# Solving Problems by Searching

Inteligencia Artificial en los Sistemas de Control Autónomo Máster Universitario en Ingeniería Industrial

Departamento de Automática





#### Objectives

- 1. Understand the role of search in AI
- 2. Describe the importance of trees in search
- 3. Express AI problems in terms of search

# 4. Apply classical search algorithms

### Bibliography

 S. Russell and P. Norvig. Chapter 3, Solving Problems by Searching. Artificial Intelligence: A Modern Approach. Pearson. 2017

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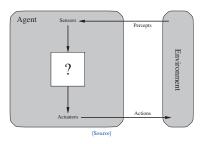
### Introduction

Introduction

### Intelligent agent

#### Agent

An agent is anything that can be viewed as perceiving its environment through sensors and acting through actuators



- Agents is a research field in AI by its own
  - ... with its own definition of agent (caution!)
- We use this term to abstract the implementation



# Motivation

Introduction

#### Early AI works were directed to

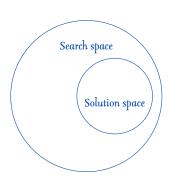
Proof of theorems, crosswords, games, ...

#### All in AI is search ...

- ... not entirely true (obviously) but more than we may imagine
- Find a good/best solution (solution space) to a problem among several potential solutions (search space)

#### Enhaustive search (or brute-force search)

- Iterate over all the potential solutions
- Unsuiteable for most real-world problems





### Types of problems

#### Types of problems depending on ...

- Knowledge
  - Observable or Non-observable or Partially observable
- Outcome
  - Deterministic or Stochastic
- Actions
  - Discrete or Continous
- Time-variance
  - Static or Dynamic

We assume static, observable, discrete and deterministic problems



### Types of problems (II)

#### Determine problem type

#### Chess



#### League of Legends



Observable or non-observable, deterministic or stochastic, discrete or continous, static or dynamic?

### Problem components (I)

#### We represent the environment as states

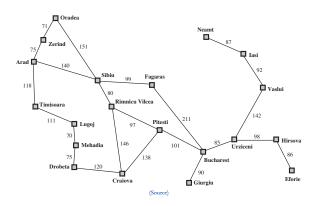
Contain the information about the world

Any problem formulation requires the following components

- Initial state: State where the search begins
- Actions: Behaviour that the agent may exhibit
- Transition model: Which states follow an action in a state (graph)
- Goal test (metas): How to determine if a state is a goal
- Path cost: Cost of a path to a state



### Problem components: Example (I)



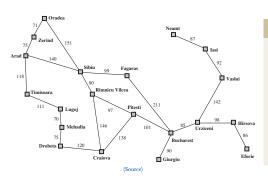
#### Problem: Move from Arad to Bucharest

On holiday in Romania; currently in Arad, flight leaves tomorrow from Bucharest Determine: Initial state, goal, states, actions, transition model, goal test and path cost



### Problem components: Example (II)

Problem formulation



#### Salution

- Initial state: Arad
- Goal: Bucharest
- States: Multiple cities
- Actions: Drive between cities
- Goal test: In Bucharest?
- Path cost: Distance



### Problem components: Search

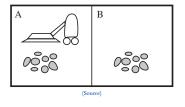
#### Search is the process of finding a solution

- A solution is a sequence of actions leading from the initial state to a goal state
- Optimal solution is a solution with the lowest cost
- ullet Example of solution: Arad o Sibiu o Fagaras o Bucharest
  - That solution is optimal?



### Toy problems (I): Vacuum world

Problem formulation

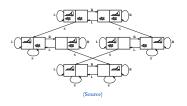


#### Problem: Clean rooms

- State? →
- Initial state? →
- Goal? →
- Actions?  $\rightarrow$
- Transition model?  $\rightarrow$
- Goal test? →
- Path cost?  $\rightarrow$



### Toy problems (I): Vacuum world



#### Problem: Clean rooms

- State?  $\rightarrow$  Dirt and location
- Initial state? ightarrow All dirt, Left
- Goal?  $\rightarrow$  No dirt, any location
- Actions?  $\rightarrow$  Left, Right, Suck
- Transition model? o See figure
- Goal test? o No dirt, any location
- Path cost?  $\rightarrow$  1 per action



