

# Adversarial Search and Game Playing

## Motivations

- Breadth-first, depth-first, hill-climbing, best-first, and A\* assume a non-hostile search space. •

The goals just sit there somewhere in the graph.

- The 8-puzzle game did not try to stop you from reaching the goal.
- But in a real 2-person game, your opponent does try to beat you and make it difficult for you to reach your goal.
- Game Trees complexity :-
  - Tic-Tac-Toe –  $9!$  – 362,880 states
  - Connect Four –  $10^{13}$  states
  - Checkers –  $10^{18}$  states
  - Chess –  $10^{50}$  states
  - Go –  $10^{170}$  states

## Typical assumptions in adversarial search

- Two agents whose actions alternate
- Utility values for each agent are the opposite of the other
  - creates the adversarial situation

- Fully observable environments
- In game theory terms:
  - “Deterministic, turn-taking, zero-sum games of perfect information”

## Search versus Games

- Search – no adversary
  - Solution is (heuristic) method for finding goal
  - can find *optimal* solution
  - Evaluation function: estimate of cost from start to goal through given node
  - Examples: path planning, scheduling activities
- Games – adversary
  - Solution is strategy (strategy specifies move for every possible opponent reply).
  - Time limits force an *approximate* solution
  - Evaluation function: evaluate “goodness” of game position
  - Examples: chess, checkers, Othello, backgammon

## Types of Games

	deterministic	chance
perfect information	chess, checkers, go, othello	backgammon monopoly
imperfect information		bridge, poker, scrabble nuclear war

- Imperfect information:- dont know all actions of other player ,don't recall all past moves of opponent or self .
- Solved games (for which best strategy is found)
  - Tic-tac-toe
  - Four In A Line
  - Checkers

## Size of search trees

- $b$  = branching factor

- $d$  = number of moves by both players
- Search tree is  $O(b^d)$
- Chess
  - $b \sim 35$
  - $D \sim 100$
  - completely impractical to search this
- Game-playing emphasizes being able to make optimal decisions in a finite amount of time
  - Somewhat realistic as a model of a real-world agent
  - Even if games themselves are artificial

## Static (Heuristic) Evaluation Functions

- An Evaluation Function:
  - estimates how good the current board configuration is for a player.
  - Typically, one figures how good it is for the player, and how good it is for the opponent, and subtracts the opponents score from the players
  - Othello: Number of white pieces - Number of black pieces
  - Chess: Value of all white pieces - Value of all black pieces

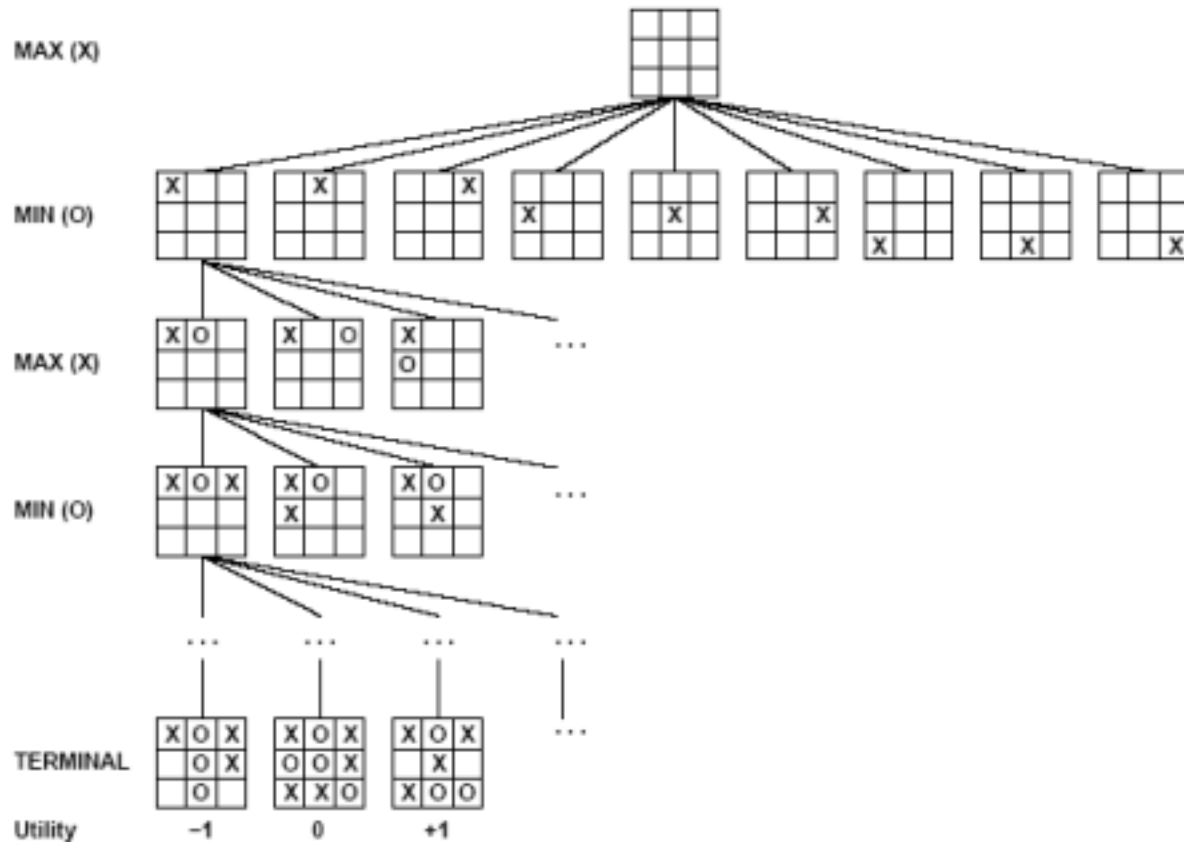
- Typical values from -infinity (loss) to +infinity (win) or  $[-1, +1]$ .
- If the board evaluation is  $X$  for a player, it's  $-X$  for the opponent
- Example:
  - Evaluating chess boards,
  - Checkers
  - Tic-tac-toe

