Key Update for OSCORE (KUDOS)

draft-ietf-core-oscore-key-update-12

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Recap of KUDOS (1/2)

- **Key Update for OSCORE (KUDOS)**
  - Renew the Master Secret and Master Salt; derive new Sender/Recipient keys
  - No change to the ID Context or Sender/Recipient ID; can achieve Forward Secrecy
  - Agnostic of the key establishment method originally used
  - Loosely inspired by Appendix B.2 of OSCORE
  - The peers update their current context CTX_OLD, deriving a new context CTX_NEW
  - Redesigned in v-10 to use a more flexible and simpler approach
    - We now explicitly define a KUDOS state machine

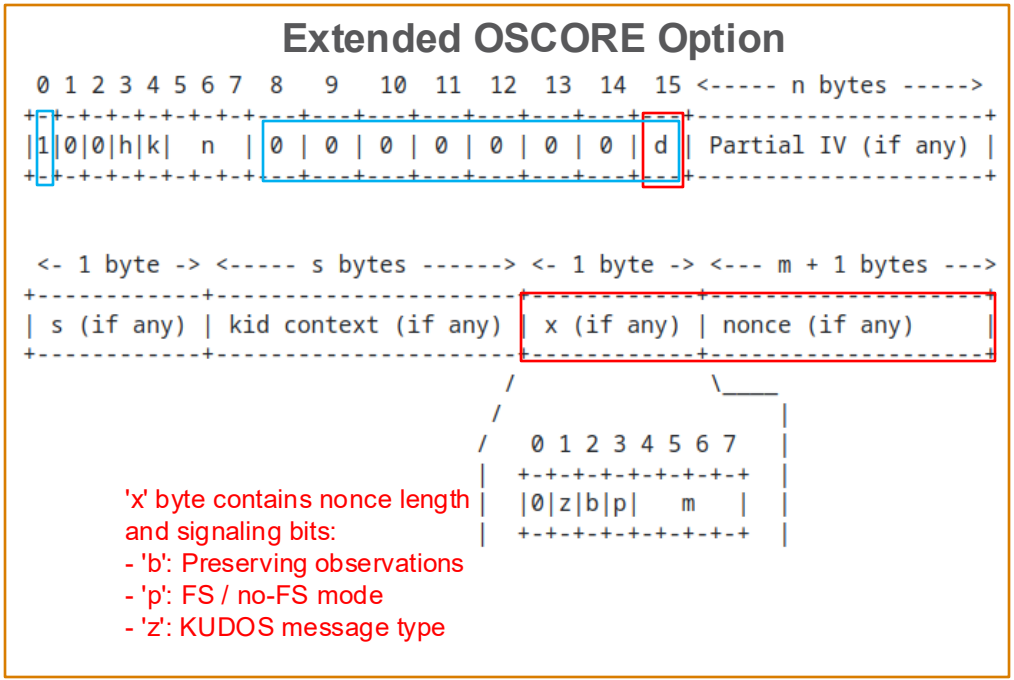
Recap of KUDOS (2/2)

- **Properties**
  - Can be initiated by either peer
  - It is robust against a peer rebooting and loss of state, avoiding the reuse of AEAD (nonce, key)
  - It typically completes in one round trip by exchanging two OSCORE-protected CoAP messages
    - The two peers achieve mutual key confirmation in a following exchange, which is protected with the newly established OSCORE Security Context
  - Flexible in terms of message flow; any CoAP message can be a KUDOS message

