

## Computer Networking I

#### **David Marchant**

Based on slides compiled by Marcos Vaz Salles, with adaptions by Vivek Shah and Michael Kirkedel Thomsen

## A quick note about me

- PhD student from Scotland
- Jeg taler ikke dansk :(
- 'David' on the Discord
- Office: HCØ, Building B, room 772-01-0-S06

# Christmas makes everything awkward

Week	Mon	Tue	Wed	Thu	Fri	Sat	Sun
50				Lecture Exercises	Cafe		
51		Lecture A4 out		Xmas			
52							
1		Lecture		Lecture Exercises A5 out	Cafe A4 in		
2		Lecture		Lecture Exercises	Cafe A5 in		
3							
4	Exam week						

# **Networking is Relevant**



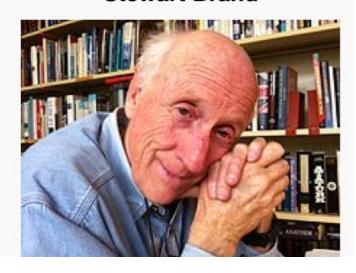
Information wants to be free because it has become so cheap to distribute, copy, and recombine... It wants to be expensive because it can be immeasurably valuable to the recipient. (1985)



Google news

WikipediA

#### Stewart Brand









#### What's the Internet: "nuts and bolts" view



PC



server



wireless laptop

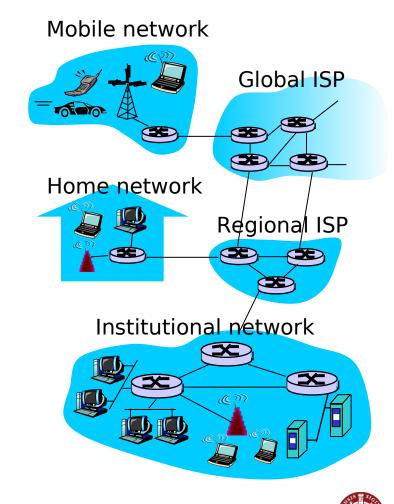


cellular handheld





- A network of networks
- Millions of connected computing devices: hosts = end systems
  - running network apps
- communication links
  - fiber, copper, radio, satellite
  - transmission rate = bandwidth
- routers: forward packets (chunks of data)





#### Protocols

- Speaking the same language
- Syntax and semantics

#### Layering

- Standing on the shoulders of giants
- A key to managing complexity

#### Resource allocation

- Dividing scarce resources among competing parties
- Memory, link bandwidth, wireless spectrum, paths

#### Naming

What to call computers, services, protocols, ...



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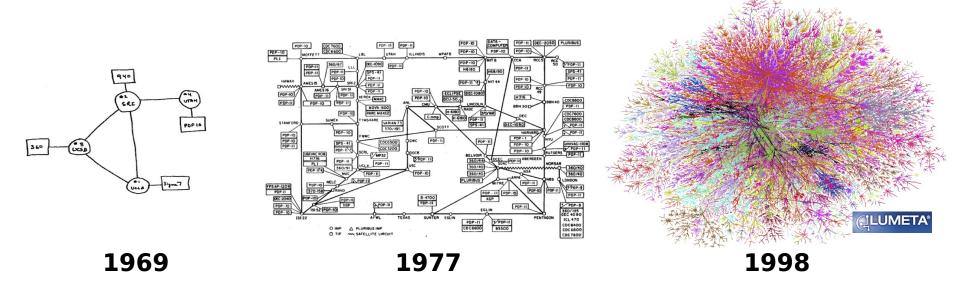
#### Naming

• What to call computers, services, protocols, ...



Source: Freedman

- Protocols
  - Speaking the same language
  - Syntax and semantics



All speak IPv4
"Internet Protocol version 4"



#### Protocol design is about tradeoffs

#### How should hosts and routers communicate?

- Standard protocol
- Fast: Machine readable in hardware at line rates

#### Browsers, web servers, and proxies?

- Can be slower: software readable
- Human readable
- Extensible and forward-compatible
- Not everybody might be familiar with extensions



### **IPv4 Packet**



4-bit	oit Header Service (TOS)		16-bit Total Length (Bytes)			
16-bit Identification			3-bit Flags	13-bit Fragment Offset		
8-bit Time to Live (TTL) 8-bit Protoco		8-bit Protocol	16-bit Header Checksum			
32-bit Source IP Address						
32-bit Destination IP Address						
Options (if any)						
Payload						

