

IFDS Framework

CMPUT 497/500 Foundations of Program Analysis

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Previously

- Inter-Procedural Data-Flow
- Inherited vs Synthesized Analysis Info
- Caller-Callee Relationships
- Valid/Invalid Paths
- Staircase of Calls and Returns
- Demand-Driven Analysis
- Types of Contexts
- Important Language Features

Interprocedural Finite Distributive Subset Problems

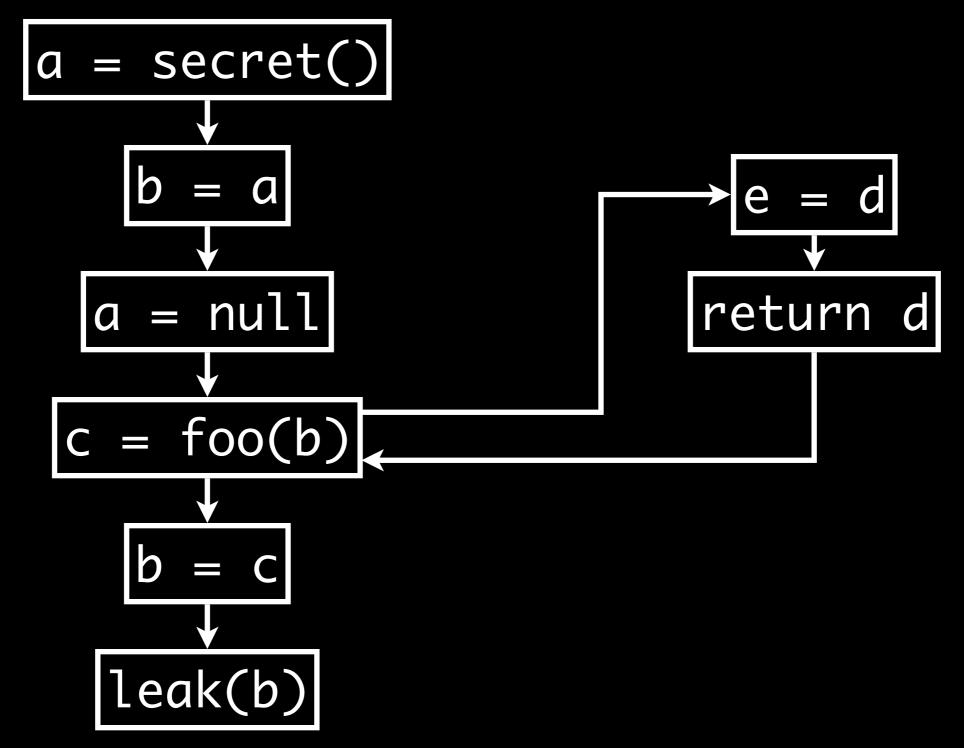
Thomas Reps, Susan Horwitz, and Mooly Sagiv. Precise Interprocedural Dataflow Analysis via Graph Reachability. ACM SIGPLAN-SIGACT Symposium on Principles of Programming Languages (POPL '95), pages 49–61.

Interprocedural Finite Distributive Subset Problems

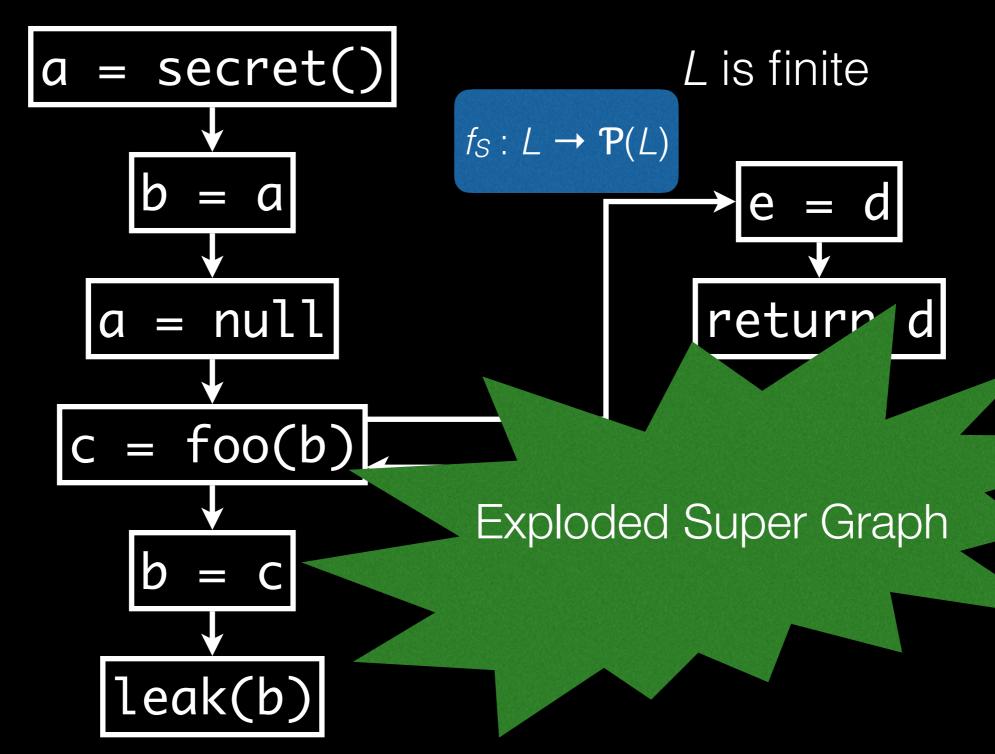
Interprocedural

```
main() {
    a = secret();
    b = a;
    a = null;
    c = foo(b);
    b = c;
    leak(b);
}
```

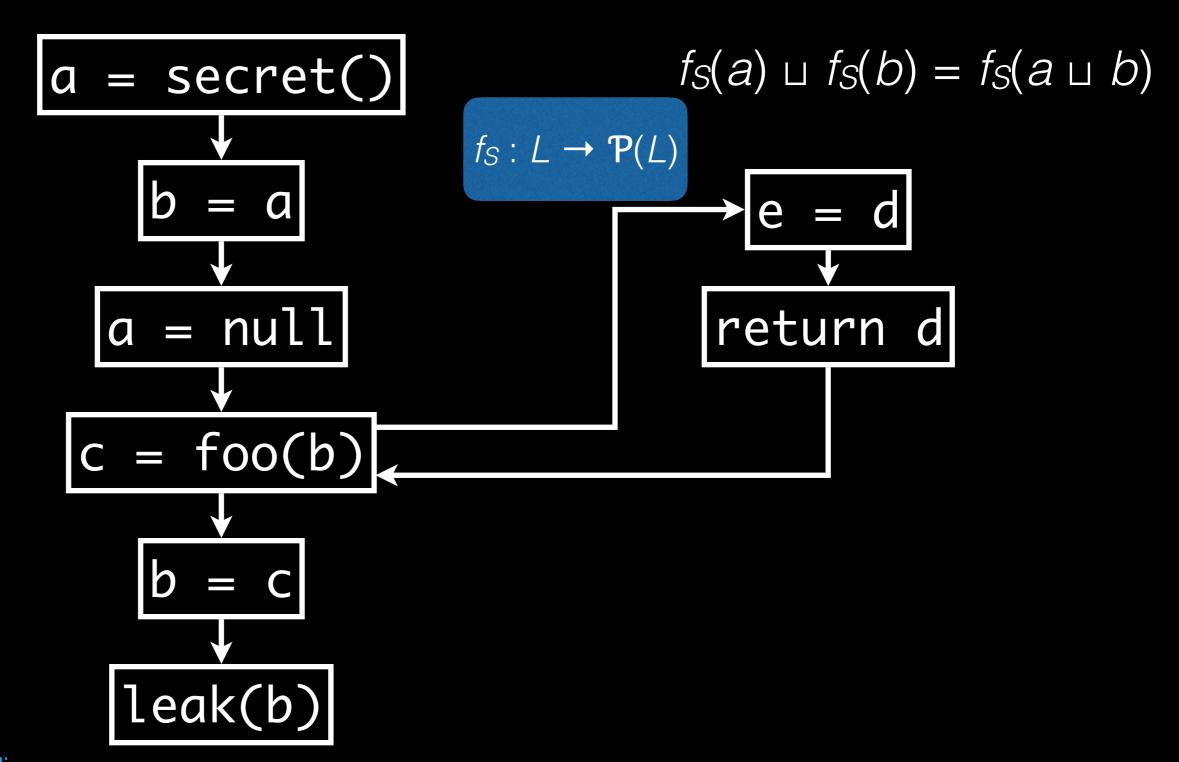
Interprocedural



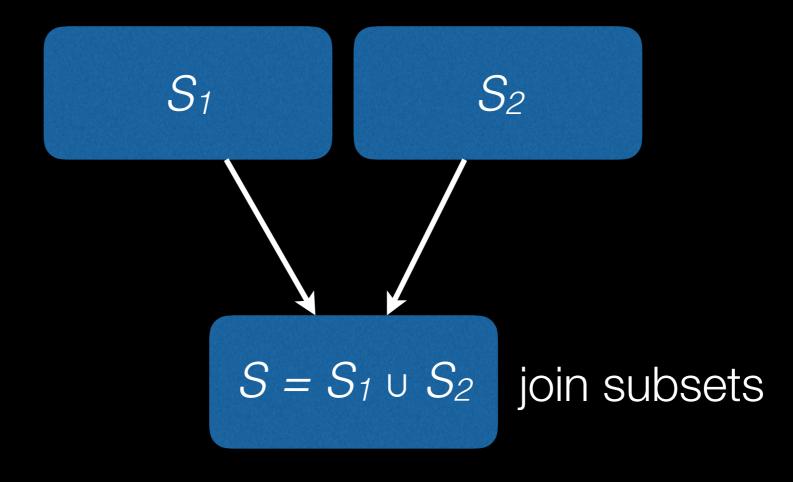
Finite



Distributive



Subset



alias taint typestate

IFDS

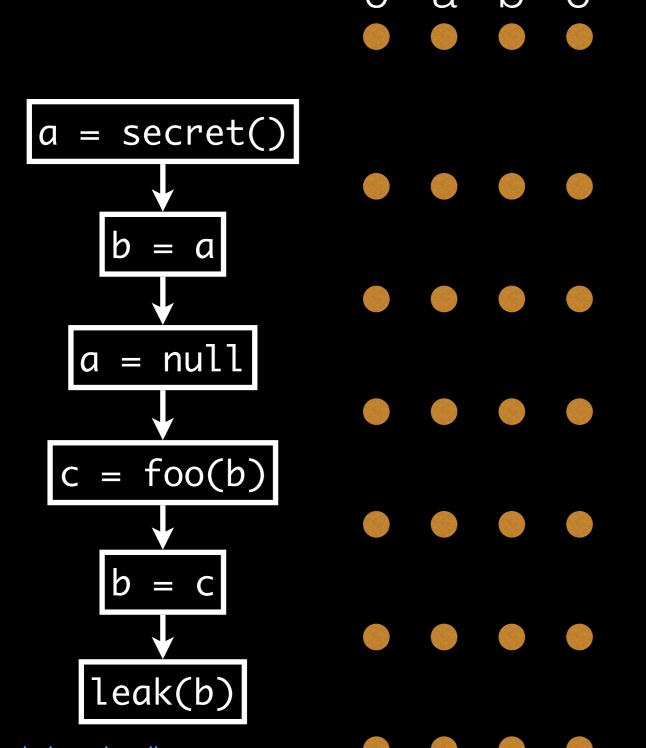
constant propagation

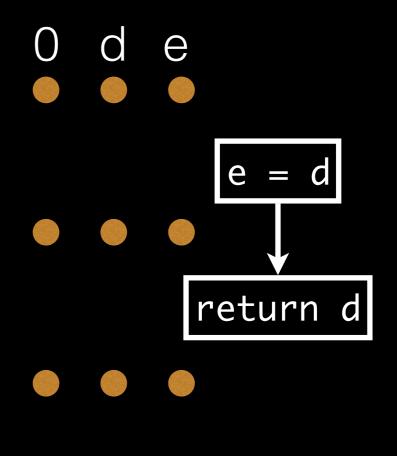
reaching definitions

uninitialized variables

Graph Reachability

Graph Reachability

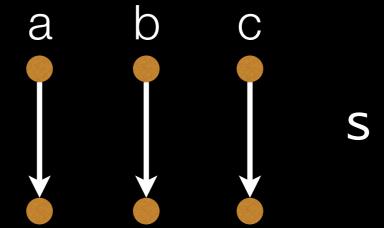




- each node is a fact
- fact holds at a statement => node is reachable in ESG

Examples of Flow Functions

Identity Flow Functions

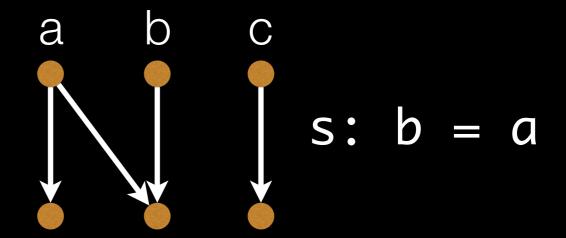


out(s) = facts after s

every fact is reachable iff it was previously reachable

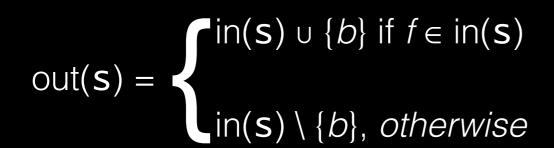
Flow Functions

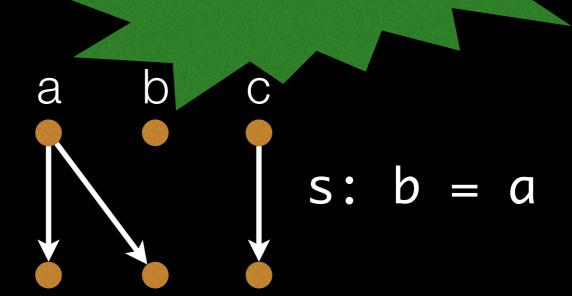
out(s) =
$$\begin{cases} in(s) \cup \{b\} & \text{if } f \in in(s) \\ in(s), & \text{otherwise} \end{cases}$$



every fact is reachable iff it was previously reachable, and b is also reachable if a was reachable

Flow Functions

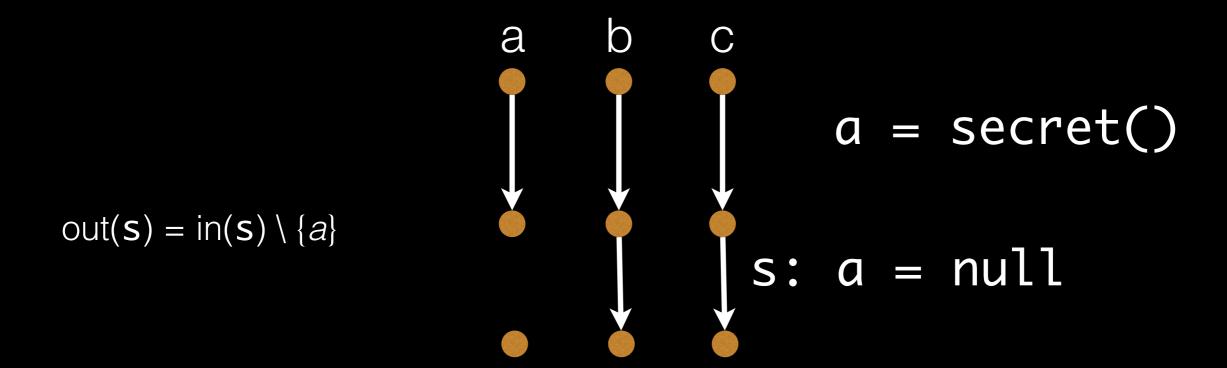




Taint Analysis

every fact except *b* is reachable iff it was previously reachable; *b* is reachable if *a* was reachable

