#### 10-708: Probabilistic Graphical Models 10-702, Spring 2017

# Homework 1

Lecturer: Eric Xing Name: Yuan Liu(yuanl4)

# 1 Bayesian Network Solution

#### 1.1 I-map

**Solution:** (Koller and Friedman Textbook Page 62)

Let  $X_1, ..., X_n$  is a topological ordering of the variables in  $\mathbb{X}$  relative to  $\mathbb{G}$ . We can first use the chain rule to represent  $\mathbb{P}$ :

$$P(X_1, ..., X_n) = \prod_{i=1}^{n} P(X_i | X_1, ..., X_{i-1})$$

Now, consider one of the factors  $P(X_i|X_1,...,X_{i-1})$ . Because  $\mathbb{G}$  is a I-map for P, we have

$$(X_i \perp NonDescendants_{x_i} | Pa_{x_i}^{\mathbb{G}}) \in I(\mathbb{P})$$

By the topological ordering assumption, all of  $X_i$ 's parents are in the set  $X_1, ..., X_{i-1}$ . Furthermore, none of  $X_i$ 's descendants can possibly be in the set. Hence,

$$\{X_1, ..., X_{i-1}\} = Pa_{X_i} \cup Z$$

where  $Z \subset NonDescendants_{X_i}$ . In addition to it, we already know  $(X_i \perp NonDescendants_{X_i} | Pa_{x_i}^{\mathbb{G}}) \in I(\mathbb{P})$ , so it follows that  $(X_i \perp Z | Pa_{x_i})$ . Hence we have that

$$P(X_i|X_1,...,X_{i-1}) = P(X_i|Pa_{X_i})$$

Applying this transformation to all of the factors in the chain rule decomposition, the result follows.

#### 1.2 D-separation

- $B \perp G|A$ : True. Since H is unobserved, and  $G \rightarrow H \leftarrow F$  is a v-structure, influence cannot flow through H.
- $C \perp D|F$ : False. Influence can flow along the path  $C \leftarrow B \rightarrow D$ .
- $C \perp D|A$ : False. Influence can flow along the path  $C \leftarrow B \rightarrow D$ .
- $H \perp B | C, F$ : True. Since F is observed and  $H \leftarrow F \leftarrow C$  or  $H \leftarrow F \leftarrow E$  are neither v-structure, then H can not flow through F.

# 1.3 I-Equivalence

•  $G_1$  and  $G_2$  has the same skeleton is a necessary, but not sufficient condition for  $G_1$  and  $G_2$  to be I-equivalent.

**Not sufficient part proof:** The v-structure and common parent structure have the same skeleton, but they are not I-equivalent.

**Necessary part proof:** Proof by contradiction. Assume  $G_1$  and  $G_2$  are I-equivalent, but they have different skeleton. In this situation, we can find a trail in one network that does not exist in the other. Assume this trail is  $X_1 \rightleftharpoons X_2 \rightleftharpoons ... \rightleftharpoons X_i$  in  $G_1$ . Given all the v-structures in this trail, like  $A_k \to B_k \leftarrow C_k$ , let all the  $\{B_k\}_k$  observed, and left others unobserved. Then  $X_1 \perp X_2 | \{B_k\}_k$  in  $G_2$ , but  $X_1 \not \perp X_2 | \{B_k\}_k$  in  $G_1$ 

•  $G_1$  and  $G_2$  has the same skeleton and same v-structures, is a sufficient, but not necessary condition for  $G_1$  and  $G_2$  to be I-equivalent.

Sufficient part proof: (http://www.wisdom.weizmann.ac.il/pgm/Ex/ps3\_sol.pdf) First we assume that  $(\mathbb{X} \perp \mathbb{Y} | \mathbb{Z}) \in I(G_1)$  and we show that  $(\mathbb{X} \perp \mathbb{Y} | \mathbb{Z}) \in I(G_2)$ . If two graphs have the same skeleton, then they have the same trails. Lets look on the trails that between  $X \in \mathbb{X}$  and  $Y \in \mathbb{Y}$  that given  $\mathbb{Z}$  is inactive in  $G_1$ , and then show this trail in  $G_2$  is in active too. Consider two cases:

