THE CONSTRAINT PROGRAMMER'S TOOLBOX

Christian Schulte, SCALE, KTH & SICS

Constraint Programming

What is constraint programming?

Sudoku is constraint programming

...is constraint programming!

			2		5			
	9					7	3	
		2			9		6	
2						4		9
				7				
6		9						1
	8		4			1		
	6	3					8	
			6		8			

Assign blank fields digits such that:
 digits distinct per rows, columns, blocks

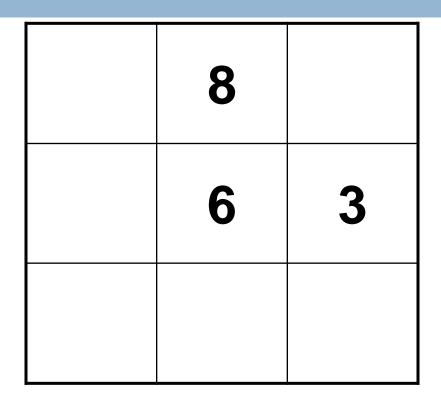
			2		5			
	9					7	3	
		2			9		6	
2						4		9
				7				
6		9						1
	8		4			1		
	6	3					8	
			6		8			

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	8		4			1		
	6	3					8	
			6		8			_

Assign blank fields digits such that:
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Block Propagation



No field in block can take digits 3,6,8

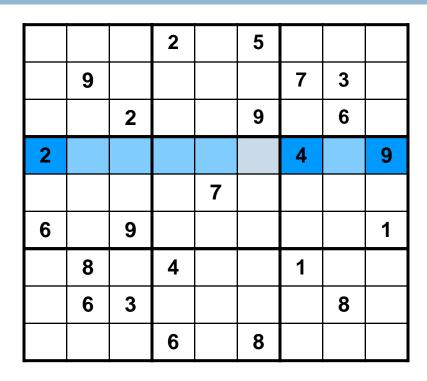
Block Propagation

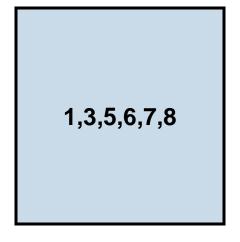
1,2,4,5,7,9	8	1,2,4,5,7,9
1,2,4,5,7,9	6	3
1,2,4,5,7,9	1,2,4,5,7,9	1,2,4,5,7,9

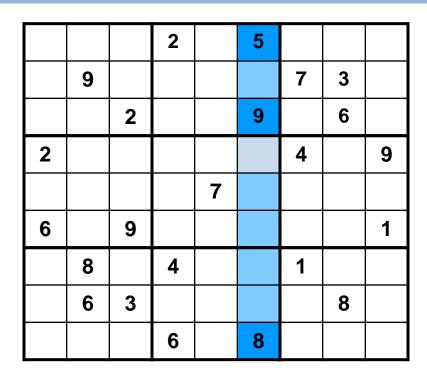
- No field in block can take digits 3,6,8
 - propagate to other fields in block
- Rows and columns: likewise

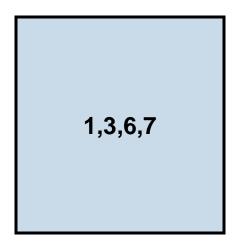
			2		5			
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	6	3					8	
			6		8			

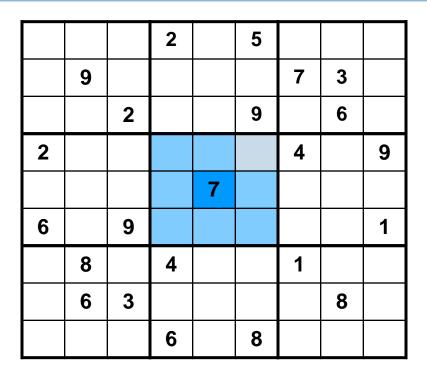
1,2,3,4,5,6,7,8,9

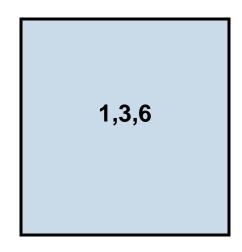












Iterated Propagation

			2		5			
	9					7	3	
		2			9		6	
2						4		9
				7				
6		9						1
	8		4			1		
	6	3					8	
			6		8			

- Iterate propagation for rows, columns, blocks
- What if no assignment: search... later

Sudoku is Constraint Programming

			2		5			
	9					7	3	
		2			9		6	
2						4		9
				7				
6		9						1
	8		4			1		
	6	3					8	
			6		8			

- Variables: fields
 - take values: digits
 - maintain set of possible values
- Constraints: distinct
 - relation among values for variables
- Modeling: variables, values, constraints
- Solving: propagation, search

Constraint Programming

- Variable domains
 - finite domain integer, finite sets, multisets, intervals, ...
- Constraints
 - distinct, arithmetic, scheduling, graphs, ...
- Solving
 - propagation, search, ...
- Modeling
 - variables, values, constraints, heuristics, symmetries, ...

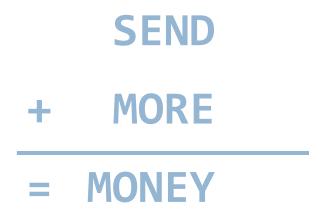
This Talk...

- Key concepts
 - constraint propagation
 - search
- The Constraint Programmer's Toolbox...
- Some few tools
 - global constraints: distinct reconsidered
 - branching heuristics: bin packing
 - user-defined constraints: personnel rostering
- Summary
 - essence of constraint programming and (very few) resources

17 Key Concepts

Running Example: SMM

Find distinct digits for letters such that



Constraint Model for SMM

Variables: $S,E,N,D,M,O,R,Y \in \{0,...,9\}$ Constraints: distinct(S,E,N,D,M,O,R,Y) $1000 \times S + 100 \times E + 10 \times N + D$ 1000×M+100×O+10×R+E + $= 10000 \times M + 1000 \times O + 100 \times N + 10 \times E + Y$ S≠0 M≠0

Solving SMM

Find values for variables

such that

all constraints satisfied

Finding a Solution

- Compute with possible values
 - rather than enumerating assignments
- Prune inconsistent values
 - constraint propagation

- Search
 - branch: define shape of search tree
 - explore: explore search tree for solution

Constraint Propagation

constraint store
propagators
constraint propagation

Constraint Store

$$x \in \{1,2,3,4\} \ y \in \{1,2,3,4\} \ z \in \{1,2,3,4\}$$

Maps variables to possible values

Constraint Store

finite domain constraints

$$x \in \{1,2,3,4\} \ y \in \{1,2,3,4\} \ z \in \{1,2,3,4\}$$

- Maps variables to possible values
 - other domains: finite sets, float intervals, graphs, ...

