20 Holdout Classes Among Skelet's 43 Hardly Non-Regular Turing Machines

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Abstract

github.com/danbriggs/Turing/blob/master/doc/record7-19-20.txt is where the reader is encouraged to go to get the state diagrams.

Introduction

In 2003, Skelet (Georgi Georgiev) released a 6218-line uncommented Pascal program "bbfind" that goes through the approximately 150 million (by the program's count) essentially distinct 5-state Turing machines and for each one attempts to either prove that it never halts when run on a blank tape or run it until it does. After a run of about two weeks, only 164 machines had stumped the program. Of these, 54 began with writing 0 to the tape, and so their runs were equivalent with an offset in number of steps to machines among the other 110. These 110 machines included 67 which were found to be "shift recursive" and so easily provable never to halt by hand.

Skelet dubbed the 43 remaining machines "Hardly non-regular," and although his work has never been independently verified, his list has remained the starting point for anyone interested in finishing the Busy Beaver of 5 ever since.

The 27th of these 43 machines was classified as "BL_2" by Skelet, and the other 42 were not classified. Note that for this reason, the machines I refer to as #s 28–43 are often referred to as 27–42 by other authors—including myself, previously.

From this list of 43, #2, #5, and #6 are equivalent with an offset in step number, #13 is equivalent to #12, #29 is equivalent to #23, #s 21, 28, and 39 are equivalent with an offset, and #41 is equivalent with an offset to #30; thus we are left with 36 essentially distinct Turing machines to study.

Of these 36, #s 16, 24, and 38 reach a phase (of at least 1 trillion steps) of equivalent action with an offset, but with several bits beyond what seems to be the working tape that are different; thus they have not been proved equivalent, but deserve to be talked about together. The situation is the same with machines #19 & #42. So instead of 36 machines, we talk of 33 different studies to be made.

Of these, I (and others independently) have discovered that machines 2 (and so 5 and 6), 14, 18, and 25 are trivially proved never to halt. I also proved machines #8, #9, #10, #11, and #12 (and so #13) never to halt. Machines #21 (and so #28 and #39), #31 and #32 were proved by univerz (Pavel Kropitz). univerz in fact showed the non-halting of twelve machines, but we need only reference these three for our reduction in number.

Thus there are 21 studies that need to be performed; thus this document is organized into 21 sections. Machine #4 was seen to be trivial in the course of producing this document, hence the title "20 Holdout Classes."

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The cleanest way to study HNR#3 is probably to inspect moments when the tape head is on the right; at these moments, the whole tape tends to be comprised of long runs of 1s with just a very few isolated 0s—often just one 0—in between. Ignoring step numbers when the tape head is on the right after it had also recently—within the past 10 steps—been at the right, we find that it achieves this configuration in state A with bit 1, state C with bit 1, state D with bit 1, or state C with bit 0—these results are for when the machine has reached the rightmost bit of the swath ever accessed, not just the current nonzero swath.

When it is in state A and there is only one 0 in between, it seems that both swaths of 1s, excluding the 1 at the tape head, may always be of even length. In order to get a handle on what it does by the time it reaches the right again afterwards, after leaving the right for at least 10 steps, I simulated runs starting from this type of configuration with two swaths of 1s of all possible even lengths up to 118, and recorded how many steps it took to get back to the right, how many swaths of 1s there were afterwards, how long the leftmost and rightmost of these were, and what state it ended up in—this time, "on the right" meaning at the rightmost 1. Tables 1, 2, 3, 4, and 5 give this data for swaths of 1s of even length ≤ 28; Table 6 shows moments when the tape head is on the right when started on a blank tape.

43	51	85	93	133	141	187	195	247	255	313	321	385	393	463
27	33	39	45	51	57	63	69	75	81	87	93	99	105	111
119	171	179	237	245	309	317	387	395	471	479	561	569	657	665
63	69	75	81	87	93	99	105	111	117	123	129	135	141	147
273	279	285	291	297	303	309	315	321	327	333	339	345	351	357
105	111	117	123	129	135	141	147	153	159	165	171	177	183	189
275	345	353	429	437	519	527	615	623	717	725	825	833	939	947
153	159	165	171	177	183	189	195	201	207	213	219	225	231	237
527	579	587	645	653	717	725	795	803	879	887	969	977	1065	1073
207	213	219	225	231	237	243	249	255	261	267	273	279	285	291
461	549	557	651	659	759	767	873	881	993	1001	1119	1127	1251	1259
267	273	279	285	291	297	303	309	315	321	327	333	339	345	351
849	855	861	867	873	879	885	891	897	903	909	915	921	927	933
333	339	345	351	357	363	369	375	381	387	393	399	405	411	417
677	783	791	903	911	1029	1037	1161	1169	1299	1307	1443	1451	1593	1601