"Killing" Facts

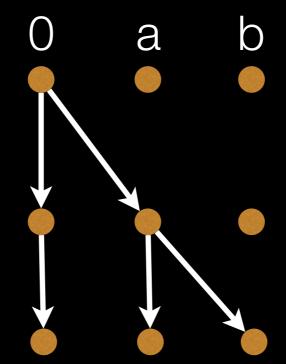


a is not reachable, even if it was before

Generating Facts

$$\operatorname{out}(\mathbf{s}) = \operatorname{in}(\mathbf{s}) \cup \{a\}$$

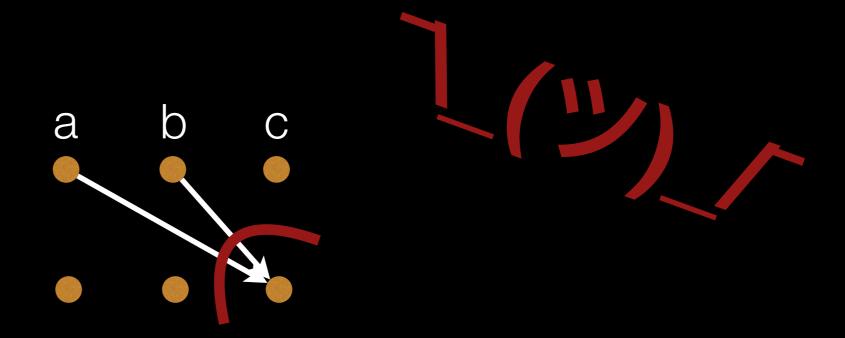
$$\mathsf{out}(\mathsf{s}) = \mathsf{in}(\mathsf{s}) \cup \{b\}$$



$$b = a$$

0 is the tautological fact => always reachable

Non-Distributive Flow Functions



e.g., full constant propagation c = a + b

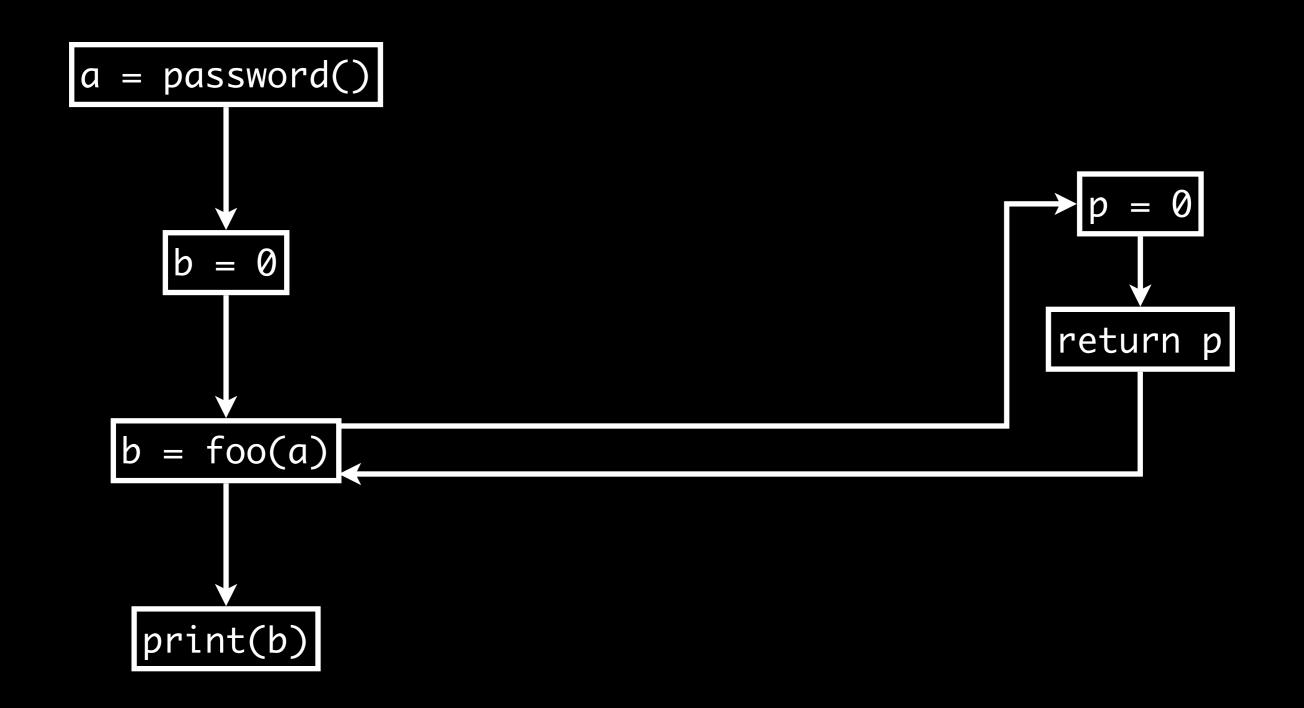
```
void main() {
    int a = password();
    int b = 0;

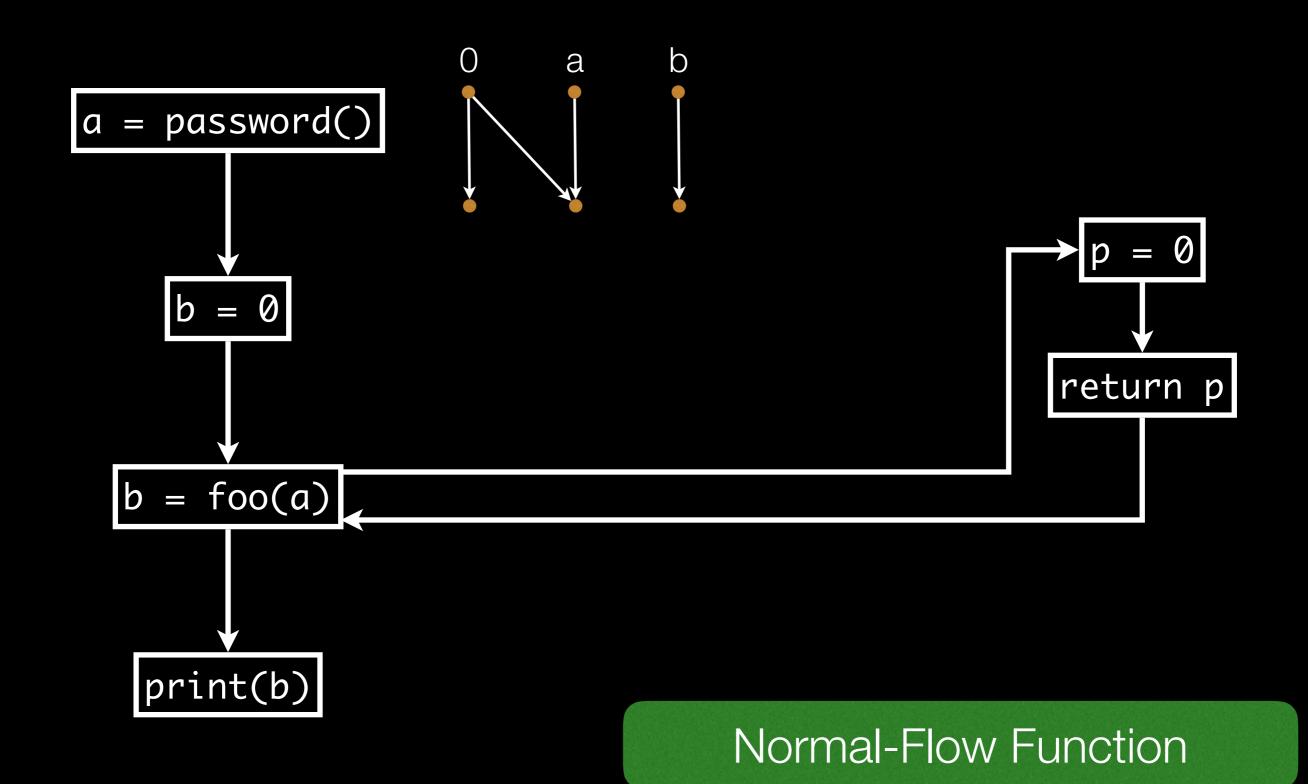
b = foo(a);

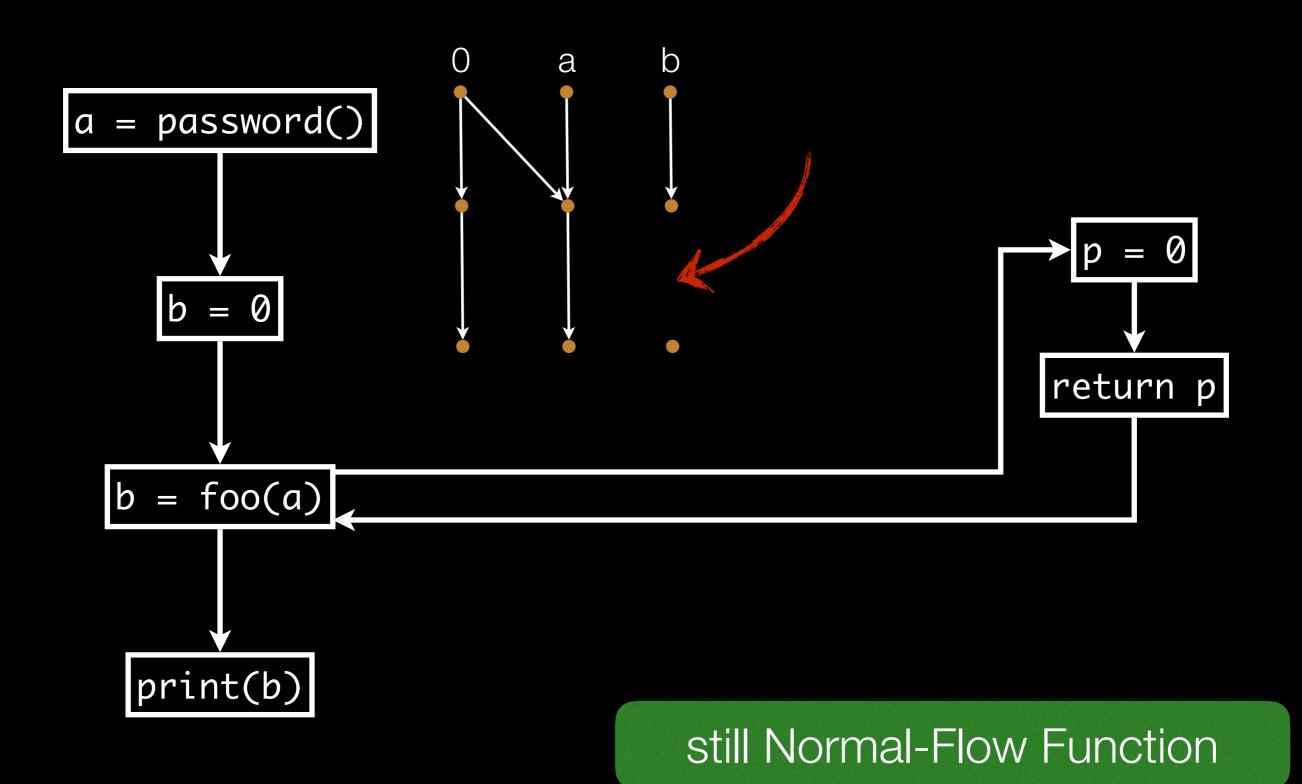
print(b);
}
```

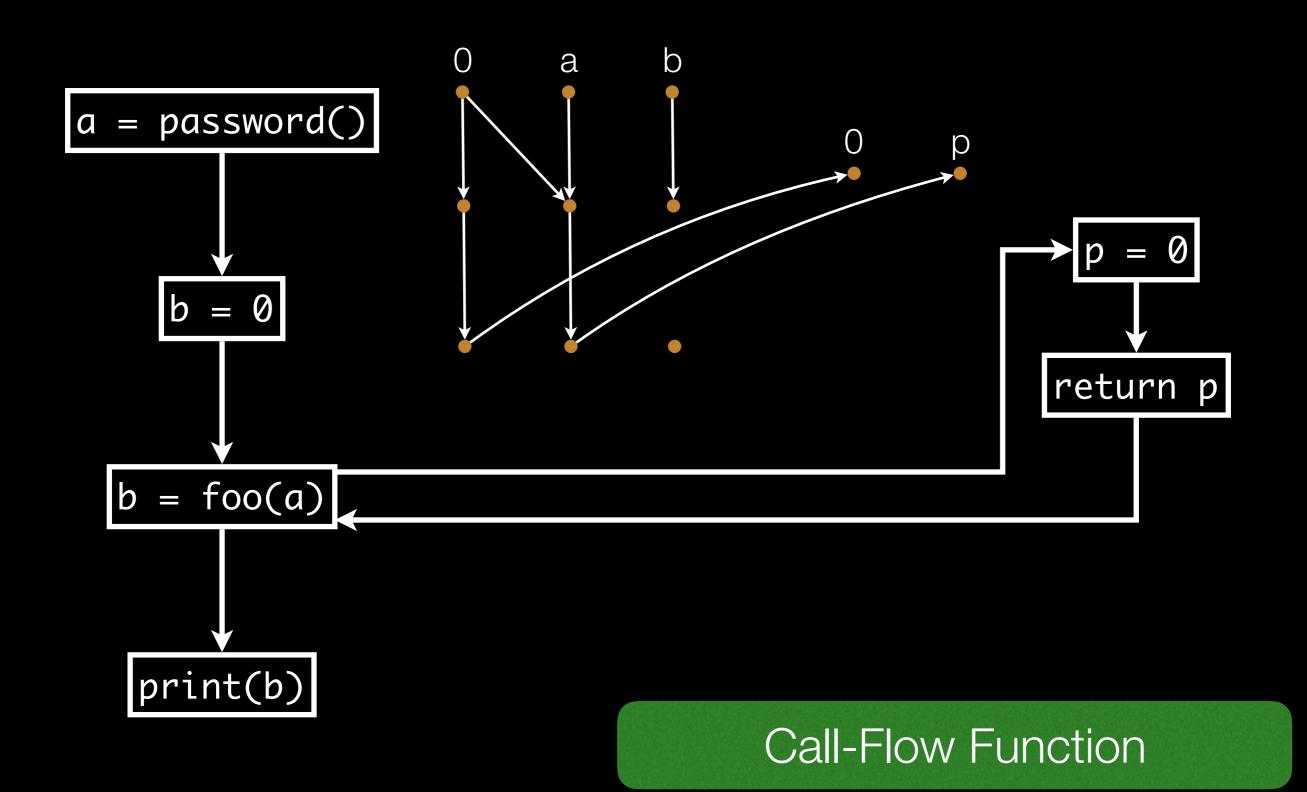
```
int foo(int p) {
   p = 0;

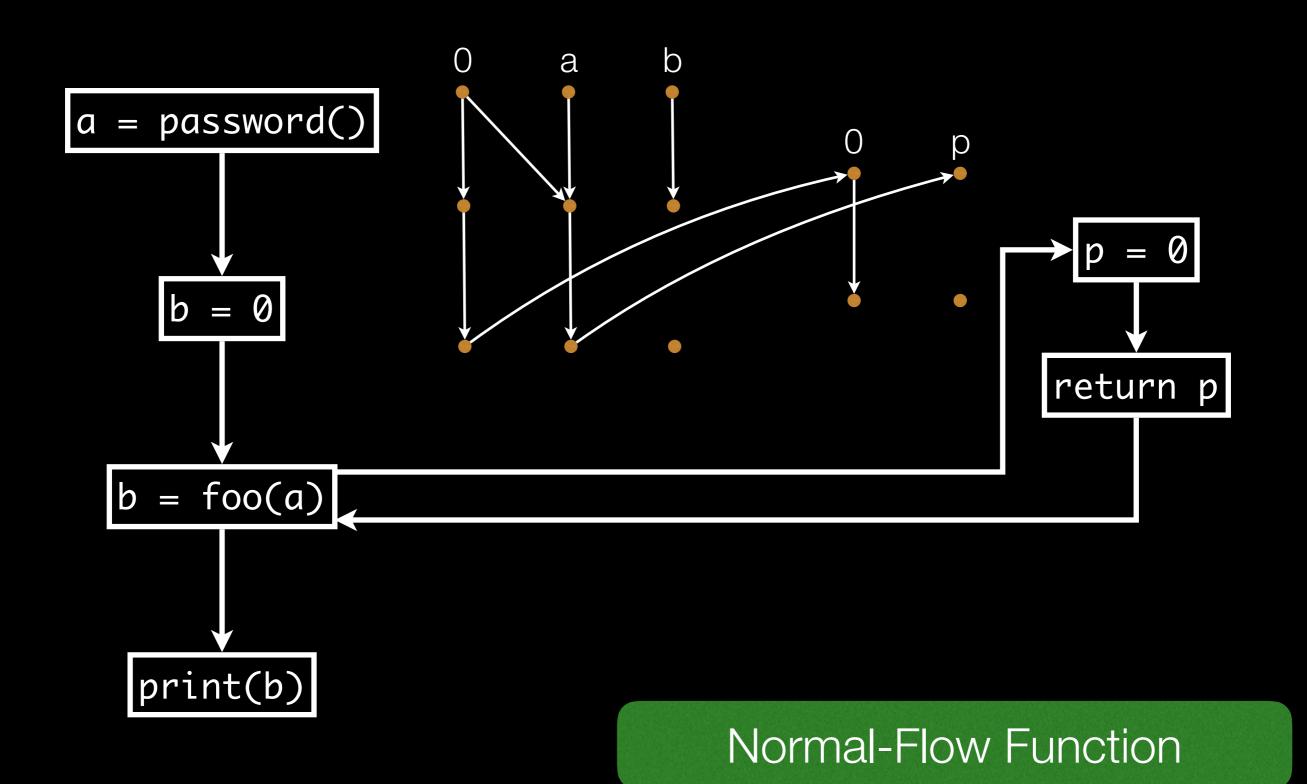
return p;
}
```

