

IFDS Framework

CMPUT 620 — Static Program Analysis

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September 26, 2017 GSB 8-59

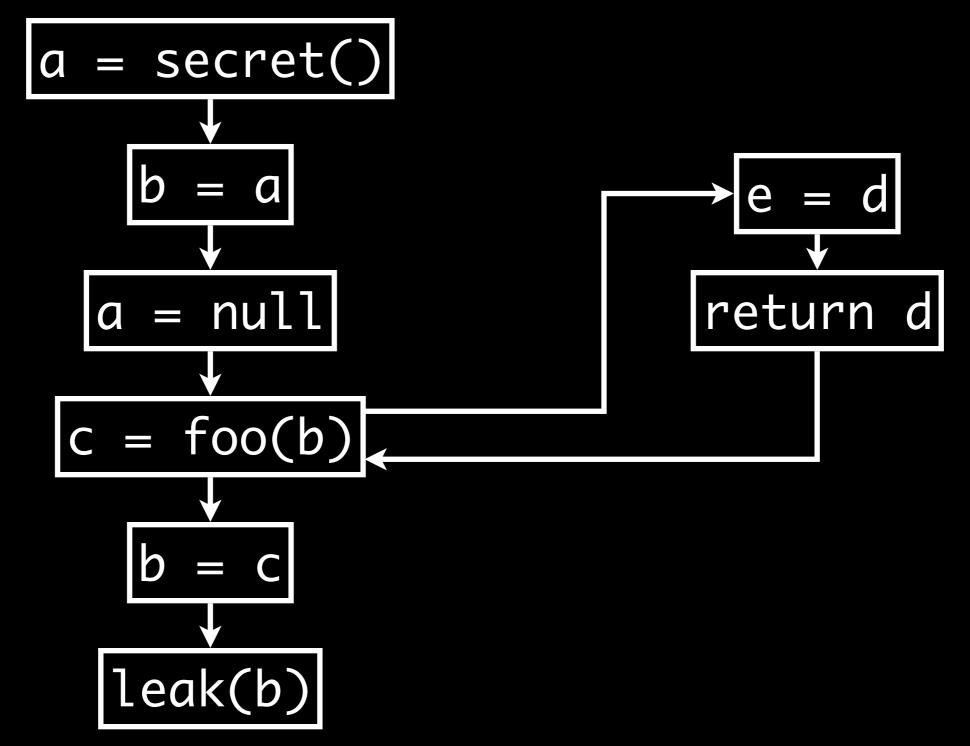
Interprocedural Finite Distributive Subset Problems

Thomas Reps, Susan Horwitz, and Mooly Sagiv. Precise Interprocedural Dataflow Analysis via Graph Reachability. ACM SIGPLAN-SIGACT Symposium on Principles of Programming Languages (POPL '95), pages 49–61.

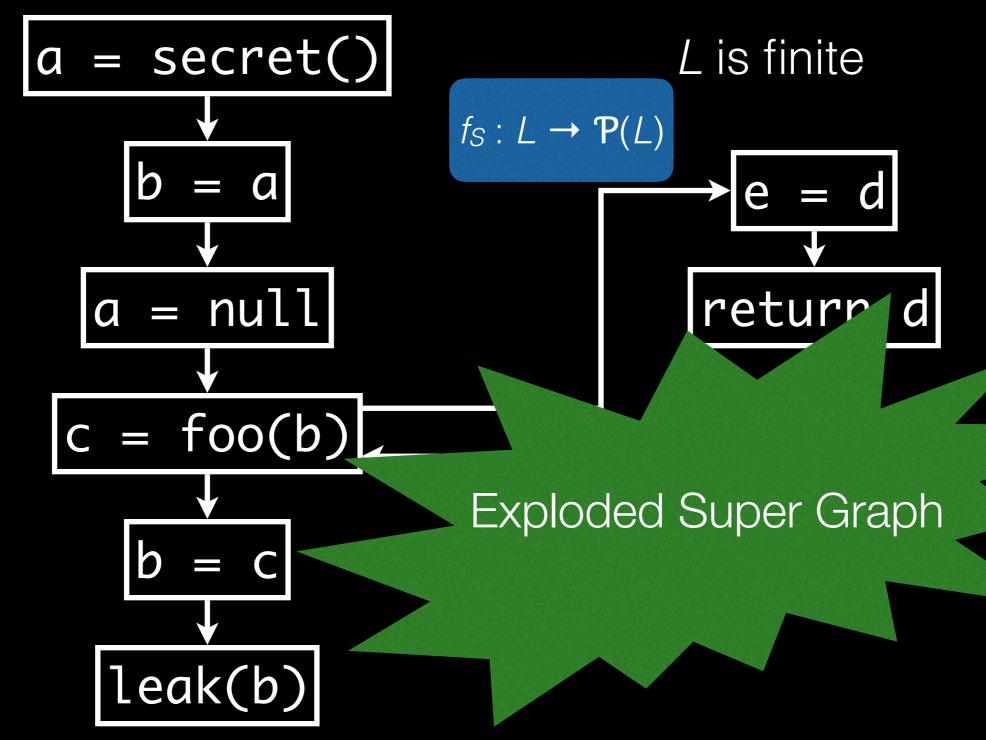
Interprocedural Finite Distributive Subset Problems

Interprocedural

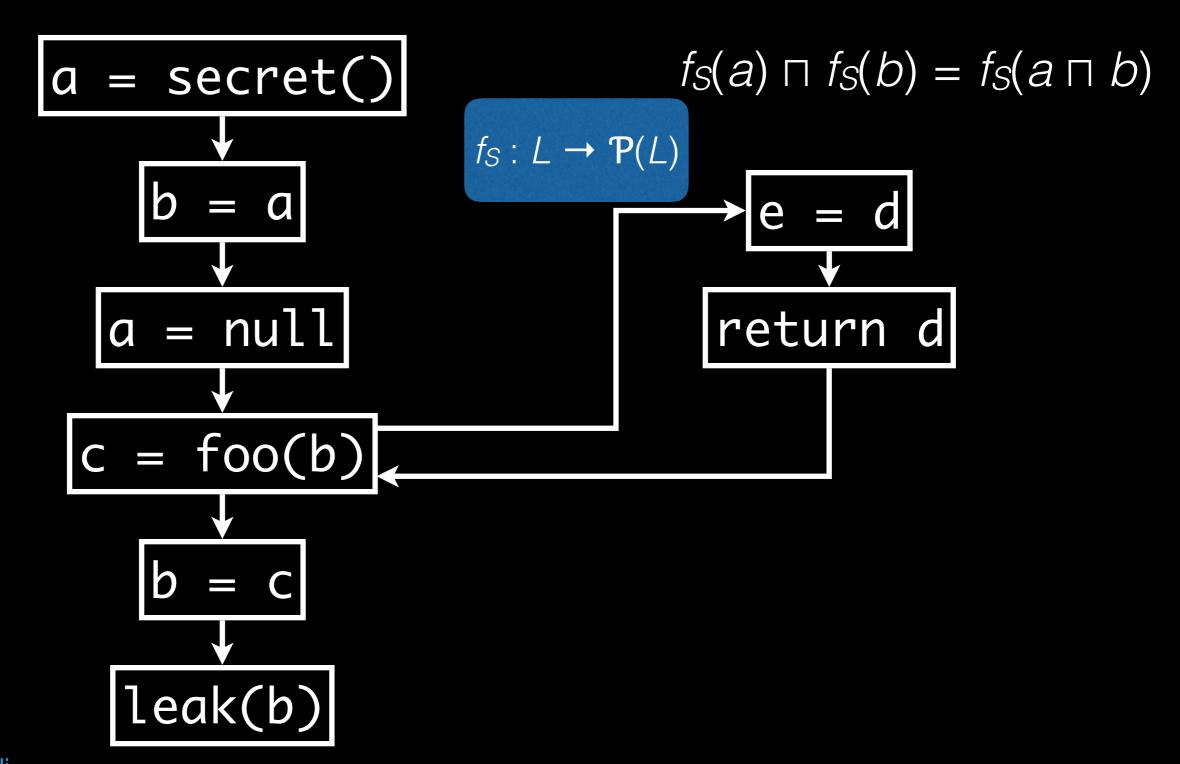
Interprocedural



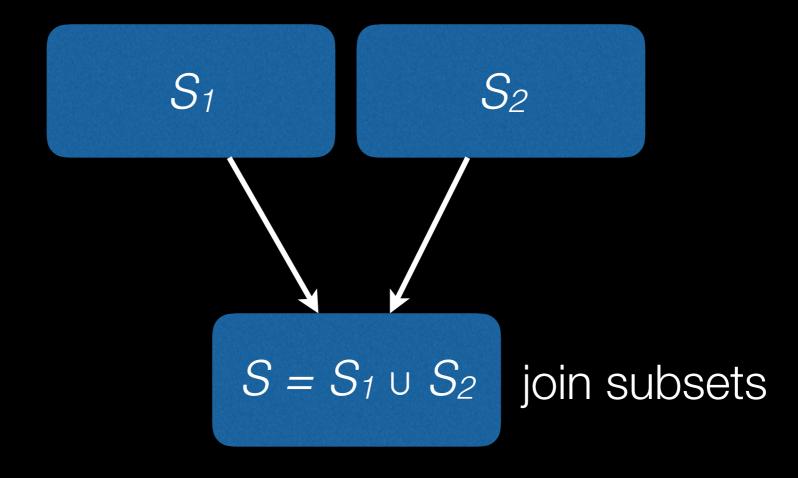
Finite



Distributive



Subset



alias taint typestate

IFDS

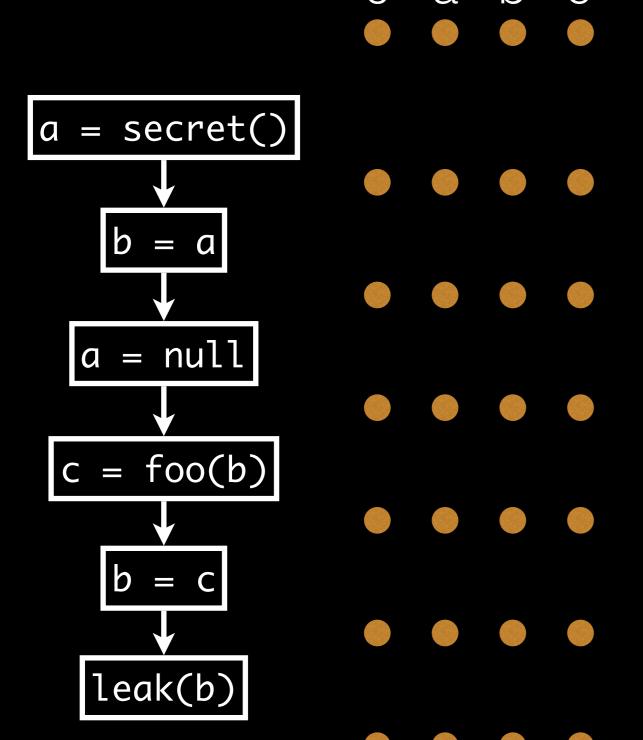
constant propagation

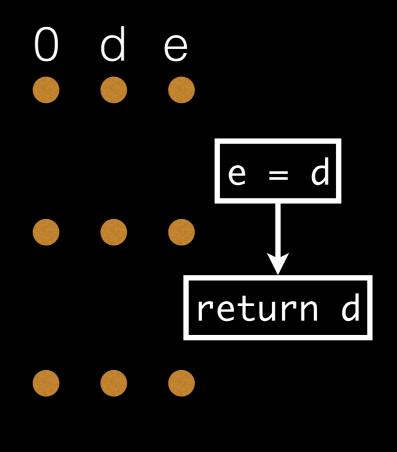
reaching definitions

uninitialized variables

Graph Reachability

Graph Reachability

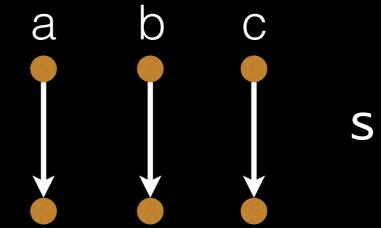




- each node is a fact
- fact holds at a statement => node is reachable in ESG

Examples of Flow Functions

Identity Flow Functions

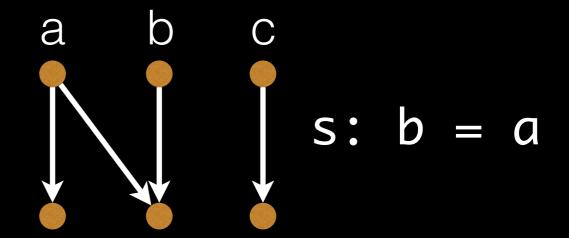


out(s) = facts after s

every fact is reachable iff it was previously reachable

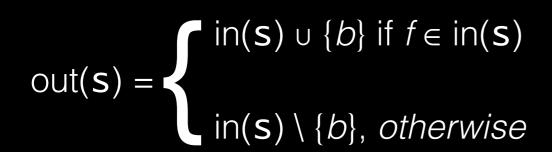
Flow Functions

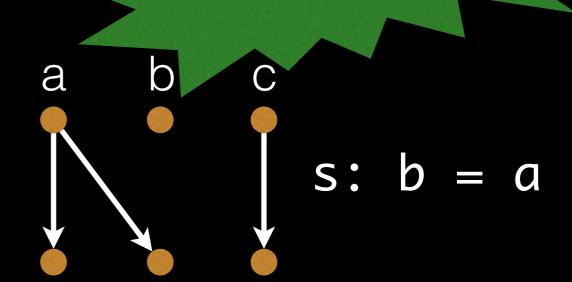
out(s) =
$$\begin{cases} in(s) \cup \{b\} & \text{if } f \in in(s) \\ in(s), & \text{otherwise} \end{cases}$$



every fact is reachable iff it was previously reachable, and b is also reachable if a was reachable

Flow Functions



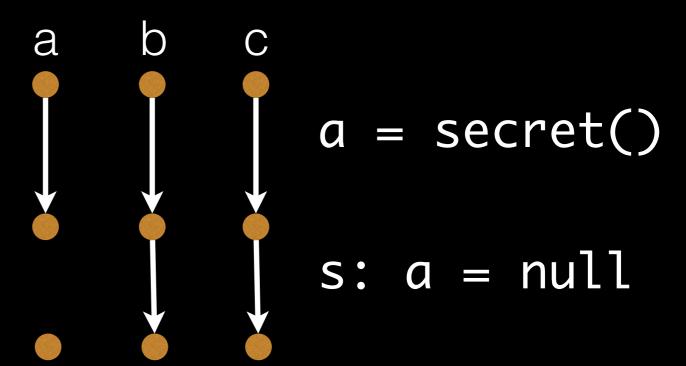


Taint Analysis

every fact except *b* is reachable iff it was previously reachable; *b* is reachable if *a* was reachable

"Killing" Facts

$$\operatorname{out}(\mathbf{s}) = \operatorname{in}(\mathbf{s}) \setminus \{a\}$$



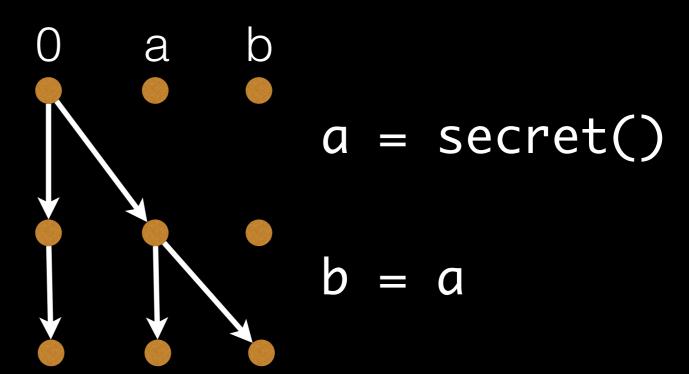
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a is not reachable, even if it was before

