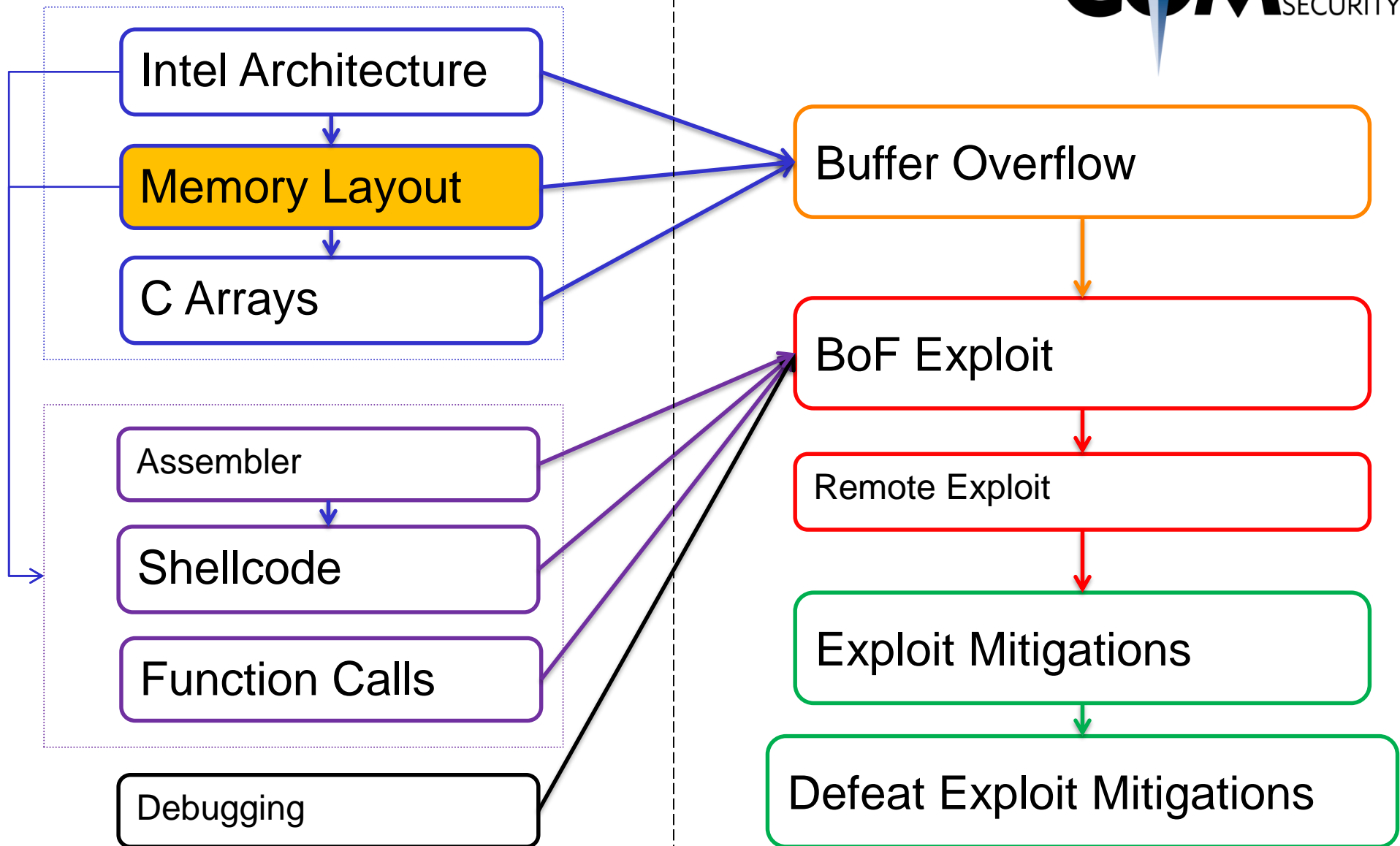


Memory Layout

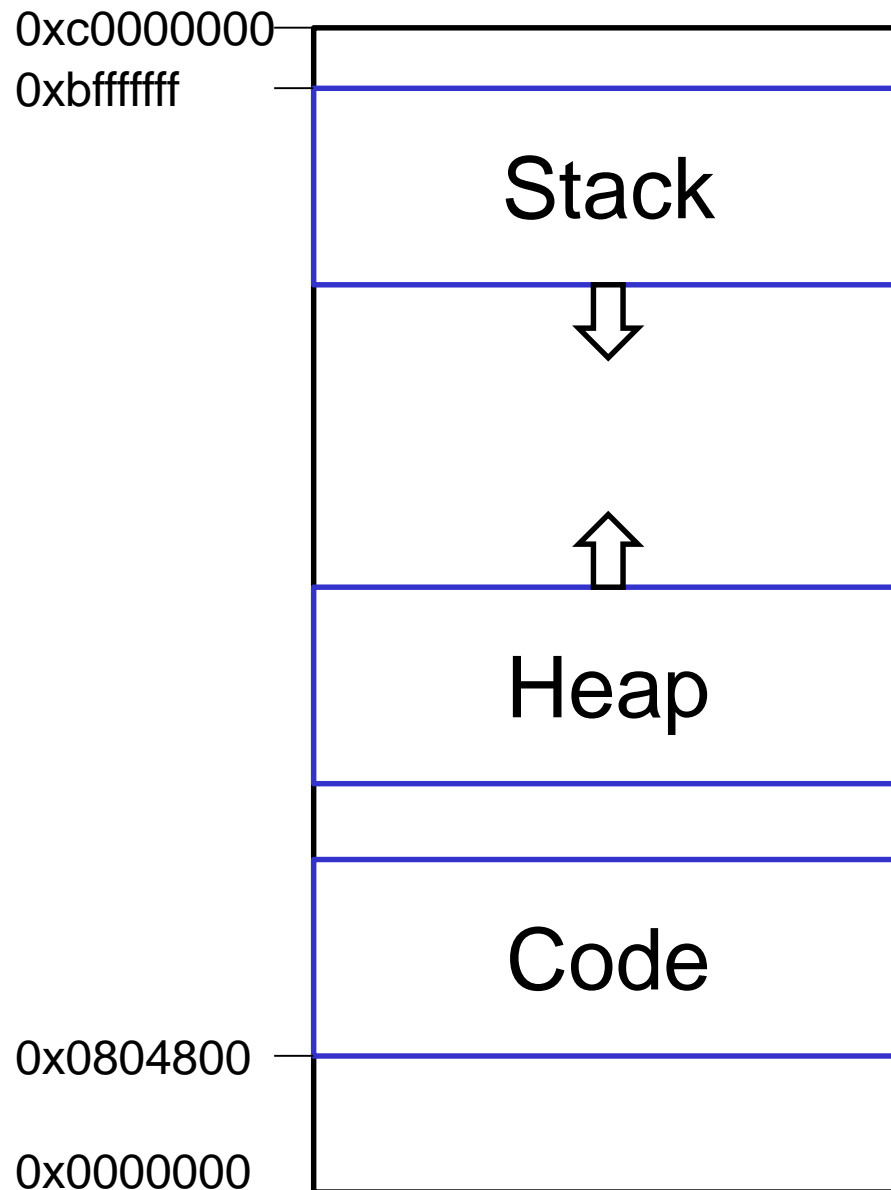
