6.851	Lecture 20	May 3, 2012

TODAY: Dynamic graphs II
- fully vs. partially dynamic
- Euler-tour trees
- $O(1)$ decremental connectivity in trees
- $O(\lg^2 n)$ fully dynamic connectivity
- survey

Dynamic connectivity:
maintain undirected graph subject to
- insert/delete edges or vertices (with no edges)
- connectivity(v, w): is there a $v \to w$ path?
  or (): is the graph connected? (≈ same)

Dynamic graph problems: characterized by updates
- fully dynamic: insert & delete ← default
- partially dynamic:
  - incremental: just insert
  - decremental: just delete

Dynamic connectivity results:

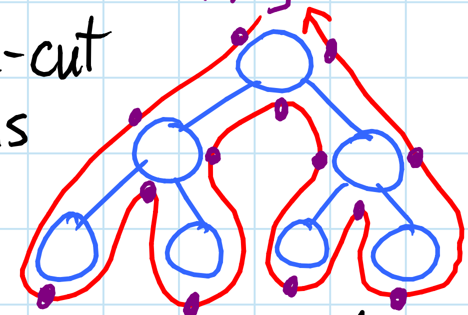
- trees: $O(\lg n)$ → link-cut [L19] & Euler tour trees → TODAY
  - decremental: $O(1)$ amortized → TODAY
- plane graphs: $O(\lg n)$
  embedded planar [Eppstein, Galil, Italiano, Spencer – JCSS 1996]
- general graphs, amortized:
  - $\boxed{\text{OPEN}}$: $O(\lg n)$ update & query
  - $O(\lg^2 n)$ update, $O(\frac{\lg n}{\lg\lg n})$ query → TODAY
    [Holm, de Lichtenberg, Thorup – J.ACM 2001]
  - $O(\lg n (\lg\lg n)^3)$ update, $O(\frac{\lg n}{\lg\lg\lg n})$ query
    [Thorup – STOC 2000]
  - incremental: $O(\alpha(m,n))$ via union-find [Tarjan – JACM 1975]
  - decremental: $O(m \lg n + n \text{ polylg } n + \#\text{queries})$ total
    [Thorup – J.ACM 1999]
- worst case: (general graphs)
  - $\boxed{\text{OPEN}}$: polylg update & query
  - $O(\sqrt{n})$ update, $O(1)$ query
    [Eppstein, Galil, Italiano, Nissenweig – J.ACM 1997]
  - incremental: $\Theta(x)$ updates $\Rightarrow$ $\Theta(\frac{\lg n}{\lg x})$ queries
    [Alstrup, Ben-Amram, Rauhe – STOC 1999]
- lower bounds: $\Omega(\lg n)$ update or query
  - $O(x \lg n)$ update $\Rightarrow \Omega(\frac{\lg n}{\lg x})$ query
  - $O(x \lg n)$ query $\Rightarrow \Omega(\frac{\lg n}{\lg x})$ update

  even for paths!

  [L21] & [Pătrașcu & Demaine – STOC 2004/SICOMP 2006]

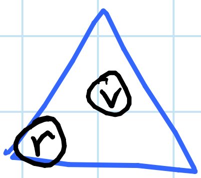
→ points on trade-off curve
  - $\boxed{\text{OPEN}}$: $o(\lg n)$ update & polylg n query?

Euler-tour trees: [Henzinger & King – STOC 1995]

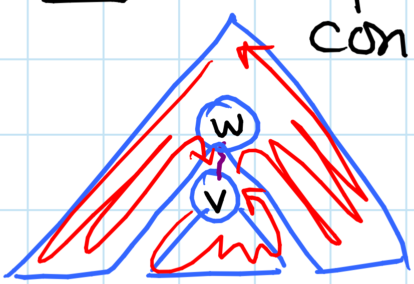


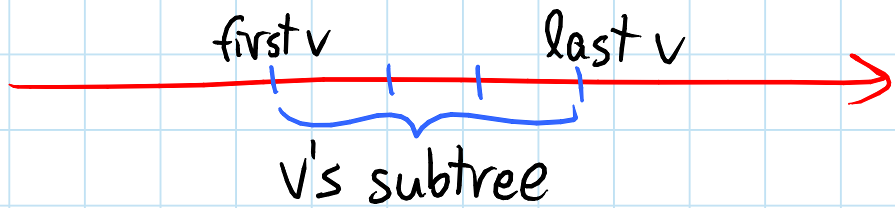
- simpler dynamic trees than link-cut
- aggregates over subtrees, vs. paths
- Euler tour [L15] around tree
  - visits each edge twice
- store node visits by Euler tour in balanced BST
- each node stores pointers to first & last visits

