

Inter-Procedural Analysis

CMPUT 497/500 Foundations of Program Analysis

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Previously

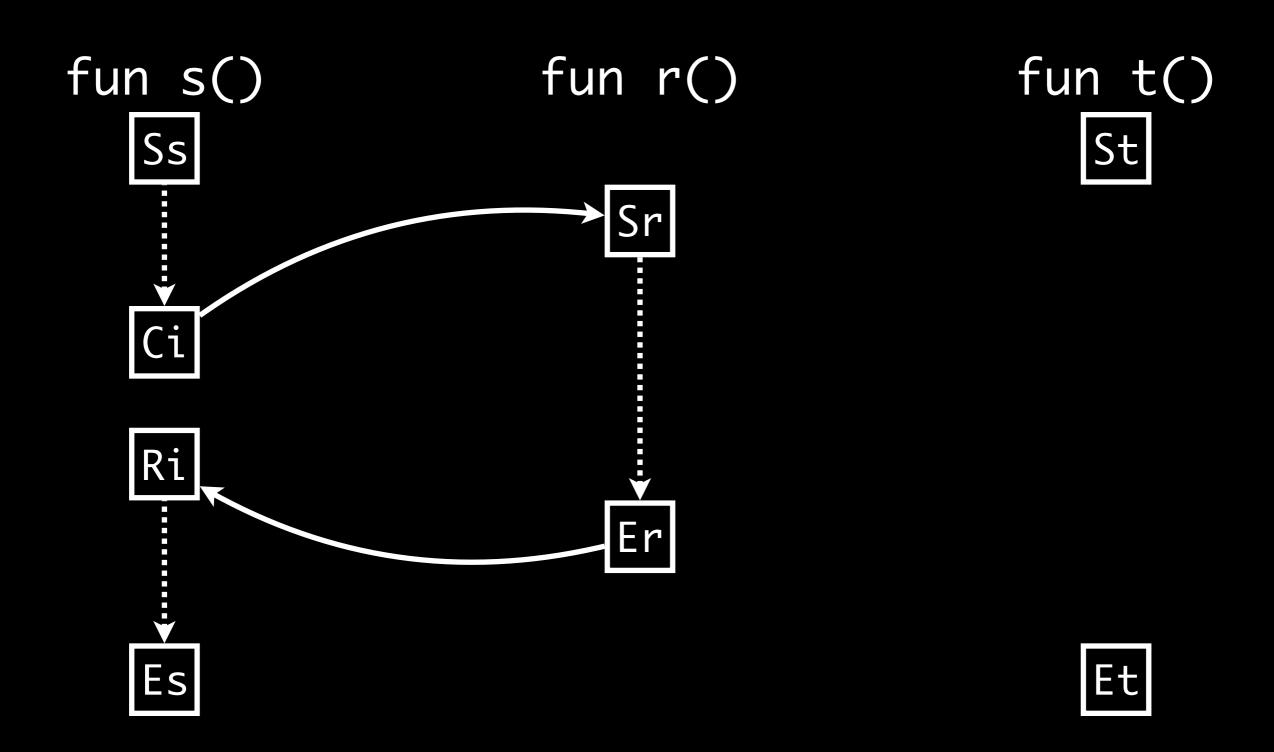
- Points-to
- Aliases
- Must and May analyses
- Incomplete Programs
- Weak vs Strong Updates
- Access Paths
- Distributivity

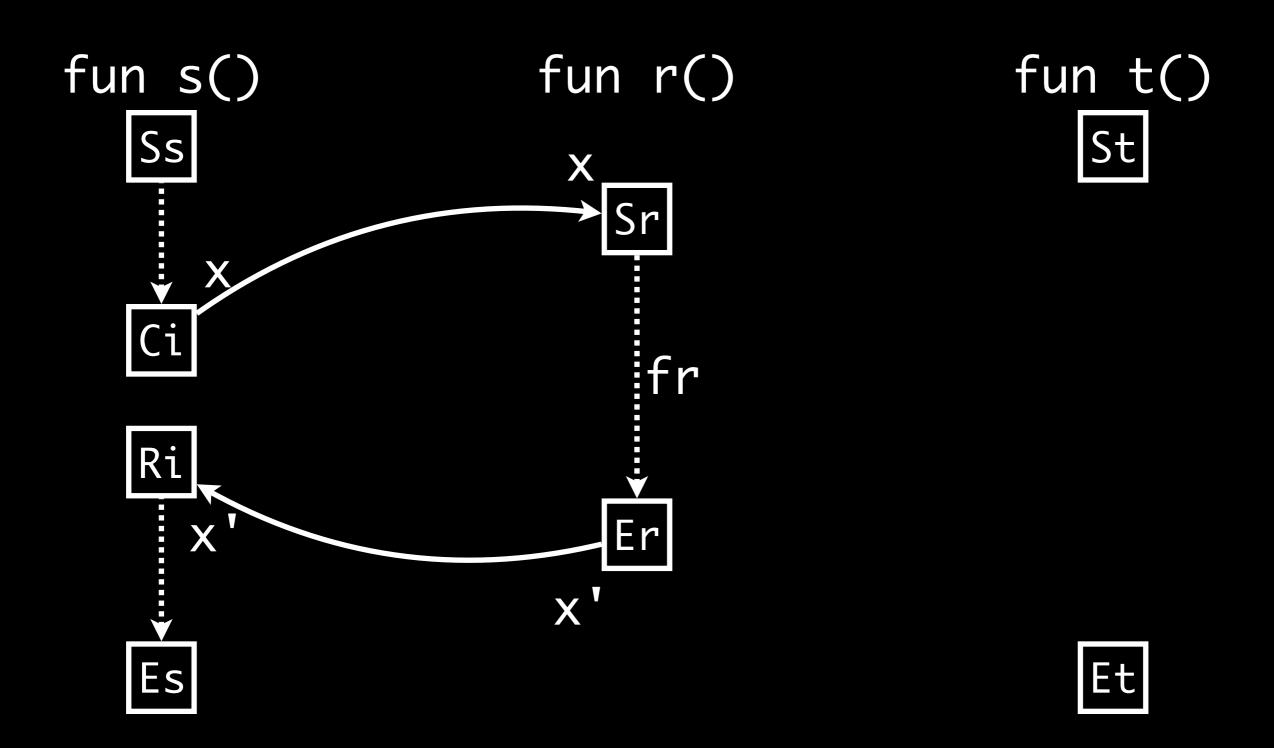
- Beyond procedure boundaries
- Model the effects of
 - calls in the callers, and
 - calling contexts in the callees

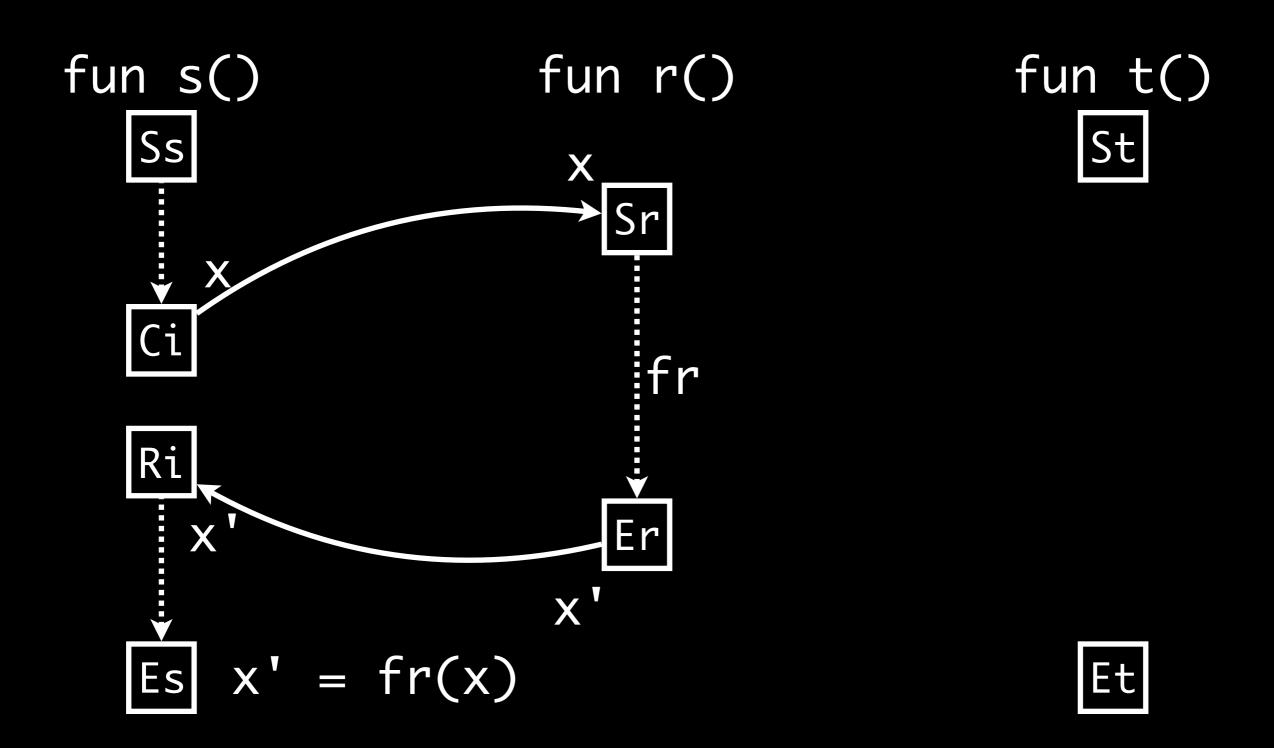
- Approaches
 - Generic: Call-strings approach, functional approach
 - Problem specific: Alias analysis,
 Points-to analysis,
 Partial redundancy elimination,
 Constant propagation

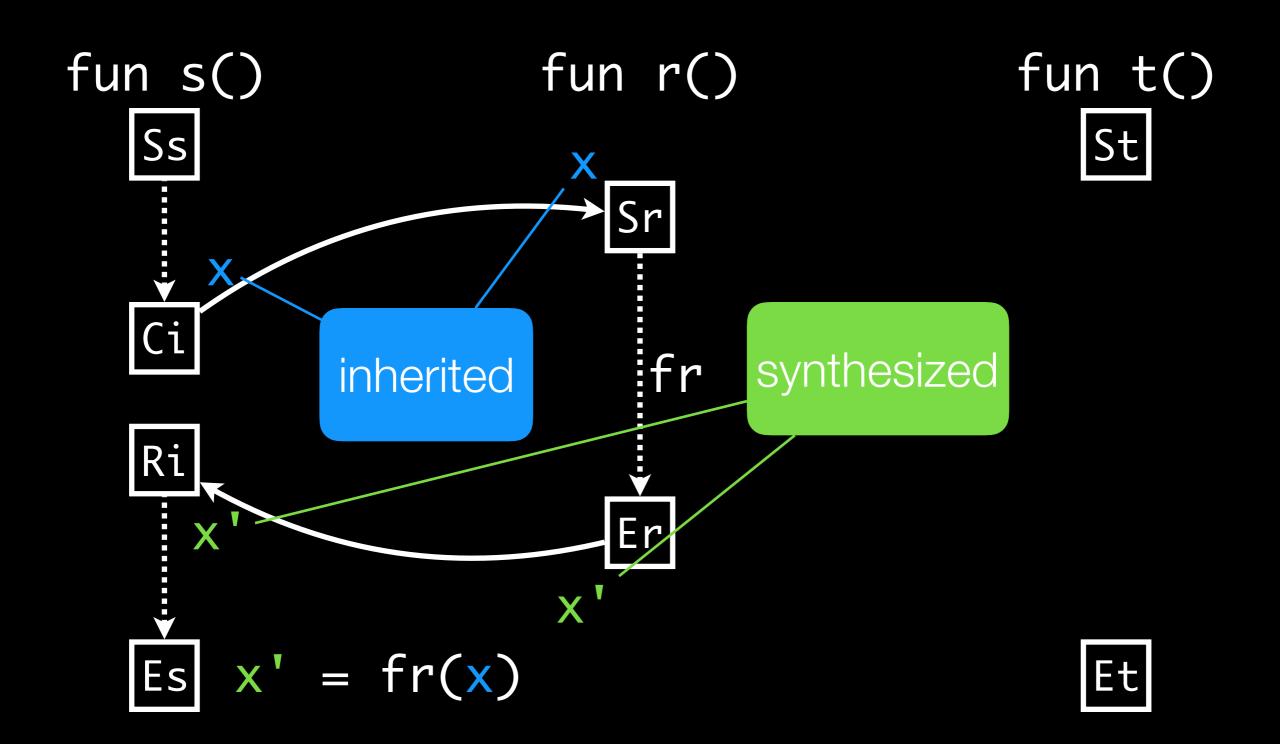
fun s() fun r() fun t()

