# Secure Origin Fallback Mechanism draft-rescorla-callerid-fallback (to be submitted)

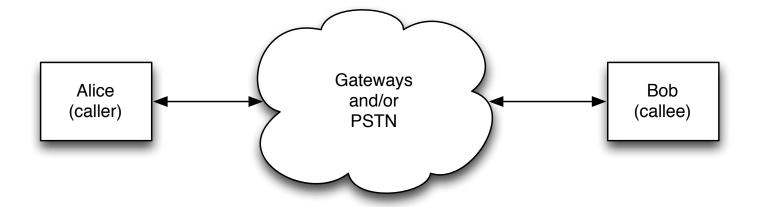
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Not-so-secret workshop May 31, 2013

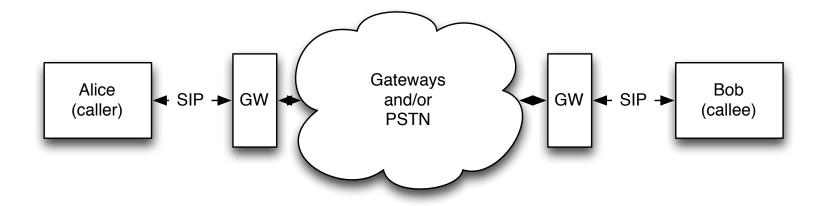
## **Overview**

- RFC 4474 can provide secure origin information
- But SBCs and/or gateways break 4474
  - Change headers
  - Recreate call entirely
- Need to provide source authentication that can survive this
- Basic idea: "Call Detail Service" that validates existence of a call

# **Basic Setting**



# **Alternate Setting**



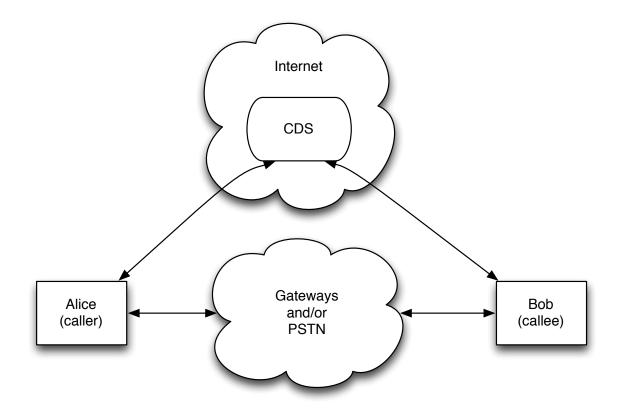
# Assumptions (see Jon's talk)

- Endpoints are programmable
  - User has a smartphone, softphone, etc.
  - User has a dumb phone but is serviced by a programmable gateway
- Very restricted VoIP channel between endpoints
  - Effectively just a PSTN call
  - Caller cannot reliably control caller-id information (CIN field)
- Each caller ID (e.g., E.164) is associated with cryptographic credentials
  - Usable for encryption, authentication, etc.

### **Credentials**

- This assumes that each phone number is associated with credentials
- Requirements
  - Bind an E.164 number to key(s)
  - Suitable for both encryption and authentication
  - Possible to quickly retrieve the credentials for any number
- Example: a public key certificate with the E.164 number as subject

# **System Architecture**



## **Call Flow**

Alice 1.111.1111

Call Detail Service

Bob 1.222.222.2222

Authenticate as 1.111.111.1111

Store E(1.222.222.1.111.111.111)

Call from 1.111.1111

Retrieve CDR from 1.111.1111

E(1.222.222.1.111.111.1111)

[Call from 1.111.111.1111]

\* Note: E(X,Y) means encrypt "Y" for "X"

### **Caller Behavior**

- Look up callee's credentials (may be cached)
- Sign and encrypt CDR for callee\*
- Contact the CDS
  - Authenticate as the caller
  - Store encrypted CDR
- Initiate call to callee

<sup>\*</sup>Special formats needed; must not contain recipient's identity in the clear.

