

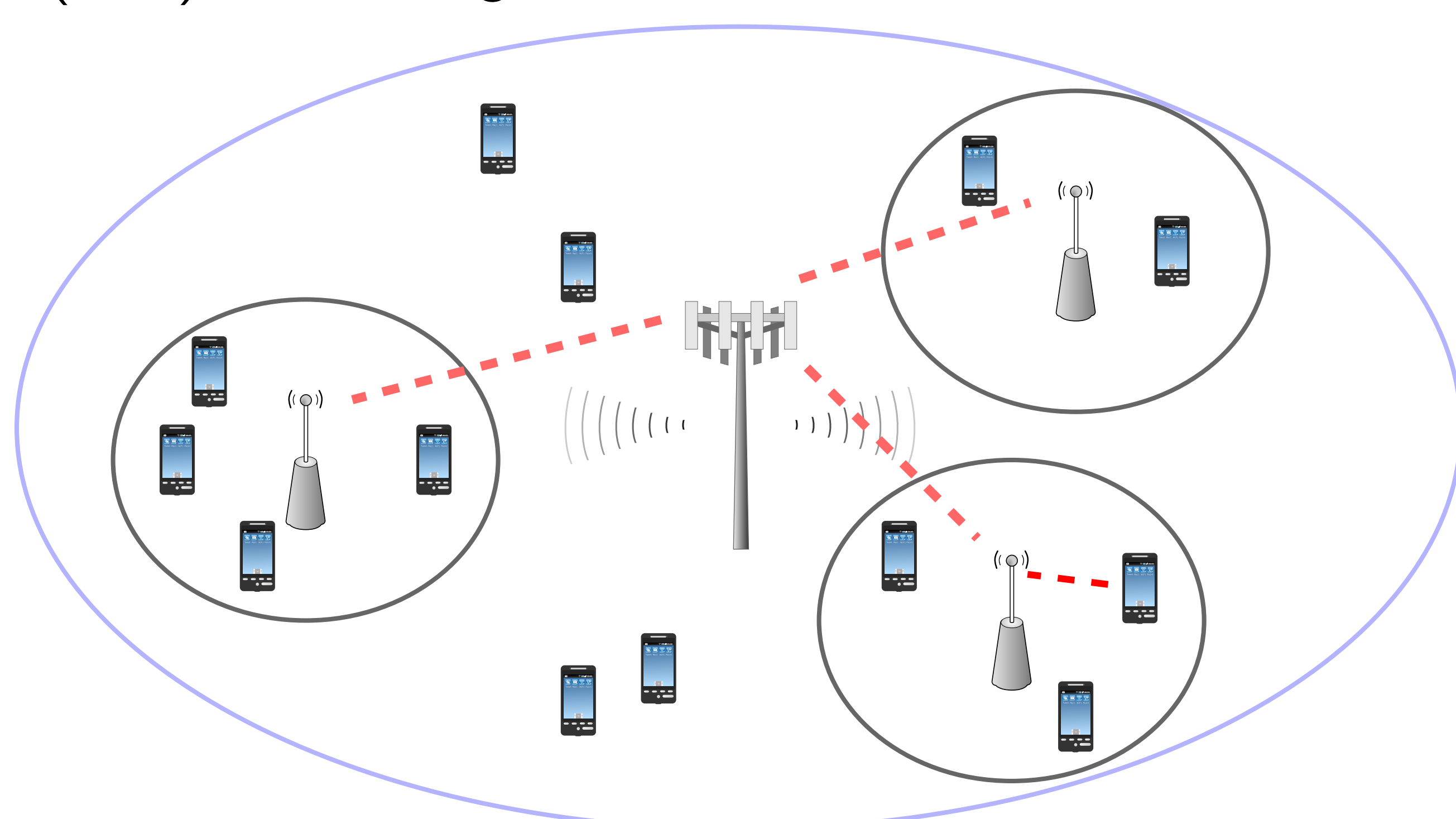
Transcoding: A New Strategy for Relay Channels

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Rise of Multi-hop, Low-latency Communication