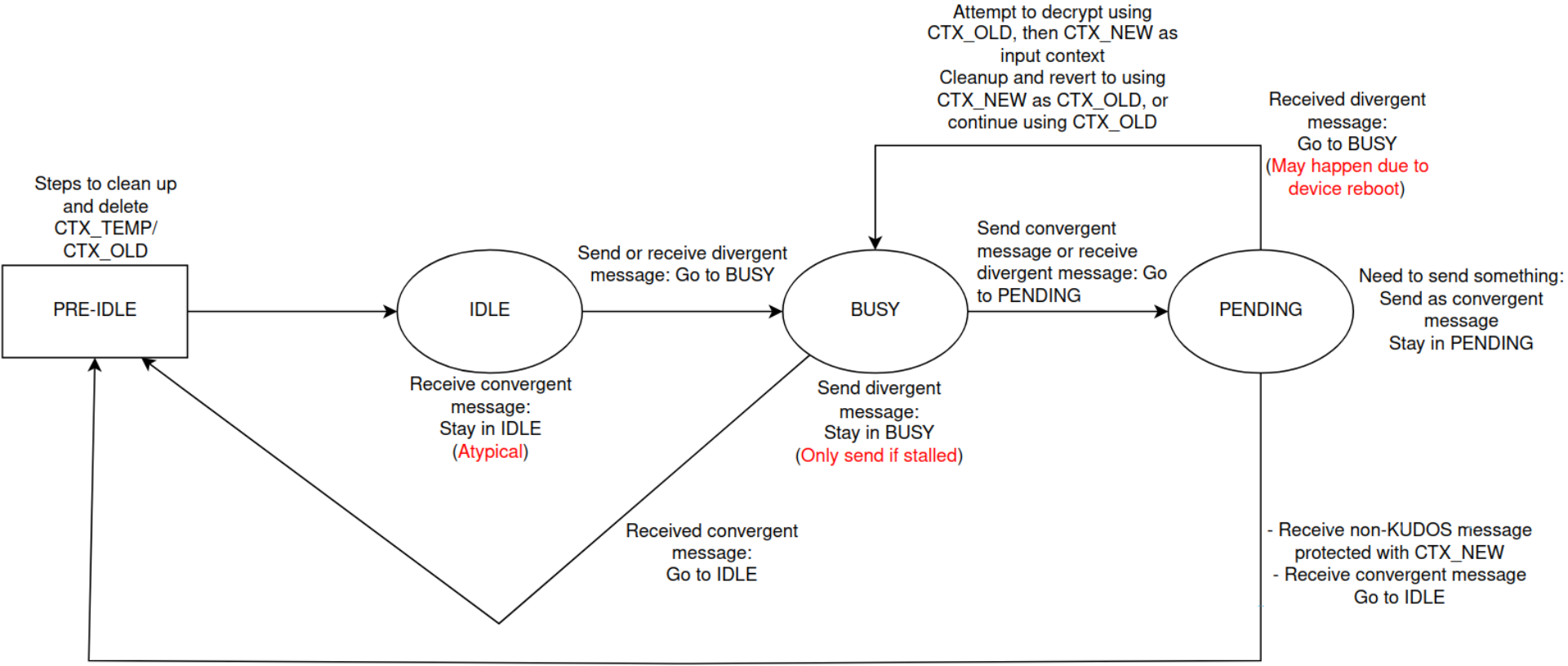
- **KUDOS message types**
  - Divergent message: The sender peer is offering its own nonce in the message and waiting to receive the other peer's nonce.
  - Convergent message: The sender peer is offering its own nonce in the message, has received the other peer's nonce, and is going to wait for key confirmation

Rekeying Procedure

- **Key Update for OSCORE (KUDOS)**
  - Message exchange to share two nonces N1 and N2
    - Decoupled from request/response and client/server concepts
  - Nonces are placed in new fields in OSCORE CoAP option
  - *UpdateCtx()* function for deriving new OSCORE Security Context using the two nonces, two ‘x’ bytes and CTX_OLD
  - Two modes
    - FS mode providing forward secrecy
    - No-FS mode for very constrained devices
  - No change of OSCORE identifiers
  - Expected to complete in 1 round trip

