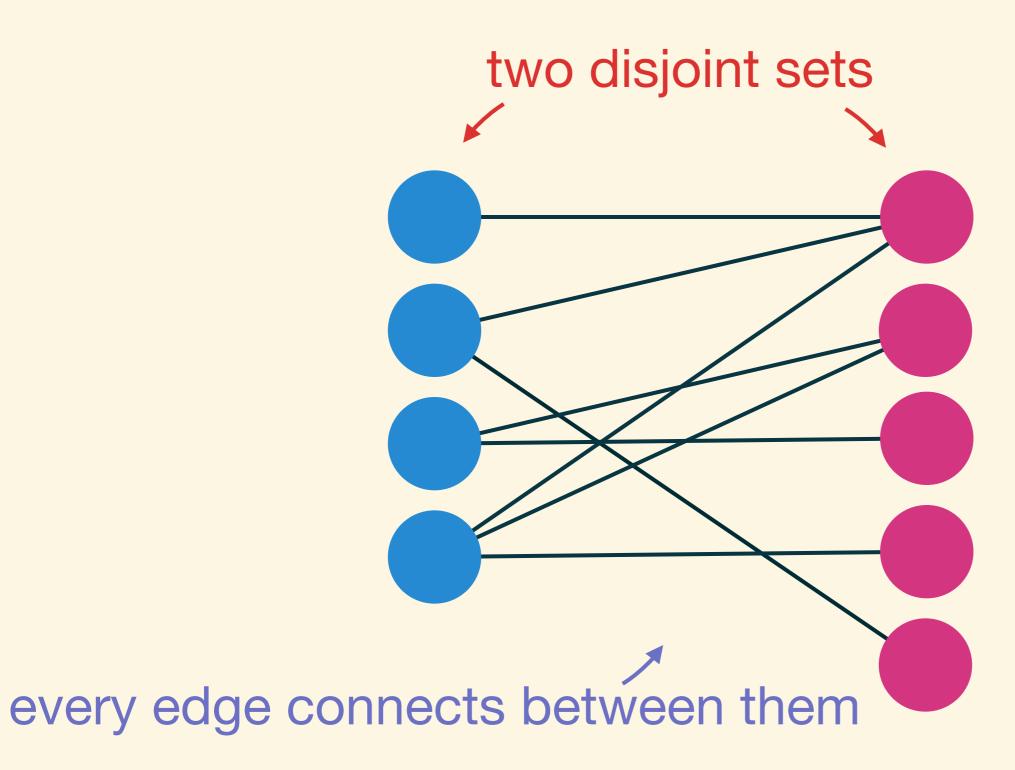
Bipartite Matching

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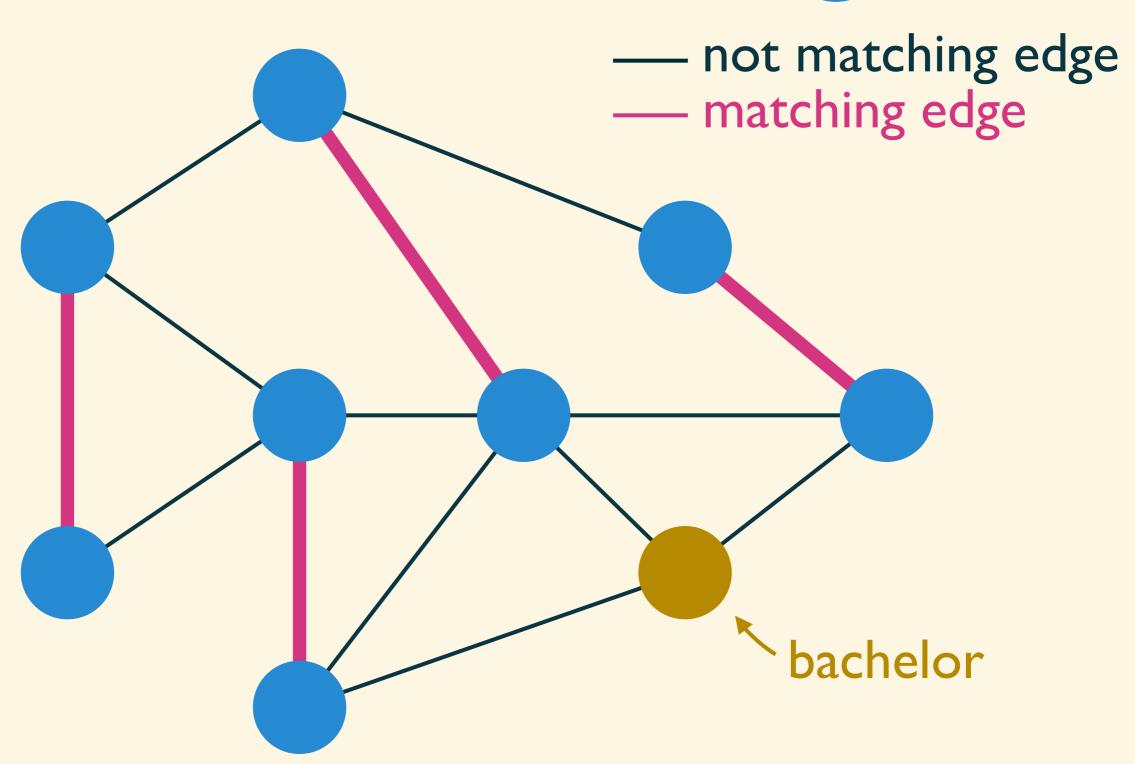