Table 1: Number of steps taken to once again arrive at the rightmost 1 bit after leaving for at least 10 steps starting with configuration $1^{2m}01^{2n}i\ A$ for $0 \leq m, n \leq 14$. Here m is the row number and n is the column number, both indexed from 0.

2	2	2	2	2	2	2	2	2	2	2	2	2	2	2
1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
2	2	2	2	2	2	2	2	2	2	2	2	2	2	2
2	2	2	2	2	2	2	2	2	2	2	2	2	2	2
3	3	3	3	3	3	3	3	3	3	3	3	3	3	3
2	2	2	2	2	2	2	2	2	2	2	2	2	2	2
2	2	2	2	2	2	2	2	2	2	2	2	2	2	2
2	2	2	2	2	2	2	2	2	2	2	2	2	2	2
3	3	3	3	3	3	3	3	3	3	3	3	3	3	3
2	2	2	2	2	2	2	2	2	2	2	2	2	2	2
2	2	2	2	2	2	2	2	2	2	2	2	2	2	2
2	2	2	2	2	2	2	2	2	2	2	2	2	2	2
3	3	3	3	3	3	3	3	3	3	3	3	3	3	3
2	2	2	2	2	2	2	2	2	2	2	2	2	2	2
2	2	2	2	2	2	2	2	2	2	2	2	2	2	2

Table 2: Number of swaths of 1s produced upon arriving at the rightmost 1 bit after leaving for at least 10 steps starting with configuration $1^{2m}01^{2n}i\ A$ for $0 \leq m, n \leq 14$. Here m is the row number and n is the column number; both are indexed from 0. Note there seems to be no dependence on the number of 1s to the right of the 0.

A	C	A	C	A	C	A	C	A	C	A	C	A	C	A
D	D	D	D	D	D	D	D	D	D	D	D	D	D	D
C	A	C	A	C	A	C	A	C	A	C	A	C	A	C
D	D	D	D	D	D	D	D	D	D	D	D	D	D	D
D	D	D	D	D	D	D	D	D	D	D	D	D	D	D
D	D	D	D	D	D	D	D	D	D	D	D	D	D	D
C	A	C	A	C	A	C	A	C	A	C	A	C	A	C
D	D	D	D	D	D	D	D	D	D	D	D	D	D	D
C	A	C	A	C	A	C	A	C	A	C	A	C	A	C
D	D	D	D	D	D	D	D	D	D	D	D	D	D	D
C	A	C	A	C	A	C	A	C	A	C	A	C	A	C
D	D	D	D	D	D	D	D	D	D	D	D	D	D	D
D	D	D	D	D	D	D	D	D	D	D	D	D	D	D
D	D	D	D	D	D	D	D	D	D	D	D	D	D	D
C	A	C	A	C	A	C	A	C	A	C	A	C	A	C

Table 3: State of the machine upon arriving at the rightmost 1 bit after leaving for at least 10 steps starting with configuration $1^{2m}01^{2n}i\ A$ for $0 \leq m, n \leq 14$. Here m is the row number and n is the column number, both indexed from 0.

2	2	4	4	6	6	8	8	10	10	12	12	14	14	16
10	12	14	16	18	20	22	24	26	28	30	32	34	36	38
2	4	4	6	6	8	8	10	10	12	12	14	14	16	16
2	2	2	2	2	2	2	2	2	2	2	2	2	2	2
4	4	4	4	4	4	4	4	4	4	4	4	4	4	4
4	4	4	4	4	4	4	4	4	4	4	4	4	4	4
4	6	6	8	8	10	10	12	12	14	14	16	16	18	18
6	6	6	6	6	6	6	6	6	6	6	6	6	6	6
4	4	4	4	4	4	4	4	4	4	4	4	4	4	4
8	8	8	8	8	8	8	8	8	8	8	8	8	8	8
6	8	8	10	10	12	12	14	14	16	16	18	18	20	20
10	10	10	10	10	10	10	10	10	10	10	10	10	10	10
10	10	10	10	10	10	10	10	10	10	10	10	10	10	10
12	12	12	12	12	12	12	12	12	12	12	12	12	12	12
8	10	10	12	12	14	14	16	16	18	18	20	20	22	22

Table 4: Length of the leftmost swath of 1s produced upon arriving at the rightmost 1 bit after leaving for at least 10 steps starting with configuration $1^{2m}01^{2n}i\ A$ for $0 \leq m, n \leq 14$. Here m is the row number and n is the column number, both indexed from 0.

