USING JACOBI PLANE ROTATIONS IN R

JAN DE LEEUW

ABSTRACT. Matrix techniques for various types of diagonalizations of matrices and three-dimensinal arrays are implemented in \mathbb{R} using Jacobi plane rotations.

1. Introduction

Many problems in multivariate analysis can be formulated as optimizing a function f(K) over the rotation matrices of order n, i.e. the square matrices that satisfy K'K = KK' = I. Or, more generally, to optimize a function $f(K_1, \dots, K_n)$ over a number of rotation matrices.

Recently, much interesting theory has been developed on how to design gradient type optimization methods for such problems, using the differential geometry of Grassman and Stiefel manifolds [Edelman et al., 1998]. (Edelman, Book – gradient methods).

In this paper we go a different route, however. We use the rotation matrix version of one-dimensional coordinate-wise optimization, by building up the optimal rotation matrices iteratively from a sequence of one-parameter plane rotations.

2. PLANE ROTATIONS

Two-by-two rotation matrices can be written in the one-parameter form

(1)
$$K(\theta) = \begin{bmatrix} \cos(\theta) & \sin(\theta) \\ -\sin(\theta) & \cos(\theta) \end{bmatrix}.$$

Note that $K(\theta)' = K(-\theta)$ and K(0) = I.

Jacobi plane rotations of order n are the matrices $K_{ij}(\theta)$, which are equal to the identity matrix of order n, but with elements (i, j), (j, j), (i, j) and (j, i) replaced

Date: Wednesday 10th December, 2008 — 18h 5min — Typeset in TIMES ROMAN.

by the four elements of $K(\theta)$. We use s_{θ} and c_{θ} as abbreviations for $\sin(\theta)$ and $\cos(\theta)$. Thus

$$\{K_{ij}(\theta)\}_{ii} = c_{\theta},$$

$$\{K_{ij}(\theta)\}_{ij} = s_{\theta},$$

$$\{K_{ij}(\theta)\}_{ji} = -s_{\theta},$$

$$\{K_{ij}(\theta)\}_{ij} = c_{\theta}.$$

Now
$$K_{ij}(\theta)' = K_{ii}(\theta) = K_{ij}(-\theta)$$
, and again $K_{ij}(0) = I$.

Suppose X is a rectangular matrix with n rows and m columns. Then $\tilde{X} = K_{ij}(\theta)X$ differs from X only in row i and row j. We have

(3a)
$$\tilde{x}_{\alpha\beta} = \begin{cases} c_{\theta}x_{i\beta} + s_{\theta}x_{j\beta} & \text{if } \alpha = i, \\ -s_{\theta}x_{i\beta} + c_{\theta}x_{j\beta} & \text{if } \alpha = j, \\ x_{\alpha\beta} & \text{otherwise.} \end{cases}$$

Similarly $\tilde{X} = XK_{k\ell}(\xi)$ differs from X only in columns k and ℓ , with

(3b)
$$\tilde{x}_{\alpha\beta} = \begin{cases} c_{\xi}x_{\alpha j} - s_{\xi}x_{\alpha\ell} & \text{if } \beta = j, \\ s_{\xi}x_{\alpha j} + c_{\xi}x_{\alpha\ell} & \text{if } \beta = \ell, \\ x_{\alpha\beta} & \text{otherwise.} \end{cases}$$

Equations (3a) and (3b) are one-sided left and right Jacobi transformations of X. The two-sided transform is defined by $\tilde{X} = K_{ij}(\theta)XK_{k\ell}(\xi)$. We find

$$\begin{cases}
c_{\xi}x_{\alpha k} - s_{\xi}x_{\alpha \ell} & \text{if } \beta = k \text{ and } \alpha \neq i, j, \\
s_{\xi}x_{\alpha k} + c_{\xi}x_{\alpha \ell} & \text{if } \beta = \ell \text{ and } \alpha \neq i, j, \\
c_{\theta}x_{i\beta} + s_{\theta}x_{j\beta} & \text{if } \alpha = i \text{ and } \beta \neq k, l, \\
-s_{\theta}x_{i\beta} + c_{\theta}x_{j\beta} & \text{if } \alpha = j \text{ and } \beta \neq k, l, \\
c_{\theta}c_{\xi}x_{ik} + s_{\theta}c_{\xi}x_{jk} - c_{\theta}s_{\xi}x_{i\ell} - s_{\theta}s_{\xi}x_{j\ell} & \text{if } \alpha = i \text{ and } \beta = k, \\
c_{\theta}s_{\xi}x_{ik} + s_{\theta}s_{\xi}x_{jk} + c_{\theta}c_{\xi}x_{i\ell} + s_{\theta}c_{\xi}x_{j\ell} & \text{if } \alpha = i \text{ and } \beta = \ell, \\
-s_{\theta}c_{\xi}x_{ik} + c_{\theta}c_{\xi}x_{jk} + s_{\theta}s_{\xi}x_{i\ell} - c_{\theta}s_{\xi}x_{j\ell} & \text{if } \alpha = j \text{ and } \beta = k, \\
-s_{\theta}s_{\xi}x_{ik} + c_{\theta}s_{\xi}x_{jk} - s_{\theta}c_{\xi}x_{i\ell} + c_{\theta}c_{\xi}x_{j\ell} & \text{if } \alpha = j \text{ and } \beta = \ell, \\
x_{\alpha\beta} & \text{otherwise.}
\end{cases}$$

Of course the one-sided transforms can be recovered by setting either $\theta=0$ or $\xi=0$.

3. CYCLES

In our algorithms we update the current rotation matrix K of order n by applying $\frac{1}{2}n(n-1)$ plane rotations $K_{ij}(\theta)$ for all i < j, where θ is chosen to optimize some criterion. In all our applications the criterion is a sum of squares of certain matrix or array elements.

Usually we update from the left, i.e. we set $\tilde{K} = K_{ij}(\hat{\theta})K$, with the optimally chosen $\hat{\theta}$. If the function we are optimizing contains a term KX, for some fixed matrix X, then of course $\tilde{X} = K_{ij}(\hat{\theta})KX$. It is usually convenient to update X in situ, which can be done by just a few multiplications and additions. Because the criterion is a sum of squares, the optimal $\sin(\hat{\theta})$ and $\cos(\hat{\theta})$ can usually be found by solving a 2×2 eigenvalue problem for the eigenvector corresponding with the smallest or largest eigenvalue.

In the symmetric matrix case we use two-sided updates. Because of symmetry we still only deal with a one parameter problem, and trigonometric identities can be used to again reduce finding the optimal plane rotation to solving a 2×2 eigenvalue problem. Each cycle of the algorithm again passes through all $\frac{1}{2}n(n-1)$ in order.

