)$	$(\langle\langle\mathbf{0}\rangle,\langle\mathbf{0}\rangle\rangle)$
type-synonym $issese$ =	$(\langle\mathbf{0}\rangle \Rightarrow \langle\mathbf{0}\rangle \Rightarrow io)$	$(\uparrow\langle\langle\mathbf{0}\rangle,\langle\mathbf{0}\rangle\rangle)$
type-synonym $ssee$ =	$(\langle\mathbf{0}\rangle \Rightarrow \mathbf{0} \Rightarrow bool)$	$(\langle\langle\mathbf{0}\rangle,\mathbf{0}\rangle)$
type-synonym $issee$ =	$(\langle\mathbf{0}\rangle \Rightarrow \mathbf{0} \Rightarrow io)$	$(\uparrow\langle\langle\mathbf{0}\rangle,\mathbf{0}\rangle)$
type-synonym $isisee$ =	$(\uparrow\langle\mathbf{0}\rangle \Rightarrow \mathbf{0} \Rightarrow io)$	$(\uparrow\langle\uparrow\langle\mathbf{0}\rangle,\mathbf{0}\rangle)$
type-synonym $isiseise$ =	$(\uparrow\langle\mathbf{0}\rangle \Rightarrow \uparrow\langle\mathbf{0}\rangle \Rightarrow io)$	$(\uparrow\langle\uparrow\langle\mathbf{0}\rangle,\uparrow\langle\mathbf{0}\rangle\rangle)$
type-synonym $isiseisise$=	$(\uparrow\langle\mathbf{0}\rangle \Rightarrow \uparrow\langle\uparrow\langle\mathbf{0}\rangle\rangle \Rightarrow io)$	$(\uparrow\langle\uparrow\langle\mathbf{0}\rangle,\uparrow\langle\uparrow\langle\mathbf{0}\rangle\rangle\rangle)$

2.2 Definitions

2.2.1 Logical Operators as Truth-Sets

abbreviation *mnot* $:: io \Rightarrow io$ ($\neg$-[52]53)
  where $\neg\varphi \equiv \lambda w.\ \neg(\varphi\ w)$
abbreviation *negpred* $:: \langle \mathbf{0}\rangle \Rightarrow \langle \mathbf{0}\rangle$ ($\rightharpoondown$-[52]53)
  where $\rightharpoondown\Phi \equiv \lambda x.\ \neg(\Phi\ x)$
abbreviation *mnegpred* $:: \uparrow\langle \mathbf{0}\rangle \Rightarrow \uparrow\langle \mathbf{0}\rangle$ ($\rightharpoondown$-[52]53)
  where $\rightharpoondown\Phi \equiv \lambda x.\lambda w.\ \neg(\Phi\ x\ w)$
abbreviation *mand* $:: io \Rightarrow io \Rightarrow io$ (**infixr**$\wedge$51)
  where $\varphi\wedge\psi \equiv \lambda w.\ (\varphi\ w)\wedge(\psi\ w)$
abbreviation *mor* $:: io \Rightarrow io \Rightarrow io$ (**infixr**$\vee$50)
  where $\varphi\vee\psi \equiv \lambda w.\ (\varphi\ w)\vee(\psi\ w)$
abbreviation *mimp* $:: io \Rightarrow io \Rightarrow io$ (**infixr**$\rightarrow$49)
  where $\varphi\rightarrow\psi \equiv \lambda w.\ (\varphi\ w)\longrightarrow(\psi\ w)$
abbreviation *mequ* $:: io \Rightarrow io \Rightarrow io$ (**infixr**$\leftrightarrow$48)
  where $\varphi\leftrightarrow\psi \equiv \lambda w.\ (\varphi\ w)\longleftrightarrow(\psi\ w)$
abbreviation *xor*$:: bool \Rightarrow bool \Rightarrow bool$ (**infixr**$\oplus$50)
  where $\varphi\oplus\psi \equiv (\varphi\vee\psi) \wedge \neg(\varphi\wedge\psi)$
abbreviation *mxor* $:: io \Rightarrow io \Rightarrow io$ (**infixr**$\bigoplus$50)
  where $\varphi\bigoplus\psi \equiv \lambda w.\ (\varphi\ w)\oplus(\psi\ w)$

2.2.2 Possibilist Quantification

abbreviation *mforall* $:: ('t \Rightarrow io) \Rightarrow io$ ($\forall$)
  where $\forall\,\Phi \equiv \lambda w.\forall x.\ (\Phi\ x\ w)$
abbreviation *mexists* $:: ('t \Rightarrow io) \Rightarrow io$ ($\exists$)
  where $\exists\,\Phi \equiv \lambda w.\exists x.\ (\Phi\ x\ w)$

abbreviation *mforallB* $:: ('t \Rightarrow io) \Rightarrow io$ (**binder**$\forall$ [8]9) — Binder notation
  where $\forall\,x.\ \varphi(x) \equiv \forall\,\varphi$
abbreviation *mexistsB* $:: ('t \Rightarrow io) \Rightarrow io$ (**binder**$\exists$ [8]9)
  where $\exists\,x.\ \varphi(x) \equiv \exists\,\varphi$

2.2.3 Actualist Quantification

The following predicate is used to model actualist quantifiers by restricting the domain of quantification at every possible world. This standard technique has been referred to as *existence relativization* ([11] p. 106), highlighting the fact that this predicate can be seen as a kind of meta-logical 'existence predicate' telling us which individuals *actually* exist at a given world. Note that since this is a meta-logical concept it will never appear in our object language.

consts *Exists*$::\uparrow\langle \mathbf{0}\rangle$ (*existsAt*)

abbreviation *mforallAct* $:: \uparrow\langle\uparrow\langle \mathbf{0}\rangle\rangle$ ($\forall^E$)
  where $\forall^E\Phi \equiv \lambda w.\forall x.\ (existsAt\ x\ w)\longrightarrow(\Phi\ x\ w)$

abbreviation *mexistsAct* :: $\uparrow\langle\uparrow\langle\mathbf{0}\rangle\rangle$ ($\exists^E$)
where $\exists^E\Phi \equiv \lambda w.\exists x.\ (existsAt\ x\ w) \wedge (\Phi\ x\ w)$

abbreviation *mforallActB* :: $\uparrow\langle\uparrow\langle\mathbf{0}\rangle\rangle$ (**binder**$\forall^E$[8]9) — binder notation
where $\forall^E x.\ \varphi(x) \equiv \forall^E\varphi$
abbreviation *mexistsActB* :: $\uparrow\langle\uparrow\langle\mathbf{0}\rangle\rangle$ (**binder**$\exists^E$[8]9)
where $\exists^E x.\ \varphi(x) \equiv \exists^E\varphi$

2.2.4 Modal Operators

consts *aRel*::$i\Rightarrow i\Rightarrow bool$ (**infixr** r 70) — accessibility relation r

abbreviation *mbox* :: $io\Rightarrow io$ ($\Box$-[52]53)
where $\Box\varphi \equiv \lambda w.\forall v.\ (w\ r\ v)\longrightarrow(\varphi\ v)$
abbreviation *mdia* :: $io\Rightarrow io$ ($\Diamond$-[52]53)
where $\Diamond\varphi \equiv \lambda w.\exists v.\ (w\ r\ v)\wedge(\varphi\ v)$

2.2.5 *Extension-of* Operator

The IHOML ↓ operator is embedded as a predicate applying to (world-dependent) atomic formulas whose first argument is a *relativized term* (i.e. a non-rigid term). This approach can be contrasted with the one taken in Fitting's book, where ↓ is an operator which, applied to a (rigid) intensional term, gives us a new (non-rigid) extensional term ([10] p. 93, for more details). Also note that, depending on the types involved, we had to define this operator differently (empha-d) to ensure type correctness. Nevertheless, in both approaches the essence of the *Extension-of* operator remains the same: a term of the form $\downarrow\varphi$ behaves as a non-rigid term, whose denotation at a given possible world corresponds to the extension of the original intensional term φ at that world.

(a) Predicate φ takes an (intensional) individual concept as argument:

abbreviation *mextIndiv*::$\uparrow\langle\mathbf{0}\rangle\Rightarrow\uparrow\mathbf{0}\Rightarrow io$ (**infix** $\downharpoonleft$ 60)
where $\varphi \downharpoonleft c \equiv \lambda w.\ \varphi\ (c\ w)\ w$

(b) Predicate φ takes an intensional predicate as argument:

abbreviation *mextPredArg*::$(('t\Rightarrow io)\Rightarrow io)\Rightarrow('t\Rightarrow io)\Rightarrow io$ (**infix** $\downarrow$ 60)
where $\varphi \downarrow P \equiv \lambda w.\ \varphi\ (\lambda x\ u.\ P\ x\ w)\ w$

(c) Predicate φ takes an extensional predicate as argument:

abbreviation *extPredArg*::$(('t\Rightarrow bool)\Rightarrow io)\Rightarrow('t\Rightarrow io)\Rightarrow io$ (**infix** $\downarrow$ 60)
where $\varphi \downarrow P \equiv \lambda w.\ \varphi\ (\lambda x.\ P\ x\ w)\ w$

(d) Predicate φ takes an extensional predicate as *first* argument:

abbreviation *extPredArg1*::$(('t\Rightarrow bool)\Rightarrow 'b\Rightarrow io)\Rightarrow('t\Rightarrow io)\Rightarrow 'b\Rightarrow io$ (**infix** $\downarrow_1$ 60)
where $\varphi \downarrow_1 P \equiv \lambda z.\ \lambda w.\ \varphi\ (\lambda x.\ P\ x\ w)\ z\ w$

2.2.6 Equality

abbreviation *meq* :: $'t \Rightarrow 't \Rightarrow io$ (**infix**$\approx$ *60*) — normal equality (for all types)
where $x \approx y \equiv \lambda w.\ x = y$
abbreviation *meqC* :: $\uparrow\langle\uparrow\mathbf{0},\uparrow\mathbf{0}\rangle$ (**infixr**$\approx^C$ *52*) — eq. for individual concepts
where $x \approx^C y \equiv \lambda w.\ \forall v.\ (x\ v) = (y\ v)$
abbreviation *meqL* :: $\uparrow\langle\mathbf{0},\mathbf{0}\rangle$ (**infixr**$\approx^L$ *52*) — Leibniz eq. for individuals
where $x \approx^L y \equiv \forall \varphi.\ \varphi(x) \rightarrow \varphi(y)$

2.2.7 Meta-logical Predicates

abbreviation *valid* :: $io \Rightarrow bool$ ($\lfloor\text{-}\rfloor$ [*8*]) **where** $\lfloor\psi\rfloor \equiv \forall w.(\psi\ w)$
abbreviation *satisfiable* :: $io \Rightarrow bool$ ($\lfloor\text{-}\rfloor^{sat}$ [*8*]) **where** $\lfloor\psi\rfloor^{sat} \equiv \exists w.(\psi\ w)$
abbreviation *countersat* :: $io \Rightarrow bool$ ($\lfloor\text{-}\rfloor^{csat}$ [*8*]) **where** $\lfloor\psi\rfloor^{csat} \equiv \exists w.\neg(\psi\ w)$
abbreviation *invalid* :: $io \Rightarrow bool$ ($\lfloor\text{-}\rfloor^{inv}$ [*8*]) **where** $\lfloor\psi\rfloor^{inv} \equiv \forall w.\neg(\psi\ w)$

2.3 Verifying the Embedding

The above definitions introduce modal logic K with possibilist and actualist quantifiers, as evidenced by the following tests:

Verifying K Principle and Necessitation:

lemma *K*: $\lfloor(\Box(\varphi \rightarrow \psi)) \rightarrow (\Box\varphi \rightarrow \Box\psi)\rfloor$ **by** *simp* — K schema
lemma *NEC*: $\lfloor\varphi\rfloor \Longrightarrow \lfloor\Box\varphi\rfloor$ **by** *simp* — necessitation

Local consequence implies global consequence (we will use this lemma often):

lemma *localImpGlobalCons*: $\lfloor\varphi \rightarrow \xi\rfloor \Longrightarrow \lfloor\varphi\rfloor \longrightarrow \lfloor\xi\rfloor$ **by** *simp*

But global consequence does not imply local consequence:

lemma $\lfloor\varphi\rfloor \longrightarrow \lfloor\xi\rfloor \Longrightarrow \lfloor\varphi \rightarrow \xi\rfloor$ **nitpick oops** — countersatisfiable

Barcan and Converse Barcan Formulas are satisfied for standard (possibilist) quantifiers:

lemma $\lfloor(\forall x.\Box(\varphi\ x)) \rightarrow \Box(\forall x.(\varphi\ x))\rfloor$ **by** *simp*
lemma $\lfloor\Box(\forall x.(\varphi\ x)) \rightarrow (\forall x.\Box(\varphi\ x))\rfloor$ **by** *simp*

(Converse) Barcan Formulas not satisfied for actualist quantifiers:

lemma $\lfloor(\forall^E x.\Box(\varphi\ x)) \rightarrow \Box(\forall^E x.(\varphi\ x))\rfloor$ **nitpick oops** — countersatisfiable
lemma $\lfloor\Box(\forall^E x.(\varphi\ x)) \rightarrow (\forall^E x.\Box(\varphi\ x))\rfloor$ **nitpick oops** — countersatisfiable

Note that we have just made use of *Nitpick* for the first time here. *Nitpick* is a (counter-)model finder for Isabelle/HOL. In the lemmas above, *Nitpick* has found a model satisfying all axioms while falsifying the given formula. This means, the formula is not valid (i.e. is countersatisfiable).