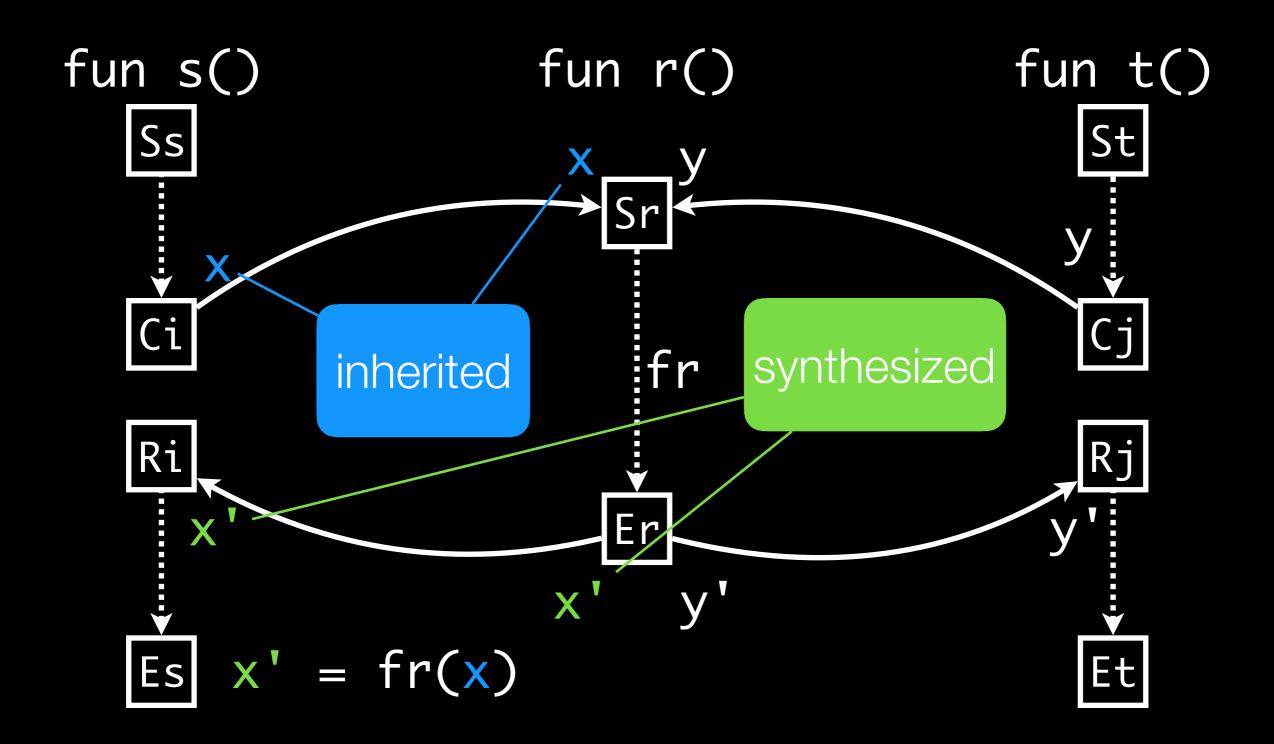
fun t() fun s() fun r() Er

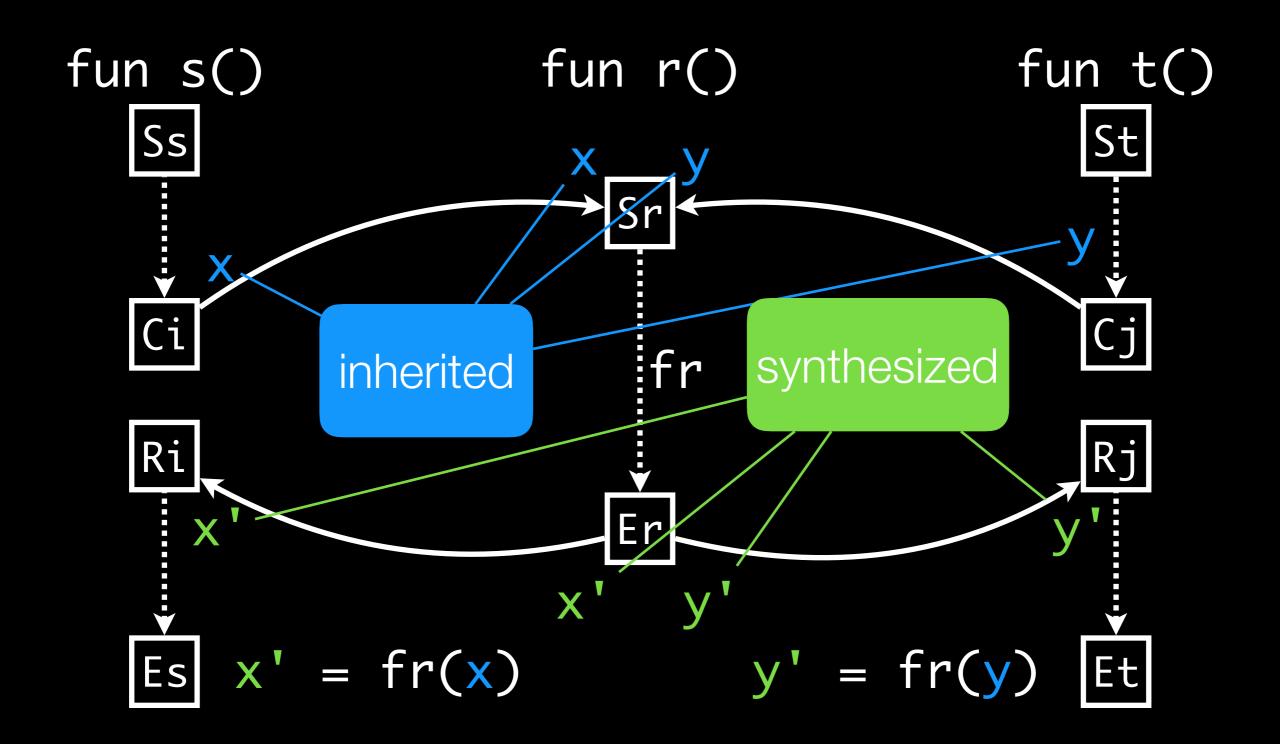












Inherited vs Synthesized Analysis Information

Inherited Analysis Information

- Answering questions about formal parameters and global variables:
 - Which variables carry constant values?
 - Which variables aliased with cac other?
 - Which locations can a point point to?

Synthesized Analysis Information

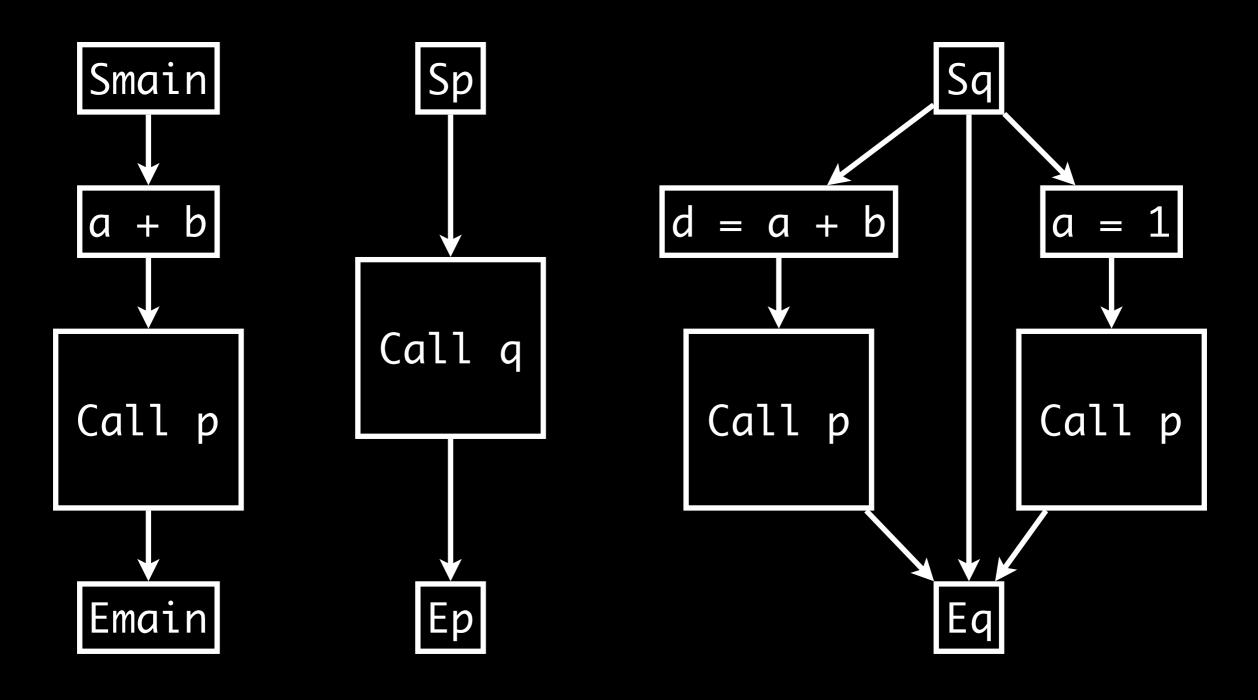
- Answering questions about sideeffects of a procedure call:
 - Which local/global/formal variables are defined in a salee?
 - Which local/global/form.

 Variables are used by a camee?

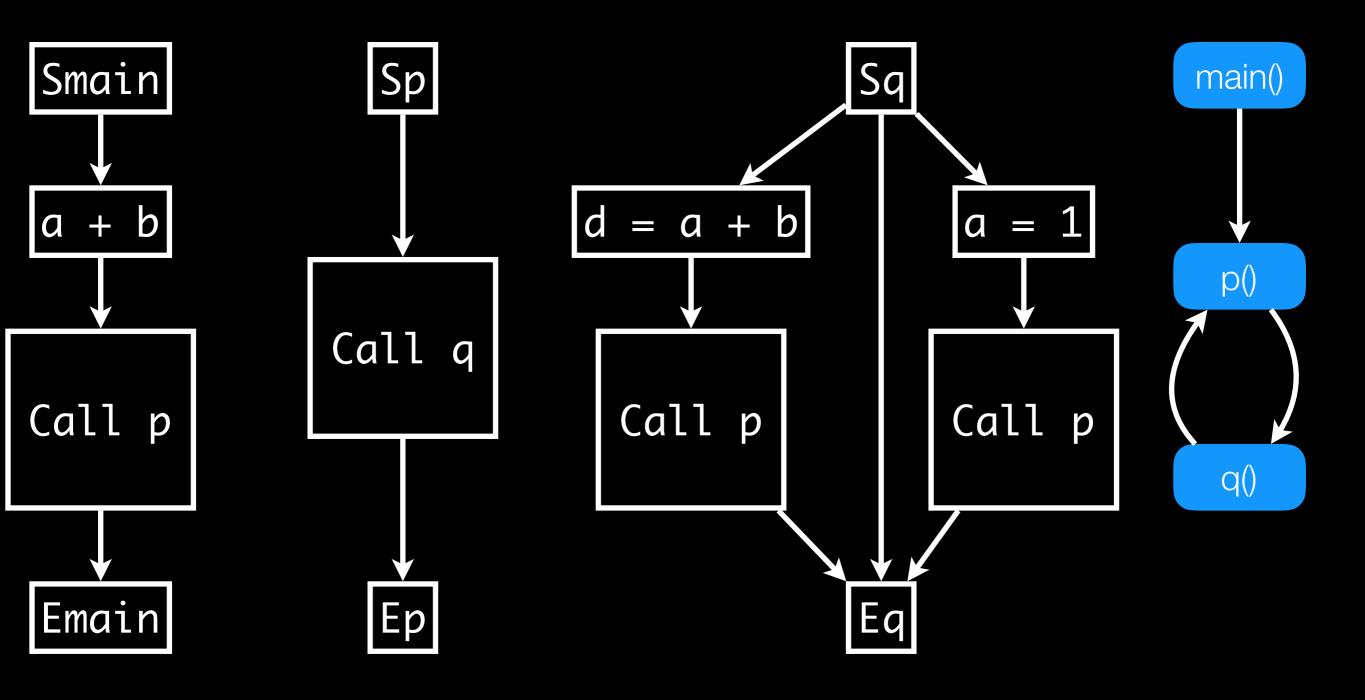
Inter-Procedural Control-Flow Graph (ICFG)

aka "program super-graph"