In the literature there are many variations on how to cycle through the plane rotations. We could choose the largest element contributing to the sum of squares, or we only decide to rotate if an element is above a certain threshold (and then lower the threshold in further cycles). This may be important for really large examples, but our \mathbb{R} code is intended for fairly small ones. In future versions of our work we will optimize by code by translating the loops in the core of the algorithm to \mathbb{C} .

4. APPLICATIONS

4.1. **Eigenvalues.** In the classical eigenvalue problem we have a symmetric matrix A, and we want to find a rotation matrix K such that K'AK is diagonal. In the classical Jacobi method we build up K by successive plane rotations. Thus the matrix X is symmetric, and we use the symmetric two-sided transform

$$\tilde{A} = K_{ij}(\xi)' A K_{ij}(\xi)$$

In the formulas from the previous section we have $\theta = -\xi$. Using this, and the symmetry of A, we find

$$\tilde{a}_{\alpha\beta} = \begin{cases}
c_{\xi}a_{\alpha i} - s_{\xi}a_{\alpha j} & \text{if } \beta = i \text{ and } \alpha \neq i, j, \\
s_{\xi}a_{\alpha i} + c_{\xi}a_{\alpha j} & \text{if } \beta = j \text{ and } \alpha \neq i, j, \\
c_{\xi}a_{i\beta} - s_{\xi}a_{j\beta} & \text{if } \alpha = i \text{ and } \beta \neq i, j, \\
s_{\xi}a_{i\beta} + c_{\xi}a_{j\beta} & \text{if } \alpha = j \text{ and } \beta \neq i, j, \\
c_{\xi}^{2}a_{ii} - 2s_{\xi}c_{\xi}a_{ij} + s_{\xi}^{2}a_{jj} & \text{if } \alpha = i \text{ and } \beta = i, \\
s_{\xi}c_{\xi}(a_{ii} - a_{jj}) + (c_{\xi}^{2} - s_{\xi}^{2})a_{ij} & \text{if } \alpha = i \text{ and } \beta = j, \\
s_{\xi}c_{\xi}(a_{ii} - a_{jj}) + (c_{\xi}^{2} - s_{\xi}^{2})a_{ij} & \text{if } \alpha = j \text{ and } \beta = i, \\
s_{\xi}^{2}a_{ii} + 2s_{\xi}c_{\xi}a_{ij} + c_{\xi}^{2}a_{jj} & \text{if } \alpha = j \text{ and } \beta = j, \\
a_{\alpha\beta} & \text{otherwise.}
\end{cases}$$

Now

$$(c_{\xi}a_{\alpha i} - s_{\xi}a_{\alpha j})^{2} + (s_{\xi}a_{\alpha i} + c_{\xi}a_{\alpha j})^{2} = a_{\alpha i}^{2} + a_{\alpha j}^{2},$$
$$(c_{\xi}a_{i\beta} - s_{\xi}a_{j\beta})^{2} + (s_{\xi}a_{i\beta} + c_{\xi}a_{j\beta})^{2} = a_{i\beta}^{2} + a_{i\beta}^{2},$$

and thus

$$\sum_{lpha < eta} ilde{a}_{lpha eta}^2 = \sum_{lpha < eta} a_{lpha eta}^2 - a_{ij}^2 + ilde{a}_{ij}^2.$$

We minimize the sum of squares of the off-diagonal elements of \tilde{A} by solving

$$\tilde{a}_{ij} = s_{\xi} c_{\xi} (a_{ii} - a_{jj}) + (c_{\xi}^2 - s_{\xi}^2) a_{ij} = \sin(2\xi) d_{ij} + \cos(2\xi) a_{ij} = 0,$$

with $d_{ij} = \frac{1}{2}(a_{ii} - a_{jj})$. One solution for the vector $(\sin(2\xi), \cos(2\xi))$ is

$$u = \frac{1}{\sqrt{a_{ij}^2 + d_{ij}^2}} \begin{bmatrix} a_{ij} \\ -d_{ij} \end{bmatrix}.$$

We now solve for any s_{ξ} and c_{ξ} such that $2s_{\xi}c_{\xi}=u_1$, $c_{\xi}^2-s_{\xi}^2=u_2$, and $c_{\xi}^2+s_{\xi}^2=1$. Thus

$$c_{\xi} = \sqrt{\frac{1+u_2}{2}},$$

$$s_{\xi} = \mathbf{sign}(u_1)\sqrt{\frac{1-u_2}{2}}.$$

Note that we can also update the matrix of eigenvectors K by starting with K = I, and by updating to \tilde{K} after each plane rotation using

(6)
$$\tilde{k}_{\alpha\beta} = \begin{cases} c_{\xi}k_{\alpha i} - s_{\xi}k_{\alpha j} & \text{if } \beta = i, \\ s_{\xi}k_{\alpha i} + c_{\xi}k_{\alpha j} & \text{if } \beta = j, \\ k_{\alpha\beta} & \text{otherwise.} \end{cases}$$

The code for the algorithm is given in the Appendix. Note that we have implemented the cyclic Jacobi method, without any searching for a largest pivot element. In applying our algorithm to various matrices, we do see the fast quadratic convergence. In this form, however, our method is not intended to be competitive with the eigen () routine in R, which is optimized compiled FORTRAN code from LAPACK.

The comparison with eigen () will become more interesting by rewriting critical sections in \underline{c} , although it is well-established that Jacobi is slower than the combination of Givens-Householder tri-diagonalization and inverse iteration. But Jacobi has the advantage that it is more easily parallelized [Sameh, 1971; Pourzandi and Tourancheau, 1995], and we intend to use OpenMP to see if we can make these routines competitive on SMP machines.

4.2. **Singular Values.** The existence theorem for the singular value decomposition says that for any rectangular X there exist rotation matrices K and L such that $\tilde{X} = KXL$ is diagonal. We can compute the singular value decomposition by one-sided Jacobi rotations, minimizing the sum of squares of the off-diagonal elements. We can first pivot through all left planar rotations, then pivot through all one-sided right planar rotations, then do the left ones again, and so on. Note that this algorithm can also be applied to symmetric matrices, in which case it gives us an alternative method to compute eigenvalues and eigenvectors.

Instead of minimizing the sum of squares of the off-diagonal elements we may as well maximize the sum of squares of the diagonal elements. For a left Jacobi transformation $K_{ij}(\theta)$ this means maximizing

$$(\cos(\theta)x_{ii} + \sin(\theta)x_{ji})^2 + (-\sin(\theta)x_{ij} + \cos(\theta)x_{jj})^2.$$

Define a 2×2 matrix V with

$$v_{11} = x_{ij}^2 + x_{ji}^2,$$

$$v_{12} = v_{21} = x_{ii}x_{ji} - x_{jj}x_{ij},$$

$$v_{22} = x_{ii}^2 + x_{jj}^2,$$

and a two element vector u with $u_1 = \sin(\theta)$ and $u_2 = \cos(\theta)$. Then we must maximize u'Vu over u'u = 1. Thus \hat{u} is the normalized eigenvector corresponding with the largest eigenvalue of V. There is no need to actually compute the optimal θ because u has all the information we need.

