Course Notes
for
MS4111
Discrete Mathematics 1

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CHAPTER 3 Predicate Logic

# 3.1 Predicates and free variables

In chapter 2 we studied the so called Propositional Logic.

Propositional logic does not provide enough instruments to deal with mathematical statements and proofs: we need to generalize the concept of proposition by introducing statements which contain variables.

# Example 3.1 The expression

 $x \ge 1$ 

contains the variable x and the truth value of it depends on the value of x, for example it is true if x = 3 and it is false if x = 0. This is an example of a (logical) predicate.

**Definition 3.1** Let p(x) be a statement with variable x and let D be a set. p is a predicate or propositional function (with respect to D) if for each  $x \in D$ , p(x) is a proposition. We call D the domain of discourse of D.

NOTE: Predicates can involve several variables! Let us see some examples of predicates.

Example 3.2 1) p(x): x is an honest man.

- 2) q(x,y): x is a man capable of doing job y.
- 3) r(x, y, z): The product of number x with number y is z.

Remark 3.12 A predicate p(x) in one variable is often expressing a property of x.

**Remark 3.13** A predicate q(x, y) in two variables is often expressing a relation between x and y.

Remark 3.14 A predicate containing no variables is a proposition.

Remark 3.15 Whenever we replace the variables of a predicate with constants, the predicate becomes a proposition.

The variables involved in a predicate are called free variables.

3.2 Predicates and bound variables: quantifiers

# 3.2.1 Universal quantifier

**Definition 3.2** Given a predicate p(x), consider the statement

for every 
$$x$$
,  $p(x)$ . (3.1)

The above statement is called the universally quantified statement and it is denoted by

$$\forall x, \quad p(x), \tag{3.2}$$

where the symbol  $\forall$  is called the universal quantifier. (3.1) (or (3.2)) is defined in the following way: it is true if p(x) is true for every x and it is false if p(x) is false for at least one x.

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(3.2) is red:
for every x, p(x)
for any x, p(x)
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for all x, p(x)

**Example 3.3** p(x): x is an honest man.

The universally quantified statement is

 $\forall x, p(x)$ : "For all x, x is an honest man", or, "Every man is honest".

Remark 3.16 The universally quantified statement

$$\forall x, p(x)$$

is false if p(x) is false for at least one x. A value x that makes p(x) false is called a counterexample to

$$\forall x, p(x).$$

Exercise 3.1 The statement

$$\forall x \in \mathbf{R}, \quad x^2 - 1 > 0$$

is false since if x = 1, proposition

$$1^2 - 1^2 > 0$$

is false. The value 1 is a counterexample to the statement

$$\forall x \in \mathbf{R}, \quad x^2 - 1 > 0.$$

# 3.2.2 Existential quantifier

**Definition 3.3** Given a predicate p(x), consider the statement

for some 
$$x$$
,  $p(x)$ . (3.3)

The above statement is called the existentially quantified statement and it is denoted by

$$\exists x, \quad p(x), \tag{3.4}$$

where the symbol  $\exists$  is called the existential quantifier. (3.3) (or (3.4)) is defined in the following way: it is true if p(x) is true for at least one x and it is false if p(x) is false for every x.

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(3.4) is red:

for some x, p(x)

for at least one x, p(x)

there exists x such that, p(x)
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**Example 3.4** p(x): x is an honest man.

The existentially quantified statement is

 $\exists x, p(x)$ : "There exists at least one x, such that x is an honest man" or, "There exists at least one honest man".

Remark 3.17 By introducing quantifiers in a predicate p(x), the variable x becomes silent and we say that x becomes a bound variable. Moreover the predicate becomes a proposition.

**Note:** If there is more than one quantifier in a predicate, then the order of the quantifiers is important!

# Example 3.5 Consider the predicate

q(x,y): x is a man capable of doing job y.

The quantified statement

$$\forall y, \exists x: q(x,y)$$

is red "For any job y, there is a man x who is capable of doing job y".

The quantified statement

$$\exists x, \ \forall y: \quad q(x,y)$$

is red "There is a man x that is capable of doing any job y".

In general, given a predicate p(x), we can restrict the range of variation of the free variable to the domain of discourse D

$$\forall x \in D, \quad p(x)$$

and the above universally quantified statement is true if p(x) is true for all  $x \in D$  and it is false if p(x) is false for at least one  $x \in D$ . We can similarly define

$$\exists x \in D, \quad p(x)$$

which is true if p(x) is true for at least one  $x \in D$  and it is false if p(x) is false for all  $x \in D$ .