Implement constraints

$$distinct(x_1, ..., x_n)$$

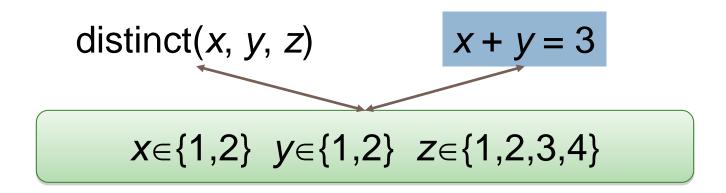
$$x + 2 \times y = z$$

schedule(
$$t_1, ..., t_n$$
)

distinct(x, y, z)
$$x + y = 3$$

 $x \in \{1,2,3,4\} \ y \in \{1,2,3,4\} \ z \in \{1,2,3,4\}$

- Strengthen store by constraint propagation
 - prune values in conflict with implemented constraint



- Strengthen store by constraint propagation
 - prune values in conflict with implemented constraint

distinct(x, y, z)
$$x + y = 3$$

 $x \in \{1,2\} \ y \in \{1,2\} \ z \in \{3,4\}$

- Iterate propagator execution until fixpoint
 - no more pruning possible

Propagation for SMM

Results in store

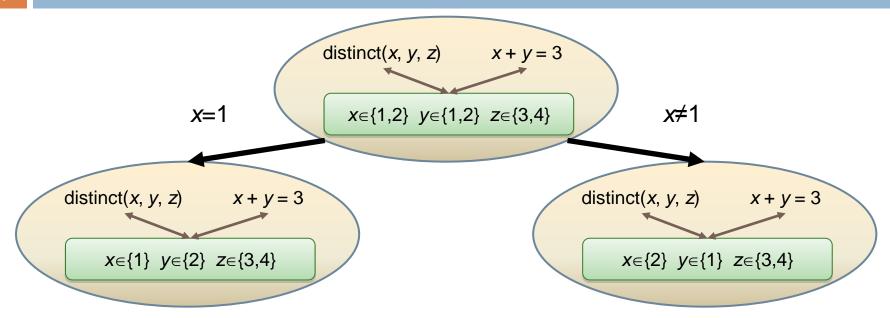
```
S \in \{9\} E \in \{4,...,7\} N \in \{5,...,8\} D \in \{2,...,8\} M \in \{1\} O \in \{0\} R \in \{2,...,8\} Y \in \{2,...,8\}
```

- Propagation alone not sufficient!
 - decompose into simpler sub-problems
 - branching and exploration for search

Search

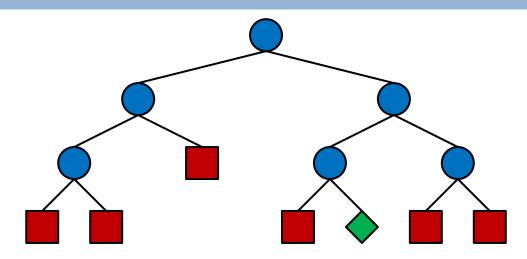
branching exploration

Branching



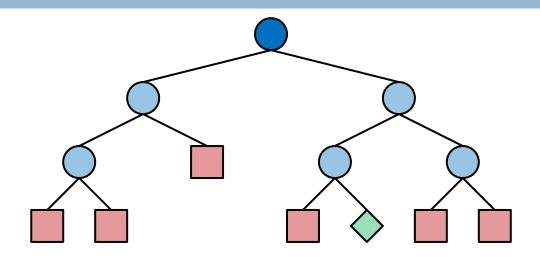
- Create subproblems with additional constraints
 - enables further propagation
 - defines search tree

Heuristic Branching

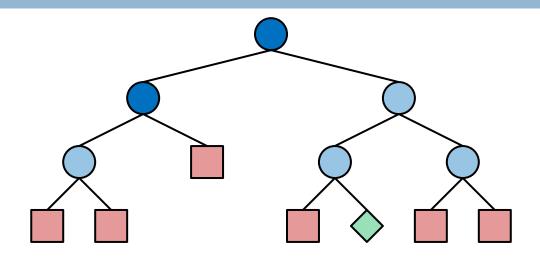


- Example branching
 - pick variable X
 - pick value
 - branch with x = n
 - and
- (at least two values) (from domain of x)
- $x \neq n$

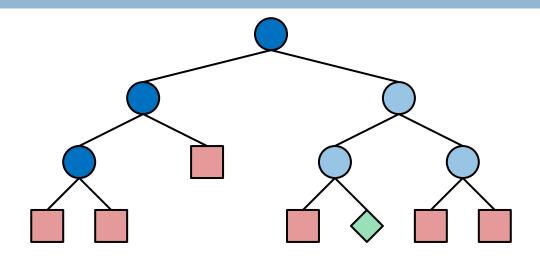
- Heuristic needed
 - which variable to select?
 - which value to select?



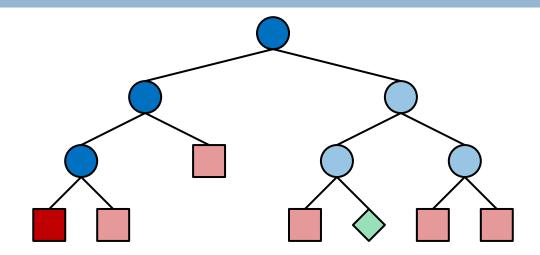
- Heuristic branching
 - defines tree shape
- Exploration of search tree
 - orthogonal aspect: DFS,



