Predicate Logic

Introduction

Statement	Translation
John walks	R_j
Bob walks	$R\,b$ ← object

↑ property

Synthax

- Symbols for constants (objects) ⟶ a, b, c
- Symbols vor variables ⟶ x, y, z
- Symbols for predicates ⟶ A, B, C
- Logical operators $\neg, \wedge, \vee, \rightarrow, \leftrightarrow$
- Quantifiers $\forall x$, $\exists x$

Syllogisms

- All A are B : $\forall x\,(Ax \rightarrow Bx)$
- Some A are B: $\exists x\,(Ax \wedge Bx)$
- All A are net B: $\forall x\,(Ax \rightarrow \neg Bx)$ $(\neg \exists x\,($

Relations

- John sees Mary ⟶ $S\,j\,m$
- John gives Mary the book ⟶ $G\,j\,m\,b$

Complex quantifier patterns

- Everyone sees someone $\forall x\,\exists y\,(Sxy)$
- Everyone is seen by someone $\forall x\,\exists y\,(Syx)$

Examples

Lxy	⟶	x loves y
Gx	⟶	x is girl
Bx	⟶	x is boy

Every boy loves a girl

$\forall x\,(Bx \rightarrow \exists y\,(Gy \wedge Lxy))$

Every girl who loves all boys does not love every girl:

$$\forall x\,((Gx \wedge \varphi(x)) \longrightarrow \psi(x))$$

Intuitive validities

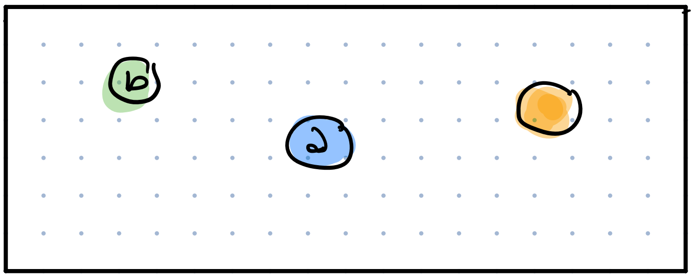
- $\neg \forall x \, \varphi x \equiv \exists x \, \neg \varphi x$

  any predicate

- $\neg \exists x \, \varphi x \equiv \forall x \, \neg \varphi x$
- $\forall x \, \varphi x \equiv \neg \exists x \, \neg \varphi x$
- $\exists x \, \varphi x \equiv \neg \forall x \, \neg \varphi x$
- $\neg \forall x (\varphi x \to \psi x) \equiv \exists x \, \neg (\varphi x \to \psi x)$

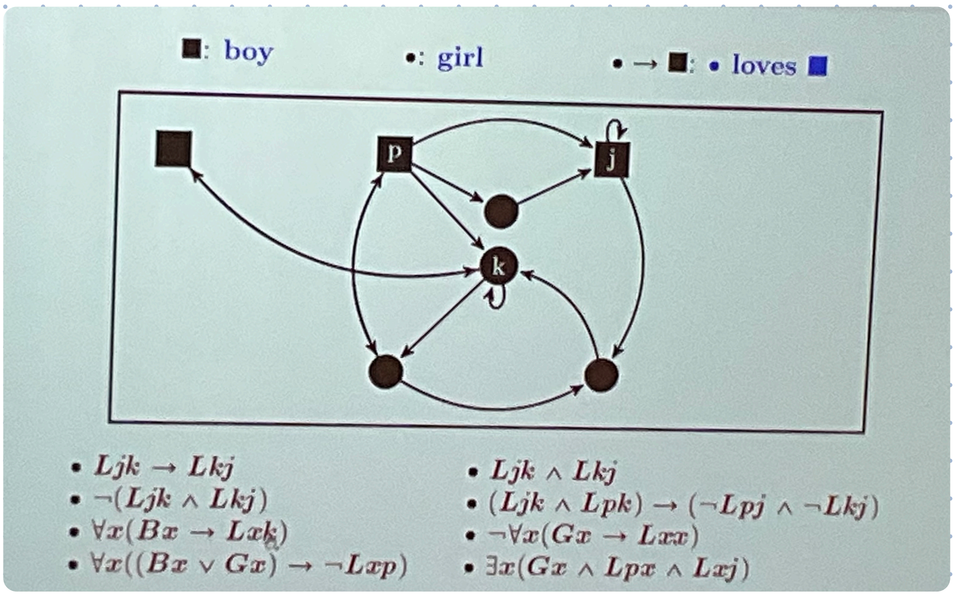
  $\equiv \exists x (\varphi x \wedge \neg \psi x)$
- $\forall x (\varphi x \wedge \psi x) \equiv \forall x \, \varphi x \wedge \forall x \, \psi x$

Evaluating formulas



Color is a property

- Ba — true (a is blue)
- $\exists x \, Sx \vee Cb$ (x is either square or circle)
- $Ra \to Sb$ true (a is not red)



Language

- Term t is a variable / constant
- Formula
  1. $P t_1 t_2$
  2. $\varphi \wedge \psi$
  3. $\forall x \varphi$

Substitution

- Inside a term → replacing the occurences of the variable y for the term t inside s

  $(s)^y_t$

  For a constant: $(c)^y_t := c$ $\qquad (a)^x_c = a$

  For a variable: $\begin{cases} (x)^y_t := x \\ (y)^y_t := t \end{cases}$

- Inside a formula

$$(Pt_1 \cdots t_n)^y_t := P(t_1)^y_t \cdots (t_n)^y_t$$
$$(\neg\varphi)^y_t := \neg(\varphi)^y_t$$
$$(\varphi \wedge \psi)^y_t := (\varphi)^y_t \wedge (\psi)^y_t$$
$$(\varphi \vee \psi)^y_t := (\varphi)^y_t \vee (\psi)^y_t$$
$$(\varphi \rightarrow \psi)^y_t := (\varphi)^y_t \rightarrow (\psi)^y_t$$
$$(\varphi \leftrightarrow \psi)^y_t := (\varphi)^y_t \leftrightarrow (\psi)^y_t$$

$$\begin{cases} (\forall x \varphi)^y_t := \forall x (\varphi)^y_t \\ (\forall y \varphi)^y_t := \forall y \varphi \end{cases}$$

$$\begin{cases} (\exists x \varphi)^y_t := \exists x (\varphi)^y_t \\ (\exists y \varphi)^y_t := \exists y \varphi \end{cases}$$

Models

A model is a tuple $\mathcal{M} = \langle D, I, g \rangle$

- D is the domain → non-empty collection of obj
- I interpretation function → assigns to each symbol
  - a relation