Generating Facts

$$\mathsf{out}(\mathsf{s}) = \mathsf{in}(\mathsf{s}) \cup \{a\}$$

$$\mathsf{out}(\mathsf{s}) = \mathsf{in}(\mathsf{s}) \cup \{b\}$$



0 is the tautological fact => always reachable

Non-Distributive Flow Functions



e.g., full constant propagation c = a + b

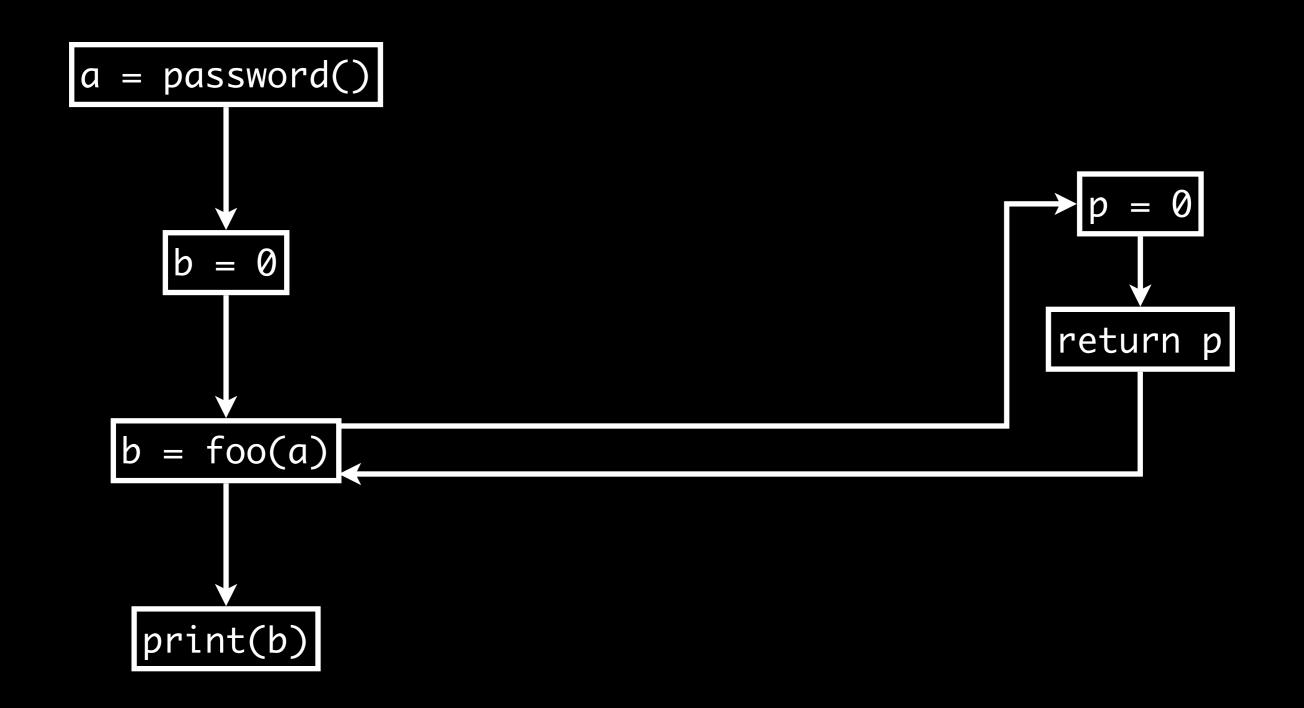
```
void main() {
    int a = password();
    int b = 0;

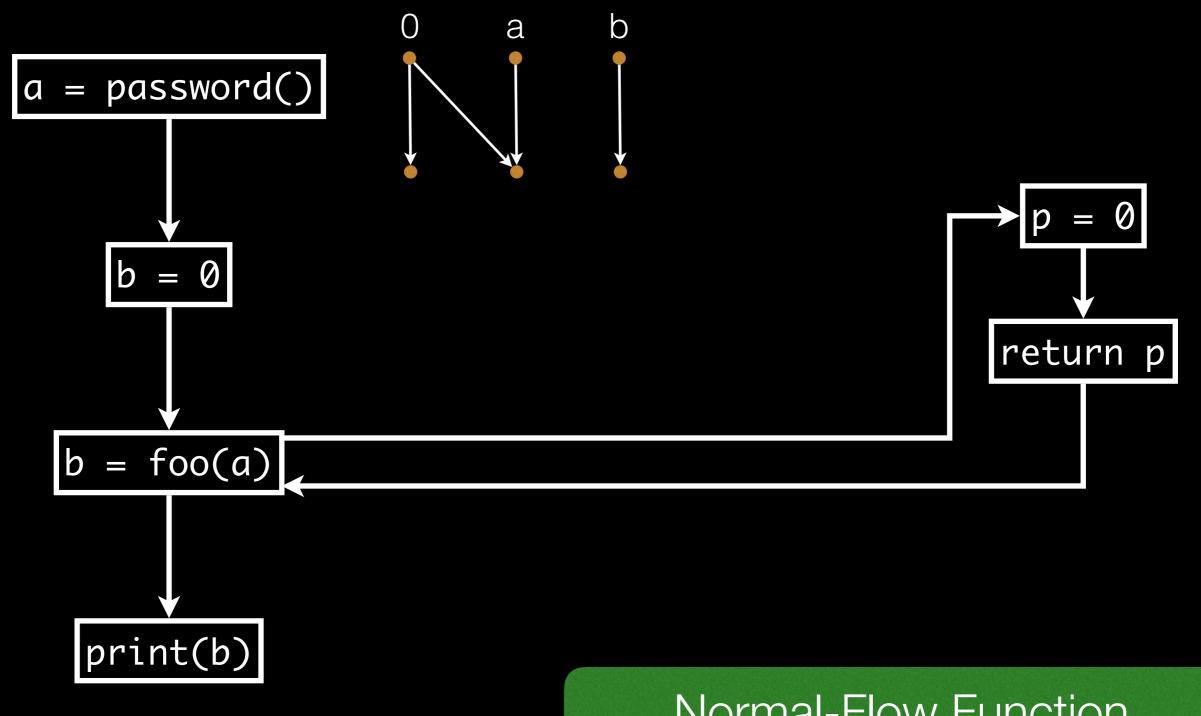
b = foo(a);

print(b);
}
```

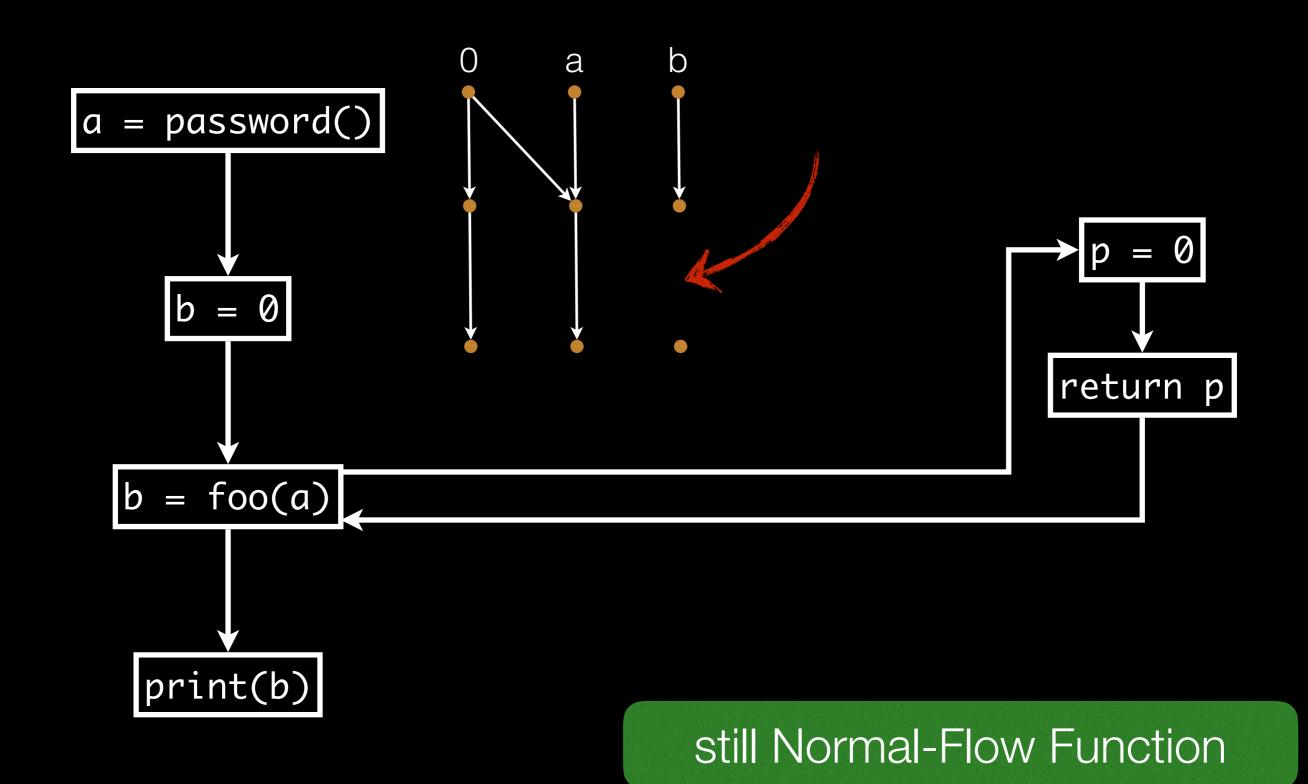
```
int foo(int p) {
   p = 0;

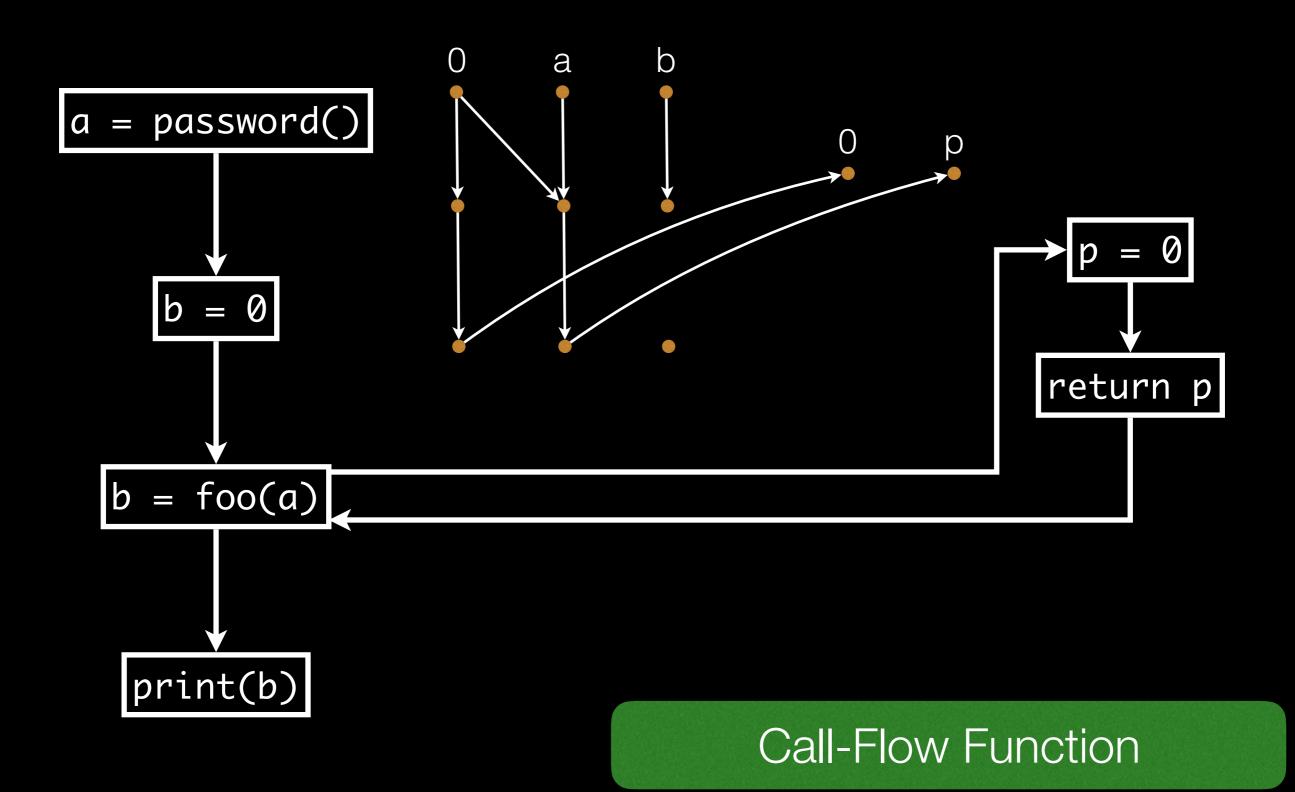
   return p;
}
```

