

# Elementary Cellular Automata as Multiplicative Automata

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## Software

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## Summary

Elementary cellular automata (ECA) are a set of simple binary programs in the form of truth tables called Wolfram codes that produce complex output when done repeatedly in parallel, and quaternions are frequently used to represent 3D space and its rotations in computer graphics. Both are well-studied subjects, this Java library puts them together in a new way. This project changes classical additive cellular automata into multiplicative automata ([Wolfram, 2002, p. 861](#)) via permutations, hypercomplex numbers, and pointer arrays. Valid solutions extend the binary ECA to complex numbers, produce a vector field, make an algebraic polynomial, and generate some very interesting fractals.

The code repository is at <https://github.com/dmcki23/MultiplicativeECA>.

## Statement of Need

Very loosely analogous to DeMorgan's law in Boolean algebra, the main algorithm produces several multiplicative versions of any given standard additive binary Wolfram code up to 32 bits and is written to support user supplied complex input at row 0 with choice of type of multiplication tables and partial product tables among other parameters. An algebraic polynomial of the automata that works with real and complex numbers is produced, and the hypercomplex 5-factor identity solution allows for the complex extension of any binary cellular automata. The GUI, though not required, allows for visual exploration of solutions with easy access to various parameters. The Java this is written in is designed to integrate well in other programs, such as Mathematica's JLink or Matlab, and is documented with Javadoc. The Cayley-Dickson and Fano construction libraries may be of value to the open source community as well.

There are other cellular automata implementations, Mathematica ([Inc., n.d.](#)), CellPyLib ([Antunes, 2021](#)), a JOSS Python project from three years ago, and books that cover related territory ([Ceccherini-Silberstein & Coornaert, 2023](#)). What sets this library apart is the conversion of any existing binary automata rule from additive to multiplicative and its extension to complex numbers without restrictions such as linearity of the rule.

## Functions

### Hypercomplex unit vector implementation

	Negative sign bit	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
Complex	2's place	1	i	-1	-i												
Quaternions	4's place	1	i	j	k	-1	-i	-j	-k								
Octonions	8's place	1	e1	e2	e3	e4	e5	e6	e7	-1	-e1	-e2	-e3	-e4	-e5	-e6	-e7

The Cayley-Dickson (CD) and Fano support classes are discussed in greater detail in the readme and the documentation, they along with the Galois class provide sets of multiplication tables to be compared with cellular automata. The CD multiplication implementation permutes the steps of splitting and recombining hypercomplex numbers to increase the scope of the CD equation,  $(a, b)x(c, d) = (ac - d * b, da + bc*)$ , where  $*$  is the conjugate. It verifies itself by producing the symmetric group of its degree when interacting with other CD multiplications. The Fano library octonions produce a triplet that is a linear match to the CD octonions as triplets{0} when the up and down recursion factoradics are equal, and produce the triplet set of John Baez's Fano plane as triplets{10}. (Baez, 2001).

The main algorithm uses a set of permutations operating on cellular automata input, each permutation permuting the neighborhood, becoming a factor, with four kinds of multiplications. The multiplication tables are input as 2D but used as N-D, where  $N = \text{numFactors}$ .

	<b>Multiplications A</b>	<b>Multiplications B</b>	<b>Multiplications C</b>	<b>Multiplications D</b>
<b>Type</b>	Hypercomplex or finite, brute-force of all the permutations of that number of factors	Cartesian product summed by a hypercomplex or finite partial product table	Complex product	Permutation composition
<b>Size</b>	Wolfram code length = L	Size of neighborhood, $\log_2(L)$	Size of neighborhood, $\log_2(L)$	Size of neighborhood, $\log_2(L)$
<b>Function</b>	Validates permutation group, reproducing the Wolfram code as a pointer array	Applies a valid solution to a user given complex neighborhood	Like B, but does the normalization before the multiplication	Orders the cell's neighborhood vector from (B), post multiplication, pre normalization
<b>Scope</b>	Entire Wolfram code, every possible binary neighborhood	Single given input neighborhood	Single given input neighborhood	Single given input neighborhood
<b>Produces</b>	Set of permutations that changes the additive automata to multiplicative, with the given multiplication table	Polynomial	Output visually similar to B	Vector
<b>Data type</b>	Binary	Complex	Complex	Discrete permutation
<b>Base 2 sum of neighborhood</b>	Construction of factors, pre multiplication	Normalization, post multiplication	Construction of factors, pre multiplication	n/a
<b>N-th root in normalization</b>	n/a	N = size of neighborhood	N = number of factors	n/a