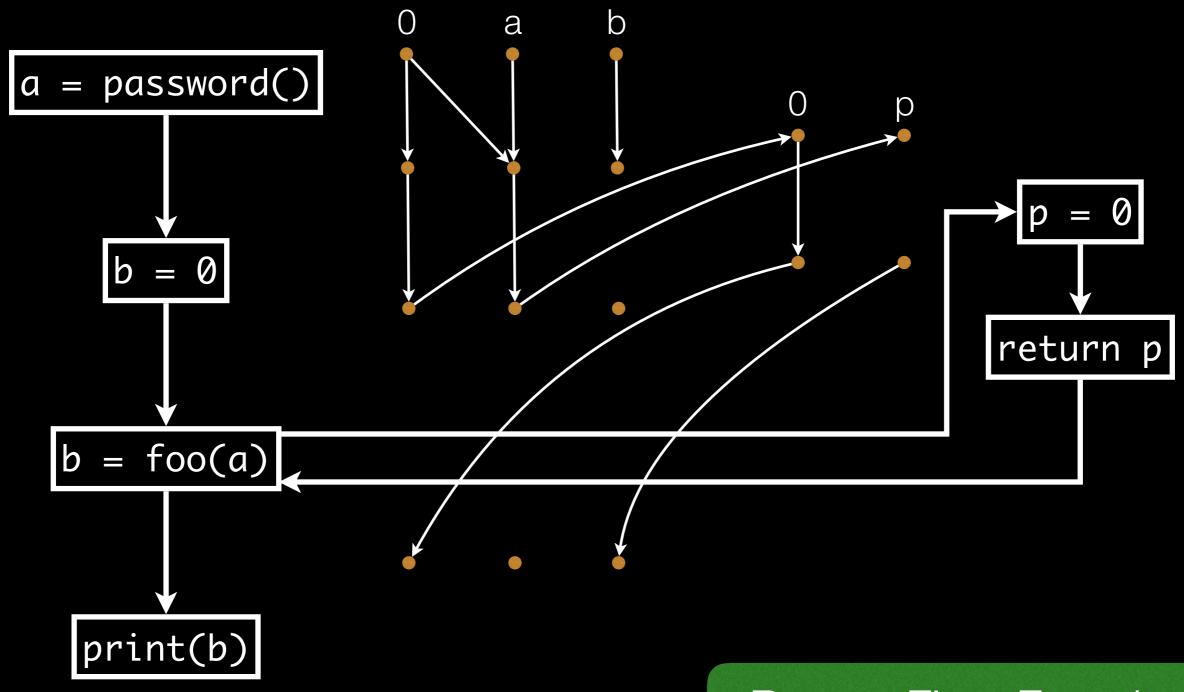




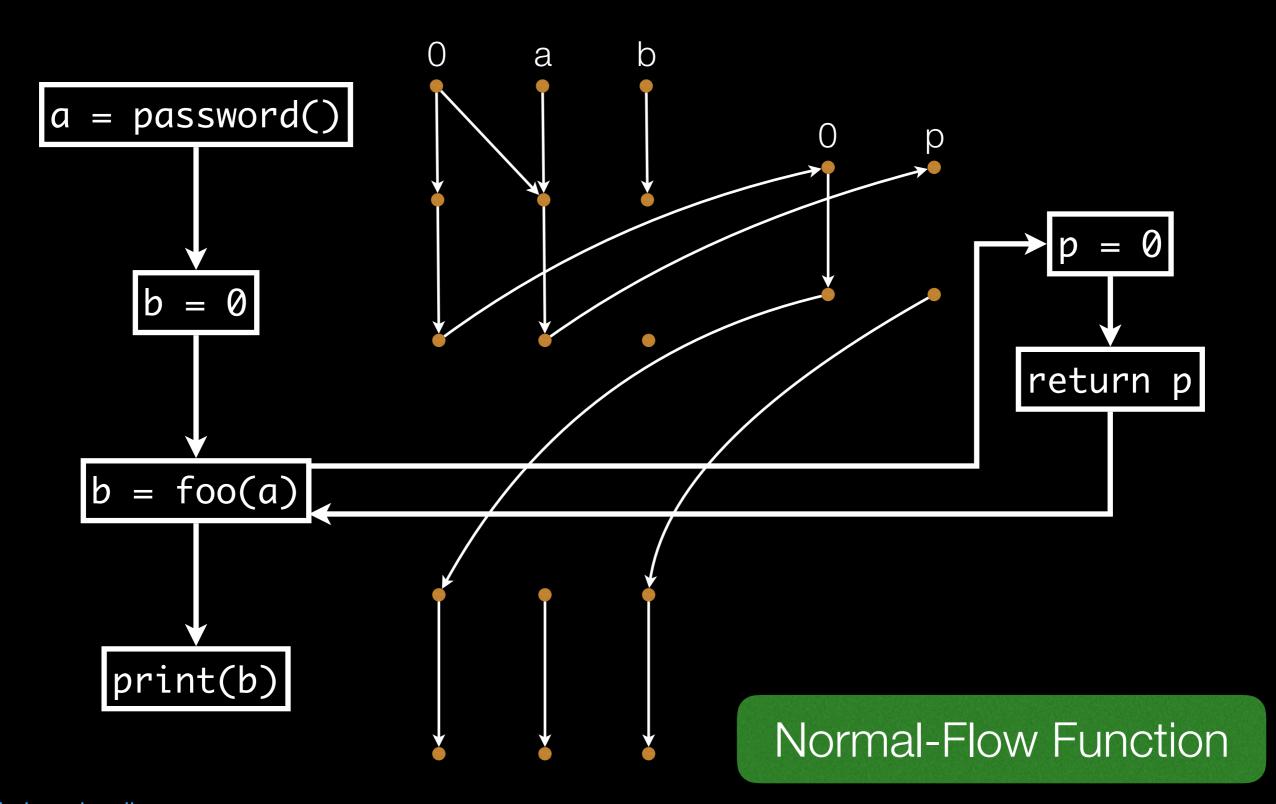


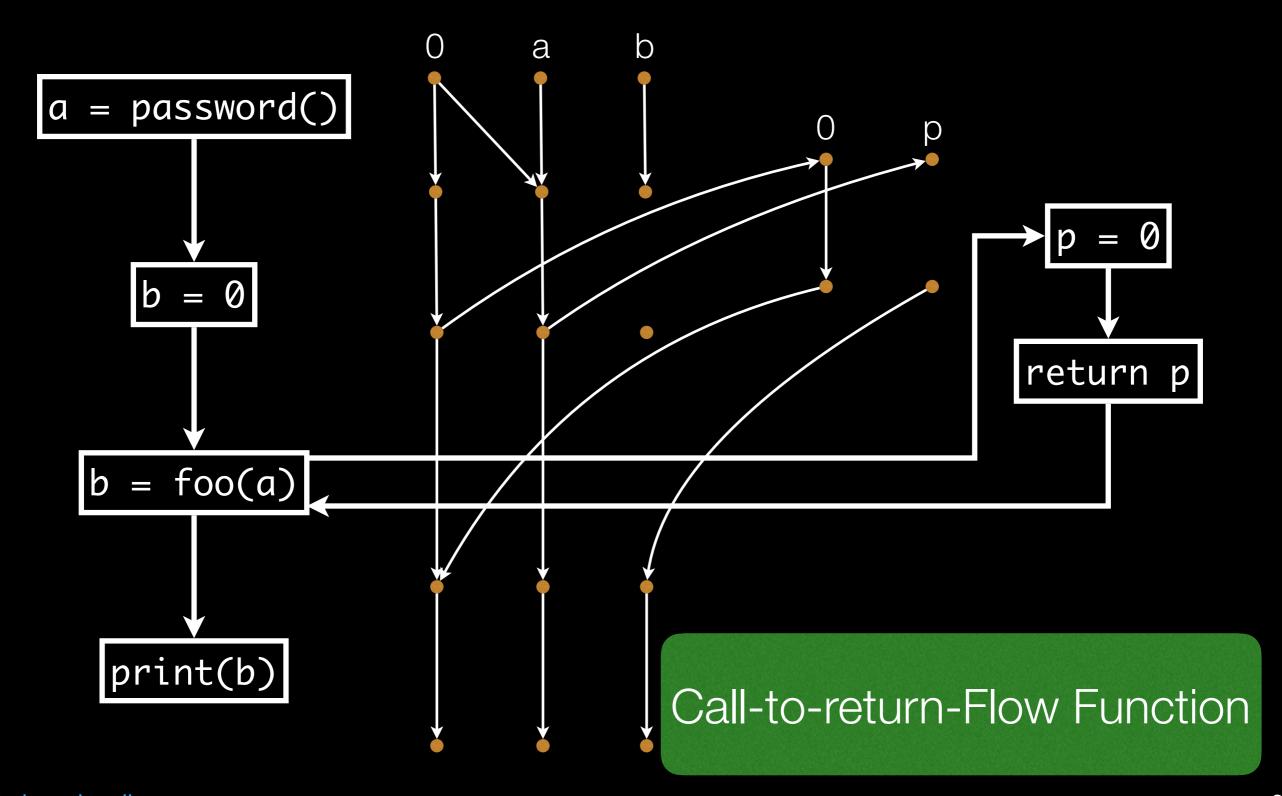






Return-Flow Function





Flow Functions

Call-Flow Function

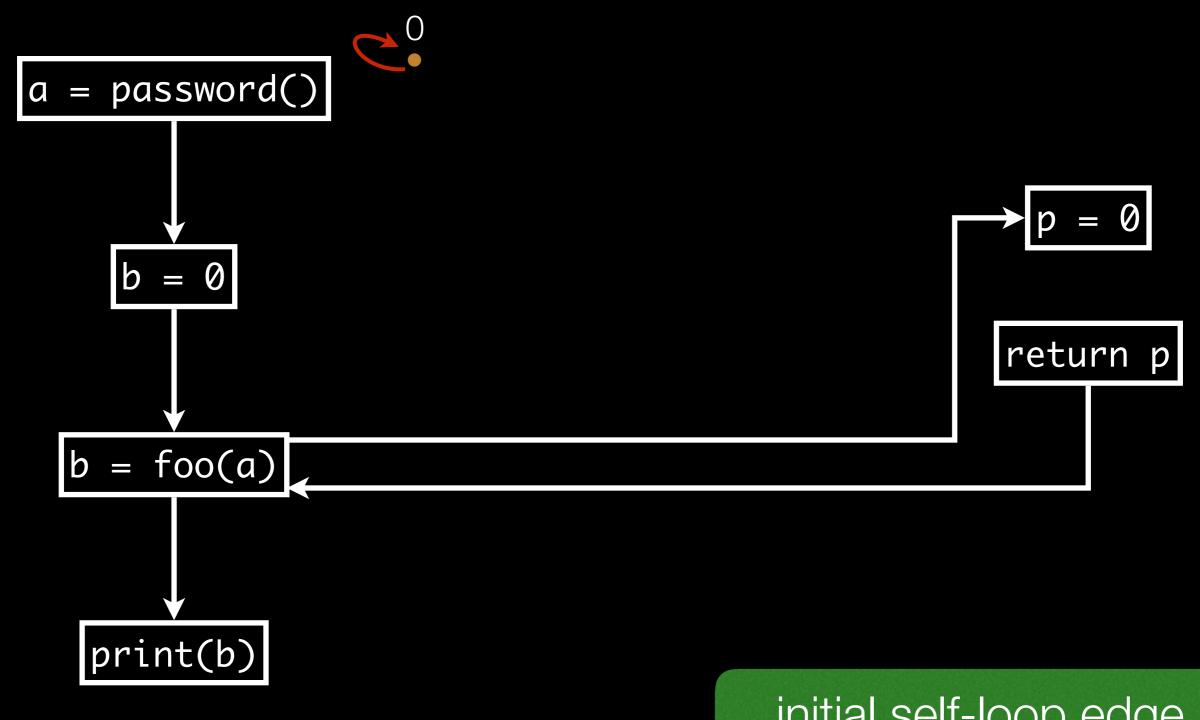
Normal-Flow Function

Return-Flow Function

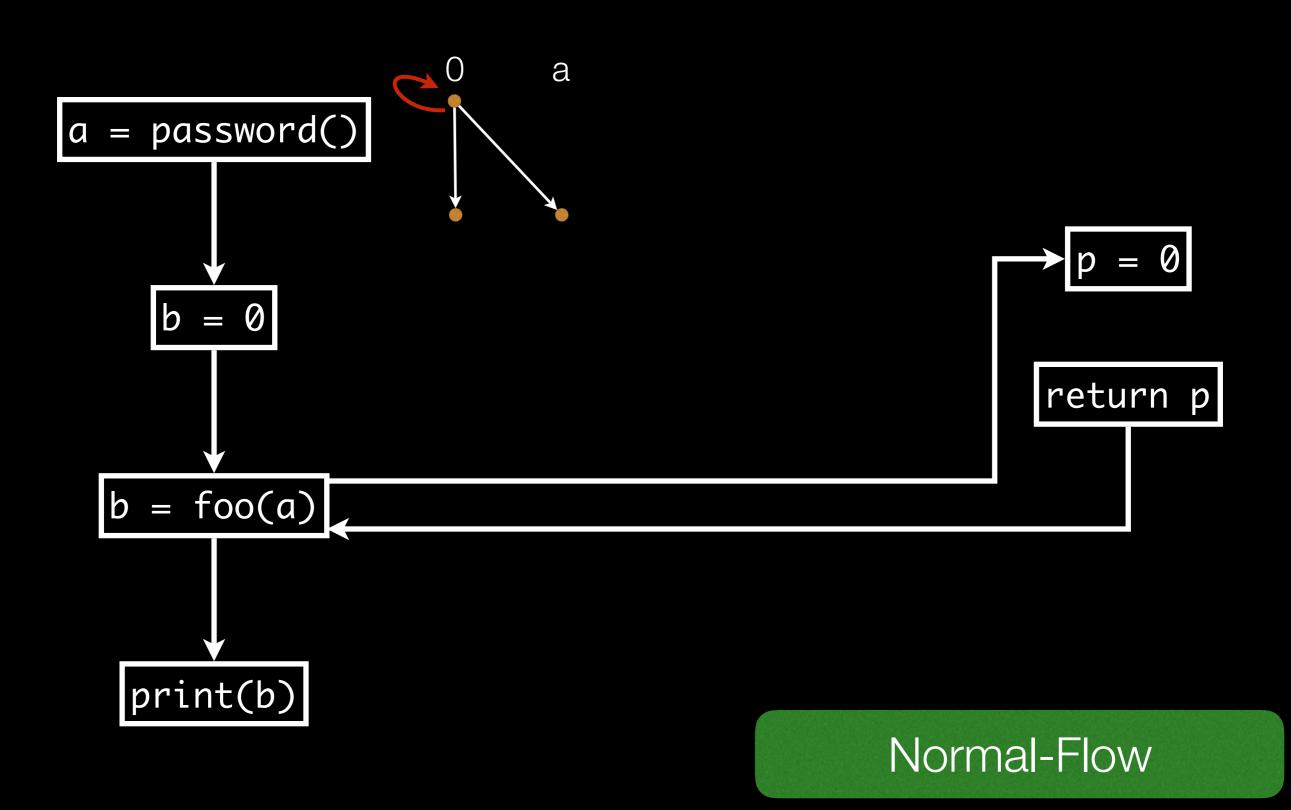
id

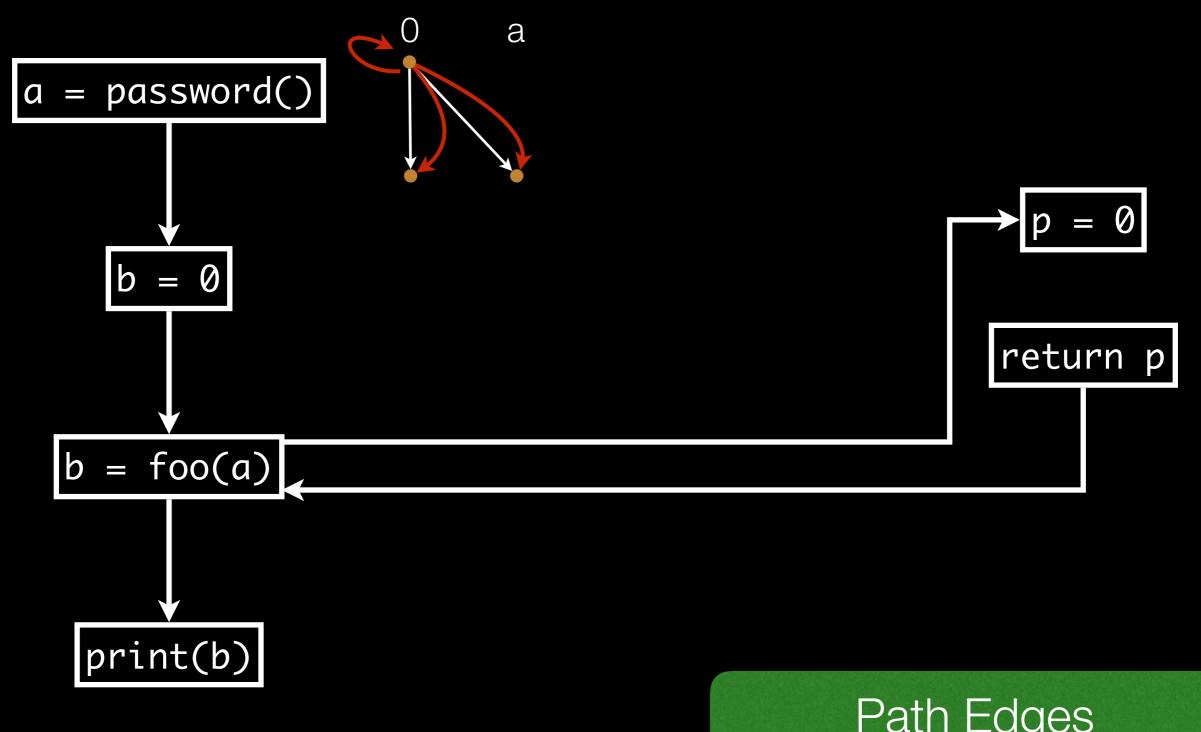
Call-to-return-Flow Function



