PARALLEL PROGRAMMING

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01. INTRO

MOTIVATION TIME, MONEY, COMPLEX PROBLEMS

· MOOR'S LAW LIMITS POWER CONSUMPTION

PARALLELIZATION:

* AUTOMATIC - NOT FEASIBLE TOOLS CAN'T EXTRACT ALL PARALLELISM BY HAND TOOLS ONLY COMPILE

TYPES OF PARALLELISM

FLYNN'S TAXONOMY : INSTRUCTION, DATA SISD, MISD, SIMD, MIMD

- · LEVEL
 - · BIT
 - · INSTRUCTION
 - · TASK PARALLEL TASK GRAPH, PIPELINE, COMMUNICATION

PARALLEL PROGRAMMING

- · NEW PROGRAMMING LANGUAGES : COMPILERS, NEW LANGUAGE
- · EXTENSIONS : EASY DESIGN, COMPILER : ONLY SOME TYPE OF PARALLELISM

02. PARALLEL PATTERNS 1

DEPENDENCIES EXECUTION ASSUMPTION, SEQUENTIAL CONSISTENCY

- TRUE (FLOW) DEPENDENCE RAW $out(S_1)$ Ç $in(S_2)$ † \mathcal{E} S_1 d S_2

- · LOOP-CARRIED DEPENDENCE
 - · CAN PREVENT LOOP ITERATION PARALLELIZATION
 - · LEXICALLY FORWARD / BACKWARD

PARALLEL PATTERN = RECURRING COMBINATION OF TASK THAT SOLVES A SPECIFIC PROBLEM IN PARALLEL ALGORITHM DESIGN

NESTING PATTERN -> ALLOWS OTHER PATTERNS TO BE COMPOSED IN A HIERARCHY

· PATIERN DIAGRAM VISUALIZES PATTERN IDEA

L> TASK BLOCK = LOCATION OF GENERAL CODE IN AN ALGORITHM CAN BE ANOTHER PATTERN

SERIAL CONTROL PATTERNS: SEQUENCE, SELECTION, ITERATION, RECURSION (+ NESTING FOR COMBINATION)

PARALLEL CONTROL PATTERN EXTEND SERIAL, RELAXED ASSUMPTIONS

- FORK-JOIN FLOW FORKS INTO MULTIPLE PARALLEL FLOWS, THAT REJOIN LATER (SYNC # BARRIER)
- MAP ELEMENTAL FUNCTION PERFORMED OVER EVERY ELEMENT OF A COLLECTION
- * STENCIL ELEMENTAL FUNCTION ACCESSES SET OF NEIGHBORS GENERALIZATION OF MAP
- REDUCTION COMBINER FUNCTION (ASSOCIATIVE) COMBINES EVERY ELEMENT IN A COLLECTION
- . SCAN COMPUTES ALL PARTIAL REDUCTIONS OF A COLLECTION
- RECURRENCE COMPLEX VERSION OF MAP (LOOP ITERATION CAN HAVE DEPENDENCIES)

SERIAL DATA MANAGEMENT PATTERNS

- · RANDOM READ / WRITE POINTERS REFER TO MEMORY ADDRESSES. ALIASING CAN CAUSE PROBLEMS
 - STACK ALLOCATION FOR DYNAMIC ALLOCATION, ARBITRARY AMOUNT IN CONSTANT TIME, PRESERVES LOCALITY :
 - . HEAP ALLOCATION WHEN DATA CAN'T BE ALLOCATED IN LIFO. SLOWER, MORE COMPLEX .
 - . OBJECTS LANGUAGE CONSTRUCTS: DATA+ CODE TO MANAGE THEM. PART OF CLASS OBJECTS

PARALLEL DATA MANAGEMENT PATTERNS

- PACK ELIMINATE UNUSED SPACE IN A COLLECTION + UNPACK
- PIPELINE PRODUCER CONSUMER TASK CONNECTION
- GEOMETRIC DECOMPOSITION ARRANGE DATA INTO SUBCOLLECTIONS (DON'T MOVE DATA) · GATHER READ COLLECTION OF DATA, GIVEN COLLECTION OF INDICES
- · SCATTER WRITE OUTPUT AT GIVEN INDEX . POSSIBLE RACE CONDITIONS!
- OTHERS: SUPERSCALAR SEQUENCES, FUTURES, SPECULATIVE SELECTION, WORKPILE