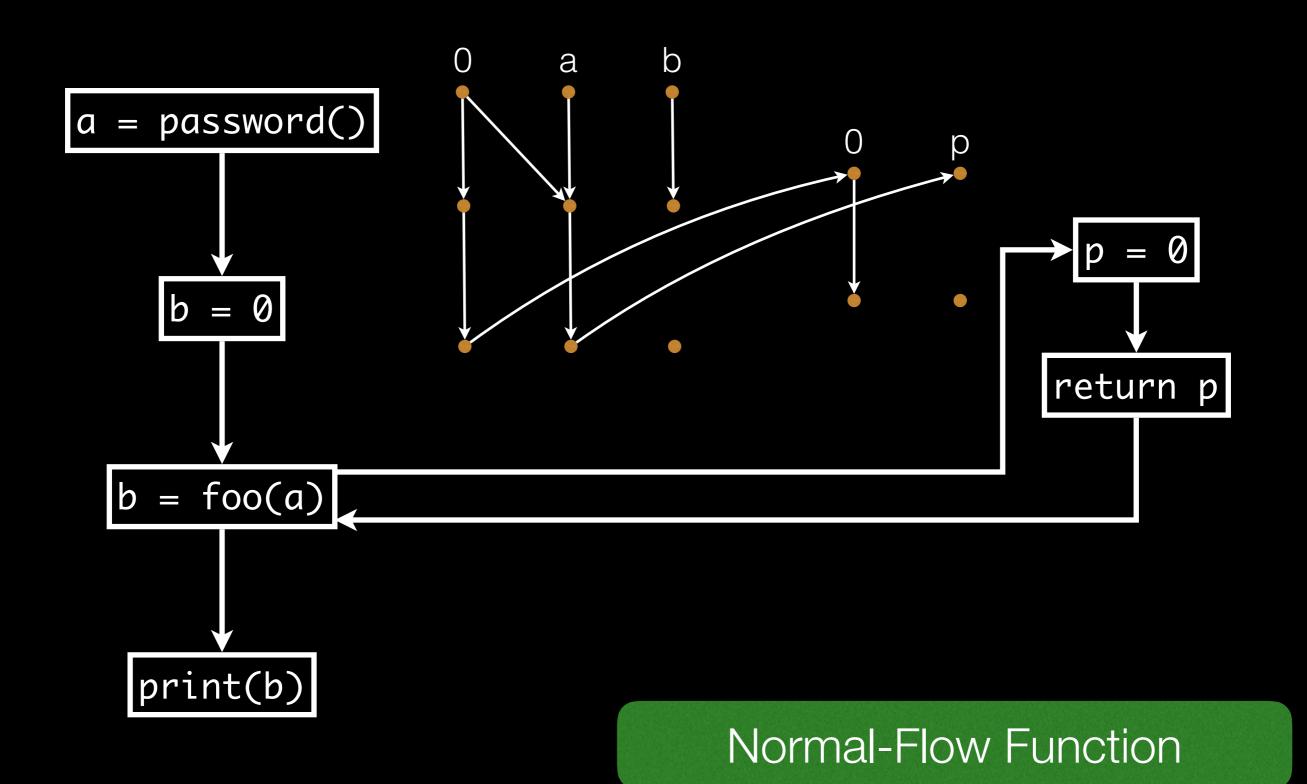


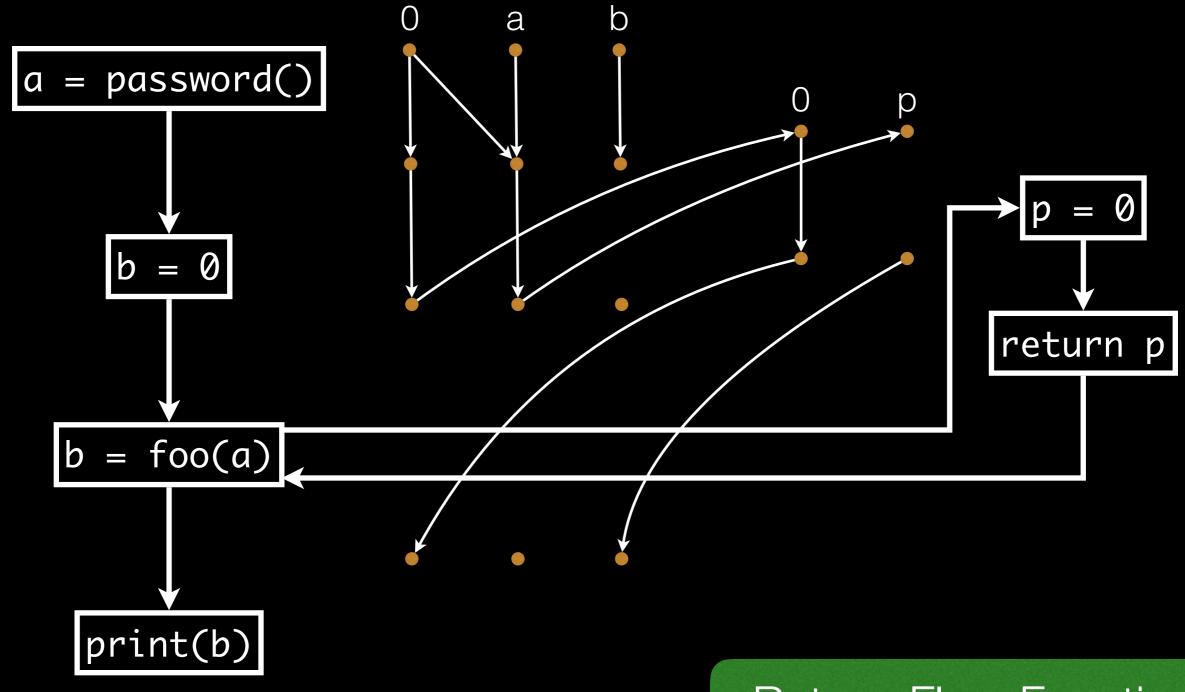


Normal-Flow Function

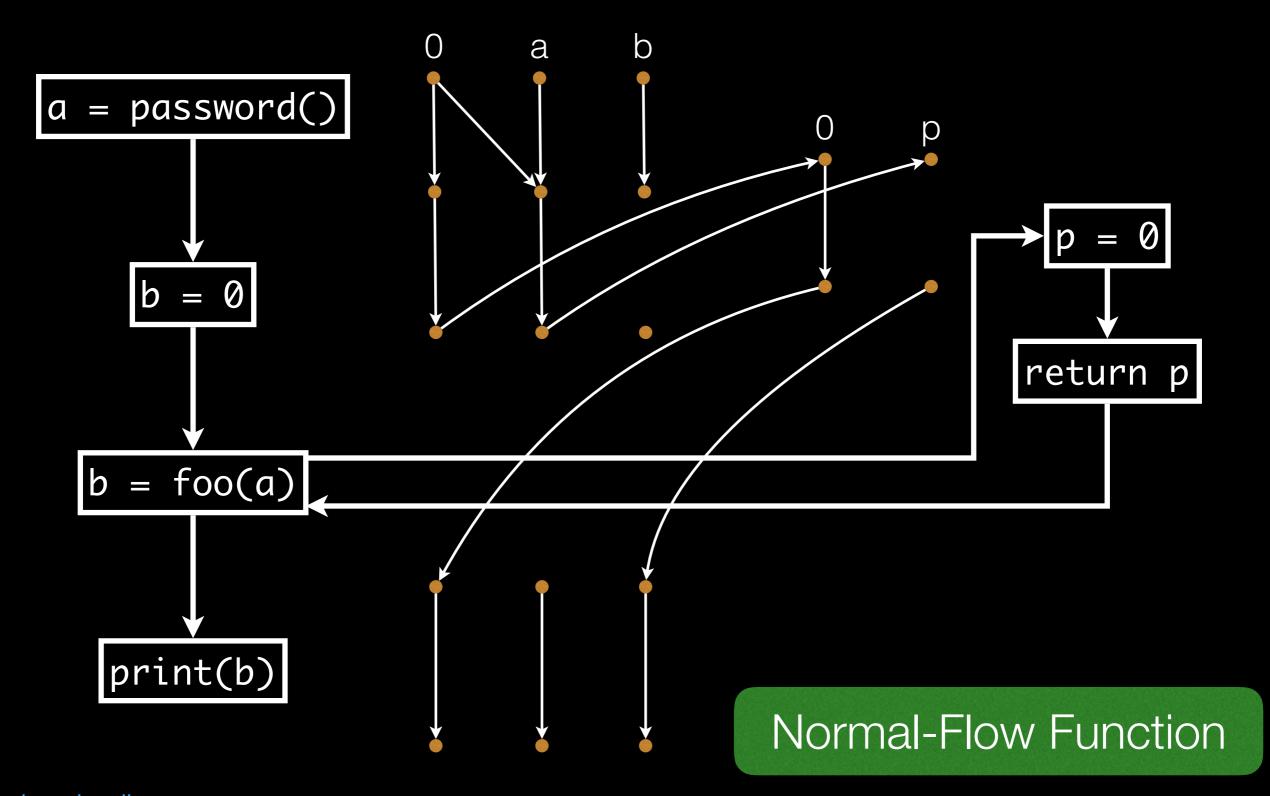


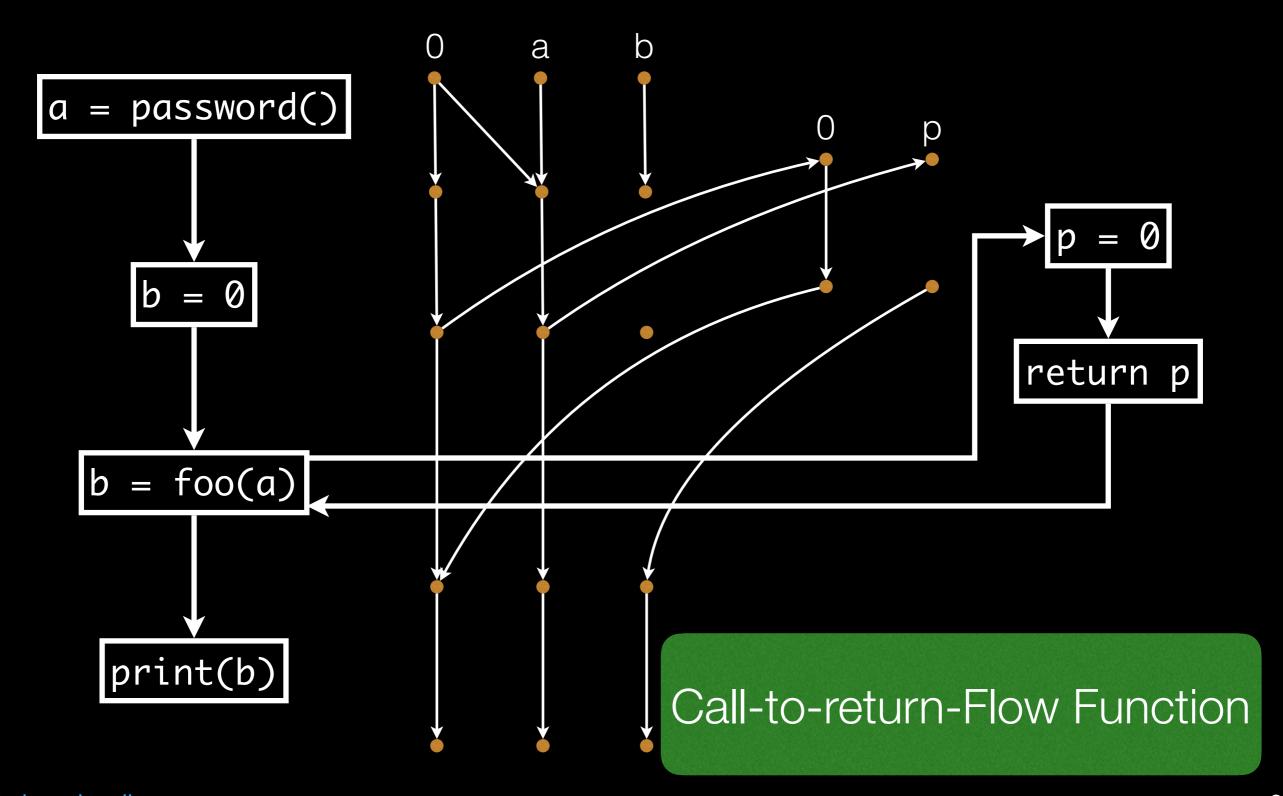






Return-Flow Function





Flow Functions

Call-Flow Function

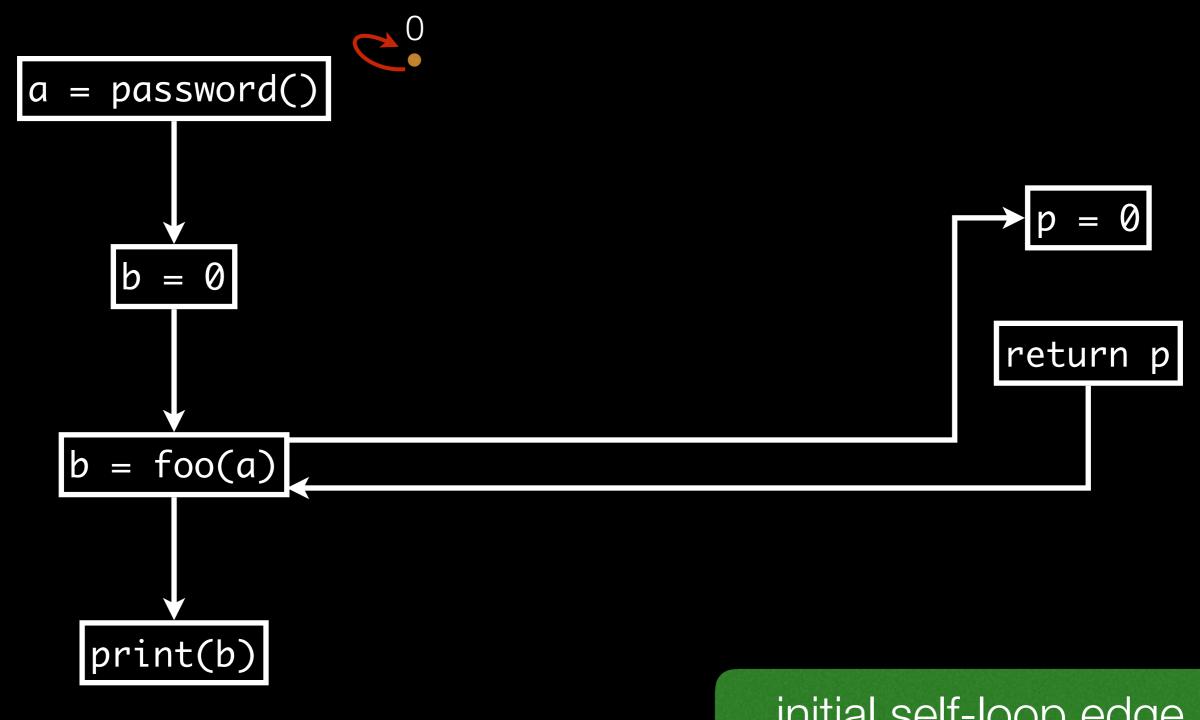
Normal-Flow Function

Return-Flow Function

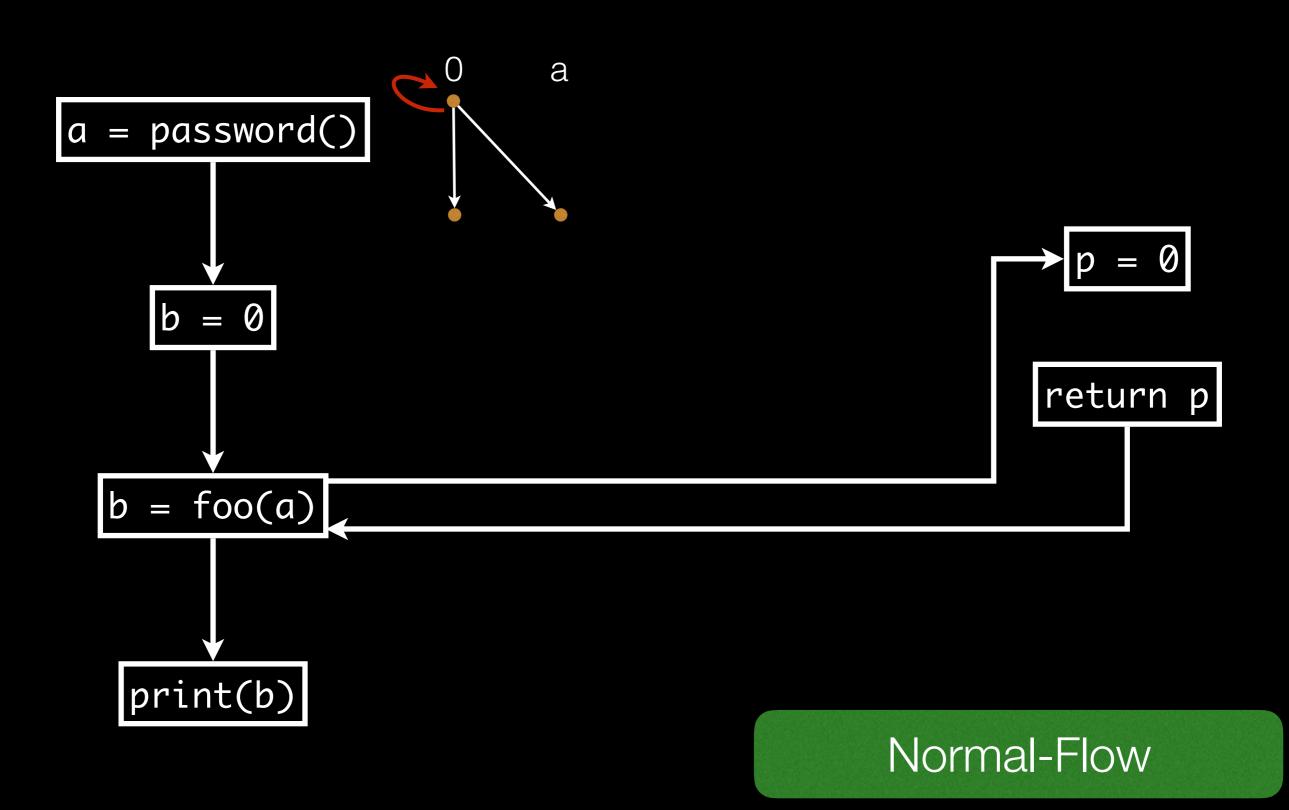


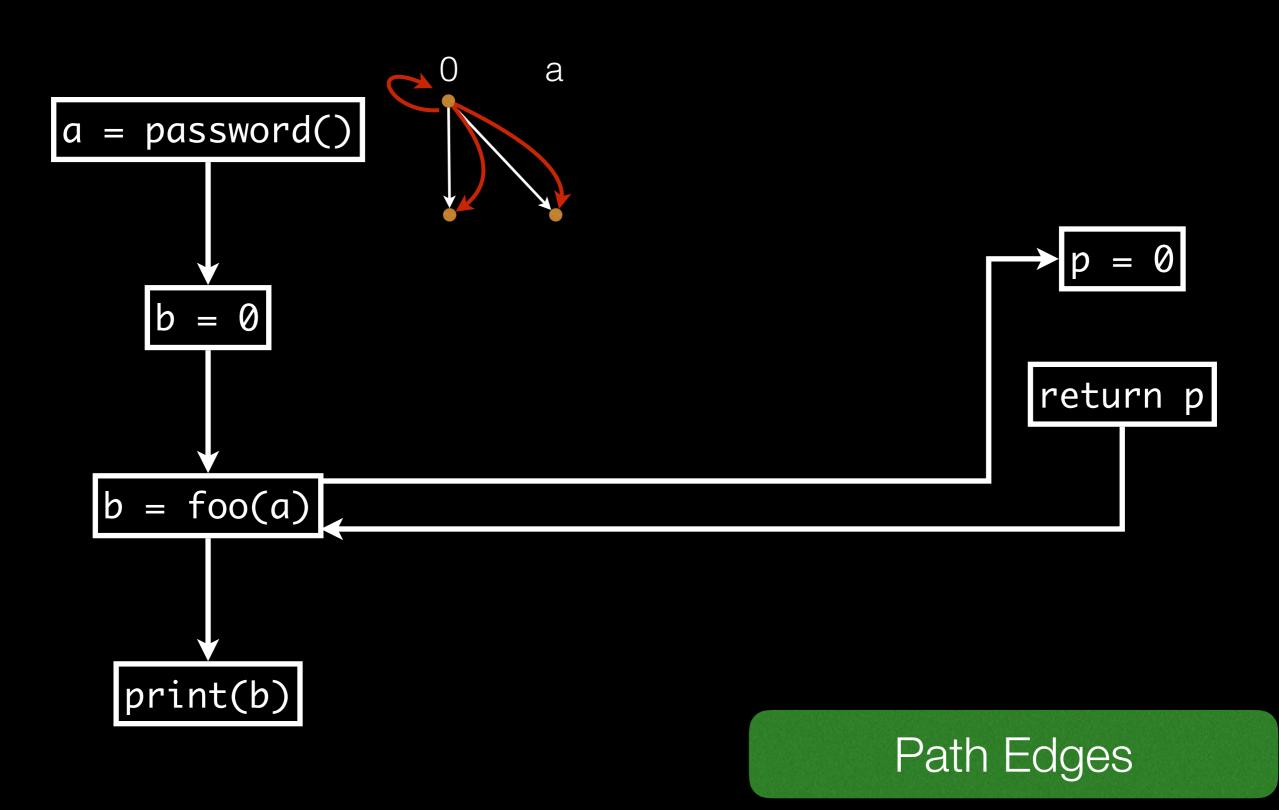