20-byte header

#### **Example: HyperText Transfer Protocol**

GET /courses/archive/spr09/cos461/ HTTP/1.1

Host: www.cs.princeton.edu

User-Agent: Mozilla/4.03

**CRLF** 

Request

#### HTTP/1.1 200 OK

Date: Mon, 2 Feb 2009 13:09:03 GMT

Server: Netscape-Enterprise/3.5.1

Last-Modified: Mon, 21 Feb 2009 11:12:23 GMT

Content-Length: 42

**CRLF** 

Site under construction





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Source: Freedman

#### Layering = Functional Abstraction

- Sub-divide the problem
  - Each layer relies on services from layer below
  - Each layer exports services to layer above
- Interface between layers defines interaction
  - Hides implementation details
  - Layers can change without disturbing other layers

**Application** 

Application-to-application channels

Host-to-host connectivity

Link hardware



**Application Layer** 

Transport Layer

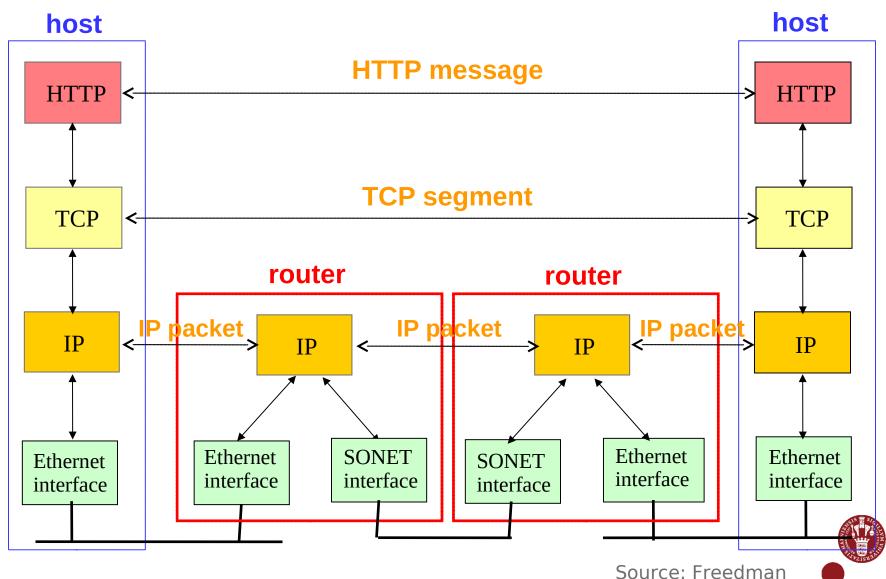
**Network Layer** 

Link Layer

**Physical Layer** 

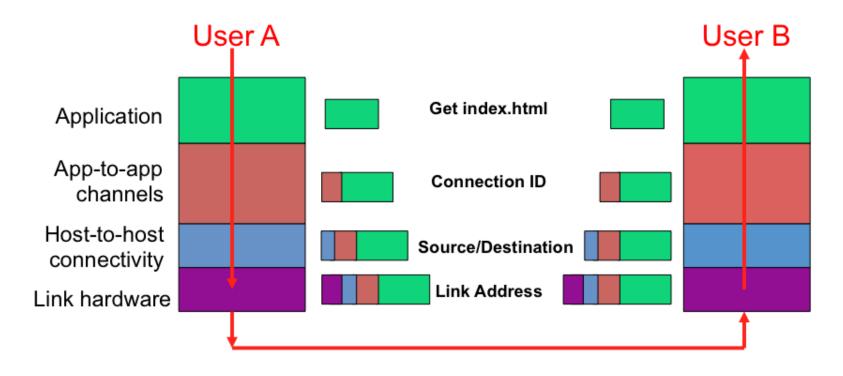


#### IP Suite: End Hosts vs. Routers



### **Layer Encapsulation in HTTP**

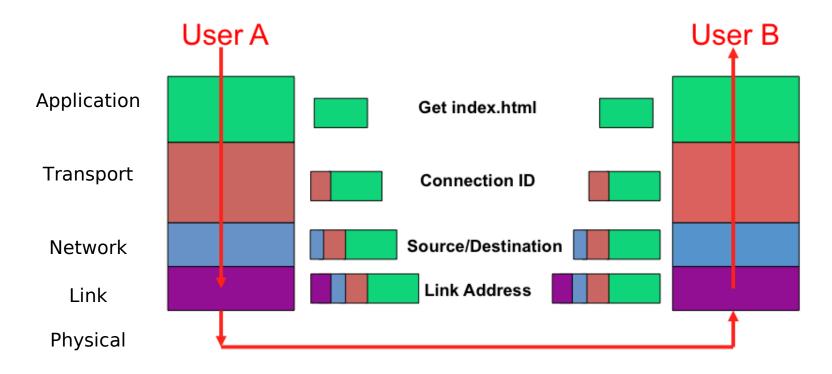






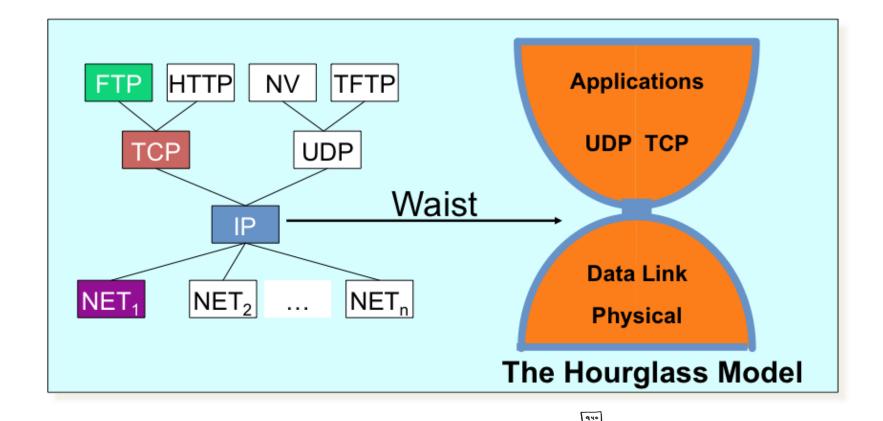
### **Layer Encapsulation in HTTP**







#### The Internet Protocol Suite



The waist facilitates interoperability



Source: Freedman

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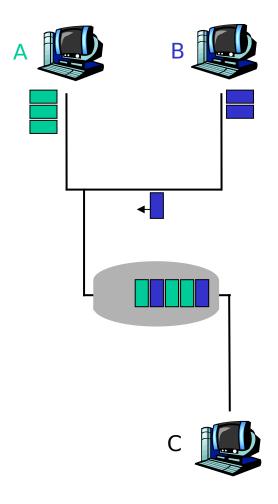


Source: Freedman

## **Network Core: Packet Switching**

## Data broken up into smaller 'packets'

- Can be sent seperately,
   each using full bandwidth
- Full bandwidth used on a single packet