3.3 Rules for negating quantifiers

# 3.3.1 Generalized De Morgan Laws (for logic)

The rules for negating quantifiers is given by the following important theorem.

Theorem 3.1 Generalized De Morgan Laws (for logic)

(a) 
$$\left(\forall x, \ p(x)\right) \equiv \exists x, \ \overline{p(x)}$$

**b**) 
$$\left( \exists x, \ p(x) \right) \equiv \forall x, \ \overline{p(x)}$$

#### Proof.

a) We want to prove that

$$\overline{\left(\forall x, \ p(x)\right)}$$
 and  $\exists x, \ \overline{p(x)}$ 

have the same truth values.

If  $(\forall x, p(x))$  is TRUE, then  $\forall x, p(x)$  is FALSE, then p(x) is

FALSE for at least one x, then  $\overline{p(x)}$  is TRUE for at least one x, i.e.  $\exists x, \ \overline{p(x)}$  is TRUE.

If  $(\forall x, p(x))$  is FALSE, then  $\forall x, p(x)$  is TRUE, then it is NOT TRUE that  $\exists x, \overline{p(x)}$ , i.e.  $\exists x, \overline{p(x)}$  is FALSE.

b) We want to prove that

$$\overline{\left(\exists x, \ p(x)\right)}$$
 and  $\forall x, \ \overline{p(x)}$ 

have the same truth values.

If  $(\exists x, p(x))$  is TRUE, then  $\exists x, p(x)$  is FALSE, then p(x) is

FALSE for every x, then  $\overline{p(x)}$  is TRUE for every x, i.e.  $\forall x, \ \overline{p(x)}$  is TRUE.

If  $(\exists x, p(x))$  is FALSE, then  $\exists x, p(x)$  is TRUE, then it is NOT

TRUE that p(x) is false for any x, i.e.  $\forall x, p(x)$  is FALSE.

### 3.3.2 Negation of quantified statement restricted to D

The Generalized De Morgan Laws (for logic) can be easily adapted to the case in which the quantified statements are restricted to a domain of discourse D in the following way

$$\mathbf{a'}$$
)  $\left(\forall x \in D, \ p(x)\right) \equiv \exists x \in D, \ \overline{p(x)}$ 

$$\mathbf{b'}$$
)  $\left(\exists x \in D, \ p(x)\right) \equiv \forall x \in D, \ \overline{p(x)}$ 

# 3.3.3 Final remarks on the Generalized De Morgan Laws

Remark 3.18 Given n propositions

$$p_1, p_2, \ldots, p_n$$

consider the compound proposition

$$p_1 \wedge p_2 \wedge \cdots \wedge p_n. \tag{3.5}$$

(3.5) is TRUE when  $p_i$  is true, for every i = 1, ..., n. The universally quantified proposition

$$\forall x, \ p(x) \tag{3.6}$$

generalizes (3.5) because it is TRUE when p(x) is true for any x.

# Remark 3.19 Given n propositions

$$p_1, p_2, \ldots, p_n$$

consider the compound proposition

$$p_1 \vee p_2 \vee \dots \vee p_n. \tag{3.7}$$

(3.5) is TRUE when there exists i, with  $1 \le i \le n$  such that  $p_i$  is true. The existentially quantified proposition

$$\exists x, \ p(x) \tag{3.8}$$

generalizes (3.7) because it is TRUE when p(x) is true for some x.

Remark 3.20 From the De Morgan Laws for logic we have

$$\overline{p_1 \wedge p_2 \wedge \dots \wedge p_n} \equiv \overline{p}_1 \vee \overline{p_2} \vee \dots \vee \overline{p_n}. \tag{3.9}$$

From the generalized De Morgan Laws we have

$$\overline{\forall x, p(x)} \equiv \exists x, \overline{p(x)}$$
(3.10)

and because of remarks 3.18 and 3.19 it makes sense to call (3.10)

Generalized De Morgan Law (it generalizes (3.9)).