SEARCH, SEGMENTATION, EXPAND, CATEGORY REDUCTION

- MAP "FOR EACH" LOOP, EACH ITERATION INDEPENDENT ~> CAN BE APPLIED WITHOUT KNOWLEDGE OF NEIGHBORS · SERIAL VS PARALLEL
 - - INDEPENDENCE + NO SHARED STATE OTHERWISE RACES, UNDEFINED BEHAVIOR >
 - · OPTIMIZATIONS
 - · SEQUENCES OF MAPS NO WRITE ALL INTERMEDIATE STAGES TO MEM!
 - CODE FUSION : ARITHMETIC INTENSITY : MEMORY USAGE
 - · CACHE FUSION BREAK WORK INTO BLOCKS
 - · STENCIL MAP+NEIGHBORS
 - · WORKPILE WORK ITEMS ADDED TO MAP WHILE IN PROGRESS
 - * DIVIDE AND CONQUER BASE CASE SOLVED SERIALLY
 - EXAMPLE:SCALED VECTOR ADDITION (SAXPY) y = ax+y

03. PARALLEL PATTERNS 2

COLLECTIVES DEAL WITH A COLLECTION AS A WHOLE (REDUCE, SCAN, PARTITION, SCATIER, GATHER)

REDUCE COMBINE COLLECTION OF EL INTO ONE SUMMARY VALUE

- . COMBINER FUNCTION ASSOCIATIVE, ELEMENTS PAIRWISE
- · VECTORIZATION, TILING
- . PRECISION WITH FLOAT "
- EXAMPLE : DOT PRODUCT MAP (*) + REDUCE (+)
- . EXAMPLE: MERGE SORT MAP-REDUCE
 - RIGHT BIASED " O(m2), INSERTION SORT
 - TREE-SHAPED SORT O(Mlogm) ; NOT A LOT TO PARALLELIZE

SCAN GENERATES NEW SEQUENCE WITH PARTIAL REDUCTIONS

- INCLUSIVE / EXCLUSIVE
- · EXAMPLE: RADIX SORT, QUICKSORT
- PARALLELIZATION : UP SWEEP + DOWN SWEEP, 3 PHASE SCAN WITH TILING

04. PARALLEL PATTERNS 3

DATA MOVEMENT LIMITS PERFORMANCE MORE THAN COMPUTATION

- · ACROSS MEMORY LOCALITY, DATA ORGANIZATION
- · ACROSS NETWORKS NUMBER OF MESSAGES

GATHER READ INPUT COLLECTION AT INPUT INDICES -> WRITE OUTPUT COLLECTION

- . "MAP + RANDOM READ"
- * OUTPUT EL. TYPE = SOURCE EL. TYPE
- * OUTPUT SHAPE = INDEX COLLECTION SHAPE

- SPECIAL CASES
 - SHIFT + DEFAULT VALUE, DUPLICATE, ROTATE
 - - 21P example: REAL+IMAGINARY PARTS -> COMPLEX NUMBERS ~> COMBINES IN PAIRS

• UNZIP SCATTER

- . "MAP + RANDOM WRITES"
- . COLLISIONS PARALLEL WRITES TO THE SAME LOCATION
 - L> UNDEFINED RESULT => NEED RULES
- COLLISION RESOLUTION
 - · ATOMIC SCATTER NON-DETERMINISTIC
 - PERMUTATION SCATTER ILLEGAL COLLISIONS -> OUTPUT PERMUTATION OF INPUT (=> GATHER)

 - MERGE SCATTER ASSOCIATIVE AND COMMUTATIVE OPERATIONS TO MERGE

· PRIORITY SCATTER PRIORITY BASED ON ITS POSITION

. BOOL ARRAY, ~ CUMSUM, WRITE OUT PUT WITH OFFSET!

UNPACK GENERALIZATION OF PACK

- · SPLIT ELEMENTS MOVED LEFT/RIGHT, NO INFO LOST
- UNSPLIT
- · BIN SPLIT WITH MORE CATEGORIES MAP + PACK
- · EXPAND EACH EL CAN PRODUCE ANY NUMBER OF EL + RES FUSED TOGETHER

PACK TO ELIMINATE UNUSED ELEMENTS FROM MEM + MOVED TO CONTIGUOUS MEM

- PARALLELIZING ALGORITHMS DIVIDING DATA INTO SECTION
 - . DIVIDE AND CONQUER
- . FORK-JOIN

- · GEOMETRIC DECOMPOSITION
- · PARTITIONING NON OVERLAPPING, EQUAL-SIZED REGIONS
- SEGMENTATION NON OVERLAPPING, NON-UNIFORM REGIONS

Ans ARRAY OF STRUCTURES

SO A STRUCTURE OF ARRAYS BETTER FOR VECTORIZATION

05. PTHREADS, OPENMP 1

- 1 DESIGN PARALLEL ALGORITHM UNDERSTAND DATA DEPENDENCIES
- 2. DESIGN PARALLEL PROGRAM ANALIZE TARGET ARCHITECTURE, CHOOSE LANGUAGE, ANALYZE COMMUNICATIONS