- findroot(v): start at first visit to v in BST
  walk up to root of BST
  walk left to min of BST
  → first visit of root node

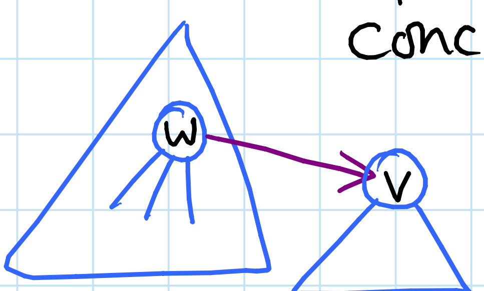


- cut(v): split BST at v's first & last visits
  concatenate "before" & "after" trees





- link(v, w): split w's BST before w's last visit
  concatenate "before last w",
  new single w,
  v's BST, and
  "after last w"



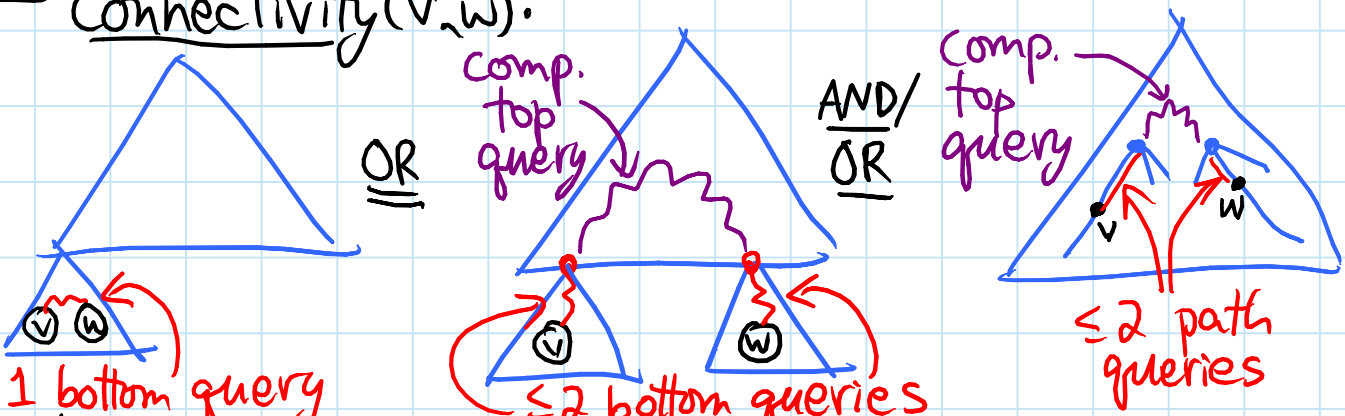
- connectivity(v, w): findroot(v) $\stackrel{?}{=}$ findroot(w)
- subtree aggregate(v): (min/max/sum/etc.)
  range query in BST between first & last visit
- $O(\lg n)$ time/op.

Decremental connectivity in a rooted tree: [Alstrup, Secher, Spork – IPL 1997]

$O(1)$ amortized, assuming all $n-1$ edges deleted

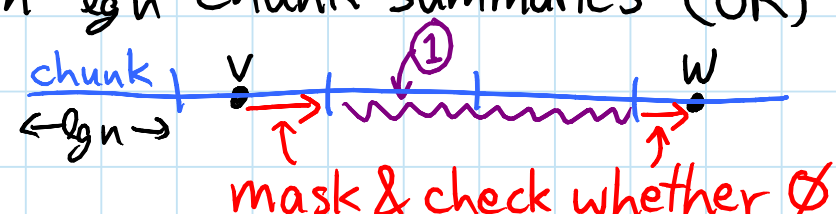
① $O(\lg n)$ via link-cut or Euler-tour trees
(simpler way: maintain explicit node→component id & relabel the smaller side)