Bipartite Graph

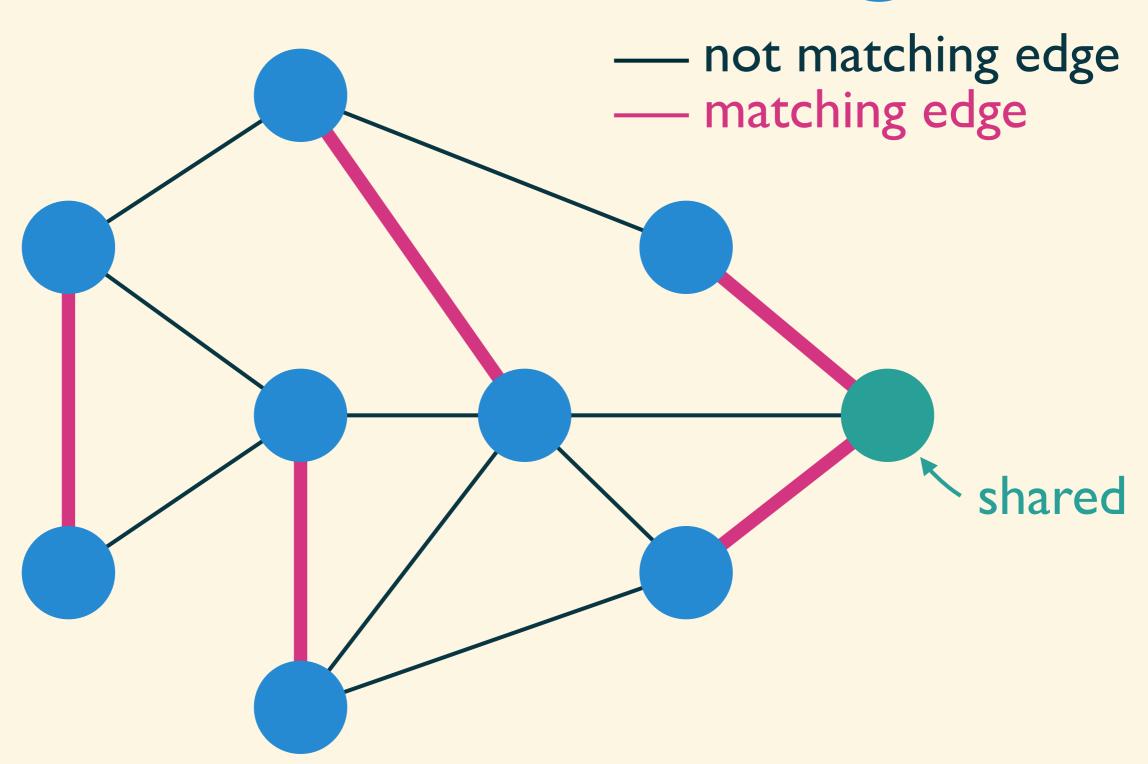




Valid Matching

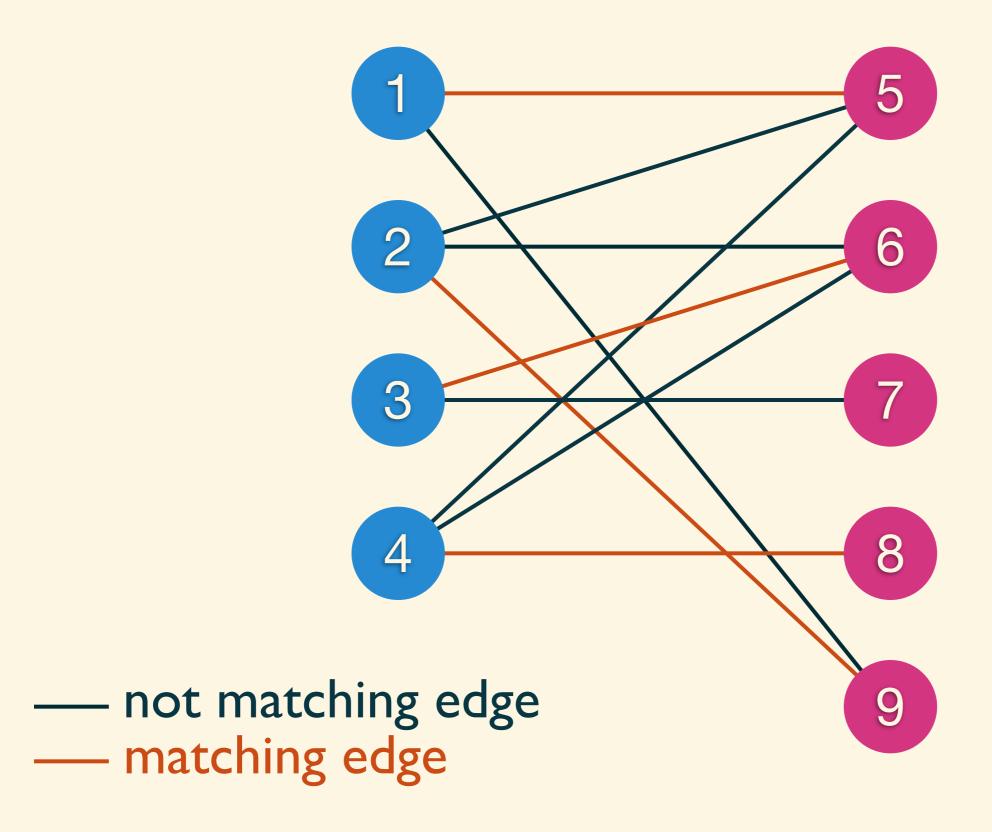


Invalid Matching



Bipartite Matching

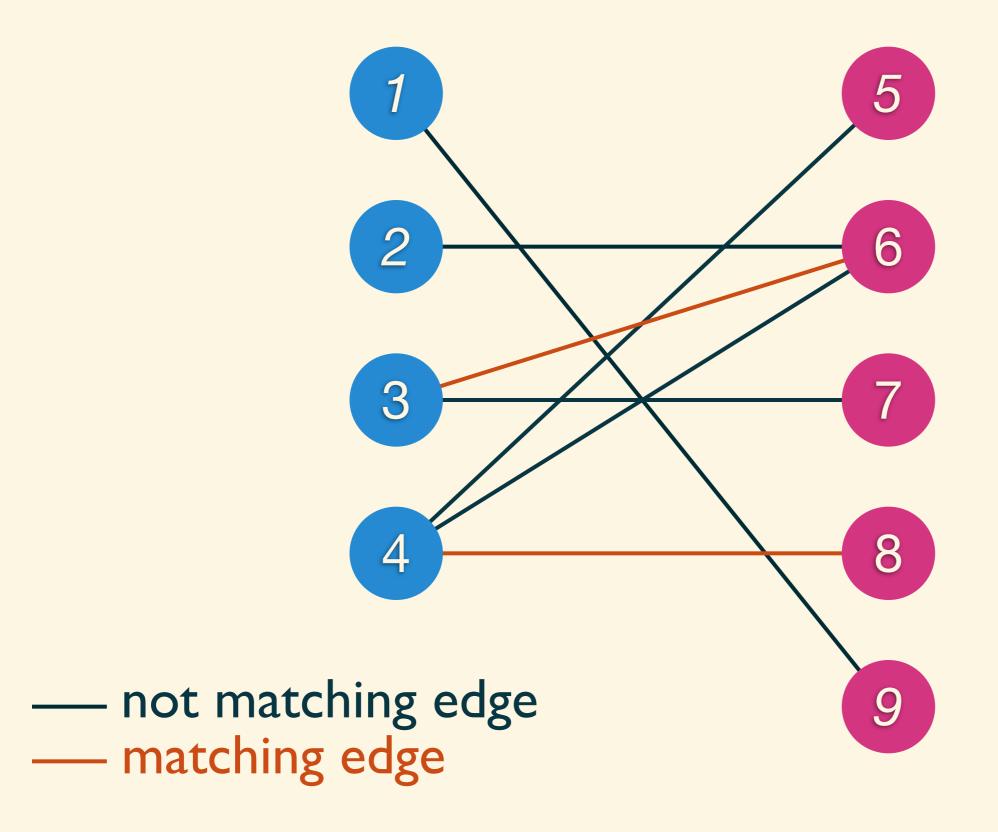




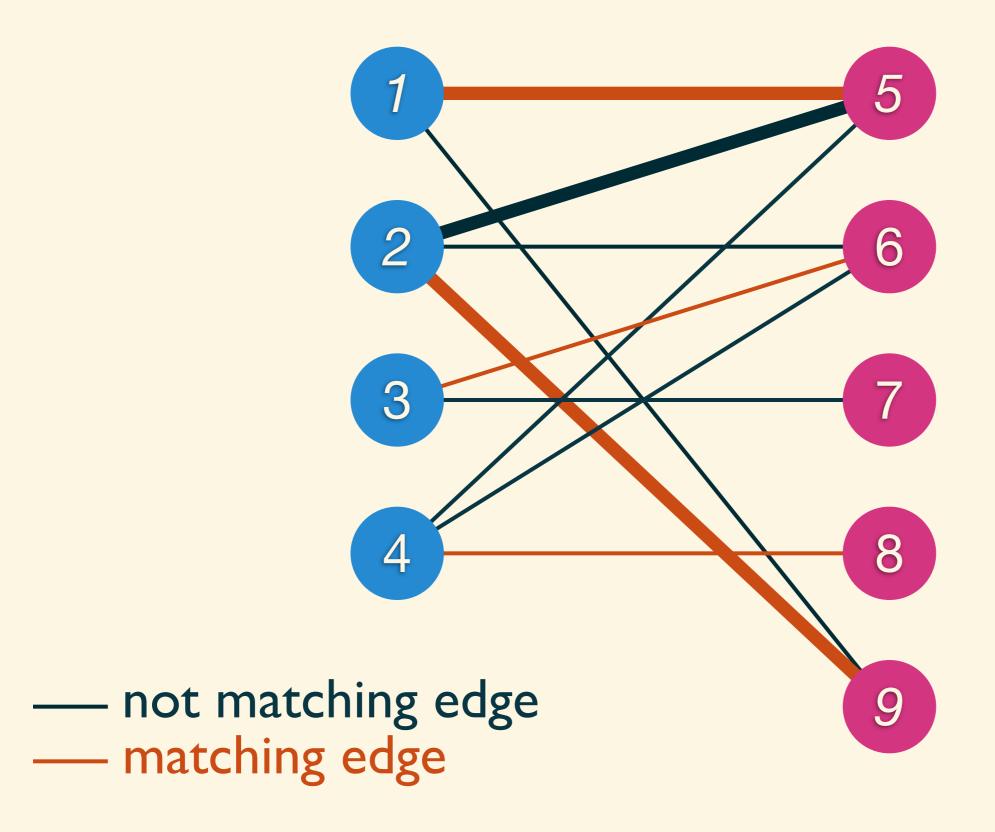
Alternating Path

A path in which the edges belong alternatively to the matching and not to the matching.

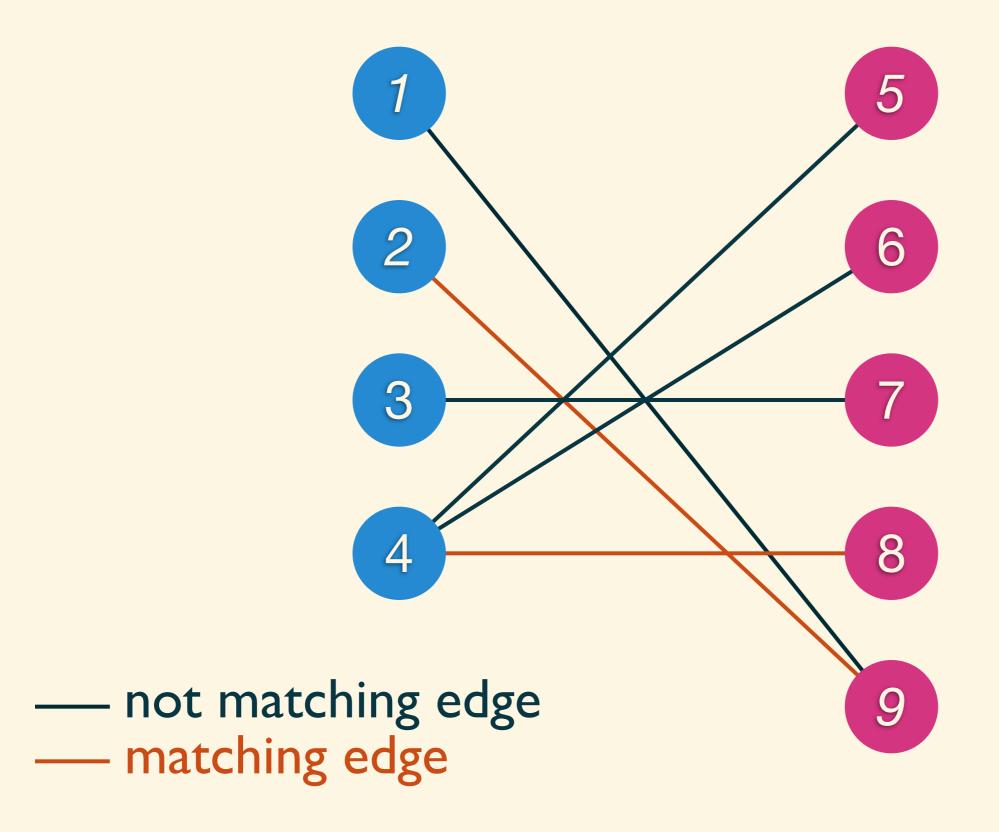
[1 - 5 - 2 - 9] is an alternating path.



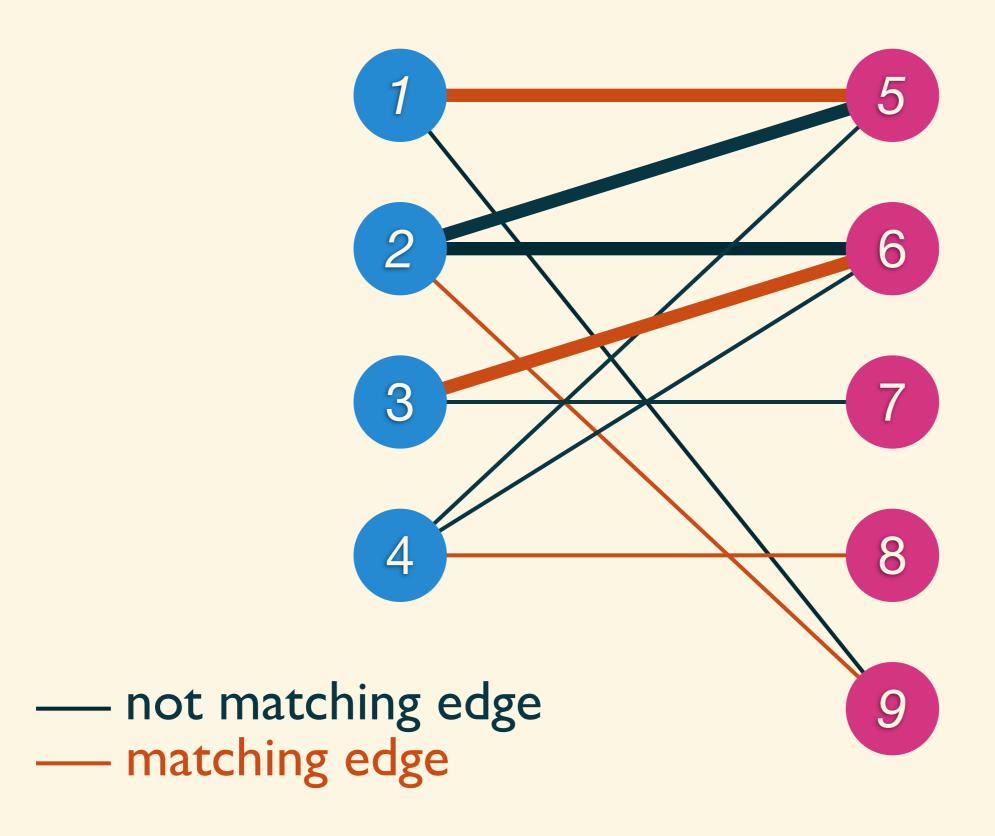
[1 - 5 - 2 - 9] is an alternating path.



[1 - 5 - 2 - 6 - 3] is **not** an alternating path.



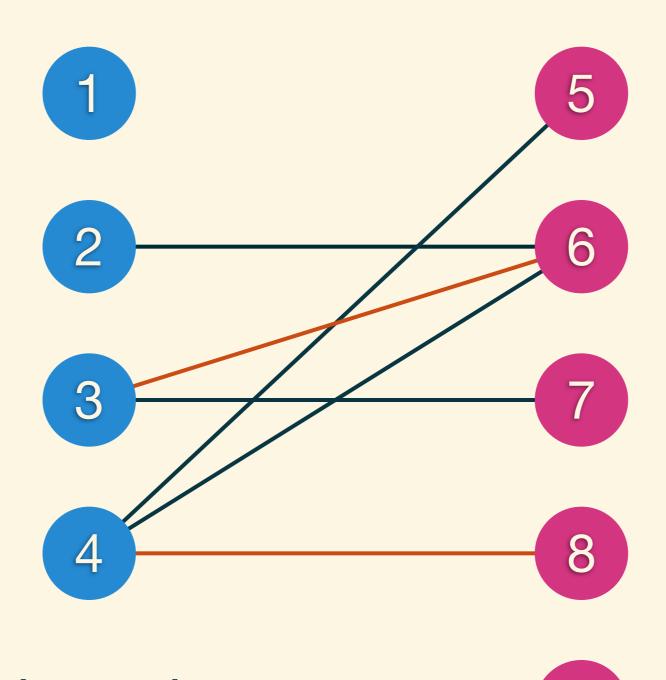
[1 - 5 - 2 - 6 - 3] is **not** an alternating path.



Alternating Cycle

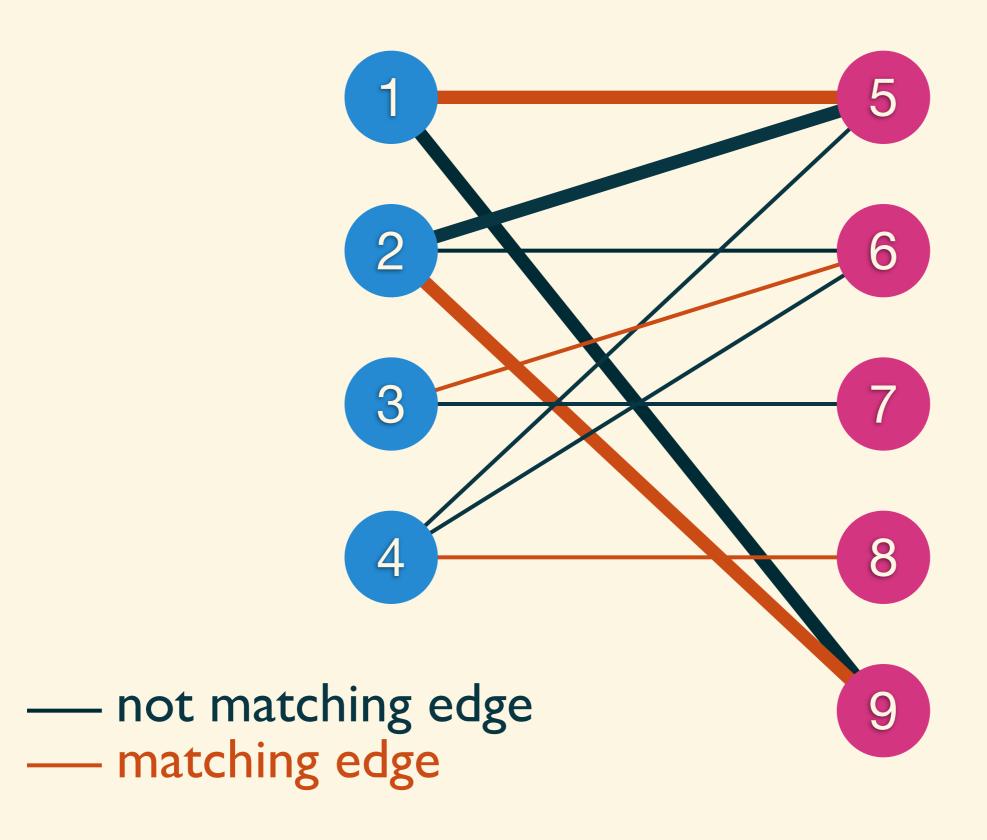
A cycle in which the edges belong alternatively to the matching and not to the matching.

[1 - 5 - 2 - 9 - 1] is an alternating cycle.

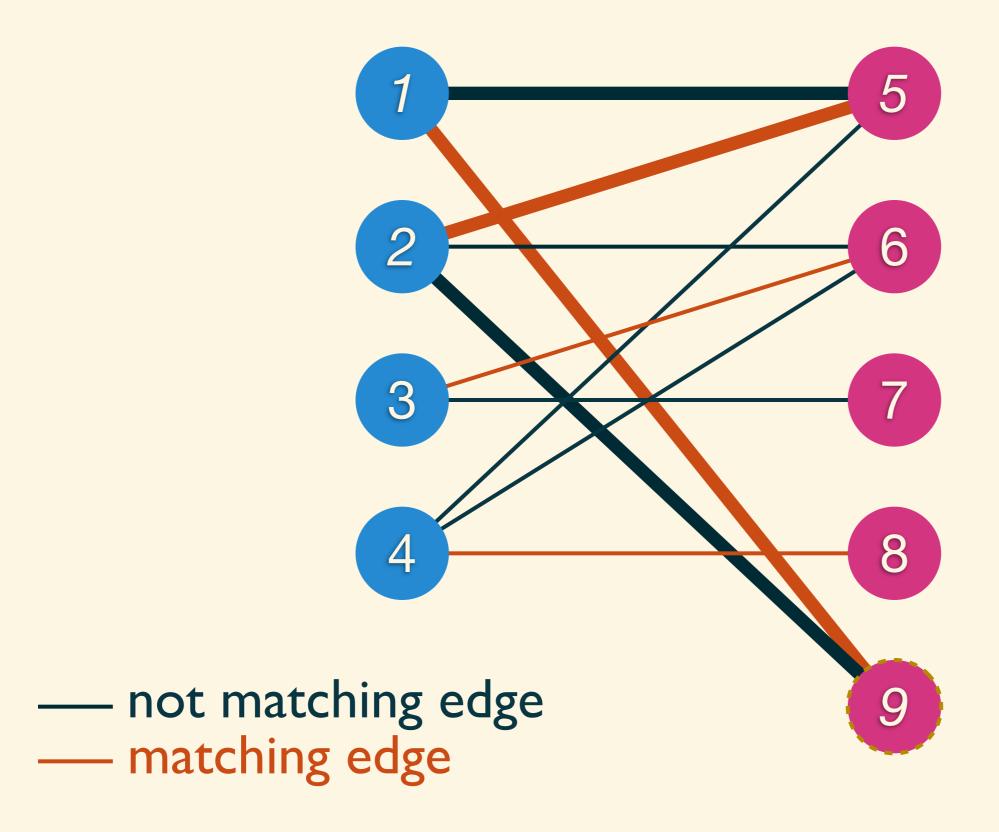


- not matching edge
- matching edge

[1 - 5 - 2 - 9 - 1] is an alternating cycle.



