

Graph Decomposition (2)

— BFS and its Applications

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Outline

BFS: Algorithm

BFS: Properties

BFS: Applications

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Problem

Problem: Shortest Path

Given an undirected graph $G = (V, E)$ and a source vertex s , to compute distance $\delta(s, u)$ for each vertex u .

$\delta(s, u) \equiv \#$ of edges of the shortest path between s and u .

BFS on Undirected Graph

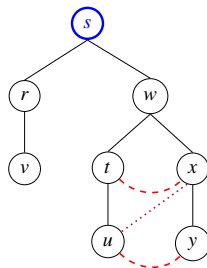
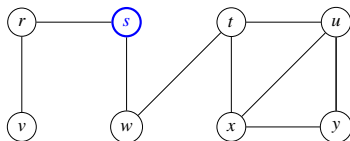
BFS with source vertex s :

- ▶ exploring every vertex u reachable from s
- ▶ computing $\delta(s, u)$, for all reachable u

BFS as a framework:

- ▶ Prim's MST algorithm
- ▶ Dijkstra's SSSP algorithm

A Physical Algorithm of BFS (Phys-BFS)



Edge Properties of Phys-BFS

$$(u, v) \in E \Rightarrow d(u) \leq d(v) \leq d(u) + 1$$

A Parallel Algorithm of BFS (Para-BFS)

graph: network of computers

vertex: computer

edge: network connections

To disseminate a computer virus from computer s .

Color Properties in Para-BFS

$$\text{states of computers} = \begin{cases} \text{WHITE} & \text{if healthy} \\ \text{GRAY} & \text{if infected} \end{cases}$$

A Parallel Algorithm of BFS (Para-BFS)

Algorithm 1 A Parallel Algorithm of BFS (Para-BFS).

```
procedure PARA-BFS( $G, s$ )  
  for all  $u \in V$  do  
     $\text{color}[u] \leftarrow \text{WHITE}$   
     $d[u] \leftarrow \infty$   
  
   $\text{color}[s] \leftarrow \text{GRAY}$   
   $d[s] \leftarrow \infty$   
  
   $Q \leftarrow \{s\}$   
  
  while  $Q \neq \emptyset$  do  
     $u \leftarrow \text{Deq}(Q)$   
    for all  $(u, v) \in E$  do  
      if  $\text{color}[v] = \text{WHITE}$  then  
         $\text{color}[v] \leftarrow \text{GRAY}$   
         $d[v] \leftarrow d[u] + 1$   
         $\text{Enq}(Q, v)$   
     $\text{color}[u] \leftarrow \text{BLACK}$ 
```

Color Properties of Para-BFS

States of vertices:

WHITE: undiscovered

GRAY: discovered but

BLACK: discovered and all its neighbors has been discovered

$$\text{WHITE} \Rightarrow \text{GRAY} \Rightarrow \text{BLACK}$$

1. $(u, v) \in E$, u is **BLACK** $\Rightarrow v$ is either **BLACK** or **GRAY**
2. **GRAY** vertex may have adjacent **WHITE** vertices
 \Rightarrow “frontier” between discovered and undiscovered vertices
3. Invariant: at any time, all **GRAY** vertices are in the Queue

Correctness Proof Para-BFS

BFS Algorithm

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Color Properties

Queue Properties

Correctness Proof

Edges Properties

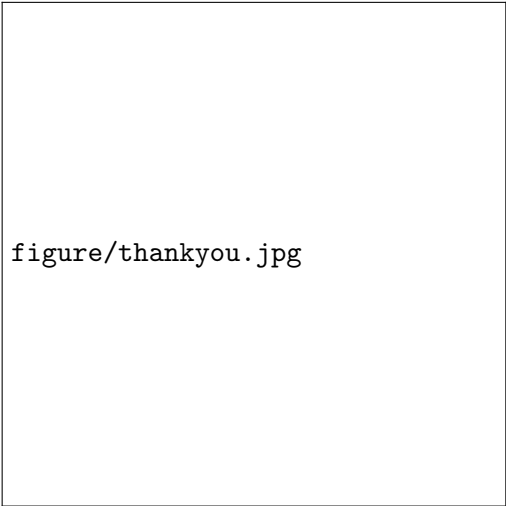
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Testing Bipartiteness



figure/thankyou.jpg