# Computational topology: Lecture 4

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Simplicial Complexes

Cellular Complexes

Chain Complexes

# Simplicial Complexes

### Join operation

The join of two sets  $P,Q \subset \mathbb{E}^n$  is the set of convex combination of their points:

$$PQ = \{\alpha \mathbf{x} + \beta \mathbf{y} | \mathbf{x} \in P, \mathbf{y} \in Q\},\$$

where  $\alpha, \beta \in \mathbb{R}$ ,  $\alpha, \beta \geq 0$  and  $\alpha + \beta = 1$ .

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A simplicial complex, often simply denoted as complex, is a triangulation  $\Sigma$  that verifies the following conditions:

- **1** if  $\sigma \in \Sigma$ , then any face of  $\sigma$  belongs to  $\Sigma$ ;
- ② if  $\sigma, \tau \in \Sigma$ , then either  $\sigma \cap \tau = \emptyset$ , or  $\sigma \cap \tau$  is a face of both  $\sigma$  and  $\tau$ .

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With some abuse of language, we call (combinatorial) s-skeletons the sets  $\mathcal{K}_s$  ( $s \leq d$ ). Geometric carrier  $|\Sigma|$ , also called the support space, is the point-set union of simplices in  $\Sigma$ .

#### Implementation

How?

help?> Lar.simplexFacets

#### Orientation

The ordering of the 0-skeleton of a simplex implies an orientation of it. The simplex can be oriented according to the even or odd permutation class of its 0-skeleton.

The two opposite orientations of a simplex will be denoted as  $+\sigma$  and  $-\sigma$ .

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#### It is assumed that:

- The two orientations of a simplex may represent its relative interior and exterior;
- 2 the two orientations of an orientable simplicial complex analogously represent the relative interior and exterior of the complex, respectively;
- the boundary of a complex maintains the same orientation of the complex.

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The volume associated with an orientation of a simplex (or complex) is positive, while the one associated with the opposite orientation has the same absolute value and opposite sign.

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$$+\sigma_3 = \langle \mathbf{v}_0, \mathbf{v}_1, \mathbf{v}_2, \mathbf{v}_3 \rangle$$
  
 $-\sigma_3 = \langle \mathbf{v}_1, \mathbf{v}_0, \mathbf{v}_2, \mathbf{v}_3 \rangle$ 

### Implementation

How?

help?> Lar.surface

The oriented facets  $\sigma_{d-1,(i)}$   $(0 \le i \le d)$  of the oriented d-simplex  $\sigma_d = +\langle \mathbf{v}_0, \mathbf{v}_1, \dots, \mathbf{v}_d \rangle$  are obtained by removing the i-th vertex  $\mathbf{v}_i$  from the 0-skeleton of  $\sigma_d$ :

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$$\sigma_{d-1,(i)} = (-1)^{i} (\sigma_{d} - \langle \mathbf{v}_{i} \rangle), \qquad 0 \le i \le d.$$
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help?> Lar.simplexFacets

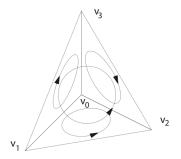


Figure 1: Coherent orientation of the facets of a 3-simplex

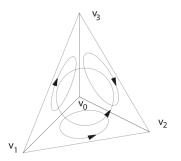


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### Implementation? See:

function Lar.simplexFacets(simplices::Cells)::Cells

## Simplicial prism

The prism over a simplex  $\sigma_d = \langle \mathbf{v}_0, \dots, \mathbf{v}_d \rangle$ , defined as the set  $P_{d+1} := \sigma_d \times [a, b]$ , with  $[a, b] \subset \mathbb{E}$ , will be called simplicial (d+1)-prism.

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$$\mathcal{K}_{d+1} = \{\sigma_{d+1,(i)} = (-1)^{id} \langle \mathbf{v}_i^a, \mathbf{v}_{i+1}^a, \dots, \mathbf{v}_d^a, \mathbf{v}_0^b, \mathbf{v}_1^b, \dots, \mathbf{v}_i^b \rangle | 0 \leq i \leq d \}$$

where  $\mathbf{v}_i^a = (\mathbf{v}_i, a)$  and  $\mathbf{v}_i^b = (\mathbf{v}_i, b)$ .

### Exercise: extract the boundary of implicial prism

#### extract the boundary of the simplicial cube

• write dimension-independent code

#### Hint

See and modify the code of Lar.simplexGrid

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#### extract the boundary of the simplicial cube

- write dimension-independent code
- apply to a grid of such decomposed cubes

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# Cellular Complexes

## Preliminary definitions

Let X be a topological space, and  $\Lambda(X) = \bigcup \Lambda_p$   $(p \in \{0, 1, ..., d\})$  a partition of X, with  $\Lambda_p$  a set of  $\{\text{(relatively) open}\}$ , connected, and manifold p-cells.

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Let X be a topological space, and  $\Lambda(X) = \bigcup \Lambda_p \ (p \in \{0, 1, \dots, d\})$  a partition of X, with  $\Lambda_p$  a set of  $\{\text{(relatively) open}\}$ , connected, and manifold p-cells.

