

# Introduction to complex networks

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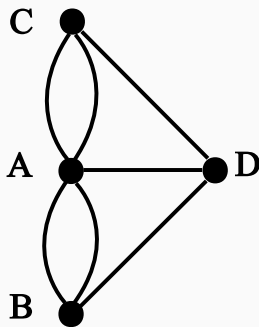
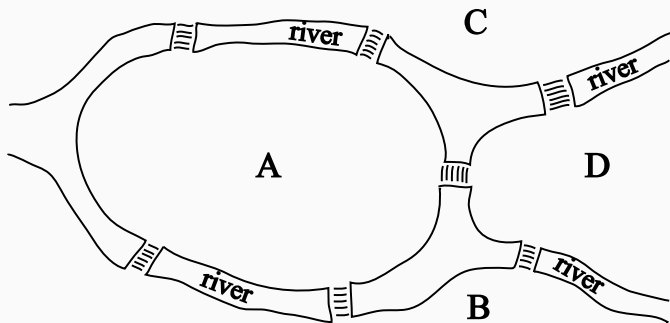
# Definitions

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# Graph theory

1736 Euler writes the first paper

1959 Erdős and Rényi introduce probabilistic methods

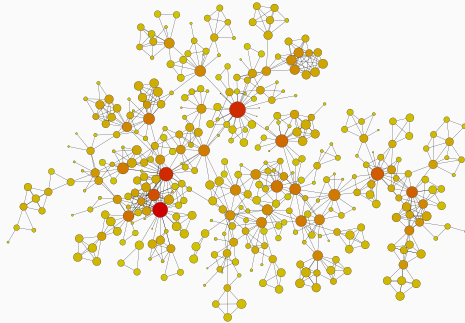


From Junker and Schreiber (2008)

# Network theory

1970s Sociologists ('social network analysis')

2000s Physicists and Computer Scientists



# Graphs and their representation

A graph  $\mathcal{G} = (\mathcal{V}, \mathcal{E})$  consists of:

1. a set of **vertices**  $\mathcal{V} = \{v_1, \dots, v_n\}$
2. a binary relation  $\mathcal{E} \subseteq \mathcal{V}^2$  that represents **edges**

# Graphs and their representation

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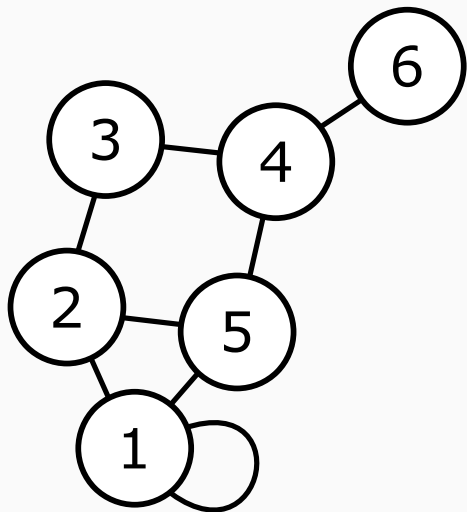
1. a set of **vertices**  $\mathcal{V} = \{v_1, \dots, v_n\}$
2. a binary relation  $\mathcal{E} \subseteq \mathcal{V}^2$  that represents **edges**

## Alternative representation of $\mathcal{E}$

In terms of the **adjacency matrix**  $\mathbf{A} \in \{0, 1\}^{n \times n}$ , with elements

$$a_{ij} = \begin{cases} 1 & \text{if } (v_i, v_j) \in \mathcal{E} \\ 0 & \text{otherwise} \end{cases}$$

## Example



$$\mathcal{V} = \{v_1, \dots, v_6\}$$

$$\mathcal{E} = \{(v_1, v_1), (v_1, v_2), \dots\}$$

$$\mathbf{A} = \begin{pmatrix} 1 & 1 & 0 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \end{pmatrix}$$



# Weights

- Assigned to elements of  $\mathcal{E}$
- Represent the **strength** of the relationship

## Examples

- Signed edges:  $\mathbf{A} \in \{-1, 0, 1\}^{n \times n}$
- Weights on the unit interval:  $\mathbf{A} \in [0, 1]^{n \times n}$
- Weights on the positive real line:  $\mathbf{A} \in [0, \infty)^{n \times n}$

# Properties of $\mathcal{E}$

## Irreflexivity

$$\forall v_i \in \mathcal{V}. (v_i, v_i) \notin \mathcal{E}$$

→ No 'self-loops' allowed

## Symmetry

$$\forall v_i, v_j \in \mathcal{V}. (v_i, v_j) \in \mathcal{E} \Rightarrow (v_j, v_i) \in \mathcal{E}$$

→ Direction of the relationship does not matter ( $\mathcal{G}$  is **undirected**)