Reverse the alternating cycle, cardinality won't change.

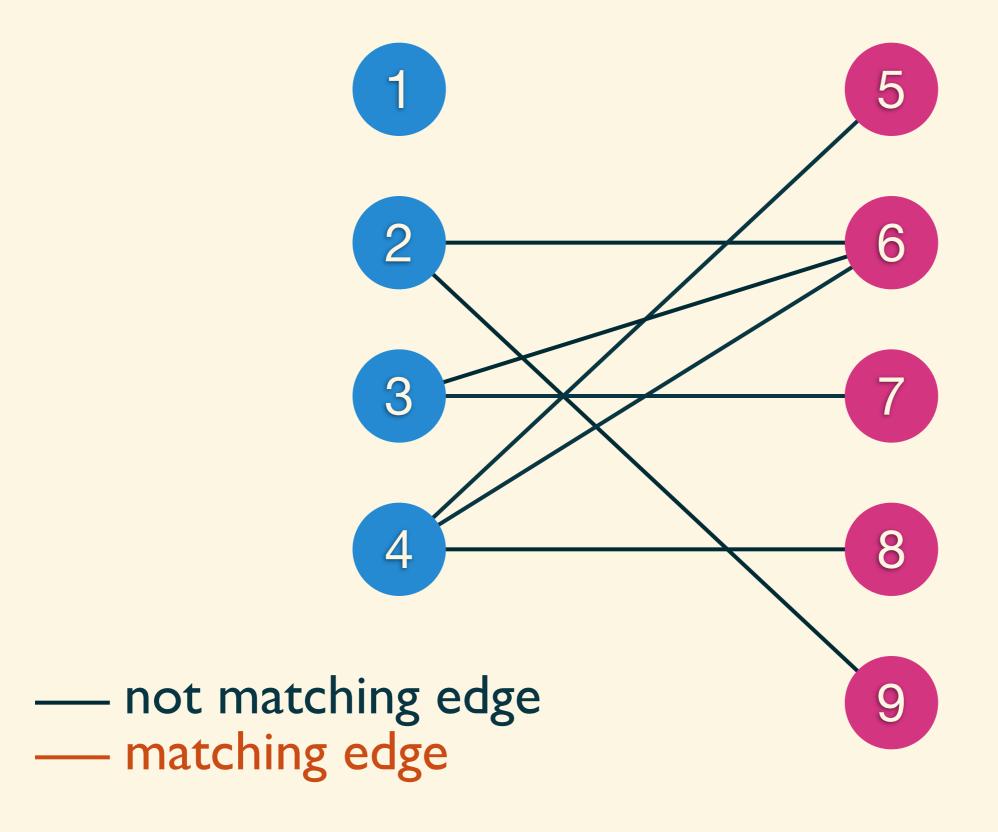


Maximum Matching Augmenting Path Algorithm

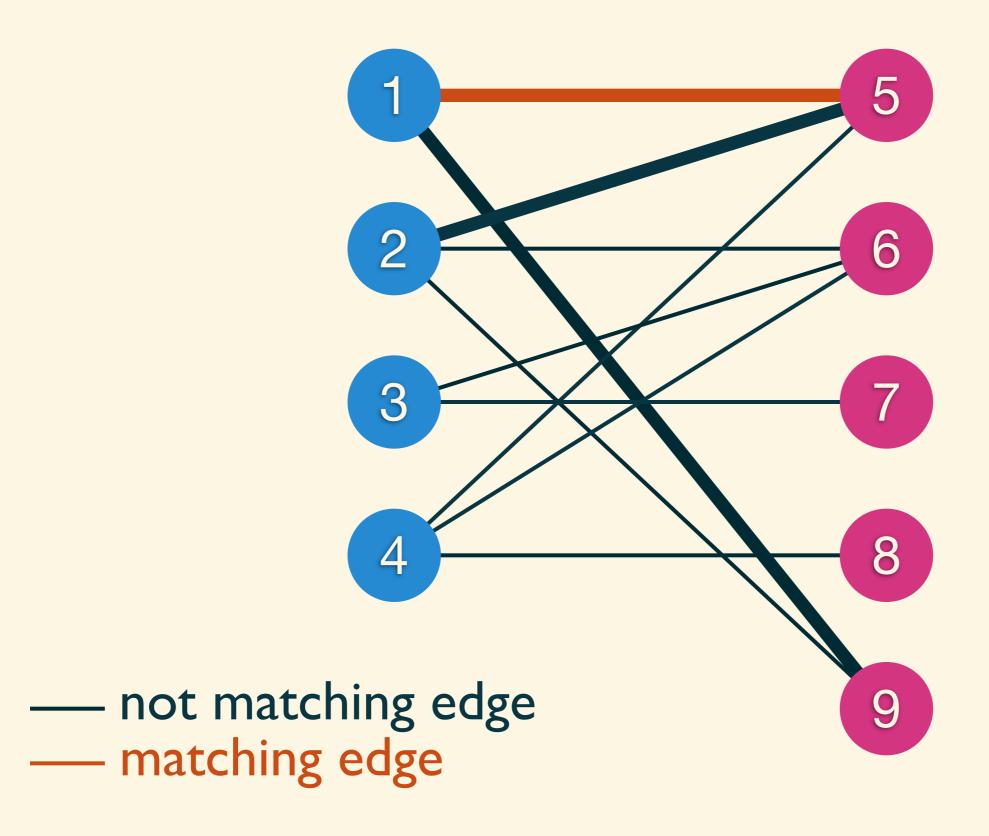
Augmenting Path

An alternating path that starts from and ends on unmatched vertices.

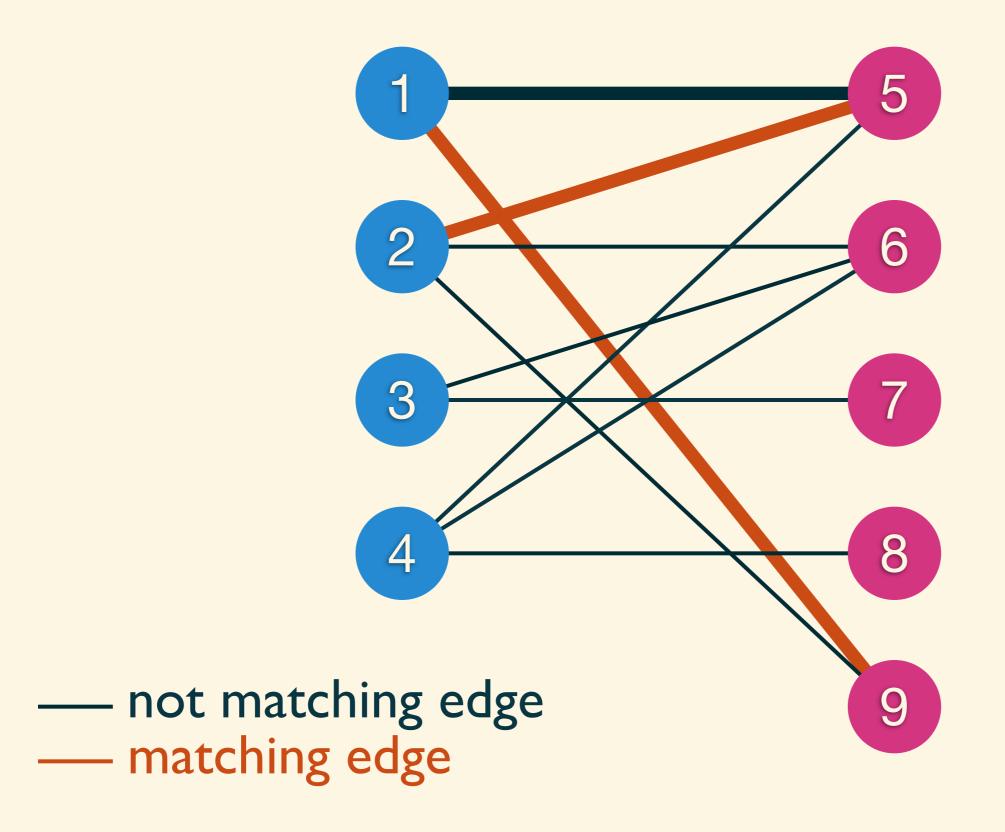
[2 - 5 - 1 - 9] is an augmenting path.



[2 - 5 - 1 - 9] is an augmenting path.



Reverse the augmenting path, cardinality will increase.



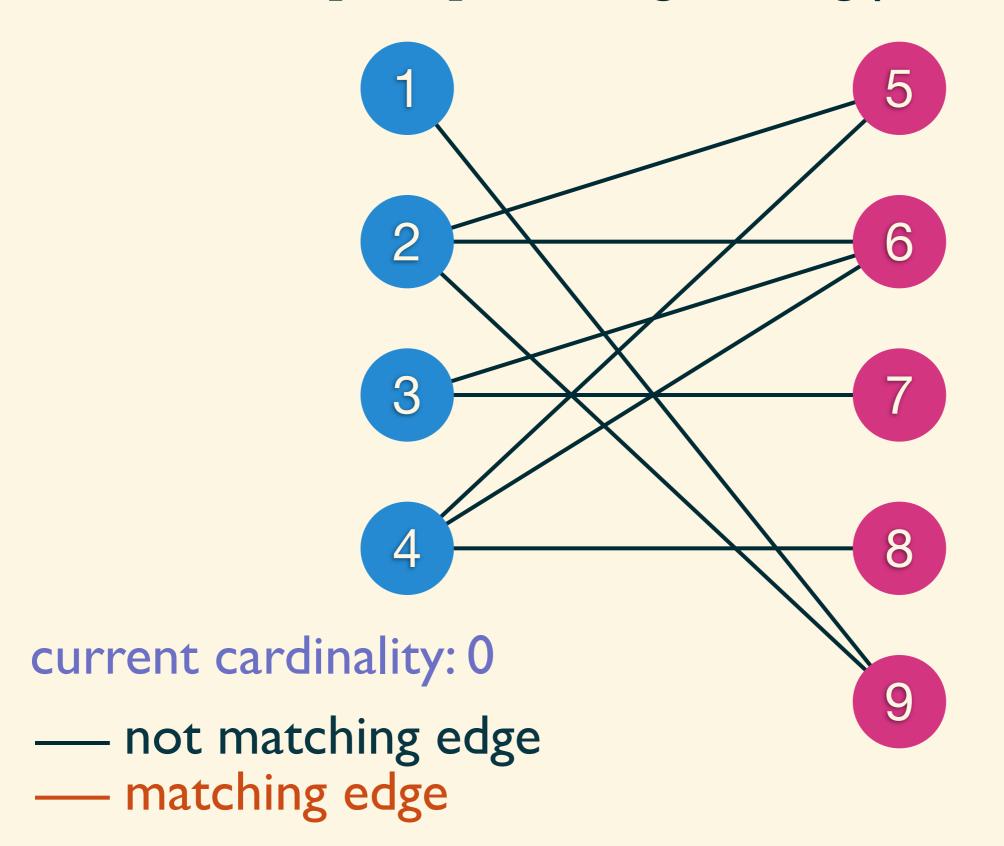
Augmenting Path

- The length of an augmenting path always is odd.
- Reversing an augmenting path will increase the cardinality.
- If there is no augmenting path, the cardinality is maximum.

