Minimum Spanning Tree (MST)

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May 26, 2017 – May 31, 2017

Minimum Spanning Tree (MST)

A generic MST algorithm

Cut property (strong)

Cut property (strong)

- Graph $G = (V, E)$
- X is some part of an MST T of G
- Any cut $(S, V \setminus S)$ *s.t.* X does not cross $(S, V \setminus S)$
- Let e be a lightest edge across $(S, V \setminus S)$

Then $X \cup \{e\}$ is some part of an MST T' of G.

Proof.

Exchange argument

□

Cut property (strong)

Correctness of Prim's and Kruskal's algorithms.

Cut property (weak)

Cut Property [Problem: 3.6.18 (a)]

- Graph $G = (V, E)$
- Any cut $(S, V \setminus S)$ where $S, V - S \neq \emptyset$
- Let $e = (u, v)$ be *a* minimum-weight edge across $(S, V \setminus S)$

Then e must be in *some* MST of G.

$$\text{“a”} \to \text{“the”} \implies \text{“some”} \to \text{“any”}$$

Applications of cut property

Application of cut property (Problem 6.10)

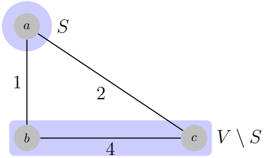
(3) (Problem 6.10–3) $e \in G$ is a lightest edge $\implies$ $e \in \exists$ MST of G

(4) $e \in G$ is the unique lightest edge $\implies$ $e \in \forall$ MST of G

Applications of cut property

Wrong divide&conquer algorithm for MST (Problem 6.14)

- $G = (V, E, w)$
- $(V_1, V_2) : ||V_1| - |V_2|| \le 1$
- $T_1 + T_2 + \{e\}$: e is a lightest edge across (V_1, V_2)



Cycle property (weak)

Cycle property (Problem 6.13–2)

- $G = (V, E, w)$
- Let C be any cycle in G
- $e = (u, v)$ is *a* maximum-weight edge in C

Then $\exists$ MST T of $G : e \notin T$.

"a" $\rightarrow$ "the" $\implies$ "some" $\rightarrow$ "any"

Applications of cycle property

Anti-Kruskal algorithm (Problem 6.13–3)

Reverse-delete algorithm (wiki)

$$O(m \log n (\log \log n)^3)$$

Proof.

Invariant: If F is the set of edges remained at the end of the while loop, then there is some MST that are a subset of F. □

Reference

- "On the Shortest Spanning Subtree of a Graph and the Traveling Salesman Problem" by Kruskal, 1956.

Application of cycle property

(Problem 6.13–1)

(1) $e \notin$ any cycle of $G \implies e \in \forall$ MST

By contradiction.

Application of cycle property

(Problem 6.10)

(1) $|E| > |V| - 1$, e is the unique max-weight edge of $G \implies e \notin \forall$ MST

(2) $\exists C \subseteq G$, e is the unique max-weight edge of $G \implies e \notin \forall$ MST

(5) Cycle $C \subseteq G$, $e \in C$ is the unique lightest edge of $G \implies e \in \forall$ MST

Unique MST

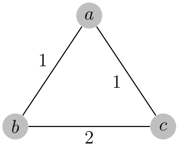
Unique MST (Problem 6.12–1)

Distinct weights $\implies$ unique MST.

Unique MST

Unique MST (Problem 6.12–2)

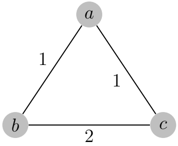
Unique MST $\not\Longrightarrow$ Equal weights.



Unique MST

Unique MST (Problem 6.12–3)

Unique MST $\not\Longrightarrow$ Minimum-weight edge across any cut is unique.



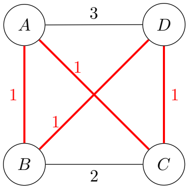
Theorem

Minimum-weight edge across any cut is unique $\Longrightarrow$ *Unique MST.*

Unique MST

Unique MST (Problem 6.12–3)

Unique MST $\not\Longrightarrow$ Maximum-weight edge in any cycle is unique.



Theorem (Conjecture)

Maximum-weight edge in any cycle is unique $\Longrightarrow$ *Unique MST.*

Unique MST

Unique MST (Problem 6.12–4)

Decide whether a graph has a unique MST?

Modify an MST by exchange argument?

Minimum Spanning Tree (MST)

Prim *vs.* Kruskal (Problem 6.4)

$$T(n,m) = O(nT(\textsf{getMin}) + nT(\textsf{deleteMin}) + mT(\textsf{decreaseKey}))$$

MST on special graphs (Problem 6.3)

Prim on special graphs (Problem 6.1)

Prim on special graphs (Problem 6.2)

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THANK YOU