## Game Setup

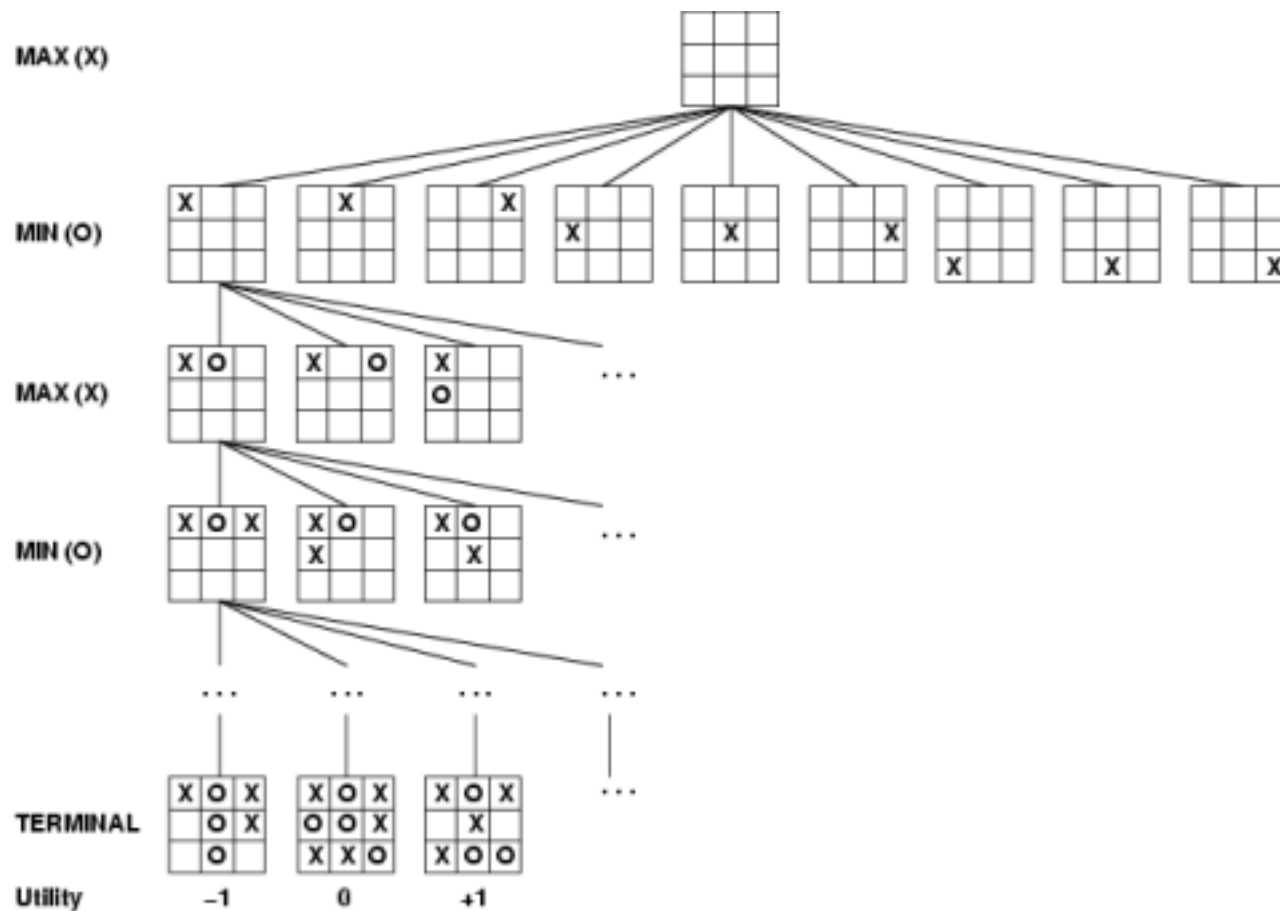
- Two players: MAX and MIN
- MAX moves first and they take turns until the game is over
  - Winner gets award, loser gets penalty.
- Games as search:
  - Initial state: e.g. board configuration of chess
  - Successor function: list of (move,state) pairs specifying legal moves.
  - Terminal test: Is the game finished?
  - Utility function: Gives numerical value of terminal states. E.g. win (+1), lose (-1) and draw (0) in tic-tac-toe or chess

- MAX uses search tree to determine next move.