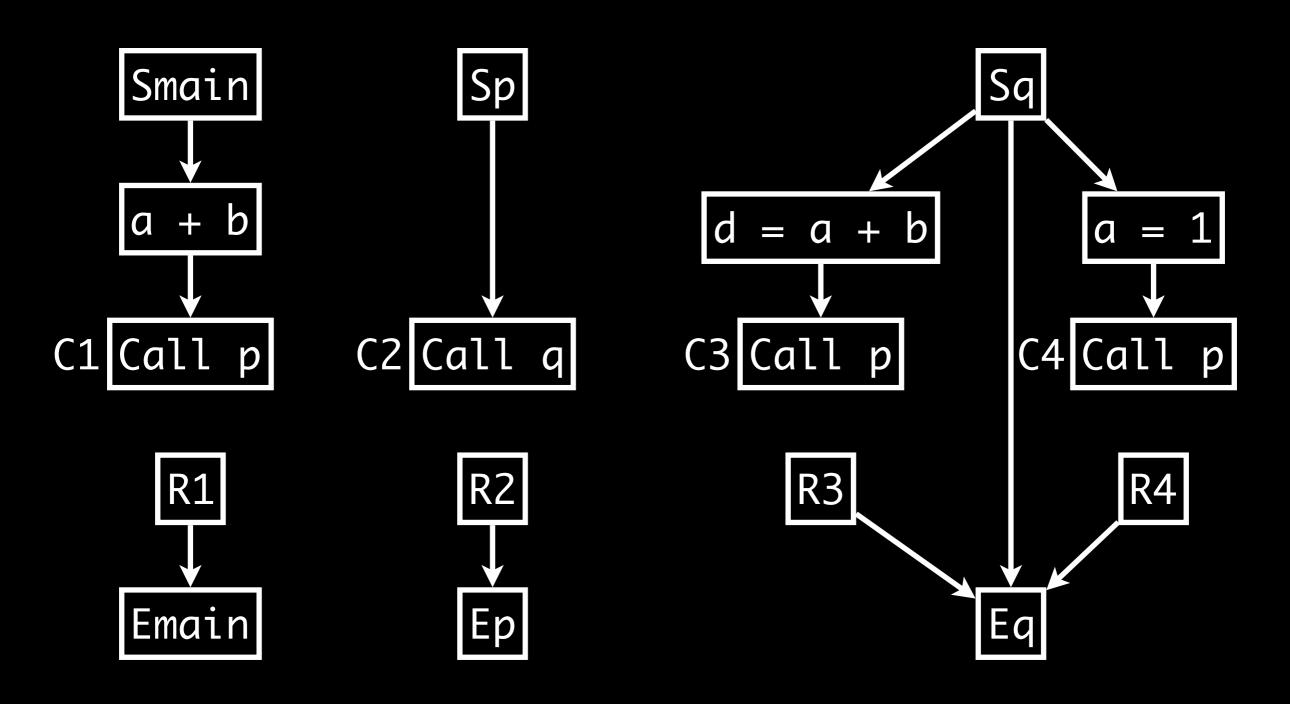
Procedure Space

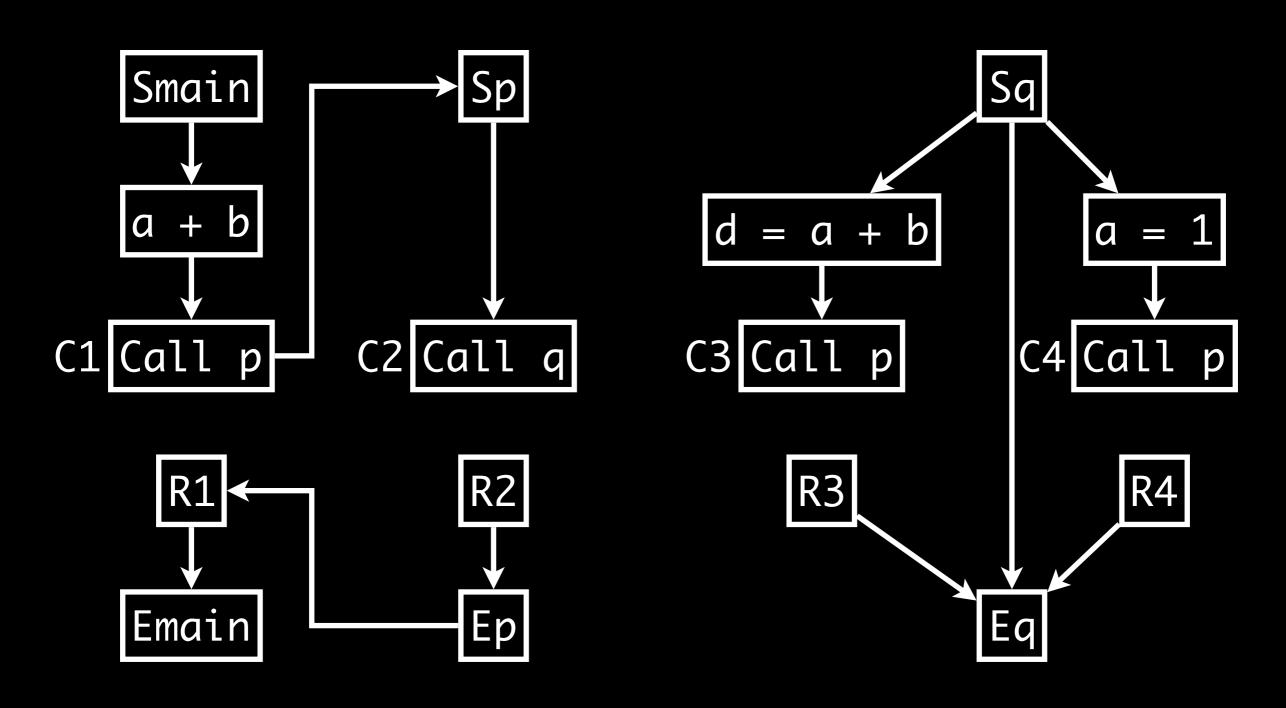


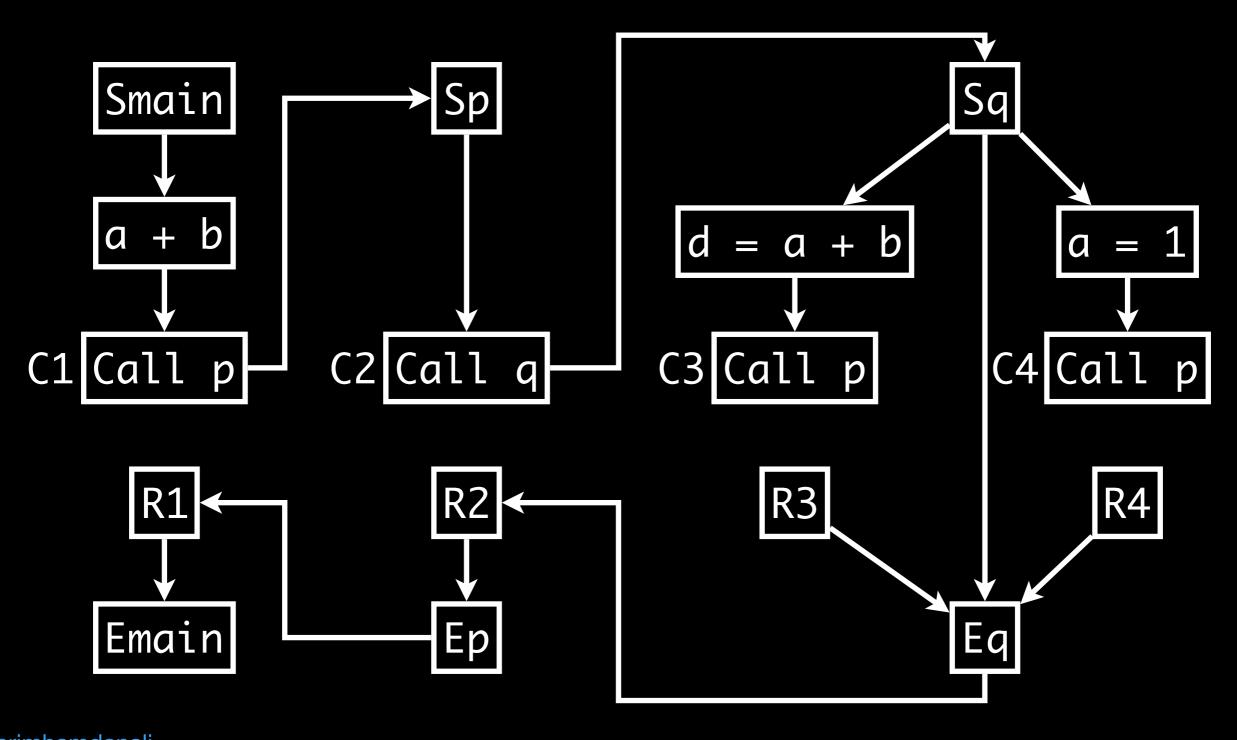
Call Graph

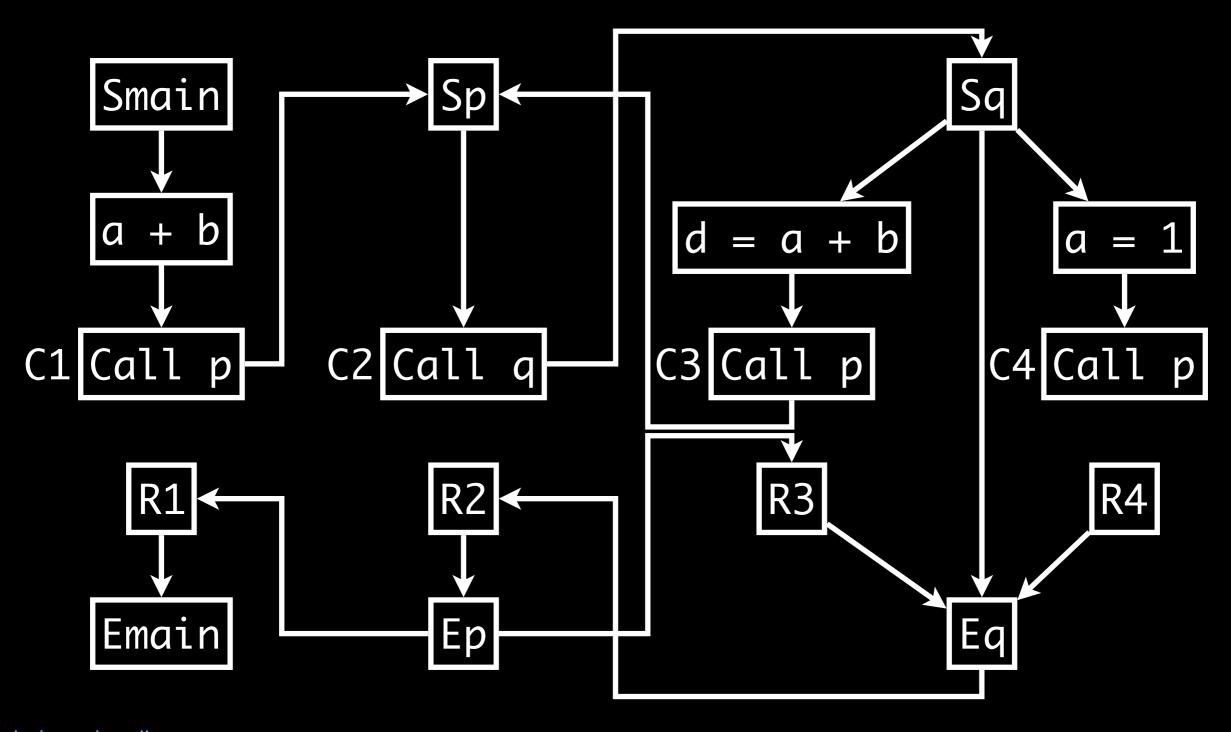


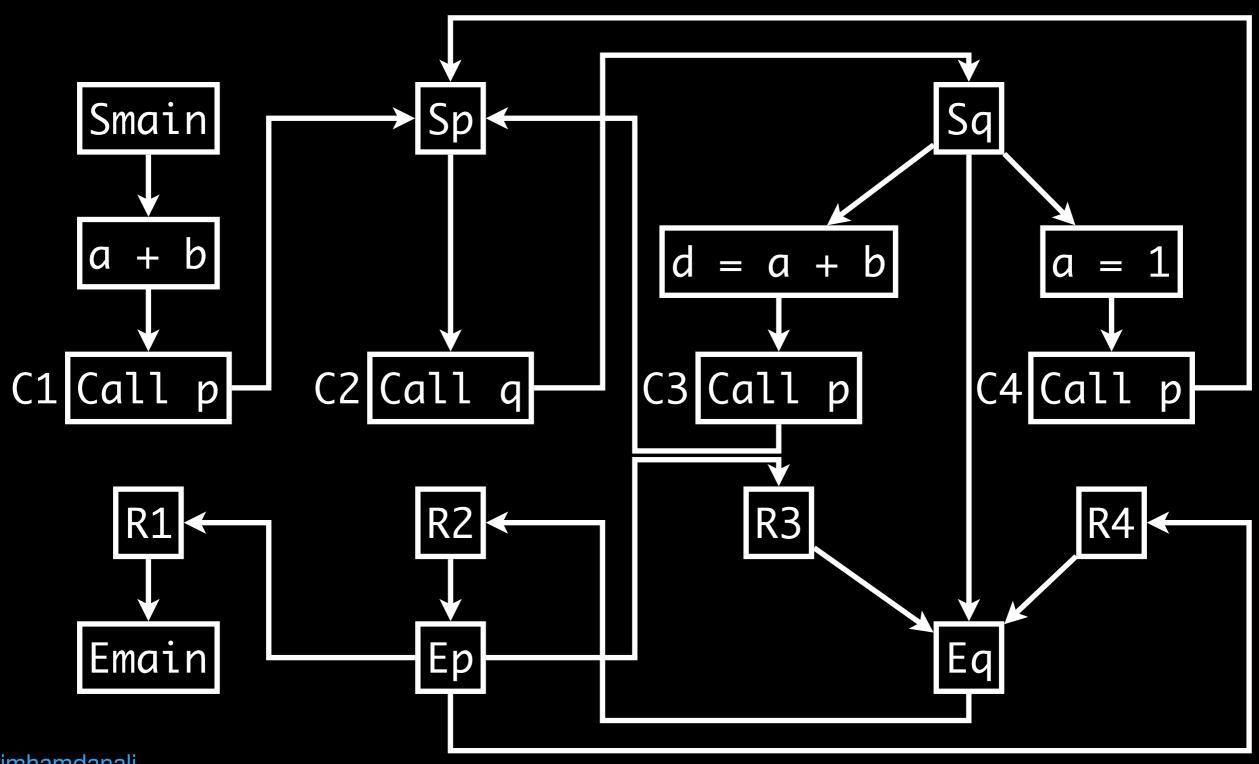
Introduce Return Sites

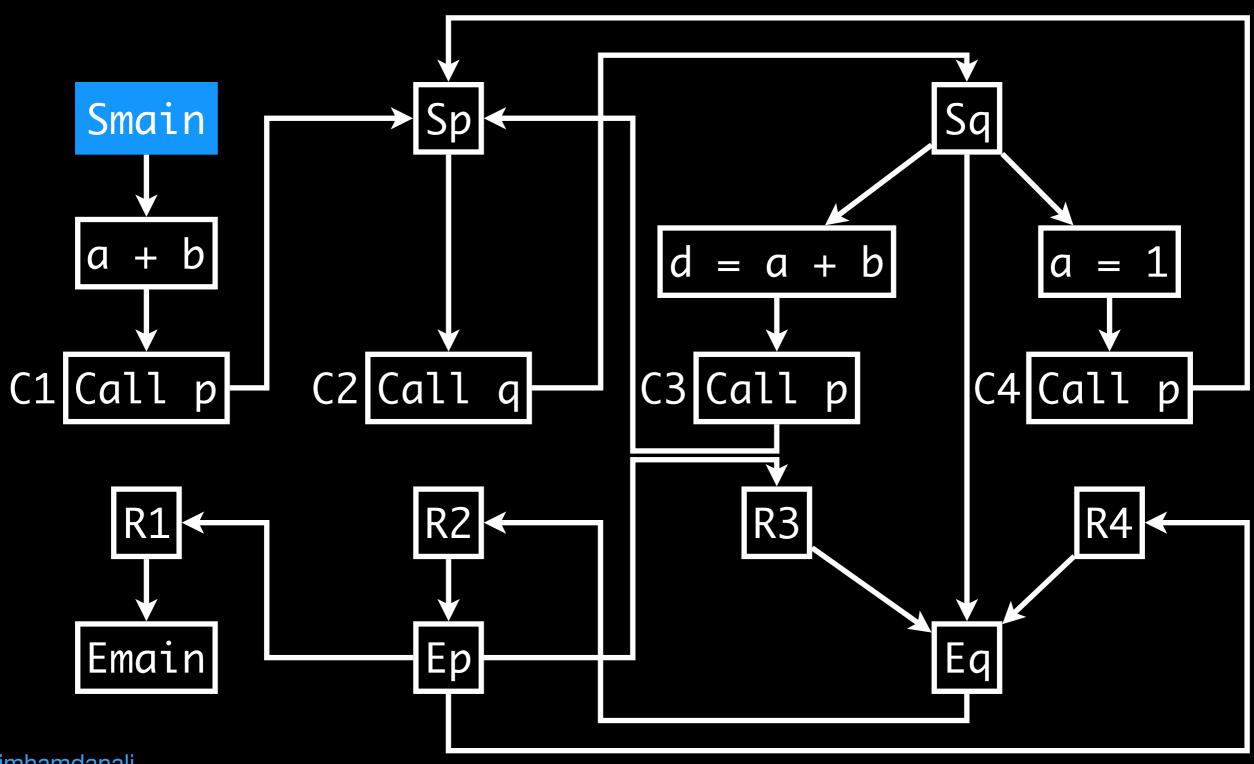


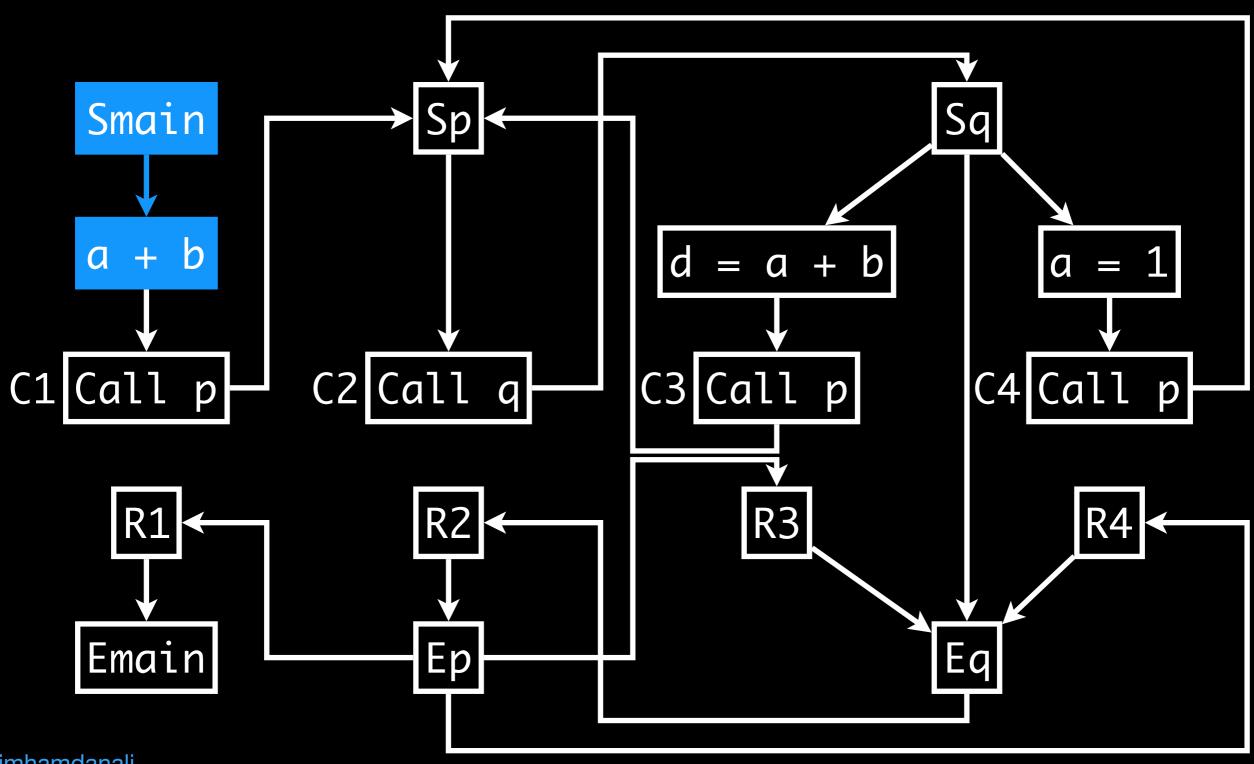