② leaf trimming: cut below maximally deep nodes with $> \lg n$ descendants
$\Rightarrow$ top tree has $O(\frac{n}{\lg n})$ leaves/branching nodes
– use ① on compressed top tree
– connectivity(v, w):



③ bottom tree in $O(1)$: (worst case, fully dynamic)
– store bit vector of which edges don't exist
– preprocess mask of each node's ancestors (1 word)
– XOR masks for v & w, mask, check whether $\emptyset$

④ path in $O(1)$ amortized:
– split path into $\frac{n}{\lg n}$ chunks of length $\lg n$
– store each chunk as bit vector (1 word)
– use ① on $\frac{n}{\lg n}$ chunk summaries (OR)
– query:



$O(\lg^2 n)$ dynamic connectivity: [Holm et al. — J.ACM 2001]

- idea:
  - store spanning forest with Euler-tour trees
  - hierarchically divide connected components
  - → $O(\lg n)$ levels of spanning forests

→ charging mechanism
- level of edge starts at $\lg n$, only decreases → $\emptyset$
- G_i = subgraph of edges at level $\le i$

$\Rightarrow G_{\lg n} = G$
- INVARIANT 1: every conn. comp. of G_i has $\le 2^i$ vxs.
- F_i = spanning forest of G_i
  - store using Euler-tour tree DS

$\Rightarrow F_{\lg n}$ is desired spanning forest of G
- INVARIANT 2: $F_\emptyset \subseteq F_1 \subseteq \cdots \subseteq F_{\lg n}$ i.e. $F_i = F_{\lg n} \cap G_i$
  i.e. $F_{\lg n}$ is a min. spanning forest w.r.t. level

insert($e = (v, w)$): $O(\lg n)$
- add e to v & w incidence lists
- $e.\text{level} = \lg n$
- if v & w disconnected in $F_{\lg n}$:
  - add e to $F_{\lg n}$ (link)
    (reroot to make v root via cyclic shift)

connectivity: $O\left(\frac{\lg n}{\lg\lg n}\right)$
- make $F_{\lg n}$ B-trees with branching factor $\Theta(\lg n)$

$\Rightarrow \underbrace{O\left(\frac{\lg n}{\lg\lg n}\right)}_{\text{depth}}$ findroot & $\underbrace{O\left(\frac{\lg^2 n}{\lg\lg n}\right)}_{\text{depth} \cdot \text{branching factor}}$ update delete($e=(v,w)$):

- remove e from v & w incidence lists
- if e is in $F_{\lg n}$: (v.parent = w or vice versa)

  $\lg^2 n \rightarrow$ - delete e from $F_{e.level}, \ldots, F_{\lg n}$ (cut)
  - look for replacement edge to reconnect v & w
    - can't be any edges with level < e.level
    - $\Rightarrow$ find min. possible level $\geq$ e.level [Invariant 2]
  - for i = e.level, ..., $\lg n$:
    - let T_v & T_w be trees of F_i containing v & w resp.
    - relabel so that $|T_v| \leq |T_w|$
    - Invariant 1 (before) $\Rightarrow |T_v| + |T_w| \leq 2^i \Rightarrow |T_v| \leq 2^{i-1}$
    - $\Rightarrow$ can "afford" to push all of T_v down to level $i-1$
    - for each level-i edge $e' = (x, y)$ with $x \in T_v$:
      - if $y \in T_w$: add e' to $F_i, F_{i+1}, \ldots, F_{\lg n}$

        return (found replacement)
      - else ($y \in T_v$): e'.level $= i-1$ $\leftarrow$ charge
- Euler-tour tree augmentation:
  - subtree sizes to test $|T_v|$ vs. $|T_w|$ in $O(1)$
  - for each node v in tree of F_i:

    does v's subtree contain any nodes incident to level-i edges?

