This document includes solution hints, but **NOT** necessarily the full solution!

Part I Evaluation Page Part A

| Name: | Signature: | |
|-------|------------|--|

1 Evaluation

This part of the exam has 180 questions, with a total of 310 points and 6 bonus points.

| Part | Max. Points | Scored Points |
|------|-------------|---------------|
| A | 60 | |
| В | 180 | |
| | Total | |

| Points: | 240-220 | 219-201 | 200-182 | 181-162 | 161-144 | 143-0 |
|---------|---------|---------|---------|---------|---------|-------|
| Grade: | A | В | С | D | Е | F |
| Score: | | | | | | |

2 Evaluation Part A

| Page | Points | Bonus Points | Score |
|--------|--------|--------------|-------|
| 5 | 5 | 0 | |
| 6 | 7 | 2 | |
| 7 | 12 | 0 | |
| 8 | 16 | 0 | |
| 9 | 5 | 1 | |
| 10 | 9 | 0 | |
| 11 | 4 | 0 | |
| 12 | 10 | 0 | |
| 13 | 12 | 0 | |
| 14 | 8 | 0 | |
| 15 | 6 | 0 | |
| 16 | 3 | 0 | |
| 17 | 4 | 0 | |
| 18 | 4 | 0 | |
| 19 | 2 | 0 | |
| 20 | 1 | 0 | |
| 21 | 2 | 2 | |
| 22 | 8 | 0 | |
| 23 | 3 | 0 | |
| 24 | 2 | 0 | |
| Total: | 123 | 5 | |

| Page | Points | Bonus Points | Score |
|--------|--------|--------------|-------|
| 26 | 6 | 0 | |
| 27 | 5 | 0 | |
| 29 | 3 | 0 | |
| 30 | 5 | 0 | |
| 31 | 2 | 0 | |
| 32 | 5 | 0 | |
| 33 | 12 | 0 | |
| 34 | 7 | 0 | |
| 35 | 13 | 0 | |
| 36 | 9 | 0 | |
| 37 | 13 | 0 | |
| 38 | 8 | 0 | |
| 39 | 8 | 0 | |
| 40 | 6 | 0 | |
| 41 | 7 | 0 | |
| 42 | 4 | 0 | |
| 43 | 8 | 0 | |
| 44 | 10 | 0 | |
| 45 | 11 | 0 | |
| 46 | 10 | 0 | |
| 47 | 8 | 0 | |
| 48 | 8 | 0 | |
| 49 | 7 | 1 | |
| 50 | 12 | 0 | |
| Total: | 187 | 1 | |

Part II

Rules

Answer the questions within the space provided. If you do not have enough space, you can use the backside of the sheet. In that case clearly indicate that your answer continues on the backside.

3 Supporting Materials

This is an examination in writing, without the usage of any electronic devices, except a scientific pocket calculator. No restriction on the model of calculator that may be used, but no device with communication capability shall be accepted as a calculator. All other electronic devices are prohibited. Writing paper is available, writing instruments (pencil, pens, etc) have to be organized by the student.

- Part A: Without any supporting material, with calculator.
- Part B: With a self written summary (format A4, 8 sheets or up to 16 pages), with calculator

4 Procedure

1. Duration:

Part A: 1 hour = 60 minutes = 60 points.

(short break)

Part B: 3 hours = 180 minutes = 180 points.

- 2. Sign the first page in the provided space. With this you certify that you are only using permitted support material and you are complying to the rules.
- 3. Write your name on any detached or additional paper sheets. Sheets without a name will not be evaluated.
- 4. Use the provided paper for your solutions. Use the provided space in the forms and tables. If needed use scratch paper. Document your way to your solution as appropriate.
- 5. Each question has a defined number of maximum points associated.
- 6. If a question is unclear, make reasonable assumptions. Document your assumptions and provide a rationale.
- 7. Write clearly and legibly. Unclear or multiple solutions will not be evaluated.
- 8. There is a short break between part A and B. You have to sign into a list for a needed break during the examination parts. Only one person can leave the room for a short time.
- 9. If something is unclear, ask your supervisor in the room.

5 Time Management

Read first all questions. Make sure you distribute your available time to all the questions. To reduce disturbance, ask questions in the first 15 minutes of the exam period.

6 Multiple-Choice Questions

- 1. Try to answer all questions if possible. If you are not sure, choose the answer which seems the best one.
- 2. For the questions of type \bigcirc : Choose **exactly one** option with \otimes (or $\sqrt{\ }$), which you think is the best match. With a correct answer you get the given number of points for that question.
- 3. For the questions of type ±: After a question or possibly incomplete sentence there are four answers or extensions. Evaluate each of them if they are true or false and mark them accordingly with '+' (true) or '-' (false). Independent if the question is formulated grammatically in singular or plural, it is possible that 0, 1, 2, 3, 4 of the choices are true. For three correct answers out of four you receive half of the points.
- 4. Wrong answers will have no penalty. Each question which has no answer is treated like a wrong answers and will be evaluated with zero points.
- 5. If you are changing your mind: cross out your old answer and clearly mark which answer is the new one.

May Dilbert be with you! ©

| Question 1 |
|---|
| $\sqrt{\text{Learning summary with 5 questions.}}$ |
| Sumo robot PCB capacitor. |
| Collection of slides. |
| Line sensor capacitance. |
| Tips for students in next semester. |
| Question 2 |
| O Variable Capacity System. |
| ○ Volatile Control Status. |
| $\sqrt{\text{Version Control System}}$. |
| O Variable Computer Software. |
| ○ Volatile Client Storage. |
| Question 3 |
| $\sqrt{\mathrm{ARM~Cortex-M0}}+$ |
| ○ Freescale HCS08 |
| \bigcirc MMA8780Q |
| ○ ARM Cortex M4 |
| |
| Question 4 |
| . gitignore readme.txt list.txt src\rotor.c src\rotor.h obj\rotor.o obj\rotor.txt |
| The .gitignore file has following content: |
| |

In above directory listing, strike through the files which are ignored.

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Reached: _____

Solution: not ignored are: .gitignore, readme.txt, src/main.c, src/main.h Explain in a single sentence what each of the following basic VCS actions mean in Git: (a) Committing $\frac{1}{2}$ **Solution:** Putting a change into the local repository. (b) Reverting $[1/_{2}]$ **Solution:** Undo a local change. (c) Pushing $[\frac{1}{2}]$ **Solution:** Moving a local change into the remote repository. (d) Cloning $[1/_{2}]$ **Solution:** Copy a repository and create a new local one. What is the fundamental difference between SVN and Git? Solution: SVN: centralized VCS, Git: distributed VCS. Explain the difference between the optimistic and pessimistic approach in a VCS. Explain it with an example. **Solution:** Optimistic: assumes that two developers do not work on the same file, so system potentially allows conflicts. Conflicts have to be resolved later. Pessimistic: assumes that conflict will happen, and whenever a developer wants to edit a file, the file gets locked so no conflict can occur. This could be your bonus question you have submitted... ②. 8. ves or no? [1] (a) Provide an example for a hard real-time system: **Solution:** Air bag, pacemaker or aircraft control system. (b) Provide an example for a *soft* real-time system: [1]