Well known relations between meta-logical notions:

lemma $\lfloor\varphi\rfloor \longleftrightarrow \neg\lfloor\varphi\rfloor^{csat}$ **by** *simp*
lemma $\lfloor\varphi\rfloor^{sat} \longleftrightarrow \neg\lfloor\varphi\rfloor^{inv}$ **by** *simp*

Contingent truth does not allow for necessitation:

lemma $\lfloor \Diamond\varphi \rfloor \longrightarrow \lfloor \Box\varphi \rfloor$ **nitpick oops** — countersatisfiable
lemma $\lfloor \Box\varphi \rfloor^{sat} \longrightarrow \lfloor \Box\varphi \rfloor$ **nitpick oops** — countersatisfiable

Modal collapse is countersatisfiable:

lemma $\lfloor \varphi \rightarrow \Box\varphi \rfloor$ **nitpick oops** — countersatisfiable

2.4 Useful Definitions for Axiomatization of Further Logics

The best known normal logics (*K4, K5, KB, K45, KB5, D, D4, D5, D45, ...*) can be obtained by combinations of the following axioms:

abbreviation M
 where $M \equiv \forall \varphi.\ \Box\varphi \rightarrow \varphi$
abbreviation B
 where $B \equiv \forall \varphi.\ \varphi \rightarrow \Box\Diamond\varphi$
abbreviation D
 where $D \equiv \forall \varphi.\ \Box\varphi \rightarrow \Diamond\varphi$
abbreviation IV
 where $IV \equiv \forall \varphi.\ \Box\varphi \rightarrow \Box\Box\varphi$
abbreviation V
 where $V \equiv \forall \varphi.\ \Diamond\varphi \rightarrow \Box\Diamond\varphi$

Because the embedding is of a semantic nature, it is more efficient to instead make use of the well-known *Sahlqvist correspondence*, which links axioms to constraints on a model's accessibility relation (e.g. reflexive, symmetric, etc. whose definitions are not shown here). We show below that axioms M, B, D, IV, V impose reflexivity, symmetry, seriality, transitivity and euclideanness respectively.

lemma *reflexive aRel* $\Longrightarrow \lfloor M \rfloor$ **by** *blast* — aka T
lemma *symmetric aRel* $\Longrightarrow \lfloor B \rfloor$ **by** *blast*
lemma *serial aRel* $\Longrightarrow \lfloor D \rfloor$ **by** *blast*
lemma *transitive aRel* $\Longrightarrow \lfloor IV \rfloor$ **by** *blast*
lemma *preorder aRel* $\Longrightarrow \lfloor M \rfloor \wedge \lfloor IV \rfloor$ **by** *blast* — S4: reflexive + transitive
lemma *equivalence aRel* $\Longrightarrow \lfloor M \rfloor \wedge \lfloor V \rfloor$ **by** *blast* — S5: preorder + symmetric

lemma *reflexive aRel* $\wedge$ *euclidean aRel* $\Longrightarrow \lfloor M \rfloor \wedge \lfloor V \rfloor$ **by** *blast* — S5

Using these definitions, we can derive axioms for the most common modal logics. Thereby we are free to use either the semantic constraints or the related *Sahlqvist* axioms. Here we provide both versions. In what follows we use the semantic constraints for improved performance.

3 Textbook Examples

In this section we verify that our embedded logic works as intended by proving the examples provided in the book. In many cases, for good mesure, we consider further theorems derived from the original ones. We were able to confirm that all results (proofs or counterexamples) agree with Fitting's claims.

3.1 Modal Logic - Syntax and Semantics (Chapter 7)

Note: In what follows, we will call a term *relativized* if it is of the form $\downarrow\varphi$ (i.e. an intensional term preceded by the *extension-of* operator), otherwise it is *non-relativized*. Note that all and only relativized terms are non-rigid.

3.1.1 Considerations Regarding $\beta\eta$-redex (p. 94)

$\beta\eta$-redex is valid for non-relativized (intensional or extensional) terms (because they designate rigidly):

lemma $\lfloor((\lambda\alpha.\ \varphi\ \alpha)\ \ (\tau{::}\uparrow\mathbf{0})) \leftrightarrow (\varphi\ \ \tau)\rfloor$ **by** *simp*
lemma $\lfloor((\lambda\alpha.\ \varphi\ \alpha)\ \ (\tau{::}\mathbf{0})) \leftrightarrow (\varphi\ \ \tau)\rfloor$ **by** *simp*
lemma $\lfloor((\lambda\alpha.\ \Box\varphi\ \alpha)\ (\tau{::}\uparrow\mathbf{0})) \leftrightarrow (\Box\varphi\ \tau)\rfloor$ **by** *simp*
lemma $\lfloor((\lambda\alpha.\ \Box\varphi\ \alpha)\ (\tau{::}\mathbf{0})) \leftrightarrow (\Box\varphi\ \tau)\rfloor$ **by** *simp*

$\beta\eta$-redex is valid for relativized terms as long as no modal operators occur inside the predicate abstract:

lemma $\lfloor((\lambda\alpha.\ \varphi\ \alpha)\ \downharpoonleft(\tau{::}\uparrow\mathbf{0})) \leftrightarrow (\varphi\ \downharpoonleft\tau)\rfloor$ **by** *simp*

$\beta\eta$-redex is non-valid for relativized terms when modal operators are present:

lemma $\lfloor((\lambda\alpha.\ \Box\varphi\ \alpha)\ \downharpoonleft(\tau{::}\uparrow\mathbf{0})) \leftrightarrow (\Box\varphi\ \downharpoonleft\tau)\rfloor$ **nitpick oops** — countersatisfiable
lemma $\lfloor((\lambda\alpha.\ \Diamond\varphi\ \alpha)\ \downharpoonleft(\tau{::}\uparrow\mathbf{0})) \leftrightarrow (\Diamond\varphi\ \downharpoonleft\tau)\rfloor$ **nitpick oops** — countersatisfiable

Example 7.13, p. 96:

lemma $\lfloor(\lambda X.\ \Diamond\exists\, X)\ \ (P{::}\uparrow\langle\mathbf{0}\rangle) \rightarrow \Diamond((\lambda X.\ \exists\, X)\ \ P)\rfloor$ **by** *simp*
lemma $\lfloor(\lambda X.\ \Diamond\exists\, X)\ {\downarrow}(P{::}\uparrow\langle\mathbf{0}\rangle) \rightarrow \Diamond((\lambda X.\ \exists\, X)\ {\downarrow}P)\rfloor$
nitpick[*card 't=1*, *card i=2*] **oops** — nitpick finds same counterexample as book

with other types for P:

lemma $\lfloor(\lambda X.\ \Diamond\exists\, X)\ \ (P{::}\uparrow\langle\uparrow\mathbf{0}\rangle) \rightarrow \Diamond((\lambda X.\ \exists\, X)\ \ P)\rfloor$ **by** *simp*
lemma $\lfloor(\lambda X.\ \Diamond\exists\, X)\ {\downarrow}(P{::}\uparrow\langle\uparrow\mathbf{0}\rangle) \rightarrow \Diamond((\lambda X.\ \exists\, X)\ {\downarrow}P)\rfloor$
nitpick[*card 't=1*, *card i=2*] **oops** — countersatisfiable
lemma $\lfloor(\lambda X.\ \Diamond\exists\, X)\ \ (P{::}\uparrow\langle\langle\mathbf{0}\rangle\rangle) \rightarrow \Diamond((\lambda X.\ \exists\, X)\ \ P)\rfloor$ **by** *simp*
lemma $\lfloor(\lambda X.\ \Diamond\exists\, X)\ {\downarrow}(P{::}\uparrow\langle\langle\mathbf{0}\rangle\rangle) \rightarrow \Diamond((\lambda X.\ \exists\, X)\ {\downarrow}P)\rfloor$
nitpick[*card 't=1*, *card i=2*] **oops** — countersatisfiable
lemma $\lfloor(\lambda X.\ \Diamond\exists\, X)\ \ (P{::}\uparrow\langle\uparrow\langle\mathbf{0}\rangle\rangle)\rightarrow \Diamond((\lambda X.\ \exists\, X)\ \ P)\rfloor$ **by** *simp*
lemma $\lfloor(\lambda X.\ \Diamond\exists\, X)\ {\downarrow}(P{::}\uparrow\langle\uparrow\langle\mathbf{0}\rangle\rangle)\rightarrow \Diamond((\lambda X.\ \exists\, X)\ {\downarrow}P)\rfloor$

nitpick[*card 't=1*, *card i=2*] **oops** — countersatisfiable

Example 7.14, p. 98:

lemma $\lfloor(\lambda X.\ \Diamond\exists\, X)\ {\downarrow}(P{::}{\uparrow}\langle\mathbf{0}\rangle) \rightarrow (\lambda X.\ \exists\, X)\ {\downarrow}P\rfloor$ **by** *simp*
lemma $\lfloor(\lambda X.\ \Diamond\exists\, X)\ \ (P{::}{\uparrow}\langle\mathbf{0}\rangle) \rightarrow (\lambda X.\ \exists\, X)\ \ P\rfloor$
nitpick[*card 't=1*, *card i=2*] **oops** — countersatisfiable

with other types for P:

lemma $\lfloor(\lambda X.\ \Diamond\exists\, X)\ {\downarrow}(P{::}{\uparrow}\langle{\uparrow}\mathbf{0}\rangle) \rightarrow (\lambda X.\ \exists\, X)\ {\downarrow}P\rfloor$ **by** *simp*
lemma $\lfloor(\lambda X.\ \Diamond\exists\, X)\ \ (P{::}{\uparrow}\langle{\uparrow}\mathbf{0}\rangle) \rightarrow (\lambda X.\ \exists\, X)\ \ P\rfloor$
nitpick[*card 't=1*, *card i=2*] **oops** — countersatisfiable
lemma $\lfloor(\lambda X.\ \Diamond\exists\, X)\ {\downarrow}(P{::}{\uparrow}\langle\langle\mathbf{0}\rangle\rangle) \rightarrow (\lambda X.\ \exists\, X)\ {\downarrow}P\rfloor$ **by** *simp*
lemma $\lfloor(\lambda X.\ \Diamond\exists\, X)\ \ (P{::}{\uparrow}\langle\langle\mathbf{0}\rangle\rangle) \rightarrow (\lambda X.\ \exists\, X)\ \ P\rfloor$
nitpick[*card 't=1*, *card i=2*] **oops** — countersatisfiable
lemma $\lfloor(\lambda X.\ \Diamond\exists\, X)\ {\downarrow}(P{::}{\uparrow}\langle{\uparrow}\langle\mathbf{0}\rangle\rangle){\rightarrow}\ (\lambda X.\ \exists\, X)\ {\downarrow}P\rfloor$ **by** *simp*
lemma $\lfloor(\lambda X.\ \Diamond\exists\, X)\ \ (P{::}{\uparrow}\langle{\uparrow}\langle\mathbf{0}\rangle\rangle){\rightarrow}\ (\lambda X.\ \exists\, X)\ \ P\rfloor$
nitpick[*card 't=1*, *card i=2*] **oops** — countersatisfiable

Example 7.15, p. 99:

lemma $\lfloor\Box(P\ \ (c{::}{\uparrow}\mathbf{0})) \rightarrow (\exists\, x{::}{\uparrow}\mathbf{0}.\ \Box(P\ x))\rfloor$ **by** *auto*

with other types for P:

lemma $\lfloor\Box(P\ \ (c{::}\mathbf{0})) \rightarrow (\exists\, x{::}\mathbf{0}.\ \Box(P\ x))\rfloor$ **by** *auto*
lemma $\lfloor\Box(P\ \ (c{::}\langle\mathbf{0}\rangle)) \rightarrow (\exists\, x{::}\langle\mathbf{0}\rangle.\ \Box(P\ x))\rfloor$ **by** *auto*

Example 7.16, p. 100:

lemma $\lfloor\Box(P\ {\downharpoonleft}(c{::}{\uparrow}\mathbf{0})) \rightarrow (\exists\, x{::}\mathbf{0}.\ \Box(P\ x))\rfloor$
nitpick[*card 't=2*, *card i=2*] **oops** — counterexample with two worlds found

Example 7.17, p. 101:

lemma $\lfloor\forall\, Z{::}{\uparrow}\mathbf{0}.\ (\lambda x{::}\mathbf{0}.\ \ \Box((\lambda y{::}\mathbf{0}.\ \ x \approx y)\ {\downharpoonleft}Z))\ {\downharpoonleft}Z\rfloor$
nitpick[*card 't=2*, *card i=2*] **oops** — countersatisfiable
lemma $\lfloor\forall\, z{::}\mathbf{0}.\ \ (\lambda x{::}\mathbf{0}.\ \ \Box((\lambda y{::}\mathbf{0}.\ \ x \approx y)\ \ \ z))\ z\rfloor$ **by** *simp*
lemma $\lfloor\forall\, Z{::}{\uparrow}\mathbf{0}.\ (\lambda X{::}{\uparrow}\mathbf{0}.\ \Box((\lambda\, Y{::}{\uparrow}\mathbf{0}.\ X \approx\ Y)\ \ \ Z))\ Z\rfloor$ **by** *simp*

3.1.2 Exercises (p. 101)

For Exercises 7.1 and 7.2 see variations on Examples 7.13 and 7.14 above.