49 Multiplications A, additive to multiplicative  
 50  $r$  = specific Wolfram code  
 51  $n$  = binary neighborhood =  $1columnZero + 2columnOne + 4columnTwo$ , points to its value in  
 52  $r$   
 53  $h$  = hypercomplex unit vector from binary  
 54  $H$  = inverse of  $h$ , binary value from hypercomplex unit vector  
 55  $p$  = a permutation of the neighborhood  
 56 using hypercomplex multiplication, a valid permutation set produces:  
 57  $WolframCode(r, n) = WolframCode(r, H(h(p(n)) \cdot h(p(n)) \cdot \dots \cdot h(p(n)) \dots numFactors) )$   
 58

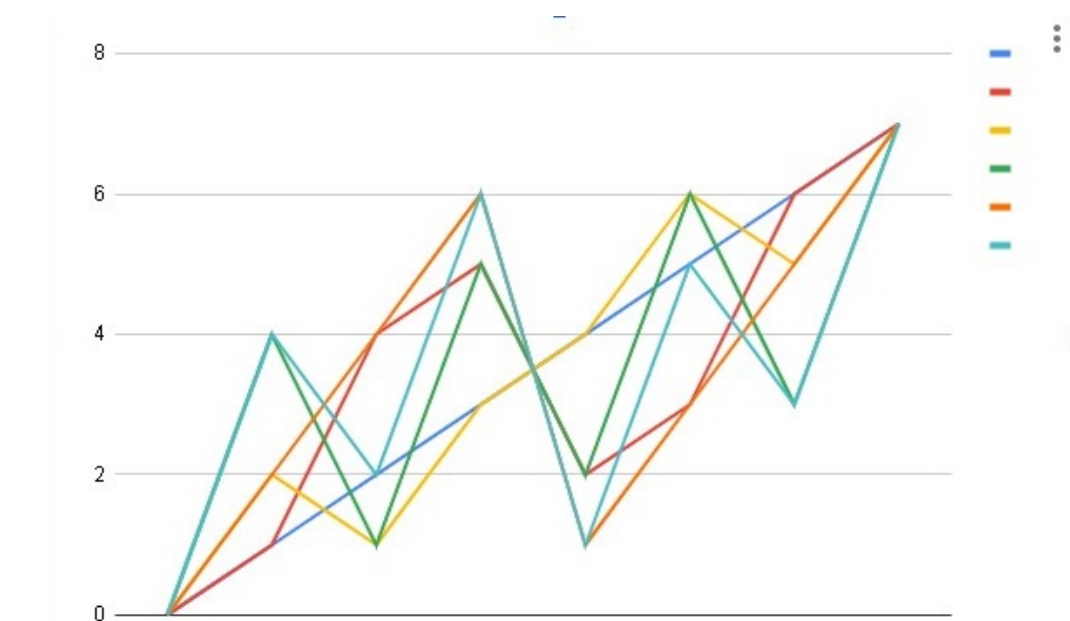
59 The first set of multiplications, column A, brute forces all possible sets of permutations on all  
 60 possible binary neighborhoods of the Wolfram code. A permutation in the set rearranges the  
 61 columns of the input neighborhood, these become a set of factors. A valid set of permutations  
 62 is one that, for all possible input neighborhoods, the set of constructed factors using the  
 63 permuted neighborhoods always multiplies out to a value that points to an equal value within  
 64 the Wolfram code. The set of multiplication results is a pointer array that reproduces the  
 65 original Wolfram code for every possible binary neighborhood.

66 Identity solutions of 5 factors using all zero permutations exist for Wolfram codes up to 32 bits  
 67 in this library using hypercomplex numbers and Galois addition. Galois multiplication takes a  
 68 mix of numbers of factors to get the identity multiplication result array, there is a function in  
 69 the GaloisField class that provides it. The factors constructed are a loose diagonal through the  
 70 multidimensional multiplication table, starting at the origin and ending at the opposite corner  
 71 while zig-zagging. The path lengths of each factor and the result are included in ValidSolution  
 72 results.