Linux Userspace Process Memory Layout



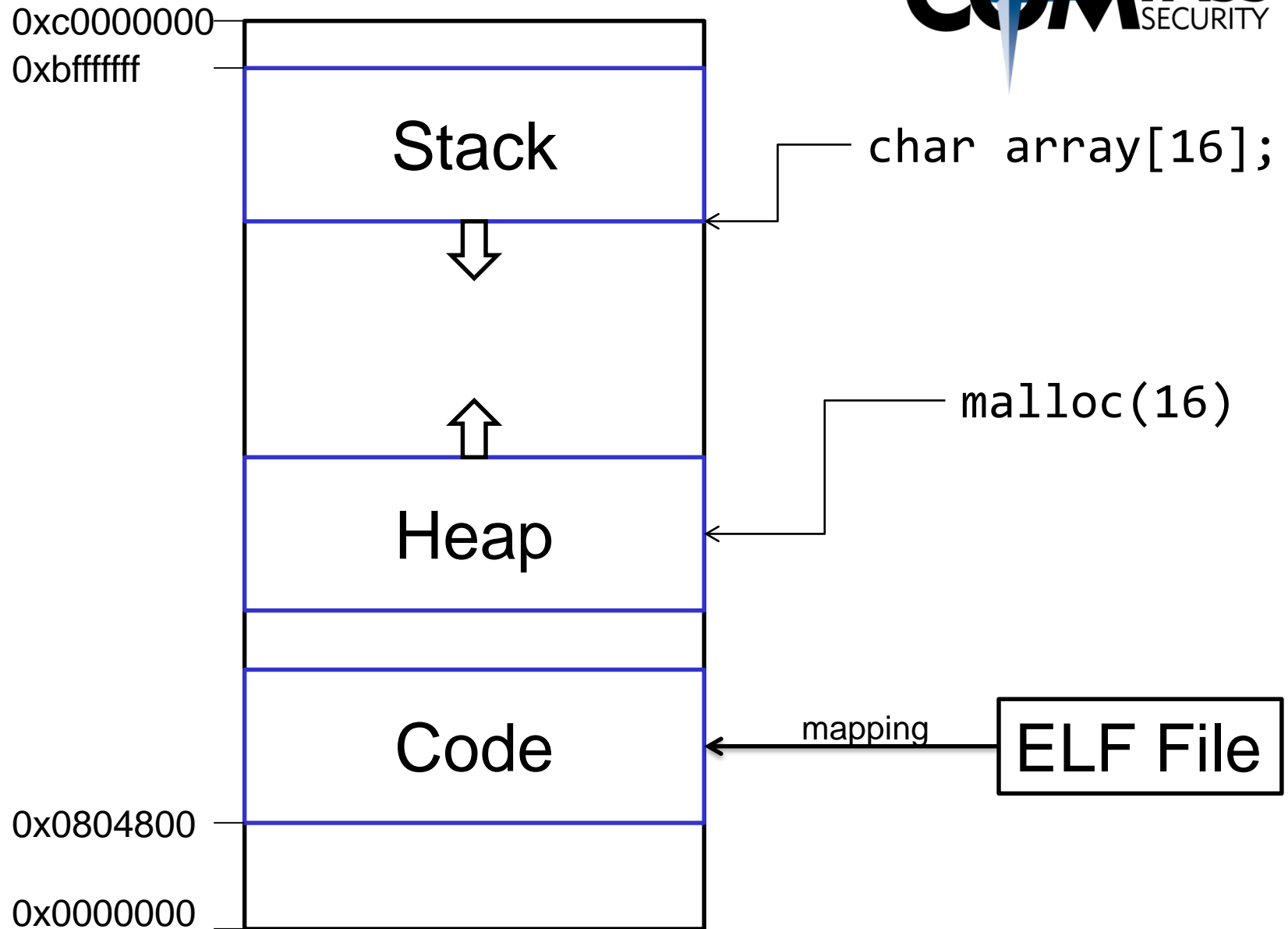
Userspace Memory Layout

In x32