Exercise 7.3:

lemma $\lfloor\Diamond\exists\,(P{::}{\uparrow}\langle\mathbf{0}\rangle) \rightarrow (\exists\, X{::}{\uparrow}\mathbf{0}.\ \Diamond(P\ {\downharpoonleft}X))\rfloor$ **by** *auto*

Exercise 7.4:

lemma $\lfloor\Diamond(\exists\, x{::}\mathbf{0}.\ (\lambda\, Y.\ \ Y\ x)\ {\downarrow}(P{::}{\uparrow}\langle\mathbf{0}\rangle)) \rightarrow (\exists\, x.\ (\lambda\, Y.\ \Diamond(Y\ x))\ {\downarrow}P)\rfloor$
nitpick[*card 't=1*, *card i=2*] **oops** — countersatisfiable

For Exercise 7.5 see Example 7.17 above.

3.2 Miscellaneous Matters (Chapter 9)

3.2.1 Equality Axioms (Subsection 1.1)

Example 9.1:

lemma $\lfloor((\lambda X.\ \Box(X\ \downharpoonleft(p{::}\uparrow\mathbf{0})))\ \downarrow(\lambda x.\ \Diamond(\lambda z.\ z \approx x)\ \downharpoonleft p))\rfloor$
by *auto* — using normal equality
lemma $\lfloor((\lambda X.\ \Box(X\ \downharpoonleft(p{::}\uparrow\mathbf{0})))\ \downarrow(\lambda x.\ \Diamond(\lambda z.\ z \approx^L x)\ \downharpoonleft p))\rfloor$
by *auto* — using Leibniz equality
lemma $\lfloor((\lambda X.\ \Box(X\ \ (p{::}\uparrow\mathbf{0})))\ \downarrow(\lambda x.\ \Diamond(\lambda z.\ z \approx^C x)\ p))\rfloor$
by *simp* — using equality as defined for individual concepts

3.2.2 Extensionality (Subsection 1.2)

In the book (p. 118), extensionality is assumed (globally) for extensional terms. Whereas Fitting introduces following extensionality principles as axioms, they are already implicit in Isabelle/HOL:

lemma *EXT*: $\forall\,\alpha{::}\langle\mathbf{0}\rangle.\ \forall\,\beta{::}\langle\mathbf{0}\rangle.\ (\forall\,\gamma{::}\mathbf{0}.\ (\alpha\ \gamma \longleftrightarrow \beta\ \gamma)) \longrightarrow (\alpha = \beta)$ **by** *auto*
lemma *EXT-set*: $\forall\,\alpha{::}\langle\langle\mathbf{0}\rangle\rangle.\ \forall\,\beta{::}\langle\langle\mathbf{0}\rangle\rangle.\ (\forall\,\gamma{::}\langle\mathbf{0}\rangle.\ (\alpha\ \gamma \longleftrightarrow \beta\ \gamma)) \longrightarrow (\alpha = \beta)$
by *auto*

3.2.3 *De Re* and *De Dicto* (Subsection 2)

De re is equivalent to *de dicto* for non-relativized (extensional or intensional) terms:

lemma $\lfloor\forall\,\alpha.\ ((\lambda\beta.\ \Box(\alpha\ \beta))\ (\tau{::}\mathbf{0}))\ \ \leftrightarrow \Box((\lambda\beta.\ (\alpha\ \beta))\ \tau)\rfloor$ **by** *simp*
lemma $\lfloor\forall\,\alpha.\ ((\lambda\beta.\ \Box(\alpha\ \beta))\ (\tau{::}\uparrow\mathbf{0}))\ \leftrightarrow \Box((\lambda\beta.\ (\alpha\ \beta))\ \tau)\rfloor$ **by** *simp*
lemma $\lfloor\forall\,\alpha.\ ((\lambda\beta.\ \Box(\alpha\ \beta))\ (\tau{::}\langle\mathbf{0}\rangle))\ \leftrightarrow \Box((\lambda\beta.\ (\alpha\ \beta))\ \tau)\rfloor$ **by** *simp*
lemma $\lfloor\forall\,\alpha.\ ((\lambda\beta.\ \Box(\alpha\ \beta))\ (\tau{::}\uparrow\langle\mathbf{0}\rangle))\ \leftrightarrow \Box((\lambda\beta.\ (\alpha\ \beta))\ \tau)\rfloor$ **by** *simp*

De re is not equivalent to *de dicto* for relativized (intensional) terms:

lemma $\lfloor\forall\,\alpha.\ ((\lambda\beta.\ \Box(\alpha\ \beta))\ \downharpoonleft(\tau{::}\uparrow\mathbf{0}))\ \leftrightarrow \Box(\ (\lambda\beta.\ (\alpha\ \beta))\ \downharpoonleft\tau)\rfloor$
nitpick[*card 't=2*, *card i=2*] **oops** — countersatisfiable
lemma $\lfloor\forall\,\alpha.\ ((\lambda\beta.\ \Box(\alpha\ \beta))\ \downarrow(\tau{::}\uparrow\langle\mathbf{0}\rangle))\ \leftrightarrow \Box(\ (\lambda\beta.\ (\alpha\ \beta))\ \downarrow\tau)\rfloor$
nitpick[*card 't=1*, *card i=2*] **oops** — countersatisfiable

Proposition 9.6 - If we can prove one side of the equivalence, then we can prove the other (p. 120):

abbreviation *deDictoImplDeRe*$::\uparrow\mathbf{0}\Rightarrow io$
where *deDictoImplDeRe* $\tau \equiv \forall\,\alpha.\ \Box((\lambda\beta.\ (\alpha\ \beta))\ \downharpoonleft\tau) \rightarrow ((\lambda\beta.\ \Box(\alpha\ \beta))\ \downharpoonleft\tau)$
abbreviation *deReImplDeDicto*$::\uparrow\mathbf{0}\Rightarrow io$
where *deReImplDeDicto* $\tau \equiv \forall\,\alpha.\ ((\lambda\beta.\ \Box(\alpha\ \beta))\ \downharpoonleft\tau) \rightarrow \Box((\lambda\beta.\ (\alpha\ \beta))\ \downharpoonleft\tau)$
abbreviation *deReEquDeDicto*$::\uparrow\mathbf{0}\Rightarrow io$
where *deReEquDeDicto* $\tau \equiv \forall\,\alpha.\ ((\lambda\beta.\ \Box(\alpha\ \beta))\ \downharpoonleft\tau) \leftrightarrow \Box((\lambda\beta.\ (\alpha\ \beta))\ \downharpoonleft\tau)$

abbreviation *deDictoImplDeRe-pred*$::('t\Rightarrow io)\Rightarrow io$
where *deDictoImplDeRe-pred* $\tau \equiv \forall\,\alpha.\ \Box((\lambda\beta.\ (\alpha\ \beta))\ \downarrow\tau) \rightarrow ((\lambda\beta.\ \Box(\alpha\ \beta))\ \downarrow\tau)$

abbreviation $deReImplDeDicto\text{-}pred{::}('t{\Rightarrow}io){\Rightarrow}io$
where $deReImplDeDicto\text{-}pred\ \tau \equiv \forall\,\alpha.\ ((\lambda\beta.\ \Box(\alpha\ \beta))\ {\downarrow}\tau) \to \Box((\lambda\beta.\ (\alpha\ \beta))\ {\downarrow}\tau)$
abbreviation $deReEquDeDicto\text{-}pred{::}('t{\Rightarrow}io){\Rightarrow}io$
where $deReEquDeDicto\text{-}pred\ \tau \equiv \forall\,\alpha.\ ((\lambda\beta.\ \Box(\alpha\ \beta))\ {\downarrow}\tau) \leftrightarrow \Box((\lambda\beta.\ (\alpha\ \beta))\ {\downarrow}\tau)$

We can prove local consequence:

lemma *AimpB*: $\lfloor deReImplDeDicto\ (\tau{::}{\uparrow}\mathbf{0}) \to deDictoImplDeRe\ \tau\rfloor$
by *force* — for individuals
lemma *AimpB-p*: $\lfloor deReImplDeDicto\text{-}pred\ (\tau{::}{\uparrow}\langle\mathbf{0}\rangle) \to deDictoImplDeRe\text{-}pred\ \tau\rfloor$
by *force* — for predicates

And global consequence follows directly (since local consequence implies global consequence, as shown before):

lemma $\lfloor deReImplDeDicto\ (\tau{::}{\uparrow}\mathbf{0})\rfloor \longrightarrow \lfloor deDictoImplDeRe\ \tau\rfloor$
using *AimpB* **by** (*rule localImpGlobalCons*) — for individuals
lemma $\lfloor deReImplDeDicto\text{-}pred\ (\tau{::}{\uparrow}\langle\mathbf{0}\rangle)\rfloor \longrightarrow \lfloor deDictoImplDeRe\text{-}pred\ \tau\rfloor$
using *AimpB-p* **by** (*rule localImpGlobalCons*) — for predicates

3.2.4 Rigidity (Subsection 3)

(Local) rigidity for intensional individuals:

abbreviation $rigidIndiv{::}{\uparrow}\langle{\uparrow}\mathbf{0}\rangle$ **where**
$rigidIndiv\ \tau \equiv (\lambda\beta.\ \Box((\lambda z.\ \beta \approx z)\ {\downharpoonleft}\tau))\ {\downharpoonleft}\tau$

(Local) rigidity for intensional predicates:

abbreviation $rigidPred{::}('t{\Rightarrow}io){\Rightarrow}io$ **where**
$rigidPred\ \tau \equiv (\lambda\beta.\ \Box((\lambda z.\ \beta \approx z)\ {\downarrow}\tau))\ {\downarrow}\tau$

Proposition 9.8 - An intensional term is rigid if and only if the *de re/de dicto* distinction vanishes. Note that we can prove this theorem for local consequence (global consequence follows directly).

lemma $\lfloor rigidIndiv\ (\tau{::}{\uparrow}\mathbf{0}) \to deReEquDeDicto\ \tau\rfloor$ **by** *simp*
lemma $\lfloor deReImplDeDicto\ (\tau{::}{\uparrow}\mathbf{0}) \to rigidIndiv\ \tau\rfloor$ **by** *auto*
lemma $\lfloor rigidPred\ (\tau{::}{\uparrow}\langle\mathbf{0}\rangle) \to deReEquDeDicto\text{-}pred\ \tau\rfloor$ **by** *simp*
lemma $\lfloor deReImplDeDicto\text{-}pred\ (\tau{::}{\uparrow}\langle\mathbf{0}\rangle) \to rigidPred\ \tau\rfloor$ **by** *auto*

3.2.5 Stability Conditions (Subsection 4)

axiomatization where
S5: *equivalence aRel* — using Sahlqvist correspondence for improved performance

Definition 9.10 - Stability conditions come in pairs:

abbreviation $stabilityA{::}('t{\Rightarrow}io){\Rightarrow}io$ **where** $stabilityA\ \tau \equiv \forall\,\alpha.\ (\tau\ \alpha) \to \Box(\tau\ \alpha)$
abbreviation $stabilityB{::}('t{\Rightarrow}io){\Rightarrow}io$ **where** $stabilityB\ \tau \equiv \forall\,\alpha.\ \Diamond(\tau\ \alpha) \to (\tau\ \alpha)$

Proposition 9.10 - In an *S5* modal logic both stability conditions are equivalent.

The last proposition holds for global consequence:

lemma $\lfloor stabilityA\ (\tau{::}\uparrow\langle\mathbf{0}\rangle)\rfloor \longrightarrow \lfloor stabilityB\ \tau\rfloor$ **using** *S5* **by** *blast*
lemma $\lfloor stabilityB\ (\tau{::}\uparrow\langle\mathbf{0}\rangle)\rfloor \longrightarrow \lfloor stabilityA\ \tau\rfloor$ **using** *S5* **by** *blast*

But it does not hold for local consequence:

lemma $\lfloor stabilityA\ (\tau{::}\uparrow\langle\mathbf{0}\rangle) \rightarrow stabilityB\ \tau\rfloor$
nitpick[*card 't=1*, *card i=2*] **oops** — countersatisfiable
lemma $\lfloor stabilityB\ (\tau{::}\uparrow\langle\mathbf{0}\rangle) \rightarrow stabilityA\ \tau\rfloor$
nitpick[*card 't=1*, *card i=2*] **oops** — countersatisfiable

Theorem 9.11 - A term is rigid if and only if it satisfies the stability conditions. Note that we can prove this theorem for local consequence (global consequence follows directly).

theorem $\lfloor rigidPred\ (\tau{::}\uparrow\langle\mathbf{0}\rangle) \leftrightarrow (stabilityA\ \tau \wedge stabilityB\ \tau)\rfloor$ **by** *meson*
theorem $\lfloor rigidPred\ (\tau{::}\uparrow\langle\uparrow\mathbf{0}\rangle) \leftrightarrow (stabilityA\ \tau \wedge stabilityB\ \tau)\rfloor$ **by** *meson*
theorem $\lfloor rigidPred\ (\tau{::}\uparrow\langle\uparrow\langle\mathbf{0}\rangle\rangle) \leftrightarrow (stabilityA\ \tau \wedge stabilityB\ \tau)\rfloor$ **by** *meson*

4 Gödel's Argument, Formally

"Gödel's particular version of the argument is a direct descendent of that of Leibniz, which in turn derives from one of Descartes. These arguments all have a two-part structure: prove God's existence is necessary, if possible; and prove God's existence is possible." [10] p. 138.