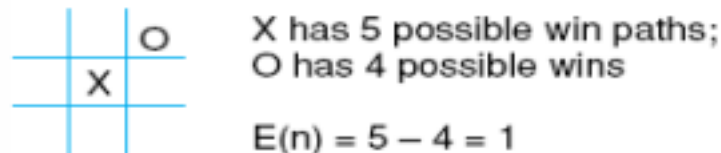
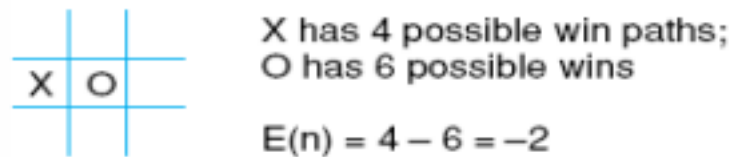
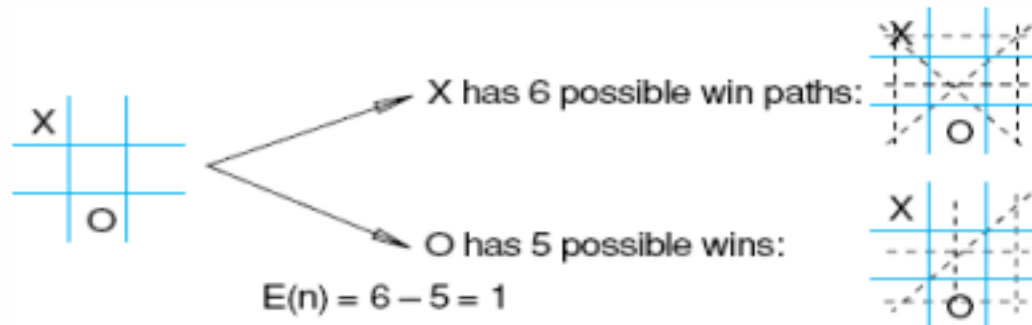
## Partial Game Tree for Tic-Tac-Toe



**Game tree (2-player, deterministic, turns)**



How  
do we search this tree to find the optimal move?  
Heuristic measuring for adversarial tic-tac-toe



Heuristic is  $E(n) = M(n) - O(n)$

where  $M(n)$  is the total of My possible winning lines

$O(n)$  is total of Opponent's possible winning lines

$E(n)$  is the total Evaluation for state  $n$

Maximize  $E(n)$

$E(n) =$

0 when my opponent and I have equal number of possibilities.

**Minimax strategy**



- Find the optimal *strategy* for MAX assuming an infallible MIN opponent
  - Need to compute this all the down the tree
- Assumption: Both players play optimally!
- Given a game tree, the optimal strategy can be determined by using the minimax value of each node:

**MINIMAX-VALUE( $n$ ) =**  
**{ UTILITY( $n$ ) If  $n$  is a terminal OR**  
**max (MINIMAX-VALUE( $s$ )) from all successors ( $s$ ) If  $n$  is a max node OR**  
**min (MINIMAX-VALUE( $s$ ) from all successors ( $s$ ) If  $n$  is a min node }**

**Simple game for example:**  
**Minimax decision**

MAX (player)

MIN  
(opponent)