- i The trail in  $G_1$  is inactive because some of the nodes on the trail that are not in v-structure are observed in  $\mathbb{Z}$ . Then clearly these nodes also blocks the trail in  $G_2$
- ii Otherwise, all nodes on the trail that are not in a v-structure are not observed (not in  $\mathbb{Z}$ ), but then for some v-structure  $V_{i-1}, V_i, V_{i+1}$  on the trail, none of the of  $V_i$  are observed. That is for every node V such that there is a directed path in  $G_1$  from  $V_i$  to V, all the nodes on the path are not observed. Consider such a directed path from such a  $V_i$  to some V in  $G_1$ , then in  $G_2$  this trail must also be directed the same, from  $V_i$  to V, because other wise it introduces a v-structure that is not in  $G_1$ (either on this path itself, or with respect to  $V_{i-1} \to V_i$ ), and clearly all it's nodes are not observed too. There fore this path also inactivates the trail between  $\mathbb{X}$  and  $\mathbb{Y}$ (given  $\mathbb{Z}$ ) in  $G_2$

Thus,  $(\mathbb{X} \perp \mathbb{Y}|\mathbb{Z}) \in I(G_1)$  implies that every trail between any pair  $X \in \mathbb{X}$  and  $Y \in \mathbb{Y}$  that given  $\mathbb{Z}$  is inactive in  $G_1$ , and showed for each such a trail that it is inactive in  $G_2$ . Since the two graphs have the same trails, then  $(\mathbb{X} \perp \mathbb{Y}|\mathbb{Z}) \in I(G_2)$ . Hence  $I(G_1) \subset I(G_2)$ , and by symmetry we also get that  $I(G_2) \subset I(G_1)$ , therefore  $I(G_1) = I(G_2)$ 

Not necessary part proof: Consider complete graphs over a set of variables. Recall that a complete graph is one to which we cannot add additional arcs without causing cycles. Such graphs encode the empty set of conditional independence assertions. Thus, any two complete graphs are I-equivalent. Although they have the same skeleton, they invariably have different v-structures.

# 2 Independence & Equivalence Testing Solution

### 2.1 D-separation and Independence

(1) First, prove that D-separation implies conditional independence, or mathematically:

$$desp(X, Y|Z) \Rightarrow X \perp Y|Z$$

**Solution** (Koller and Friedman Textbook Page 137) Let  $U = X \cup Y \cup Z$ . Let  $U^* = U \cup Ancestors_U$ . Let  $G_{U^*} = G^+[U]$  be the induced graph over  $U^*$ , and let H be the moralized graph  $M[G_{U^*}]$ . Let  $P_{U^*}$  be the Bayesian network distribution defined over  $G_{U^*}$  in the obvious way: the CPD for any variable in  $U^*$  is the same as in original distribution  $P_B$ . Because  $U^*$  is upwardly closed, all variables used in these CPDs are in  $U^*$ .

Now, consider an independence assertion  $(X \perp Y|Z) \in I(G)$ ; we want to prove that  $P \models (X \perp Y|Z)$ . By definition, if  $(X \perp Y|Z) \in I(G)$ , we have that desp(X,Y|Z). It follows that  $sep_H(X,Y|Z)$ , and hence that  $(X \perp Y|Z) \in I(H)$ .  $P_{U^*}$  is a Gibbs distribution over H, and hence,  $P_{U^*} \models (X \perp Y|Z)$ . We know  $P_{U^*}(U^*)$  is the same as  $P_B(U^*)$ . Hence, it follows that  $P_B \models (X \perp Y|Z)$ 

- (2) Prove the completeness of the D-separation property.
- (3) Provide a pseudo-code and describe in detail an algorithm for testing D-separation.

**Solution** (Koller and Friedman Textbook Page 75) We use the following definition of active trail to find all the nodes that can be reached by X.

- Whenever we have a v-structure  $X_{i-1} \to X_i \leftarrow X_{i+1}$ , then  $X_i$  or one of its descendants are in Z.
- no other node along the trail is in Z.

So we first need to find all the ancestors of Z, and put them into A.(Using broad first search) Then we use broad-first search to find all the linked node.

#### 2.2 I equivalence

(1) Proof: The two BNs, G and G', are I-equivalent if both graphs have the same set of trails and a trail is active in G if and only if it is active in G'.

**Solution:** First we assume that  $(\mathbb{X} \perp \mathbb{Y} | \mathbb{Z}) \in I(G_1)$  and we show that  $(\mathbb{X} \perp \mathbb{Y} | \mathbb{Z}) \in I(G_2)$ .