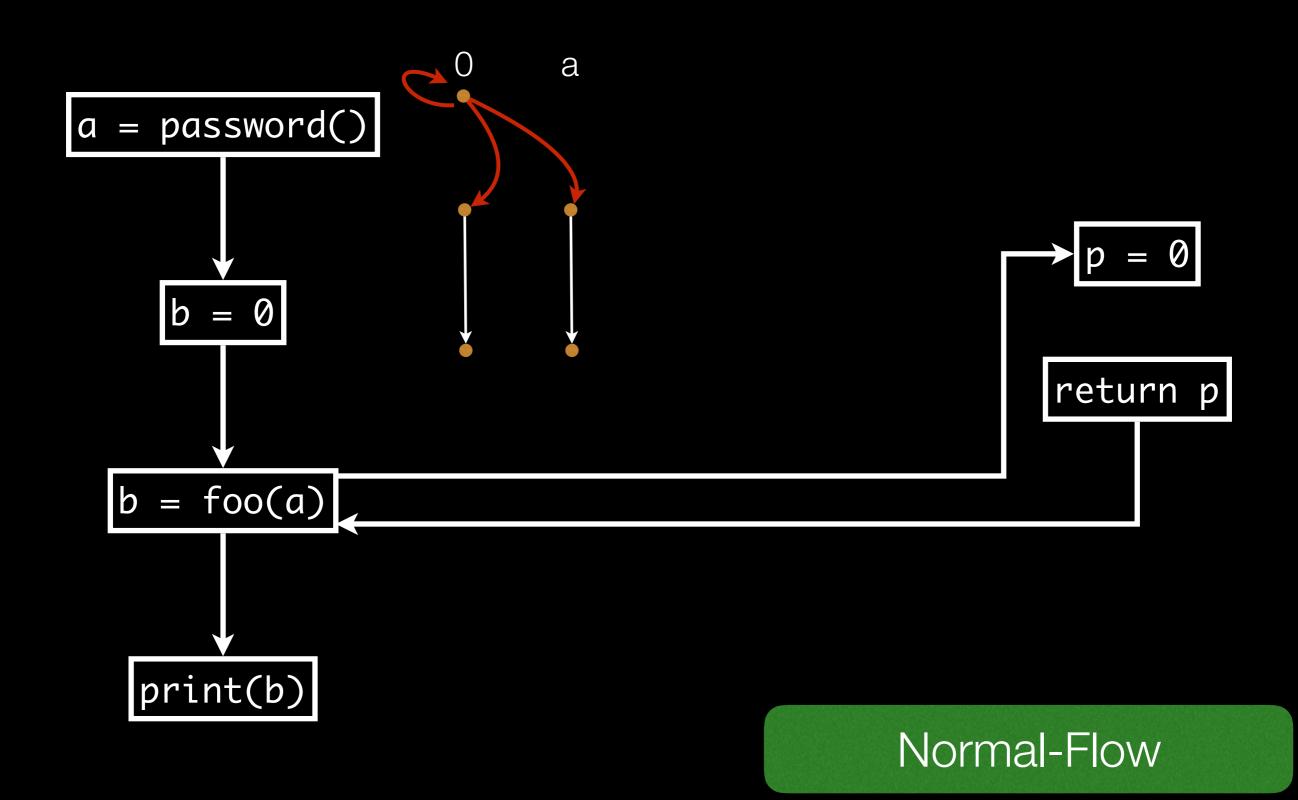


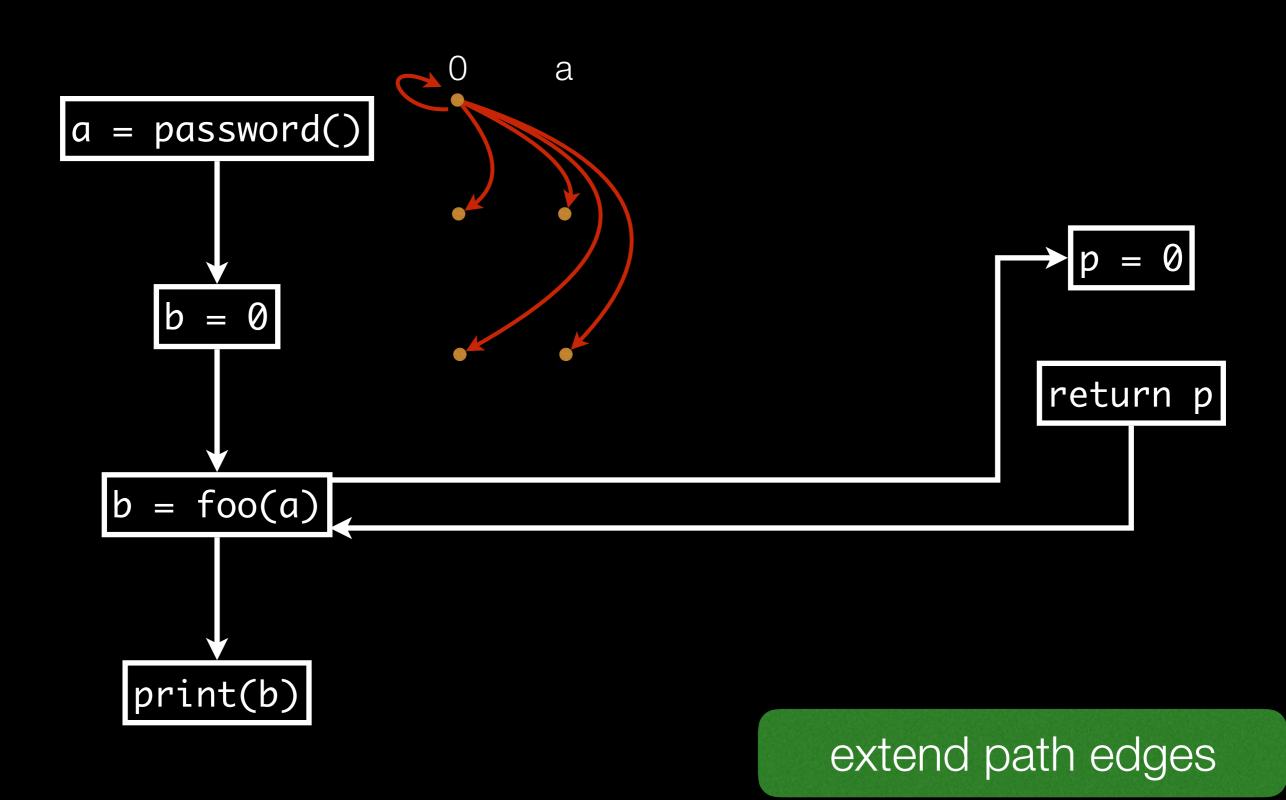
initial self-loop edge

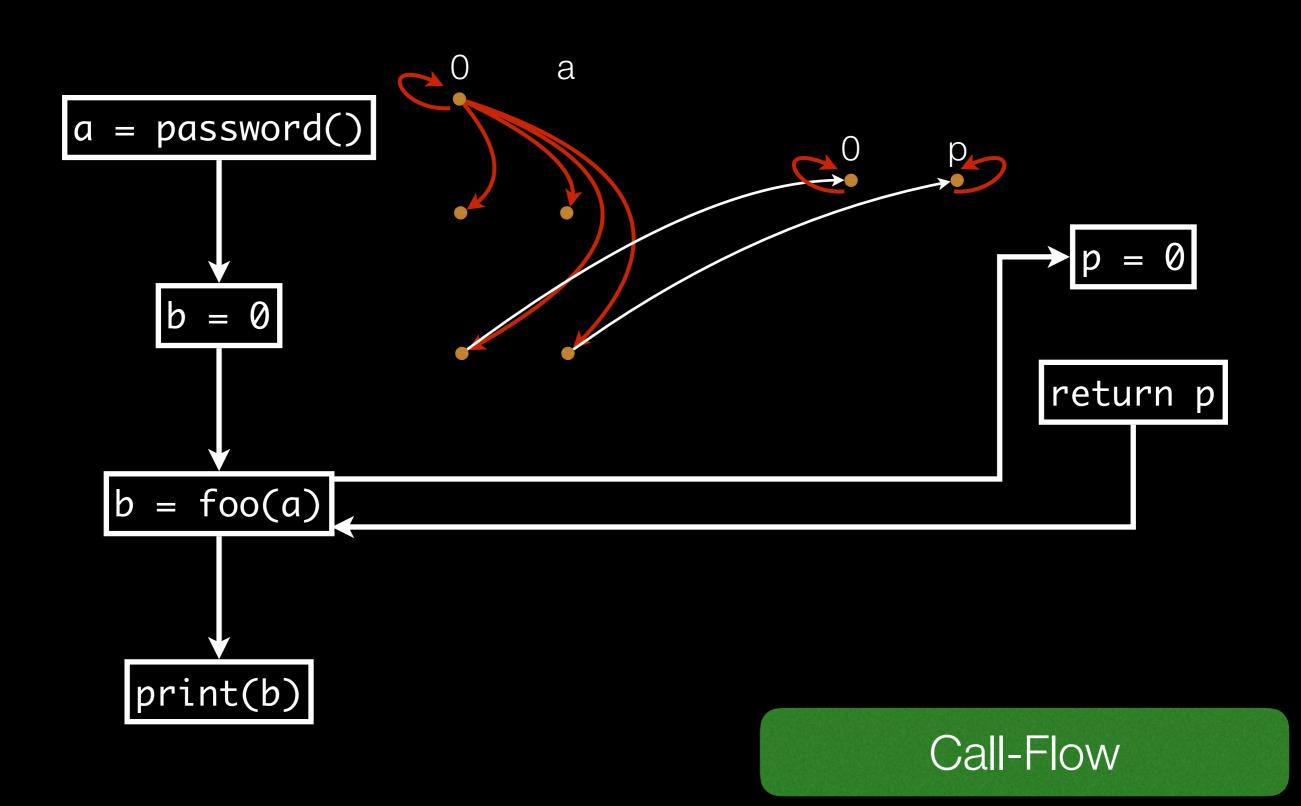


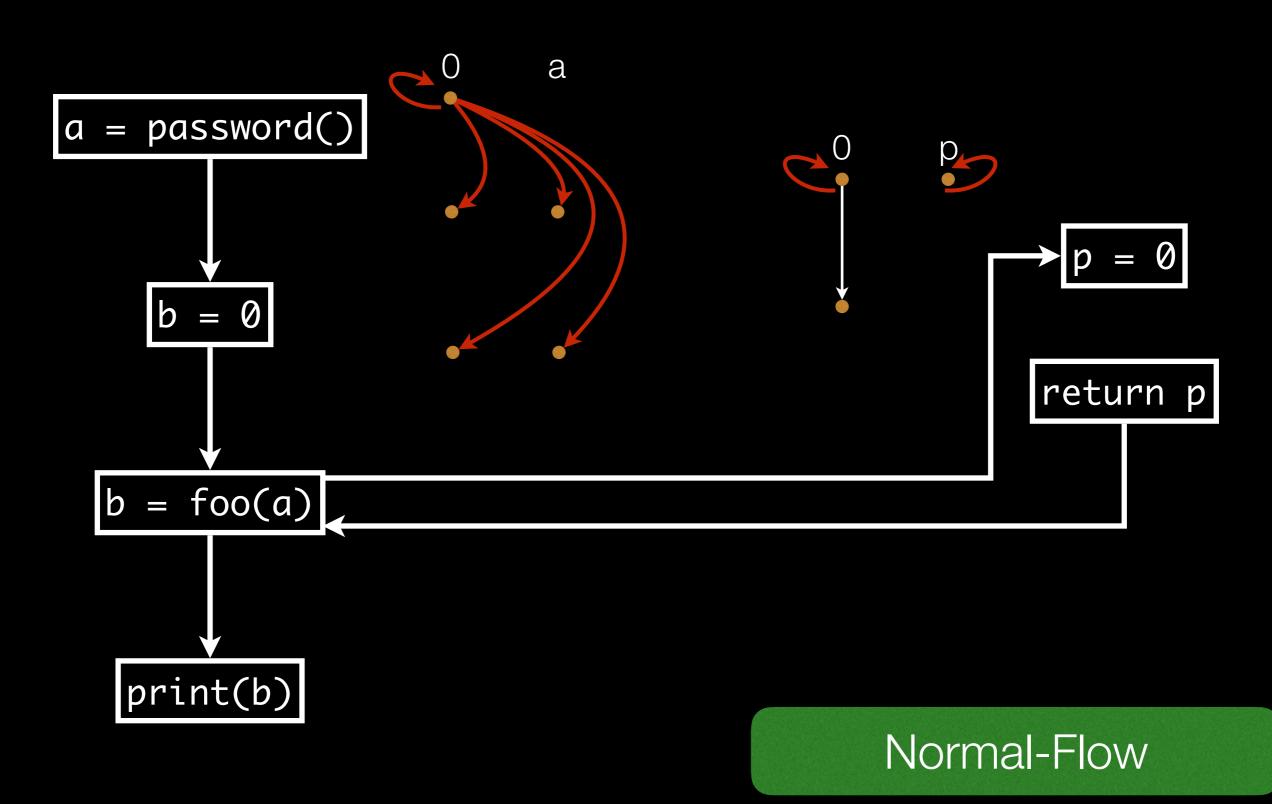


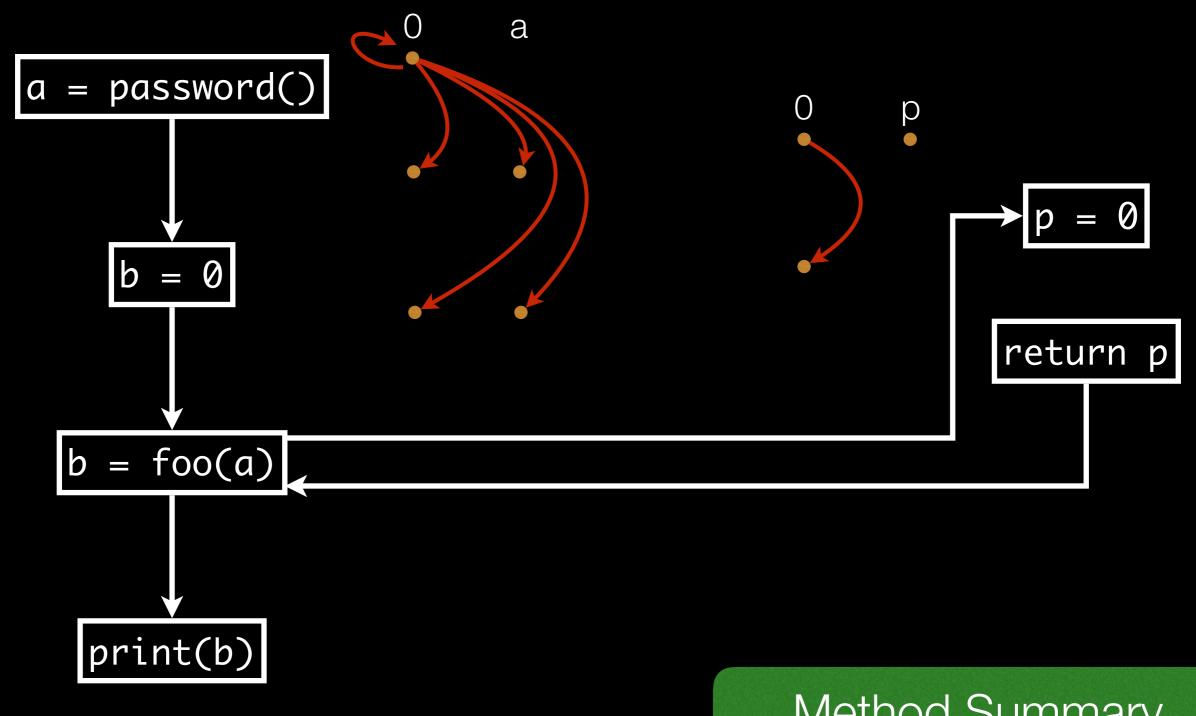
Path Edges



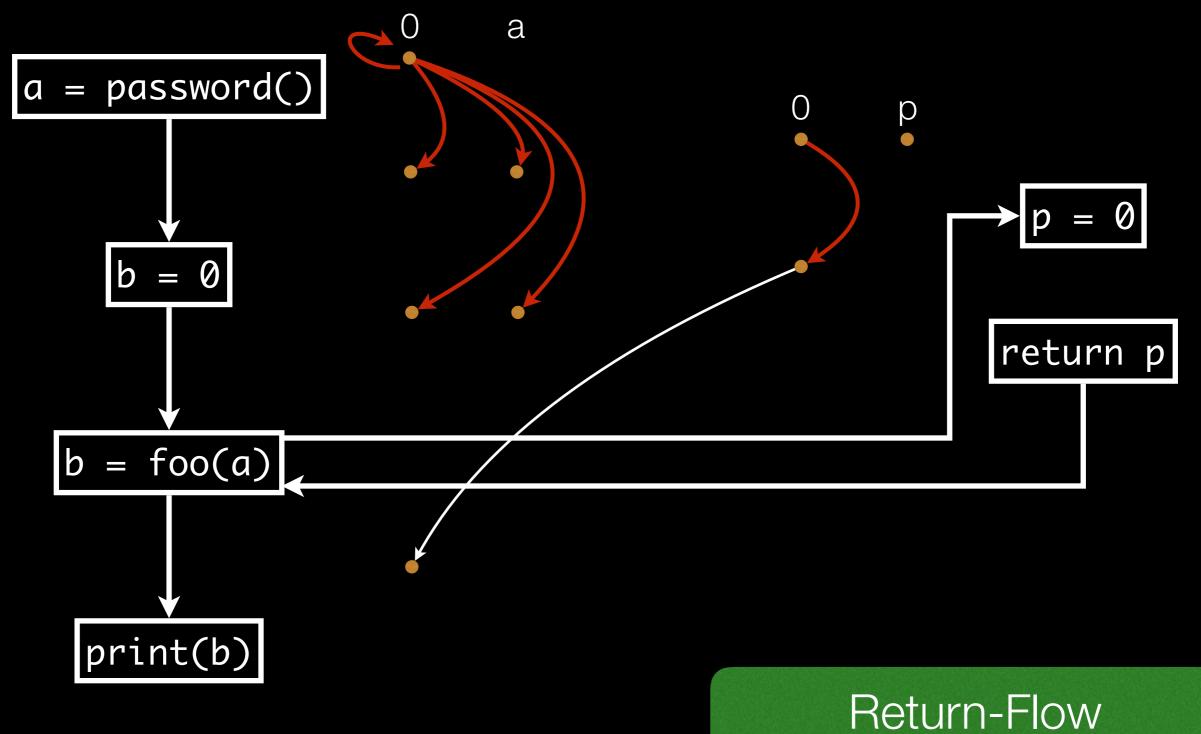


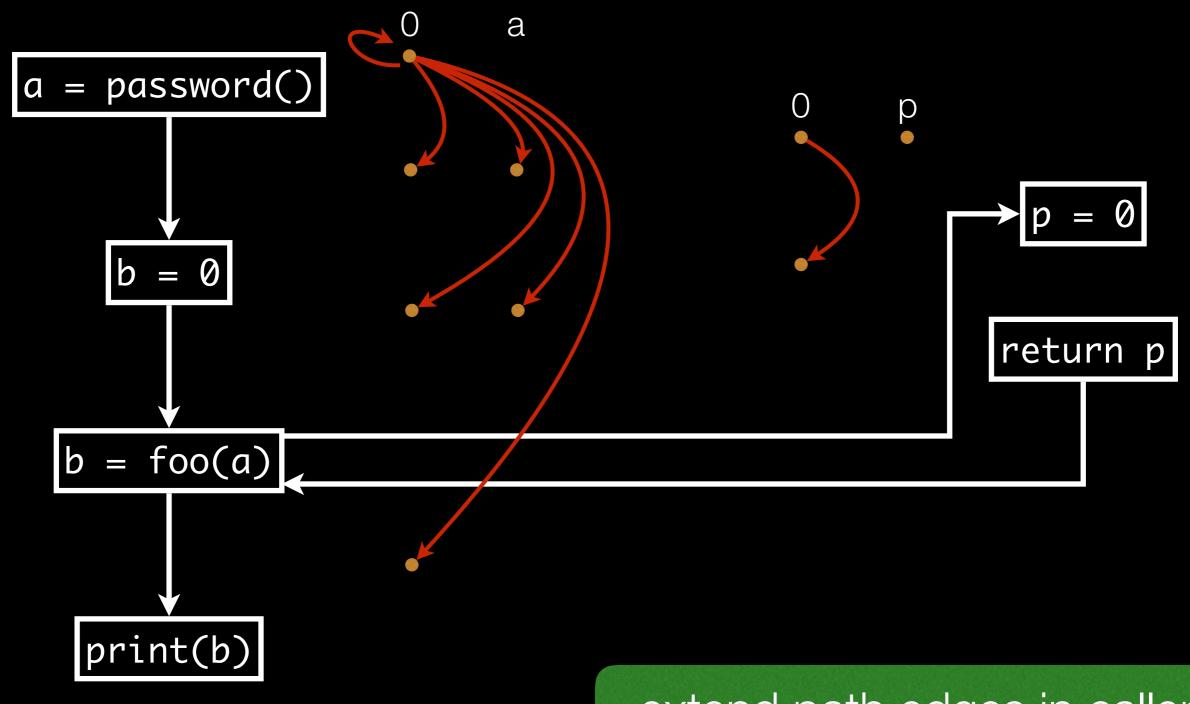




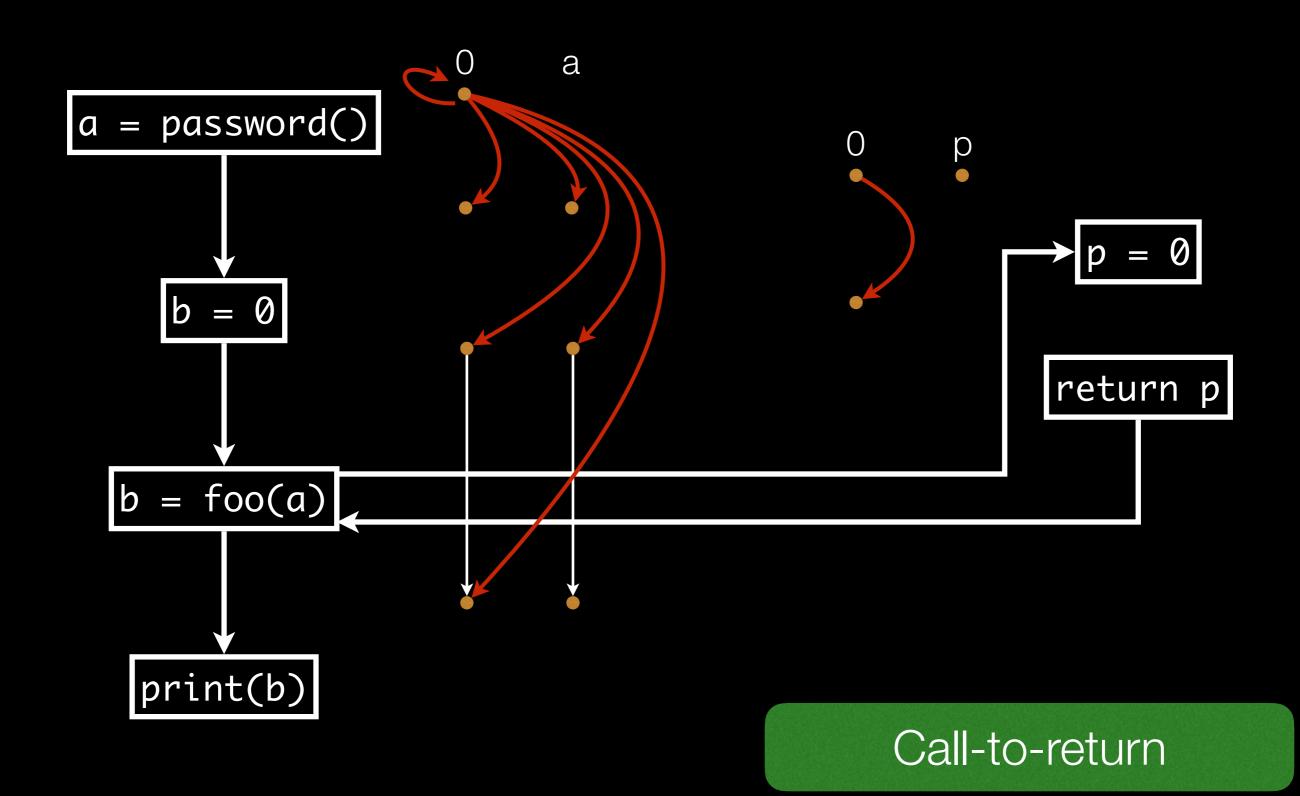


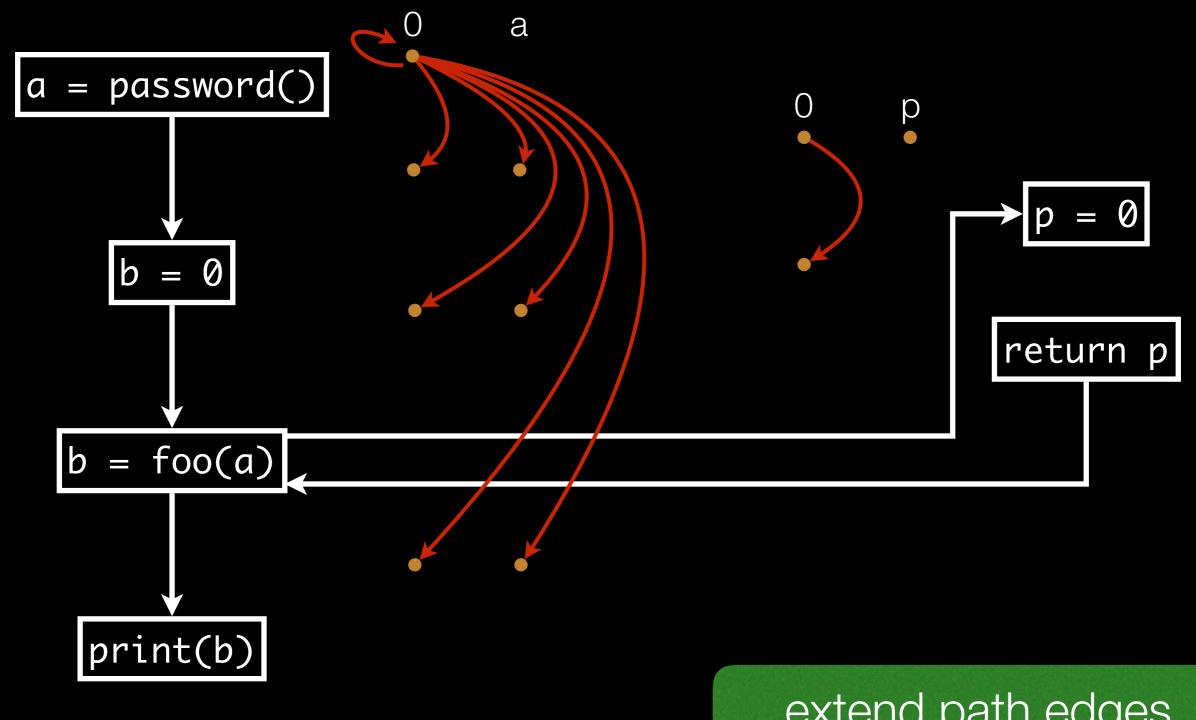
Method Summary