5	7	7	9	9	11	11	13	13	15	15	17	17	19	19
10	12	14	16	18	20	22	24	26	28	30	32	34	36	38
13	13	15	15	17	17	19	19	21	21	23	23	25	25	27
12	14	16	18	20	22	24	26	28	30	32	34	36	38	40
4	6	8	10	12	14	16	18	20	22	24	26	28	30	32
14	16	18	20	22	24	26	28	30	32	34	36	38	40	42
19	19	21	21	23	23	25	25	27	27	29	29	31	31	33
16	18	20	22	24	26	28	30	32	34	36	38	40	42	44
13	13	15	15	17	17	19	19	21	21	23	23	25	25	27
18	20	22	24	26	28	30	32	34	36	38	40	42	44	46
25	25	27	27	29	29	31	31	33	33	35	35	37	37	39
20	22	24	26	28	30	32	34	36	38	40	42	44	46	48
4	6	8	10	12	14	16	18	20	22	24	26	28	30	32
22	24	26	28	30	32	34	36	38	40	42	44	46	48	50
31	31	33	33	35	35	37	37	39	39	41	41	43	43	45

Table 5: Length of the rightmost swath of 1s produced upon arriving at the rightmost 1 bit after leaving for at least 10 steps starting with configuration $1^{2m}01^{2n}i\ A$ for $0 \leq m, n \leq 14$. Here m is the row number and n is the column number, both indexed from 0.

43	A	1	6	$1^{2}01^{4}i$
83	D	0	13	$1^{13}i$
268	A	1	18	$1^{8}01^{10}i$
572	D	2	30	$1^{4}01^{13}01^{13}i$
923	A	1	34	$1^{14}01^{20}i$
1137	D	1	41	$1^{6}01^{35}i$
1300	D	1	47	$1^{2}01^{45}i$
1457	D	0	53	$1^{53}i$
2512	A	1	58	$1^{28}01^{30}i$
4262	A	1	68	$1^{24}01^{44}i$
5244	D	2	80	$1^{10}01^{23}01^{47}i$
6881	C	1	84	$1^{36}01^{48}i$
7035	C	1	84	$1^{38}01^{46}o$
7731	D	1	89	$1^{18}01^{71}i$
8146	D	1	95	$1^{8}01^{87}i$
8675	D	2	106	$1^{4}01^{13}01^{89}i$
11762	A	1	110	$1^{52}01^{58}i$
16410	A	1	120	$1^{44}01^{76}i$
21486	C	1	130	$1^{50}01^{80}i$
21736	C	1	130	$1^{52}01^{78}o$
27456	C	1	138	$1^{52}01^{86}i$
27724	C	1	138	$1^{54}01^{84}o$
28930	D	1	143	$1^{26}01^{117}i$
29609	D	1	149	$1^{12}01^{137}i$
35900	C	1	158	$1^{72}01^{86}i$
36168	C	1	158	$1^{74}01^{84}o$
38004	D	1	163	$1^{36}01^{127}i$
46117	A	1	172	$1^{74}01^{98}i$
48005	D	1	179	$1^{36}01^{143}i$
57450	A	1	188	$1^{82}01^{106}i$
59656	D	1	195	$1^{40}01^{155}i$
61763	D	2	206	$1^{16}01^{33}01^{157}i$
71594	A	1	210	$1^{98}01^{112}i$
74478	D	1	217	$1^{48}01^{169}i$
86025	C	2	231	$1^{10}01^{109}01^{112}i$
86371	C	2	231	$1^{10}01^{111}01^{110}o$
99063	A	1	234	$1^{68}01^{166}i$
114997	A	1	244	$1^{102}01^{142}i$
118151	D	1	251	$1^{50}01^{201}i$
119604	D	1	257	$1^{24}01^{233}i$
121147	D	2	268	$1^{10}01^{23}01^{235}i$
137166	C	1	272	$1^{130}01^{142}i$
137602	C	1	272	$1^{132}01^{140}o$
159556	A	1	280	$1^{104}01^{176}i$
167144	D	2	292	$1^{40}01^{73}01^{179}i$
184393	C	1	296	$1^{132}01^{164}i$
184895	C	1	296	$1^{134}01^{162}o$
189755	D	1	301	$1^{66}01^{235}i$
191778	D	1	307	$1^{32}01^{275}i$
193873	D	3	323	$1^{4}01^{23}01^{19}01^{277}i$
214312	A	2	327	$1^{4}01^{165}01^{158}i$
240390	C	1	332	$1^{86}01^{246}i$
241138	C	1	332	$1^{88}01^{244}o$
247234	D	2	342	$1^{34}01^{63}01^{245}i$
272173	A	1	346	$1^{160}01^{186}i$
286901	D	3	363	$1^{16}01^{75}01^{83}01^{189}i$
307178	A	2	367	$1^{16}01^{173}01^{178}i$
338230	C	1	372	$1^{108}01^{264}i$
339032	C	1	372	$1^{110}01^{262}o$
342914	D	1	377	$1^{54}01^{323}i$
344841	D	1	383	$1^{26}01^{357}i$
346240	D	1	389	$1^{12}01^{377}i$
380071	C	1	398	$1^{192}01^{206}i$
380699	C	1	398	$1^{194}01^{204}o$
389825	D	1	403	$1^{96}01^{307}i$
396816	D	3	419	$1^{10}01^{49}01^{51}01^{309}i$
429039	A	2	423	$1^{10}01^{207}01^{206}i$
471335	C	1	428	$1^{116}01^{312}i$
472281	C	1	428	$1^{118}01^{310}o$
476709	D	1	433	$1^{58}01^{375}i$
478906	D	1	439	$1^{28}01^{411}i$
522833	A	1	448	$1^{214}01^{234}i$
533739	D	1	455	$1^{106}01^{349}i$
537694	D	1	461	$1^{52}01^{409}i$