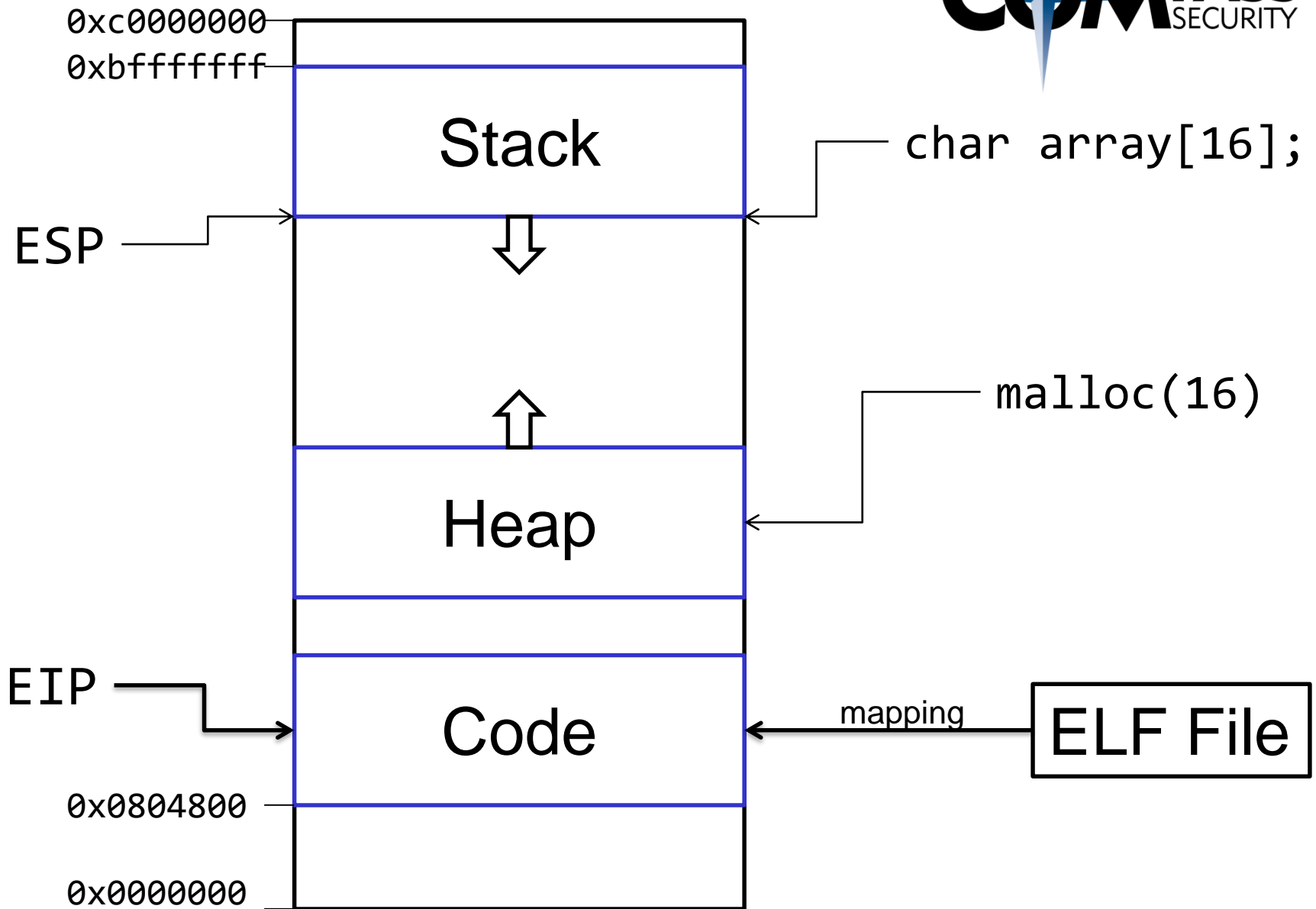
x32 Memory Layout



x32 Memory Layout



x32 Memory Layout



Memory regions:

Stack

- ✦ There's one contiguous memory region containing the stack for the process
- ✦ LIFO – Last in, First Out
- ✦ Contains function **local variables**
- ✦ Also contains: **Saved Instruction Pointer (SIP)**
- ✦ Current function adds data to the top (bottom) of the stack

Heap

- ✦ There's one contiguous memory region containing the heap
- ✦ Memory allocator returns specific pieces of the memory region
- ✦ For **malloc()**
- ✦ Also contains: heap management data

Memory regions:

Code

- ✦ Compiled program code

ELF Format

How do programs on disk look like

Programs are stored in ELF files

ELF: Executable and Linkable Format

- ✦ Previously: "a.out" (Linux 1.2)
- ✦ Like COFF, PE (EXE), COM, ...

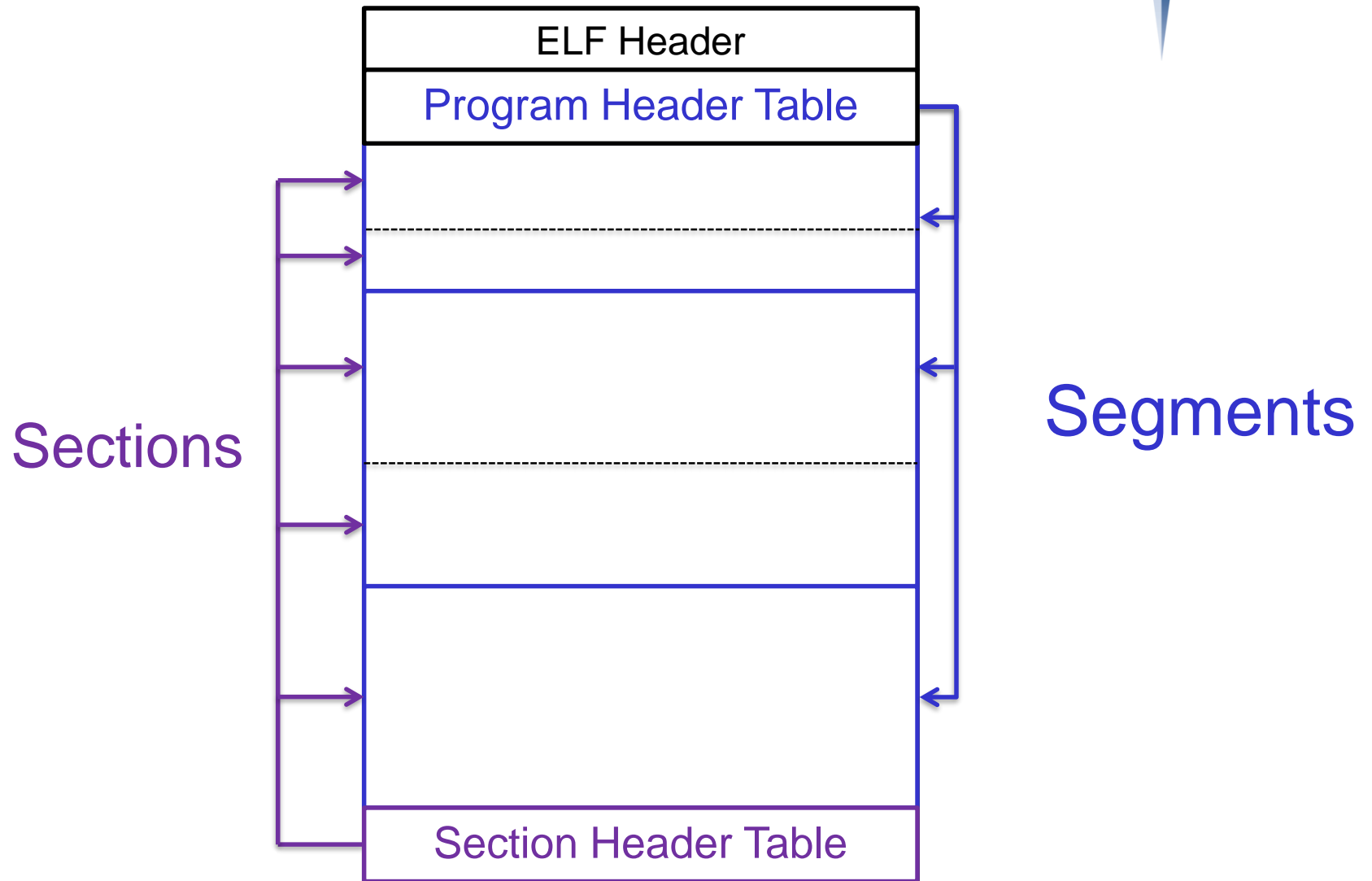
ELF types:

- ✦ ET_EXEC: Executable File
- ✦ ET_REL: Relocatable File
- ✦ ET_DYN: Shared Object File

ELF "views":