### **CDS** Behavior

- Only store credentials from authorized callers
  - This prevents spamming the CDS
- Provide CDR to any responder
- What if no CDR exists?
  - Generate a random CDR(s)
  - This helps mask the calling rate
    - \* Though not so well for high-rate callers such as call centers

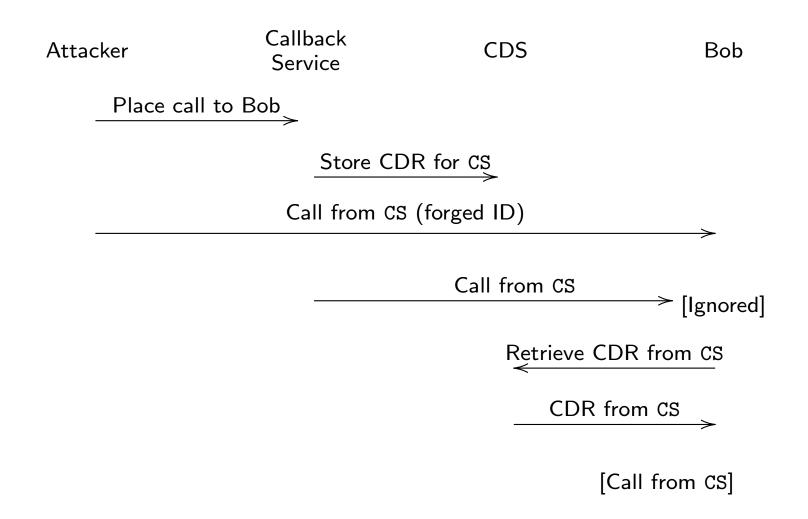
## **Callee Behavior**

- Retrieve encrypted CDR from CDS using claimed caller number
- Decrypt CDR using private key
- Verify CDR signature matches caller's alleged identity
- Check timestamp for relevance (replay prevention)

## What are the security guarantees?

- There exists a relatively recent call from caller to callee
  - Assuming credentials not compromised, etc.
- No guarantee that it is this call
- This defends against robocalling but not MITM attacks

## **Substitution Attack**



## **Privacy Properties I: Off-path Attackers**

- Cannot determine anything about who is calling who
- Cannot determine how many calls a callee is getting
- Limited information about how many calls a caller is generating
  - By polling caller number
  - Can tell if it is more than the minimum number of fake CDRs the CDS generates

## Privacy Properties II: On-path Attackers (to CDS)

- Cannot directly tell who is calling who
  - Assuming communications to CDS are encrypted
- If call volumes are low, can do traffic analysis
  - Alice and Bob both contacted the CDS within a few seconds
- Can measure call volumes for caller and callee
  - Unless they are hidden behind some kind of proxy (Tor, etc.)

# **Privacy Properties III: CDS**

- Anything an on-path attacker can do
- Can directly measure caller's call volume

### Federated CDSs

- Don't need to have one giant CDS
- Each user can select their own CDS
  - Indicated in their credentials?
  - Or delegated from the master CDS?
- This does not need to be exact
  - Callers can fall back (or be bounced) to master CDS during transitions

### What about the Credential Service?

- All callers and callees need to have credentials
- Must be possible for any caller to get callee credentials
  - Quickly
  - Somewhat privately
  - Possible design approaches
    - \* Pre-fetch plus pub-sub
    - \* Caching servers/proxies (a la DNS)
- Caller can provide the callee with his credentials

## How important is credential timeliness?

- Attacker has caller's credentials
  - Can forge calls from attacker
- Attacker has callee's credentials
  - Can poll for calls to callee
  - But probably only for a small number of callers
- Compromise versus reassignment?
  - Can we not reassign during credential validity window?
  - This lets us make validity windows longer
  - Doesn't do anything for compromise
- What is the minimum detection time?

## Why not...?

- Insert a correlation token in the caller's number
  - Could use it to store CDR
  - Assumed not to be possible
- Store under a hash of the caller's number
  - Better privacy but requires sending more traffic from CDS  $\rightarrow$  callee
- Store under a hash of the caller + callee's number
  - May make privacy situation worse
  - Unless hash is shorter than either number
  - Very weak if either side is known
- Is there a practical Private Information Retrieval (PIR) protocol we can use here?

# **Questions?**