



Chapter 10: Storage and File Structure

Database System Concepts, 6th Ed.

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File Organization, Record Organization and Storage Access



File Organization

- The database is stored as a collection of *files*. Each file is a sequence of *records*. A record is a sequence of fields.
 - One approach:
 - assume record size is fixed
 - each file has records of one particular type only
 - different files are used for different relations
- This case is easiest to implement; will consider variable length records later.



File Organization, Record Organization and Storage Access



Fixed-Length Records

■ Simple approach:

- Store record i starting from byte $n * (i - 1)$, where n is the size of each record.
- Record access is simple but records may cross blocks
 - ▶ Modification: do not allow records to cross block boundaries

■ Deletion of record i : alternatives:

- move records $i + 1, \dots, n$ to $i, \dots, n - 1$
- move record n to i
- do not move records, but link all free records on a *free list*

record 0	10101	Srinivasan	Comp. Sci.	65000
record 1	12121	Wu	Finance	90000
record 2	15151	Mozart	Music	40000
record 3	22222	Einstein	Physics	95000
record 4	32343	El Said	History	60000
record 5	33456	Gold	Physics	87000
record 6	45565	Katz	Comp. Sci.	75000
record 7	58583	Califieri	History	62000
record 8	76543	Singh	Finance	80000
record 9	76766	Crick	Biology	72000
record 10	83821	Brandt	Comp. Sci.	92000
record 11	98345	Kim	Elec. Eng.	80000



Deleting record 3 and compacting

record 0	10101	Srinivasan	Comp. Sci.	65000
record 1	12121	Wu	Finance	90000
record 2	15151	Mozart	Music	40000
record 4	32343	El Said	History	60000
record 5	33456	Gold	Physics	87000
record 6	45565	Katz	Comp. Sci.	75000
record 7	58583	Califieri	History	62000
record 8	76543	Singh	Finance	80000
record 9	76766	Crick	Biology	72000
record 10	83821	Brandt	Comp. Sci.	92000
record 11	98345	Kim	Elec. Eng.	80000



Deleting record 3 and moving last record

record 0
record 1
record 2
record 11
record 4
record 5
record 6
record 7
record 8
record 9
record 10

	10101	Srinivasan	Comp. Sci.	65000
record 1	12121	Wu	Finance	90000
record 2	15151	Mozart	Music	40000
record 11	98345	Kim	Elec. Eng.	80000
record 4	32343	El Said	History	60000
record 5	33456	Gold	Physics	87000
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Free Lists

- Store the address of the first deleted record in the file header.
- Use this first record to store the address of the second deleted record, and so on
- Can think of these stored addresses as **pointers** since they “point” to the location of a record.
- More space efficient representation: reuse space for normal attributes of free records to store pointers. (No pointers stored in in-use records.)

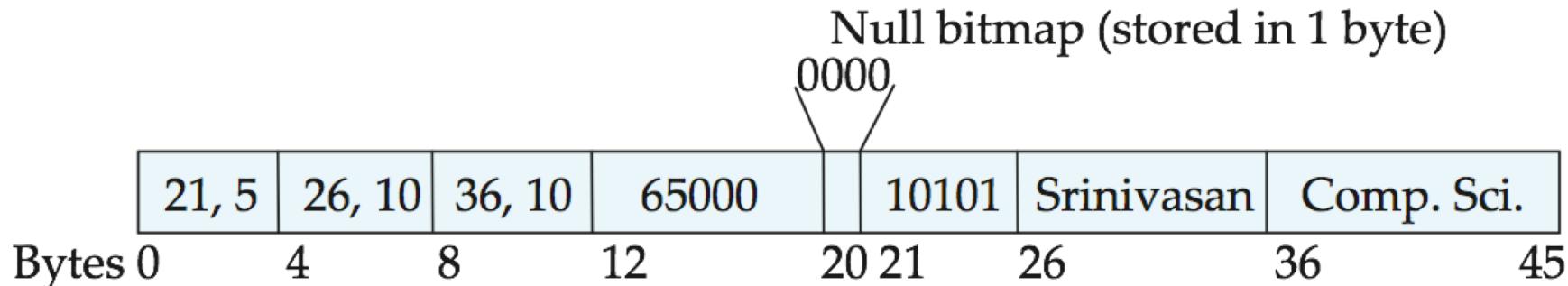
header				
record 0	10101	Srinivasan	Comp. Sci.	65000
record 1				
record 2	15151	Mozart	Music	40000
record 3	22222	Einstein	Physics	95000
record 4				
record 5	33456	Gold	Physics	87000
record 6				
record 7	58583	Califieri	History	62000
record 8	76543	Singh	Finance	80000
record 9	76766	Crick	Biology	72000
record 10	83821	Brandt	Comp. Sci.	92000
record 11	98345	Kim	Elec. Eng.	80000

The diagram illustrates the use of free lists. It shows a table of student records with four columns: ID, Name, Major, and Grade. The fifth row, labeled 'record 4', is highlighted in light blue. Arrows point from the 'Major' and 'Grade' fields of each record in rows 1 through 3 to the corresponding fields of record 4. This indicates that record 4 is being reused to store pointers to the previous records. A final arrow points from the end of record 4's data to a horizontal line, which then points to a ground symbol, signifying the end of the list.



