

PAUL'S ANSWERS TO LOGIC BONUS ASSIGNMENT 1.

20/04/2023.

1 a)  $\neg(\neg p \vee \neg(q \wedge p))$

$p\ q$	$\neg(\neg p \vee \neg(q \wedge p))$						
1 1	1	0	0	0	1	1	1
1 0	0	0	1	1	0	0	1
0 1	0	1	2	1	1	0	0
0 0	0	1	2	1	0	0	0

↑

Not a contradiction, not a tautology, but is satisfiable  
(when  $p=1, q=1$ )

Not a contradiction, not a tautology, but is satisfiable (when both p and q are true).

1 B)  $(\neg a \vee \neg b) \leftrightarrow \neg(a \wedge b)$

$a\ b$	$(\neg a \vee \neg b)$	$\leftrightarrow$	$\neg(a \wedge b)$
1 1	0	0	0
1 0	0	1	1
0 1	1	1	0
0 0	1	1	1

TAUTOLOGY, NOT A CONTRADICTION, Satisfiable.

1 C)  $\neg(a \rightarrow b) \rightarrow (a \rightarrow b)$

$a\ b$	$\neg(a \rightarrow b)$	$\rightarrow$	$(a \rightarrow b)$
1 1	0	1	1
1 0	1	0	0
0 1	0	1	1
0 0	0	1	0

NOT a Tautology  
NOT a Contradiction,  
But is satisfiable.

$$1 D) \neg b \rightarrow (a \rightarrow b)$$

True when  $b = 1$ , regardless of  $a$ 's value.

False when  $b = 0, a = 1$

True when  $b = 0, a = 0$

So Not a TAUTOLGY, not a CONTRADICTION, But Satisfiable.

$$2 A) a \models (b \wedge c) \rightarrow (a \rightarrow b)$$

$a \ b \ c$	$a \models (b \wedge c) \rightarrow (a \rightarrow b)$
1 1 1	1 : 1 : 1 : 1 : 1
1 1 0	1 : 1 : 0 : 1 : 1
1 0 1	1 : 1 : 0 : 1 : 0
1 0 0	1 : 1 : 0 : 1 : 0
0 1 1	:
0 1 0	:
0 0 1	:
0 0 0	:

VALID

INFERENCE.

$$2 B) \{p \vee q, q \vee r\} \models (p \wedge r) \rightarrow \neg q$$

$p \ q \ r$	$(p \vee q) \models (q \vee r) \models (p \wedge r) \models \neg q$
1 1 1	1 : 1 : 0 : 1 : 0 : 0
1 1 0	1 : 1 : 1 : 0 : 1 : 0
1 0 1	1 : 1 : 1 : 1 : 1 : 1
1 0 0	1 : 0 : 1 : 0 : 1 : 1
0 1 1	1 : 1 : 1 : 0 : 1 : 0
0 1 0	1 : 1 : 0 : 0 : 1 : 0
0 0 1	0 : 1 : 0 : 1 : 1 : 1
0 0 0	0 : 0 : 0 : 1 : 1 : 1

NOT VALID ( $p=1, q=1, r=1$ )

$$2C \quad \{ p \rightarrow q, q \rightarrow r \} \models p \rightarrow r$$

$p \ q \ r$	$p \rightarrow q$	$q \rightarrow r$	$\models$	$p \rightarrow r$
1 1 1	1	1	1	1
1 1 0	1	0		0
1 0 1	0	1		1
1 0 0	0	1		0
0 1 1	1	1	1	1
0 1 0	1	0		1
0 0 1	1	1	1	1
0 0 0	1	1	1	1

VALID INFERENCE ☺

3.

Propositional logic:

$$a \leftrightarrow b$$

$$b \leftrightarrow \neg c$$

$$\ast = c \leftrightarrow (a \wedge b \wedge \neg c) \vee (a \wedge \neg b \wedge c) \vee (\neg a \wedge b \wedge c) = \ast$$

No configuration possible:

if  $a$  is true,  $b$  must be true, so  $c$  must not be true, but it is.