- ▶ Number of wireless devices continues to grow
- ▶ Many “new” devices will be low-power Internet of Things (IoT) devices
- ▶ Potentially **no direct connection to base station**
- ▶ Cellular: Move to **small cell networks**, in-band wireless (self-)backhauling

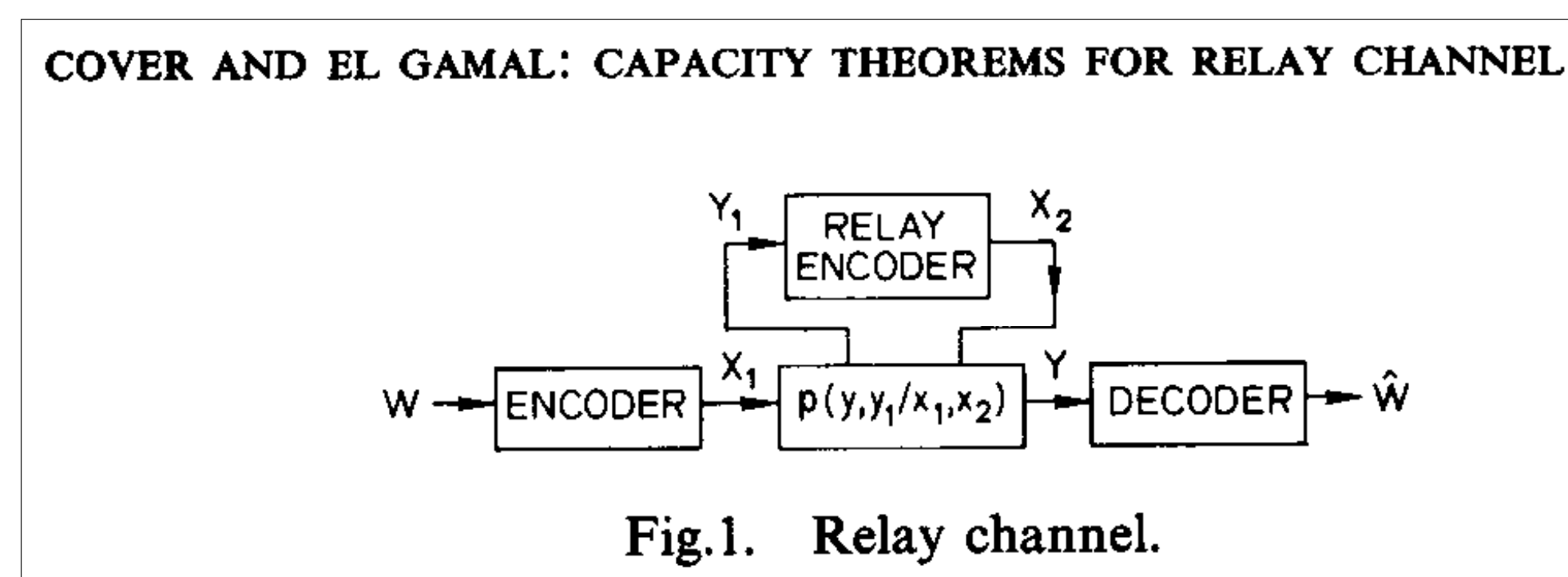


Small cell network with wireless backhaul

- ▶ Additionally: Growing focus on **low latency**
- ▶ Sub-1ms latency in IMT-2020

