

Bottom-up Chart Parsing: the CKY algorithm

Data Structures and Algorithms for Computational Linguistics III (ISCL-BA-07)

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Seminar für Sprachwissenschaft

Winter Semester 2022/23

version: 19b769a-20221012-20

Parsing so far

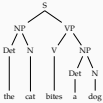
- Parsing is the task of automatic syntactic analysis
- For most practical purposes, context-free grammars are the most useful formalism for parsing
 - Top-down: begin with the start symbol, try to produce the input string to be parsed
 - Bottom up: begin with the input, and try to reduce it to the start symbol
- Both strategies can be cast as search with backtracking
- Backtracking parsers are inefficient: they recompute sub-trees multiple times

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Introduction CSF CS3

Bottom-up parsing as search



S → NP VP
NP → Det N
VP → V NP
VP → V
Det → a
Det → the
N → cat
N → dog
V → bites
N → bites

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Introduction CSF CS3

Dealing with ambiguity

I saw her duck

S → NP VP
NP → Prm N
NP → Prm
VP → V NP
VP → V
VP → V S
N → duck
V → duck
V → saw
Prm → I
Prm → she
Prm → her

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Dealing with ambiguity



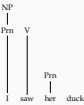
S → NP VP
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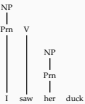
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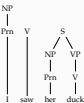
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Introduction CSF CS3

Dealing with ambiguity



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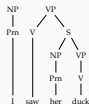
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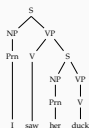
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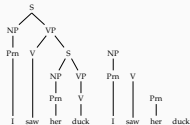
$S \rightarrow NP VP$
 $NP \rightarrow Pm N$
 $NP \rightarrow Pm$
 $VP \rightarrow V NP$
 $VP \rightarrow V$
 $VP \rightarrow V S$
 $N \rightarrow duck$
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 $Pm \rightarrow I$
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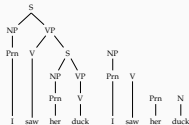
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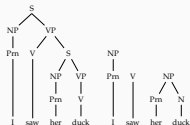
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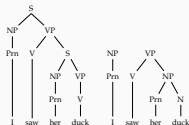
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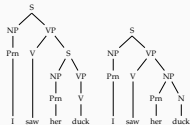
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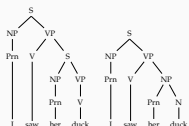
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 $Pm \rightarrow she$
 $Pm \rightarrow her$

How to represent multiple parses

parse forest grammar



$S_{2,4} \rightarrow NP_{0,1} VP_{1,4}$
 $NP_{2,1} \rightarrow Pm_{2,1}$
 $NP_{0,1} \rightarrow I_{0,1}$
 $VP_{1,4} \rightarrow V_{1,2} S_{2,4}$
 $V_{1,2} \rightarrow saw_{1,2}$
 $S_{2,4} \rightarrow Pm_{2,3} V_{3,4}$
 $V_{3,4} \rightarrow duck_{3,4}$
 $VP_{1,4} \rightarrow V_{1,2} NP_{2,4}$
 $NP_{2,4} \rightarrow Pm_{2,3} N_{3,4}$

CKY algorithm

- The CKY (Cocke-Kasami-Younger) parsing algorithm is a dynamic programming algorithm
- It processes the input *bottom up*, and saves the intermediate results on a *chart*
- Time complexity for recognition is $O(n^3)$
- Space complexity is $O(n^2)$
- It requires the CFG to be in *Chomsky normal form* (CNF) (can somewhat be relaxed, but not common)

Chomsky normal form (CNF)

- A CFG is in CNF, if the rewrite rules are in one of the following forms
 - $A \rightarrow BC$
 - $A \rightarrow a$
 where A, B, C are non-terminals and a is a terminal
- Any CFG can be converted to CNF
- Resulting grammar is *not* equivalent to the original grammar:
 - it generates/accepts the same language
 - but the derivations are different

Converting to CNF: example

$S \rightarrow NP VP$
 $S \rightarrow Aux NP VP$
 $NP \rightarrow the N$
 $VP \rightarrow V NP$
 $VP \rightarrow V$
 $N \rightarrow cat$
 $N \rightarrow dog$
 $V \rightarrow bites$
 $N \rightarrow bites$