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| | Solution: Video streamer. |
|-------------------|---|
| | n 10 |
| | ition: Waiting for hardware state, waiting for semaphore/resource, runaway, stack overflow. |
| | n 11 |
| (a) | Name some benefits implementing a state machine: |
| | Solution: well structured, easy to implement, reusable design pattern. |
| (b) | What should be the first steps when implementing a state machine? |
| | Solution: Define all states with input/output and transitions. Draw a diagram. |
| | n 12Points: [5] TRO we implemented an 'Events' driver. |
| (a) | Why did we implement it as an array of bits? |
| | Solution: To save RAM. |
| (b) | What is the fundamental disadvantage of such an array of bits? |
| | Solution: Costs runtime performance for bit manipulation. |
| , | It implements critical section (e.g. to set an event bit) with EnterCritical() and ExitCritical(). Under which conditions such a critical section would not be required? |
| | Solution: If there is not a possiblity that the operation gets interrupted, or if the operation is atomic. |
| (d) | List reasons why an interrupt service routine should use such an Event module: |
| | Solution: Reduce interrupt latency time, only setting a bit and let the main application do the heavy lifting. |
| | n 13Points: [2] TOS can be either pre-emptive or cooperative: Explain the difference: |
| prio | tion: Pre-emptive: Always runs the highest available task. Tasks of identical rity share CPU time Cooperative: Context switches only occur if a task blocks, eplicitly calls yield. |
| \mathbf{uestio} | n 14Points: [4] |

[2](a) In an RTOS, each task can be in one of 5 fundamental states: List them: Solution: New, Ready, Running, Waiting, Stopped. (b) What's the purpose of the scheduler in an RTOS? [2]**Solution:** To determine which tasks gets executed next to minimize waiting time. Provide a short definition of the term Interrupt Latency, and which factors/aspects are contributing to it: **Solution:** Time between the event itself and until the ISR executes. Factors are stopping/finishing the current interrupt, interrupt destination calculation/arbitration, pushing state and diverting to the ISR. Question 16......Points: [2] Provide an example of a typical *Reactive System*, and explain why this is a reactive system: **Solution:** Airbag, it reacts on external events. (a) A PWM signal on a H-Bridge is labeled as low active. Explain what this means [2]and how this impacts the speed of a DC motor: **Solution:** low active means that the motor is active when the signal is low. It means for the PWM duty cycle: the longer in the low state, the higher the voltage, the faster the motor turns. (b) Draw a timing diagram for that PWM signal: the PWM period is 5 ms, and [2]the motor shall at 20% speed. Indicate how many milliseconds the signal is high and low. **Solution:** [Timing diagram drawn here, with a frequency of 5 ms and 20% low duty cycle (1 ms low, 4 ms high)]. Given the following program: #define ADC CONFIG (*(volatile uint8 t*)0x123) static void Interrupt(void) { uint8 t i; while (ADC CONFIG & $^{\sim}0x10$); for (i = 0; i < 10; i ++) {

asm("nop");

This program is using Interrupt synchronization. O Gadfly synchronization. Realtime synchronization. $\sqrt{\text{Realtime and Gadfly synchronization}}$. O No synchronization. Your Eclipse project stores the make files, object files and the final (binary) application file in a sub folder inside your project. Are you going to store this folder and files in a version control system? Justify your answer: Solution: No, as the content of this folder is generated. It does not make sense to store derived content in a version control system, as it can be generated from the sources. A tourist walked into a pet shop and was looking at the animals on display. Which monkey is the most expensive one? $\sqrt{\text{The monkey which programs nothing.}}$ O The monkey which programs in C. \bigcirc The monkey which programs in C++. \bigcirc The monkey which programs in Java and C++. \bigcirc The monkey which programs in C, C++ and Java. Question 21......Points: [3] Given the source of a PID control loop implementation. Identify in the source lines for the P, I and D part: mark them clearly and label it with P, I and D. Mark/circle this in the following source listing: #define max 0×33 ff static int 32 t old =0, b=0; void PID Control(void) { int32 t f, s, a; v = 0;

#define max 0x33ff
static int32_t old=0, b=0;
void PID_Control(void) {
 int32_t f, s, a;

 v = 0;
 f = should-actual;
 a = f-old;
 old = f;
 v += a/10;
 v += f*35;
 b += f;
 if (b > max) { b = max; }
 v += b/4;

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```
setAcuator(v);
   Solution: D, then P, then I with anti-windup.
Question 22......Points: [3]
   Processor Expert components are using the concept of Methods, Properties and
   Events. What would you expect for an ADC (Analog to Digital Converter) com-
   ponent?
    (a) 2 typical Methods for an ADC component:
                                                                             [1]
       Solution: Measure(), SetChannel(), ...
                                                                             [1]
   (b) 2 typical Properties for of an ADC component:
       Solution: Pin, sampling time, channel
    (c) 2 typical Events for an ADC component:
                                                                             [1]
       Solution: OnSamplingStart(), OnConversionEnd(), OnError(), ...
Given following variable definition:
   static char *string = "hello";
   What is the difference between the two following usages
   sizeof (string)
   strlen (string)
   in respect to the result and the expected code generated?
   Solution: size of gives the size in memory, which is here the size of a pointer (2 or 4
   bytes, depending on the machine), while strlen() the length without the zero byte.
   size of() is calculated at compile time (constant), while strlen() is a library routine
   call.
Question 24......Points: [3]
   Given following interface implementation for a DC motor driver:
```

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```
/* motor.h */
#include "LED.h" /* LED interface */
#include "PWM.h" /* PWM interface */
static uint16_t MOT_motorSpeed;

void MOT_Init(void);
/* end of motor.h */
```

Identify three issues with such an interface implementation (issues which could lead to linker/compiler failure, or things which are not considered as good programming style):

Solution: No #ifndef...#define, not necessary includes, static definition in header file.

Question 25......Points: [2]

Given following program:

```
#define DEC(i) {int b=0; i--;}

void main(void) {
  int a = 5, b = 5;
  DEC(a);
  DEC(b);
  printf("a is %d, b is %d\n", a, b);
}
```

What is the output of printf()?

```
Solution: a is 4, b is 5
```

Given the *Mealy Sequential State Machine* in Figure 1 with five states, two input values and two output values:

(a) The state machine in Figure 1 is in the state 'B'. Determine the *input sequence* in order to generate the following *output*:
1, 1, 1, 0

```
Solution: B: input1 output1 -> C
C: input0 output1 -> C
C: input0 output1 -> C
C: input0 output1 -> B
Solution: 1001
```

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[2]

[1]

[2]

[2]

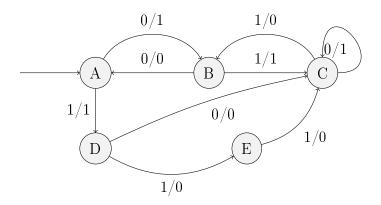


Figure 1: Mealy Sequential State Machine

(b) The State Machine in Figure 1 is not complete and has an undefined transition from one state to another: fix this with a solution in Figure 1.

Solution: Node E needs to have another outgoing arrow with 0/? (e.g. to node E).

Write in C code a single statement which only sets bit number 0 in PORTA and let the other bits untouched (bit number 0 is the least significant bit):

Solution: PORTA |= (1 < 0);

(a) Provide an example of a typical *Transforming System*, and explain why this is a Transforming System:

Solution: Network router, it transform packes and distributes them.

(b) Explain why Optimized Memory Usage is a typical attribute for a Transforming System:

Solution: Such systems transform an input stream into an output stream, and this usually involves larger amount of memory for buffering and transforming. As memory is expensive, such systems need to be optimized for this.