- Packets from different sources are interspersed
- Resources only used when needed
- Unreliable transfer





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Source: Freedman

#### A Programmers View of the Internet

- 1. Hosts are mapped to a set of 32-bit IP addresses
  - **1**28.2.203.179
- 2. The set of IP addresses is mapped to a set of identifiers called Internet domain names
  - 128.2.217.3 is mapped to www.cs.cmu.edu
- 3. A process on one Internet host can communicate with a process on another Internet host over a *connection*



#### IPv4 and IPv6

- The original Internet Protocol, with its 32-bit addresses, is known as Internet Protocol Version 4 (IPv4)
- 1996: Internet Engineering Task Force (IETF) introduced Internet Protocol Version 6 (IPv6) with 128-bit addresses
  - Intended as the successor to IPv4
- As of 2021, majority of Internet traffic still carried by IPv4
  - Only 30-35% of users access Google services using IPv6.
- We will focus on IPv4, but will show you how to write networking code that is protocol-independent.



## IPv4, IPv6 and beyond

IPv0	Reserved			
IPv1-3	Initial attempts, became IPv4 and TCP('73-'79)			
IPv4	First actually deployed protocol ('81)			
5	Experiment ('90 -'95) into video streaming protocol			
IPv6	Improvement on IPv4 ('95)			
IPv7	Incorrectly assigned when IPv6 was assumed to already be in use('88)			
8	PIP, obsolete replacement for IPv4 ('94)			
9	TUBA, obsolete replacement for IPv4 ('92),			
IPv9	April Fools Joke by IETF ('94),			
IPv9	Chinese research project, seemingly abandoned ('04)			
IPv11-14	Unassigned			
IPv15	Reserved			