### Toy problems (II): 8-puzzle





Start State

Goal State

(Source)

#### Problem: Solve 8-puzzle

- State?  $\rightarrow$
- Initial state?  $\rightarrow$
- Goal?  $\rightarrow$
- Actions?  $\rightarrow$
- Transition model?  $\rightarrow$
- Goal test? →
- Path cost?  $\rightarrow$



### Toy problems (II): 8-puzzle





Start State Goal State

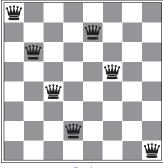
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#### Problem: Solve 8-puzzle

- State?  $\rightarrow$  Location of tiles 9!/2 = 181,440 states
- Initial state?  $\rightarrow$  Any
- Goal?  $\rightarrow$  See figure
- Actions?  $\rightarrow$  Left, Right, Up, Down
- Transition model?  $\rightarrow$  Complex graph
- Goal test?  $\rightarrow$  Goal state
- Path cost?  $\rightarrow$  1 per move



### Toy problems (III): 8-queens



(Source)

State?

Initial state?

Goal?

Actions?

Transition model?

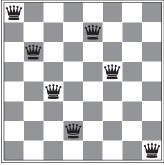
Goal test?

Path cost?



### Toy problems (III): 8-queens

Problem formulation



(Source)

# Problem: Place 8 queens no queen attacks any other

- State?  $\rightarrow$  Any arrangement of 0 to 8 queens
- Initial state? → Empty board
- Goal?  $\rightarrow$  See figure
- Actions? o Add queen to empty square
- Transition model?  $\rightarrow$  Complex graph
- Goal test? ightarrow 8 queens on board, none attacked
- Path cost?  $\rightarrow$  1 per move



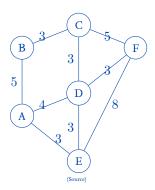
Travelling Salesman Problem (TSP)

### TSP formulation

A travelling salesman must visit a set of cities only one time each. Find the shortest route.

TSP is a very big problem in AI!

- First formulated in 1930 and still a hot research topic!
- NP-hard problem
- Many real world applications



#### In general ...

- Each problem has a search graph, or state space
- Searching means finding a path from the initial state to a goal state

#### Basic idea

- Explore search space
- Generate a search tree (i.e., expanding nodes)

A search strategy is defined by picking the order of node exansion

- Uninformed search: Only uses the problem definition
- Informed search: Uses problem-specific knowledge



### Search strategy (II)

#### Search strategies are evaluated along the following dimensions

- Completeness
- Time complexity
- Space complexity
- Optimality

Time and space are measured in terms of

- b: Maximum branching factor
- d: Depth of the least-cost solution
- m: Maximum depth of the state space



#### Uninformed search algorithms

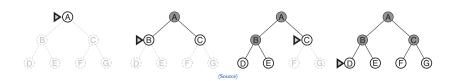
- Breadth-first search (búsqueda en anchura)
- Uniform-cost search (búsqueda de coste uniforme)
- Depth-first search (búsqueda en profundidad)
- Depth-limited search (búsqueda en profundidad limitada)
- Iterative deepening search (búsqueda de profundización iterativa)



### Breadth-first search (I)

#### Expand shallowest unexpanded node

• Implemented with a FIFO queue (First-In First-Out)





### Breadth-first search (II)

Depth	Nodes	Time		Memory	
2	110	1.1	milliseconds	107	kilobytes
4	11,110	111	milliseconds	10.6	megabytes
6	$10^{6}$	11	seconds	1	gigabytes
8	$10^{8}$	19	minutes	103	gigabytes
10	$10^{10}$	31	hours	10	terabytes
12	$10^{12}$	129	days	1	petabytes
14	$10^{14}$	35	years	99	petabytes
16	$10^{16}$	3,500	years	10	exabytes

Figure 3.13 Time and memory requirements for breadth-first search. The numbers shown assume branching factor b=10;100,000 nodes/second; 1000 bytes/node.