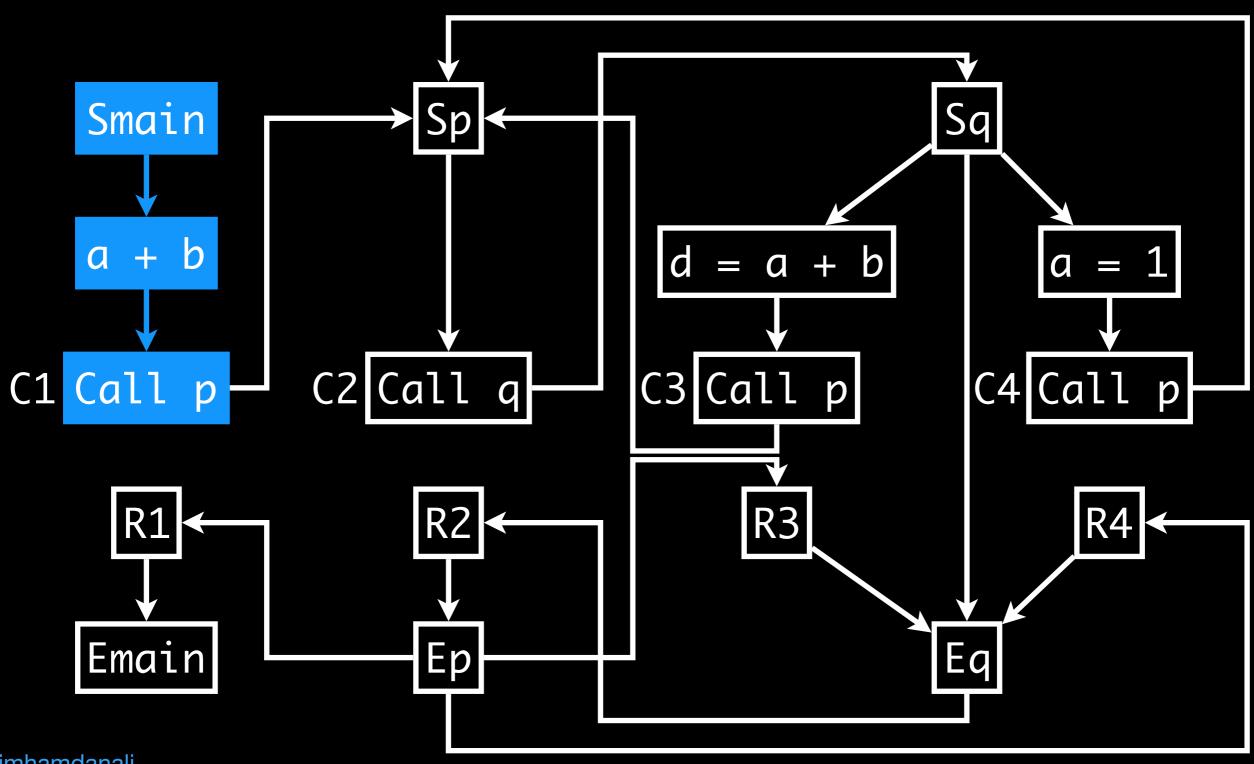


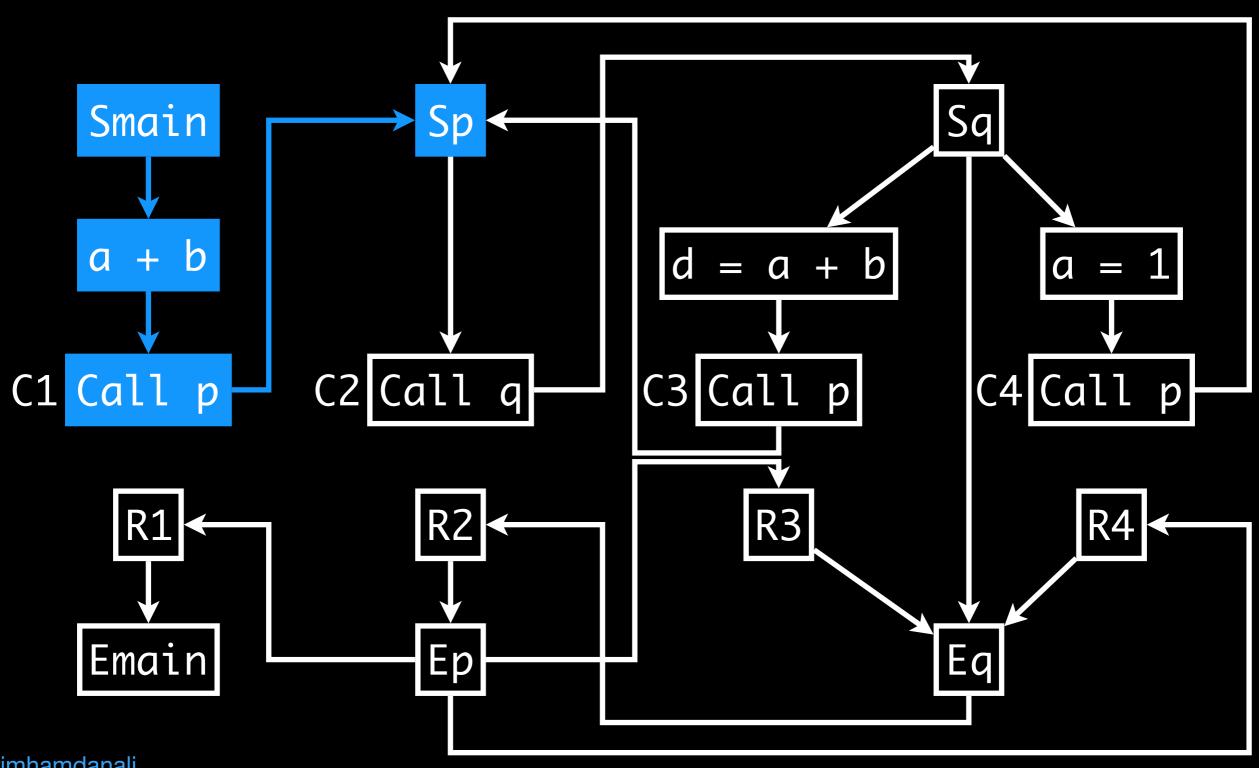


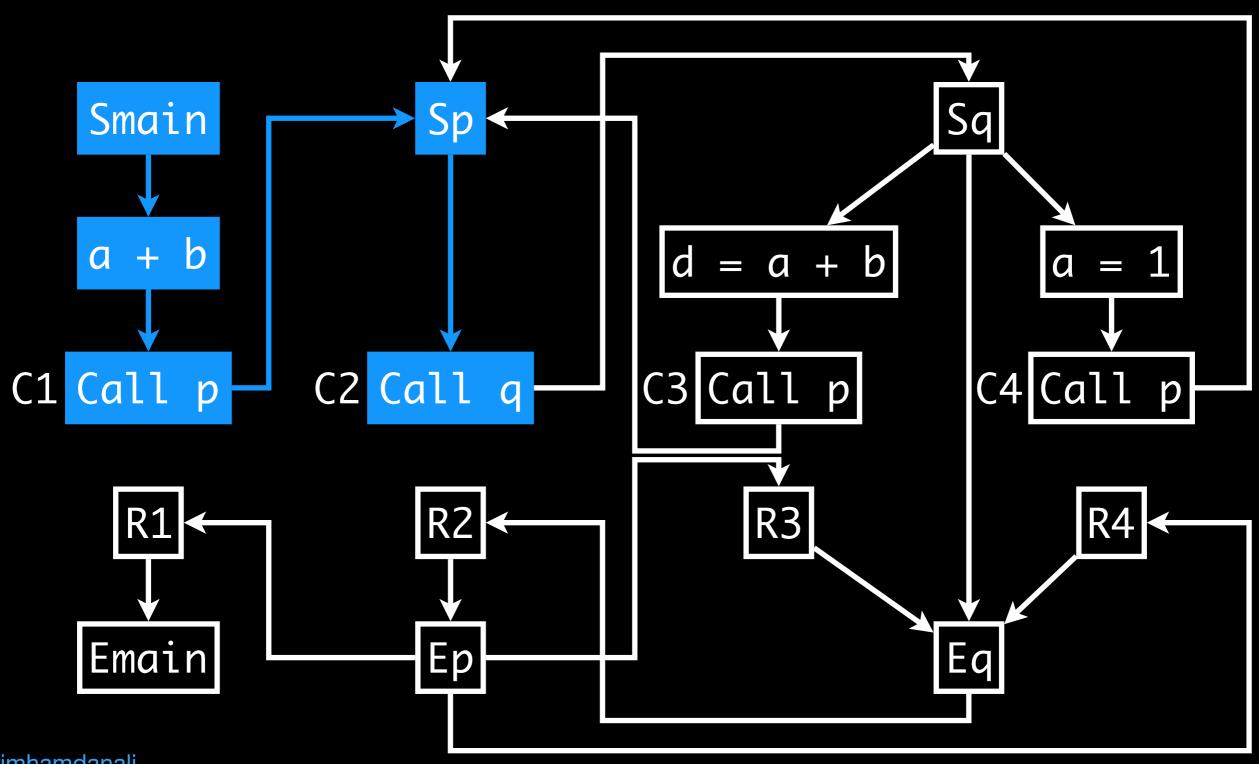


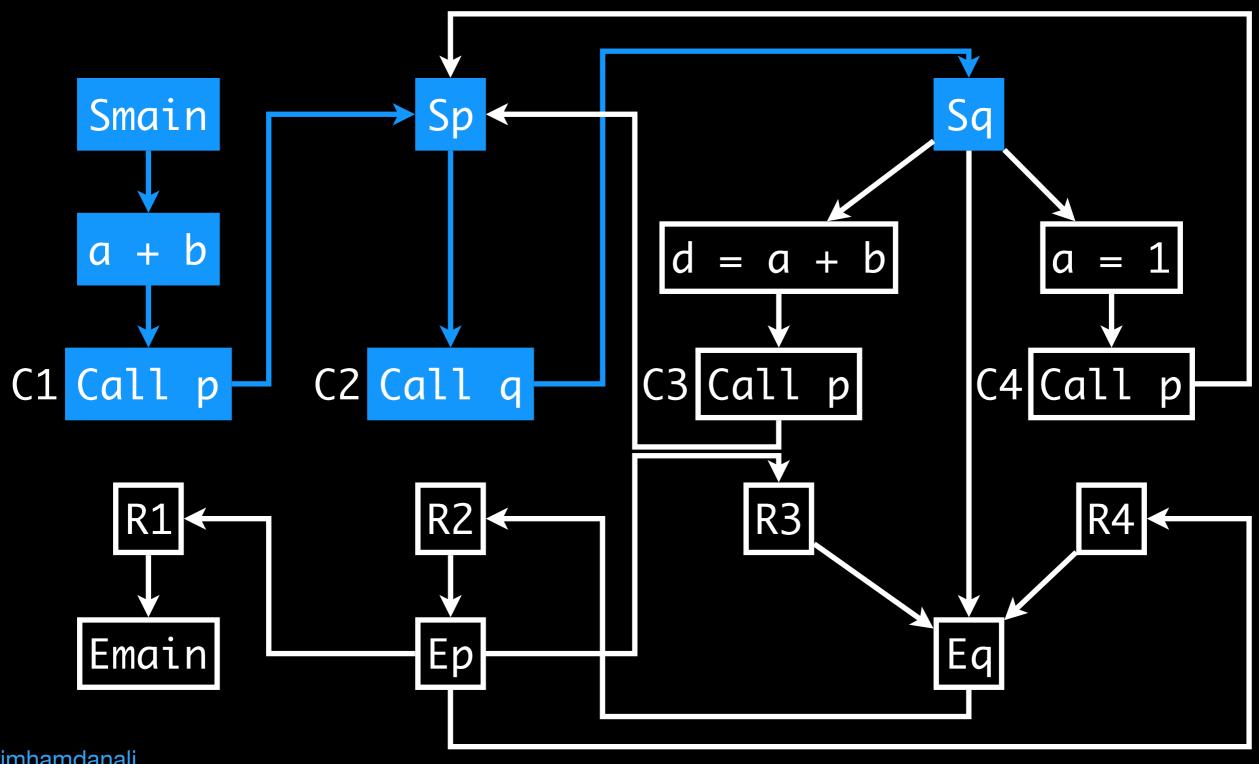


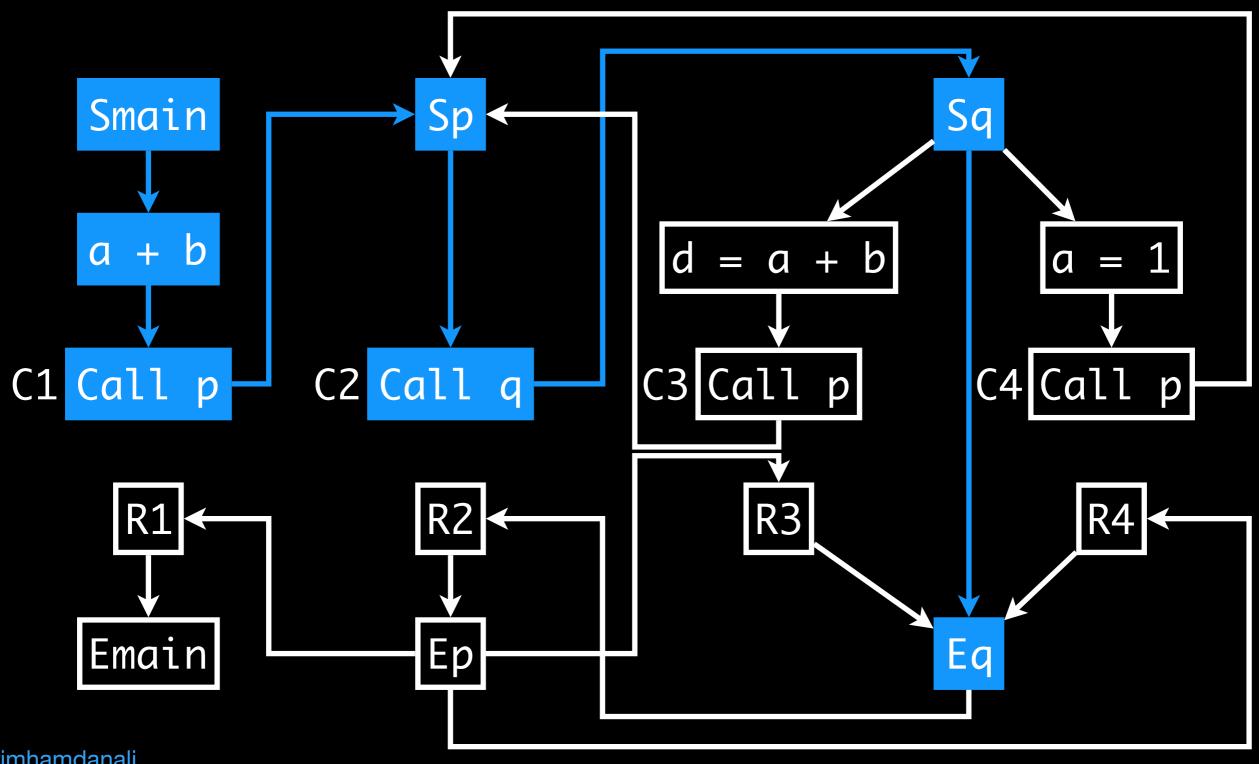


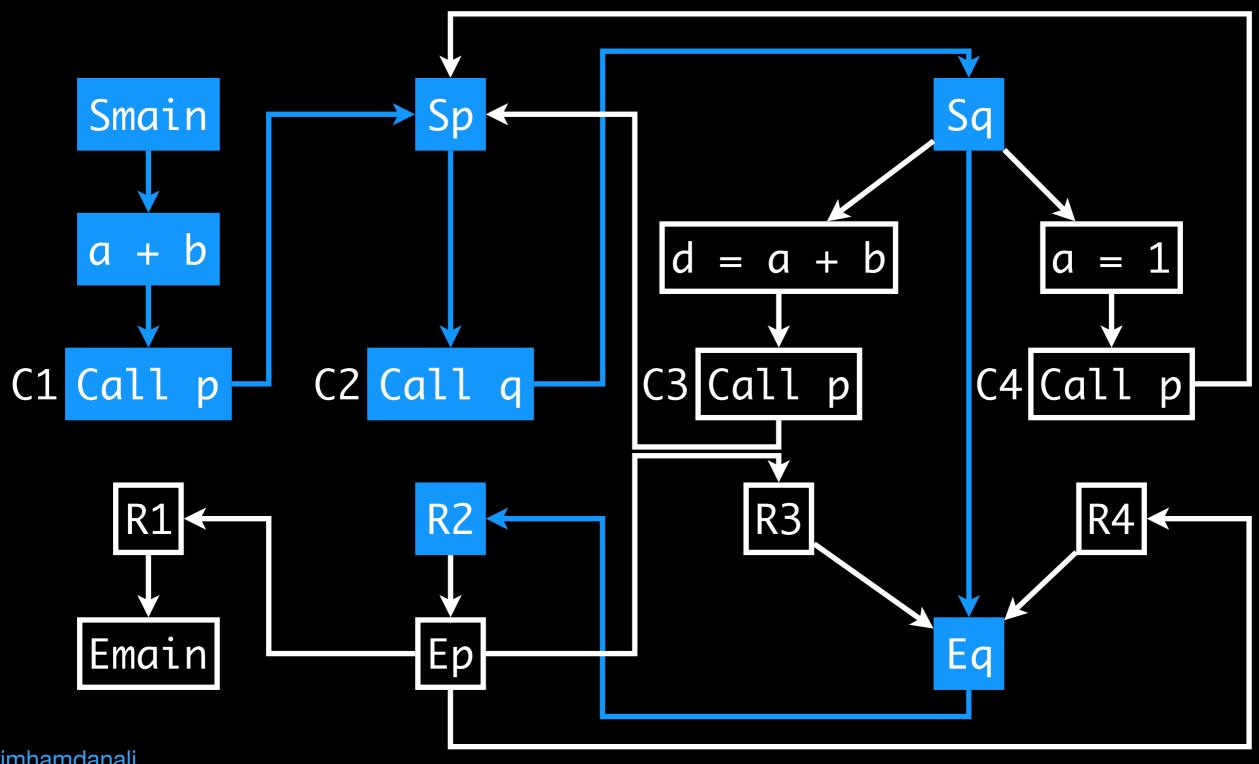


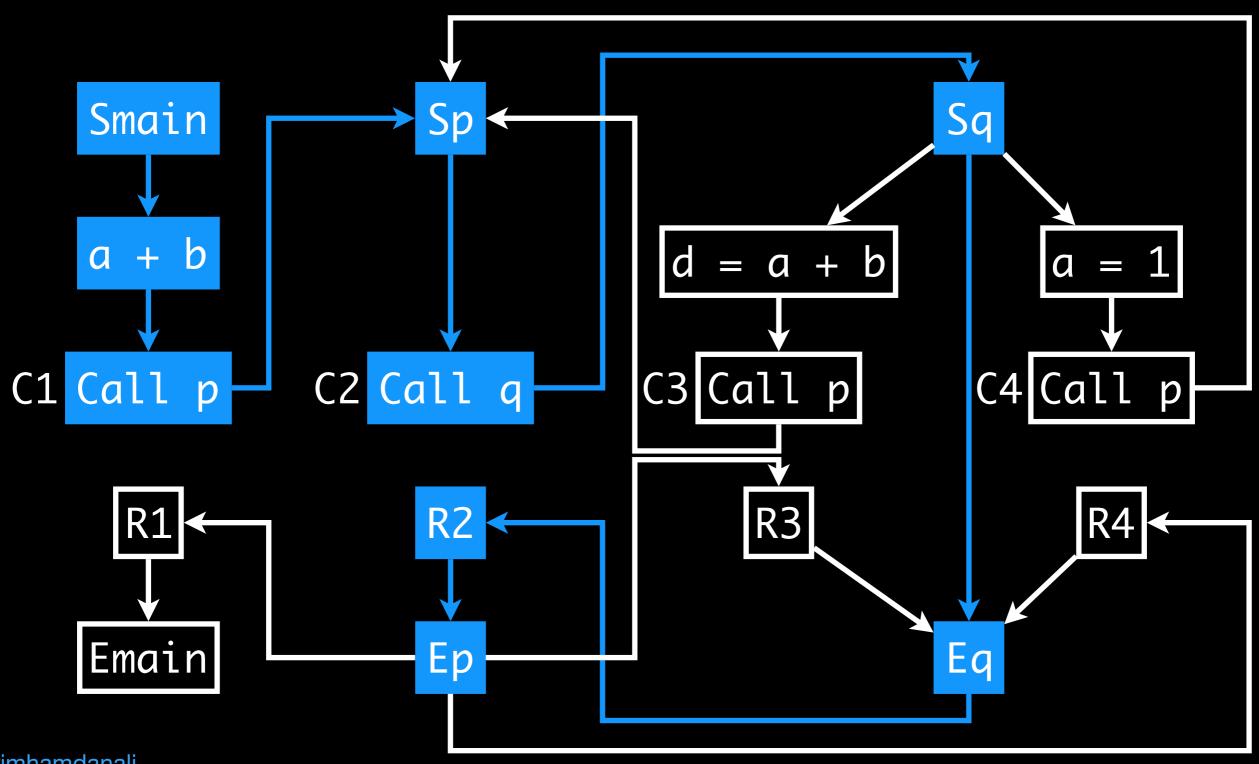


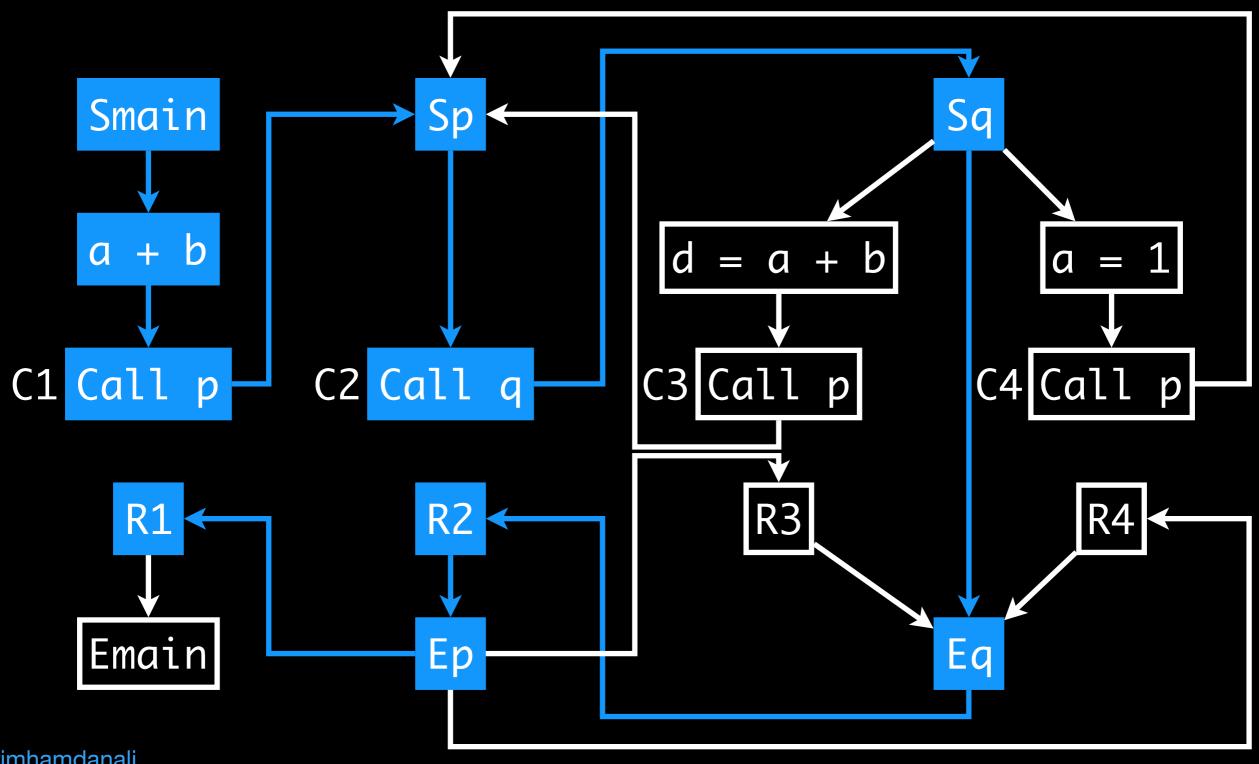