Explain the reason why some processors push all their core registers onto the interrupt stack, and some only push a subset of the registers:

Solution: In order not to increase the interrupt latency in case there are many registers, and as well to reduce the chance for stack overflow.

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| $\mathbf{Q}\mathbf{u}$ | estic | on 30 | |
|------------------------|-------|---|----------|
| | (a) | List 2 typical reactive systems: | [1] |
| | | Solution: Airbag, ABS | |
| | (b) | List 2 typical interactive systems: | [1] |
| | | Solution: PDA, ticket selling machine | |
| | (c) | List 2 typical transformative systems: | [1] |
| | | Solution: network router, encryption engine | |
| | | | Total: 3 |
| Qu | | on 31 | |
| | Pro | ution: Timing might be difficult to calculate. Everything must be deterministic. blem with stack consumption, timing problems, missed interrupts. Interrupts the not be acknowledged. | |
| Qu | List | on 32 | |
| | on | ution: Stop actual instruction (or undo, or finish), calculate new address based vector, store status on stack, branch to ISR and context switch. Do the same ag in reverse order to return from the ISR. | |
| Qu | | on 33 | |
| | iste | ution: Latency time depends on the speed of the CPU, the amount of data/regrs/stack to be changed for the context switch, and the entry time inside the ISR il the ISR routine can start. | |
| Qu | | on 34 | |
| | Sol | ution: The interrupts are disabled. | |
| Qu | | on 35 | |
| | | 35. <u>No</u> | |
| | | | |

Solution: Yes, an interrupt with lower interrupt priority number (higher interrupt priority number) can be interrupted by an interrupt source with higher priority (lower interrupt priority number).

| <pre>typedef signed short MyType; static unsigned char myVar[3]; typedef enum { RED=5, GREEN, YELLOW } Color</pre> | rs; | | |
|--|-------------------------|-------|--------|
| | | | |
| (a) What gives sizeof(MyType) for the FRDM | ${ m board/projec}$ | et: | |
| | (a) | 2 | |
| (b) What gives sizeof(MyType) for the Robot b | ${ m poard/projec}$ | t? | |
| | (b) | 2 | |
| (c) What gives sizeof(MyVar) for the FRDM b | $\mathrm{oard/project}$ | ? | |
| | (c) | 3 | |
| (d) What gives sizeof(MyVar) for the Robot bo | m pard/project' | ? | |
| | (d) | 3 | |
| (e) Which value has YELLOW? | | | |
| | (e) | 7 | |
| | | | Τ |
| estion 37 | ••••• | Point | s: [3] |
| <pre>uint16_t abcd[16]; uint8_t buf[10]; static uint16_t values[3];</pre> | | | |
| | | | |
| (a) Determine the value of following expression: sizeof("abcd") | | | |
| | (a) | 5 | |
| | | | |
| (b) Determine the value of following expression: sizeof(buf) | | | |
| | (b) | 10 | |

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Reached: _____ from 8 points

sizeof(values)

(c) _____6

Question 38......Points: [3] You have to implement the graph in Figure 2 with Doxygen/Dot.

Fill the gaps ('_____') in following code:

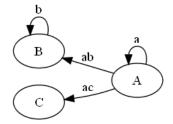


Figure 2: Doxygen ABC

```
dot
digraph example_dot_graph {
  node [];

rankdir = RL ______;
A __ [label = "A"] _____;
B__ [label = "B" _____];
C__ [label = "C"_____];
A__ -> _A __[label="a"];
A -> B ____[label="ab"];
A -> C ____[label="ac"];
B -> B [label="b"];
}
\enddot
```

Question 39......Points: [3]

Consider following doxygen source:

```
\dot
digraph example_dot_graph {
   node [shape=triangle];
```

```
rankdir=RL;
A    [style=filled,label="A"];
B    [style=filled,label="B"];
C    [style=filled,label="C"];
A -> A [label="a/b"];
A -> B -> C -> A -> C;
B -> B [label="b/c"];
}
\enddot
```

This produces the following graph:

 $\ominus \pm$ Solution is Figure 3.

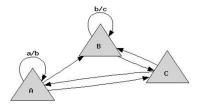


Figure 3: Dot Graph A

 $\ominus \pm$ Solution is Figure 4.

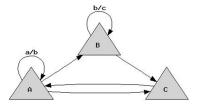


Figure 4: Dot Graph B

 $\ominus \pm$ Solution is Figure 5.

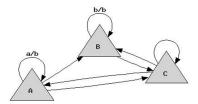


Figure 5: Dot Graph C

 $\ominus \pm$ Solution is Figure 6.

1. Linked Folder

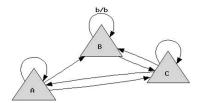


Figure 6: Dot Graph D

- 2. Linked Files
- 3. Virtual Group

List pros and cons for each approach:

Solution: Linked folder: Pros: new files in that folder automatically get added to the project. Cons: you get all or nothing.

Linked Files: Pros: You can decide for each remote file if it is included or not. Cons: you need to do this for every file.

Virtual Group: Pro: arbitrary group of files. Cons: No physical folder you can use for the build tools settings.

A hard realtime system or a soft realtime system: which do you consider easier to implement and test? List one pro and one cons for each:

Solution: A hard realtime system is probably harder to implement, but it is easier to test, as if you can make it fail a deadline, it is clear that it fails. A soft realtime system is probably easier to implement as it does not have to stick to hard deadlines, but it will be more difficult to test.

Consider following source:

```
#define MACRO(a,b) a = j \setminus a
```

Write down the text which would be produced by the preprocessor of the compiler, if you call the MACRO as following:

```
MACRO(i,5);
```

Solution:

```
MACRO(i, 5);
i = j = 5;
```

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Implement a function with doxygen comments which creates a doxygen output as in Figure 7.

```
uint8_t MyTest ( uint16_t val )

Implements a test routine.

Parameters:
    [in] val Input value

Returns:
    Error code

Todo:
    Implement function

{
    return 0;
}
```

Figure 7: Doxygen for MyTest()

Solution:

```
/*!

* | brief Implements a test routine

* | param[in] val Input value

* | return Error code

*/

uint8_t MyTest(uint16_t val) {

/*! | todo Implement function */

return 0;
}
```

Given following C source:

```
#define MACRO(var, mask1, mask2) \
    (var = (var & (~(uint8_t)(mask1))) | (uint8_t)(mask2))
    static uint8_t var;

void foo(void) {
    var = 0x22;
    MACRO(var, 16, 0x13);
}
```

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What is the value of var after execution of foo()?

44. <u>0x33 oder 51</u>

Solution:

```
 \begin{array}{l} var = 0x22\,;\\ (var = (var \& (~(uint8\_t)(16))) \mid (uint8\_t)(0x13))\,;\\ var = (var \& (~(uint8\_t)(0x10))) \mid (uint8\_t)(0x13)\,;\\ =>> clear & bits & in & mask1\,, & set & bits & in & mask2\\ =>> clear & bit & 0x10\,, & set & bits & 0x13\\ =>> clear & has & no & effect\,, & so & it & is & 0x22 \mid 0x13 ==> 0x33 \\ \end{array}
```

Which sequence is the correct one to configure a keyboard interrupt?