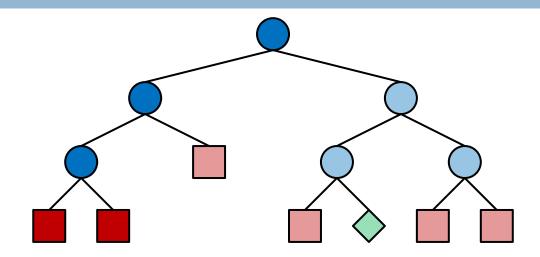
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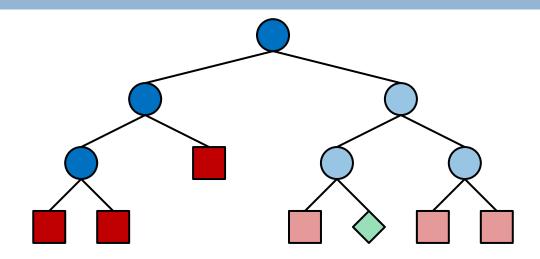
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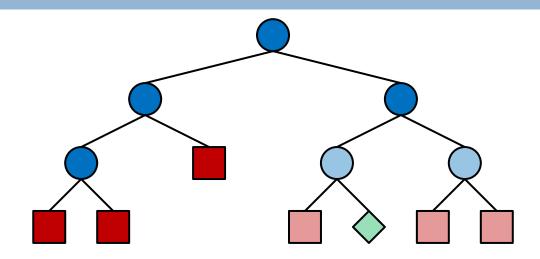
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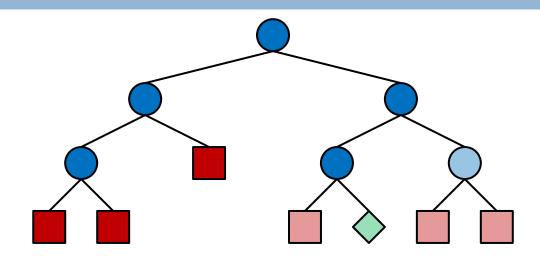
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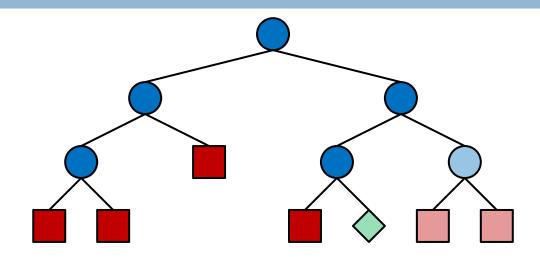
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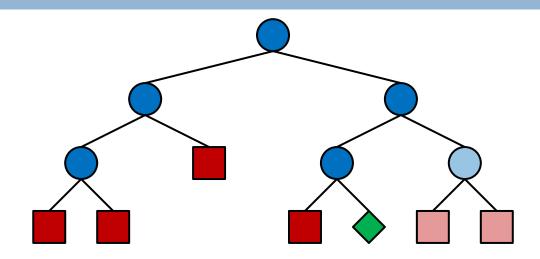
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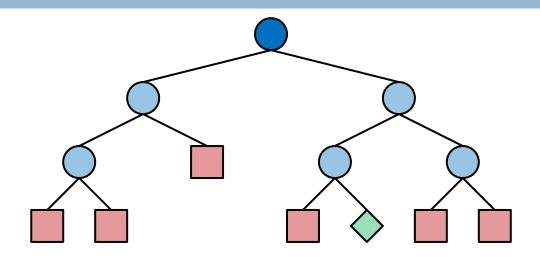
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- Heuristic branching
 - defines tree shape
- Exploration of search tree
 - orthogonal aspect: DFS, BFS, IDFS, LDS, parallel, ...

Summary

Modeling

- variables with domain
- constraints to state relations
- branching strategy
- in real: an array of modeling techniques...

Solving

- constraint propagation
- branching
- search tree exploration
- in real: an array of solving techniques...

The Constraint Programmer's Toolbox

Widely Applicable

- Timetabling
- Scheduling
- Personnel and crew rostering
- Resource allocation
- Workflow planning and optimization
- Gate allocation at airports
- Sports-event scheduling
- Railroad: track allocation, train allocation, schedules
- Automatic composition of music
- Genome sequencing
- Frequency allocation
- ...

Current Interest: Constraint-based Code Generation

- From
 - input program (in intermediate representation)
 - hardware architecture description generate constraint problem
- Solution = executable program code
 - simplicity: avoid bugs, ...
 - flexibility: architecture change, ...
 - quality: possibly optimal, ...
- Ongoing project 2010 2015
 - funded by Ericsson and Vetenskapsrådet

Problems Are Hard

- The problems are NP hard
 - no efficient algorithm is likely to exist
- Tremendously difficult to
 - always solve any problem instance
 - scale to large instances
 - have single silver bullet method
- Property of problems...
 - ...not of method
 - ...hence no silver bullet

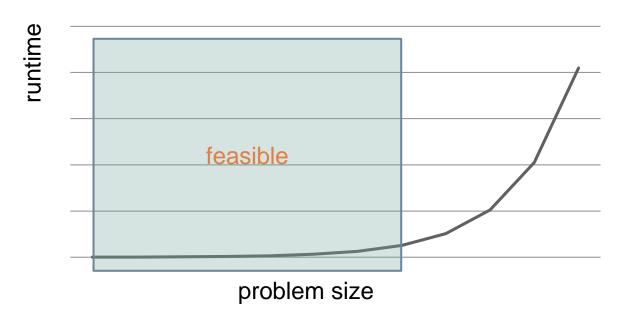
Why Is a Toolbox Needed?

- Initial model: model to capture problem
 - correctness
- Improved model: model to solve problem
 - robustness and scalability
 - often difficult
- Tools in the toolbox are needed for...
 - ...modeling to solve problems

Parts of the Toolbox

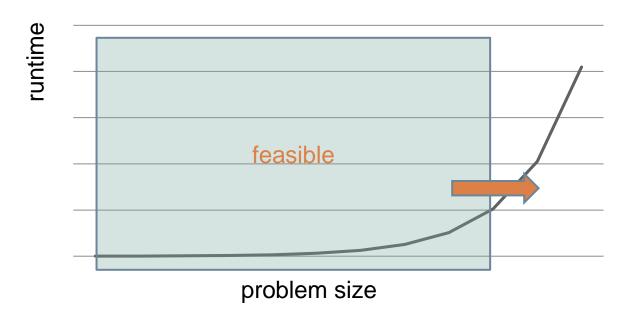
- "Global" constraints
 - capture structure during modeling
 - provide strong constraint propagation
- Search heuristics
 - application specific
- Symmetries and dominance relations
 - reduce size of search space
- Propagation-boosting constraints
- Randomized restarts during search
 - including no-goods from restarts
- ...

The Best We Can Hope for...



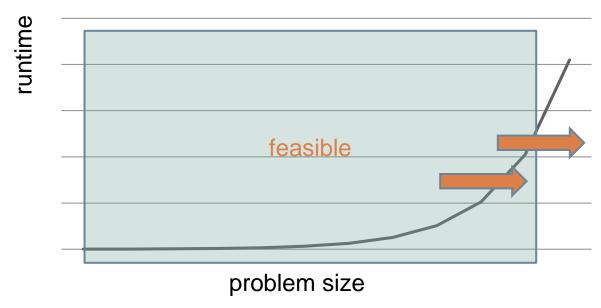
- Exponential growth in runtime
- Without using tools

The Best We Can Hope for...



