Solving Problems by Searching

Inteligencia Artificial en los Sistemas de Control Autónomo Máster en Ciencia y Tecnología desde el Espacio

Departamento de Automática





Objectives

- 1. Understand the role of search in AI
- 2. Describe the importance of trees in search
- 3. Express AI problems in terms of search

4. Apply classical search algorithms

Bibliography

 S. Russell and P. Norvig. Chapter 3, Solving Problems by Searching. Artificial Intelligence: A Modern Approach. Pearson. 2017

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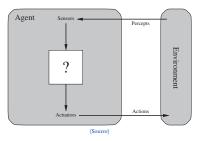
Introduction

Introduction

Intelligent agent

Agent

An agent is anything that can be viewed as perceiving its environment through sensors and acting through actuators



- Agents is a research field in AI by its own
 - ... with its own definition of agent (caution!)
- We use this term to abstract the implementation



Motivation

Introduction

Early AI works were directed to

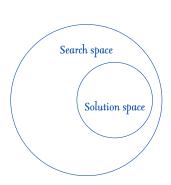
Proof of theorems, crosswords, games, ...

All in AI is search ...

- ... not entirely true (obviously) but more than we may imagine
- Find a good/best solution (solution space) to a problem among several potential solutions (search space)

Enhaustive search (or brute-force search)

- Iterate over all the potential solutions
- Unsuiteable for most real-world problems





Types of problems

Types of problems depending on ...

- Knowledge
 - Observable or Non-observable or Partially observable
- Outcome
 - Deterministic or Stochastic
- Actions
 - Discrete or Continous
- Time-variance
 - Static or Dynamic

We assume static, observable, discrete and deterministic problems



Types of problems (II)

Determine problem type

Chess



League of Legends



Observable or non-observable, deterministic or stochastic, discrete or continous, static or dynamic?

Problem formulation

Problem components (I)

We represent the environment as states

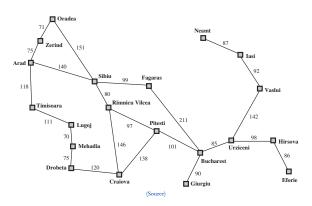
Contain the information about the world

Any problem formulation requires the following components

- Initial state: State where the search begins
- Actions: Behaviour that the agent may exhibit
- Transition model: Which states follow an action in a state (graph)
- Goal test (metas): How to determine if a state is a goal
- Path cost: Cost of a path to a state



Problem components: Example (I)



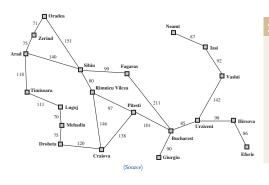
Problem: Move from Arad to Bucharest

On holiday in Romania; currently in Arad, flight leaves tomorrow from Bucharest Determine: Initial state, goal, states, actions, transition model, goal test and path cost



Problem formulation

Problem components: Example (II)



Solution

- Initial state: Arad
- Goal: Bucharest
- States: Multiple cities
- Actions: Drive between cities
- Goal test: In Bucharest?
- Path cost: Distance



Problem components: Search

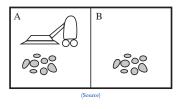
Problem formulation

Search is the process of finding a solution

- A solution is a sequence of actions leading from the initial state to a goal state
- Optimal solution is a solution with the lowest cost
- ullet Example of solution: Arad o Sibiu o Fagaras o Bucharest
 - That solution is optimal?



Toy problems (I): Vacuum world



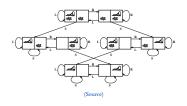
Problem: Clean rooms

- State? →
- Initial state? →
- Goal? \rightarrow
- Actions? \rightarrow
- Transition model? \rightarrow
- Goal test? \rightarrow
- Path cost? \rightarrow



Toy problems (I): Vacuum world

Problem formulation 0000000000000



- State? \rightarrow Dirt and location
- Initial state? \rightarrow All dirt, Left
- Goal? \rightarrow No dirt, any location
- Actions? → Left, Right, Suck
- Transition model? \rightarrow See figure
- Goal test? \rightarrow No dirt, any location
- Path cost? \rightarrow 1 per action