Lets look on the trail that between  $X \in \mathbb{X}$  and  $Y \in \mathbb{Y}$  that given  $\mathbb{Z}$  is inactive in  $G_1$ , and because of the assumption this trail in  $G_2$  is inactive too.

Thus,  $(\mathbb{X} \perp \mathbb{Y}|\mathbb{Z}) \in I(G_1)$  implies that every trail between any pair  $X \in \mathbb{X}$  and  $Y \in \mathbb{Y}$  that given  $\mathbb{Z}$  is inactive in  $G_1$ , and showed for each such a trail that it is inactive in  $G_2$ . Since the two graphs have the same trails, then  $(\mathbb{X} \perp \mathbb{Y}|\mathbb{Z}) \in I(G_2)$ . Hence  $I(G_1) \subset I(G_2)$ , and by symmetry we also get that  $I(G_2) \subset I(G_1)$ , therefore  $I(G_1) = I(G_2)$ 

(2) Prove that a minimal active trail may only contain a triangle of the following form:

#### **Algorithm 1** D-separation

```
Input: Bayesian network graph: G; Source variables: X; Evidence variables: Z;
Output: Reachable Set R
 1: //Step one: find all the ancestors of Z
 2: L \leftarrow Z //Node to be visited
 3: A \leftarrow \emptyset //Ancestors of Z
 4: //Broad first search
 5: while L \neq \emptyset do
        Select some Y from L
 6:
         L \leftarrow L - \{Y\}
 7:
        if Y \notin A then
 8:
 9:
             L \leftarrow L \cup Pa_Y
10:
         end if
         A \leftarrow A \cup \{Y\}
11:
12: end while
          //Step two:traverse active trails starting from X
13: L \leftarrow \{(X,\uparrow)\} //(Node, direction) to be visited
14: V \leftarrow \emptyset //(Node, direction) marked as visited
15: R \leftarrow \emptyset //Nodes reachable via active trail
    while / do/Select some (Y,d) from L
17:
         L \leftarrow L - (Y, d)
        if (Y,d) \notin V then
18:
             if Y \notin Z then
19:
                 R \leftarrow R \cup \{Y\} // Y is reachable
20:
             end if
21:
             V \leftarrow V \cup (Y, d) // \operatorname{Mark}(Y, d) as visited
22:
             if d = \uparrow and Y \notin Z then //Trail up through Y active if Y not in Z
23:
                 for each Z \in Pa_Y do //Y's parents to be visited from bottom
24:
                     L \leftarrow L \cup \{(Z,\uparrow)\}
25:
                 end for
26:
                 for each Z \in Ch_Y do //Y's children to be visited from top
27:
                     L \leftarrow L \cup \{(Z,\downarrow)\}
28:
                 end for
29:
             else if d = \downarrow then
30:
                 if Y \notin Z then //Downward trails to Y's children are active
31:
                     for each Z \in Ch_Y do
32:
                         L \leftarrow L \cup \{(Z,\downarrow)\} //Y's children to be visited from top
33:
                     end for
34:
                 end if
35:
                 if Y \in A then //v-structure trails are active
36:
                     for each Z \in Pa_Y do
37:
                         L \leftarrow L \cup \{(Z,\uparrow)\} // Y's parents to be visited from bottom
38:
                     end for
39:
                 end if
40:
             end if
41:
         end if
42:
43: end while
44: Return R
```

**Solution** Let first consider the trail:  $X_{i-1} \rightleftharpoons X_i \rightleftharpoons ... \rightleftharpoons X_{i+1}$ . There are four following situations:

- $-X_{i-1} \to X_i \to X_{i+1}$ . Because these three variables form a triangle. There are two sub-cases in this situation:
  - \*  $X_{i-1} \to X_{i+1}$  In this situation,  $X_{i-1} \to X_{i+1}$  will be an active trail, which breaks the minimal assumption.
  - \*  $X_{i-1} \leftarrow X_{i+1}$  In this situation, there will be a cycle among these three variables.
- $-X_{i-1} \leftarrow X_i \leftarrow X_{i+1}$ . Because these three variables form a triangle. There are two sub-cases in this situation:
  - \*  $X_{i-1} \to X_{i+1}$  In this situation, there will be a cycle among these three variables.
  - \*  $X_{i-1} \leftarrow X_{i+1}$  In this situation,  $X_i \rightarrow X_{i-1} \leftarrow X_{i+1}$  is a v-structure. So there must be  $X_{i-2} \leftarrow X_{i-1}$  to make the whole trail active. It follows  $X_{i-2} \leftarrow X_{i-1} \leftarrow X_{i+1}$  is active, which breaks the minimal assumption.
- $-X_{i-1} \to X_i \leftarrow X_{i+1}$ . Because these three variables form a triangle. There are two sub-cases in this situation:
  - \*  $X_{i-1} \to X_{i+1}$  In this situation,  $X_{i-2} \rightleftharpoons X_{i-1} \to X_{i+1}$  will be an active trail, which breaks the minimal assumption.
  - \*  $X_{i-1} \leftarrow X_{i+1}$  In this situation,  $X_{i-1} \rightarrow X_i \leftarrow X_{i+1}$  is a v-structure, so there must be  $X_{i-2} \leftarrow X_{i-1}$ . In this situation  $X_{i-2} \leftarrow X_{i-1} \leftarrow X_{i+1}$  breaks the minimal assumption.
- $-X_{i-1} \leftarrow X_i \rightarrow X_{i+1}$  This is the situation we want to proof.
- (3) Two Bayes networks, G and G', are I-equivalent if and only if G and G' have the same skeletons and the same set of immoralities.

#### Solution:

**Part one:** Prove that if G and G' have the same skeletons and the same set of immoralities then they are I-equivalent. So we need to prove if  $(X \perp Y|Z) \in I(G)$ , then  $(X \perp Y|Z) \in I(G')$ . Proof by contradiction.

Assume  $(X \perp Y|Z) \in I(G)$  and  $(X \perp Y|Z) \notin I(G')$ . So there is a minimal active trail between X and Y given Z in G', but blocked in G. For example  $X_{i-1} \rightleftharpoons X_i \rightleftharpoons X_{i+1}$  is blocked in G. There are two situations it is blocks:

Situation one: In  $X_{i-1}, X_i, X_{i+1}, X_i$  is observed and they form cascade or common parent structure in G, while they form v-structure in G'. Because they forms v-structure in G' and the trail is minimal, there is no link between  $X_{i-1}$  and  $X_{i+1}$  (Problem 2's result). Thus this v-structure is a immorality in G'. So it should also be a immorality in G. This is a contradiction.

Situation two:  $X_{i-1}, X_i, X_{i+1}$  are all unobserved and they form a v-structure in G, while in G' they don't form a v-structure. Thus  $X_{i-1}, X_i, X_{i+1}$  can not form a immorality, and there will be a edge between  $X_{i-1}$  and  $X_{i+1}$ . It violates the assumption that  $X_{i-1}$  and  $X_{i+1}$  is blocked in G.

**Part two:** Prove that two networks G and G' that induce the same conditional independence assumptions must have the same skeleton and the same immoralities. By Problem 1.3 we already proved that G and G' have the same skeleton, thus we only need to prove they have the same immoralities. Proof by contradiction:

Two networks G and G' are I-equivalent, have the same skeleton but different immoralities. For example,  $X_{i-1}$ ,  $X_i$ ,  $X_{i+1}$  forms a immorality in G, but they form a cascade or common parent structure in G'. Given  $X_i$  observed so the trail  $X_{i-1} \to X_i \leftarrow X_{i+1}$  is active in G, where in G' this trail is inactive. This situation violates the Proposition 1, we proved.

# Algorithm 2 I-equivalent

```
Input: Bayesian network graph's direct edge set (X,Y): G1, G2;
Output: Whether they are I-equivalent: Flag
 2: function isIEquivalent(G1,G2)
    //Whether two graphs have the same skeleton
       if isSameSkeleton(G1,G2)=False then
 4:
 5:
           return False
       end if
 6:
    //Get immoralities of each graph
       M1 \leftarrow getImmoralities(G1)
       M2 \leftarrow qetImmoralities(G2)
 9:
       if M1 != M2 then
10:
           return False
11:
       end if
12:
       return True
13:
   end function
14:
15:
    // Whether two graphs have the same skeleton
17: function isSameSkeleton(G1,G2)
       if G1.size()!=G2.size() then
18:
           return False
19:
20:
       end if
       for (X,Y) \in G1 do
21:
           if (X,Y) \notin G2 and (Y,X) \notin G2 then
22:
              return False
23:
           end if
24:
       end for
25:
       return True
26:
27: end function
28:
   //Output the immoralities in the graph
29:
30: function GETIMMORALITIES(G)
31: //Get the node Set V
32:
       V \leftarrow \emptyset
       for (X,Y) \in G do
33:
           V \leftarrow V \cup \{X, Y\}
34:
       end for//Iteration all there node triple to find the immoralities
35:
       M \leftarrow \emptyset //The set of all the immoralities
36:
       for X \in V do
37:
           for Y \in V do
38:
              for Z \in V do
39:
                  if X,Y,Z are distinct and (X,Y) \in G and (Z,Y) \in G then
40:
                      if (X,Z) \notin G and (Z,X) \notin G then
41:
                         M = M \cup (X, Y, Z)
42:
                      end if
43:
                  end if
44:
               end for
45:
           end for
46:
47:
       end for
       return M
48:
49: end function
```