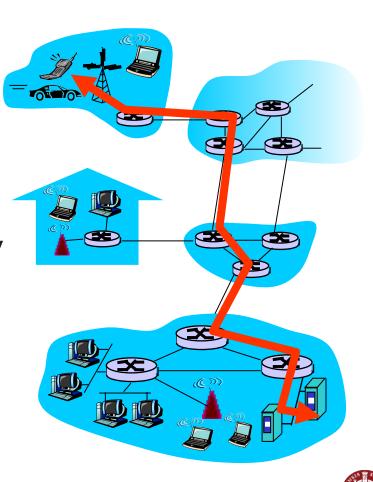


# Using the network



## **Network Core: Circuit Switching**

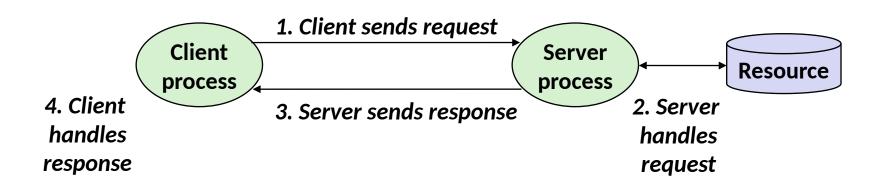
- To send a message we can use libaries at the application layer
- Underlying layers will take care of most of the work
- Can be done in (almost) any lanaguage. C is difficult though, so we'll use Python





## **A Client-Server Transaction**

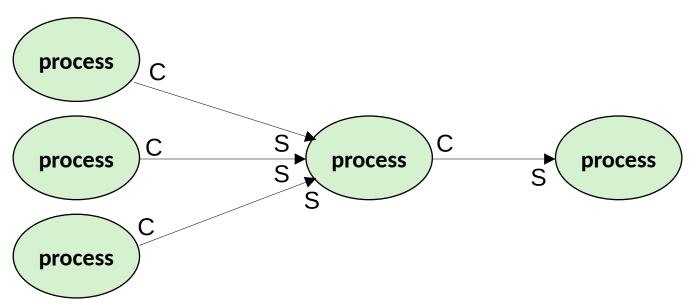
- Most network applications are based on the client-server model:
  - A server process and one or more client processes
  - Server manages some resource
  - Server provides service by manipulating resource for clients
  - Server activated by request from client (vending machine analogy)



Note: clients and servers are processes running on hosts (can be the same or different hosts)

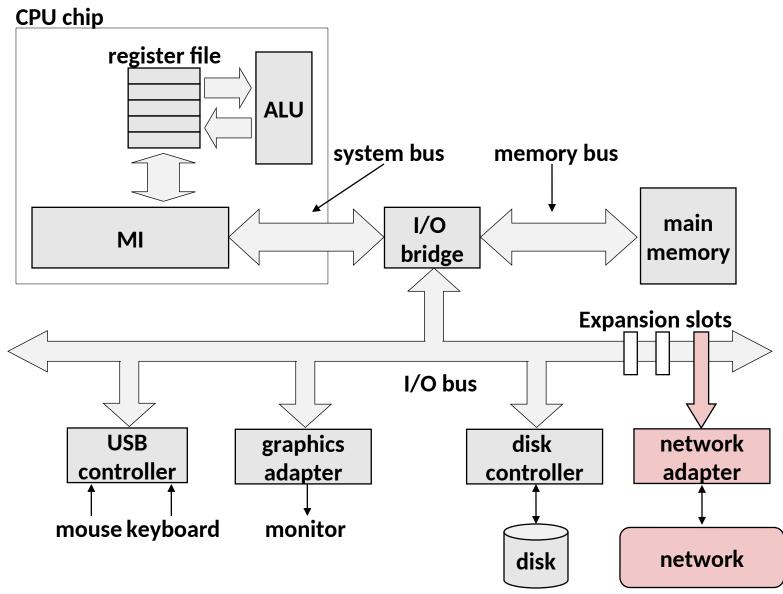
## **A Client-Server Transaction**

- Most network applications are based on the client-server model:
  - There can be multiple clients connected to a single server
  - Individual hosts can act as both *clients* and *servers* at the same time



Note: clients and servers are processes running on hosts (can be the same or different hosts)

## Hardware Organization of a Network Host



#### **Internet Connections**

- Clients and servers communicate by sending streams of bytes over connections. Each connection is:
  - Point-to-point: connects a pair of processes.
  - Full-duplex: data can flow in both directions at the same time,
  - Reliable: stream of bytes sent by the source is eventually received by the destination in the same order it was sent.
- A socket is an endpoint of a connection
  - Socket address is an IPaddress:port pair
- A port is a 16-bit integer that identifies a process:
  - **Ephemeral port:** Assigned automatically by client kernel when client makes a connection request.
  - Well-known port: Associated with some service provided by a server (e.g., port 80 is associated with Web servers)

## **Well-known Ports and Service Names**

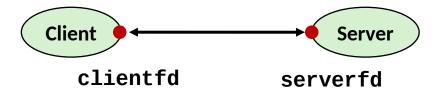
- Popular services have permanently assigned well-known ports and corresponding well-known service names:
  - echo server: 7/echo
  - ssh servers: 22/ssh
  - email server: 25/smtp
  - Web servers: 80/http
- Mappings between well-known ports and service names is contained in the file /etc/services on each Linux machine.