- Enable Keyboard Interrupts;
 - Set Port direction as input;
 - Enable Pull-Up Resistors;
 - Acknowledge Pending Interrupt;
- Acknowledge Pending Interrupt;
 - Set Port direction register as input;
 - Enable Keyboard Interrupts;
 - Enable Pull-Up Resistors;
- $\sqrt{\text{Set Port direction register as input;}}$
 - Enable Pull-Up Resistors:
 - Acknowledge Pending Interrupt;
 - Enable Keyboard Interrupts;
- Carry Enable Pull-Up Resistors;
 - Enable Keyboard Interrupts;
 - Acknowledge Pending Interrupt;
 - Set Port direction register as input;

For the implementation of a driver for an interrupt hardware following has to be considered:

- \ominus ± Interrupts have to be enabled globally during the driver initialization.
- \oplus \pm The driver shall reset the device interrupt flag during initialization.
- \oplus \pm After a power-on reset, it might be necessary to wait a certain time until the hardware signals have stabilized.
- \ominus ± The interrupt handler shall be as efficient as possible in order to increase the interrupt latency time.

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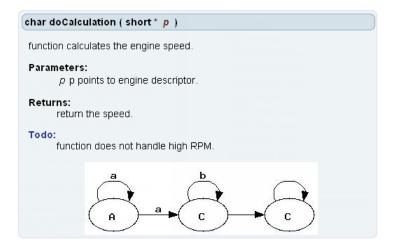


Figure 8: Doxygen Dokumenation

Given the doxygen graph in Figure 8. This figure has been created with following doxygen extract:

```
\dot digraph a_graph {
        node[],rankdir=RL; A,B[label="B"],C;
        A \rightarrow A \rightarrow B[label="a"];
        B->B[label="b"]; B->C->C;
      }
      \enddot
     \dot digraph b_graph {
       node[],rankdir=LR; A,B[label="C"],C;
        A \rightarrow A \rightarrow B[label="a"];
        B->B[label="b"]; B->C->C;
      }
     \enddot
     \dot digraph c_graph {
       node[],rankdir=RL; A,B,C;
        A \rightarrow A \rightarrow B[label="b"];
        B->B[label="a"]; B->C;
      }
     \enddot
     \dot digraph d_graph {
       node[],rankdir=LR; A[label="B"],B,C;
        A \rightarrow A \rightarrow B[label="b"];
        B->B[label="a"]; B->C;
      }
     \enddot
     \dot digraph e_graph {
\bigcirc
        node[],rankdir=LR; A,B,C[label="B"];
        B->A->B[label="b"];
        B->B[label="a"]; B->C->A;
```

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} \enddot

```
void main(void) {
unsigned char *src = (unsigned char*) 0x100, buffer [0x100], i;
for (i = 0; i < 100; i++) {
   buffer [i] = *src;
}
}</pre>
```

For the above program, following applies:

- \ominus ± It reads the values from the address 256 and 512 and stores it in a buffer.
- \oplus ± It reads 100 times the value at the address 0x100 and stores the values one after each other in a buffer.
- \ominus ± At termination of the program, the whole buffer is filled with the values from address 0x100.
- \oplus ± With disabled interrupts, the program behaves in a deterministic way.

Solution: 3rd Answer: only part of the buffer ([0]..[99] is filled with *0x100, but the buffer has the size of [0x100].

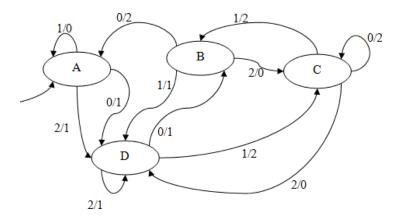


Figure 9: Mealy Machine

(a) The machine in Figure 9 is currently in state 'C'. Determine the output sequence for following input values: 0, 1, 0, 1, 1, 0

[1]

[6]

Solution: 2, 2, 2, 0, 0, 1

(b) Given following Mealy program:

```
typedef enum {A=0, B, C, D} States;
void Run(void) {
   char j, i = 0

for (;;) {
    j = Input();
    Output(tbl[i][j][1]);
    i = tbl[i][j][0];
   }
}
```

To implement the machine in Figure 9, complete the initialization of table tbl:

```
const char tbl[4][3][2] = {
```

Solution:

For realtime systems following applies:

- \ominus ± Realtime systems have to have reaction times below 1 ms in order to be realtime compliant.
- \oplus \pm For a realtime system not the average system load matters, but the highest possible system load.
- \ominus ± Hard realtime systems are more difficult to verify, because the realtime conditions are not exactly specified.
- \ominus ± A system can be a realtime system, if it is using true random number generator for its decision instead of a pseudo random number generator.

Given following program:

```
char buf[0x100];
int i,j;

static void test(void) {
  for(i=0; i<sizeof(buf); i++) {</pre>
```

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```
CFG = 0x80; PORTB = 4;

buf [i] = PORTA;

PORTB = 0;

}
```

For this program following applies:

- $\ominus \pm$ Implements an interrupt synchronization.
- $\ominus \pm$ Implements a gadfly synchronization.
- $\ominus \pm$ Implements a realtime synchronization.
- $\oplus \pm$ None of above.

Question 53.......Points: [1]

For all reentrant functions in C, following has to apply:

- \ominus ± The function shall not be recursive.
- \ominus ± The function shall not be called from an ISR.
- \oplus \pm The access to shared data has to be protected from mutual access.
- \oplus ± The function shall not modify itself (self modifying code).

Question 54.......Points: [1]

The following program gets compiled for the FRDM board with default compiler options:

```
static char ch@0x10;
void foo(void) {
    static char i, j=4;
    volatile char v;
    v = i;
    v = j;
    ch = v;
}
```

Following applies:

- \ominus ± The variables i, j and v are allocated on the stack.
- \oplus \pm The compiler cannot optimize the two assignments to v because of volatile.
- \ominus ± At execution time of foo(), the variable v gets initialized with a value of 4.
- \oplus ± After execution of foo(), the memory at address 0x10 will have a value of 4.

For the interrupt system of the ARM Cortex following applies:

- ⊕ ± The interrupt latency is the sum of execution time of the current instruction, pushing of the registers, calculating the ISR PC address and the branching to the ISR itself.
- \ominus ± With 'masking the interrupts' we are enabling the interrupts.
- \oplus ± In order for the ISR program to return to the interrupted program, the return address of the interrupted program is stored on the stack by the hardware.
- \oplus \pm In order to reduce the interrupt latency time, the core can decide not to push all registers on the stack.

The diagram in Figure 10 shows an interrupt system with multiple interrupts (IRQ1 and IRQ2) and the corresponding interrupt service routines (ISR) #1 and #2). The lines on the time axis denote the execution time boundaries of the instructions. Following applies:

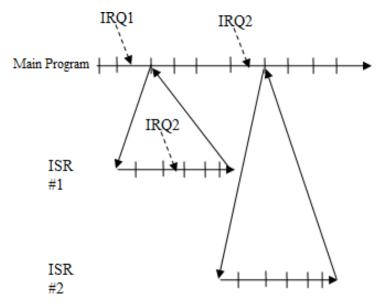


Figure 10: Interrupts

- At the beginning of ISR #1 all interrupts get disabled, and at the end of ISR #1 the flag for IRQ1 gets acknowledged.
- \bigcirc The main program has at the beginning all interrupts disabled and has the IRQ1 flag acknowledged. After execution of ISR #1 the main program enables all interrupts.
- $\sqrt{\text{ISR } #1 \text{ turns off all interrupts at the beginning.}}$ At the end of ISR #1 it acknowledged the IRQ1 and IRQ2 flag and enables all interrupts again.
- \bigcirc At the beginning of ISR #1 the flags for IRQ1 and IRQ2 are acknowledged. All interrupts get disabled at the end of ISR #1.
- \bigcirc ISR #1 has not acknowledged the IRQ1 flag. ISR #2 acknowledged the flags for IRQ1 and IRQ2 at the beginning of ISR #2.

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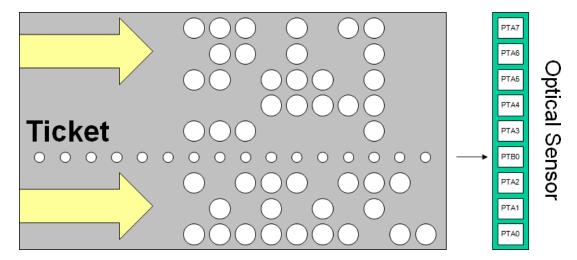


Figure 11: Parking Ticket

Question 57.......Points: [5]

A punched paper ticket is used in a parking system. The punched paper ticket is using following format for each data line in Figure 11:

- 1 guidance bit (small holes)
- 8 data bits (large holes)

The punched paper tape gets pulled into the machine with constant speed of 50 ms for each data line. The data lines are scanned with an optical sensor, and the sensor digital output is attached to the port of a microcontroller. The state of the sensor/holes is available on the microcontroller PORTA, bit 0 to 7:

- Value of bit is 0: no hole, light does not go through
- Value of bit is 1: hole, light goes through

The state of the guidance hole is available on bit 0 of PORTB. The bit 0 of PORTB is configured to raise an keyboard interrupt on falling edge.

Given following program:

```
extern WaitMs(unsigned int ms); /* wait for the given ms */
unsigned char buffer[16]; /* contains the data read */

void Read(void) {
    uint8_t i;

    for (i = 0; i < size of (buffer); i++) {
        WaitMs(50);
        buffer[i] = PORTA;
    }
}
interrupt KBI(void) {
    /* Guidance Hole Sensor */
    AcknowledgeKBI();
    DisableInterrupts();
    Read();</pre>
```

```
EnableInterrupts
}

void main(void) {
    for (;;);
}
```

- (a) Which synchronization method is used for the detection of insertion of the [1]parking ticket? Combination of interrupt and realtime synchronization. Interrupt synchronization. Realtime synchronization. Combination of gadfly synchronization and realtime synchronization. O Gadfly synchronization. [1] (b) Which synchronization method is used for the synchronization on the **first data** hole? Combination of interrupt and realtime synchronization. ○ Interrupt synchronization. $\sqrt{\text{Realtime synchronization}}$. Combination of gadfly synchronization und realtime synchronization. O Gadfly synchronization. [3] (c) Implement a new function ReadGadfly() which does the same as Read(), but
 - uses a gadfly synchronization method.

void ReadGadfly(void) { ...

Solution:

```
void ReadGadfly(void) {
    uint8_t i;

for(i=0; i < sizeof(buffer); i++) {
    while((PORTB&1)==0);
    buffer[i] = PORTA;
    while((PORTB&1)==1);
} /* end for */
} /* end ReadGadfly */</pre>
```

Total: 5

Question 58......Points: [1]

Given following program:

```
double power(double x, int exp) {
  if (exp<=0) return 1;</pre>
```

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Question 58 continuous on the next page...

Reached: _____

```
return(x*power(x, exp-1));
}
```

Evaluate following:

- \ominus ± In order to have this program reentrant, it is sufficient that **x** and **exp** are variables on a hardware stack.
- \oplus ± It depends on the compiler and the generated code, if this program is reentrant or not.
- \ominus ± The program is reentrant if it is called from an interrupt service routine only.
- \ominus ± The recursive implementation of this program ensures that it is reentrant.

Solution: Notice that the compiler routines for float operations (runtime routines) might not be reentrant. And not every hardware architecture is using a hardware stack for local variables, e.g. some PIC controlles or ST5 do not have a hardware stack (variables are on a software stack which prevent recursion or reentrant code).

| Program | Main Priority | Sub Priority | Time |
|---------|---------------|--------------|-------------|
| HP | 0 | 0 | 5 ms |
| UP1 | 1 | 1 | $2 \mu s$ |
| UP2 | 1 | 2 | $3 \ \mu s$ |
| UP3 | 2 | 1 | $5 \ \mu s$ |
| UP4 | 2 | 2 | $2 \mu s$ |

Table 1: Interrupt System

The timing required for a context switch is given in table 2, which is illustrated in Figure 12.

| Context Switch | Time |
|--|-------------|
| Total time for the interrupt, switch to a new program and starting execution | $1 \ \mu s$ |
| of the waiting program | |
| Total time for the interrupt, switch to the interrupted program, immediate | $1 \ \mu s$ |
| interruption of this program and switching and starting execution of the | |
| waiting program | |

Table 2: Context Switch Timing

The interrupt system is using following rules (as used in the lecture):

$$if(MP(s) \le MP(fn)) \to ws = ws \cup s$$
 (1)

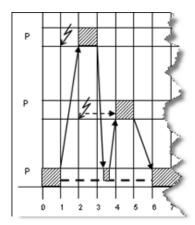


Figure 12: Example Context Switch

$$if(MP(s) > MP(fn)) \to INT(fn)$$
 (2)

$$if(MP(s) \le MP(fn)) \to ws = ws \cup s$$
 (3)

$$if(SP(s) > SP(fn)) \to ws = ws \cup s$$
 (4)

$$if(ws \neq \{\}) \rightarrow fn(MAX(SP(ws)))$$
 (5)

$$if(MP(ws) > MP(in)) \rightarrow fn = in \rightarrow ws = in$$
 (6)

The programs run according following information:

- 1. At the time 0 μ s HP starts.
- 2. At the time 2 μ s an interrupt for UP1 is raised.
- 3. At the time 4 μ s an interrupt for UP4 is raised
- 4. At the time 6 μ s an interrupt for UP2 is raised
- 5. At the time 9 μ s an interrupt for UP3 is raised
- 6. At the time 23 μ s an interrupt for UP2 is raised

Show the sequence of programs and interrupts in Figure 13. Use the same notation as in Figure 12 for interrupts (Exception, Pending), program switches, program (aktive, suspended).

Solution: See Figure below.

- 1. A signal with a higher main priority interrupts always a running program.
- 2. A signal with same main priority interrupts never.
- 3. A signal with higher sub priority has to wait for an already started program with same main priority, even if it had been interrupted.
- 4. If there are multiple signals waiting with the same main priority, then the sub priority decides which one will be handled first. With sub priorities it is possible to influence the sequence of execution and you are not depending on the signal raise time.

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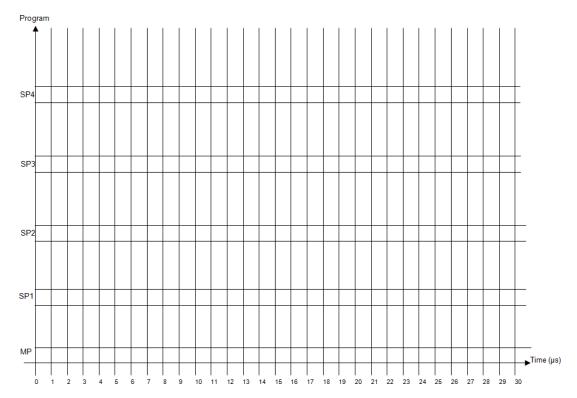


Figure 13: Program Timing

5. If there is a pending signal with a higher main priority than the last interrupted program, then the control goes back to that interrupted program, but gets interrupted directly again.

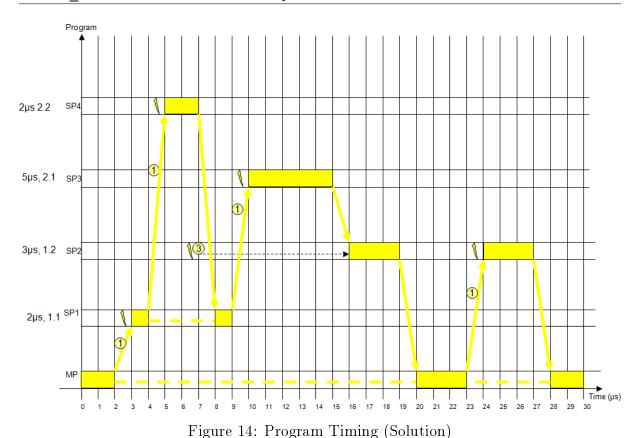
```
#define FUNC(a,b) i+a+b
int foo(int i, int j) {
   return FUNC(i,j);
}
```

Determin the return value for foo(5,6);:

60. _____<u>16</u>

