

## INTRODUCTION

A preemptible kernel creates new locking issues. The issues are the same as those under SMP: concurrency and reentrancy. Thankfully, the Linux preemptible kernel model leverages existing SMP locking mechanisms. Thus, the kernel requires explicit additional locking for very few additional situations.

This document is for all kernel hackers. Developing code in the kernel requires protecting these situations.

**RULE #1:** Per-CPU data structures need explicit protection

Two similar problems arise. An example code snippet:

```
struct this_needs_locking tux[NR_CPUS];
tux[smp_processor_id()] = some_value;
/* task is preempted here... */
something = tux[smp_processor_id()];
```

First, since the data is per-CPU, it may not have explicit SMP locking, but require it otherwise. Second, when a preempted task is finally rescheduled, the previous value of `smp_processor_id` may not equal the current. You must protect these situations by disabling preemption around them.

You can also use `put_cpu()` and `get_cpu()`, which will disable preemption.

**RULE #2:** CPU state must be protected.

Under preemption, the state of the CPU must be protected. This is arch-dependent, but includes CPU structures and state not preserved over a context switch. For example, on x86, entering and exiting FPU mode is now a critical section that must occur while preemption is disabled. Think what would happen if the kernel is executing a floating-point instruction and is then preempted. Remember, the kernel does not save FPU state except for user tasks. Therefore, upon preemption, the FPU registers will be sold to the lowest bidder. Thus, preemption must be disabled around such regions.

Note, some FPU functions are already explicitly preempt safe. For example, `kernel_fpu_begin` and `kernel_fpu_end` will disable and enable preemption. However, `math_state_restore` must be called with preemption disabled.

**RULE #3:** Lock acquire and release must be performed by same task

A lock acquired in one task must be released by the same task. This means you can't do oddball things like acquire a lock and go off to play while another task releases it. If you want to do something like this, acquire and release the task in the same code path and have the caller wait on an event by the other task.

## SOLUTION

Data protection under preemption is achieved by disabling preemption for the duration of the critical region.

|  |   |
|--|---|
| <code>preempt_enable()</code>            | decrement the preempt counter             |
| <code>preempt_disable()</code>           | increment the preempt counter             |
| <code>preempt_enable_no_resched()</code> | decrement, but do not immediately preempt |
| <code>preempt_check_resched()</code>     | if needed, reschedule                     |
| <code>preempt_count()</code>             | return the preempt counter                |

The functions are nestable. In other words, you can call `preempt_disable` *n*-times in a code path, and preemption will not be reenabled until the *n*-th call to `preempt_enable`. The preempt statements define to nothing if preemption is not enabled.

Note that you do not need to explicitly prevent preemption if you are holding any locks or interrupts are disabled, since preemption is implicitly disabled in those cases.

But keep in mind that 'irqs disabled' is a fundamentally unsafe way of disabling preemption – any `spin_unlock()` decreasing the preemption count to 0 might trigger a reschedule. A simple `printk()` might trigger a reschedule. So use this implicit preemption-disabling property only if you know that the affected codepath does not do any of this. Best policy is to use this only for small, atomic code that you wrote and which calls no complex functions.

Example:

```

cpucache_t *cc; /* this is per-CPU */
preempt_disable();
cc = cc_data(searchp);
if (cc && cc->avail) {
    __free_block(searchp, cc_entry(cc), cc->avail);
    cc->avail = 0;
}
preempt_enable();
return 0;

```

Notice how the preemption statements must encompass every reference of the critical variables. Another example:

```

int buf[NR_CPUS];
set_cpu_val(buf);
if (buf[smp_processor_id()] == -1) printf(KERN_INFO "wee!\n");
spin_lock(&buf_lock);
/* ... */

```

preempt-locking.txt

This code is not preempt-safe, but see how easily we can fix it by simply moving the spin\_lock up two lines.

## PREVENTING PREEMPTION USING INTERRUPT DISABLING

It is possible to prevent a preemption event using `local_irq_disable` and `local_irq_save`. Note, when doing so, you must be very careful to not cause an event that would set `need_resched` and result in a preemption check. When in doubt, rely on locking or explicit preemption disabling.

Note in 2.5 interrupt disabling is now only per-CPU (e.g. local).

An additional concern is proper usage of `local_irq_disable` and `local_irq_save`. These may be used to protect from preemption, however, on exit, if preemption may be enabled, a test to see if preemption is required should be done. If these are called from the `spin_lock` and `read/write lock` macros, the right thing is done. They may also be called within a spin-lock protected region, however, if they are ever called outside of this context, a test for preemption should be made. Do note that calls from interrupt context or bottom half/ tasklets are also protected by preemption locks and so may use the versions which do not check preemption.