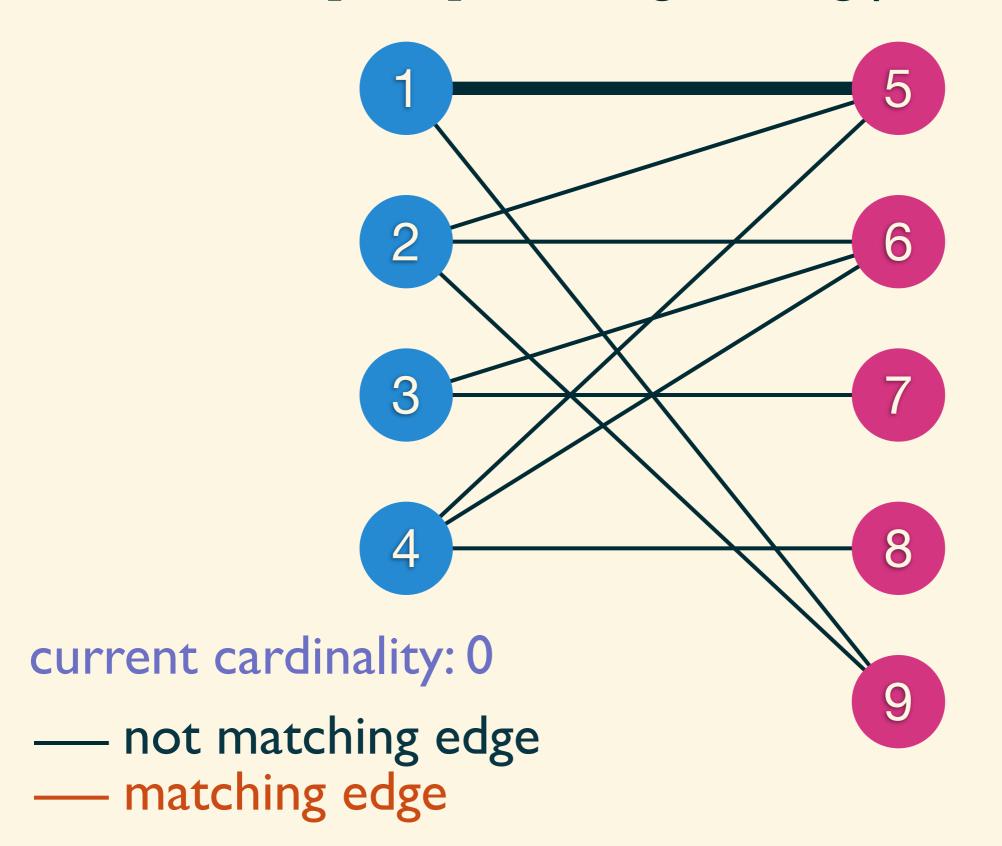
Algorithm

- 1. Try to build augmenting paths from all vertices in one side.
- 2. Travel on graph.
- 3. If an augmenting path exists, reverse the augmenting path to increase cardinality.
- 4. If no augmenting path exists, ignore this vertex.
- 5. Repeat above step until there no augmenting path exists.

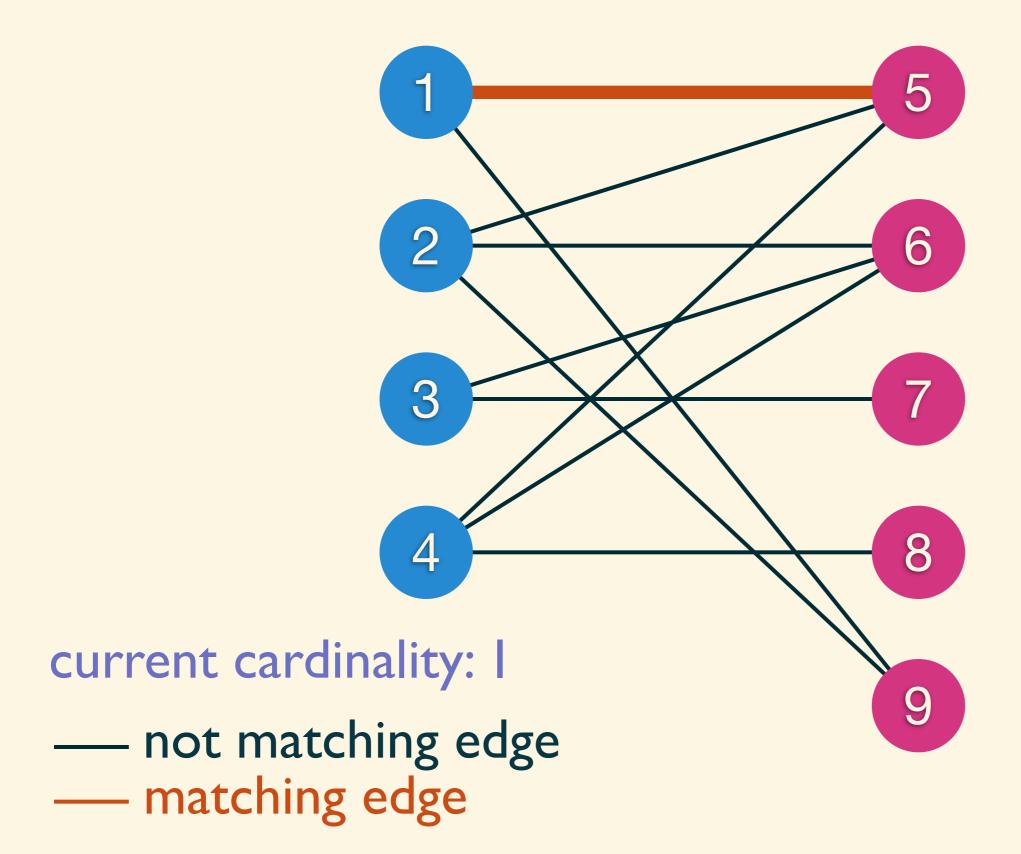
[I - 5] is an augmenting path.



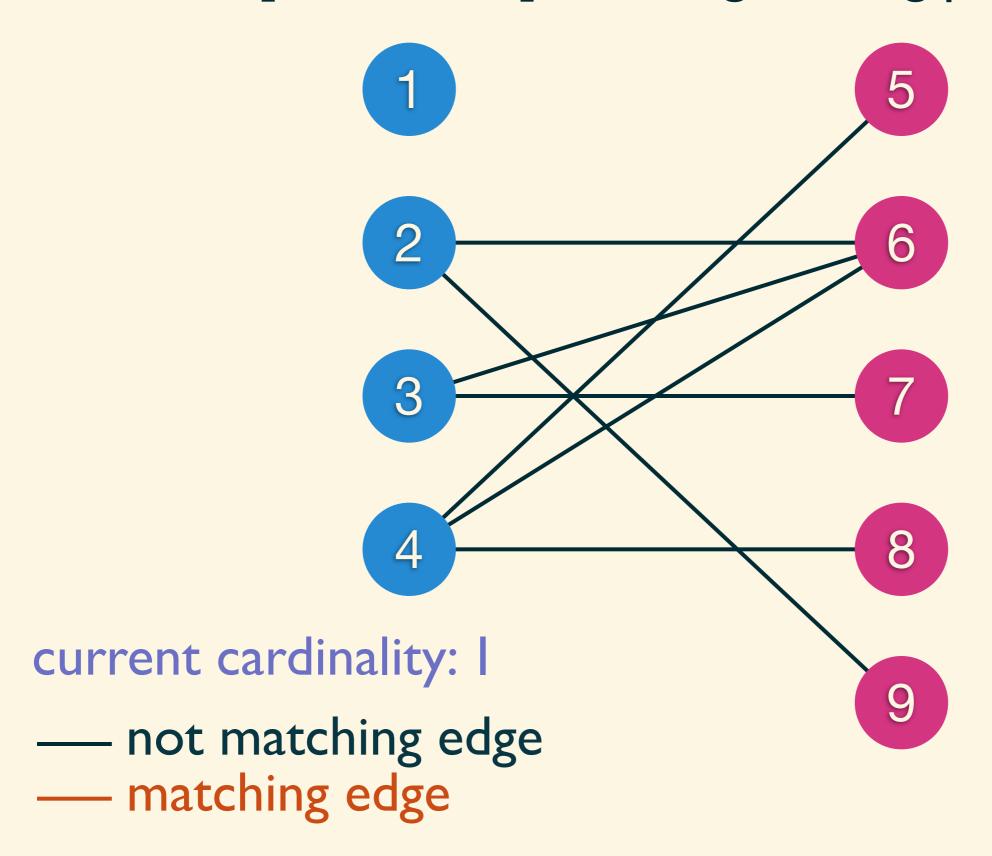
[I - 5] is an augmenting path.



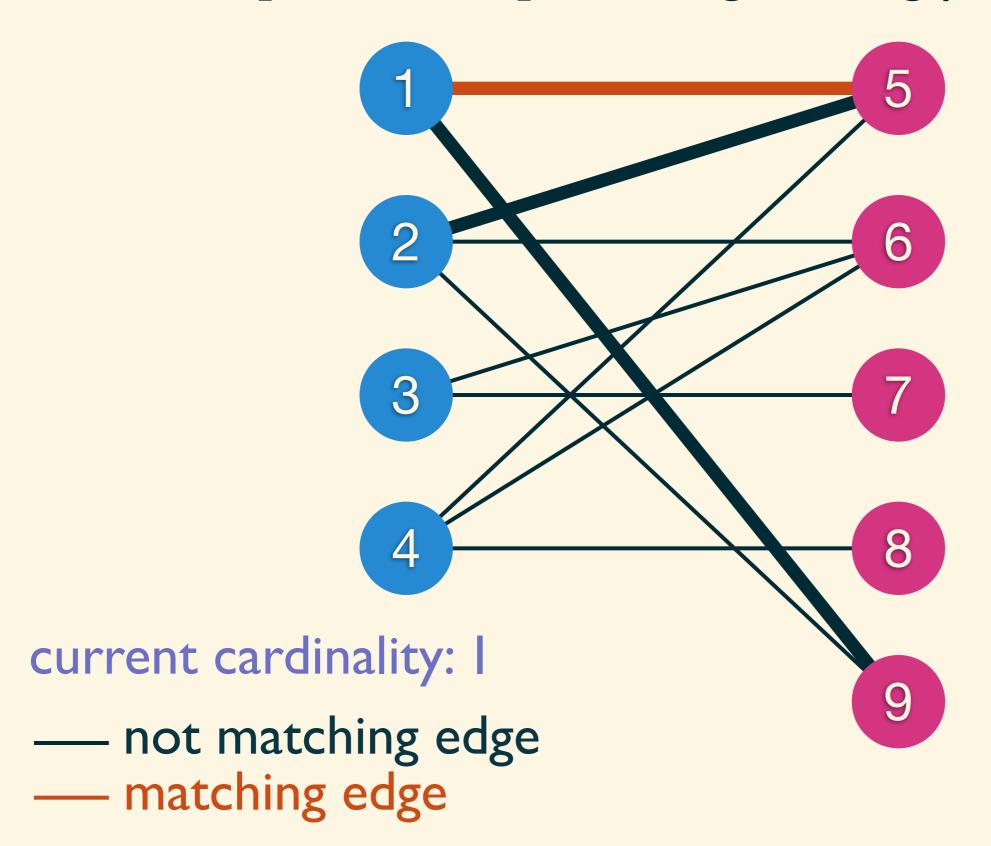
Reverse it!



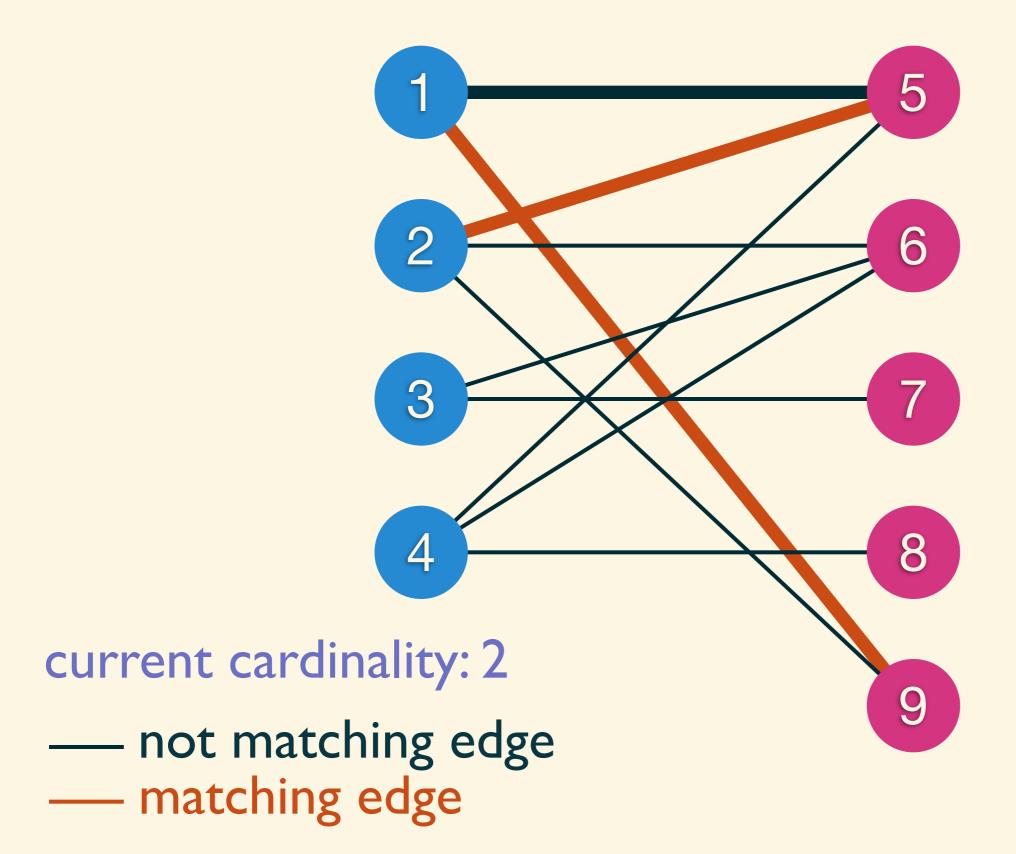
[2 - 5 - I - 9] is an augmenting path.