For a right Jacobi rotation $K_{ij}(\theta)$ we minimize

$$(\cos(\theta)x_{ii} - \sin(\theta)x_{ij})^2 + (\sin(\theta)x_{ji} + \cos(\theta)x_{jj})^2,$$

and we compute the largest eigenvalue and corresponding eigenvector of

$$v_{11} = x_{ij}^2 + x_{ji}^2,$$

$$v_{12} = v_{21} = x_{jj}x_{ji} - x_{ii}x_{ij},$$

$$v_{22} = x_{ii}^2 + x_{jj}^2.$$

The code for the algorithm in the Appendix, in the function jSVD(). There are various ways to deal with the fact that X is not square, and that consequently some of the singular vectors correspond with zero singular values. An obvious, although somewhat wasteful, way is to make X square by appending rows or columns with zeroes.

In our numerical experiments with jSVD() we find that the algorithm exhibits slow linear or even sublinear convergence. It performs reliably, but it is dreadfully slow. We are far better off applying the quadratically convergent Jacobi eigenvalue algorithm to X'X or XX', or even to

$$\begin{bmatrix} I & X \\ X' & I \end{bmatrix}.$$

On the other hand the algorithm can be generalized very simply to approximate simultaneous diagonalization of a number of rectangular matrices (see 4.4 below).

4.3. **Simultaneous Diagonalization.** If there are m symmetric matrices A_v they generally cannot be simultaneously diagonalized by an orthogonal transformation K. We can find K such that $K'A_vK$ is diagonal for all v if and only if the A_v commute in pairs, i.e. if and only if $A_vA_\mu = A_\mu A_v$ for all v, μ . We can find K, however, such that the transformed $K'A_vK$ are as diagonal as possible in the least squares sense. For this we minimize the sum of squares of all off-diagonal elements.

As Subsection 4.1 shows, for $K_{ij}(\xi)$ we must choose the plane rotation angle ξ such that

$$\sum_{v=1}^{m} \left[\frac{1}{2} \sin(2\xi) (a_{iiv} - a_{jjv}) + \cos(2\xi) a_{ijv} \right]^{2}$$

is minimized. Let $d_{ijv} = \frac{1}{2}(a_{iiv} - a_{jjv})$. Define a 2 × 2 matrix V with

$$v_{11} = \sum_{v=1}^{m} d_{ijv}^{2},$$

$$v_{12} = v_{21} = \sum_{v=1}^{m} d_{ijv} a_{ijv},$$

$$v_{22} = \sum_{v=1}^{m} a_{ijv}^{2},$$

and a two element vector u with $u_1 = \sin(2\xi)$ and $u_2 = \cos(2\xi)$. Then we must minimize u'Vu over u'u = 1, and thus the optimal u is any eigenvector corresponding with the smallest eigenvalue of V. Again, as in the Jacobi eigenvector method,

$$c_{\xi} = \sqrt{\frac{1+u_2}{2}},$$

$$s_{\xi} = \mathbf{sign}(u_1)\sqrt{\frac{1-u_2}{2}}.$$

This solution was first derived by De Leeuw and Pruzansky [1978], although in an unnecessarily complicated way. It was implemented in a convenient FORTRAN routine by Clarkson [1988]. Ten Berge [1984] has shown, in an interesting paper, that Kaiser's VARIMAX rotation method [Kaiser, 1958] can be formulated as a simultaneous diagonalization problem, and that the algorithm proposed by Kaiser is the same as the one in De Leeuw and Pruzansky [1978]. The implementation of the function <code>jSimDiag()</code> in the Appendix converges quite rapidly.

4.4. **Tucker Models.** Suppose X_1, \dots, X_m are rectangular matrices of the same dimensions. The problem is to find K and L such that $\tilde{X}_j = KX_jL$ are as diagonal as

possible in the least squares sense. But for this we can straightforwardly use all the results from 4.2. This gives an orthogonal version of the TUCKER-2 model [Kroonenberg and De Leeuw, 1980]. The function <code>jSimSVD()</code> is in the Appendix.

A more general problem is to transform the three-dimensional array $X = \{x_{ijk}\}$, using three rotation matrices K, L and M, by

$$\tilde{x}_{abc} = \sum_{p} \sum_{q} \sum_{r} k_{ap} l_{bq} m_{cr} x_{pqr}.$$

We can define different criteria to optimize. An interesting one is to maximize the sum of squares of the body diagonal, i.e. of the elements for which a=b=c. This fits an orthonormal version of the INDSCAL/PARAFAC model. It is implemented in <code>jTucker3Diag()</code>. Another option is to maximize the sum of squares of the leading principal block with $a \le A, b \le B$ and $c \le C$. This equivalent to fitting the original TUCKER-3 model [Kroonenberg and De Leeuw, 1980]. The function <code>jTucker3Block()</code> is also given in the Appendix. Both approaches are conceptually straightforward generalizations of Principal Component Analysis and the Singular Value Decomposition.

All techniques discussed in this section can be generalized to arrays in more than three dimensions. In \mathbb{R} this is most easily done by using the machinery developed by (APL). The code will be presented in a subsequent publication.

4.5. **PREHOM.** In De Leeuw [1982]; Bekker [1982] and in Bekker and De Leeuw [1988] a variation of multiple correspondence analysis using Jacobi plane rotations is proposed. The matrix to be analyzed is the Burt matrix [Burt, 1950] of m categorical variables. Variable j has k_j categories. If $G = \begin{bmatrix} G_1 & \cdots & G_m \end{bmatrix}$ are the concatenated indicator matrices (dummies) of the variables, with G_j of dimension $n \times k_j$, then the Burt matrix is C = G'G. Matrix C has submatrices $C_{k\ell}$ of dimension $k_j \times k_\ell$. If $j \neq \ell$ then $C_{j\ell}$ contains the bivariate marginals (cross table) of variables j and ℓ . If $j = \ell$ then $C_{j\ell}$ is a diagonal matrix with the univariate marginals of variable j on the diagonal. Also define $D = \operatorname{diag}(C)$, and the scaled Burt matrix $A = D^{-\frac{1}{2}}CD^{-\frac{1}{2}}$.

Multiple correspondence analysis can be defined as diagonalization of the matrix A. Thus we could use the classical Jacobi method of Subsection 4.1. In De Leeuw [1982] a three-step approximate diagonalization of the scaled Burt matrix is proposed, which is more revealing from a data analysis pont of view. In the first step

Jacobi plane rotations are used to approximately diagonalize all of the $k_j \times k_\ell$ submatrices $A_{j\ell}$ simultaneously. Note that the diagonal submatrices A_{jj} are equal to the identity matrix of order k_j , and are thus already diagonal. After the Jacobi step has finished we permute rows and columns to collect the diagonal elements of the $A_{j\ell}$ into a direct sum of diagonal blocks. If all k_j are equal to k, then there are k blocks of order k. Each diagonal block is a correlation matrix. The first block corresponds to the k0,1 elements of all k1,2 the second block to the k2,2 elements, and so on. If some variables have fewer categories, they will not occur in the later blocks. The third step of the approximate diagonalization computes the eigenvectors of the diagonal blocks, and uses them to diagonalize these blocks.