- Exponential growth in runtime
- With propagation and heuristic search

The Best We Can Hope for...



- Exponential growth in runtime
- With propagation, heuristic search, symmetry breaking, restarts, ...

Capturing Structure

distinct reconsidered

Naïve Is Not Good Enough

- \square distinct(x, y, z)
 - naïve decomposition: $x \neq y$ and $x \neq z$ and $y \neq z$
 - propagates only as soon as x, y, or z assigned

- $\square x \in \{1,2,3\}, y \in \{1,2\}, z \in \{1,2\}$
 - should propagate $x \in \{3\}$
- $\Box x \in \{1,2\}, y \in \{1,2\}, z \in \{1,2\}$
 - should exhibit failure without search

Strong Propagation Idea

- \square distinct($x_0, ..., x_4$)
 - $x_0 \in \{0,1,2\} \ x_1 \in \{1,2\} \ x_2 \in \{1,2\} \ x_3 \in \{2,4,5\} \ x_4 \in \{5,6\}$
- Collect all solutions (permutations)
 - $x_0=0$ $x_1=1$ $x_2=2$ $x_3=4$ $x_4=5$
 - $x_0=0$ $x_1=1$ $x_2=2$ $x_3=4$ $x_4=6$
 - $x_0=0$ $x_1=1$ $x_2=2$ $x_3=5$ $x_4=6$
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 - $x_0=0$ $x_1=2$ $x_2=1$ $x_3=5$ $x_4=6$
- Collect values from solutions
 - $x_0 \in \{0\}$ $x_1 \in \{1,2\}$ $x_2 \in \{1,2\}$ $x_3 \in \{4,5\}$ $x_4 \in \{5,6\}$

Strong Propagation Idea

- \square distinct($x_0, ..., x_4$)
 - $x_0 \in \{0,1,2\} \ x_1 \in \{1,2\} \ x_2 \in \{7\}$

infeasible: all permutations!

- Collect all solutions (permutation)
 - $x_0=0$ $x_1=1$ $x_2=2$ $x_3=4$ $x_4=5$
 - $x_0=0$ $x_1=1$ $x_2=2$ $x_3=4$ $x_4=6$
 - $x_0=0 x_1=1 x_2=2 x_3=5 x_4=6$
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Strong Propagation Idea

- \square distinct($x_0, ..., x_4$)
 - $x_0 \in \{0,1,2\} \ x_1 \in \{1,2\} \ x_2 \in \{1,2\}$
- Characterize all solution

$$x_0=0$$
 $x_1=1$ $x_2=2$ $x_3=4$ $x_4=5$

$$x_0=0$$
 $x_1=1$ $x_2=2$ $x_3=4$ $x_4=6$

$$x_0=0$$
 $x_1=1$ $x_2=2$ $x_3=5$ $x_4=6$

$$x_0=0$$
 $x_1=2$ $x_2=1$ $x_3=4$ $x_4=5$

$$x_0=0$$
 $x_1=2$ $x_2=1$ $x_3=4$ $x_4=6$

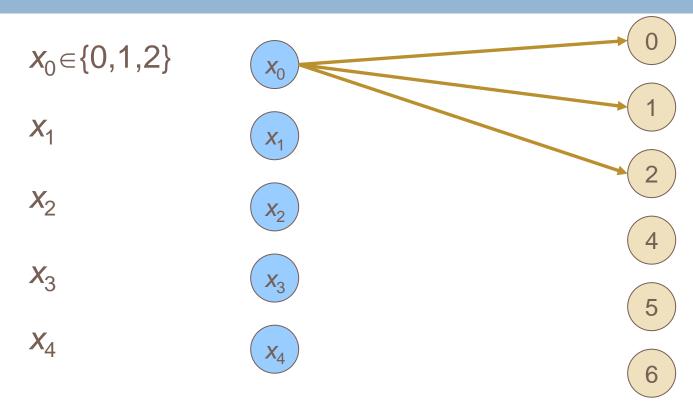
$$x_0=0$$
 $x_1=2$ $x_2=1$ $x_3=5$ $x_4=6$

translate into simple graph problem!

[Régin. A Filtering Algorithm for Constraints of Difference in CSPs. AAAI 1994]

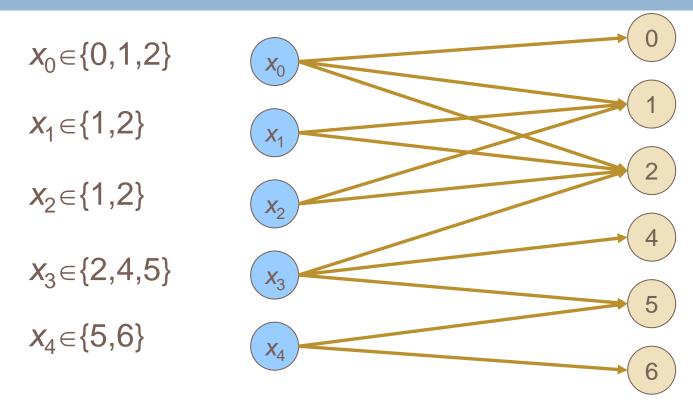
- Collect values from solutions
 - $x_0 \in \{0\}$ $x_1 \in \{1,2\}$ $x_2 \in \{1,2\}$ $x_3 \in \{4,5\}$ $x_4 \in \{5,6\}$

Variable Value Graph



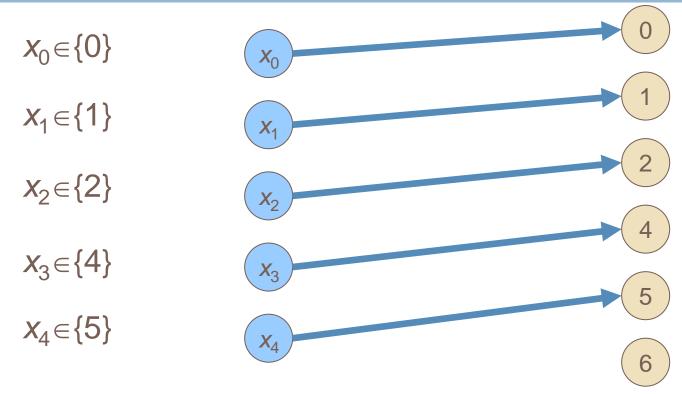
- Translates propagation into graph problem
 - variable nodes → value nodes