In the same way, if we compare the De Morgan Law for logic

$$\overline{p_1 \vee p_2 \vee \dots \vee p_n} \equiv \overline{p}_1 \wedge \overline{p_2} \wedge \dots \wedge \overline{p_n}. \tag{3.11}$$

with the generalized De Morgan law

$$\overline{\exists x, p(x)} \equiv \forall x, \overline{p(x)},$$
(3.12)

it makes sense to call (3.12) Generalized De Morgan Law because of remarks 3.18 and 3.19 (it generalizes (3.11)).

# 3.4 Revision Exercises

Exercise 3.2 Prove that

$$\forall x \in \mathbf{R}, \quad x > 1 \implies x + 1 > 1 \tag{3.13}$$

is true.

Answer. Let us denote

$$p(x): x > 1 \implies x + 1 > 1.$$

We want to prove that p(x) is true for any real number x.

1) If x > 1 is false then p(x) is true because of the false hypothesis (independently of the truth value of x + 1 > 1).

2) If x > 1 is true, then we have

$$x + 1 > x > 1$$

therefore x + 1 > 1 is true, therefore

$$p(x): x > 1 \implies x + 1 > 1$$

is true because of type  $T \Rightarrow T$ .

Therefore we proved that p(x) is true for any real number x, therefore

$$\forall x \in \mathbf{R}, \quad x > 1 \implies x + 1 > 1$$

is true.

# Exercise 3.3 Prove that

 $\exists positive integer n, \quad nis prime \Rightarrow n+1, n+2, n+3, n+4 \quad are not prime.$ (3.14)

is true.

Answer. Let us denote

 $p(n): n \text{ is prime } \Rightarrow n+1, n+2, n+3, n+4 \text{ are not prime}$ 

We just need to find a positive integer n for which p(n) is true. Take

$$n = 23,$$

we obtain the proposition

 $p(23): 23 \text{ is prime} \implies 24, 25, 26, 27 \text{ are not prime},$ 

which is a true proposition since it is of type  $T \Rightarrow T$ .

**Note:** p(n) in the above exercise is not true for every values of n, if we take for example n=2, we obtain

 $p(2): 2 \text{ is prime} \implies 3, 4, 5, 6 \text{ are not prime},$ 

which is of type  $T \Rightarrow F$  and therefore it is false.

#### Exercise 3.4 Prove that

$$\exists x \in \mathbf{R}, \quad \frac{1}{x^2 + 1} > 1.$$
 (3.15)

is false.

Answer. If we denote

$$p(x): \frac{1}{x^2+1} > 1,$$

we want to prove that p(x) is false for any  $x \in \mathbf{R}$ , i.e. we want to prove that

$$\frac{1}{x^2+1} > 1$$

is false for any  $x \in \mathbf{R}$ , i.e. we want to prove that  $\overline{p(x)}$  is true for any  $x \in \mathbf{R}$  which means that

$$\frac{1}{x^2+1} \le 1$$

is true for any  $x \in \mathbf{R}$ . In other words we are going to prove that the following statement

$$\forall x \in \mathbf{R}, \quad \frac{1}{x^2 + 1} \le 1 \tag{3.16}$$

is true. Let x be a real number, then  $x^2 \ge 0$ , therefore

$$x^2 + 1 \ge 1,$$

therefore

$$\frac{1}{x^2+1} \le 1,$$

which proves that statement (3.16) is true and therefore that the statement

$$\exists x \in \mathbf{R}, \quad \frac{1}{x^2 + 1} > 1$$

is false.

**Note:** In the above exercise we had p(x):  $\frac{1}{x^2+1} > 1$  and we proved that

$$\exists x, \ p(x)$$

is false by proving that

$$\forall x, \ \overline{p(x)}$$

is true: this method is correct because of the Generalized De Morgan Laws

$$\overline{\exists x, p(x)} \equiv \forall x, \ \overline{p(x)},$$

therefore to prove that  $\exists x, p(x)$  is false, we proved that its negation  $\overline{\exists x, p(x)}$  is true by making use of the Generalized De Morgan Law and proving that  $\forall x, \overline{p(x)}$  is true.

Exercise 3.5 Prove that

$$\forall x \in \mathbf{R}, \exists y \in \mathbf{R}, \quad x + y = 0. \tag{3.17}$$

is true.