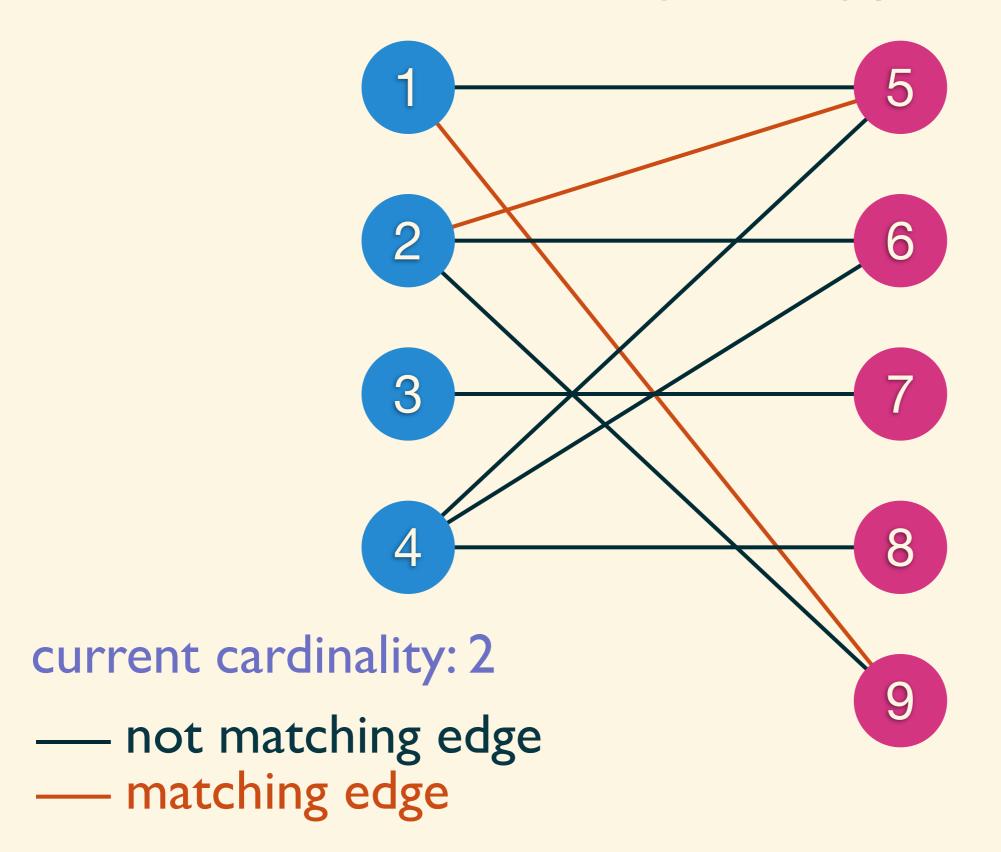
[2 - 5 - I - 9] is an augmenting path.



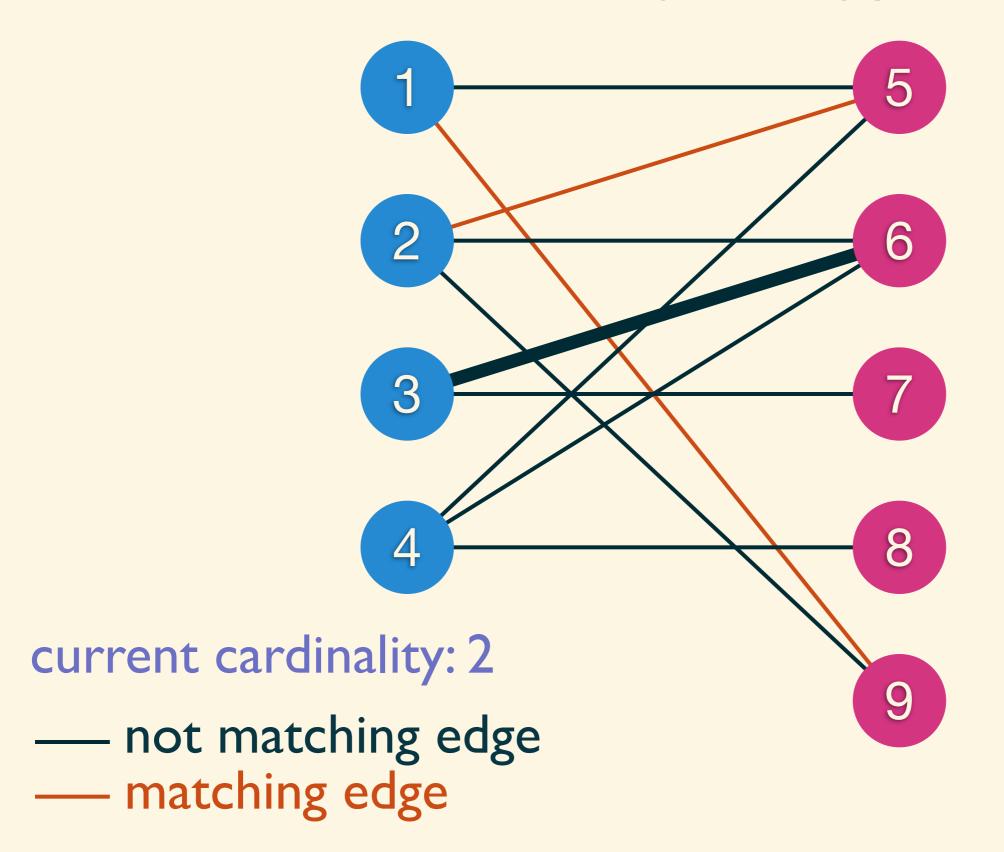
Reverse it!



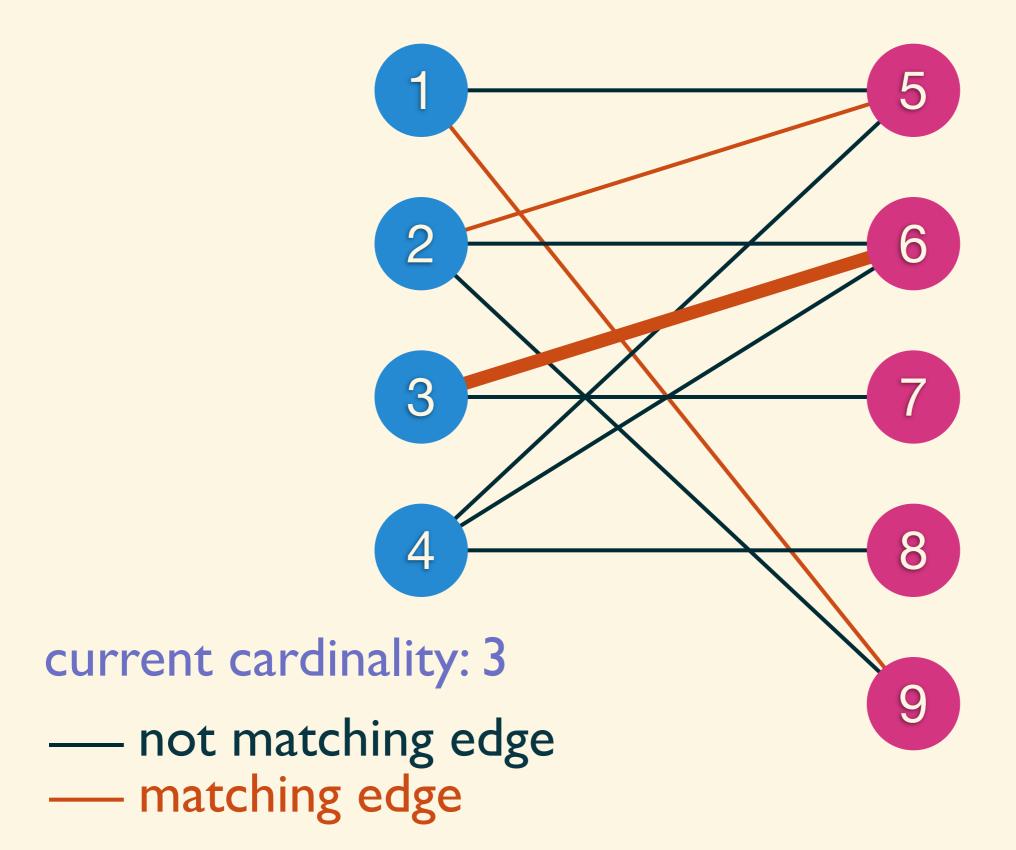
[3 - 6] is an augmenting path.



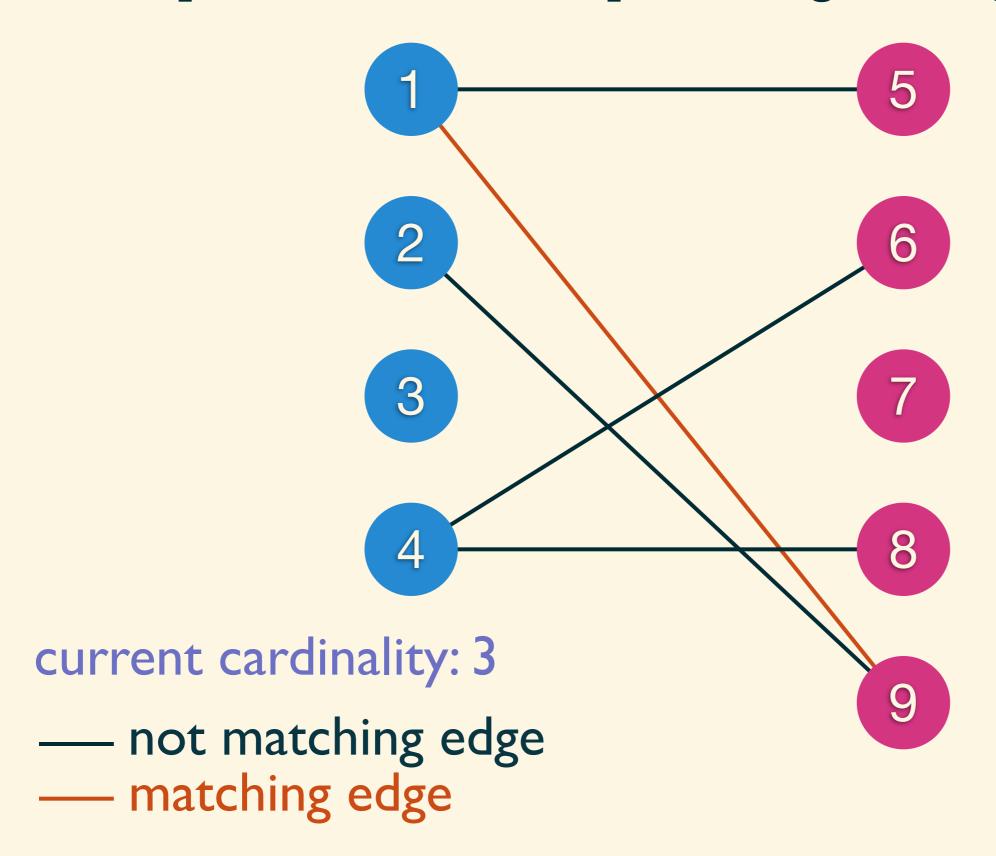
[3 - 6] is an augmenting path.



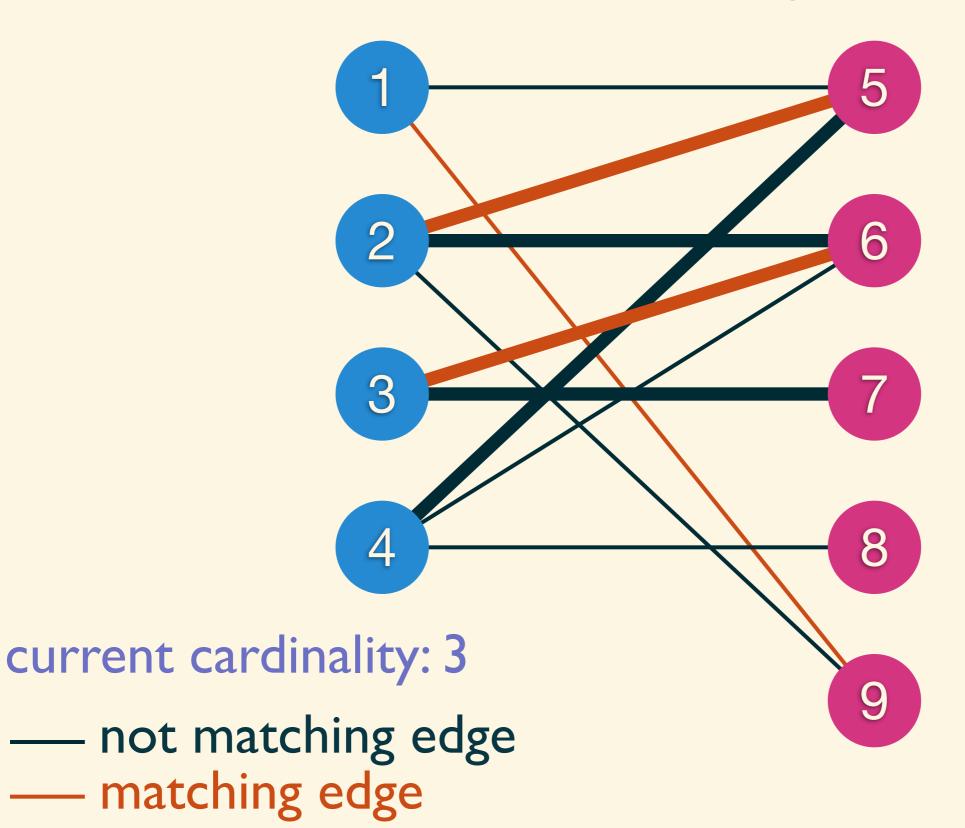
Reverse it!