Table 6: The step numbers when the tape head was at the rightmost bit of the swath ever accessed after leaving the right for at least 10 steps; what state the machine was in; how many intermediate zeros there were; how many 1s there were other than at the tape head; what the tape looked like. Run started on a blank tape.

Repeating the analysis done in Tables 1, 2, 3, 4, and 5, but with the machine starting in state C instead of A gives much simpler data: there are always two swaths; the number of steps and lengths of the left and right swaths are linear functions of the lengths of the original swaths; the machine ends up in state D as long as the right swath has at least six 1s. The machine starting in state D is exactly analogous to starting in state A, but the tape head is one bit further to the right and the step number is advanced by three. C with tape head at a 0 bit always comes from D with tape head at a 1 bit but otherwise of the same form.

Thus, it would suffice to understand the patterns shown in Tables 1, 2, 3, 4, and 5 to predict the action of HNR#3 from situations consisting of at most two swaths of 1s. However, situations with more than two swaths as shown in Table 6 may have to be understood separately, and could increase the difficulty of showing that HNR#3 never halts or accelerating it until it does.

Out of the first billion steps, the maximum number of swaths of 1s when the tape head is on the right is 6, and this first occurs at step 6333029, when the machine is in state D, and the tape is $1^4\ 0\ 1^{23}\ 0\ 1^{53}\ 0\ 1^{197}\ 0\ 1^{259}\ 0\ 1^{550}$.

An argument backtracking from a halt state may be capable of sidestepping all this detailed analysis; it is clear from the machine diagram

	A	B	C	D	E
0	C1L	H1L	D1R	A1R	C0L
1	A0R	E1L	B0L	C1R	D1L

that 01o A and 01o E are the only ways to achieve halting; these must come from 0o0 D, 010i B, respectively; the former is a dead-end, whereas the latter comes from 0101i C; unfortunately, this branches into three possibilities: 01011o A, 010o1 D, and 01011o E. These respectively come from 0101o0 D, 0100i E, and 010110i B; these, from 010i00 C, 01001i B, and 0101101i C;