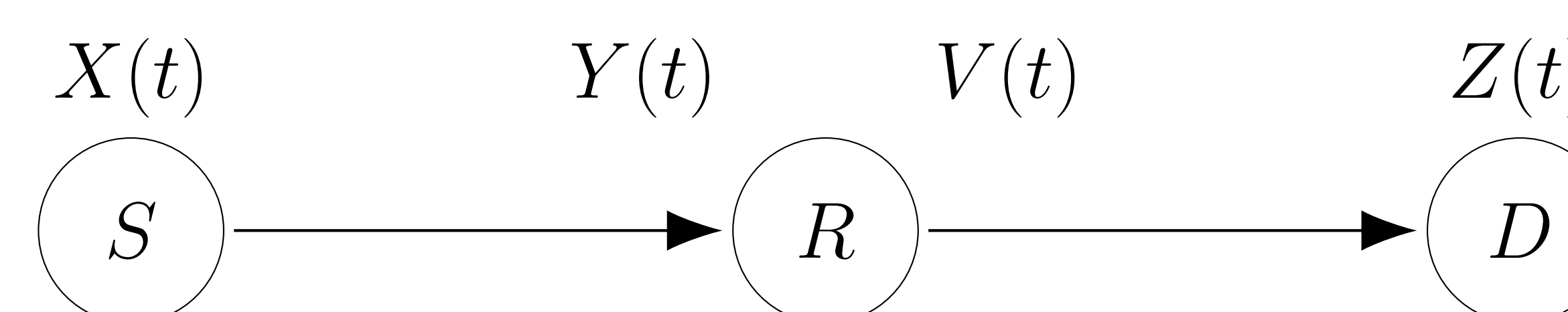
The Relay Channel: A Classic Problem in Information Theory

- ▶ Lots of focus from industry & academia since 1970s



- ▶ Many different design philosophies: Compress-&-Forward [1], Hash-&-Forward [2], Compute-&-Forward [3], Noisy Network Coding [4]
- ▶ De-facto industry standard: Decode-&-Forward (DF) [1]
- ▶ **Traditional schemes focus on capacity rather than delay performance** (\Rightarrow long block lengths)
- ▶ Our interest: Low latency, short, **finite block lengths**
- ▶ Amplify-&-Forward (AF) gives best delay performance but suffers throughput loss due to noise build-up

Separated Two-hop Relay Channel

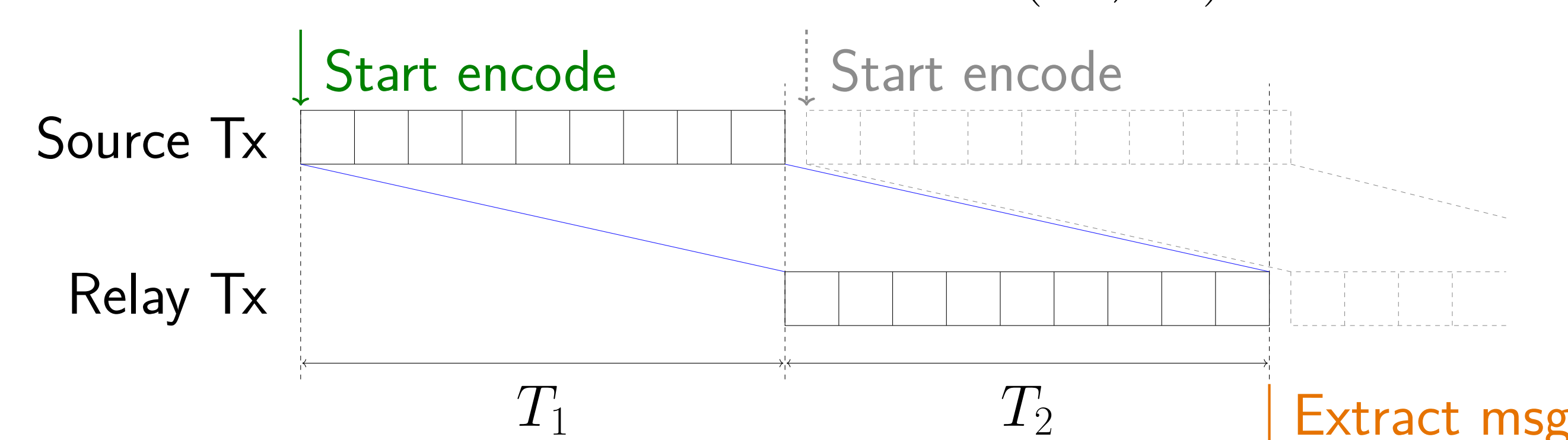


- ▶ No connection between source and destination
- ▶ “Degraded” relay channel, DF achieves capacity (with infinite delay)

Decode-&-Forward

- ▶ Error control through **coding at source and relay**
- ▶ End-to-end delay: $T = T_1 + T_2$
- ▶ Pipelined coding rate:

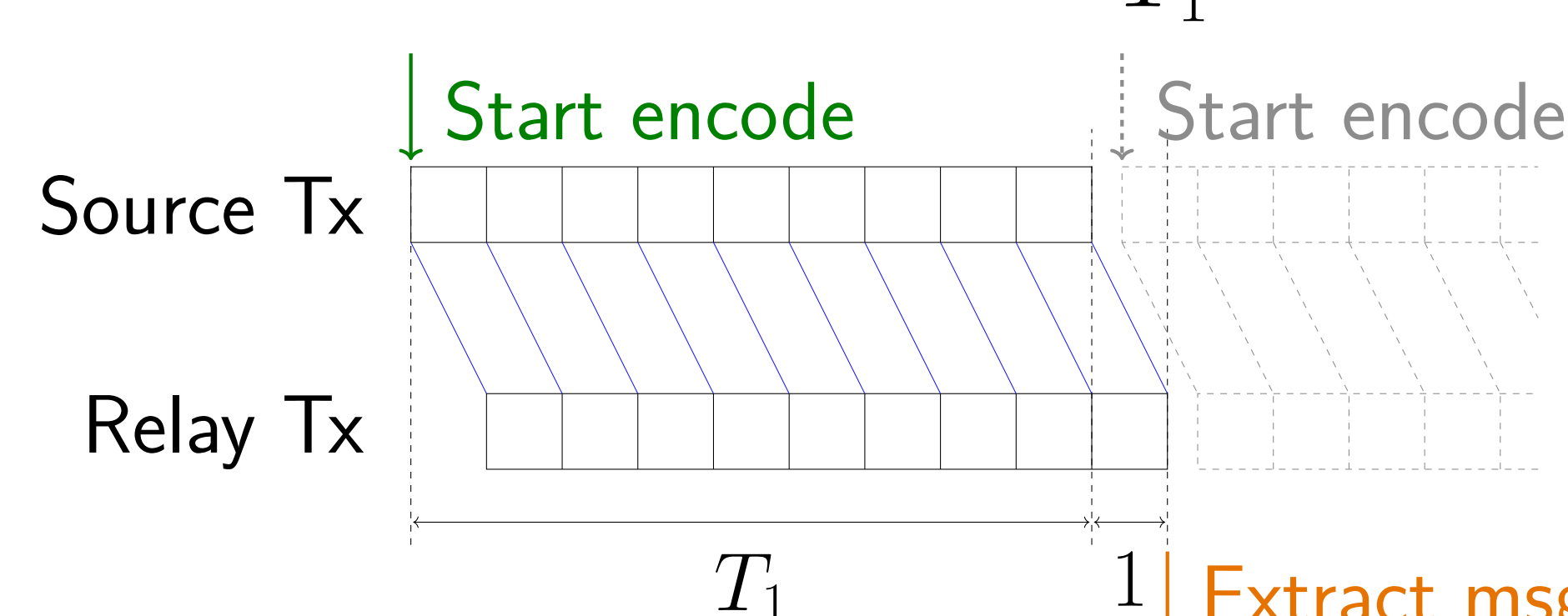
$$R_{DF}(T, \epsilon) = \frac{\log M}{\max(T_1, T_2)}$$



Amplify-&-Forward

- ▶ No error correction at relay \Rightarrow **noise accumulation**
- ▶ End-to-end delay: $T = T_1 + 1$
- ▶ Pipelined coding rate:

$$R_{AF}(T, \epsilon) = \frac{\log M}{T_1}$$



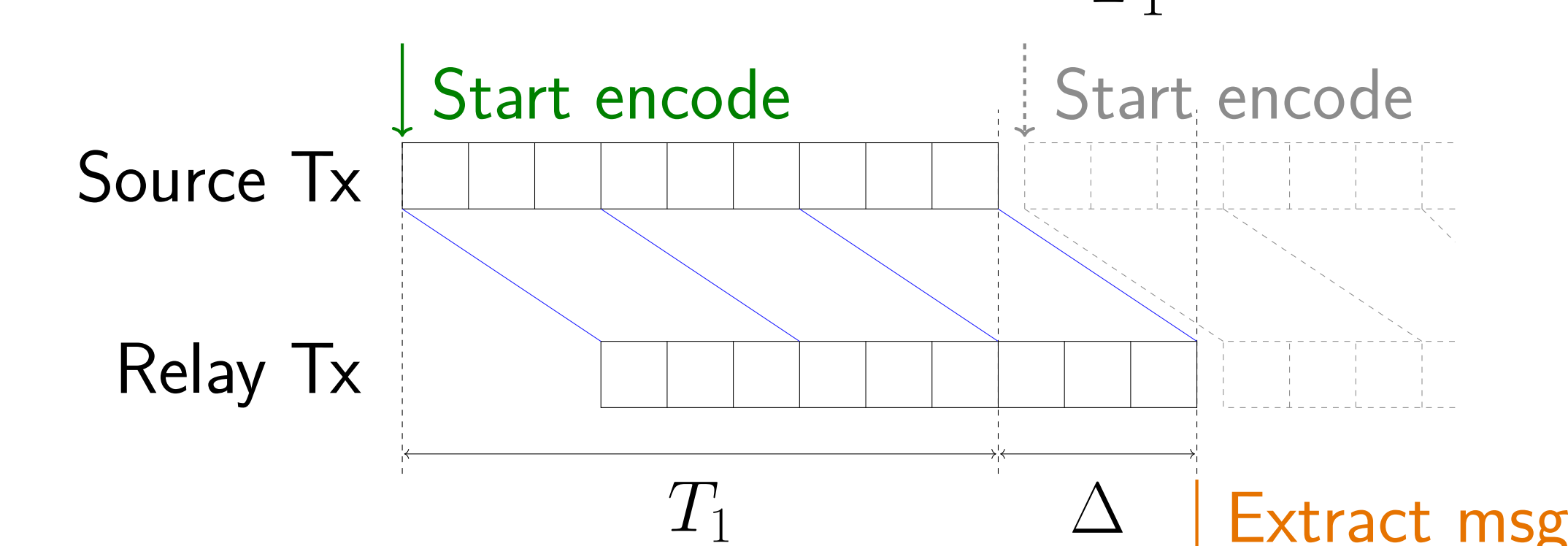
Transcoding as “middle ground” between AF and DF

- ▶ Can be viewed as “smart” AF with error protection
- ▶ Improved coding rate in low latency regime