4.1 Part I - God's Existence is Possible

We divide Gödel's Argument as presented in Fitting's textbook (ch. 11) in two parts. For the first one, while Leibniz provides some kind of proof for the compatibility of all perfections, Gödel goes on to prove an analogous result: *(T1) Every positive property is possibly instantiated*, which together with *(T2) God is a positive property* directly implies the conclusion. In order to prove *T1* Gödel assumes *A2: Any property entailed by a positive property is positive.*

We are currently contemplating a follow-up analysis of the philosophical implications of these axioms, which encompasses some criticism of the notion of *property entailment* used by Gödel throughout the argument.

4.1.1 General Definitions

abbreviation $existencePredicate{::}\uparrow\langle\mathbf{0}\rangle$ $(E!)$
  where $E!\ x \ \equiv \lambda w.\ (\exists^E y.\ y{\approx}x)\ w$ — existence predicate in object language
lemma $E!\ x\ w \longleftrightarrow existsAt\ x\ w$
  by *simp* — safety check: $E!$ correctly matches its meta-logical counterpart

consts $positiveProperty{::}\uparrow\langle\uparrow\langle\mathbf{0}\rangle\rangle$ $(\mathcal{P})$ — Positiveness/Perfection

Definitions of God (later shown to be equivalent under axiom *A1b*):

abbreviation $God{::}\uparrow\langle\mathbf{0}\rangle$ (G) **where** $G \equiv (\lambda x.\ \forall\, Y.\ \mathcal{P}\ Y \rightarrow\ Y\, x)$
abbreviation $God\text{-}star{::}\uparrow\langle\mathbf{0}\rangle$ $(G*)$ **where** $G* \equiv (\lambda x.\ \forall\, Y.\ \mathcal{P}\ Y \leftrightarrow\ Y\, x)$

Definitions needed to formalise *A3*:

abbreviation $appliesToPositiveProps{::}\uparrow\langle\uparrow\langle\uparrow\langle\mathbf{0}\rangle\rangle\rangle$ (pos) **where**
  $pos\ Z \equiv\ \forall\, X.\ Z\, X \rightarrow \mathcal{P}\ X$
abbreviation $intersectionOf{::}\uparrow\langle\uparrow\langle\mathbf{0}\rangle,\uparrow\langle\uparrow\langle\mathbf{0}\rangle\rangle\rangle$ $(intersec)$ **where**
  $intersec\ X\ Z \equiv\ \Box(\forall\, x.(X\, x \leftrightarrow (\forall\, Y.\ (Z\ Y) \rightarrow (Y\, x))))$ — quantifier is possibilist

abbreviation $Entailment{::}\uparrow\langle\uparrow\langle\mathbf{0}\rangle,\uparrow\langle\mathbf{0}\rangle\rangle$ (**infix** $\Rrightarrow$ *60*) **where**
  $X \Rrightarrow\ Y \equiv\ \Box(\forall^E z.\ X\ z \rightarrow\ Y\ z)$

4.1.2 Axioms

axiomatization where
  $A1a{:}\lfloor\forall\, X.\ \mathcal{P}\ (\rightharpoondown X) \rightarrow \neg(\mathcal{P}\ X)\ \rfloor$ **and** — axiom 11.3A

A1b:$\lfloor\forall X.\ \neg(\mathcal{P}\ X) \rightarrow \mathcal{P}\ (\rightharpoondown X)\rfloor$ **and** — axiom 11.3B
A2: $\lfloor\forall X\ Y.\ (\mathcal{P}\ X \wedge (X \Rrightarrow Y)) \rightarrow \mathcal{P}\ Y\rfloor$ **and** — axiom 11.5
A3: $\lfloor\forall Z\ X.\ (pos\ Z \wedge intersec\ X\ Z) \rightarrow \mathcal{P}\ X\rfloor$ — axiom 11.10

lemma *True* **nitpick**[*satisfy*] **oops** — model found: axioms are consistent

lemma $\lfloor D\rfloor$ **using** *A1a A1b A2* **by** *blast* — axioms already imply D axiom
lemma $\lfloor D\rfloor$ **using** *A1a A3* **by** *metis*

4.1.3 Theorems

lemma $\lfloor\exists X.\ \mathcal{P}\ X\rfloor$ **using** *A1b* **by** *auto*
lemma $\lfloor\exists X.\ \mathcal{P}\ X \wedge\ \Diamond\exists^E\ X\rfloor$ **using** *A1a A1b A2* **by** *metis*

Being self-identical is a positive property:

lemma $\lfloor(\exists X.\ \mathcal{P}\ X \wedge\ \Diamond\exists^E\ X) \rightarrow \mathcal{P}\ (\lambda x\ w.\ x = x)\rfloor$ **using** *A2* **by** *fastforce*

Proposition 11.6

lemma $\lfloor(\exists X.\ \mathcal{P}\ X) \rightarrow \mathcal{P}\ (\lambda x\ w.\ x = x)\rfloor$ **using** *A2* **by** *fastforce*

lemma $\lfloor\mathcal{P}\ (\lambda x\ w.\ x = x)\rfloor$ **using** *A1b A2* **by** *blast*
lemma $\lfloor\mathcal{P}\ (\lambda x\ w.\ x = x)\rfloor$ **using** *A3* **by** *metis*

Being non-self-identical is a negative property:

lemma $\lfloor(\exists X.\ \mathcal{P}\ X\ \wedge \Diamond\exists^E\ X) \rightarrow\ \mathcal{P}\ (\rightharpoondown (\lambda x\ w.\ \neg x = x))\rfloor$
using *A2* **by** *fastforce*

lemma $\lfloor(\exists X.\ \mathcal{P}\ X) \rightarrow\ \mathcal{P}\ (\rightharpoondown (\lambda x\ w.\ \neg x = x))\rfloor$ **using** *A2* **by** *fastforce*
lemma $\lfloor(\exists X.\ \mathcal{P}\ X) \rightarrow\ \mathcal{P}\ (\rightharpoondown (\lambda x\ w.\ \neg x = x))\rfloor$ **using** *A3* **by** *metis*

Proposition 11.7

lemma $\lfloor(\exists X.\ \mathcal{P}\ X) \rightarrow \neg\mathcal{P}\ ((\lambda x\ w.\ \neg x = x))\rfloor$ **using** *A1a A2* **by** *blast*
lemma $\lfloor\neg\mathcal{P}\ (\lambda x\ w.\ \neg x = x)\rfloor$ **using** *A1a A2* **by** *blast*

Proposition 11.8 (Informal Proposition 1) - Positive properties are possibly instantiated:

theorem *T1*: $\lfloor\forall X.\ \mathcal{P}\ X \rightarrow \Diamond\exists^E\ X\rfloor$ **using** *A1a A2* **by** *blast*

Proposition 11.14 - Both defs (*God/God**) are equivalent. For improved performance we may prefer to use one or the other:

lemma *GodDefsAreEquivalent*: $\lfloor\forall x.\ G\ x \leftrightarrow G{*}\ x\rfloor$ **using** *A1b* **by** *force*

Proposition 11.15 - Possibilist existence of *God* directly implies *A1b*:

lemma $\lfloor\exists\ G{*} \rightarrow (\forall X.\ \neg(\mathcal{P}\ X) \rightarrow \mathcal{P}\ (\rightharpoondown X))\rfloor$ **by** *meson*

Proposition 11.16 - *A3* implies *P(G)* (local consequence):

lemma *A3implT2-local*: $\lfloor(\forall Z\ X.\ (pos\ Z \wedge intersec\ X\ Z) \rightarrow \mathcal{P}\ X) \rightarrow \mathcal{P}\ G\rfloor$

proof –
{
fix w
have *1*: *pos* $\mathcal{P}$ w **by** *simp*
have *2*: *intersec* G $\mathcal{P}$ w **by** *simp*
{
assume $(\forall Z\ X.\ (pos\ Z \wedge intersec\ X\ Z) \rightarrow \mathcal{P}\ X)\ w$
hence $(\forall X.\ ((pos\ \mathcal{P}) \wedge (intersec\ X\ \mathcal{P})) \rightarrow \mathcal{P}\ X)\ w$ **by** (*rule allE*)
hence $(((pos\ \mathcal{P}) \wedge (intersec\ G\ \mathcal{P})) \rightarrow \mathcal{P}\ G)\ w$ **by** (*rule allE*)
hence *3*: $((pos\ \mathcal{P} \wedge intersec\ G\ \mathcal{P})\ w) \longrightarrow \mathcal{P}\ G\ w$ **by** *simp*
hence *4*: $((pos\ \mathcal{P}) \wedge (intersec\ G\ \mathcal{P}))\ w$ **using** *1 2* **by** *simp*
from *3 4* **have** $\mathcal{P}\ G\ w$ **by** (*rule mp*)
}
hence $(\forall Z\ X.\ (pos\ Z \wedge intersec\ X\ Z) \rightarrow \mathcal{P}\ X)\ w\ \longrightarrow \mathcal{P}\ G\ w$ **by** (*rule impI*)
}
thus *?thesis* **by** (*rule allI*)
qed

A3 implies $P(G)$ (as global consequence):

lemma *A3implT2-global*: $\lfloor \forall Z\ X.\ (pos\ Z \wedge intersec\ X\ Z) \rightarrow \mathcal{P}\ X \rfloor \longrightarrow \lfloor \mathcal{P}\ G \rfloor$
using *A3implT2-local* **by** (*rule localImpGlobalCons*)

Being Godlike is a positive property. Note that this theorem can be axiomatized directly, as noted by Dana Scott (see [10] p. 152). We will do so for the second part.

theorem *T2*: $\lfloor \mathcal{P}\ G \rfloor$ **using** *A3implT2-global A3* **by** *simp*

Theorem 11.17 (Informal Proposition 3) - Possibly God exists:

theorem *T3*: $\lfloor \Diamond \exists^E\ G \rfloor$ **using** *T1 T2* **by** *simp*

4.2 Part II - God's Existence is Necessary if Possible

We show here that God's necessary existence follows from its possible existence by adding some additional (potentially controversial) assumptions including, among others, an *essentialist* premise and the emphS5 axioms. Further results like monotheism and the rejection of free will (*modal collapse*) are also proved. A more detailed analysis of these rather philosophical issues is foreseen as follow-up work.