4

Machine 4 seems to be a trivial one, the proof just needs to be carried out. It counts up in binary using 11 as 0 and 10 as 1, with less significant digits to the left, on the left side of the tape, with an additional o11 at the left end in C. Then after a number of steps increasing by a constant difference of 12 modified by four times the location in the number of the 0 to be turned into a 1 (increasing when that bit is *closer*), it produces the next number, written in the same binary scheme, with an extra 0 as padding on the right before its bookend 011. For example,

```
 882  C            o11101011101111111111111011  11.  Then after 145,
1027  C          o11111110101111111111111111011  12.  Then after 165 (+20 3/1),
1192  C        o11101110101111111111111111011  13.  Then after 173 (+8 1/2),
1365  C      o11111010101111111111111111111011  14.  Then after 189 (+16 2/1),
1554  C    o11101010101111111111111111111111011  15.  Then after 185 (-4 1/5),
1739  C  o11111111110111111111111111111111011  16.  Then after 213 (+28).
```

Table 7: Machine 4. The numbers before and after the slash portray which bit in is to be turned from a 0 to a 1 previously/currently. The difference between the two numbers determines the multiple of 4 offset from 12 more the step difference will be above the previous step difference.

7

33	A	$1^{2}0^{1}1^{2}$
57	C	1^{10}
188	D	$1^{9}0^{1}1^{6}$
276	C	$1^{20}0^{1}1^{2}$
355	C	1^{28}
746	A	$1^{19}0^{1}1^{14}$
810	A	$1^{16}0^{1}1^{16}$
1602	D	$1^{19}0^{1}1^{23}0^{1}1^{4}$
2234	A	$1^{35}0^{1}1^{16}$
2346	A	$1^{32}0^{1}1^{18}$
2640	C	$1^{48}0^{1}1^{8}$
3049	C	$1^{50}0^{1}1^{13}0^{1}1^{4}$
4426	D	$1^{39}0^{1}1^{32}$
5816	C	$1^{42}0^{1}1^{19}0^{1}1^{23}0^{1}1^{4}$
7139	D	$1^{41}0^{1}1^{47}0^{1}1^{4}$
9473	D	$1^{69}0^{1}1^{28}$
12905	A	$1^{65}0^{1}1^{42}$
13107	A	$1^{62}0^{1}1^{44}$
17183	A	$1^{73}0^{1}1^{42}$
17409	A	$1^{70}0^{1}1^{44}$
21971	A	$1^{77}0^{1}1^{46}$
22209	A	$1^{74}0^{1}1^{48}$
26915	A	$1^{65}0^{1}1^{61}0^{1}1^{10}$
27117	A	$1^{62}0^{1}1^{63}0^{1}1^{10}$
31493	D	$1^{95}0^{1}1^{44}$
37941	D	$1^{89}0^{1}1^{60}$
45261	A	$1^{99}0^{1}1^{60}$
45565	A	$1^{96}0^{1}1^{62}$
47041	C	$1^{134}0^{1}1^{30}$
47840	C	$1^{156}0^{1}1^{14}$
48453	C	$1^{170}0^{1}1^{6}$
49018	C	$1^{180}0^{1}1^{2}$
49577	C	1^{188}
57768	A	$1^{99}0^{1}1^{94}$
58072	A	$1^{96}0^{1}1^{96}$
64426	C	$1^{98}0^{1}1^{51}0^{1}1^{49}0^{1}1^{10}$
70891	D	$1^{101}0^{1}1^{101}0^{1}1^{10}$
81739	D	$1^{153}0^{1}1^{64}$
93763	A	$1^{113}0^{1}1^{97}0^{1}1^{23}0^{1}1^{4}$
94109	A	$1^{110}0^{1}1^{99}0^{1}1^{23}0^{1}1^{4}$
105757	D	$1^{155}0^{1}1^{81}0^{1}1^{4}$
120999	A	$1^{161}0^{1}1^{84}$
121489	A	$1^{158}0^{1}1^{86}$
123997	C	$1^{208}0^{1}1^{42}$
125270	C	$1^{236}0^{1}1^{20}$
141651	D	$1^{141}0^{1}1^{124}$
162381	A	$1^{173}0^{1}1^{102}$
162907	A	$1^{170}0^{1}1^{104}$
170467	C	$1^{172}0^{1}1^{73}0^{1}1^{40}$
186702	A	$1^{161}0^{1}1^{128}$
187192	A	$1^{158}0^{1}1^{130}$
191812	C	$1^{230}0^{1}1^{64}$
211699	A	$1^{151}0^{1}1^{135}0^{1}1^{23}0^{1}1^{4}$
212159	A	$1^{148}0^{1}1^{137}0^{1}1^{23}0^{1}1^{4}$
232791	A	$1^{213}0^{1}1^{99}0^{1}1^{4}$
233437	A	$1^{210}0^{1}1^{101}0^{1}1^{4}$
259509	D	$1^{207}0^{1}1^{112}$
284371	D	$1^{163}0^{1}1^{149}0^{1}1^{22}$
309195	A	$1^{233}0^{1}1^{106}$
309901	A	$1^{230}0^{1}1^{108}$
340719	A	$1^{205}0^{1}1^{142}$
341341	A	$1^{202}0^{1}1^{144}$
370731	A	$1^{177}0^{1}1^{155}0^{1}1^{28}$
371269	A	$1^{174}0^{1}1^{157}0^{1}1^{28}$
399339	D	$1^{245}0^{1}1^{118}$
403579	C	$1^{312}0^{1}1^{58}$
405584	C	$1^{348}0^{1}1^{28}$
438615	D	$1^{203}0^{1}1^{182}$
446809	C	$1^{302}0^{1}1^{90}$
449912	C	$1^{354}0^{1}1^{44}$
487277	A	$1^{219}0^{1}1^{188}$
487941	A	$1^{216}0^{1}1^{190}$
496785	C	$1^{318}0^{1}1^{94}$
500104	C	$1^{372}0^{1}1^{46}$
501965	C	$1^{402}0^{1}1^{22}$
503430	C	$1^{420}0^{1}1^{10}$
504787	C	$1^{432}0^{1}1^{4}$
546044	D	$1^{227}0^{1}1^{218}$
557280	C	$1^{344}0^{1}1^{108}$
608971	D	$1^{261}0^{1}1^{200}$