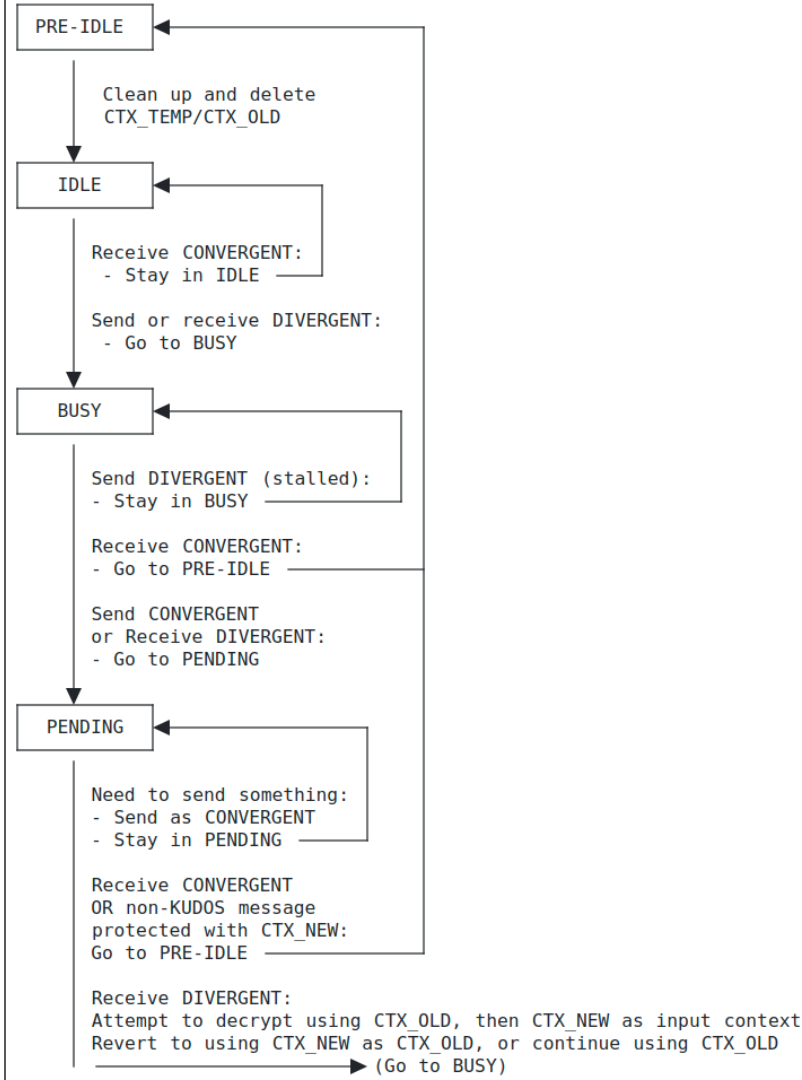


KUDOS State Machine



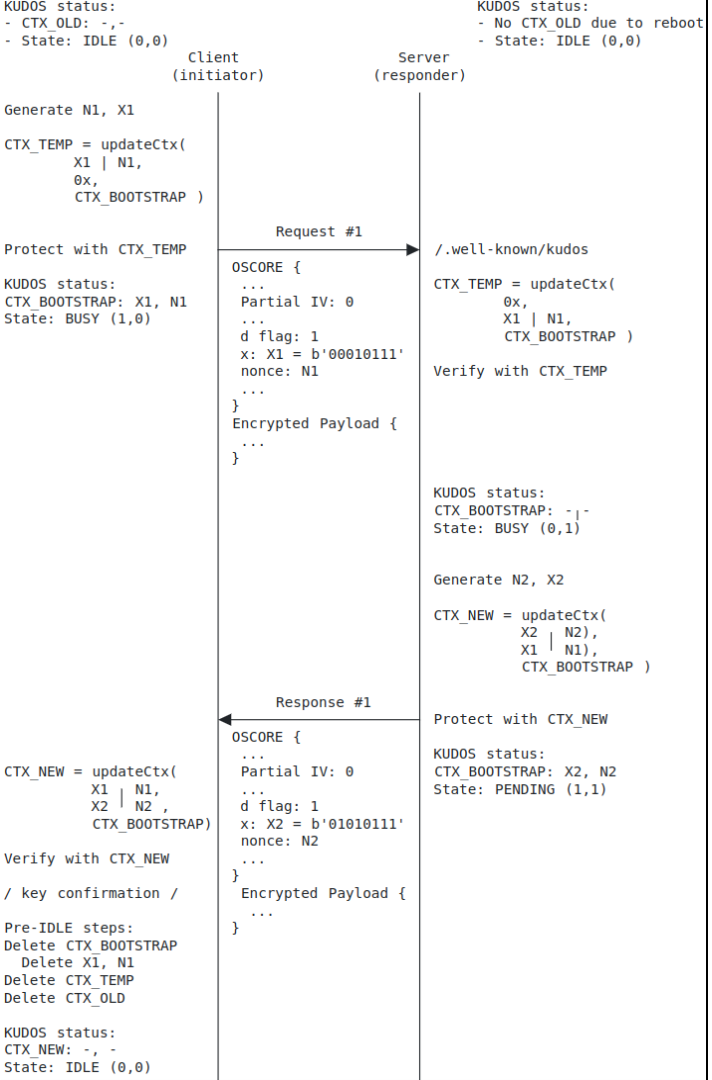
Changes for v-12

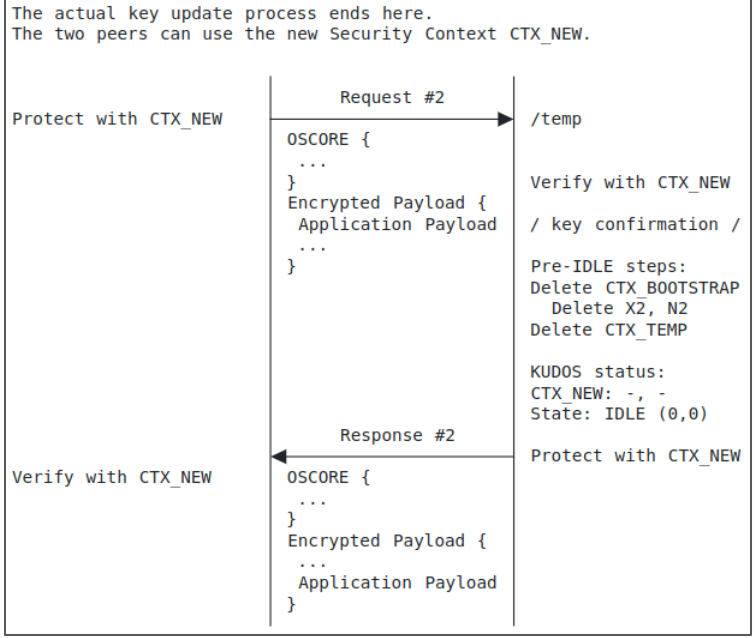
- **Added state machine as ASCII figure**
  - Now in Appendix B
  - Aligned with was shown in the previous slide
  - The state machine is described in Section 4.3.3
- **Editorial improvements**