Toy problems (II): 8-puzzle





Start State

Goal State

(Source)

Problem: Solve 8-puzzle

- State? →
- Initial state? \rightarrow
- Goal? \rightarrow
- Actions? \rightarrow
- Transition model? \rightarrow
- Goal test? →
- Path cost? \rightarrow



Toy problems (II): 8-puzzle





Goal State

Start State

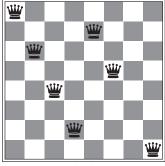
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Problem: Solve 8-puzzle

- State? \rightarrow Location of tiles 9!/2 = 181,440 states
- Initial state? \rightarrow Any
- Goal? \rightarrow See figure
- Actions? \rightarrow Left, Right, Up, Down
- Transition model? \rightarrow Complex graph
- Goal test? → Goal state
- Path cost? \rightarrow 1 per move



Toy problems (III): 8-queens



(Source)

State?

Initial state?

Goal?

Actions?

Transition model?

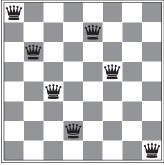
Goal test?

Path cost?



Toy problems (III): 8-queens

Problem formulation 00000000000000



(Source)

- State? \rightarrow Any arrangement of o to 8 queens
- Initial state? → Empty board
- Goal? \rightarrow See figure
- Actions? o Add queen to empty square
- Transition model? → Complex graph
- Goal test? \rightarrow 8 queens on board, none attacked
- Path cost? \rightarrow 1 per move



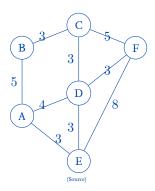
Travelling Salesman Problem (TSP)

TSP formulation

A travelling salesman must visit a set of cities only one time each. Find the shortest route.

TSP is a very big problem in AI!

- First formulated in 1930 and still a hot research topic!
- NP-hard problem
- Many real world applications



In general ...

- Each problem has a search graph, or state space
- Searching means finding a path from the initial state to a goal state

Basic idea

- Explore search space
- Generate a search tree (i.e., expanding nodes)

A search strategy is defined by picking the order of node exansion

- Uninformed search: Only uses the problem definition
- Informed search: Uses problem-specific knowledge



Search strategy (II)

Search strategies are evaluated along the following dimensions

- Completeness
- Time complexity
- Space complexity
- Optimality

Time and space are measured in terms of

- b: Maximum branching factor
- d: Depth of the least-cost solution
- m: Maximum depth of the state space



Uninformed search

Uninformed search algorithms

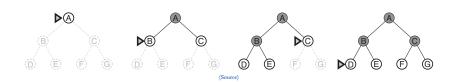
- Breadth-first search (búsqueda en anchura)
- Uniform-cost search (búsqueda de coste uniforme)
- Depth-first search (búsqueda en profundidad)
- Depth-limited search (búsqueda en profundidad limitada)
- Iterative deepening search (búsqueda de profundización iterativa)



Breadth-first search (I)

Expand shallowest unexpanded node

• Implemented with a FIFO queue (First-In First-Out)





Uninformed search

Breadth-first search (II)

Depth	Nodes	Time	Memory
2	110	1.1 milliseconds	107 kilobytes
4	11,110	111 milliseconds	10.6 megabytes
6	10^{6}	11 seconds	1 gigabytes
8	108	19 minutes	103 gigabytes
10	10^{10}	31 hours	10 terabytes
12	10^{12}	129 days	1 petabytes
14	10^{14}	35 years	99 petabytes
16	10^{16}	3,500 years	10 exabytes

Figure 3.13 Time and memory requirements for breadth-first search. The numbers shown assume branching factor b=10;100,000 nodes/second; 1000 bytes/node.