```
return FUNC(i,j);
return i+i+j;
return 5+5+6;
return 16;
```

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Question 61......Points: [3]

Given following source code:

```
uint16_t abcd[8];
uint32_t buf[10];
static uint16_t values[2];
```

(a) Determine the value of following expression: sizeof("abcd")

(a) _____<u>5</u>

[1]

[1]

[1]

(b) Determine the value of following expression: sizeof(buf)

(b) ______

(c) Determine the value of following expression: sizeof(values)

(c) <u>4</u>

Consider following program:

unsigned char data@0x10;

```
#define p_data ((unsigned char*)0x10)
void foo(void) {
   data = 10;
   *p_data = 17;
   data++;
}
```

After execution of foo(), following applies:

```
\ominus \pm \text{ data} == 11, *p\_\text{data} == 17, p\_\text{data} == 10
\ominus \pm \text{ data} == 10, *p\_\text{data} == 18, p\_\text{data} == 17
\ominus \pm \text{ data} == 10, *p\_\text{data} == 17, p\_\text{data} == 0x10
\ominus \pm \text{ data} == 17, *p\_\text{data} == 17, p\_\text{data} == 16
\ominus \pm \text{ data} == 18, *p\_\text{data} == 18, p\_\text{data} == 16
```

Consider following program:

```
void delay(void) {
  uint8_t i;
  for(i=0;i<50;i++);
}</pre>
```

This program

- $\ominus \pm$ always waits for 50 ms
- \oplus ± can be optimized by a smart compiler to a function which only contains a return; statement
- \oplus ± will wait for a certain time which is depending on the speed of the microcontroller used
- $\ominus \pm$ will never terminate

Consider following Mealy Sequential State Machine with 3 states and two inputs:

```
typedef enum {A, B, C, D, E, } States;
const char tbl[3][2][2] =
{ {{A,0}, {B,1}},
    {{C,3}, {A,4}},
    {{C,0}, {B,5}}
};

void Run(void) {
    char j, i = 0

for (;;) {
        j = Input();
        Output(tbl[i][j][1]);
        i = tbl[i][j][0];
}
```

[1]

[1]

```
}
```

- (a) Given following sequence of Input() values: 0, 1, 0, 1, 1, 1. Determine the sequence of Output() values:
 - (a) ______0, 1, 3, 5, 4, 1
- (b) Draw the corresponding state diagram:

Solution: See diagram.

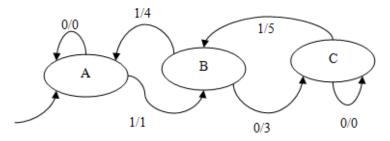


Figure 15: Solution Mealy Diagram Drawing

Total: 2

Consider following program:

```
void main (void) {
```

```
char buf [0x100];
int i, j;
PORTB = 0;
for (i = 0; i < sizeof(buf); i++) {
  CFG = 0x80; PORTB = 4;
  while (CFG! = 0);
  buf[i] = PORTA;