EVALUATION - SEQUENTIAL ALGORITHMS - TIME, MEMORY COMPLEXITY PARALLEL ALGORITHMS -> TIME, RESOURCE COMPLEXITY

QUALITATIVE EVALUATION: DAG, WORK, SPAN, PARALLELISM

PP MODELS SHARED MEMORY, THREADS, MESSAGE PASSING/DISTRIBUTED MEM, DATA PARALLEL

- VERILOG/VHDL :: COMMUNICATION+MEM CONTROL, NO OVERHEAD :: SPECIFIC HW, DIFFICULT TO LEARN, NO DIFF ARCH
- MPI : DIFFERENT ARCH, SYNCH + DATA COMM EXPLICITLY MANAGED 💢 COMM OVERHEAD, MORE DIFFICULT
- 🍍 PTHREAD 🖰 DIFFERENT ARCH, EXPLICIT PARALLELISM, APP CONTROL 💢 TASK MANAGEMENT CH, NOT SCALABLE, LOW-LEVEL API
- OPENMP :: EASY LEARN, SCALABLE 💢 ONLY SH-MEM HOMOGENEOUS SYS, SMALL TASK INTERACTION
- OPENCL : TARGET INDEPENDENT, HIDES ARCH DETAILS : DIFFICULT TO PROGRAM+OBTAIN BEST PERFORMANCE
- APACHE SPARK : API FOR DIFFERENT LANG, NO EXPLICIT PARALL/COMM : ONLY FOR BIG DATA APP, NO GPU FULL SUPPORT MIX OPENMP+ CUDA, MPI+OPENMP, OPENCL+VERILOG/VHDL

CUDA : GPU, EASY KERNELWRITING, OPTIMIZED LIBRARIES : ONLY NVIDIA GPU, DIFFICULT MASSIVE PAR. + KERNEL OPT

PROCESS INFO ABOUT RESOURCES, EXEC STATE

THREAD INDEPENDENT STREAM OF INSTR WITHIN A PROCESS

· CAN EXECUTE AT THE SAME TIME

- . HAS LOCAL RESOURCES, CAN ACCESS SHARED PROCESS RESOURCES
- . IMPLICIT COMMUNICATION (SHARED VAR) -> EXPLICIT SYNCHRONIZATION TO AVOID DATA RACE

THREADED PROGRAM

- IF ORGANIZABLE IN DISCRETE CONCURRENT TASKS, THREAD SAFE
- · MODELS: MANAGER/WORKERS, PIPELINE
- · STANDARD IMPLEMENTATIONS: PTHREAD, OPENMP

PTHREADS

THREAD MANAGEMENT CREATE, JOIN / DETATCH, BARRIER INIT / WAIT, CANCEL / EXIT

- · SYNCHRONIZATION MUTEX, CONDITION VARIABLES
- examples CALCULATION OF T (MONTECARLO), MULTIRATE BAND-PASS FILTER, PRODUCER-CONSUMER OPENMP API FOR MULTI-THREADED SHARED-MEM PROGRAMMING
 - · BASED ON FORK-JOIN MASTER + CONCURRENT SLAVES
 - · LANG EXTENSIONS: PRAGMA PREPROCESSOR DIRECTIVES, IDENTIFY TASKS. IMPLIED BARRIER AT THE END
 - PARALLEL [NUM THREADS, IF]
 - · WORK SHARING: FOR, SECTIONS, SINGLE/MASTER, TASK

OG. OPENMP 2

OPENMP LANGUAGE EXTENSIONS

- · SYNCHRONIZATION: CRITICAL [NAME], BARRIER, ATOMIC
- DAT A ENVIRONMENT : DATA SCOPE ATTRIBUTE CLAUSES (PRIVATE, SHARED, DEFAULT, REDUCTION)
 - · SEQUENTIAL MEM CONSISTENCY LOAD/STORE IN ORIGINAL SEQUENTIAL ORDER
 - RELAXED MEM CONSISTENCY EACH THREAD TMP VIEW. FLUSH -> TO HAVE SAME GLOBAL VIEW
- · RUNTIME FUNCTIONS GET NUM THREADS, GET THREAD NUM, SET NUM THREADS, GET WITIME, GET WITICK
- · examples: CALCULATION OF TI (INTEGRAL)

NESTED PARALLELISM PARALLEL, WORKSHARING ~> BUT ID RESTARTS FROM O!
THREAD CANCELLATION CANCEL + CANCELLATION POINT

O7. OPENMP 3

OPEN MP TASK = BLOCK OF CODE IN A PARALLEL REGION, THAT CAN BE EXECUTED SIMULTANEOUSLY WITH OTHER TASKS IN THE SAME REGION