3 12 8 2 4 6 14 5 2

**Simple game for  
example: Minimax  
decision**

	MIN
MAX (player)	(opponent)
	3

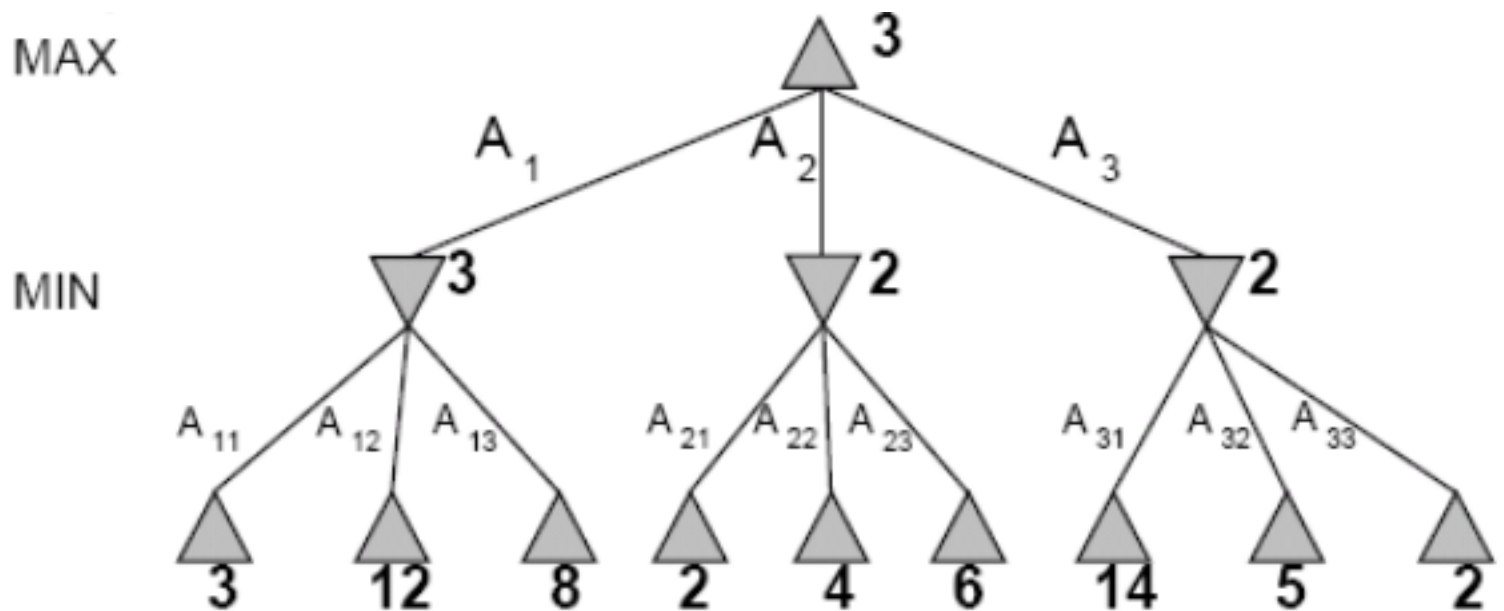
3 2 2

3 12 8 2 4 6 14 5 2

## **Example 2:- Two-Ply Game Tree**

**Minimax maximizes the utility for the worst-case outcome for max**

*The minimax decision*



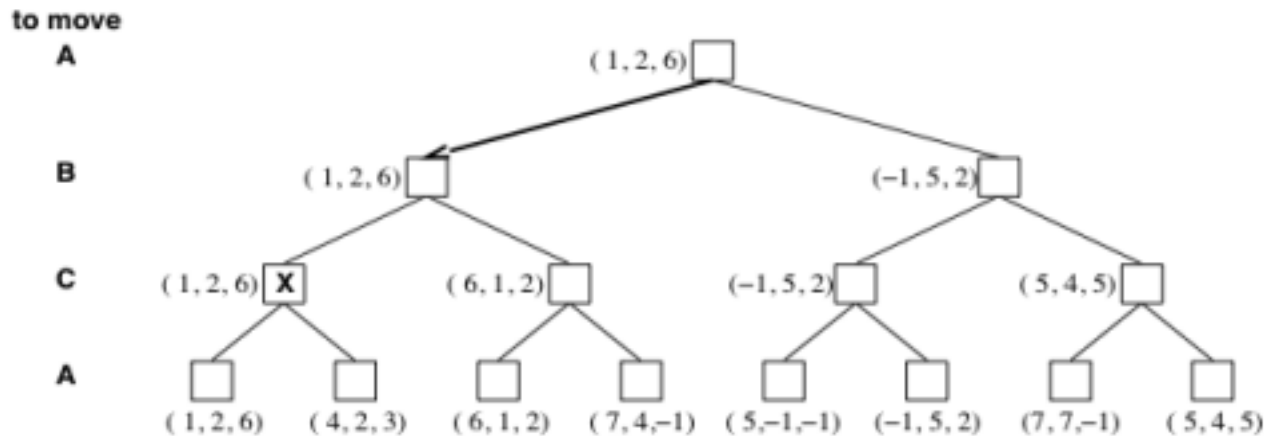
## What if MIN does not play optimally?

- Definition of optimal play for MAX assumes MIN plays optimally:
  - maximizes worst-case outcome for MAX
- But if MIN does not play optimally, MAX will do even better

