What is hwpoison?

Upcoming Intel CPUs have support for recovering from some memory errors (`MCA recovery'). This requires the OS to declare a page "poisoned", kill the processes associated with it and avoid using it in the future.

This patchkit implements the necessary infrastructure in the VM.

To quote the overview comment:

- * High level machine check handler. Handles pages reported by the * hardware as being corrupted usually due to a 2bit ECC memory or cache * failure.
- * This focusses on pages detected as corrupted in the background.

 * When the current CPU tries to consume corruption the currently

 * running process can just be killed directly instead. This implies

 * that if the error cannot be handled for some reason it's safe to

 * just ignore it because no corruption has been consumed yet. Instead
- * when that happens another machine check will happen.
- * Handles page cache pages in various states. The tricky part
 * here is that we can access any page asynchronous to other VM
 * users, because memory failures could happen anytime and anywhere,
 * possibly violating some of their assumptions. This is why this code
 * has to be extremely careful. Generally it tries to use normal locking
 * rules, as in get the standard locks, even if that means the
 * error handling takes potentially a long time.
- * Some of the operations here are somewhat inefficient and have non
 * linear algorithmic complexity, because the data structures have not
 * been optimized for this case. This is in particular the case
 * for the mapping from a vma to a process. Since this case is expected
 * to be rare we hope we can get away with this.

The code consists of a the high level handler in mm/memory-failure.c, a new page poison bit and various checks in the VM to handle poisoned pages.

The main target right now is KVM guests, but it works for all kinds of applications. KVM support requires a recent qemu-kvm release.

For the KVM use there was need for a new signal type so that KVM can inject the machine check into the guest with the proper address. This in theory allows other applications to handle memory failures too. The expection is that near all applications won't do that, but some very specialized ones might.

There are two (actually three) modi memory failure recovery can be in:

vm. memory failure recovery sysctl set to zero:

All memory failures cause a panic. Do not attempt recovery. (on x86 this can be also affected by the tolerant level of the MCE subsystem)

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early kill

(can be controlled globally and per process)
Send SIGBUS to the application as soon as the error is detected
This allows applications who can process memory errors in a gentle
way (e.g. drop affected object)
This is the mode used by KVM gemu.

late kill

Send SIGBUS when the application runs into the corrupted page. This is best for memory error unaware applications and default Note some pages are always handled as late kill.

User control:

vm.memory_failure_recovery See sysctl.txt

vm.memory_failure_early_kill Enable early kill mode globally

PR MCE KILL

Set early/late kill mode/revert to system default argl: PR_MCE_KILL_CLEAR: Revert to system default argl: PR_MCE_KILL_SET: arg2 defines thread specific mode PR_MCE_KILL_EARLY: Early kill PR_MCE_KILL_LATE: Late kill PR_MCE_KILL_DEFAULT: Use system global default

PR MCE KILL GET

return current mode

Testing:

madvise(MADV_HWPOISON,)
(as root)
Poison a page in the process for testing

hwpoison-inject module through debugfs

/sys/debug/hwpoison/

corrupt-pfn

Inject hwpoison fault at PFN echoed into this file. This does some early filtering to avoid corrupted unintended pages in test suites.

unpoison-pfn

Software-unpoison page at PFN echoed into this file. This way a page can be reused again.

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This only works for Linux injected failures, not for real memory failures.

Note these injection interfaces are not stable and might change between kernel versions

corrupt-filter-dev-major
corrupt-filter-dev-minor

Only handle memory failures to pages associated with the file system defined by block device major/minor. -1U is the wildcard value. This should be only used for testing with artificial injection.

corrupt-filter-memcg

Limit injection to pages owned by memgroup. Specified by inode number of the memcg.

Example:

mkdir /cgroup/hwpoison

usemem -m 100 -s 1000 & echo jobs -p > /cgroup/hwpoison/tasks

memcg_ino=\$(ls -id /cgroup/hwpoison | cut -f1 -d' ')
echo \$memcg_ino > /debug/hwpoison/corrupt-filter-memcg

page-types -p pidof init --hwpoison # shall do nothing page-types -p pidof usemem --hwpoison # poison its pages

corrupt-filter-flags-mask corrupt-filter-flags-value

When specified, only poison pages if ((page_flags & mask) == value). This allows stress testing of many kinds of pages. The page_flags are the same as in /proc/kpageflags. The flag bits are defined in include/linux/kernel-page-flags.h and documented in Documentation/vm/pagemap.txt

Architecture specific MCE injector

x86 has mce-inject, mce-test

Some portable hwpoison test programs in mce-test, see blow.

References:

http://halobates.de/mce-lc09-2.pdf Overview presentation from LinuxCon 09

git://git.kernel.org/pub/scm/utils/cpu/mce/mce-inject.git 第 3 页

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x86 specific injector

Limitations:

Not all page types are supported and never will. Most kernel internal objects cannot be recovered, only LRU pages for now.
Right now hugepage support is missing.

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