- PIPELINE SECTIONS & COMMUNICATION AND WAD IMBALANCE PROBLEM
 TASKS & CAN EXECUTE WHENEVER INPUT IS READY
- · TASKWAIT /TASKGROUP
- DEPEND IN OUT
 PRIORITY IMPORTANT TASKS TO BE EXECUTED MORE FREQUENTLY
- TASKLOOP

SIMD VECTORIZATION INSTRUCTIONS ON BLOCK OF DATA (VECTORS)

- . COMPILER ISSUES LOOP DEPENDENCIES, POINTER ALIASES, DATA ALIGNMENT, LOOP BOUNDS
- . VECTOR REGISTERS LOOP CHUNK SHOULD FIT IT
- · simulen PREFERRED SIZE (PERFORMANCE)

HETEROGENEOUS ARCHITECTURES

- # pragma omp target [map(...)]
 - · CODE TO ACCELERATOR, IF PRESENT
 - · HOST THREAD BLOCKED
 - · map copies variables in target memory (17's shared) + to, from, togrom
- BATCH PROCESSING/COOPERATIVE PATIERNS

08. MPI 1

SHARED MEM ARCHITECTURES : HARD TO BUILD : EASY TO PROGRAM (OPENMP, C++ LIBRARY)

DISTRIBUTED MEM ARCHITECTURE REQUIRES EXPLICIT COMMUNICATION : SYNCHRONIZATION

MPI (MESSAGE PASSING INTERFACE) STANDARD THAT DEFINES INTERFACES FOR DISTRIBUTED MEM COMMUNICATIONS

- · OUTCOME : PDF DOCUMENT
- · IMPLEMENTATIONS OPENMPI, MPICH
 - · libmpi. SO IMPLEMENTATION LIBRARY
 - · mpicc COMPILER
 - · mpiexec LAUNCHER
 - · mpirum FORK+EXEC CODE
- . INIT + THREAD USAGE
- · FINALIZE
- . COMMUNICATOR CONTEXT

POINT TO POINT COMMUNICATION

- · SEND, RECV (STATUS), PROBE
- · WAIT, TEST, CANCEL

09. MPI 2

COLLECTIVE COMMUNICATIONS

- · COMMUNICATOR HANDLING DUP, SPLIT, FREE
- · SYNCHRONIZATION BARRIER
- DAT A TRANSFER BROADCAST, GATHER, SCATTER REDUCTIONS REDUCE (MAX, SUM, LAND, MAXLOC)

10 HALIDE

POWER = OP / SECONDS · TOULES /OP

GENERAL-PURPOSE PROCESSORS NOT ENERGY EFFICIENT

HETEROGENEOUS PROCESSING

- GPU
- . DSP DIGITAL SIGNAL PROCESSORS
- . ASIC DOMAIN-SPECIFIC HW EXTREME SPECIALIZATION
- · FPGA

CHALLENGES ___ SW DESIGN ALGO THAT CAN BE DECOMPOSED AND MAP WELL ON HW

HW CHOOSE RESOURCES, CHIP AREA FOR EACH FUNCTIONALITY

ENERGY-EFFICIENT COMPUTING REDUCE COMPUTATION, DATA TRANSFER , CHOOSE BEST COMPUTE UNIT

DSL PRODUCTIVITY + PERFORMANCE, NO GENERALITY

IMAGE PROCESSING DEMANDS OPTIMIZATION -> FASTER ALGO, HW, MORE EFFICIENT USE OF HW

- · C++ WITH MULTITHREADING
- · CUDA/OPENCL
- · OPTIMIZED LIBRARIES

HALIDE DSL EMBEDDED IN C++ + SPECIALIZED COMPILER

- · DECOUPLE ALGO FROM SCHEDULE
- X ALL COMPUTATIONS OVER REGULAR GRIDS, MANUALLY SPECIFY SCHEDULE (AUTOSCHEDULER CAN SUGGEST)
- · CAN OPTIMIZE WITH
 - ORDER OF VALUES COMPUTED SERIAL, VECTORIZE, TILING
 - . WHEN COMPUTE RESULTS eg: COMPUTE ALL/SUBSET OF PRODUCER, THEN PASS TO CONSUMER
- · SCHEDULING ACROSS STAGES DRIVEN BY CONSUMER
- · examples: BLUR, BRIGHTEN

11. PARALLEL PATTERNS APPLIED

PP MODELS OFFER API FOR PATTERNS

- · MAP # pragma omp parallel for
- REDUCTION # pragma omp for reduction (+: van)
- · WORKPILE example: TREE SEARCH
- SCAN + 3 PHASE TILING APPROACH
- + SORTING