## Multiplayer games

- Games allow more than two players
- Single minimax values become vectors

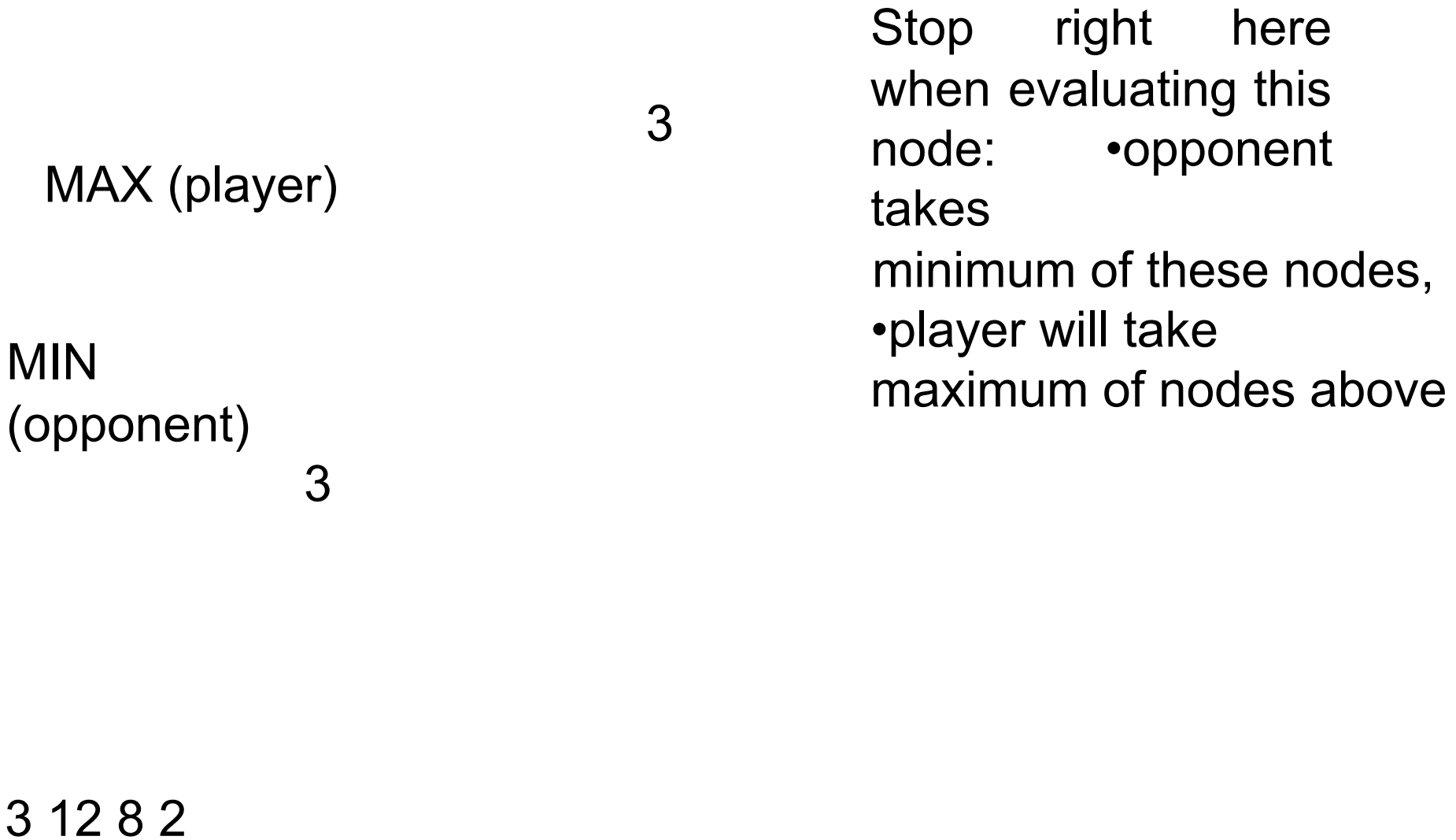
- standard minimax analysis assumes that each player operates to maximize only their own utility
- If game is not zero-sum (i.e.,  $\text{utility}(A) = -\text{utility}(B)$ ) then alliances can be useful even with 2 players
  - e.g., both cooperate to maximum the sum of the utilities



## Practical problem with minimax search

- Number of game states is exponential in the number of moves.
  - Solution: Do not examine every node
    - => pruning
      - Remove branches that do not influence final decision