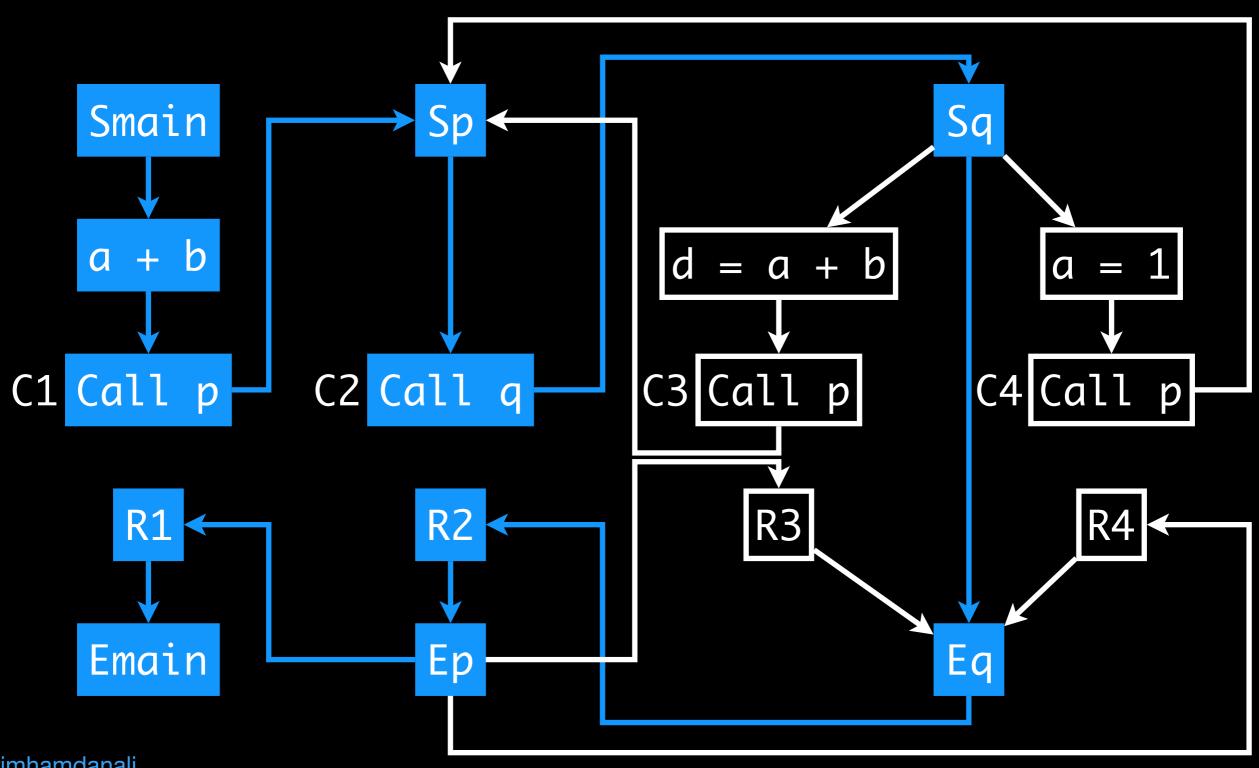


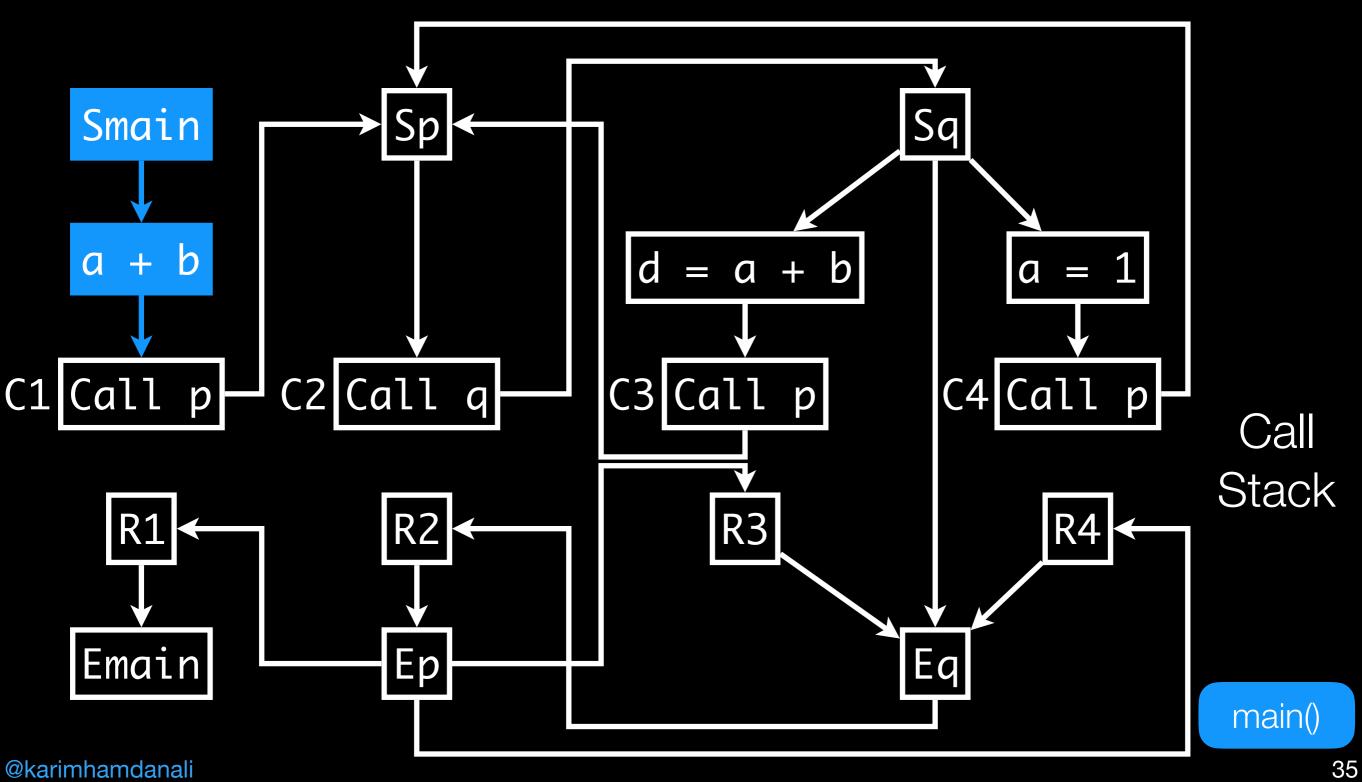


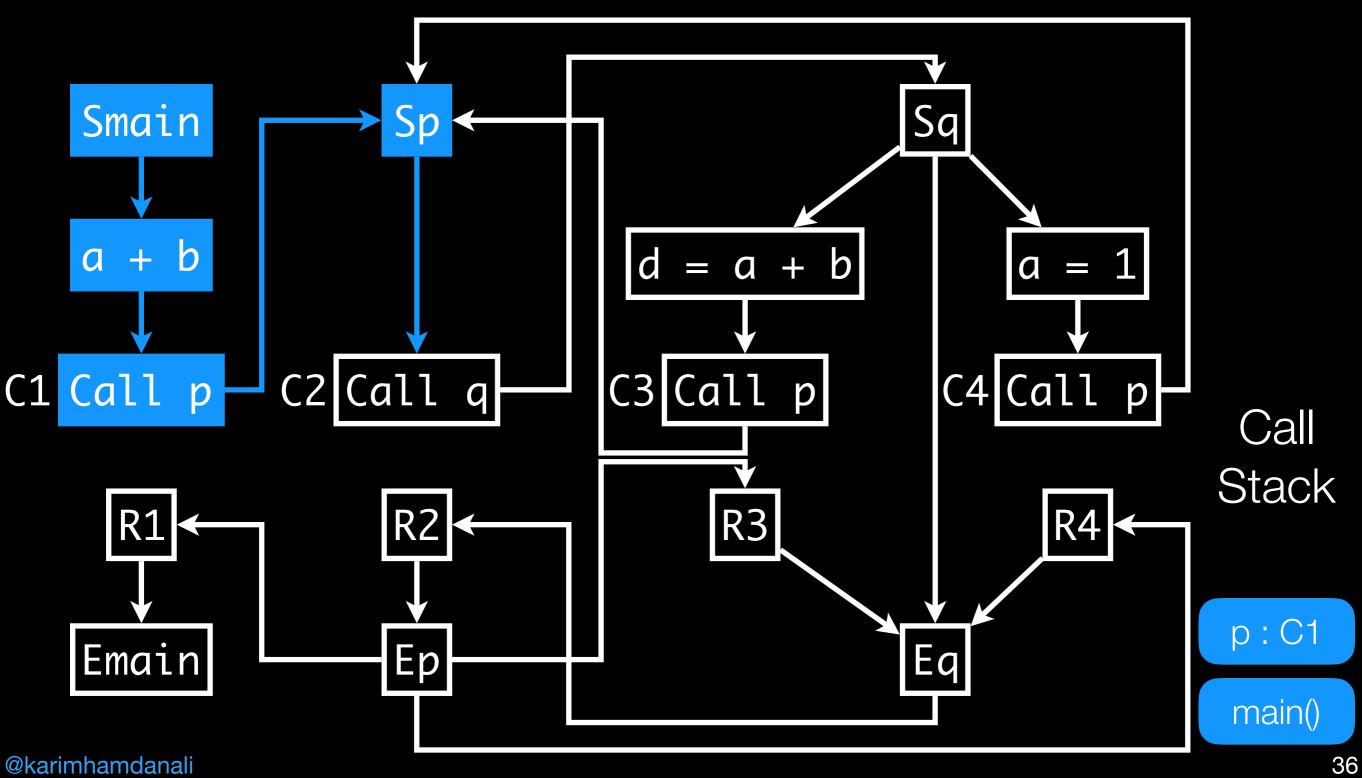


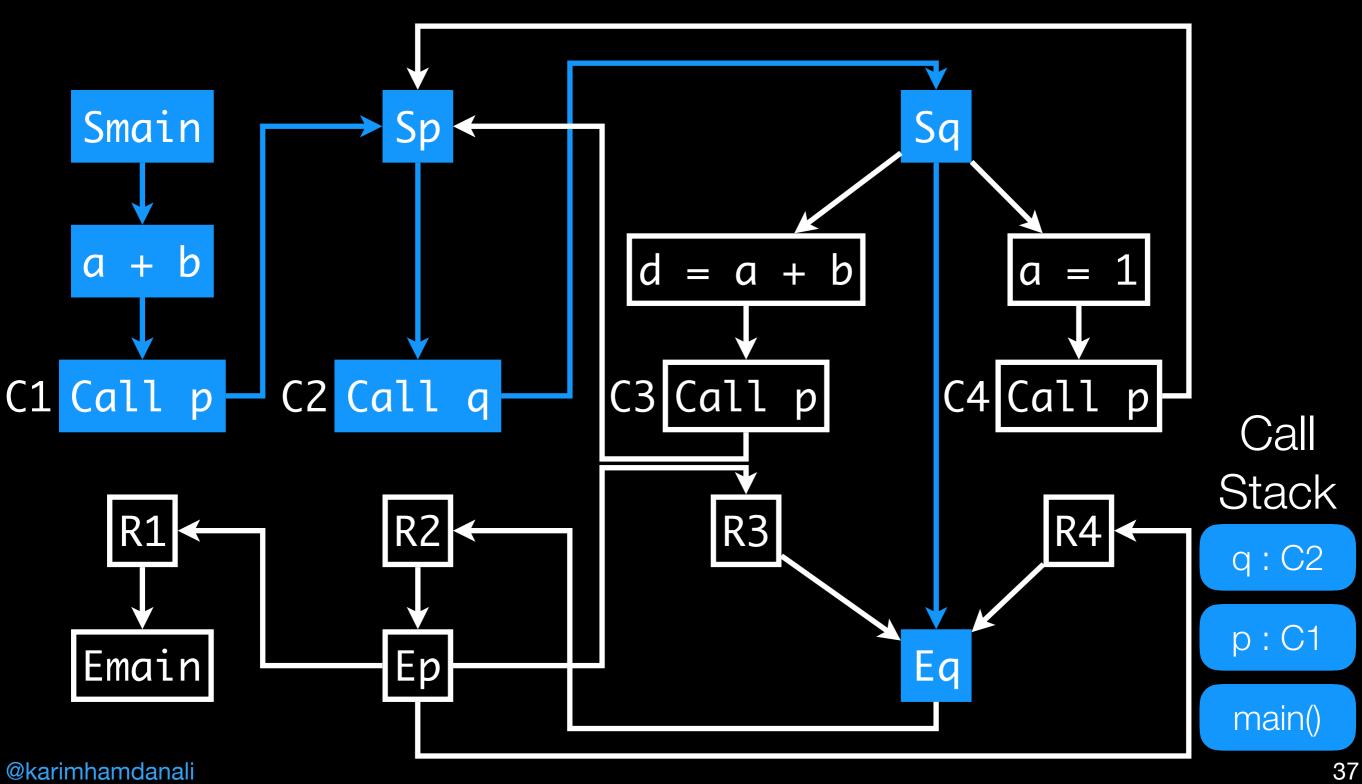


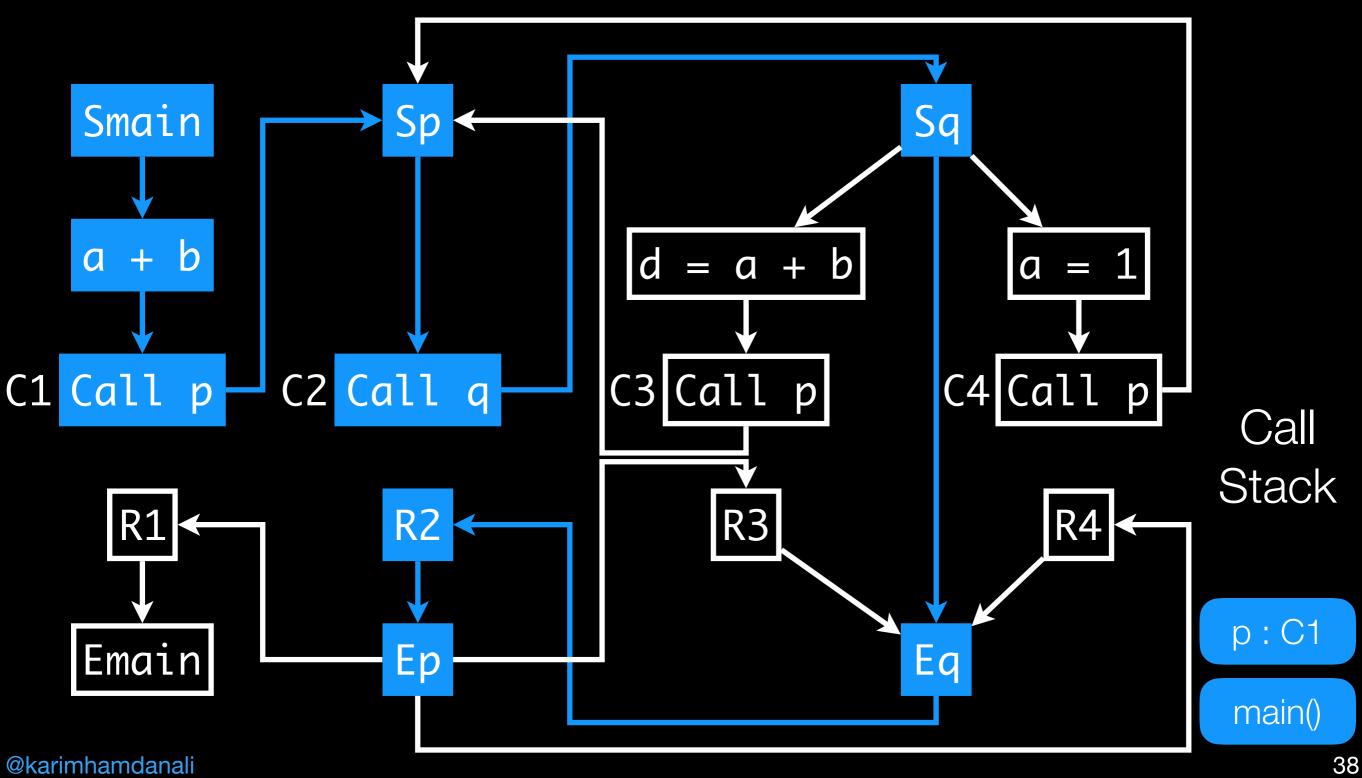


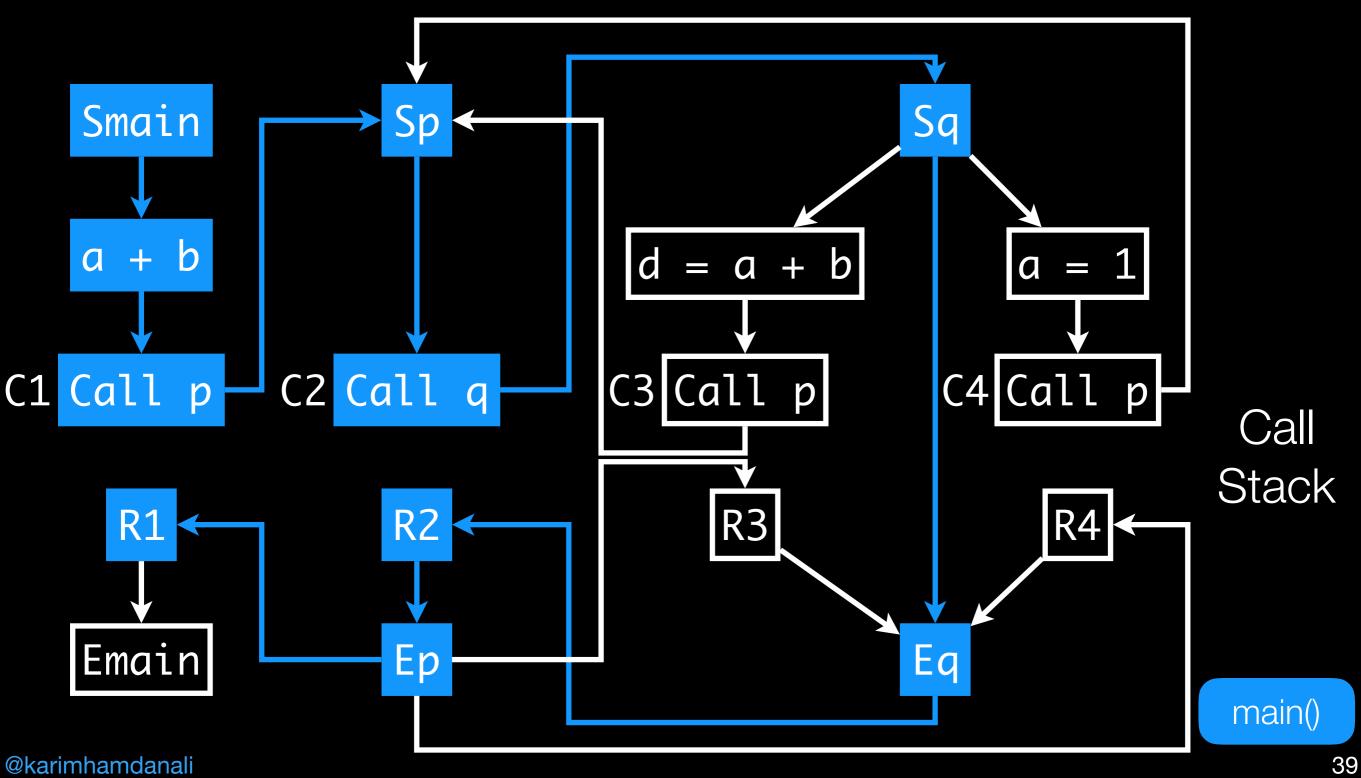


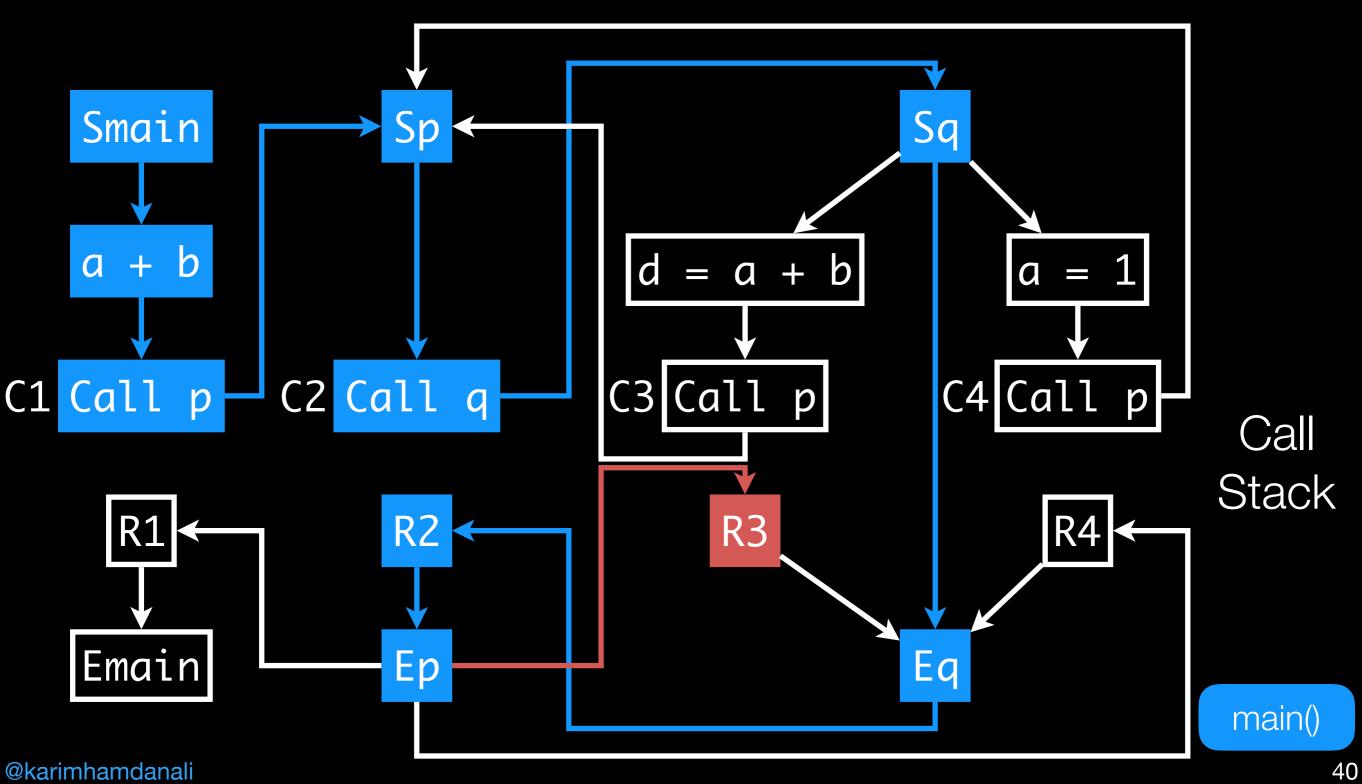


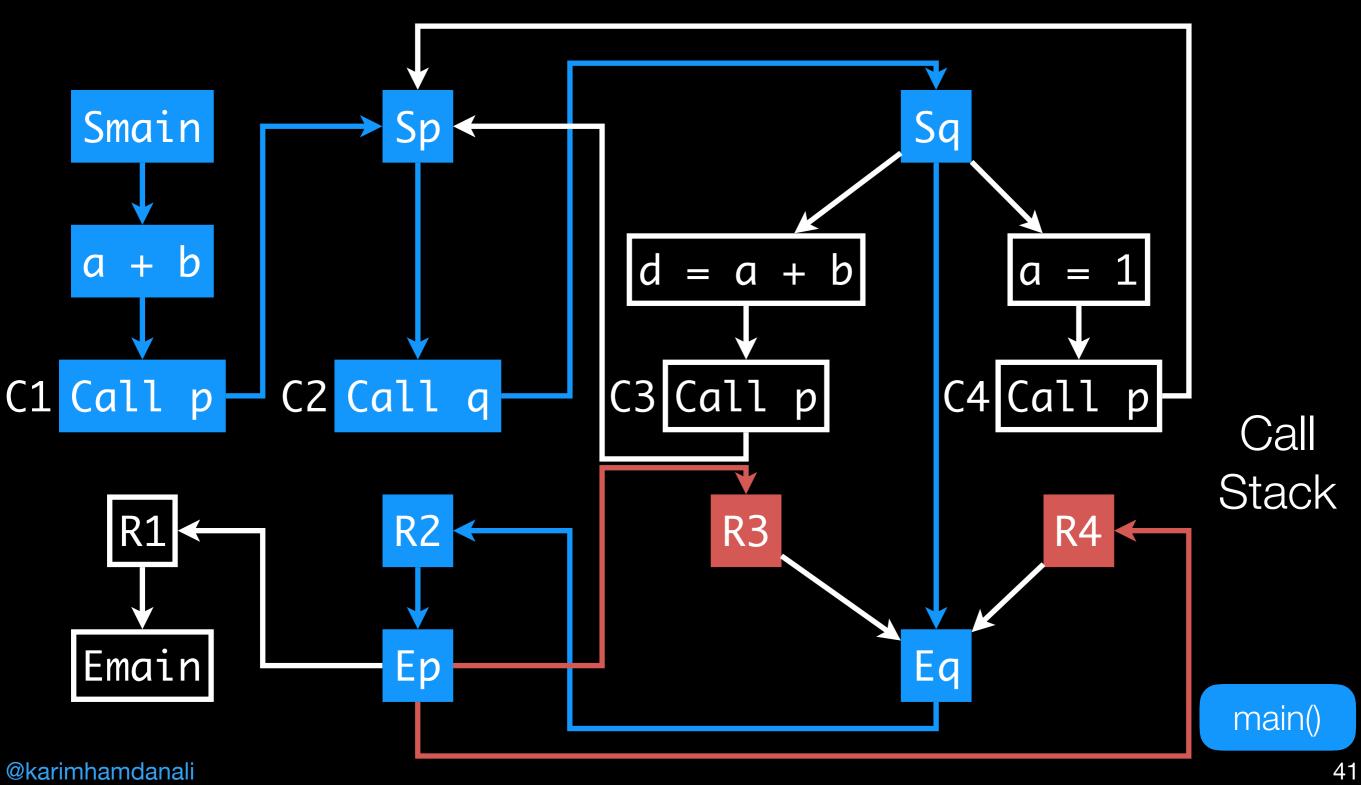