Call-to-return-Flow Function

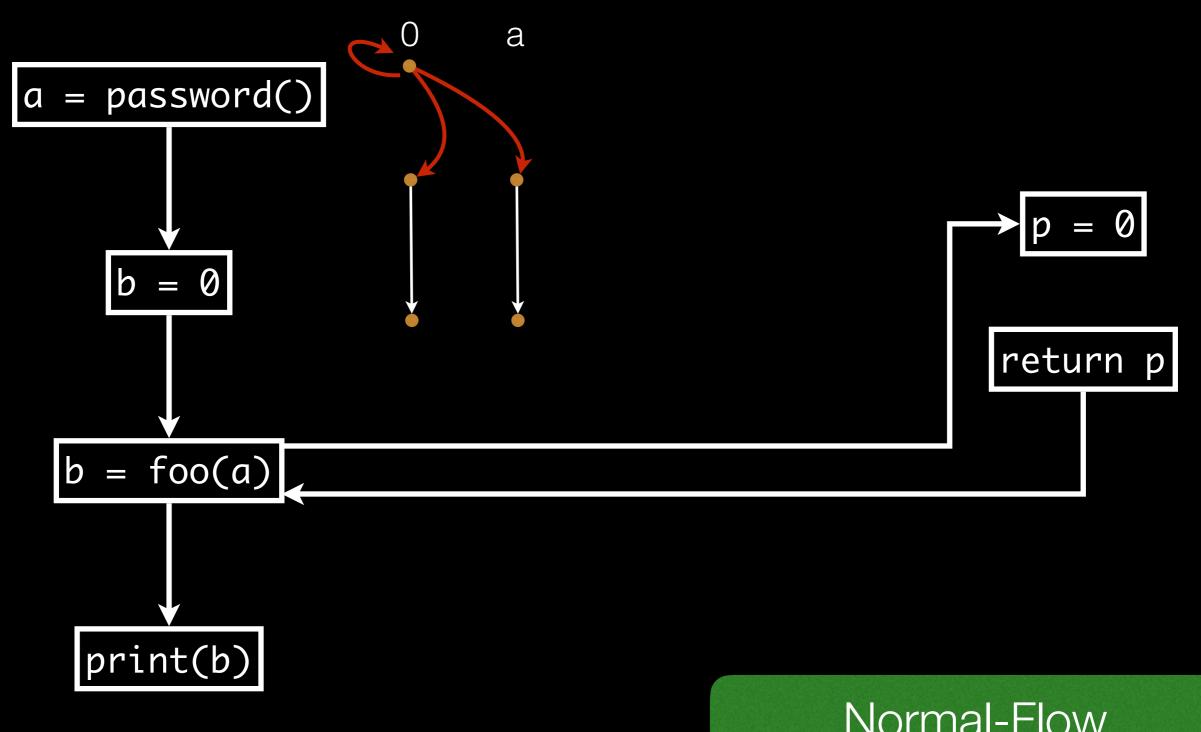




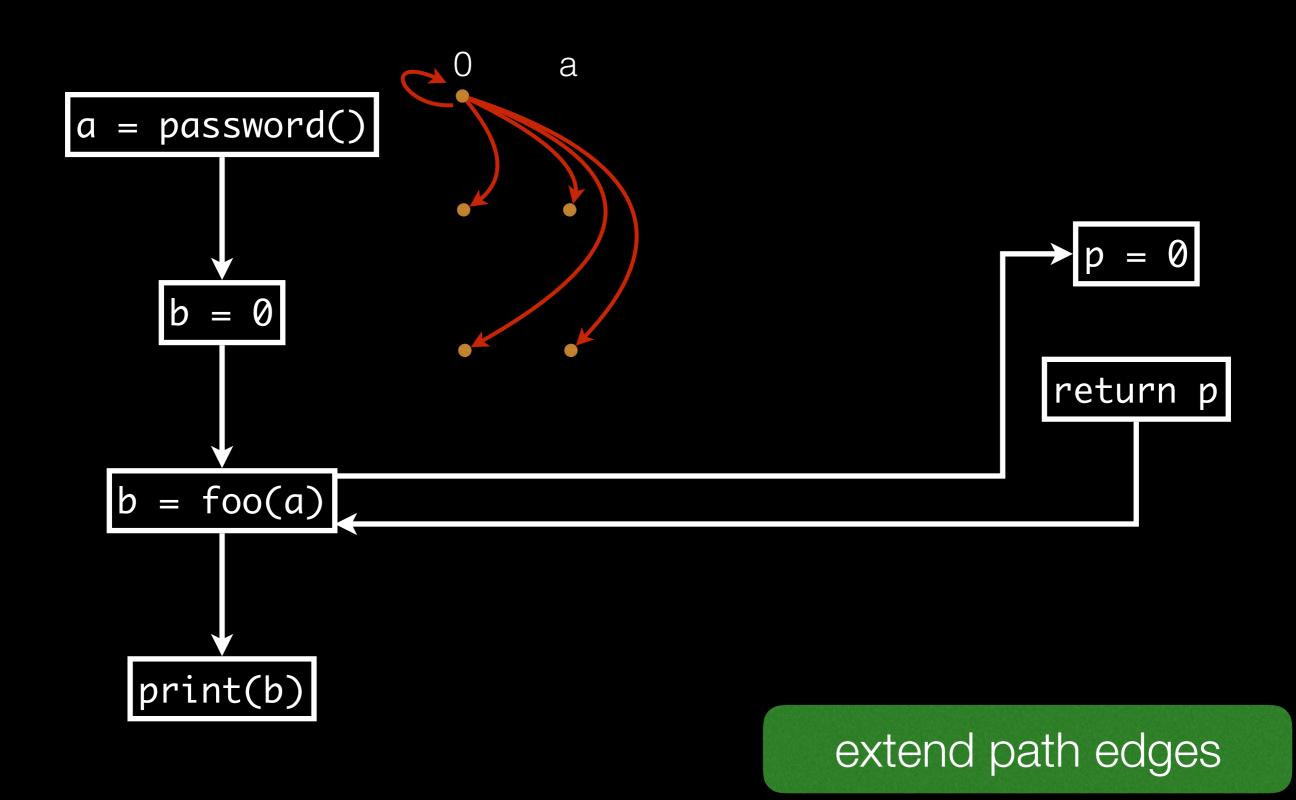
initial self-loop edge

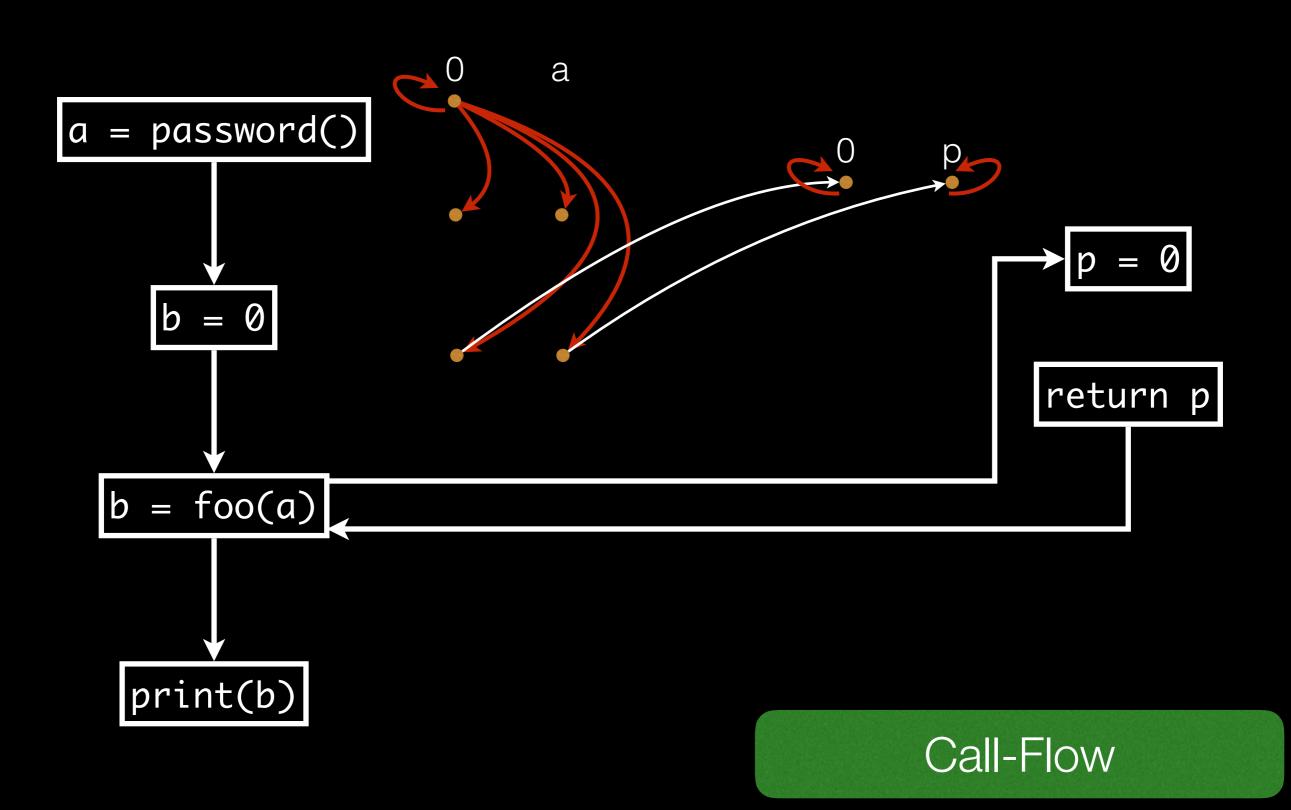


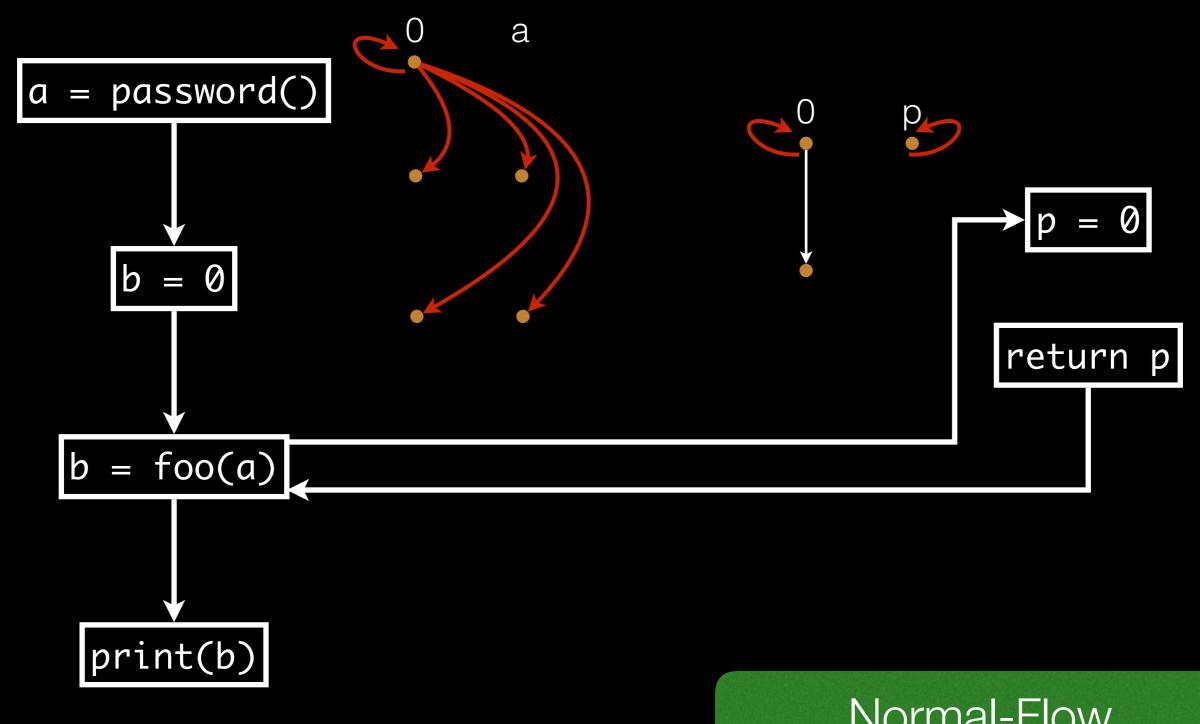




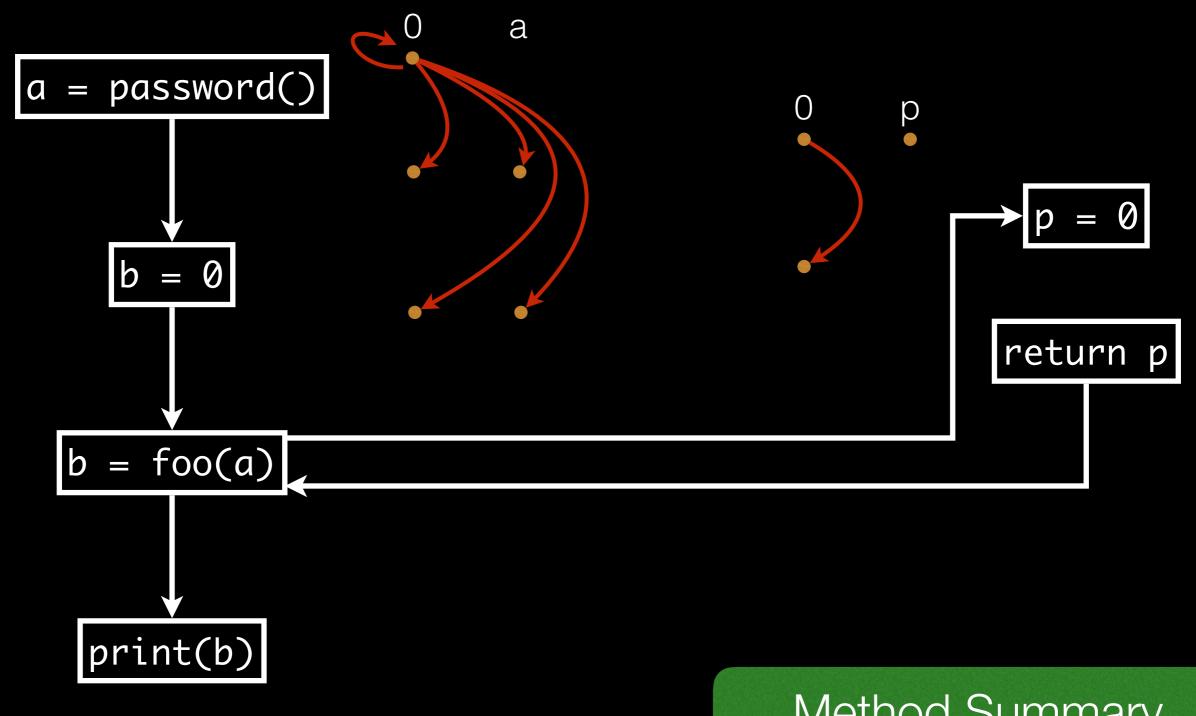
Normal-Flow



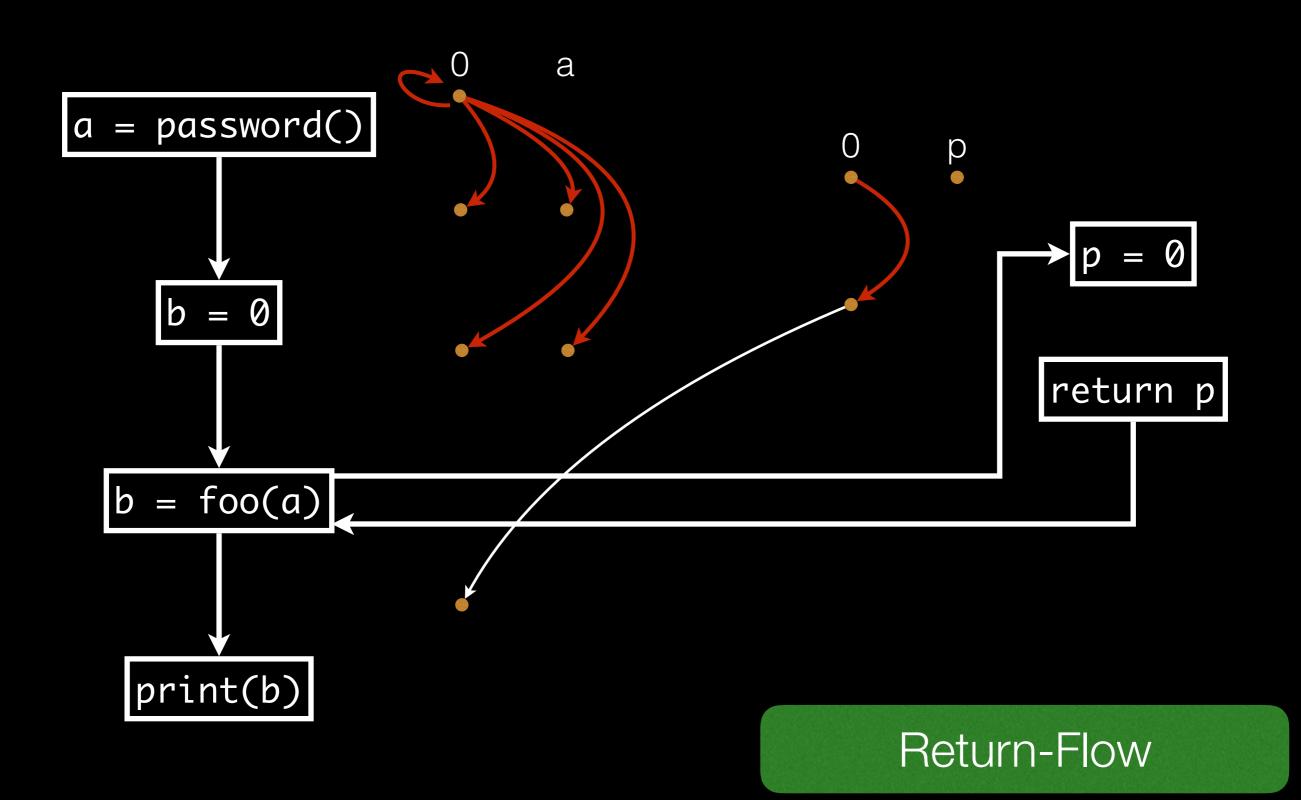


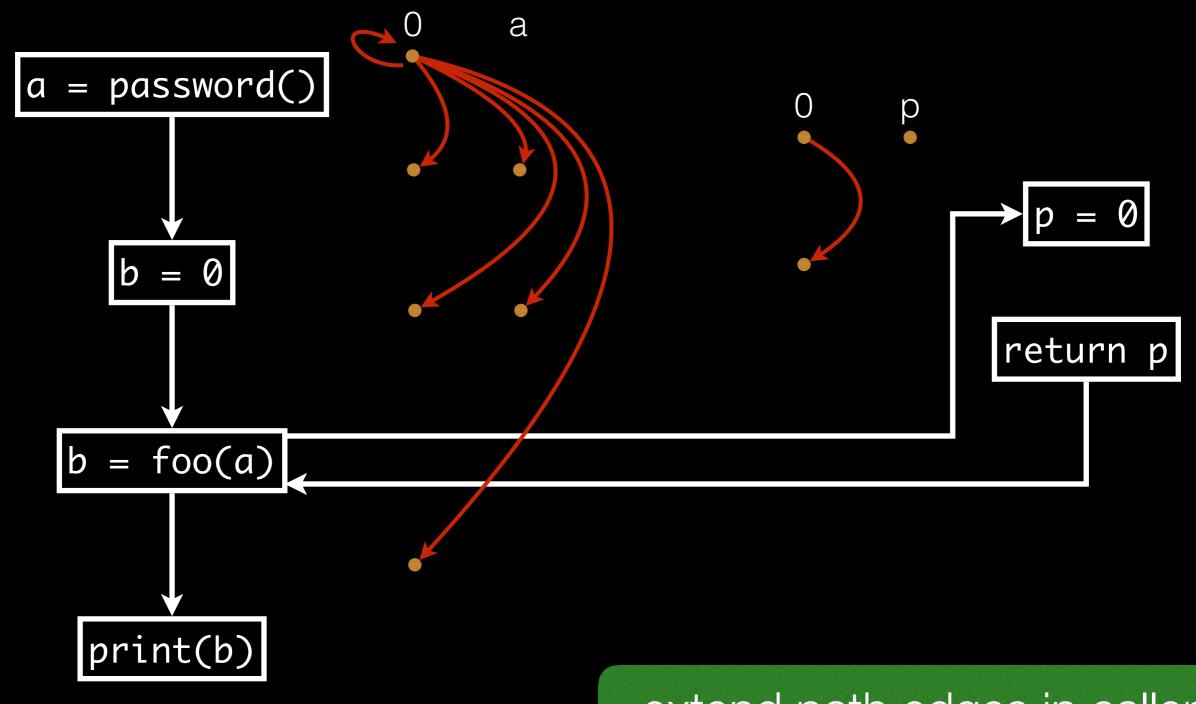


Normal-Flow

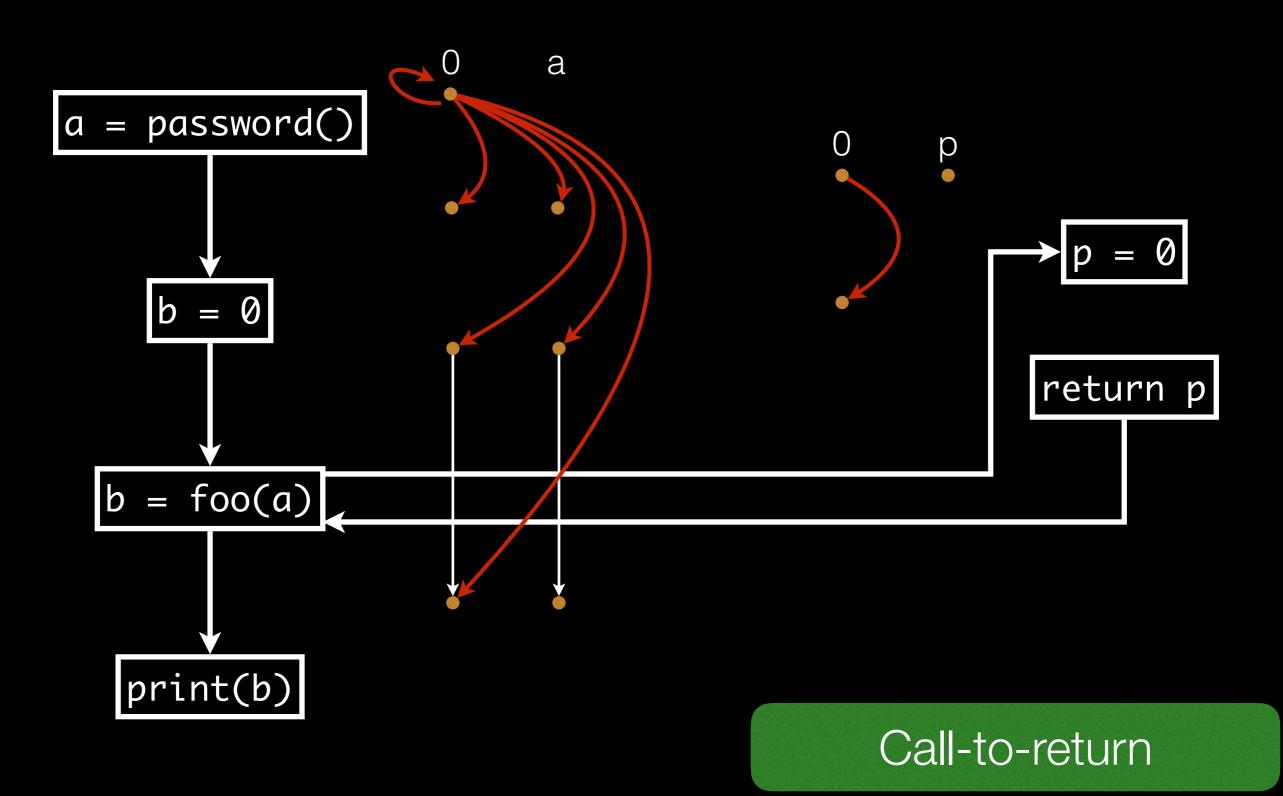