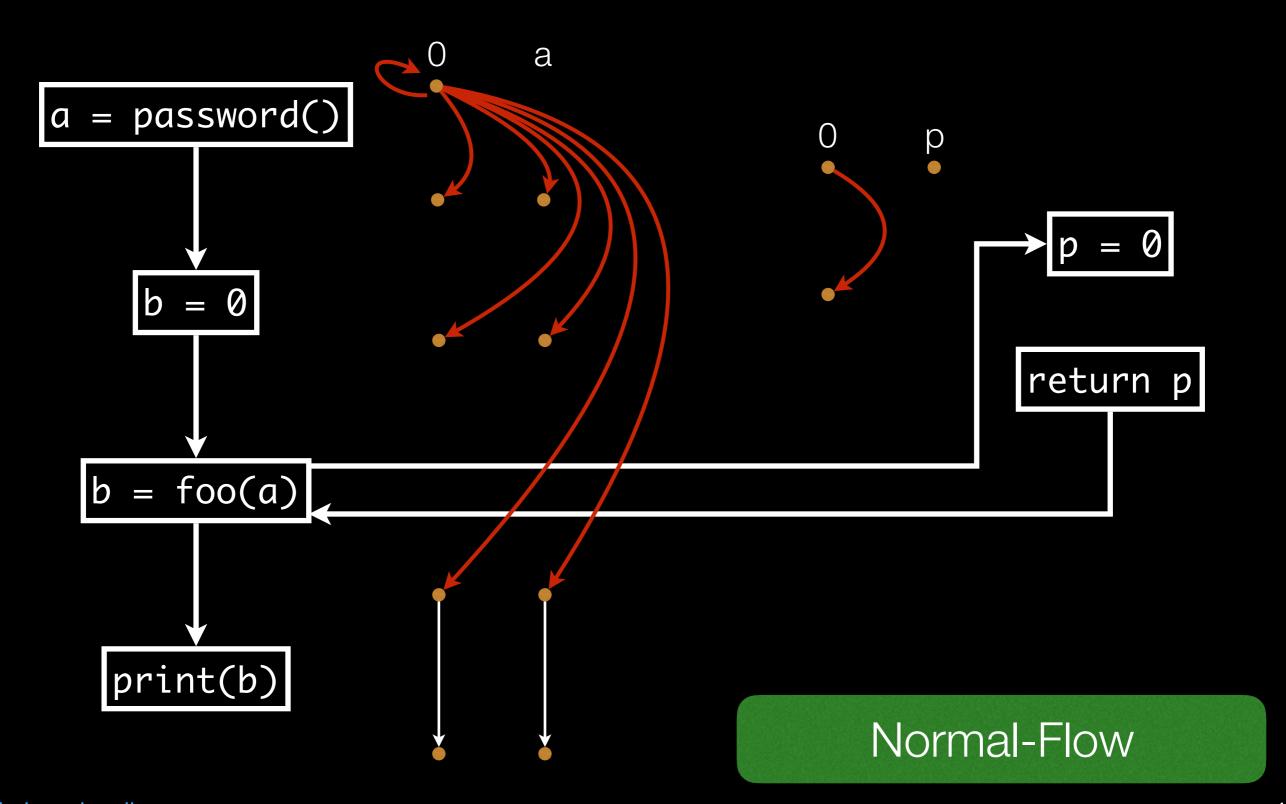


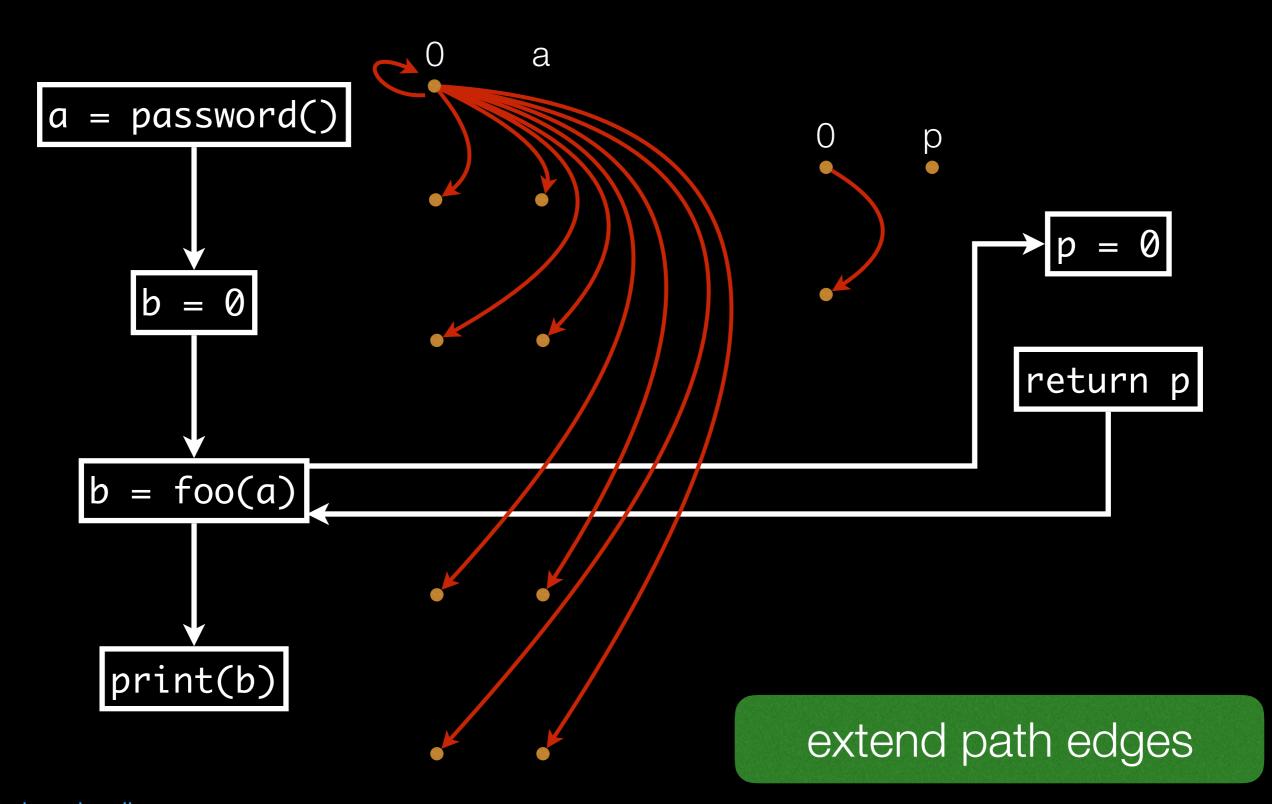
extend path edges in caller

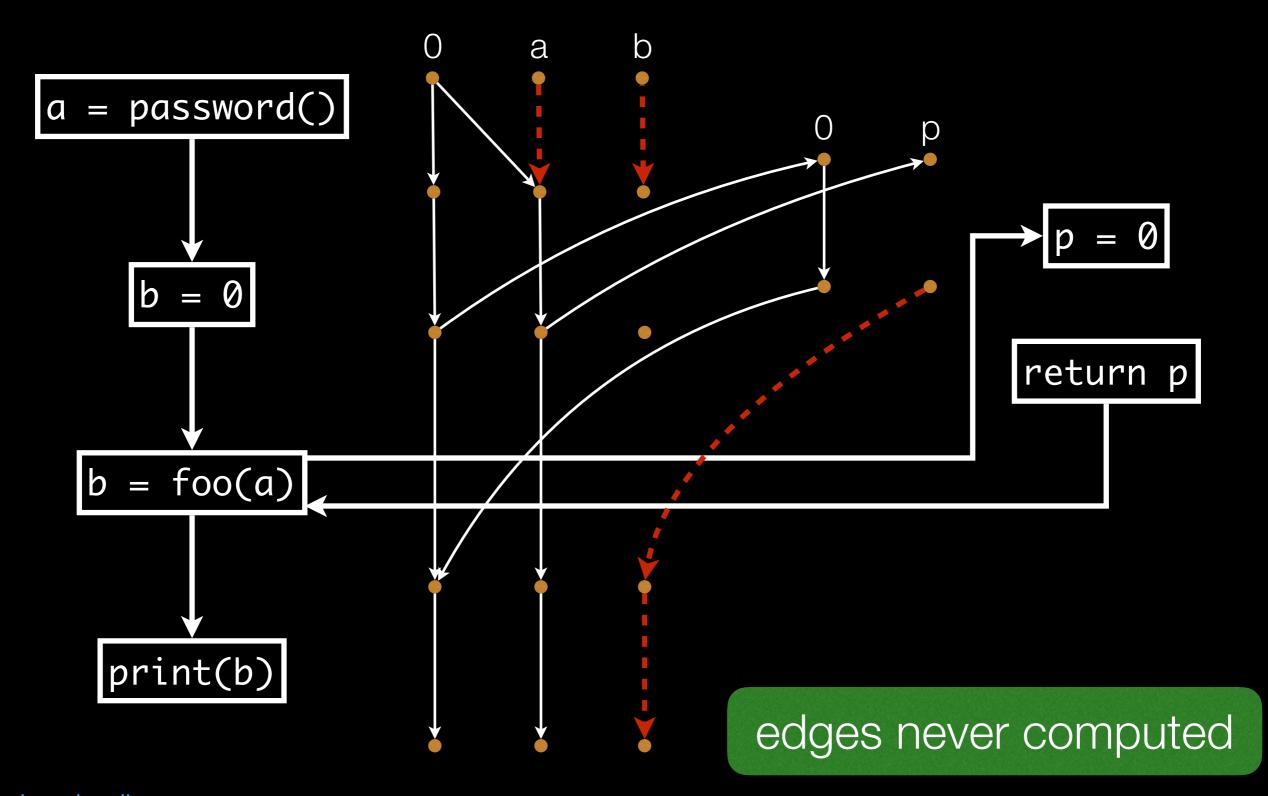




extend path edges







Tidbits!

Field-Sensitivity

a.f

a.*

A.f

Field-Sensitive

Field-Insensitive

Field-Based

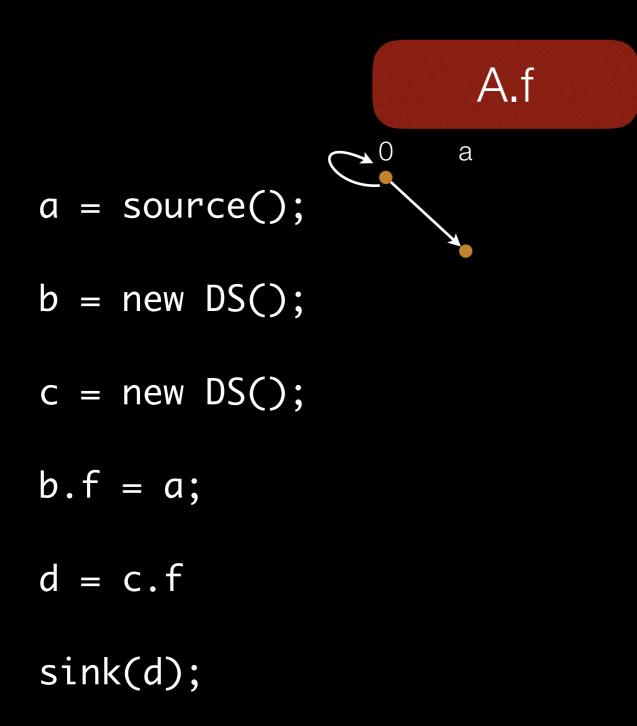
A.f

$$b = new DS();$$

$$c = new DS();$$

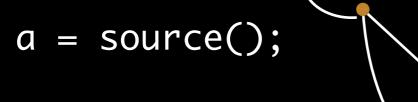
$$b.f = a;$$