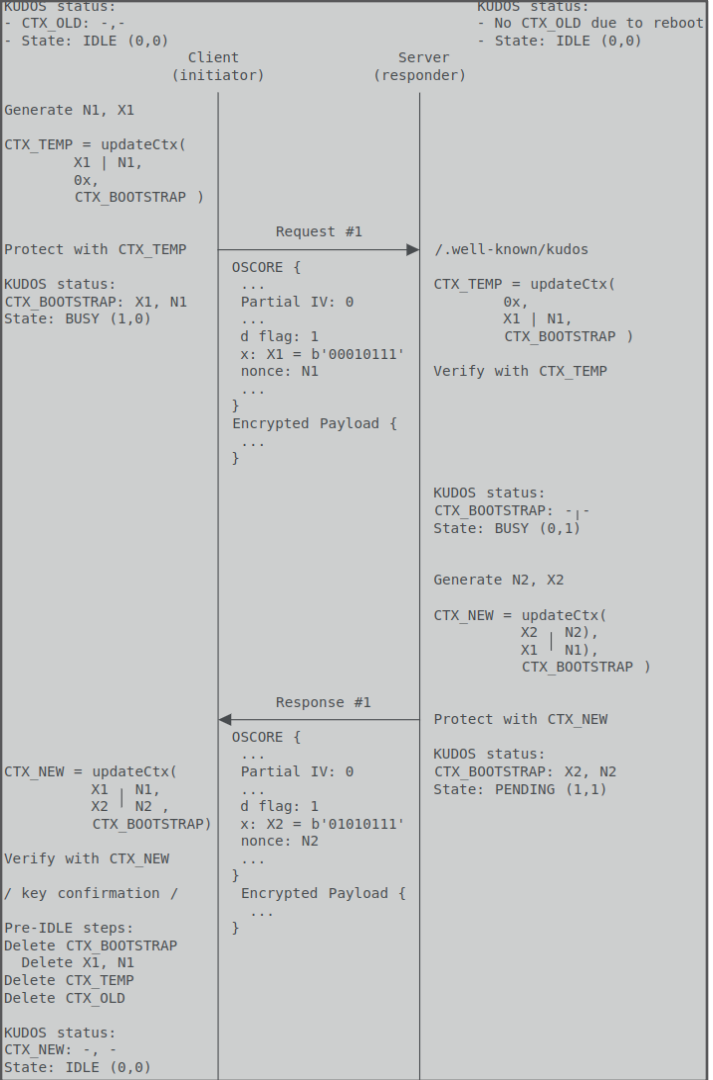
Changes for v-12

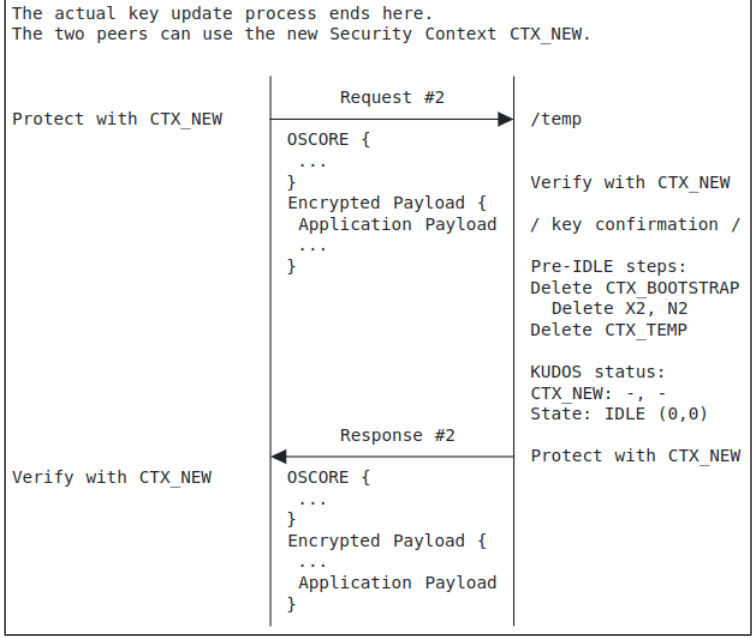
- **Multiple additional message flow examples as appendixes**
  - KUDOS Execution Initiated with a Request Message, with Non-capable Server that has Rebooted
    - *The peers have run KUDOS prior to this KUDOS execution and have learned that they must from now on run KUDOS only in no-FS mode*
  - KUDOS Execution Initiated with a Request Message, where the Client Executes KUDOS again after the first Execution
    - *A second KUDOS execution is started by the client immediately after a successful KUDOS key update*
  - KUDOS Execution Initiated with a Request Message, where KUDOS Response #1 is Lost
    - *The server's first response is dropped by the network; thus the client retries and both sides end up deriving the same a CTX_NEW*
  - KUDOS Execution Completed using two Request Messages