(Source)

Properties of breadth-first search

- Completeness: Yes
- Time complexity: O(b^{d+1})
- Space complexity: O(b^{d+1})
- Optimality: Yes (if cost = 1 per step)

Space is the biggest problem (more than time)



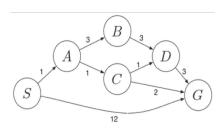
Uninformed search

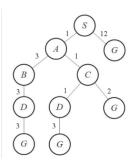
Uninformed search

Uniform-cost search (I)

Special case of breadth-first search

- Expand least-cost unexpanded node
- The queue is sorted by cost





Uniform-cost search (II)

```
Initialization: \{[S, 0]\}
Iter. I: \{[S \rightarrow A, 1], [S \rightarrow G, 12]\}
Iter. 2: \{[S \rightarrow A \rightarrow C, 2], [S \rightarrow A \rightarrow B, 4], [S \rightarrow G, 12]\}
Iter. 3:
\{[S \rightarrow A \rightarrow C \rightarrow D, 3], [S \rightarrow A \rightarrow C \rightarrow G, 4], [S \rightarrow A \rightarrow B \rightarrow D, 7], [S \rightarrow G, 12]\}
Iter. 4: \{[S \rightarrow A \rightarrow C \rightarrow D \rightarrow G, 6], [S \rightarrow A \rightarrow C \rightarrow G, 4], [S \rightarrow A \rightarrow B \rightarrow D, 7], [S \rightarrow G, 4], [S
G. 12]
Iter. 5: \{[S \rightarrow A \rightarrow C \rightarrow G, 4], [S \rightarrow A \rightarrow C \rightarrow D \rightarrow G, 10], [S \rightarrow G, 12]\}
Solution: S \to A \to C \to G
```

Uniform-cost search (III)

Properties

- Completeness: Yes, if step cost $\geq \epsilon$
- Time complexity: $O(b^{\lceil C^*/\epsilon \rceil})$, where C^* is the cost of the optimal solution
- Space complexity: $O(b^{\lceil C^*/\epsilon \rceil})$
- Optimality: Yes

Space is the biggest problem (more than time)

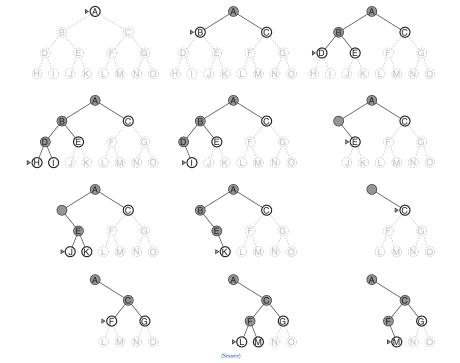


Depth-first search (I)

Expand deepest unexpanded node

• Implemented with a LIFO stack





Uninformed search

Depth-first search (III)

Properties of depth-first search

- Completeness: No, fail in infinite-depth spaces or spaces with loops
- Time complexity: $O(b^m)$, (terrible if m >> d)
- Space complexity: O(bm)
- Optimality: No



Depth-first search with depth limit L

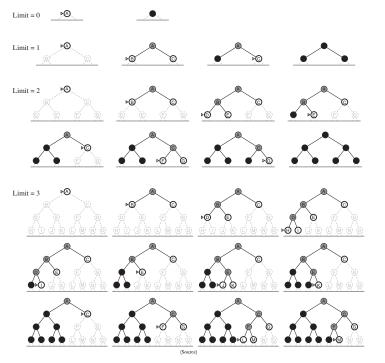
• Nodes at depth L are not expanded



Uninformed search Iterative deepening depth-first search (I)

Depth-limited search where gradually increases \boldsymbol{L}





Iterative deepening depth-first search (III)

Properties

- Completeness: Yes
- Time complexity: O(b^d)
- Space complexity: O(bd)
- Optimality: Yes if step cost = 1



Comparison of uninformed search algorithms

Criterion	Breadth-	Uniform-	Depth-	Depth-	Iterative
	First	Cost	First	Limited	Deepening
Complete	Yes*	Yes*	No	Yes, if $l \ge d$	Yes
Time	b^{d+1}	$b^{\lceil C^*/\epsilon ceil}$	b^m	$\mathbf{b^l}$	\mathbf{b}^{d}
Space	b^{d+1}	$b^{\lceil \mathrm{C}^*/\epsilon ceil}$	bm	bl	bd
Optimal	Yes^*	Yes	No	No	Yes^*



Introduction (I)