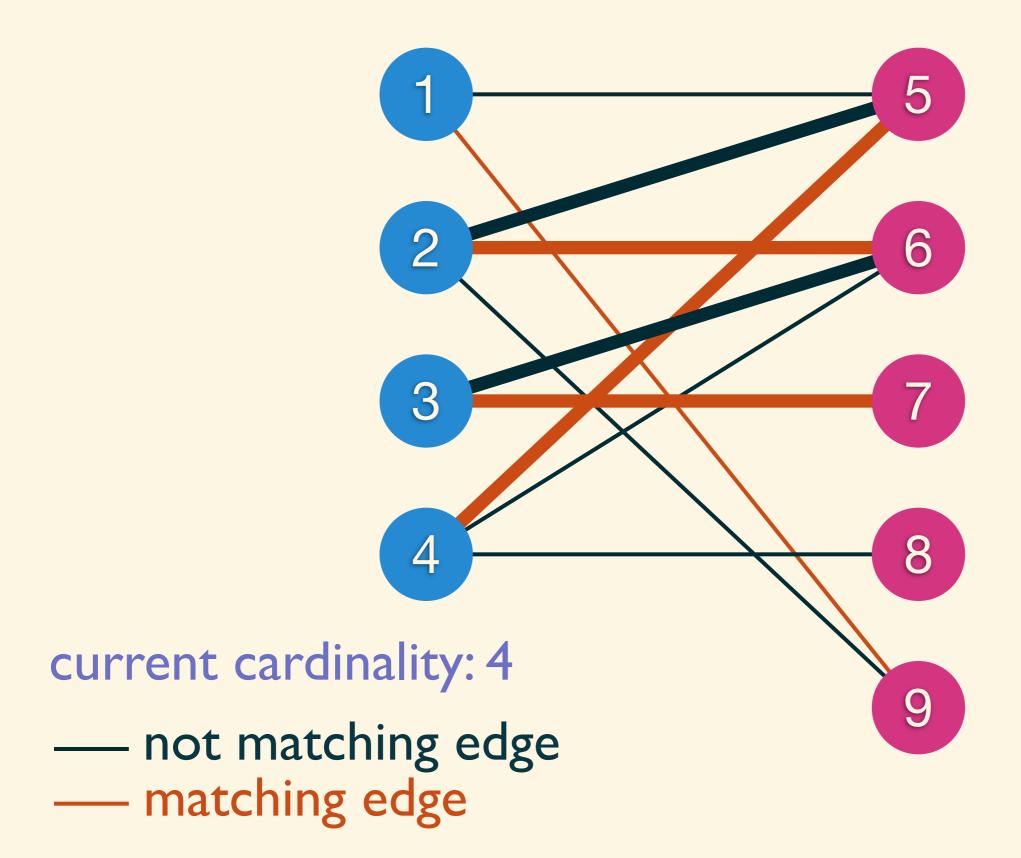
[4 - 5 - 2 - 6 - 3 - 7] is an augmenting path.



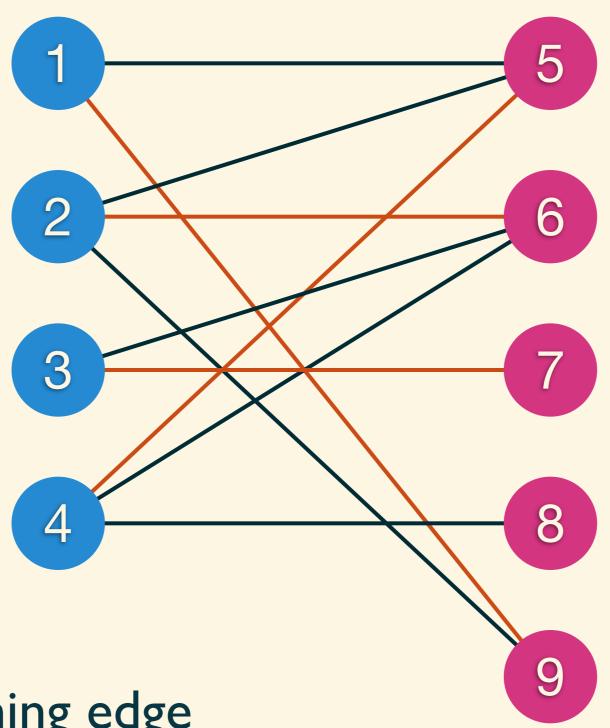
[4 - 5 - 2 - 6 - 3 - 7] is an augmenting path.



Reverse it!



Max cardinality is 4.



not matching edgematching edge

Time Complexity

$$O(V \times E)$$

V: number of vertices

E: number of edges

Source Code

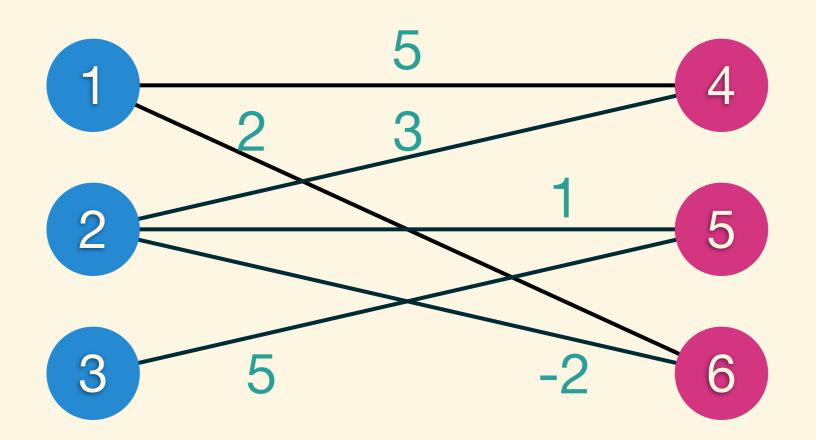
```
// find an augmenting path
bool find_aug_path( int x ) {
    for ( int i = 0; i < ( int )vertex[ x ].size(); ++i ) {</pre>
        int next = vertex[ x ][ i ];
        // not in augmenting path
        if ( !visit[ next ] ) {
            // setup this vertex in augmenting path
            visit \lceil \text{next} \rceil = 1;
             * If this vertex is a unmatched vertex, reverse
augmenting path and return.
             * If this vector is a matched vertex, try to reverse
augmenting path and continue find an unmatched vertex.
            if ( lnk2[ next ] == -1 || find_aug_path( lnk2[ next ] ) )
                lnk1[x] = next, lnk2[next] = x;
                return 1;
    return 0;
```

Practice Now

UVa 670 - The dog task

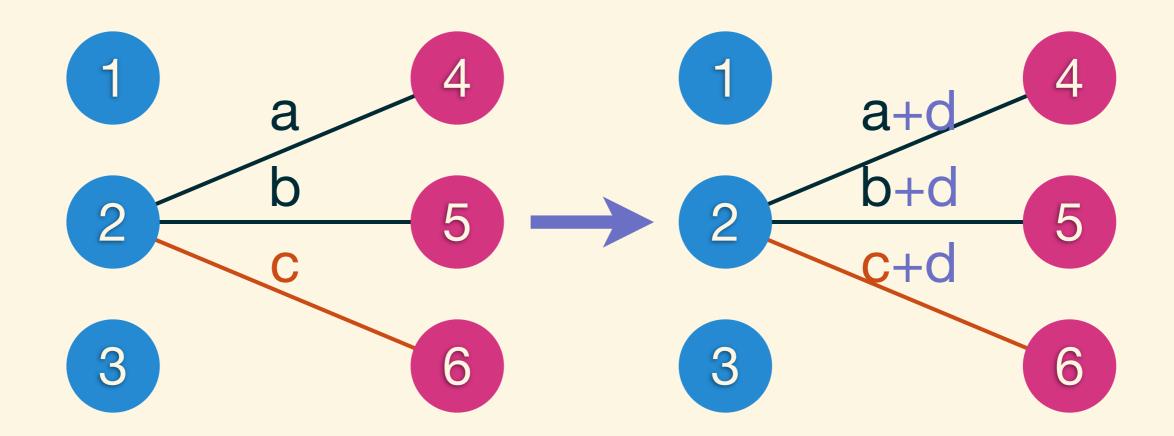
Maximum Weight Matching Hungarian Algorithm

(Kuhn-Munkres Algorithm)



weight adjustment

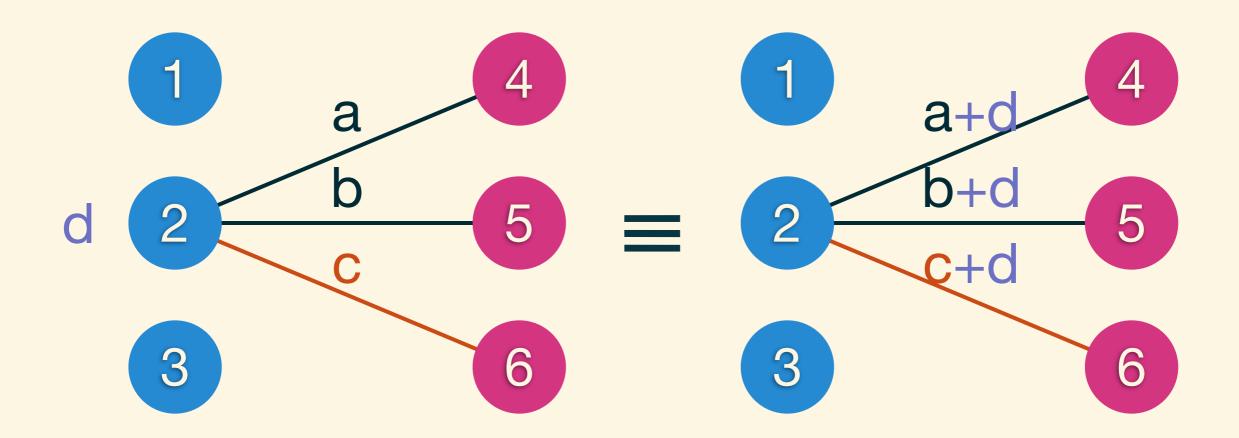
If add (subtract) some value on all edges connected with vertex X, the maximum matching won't be effected.



— matching edge — not matching edge

vertex labeling

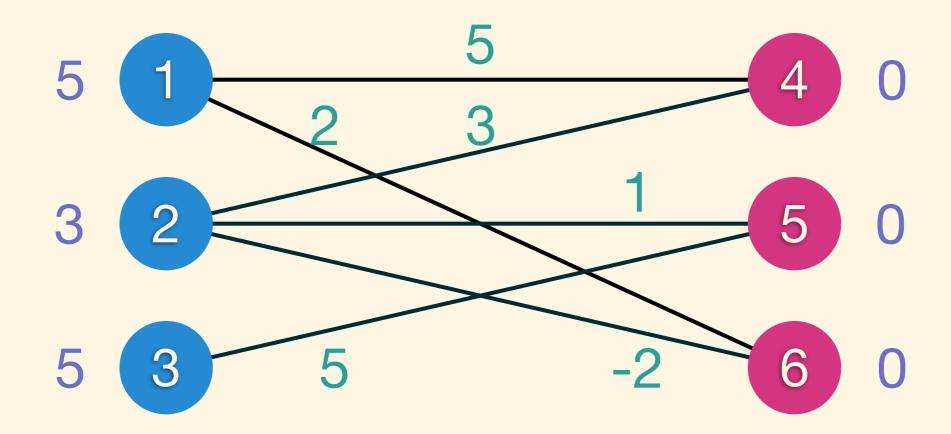
For convenient, add a variable on vertices to denote the value add (subtract) on edges connected with them.



— matching edge — not matching edge

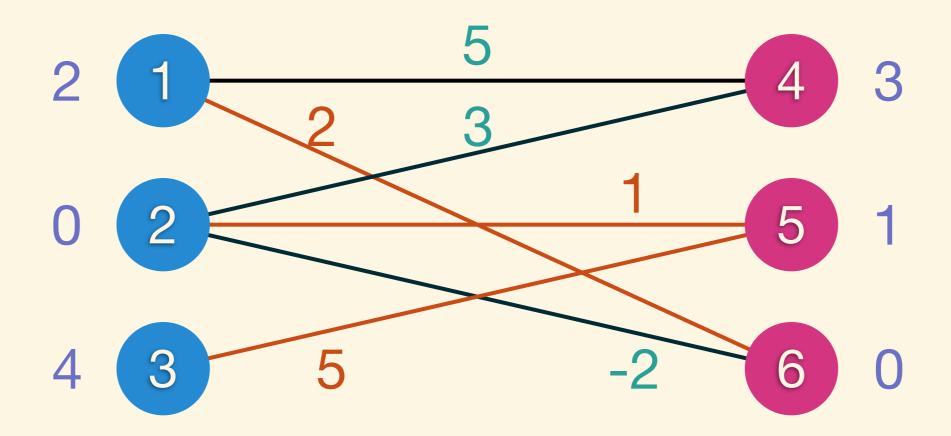
vertex labeling

 $l(x)+l(y)\geq w(x,y)$, for each edge(x,y)

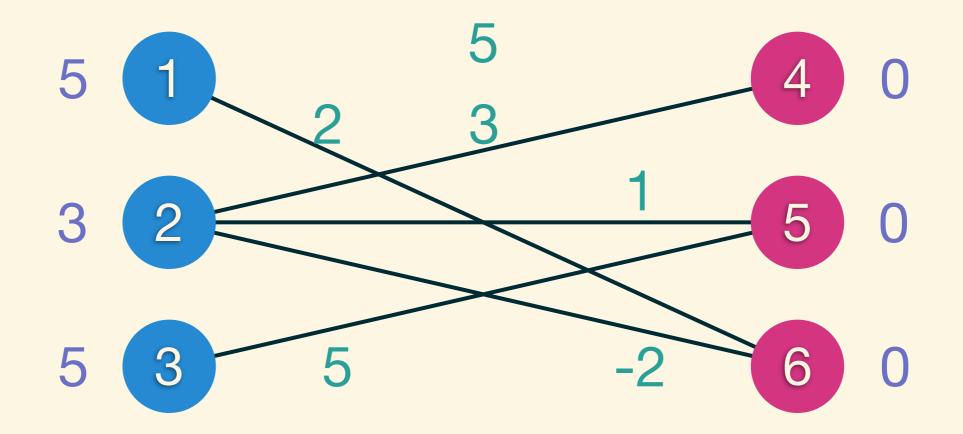


minimize vertex labeling

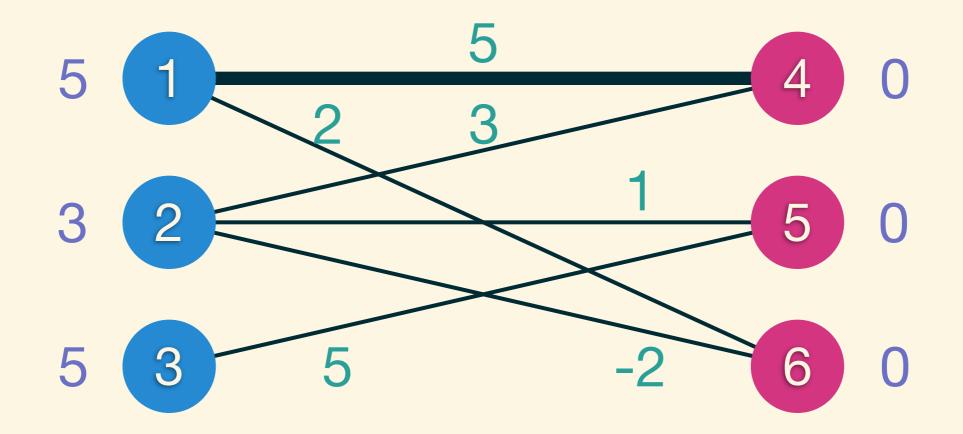
Admissible Edge: l(x)+l(y)=w(x,y)