Table 8: Machine 7, tapehead at leftmost bit. Perhaps analysis of HNR#3 will serve HNR#7 as well.

15

31167	C	$i1^{0}$01111111010111111101010111110101110101010111011111111010101010101011	diff off
31203	C	$i1^{4}$01110101110111111101010111110101110101010111011111111010101010101011	+36 0
31267	C	$i1^{8}$01110111110111111101010111110101110101010111011111111010101010101011	+64 -8
31407	C	$i1^{12}$01110111011111111101010111110101110101010111011111111010101010101011	+140 -4
31703	C	$i1^{16}$01110111010101010110101111101011101010101110111111101010101010101011	+296 8
32271	C	$i1^{20}$01110111010110101110101111101011101010101110111111101010101010101011	+568 -8
33415	C	$i1^{24}$01110111010111110111010111110101110101010111011111111010101010101011	+1144 -8
35711	C	$i1^{28}$01110111010111111111010111110101110101010111011111111010101010101011	+2296 -8
40315	C	$i1^{32}$01110111010111111110110111110101110101010111011111111010101010101011	+4604 -4
49527	C	$i1^{36}$01110111010111111110101111110101110101010111011111111010101010101011	+9212 -4
67963	C	$i1^{40}$01110111010111111110101010101110111010101011101111111010101010101011	+18436 4
104819	C	$i1^{44}$01110111010111111110101011011101110101010111011111111010101010101011	+36856 -8
178539	C	$i1^{48}$01110111010111111110101011111101110101010111011111111010101010101011	+73720 -8
325991	C	$i1^{52}$01110111010111111110101011111011110101010111011111111010101010101011	+147452 -4
620903	C	$i1^{56}$01110111010111111110101011111010101101010111011111111010101010101011	+294912 0
1210719	C	$i1^{60}$01110111010111111110101011111010111101010111011111111010101010101011	+589816 -8
2390363	C	$i1^{64}$01110111010111111110101011111010111011010111011111111010101010101011	+1179644 -4
4749655	C	$i1^{68}$01110111010111111110101011111010111010110111011111111010101010101011	+2359292 -4
9468243	C	$i1^{72}$01110111010111111110101011111010111010101111011111111010101010101011	+4718588 -4
18905427	C	$i1^{76}$01110111010111111110101011111010111010101011111111111010101010101011	+9437184 0
37779787	C	$i1^{80}$01110111010111111110101011111010111010101011111111111010101010101011	+18874360 -8
75528531	C	$i1^{84}$01110111010111111110101011111010111010101011101010101110101010101011	+37748744 8
151025995	C	$i1^{88}$01110111010111111110101011111010111010101011101110101110101010101011	+75497464 -8
302020931	C	$i1^{92}$01110111010111111110101011111010111010101011101111101110101010101011	+150994936 -8
604010811	C	$i1^{96}$01110111010111111110101011111010111010101011101111111110101010101011	+301989880 -8
1207990583	C	$i1^{100}$01110111010111111110101011111010111010101011101111111011010101011	+603979772 -4
2415950131	C	$i1^{104}$01110111010111111110101011111010111010101011101111111010111010101011	+1207959548 -4
4831869231	C	$i1^{108}$01110111010111111110101011111010111010101011101111111010101110101011	+2415919100 -4
38654736675	C	$i1^{120}$01110111010111111110101011111010111010101011101111111010101010101111	+19327352828 -4