4.2.1 General Definitions

abbreviation *existencePredicate*::$\uparrow\langle\mathbf{0}\rangle$ $(E!)$ **where**
$E!\ x\ \equiv (\lambda w.\ (\exists^E y.\ y{\approx}x)\ w)$
consts *positiveProperty*::$\uparrow\langle\uparrow\langle\mathbf{0}\rangle\rangle$ $(\mathcal{P})$

abbreviation *God*::$\uparrow\langle\mathbf{0}\rangle$ (G) **where** $G \equiv (\lambda x.\ \forall Y.\ \mathcal{P}\ Y \rightarrow Y\ x)$

abbreviation *God-star*::$\uparrow\langle\mathbf{0}\rangle$ ($G*$) **where**
$G* \equiv (\lambda x.\ \forall Y.\ \mathcal{P}\ Y \leftrightarrow Y\ x)$
abbreviation *Entailment*::$\uparrow\langle\uparrow\langle\mathbf{0}\rangle,\uparrow\langle\mathbf{0}\rangle\rangle$ (**infix** $\Rrightarrow$ *60*) **where**
$X \Rrightarrow Y \equiv \Box(\forall^E z.\ X\ z \rightarrow Y\ z)$

4.2.2 Results from Part I

Note that the only use Gödel makes of axiom A3 is to show that being Godlike is a positive property (*T2*). We follow therefore Scott's proposal and take (*T2*) directly as an axiom:

axiomatization where
A1a:$\lfloor\forall X.\ \mathcal{P}\ (\rightharpoondown X) \rightarrow \neg(\mathcal{P}\ X)\ \rfloor$ **and** — axiom 11.3A
A1b:$\lfloor\forall X.\ \neg(\mathcal{P}\ X) \rightarrow \mathcal{P}\ (\rightharpoondown X)\rfloor$ **and** — axiom 11.3B
A2: $\lfloor\forall X\ Y.\ (\mathcal{P}\ X \wedge (X \Rrightarrow Y)) \rightarrow \mathcal{P}\ Y\rfloor$ **and** — axiom 11.5
T2: $\lfloor\mathcal{P}\ G\rfloor$ — proposition 11.16

lemma *True* **nitpick**[*satisfy*] **oops** — model found: axioms are consistent

lemma $\lfloor D\rfloor$ **using** *A1a A1b A2* **by** *blast* — axioms already imply D axiom

lemma *GodDefsAreEquivalent*: $\lfloor\forall x.\ G\ x \leftrightarrow G*\ x\rfloor$ **using** *A1b* **by** *fastforce*

theorem *T1*: $\lfloor\forall X.\ \mathcal{P}\ X \rightarrow \Diamond\exists^E\ X\rfloor$
using *A1a A2* **by** *blast* — positive properties are possibly instantiated
theorem *T3*: $\lfloor\Diamond\exists^E\ G\rfloor$ **using** *T1 T2* **by** *simp* — God exists possibly

4.2.3 Axioms

$\mathcal{P}$ satisfies the so-called stability conditions (see [10] p. 124), which means it designates rigidly (note that this makes for an *essentialist* assumption).

axiomatization where
A4a: $\lfloor\forall X.\ \mathcal{P}\ X \rightarrow \Box(\mathcal{P}\ X)\rfloor$ — axiom 11.11
lemma *A4b*: $\lfloor\forall X.\ \neg(\mathcal{P}\ X) \rightarrow \Box\neg(\mathcal{P}\ X)\rfloor$ **using** *A1a A1b A4a* **by** *blast*

abbreviation *rigidPred*::$('t{\Rightarrow}io){\Rightarrow}io$ **where**
rigidPred $\tau \equiv (\lambda\beta.\ \Box((\lambda z.\ \beta \approx z)\ {\downarrow}\tau))\ {\downarrow}\tau$

lemma $\lfloor$*rigidPred* $\mathcal{P}\rfloor$
using *A4a A4b* **by** *blast* — $\mathcal{P}$ is therefore rigid

lemma *True* **nitpick**[*satisfy*] **oops** — model found: so far all axioms A1-4 consistent

4.2.4 Theorems

abbreviation *essenceOf*::$\uparrow\langle\uparrow\langle\mathbf{0}\rangle,\mathbf{0}\rangle$ ($\mathcal{E}$) **where**
$\mathcal{E}\ Y\ x \equiv (Y\ x) \wedge (\forall Z.\ Z\ x \rightarrow Y \Rrightarrow Z)$
abbreviation *beingIdenticalTo*::$\mathbf{0}{\Rightarrow}\uparrow\langle\mathbf{0}\rangle$ (*id*) **where**

$id\ x\ \equiv (\lambda y.\ y{\approx}x)$ — note that id is a rigid predicate

Theorem 11.20 - Informal Proposition 5

theorem *GodIsEssential*: $\lfloor\forall x.\ G\ x \rightarrow (\mathcal{E}\ G\ x)\rfloor$ **using** *A1b A4a* **by** *metis*

Theorem 11.21

theorem $\lfloor\forall x.\ G{*}\ x \rightarrow (\mathcal{E}\ G{*}\ x)\rfloor$ **using** *A4a* **by** *meson*

Theorem 11.22 - Something can have only one essence:

theorem $\lfloor\forall X\ Y\ z.\ (\mathcal{E}\ X\ z \wedge \mathcal{E}\ Y\ z) \rightarrow (X \Rrightarrow Y)\rfloor$ **by** *meson*

Theorem 11.23 - An essence is a complete characterization of an individual:

theorem *EssencesCharacterizeCompletely*: $\lfloor\forall X\ y.\ \mathcal{E}\ X\ y \rightarrow (X \Rrightarrow (id\ y))\rfloor$
proof (*rule ccontr*)
  assume $\neg\ \lfloor\forall X\ y.\ \mathcal{E}\ X\ y \rightarrow (X \Rrightarrow (id\ y))\rfloor$
  hence $\exists w.\ \neg((\ \forall X\ y.\ \mathcal{E}\ X\ y \rightarrow X \Rrightarrow id\ y)\ w)$ **by** *simp*
  then obtain w **where** $\neg((\ \forall X\ y.\ \mathcal{E}\ X\ y \rightarrow X \Rrightarrow id\ y)\ w)$ **..**
  hence $(\exists X\ y.\ \mathcal{E}\ X\ y \wedge \neg(X \Rrightarrow id\ y))\ w$ **by** *simp*
  hence $\exists X\ y.\ \mathcal{E}\ X\ y\ w \wedge (\neg(X \Rrightarrow id\ y))\ w$ **by** *simp*
  then obtain P **where** $\exists y.\ \mathcal{E}\ P\ y\ w \wedge (\neg(P \Rrightarrow id\ y))\ w$ **..**
  then obtain a **where** 1: $\mathcal{E}\ P\ a\ w \wedge (\neg(P \Rrightarrow id\ a))\ w$ **..**
  hence 2: $\mathcal{E}\ P\ a\ w$ **by** (*rule conjunct1*)
  from 1 **have** $(\neg(P \Rrightarrow id\ a))\ w$ **by** (*rule conjunct2*)
  hence $\exists x.\ \exists z.\ w\ r\ x \wedge\ existsAt\ z\ x \wedge P\ z\ x \wedge \neg(a = z)$ **by** *blast*
  then obtain $w1$ **where** $\exists z.\ w\ r\ w1 \wedge\ existsAt\ z\ w1\ \wedge P\ z\ w1\ \wedge \neg(a = z)$ **..**
  then obtain b **where** 3: $w\ r\ w1 \wedge\ existsAt\ b\ w1\ \wedge P\ b\ w1\ \wedge \neg(a = b)$ **..**
  hence $w\ r\ w1$ **by** *simp*
  from 3 **have** $existsAt\ b\ w1$ **by** *simp*
  from 3 **have** $P\ b\ w1$ **by** *simp*
  from 3 **have** 4: $\neg(a = b)$ **by** *simp*
  from 2 **have** $P\ a\ w$ **by** *simp*
  from 2 **have** $\forall Y.\ Y\ a\ w \longrightarrow ((P \Rrightarrow Y)\ w)$ **by** *auto*
  hence $(\dot\neg(id\ b))\ a\ w \longrightarrow (P \Rrightarrow (\dot\neg(id\ b)))\ w$ **by** (*rule allE*)
  hence $\neg(\dot\neg(id\ b))\ a\ w \vee ((P \Rrightarrow (\dot\neg(id\ b)))\ w)$ **by** *blast*
  then show *False* **proof**
    assume $\neg(\dot\neg(id\ b))\ a\ w$
    hence $a = b$ **by** *simp*
    thus *False* **using** 4 **by** *auto*
    next
    assume $((P \Rrightarrow (\dot\neg(id\ b)))\ w)$
    hence $\forall x.\ \forall z.\ (w\ r\ x \wedge existsAt\ z\ x \wedge P\ z\ x) \longrightarrow (\dot\neg(id\ b))\ z\ x$ **by** *blast*
    hence $\forall z.\ (w\ r\ w1\ \wedge existsAt\ z\ w1\ \wedge P\ z\ w1) \longrightarrow (\dot\neg(id\ b))\ z\ w1$
      by (*rule allE*)
    hence $(w\ r\ w1\ \wedge existsAt\ b\ w1\ \wedge P\ b\ w1) \longrightarrow (\dot\neg(id\ b))\ b\ w1$ **by** (*rule allE*)
    hence $\neg(w\ r\ w1\ \wedge existsAt\ b\ w1\ \wedge P\ b\ w1) \vee (\dot\neg(id\ b))\ b\ w1$ **by** *simp*
    hence $(\dot\neg(id\ b))\ b\ w$ **using** 3 **by** *simp*
    hence $\neg(b{=}b)$ **by** *simp*
    thus *False* **by** *simp*

  qed
qed

Definition 11.24 - Necessary Existence (Informal Definition 6):

abbreviation *necessaryExistencePred*::$\uparrow\langle\mathbf{0}\rangle$ (*NE*)
  where $NE\ x \ \equiv (\lambda w.\ (\forall\, Y.\ \ \mathcal{E}\ Y\, x \rightarrow \Box\exists^{E}\ Y)\ w)$

Axiom 11.25 (Informal Axiom 5)

axiomatization where
 A5: $\lfloor\mathcal{P}\ NE\rfloor$

lemma *True* **nitpick**[*satisfy*] **oops** — model found: so far all axioms consistent

Theorem 11.26 (Informal Proposition 7) - Possibilist existence of God implies necessary actualist existence:

theorem *GodExistenceImpliesNecExistence*: $\lfloor\exists\ \ G \rightarrow\ \ \Box\exists^{E}\ G\rfloor$
proof –
{
  fix w
  {
    assume $\exists x.\ G\ x\ w$
    then obtain g **where** *1*: $G\ g\ w$ **..**
    hence $NE\ g\ w$ **using** *A5* **by** *auto* — axiom 11.25
    hence $\forall\, Y.\ (\mathcal{E}\ \ Y\ g\ w)\ \longrightarrow (\Box\exists^{E}\ Y)\ w$ **by** *simp*
    hence *2*: $(\mathcal{E}\ \ G\ g\ w)\ \longrightarrow (\Box\exists^{E}\ G)\ w$ **by** (*rule allE*)
    have $(\forall x.\ G\ x \rightarrow (\mathcal{E}\ \ G\ x))\ w$ **using** *GodIsEssential*
      by (*rule allE*) — GodIsEssential follows from Axioms 11.11 and 11.3B
    hence $(G\ g \rightarrow (\mathcal{E}\ \ G\ g))\ w$ **by** (*rule allE*)
    hence $G\ g\ w \longrightarrow \mathcal{E}\ \ G\ g\ w$ **by** *simp*
    from *this 1* **have** *3*: $\mathcal{E}\ \ G\ g\ w$ **by** (*rule mp*)
    from *2 3* **have** $(\Box\exists^{E}\ G)\ w$ **by** (*rule mp*)
  }
  hence $(\exists x.\ G\ x\ w) \longrightarrow (\Box\exists^{E}\ G)\ w$ **by** (*rule impI*)
  hence $((\exists x.\ G\ x) \rightarrow\ \ \Box\exists^{E}\ G)\ w$ **by** *simp*
}
 thus *?thesis* **by** (*rule allI*)
qed

Modal collapse is countersatisfiable (unless we introduce S5 axioms):

lemma $\lfloor\forall\Phi.(\Phi \rightarrow (\Box\ \Phi))\rfloor$ **nitpick oops**

Below we axiomatize semantic frame conditions for some modal logics. Taken all together they make for an equivalence relation and therefore an *S5* logic (via *Sahlqvist correspondence*). We prefer to introduce them individually in order to get more detailed information about their relevance.

axiomatization where
 refl: *reflexive aRel* **and**
 tran: *transitive aRel* **and**

symm: *symmetric aRel*

lemma *True* **nitpick**[*satisfy*] **oops** — model found: axioms still consistent

Using an *S5* logic, *modal collapse* ($\lfloor\forall\,\Phi.(\Phi \rightarrow (\Box\ \Phi))\rfloor$) is actually valid (see 'More Objections' some pages below)

We prove some useful inference rules:

lemma *modal-distr*: $\lfloor\Box(\varphi \rightarrow \psi)\rfloor \Longrightarrow \lfloor(\Diamond\varphi \rightarrow \Diamond\psi)\rfloor$ **by** *blast*
lemma *modal-trans*: $(\lfloor\varphi \rightarrow \psi\rfloor \wedge \lfloor\psi \rightarrow \chi\rfloor) \Longrightarrow \lfloor\varphi \rightarrow \chi\rfloor$ **by** *simp*

Theorem 11.27 - Informal Proposition 8. Note that only symmetry and transitivity for the accessibility relation are needed. Nevertheless we already rely on an S5 logic here, since our axioms imply D and therefore seriality (via *Sahlqvist correspondence*).

theorem *possExistenceImpliesNecEx*: $\lfloor\Diamond\exists\ \ G \rightarrow \Box\exists^{\,E}\ G\rfloor$ — local consequence
proof –
  have $\lfloor\exists\ \ G \rightarrow \Box\exists^{\,E}\ G\rfloor$ **using** *GodExistenceImpliesNecExistence*
    by *simp* — follows from Axioms 11.11, 11.25 and 11.3B
  hence $\lfloor\Box(\exists\ \ G \rightarrow \Box\exists^{\,E}\ G)\rfloor$ **using** *NEC* **by** *simp*
  hence *1*: $\lfloor\Diamond\exists\ \ G \rightarrow \Diamond\Box\exists^{\,E}\ G\rfloor$ **by** (*rule modal-distr*)
  have *2*: $\lfloor\Diamond\Box\exists^{\,E}\ G \rightarrow \Box\exists^{\,E}\ G\rfloor$ **using** *symm tran* **by** *metis* — frame conditions
  from *1 2* **have** $\lfloor\Diamond\exists\ \ G \rightarrow \Diamond\Box\exists^{\,E}\ G\rfloor \wedge \lfloor\Diamond\Box\exists^{\,E}\ G \rightarrow \Box\exists^{\,E}\ G\rfloor$ **by** *simp*
  thus *?thesis* **by** (*rule modal-trans*)
qed

lemma *T4*: $\lfloor\Diamond\exists\ \ G\rfloor \longrightarrow \lfloor\Box\exists^{\,E}\ G\rfloor$ **using** *possExistenceImpliesNecEx*
    by (*rule localImpGlobalCons*) — global consequence