    $\Rightarrow$ can find next level-i edge incident to $x \in T_v$ in $O(\lg n)$ time

    (successor, skipping over empty subtrees)

$\Rightarrow$ time: $O(\lg^2 n + \#\text{charges} \cdot \lg n)$

- each inserted edge charged $\leq \lg n$ times

k-connectivity: vertex or edge

- disjoint paths between pairs of vertices:
  - 2-edge: $O(\lg^4 n)$ – 2-vertex: $O(\lg^5 n)$ [Holm et al. – JACM 2001]
  - $\boxed{\text{OPEN}}$: polylg n for $k = O(1)$? $k = $ polylg n?
  - planar decremental: $O(\lg^2 n)$ 3-edge-conn. [Giammaresi & Italiano – Algorithmica 1996]
  - worst case: [Eppstein et al. – J.ACM 1997]
    - 2-edge-conn.: $O(\sqrt{n})$ – 2-vertex-conn.: $O(n)$
    - 3-edge-conn.: $O(n^{2/3})$ – 3-vertex-conn.: $O(n)$
    - $k=4$: $O(n \cdot \alpha(n))$
    - $O(1)$-edge-conn.: $O(n \lg n)$
- whole graph ≈ min cut = max flow
  - $O(\text{polylg } n)$-edge-conn. (& min cut up to that size): $O(\sqrt{n} \cdot \text{polylg } n)$ [Thorup – STOC 2001]
  - $\boxed{\text{OPEN}}$: polylg n for $k = O(1)$? $k = $ polylg n?

Minimum spanning forest: (MST on each conn. comp. as dyn. tree)

- $O(\lg^4 n)$ update [Holm, de Lichtenberg, Thorup – J.ACM 2001]
- worst case: $O(\sqrt{n})$ update [Eppstein et al. – J.ACM 1997]
- plane graphs: $O(\lg n)$ [Eppstein et al. – J.ACM 1992]
- can use to solve bipartiteness: is graph 2-colorable?

Planarity testing: insert e or report planarity violation

- $O(n^{2/3})$ [Galil, Italiano, Sarnak – J.ACM 1997]
- plane (fix embedding): $O(\lg^2 n)$ [Eppstein et al. – J.ACM 1997]
- incremental: $O(\alpha(m,n))$ amortized [la Poutré – STOC 1994]

Directed graphs:

Transitive closure: is there a $v \to w$ directed path?

- bulk update: insert/delete vertex & incident edges
- $O(n^2)$ am. bulk update, $O(1)$ worst-case query
  [Demetrescu & Italiano — FOCS 2000; Roditty — SODA 2003]
- same, worst case [Sankowski — FOCS 2004]
- optimal if explicitly storing trans. closure matrix
- $\boxed{\text{OPEN}}$: $o(n^2)$ worst-case update?
- $O(m\sqrt{n} \cdot t)$ am. bulk update, $O(\sqrt{n}/t)$ w.c. query
  for any $t = O(\sqrt{n})$ [Roditty & Zwick — FOCS 2002]
- $O(m + n \lg n)$ am. bulk update, $O(n)$ w.c. query
  [Roditty & Zwick — STOC 2004]
- $\boxed{\text{OPEN}}$: full trade-off: update · query $= O(mn)$ or $O(n^2)$
- acyclic: $O(n^{1.575} \cdot t)$ update, $O(n^{0.575}/t)$ query, $t = O(\sqrt{n})$
- decremental: $O(n)$ am. update, $O(1)$ w.c. query
  [Demetrescu & Italiano — FOCS 2000]

All-pairs shortest paths: weight of shortest $v \to w$ path

- $O(n^2(\lg n + \lg^2(1 + \frac{m}{n})))$ am. bulk update, $O(1)$ w.c. query
  [Thorup — SWAT 2004] improving [Demetrescu & Italiano — STOC 2003]
- $\boxed{\text{OPEN}}$: $O(n^2)$ or $o(n^2)$ update, even undirected graphs?
- $O(n^{2.75})$ w.c. update, $O(1)$ query [Thorup — STOC 2005]
- unweighted: $O(m\sqrt{n} \cdot \text{polylg } n)$ am. update,
  $O(n^{3/4})$ w.c. query [Roditty & Zwick — ESA 2004]
- undirected, unweighted & $(1+\varepsilon)$-approx.: [Roditty & Zwick
  $O(\sqrt{m}\, n \cdot t)$ am. update, $O(\sqrt{m}/t)$ w.c. query, $t = O(\sqrt{n})$ — FOCS 2004]
- static & $(1+\varepsilon)$-approx.: distance oracles [...]