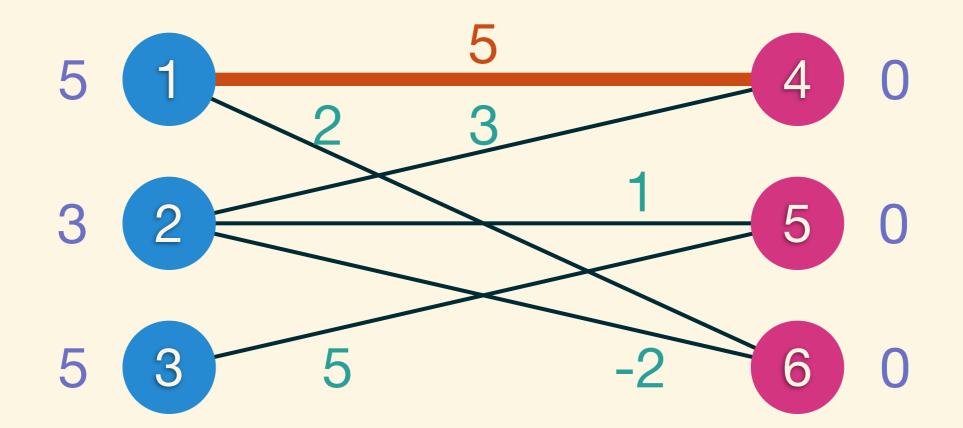
convert maximum weight problem to minimum vertex labeling problem



[1 - 4] is an augment path with admissible edges.

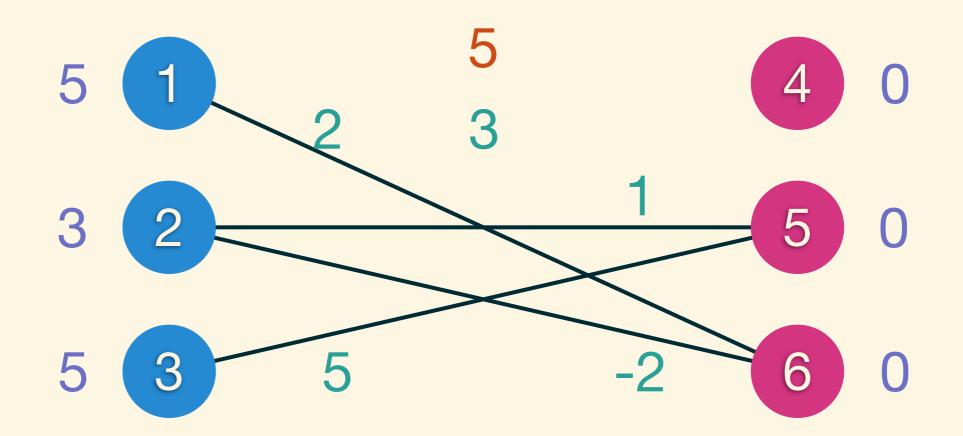


[1 - 4] is an augment path with admissible edges.

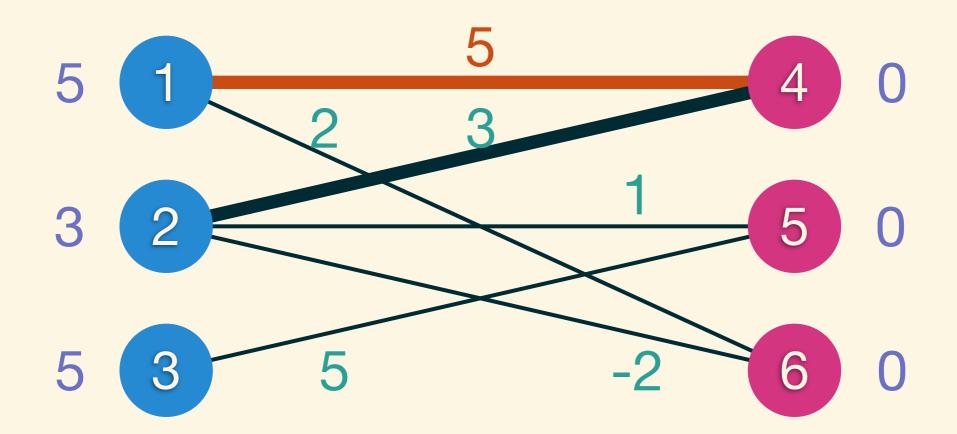


Reverse it!

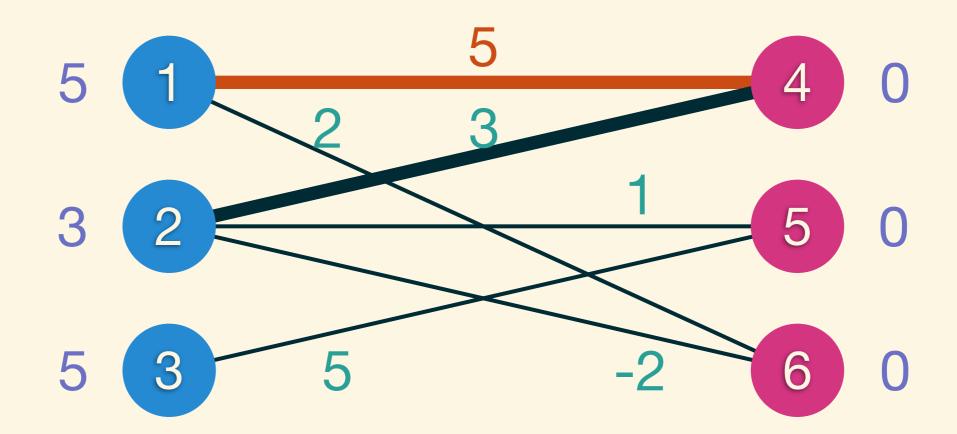
Current Weight = 5



No augmenting path found start from vertex 2.



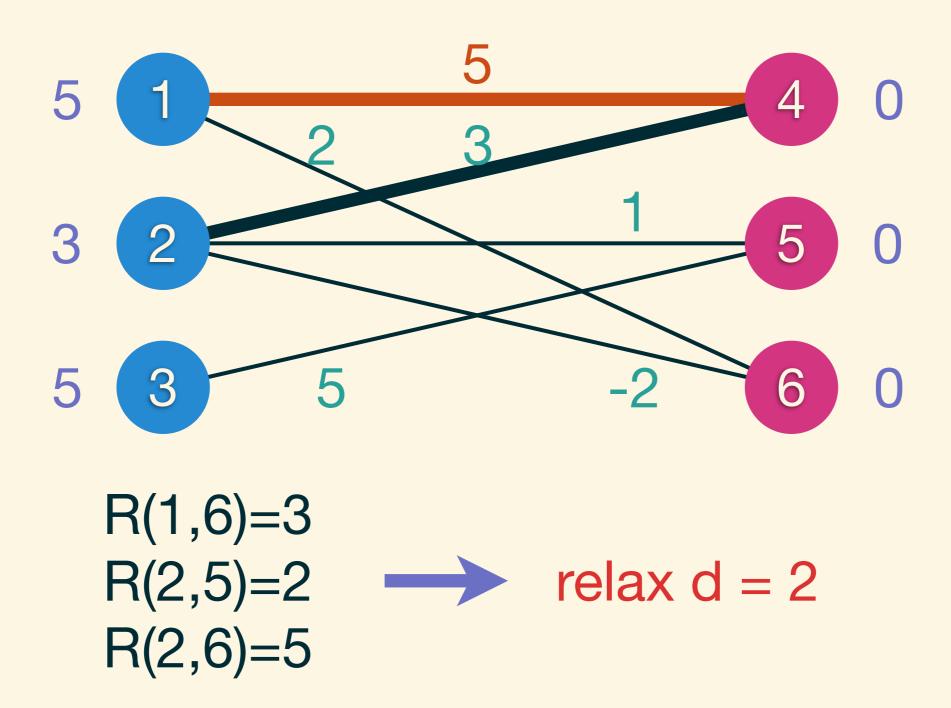
No augmenting path found start from vertex 2.

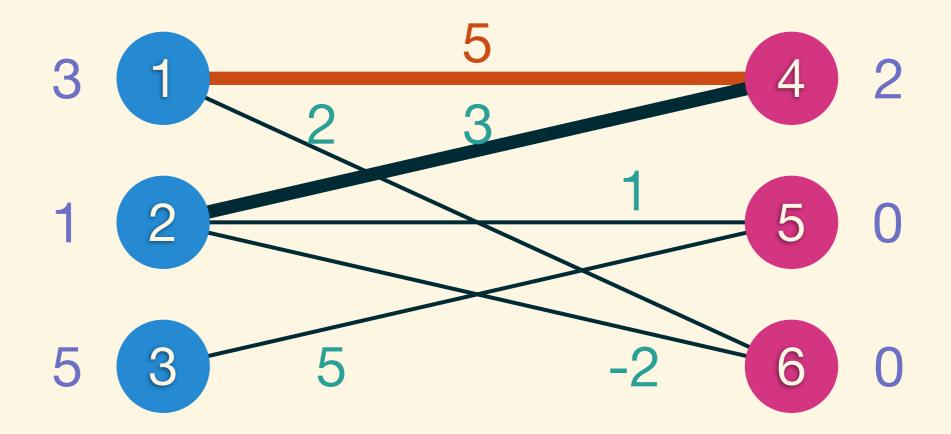


relax value d = min(l(x)+l(y)-w(x,y))

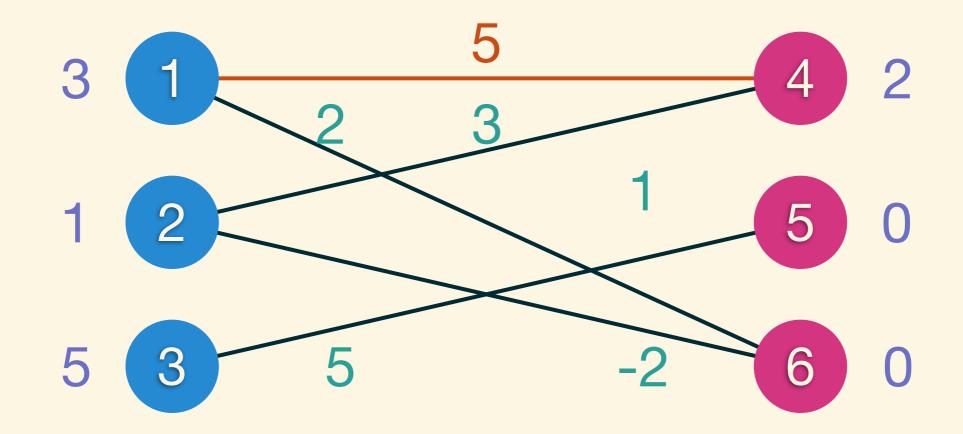
x: vertex in alternating path

y: vertex y not in alternating path

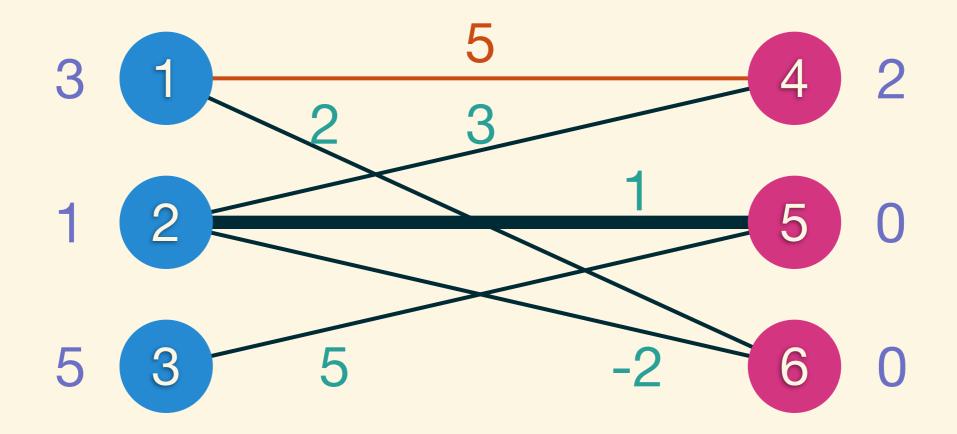




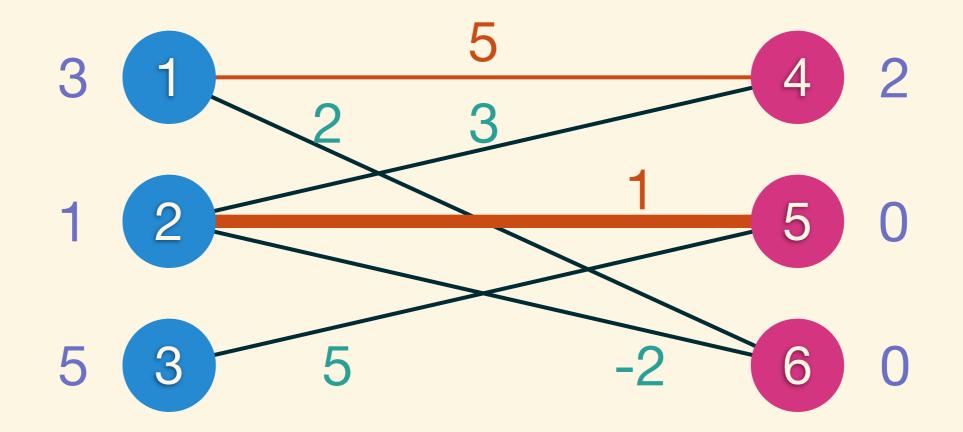
l(x) subtract value d.l(y) add value d.