# Alpha-Beta Pruning: Example



# Alpha-Beta Example



**Do DF-search until first**

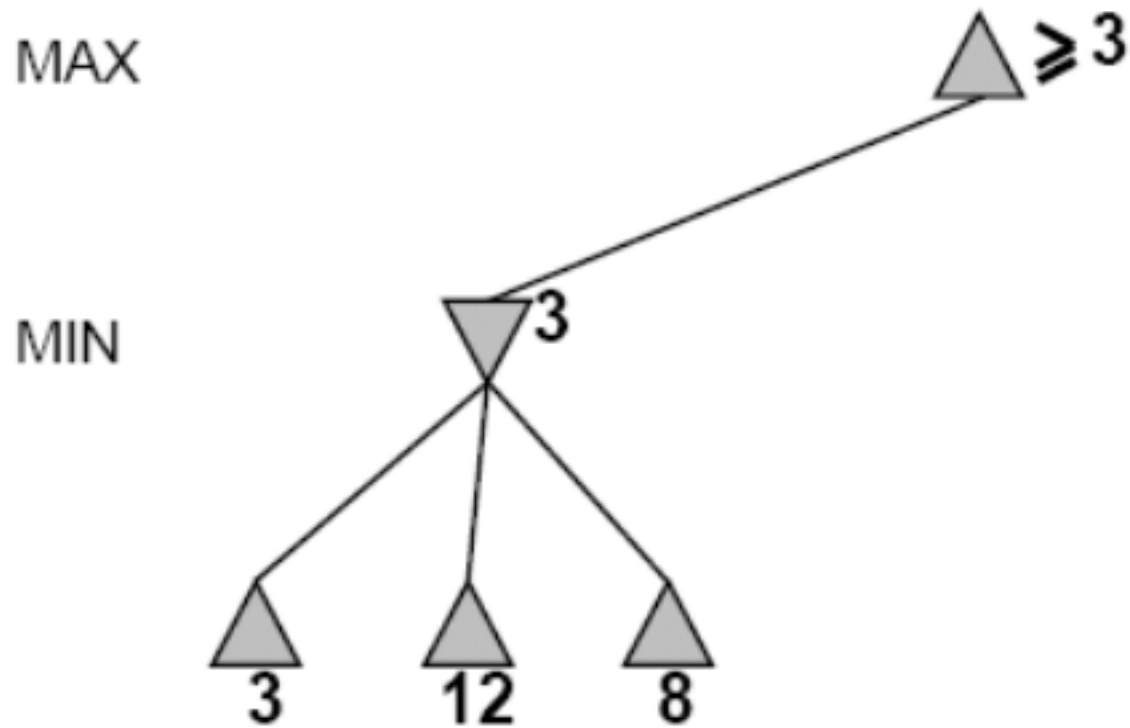
**leaf**

*Range of possible values*

$[-\infty, +\infty]$

$[-\infty, +\infty]$

## Alpha-Beta Example (continued)





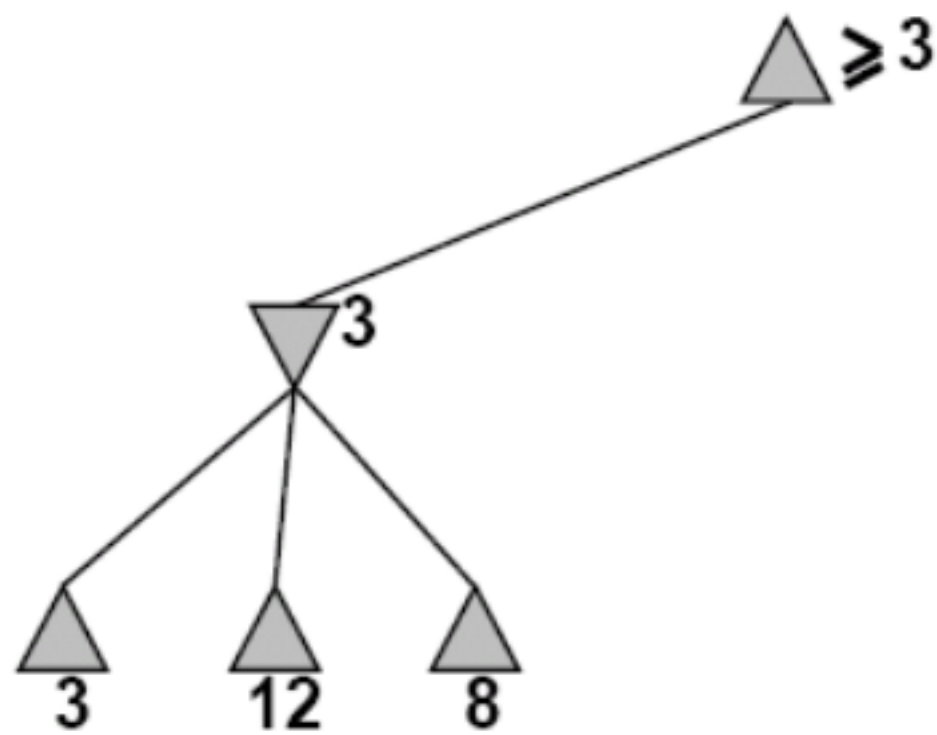
$[-\infty, 3]$

$[-\infty, +\infty] \leq 3$

**Alpha-Beta Example (continued)**

MAX

MIN



$$[-\infty, 3]$$

$$[-\infty, +\infty] \leq 3$$



ple (continued)

$$[3, +\infty]$$

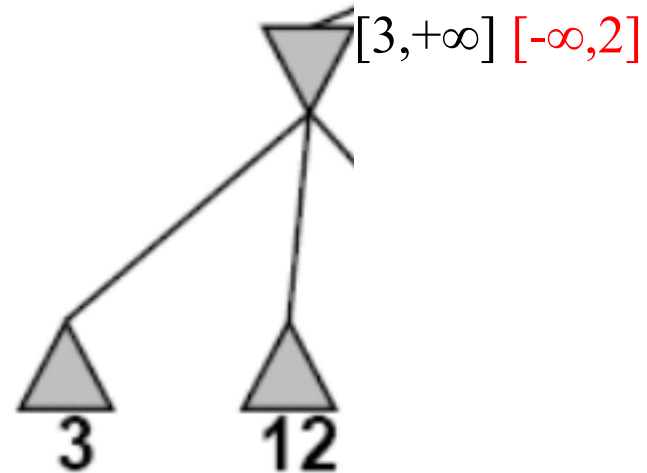
$$[3, 3]$$

# Alpha-Beta Example

(continued)

MAX

MIN



$[3,3]$

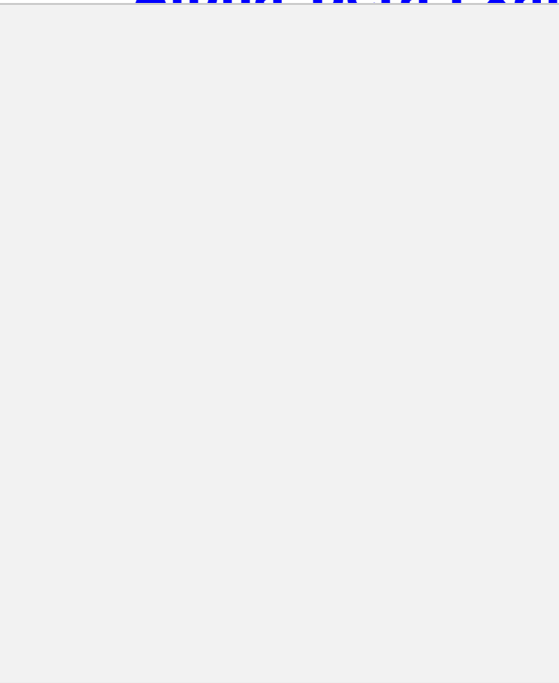
*This node is worse  
for MAX*

## Alpha-Beta Pruning: Concept

There are two rules for terminating search:

- Search can be stopped below any MIN node having a beta value less than or equal to the alpha value of any of its MAX ancestors.
- Search can be stopped below any MAX node having an alpha value greater than or equal to the beta value of any of its MIN ancestors.