$S \rightarrow Aux NP VP$
 $S \rightarrow Aux NP VP \Rightarrow S \rightarrow Aux X$
 $NP \rightarrow the N \Rightarrow NP \rightarrow X N$
 $NP \rightarrow the N \Rightarrow X \rightarrow the$
 $VP \rightarrow V$
 $VP \rightarrow V \Rightarrow VP \rightarrow bites$

Converting to CNF

- Eliminate the ϵ rules: if $A \rightarrow \epsilon$ is in the grammar
 - replace any rule $B \rightarrow \alpha A \beta$ with two rules
 - $B \rightarrow \alpha \beta$
 - $B \rightarrow \alpha A' \beta$
 - add $A' \rightarrow \alpha$ for all α (except ϵ) whose LHS is A
 - repeat the process for newly created ϵ rules
 - remove the rules with ϵ on the RHS (except $S \rightarrow \epsilon$)
- Eliminate unit rules: for a rule $A \rightarrow B$
 - Replace the rule with $A \rightarrow a_1 | \dots | a_n$, where a_1, \dots, a_n are all RHS or rule B
 - Remove the rule $A \rightarrow B$
 - Repeat the process until no unit rules remain
- Binarize all the non-binary rules with non-terminal on the RHS: for a rule $A \rightarrow X_1 X_2 \dots X_n$:
 - Replace the rule with $A \rightarrow A_1 X_1 \dots X_{n-1}$ and add $A_1 \rightarrow X_2 X_3$
 - Repeat the process until all new rules are binary

CKY demonstration

an ambiguous example



CKY demonstration

an ambiguous example



CKY demonstration

an ambiguous example



CKY demonstration

an ambiguous example



CKY demonstration

an ambiguous example



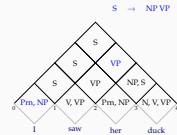
CKY demonstration

an ambiguous example



CKY demonstration

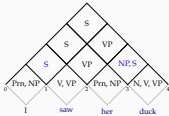
an ambiguous example



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 NP \rightarrow Prm N
 VP \rightarrow V NP
 VP \rightarrow V S
 N \rightarrow duck
 VP \rightarrow duck | saw
 V \rightarrow duck | saw
 Prm \rightarrow I | she | her
 NP \rightarrow I | she | her

CKY demonstration

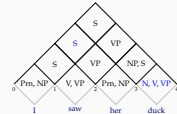
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CKY demonstration

an ambiguous example



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 VP \rightarrow duck | saw
 V \rightarrow duck | saw
 Prm \rightarrow I | she | her
 NP \rightarrow I | she | her

CKY demonstration: the chart

our chart is a 2D array

NP, Prn	S	S	S
	V, VP	VP	VP
		Prn	NP, S
			V, N, NP

Space complexity is $O(n^2)$.

CKY demonstration: the chart

our chart is a 2D array – this is more convenient for programming

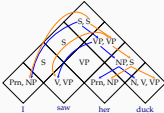
S			
S	VP		
S	VP	NP, S	
NP, Prn	V, VP	Prn, NP	V, N, NP

Space complexity is $O(n^2)$.

Parsing vs. recognition

- We went through a recognition example
- Note that the algorithm is not directional: it takes the complete input
- Recognition accepts or rejects a sentence based on a grammar
- For parsing, we want to know the derivations that yielded a correct parse
- To recover parse trees, we
 - follow the same procedure as recognition
 - add back links to keep track of the derivations

Chart parsing example (CKY parsing)



The chart stores a *parse forest* efficiently.

Summary

- + CKY avoids re-computing the analyses by storing the earlier analyses (of sub-spans) in a table
- It still computes lower level constituents that are not allowed by the grammar
- CKY requires the grammar to be in CNF
- CKY has $O(n^3)$ recognition complexity
- For parsing we need to keep track of backlinks
- CKY can efficiently store all possible parses in a chart
- Enumerating all possible parses have exponential complexity (worst case)
- Suggested reading: **Jurafsky 2009**

Next:

- Top-down chart parsing: Earley algorithm
- Suggested reading:
 - **Jurafsky 2009**
 - **Grune 2008**

Acknowledgments, references, additional reading material