

# Zero-Jitter Task Chains via Algebraic Rings

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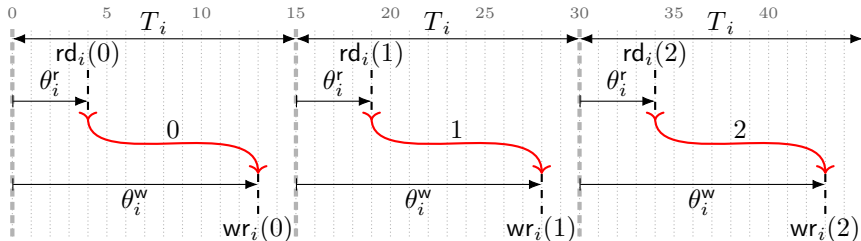
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## Model of a LET task $\tau_i$



- A LET task  $\tau_i$  is composed by *periodic jobs* (curvy arrows)
  - in LET, we only care of the read and write instants

$\theta_i^r$  read phasing of  $\tau_i$ , relative to the period  $T_i$  (aka the offset)

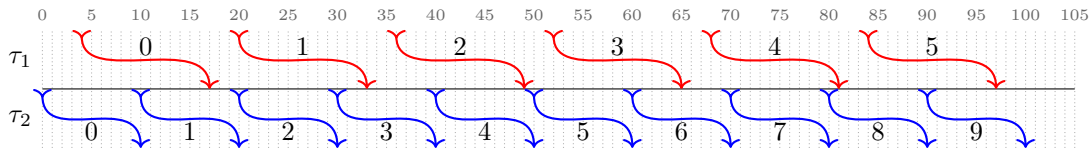
$\theta_i^w$  write phasing of  $\tau_i$ , rel. to period  $T_i$  ( $\theta_i^w - \theta_i^r$  is aka the deadline)

$rd_i(j) = j T_i + \theta_i^r$  read instant of job  $j$  of  $\tau_i$

$wr_i(j) = j T_i + \theta_i^w$  write instant of job  $j$  of  $\tau_i$



## Chain of 2 LET tasks

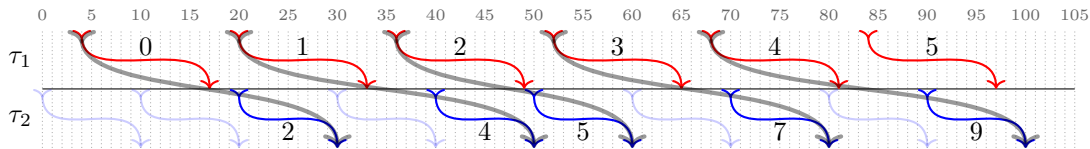


- Chain of LET tasks

- each task reads data written by the previous one
- data is on **shared memory**
  - ★  $\tau_1$  may over-write data before  $\tau_2$  reads (if  $T_1 < T_2$ )
  - ★  $\tau_2$  may read again the same data (if  $T_2 < T_1$ )



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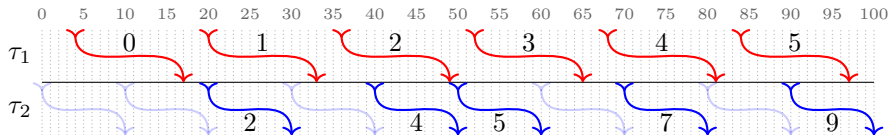
- Chains of jobs have **variable** read-to-write delay

	rd→wr delay		
$(j_1, j_2)$	$rd_1(j_1)$	$wr_2(j_2)$	$wr_2(j_2) - rd_1(j_1)$
(0, 2)	4	30	$26 = 13 + \mathbf{3} + 10$
(1, 4)	20	50	$30 = 13 + \mathbf{7} + 10$
(2, 5)	36	60	$24 = 13 + \mathbf{1} + 10$
(3, 7)	52	80	$28 = 13 + \mathbf{5} + 10$
(4, 9)	68	100	$32 = 13 + \mathbf{9} + 10$
...	...	...	...



## Variability of the input-output delay

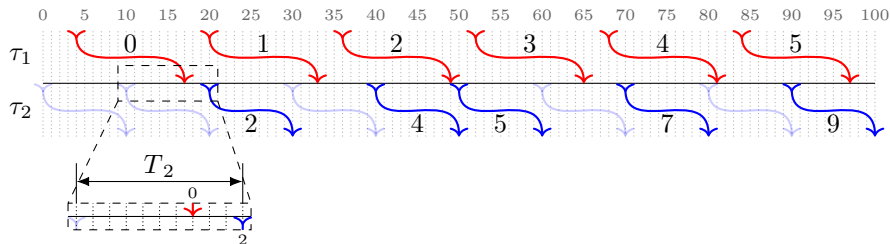
- The source of variability is  $rd_2(j_2) - wr_1(j_1)$ . In the example: 3, 7, 1, 5, 9, 3, 7, ...