73 Permutations of 3 bit neighborhoods

Permutation: 0, [0, 1, 2, 3, 4, 5, 6, 7]  
 Permutation: 1, [0, 1, 4, 5, 2, 3, 6, 7]  
 Permutation: 2, [0, 2, 1, 3, 4, 6, 5, 7]  
 Permutation: 3, [0, 4, 1, 5, 2, 6, 3, 7]  
 Permutation: 4, [0, 2, 4, 6, 1, 3, 5, 7]  
 Permutation: 5, [0, 4, 2, 6, 1, 5, 3, 7]

74  
 75 Flattened path through a six dimensional multiplication table  
 76 Six factors, permutation set = {0,1,2,3,4,5}



77

78

79 Multiplications B and C apply a valid solution from the first set of multiplications to any given  
80 individual neighborhood with binary, non-negative real, and complex values. Multiplication B is  
81 the Cartesian product of the permuted neighborhoods, using a closed partial product table to  
82 generate a polynomial. Multiplication C does the binary sum of complex neighborhood, then  
83 multiplies as complex. Both B and C take the  $n$ -th root of the result, with  $n = \text{numColumns}$   
84 and  $n = \text{numFactors}$ , respectively. Multiplications B and C both include a binary weighted  
85 sum of the neighborhood, same as the construction of the factors from A, though B and C  
86 use complex. B, as part of the normalization and C as the construction. Multiplication C is  
87 the permutation composition product. B, just before the normalization is a neighborhood of  
88 multiplication results, with each column of it being a unit vector coefficient. This multiplication  
89 result neighborhood is permuted by the inverse of the permutation composition product to  
90 properly order the output vector.

91 Control Panel

ECA rule

54

Multiplication Table to use

Permuted Cayley-Dickson

Specific solution to use

0

Degree: 2 = quaternions, 3 = octonions, etc., if applicable

2

Number of factors to use

5

Number of rows in the ECA, 1 row = 3 bit neighborhood, 2 rows = 5 bit neighborhood

1

Partial product table, size = places x places

Galois addition, XOR, 3x3

Keeps functions from running longer than the user want, in seconds

30

the calculate button produces all solutions for the chosen parameters

Refresh

this button re-randomizes and displays the ECA rule with the particular solution number chosen

Display specific solution

Deep search using above parameters

Start deep search

Width of random input 200

Number of factors in logic gate search

5

Logic gate, AND = 8, OR = 14, XOR = 6, etc

6: XOR

Logic gate solution:

Which multiplication table to use

XOR

Partial product table

Galois addition, XOR, 2x2

Refresh logic gate solutions

Refresh

Display specific logic gate solution

Display specific solution

Search all logic gates for solutions and crossreference gates that have solutions in common

Deep logic gate search

Table Display Degree, 2 = quaternions, 3 = octonions, 4 = sedonions

2

Cayley-Dickson permutation number, (cdz, \_\_), down in recursion

0

Cayley-Dickson permutation= number, (\_\_ , cdo), up in recursion

0

Fano plane octonions

0

Galois Field, Prime

2

Galois Field, Power

1

Length of permutations

4

Refresh permuted Cayley-Dickson solutions

Refresh

Display tables with above parameters

Display specific tables

Compare Fano-generated octonions with permuted Cayley-Dickson octonions

Fano/CD Compare

Compare permuted CD with permuted CD

CD v CD

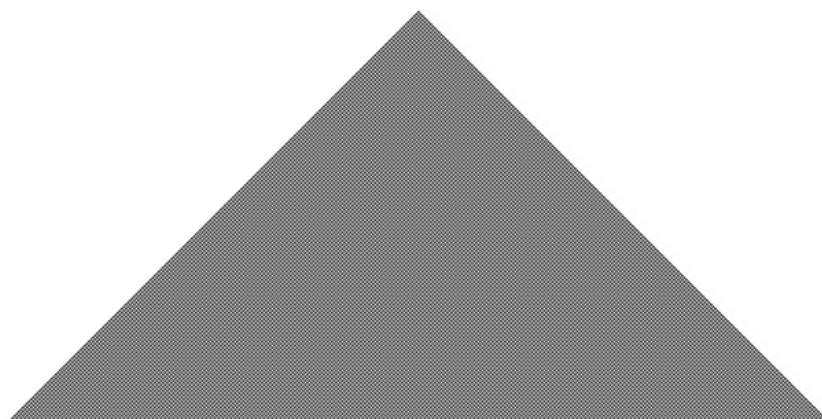
Picks a random Wolfram code with 5 factors, identify solution

