Finite state automata

Data Structures and Algorithms for Computational Linguistics III
(ISCL-BA-07)

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Why study finite-state automata?

- Finite-state automata are efficient models of computation
- There are many applications
  - Electronic circuit design
  - Workflow management
  - Games
  - Pattern matching
  - ...

  But more importantly ;-)
  - Tokenization, stemming
  - Morphological analysis
  - Spell checking
  - Shallow parsing/chunking
  - ...

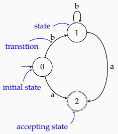
Finite-state automata (FSA)

- A finite-state machine is in one of a finite-number of states in a given time
- The machine changes its state based on its input
- Every regular language is generated/recognized by an FSA
- Every FSA generates/recognizes a regular language
- Two flavors:
  - *Deterministic finite automata* (DFA)
  - *Non-deterministic finite automata* (NFA)

  Note: the NFA is a superset of DFA.

FSA as a graph

- An FSA is a directed graph
- States are represented as nodes
- Transitions are labeled edges
- One of the states is the *initial state*
- Some states are accepting states



DFA: formal definition

Formally, a finite state automaton, M, is a tuple $(\Sigma, Q, q_0, F, \Delta)$ with

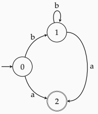
- Σ is the alphabet, a finite set of symbols
- Q a finite set of states
- q_0 is the start state, $q_0 \in Q$
- F is the set of final states, $F \subseteq Q$
- Δ is a function that takes a state and a symbol in the alphabet, and returns another state ($\Delta : Q \times \Sigma \to Q$)

At any state and for any input,
a DFA has a single well-defined action to take.

DFA: formal definition

an example

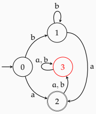
$$\begin{aligned}
\Sigma &= \{a, b\} \\
Q &= \{q_0, q_1, q_2\} \\
q_0 &= q_0 \\
F &= \{q_2\} \\
\Delta &= \{(q_0, a) \to q_2, \quad (q_0, b) \to q_1, \\
&\quad\ \ (q_1, a) \to q_2, \quad (q_1, b) \to q_1\}
\end{aligned}$$



Another note on DFA

error or sink state

- Is this FSA deterministic?
- To make all transitions well-defined, we can add a sink (or error) state
- For brevity, we skip the explicit error state
  - In that case, when we reach a dead end, recognition fails

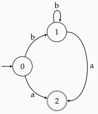


DFA: the transition table

transition table

state	*symbol* a	*symbol* b
→**0**	2	1
1	2	1
***2**	∅	∅

→ marks the start state
* marks the accepting state(s)

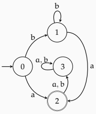


DFA: the transition table

transition table

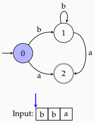
state	*symbol* a	*symbol* b
→**0**	2	1
1	2	1
***2**	3	3
3	3	3

→ marks the start state
* marks the accepting state(s)



DFA recognition

1. Start at q_0
2. Process an input symbol, move accordingly
3. Accept if in a final state at the end of the input



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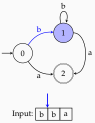

DFA recognition

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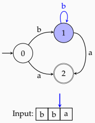

DFA recognition

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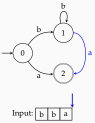

DFA recognition

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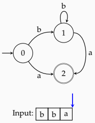
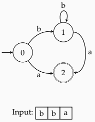


DFA recognition

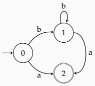
1. Start at q_0
2. Process an input symbol, move accordingly
3. Accept if in a final state at the end of the input

- What is the complexity of the algorithm?
- How about inputs:
  - bbbb
  - aa



A few questions

- What is the language recognized by this FSA?
- Can you draw a simpler DFA for the same language?
- Draw a DFA recognizing strings with even number of 'a's over $\Sigma = \{a, b\}$



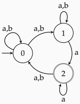
Non-deterministic finite automata

Formal definition

A non-deterministic finite state automaton, M, is a tuple $(\Sigma, Q, q_0, F, \Delta)$ with

- Σ is the alphabet, a finite set of symbols
- Q a finite set of states
- q_0 is the start state, $q_0 \in Q$
- F is the set of final states, $F \subseteq Q$
- Δ is a function from (Q, Σ) to $P(Q)$, power set of Q ($\Delta : Q \times \Sigma \to P(Q)$)