PORTB = 0;
```

| The follow | ving applies: |
|-----------------|---|
| $\ominus \pm$ | It implements an interrupt synchronization. |
| \oplus \pm | It implements a Gadfly synchronization. |
| \ominus \pm | It implements a Realtime synchronization. |
| $\ominus \pm$ | It implements no synchronization. |
| = | |
| \oplus \pm | A function which modifies its own code is not reentrant. |
| $\ominus \pm$ | A function which calls an interrupt service routine is not reentrant. |
| $\ominus \pm$ | Recursive functions are always reentrant. |
| | Interrupt service routines are always reentrant if they do not call another routine. |
| - - | |
| \oplus \pm | relative short answer times are typical for interactive systems. |
| ⊕ ± | reactive systems are common in systems which do measurement and control. |
| # ± | transforming systems are typically optimized for high throughput. |
| \oplus \pm | an example for an transforming system could be a network router. |
| | |
| $\ominus \pm$ | Realtime means to produce a result as fast as possible. |
| | A computer is realtime, if is is able to produce at average system load the correct result as fast as possible. |
| $\ominus \pm$ | For realtime it is sufficient to have an accurate timing system. |
| $\ominus \pm$ | An RTOS is required for a realtime system. |
| | |
| $\ominus \pm$ | A reentrant function shall not be interrupted. |
| | Interrupt functions does not have to be reentrant, but all functions called from that interrupt routine. |
| | A function which modify itself is reentrant, as long the self modification happens with disable interrupts. |
| \oplus \pm | On the ARM Cortex-M0+/M4F the usage of local stack variables does not |

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violate reentrancy.

Explain in a few words the reasons why a switch (like a button) needs a resistor. Illustrate it with a small drawing.

Solution: Without a resistor, the signal is undetermined or open. It is needed to define the signal to a defined level (low or high) if the switch is not closed.

List important points to be considered for the implementation of an ISR:

Solution: As short and as fast as possible. Need to acknowledge the interrupt at the beginning (or additionally at the end. Need to care about reentrancy, and that shared functions are reentrant and shared variables are protected.

Question 72......Points: [2]

Explain two different ways how a microcontroller can implement interrupts:

Solution: Using a vector table: the interrupt source gets translated into a vector number which is used as index into a vector table. The other way is that the processor directly jumps to an address and executes the code there (stub based approach).

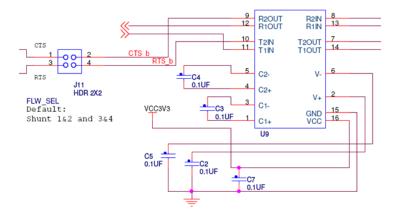


Figure 16: IC U9

- Motor H-Bridge IC.
- O Quadratur Decoder IC.
- Analog/Digital Converter IC.
- $\sqrt{\text{RS-232}}$ Level Shifter IC.
- Accelerometer IC.

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| Question 74 |
|---|
| What happens if two developers work at the same project but in different files or at different parts in one file and commit on Git? |
| Solution: Git merge automatically the changes together. |
| Question 75 |
| Solution: SVN: Centralized VCS, Git: Distributed VCS. |
| Question 76 |
| Solution: There will be a conflict which cannot be resolved by the VCS. The developers have to solve this problem manually (solving the conflict). |
| Question 77 |
| Solution: Every file is load in the repository except those are mentioned in the .gitignore. |
| Question 78 |
| Solution: It's not necessary, because the documentation is generated out of the code. Apart from that, there would be too many unnecessarily data on the repository. |
| Question 79 |
| Solution: TS: Video encoder, RS: PID Controller, IS: ticket selling machine |
| Question 80 |
| #define CALC1 (2+5) #define CALC2 (5*3) |

| Solution: To ensure proper usage in other macros or in calculations. $10*CALC1$ is not the same as $(10*2)+5$ |
|--|
| Question 81 |
| Solution: Advantages: Faster code, smaller code. Disadvantages: Interface, Encapsulation, Debugging |
| |
| Question 82 |
| Solution: |
| |
| #ifndefLED_H_ |
| #defineLED_H_ |
| <pre>/* content of header file */ #endif</pre> |
| #GHUII |
| Question 83 |
| Solution: In a header file with *.h extension. Using extern for variable declaration, e.g. extern int LED_global; I use a good name with prefix because the name is visible in the whole project. |
| Question 84 |
| Solution: Yes, any text files can be included (as long as understood by the compiler). For example I can have an array of hex values in an array and then include it as bitmap.txt. |
| Question 85 |
| Solution: HTML, RTF, PDF, LaTeX |
| Question 86 |

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| Solution: Doxygen extracts the documentation from the source file and from the comments embedded in the files. |
|---|
| Question 87 |
| Solution: It blocks further execution. |
| Question 88 |
| Solution: Each ISR needs to ensure that the other ISR is not executed. This can be with priorities, or with disabling the interrupt for the other ISR. |
| Question 89 |
| Solution: Each ISR needs to ensure that the other ISR is not executed. This can be with priorities, or with disabling the interrupt for the other ISR. |
| Question 90 |
| Solution: Disable or mask them, or have a dummy ISR. |
| Question 91 |
| Solution: Text file with code (encoded as ASCII in a defined format). |
| Question 92 |
| Solution: Synchronous (periodic timer interrupt, periodic task output) and asynchronous (button pressed, transceiver packet received). |
| Question 93 |
| Solution: Marking the end of a list. Additionally it is used to calculate the size of the array or to compute the number of items. |
| Question 94 |

| Solution: Advantage: simple. Disadvantages: long if-elsif-else/switch in the event handler, order of events has impact about priority, need mutual exclusion for shared data. |
|--|
| Question 95 |
| Solution: ARM Cortex-M0+ |
| Question 96 |
| Solution: Main Stack Pointer (MSP) and Process Stack Pointer (PSP). |
| Question 97 |
| Solution: Absolute (e.g. 13:50) or relative (e.g. after 50 ms). |
| Question 98 |
| |
| Solution: 8 MHz crystal oscillator. |
| |
| Solution: 8 MHz crystal oscillator. Question 99 |
| Solution: 8 MHz crystal oscillator. Question 99 |
| Solution: 8 MHz crystal oscillator. Question 99 |
| Solution: 8 MHz crystal oscillator. Question 99 |
| Solution: 8 MHz crystal oscillator. Question 99 |

| Solution: functional, hierarchical, else-if state machine. |
|--|
| Question 103 |
| Solution: if-elsif-else, switch, table |
| Question 104 |
| Solution: determined by current state and current output. |
| Question 105 |
| Solution: Handle multiple 'interrupt like things' with just one timer. |
| Question 106 |
| Solution: yes. |
| Question 107 |
| Solution: To avoid multiple interrupts, and to have reached a stable state. |
| Question 108 |
| Solution: hardware (R-C) or software (timer, delay). |
| Question 109 |
| Solution: USB is a much more complex protocol, and therefore, has a bigger overhead. If the USB performance is not needed, a simpler UART connection is more efficient. At the Freedom board, there is an extra MCU to convert the USB to a UART protocol. This means that there is an extra MCU to take over the USB overhead. |
| Question 110 |

Solution: Serial means that it sends the data bits after bits on a single data line (as a sample Rx or Tx line). Asynchronous means that there is no clock supported to read the data. The start of the data has to be detected by the protocol (start

bits, stop bits). As which device class does the Freedom board enumerate at the PC, and what is the purpose of it? Solution: CDC (Communication Device Class), the PC opens a virtual COM port for the device. This enables to have a connection, which acts like a UART connection. Furthermore, it also enumerates as Mass Storage Device, this mode can be used to update the firmware. What's the command parser table? **Solution:** It's a list with function pointers. In this list the parser choses the method which is going to be executed. What's the difference between Memory Scheme 1 and 2 in FreeRTOS? Solution: Memory Scheme 1 only allocates Memory. It's not possible to delete Tasks. With Scheme 2 it's possible to free space and reuse. There would be another Scheme where you can merge freed blocks. Provide a good example how FreeRTOS queues can be used between multiple tasks: **Solution:** To send messages between tasks. What is the difference between xQueueReceive() and xQueuePeek() in FreeRTOS? Solution: xQueuePeek() only checks if there is an item in the queue. xQueue-Receive() does remove the item. What are the H-bridges of the motors needed for?

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Reached:

| Solution: They are used to set the direction of the engine by simply twisting the two wires of the motor. |
|--|
| Question 117 |
| Solution: There is only one Timer for all the different PWM-channels and therefore the frequency can only set for this common timer. The different channels only have a separate value register which all compare with the same common timer. |
| Question 118 |
| Solution: The coast modus of the engines is only possible if "Mode 0" is used on the motor drivers. On the SUMO robot, this pin is not routet to the FRDM board and just hardware wise set to HIGH (Mode 1). |
| Question 119 |
| Solution: Measure the accelerometer value for a known value, e.g. 0g or 1g. Compare it with the value it should report. The difference is the offset correction. |
| Question 120 |
| Solution: Typical interfaces are analog with ADC and digital with SPI/I ² C. |
| Question 121 |
| Solution: More write cycles, durable, defined read/write times. |
| Question 122 |
| Solution: Add a checksum for the blocks. |
| Question 123 |

Solution: The Binary Code several bit changes between the counting steps are possible whereas the Grey Code changes only one bit between each step. It's possible to build the Grey Code recursive and it is permutable.

Solution: Interrupts, Sampling, input capture, dedicated quadrature peripheral/IC