$a$  is false,  $b$  must be false, so  $c$  must be true, but it is not

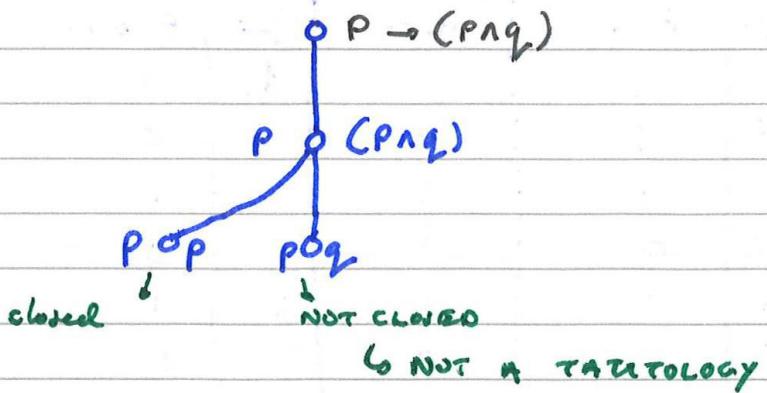
This problem is not satisfiable, it is a paradox.

$a \ b \ c$	$a \leftrightarrow b$	$b \leftrightarrow \neg c$	$\ast$
1 1 1	1	0	0
1 1 0	1	1	0
1 0 1	0	1	1
1 0 0	0	0	1

$a \ b \ c$	$a \leftrightarrow b$	$b \leftrightarrow \neg c$	$\ast$
0 1 1	0	0	1
0 1 0	0	1	1
0 0 1	1	1	1
0 0 0	1	0	1

4. Tabular method  $\rightarrow$  TAUTOLOGIES  $\rightarrow$  Logic in Action chp 8 Hickey  
ex 5 tick 60 + 71

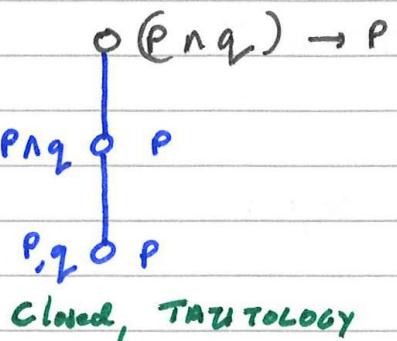
a)



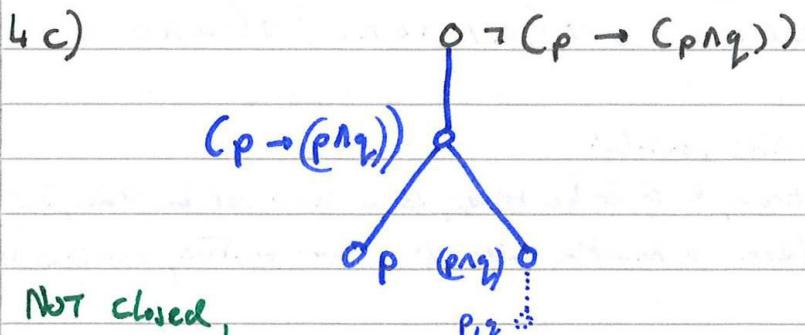
Counterexample:

$$P = 1, q = 0$$

4b)



4c)

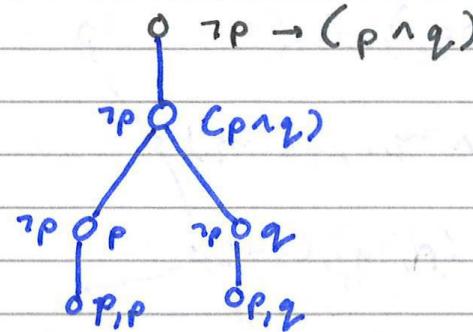


Counterexamples:

$$P = 0, q = 0 \text{ or } 1.$$

$$P = 1, q = 1.$$

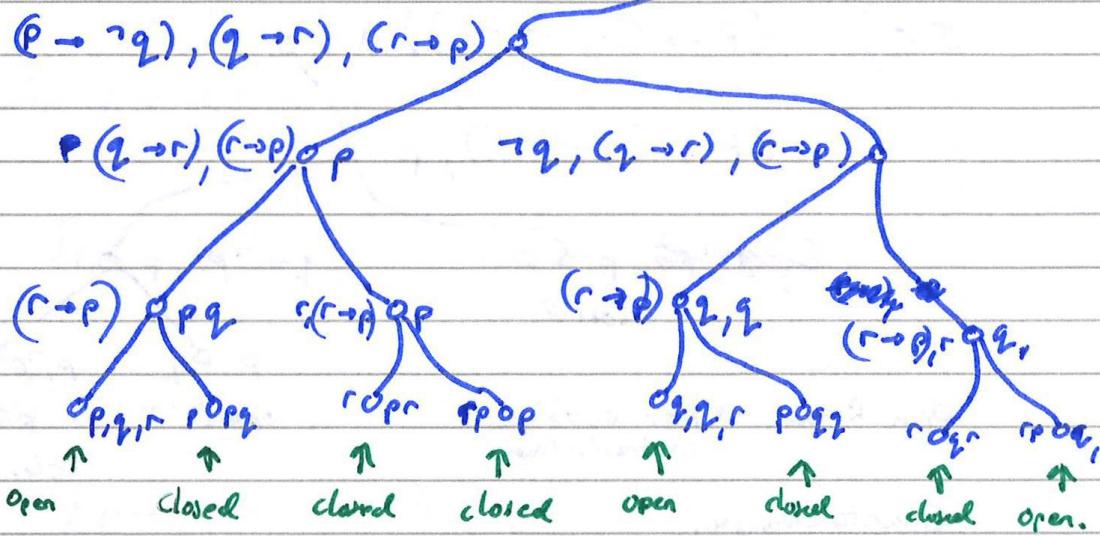
4 D



NOT CLOSED, NOT A TAUTOLOGY; Counterexample:  $p=0, q=0 \text{ or } 1.$

5 → Using the tableau method to ~~find~~ check for contradictions.

$$5a) (p \rightarrow \neg q) \wedge (q \rightarrow r) \wedge (r \rightarrow p) \circ$$



3 open branches → Not a contradiction!

Counterexamples:  $p=q=0, r=0,$

$p=1, q=0, r=0,$

$p=1, q=0, r=1.$

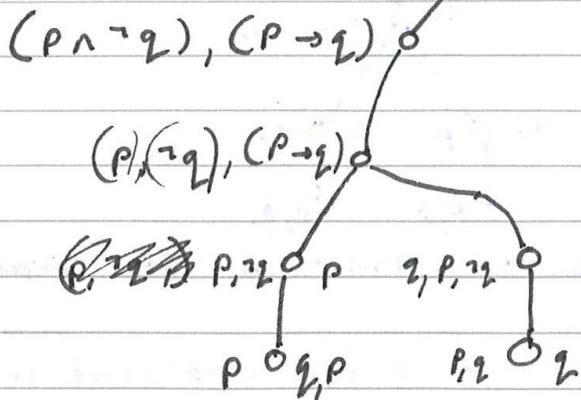
How did I find these? Look at the open branches!

$\circ p, q, r \Rightarrow r=0, q=0, r=0.$

$\circ q, r \Rightarrow q=0, r=0, p=0 \text{ or } 1$

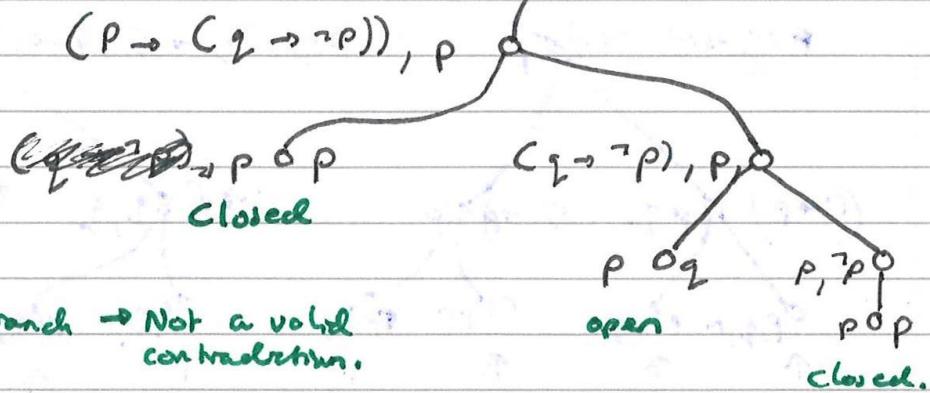
$\circ p, q \Rightarrow$

$$5B) (P \wedge \neg q) \wedge (P \rightarrow q) \circ$$



All are closed  
 $\Rightarrow$  Contradiction.

$$5C) (P \rightarrow (q \rightarrow \neg p)) \wedge P \circ$$



Open Branch  $\rightarrow$  Not a valid contradiction.