Corollary 11.28 - Necessary (actualist) existence of God (for both definitions):

lemma *GodNecExists*: $\lfloor\Box\exists^{\,E}\ G\rfloor$ **using** *T3 T4* **by** *metis*
lemma *God-starNecExists*: $\lfloor\Box\exists^{\,E}\ G*\rfloor$
  using *GodNecExists GodDefsAreEquivalent* **by** *simp*

4.2.5 Monotheism

Monotheism for non-normal models (with Leibniz equality) follows directly from God having all and only positive properties:

theorem *Monotheism-LeibnizEq*: $\lfloor\forall\,x.\ G\ x \rightarrow (\forall\,y.\ G\ y \rightarrow (x\ \approx^{L}\ y))\rfloor$
  using *GodDefsAreEquivalent* **by** *simp*

Monotheism for normal models is trickier. We need to consider some previous results (p. 162):

lemma *GodExistenceIsValid*: $\lfloor\exists^{\,E}\ G\rfloor$ **using** *GodNecExists refl*
  by *auto* — frame reflexivity is explicitly required by the solver

Proposition 11.29:

theorem *Monotheism-normalModel*: $\lfloor \exists x. \forall y.\ G\ y \leftrightarrow x \approx y \rfloor$
proof –
{
fix w
have $\lfloor \exists^E\ G \rfloor$ **using** *GodExistenceIsValid* **by** *simp* — follows from corollary 11.28

hence $(\exists^E\ G)\ w$ **by** (*rule allE*)
then obtain g **where** *1*: $existsAt\ g\ w \wedge G\ g\ w$ **..**
hence *2*: $\mathcal{E}\ G\ g\ w$ **using** *GodIsEssential* **by** *blast* — follows from ax. 11.11/11.3B

{
fix y
have $G\ y\ w \longleftrightarrow (g \approx y)\ w$ **proof**
assume $G\ y\ w$
hence *3*: $\mathcal{E}\ G\ y\ w$ **using** *GodIsEssential* **by** *blast*
have $(\mathcal{E}\ G\ y \rightarrow (G \Rrightarrow id\ y))\ w$ **using** *EssencesCharacterizeCompletely*
by *simp* — follows from theorem 11.23
hence $\mathcal{E}\ G\ y\ w \longrightarrow ((G \Rrightarrow id\ y)\ w)$ **by** *simp*
from *this 3* **have** $(G \Rrightarrow id\ y)\ w$ **by** (*rule mp*)
hence $(\Box(\forall^E z.\ G\ z \rightarrow z \approx y))\ w$ **by** *simp*
hence $\forall x.\ w\ r\ x \longrightarrow ((\forall z.\ (existsAt\ z\ x \wedge G\ z\ x) \longrightarrow z = y))$ **by** *auto*
hence $w\ r\ w \longrightarrow ((\forall z.\ (existsAt\ z\ w \wedge G\ z\ w) \longrightarrow z = y))$ **by** (*rule allE*)
hence $\forall z.\ (w\ r\ w \wedge existsAt\ z\ w \wedge G\ z\ w) \longrightarrow z = y$ **by** *auto*
hence *4*: $(w\ r\ w \wedge existsAt\ g\ w \wedge G\ g\ w) \longrightarrow g = y$ **by** (*rule allE*)
have $w\ r\ w$ **using** *refl*
by *simp* — using frame reflexivity (Axiom M)
hence $w\ r\ w \wedge (existsAt\ g\ w \wedge G\ g\ w)$ **using** *1* **by** (*rule conjI*)
from *4 this* **have** $g = y$ **by** (*rule mp*)
thus $(g \approx y)\ w$ **by** *simp*
next
assume $(g \approx y)\ w$
from *this 2* **have** $\mathcal{E}\ G\ y\ w$ **by** *simp*
thus $G\ y\ w$ **by** (*rule conjunct1*)
qed
}
hence $\forall y.\ G\ y\ w \longleftrightarrow (g \approx y)\ w$ **by** (*rule allI*)
hence $\exists x.\ (\forall y.\ G\ y\ w \longleftrightarrow (x \approx y)\ w)$ **by** (*rule exI*)
hence $(\exists x.\ (\forall y.\ G\ y \leftrightarrow (x \approx y)))\ w$ **by** *simp*
}
thus *?thesis* **by** (*rule allI*)
qed

Corollary 11.30:

lemma *GodImpliesExistence*: $\lfloor \forall x.\ G\ x \rightarrow E!\ x \rfloor$
using *GodExistenceIsValid Monotheism-normalModel* **by** *metis*

4.2.6 Positive Properties are Necessarily Instantiated

lemma *PosPropertiesNecExist*: $\lfloor \forall Y.\ \mathcal{P}\ Y \rightarrow \Box\exists^E\ Y \rfloor$ **using** *GodNecExists A4a*

by *meson* — proposition 11.31: follows from corollary 11.28 and axiom A4a

4.2.7 More Objections

In this section Fitting discusses the objection raised by Sobel [17], who argues that Gödel's axiom system is so strong it implies that whatever is the case is so necessarily, i.e. the modal system collapses ($\varphi \longrightarrow \Box\varphi$). The *modal collapse* has been philosophically interpreted as implying the absence of free will.

We start by proving an useful FOL lemma:

lemma *useful*: $(\forall x.\ \varphi\ x \longrightarrow \psi) \Longrightarrow ((\exists x.\ \varphi\ x) \longrightarrow \psi)$ **by** *simp*

After introducing the S5 axioms *modal collapse* becomes valid (pp. 163-4):

lemma *ModalCollapse*: $\lfloor\forall\,\Phi.(\Phi \rightarrow (\Box\ \Phi))\rfloor$
proof –
 {
 fix w
  {
  fix Q
  have $(\forall x.\ G\ x \rightarrow (\mathcal{E}\ G\ x))\ w$ **using** *GodIsEssential*
   by (*rule allE*) — follows from Axioms 11.11 and 11.3B
  hence $\forall x.\ G\ x\ w \longrightarrow \mathcal{E}\ G\ x\ w$ **by** *simp*
  hence $\forall x.\ G\ x\ w \longrightarrow (\forall Z.\ Z\ x \rightarrow \Box(\forall^E z.\ G\ z \rightarrow Z\ z))\ w$ **by** *force*
  hence $\forall x.\ G\ x\ w \longrightarrow ((\lambda y.\ Q)\ x \rightarrow \Box(\forall^E z.\ G\ z \rightarrow (\lambda y.\ Q)\ z))\ w$ **by** *force*
  hence $\forall x.\ G\ x\ w \longrightarrow (Q \rightarrow \Box(\forall^E z.\ G\ z \rightarrow Q))\ w$ **by** *simp*
  hence *1*: $(\exists x.\ G\ x\ w) \longrightarrow ((Q \rightarrow \Box(\forall^E z.\ G\ z \rightarrow Q))\ w)$ **by** (*rule useful*)
  have $\exists x.\ G\ x\ w$ **using** *GodExistenceIsValid* **by** *auto*
  from *1 this* **have** $(Q \rightarrow \Box(\forall^E z.\ G\ z \rightarrow Q))\ w$ **by** (*rule mp*)
  hence $(Q \rightarrow \Box((\exists^E z.\ G\ z) \rightarrow Q))\ w$ **using** *useful* **by** *blast*
  hence $(Q \rightarrow (\Box(\exists^E z.\ G\ z) \rightarrow \Box Q))\ w$ **by** *simp*
  hence $(Q \rightarrow \Box Q)\ w$ **using** *GodNecExists* **by** *simp*
  }
 hence $(\forall\,\Phi.\ \Phi \rightarrow \Box\ \Phi)\ w$ **by** (*rule allI*)
 }
 thus *?thesis* **by** (*rule allI*)
qed

5 Fitting's Solution

In this section we consider Fitting's solution to the objections raised in his discussion of Gödel's Argument pp. 164-9, especially the problem of *modal collapse*, which has been metaphysically interpreted as implying a rejection of free will. Since we are generally commited to the existence of free will (in a pre-theoretical sense), such a result is philosophically unappealing and rather seen as a problem in the argument's formalisation.

This part of the book still leaves several details unspecified and the reader is thus compelled to fill in the gaps. As a result, we came across some premises and theorems allowing for different formalisations and therefore leading to disparate implications. Only some of those cases are shown here for illustrative purposes. The options we have chosen here are such that they validate the argument and we assume they correspond to Fitting's ideas.

5.1 General Definitions

The following technical definition is needed only for type correctness. It is used to convert extensional objects into rigid intensional ones:

abbreviation *trivialExpansion*::*bool*⇒*io* ((|-|)) **where** $(|\varphi|) \equiv \lambda w.\ \varphi$

The following is an existence predicate for our object-language. (We have previously shown it is equivalent to its meta-logical counterpart.)

abbreviation *existencePredicate*::$\uparrow\langle\mathbf{0}\rangle$ ($E!$) **where**
$E!\ x \ \equiv (\lambda w.\ (\exists^E y.\ y{\approx}x)\ w)$

consts *positiveProperty*::$\uparrow\langle\langle\mathbf{0}\rangle\rangle$ ($\mathcal{P}$)
abbreviation *God*::$\uparrow\langle\mathbf{0}\rangle$ (G) **where** $G \equiv (\lambda x.\ \forall\, Y.\ \mathcal{P}\ \ Y \rightarrow (|\,Y\ x|))$
abbreviation *God-star*::$\uparrow\langle\mathbf{0}\rangle$ ($G*$) **where** $G* \equiv (\lambda x.\ \forall\, Y.\ \mathcal{P}\ \ Y \leftrightarrow (|\,Y\ x|))$

abbreviation *Entailment*::$\uparrow\langle\langle\mathbf{0}\rangle,\langle\mathbf{0}\rangle\rangle$ (**infix** $\Rrightarrow$ *60*) **where**
$X \Rrightarrow Y \equiv \ \Box(\forall^E z.\ (|X\ z|) \rightarrow (|\,Y\ z|))$

5.2 Part I - God's Existence is Possible

axiomatization where
A1a:$\lfloor\forall\, X.\ \mathcal{P}\ (\rightharpoondown X) \rightarrow \neg(\mathcal{P}\ \ X)\ \rfloor$ **and** — axiom 11.3A
A1b:$\lfloor\forall\, X.\ \neg(\mathcal{P}\ \ X) \rightarrow \mathcal{P}\ \ (\rightharpoondown X)\rfloor$ **and** — axiom 11.3B
A2: $\lfloor\forall\, X\ Y.\ (\mathcal{P}\ \ X \wedge (X \Rrightarrow\ Y)) \rightarrow \mathcal{P}\ \ Y\rfloor$ **and** — axiom 11.5
T2: $\lfloor\mathcal{P}\ {\downarrow}G\rfloor$ — proposition 11.16 (modified)

lemma *True* **nitpick**[*satisfy*] **oops** — model found: axioms are consistent

lemma *GodDefsAreEquivalent*: $\lfloor\forall\, x.\ G\ x \leftrightarrow G*\ x\rfloor$ **using** *A1b* **by** *fastforce*

T1 (Positive properties are possibly instantiated) can be formalised in two different ways:

theorem *T1a*: $\lfloor\forall X{::}\langle\mathbf{0}\rangle.\ \mathcal{P}\ X \rightarrow \Diamond(\exists^{E} z.\ (\!|X\ z|\!))\rfloor$
using *A1a A2* **by** *blast* — this is the one used in the book
theorem *T1b*: $\lfloor\forall X{::}\uparrow\langle\mathbf{0}\rangle.\ \mathcal{P}\ \downarrow X \rightarrow \Diamond(\exists^{E} z.\ X\ z)\rfloor$
nitpick oops — this one is also possible but not valid so we won't use it

Some interesting (non-)equivalences:

lemma $\lfloor\Box\exists^{E}\ (Q{::}\uparrow\langle\mathbf{0}\rangle) \leftrightarrow \Box(\exists^{E} \downarrow Q)\rfloor$ **by** *simp*
lemma $\lfloor\Box\exists^{E}\ (Q{::}\uparrow\langle\mathbf{0}\rangle) \leftrightarrow ((\lambda X.\ \Box\exists^{E}\ X)\ Q)\rfloor$ **by** *simp*
lemma $\lfloor\Box\exists^{E}\ (Q{::}\uparrow\langle\mathbf{0}\rangle) \leftrightarrow ((\lambda X.\ \Box\exists^{E} \downarrow X)\ Q)\rfloor$ **by** *simp*
lemma $\lfloor\Box\exists^{E}\ (Q{::}\uparrow\langle\mathbf{0}\rangle) \leftrightarrow ((\lambda X.\ \Box\exists^{E}\ X) \downarrow Q)\rfloor$ **nitpick oops** — not equivalent!