- ✦ Sections
- ✦ Segments

\$ readelf -l <binary>



Program Headers:

| Type | Offset FileSiz | VirtAddr MemSiz | PhysAddr Flags Align |
|------------------|--------------------------|------------------------------|--|
| PHDR | 0x00000040 0x000001c0 | 0x0000400040 0x00000001c0 | 0x0000000000400040 R E 8 |
| INTERP | 0x00000200 0x0000001c | 0x0000400200 0x000000001c | 0x0000000000400200 R 1 |
| LOAD | 0x00000000 0x00000b24 | 0x0000400000 0x0000000b24 | 0x0000000000400000 R E 200000 |
| LOAD | 0x00000b28 0x00000270 | 0x0000600b28 0x0000000278 | 0x0000000000600b28 RW 200000 |
| DYNAMIC | 0x00000b40 0x000001e0 | 0x0000600b40 0x00000001e0 | 0x0000000000600b40 RW 8 |
| NOTE | 0x0000021c 0x00000044 | 0x000040021c 0x0000000044 | 0x000000000040021c R 4 |
| GNU_EH_FRAME | 0x000009ac 0x00000044 | 0x00004009ac 0x0000000044 | 0x00000000004009ac R 4 |
| GNU_STACK | 0x00000000 0x00000000 | 0x0000000000 0x0000000000 | 0x0000000000000000 RW 10 |

```
$ readelf -l challenge0
```

Section to Segment mapping:

Segment Sections...

00

01 .interp

02

.interp .note.ABI-tag .note.gnu.build-id .gnu.hash

.dynsym .dynstr .gnu.version .gnu.version_r

.rela.dyn .rela.plt .init **.plt .text .fini .rodata**

.eh_frame_hdr .eh_frame

03

.init_array .fini_array .jcr .dynamic **.got .got.plt**

.data .bss

04

.dynamic

05

.note.ABI-tag .note.gnu.build-id

06

.eh_frame_hdr

07

Sections:

- ✦ .text: Executable instructions
- ✦ .bss: Uninitialized data (usually the heap)
- ✦ .data: initialized data
- ✦ .rodata: Read-Only data
- ✦ .got: Global Offset Table
- ✦ .plt: Procedure Linkage Table
- ✦ .init/.fini: Initialization instructions (“glibc”)

Program Headers:

| Type | Offset | PhysAddr |
|----------------|---------------------|--------------------|
| | FileSiz | Flags Align |
| (02) LOAD | 0x0000000000000000 | 0x0000000000400000 |
| | 0x00000000000000b24 | R E 200000 |
| (03) LOAD | 0x00000000000000b28 | 0x0000000000600b28 |
| | 0x00000000000000270 | RW 200000 |
| (07) GNU_STACK | 0x0000000000000000 | 0x0000000000000000 |
| | 0x0000000000000000 | RW 10 |

02 .init .plt .text .fini .rodata

03 .got .got.plt .data .bss

07

ELF Loader

Compass Security Schweiz AG
Werkstrasse 20
Postfach 2038
CH-8645 Jona

Tel +41 55 214 41 60
Fax +41 55 214 41 61
team@csnc.ch
www.csnc.ch

| |
|----------------------|
| ELF Header |
| Program Header Table |
| .plt |
| .text |
| .init |
| .got |
| .data |
| .bss |
| |
| Section Header Table |