# 3 Exact Inference

# 3.1 Variable Elimination

$$\alpha_{t+1}(j) = p(z_{t+1} = j | x_{1,...,t+1}) \propto \sum_{i} p(z_{t+1} = j | z_{t} = i) p(x_{t+1} | z_{t+1} = j) p(z_{t} = i | x_{1,...,t})$$

$$\propto p(x_{t+1} | z_{t+1} = j) \sum_{i} p(z_{t+1} = j | z_{t} = i) \alpha_{t}(i)$$

$$\beta_{t-1}(j) = P(x_{t,...,T} | z_{t-1} = j) \propto \sum_{i} P(x_{t} | z_{t} = i) P(z_{t} = i | z_{t-1} = j) P(x_{t+1,...,T} | z_{t} = i)$$

$$\propto \sum_{i} p(x_{t} | z_{t} = i) p(z_{t} = i | z_{t-1} = j) \beta_{t}(i)$$

Then because

$$p(z_t = j | x_1, ..., T) \propto \alpha_t(j) \beta_t(j)$$

we can get

$$p(z_t = j | x_1, ..., T) = \frac{\alpha_t(j)\beta_t(j)}{\sum_i \alpha_t(i)\beta_t(i)}$$

# 3.2 Gaussian Belief Propagation

(1) Proof: 
$$N(x|\mu_1, \lambda_1^{-1}) \times N(x|\mu_2, \lambda_2^{-1}) = CN(x|\mu, \lambda^{-1})$$

Solution

$$N(x|\mu_1, \lambda_1^{-1}) \times N(x|\mu_2, \lambda_2^{-1}) = exp(-\frac{\lambda_1(x-\mu_1)^2}{2}) \times exp(-\frac{\lambda_2(x-\mu_2)^2}{2})$$
$$= exp(-\frac{1}{2}(\lambda_1(x-\mu_1)^2 + \lambda_2(x-\mu_2)^2)$$

$$\begin{split} \lambda_1(x-\mu_1)^2 + \lambda_2(x-\mu_2)^2 = &(\lambda_1+\lambda_2)x^2 - 2(\lambda_1\mu_1 + \lambda_2\mu_2)x + \lambda_1\mu_1^2 + \lambda_2\mu_2^2 \\ = &(\lambda_1+\lambda_2)[x^2 - \frac{2(\lambda_1\mu_1 + \lambda_2\mu_2)}{\lambda_1 + \lambda_2}x + (\frac{\lambda_1\mu_1 + \lambda_2\mu_2}{\lambda_1 + \lambda_2})^2] - \frac{(\lambda_1\mu_1 + \lambda_2\mu_2)^2}{\lambda_1 + \lambda_2} \\ + &(\lambda_1\mu_1^2 + \lambda_2\mu_2^2) \\ = &(\lambda_1+\lambda_2)[x - \frac{\lambda_1\mu_1 + \lambda_2\mu_2}{\lambda_1 + \lambda_2}]^2 + \frac{\lambda_1\lambda_2}{\lambda_1 + \lambda_2}(\mu_1 - \mu_2)^2 \end{split}$$

So we can get

$$\lambda = \lambda_1 + \lambda_2$$

$$\mu = \frac{\lambda_1 \mu_1 + \lambda_2 \mu_2}{\lambda_1 + \lambda_2}$$

$$C = \sqrt{\frac{\lambda_1 \lambda_2}{2\pi(\lambda_1 + \lambda_2)}} exp(-\frac{1}{2} \frac{\lambda_1 \lambda_2}{\lambda_1 + \lambda_2} (\mu_1 - \mu_2)^2)$$

(2) Solve for  $m(x_t)$ . Hint First solve for the messages passed into Node t's neighbor nodes from their neighbors excluding t, then use these messaged to represent  $m(x_t)$ 