$$d = c.f$$





a



$$b = new DS();$$

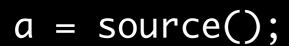
$$c = new DS();$$

$$b.f = a;$$

$$d = c.f$$



a



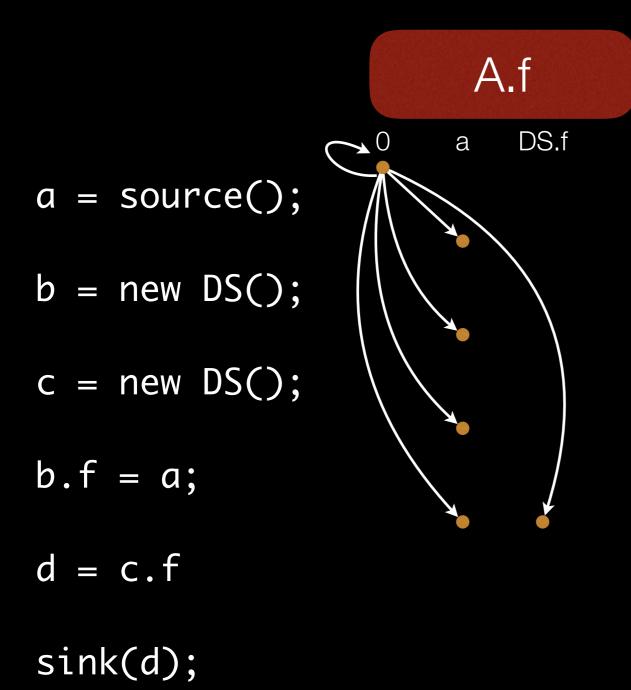
b = new DS();

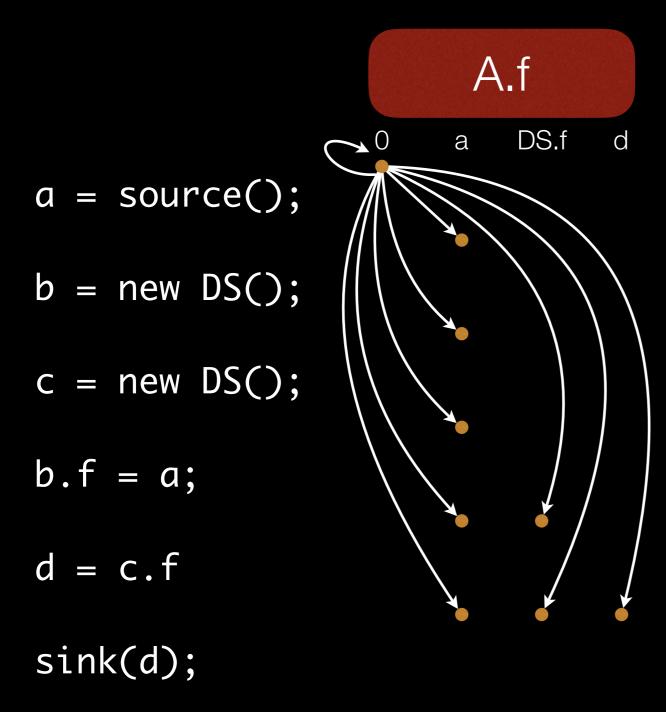
c = new DS();

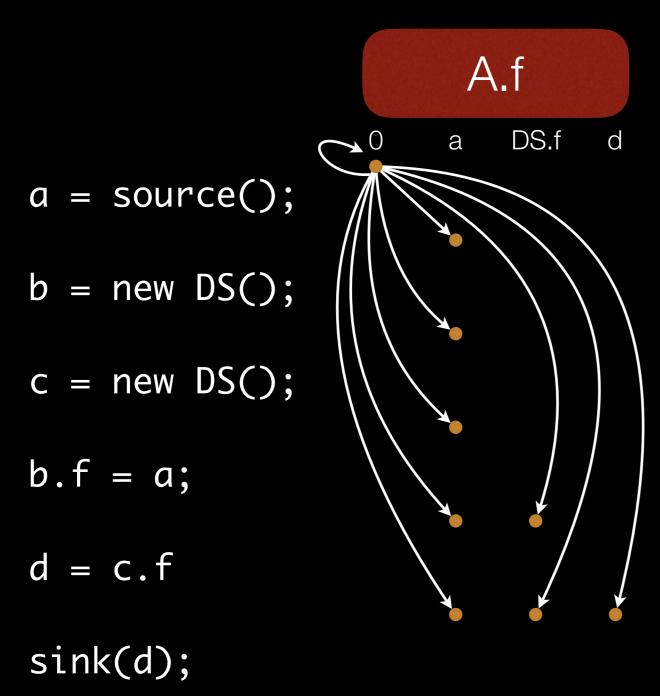
$$b.f = a;$$

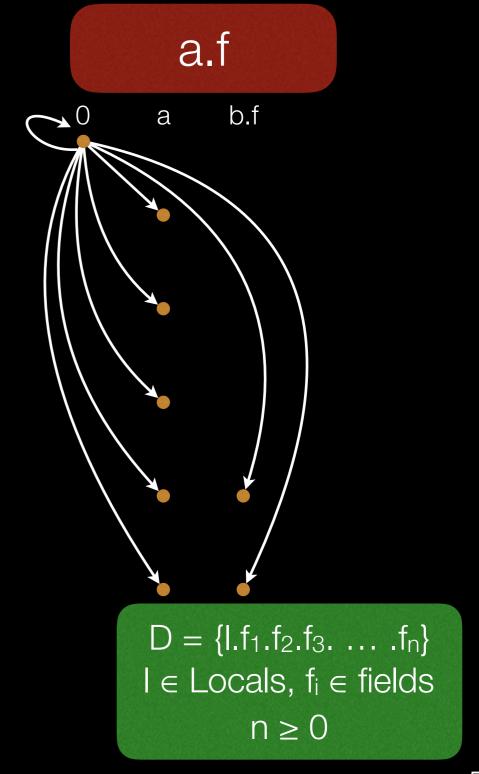
d = c.f

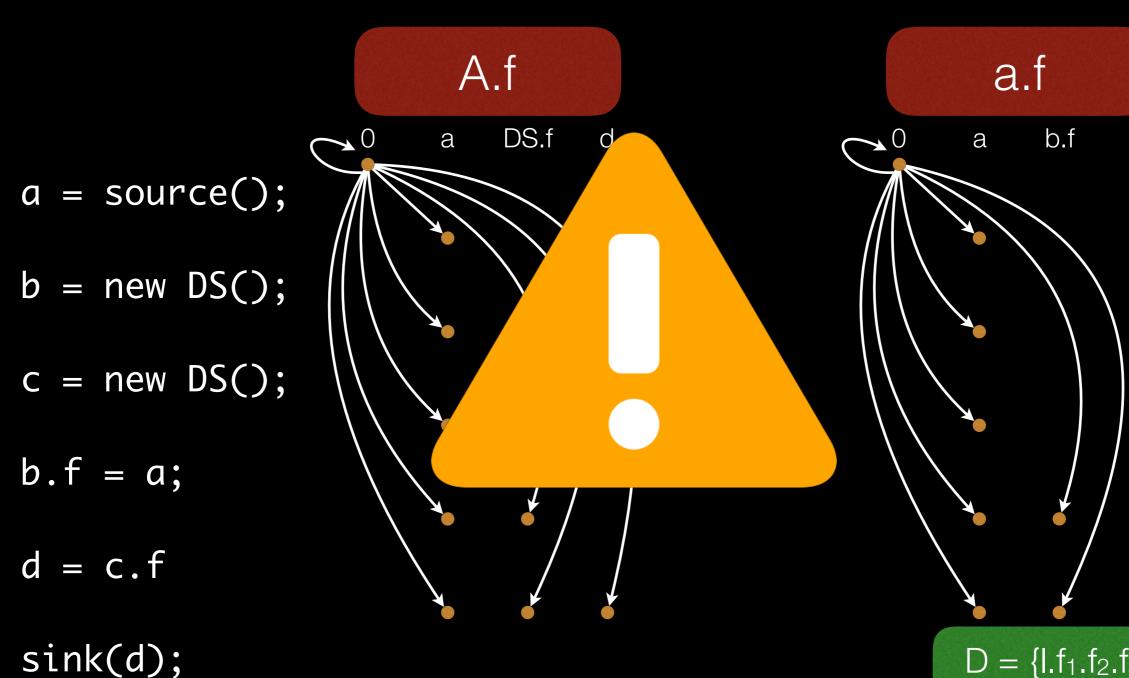
sink(d);