Counterexample:

$$P = 1, q = 0$$

Remember, left of circle = positive/1.  
 right of circle = negative/neutral/0.

$$\text{So } P \circ q \rightarrow P = 1, q = 0.$$

$\pi$   
 Counterexample.

5 What is wrong with the following formula?

$$(p \rightarrow q) \wedge (q \rightarrow p) \rightarrow (p \vee \neg q)$$

→ There is nothing wrong with the formula.

When  $p=1$  and  $q=1 \rightarrow$  True.

$p=1$  and  $q=0 \rightarrow$  True.

$p=0$  and  $q=1 \rightarrow$  True

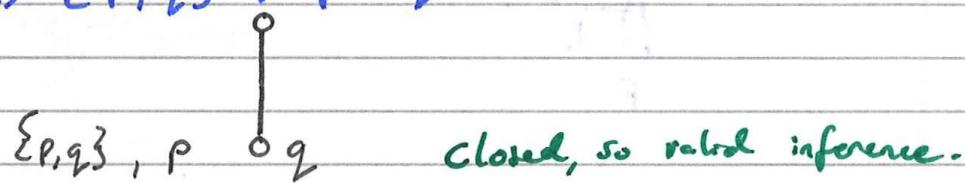
$p=0$  and  $q=0 \rightarrow$  True.

IT IS A TAUTOLGY.

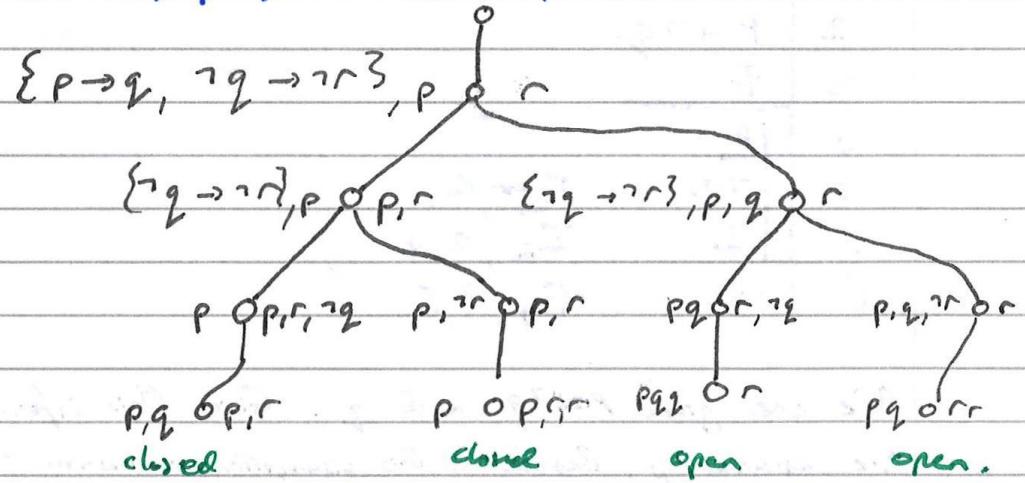
Alternative answer:  
The order in the formula  
is unclear. Is it  
 $((p \rightarrow q) \wedge (q \rightarrow p)) \rightarrow (p \vee \neg q)$   
-- OR --  
 $(p \rightarrow q) \wedge ((q \rightarrow p) \rightarrow (p \vee \neg q))$ ?

