

# Paths of Graphs

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for all  $v \in V$  do
     $\text{dist}[v] \leftarrow \infty$ 
 $\text{dist}[s] \leftarrow 0$ 

 $Q \leftarrow \text{MINPQ}(V)$ 
while  $Q \neq \emptyset$  do
     $u \leftarrow \text{DELETMIN}(Q)$ 

    for all  $(u, v) \in E \wedge v \notin Q$  do
        if  $\text{dist}[v] > \text{dist}[u] + l(u, v)$  then
             $\text{dist}[v] \leftarrow \text{dist}[u] + l(u, v)$ 
             $\text{DECREASEKEY}(Q, v)$ 

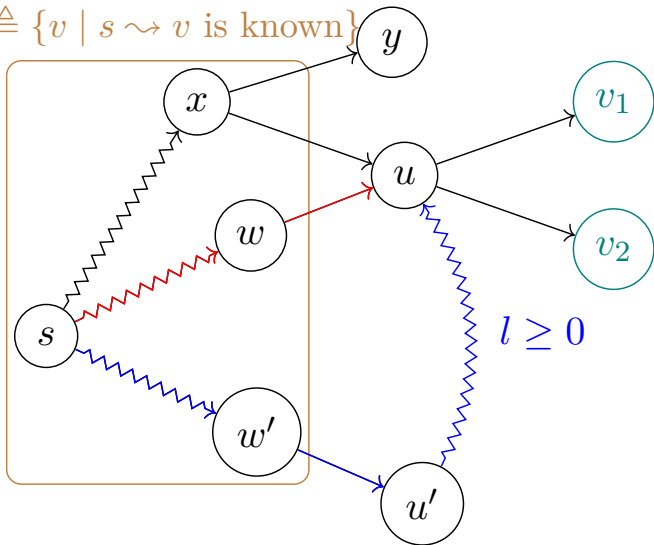
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$$O(n + (n + m) \log n) \implies O((n + m) \log n) \implies O(m \log n)$$

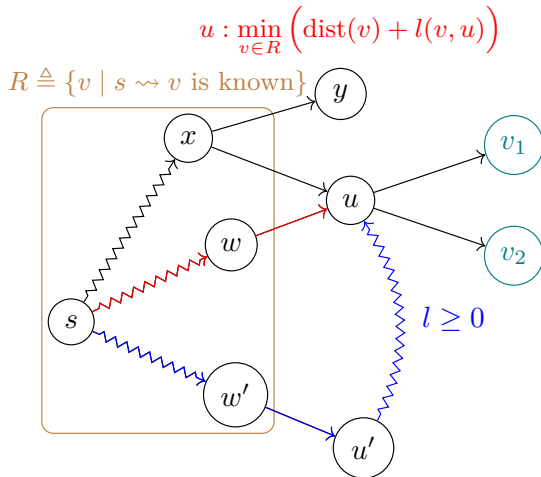
$$u : \min_{v \in R} (\text{dist}(v) + l(v, u))$$

$$R \triangleq \{v \mid s \rightsquigarrow v \text{ is known}\}$$



## Negative Edges from $s$ (Problem 11.9)

All negative edges are from  $s$ .



## Shortest Paths Through $v_0$ (Problem 13.7)

Strongly connected digraph  $G = (V, E)$ ,  $w(e) > 0$

$$v_0 \in V$$

Find shortest paths  $s \rightsquigarrow^{\text{SP}} t$  through  $v_0$ .

$$s \rightsquigarrow^{\text{SP}} v_0 \rightsquigarrow^{\text{SP}} t$$

$$\forall v : v_0 \rightsquigarrow^{\text{SP}} v$$

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**for all**  $v \in V$  **do**

$\text{dist}[v] \leftarrow \infty$

$\text{dist}[s] \leftarrow 0$

$Q \leftarrow \text{MinPQ}(V)$

**while**  $Q \neq \emptyset$  **do**

$u \leftarrow \text{DELETEMIN}(Q)$

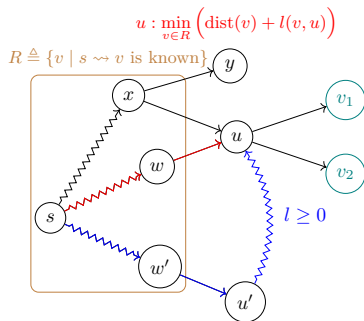
**for all**  $(u, v) \in E \wedge v \notin Q$  **do**

**if**  $\text{dist}[v] > \text{dist}[u] + l(u, v)$  **then**

$\text{dist}[v] \leftarrow \text{dist}[u] + l(u, v)$

$\text{DECREASEKEY}(Q, v)$

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## Min-Max Path (Problem 11.12)

$G = (V, E)$  : network of highways

$l_e$  : road length     $L$  : tank capacity

Given  $G$ , to compute  $\min L$  in  $O(m \log n)$  from  $s$  to  $t$ .

$Q \leftarrow \text{MinPQ}(V)$

**for all**  $v \in V$  **do**

$L[v] \leftarrow \infty$

$L[s] \leftarrow 0$

**if**  $L[v] > \max \left( L[u], l(u, v) \right)$  **then**

$L[v] \leftarrow \max \left( L[u], l(u, v) \right)$

## Max-Min Path (Problem 13.2 (1))

$G = (V, E)$  : network of oil pipelines

$c(u, v)$  : capacity of  $(u, v)$

$\text{cap}(s, t)$  :  $\max \min s \rightsquigarrow t$

Given  $s$ , to compute  $\text{cap}(s, v)$ .

$Q \leftarrow \text{MaxPQ}(V)$

**for all**  $v \in V$  **do**

$\text{cap}[v] \leftarrow -\infty$

$\text{cap}[s] \leftarrow 0$

**if**  $\text{cap}[v] < \min(\text{cap}[u], c(u, v))$  **then**

$\text{cap}[v] \leftarrow \min(\text{cap}[u], c(u, v))$



## Max-Min Path (Problem 13.2 (2))

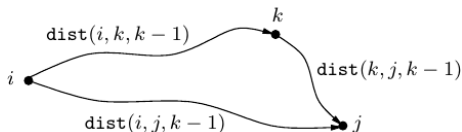
$G = (V, E)$  : network of oil pipelines

$c(u, v)$  : capacity of  $(u, v)$

$\text{cap}(s, t)$  :  $\max \min s \rightsquigarrow t$

Compute all-pair  $\text{cap}(i, j)$ .

$$\text{cap}(i, j, k) = \max \left( \text{cap}(i, j, k-1), \min(\text{cap}(i, k, k-1), \text{cap}(k, j, k-1)) \right)$$



## Most Critical Edge (Problem 11.3)

$$s, t \in V$$

$$e : E \setminus \{e\} \implies \text{dist}(s, t) \text{ increases most}$$



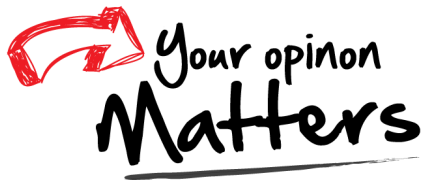
*“Most Vital Links and Nodes in Weighted Networks”, 1992*

$$O(m \log n)$$

## Bitonic Shortest Path (Problem 11.7)







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