02 Executable Segment

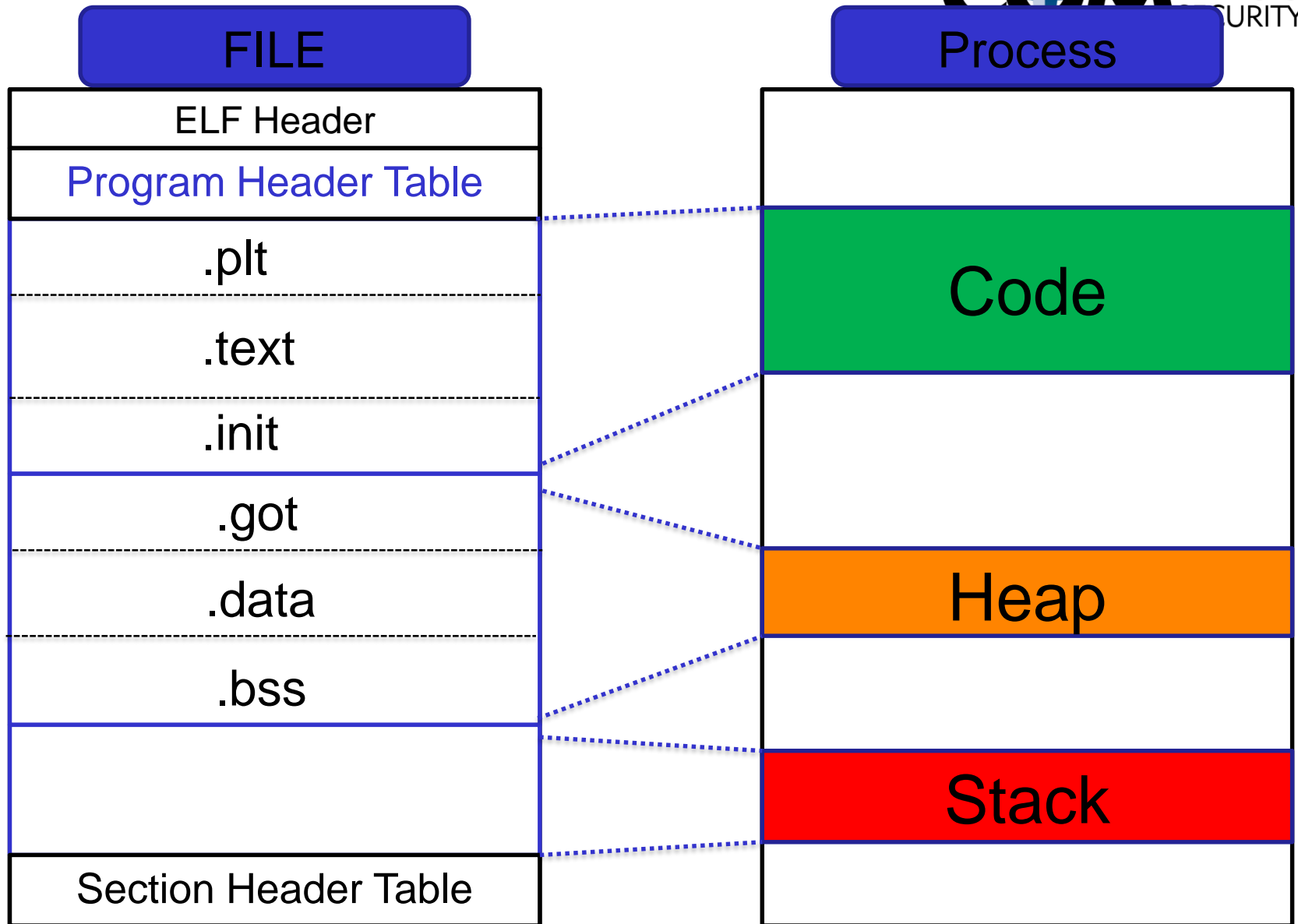
r-X

03 Data Segment

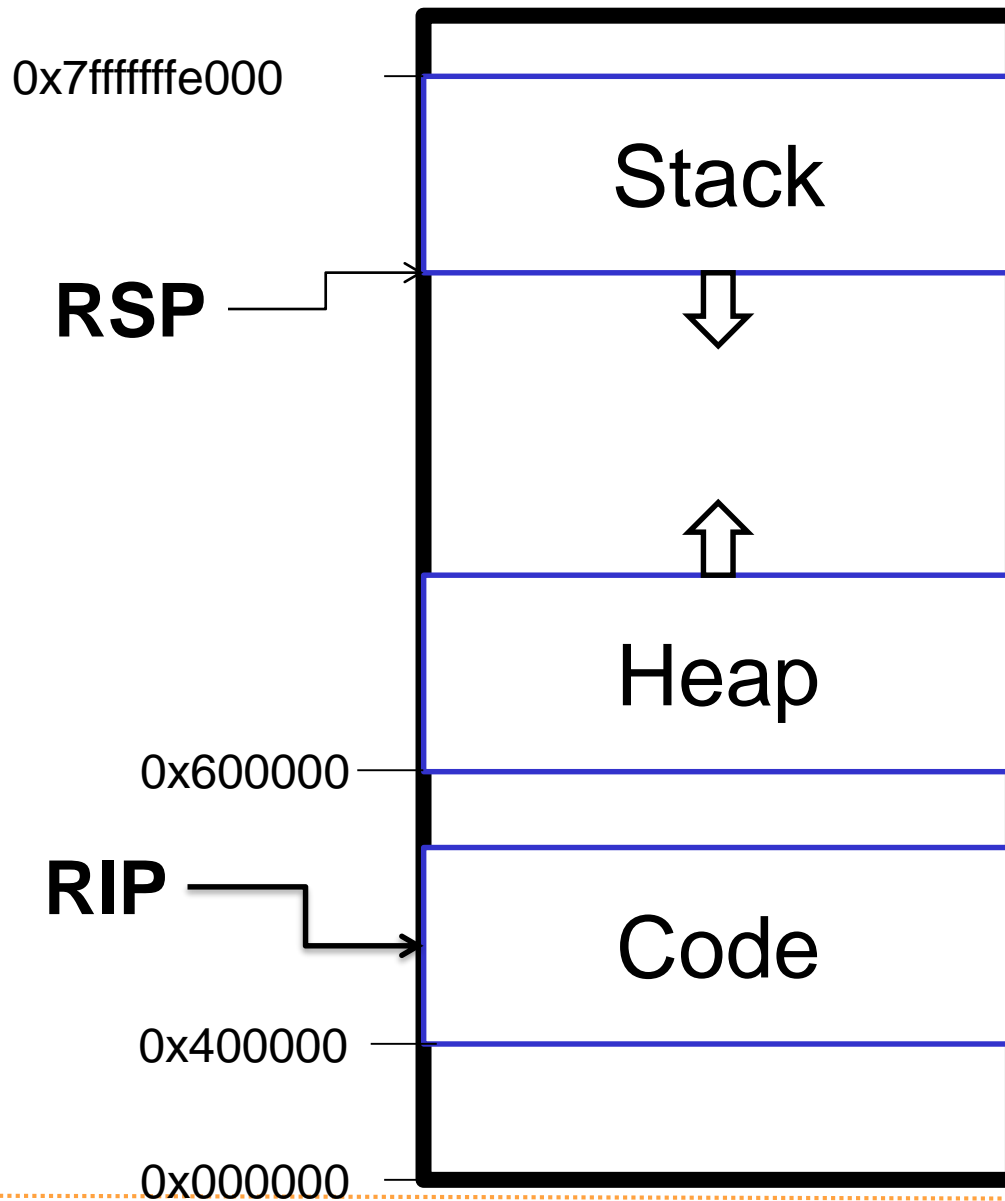
rw-

07 Stack

rw-



x64 Memory Layout



Lets do an example

some static and dynamic binary analysis

```
char *globalVar = "Global";

void main(void) {
    char stackVar[16];
    char *heapVar = (char *) malloc(4);

    printf("Global var: %p\n", globalVar);
    printf("Heap var   : %p\n", heapVar);
    printf("Stack var  : %p\n", stackVar);
}
```

Global var: **0x400654**

Heap var : **0x601010**

Stack var : **0x7fffffffefe990**

(2) **LOAD** **0x000000000000400000**

R E 200000

(3) **LOAD** **0x000000000000600b28**

RW 200000

(7) **GNU_STACK** **0x00000000000000000000**

RW 10

See it at runtime

```
# cat /proc/self/maps
```

```
00400000-0040c000 r-xp 00000000 08:01 391694 /bin/cat
```

```
0060b000-0060c000 r--p 0000b000 08:01 391694 /bin/cat
```

```
0060c000-0060d000 rw-p 0000c000 08:01 391694 /bin/cat
```

```
...
```

```
7fffffffde000-7fffffffff000 rw-p 00000000 00:00 0 [stack]