Table 9: Machine 15. Let i be the row in the table above, indexed from -1; let $x_i = 36 * 2^i$. Let d_i be the difference in step number between the ith row and the $i-1$st row, and let k_i be the place of the first 1 bit that was a 0 bit in the previous line, indexing from 0 just after the compressed portion. It is observed that $d_i - x_i$ relates to $k_i - 2i$ in the following linear fashion: $d_i - x_i = -16 + 2(k_i - 2i)$. Also, $k_i - 2i$ always seems to be 2, 4, 6, 8, or 10. Also, the positions of the new 1s, highlighted in red, seem to be injective and surjective among every other position throughout the swath other than at the very end. Furthermore, once all this work is done, only the second 0 and the last 0 have been changed, and everything else in its "workspace" is as at the beginning.

76761382320	C	$i1^{0}011111111111111111111111111101010101010111111111010111010101010111011111010111010101011111011101011101010101011111011101111101111101111$	diff
76761382396	C	$i1^{4}011101010101010101010101010111010101010111111111010111010101010111011111010111010101011111011101011101010101011111011101111101111101111$	+76
76761382460	C	$i1^{8}011101110101010101010101010111010101010111111111010111010101010111011111010111010101011111011101011101010101011111011101111101111101111$	+64
76761382596	C	$i1^{12}011101111010101010101010101110101010101111111111010111010101010111011111010111010101011111011101011101010101011111011101111101111101111$	+136
76761382876	C	$i1^{16}011101111110101010101010101110101010101111111111010111010101010111011111010111010101011111011101011101010101011111011101111101111101111$	+280
76761383444	C	$i1^{20}011101111111110101010101010111010101010111111111010111010101010111011111010111010101011111011101011101010101011111011101111101111101111$	+568
76761384588	C	$i1^{24}011101111111111101010101010111010101010111111111010111010101010111011111010111010101011111011101011101010101011111011101111101111101111$	+1144
76761386884	C	$i1^{28}011101111111111111010101010111010101010111111111010111010101010111011111010111010101011111011101011101010101011111011101111101111101111$	+2296
76761391484	C	$i1^{32}011101111111111111101010101110101010101111111111010111010101010111011111010111010101011111011101011101010101011111011101111101111101111$	+4600
76761400692	C	$i1^{36}011101111111111111111101010111010101010111111111010111010101010111011111010111010101011111011101011101010101011111011101111101111101111$	+9208
76761419116	C	$i1^{40}011101111111111111111111010111010101010111111111010111010101010111011111010111010101011111011101011101010101011111011101111101111101111$	+18424
76761455972	C	$i1^{44}011101111111111111111111111011101010101011111111010111010101010111011111010111010101011111011101011101010101011111011101111101111101111$	+36856
76761529692	C	$i1^{48}011101111111111111111111111110101010101011111111010111010101010111011111010111010101011111011101011101010101011111011101111101111101111$	+73720
76761677144	C	$i1^{52}011101111111111111111111111110110101010111111111010111010101010111011111010111010101011111011101011101010101011111011101111101111101111$	+147452
76761972052	C	$i1^{56}011101111111111111111111111101011010101111111111010111010101010111011111010111010101011111011101011101010101011111011101111101111101111$	+294908
76762561872	C	$i1^{60}011101111111111111111111111101010110101111111111010111010101010111011111010111010101011111011101011101010101011111011101111101111101111$	+589820
76763741516	C	$i1^{64}011101111111111111111111111101010101101111111111010111010101010111011111010111010101011111011101011101010101011111011101111101111101111$	+1179644
76766100808	C	$i1^{68}011101111111111111111111111101010101011111111111010111010101010111011111010111010101011111011101011101010101011111011101111101111101111$	+2359292
76770819412	C	$i1^{72}011101111111111111111111111101010101010101010101101110101010101110111110101110101011111011101011101010101011111011101111101111101111$	+4718604
76780256588	C	$i1^{76}011101111111111111111111111101010101010101101010111011101010101011101111101011101011111011101011101010101011111011101111101111101111$	+9437176
76799130948	C	$i1^{80}011101111111111111111111111101010101010111101011101110101010101110111110101110101011111011101011101010101011111011101111101111101111$	+18874360
76836879676	C	$i1^{84}011101111111111111111111111101010101010111111101111010101010101110111110101110101011111011101011101010101011111011101111101111101111$	+37748728
76912377140	C	$i1^{88}011101111111111111111111111101010101010111111111111010101010101110111110101110101011111011101011101010101011111011101111101111101111$	+75497464
77063372080	C	$i1^{92}011101111111111111111111111101010101010111111111011110101010101110111110101110101011111011101011101010101011111011101111101111101111$	+150994940
77365361968	C	$i1^{96}011101111111111111111111111101010101010111111111010101101010101110111110101110101011111011101011101010101011111011101111101111101111$	+301989888
77969341736	C	$i1^{100}011101111111111111111111111101010101010111111111010111110101010111011111010111010101011111011101011101010101011111011101111101111101111$	+603979768
79177301284	C	$i1^{104}011101111111111111111111111101010101010111111111010111011010101110111110101110101011111011101011101010101011111011101111101111101111$	+1207959548