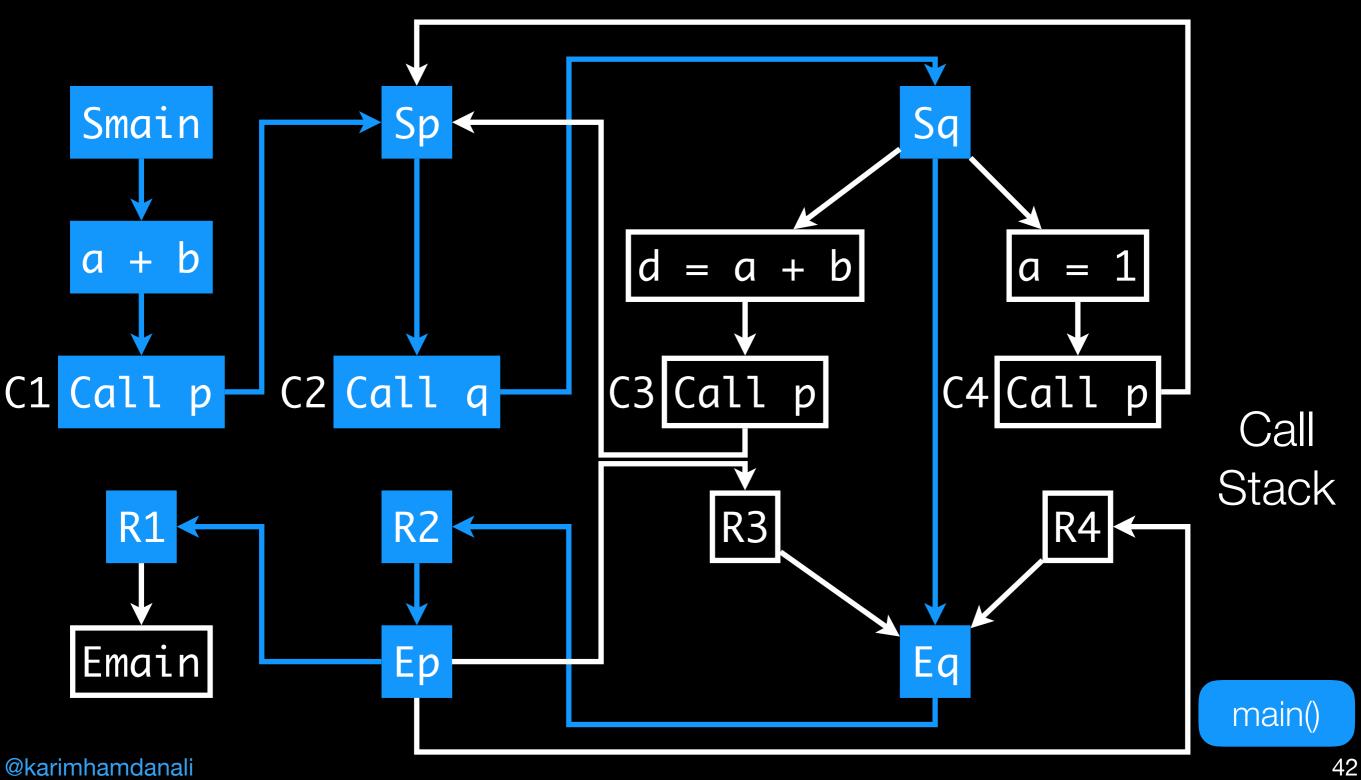


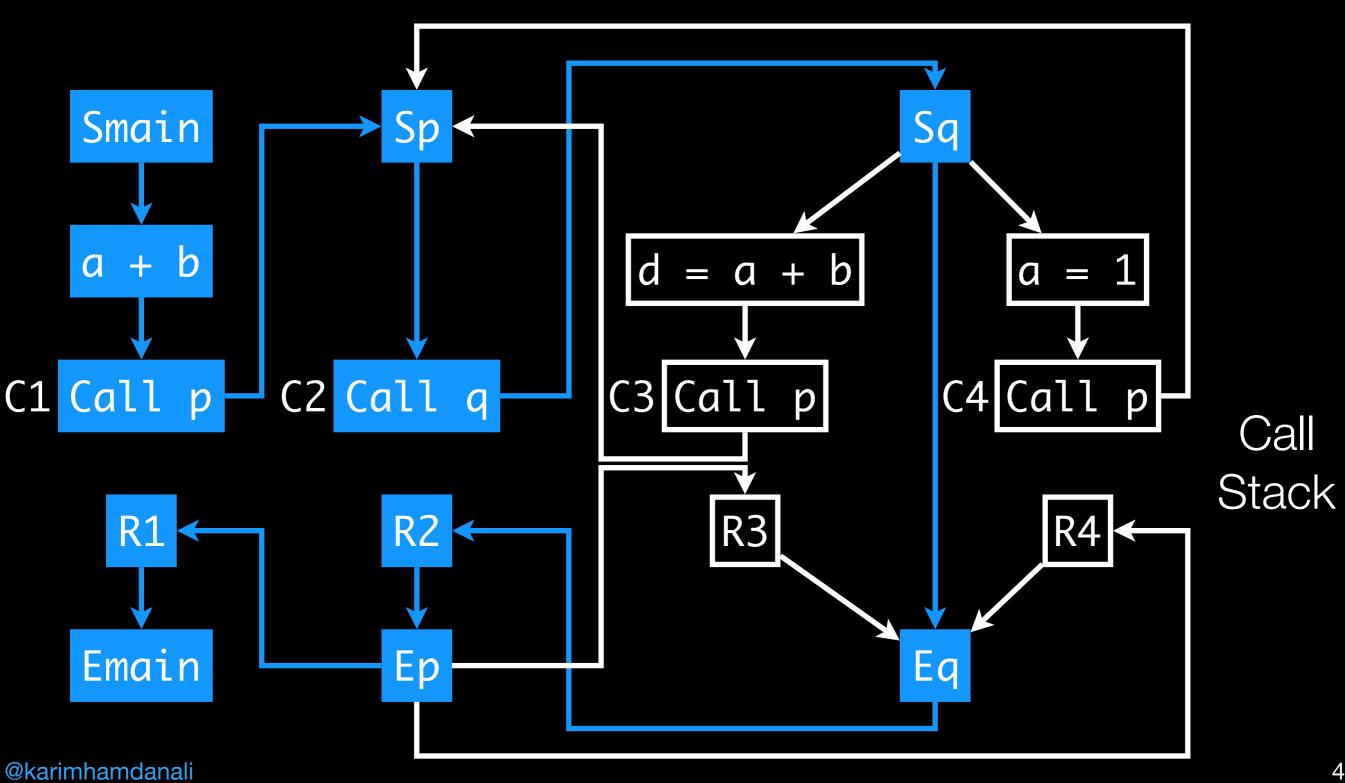






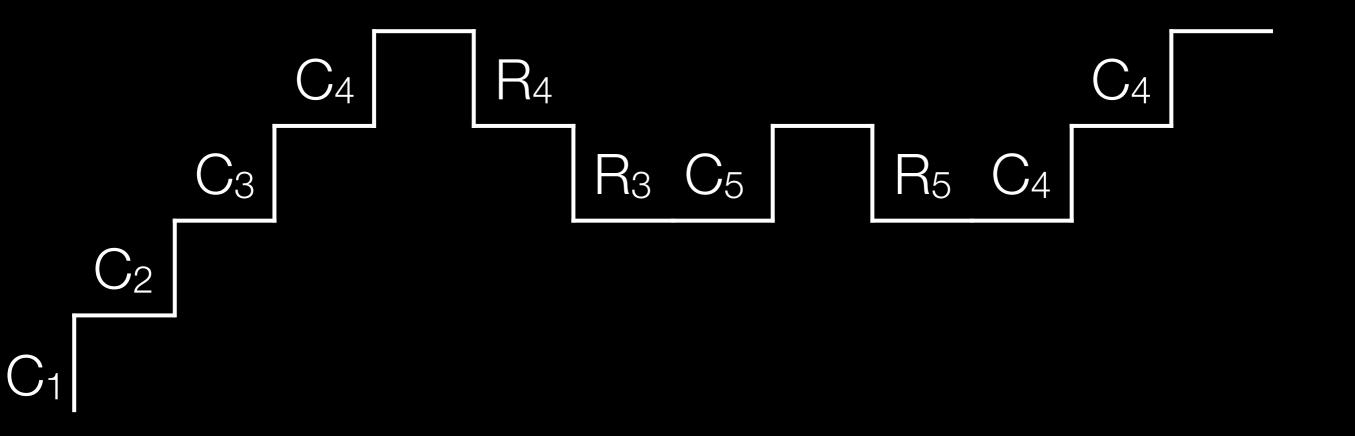




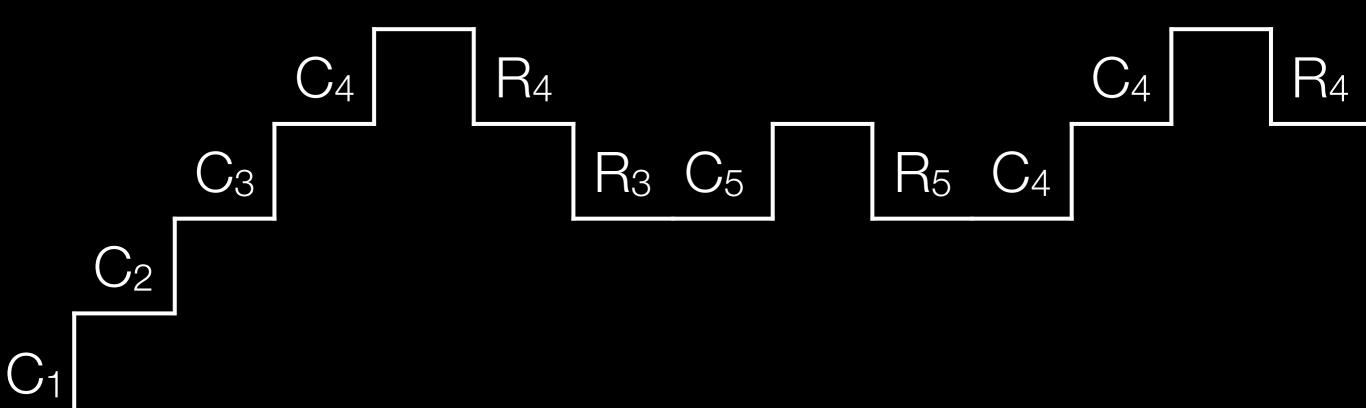


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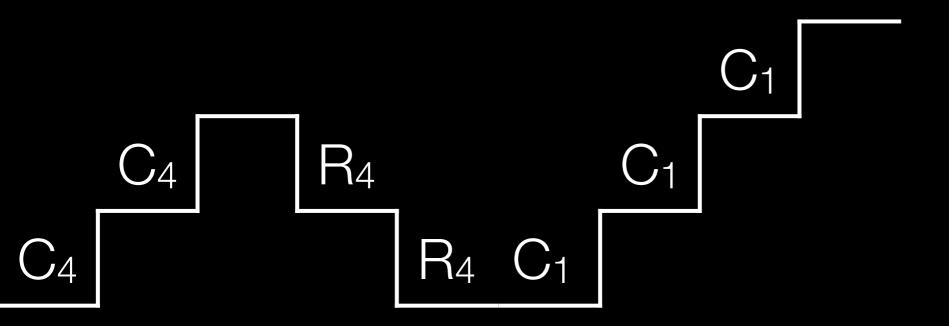
Recognizing Invalid Paths Staircase of Calls & Returns