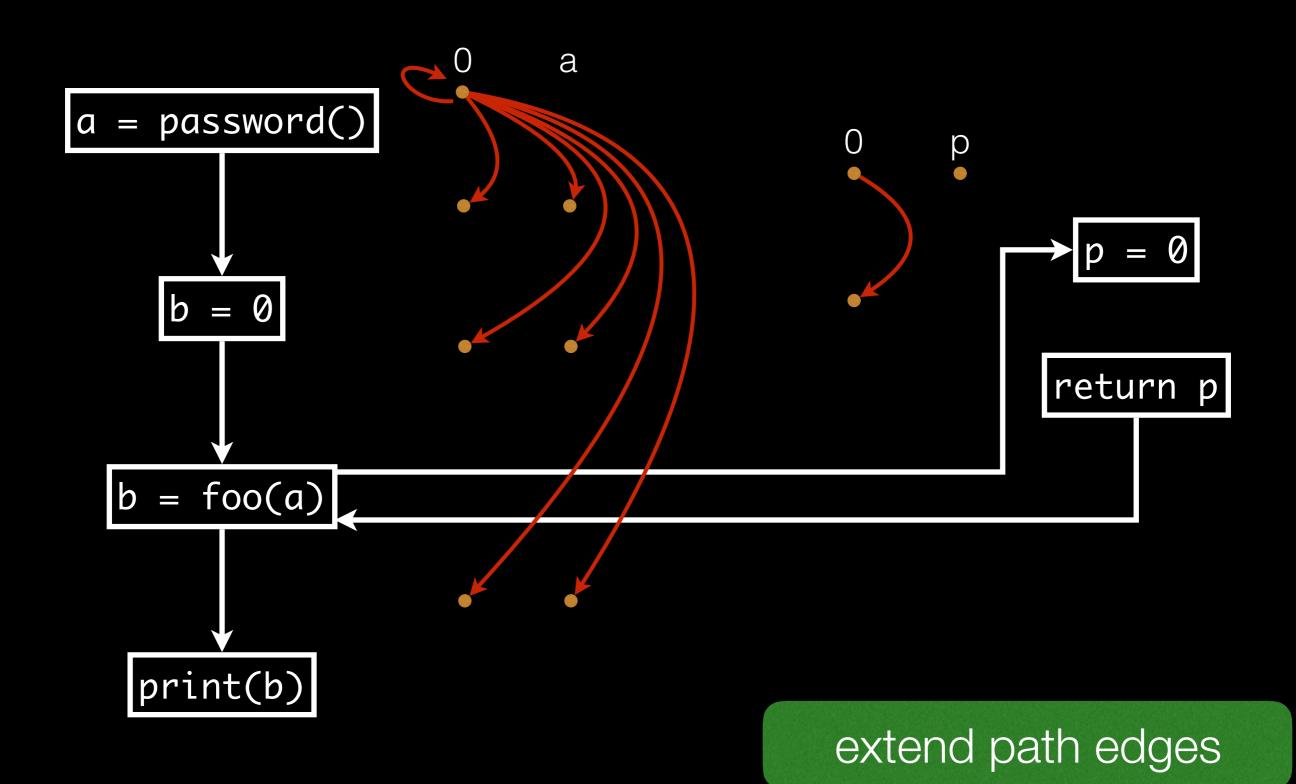
Method Summary

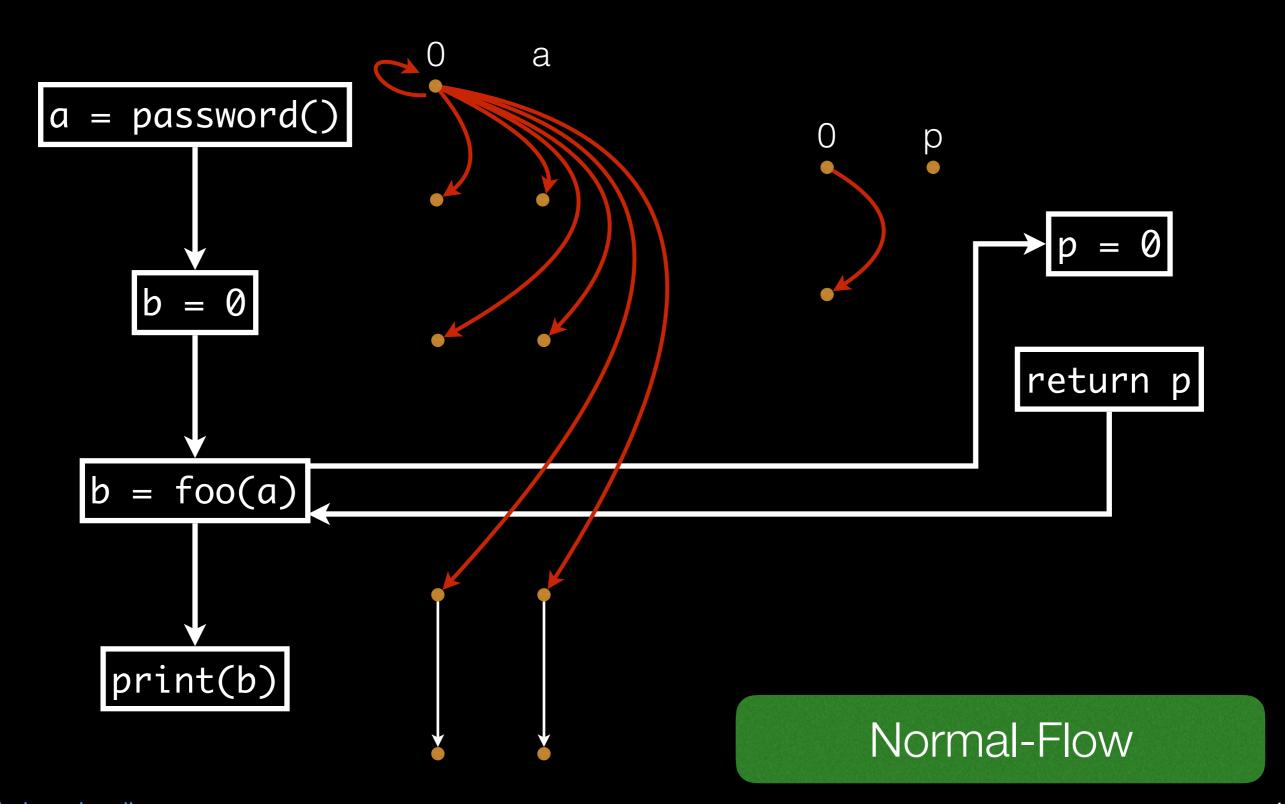


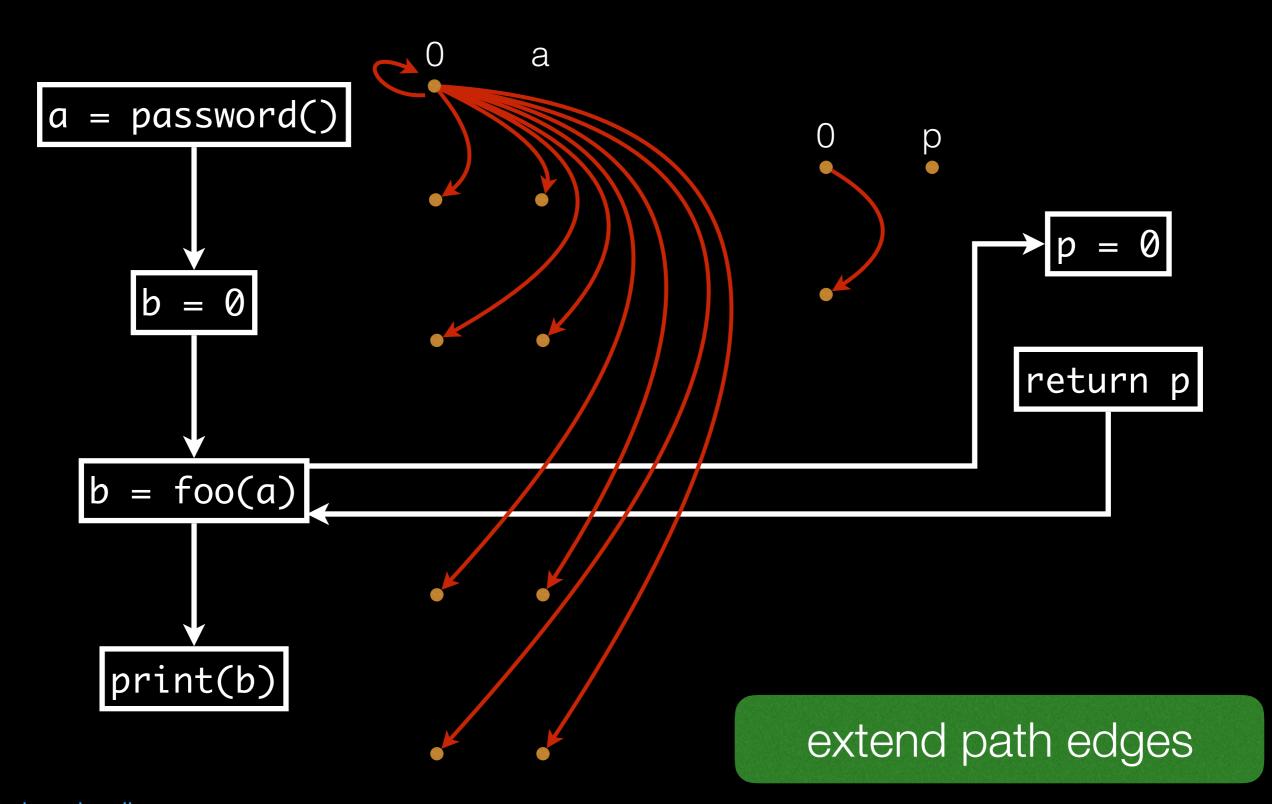


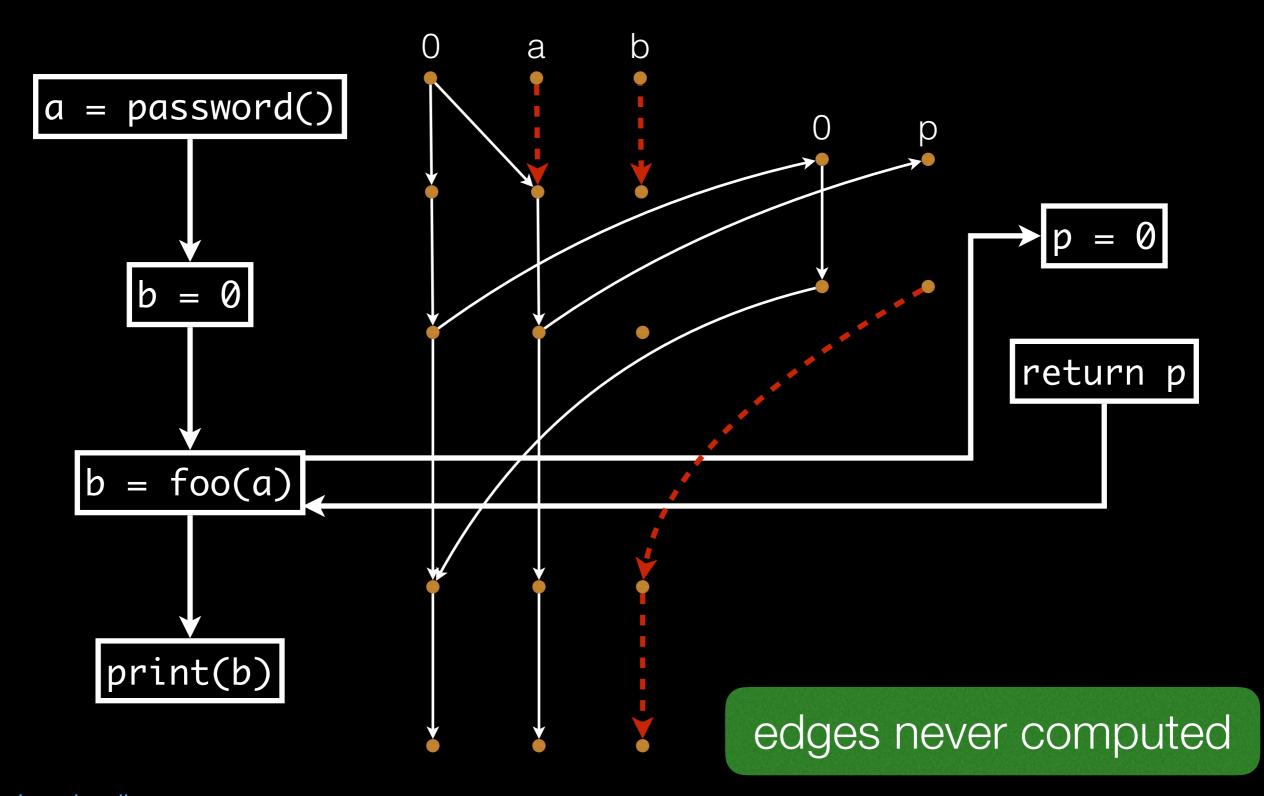
extend path edges in caller











Tidbits!

Field-Sensitivity

a.f

a.*

A.f

Field-Sensitive

Field-Insensitive

Field-Based

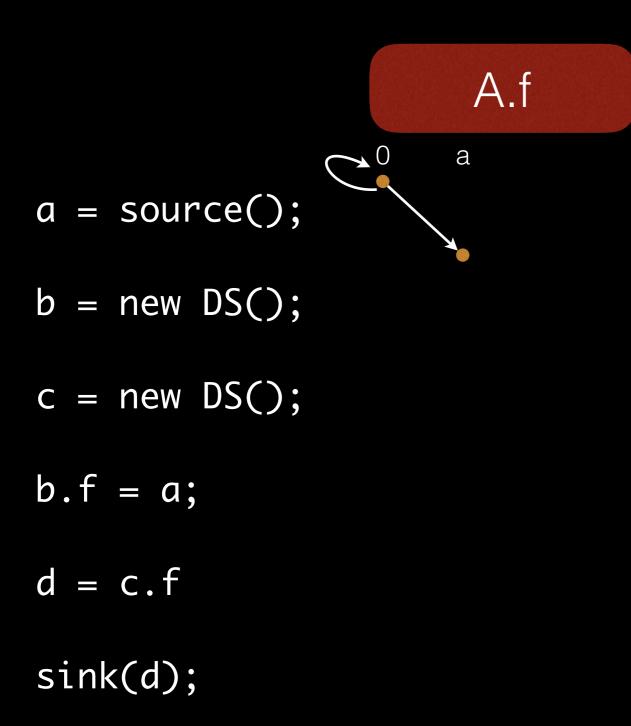
A.f

$$b = new DS();$$

$$c = new DS();$$

$$b.f = a;$$

