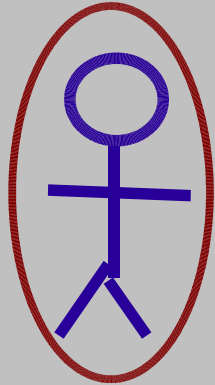


(not a mascot)



Operating System (without a name) (so far)

Benji, Jaap & Ting

Traditional “timesharing” OS

- Protect processes from each other
- Usual trick:
 - Virtualization (pretend we have >1 machine)
- Problems:
 - Context switches are slow
 - Copying buffers is slow
 - Too little security (big rooms)
 - Too much security (big walls)

What is the actual goal?

To protect resources!

```
char* vga = (char *) 0xb8000;  
strcpy("H e l l o ", vga);
```

- UNIX programmer says:
 - “Shouldn't have that address mapped.”
- ML/Lisp/have-your-pick programmer says:
 - (char *) is inherently unsafe

Loading user code in UNIX

- `fork(2)` creates a new memory space
 - `exec(2)` loads an object file into memory
 - Figure out where `main()` is and jump to it
-

Protection through virtualization

Loading user code in (unnamed)

- Check that code came out of a type-safe compiler
 - Do *not* generate a new memory map
 - Load an object file into memory
 - Figure out where main() is and jump to it
-

Protection without virtualization

Why does type safety help?

- *Theorem: “While this program is being executed, no memory outside of range R will be accessed, and control will not be transferred outside of range U .”*
- *Proof:*
 - For a correctly typed program: trivial
 - Pointers don't appear out of thin air
 - For a C program: impossible
 - Pointers do appear out of thin air

(unnamed OS) so far....



Booting made easy



Emulation made easy

```
qemu -m 16 -fda boot -fdb root
```

Next: basic device drivers

- VGA
- 8259 PIC (Interrupt controller)
- Keyboard
- Floppy disk
- Hard drive
- File system