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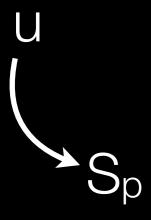
Recognizing Invalid Paths Staircase of Calls & Returns

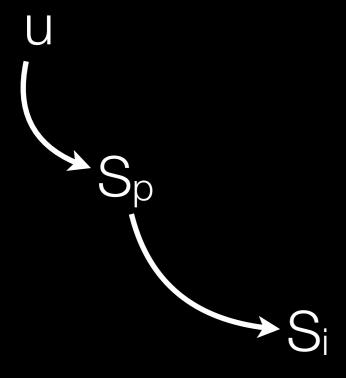


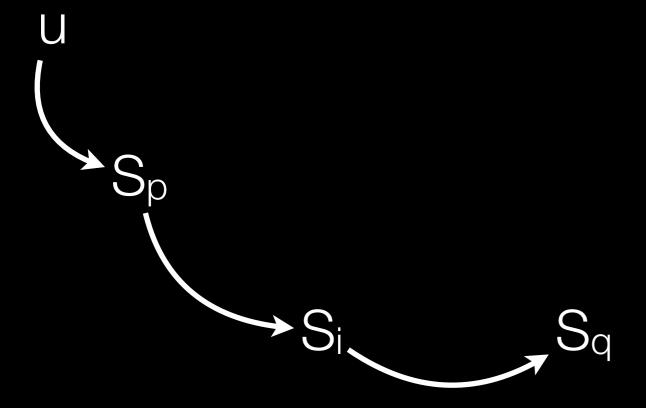
Every descending step must match a corresponding ascending step

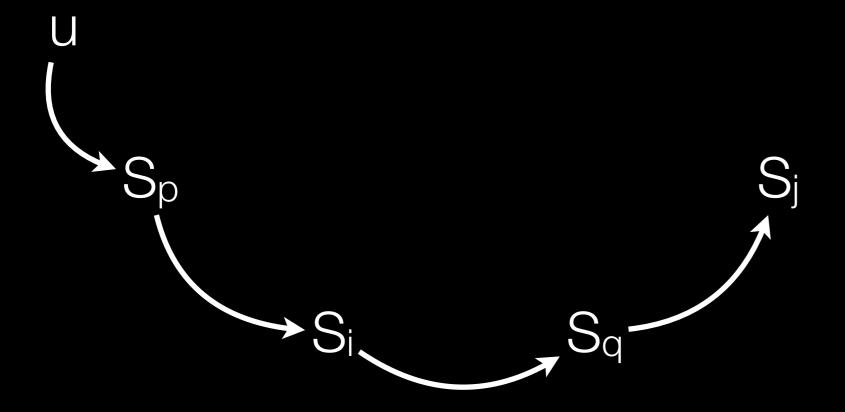
2 problems with that!