Solution: (Gaussian Belief Propagation: Theory and Application https://arxiv.org/pdf/0811.2518.pdf)

$$\begin{split} N(\mu_{s/t}, \lambda_{s/t}) &\propto \phi_s(x_s) \prod_{k \in N(s)/t} m_{ks}(x_s) \\ \lambda_{s/t} &= \lambda_{ss} + \sum_{k \in N(s)/t} \lambda_{ks} \\ \mu_{s/t} &= \lambda_{s/t}^{-1} (\lambda_{ss} \mu_{ss} + \sum_{k \in N(s)/t} \lambda_{ks} \mu_{ks}) \end{split}$$

Where  $\lambda_{ks}$  and  $\mu_{ks}$  are the precision and mean of the message  $m_{ks}(x_s)$ , and  $\lambda_{ss} \triangleq A_{ss}$ ,  $\mu_{ss} \triangleq \frac{b_s}{A_{ss}}$ . Next we calculate  $m_{st}(x_t)$ 

$$m_{st}(x_t) = \int_{x_s} \phi_{s,t}(x_s, x_t) \phi_s(x_s) \prod_{k \in N(s)/t} m_{ks}(x_s)$$

$$\propto \int_{x_s} exp(-x_s A_{st} x_t) exp(-\lambda_{s/t} (x_s^2/2 - \mu_{s/t} x_s))$$

$$= \int_{x_s} exp(-\lambda_{s/t} x_s^2/2 + (\lambda_{s/t} \mu_{s/t} - A_{st} x_t) x_i)$$

$$\propto exp((\lambda_{s/t} \mu_{s/t} - A_{st} x_t)^2/(2\lambda_{s/t}))$$

$$\propto N(\mu_{st}, \lambda_{st})$$

Where  $\mu_{st}$  and  $\mu_{st}$  are the precision and mean of the message  $m_{st}(x_t)$ .

$$\mu_{st} = \frac{\lambda_{s/t}\mu_{s/t}}{A_{st}}$$
$$\lambda_{st} = -\frac{A_{st}^2}{\lambda_{s/t}}$$

Then we can calculate  $m(x_t)$ 

$$m(x_t) = \phi_t(x_t) \prod_{s \in N(t)} m_{st}(x_t) \propto N(\mu_t, \lambda_t^{-1})$$
$$\lambda_t = \lambda_{tt} + \sum_{s \in N(t)} \lambda_{st}$$
$$\mu_t = \lambda_t^{-1} (\lambda_{tt} \mu_{tt} + \sum_{s \in N(t)} \lambda_{st} \mu_{st})$$

Where  $\lambda_{tt} = A_{tt}$ ,  $\mu_{tt} = \frac{b_t}{A_{tt}}$ 

# 4 Undirected Graphical Models

### 4.1 Ising Model & Boltzmann Machine

Solution:

$$\begin{split} \exp(\sum_{(i,j)\in E} W_{ij} x_i x_j - \sum_{i\in V} u_i x_i) &\propto \exp(\sum_{(i,j)\in E} 4W_{ij} (\frac{x_i}{2} + \frac{1}{2}) (\frac{x_j}{2} + \frac{1}{2}) - \sum_{(i,j)\in E} W_{ij} x_i - \sum_{(i,j)\in E} W_{ij} x_j - \sum_{i\in V} u_i x_i) \\ &\propto \exp[\sum_{(i,j)\in E} 4W_{ij} (\frac{x_i}{2} + \frac{1}{2}) (\frac{x_j}{2} + \frac{1}{2}) - \sum_{i\in V} (u_i + 2\sum_{j\in N(i)} W_{ij}) x_i] \\ &\propto \exp[\sum_{(i,j)\in E} 4W_{ij} (\frac{x_i}{2} + \frac{1}{2}) (\frac{x_j}{2} + \frac{1}{2}) - \sum_{i\in V} 2(u_i + 2\sum_{j\in N(i)} W_{ij}) (\frac{x_i}{2} + \frac{1}{2})] \end{split}$$

Thus we can get:

$$x'_{i} = \frac{x_{i}}{2} + \frac{1}{2} \in \{0, 1\}$$

$$W'_{ij} = 4W_{ij}$$

$$u'_{i} = 2(u_{i} + 2\sum_{j \in N(i)} W_{ij})$$

### 4.2 Determinantal Point Process