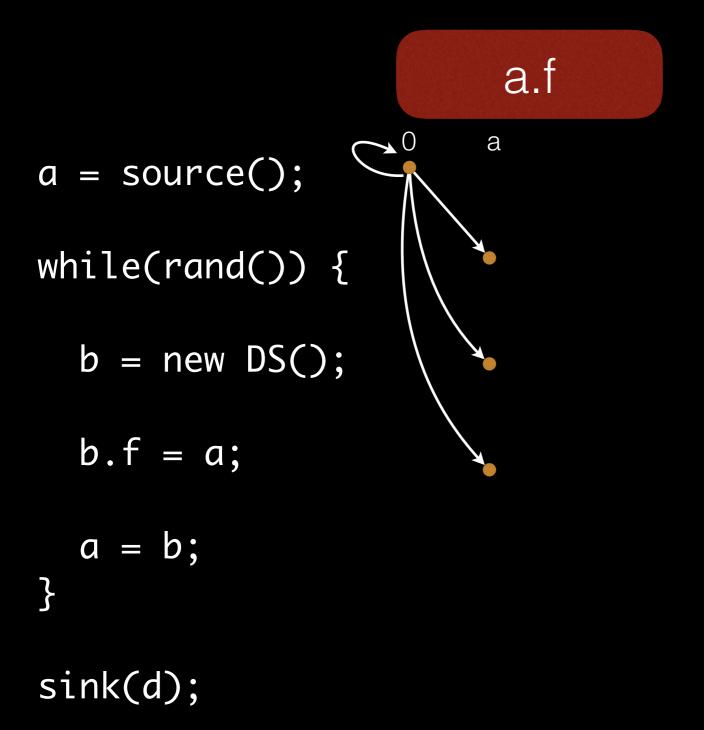


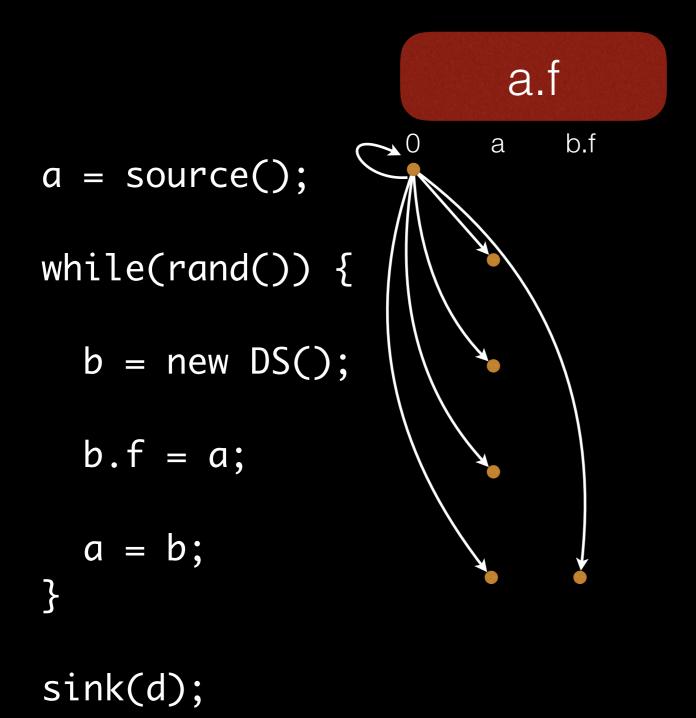


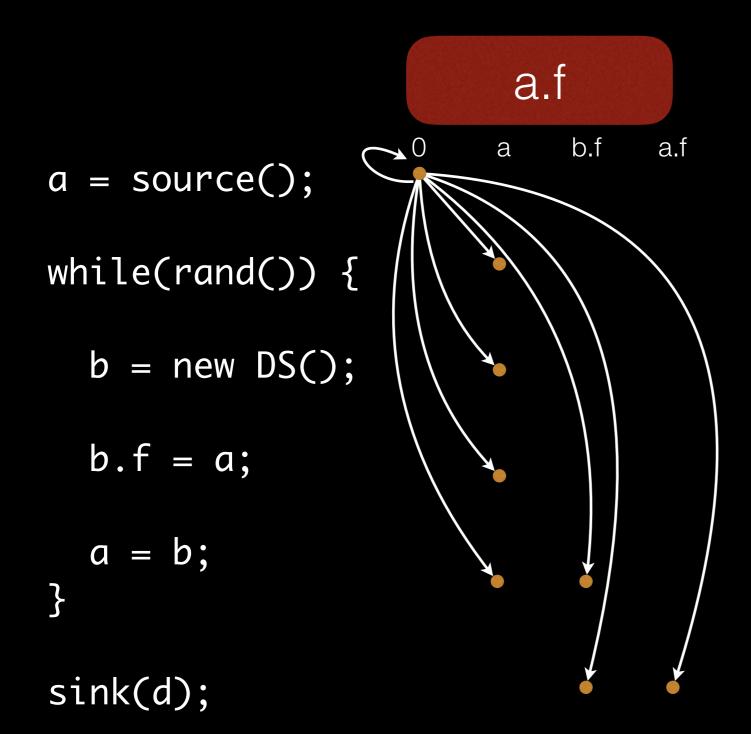
Infinite domain!!

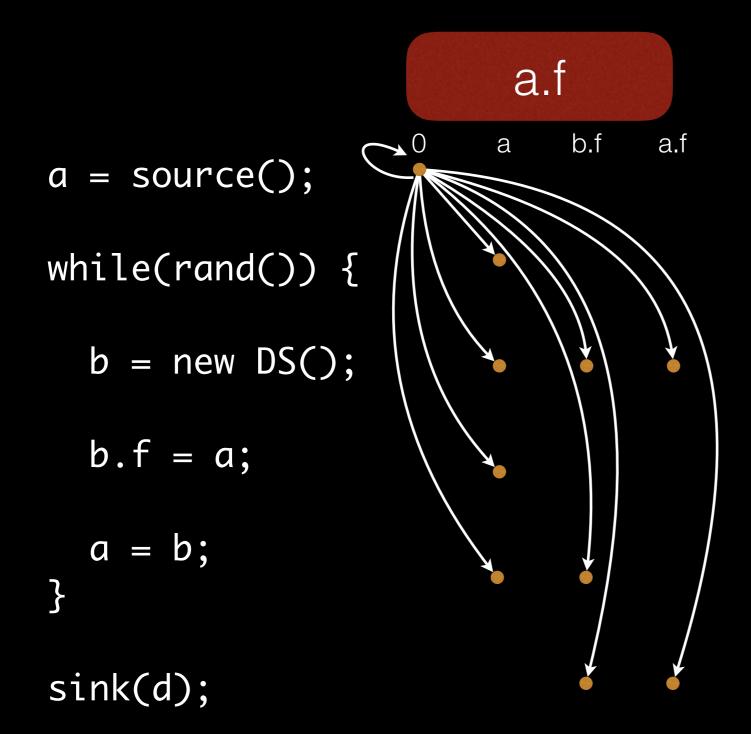
```
a.f
a = source();
while(rand()) {
  b = new DS();
  b.f = a;
  a = b;
sink(d);
```

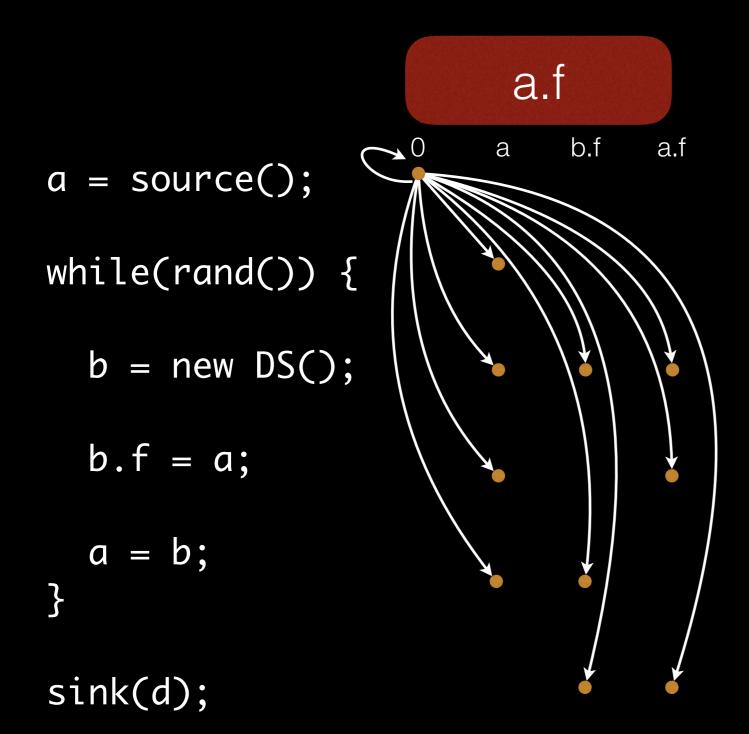
a.f a = source();while(rand()) { b = new DS();b.f = a;a = b;sink(d);