T3 (God exists possibly) can be formalised in two different ways, using a *de re* or a *de dicto* reading.

theorem *T3-deRe*: $\lfloor(\lambda X.\ \Diamond\exists^{E}\ X) \downarrow G\rfloor$ **using** *T1a T2* **by** *simp*
theorem *T3-deDicto*: $\lfloor\Diamond\exists^{E} \downarrow G\rfloor$ **nitpick oops** — countersatisfiable

From the last two theorems, we think *T3-deRe* should be the version originally implied in the book, since *T3-deDicto* is not valid (*T1b* were valid but it isn't)

lemma assumes *T1b*: $\lfloor\forall X.\ \mathcal{P} \downarrow X \rightarrow \Diamond(\exists^{E} z.\ X\ z)\rfloor$
shows *T3-deDicto*: $\lfloor\Diamond\exists^{E} \downarrow G\rfloor$ **using** *assms T2* **by** *simp*

5.3 Part II - God's Existence is Necessary if Possible

In this variant $\mathcal{P}$ also designates rigidly, as shown in the last section.

axiomatization where
A4a: $\lfloor\forall X.\ \mathcal{P}\ X \rightarrow \Box(\mathcal{P}\ X)\rfloor$ — axiom 11.11
lemma *A4b*: $\lfloor\forall X.\ \neg(\mathcal{P}\ X) \rightarrow \Box\neg(\mathcal{P}\ X)\rfloor$ **using** *A1a A1b A4a* **by** *blast*

lemma *True* **nitpick**[*satisfy*] **oops** — model found: so far all axioms consistent

abbreviation *essenceOf*::$\uparrow\langle\langle\mathbf{0}\rangle,\mathbf{0}\rangle$ ($\mathcal{E}$) **where**
$\mathcal{E}\ Y\ x \equiv (\!|Y\ x|\!) \wedge (\forall Z{::}\langle\mathbf{0}\rangle.\ (\!|Z\ x|\!) \rightarrow Y \Rrightarrow Z)$

Theorem 11.20 - Informal Proposition 5

theorem *GodIsEssential*: $\lfloor\forall x.\ G\ x \rightarrow ((\mathcal{E} \downarrow_1 G)\ x)\rfloor$ **using** *A1b* **by** *metis*

Theorem 11.21

theorem *God-starIsEssential*: $\lfloor\forall x.\ G{*}\ x \rightarrow ((\mathcal{E} \downarrow_1 G{*})\ x)\rfloor$ **by** *meson*

abbreviation *necExistencePred*:: $\uparrow\langle\mathbf{0}\rangle$ (*NE*) **where**
$NE\ x \equiv \lambda w.\ (\forall Y.\ \mathcal{E}\ Y\ x \rightarrow \Box(\exists^{E} z.\ (\!|Y\ z|\!)))\ w$

Informal Axiom 5

axiomatization where
 A5: $\lfloor \mathcal{P} \downarrow NE \rfloor$

lemma *True* **nitpick**[*satisfy*] **oops** — model found: so far all axioms consistent

Reminder: We use the down-arrow notation because it is more explicit. See (non-)equivalences above.

lemma $\lfloor \exists\ G \leftrightarrow \exists \downarrow G \rfloor$ **by** *simp*
lemma $\lfloor \exists^E\ G \leftrightarrow \exists^E \downarrow G \rfloor$ **by** *simp*
lemma $\lfloor \Box\exists^E\ G \leftrightarrow\ \Box\exists^E \downarrow G \rfloor$ **by** *simp*

Theorem 11.26 (Informal Proposition 7) - (possibilist) existence of God implies necessary (actualist) existence.

There are two different ways of formalising this theorem. Both of them are proven valid:

First version:

theorem *GodExImpliesNecEx-v1*: $\lfloor \exists \downarrow G \rightarrow\ \Box\exists^E \downarrow G \rfloor$
proof –
{
  fix w
  {
    assume $\exists x.\ G\ x\ w$
    then obtain g **where** *1*: $G\ g\ w$ **..**
    hence $NE\ g\ w$ **using** *A5* **by** *auto*
    hence $\forall Y.\ (\mathcal{E}\ \ Y\ g\ w) \longrightarrow (\Box(\exists^E z.\ (\!|Y\ z|\!)))\ w$ **by** *simp*
    hence $(\mathcal{E}\ (\lambda x.\ G\ x\ w)\ g\ w) \longrightarrow (\Box(\exists^E z.\ (\!|(\lambda x.\ G\ x\ w)\ z|\!)))\ w$ **by** (*rule allE*)
    hence *2*: $((\mathcal{E} \downarrow_1 G)\ g\ w) \longrightarrow (\Box(\exists^E\ G))\ w$ **using** *A4b* **by** *meson*
    have $(\forall x.\ G\ x \rightarrow ((\mathcal{E} \downarrow_1 G)\ x))\ w$ **using** *GodIsEssential* **by** (*rule allE*)
    hence $(G\ g \rightarrow ((\mathcal{E} \downarrow_1 G)\ g))\ w$ **by** (*rule allE*)
    hence $G\ g\ w \longrightarrow (\mathcal{E} \downarrow_1 G)\ g\ w$ **by** *simp*
    from *this 1* **have** *3*: $(\mathcal{E} \downarrow_1 G)\ g\ w$ **by** (*rule mp*)
    from *2 3* **have** $(\Box\exists^E\ G)\ w$ **by** (*rule mp*)
  }
  hence $(\exists x.\ G\ x\ w) \longrightarrow (\Box\exists^E\ G)\ w$ **by** (*rule impI*)
  hence $((\exists x.\ G\ x) \rightarrow\ \Box\exists^E\ G)\ w$ **by** *simp*
}
 thus *?thesis* **by** (*rule allI*)
qed

Second version (which can be proven directly by automated tools using last version):

theorem *GodExImpliesNecEx-v2*: $\lfloor \exists \downarrow G \rightarrow ((\lambda X.\ \Box\exists^E\ X) \downarrow G)\rfloor$
  using *A4a GodExImpliesNecEx-v1* **by** *metis*

In contrast to Gödel's argument (as presented by Fitting), the following theorems can be proven in *K* logic (note that the *S5* axioms are no longer needed):

Theorem 11.27 - Informal Proposition 8

theorem *possExImpliesNecEx-v1*: $\lfloor \Diamond\exists \downarrow G \rightarrow \Box\exists^{E} \downarrow G \rfloor$
using *GodExImpliesNecEx-v1 T3-deRe* **by** *metis*
theorem *possExImpliesNecEx-v2*: $\lfloor (\lambda X.\ \Diamond\exists^{E}\ X) \downarrow G \rightarrow ((\lambda X.\ \Box\exists^{E}\ X) \downarrow G) \rfloor$
using *GodExImpliesNecEx-v2* **by** *blast*

Corollaries:

lemma *T4-v1*: $\lfloor \Diamond\exists \downarrow G \rfloor \longrightarrow \lfloor \Box\exists^{E} \downarrow G \rfloor$
using *possExImpliesNecEx-v1* **by** *simp*
lemma *T4-v2*: $\lfloor (\lambda X.\ \Diamond\exists^{E}\ X) \downarrow G \rfloor \longrightarrow \lfloor (\lambda X.\ \Box\exists^{E}\ X) \downarrow G \rfloor$
using *possExImpliesNecEx-v2* **by** *simp*

5.4 Conclusion (*De Re* and *De Dicto* Reading)

Version I - Necessary Existence of God (*de dicto*):

lemma *GodNecExists-v1*: $\lfloor \Box\exists^{E} \downarrow G \rfloor$
using *GodExImpliesNecEx-v1 T3-deRe* **by** *fastforce* — corollary 11.28
lemma *God-starNecExists-v1*: $\lfloor \Box\exists^{E} \downarrow G* \rfloor$
using *GodNecExists-v1 GodDefsAreEquivalent* **by** *simp*
lemma $\lfloor \Box(\lambda X.\ \exists^{E}\ X) \downarrow G* \rfloor$
using *God-starNecExists-v1* **by** *simp* — *de dicto* shown here explicitly

Version II - Necessary Existence of God (*de re*)

lemma *GodNecExists-v2*: $\lfloor (\lambda X.\ \Box\exists^{E}\ X) \downarrow G \rfloor$
using *T3-deRe T4-v2* **by** *blast*
lemma *God-starNecExists-v2*: $\lfloor (\lambda X.\ \Box\exists^{E}\ X) \downarrow G* \rfloor$
using *GodNecExists-v2 GodDefsAreEquivalent* **by** *simp*

5.5 Modal Collapse

Modal collapse is countersatisfiable even in *S5*. Note that countermodels with a cardinality of one for the domain of individuals are found by *Nitpick* (the countermodel shown in the book has cardinality of two).

lemma $\lfloor \forall\, \Phi.(\Phi \rightarrow (\Box\ \Phi)) \rfloor$
nitpick[*card 't=1*, *card i=2*] **oops** — countermodel found in *K*

axiomatization where
S5: *equivalence aRel* — assume *S5* logic

lemma $\lfloor \forall\, \Phi.(\Phi \rightarrow (\Box\ \Phi)) \rfloor$
nitpick[*card 't=1*, *card i=2*] **oops** — countermodel also found in *S5*

6 Anderson's Alternative

In this last section we consider Anderson's Alternative to the objections previously shown, as exposed in the last part of the textbook (pp. 169-171)

6.1 General Definitions

abbreviation *existencePredicate*::$\uparrow\langle\mathbf{0}\rangle$ ($E!$)
  where $E!\ x\ \equiv \lambda w.\ (\exists^E y.\ y{\approx}x)\ w$

consts *positiveProperty*::$\uparrow\langle\uparrow\langle\mathbf{0}\rangle\rangle$ ($\mathcal{P}$)

abbreviation *God*::$\uparrow\langle\mathbf{0}\rangle$ (G^A) **where** $G^A \equiv \lambda x.\ \forall\, Y.\ (\mathcal{P}\ \ Y) \leftrightarrow \Box(Y\, x)$

abbreviation *Entailment*::$\uparrow\langle\uparrow\langle\mathbf{0}\rangle,\uparrow\langle\mathbf{0}\rangle\rangle$ (**infix** $\Rrightarrow$ *60*) **where**
  $X \Rrightarrow\ Y \equiv\ \Box(\forall^E z.\ X\ z \rightarrow\ Y\ z)$

6.2 Part I - God's Existence is Possible

axiomatization where
  A1a:$\lfloor\forall\, X.\ \mathcal{P}\ (\rightharpoondown X) \rightarrow \neg(\mathcal{P}\ X)\ \rfloor$ **and** — Axiom 11.3A
  A2: $\lfloor\forall\, X\ Y.\ (\mathcal{P}\ X \wedge (X \Rrightarrow\ Y)) \rightarrow \mathcal{P}\ Y\rfloor$ **and** — Axiom 11.5
  T2: $\lfloor\mathcal{P}\ G^A\rfloor$ — Proposition 11.16

lemma *True* **nitpick**[*satisfy*] **oops** — model found: axioms are consistent

theorem *T1*: $\lfloor\forall\, X.\ \mathcal{P}\ X \rightarrow \Diamond\exists^E\ X\rfloor$
  using *A1a A2* **by** *blast* — positive properties are possibly instantiated
theorem *T3*: $\lfloor\Diamond\exists^E\ G^A\rfloor$ **using** *T1 T2* **by** *simp* — God exists possibly

6.3 Part II - God's Existence is Necessary if Possible

$\mathcal{P}$ now satisfies only one of the stability conditions. But since the argument uses an *S5* logic, the other stability condition is implied. Therefore $\mathcal{P}$ becomes rigid (see p. 124 in textbook).

axiomatization where
  A4a: $\lfloor\forall\, X.\ \mathcal{P}\ X \rightarrow \Box(\mathcal{P}\ X)\rfloor$ — axiom 11.11

We again postulate our *S5* axioms:

axiomatization where
 refl: *reflexive aRel* **and**
 tran: *transitive aRel* **and**
 symm: *symmetric aRel*

lemma *True* **nitpick**[*satisfy*] **oops** — model found: so far all axioms consistent

abbreviation *rigidPred*::$('t{\Rightarrow}io){\Rightarrow}io$ **where**
 $rigidPred\ \tau \equiv (\lambda\beta.\ \Box((\lambda z.\ \beta \approx z)\ \downarrow\!\tau))\ \downarrow\!\tau$

lemma *A4b*: $\lfloor \forall X.\ \neg(\mathcal{P}\ X) \rightarrow \Box\neg(\mathcal{P}\ X)\rfloor$
  using *A4a symm* **by** *auto* — note only symmetry is needed (B axiom)
lemma $\lfloor$*rigidPred* $\mathcal{P}\rfloor$
  using *A4a A4b* **by** *blast* — $\mathcal{P}$ is therefore rigid in a B logic

Essence, Anderson Version (Definition 11.34)

abbreviation *essenceOf*::$\uparrow\langle\uparrow\langle\mathbf{0}\rangle,\mathbf{0}\rangle$ ($\mathcal{E}^A$) **where**
  $\mathcal{E}^A\ \ Y\, x \equiv (\forall Z.\ \Box(Z\ x) \leftrightarrow Y \Rrightarrow Z)$

Necessary Existence, Anderson Version (Definition 11.35)

abbreviation *necessaryExistencePred*::$\uparrow\langle\mathbf{0}\rangle$ (NE^A)
  where $NE^A\ x\ \equiv (\lambda w.\ (\forall Y.\ \ \mathcal{E}^A\ \ Y\, x \rightarrow \Box\exists^E\ Y)\ w)$

Theorem 11.36 - If g is God-like, then the property of being God-like is the essence of g.