Variable Value Graph



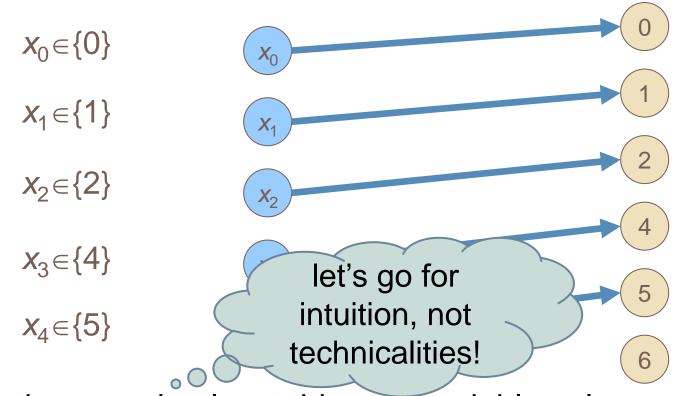
- Translates propagation into graph problem
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Graph Solution (1)



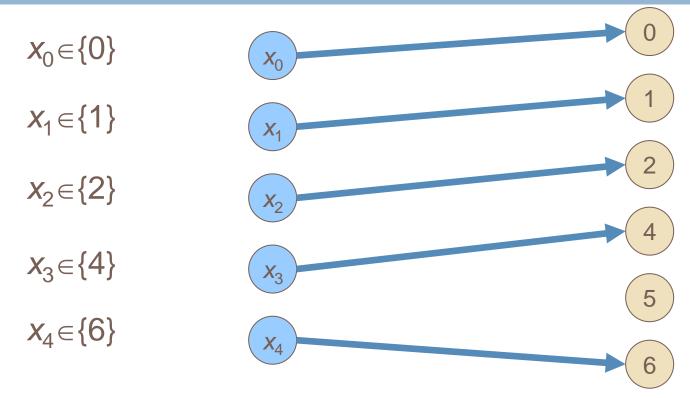
- Solutions maximal matchings in variable value graph
 - variable nodes → value nodes

Graph Solution (1)



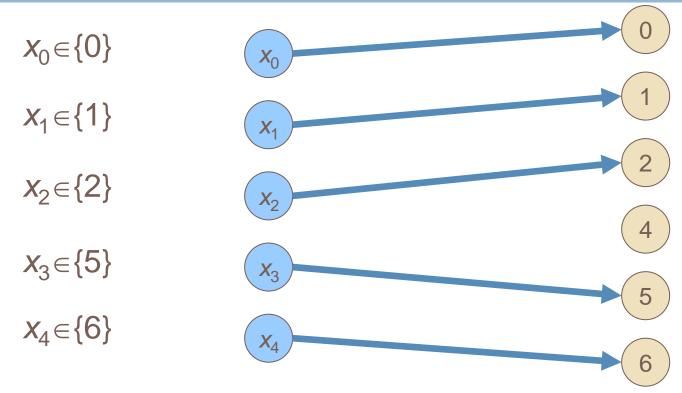
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Graph Solution (2)



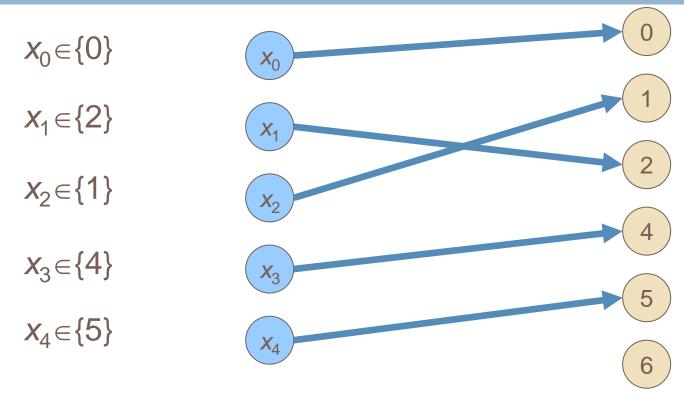
- Solutions maximal matchings in variable value graph
 - variable nodes → value nodes

Graph Solution (3)



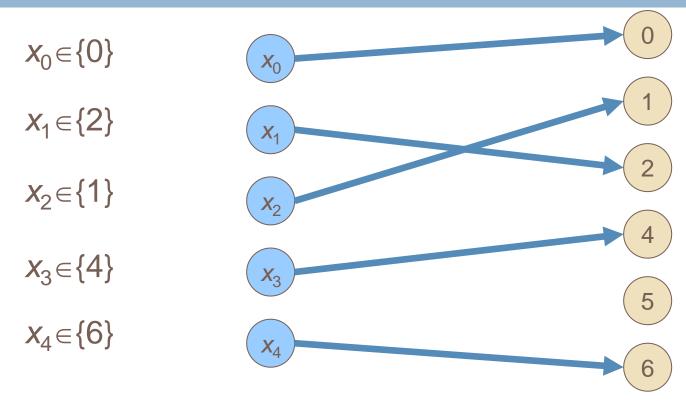
- Solutions maximal matchings in variable value graph
 - variable nodes → value nodes

Graph Solution (4)



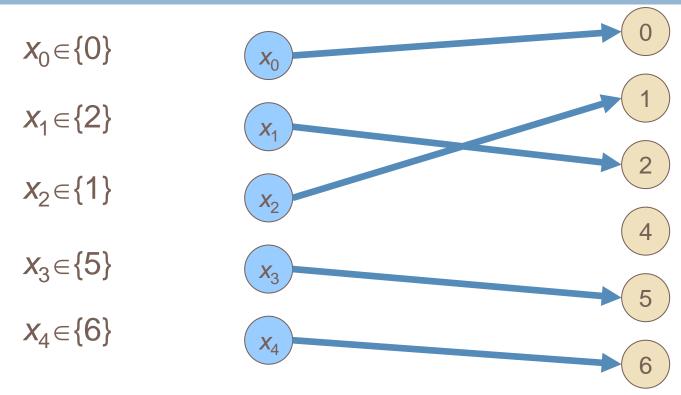
- Solutions maximal matchings in variable value graph
 - variable nodes → value nodes

Graph Solution (5)

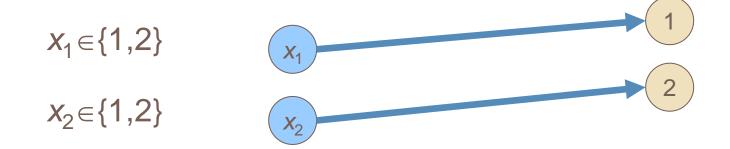


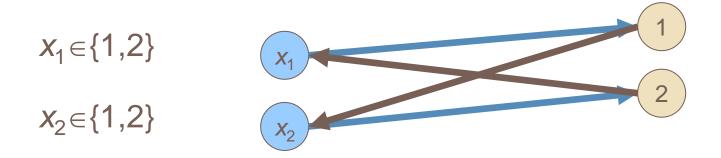
- Solutions maximal matchings in variable value graph
 - variable nodes → value nodes