## Variability of the input-output delay

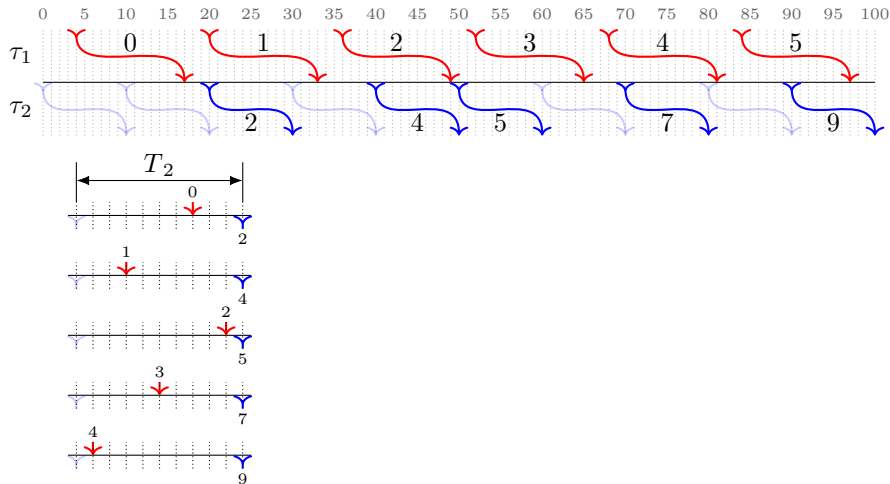
- $rd_2(j_2) - wr_1(j_1) < T_2$  ( $= 10$  in the example), always. Let's zoom in ...





## Variability of the input-output delay

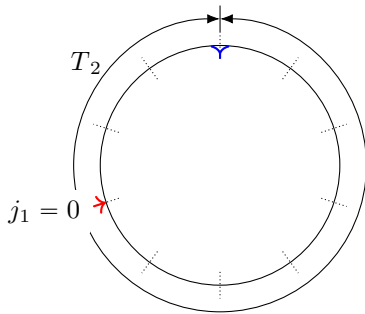
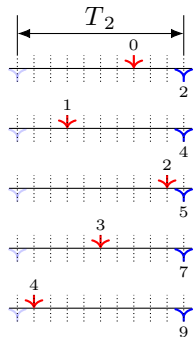
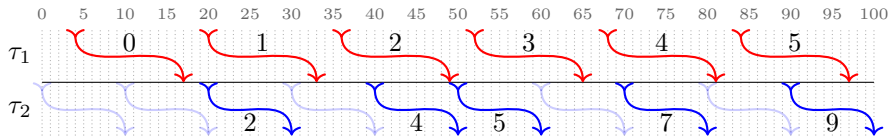
- ... and take  $rd_2(j_2)$  as reference for all  $wr_1(j_1) \rightarrow rd_2(j_2)$  delays.





## Variability of the input-output delay

- A repetitive sequence: let's bend all the  $T_2$ -long segments over the (algebraic) ring  $\mathbb{Z}/T_2\mathbb{Z}$ .

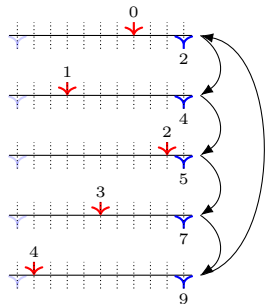
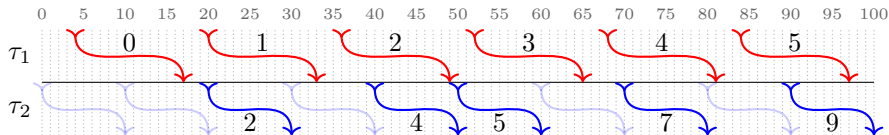




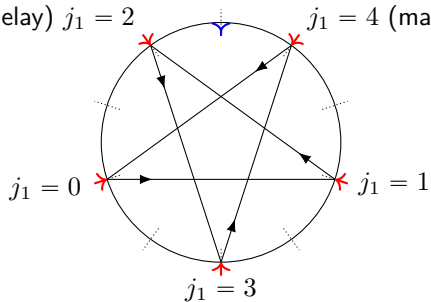


## Variability of the input-output delay

- By inverting over the ring  $\mathbb{Z}/T_2\mathbb{Z}$ , max/min delay is found **without** unrolling the schedule.



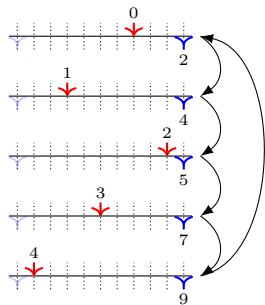
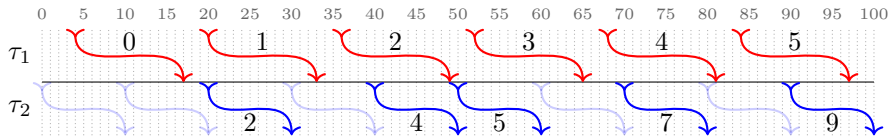
(min delay)  $j_1 = 2$  (max delay)  $j_1 = 4$





## Variability of the input-output delay

- We can **eliminate the jitter** of a 2-tasks chain, by adding a copier task.



(min delay)  $j_1 = 2$   $j_1 = 4$  (max delay)