A CW-structure on the space X is a filtration  $\emptyset = X_{-1} \subset X_0 \subset X_1 \subset \ldots \subset X_{d-1} \subset X = \bigcup_p X_p$ , such that, for each p, the skeleton  $X_p$  is homeomorphic to a space obtained from  $X_{p-1}$  by attachment of p-cells in  $\Lambda_p = \Lambda_p(X) \sim [?]$ .

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- A cellular *d*-complex *X* is orientable when its *d*-cells can be coherently oriented.
- The support space  $|\sigma|$  of a cell  $\sigma$  is its compact point-set.

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# Chain Complexes

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A p-chain of X with coefficients in G is a mapping  $c_p: X \to G$  such that, for each  $\sigma \in X_p$ , reversing a cell orientation changes the sign of the chain value:

$$c_p(-\sigma) = -c_p(\sigma).$$

aaaaa

Chain addition is defined by addition of chain values: if  $c_p^1, c_p^2$  are p-chains, then  $(c_p^1 + c_p^2)(\sigma) = c_p^1(\sigma) + c_p^2(\sigma)$ , for each  $\sigma \in X_p$ .

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Each chain can be written in a unique way as a sum of elementary chains.

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- Finally, a d-cell may be oriented as the sign of its oriented volume.

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To allow not only for chain addition, but also for linear combination of chains, coefficients should be taken from a set endowed with the structure of a field, such as  $(\mathbb{F},+,\times,0,1)$ ,

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Each chain can be written in a unique way as a linear combination of unit chains  $u \in U$ , if the outer cell is not taken into account.

Canonical representation and standard basis

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In particular,  $\#U_d = \#\Lambda_d - 1$ .

Often, with some abuse of notation, one does not distinguish between a *p*-cell and the corresponding unit *p*-chain.

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We call characteristic matrix M of a collection of subsets  $A_i \subseteq S$  (i = 1, ..., n) the binary matrix  $M = (m_{ij})$ , with  $m_{ij} = \chi_{A_i}(s_j)$ .

A matrix  $M_p$ , whose rows are indexed by unit p-chains and columns are indexed by unit 0-chains, provides a useful representation of a basis for the linear space  $C_p$ .

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While chains are mostly presented as formal sums of cells, in the actual implementation their signed coordinate vectors are used as sparse arrays, and in particular as CSC (Compressed Sparse Column) maps :  $\mathbb{N} \to \{-1,0,1\}.$ 

Implementation



Implementation

#### Implementation

How?

help?> Lar.characteristicMatrix

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The chain operators corresponding to the incidence relations are

$$VV: C_0 \to C_0$$
,

$$\mathcal{E}\mathcal{V}: \mathcal{C}_0 \to \mathcal{C}_1$$

$$\mathcal{FV}: \mathcal{C}_0 \to \mathcal{C}_2;$$

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# Characteristic matrices: Examples

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The corresponding CSR matrices are readily computed:

$$\begin{array}{l} \mathcal{V}\mathcal{V} = \mathcal{V}\mathcal{E} \circ \mathcal{E}\mathcal{V} = \mathcal{E}\mathcal{V}^{\top} \circ \mathcal{E}\mathcal{V} \Rightarrow [\mathcal{V}\mathcal{V}] = M_{1}^{t}M_{1} \\ \mathcal{V}\mathcal{E} = \mathcal{E}\mathcal{V}^{\top} \Rightarrow [\mathcal{V}\mathcal{E}] = M_{1}^{t} \\ \mathcal{V}\mathcal{F} = \mathcal{F}\mathcal{V}^{\top} \Rightarrow [\mathcal{V}\mathcal{F}] = M_{2}^{t} \\ \mathcal{E}\mathcal{V} \quad [\mathcal{E}\mathcal{V}] = M_{1} \\ \mathcal{E}\mathcal{E} = \mathcal{E}\mathcal{V} \circ \mathcal{V}\mathcal{E} = \mathcal{E}\mathcal{V} \circ \mathcal{E}\mathcal{V}^{\top} \Rightarrow [\mathcal{E}\mathcal{E}] = M_{1}M_{1}^{t} \\ \mathcal{E}\mathcal{F} = \mathcal{E}\mathcal{V} \circ \mathcal{V}\mathcal{F} = \mathcal{E}\mathcal{V} \circ \mathcal{F}\mathcal{V}^{\top} \Rightarrow [\mathcal{E}\mathcal{F}] = M_{1}M_{2}^{t} \\ \mathcal{F}\mathcal{V} \quad [\mathcal{F}\mathcal{V}] = M_{2} \\ \mathcal{F}\mathcal{E} = \mathcal{F}\mathcal{V} \circ \mathcal{V}\mathcal{E} = \mathcal{F}\mathcal{V} \circ \mathcal{E}\mathcal{V}^{\top} \Rightarrow [\mathcal{F}\mathcal{E}] = M_{2}M_{1}^{t} \\ \mathcal{F}\mathcal{F} = \mathcal{F}\mathcal{V} \circ \mathcal{V}\mathcal{F} = \mathcal{F}\mathcal{V} \circ \mathcal{F}\mathcal{V}^{\top} \Rightarrow [\mathcal{F}\mathcal{F}] = M_{2}M_{2}^{t}. \end{array}$$