Graph Solution (6)

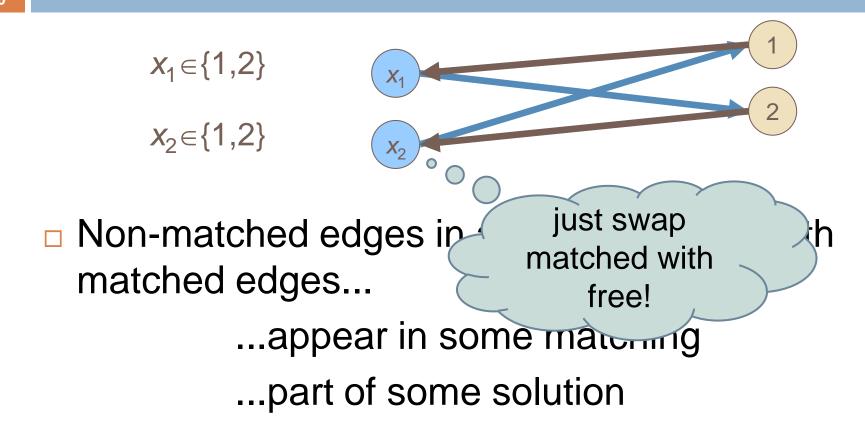


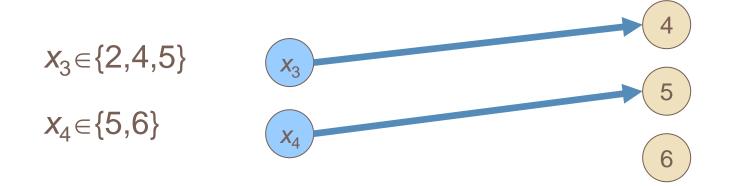
- Solutions maximal matchings in variable value graph
 - variable nodes → value nodes

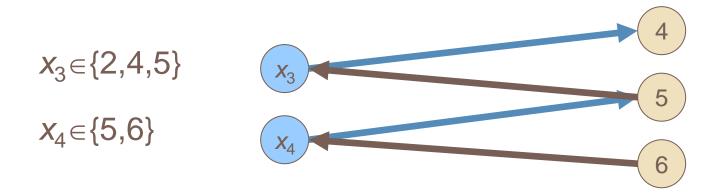




- Non-matched edges in alternating cycle with matched edges...
 - ...appear in some matching
 - ...part of some solution

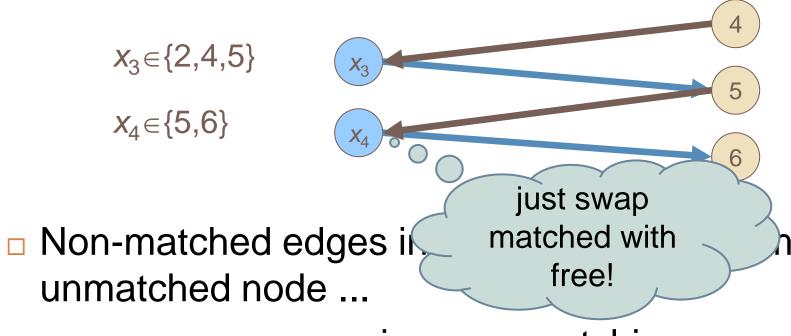






- Non-matched edges in alternating path from unmatched node ...
 - ...appears in some matching
 - ...part of some solution