```

Show Code section, and disassemble:

```
$ objdump -d ./challenge1
```

```
./challenge1:      file format elf64-x86-64
```

```
Disassembly of section .init:
```

```
0000000000400588 <_init>:
```

```
...
```

```
000000000040077f <handleData>:
```

| | | | |
|---------|-------------|------|------------------|
| 40077f: | 55 | push | %rbp |
| 400780: | 48 89 e5 | mov | %rsp,%rbp |
| 400783: | 48 83 ec 30 | sub | \$0x30,%rsp |
| 400787: | 48 89 7d d8 | mov | %rdi,-0x28(%rbp) |
| 40078b: | 48 89 75 d0 | mov | %rsi,-0x30(%rbp) |

The process of creating a process from an ELF file is called:

- ✦ “Linking and Loading”

Sections:

- ✦ Are for compiler (gcc), to link several object files together (.o)

Segments:

- ✦ Are for the loader, to create the process
- ✦ Consists of one or more sections

Recap:

- ✦ Program Code is stored in ELF Files
- ✦ ELF Files contain segments
- ✦ Segments are copied 1:1 in the memory

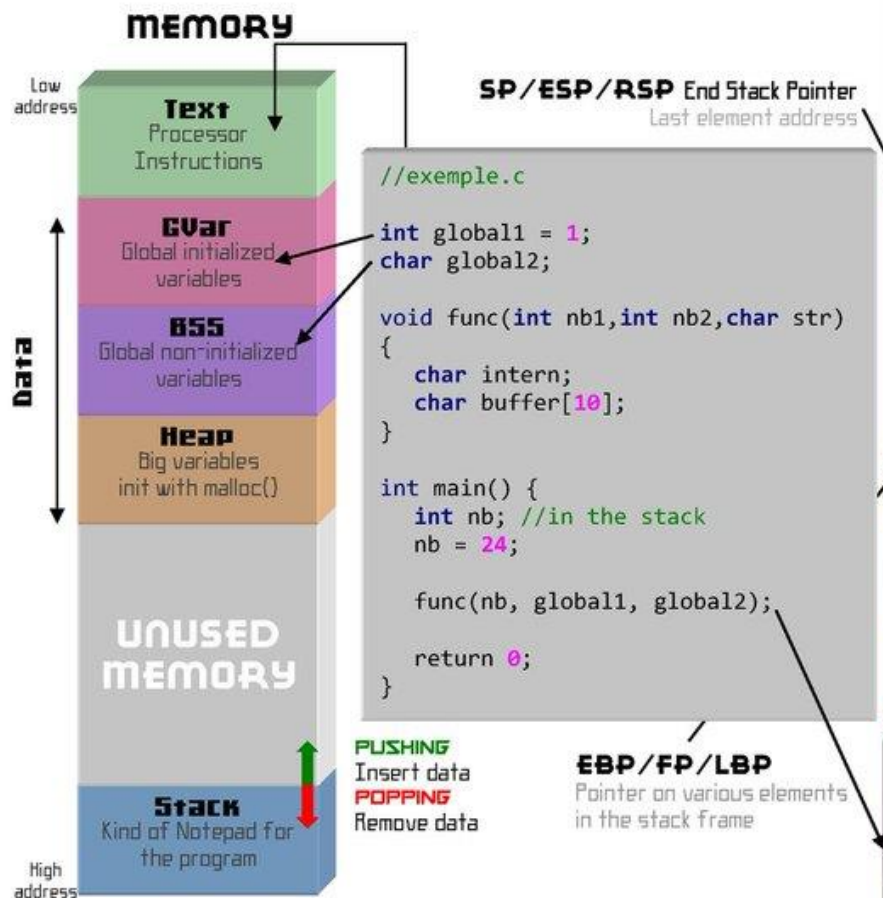
Challenges:

<https://exploit.courses>

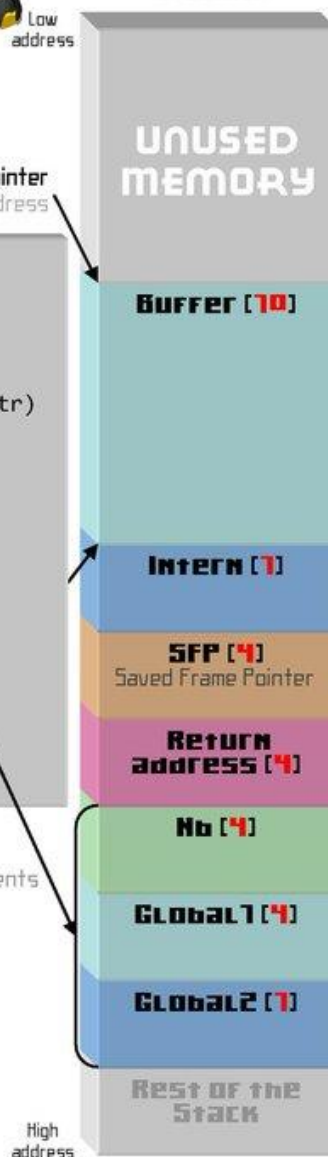
- ✦ Challenge 0: Introduction to memory layout – basic
- ✦ Challenge 1: Introduction to memory layout - advanced

memory Segmentation Cheat Sheet

0xOff.info production, inspired by bases-hacking.org



STACK



GLOBAL REGISTERS

Used for general purpose

X86

AL/AH/AX/EAX/RAX

BL/BH/BX/EBX/RBX

CL/CH/CX/ECX/RCX

DL/DH/DX/EDX/RDX

ARM

r0-r12

| 64 BITS REGISTER | | | |
|------------------|--------|-------|------|
| 63..32 | 31..16 | 15..8 | 7..0 |



SEGMENT REGISTERS
Use to easily read/write to memory

X86

| | |
|------------|--------------------|
| CS : Code | ES : Extra data #1 |
| DS : Data | FS : Extra data #2 |
| SS : Stack | GS : Extra data #3 |

MEMORY ALIGNMENT
Data must be aligned on 4,8,16... Bytes, depending on your system.

EXAMPLE
There is an 3 bytes long empty gap between intern and SFP.

BUFFER OVERFLOW

when input is longer than the allocated memory space.

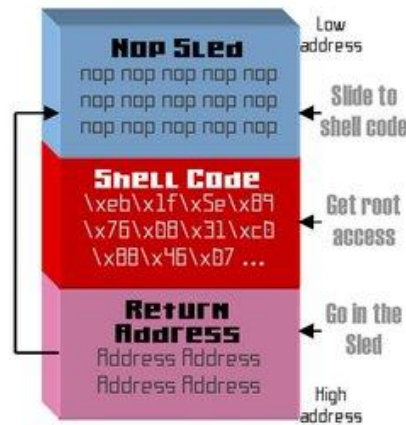
[Stack based]

Smart overwrite of return address

EXAMPLE

Put a 22 bytes long string and overwrite intern Return address.

EXPLOIT ANATOMY



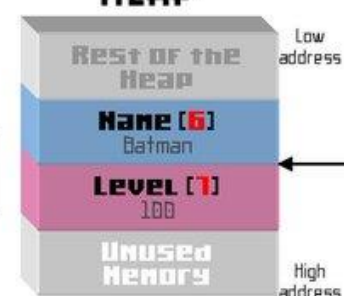
[Heap based]

Smart overwrite of others variables like file name.

EXAMPLE

You have to enter a 6 letters name in the character builder of a game. You enter **Batman\\x64**, to overwrite level variable and get max stats!

HEAP



HOW TO WEAKEN A PROGRAM TRAINING

LINUX

Disable ASLR (random memory address):

echo 0 > /proc/sys/kernel/randomize_va_space

Disable non-executable stack:

\$ gcc -z execstack ...

Disable stack protector:

\$ gcc -fno-stack-protector ...

Force 32-bits compilation mode:

\$ gcc -m32 ...