Solution: The sampling method causes a constant system load whereas the interrupt method causes a speed dependent system load. It's easier to handle a constant system load to test a system.

Solution: If you use the interrupt method the system freezes (to many interrupts). If you use the sampling method the system still works normally, but can't recognize every encoder step (errors).

Question 127......Points: [1] What are the input and output signals of QuadCounter.c?

Solution: The input signals are the two hardware signals and the output signals are the number of errors and the current position/steps.

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| Question 128 |
|---|
| Solution: It is used as a sentinel, so the code can iterate through the table until there is a NULL entry. |
| Question 129 |
| Solution: sizeof() returns the size of the string including zero byte, but we want to compare the string without it. |
| Question 130 |
| Solution: baudrate, number of data bits, stopbit, parity |
| Question 131 |
| Solution: volatile |
| Question 132 |
| Solution: Because they operate with different speeds. |
| Question 133 |
| Solution: Because an interrupt can happen any time, it must be ensured that there is no data corruption. |
| Question 134 |
| Solution: No, for example the HCS08 does not have nested interrupts. |
| Question 135 |

| Solution: Reset, NMI, HardFault |
|--|
| Question 136 |
| Solution: Yes, e.g. by a Reset or NMI as they have numerically lower interrupt priorities. |
| Question 137 |
| Solution: Measure the delta-time (duration of period) if frequency is low. Use counting the steps or periods if time intervals are short. |
| Question 138 |
| Solution: The speed can be estimated more frequently than the measurement interval. Additionally it allows an averaging of the speed. With the ring buffer a configurable and dynamic time span for the estimate can be used. |
| Question 139 |
| Solution: Is the system control value (e.g. PWM for the motor) limited in range, then the integral part can go out of this range. Then the integral sum will increase without having an additional impact on the control value, and delays the impact of the integral on the control value. Additionally the limit avoids a numerical overflow of the integral sum. |
| Question 140 |
| Solution: Nyquist/Shannon: need to sample it with at least 50 μ s. |
| Question 141 |

Reached: _____

Solution: There are several options: one way is to create a drive task and call the PID from there. Or to use a periodic interrupt to call the PID.

Why is it not possible to directly measure the output signal of the optical quadrature encoder we have used?

Solution: The signal is an analog sinus-like signal, without a 50%-50% high-low signal distribution. To get a clean quadrature signal a DAC with comparator devices are used.

Discuss the pros and cons of using either sampling or interrupt method for a quadrature signal:

Solution: Usually sampling is usually preferred as it will create a constant system load. However, this creates a high load of the system even if the wheel is not moving or only slowly moving. Using interrupts can cause problems if there are too many interrupts. It would allow a low system load if the wheel is not moving or slowly moving.

What fundamental problem exists for absolute position encoders, and how can it be solved:

Solution: Because of mechanical tolerances, multiple bits can change from one step to another. The solution is to use a Gray code/encoder, as with this only one bit changes from one sector to another. In addition the Gray code is cyclic and is therefore ideal for wheel position measurement

Can you list the main features of the MCP4728?

Solution: 12bit DA-Converter, includes an EEPROM to store DAC values and settings, I²C bus and protocol, 4 DAC output signals.

What are the special things or attributes of the Gray code?

Solution: Hamming distance of 1 (only one bit changes), the code has permutation (every code only occurs once), it is cyclic (last and first code confirm to the rules) and recursive (codes of lower order are embedded in code of higher order), and it is simple to transform a binary code into a Gray code.

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| Question 147 |
|--|
| If the robot moves with a speed of 1 m/s, and you measure the reflectance sensor with 100 Hz, what would be an estimated way distance over the white line until you detect the white sumo line in the application? |
| Solution: If robo is moving with 1 m/s and we measure with 100 Hz, then the robot will move in average 2 cm until the motors get stopped. |
| Question 148 |
| The reflectance sensor has two red LEDs to indicate if the sensor is on. For the red LEDs there is a 1K Ohm resistor in series to limit the current through the LED. But why is there another 220 Ohm resistor in series to the LED with that 1K Ohm resistor? |
| Solution: To limit the current through the photo diodes. |
| |
| Question 149 |
| Solution: This generates a digital signal we can measure, and we do not need an A/D converter pin. |
| Question 150 |
| Solution: Predictability, precise timing, speed. |
| Question 151 |
| Solution: Running multiple things in parallel, using RTOS services (queues, semaphores, mutex,), scalability of application. |
| Question 152 |
| Solution: Preemptive: the RTOS distributes the processing time, tasks get suspended by the RTOS. Non-preemptive: the task are cooperative, they pas the control back to the kernel. |
| Question 153 |
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| Solution: To maximize the CPU utilization among different tasks. |
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| Question 154 |
| Solution: The RTOS needs to adhere to strict timing and needs to produce output with given timing constraints. |
| Question 155 |
| Solution: IDLE, PRESSED, RELEASE. |
| Question 156 |
| $\textbf{Solution:} \ \ Software (state machine, time delay/low pass filter) or hardware (capacitor, Schmitt-Trigger.$ |
| Question 157 |
| Solution: Bouncing is a mechanical problem. To process the state of a bouncing push button, it needs first to stabilize. |
| Question 158 |
| Solution: Add additional events to message an inter-click. Extend the state machine so it either continues or uses different states in the state machine. |
| Question 159 |
| Solution: The fast decay principle is to revert the current from the previous movement. E.g. if the motor is turning forward, to put current into the H-Bridge to in reverse order (turning it backward). This will bring the energy stored in the inductor down fast. So the 'fast' is about how fast the current reaches zero. With 'slow decay' either the lower half or upper half of the transistors are on, allowing the inductive current to flow back. Because this takes longer than with the 'fast' method, this is called 'slow decay'. With 'fast decay', the motor coasts down the speed, while with 'slow decay' the H-Bridge using an active break. Hint: . |
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| Question 160 |
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| Solution: No, this depends if the signal is LOW or HIGH active. |
| Question 161 |
| Solution: Increase the PWM frequency. |
| Question 162 |
| Solution: A half-H Bridge has two transistors, while a full H-Bridge has 4 transistors. With a full H-Bridge the direction of a DC motor can be changed. |
| Question 163 |
| Solution: Files which can be used for multiple projects, like files of a library. |
| Question 164 |
| Solution: A lot of duplicated files, maintenance problem could occur as I need to change files in different places. |
| Question 165 |
| Solution: easier to read, re-usability, easier to understand. |
| Question 166 |
| Solution: Either with '//' or e.g. with '{PARENT-1 PROJECT_LOC} |
| Question 167 |

| Solution: Because with sharing the project the paths might be different on another machine. |
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| Question 168 |
| Solution: To debug it a preprocessor listing needs to be generated (-E for gcc). |
| Question 169 |
| Solution: With double quotes 'user' header files are included, and with <> library header files. The compiler uses two different search settings for user and library header files. |
| Question 170 |
| Solution: Using the mainpage command. |
| Question 171 |
| Solution: ???? |
| Question 172 |
| Solution: The processor is using muxing: a single pin can have different purposes, like as touch sensing pin, as a input or output port pin or as a timer channel pin. |
| Question 173 |
| Solution: That it is only possible to use the pin in one mode at a given time. And that muxing a pin might have impacts and side effects on other pins, e.g. certain functions are not possible any more. |
| Question 174 |

| Solution: printf() is not a save function and can cause a stack overflow, and is subject of security issues. Additionally it needs a lot of code space and $stack/RAM$. |
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| Question 175 |
| Solution: Because the string is terminated with zero byte. |
| Question 176 |
| static int counter = 0; |
| Now you remove the static. What is the effect? |
| Solution: The variable now gets initialized with zero at every interrupt. |
| Question 177 |
| Solution: Pros: events are all handled in a single place and centralized. Complexity is simple. Cons: can be a long switch/if-else-if, event handling depends on frequency of main loop. Cannot handle events independently/concurrently. |
| Question 178 |
| Solution: Maximum core clock of KL25Z is 48 MHz, while the K22 can run up to 120 MHz. The K22 core is a Cortex-M4F and it can run at a higher speed than the Cortex-M0+. Additionally, the bus clock can run at maximum half of the speed of the core clock. |
| Question 179 |
| Solution: Pros: more accurate, higher overall speed possible. Cons: more costs, more PCB space needed. |
| Question 180 |

Solution: The CPU is clocked by the CPU clock, the bus clock is used for data transfer and to access memory, and the system clock is used to clock the peripherals.

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End of Example Exam Part A.

Reached: _____ from 0 points