### Alpha-Beta Example (continued)

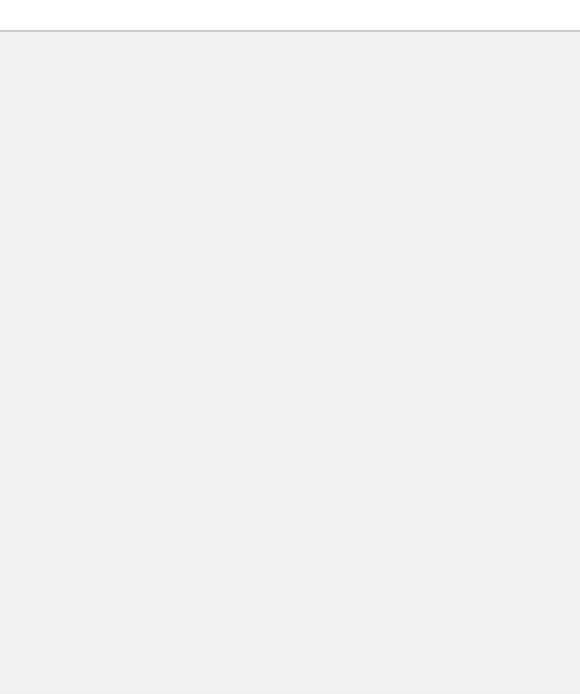


[3,14]

,

[3,3] [-∞,14] [-∞,2]

### Alpha-Beta Example (continued)



$[3,5]$

,

$[3,3]$   $[-\infty,5]$   $[-\infty,2]$

**Alpha-Beta Example (continued)**



[3,3]



[3,3]

$[-\infty, 2]$  [2,2]

**Alpha-Beta Example (continued)**





[3,3]



[3,3]

$[-\infty, 2]$  [2,2]

## General alpha-beta pruning

- Consider a node  $n$  somewhere in the tree
- If player has a better choice at – Parent node of  $n$ 
  - Or any choice point further up
- $n$  will **never** be reached in actual play.



- Hence when enough is known about  $n$ , it can be pruned.

## Alpha-beta Algorithm

- Depth first search – only considers nodes along a single path at any time

$\alpha$  = highest-value choice we have found at any choice point along the path for MAX

$\beta$  = lowest-value choice we have found at any choice point along the path for MIN

- update values of  $\alpha$  and  $\beta$  during search and prunes remaining branches as soon as the value is known to be worse than the current  $\alpha$  or  $\beta$  value for MAX or MIN

Minimax without pruning



Alpha-beta pruning example



F

G H I<sup>J</sup>

Unevaluated nodes

## Effectiveness of Alpha-Beta Search

- Worst-Case
  - branches are ordered so that no pruning takes place. In this case alpha-beta gives no improvement over exhaustive search
- Best-Case
  - each player's best move is the left-most alternative (i.e., evaluated first)
  - in practice, performance is closer to best rather than worst-case

- In practice often get  $O(b^{(d/2)})$  rather than  $O(b^d)$  – this is the same as having a branching factor of  $\sqrt{b}$ , • since  $(\sqrt{b})^d = b^{(d/2)}$ 
  - i.e., we have effectively gone from  $b$  to square root of  $b$ 
    - e.g., in chess go from  $b \sim 35$  to  $b \sim 6$
  - this permits much deeper search in the same amount of time

## Final Comments about Alpha-Beta

**Pruning** • Pruning does not affect final

results

- Entire subtrees can be pruned.
- Good move *ordering* improves effectiveness of pruning

## Iterative (Progressive) Deepening

- In real games, there is usually a time limit  $T$  on making a move

- In practice, iterative deepening search (IDS) is used – IDS runs depth-first search with an increasing depth-limit – when the clock runs out we use the solution found at the previous depth limit

## **The State of Play**

- Checkers:
  - Chinook ended 40-year-reign of human world champion Marion Tinsley in 1994.
- Chess:
  - Deep Blue defeated human world champion Garry Kasparov in a six-game match in 1997.
- Othello:
  - human champions refuse to compete against computers: they are too good.
- Go:
  - human champions refuse to compete against computers: they are

too bad

- $b > 300$  (!)

## Deep Blue

- 1957: Herbert Simon
  - “within 10 years a computer will beat the world chess champion”
- 1997: Deep Blue beats Kasparov
- Parallel machine with 30 processors for “software” and 480 VLSI processors for “hardware search”
- Searched 126 million nodes per second on average
  - Generated up to 30 billion positions per move
  - Reached depth 14 routinely
- Uses iterative-deepening alpha-beta search with transpositioning
  - Can explore beyond depth-limit for interesting moves