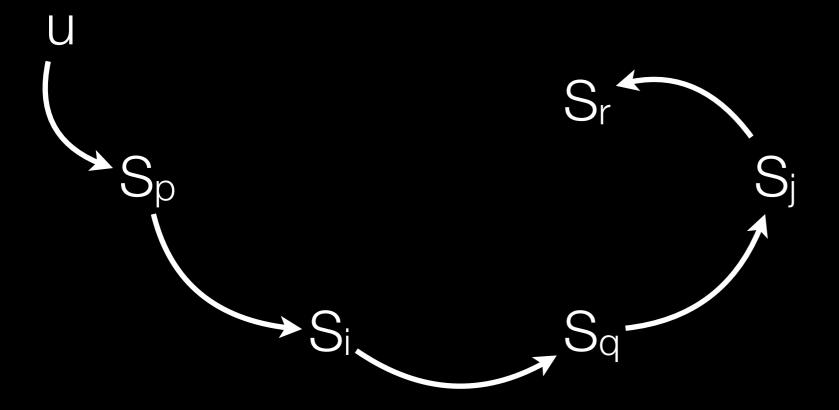
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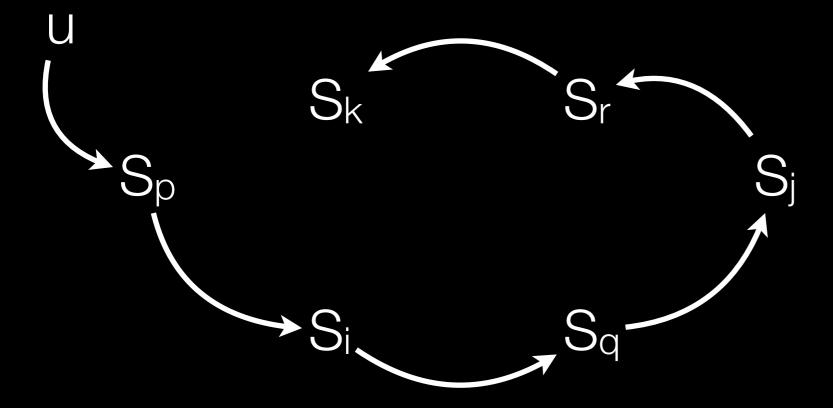


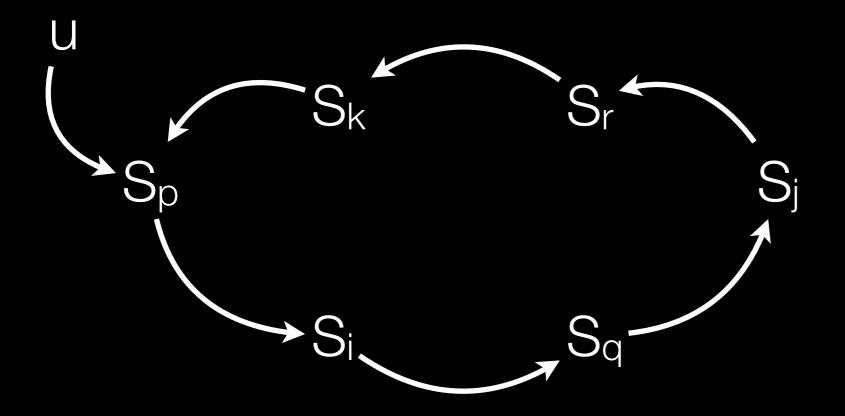




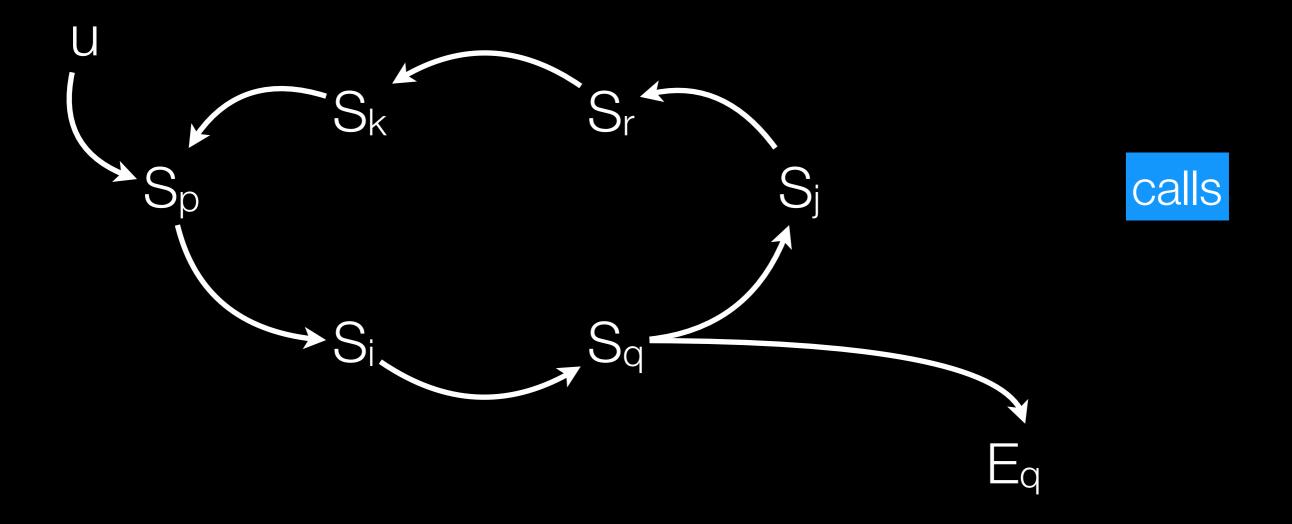


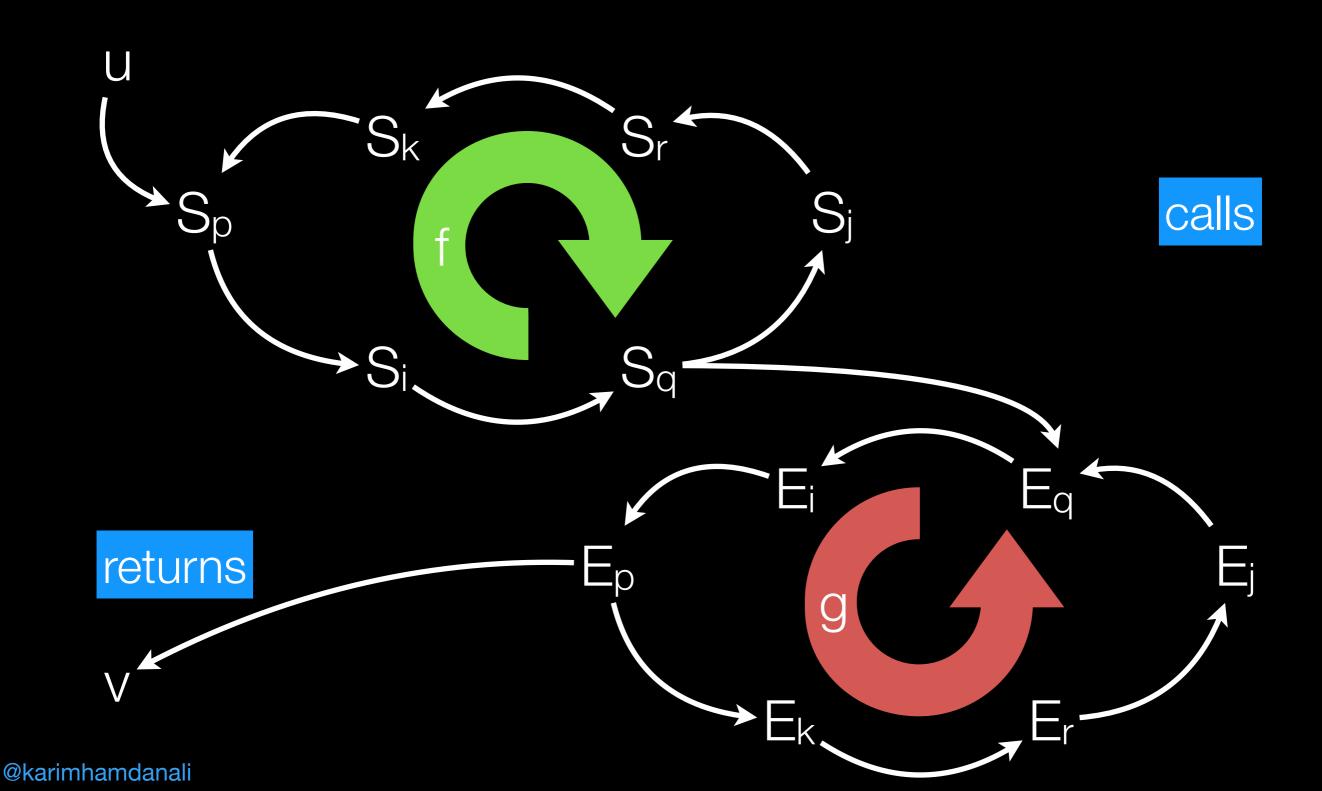


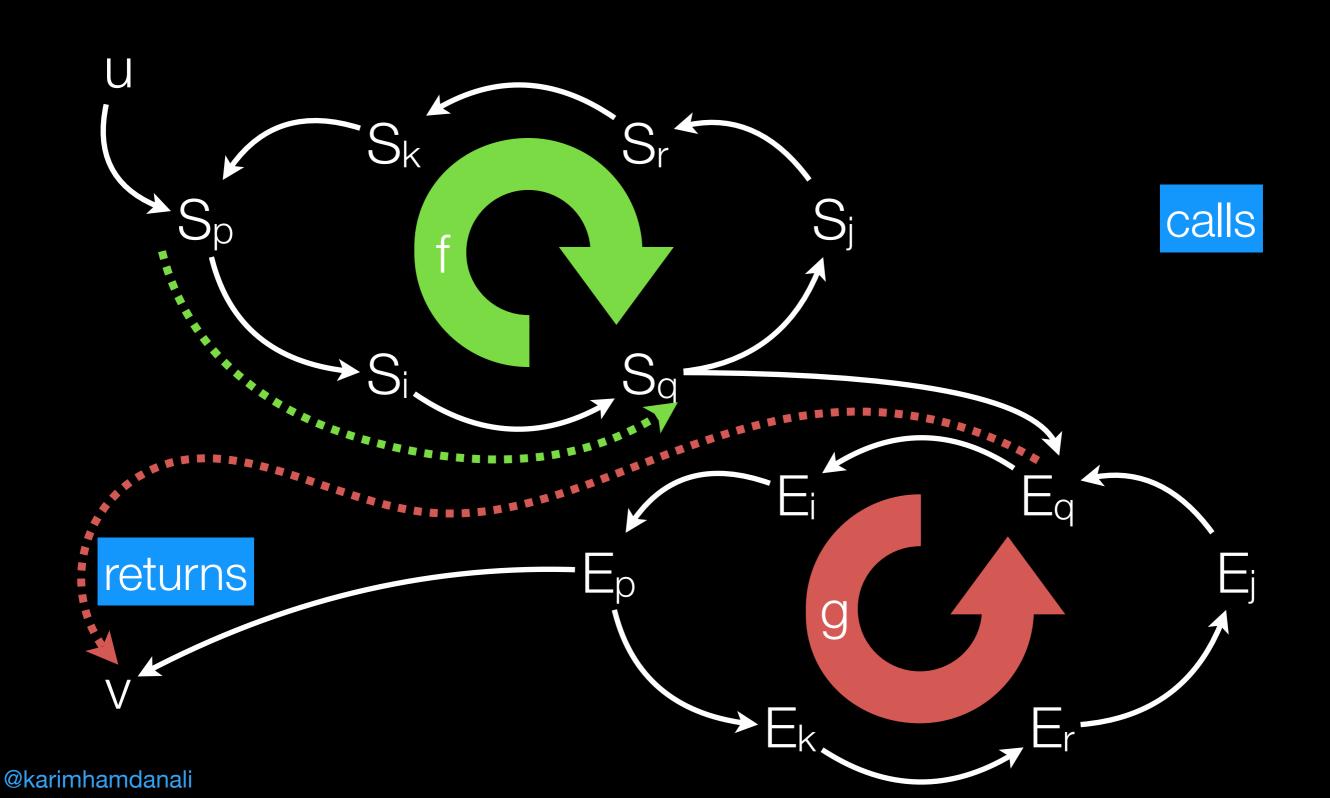












Problem # 2: Demand-Driven Analysis

```
main() {
    s = secret();
    foo(s);
    t = "123";
    foo(t);
}
foo(v) { leak(v); }

assume we search
    from foo(v)
    backwards to find
    possible inputs
    Rt Cs Rs ...
```

here: "unbalanced return" without a call must return to all possible callers

Solution: Context-Sensitive Analysis

Context-Sensitive Analysis

 Analyze the same method, depending on the context of the current call to that method

Context-Sensitive Analysis

Considerations:

 How to distinguish different contexts?

 Which contexts can be merged?

Types of Context

- A call string that encodes the methods/call sites on the current call stack
- A value context that uses the input domain values as context
- An object context that uses the currently executing object as context
- and more...

Important Language Features

Recursion

Must bound computation and contexts

 Often uses flow-insensitive analysis to over-approximate

Parameters/Return Values

- Must map actuals to formals and vice versa
- Don't propagate too much info:
 - at a call: propagate only the facts relevant to that callee
 - at a return: propagate only the facts relevant to the caller
- Question: what to do with static fields?

```
main(){
    x = source();
    y = x;
    z = foo(x);
}

foo(a) {
    b = a;
    return 0;
}
```

Aliasing

aliases might be created by callers and callees

```
main(){
    a.f.g = source();
    foo(a,b);
    leak(b.f.g);
}

foo(x,y) {
    y.f=x.f;
}
```

Virtual Dispatch

 Multiple possible call targets per call site

- Consider them all!
 - "may" or "must" analysis?
 - similar to intra-procedural branches at if-then-else constructs (combine)

Threads

- Intra-procedural analyses are typically sound despite multi-threaded execution
- Inter-procedural analyses are typically unsound if flow-sensitive!
- Flow-insensitive analyses not impacted by multi-threading
- Effective modelling of synchronization constructs is a big open research problem!

Library Dependencies

- Typically analyze an application with its dependencies
- But what about native code?
- Often need to resort to hand-crafted summaries
- Possible way out: summarization (e.g., Averroes)

Recap

 Context-sensitivity analyzes a method multiple times, once per context

 Challenges: Recursion, parameters, aliasing, virtual dispatch, threads, libraries

Next

Context sensitivity