Thus the eigenvectors of the scaled Burt matrix are approximated by the product KPL, where K is the cumulative product of the Jacobi rotations that approximately diagonalize the $A_{j\ell}$, P is the permutation matrix that permutes the elements to approximate block-diagonal form, and L is the direct sum of the eigenvectors that diagonalize the correlation matrices along the diagonal. For m variables with k categories each we use $\frac{1}{2}mk(k+m-2)$ parameters to do the approximate diagonalization, instead of the $\frac{1}{2}mk(mk-1)$ parameters for the full diagonalization.

The code for the \mathbb{R} implementation is the function jMCA () in the Appendix.

REFERENCES

- P. Bekker. Varianten van Niet-lineaire Principale Komponenten Analyse. Master's thesis, Department of Psychology, University of Leiden, 1982.
- P. Bekker and J. De Leeuw. Relation between Variants of Nonlinear Principal Component Analysis. In J.L.A. Van Rijckevorsel and J. De Leeuw, editors, *Component and Correspondence Analysis*. Wiley, Chichester, England, 1988.
- C. Burt. The Factorial Analysis of Qualitative Data. *British Journal of Statistical Psychology*, 3:166–185, 1950.
- D.B. Clarkson. Remark AS R74: A Lest Squares Version of Algorithm AS 211: The F-G Diagonalization Algorithm. *Applied Statistics*, 37:317–321, 1988.
- J. De Leeuw. Nonlinear Principal Component Analysis. In H. Caussinus et al., editor, *COMPSTAT 1982*, pages 77–86, Vienna, Austria, 1982. Physika Verlag.
- J. De Leeuw and S. Pruzansky. A New Computational Method to fit the Weighted Euclidean Distance Model. *Psychometrika*, 43:479–490, 1978.

- A. Edelman, T.A. Arias, and S.T. Smith. The Geometry of Algorithms with Orthogonality Constraints. *SIAM Journal of Matrix Analysis and Applications*, 20: 303–353, 1998.
- H.F. Kaiser. The Varimax Criterion of Analytic Rotation in Factor Analysis. *Psychometrika*, 23:187–200, 1958.
- P.M. Kroonenberg and J. De Leeuw. Principal Component Analysis of Three-Mode Data by Means of Alternating Least Squares Algorithms. *Psychometrika*, 45:69–97, 1980.
- M. Pourzandi and B. Tourancheau. Parallel Performance of Jacobi Eigenvalue Solution. *Computing Systems in Engineering*, 6:377–383, 1995.
- A.H. Sameh. On Jacobi and Jacobi-Like Algorithms for a Parallel Computer. *Mathematics of Computation*, 25(579–590), 1971.
- J.M.F. Ten Berge. A Joint Treatment of Varimax Rotation and the Problem of Diagonalizing Symmetric Matrices Simultaneously in the Least Squares Sense. *Psychometrika*, 49(347–358), 1984.