(Source)

#### Properties of breadth-first search

- Completeness: Yes
- Time complexity:  $O(b^{d+1})$
- Space complexity: O(b<sup>d+1</sup>)
- Optimality: Yes (if cost = 1 per step)

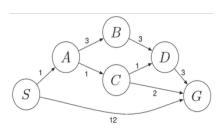
Space is the biggest problem (more than time)

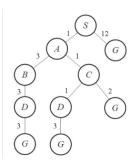


### Uniform-cost search (I)

#### Special case of breadth-first search

- Expand least-cost unexpanded node
- The queue is sorted by cost





#### Uniform-cost search (II)

```
Initialization: \{[S, 0]\}
Iter. I: \{[S \rightarrow A, 1], [S \rightarrow G, 12]\}
Iter. 2: \{[S \rightarrow A \rightarrow C, 2], [S \rightarrow A \rightarrow B, 4], [S \rightarrow G, 12]\}
Iter. 3:
\{[S \rightarrow A \rightarrow C \rightarrow D, 3], [S \rightarrow A \rightarrow C \rightarrow G, 4], [S \rightarrow A \rightarrow B \rightarrow D, 7], [S \rightarrow G, 12]\}
Iter. 4: \{[S \rightarrow A \rightarrow C \rightarrow D \rightarrow G, 6], [S \rightarrow A \rightarrow C \rightarrow G, 4], [S \rightarrow A \rightarrow B \rightarrow D, 7], [S \rightarrow G, 4], [S
G, 12
Iter. 5: \{[S \rightarrow A \rightarrow C \rightarrow G, 4], [S \rightarrow A \rightarrow C \rightarrow D \rightarrow G, 10], [S \rightarrow G, 12]\}
Solution: S \to A \to C \to G
```

### Uniform-cost search (III)

#### **Properties**

- Completeness: Yes, if step cost  $\geq \epsilon$
- Time complexity:  $O(b^{\lceil C^*/\epsilon \rceil})$ , where  $C^*$  is the cost of the optimal solution
- Space complexity:  $O(b^{\lceil C^*/\epsilon \rceil})$
- Optimality: Yes

Space is the biggest problem (more than time)

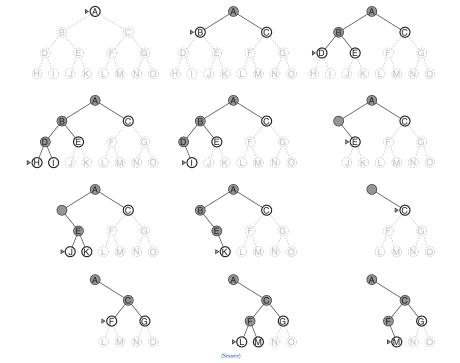


### Depth-first search (I)

Expand deepest unexpanded node

• Implemented with a LIFO stack





### Depth-first search (III)

#### Properties of depth-first search

- Completeness: No, fail in infinite-depth spaces or spaces with loops
- Time complexity:  $O(b^m)$ , (terrible if m >> d)
- Space complexity: O(bm)
- Optimality: No



Depth-limited search

Depth-first search with depth limit L

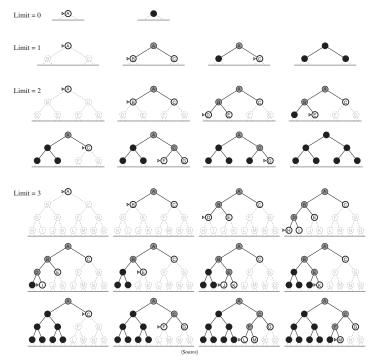
• Nodes at depth L are not expanded



Iterative deepening depth-first search (I)

Depth-limited search where gradually increases L





### Iterative deepening depth-first search (III)

#### Properties

- Completeness: Yes
- Time complexity: O(b<sup>d</sup>)
- Space complexity: O(bd)
- Optimality: Yes if step cost = 1