As shown before, this theorem's proof could be completely automatized for Gödel's and Fitting's variants. For Anderson's version however, we had to provide Isabelle with some help based on the corresponding natural-language proof given by Anderson (see [2] Theorem 2*, p. 296)

theorem *GodIsEssential*: $\lfloor\forall x.\ G^A\ x \rightarrow (\mathcal{E}^A\ \ G^A\ x)\rfloor$
proof –
{
  fix w
  {
    fix g
    {
      assume $G^A\ g\ w$
      hence *1*: $\forall Y.\ (\mathcal{P}\ \ Y\ w) \longleftrightarrow (\Box(Y\ g))\ w$ **by** *simp*
      {
        fix Q
        from *1* **have** *2*: $(\mathcal{P}\ \ Q\ w) \longleftrightarrow (\Box(Q\ g))\ w$ **by** (*rule allE*)
        have $(\Box(Q\ g))\ w \longleftrightarrow (G^A \Rrightarrow Q)\ w$ — we need to prove $\rightarrow$ and $\leftarrow$
        proof
            assume $(\Box(Q\ g))\ w$ — suppose g is God-like and necessarily has Q
            hence *3*: $(\mathcal{P}\ \ Q\ w)$ **using** *2* **by** *simp* — then Q is positive

            {
              fix u
              have $(\mathcal{P}\ \ Q\ u) \longrightarrow (\forall x.\ G^A\ x\ u \longrightarrow (\Box(Q\ x))\ u)$
                by *auto* — using the definition of God-like
              have $(\mathcal{P}\ \ Q\ u) \longrightarrow (\forall x.\ G^A\ x\ u \longrightarrow ((Q\ x))\ u)$
                using *refl* **by** *auto* — and using $\Box(\varphi\ x) \longrightarrow \varphi\ x$
            }
            hence $\forall z.\ (\mathcal{P}\ \ Q\ z) \longrightarrow (\forall x.\ G^A\ x\ z \longrightarrow Q\ x\ z)$ **by** (*rule allI*)
            hence $\lfloor\mathcal{P}\ \ Q \rightarrow (\forall x.\ G^A\ x \rightarrow Q\ x)\rfloor$
              by *auto* — if Q is positive, then whatever is God-like has Q

hence $\lfloor\Box(\mathcal{P}\ Q \to (\forall x.\ G^A\ x \to Q\ x))\rfloor$ **by** (*rule NEC*)

hence $\lfloor(\Box(\mathcal{P}\ Q)) \to \Box(\forall x.\ G^A\ x \to Q\ x)\rfloor$ **using** K **by** *auto*
hence $\lfloor(\Box(\mathcal{P}\ Q)) \to G^A \Rrightarrow Q\rfloor$ **by** *simp*
hence $((\Box(\mathcal{P}\ Q)) \to G^A \Rrightarrow Q)\ w$ **by** (*rule allE*)
hence *4*: $(\Box(\mathcal{P}\ Q))\ w \longrightarrow (G^A \Rrightarrow Q)\ w$ **by** *simp*
have $\lfloor\forall X.\ \mathcal{P}\ X \to \Box(\mathcal{P}\ X)\rfloor$ **by** (*rule A4a*) — using axiom 4
hence $(\forall X.\ \mathcal{P}\ X \to (\Box(\mathcal{P}\ X)))\ w$ **by** (*rule allE*)
hence $\mathcal{P}\ Q\ w \longrightarrow (\Box(\mathcal{P}\ Q))\ w$ **by** (*rule allE*)
hence $\mathcal{P}\ Q\ w \longrightarrow (G^A \Rrightarrow Q)\ w$ **using** *4* **by** *simp*
thus $(G^A \Rrightarrow Q)\ w$ **using** *3* **by** (*rule mp*) — $\to$ direction
next
assume *5*: $(G^A \Rrightarrow Q)\ w$ — suppose Q is entailed by being God-like
have $\lfloor\forall X\ Y.\ (\mathcal{P}\ X \wedge (X \Rrightarrow Y)) \to \mathcal{P}\ Y\rfloor$ **by** (*rule A2*)
hence $(\forall X\ Y.\ (\mathcal{P}\ X \wedge (X \Rrightarrow Y)) \to \mathcal{P}\ Y)\ w$ **by** (*rule allE*)
hence $\forall X\ Y.\ (\mathcal{P}\ X\ w \wedge (X \Rrightarrow Y)\ w) \longrightarrow \mathcal{P}\ Y\ w$ **by** *simp*
hence $\forall Y.\ (\mathcal{P}\ G^A\ w \wedge (G^A \Rrightarrow Y)\ w) \longrightarrow \mathcal{P}\ Y\ w$ **by** (*rule allE*)
hence *6*: $(\mathcal{P}\ G^A\ w \wedge (G^A \Rrightarrow Q)\ w) \longrightarrow \mathcal{P}\ Q\ w$ **by** (*rule allE*)
have $\lfloor\mathcal{P}\ G^A\rfloor$ **by** (*rule T2*)
hence $\mathcal{P}\ G^A\ w$ **by** (*rule allE*)
hence $\mathcal{P}\ G^A\ w \wedge (G^A \Rrightarrow Q)\ w$ **using** *5* **by** (*rule conjI*)
from *6 this* **have** $\mathcal{P}\ Q\ w$ **by** (*rule mp*) — Q is positive by A2 and T2
thus $(\Box(Q\ g))\ w$ **using** *2* **by** *simp*
qed
}
hence $\forall Z.\ (\Box(Z\ g))\ w \longleftrightarrow (G^A \Rrightarrow Z)\ w$ **by** (*rule allI*)
hence $(\forall Z.\ \Box(Z\ g) \leftrightarrow\ G^A \Rrightarrow Z)\ w$ **by** *simp*
hence $\mathcal{E}^A\ G^A\ g\ w$ **by** *simp*
}
hence $G^A\ g\ w \longrightarrow \mathcal{E}^A\ G^A\ g\ w$ **by** (*rule impI*)
}
hence $\forall x.\ G^A\ x\ w \longrightarrow \mathcal{E}^A\ G^A\ x\ w$ **by** (*rule allI*)
}
thus *?thesis* **by** (*rule allI*)
qed

Axiom 11.37 (Anderson's Version of 11.25)

axiomatization where
A5: $\lfloor\mathcal{P}\ NE^A\rfloor$

lemma *True* **nitpick**[*satisfy*] **oops** — model found: so far all axioms consistent

Theorem 11.38 - Possibilist existence of God implies necessary actualist existence:

theorem *GodExistenceImpliesNecExistence*: $\lfloor\exists\ G^A \to\ \Box\exists^E\ G^A\rfloor$
proof –
{
fix w
{

assume $\exists\, x.\ G^A\ x\ w$
then obtain g **where** *1*: $G^A\ g\ w$ **..**
hence $NE^A\ g\ w$ **using** *A5* **by** *blast* — axiom 11.25
hence $\forall\, Y.\ (\mathcal{E}^A\ \ Y\ g\ w) \longrightarrow (\Box\exists^{\,E}\ \ Y)\ w$ **by** *simp*
hence *2*: $(\mathcal{E}^A\ \ G^A\ \ g\ w) \longrightarrow (\Box\exists^{\,E}\ \ G^A)\ w$ **by** (*rule allE*)
have $(\forall\, x.\ G^A\ \ x \rightarrow (\mathcal{E}^A\ \ G^A\ \ x))\ w$ **using** *GodIsEssential*
by (*rule allE*) — GodIsEssential follows from Axioms 11.11 and 11.3B
hence $(G^A\ \ g \rightarrow (\mathcal{E}^A\ \ G^A\ \ g))\ w$ **by** (*rule allE*)
hence $G^A\ g\ w \longrightarrow \mathcal{E}^A\ \ G^A\ \ g\ w$ **by** *blast*
from *this 1* **have** *3*: $\mathcal{E}^A\ \ G^A\ \ g\ w$ **by** (*rule mp*)
from *2 3* **have** $(\Box\exists^{\,E}\ \ G^A)\ w$ **by** (*rule mp*)
}
hence $(\exists\, x.\ G^A\ \ x\ w) \longrightarrow (\Box\exists^{\,E}\ \ G^A)\ w$ **by** (*rule impI*)
hence $((\exists\, x.\ G^A\ \ x) \rightarrow\ \Box\exists^{\,E}\ \ G^A)\ w$ **by** *simp*
}
thus *?thesis* **by** (*rule allI*)
qed

Some useful rules:

lemma *modal-distr*: $\lfloor\Box(\varphi \rightarrow \psi)\rfloor \Longrightarrow \lfloor(\Diamond\varphi \rightarrow \Diamond\psi)\rfloor$ **by** *blast*
lemma *modal-trans*: $(\lfloor\varphi \rightarrow \psi\rfloor \wedge \lfloor\psi \rightarrow \chi\rfloor) \Longrightarrow \lfloor\varphi \rightarrow \chi\rfloor$ **by** *simp*

Anderson's Version of Theorem 11.27

theorem *possExistenceImpliesNecEx*: $\lfloor\Diamond\exists\ \ G^A \rightarrow \Box\exists^{\,E}\ \ G^A\rfloor$ — local consequence

proof –
have $\lfloor\exists\ \ G^A \rightarrow \Box\exists^{\,E}\ \ G^A\rfloor$ **using** *GodExistenceImpliesNecExistence*
by *simp* — follows from Axioms 11.11, 11.25 and 11.3B
hence $\lfloor\Box(\exists\ \ G^A \rightarrow \Box\exists^{\,E}\ \ G^A)\rfloor$ **using** *NEC* **by** *simp*
hence *1*: $\lfloor\Diamond\exists\ \ G^A \rightarrow \Diamond\Box\exists^{\,E}\ \ G^A\rfloor$ **by** (*rule modal-distr*)
have *2*: $\lfloor\Diamond\Box\exists^{\,E}\ \ G^A \rightarrow \Box\exists^{\,E}\ \ G^A\rfloor$ **using** *symm tran* **by** *metis*
from *1 2* **have** $\lfloor\Diamond\exists\ \ G^A \rightarrow \Diamond\Box\exists^{\,E}\ \ G^A\rfloor \wedge \lfloor\Diamond\Box\exists^{\,E}\ \ G^A \rightarrow \Box\exists^{\,E}\ \ G^A\rfloor$ **by** *simp*
thus *?thesis* **by** (*rule modal-trans*)
qed

lemma *T4*: $\lfloor\Diamond\exists\ \ G^A\rfloor \longrightarrow \lfloor\Box\exists^{\,E}\ \ G^A\rfloor$ **using** *possExistenceImpliesNecEx*
by (*rule localImpGlobalCons*) — global consequence

Conclusion - Necessary (actualist) existence of God:

lemma *GodNecExists*: $\lfloor\Box\exists^{\,E}\ \ G^A\rfloor$ **using** *T3 T4* **by** *metis*

6.4 Modal Collapse

Modal collapse is countersatisfiable

lemma $\lfloor\forall\,\Phi.(\Phi \rightarrow (\Box\ \Phi))\rfloor$ **nitpick oops**

7 Conclusion

In this work we presented a shallow embedding in Isabelle/HOL for an intensional higher-order modal logic (a successor of Montague/Gallin intensional logics) as introduced by M. Fitting in his textbook. We employed this logic to formalise and verify all results (theorems, examples and exercises) relevant to the discussion of Gödel's ontological argument in the last part of the book. Three different versions of this ontological argument have been considered: the first one by Gödel himself, the second one by Fitting and the last one by Anderson.

By employing an interactive theorem-prover like Isabelle, we were not only able to verify Fitting's results, but also to demonstrate axiom's consistency throughout the whole argumentation. We could prove even stronger versions of many of the theorems and find better countermodels (i.e. with smaller cardinality) than the ones presented in the book. Another interesting aspect was the possibility to explore the implications of alternative formalisations for definitions and theorems which shed light on interesting philosophical issues concerning entailment, essentialism and free will, which are currently the subject of some follow-up analysis.

The latest developments in *automated theorem proving* allow us to engage in much more experimentation during the formalisation and assessment of arguments than ever before. The potential reduction (of several orders of magnitude) in the time needed for proving or disproving theorems (compared to pen-and-paper proofs), results in almost real-time feedback about the suitability of our speculations. The practical benefits of computer-supported argumentation go beyond mere quantitative (easier, faster and more reliable proofs). The advantages are also qualitative, since it fosters a different approach to argumentation: We can now work iteratively (by 'trial-and-error') on an argument by making gradual adjustments to its definitions, axioms and theorems. This allows us to continuously expose and revise the assumptions we indirectly commit ourselves everytime we opt for some particular formalisation.

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