## **Sockets Interface**

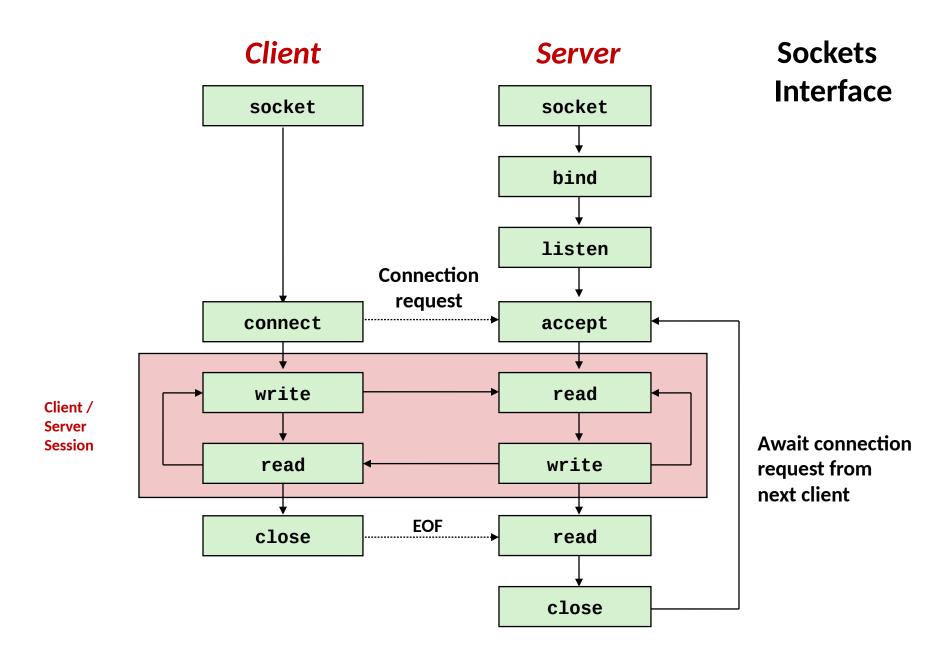
- Set of system-level functions used in conjunction with Unix I/O to build network applications.
- Created in the early 80's as part of the original Berkeley distribution of Unix that contained an early version of the Internet protocols.
- Available on all modern systems
  - Unix variants, Windows, OS X, IOS, Android, ARM

#### **Sockets**

- What is a socket?
  - To the kernel, a socket is an endpoint of communication
  - To an application, a socket is a file descriptor that lets the application read/write from/to the network
    - Remember: All Unix I/O devices, including networks, are modeled as files
- Clients and servers communicate with each other by reading from and writing to socket descriptors



The main distinction between regular file I/O and socket I/O is how the application "opens" the socket descriptors



#### Sockets Interface: socket

Clients and servers use the socket function to create a socket descriptor:

```
int socket(int domain, int type, int protocol)
```

Example:

C:

```
#include <sys/socket.h>
int socket_fd = socket(AF_INET, SOCK_STREAM, 0);
```

Indicates that we are using 32-bit IPV4 addresses

Indicates that the socket will be the end point of a connection

## **Python:**

```
from socket import *
with socket(AF_INET, SOCK_STREAM) as sock:
...
```

# Sockets Interface: listen

- By default, kernel assumes that descriptor from socket function is an active socket that will be on the client end of a connection.
- A server calls the listen function to tell the kernel that a descriptor will be used by a server rather than a client:

```
int listen(int sockfd, int backlog);
```

Converts sockfd from an active socket to a listening socket that can accept connection requests from clients.

```
C: listen(socket_fd, 10);
```

Python: sock.listen(10)

# Sockets Interface: accept

Servers wait for connection requests from clients by calling accept:

```
int accept(int listenfd, SA *addr, int *addrlen);
```

- Waits for connection request to arrive on the connection bound to listenfd, then fills in client's socket address in addr and size of the socket address in addrlen.
- Returns a connected descriptor that can be used to communicate with the client via Unix I/O routines.

```
socklen_t clientlen;
struct sockaddr_storage clientaddr;
conn_fd = accept(socket_fd, (SA *) &clientaddr, &clientlen);
```

```
Python: conn_addr = sock.accept()
```