An example NFA



transition table

state	symbol a	symbol b
→0	0,1	0,1
1	1,2	1
*2	0,2	0

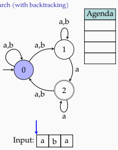
- We have nondeterminism, e.g., if the first input is a, we need to choose between states 0 or 1
- Transition table cells have *sets* of states

Dealing with non-determinism

- Follow one of the links, store alternatives, and *backtrack* on failure
- Follow all options in parallel

NFA recognition

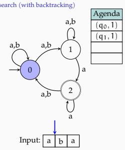
as search (with backtracking)



1. Start at q_0
2. Take the next input, place all possible actions to an *agenda*
3. Get the next action from the agenda, act
4. At the end of input
   - Accept: if in an accepting state
   - Reject: not in accepting state & agenda empty
   - Backtrack: otherwise

NFA recognition

as search (with backtracking)



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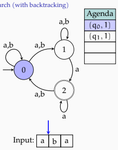
NFA recognition

as search (with backtracking)



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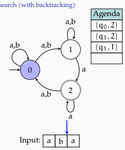
NFA recognition

as search (with backtracking)



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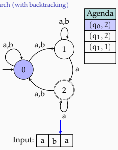
NFA recognition

as search (with backtracking)



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2. Take the next input, place all possible actions to an *agenda*
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4. At the end of input
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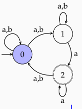
   - Reject: not in accepting state & agenda empty
   - Backtrack: otherwise

NFA recognition
as search (with backtracking)



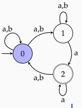
Agenda
$(q_0, 3)$
$(q_1, 3)$
$(q_1, 2)$
$(q_1, 1)$

1. Start at q_0
2. Take the next input, place all possible actions to an *agenda*
3. Get the next action from the agenda, act
4. At the end of input

Accept if in an accepting state
Reject not in accepting state & agenda empty
Backtrack otherwise

Input: | a | b | a |

NFA recognition
as search (with backtracking)



Agenda
$(q_0, 3)$
$(q_1, 3)$
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$(q_1, 1)$

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Accept if in an accepting state
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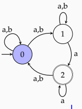
Backtrack otherwise

Input: | a | b | a |

NFA recognition
as search (with backtracking)



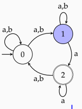
Agenda
$(q_1, 3)$
$(q_1, 2)$
$(q_1, 1)$

1. Start at q_0
2. Take the next input, place all possible actions to an *agenda*
3. Get the next action from the agenda, act
4. At the end of input

Accept if in an accepting state
Reject not in accepting state & agenda empty
Backtrack otherwise

Input: | a | b | a |

NFA recognition
as search (with backtracking)



Agenda
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$(q_1, 1)$

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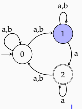
Backtrack otherwise

Input: | a | b | a |

NFA recognition
as search (with backtracking)



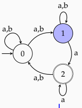
Agenda
$(q_1, 2)$
$(q_1, 1)$

1. Start at q_0
2. Take the next input, place all possible actions to an *agenda*
3. Get the next action from the agenda, act
4. At the end of input

Accept if in an accepting state
Reject not in accepting state & agenda empty
Backtrack otherwise

Input: | a | b | a |

NFA recognition
as search (with backtracking)



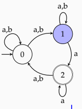
Agenda
$(q_1, 2)$
$(q_1, 1)$

1. Start at q_0
2. Take the next input, place all possible actions to an *agenda*
3. Get the next action from the agenda, act
4. At the end of input

Accept if in an accepting state
Reject not in accepting state & agenda empty
Backtrack otherwise

Input: | a | b | a |

NFA recognition
as search (with backtracking)



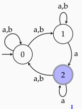
Agenda
$(q_2, 3)$
$(q_1, 3)$
$(q_1, 1)$

1. Start at q_0
2. Take the next input, place all possible actions to an *agenda*
3. Get the next action from the agenda, act
4. At the end of input

Accept if in an accepting state
Reject not in accepting state & agenda empty
Backtrack otherwise

Input: | a | b | a |

NFA recognition
as search (with backtracking)



Agenda
$(q_2, 3)$
$(q_1, 3)$
$(q_1, 1)$

1. Start at q_0
2. Take the next input, place all possible actions to an *agenda*
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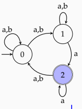
Input: | a | b | a |

NFA recognition
as search (with backtracking)