$$d = c.f$$





a



$$b = new DS();$$

$$c = new DS();$$

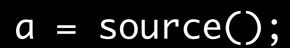
$$b.f = a;$$

$$d = c.f$$

sink(d);



a



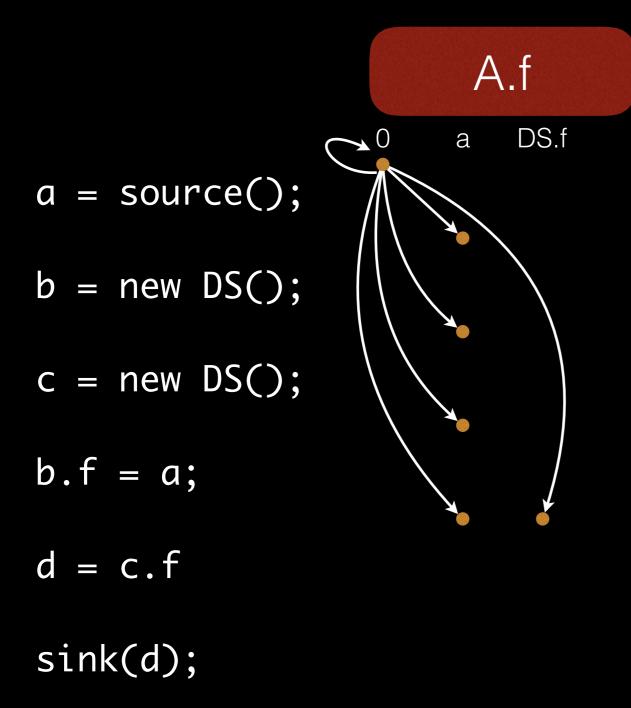
b = new DS();

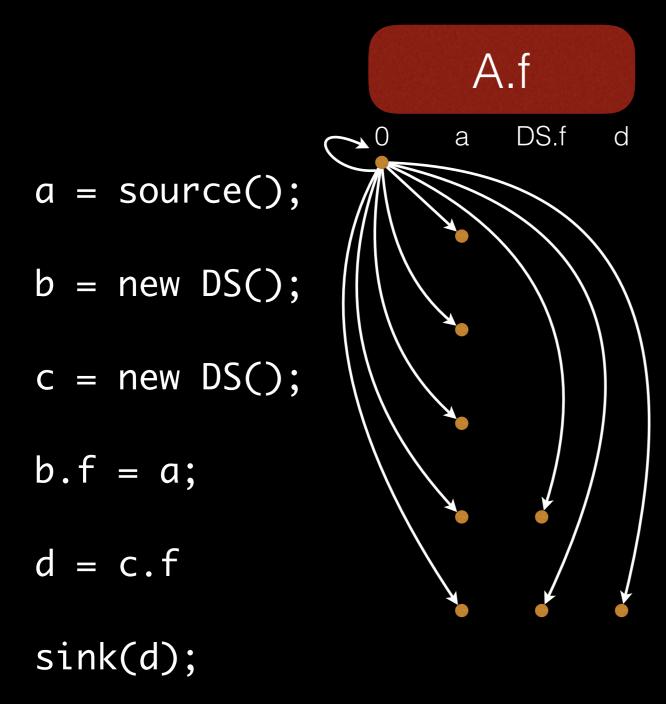
c = new DS();

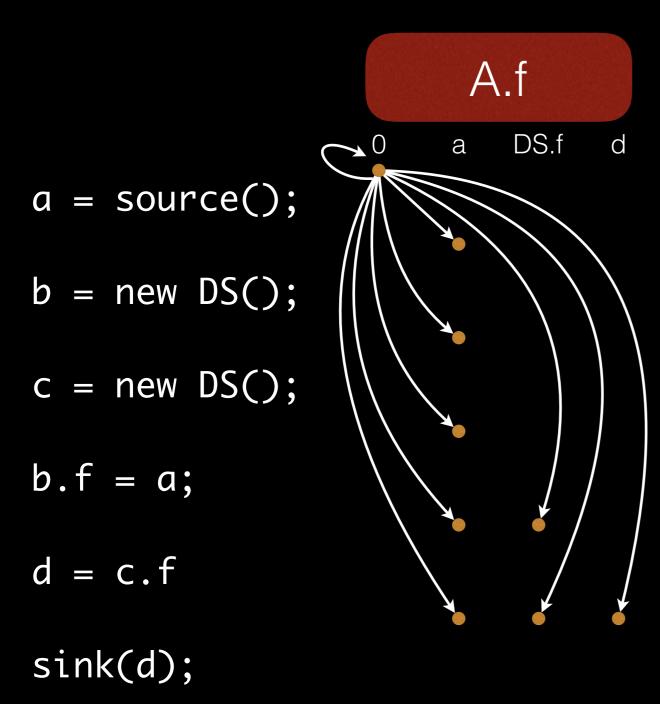
$$b.f = a;$$

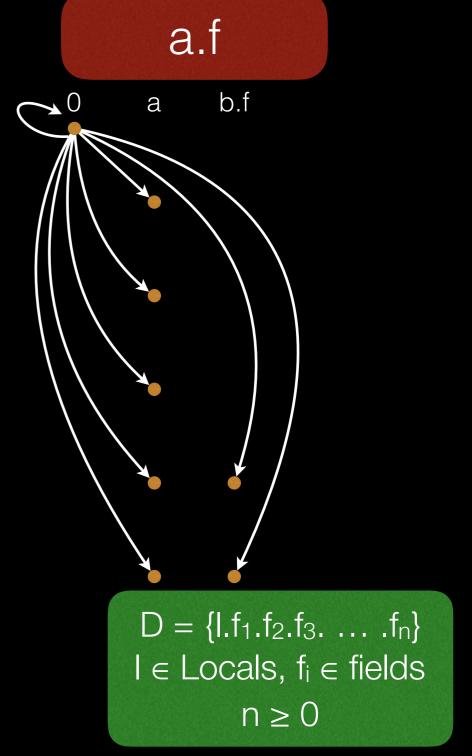
$$d = c.f$$

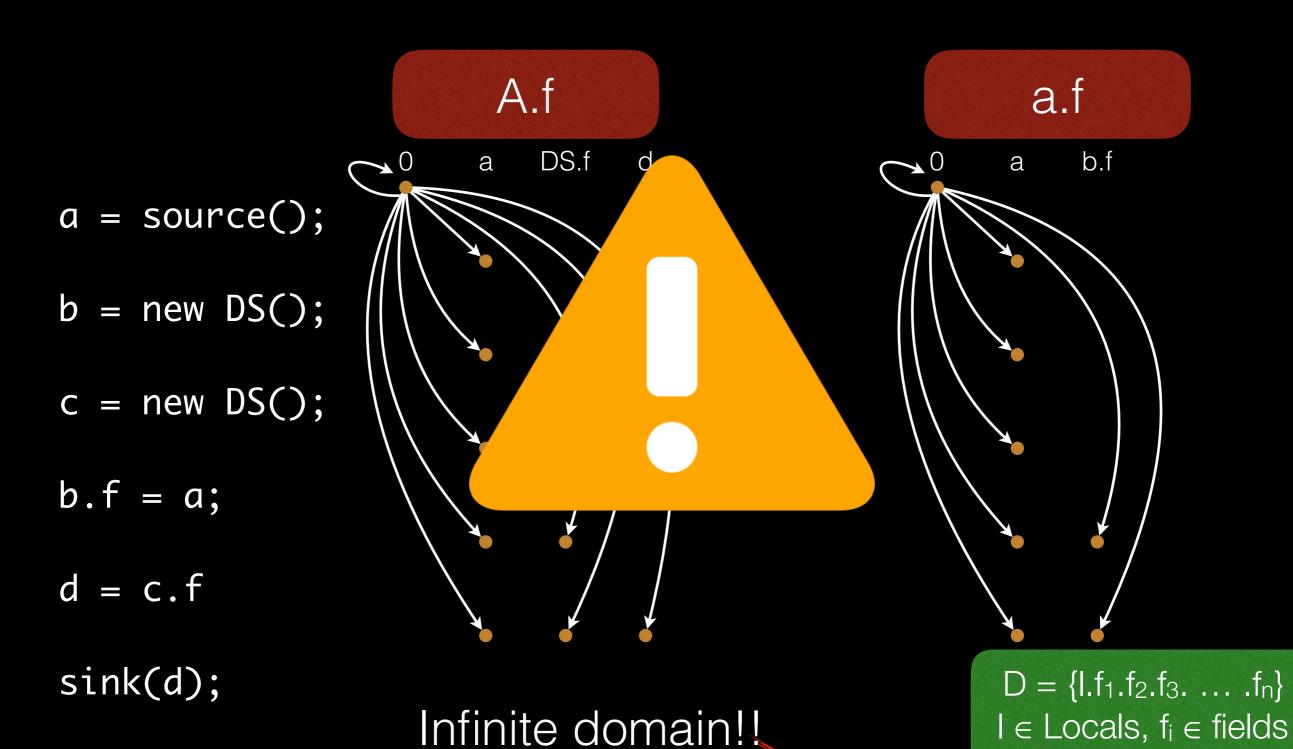
sink(d);