# Sockets Interface: connect

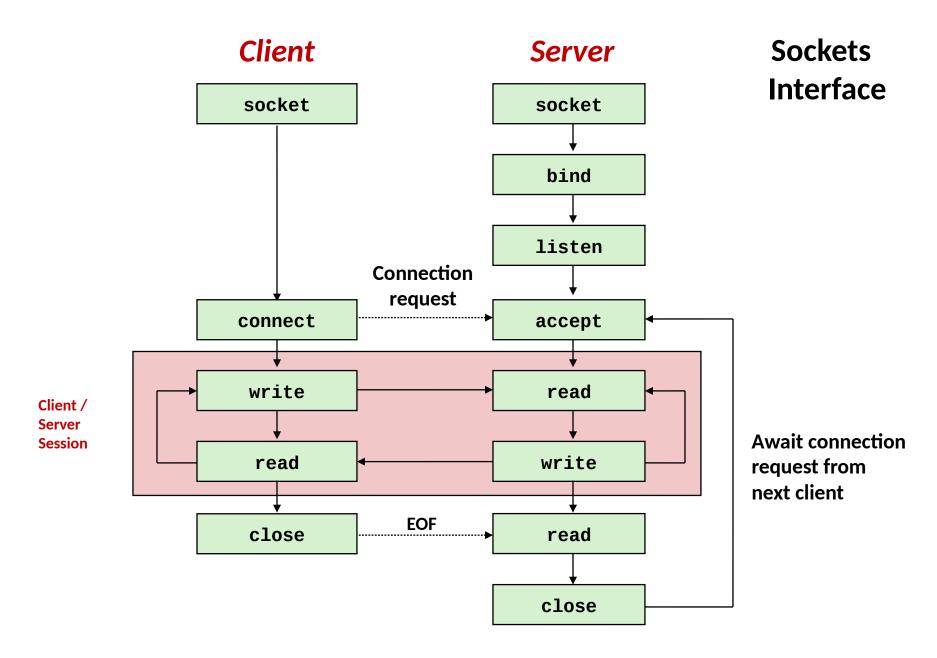
A client establishes a connection with a server by calling connect:

```
int connect(int clientfd, SA *addr, socklen_t addrlen);
```

- Attempts to establish a connection with server at socket address addr
  - If successful, then clientfd is now ready for reading and writing.
  - Resulting connection is characterized by socket pair (x:y, addr.sin\_addr:addr.sin\_port)
    - x is client address
    - y is ephemeral port that uniquely identifies client process on client host

```
struct sockaddr s_addr;
connect(socket_fd, (struct sockaddr *)&s_addr, sizeof(s_addr));
```

Python: client\_sock.connect("130.226.237.173", 56)



# **Python Example**

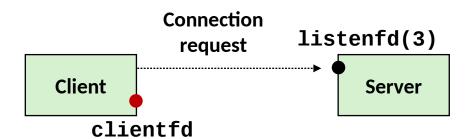
#### **Client:**

```
import socket
with socket.socket(socket.AF_INET, socket.SOCK_STREAM) as client_socket:
    client_socket.connect(("127.0.0.1", 5678))
    request = bytearray("This is a message".encode())
    client_socket.sendall(request)
    response = client_socket.recv(1024)
    print(response)
                                                  Note that
 Server:
                                                  these lines
import socket
                                                  are where our
with socket.socket(socket.AF_INET, socket.SOCK_ST Processes
                                                                  socket:
                                                  synchronise
    server_socket.bind(("127.0.0.1", 5678))
    server_socket.listen()
    while True
        connection, connection_address = server_sock!t.accept()
        with connection:
            message = connection.recv(1024)
            connection.sendall(response)
```

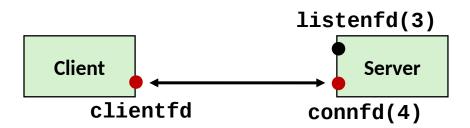
# accept Illustrated



1. Server blocks in accept, waiting for connection request on listening descriptor listenfd



2. Client makes connection request by calling and blocking in connect



3. Server returns connfd from accept. Client returns from connect. Connection is now established between clientfd and connfd

# **Connected vs. Listening Descriptors**

#### Listening descriptor

- End point for client connection <u>requests</u>
- Created once and exists for lifetime of the server