    - *Both peers independently initiate KUDOS and exchange two request messages that ultimately result in the same CTX_NEW*





Successful KUDOS Execution Initiated with a Request Message, with Non-capable Server that has Rebooted





Successful KUDOS Execution Initiated with a Request Message, with Non-capable Server that has Rebooted

Summary and next steps

- **KUDOS implementations**
  - Update the implementation in Java [1] to be aligned with the latest design
  - Update the implementation in C for Contiki-NG to be aligned with the latest design
- **Procedural: Move content relevant to SCHC from draft-ietf-schc-8824-update**
  - Considering the timeline, that content about the extended OSCORE option fits better in this draft
- **Comments and reviews are welcome!**

[1] https://github.com/rikard-sics/californium/blob/oscore_cmd/cf-oscore/

Thank you!

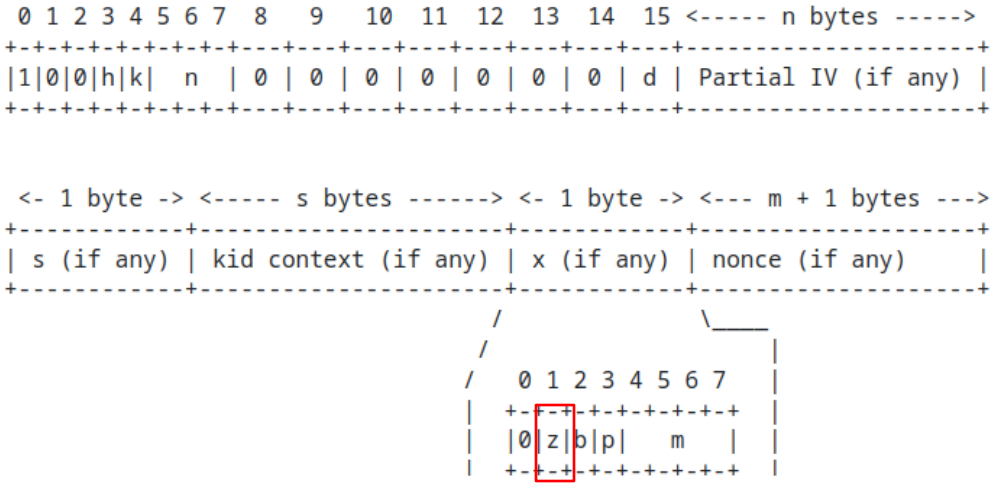
Comments/questions?