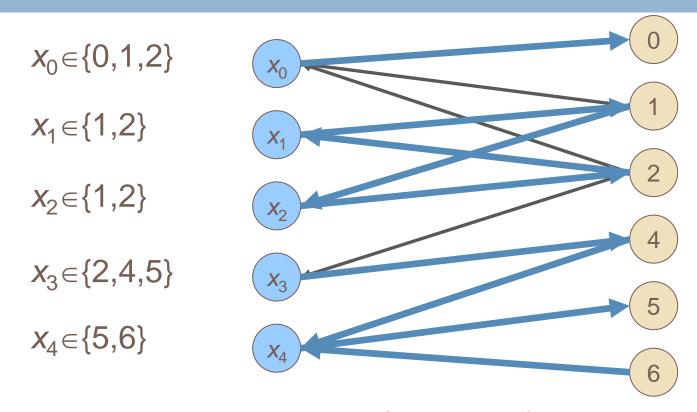
Characterizing All Solutions



...appears in some matching

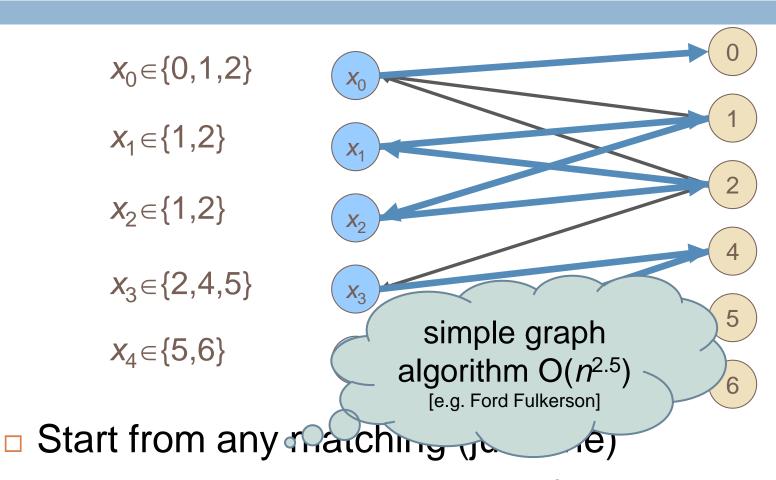
...part of some solution

Variable Value Graph



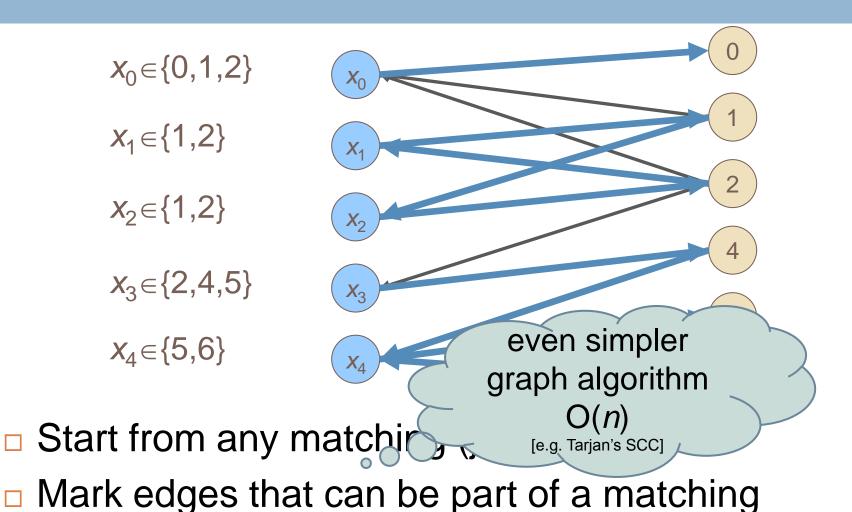
- Start from any matching (just one)
- Mark edges that can be part of a matching

Variable Value Graph

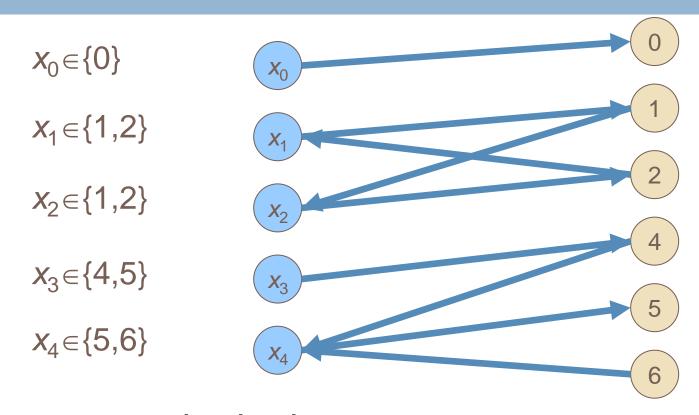


Mark edges that can be part of a matching

Variable Value Graph



Propagation, Finally!



Prune unmarked edges...

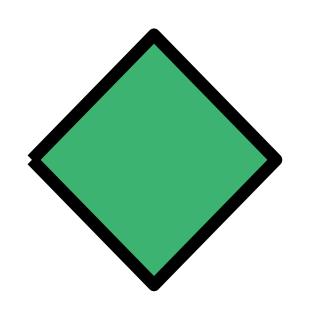
...and their corresponding values

Summary

- Constraints capture problem structure ("global")
 - ease modeling (commonly recurring structures)
 - enable solving (efficient and strong algorithms available)
- Constraints as
 - reusable
 - powerful

software components in the toolbox

SMM: Strong Propagation



```
SEND
   MORE
= MONEY
   9567
   1085
  10652
```

Branching Heuristics

bin packing

Branching Heuristics

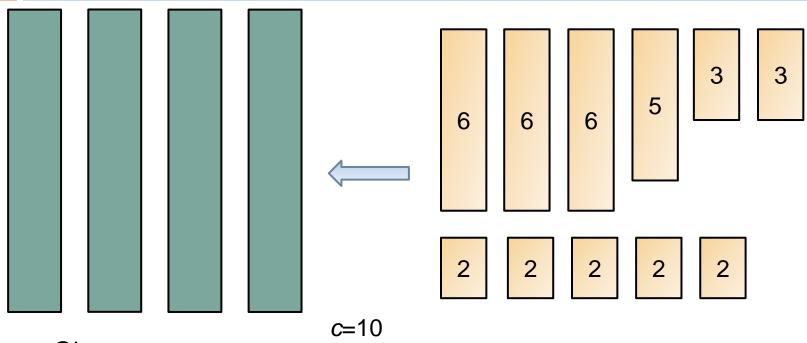
- CP advantage: programmable heuristics
 - application domain dependent: scheduling, assignment, bin-packing, ...
 - requires deep insight into problem structure
 - limited reuse even though recurring principles
- CP disadvantage: universal heuristics just emerging
 - CP solver as "black box" tool
 - ultimate goal: robust and autonomous search
 - contrast to SAT and MIP
- Here: bin packing as case study for programmable heuristics

First-Fail Principle

- Could be paraphrased as:
 - to succeed, try first where you are most likely to fail!
 - minimize cost to find out that decision is in fact wrong
 - cost = amount of search needed (depth-first search)
- Avoid thrashing
 - make wrong decision: search will have to find out
 - make many unrelated or non-difficult decisions
 - takes ages to find that decision was wrong!

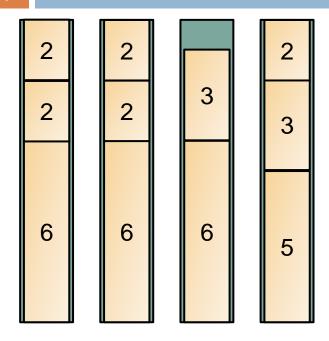
[Haralick, Elliott. Increasing tree search efficiency for constraint satisfaction problems. Artificial Intelligence, 1980]

Bin Packing Problem



- Given
 - bins of capacity c
 - \blacksquare *n* items of size s_i
- Sought
 - find least number of bins such that each item packed into bin

Bin Packing Problem



c=10

- Given
 - bins of capacity c
 - \blacksquare *n* items of size s_i
- Sought
 - find least number of bins such that each item packed into bin

Simplify Problem

- Repeat simpler problem for m such that: is it possible to pack n items into m bins?
- Restrict m by lower bound

$$I = \lceil (s_1 + \dots + s_n) / c \rceil$$

and upper bound

u =#bins from some (non-optimal) packing

- Try m between I and u: least feasible m optimal
 - even better lower bounds are known

Constraint Model: Variables

- □ Bin variable b_i for each item $i b_i \in \{1, ..., n\}$
 - into which bin is item i packed

- Load variable I_i for each bin j
 - size of items packed into bin j
- $I_j \in \{0, ..., c\}$

- Packing variable x_{ii}
 - whether item i is packed into bin j

$$x_{ij} \in \{0,1\}$$

Constraint Model: Constraints

Total load is size of all items

$$I_1 + \dots + I_m = S_1 + \dots + S_n$$

□ Load corresponds to items packed into bin j $I_{i} = s_{1} \cdot x_{1i} + ... + s_{n} \cdot x_{ni}$

□ Bin variables correspond to packing variables

$$x_{ij} = 1$$
 if and only if $b_i = j$

Constraint Model: Improved

- □ Use dedicated bin packing constraint binpacking($\langle b_1,...,b_n \rangle$, $\langle s_1,...,s_n \rangle$, $\langle I_1,...,I_m \rangle$)
 - no packing variables needed
 - much stronger propagation
- If items i and j with i<j have same size b_i ≤ b_j
 - reduce search space ("symmetry breaking")
- □ Assign large items $(s_i > c/2)$ to fixed bins
- [Shaw. A Constraint for Bin Packing. CP 2004]

How To Branch?

