

## Lab 5 Ex 1 - SDP - Davide Gallitelli S241521

The goal of this exercise is to find out which are the system calls involved when running the command:

```
wc < test.txt | grep 1
```

The sequence starts with 22 `SYS_READ` calls, as many as the characters contained in the above mentioned command.

```
5 - SYS_read [x22]
[...]
```

Then, the shell performs a `SYS_FORK` call, so that its child will execute the command. After that, the shell executes a `SYS_WAIT`, in order to wait for child termination. The `SYS_sbrk` is the system call version of the `malloc()` function, which allocates memory to the newly created process.

```
1 - SYS_fork
3 - SYS_wait
12 - SYS_sbrk
```

The shell first child can now execute the command. This command is actually made of two commands piped: the child executes a `SYS_pipe` system call to create a nameless pipe to allow for interprocess communication.

```
4 - SYS_pipe
```

The following system calls are explained in the file `sh.c`, which contains information about shell execution. When a *pipe* instruction, or related system call, is received, the shell forks twice, creating two children, which will be referred to in the following lines as *CH1* and *CH2*, to handle the two commands, `wc < test.txt` and `grep 1`. Finally, it calls for a `SYS_wait` in order to wait for children to complete execution.

```
1 - SYS_fork
1 - SYS_fork
3 - SYS_wait
```

The two children created both perform four system calls: `CLOSE`, `DUP`, `CLOSE`, `CLOSE`.

```
// First children
21 - SYS_close
```

```
10 - SYS_dup
21 - SYS_close
21 - SYS_close
// Second children
21 - SYS_close
10 - SYS_dup
21 - SYS_close
21 - SYS_close
```

CH1, the child responsible for *wc*, performs a `SYS_exec` syscall to change its shell to the *wc* command. Then, it performs a `SYS_close` syscall to close the file descriptor related to *stdout*, because it will write on the pipe, and the one related to the read from the pipe.

```
21 - SYS_close
21 - SYS_close
7 - SYS_exec
```

Also the CH2, which will execute *grep*, performs `SYS_exec` and two `SYS_close`, closing the file descriptor related to *stdin*, as it will read the output from the pipe, and the one for writing on the pipe, because it will write on *stdout*.

```
21 - SYS_close
21 - SYS_close
7 - SYS_exec
```

The first step of the *wc* command is to open the file *test.txt*, which is done by a `SYS_open` syscall. It reads 512 bytes, receiving 31 bytes (one per character). Then it reads again from the file, returning 0 bytes meaning that it reached the EOF.

```
15 - SYS_open
// read outputs of the wordcount (1) (4) (31)
5 - SYS_read // read 512 bytes
5 - SYS_read // read 0 bytes - EOF
```

Then, CH1 can write on the pipe the 8 bytes generated as output from the *wc* command, then exits.

```
16 - SYS_write
16 - SYS_write
16 - SYS_write
16 - SYS_write
16 - SYS_write
16 - SYS_write
```

```
16 - SYS_write
16 - SYS_write
2 - SYS_exit
```

After being awoken by the child termination, it waits for the other child to die.

```
3 - SYS_wait
```

CH2 can now read from the pipe by means of a `SYS_read` syscall, and writes the output to *stdout*. Then reads again from the pipe to look for the end of the pipe, and it exits.

```
5 - SYS_read
16 - SYS_write // [ 1 4 31]
5 - SYS_read
2 - SYS_exit
```

The shell child now exits and returns control to the shell.

```
2 - SYS_exit
```