Answer. If we denote

$$p(x,y): x+y=0,$$

then we want to prove that

$$\exists y \in \mathbf{R}, \quad p(x,y)$$

is true for all  $x \in \mathbf{R}$ . Take a general  $x \in \mathbf{R}$  and we want to prove that

is true for at least one  $y \in \mathbf{R}$ . If we take

$$y = -x$$

then

$$x + y = x - x = 0$$

and we found one value of y for which p(x, y) is true, i.e.

$$\exists y \in \mathbf{R}, \quad p(x,y)$$

is true and we proved it by taking a general  $x \in \mathbf{R}$ , which proves that

$$\exists y \in \mathbf{R}, \quad p(x,y)$$

is true for all  $x \in \mathbf{R}$ , i.e.

$$\forall x \in \mathbf{R}, \exists y \in \mathbf{R}, \quad x + y = 0$$

is true.

### Exercise 3.6 Prove that

$$\exists y \in \mathbf{R}, \forall x \in \mathbf{R}, \quad x + y = 0. \tag{3.18}$$

is false.

Answer. If we denote

$$p(x,y): x+y=0,$$

then we can prove that

$$\exists y \in \mathbf{R}, \ \forall x \in \mathbf{R}, \quad p(x,y)$$

is false by proving that its negation is true, but its negation can be expressed by making use of the Generalized De Morgan Law by

$$\exists y, \, \forall x, \quad p(x,y) \equiv \forall y, \overline{\forall x, \quad p(x,y)} 
 \equiv \forall y, \exists x, \, \overline{x+y=0} 
 \equiv \forall y, \exists x, \, x+y \neq 0.$$

We therefore want to prove that the statement

$$\forall y, \exists x, \ x + y \neq 0$$

is true, i.e. that the statement

$$\exists x, \ x + y \neq 0$$

is true for every y. Take a general  $y \in \mathbf{R}$  and we want to prove that

$$x + y \neq 0$$
,

for at least a value of  $x \in \mathbf{R}$ . If we take

$$x = 1 - y,$$

we obtain

$$x + y = 1 - y + y = 1 \neq 0,$$

therefore we proved that

$$x + y \neq 0$$

for at least one value of x given by x = 1 - y, i.e. we proved that

$$\exists x, \ x + y \neq 0$$

is true and it is true for any value  $y \in \mathbf{R}$ , i.e.

$$\forall y, \exists x, \ x + y \neq 0$$

is true.

Exercise 3.7 Prove that

$$\forall x \in \mathbf{R}, \ \forall y \in \mathbf{R}, \quad x^2 < y^2 \ \Rightarrow \ x < y. \tag{3.19}$$

is false.

Answer. If we denote

$$p(x,y): x^2 < y^2 \implies x < y,$$

we want to prove that

$$\forall x \in \mathbf{R}, \ \forall y \in \mathbf{R}, \quad p(x,y)$$

is true, i.e.

$$\exists x \in \mathbf{R}, \exists y \in \mathbf{R}, \quad \overline{p(x,y)}$$

is true. In other words we want to find a counterexample: take

$$x = 1,$$
  $y = -2$ 

and with the above choices we have that 1 < 4 is true and 1 < -2 is false, therefore

$$1 < 4 \Rightarrow 1 < -2$$

is false because of type  $T \Rightarrow F$ , therefore

$$\overline{1 < 4 \implies 1 < -2}$$

is true for the values x = 1 and y = -2.

Exercise 3.8 Prove that

$$\forall x \in \mathbf{R}, \ \exists y \in \mathbf{R}, \quad x^2 < y^2 \ \Rightarrow \ x < y. \tag{3.20}$$

is true.

Answer. If we denote

$$p(x,y): x^2 < y^2 \implies x < y,$$

we want to prove that

$$\exists y \in \mathbf{R}, \quad p(x,y)$$

is true for all  $x \in \mathbf{R}$ . Take a general  $x \in \mathbf{R}$ , we want to prove that

is true for at least one value of  $y \in \mathbf{R}$ . If we take y = 0, then we obtain

$$p(x,0): x^2 < 0 \implies x < 0,$$

which is of type  $F \Rightarrow T$  and therefore it is true. We proved that

$$\exists y \in \mathbf{R}, \ p(x,y)$$

is true (by taking y = 0) and this is true for any  $x \in \mathbf{R}$ , therefore we proved that

$$\forall x \in \mathbf{R}, \ \exists y \in \mathbf{R}, \quad p(x,y)$$

is true.