APPENDIX A. CODE

```
1 #
  3
                    Copyright (C) 2008 Jan de Leeuw <deleeuw@stat.ucla.edu>
                    UCLA Department of Statistics, Box 951554, Los Angeles, CA 90095-1554
  5
  6 # This program is free software; you can redistribute it and \underline{\!\!/} or modify
  7 # it under the terms of the GNU General Public License as published by
  8 # the Free Software Foundation; either version 2 of the License, or
  9 #
                   (at your option) any later version.
10
11 # This program is distributed in the hope that it will be useful,
12 # but WITHOUT ANY WARRANTY; without even the implied warranty of
                   MERCHANTABILITY or FITNESS FOR A PARTICULAR PURPOSE. See the
13
14
                    GNU General Public License for more details.
15
16
                    You should have received a copy of the GNU General Public License
17
                   along with this program; if not, write to the Free Software
                   Foundation, Inc., 675 Mass Ave, Cambridge, MA 02139, USA.
18 #
19 #
21 #
22 # version 0.0.1, 2008-12-05 Eigen and SVD
23 # version 0.1.0, 2008-12-06 Simultaneous Diagonalization
24 # version 0.2.0, 2008-12-08 Simultaneous SVD
25 # version 0.2.1, 2008-12-08 Various Bugfixes
26 # version 0.2.2, 2008-12-08 Small efficiency gains
          # version 0.3.0, 2008-12-10 PREHOM in jMCA added
27
28
          # version 1.0.0, 2008-12-10 Tucker3 Methods added
30
31 \quad \texttt{jEigen} \\ \underbrace{\texttt{s-function}}_{} (\texttt{a,eps1} \\ = \texttt{1e-6,eps2} \\ = \texttt{1e-10,itmax} \\ = \texttt{100,vectors} \\ = \texttt{TRUE,verbose} \\ = \texttt{FALSE}) \quad \{\texttt{a,eps1} \\ = \texttt{a.eps1} \\ = \texttt{a.eps2} \\ = 
32 n \leq n \leq n  (a); k \leq n  (b); itel\leq n  (c); n \leq n  (a)
33 repeat {
34
                    for (i in 1:(n-1)) for (j in (i+1):n) {
35
                             aij<-a[i,j]; bij<-abs (aij);
36
                              if (bij < eps1) next()</pre>
37
                             mx<-max(mx,bij)</pre>
38
                              am < -(a[i,i]-a[j,j])/2
39
                              u \leq c(aij, -am); u \leq u/sqrt(sum(u^2))
40
                              c \leq -sqrt((1+u[2])/2); s \leq -sign(u[1]) + sqrt((1-u[2])/2)
41
                              ss<-s^2; cc<-c^2; sc<-s*c
                              ai<u><-</u>a[i,]; aj<u><-</u>a[j,]
42
43
                              aii < -a[i,i]; ajj < -a[j,j]
44
                              a[i,] \leq a[,i] \leq c_{\underline{\star}}ai - s_{\underline{\star}}aj
45
                              a[j,] \leq a[,j] \leq s * ai + c * aj
46
                              a[i,j] \leq a[j,i] \leq 0
47
                              a[i,i]<-aii*cc+ajj*ss-2*sc*aij
48
                              a[j,j]<-ajj*cc+aii*ss+2*sc*aij
                              if (vectors) {
49
50
                                       ki \leq k[,i]; kj \leq k[,j]
51
                                       k[,i] \leq c_{\underline{*}} ki - s_{\underline{*}} kj
```

```
52
                      k[,j] \leq s ki + c kj
53
55
           ff<-sqrt (saa-sum (diag (a) ^2))</pre>
56
           if (verbose)
57
                cat("Iteration ", formatC(itel, digits=4), "maxel ", formatC(mx, width=10),"
                       loss ", formatC (ff, width=10), "\n")
           <u>if</u> ((mx < eps1) || (ff < eps2) || (itel == itmax)) <u>break()</u>
59
           itel<-itel+1; mx<-0
60
61 d \leq -diag(a); o \leq -order(d, decreasing=TRUE)
\underline{\text{if}} \text{ (vectors) } \underline{\text{return}} \text{(} \underline{\text{list}} \text{(values=d[o], vectors=k[,o]))}
63
           else return(values=d[o])
64 }
65
66 jSVD<-function(x,eps1=le-6,eps2=le-6,itmax=1000,vectors=TRUE,verbose=FALSE) {
67 n \leq -nrow(x); m \leq -ncol(x); itel\leq -1; mx \leq -0
68 kkk \leq -diag(n); lll \leq -diag(m);
69 sxx < -sum(x^2); sxm < -sqrt(sxx/(n*m))
70
     repeat {
71
           for (i in 1:(n-1)) {
72
                           <u>if</u> (i > m) <u>next()</u>
                           <u>for</u> (j in (i+1):n) {
73
74
                           xi \leq x[i,]; xj \leq x[j,]
75
                           xij \leq -ifelse(j > m, 0, x[i, j])
76
                            xjj < -ifelse(j > m, 0, x[j, j])
77
                                     xii < x[i,i]; xji < x[j,i]
78
                           mx<-max (mx, abs (xij) /sxm, abs (xji) /sxm)</pre>
                           v \leq -matrix(0,2,2)
80
                            v[1,1] \leq xij^2 + xji^2
81
                           v[1,2] \leq v[2,1] \leq xii \times xji - xjj \times xij
82
                           v[2,2]<-xii^2+xjj^2
83
                           u<-eigen(v) $vectors[,1]
84
                           x[i,] \leq u[2] \times xi + u[1] \times xj
85
                           x[j,] \leq u[2] \times xj = u[1] \times xi
86
                           if (vectors) {
87
                                ki<-kkk[i,]; kj<-kkk[j,]
88
                                kkk[i,] \leq u[2] \times ki + u[1] \times kj
89
                                 kkk[j,] \leq u[2] * kj - u[1] * ki
90
                                 }
91
92
93
           ff < -sqrt ((sxx-sum (diag(x)^2))/sxx)
94
           if (verbose)
95
                 \underline{\mathtt{cat}}\,(\texttt{"Left iteration ",}\underline{\mathtt{formatC}}\,(\mathtt{itel,digits=4})\,,\texttt{"maxel ",}\underline{\mathtt{formatC}}\,(\mathtt{mx,width})
                       =10), "loss ", <u>formatC</u>(ff, width=10), "\n")
96
           for (k in 1:(m-1)) {
97
                 \underline{if} (k > n) \underline{next}()
98
                           for (1 in (k+1):m) {
99
                           xk \leq x[,k]; xl \leq x[,1]
100
                           xlk \leq -ifelse(1 > n, 0, x[1, k])
101
                           xll \leq -ifelse(1 > n, 0, x[1, 1])
102
                           xkk \leq x[k,k]; xkl \leq x[k,1]
```

```
103
                            mx<-max (mx, abs (xkl) /sxm, abs (xlk) /sxm)</pre>
104
                            v \leq -matrix(0,2,2)
105
                            v[1,1] < -xk1^2 + x1k^2
106
                            v[1,2] \leq v[2,1] \leq xll \times xlk - xkk \times xkl
107
                            v[2,2] < -xkk^2+x11^2
108
                            u<-eigen (v) $vectors[,1]
109
                            x[,k] \leq u[2] \times xk - u[1] \times x1
110
                            x[,1] \leq u[1] \star xk + u[2] \star xl
111
                            if (vectors) {
112
                                 1k<-111[,k]; 11<-111[,1]
113
                                 111[,k] < u[2] * 1k - u[1] * 11
114
                                 111[,1] < u[1] * 1k+u[2] * 11
115
116
                             }
117
118
            ff<-sqrt ((sxx-sum(diag(x)^2))/sxx)
119
            if (verbose)