Transcoding

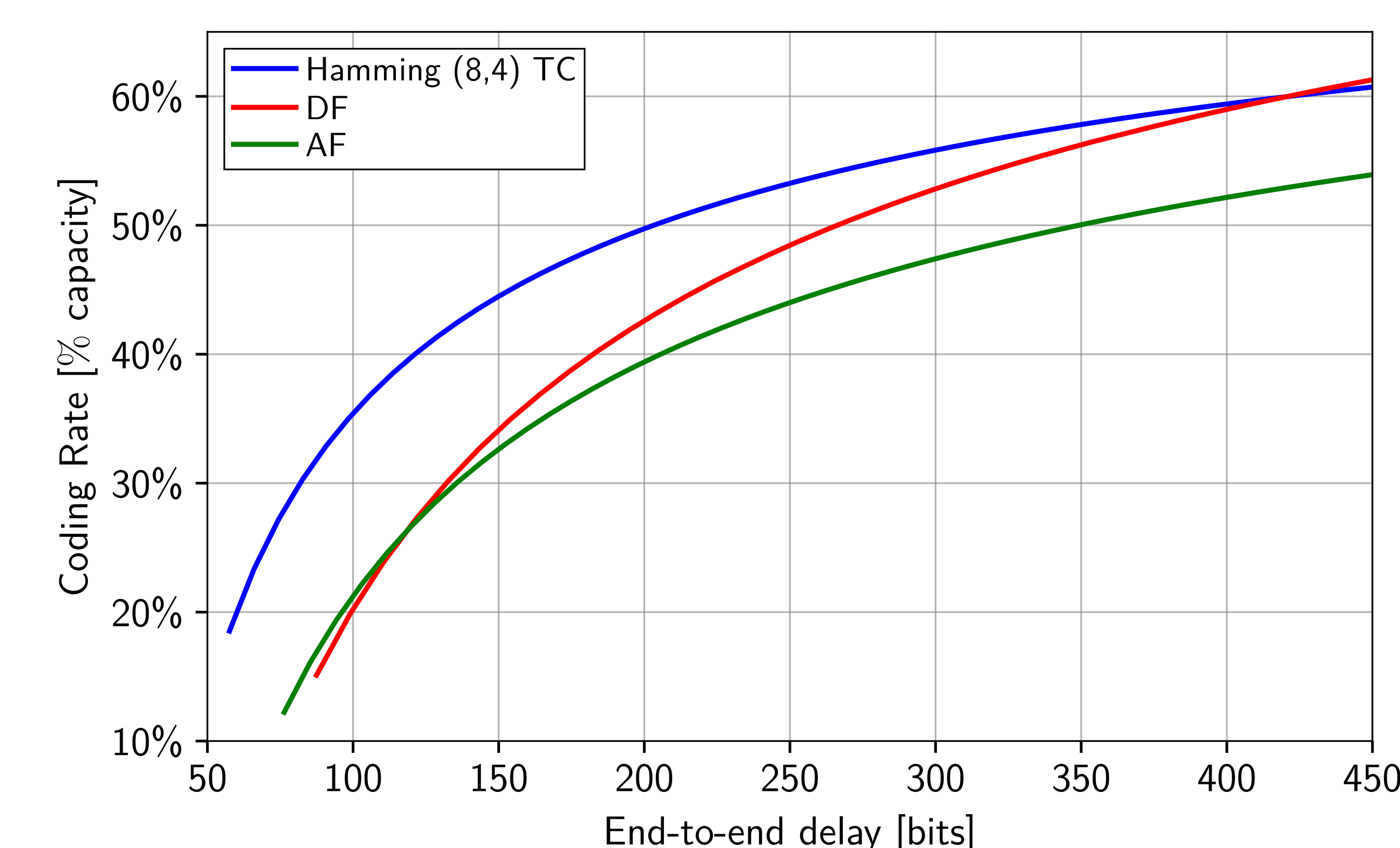
- ▶ Idea: Relay processes sub-blocks of size Δ
- ▶ Structure in sub-blocks \Rightarrow **Error control at relay**
- ▶ End-to-end delay: $T = T_1 + \Delta$
- ▶ Pipelined coding rate:

$$R_{TC}(T, \epsilon) = \frac{\log M}{T_1}$$



Example: $\Delta = 8$ [5]

- ▶ Binary symmetric channel on both hops
- ▶ Take sub-blocks from (8,4) extended Hamming Code
- ▶ Parameters: $p_1 = 0.04$, $p_2 = 0.13$, $\epsilon = 10^{-3}$
- ▶ Normal approximation [6] to evaluate rate-delay tradeoff



Ongoing Work

- ▶ Random code construction & analysis
- ▶ General transcoder design theory
- ▶ Extension to multi-hop & fading channels

- [1] T. Cover and A. E. Gamal, “Capacity theorems for the relay channel,” *IEEE Trans. Inform. Theory*, vol. 25, no. 5, pp. 572–584, Sep. 1979, ISSN: 0018-9448.
- [2] T. M. Cover and Y. H. Kim, “Capacity of a class of deterministic relay channels,” in *2007 IEEE International Symposium on Information Theory*, Jun. 2007, pp. 591–595.
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- [4] S. H. Lim, Y. H. Kim, A. E. Gamal, and S. Y. Chung, “Noisy network coding,” *IEEE Transactions on Information Theory*, vol. 57, no. 5, pp. 3132–3152, May 2011, ISSN: 0018-9448.
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- [6] Y. Polyanskiy, H. V. Poor, and S. Verdú, “Channel coding rate in the finite blocklength regime,” *IEEE Trans. Inform. Theory*, vol. 56, no. 5, pp. 2307–2359, May 2010, ISSN: 0018-9448.

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