Agenda
$(q_1, 3)$
$(q_1, 1)$

1. Start at q_0
2. Take the next input, place all possible actions to an *agenda*
3. Get the next action from the agenda, act
4. At the end of input

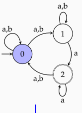
Accept if in an accepting state
Reject not in accepting state & agenda empty
Backtrack otherwise

Input: | a | b | a |

NFA recognition as search
summary

- Worst time complexity is exponential
  - Complexity is worse if we want to enumerate all derivations
- We used a stack as *agenda*, performing a depth-first search
- A queue would result in breadth-first search
- If we have a reasonable heuristic A* search may be an option
- Machine learning methods may also guide finding a fast or the best solution

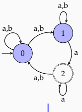
NFA recognition
parallel version



1. Start at q_0
2. Take the next input, mark all possible next states
3. If an accepting state is marked at the end of the input, accept

Input: | a | b | a | b |

NFA recognition
parallel version



1. Start at q_0
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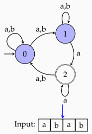
Input: | a | b | a | b |

NFA recognition
parallel version



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NFA recognition
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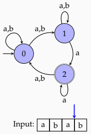

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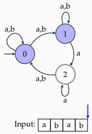

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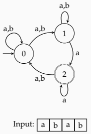
NFA recognition
parallel version

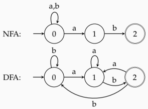


1. Start at q_0
2. Take the next input, mark all possible next states
3. If an accepting state is marked at the end of the input, accept

Note: the process is *deterministic*, and *finite-state*.

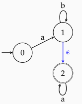
An exercise

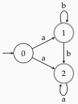
Construct an NFA and a DFA for the language over $\Sigma = \{a, b\}$ where all sentences end with ab.



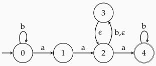
One more complication: ϵ transitions

- An extension of NFA, ϵ-NFA, allows moving without consuming an input symbol, indicated by an ϵ-transition (sometimes called a λ-transition)
- Any ϵ-NFA can be converted to an NFA





ϵ-transitions need attention



- How does the (depth-first) NFA recognition algorithm we described earlier work on this automaton?
- Can we do without ϵ transitions?

NFA–DFA equivalence

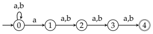
- The language recognized by every NFA is recognized by some DFA
- The set of DFA is a subset of the set of NFA (a DFA is also an NFA)
- The same is true for ϵ-NFA
- All recognize/generate regular languages
- NFA can automatically be converted to the equivalent DFA

Why do we use an NFA then?

- NFA (or ϵ-NFA) are often easier to construct
  - Intuitive for humans (cf. earlier exercise)
  - Some representations are easy to convert to NFA rather than DFA, e.g., regular expressions
- NFA may require less memory (fewer states)

A quick exercise – and a not-so-quick one

1. Construct (draw) an NFA for the language over $\Sigma = \{a, b\}$, such that 4th symbol from the end is an a



2. Construct a DFA for the same language

Summary

- FSA are efficient tools with many applications
- FSA have two flavors: DFA, NFA (or maybe three: ϵ-NFA)
- DFA recognition is linear, recognition with NFA may require exponential time
- Reading suggestion: **hopcroft1979** (and its successive editions), **jurafsky2009**

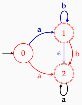
Next:

- FSA determinization, minimization
- Reading suggestion: **hopcroft1979** (and its successive editions), **jurafsky2009**

Acknowledgments, credits, references

ϵ removal

- We start with finding the ϵ-*closure* of all states
  - ϵ-closure(q_0) = {q0}
  - ϵ-closure(q_1) = {q1, q2}
  - ϵ-closure(q_2) = {q2}
- Replace each arc to each state with arc(s) to all states in the ϵ-closure of the state

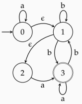


ϵ removal

a(nother) solution with the transition table

transition table

state	a	b	ϵ	ϵ^*
→0	0	∅	1	0,1,2
1	∅	1,3	2	1,2
2	3	∅	∅	2
***3**	3	1	∅	3



ϵ removal

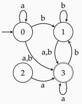
a(nother) solution with the transition table

transition table

state	a	b	ϵ	ϵ^*
→0	0	∅	1	0,1,2
1	∅	1,3	2	1,2
2	3	∅	∅	2
***3**	3	1	∅	3

⇒

state	a	b
→0	0,1,2	1,3
1	1,2,3	1,2,3
2	3	∅
***3**	3	1,2



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