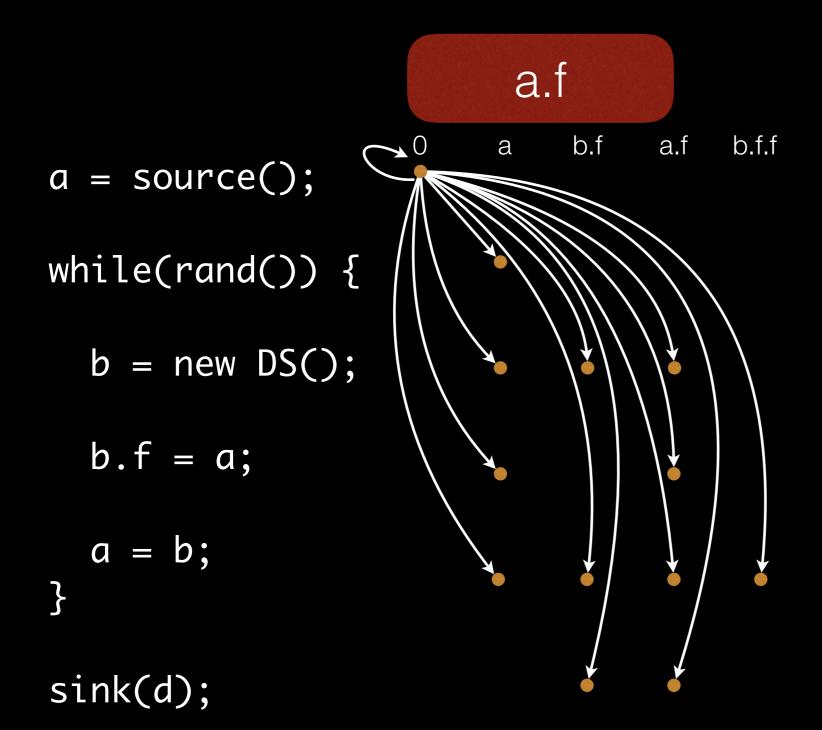


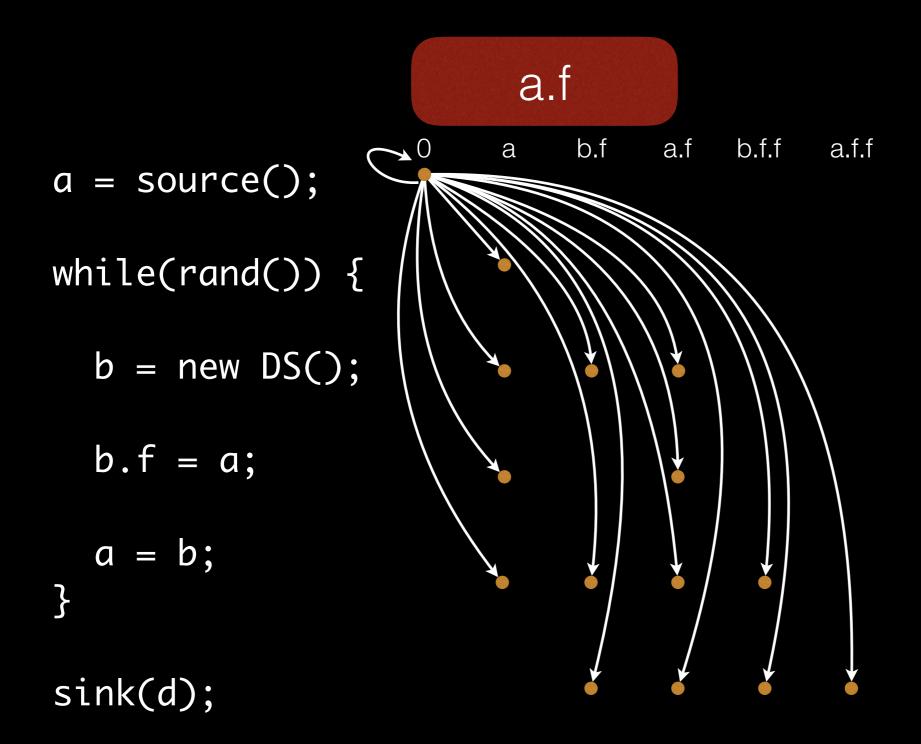


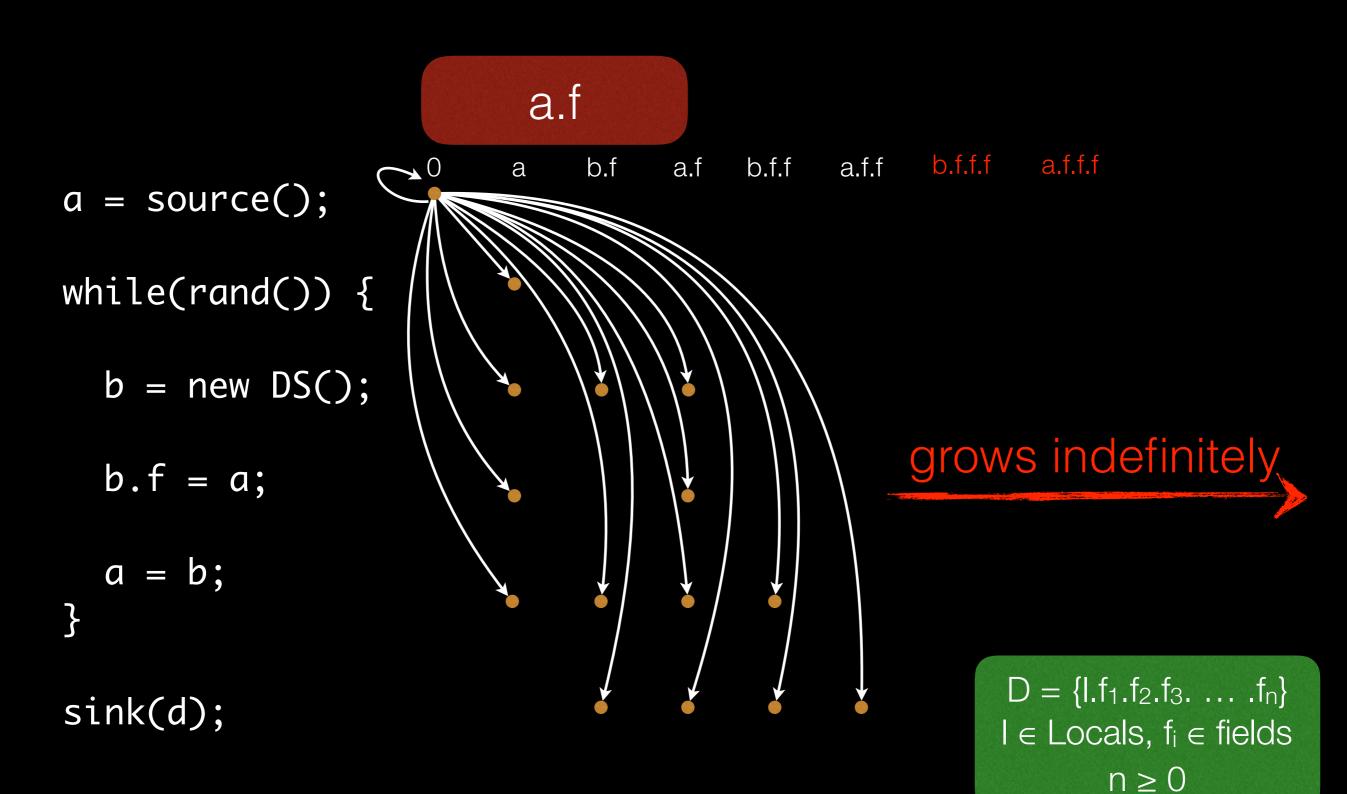


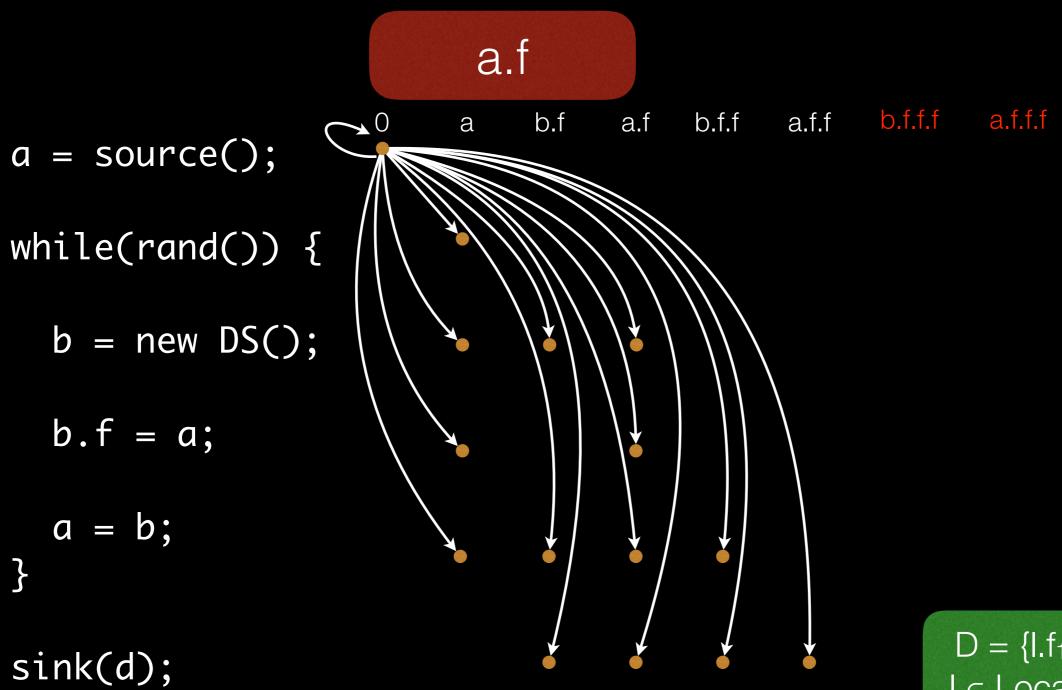


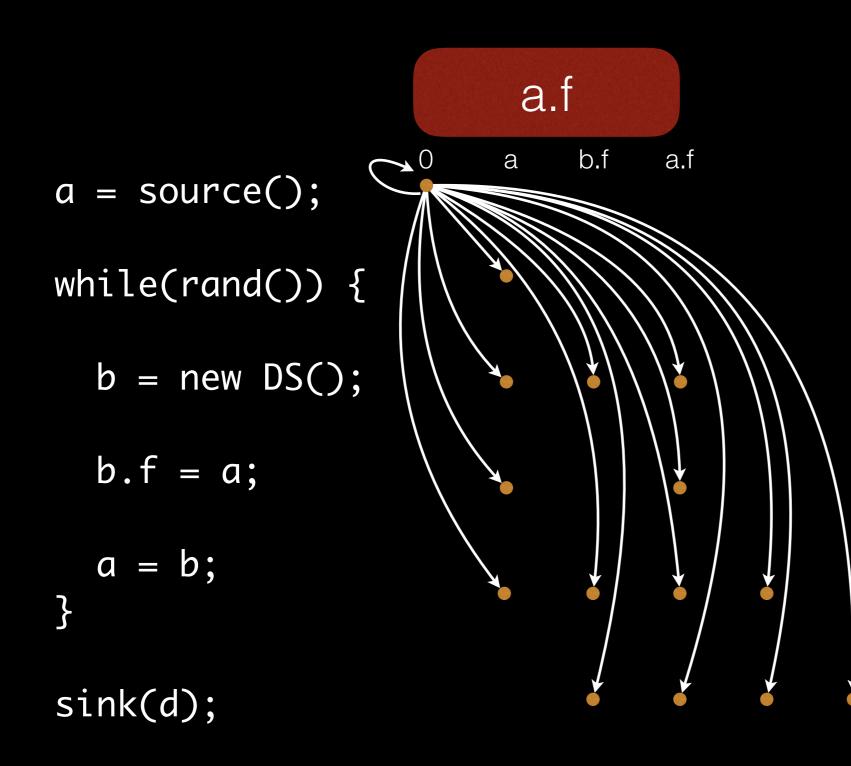








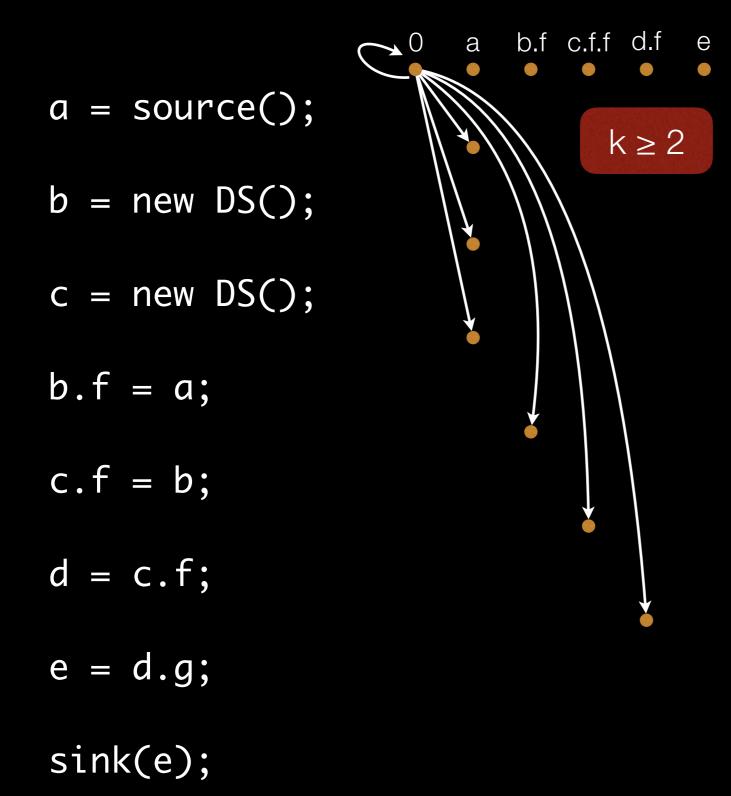


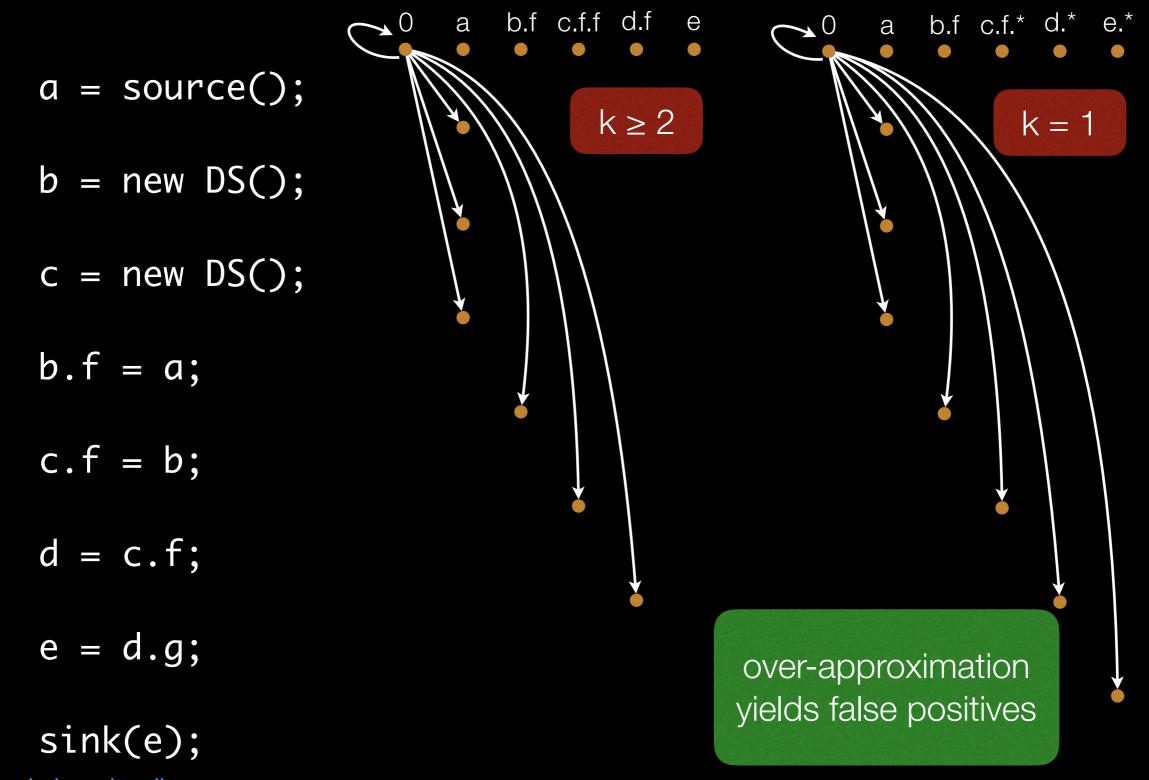


unsound domain beyond k => imprecise

k = 2

 $D = \{I.f_1.f_2.f_3.f_n\}$ $I \in Locals, f_i \in fields$ $0 \le n \le k$





Tidbits

Flow Sensitivity?

Context Sensitivity?

Tidbits

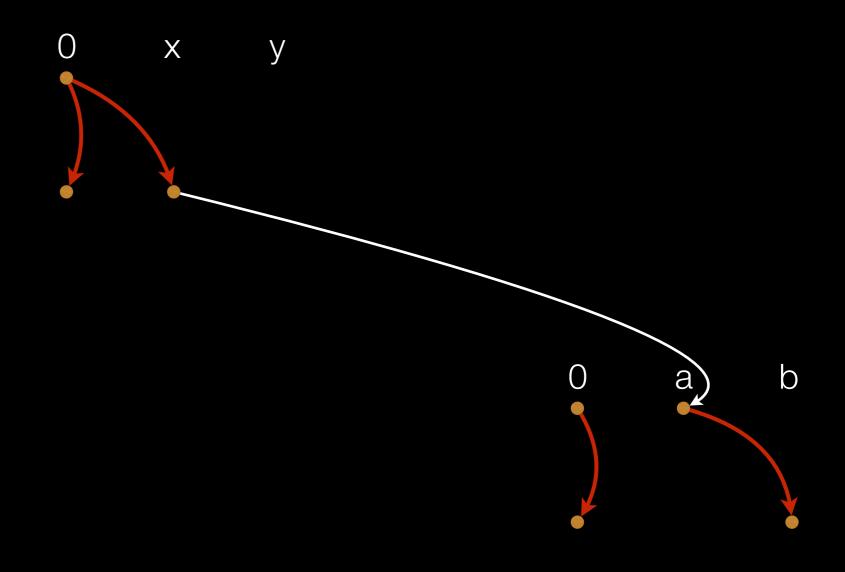
• Flow Sensitivity? Method Summaries

Context Sensitivity?

• •

call1();

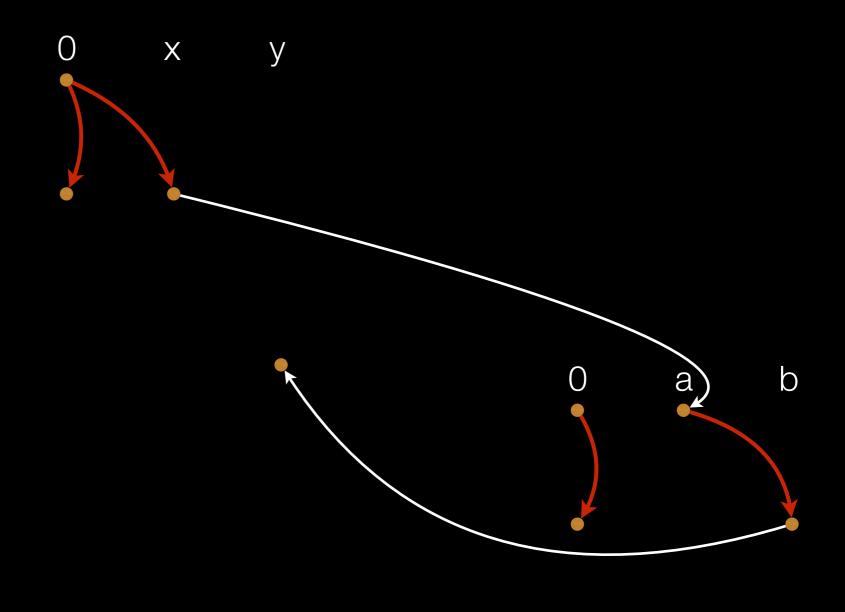
• •



• •

call1();

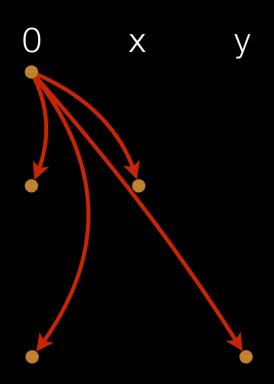
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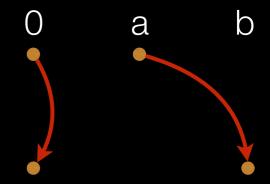


• • •

call1();

• •

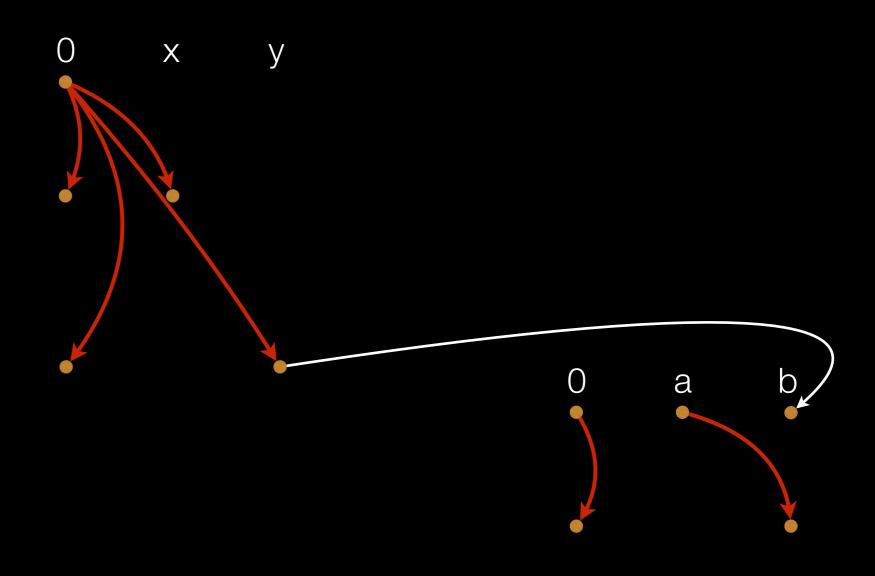




• • •

call1();

• •

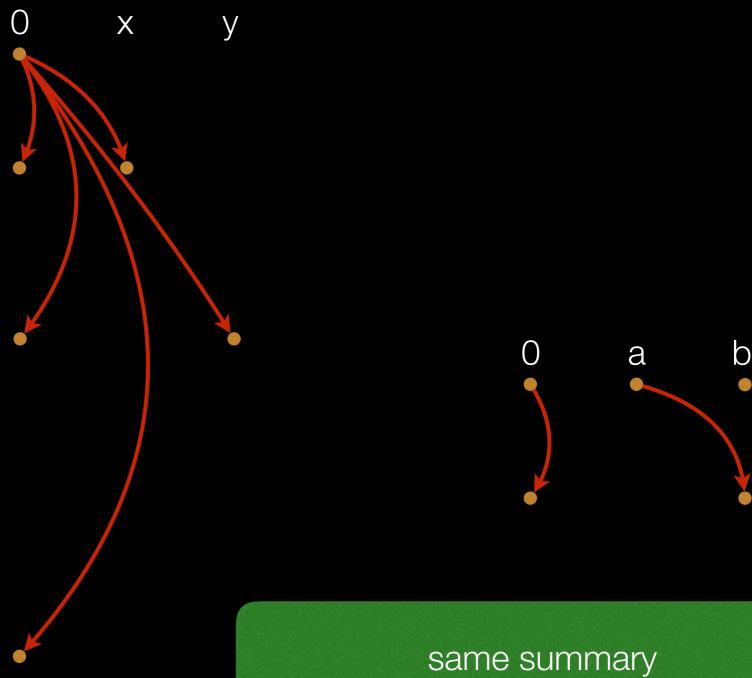


• • •

call1();

• •

call2();



different contexts => different results

Tidbits

- Flow Sensitivity
- Context Sensitivity
- Large Domains => performance problems

$$r = v + 1$$
 $0 \quad (v,0) \quad (r,1)$

In any context, if v is 0 before the call, then it is true that r is 1 after the call

```
main() {
x = inc(1);
y = inc(x);
z = inc(y);
 print(z);
      (v,0) (r,1)
```

```
inc(v) {
 r = v + 1;
 return r;
```

```
main() {
                          inc(v) {
x = inc(1);
                            r = v + 1;
y = inc(x);
                            return r;
 z = inc(y);
 print(z);
       (v,0) (r,1) (v,1) (r,2) (v,2) (r,3)
```

```
main() {
                              inc(v) {
 x = inc(1);
                                 = \vee + 1.
 V = inc(V)
           Solution: IDE instead of IFDS
          (v,0) (r,1) (v,1) (r,2) (v,2)
```

Recap

- Interprocedural Finite Distributive Subset
- Flow Functions
- Taint Analysis Example
- On-the-fly ESG
- Tidbits: field-sensitivity, contextsensitivity, large domains

Next

Context sensitivity