

<previous description obsolete, deleted>

Virtual memory map with 4 level page tables:

000000000000000000 - 00007fffffffffffff (=47 bits) user space, different per mm  
hole caused by [48:63] sign extension  
ffff800000000000 - ffff80ffffffffffff (=40 bits) guard hole  
ffff880000000000 - ffffc7ffffffffffff (=64 TB) direct mapping of all phys. memory  
ffffc80000000000 - ffffc8ffffffffffff (=40 bits) hole  
ffffc90000000000 - ffffe8ffffffffffff (=45 bits) vmalloc/ioremap space  
ffffe90000000000 - ffffe9ffffffffffff (=40 bits) hole  
fffffea00000000000 - ffffeaffffffffffffff (=40 bits) virtual memory map (1TB)  
... unused hole ...  
ffffffff80000000 - ffffffffafa0000000 (=512 MB) kernel text mapping, from phys 0  
ffffffffffa0000000 - ffffffffffffff000000 (=1536 MB) module mapping space

The direct mapping covers all memory in the system up to the highest memory address (this means in some cases it can also include PCI memory holes).

vmalloc space is lazily synchronized into the different PML4 pages of the processes using the page fault handler, with `init_level4_pgt` as reference.

Current X86-64 implementations only support 40 bits of address space, but we support up to 46 bits. This expands into MBZ space in the page tables.

-Andi Kleen, Jul 2004