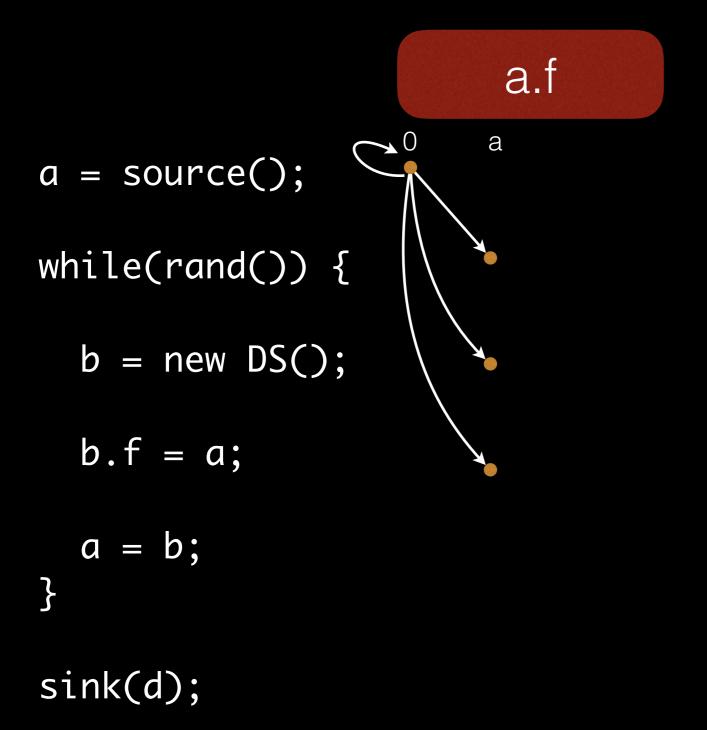


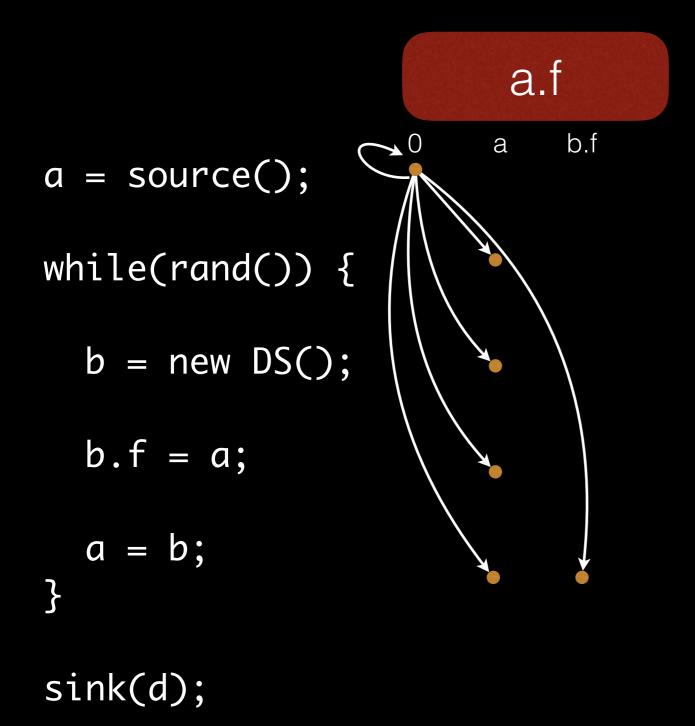
@karimhamdanali

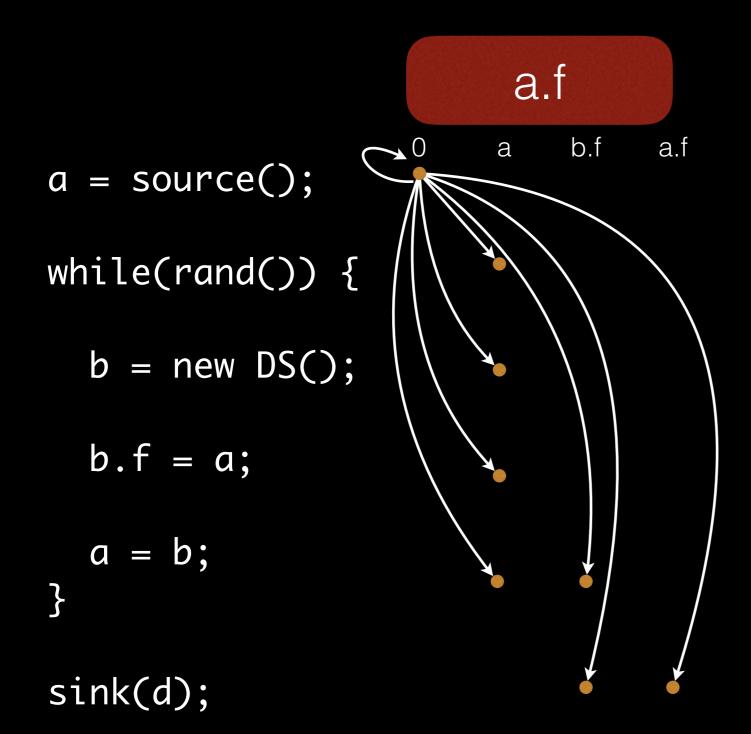
>n ≥ 0

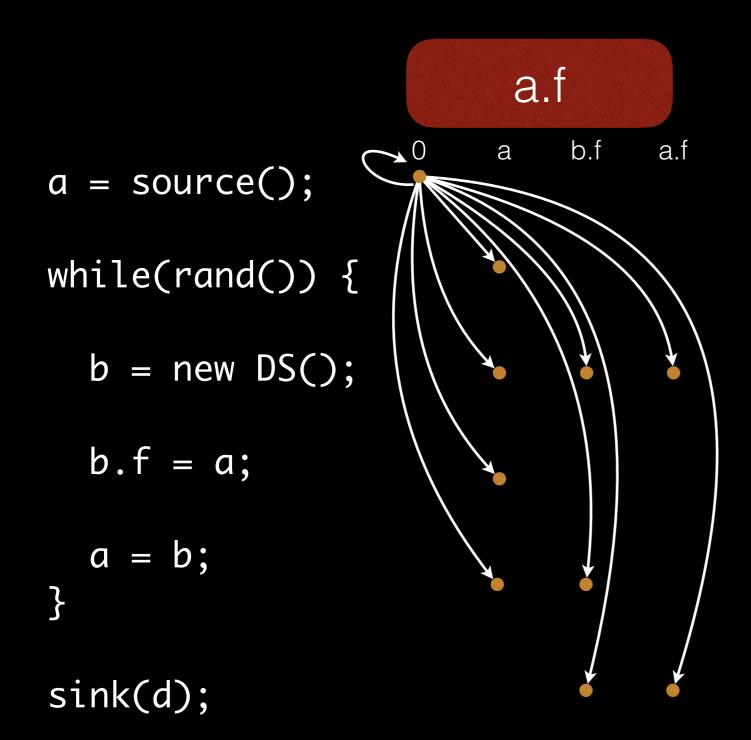
```
a.f
                       a
a = source();
while(rand()) {
  b = new DS();
  b.f = a;
  a = b;
sink(d);
```

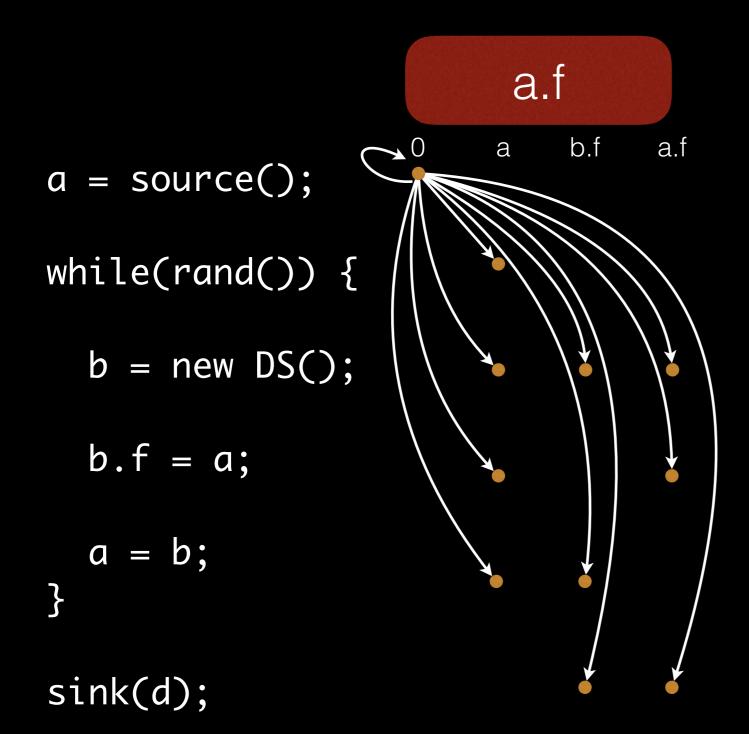
a.f a a = source();while(rand()) { b = new DS();b.f = a;a = b;sink(d);