### Comparison of uninformed search algorithms

Criterion	Breadth-	Uniform-	Depth-	Depth-	Iterative
	First	Cost	First	Limited	Deepening
Complete	Yes*	Yes*	No	Yes, if $l \ge d$	Yes
Time	$b^{d+1}$	$b^{\lceil C^*/\epsilon  ceil}$	$b^m$	$\mathbf{b^l}$	$\mathrm{b}^{\mathrm{d}}$
Space	$b^{d+1}$	$b^{\lceil \mathrm{C}^*/\epsilon  ceil}$	bm	bl	bd
Optimal	$Yes^*$	Yes	No	No	$\mathrm{Yes}^*$



# Introduction (I)

#### Use problem-specific knowledge beyond problem definition

- Best-first search (búsqueda primero el mejor)
  - Greedy best-first search (Búsqueda voraz)
  - A\* search
- Local search algorithms
  - Hill-climbing search (búsqueda en escalada)
  - Simulated annealing search (búsqueda de temple simulado)
  - Local beam search (búsqueda de haz local)
  - Genetic Algorithms



### Beast-first search

Use an evaluation function f(n) for each node

- f is a cost estimate, i.t, its "desirability"
- Expand most desirable unexpanded nodes
- n is a node

Most algorithms use a heuristic function or just heuristic (h(n))

- Estimated cost from a state to the goal
- h only depends on the state (does not consider the path)
- h is a negative, nonnegative problem-specific function
- If n is a goal node, then h(n) = 0

The choice of f determines the search strategy

- Greedy best-first search
- A\*



### Greedy best-first search (I)

It only considers the heuristic

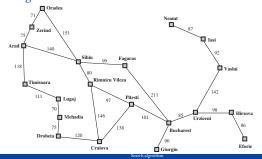
- f(n) = h(n)
- Remember: h is the estimate coast from a state to the goal

#### Greedy search

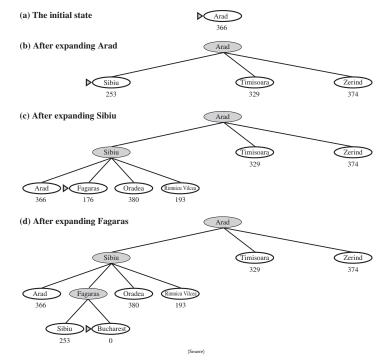
$$f(n) = h(n)$$

Example: Find a path between Arad and Bucharest

• Heuristic: Straight-line distance







Idea: avoid expanding paths that are already expensive

 It consider the path to a state (past), and its estimated cost to goal (future)

Evaluation function: f(n) = g(n) + h(n)

- g(n): Cost so far to reach node n
- h(n): Estimated cost from node n to goal
- f(n): Estimated total cost of path through n to goal

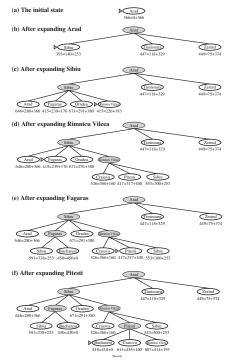
Theorem: A\* is optimal if h(n) is admisible

- A\* is admisible if it never overestimates the cost
- Example: Straight-line distance never overestimates road distance



$$f(n) = g(n) + h(n)$$





A\* (III)

#### Properties

- Completeness: Yes
- Time complexity: Exponential
- Space complexity: Keeps all nodes in memory
- Optimality: Yes



#### Local search: Introduction (I)

In many optimization problems, the path to the gal is irrelevant; the goal itself is the solution

The path is stored in memory

Solution: Keep a single "current" state and try to improve it

• Generally, moving to the neighboring state

The paths followed by the search are not retained

- The use little memory
- They can find reasonable solutions in large or even infinite state spaces



#### Local search: Introduction (II)

#### Example: n-queens

Put n queens on an n  $\times$  n board with no two queens on the same row, column or diagonal





### Hill-climbing search (I)

Its just a loop that moves in the direction of increasing value

- Ends when it reaches a peak where no neighbor has a higher value
- No search tree is kept, just a datastructure with the current node

"Like climbing Everest in thick fog with amnesia"



### Hill-climbing search (II)

#### Good for pure optimization problems

- No obvious cost function for such problems
- Objective function: How good a state is

