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The best programs are written so that computing machines can perform them quickly and so that human beings can understand them clearly. A programmer is ideally an essayist who works with traditional aesthetic and literary forms as well as mathematical concepts, to communicate the way that an algorithm works and to convince a reader that the results will be correct. — **Donald E. Knuth**

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Algorithms

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Dynamic Programming

Problem :-

Splitting a String into minimum number of palindromes.

Given a string, find the minimum no. of palidromes into which the string can be broken.

Consider the string "civic" (String is Palindrome by itself, so min no. of palindromes = 1)

The string "abcde" cannot be splitted into less than 5 palindromes (each character is a palindrome)

Consider the string "civicacaramanamaracacammacdeified".

It can be broken into 4 palindromes " civic " + " acaramanamaraca " + " cammac " + " deified "

Solution :-

Suppose the size of the string is n.

We maintain a 2D array palins [n][n] where palins [i][j] stores the minimum number of palindromes in the string str [i...j] i.e from index 'i' to index 'j' of the string. So, the final solution of breaking the string of size n into minimum no. of palindromes is given by palins [0][n-1].

Let us determine the structure of the optimal solution.

Let palins (i , j) denotes the min no. of palindromes in the string str (i ... j).

We will try to break the string at all possible positions k such that $i \le k \le j$ and sum up the minimum no. of palindromes in **palins** (i, k) and **palins** (k, j).

Thus, the structure of optimal solution is determined recursively as :

palins (i , j) = palins (i , k) + palins (k , j) for all $i \le k \le j$

Finally, palins ($\bf 0$, $\bf n$ - $\bf 1$) gives the required solution. See the implementation below.

```
#include<iostream>
     #include<cstring>
    using namespace std;
     // check if a string is palindrome, given the start and end index of the str
bool isPalindrome(char str[], int i, int j) {
 6
         while (i < j) {
   if (str[i++] != str[j--])</pre>
 8
                 return false;
 9
10
         return true;
11
12
    }
13
     /* Find the minimum no. of palindromes in the input string i.e the
14
15
         string has to broken down into min. no. of palindromes
         Suppose palins[i][j] -> min no. of palindromes in the string str[i...j] Then, palins[0][n-1] is the final solution
16
17
         Standard Recursive Solution :
18
19
         palins(i,j) = min(palins(i,k) + palins(k,j)) for i <= k < j</pre>
20
21
     int minPalindromes(char str[]) {
         int i, j, k;
22
         int substr_size; //substring size
23
24
         int n = strlen(str);
25
         // Matrix where cell (i,j) stores the minimum no. of palindromes in str[i
26
         int palins[n][n];
27
28
         // iterating over all possible substring sizes
         for (substr_size = 1; substr_size <= n; substr_size++) {
    // 'i' and 'j' tracks the start and end index of the current substring
    for (i = 0; i <= n - substr_size; i++) {</pre>
29
30
31
                 j = i + substr_size - 1;
if (isPalindrome(str, i, j)) {
   // substring is a palindrome
32
33
34
35
                     palins[i][j] = 1;
36
37
                 else {
                     // search for palindromes within the substring
38
                     // worst case (min. no of palindromes = length of the substring)
// each character is a palindrome
39
```

```
palins[i][j] = substr_size;
                       // find min no. of palindromes within substring for (k = i; k < j; k++) {
 42
 43
                          // try splitting substring at each position
int sum = palins[i][k] + palins[k+1][j];
if (sum < palins[i][j])
  palins[i][j] = sum;</pre>
 44
 45
 46
 47
 48
 49
                  }
               }
 50
 51
           \H// return min no. of palindromes in the input string
 52
 53
           return palins[0][n-1];
 54
      }
 55
       // main
 56
 57
       int main() {
           char str[] = "civicacaramanamaracacammacdeified";
//civic + acaramanamaraca + cammac + deified (4 palindromes)
 58
 59
           int min_palin = minPalindromes(str);
 60
           \verb|cout|<<"\nMinimum no. of palindromes|:: "<<min_palin;
 61
           cout<<endl;
 62
 63
           return 0;
 64
     }
      4
                                                                                          Compile & Run
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All problems on Dynamic Programming
* Find the nth term of fibonacci series
* Evaluate combination(n, r)
* Solve the Edit-Distance problem
* Longest Common Subsequence ( LCS ) problem
* Given a set of coin denominations, find the minimum number of coins required to make a change for
a target value
* Longest Increasing Subsequence ( LIS ) problem
* Unbounded Knapsack problem
* 0/1 Knapsack problem
* Splitting a string into minimum number of palindromes
```

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