https://github.com/core-wg/oscore-key-update

Backup

KUDOS Message Types

- **Two types of KUDOS messages, distinguished by the 7th least significant bit 'z' in the 'x' byte**
  - Indicates if only one or both nonces have been exchanged



Note: The z bit used to have another meaning

- **z = 0: "divergent message"**
- This message is protected with the temporary Security Context CTX_TEMP (was CTX_1)
- The sender peer is offering its own nonce in the message and waiting to receive the other peer's nonce.

- **z = 1: "convergent message"**
- This message is protected with the final Security Context CTX_NEW.
- The sender peer is offering its own nonce in the message, has received the other peer's nonce, and is going to wait for key confirmation

State Machine when in PENDING

* Upon receiving a divergent message:

  - In case of successful decryption and verification of the message using a CTX_TEMP derived from CTX_OLD:

    1. Delete CTX_NEW.

    2. Delete the pair (X, nonce) associated with the Security Context CTX_IN that was used to generate the CTX_NEW deleted at the previous step.

    3. Abort the ongoing KUDOS execution.

    4. Move to **BUSY** and enter it consistently with the reception of a divergent message.

  - Otherwise, in case of successful decryption and verification of the message using a CTX_TEMP derived from CTX_NEW:

    1. Delete the oldest CTX_TEMP.

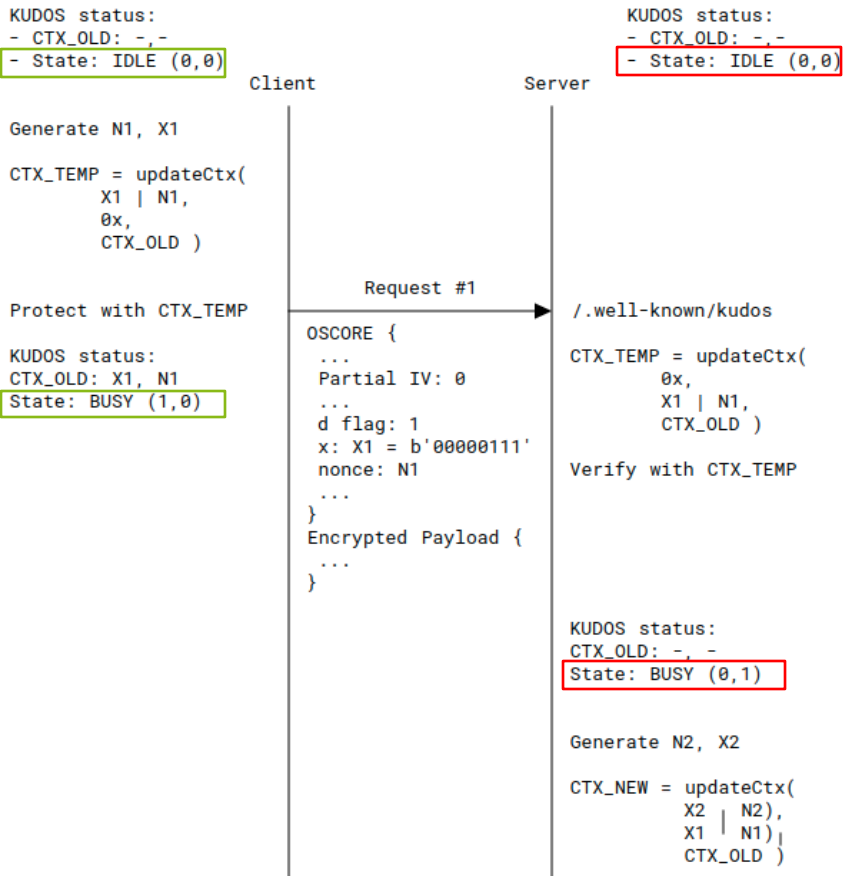
    2. Delete the Security Context that was used as CTX_IN to generate the CTX_TEMP deleted at the previous step.

    3. CTX_NEW becomes the oldest Security Context. From this point on, that Security Context is what this KUDOS execution refers to as CTX_OLD.

    4. Abort the ongoing KUDOS execution.

    5. Move to **BUSY** and enter it consistently with the reception of a divergent message.

Example Execution



Example Execution