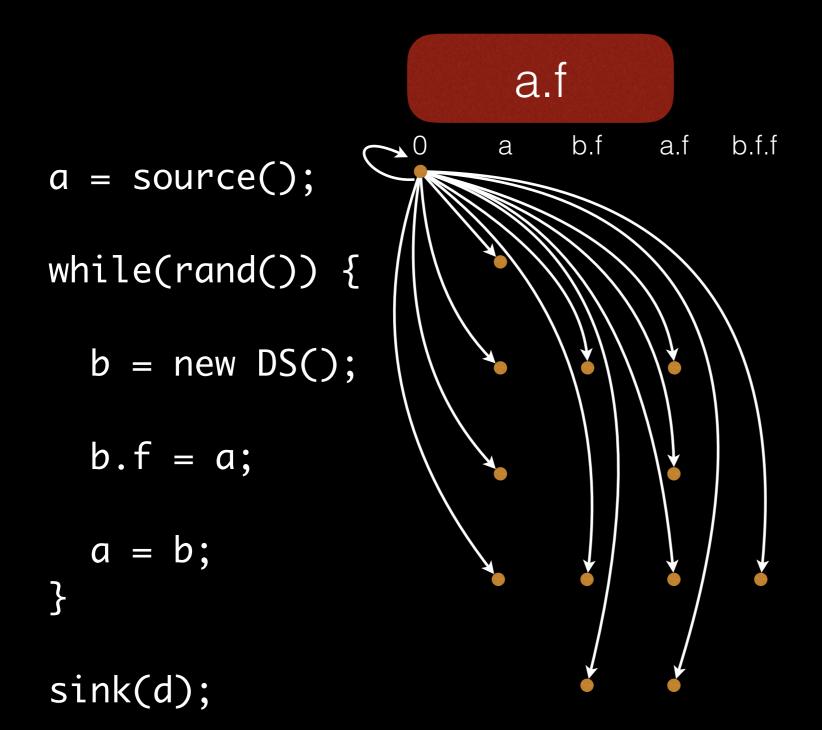


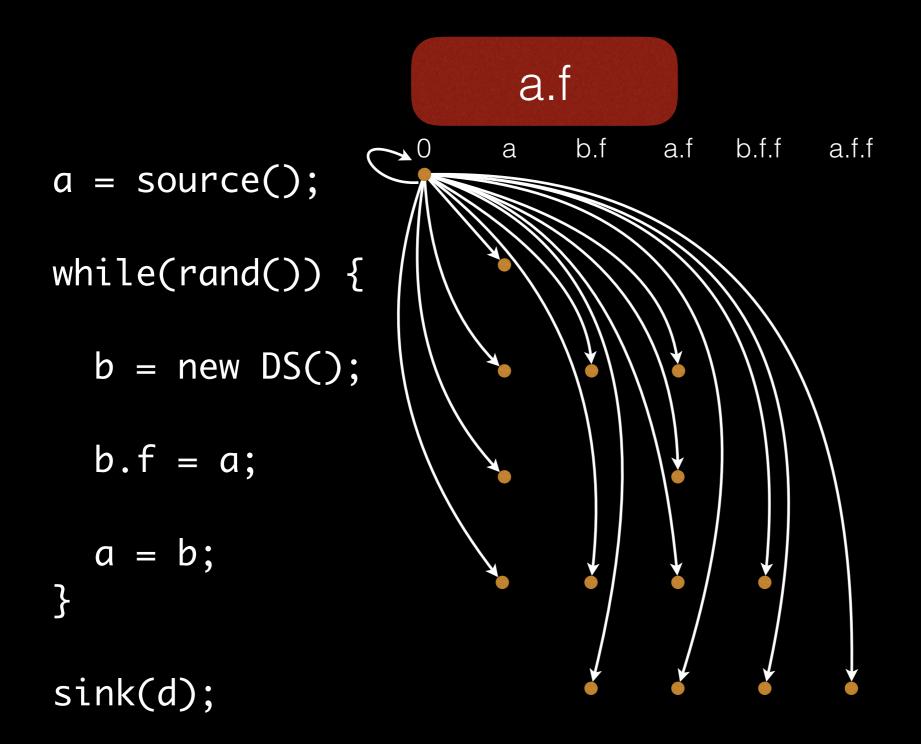


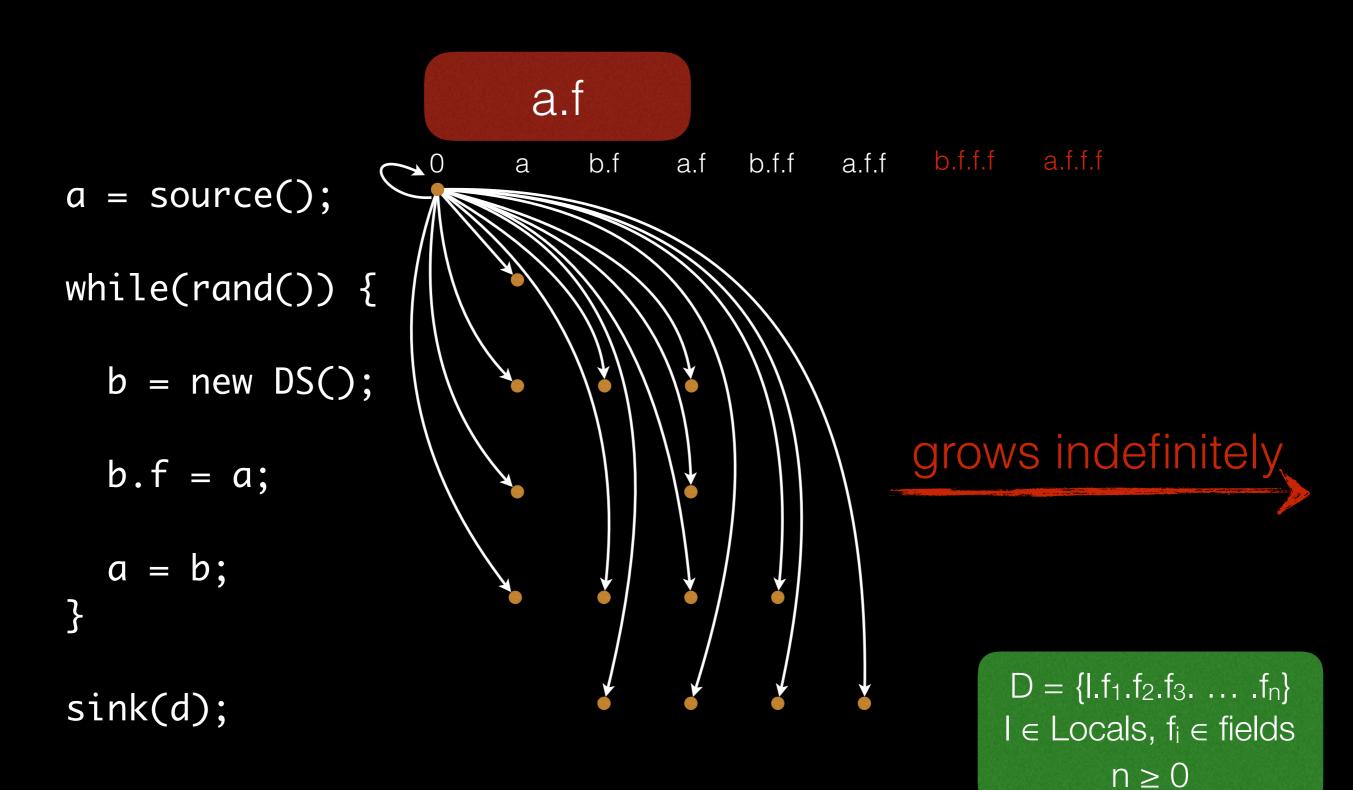


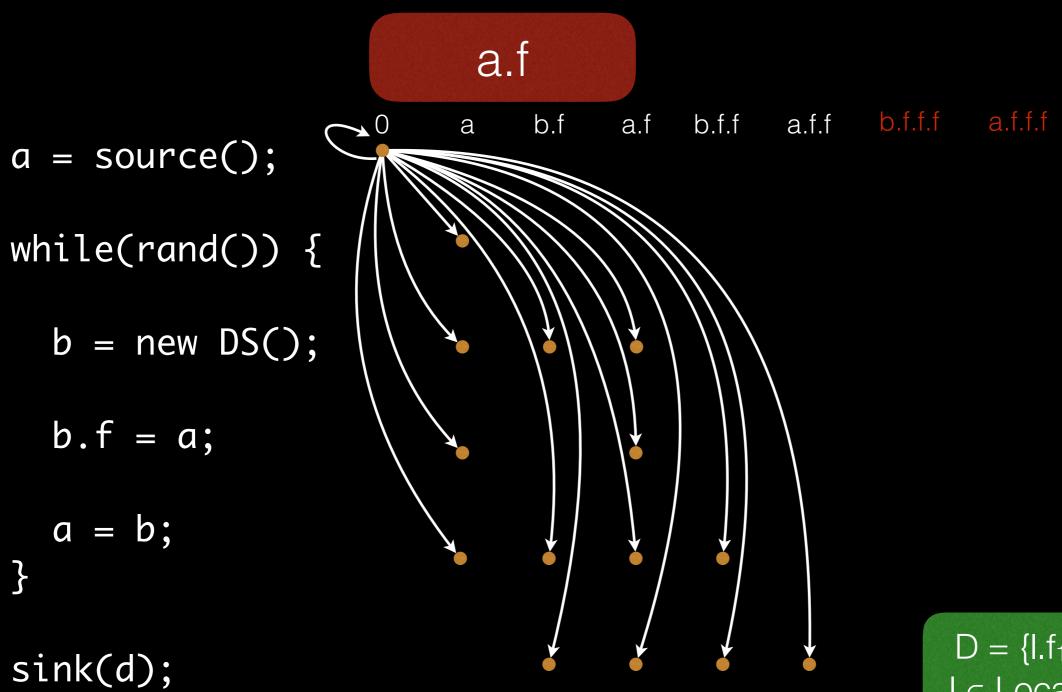


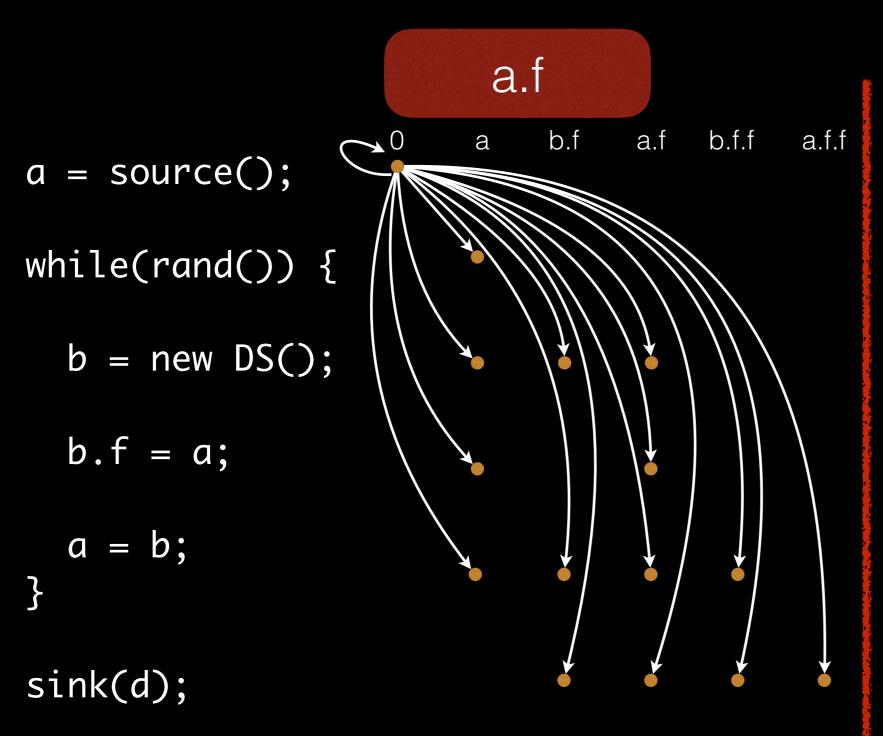




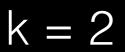


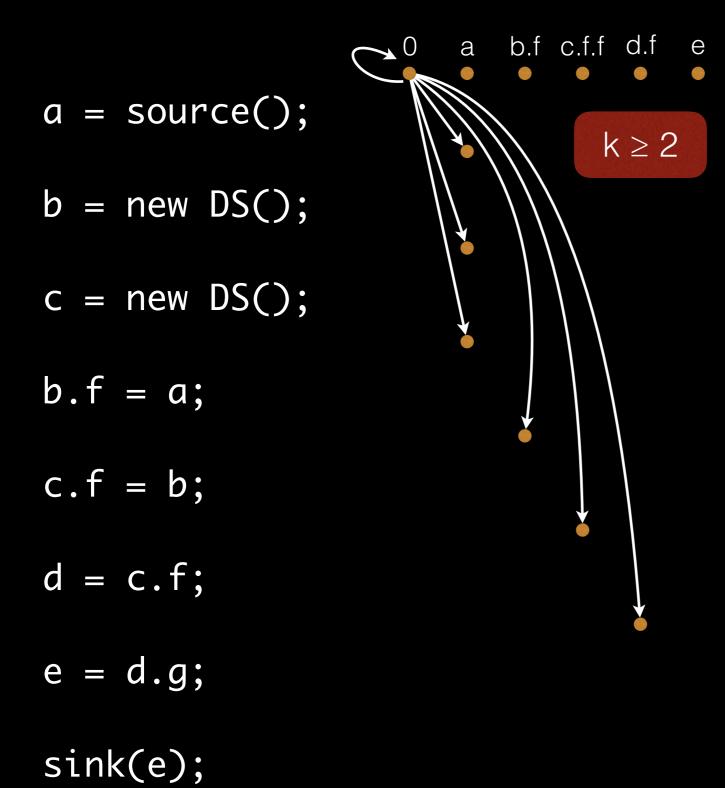


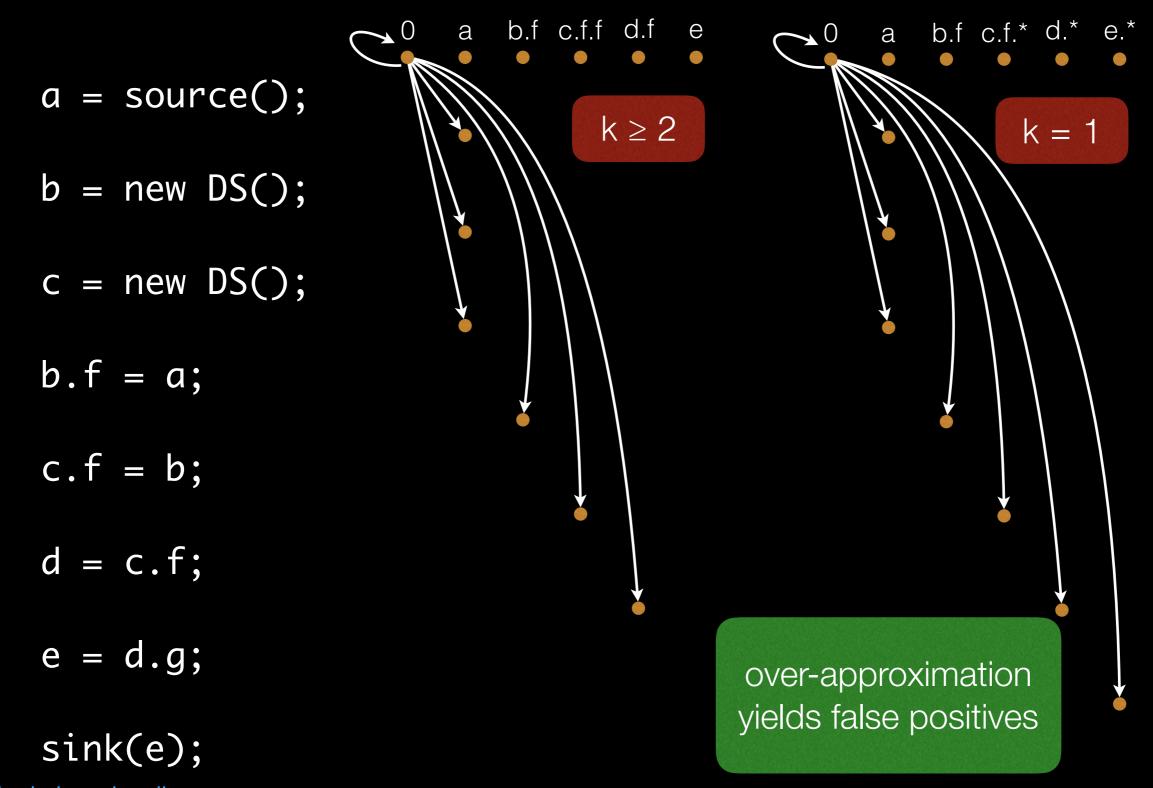




unsound domain beyond k = > imprecise







Tidbits

- Flow Sensitivity?
- Context Sensitivity?

Tidbits

Flow Sensitivity?

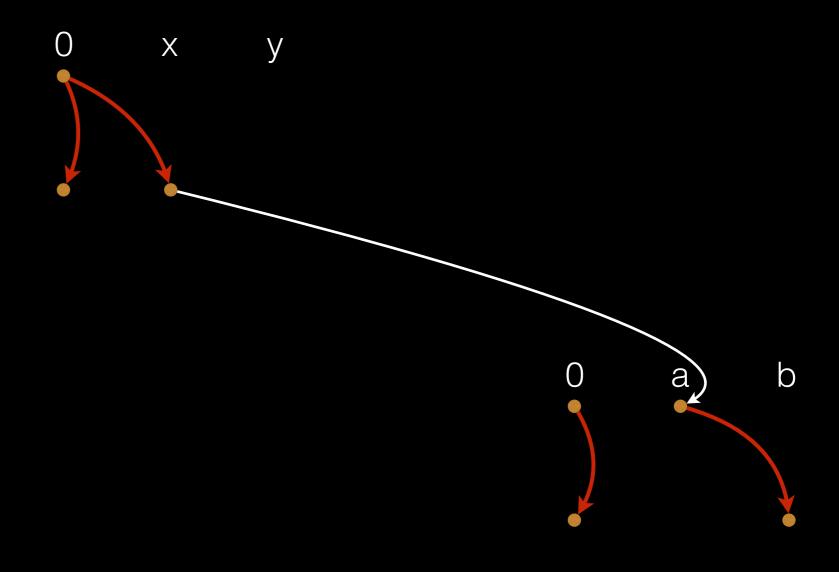
Method Summaries

Context Sensitivity?

• •

call1();

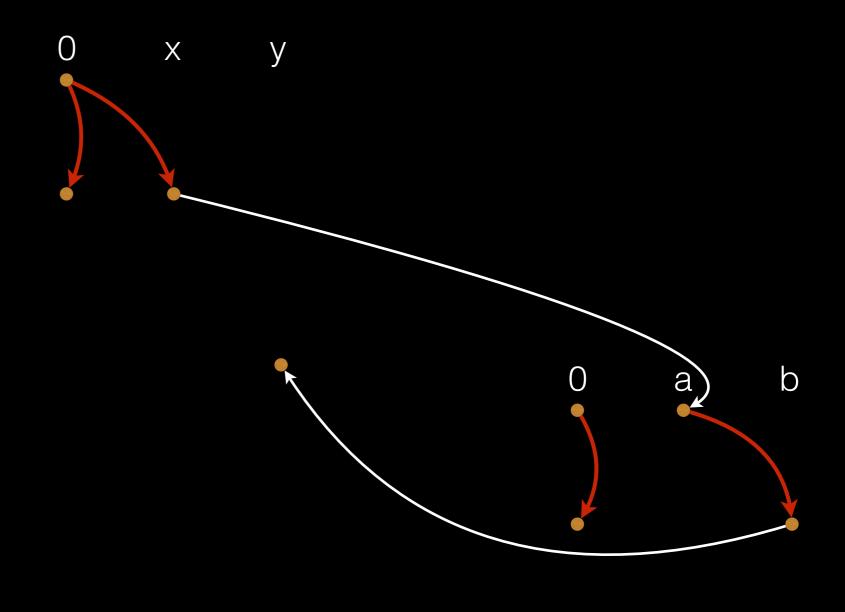
• •



• • •

call1();

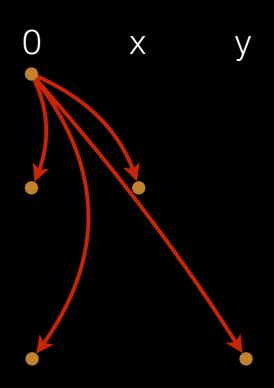
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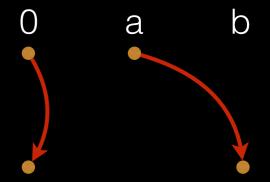


• • •

call1();

• •

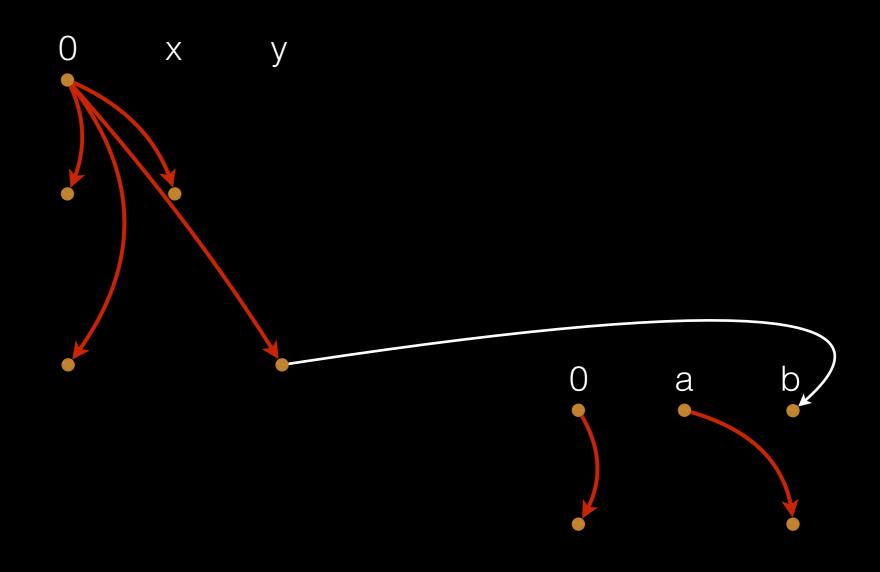




• • •

call1();

• •

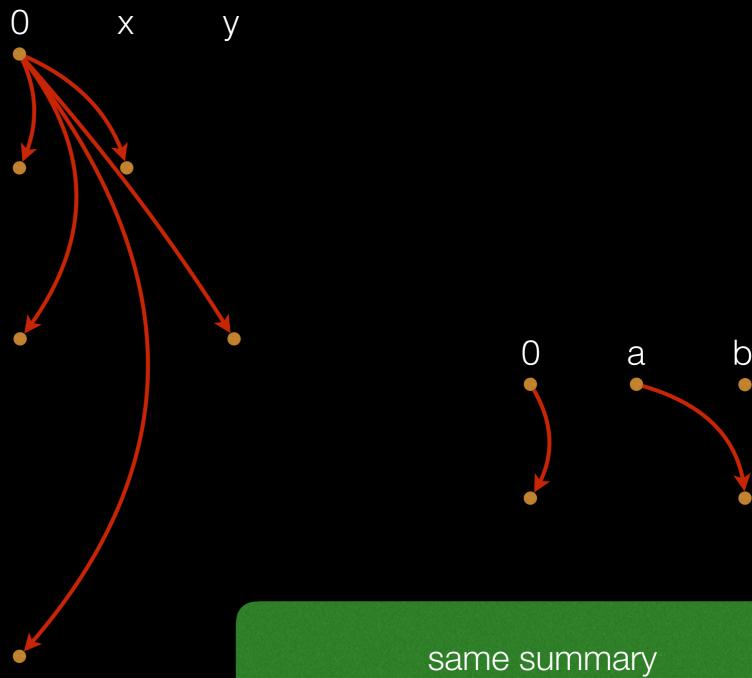


• •

call1();

• •

call2();



different contexts => different results

Tidbits

- Flow Sensitivity
- Context Sensitivity
- Large Domains => performance problems

$$r = v + 1$$
 $0 \quad (v,0) \quad (r,1)$

In any context, if v is 0 before the call, then it is true that r is 1 after the call

```
main() {
                         inc(v) {
x = inc(1);
y = inc(x);
z = inc(y);
 print(z);
      (v,0) (r,1)
```

r = v + 1;return r;

```
main() {
                          inc(v) {
x = inc(1);
                            r = v + 1;
y = inc(x);
                            return r;
 z = inc(y);
 print(z);
       (v,0) (r,1) (v,1) (r,2) (v,2) (r,3)
```

```
main() {
                              inc(v) {
 x = inc(1);
                                 = \vee + 1.
 V = inc(V)
           Solution: IDE instead of IFDS
         (v,0) (r,1) (v,1) (r,2) (v,2)
```

Recap

- Interprocedural Finite Distributive Subset
- Flow Functions
- Taint Analysis Example
- On-the-fly ESG
- Tidbits: field-sensitivity, context-sensitivity, large domains

On Thursday

- Proposal presentations
- Solving Assignments 1 and 2 in class