## Example :- Complete game tree for Nim-7

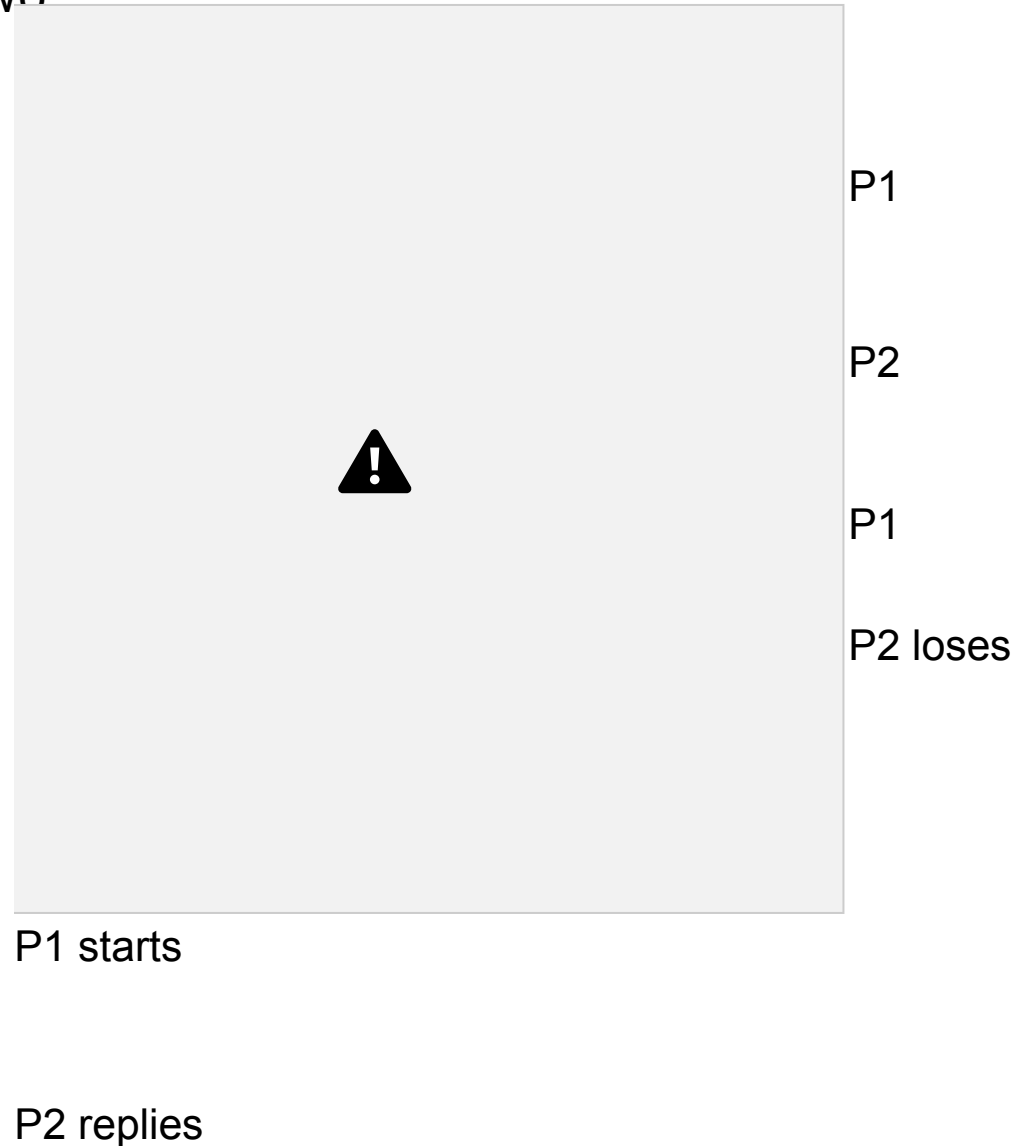


✓ 7 coins are placed on a table between the two opponents

✓ A move consists of dividing a pile of coins into two nonempty piles of different sizes

✓ For example, 6 coins can be divided into piles of 5 and 1 or 4 and 2, but not 3 and 3

✓ The first player who can no longer make a move loses the game



P2 loses P1 loses

## MIN vs. MAX in a Nim game

The best that MIN (Player1) can do is to lose unless Player2 makes a mistake.

moves first to minimize

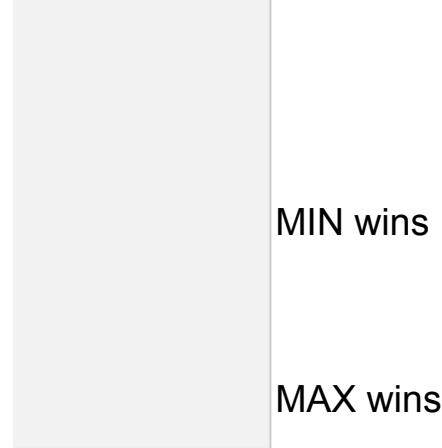
to maximize

Node score = 0  
means MIN wins.

1 means MAX wins.

Bold edges indicate forced win for MAX, Player2.

**What should be the winning strategy ?**



MIN wins

## Summary

- Game playing can be effectively modeled as a search problem
- Game trees represent alternate computer/opponent moves
- Evaluation functions estimate the quality of a given board configuration for the Max player.
- Minimax is a procedure which chooses moves by assuming that the opponent will always choose the move which is best for them

- Alpha-Beta is a procedure which can prune large parts of the search tree and allow search to go deeper
- For many well-known games, computer algorithms based on heuristic search match can out-perform human world experts.