Use problem-specific knowledge beyond problem definition

- Best-first search (búsqueda primero el mejor)
 - Greedy best-first search (Búsqueda voraz)
 - A* search
- Local search algorithms
 - Hill-climbing search (búsqueda en escalada)
 - Simulated annealing search (búsqueda de temple simulado)
 - Local beam search (búsqueda de haz local)
 - Genetic Algorithms



Informed search

Beast-first search

Use an evaluation function f(n) for each node

- f is a cost estimate, i.t, its "desirability"
- Expand most desirable unexpanded nodes
- n is a node

Most algorithms use a heuristic function or just heuristic (h(n))

- Estimated cost from a state to the goal
- h only depends on the state (does not consider the path)
- h is a negative, nonnegative problem-specific function
- If n is a goal node, then h(n) = 0

The choice of f determines the search strategy

- Greedy best-first search
- A*



Greedy best-first search (I)

It only considers the heuristic

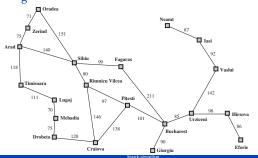
- f(n) = h(n)
- Remember: h is the estimate coast from a state to the goal

Greedy search

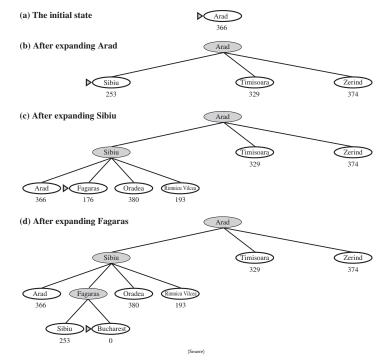
$$f(n) = h(n)$$

Example: Find a path between Arad and Bucharest

• Heuristic: Straight-line distance







$$A*(I)$$

Idea: avoid expanding paths that are already expensive

 It consider the path to a state (past), and its estimated cost to goal (future)

Evaluation function: f(n) = g(n) + h(n)

- g(n): Cost so far to reach node n
- h(n): Estimated cost from node n to goal
- f(n): Estimated total cost of path through n to goal

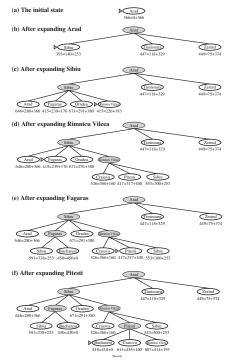
Theorem: A* is optimal if h(n) is admisible

- A* is admisible if it never overestimates the cost
- Example: Straight-line distance never overestimates road distance



$$f(n) = g(n) + h(n)$$





A* (III)

Properties

- Completeness: Yes
- Time complexity: Exponential
- Space complexity: Keeps all nodes in memory
- Optimality: Yes



Local search: Introduction (I)

In many optimization problems, the path to the gal is irrelevant; the goal itself is the solution

The path is stored in memory

Solution: Keep a single "current" state and try to improve it

• Generally, moving to the neighboring state

The paths followed by the search are not retained

- The use little memory
- They can find reasonable solutions in large or even infinite state spaces



Local search: Introduction (II)

Put n queens on an $n \times n$ board with no two queens on the same row, column or diagonal





Hill-climbing search (I)

Its just a loop that moves in the direction of increasing value

- Ends when it reaches a peak where no neighbor has a higher value
- No search tree is kept, just a datastructure with the current node

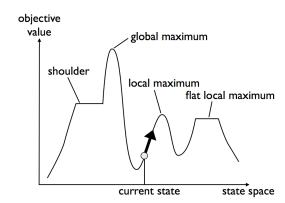
"Like climbing Everest in thick fog with amnesia"



Hill-climbing search (II)

Good for pure optimization problems

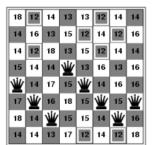
- No obvious cost function for such problems
- Objective function: How good a state is





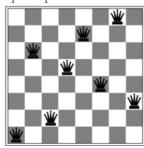
Hill-climbing search (III)

Hill climbing search: 8-queens problem



h = number of pairs of queens that are attacking each other

- h = 0 for the above state
- All successors h > 12



A local minimum (h = 1)



Hill-climbing search (IV)