120
                 cat("Right iteration ", formatC(itel, digits=4), "maxel ", formatC(mx, width
                        =10), "loss ", <u>formatC</u>(ff, width=10), "\n")
121
            \underline{\text{if}} ((mx < eps1) || (ff < eps2) || (itel == itmax)) \underline{\text{break}}()
122
            itel < -itel + 1; mx < -0
123
124
      \underline{\text{return}} (\underline{\text{list}} (d = \underline{\text{diag}} (x), u = \underline{\text{t}} (kkk), v = 111))
125
      }
126
127
      jSimDiag<-function(a,eps=1e-10,itmax=100,vectors=TRUE,verbose=FALSE) {
128
     n \leq -\dim(a)[1]; kk \leq -\dim(a)[3]; itel \leq 1; saa \leq -sum(a^2)
129 fold \leq saa = sum(apply(a, 3, function(x) sum(diag(x^2))))
130
     repeat {
131
            for (i in 1:(n-1)) for (j in (i+1):n) {
132
            ad \leq (a[i,i,]-a[j,j,])/2
133
            av \leq a[i,j,]
134
           v \leq -matrix(0,2,2)
135
           v[1,1] \leq sum(ad^2)
           v[1,2] \leq v[2,1] \leq sum(av * ad)
136
137
           v[2,2] < -sum(av^2)
138
           u<-eigen (v) $vectors[,2]
139
           c \leq -sqrt((1+u[2])/2); s \leq -sign(u[1]) + sqrt((1-u[2])/2)
140
           for (k in 1:m) {
141
                     ss<u><-</u>s^2; cc<u><-</u>c^2; sc<u><-</u>s<u>*</u>c
142
                      ai < -a[i,,k]; aj < -a[j,,k]
143
                      aii < -a[i,i,k]; ajj < -a[j,j,k]; aij < -a[i,j,k]
144
                      a[i,,k] \leq a[,i,k] \leq c_{\star}ai-s_{\star}aj
145
                      a[j,,k] < -a[,j,k] < -s * ai + c * aj
146
                      a[i,j,k] \leq a[j,i,k] \leq u[1] \times (aii-ajj) / 2 + u[2] \times aij
147
                       a[i,i,k] < -aii * cc + ajj * ss - 2 * sc * aij
148
                       a[j,j,k] \leq -ajj \times cc + aii \times ss + 2 \times sc \times aij
149
                      }
            if (vectors) {
150
151
                 ki < -kk[,i]; kj < -kk[,j]
152
                 kk[,i] \leq c_{\underline{\star}}ki - s_{\underline{\star}}kj
153
                 kk[,j] \leq -s_{\underline{\star}}ki+c_{\underline{\star}}kj
154
```

```
155
156
           f_{new} \leq s_{aa} - s_{um} (apply (a, 3, f_{unction}(x) s_{um} (d_{iag}(x^2))))
157
           if (verbose)
               cat("Iteration ", formatC(itel, digits=4), "old loss ", formatC(fold, width=10)
158
                      , "new loss ", formatC (fnew, width=10), "\n")
159
           if (((fold-fnew) < eps) || (itel == itmax)) break()</pre>
160
           itel\leq-itel+1; fold\leq-fnew
161
162
     return(list(a=a, d<-apply(a, 3, diag), k=kk))</pre>
163
164
165 jSimSVD<-function(x,eps=le-6,itmax=1000,vectors=TRUE,verbose=FALSE) {
166 n \leq -\dim(x) [1]; m \leq -\dim(x) [2]; n \max \leq -\dim(x) [3]; itel \leq -1
167 kkk < -diag(n); lll < -diag(m); sxx < -sum(x^2); fold < -Inf
169
     repeat {
170
          for (i in 1:(n-1)) {
171
               \underline{if} (i > m) \underline{next}()
172
                          for (j in (i+1):n) {
173
                          v \leq -matrix(0,2,2)
174
                         for (imat in 1:nmat) {
                              xij < -ifelse(j > m, 0, x[i, j, imat])
175
176
                              xjj < -ifelse(j > m, 0, x[j, j, imat])
177
                                             xii<-x[i,i,imat]; xji<-x[j,i,imat]</pre>
178
                              v[1,1] \leq v[1,1] + (xij^2+xji^2)
179
                              v[1,2] \leq v[2,1] \leq v[1,2] + (xii *xji - xjj *xij)
180
                              v[2,2] \leq v[2,2] + (xii^2+xjj^2)
181
182
                         u<-eigen(v) $vectors[,1]
183
                          for (imat in 1:nmat) {
184
                              xi < x[i,,imat]; xj < x[j,,imat]
185
                              x[i,,imat] \leq u[2] \times xi + u[1] \times xj
186
                              x[j,,imat] \leq u[2] \times xj - u[1] \times xi
187
188
                         if (vectors) {
189
                              ki<u><-</u>kkk[i,]; kj<u><-</u>kkk[j,]
190
                              kkk[i,] \leq u[2] * ki + u[1] * kj
191
                              kkk[j,] \leq u[2] * kj - u[1] * ki
192
193
194
195
           ss \leq sum(apply(x, 3, diag)^2); fnew \leq sqrt((sxx-ss)/sxx)
196
           if (verbose)
197
                cat(" Left iteration ", formatC(itel, digits=4), "loss ", formatC(fnew, digits
                      =6, width=10), "\n")
198
           <u>for</u> (k in 1: (m-1)) {
199
                \underline{if} (k > n) \underline{next}()
200
                          for (l in (k+1):m) {
201
                         v \leq -matrix(0,2,2)
                         for (imat in 1:nmat) {
202
203
                              xlk \leq -ifelse(1 > n, 0, x[1, k, imat])
204
                              xll \leq -ifelse(1 > n, 0, x[1, 1, imat])
205
                              xkl \leq x[k,l,imat]; xkk \leq x[k,k,imat]
```

```
206
                                 v[1,1] \leq v[1,1] + (xk1^2+x1k^2)
207
                                 v[1,2] < v[2,1] < v[1,2] + (xll \times xlk - xkk \times xkl)
208
                                 v[2,2] < -v[2,2] + (xkk^2+x11^2)
209
210
                           u<-eigen(v) $vectors[,1]
211
                           for (imat in 1:nmat) {
212
                                 xk < x[,k,imat]; xl < x[,l,imat]
213
                                 x[,k,imat] \leq u[2] \times xk - u[1] \times xl
214
                                 x[,1,imat] < -u[1] *xk+u[2] *xl
215
216
                           if (vectors) {
217
                                1k<-111[,k]; 11<-111[,1]
218
                                111[,k] < u[2] * 1k - u[1] * 11
219
                                111[,1] < u[1] * 1k+u[2] * 11
220
221
222
223
            ss \leq -sum(apply(x, 3, diag)^2); fnew \leq -sqrt((sxx-ss)/sxx)
224
            if (verbose)
225
                 cat("Right iteration ", formatC(itel, digits=4), "loss ", formatC(fnew, digits
                       =6, width=10), "\n")
226
            if (((fold - fnew) < eps) || (itel == itmax)) break()</pre>
227
            itel<-itel+1; fold<-fnew
228
229
      \underline{\text{return}} (\underline{\text{list}} (d = \underline{\text{apply}} (x, 3, \underline{\text{diag}}), u = \underline{\text{t}} (kkk), v = 111))
230
     }
231
232 jTucker3Diag<-function(a,eps=1e-6,itmax=100,vectors=TRUE,verbose=TRUE) {
233 n \leq -\dim(a)[1]; m \leq -\dim(a)[2]; k \leq -\dim(a)[3]; nmk \leq -\min(n, m, k)
234 kn < -diag(n); km < -diag(m); kk < -diag(k); ossq < -0; itel < -1
235
      repeat {
236
                 for (i in 1:(n-1)) for (j in (i+1):n) {
237
                           ai<u><-</u>a[i,,]; aj<u><-</u>a[j,,]
238
                           acc<-ass<-asc<-0
                           <u>if</u> (i <= <u>min</u>(m,k)) {
239
240
                                      acc<-acc+a[i,i,i]^2
241