### Connected descriptor

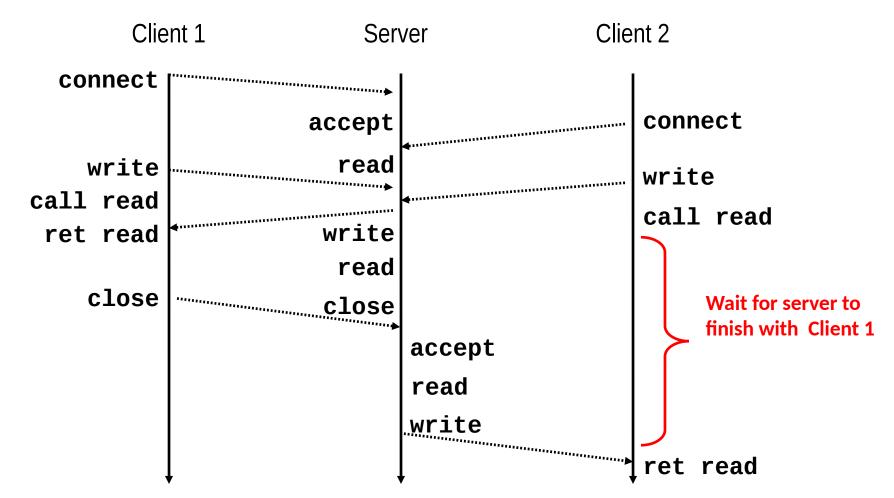
- End point of the <u>connection</u> between client and server
- A new descriptor is created each time the server accepts a connection request from a client
- Exists only as long as it takes to service client

#### Why the distinction?

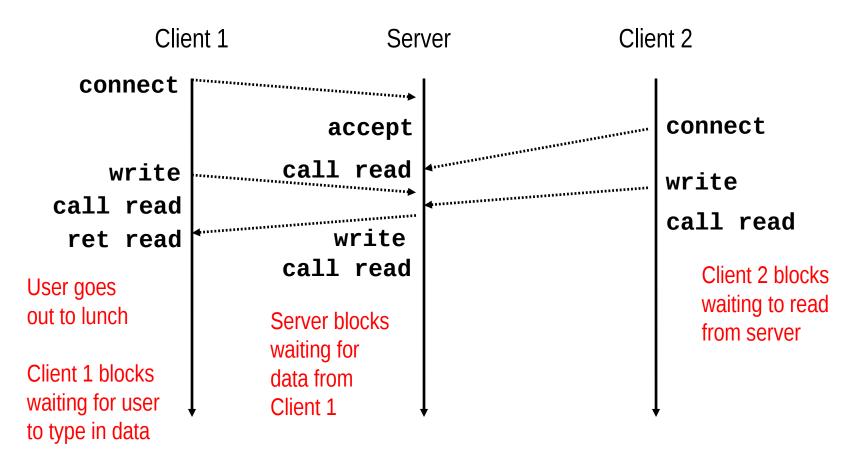
- Allows for concurrent servers that can communicate over many client connections simultaneously
  - E.g., Each time we receive a new request, we fork a child to handle the request

## **Iterative Servers**

Iterative servers process one request at a time

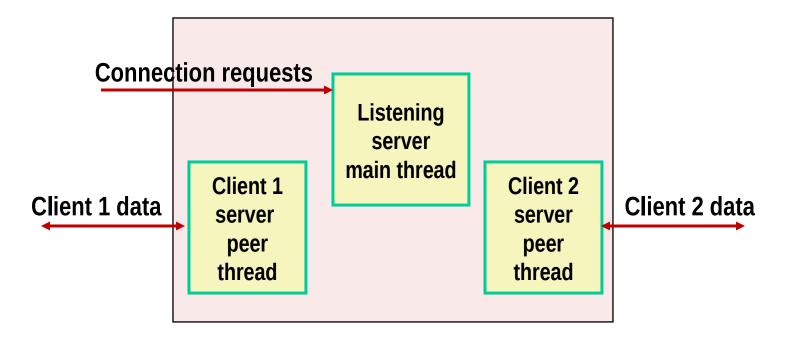


## **Fundamental Flaw of Iterative Servers**



- Solution: use concurrent servers instead
  - Concurrent servers use multiple concurrent flows to serve multiple clients at the same time

# **Thread-based Server Execution Model**



- Each client handled by individual peer thread
- Threads share all process state except TID
- Each thread has a separate stack for local variable
- We'll get into what this actually means after Christmas . . .