The algorithm gets stuck for several reasons

- Local Maximum: it is a peak that is higher than each of its neighbours, but lower than the maximum overall
- Ridges: cause a sequence of local maxima that make navigation difficult
- Plateau (flat): no ascendant exit

In the 8-queens, it gets stuck in 86 %

• If we allow lateral movements $\rightarrow 100 \,\%$

Variants

- Stochastically chooses upward movements
- Random restart (random generation of initial state)



Simulated annealing search (I)

Process of tempering or hardeing metals by heating and then cooling them gradually

• Idea: escape local maxima by allowing some "bad" moves, but gradually decrease their frequency

Combination of hill-climbing with random generation successor

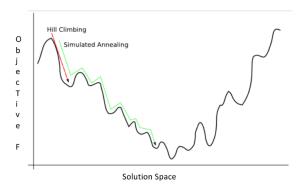
• Problem: determine its parameteters, need of experimentation

Application

- Good for problems with a large search space, optimum is surrounded by many local optima
- Widely used in VLSI layout, airline scheduling, etc



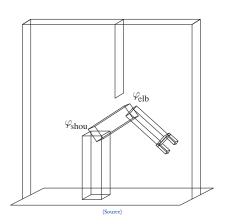
Simulated annealing search (II)





Case studies

Case study I: Robot arm with two DOF (I)

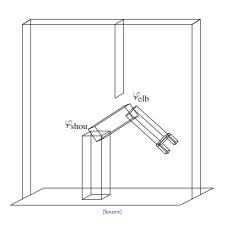


- State? \rightarrow
- Actions? \rightarrow
- Goal test? →
- Path cost? \rightarrow



Case studies

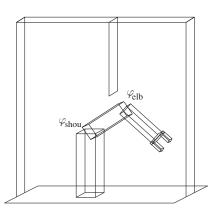
Case study I: Robot arm with two DOF (I)

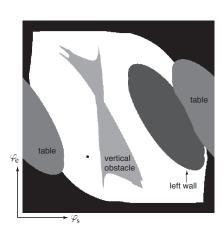


- State? \rightarrow Real-valued coordinates of robot joint angles
- Actions? → Continuous motions of joints
- Goal test? \rightarrow Complete assembly
- Path cost? → Time to complete



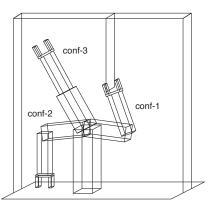
Case study I: Robot arm with two DOF (II)

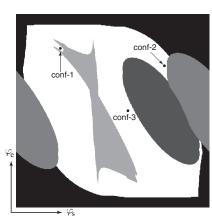




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Case study I: Robot arm with two DOF (III)

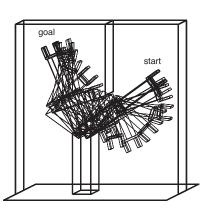


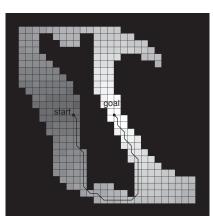


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Case study I: Robot arm with two DOF (IV)





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Case study II: 9th Global Trajectory Optimization Competition (I)

GTOC: Global Trajectory Optimization Competition

- Proposed by ESA Advanced Concepts Team
- Difficult trajectory optimization problems

GTOC 9: The Kesser Run

- 123 orbiting debris
- Remove debris
- Design multiple missions

(Video) (Solution)



Case study III: Mars orbital insertion

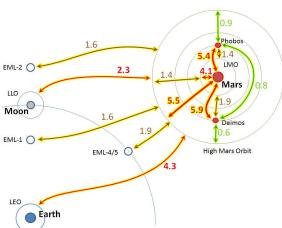


Chart by Richard Penn CC-BY, data from David Hollister hopsblog-hop.blogspot.co.uk

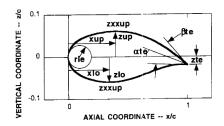


Study case

Case study IV: Transonic wing shape optimization

Problem: Design a wing shape for transonic flight

Maximize lift



Holst T.L., Pulliam T.H. (2003) Transonic Wing Shape Optimization Using a Genetic Algorithm. In: IUTAM Symposium Transsonicum IV. Fluid Mechanics and its Applications, vol 73. Springer.