                                      ass<-ass+a[j,i,i]^2
242
                                      asc \leq -asc + a[i,i,i] * a[j,i,i]
243
                                      }
244
                           \underline{if} (j <= \underline{min} (m, k)) {
245
                                     acc<-acc+a[j,j,j]^2
246
                                      ass<-ass+a[i,j,j]^2
247
                                      asc < -asc - a[j, j, j] * a[i, j, j]
248
                                      }
249
                           u \leq -eigen (matrix (c(acc, asc, asc, ass), 2, 2)) $vectors[,1]
250
                            c \le u[1]; s \le u[2]
251
                           a[i,,] \leq -c_{\underline{\star}}ai + s_{\underline{\star}}aj
252
                           a[j,,]<-c<u>*</u>aj-s<u>*</u>ai
253
                           <u>if</u> (vectors) {
254
                                 ki < -kn[i,]; kj < -kn[j,]
255
                                kn[i,]<u><-</u>c<u>*</u>ki+s<u>*</u>kj
256
                                 kn[j,] \leq c_{\underline{\star}} kj - s_{\underline{\star}} ki
257
                                     }
```

```
258
259
                  for (i in 1:(m-1)) for (j in (i+1):m) {
260
                            ai<-a[,i,]; aj<-a[,j,]
261
                             acc<u><-</u>ass<u><-</u>asc<u><-</u>0
262
                             \underline{if} (i <= \underline{min}(n,k)) {
263
                                       acc<-acc+a[i,i,i]^2
264
                                        ass < -ass + a[i, j, i]^2
265
                                        asc < -asc + a[i,i,i] * a[i,j,i]
266
267
                            \underline{if} (j <= \underline{min}(n,k)) {
268
                                       acc<-acc+a[j,j,j]^2
269
                                       ass < -ass + a[j,i,j]^2
270
                                        asc < -asc - a[j, j, j] * a[j, i, j]
271
272
                            u<-eigen (matrix (c (acc, asc, asc, ass), 2, 2)) $vectors[,1]
273
                            c<-u[1]; s<-u[2]
274
                            a[,i,]<-c*ai+s*aj
275
                            a[,j,]<mark><-</mark>c<u>*</u>aj-s<u>*</u>ai
276
                             if (vectors) {
277
                                  ki < -km[i,]; kj < -km[j,]
278
                                  km[i,] \leq c_{\underline{\star}} ki + s_{\underline{\star}} kj
279
                                  km[j,] < -c \times kj - s \times ki
280
                                       }
281
                             }
282
                 for (i in 1:(k-1)) for (j in (i+1):k) {
283
                            ai<-a[,,i]; aj<-a[,,j]
284
                            acc<-ass<-asc<-0
285
                             <u>if</u> (i <= <u>min</u>(n, m)) {
286
                                       acc<-acc+a[i,i,i]^2
287
                                       ass<-ass+a[i,i,j]^2
288
                                        asc < -asc + a[i,i,i] * a[i,i,j]
289
290
                             \underline{if} (j <= \underline{min}(n,m)) {
291
                                       acc<-acc+a[j,j,j]^2
292
                                       ass<-ass+a[j,j,i]^2
293
                                        asc < -asc - a[j, j, j] * a[j, j, i]
294
                                       }
295
                            u<-eigen (matrix (c (acc, asc, asc, ass), 2, 2)) $vectors[,1]
296
                            c \le u[1]; s \le u[2]
297
                            a[,,i]<-c<u>*</u>ai+s<u>*</u>aj
298
                            a[,,j]<u><-</u>c<u>∗</u>aj-s<u>∗</u>ai
299
                             if (vectors) {
300
                                  ki<-kk[i,]; kj<-kk[j,]
301
                                  kk[i,] \leq c_{\underline{\star}}ki+s_{\underline{\star}}kj
302
                                  kk[j,] \leq c_{\underline{\star}} k_{j} - s_{\underline{\star}} k_{i}
303
                                       }
304
305
                 nssq < 0; for (v in 1:nmk) nssq < nssq + a[v, v, v]^2
306
            if (verbose)
307
                 cat("Iteration ", formatC(itel, digits=4), "ssq ", formatC(nssq, digits=10,
                        width=15), "\n")
308
            \underline{\text{if}} (((nssq - ossq) < eps) || (itel == itmax)) \underline{\text{break}}()
309
            itel\leq-itel+1; ossq\leq-nssq
```

```
310
311 d \leftarrow rep(0, nmk); for (v in 1:nmk) <math>d[v] \leftarrow a[v, v, v]
      return(list(a=a,d=d,kn=kn,km=km,kk=kk))
312
313
314
315 jTucker3Block<-function(a,dims,eps=1e-6,itmax=100,vectors=TRUE,verbose=TRUE) {
316
      n \leq -\dim(a) [1]; m \leq -\dim(a) [2]; k \leq -\dim(a) [3]; nmk \leq -\min(n, m, k)
      p<-dims[1]; q<-dims[2]; r<-dims[3]</pre>
318 kn \leq -diag(n); km \leq -diag(m); kk \leq -diag(k); ossq \leq -0; itel \leq -1
319
      repeat {
320
                for (i in 1:(n-1)) for (j in (i+1):n) {
321
                           ai<u><-</u>a[i,,]; aj<u><-</u>a[j,,]
322
                            acc<u><-</u>ass<u><-</u>asc<u><-</u>0
323
                            \underline{if} (i <= p)
324
                                       for (u in 1:g) for (v in 1:r) {
325
                                                 acc<-acc+a[i,u,v]^2
326
                                                 ass < -ass + a[j,u,v]^2
327
                                                 asc < -asc + a[i,u,v] * a[j,u,v]
328
329
                            \underline{if} (j <= p)
330
                                       for (u in 1:q) for (v in 1:r) {
331
                                                 acc<-acc+a[j,u,v]^2
332
                                                 ass < -ass + a[i,u,v]^2
333
                                                 asc < -asc - a[j, u, v] * a[i, u, v]
334
                                                 }
335
                            u<-eigen (matrix (c (acc, asc, asc, ass), 2, 2)) $vectors[,1]
336
                            c \le u[1]; s \le u[2]
337
                            a[i,,] < -c * ai + s * aj
338
                            a[j,,]<u><-</u>c<u>∗</u>aj-s<u>∗</u>ai
339
                            if (vectors) {
340
                                 ki<-kn[i,]; kj<-kn[j,]
341
                                 kn[i,] < -c * ki + s * kj
342
                                 kn[j,] \leq c_{\underline{\star}}kj - s_{\underline{\star}}ki
343
                                      }
344
                            }
345
                 for (i in 1:(m-1)) for (j in (i+1):m) {
346
                           ai<u><-</u>a[,i,]; aj<u><-</u>a[,j,]
347
                            acc<u><-</u>ass<u><-</u>asc<u><-</u>0
348
                            \underline{if} (i <= \underline{q})
349
                                       for (u in 1:p) for (v in 1:r) {
350