- Use socketserve library
- Two parts: Server class, and Handler class
- Server is a class to run continuously, handle errors, and spawn new threads to respond to each new request
- Handler is what actually reads in new messages and processes them

#### **Iterative Server:**

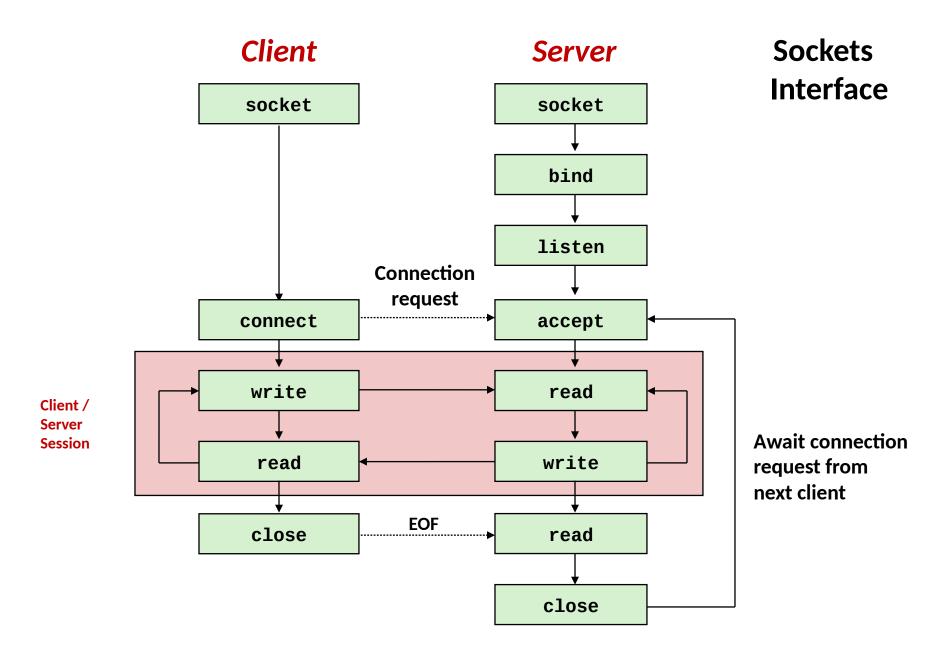
```
import socket
with socket.socket(socket.AF_INET, socket.SOCK_STREAM) as server_socket:
    server_socket.bind(("127.0.0.1", 5678))
    server_socket.listen()
    while True:
        connection, connection_address = server_socket.accept()
        with connection:
            message = connection.recv(1024)
            connection.sendall(response)
```

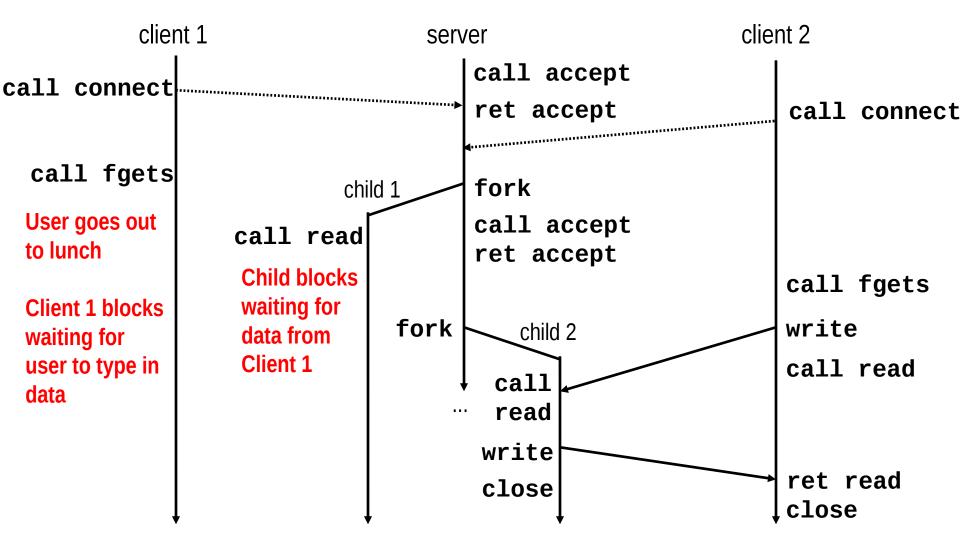
#### **Concurrent Server:**

```
from socketserver import ThreadingTCPServer, StreamRequestHandler

class MyHandler(StreamRequestHandler):
    def handle(self) -> None:
        message = self.request.recv(1024)
        self.request.sendall(message)

with ThreadingTCPServer(("127.0.0.1", 5678), MyHandler) as my_server:
    my_server.serve_forever()
```





- Demonstrated in exercises for this week
- This allows for Non-blocking servers. Very good for performance as all clients can get respodned too without a huge queue.
- We have not really talked about threading yet (and might not till after you've finished the networking assignment), so don't get too concerned with the details just yet.