Random Wolfram Code

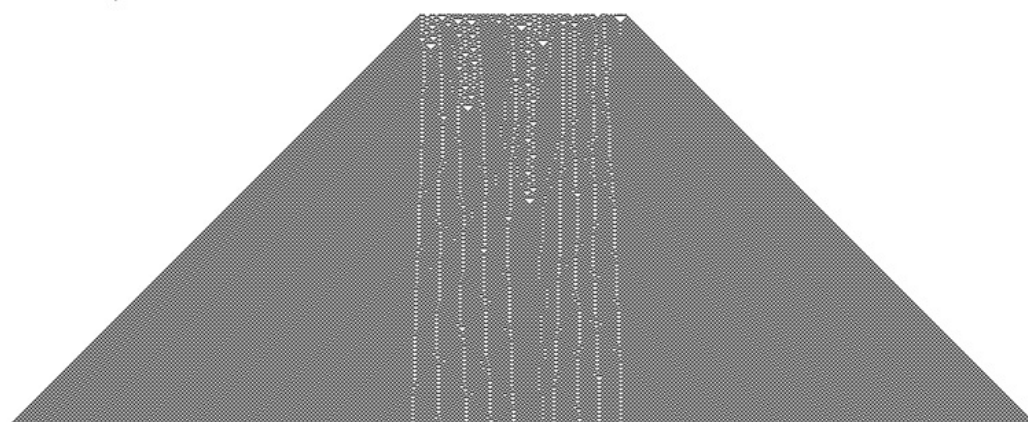
ECA 54, binary and non negative real



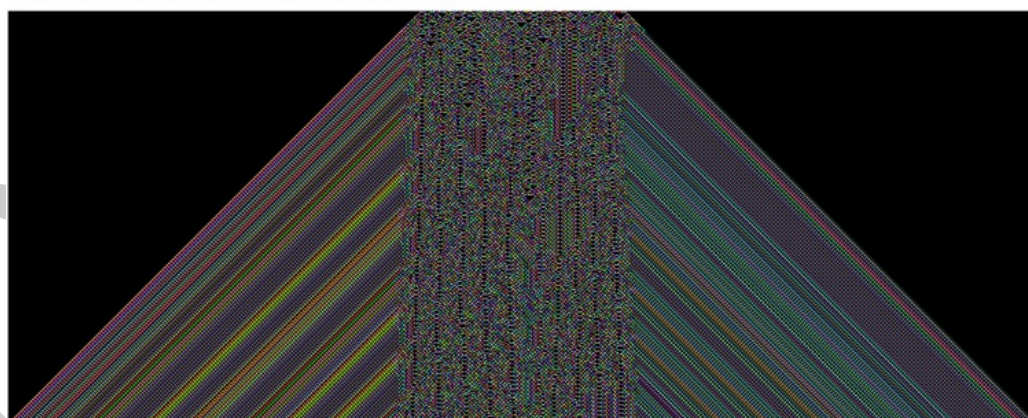
Single bit initial input:



Random initial input:



Same random initial input with solution applied to random  $(0,1)$  non negative real:



95

96

97 ECA 54, solution parameters, including polynomial

```
ValidSolution
Wolfram code: [0, 1, 1, 0, 1, 1, 0, 0]
Permutation: 0 Permuted Axis: [0, 1, 2, 3, 4, 5, 6, 7]
times
Permutation: 0 Permuted Axis: [0, 1, 2, 3, 4, 5, 6, 7]
times
Permutation: 0 Permuted Axis: [0, 1, 2, 3, 4, 5, 6, 7]
times
Permutation: 0 Permuted Axis: [0, 1, 2, 3, 4, 5, 6, 7]
times
Permutation: 0 Permuted Axis: [0, 1, 2, 3, 4, 5, 6, 7]
times
-----
Equals: [0, 1, 2, 3, 4, 5, 6, 7]
Apply Wolfram code to multiplication result
Equals: [0, 1, 1, 0, 1, 1, 0, 0]
Original Wolfram code: [0, 1, 1, 0, 1, 1, 0, 0]

Permutation composition product: 0, inverse: 0

Multiplication table type: 0
2D multiplication table used:
[0, 1, 2, 3, 4, 5, 6, 7]
[1, 4, 7, 2, 5, 0, 3, 6]
[2, 3, 4, 5, 6, 7, 0, 1]
[3, 6, 1, 4, 7, 2, 5, 0]
[4, 5, 6, 7, 0, 1, 2, 3]
[5, 0, 3, 6, 1, 4, 7, 2]
[6, 7, 0, 1, 2, 3, 4, 5]
[7, 2, 5, 0, 3, 6, 1, 4]

numFactors: 5 numBits: 3

1*((a^5)*(b^0)*(c^0)) + 20*((a^3)*(b^1)*(c^1)) + 10*((a^2)*(b^3)*(c^0)) + 10*((a^2)*(b^0)*(c^3)) + 30*((a^1)*(b^2)*(c^2)) +
5*((a^0)*(b^4)*(c^1)) + 5*((a^0)*(b^1)*(c^4))

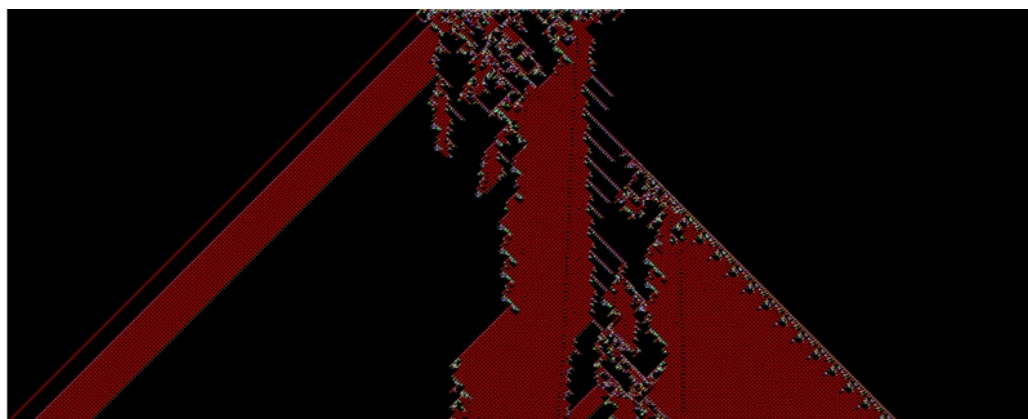
5*((a^4)*(b^1)*(c^0)) + 10*((a^3)*(b^0)*(c^2)) + 30*((a^2)*(b^2)*(c^1)) + 5*((a^1)*(b^4)*(c^0)) + 20*((a^1)*(b^1)*(c^3)) +
10*((a^0)*(b^3)*(c^2)) + 1*((a^0)*(b^0)*(c^5))

5*((a^4)*(b^0)*(c^1)) + 10*((a^3)*(b^2)*(c^0)) + 30*((a^2)*(b^1)*(c^2)) + 20*((a^1)*(b^3)*(c^1)) + 5*((a^1)*(b^0)*(c^4)) +
1*((a^0)*(b^5)*(c^0)) + 10*((a^0)*(b^2)*(c^3))
```

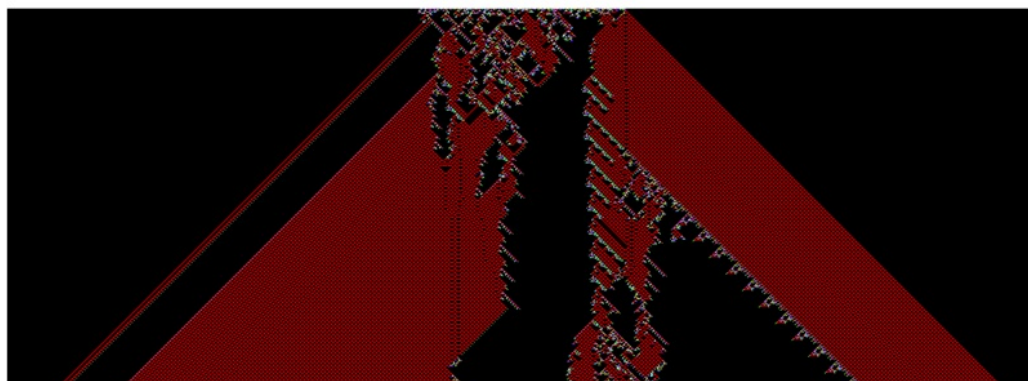
98

99

100 ECA 54, solution output, complex



Complex part



101

102

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103

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