7. Use the tableau method to check if these inferences are valid.

A)  $\{p, q\} \models p \rightarrow q$

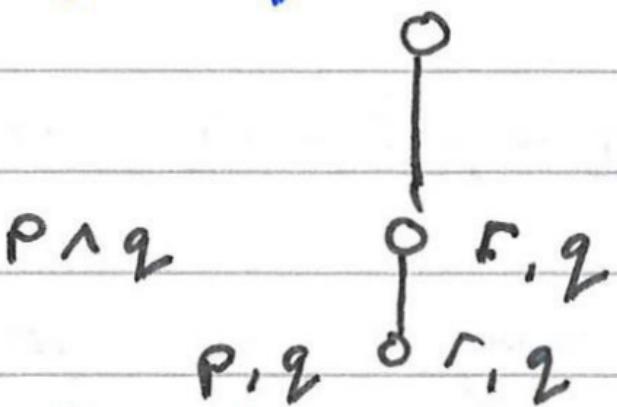


B)  $\{p \rightarrow q, \neg q \rightarrow \neg r\} \models p \rightarrow r$



Not a valid inference.

7. c)  $\{p \wedge q, \neg q\} \models r$



Closed, so valid inference.

PAUL'S SOLUTIONS TO LOGIC BONUS 1 2023.

8 A) From  $\neg\neg p$ , derive  $p$   $\neg\neg p \models p$

1	$\neg\neg p$	given
2	$\neg p$	assumption
3	$\perp$	
4	$p$	$E_{\neg} (2, 3)$ $I_{\neg} (2, 3)$

8 B) from  $p \rightarrow \neg q$  and  $q$ , derive  $\neg p$

1	$p \rightarrow \neg q$	$\neg q$	given
2	$q$		given
3	$p$		Assumption
4	$\neg q$		$E_{\neg} (2, 3)$
5	$\perp$		$I_{\perp} (2, 4)$
6	$\neg p$		$I_{\neg} (3, 5)$

Explanation: We are given  $p \rightarrow \neg q$  and  $q$ . Given this information, we assume  $p$ . Based on this assumption, we assume  $\neg q$ . Since  $p \rightarrow \neg q$  must hold, however, this contradicts line 2, where we assumed to be true. Thus, we derive  $\neg p$  to be true.

8 C) Derive  $p \vee \neg p$  (so  $\models p \vee \neg p$ )

1	$\neg p \vee p$	Assume Contrary.
2	$p$	Assume
3	$p \vee \neg p$	$E_{\vee} (2)$
4	$\perp$	$I_{\perp} (1, 4)$
5	$\neg p$	Assume, $I_{\neg} (2, 4)$
6	$\neg p \vee \neg p$	$E_{\vee} (5)$
7	$\perp$	$I_{\perp} (1, 6)$
8	$p \vee \neg p$	$I_{\vee} (1, 7)$

from  $P \vee \neg q$  derive  $q \rightarrow p$

8 D)	$\frac{P \vee \neg q}{\frac{\frac{P}{\frac{q \rightarrow p}{\frac{\neg q}{\frac{q}{\perp}}}}}{\frac{P}{q \rightarrow p}}}$	assume $\neg q \rightarrow (2)$ $\neg q \rightarrow (3, 4)$ $\neg q \rightarrow (5, 6)$ $\neg q \rightarrow (7, 8)$ $\neg q \rightarrow (1, 2, 3, 4, 5)$
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8 E) from  $P \rightarrow q$  derive  $\neg q \rightarrow \neg p$ .

1.	$\frac{P \rightarrow q}{\frac{\neg q}{\frac{P}{\perp}}}$	Given
2		Assume
3		Assume
4		$E \rightarrow (2, 3)$
5		$I \rightarrow (2, 4)$
6		$I \rightarrow (3, 5)$
7	$\neg q \rightarrow \neg p$	$I \rightarrow (2, 6)$

8 F)	$\frac{1}{\frac{\frac{\neg p \rightarrow \neg q}{\frac{q}{\frac{\neg q}{\frac{\neg p}{\perp}}}}}{\frac{q \rightarrow p}{\frac{\neg p \rightarrow \neg q}{\frac{\neg q}{\frac{\neg p}{\perp}}}}}}$	mp (1, 3) $I \perp (2, 4)$ $E \rightarrow (3-5.)$ $I \rightarrow (2, 6)$
8	$(\neg p \rightarrow \neg q) \rightarrow (q \rightarrow p)$	$I \rightarrow (1-7)$