                                                acc < -acc + a[u,i,v]^2
351
                                                 ass < -ass + a[u, j, v]^2
352
                                                 asc < -asc + a[u,i,v] * a[u,j,v]
353
354
                            \underline{if} (j <= \underline{q})
355
                                       for (u in 1:p) for (v in 1:r) {
356
                                                 acc < -acc+a[u,j,v]^2
357
                                                  ass < -ass + a[u,i,v]^2
358
                                                  asc < -asc - a[u,i,v] * a[u,j,v]
359
                                                  }
360
                            u \leq -eigen (matrix (c(acc, asc, asc, ass), 2, 2)) $\square vectors[,1]
361
                            c \le u[1]; s \le u[2]
362
                            a[,i,] \leq c_{\underline{\star}} ai + s_{\underline{\star}} aj
```

```
363
                              a[,j,]<u><-</u>c<u>*</u>aj-s<u>*</u>ai
364
                              if (vectors) {
365
                                    ki \leq km[i,]; kj \leq km[j,]
366
                                    km[i,] \leq c_{\underline{\star}}ki + s_{\underline{\star}}kj
367
                                    km[j,] < -c *kj - s *ki
368
                                         }
369
370
                  for (i in 1:(k-1)) for (j in (i+1):k) {
371
                             ai<-a[,,i]; aj<-a[,,j]
372
                              acc<-ass<-asc<-0
373
                              if (i <= r)</pre>
374
                                         for (u in 1:p) for (v in 1:q) {
375
                                                    acc<-acc+a[u,v,i]^2
376
                                                     ass < -ass + a[u, v, j]^2
377
                                                     asc < -asc + a[u, v, i] * a[u, v, j]
378
379
                              <u>if</u> (j <= r)
380
                                         \underline{\text{for}} (u in 1:p) \underline{\text{for}} (v in 1:q) {
381
                                                     acc < -acc + a[u, v, j]^2
382
                                                     ass < -ass + a[u, v, i]^2
383
                                                     asc < -asc - a[u, v, i] * a[u, v, j]
384
                                                     }
                              u<-eigen(matrix(c(acc,asc,asc,ass),2,2))$vectors[,1]
385
386
                              c \le u[1]; s \le u[2]
387
                              a[,,i]<mark><-</mark>c<u>∗</u>ai+s<u>∗</u>aj
388
                              a[,,j]<mark><-</mark>c<u>*</u>aj-s<u>*</u>ai
389
                              if (vectors) {
390
                                    ki < -kk[i,]; kj < -kk[j,]
391
                                    kk[i,] \leq c_{\star}ki + s_{\star}kj
                                    kk[j,]<u><-</u>c<u>∗</u>kj-s<u>∗</u>ki
392
393
394
395
                  nssq \leq 0; for (i in 1:p) for (j in 1:q) for (l in 1:r) nssq \leq nssq + a[i,j,l]
                        ]^2
396
             if (verbose)
397
                  cat("Iteration ", formatC(itel, digits=4), "ssq ", formatC(nssq, digits=10,
                         width=15), "\n")
398
             \underline{if} (((nssq - ossq) < eps) || (itel == itmax)) \underline{break}()
399
            itel<-itel+1; ossq<-nssq
400
             }
401
      d<u><-</u>a[1:p,1:<u>q</u>,1:r]
      return(list(a=a,d=d,kn=kn,km=km,kk=kk))
403
      }
404
405
      jMCA<-function(burt,k,eps=1e-6,itmax=500,verbose=TRUE,vectors=TRUE) {</pre>
406
      m \leq -length(k); burt\leq -m \pm m \pm burt / sum(burt); sk \leq -sum(k)
407
      \label{eq:dbdef} db \underline{<-\text{diag}} \, (\text{burt}) \, ; \; 11 \underline{<-} kk \underline{<-} ww \underline{<-\text{diag}} \, (\text{sk}) \, ; \; \text{itel} \underline{<-}1 ; \; \text{ossq} \underline{<-}0
      klw \leq 1 + cumsum(c(0,k))[1:m]; kup \leq -cumsum(k)
409
      ind<-lapply(1:m, function(i) klw[i]:kup[i])</pre>
410
      sburt<-burt/sqrt (outer(db, db))</pre>
411
      <u>for</u> (i in 1:m)
412
           kk[ind[[i]],ind[[i]]] \leq -t (svd(sburt[ind[[i]],]) \leq u)
413 kbk \leq -kk \frac{% * %}{s} sburt \frac{% * %t}{s} (kk)
```

```
414 <u>for</u> (i in 1:m) <u>for</u> (j in 1:m)
415
         ww[ind[[i]],ind[[j]]]<-ifelse(outer(1:k[i],1:k[j],"=="),1,0)
416 <u>repeat</u> {
417
           <u>for</u> (l in 1:m) {
418
                <u>if</u> (k[1] == 2) <u>next()</u>
419
                li<u><-</u>ind[[1]]
420
                for (i in (klw[l]+1):(kup[l]-1)) for (j in (i+1):kup[l]) {
421
                     bi \leq kbk[i,-li]; bj \leq kbk[j,-li]
422
                     wi<-ww[i,-li]; wj<-ww[j,-li]
423
                    acc<-sum(wi*bi^2)+sum(wj*bj^2)
424
                    acs<-sum((wi-wj)*bi*bj)
425
                     ass<-sum(wi*bj^2)+sum(wj*bi^2)
426
                    u<-eigen (matrix (c (acc, acs, acs, ass), 2, 2)) §vectors[,1]
427
                    c < -u[1]; s < -u[2]
428
                    kbk[-li,i] < -kbk[i,-li] < -c *bi + s *bj
429
                    kbk[-li,j]<-kbk[j,-li]<-c*bj-s*bi
430
                     if (vectors) {
431
                           ki \leq -kk[i,li]; kj \leq -kk[j,li]
432
                           kk[i,li] \leq c_{\star}ki+s_{\star}kj
433
                           kk[j,li]<u><-</u>c<u>∗</u>kj-s<u>∗</u>ki
434
                           }
435
436
                }
437
           nssq<-sum (ww*kbk^2)-sum (diag(kbk)^2)
438
           <u>if</u> (verbose)
439
                cat("Iteration ", formatC(itel, digits=4), "ssq ", formatC(nssq, digits=10,
                      width=15), "\n")
440
           if (((nssq - ossq) < eps) || (itel == itmax)) break()
441
           itel<-itel+1; ossq<-nssq
442
443 kl<-unlist(sapply(k, function(i) 1:i))
444 pp<-ifelse(outer(1:sk,order(kl),"=="),1,0)
445 pkbkp<u><-t</u> (pp) <u>%*%</u>kbk<u>%*%</u>pp
446 pk<-t (pp) % * % kk
     km<-as.vector(table(kl)); nm<-length(km)</pre>
447
448 klw < -1 + cumsum (c(0,km))[1:nm]; kup < -cumsum (km)
449 <u>for</u> (i in 1:<u>length</u>(km)) {
450
              <u>if</u> (km[i]==1) <u>next</u>()
451
               ind<-klw[i]:kup[i]
452
               ll[ind,ind] <-eigen (pkbkp[ind,ind]) $vectors</pre>
454 lpkbkpl<-t (11) % * % pkbkp% * % 11
455 lpk<u><-t</u>(ll) <u>%*%</u>pk
 \frac{456}{\text{return}} \left( \frac{\text{list}}{\text{kbk} + \text{kbk}}, \text{pkbkp} = \text{pkbkp}, \text{lpkbkpl} = \text{lpkbkpl}, \text{kk} + \underline{\textbf{t}}(\text{kk}), \text{pp} = \text{pp}, \text{ll} = \text{ll}, \text{kpl} + \underline{\textbf{t}}(\text{lpk}) \right) \right) 
457 }
```

DEPARTMENT OF STATISTICS, UNIVERSITY OF CALIFORNIA, LOS ANGELES, CA 90095-1554

E-mail address, Jan de Leeuw: deleeuw@stat.ucla.edu

URL, Jan de Leeuw: http://gifi.stat.ucla.edu