#### Searching and Selection

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- Selection
- Searching

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$$V_1(n) = n - 1$$
  
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#### Reference

"The Art of Computer Programming, Vol 3: Sorting and Searching" by Donald E. Knuth.

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"Does your algorithm need to find the 1st and the 2nd elements?"

"YES!"

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"YFS!"

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"Does your algorithm need to find the 1st and the 2nd elements?"

"YES!"

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"NO!"

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"YES!"

"Do all algorithms have to find the 1st and the 2nd elements?"

"NO!"

#### References

"Selecting the Top Three Elements" by Aigner, 1982.



#### The largest k elements (Problem 3.5)



### Close to median (Problem 3.6)



#### Medians of sorted arrays (Problem 3.7)

#### Weighted median (Problem 3.9)

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#### $\max / \min$ differences (Problem 4.5)

- (a) unsorted;  $\max |x y|$ ; O(n)
- (b) sorted;  $\max |x y|$ ; O(1)
- (c) unsorted;  $\min |x y|$ ;  $O(n \log n)$
- (d) sorted;  $\min |x y|$ ; O(n)

- ightharpoonup M: matrix  $m \times n$
- row: increasing from left to right
- col: increasing from top to down
- ▶ Is  $x \in M$ ?

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Divide and conquer.

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Always checking the lower left corner.



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Always checking the lower left corner.

$$T(m,n) = m + n - 1$$

Assume  $M: n \times n$ 

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 $W(n) \geq 2n-1$  by adversary argument!

$$i+j \le n-1 \implies x > M_{ij}$$
  
 $i+j > n-1 \implies x < M_{ij}$ 

#### Local