- \square Branch over the bin variables b_i
 - that is: assign items to bins
- Which item to pick first: largest!
- Which bin to pick first: tightest!
 - best fit (least slack)!
- "Easy" to express with standard heuristics...
 - ...can programming do more?

Programming Heuristic

Avoid search

- perfect fit of item i to bin b: assign i to b (no search)
- all bins have same slack: assign i to some b

Learn from failure

- try to assign item i to bin b
- if search fails: no other item j with $s_i = s_j$ can go to b
- if search fails: item i cannot go to bin with same slack (also for items j with $s_i = s_i$)
- "symmetry breaking during search"
- known as CDBF: complete decreasing best-fit

[Gent, Walsh. From approximate to optimal solutions: constructing pruning and propagation rules. IJCAI 1997.]

Local Reasoning

beauty and curse of constraint programming

		11	4		
	5 14			10	
17					3
6			3		
	10				
		3			

		11	4		
	5 14			10	
17					3
6			3		
	10				
		3			

- Fields take digits
- Hints describe
 - for row or column
 - digit sum must be hint
 - digits must be distinct

		11	4		
	5 14			10	
17					3
6			3		1
	10				2
		3			

For hint 31 + 2

		11	4		
	5 14			10	
17					3
6			3		2
	10				1
		3			

For hint 3

$$1 + 2$$

or

$$2 + 1$$

		11	4		
	5 14			10	
17					3
6			3	1	3
	10				
		3			

For hint 41 + 3

		11	4		
	5 14			10	
17					3
6			3	3	1
	10				
		3			

For hint 41 + 3or3 + 1

		11	4		
	5 14			10	
17					3
6			3	3	1
	10				2
		3			

- For hint 3
 - 1 + 2
- For hint 4

$$1 + 3$$

Kakuro Solution

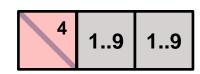
		11	4		
	5 14	2	3	10	
17	9	5	1	2	3
6	5	1	3	3	1
	10	3	1	4	2
		3	2	1	

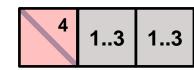
Modeling and Solving Kakuro

- Obvious model: for each hint
 - distinct constraint
 - sum constraint
- □ Good case... (?)
 - few variables per hint
 - few values per variable
- Let's try it...
 - 22×14, 114 hints: 9638 search nodes, 2min 40sec
 - 90×124, 4558 hints: ? search nodes, ? minutes years? centuries? eons?

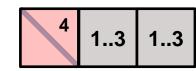
Local Reasoning: Decomposition

- Possible values
 - = all digits
- Propagating sum = 4
 - in isolation!
- Propagating distinct
 - in isolation!
- Propagating both
 - in combination!
 - but how?
 - where is the tool (constraint) for it?





all solutions: $\langle 1,3 \rangle \langle 2,2 \rangle \langle 3,1 \rangle$



all solutions: $\langle 1,2 \rangle \langle 1,3 \rangle \langle 2,1 \rangle$ $\langle 2,3 \rangle \langle 3,1 \rangle \langle 3,2 \rangle$



all solutions: $\langle 1,3 \rangle \langle 3,1 \rangle$

Failing for Kakuro...

- Beauty of constraint programming
 - local reasoning
 - propagators are independent
 - variables as simple communication channels
- Curse of constraint programming
 - local reasoing
 - propagators are independent
 - variables as simple communication channels

User-defined Constraints

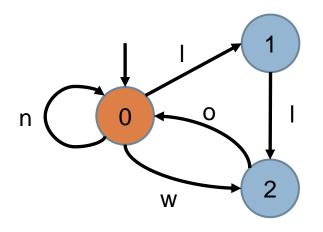
personnel rostering Kakuro reconsidered

Modeling Rostering: User-defined

- Personnel rostering: example (nonsensical)
 - one day off (o) after weekend shift (w)
 - one day off (o) after two consectuive long shifts (l)
 - normal shifts (n)
- Infeasible to implement propagator for everchanging rostering constraints
- User-defined constraints: describe legal rosters by regular expression
 - (wo | llo | n)*

Regular Constraint

(wo | Ilo | n)*



regular($x_1, ..., x_n, r$)

- $x_1 \dots x_n$ word in r
- □ or, accepted by DFA d for r

- Propagation idea: maintain all accepting paths in DFA
 - from start state (0) to a final state (0): solutions!
 - symbols on transitions comply with variable values

[Pesant. A Regular Language Membership Constraint for Finite Sequences of Variables. CP 2004]

Kakuro Reconsidered

- Real model: for each hint
 - one regular constraint combining distinct and sum
 - example: regular expression for hint 5 with two fields 14 | 23 | 32 | 41
 - precompute when model is setup
- Good case...
 - few solutions for combined constraint
- Let's try again (precomputation time included)
 - 22x14, 114 hints: 0 search nodes, 28 msec
 - 90×124, 4558 hints: 0 search nodes, 345 msec

Summary

- User-defined constraints
 - high degree of flexibility
 - efficient and perfect propagation
 - limited to medium-sized constraints
 - even better methods than regular known

- Kakuro: decomposition is harmful [again]
 - capture essential structure by few constraints
 - best by single constraint

108 Summary

Essence

- Constraint programming is about...
 - ...local reasoning exploiting structure
 - ...an array of modeling tools for solving
- Strength
 - simplicity, compositionality, exploiting structure
 - rich toolbox of techniques
- Challenges
 - lack of global picture during search
 - difficult to find global picture due to rich structure

Resources

Overview

- Rossi, Van Beek, Walsh, eds. Handbook of Constraint Programming, Elsevier, 2006 (around 950 pages).
- National perspective
 - Flener, Carlsson, Schulte. Constraint Programming in Sweden, *IEEE Intelligent Systems*, pages 87-89. IEEE Press, March/April, 2009.
 - SweConsNet: Swedish network for people interested in constraints. Yearly workshops, see:

www.it.uu.se/research/SweConsNet/

- Advanced (ID2204)/graduate (ID3005) course
 - taught by me
 - period 4, 2014