Variable-Length Records

- Variable-length records arise in database systems in several ways:
 - Storage of multiple record types in a file.
 - Record types that allow variable lengths for one or more fields such as strings (**varchar**)
 - Record types that allow repeating fields (used in some older data models).
- Attributes are stored in order
- Variable length attributes represented by fixed size (offset, length), with actual data stored after all fixed length attributes
- Null values represented by null-value bitmap





Organization of Records in Files

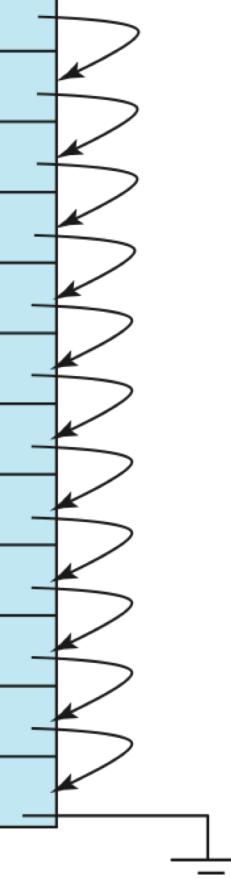
- **Heap** – a record can be placed anywhere in the file where there is space
- **Sequential** – store records in sequential order, based on the value of the search key of each record
- **Hashing** – a hash function computed on some attribute of each record; the result specifies in which block of the file the record should be placed
- Records of each relation may be stored in a separate file. In a **multitable clustering file organization** records of several different relations can be stored in the same file
 - Motivation: store related records on the same block to minimize I/O



Sequential File Organization

- Suitable for applications that require sequential processing of the entire file
- The records in the file are ordered by a search-key

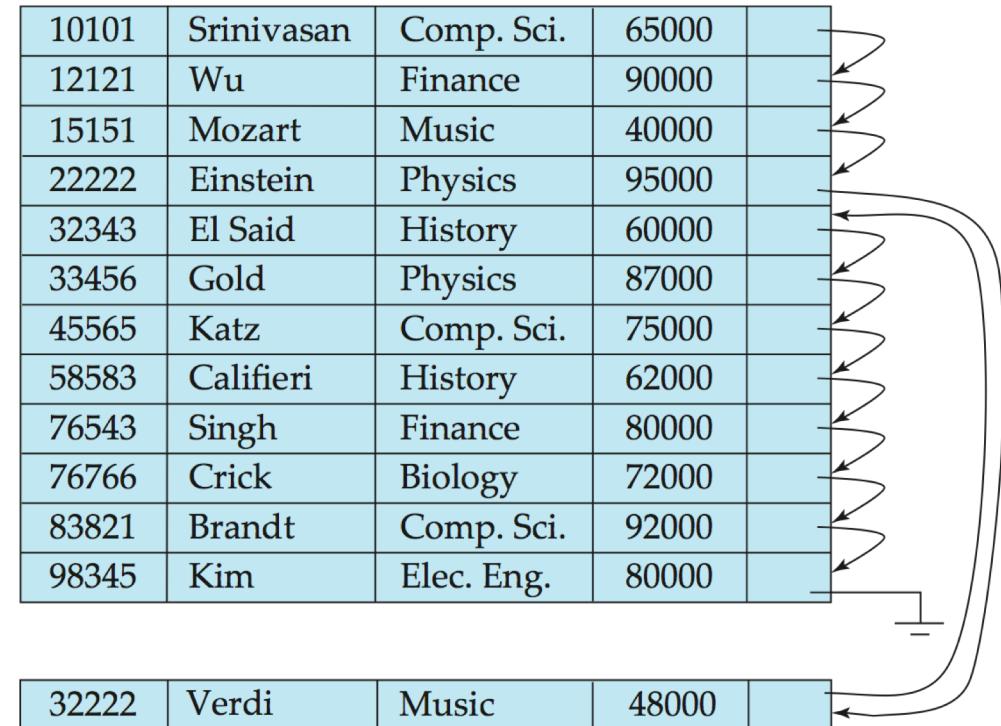
10101	Srinivasan	Comp. Sci.	65000	
12121	Wu	Finance	90000	
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76766	Crick	Biology	72000	
83821	Brandt	Comp. Sci.	92000	
98345	Kim	Elec. Eng.	80000	





Sequential File Organization (Cont.)

- Deletion – use pointer chains
- Insertion – locate the position where the record is to be inserted
 - if there is free space insert there
 - if no free space, insert the record in an **overflow block**
 - In either case, pointer chain must be updated
- Need to reorganize the file from time to time to restore sequential order





Storage Access

- A database file is partitioned into fixed-length storage units called **blocks**. Blocks are units of both storage allocation and data transfer.
- Database system seeks to minimize the number of block transfers between the disk and memory. We can reduce the number of disk accesses by keeping as many blocks as possible in main memory.
- **Buffer** – portion of main memory available to store copies of disk blocks.
- **Buffer manager** – subsystem responsible for allocating buffer space in main memory.