Table 10: The machine's behavior after returning to the left at step 76761382305 B. It has set up a new task for itself in that it has written new stuff further to the right of where it had stuff previously, and it has gone back to having no 1s to the left of the workspace. It repeats its prior behavior, again taking nearly exactly $36 * 2^{n/4}$ steps to install a new 1 and return to the left, where n is the number of 1s after the i on the left (the first difference is significantly greater than 36 since the first new 1 to install is so far to the right), again injectively and bijectively ostensibly until $n/4$ is at or near 93, and then when $n = 94$ it would likely either write a new workspace for itself or do something new entirely. So in order to continue analyzing this machine effectively, an acceleration should be written that is able to automate away these $36 * 2^{93}$ or so steps. After all this work is done, seemingly only the second 0 and the last 0 would be altered, all the rest of the bits in the workspace becoming as they were before the work started.

Ostensibly the pack of 1s on the left is telling the machine to go to at least the $\frac{n}{4} + 2$nd even bit after the pack (indexing from 0 at the first bit after the pack) and turn the first 0 bit encountered from that point on into a 1, turning alternate 1 bits that it encounters between that index and that bit back into 0s. The most primitive way that a Turing machine could carry out a task like this of course takes exponential time in the number of 1 bits on the left, matching the pattern in step differences that we see, and this would also explain why the bits of the working tape end up nearly exactly as they were before the work began once the machine has accomplished this for each bit.

Conjecture. *For any string $x = a_0 \dots a_m$ with each of the a_j being* 10 *or* 11*, and any nonnegative integer $n \leq m$, let j be the least index $j \geq n$ such that $a_j = 10$; suppose such index exists. Then for any string y,*

$$0^\infty i 1^{4n} 011xy \ \mathtt{C} \vdash \ 0^\infty i 1^{4n+4} 011\overline{x}y \ \mathtt{C} \qquad 36 \cdot 2^n - 8 + 4(j-n),$$

where $\overline{x} = a_0 a_1 \dots a_{n-1} \overline{a_n}\,\overline{a_{n+1}} \dots \overline{a_j} a_{j+1} \dots a_m$, for $a = 10$, $\overline{a} = 11$, and vice versa.

The conjecture has been verified to produce each subsequent line of either of the two tables above from any given current line using the Python script found at X. Assuming the conjecture is true, one obtains

356526731314189519247709158424 C
i111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111/
111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111/
111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111/
111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111111/
101010101010111111111010111010101010111011111010111010101111101110101110101010101111011101111101111/
101110101010101010101010101010101010101010101010101010101010101010101111,

i.e., the machine runs for at least 356 octillion or $3.56 \cdot 10^{29}$ steps.

16∼24∼38

17

19∼42

20

22

23=29

26

BL_2

30=41

33

34

35

36

37

40

43