[2 - 5] is an augment path with admissible edges.

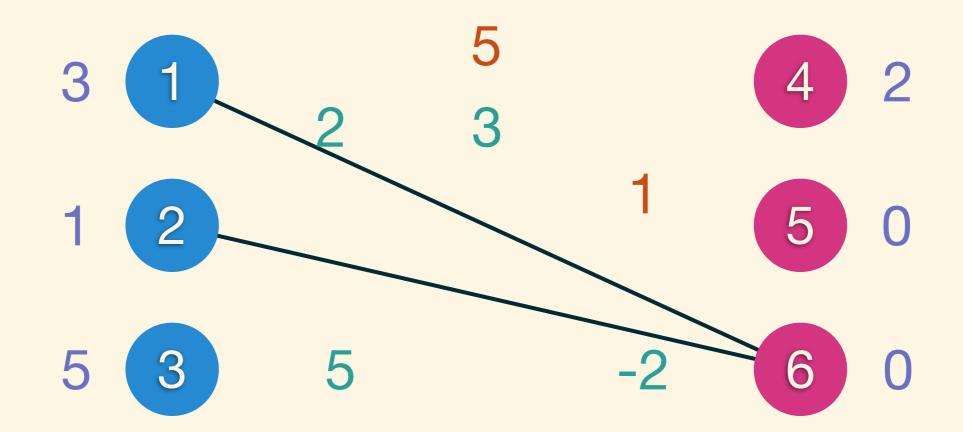


[2 - 5] is an augment path with admissible edges.

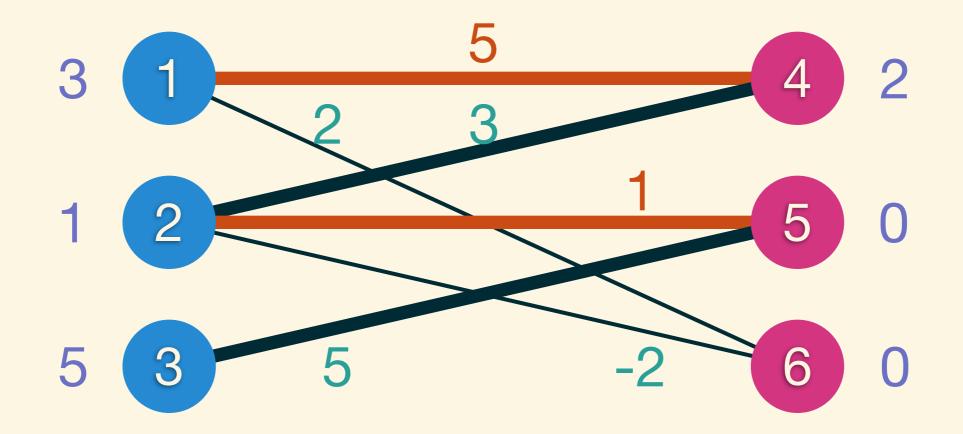


Reverse it!

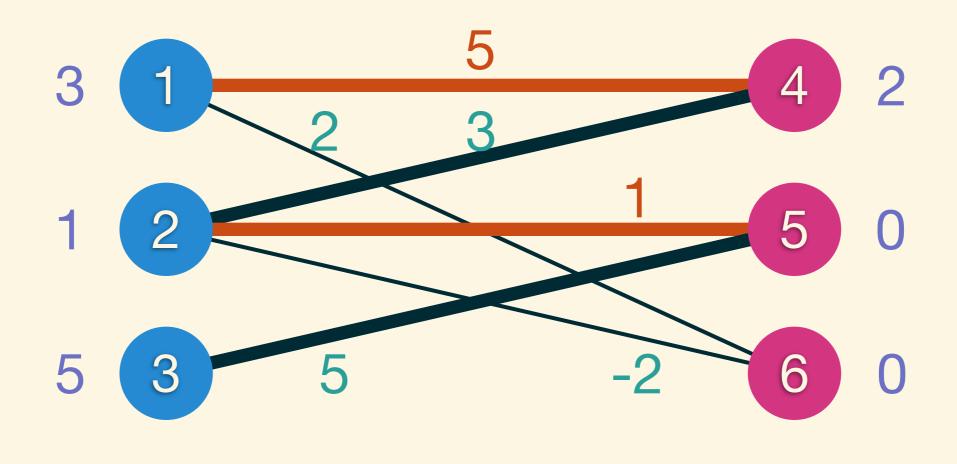
Current Weight = 6



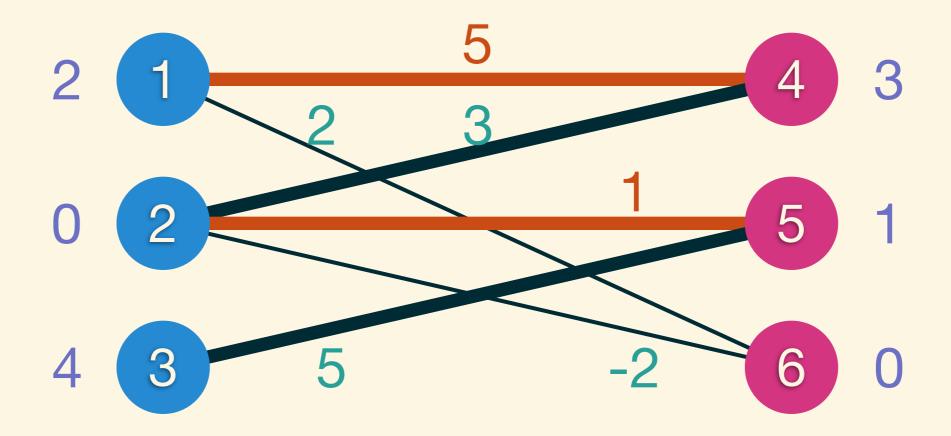
No augmenting path found start from vertex 3.



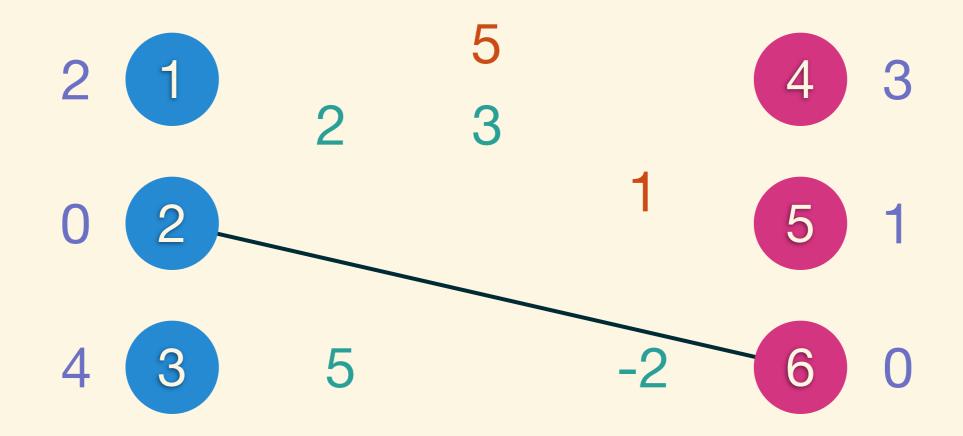
No augmenting path found start from vertex 3.



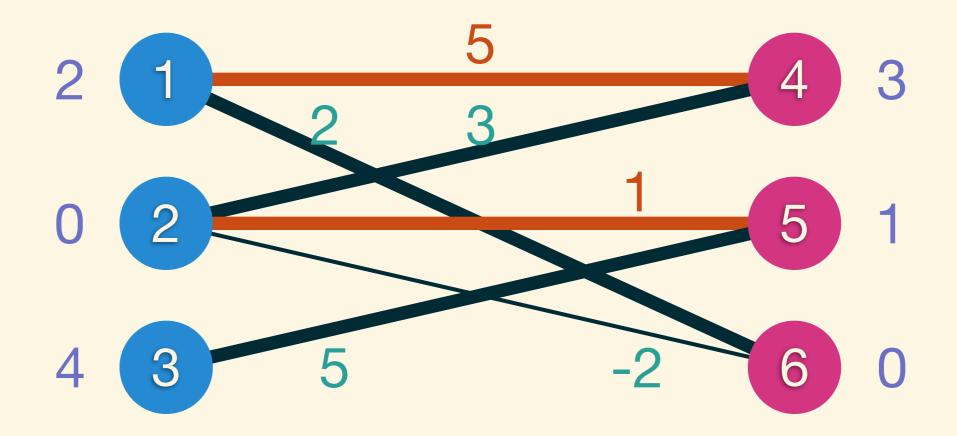
$$R(1,6)=1$$
 $R(2,6)=3$
relax d = 1



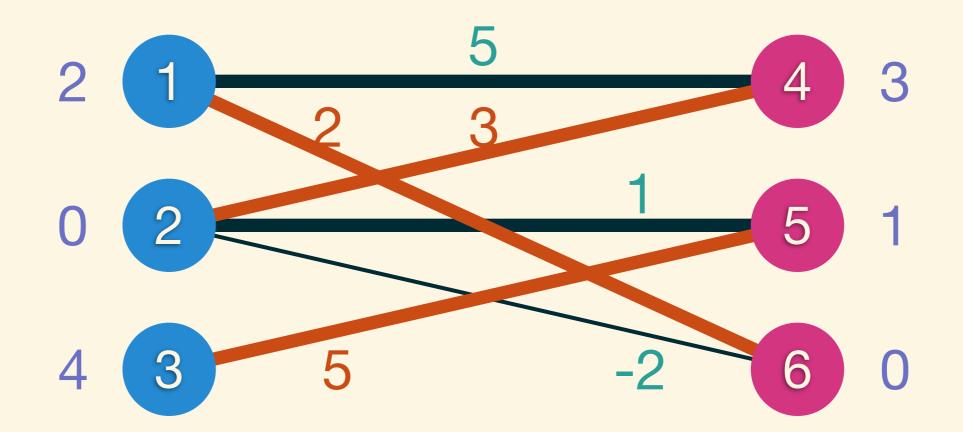
- l(x) subtract value d.
- l(y) add value d.



[3 - 5 - 2 - 4 - 1 - 6] is an augment path with admissible edges.



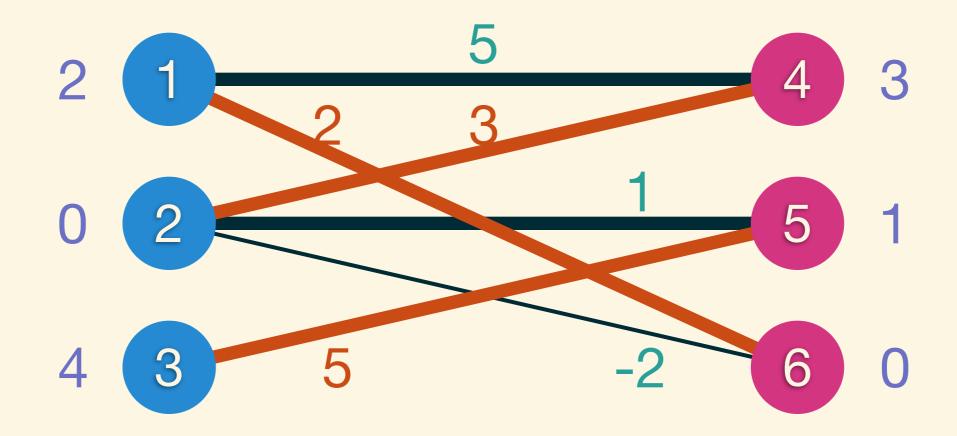
[3 - 5 - 2 - 4 - 1 - 6] is an augment path with admissible edges.



Reverse it!

Current Weight = 10

Maximum Weight Matching = 10



Algorithm

- 1. Initial vertex labeling to fit $l(x)+l(y)\geq w(x,y)$
- 2. Find augmenting paths composed with admissible edges from all vertices in one side.
- 3. If no augmenting path exists, adjust vertex labeling until the augmenting path found.
- 4. If augmenting path exists, continue to find next augmenting path.
- 5. Repeat above step until there no augmenting path exists.
- 6. Calculate the sum of weight on matching edges.

Practice Now

POJ 2195 - Going Home

Problem List

UVa 670
UVa 753
UVa 10080
UVa 10092
UVa 10243
UVa 10418
UVa 10984
POJ 3565

Reference

- http://en.wikipedia.org/wiki/Bipartite_graph
- http://en.wikipedia.org/wiki/Matching (graph_theory)
- http://www.flickr.com/photos/marceau_r/5244129689
- http://www.sanfilippo-chianti.it/offerta-svalentino.html
- http://www.csie.ntnu.edu.tw/~u91029/Matching.html

Thank You for Your Listening.