→  $\mathbf{A} = \mathbf{A}^\top$  is also symmetric

# Properties of vertices

## Neighbourhood

$$\mathcal{N}(v_i) = \{v_j : (v_i, v_j) \in \mathcal{E}\} \subseteq \mathcal{V}$$

→ **Set** of vertices to which  $v_i$  is connected by edges

## Degree

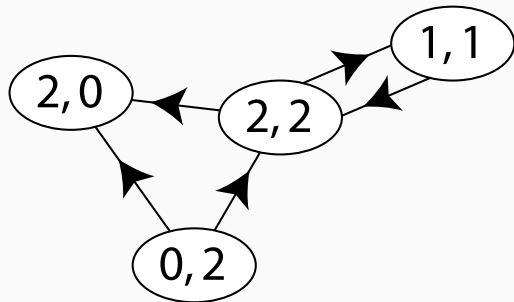
$$k_i = k(v_i) = |\mathcal{N}(v_i)|$$

→ **Number** of vertices to which  $v_i$  is connected by edges

# In- and out-degree

$$k_i^{\text{in}} = |\{(v_j, v_i) \mid (v_j, v_i) \in \mathcal{E}\}|$$

$$k_i^{\text{out}} = |\{(v_i, v_j) \mid (v_i, v_j) \in \mathcal{E}\}|$$



# Paths

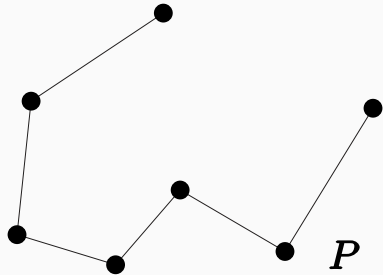
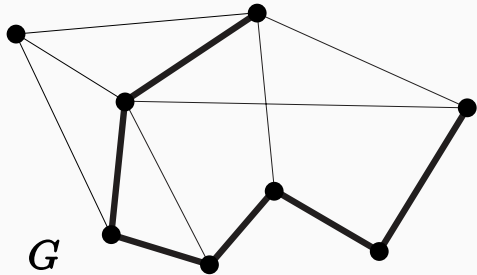
A **path** on  $\mathcal{G} = (\mathcal{V}, \mathcal{E})$  is a graph  $\mathcal{P} = (\mathcal{V}_{\mathcal{P}}, \mathcal{E}_{\mathcal{P}})$  with:

1.  $\mathcal{V}_{\mathcal{P}} = \{v_{(0)}, \dots, v_{(l)}\} \subseteq \mathcal{V}$
2.  $\mathcal{E}_{\mathcal{P}} = \{(v_{(0)}, v_{(1)}), \dots, (v_{(l-1)}, v_{(l)})\} \subseteq \mathcal{E}$

**Ends** Vertices  $v_{(0)}$  and  $v_{(l)}$

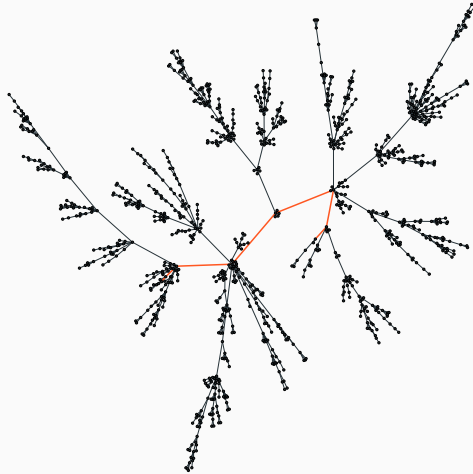
**Length** Number of edges  $|\mathcal{E}_{\mathcal{P}}| = l$

# Paths



From Diestel (2010)

# Shortest paths



# Local properties

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Which are the most important  
(or 'central') vertices?

## **Degree centrality**

How many connections does a vertex have to other vertices?

Which are the most important  
(or 'central') vertices?

## **Closeness centrality**

How 'far' (in terms of shortest paths) is a vertex with respect to all other vertices?

Which are the most important  
(or 'central') vertices?

## **Betweenness centrality**

How many times does a vertex act as a 'bridge' along the shortest paths between all pairs of vertices?

Which are the most important  
(or ‘central’) vertices?

## **Eigenvector centrality (PageRank)**

How many times do we stumble upon a vertex while randomly walking on the graph?

# Local clustering coefficient

$$C_i = C(v_i) = \frac{|\{(v_j, v_k) : \{(v_i, v_j), (v_i, v_k), (v_j, v_k)\} \subseteq \mathcal{E}\}|}{k_i(k_i - 1)/2} \in [0, 1]$$

- Probability that a pair of neighbours of  $v_i$  are connected
  - Empirically:  $k_i \nearrow, C_i \searrow$
- Can be used to identify ‘structural holes’

# Large-scale structure

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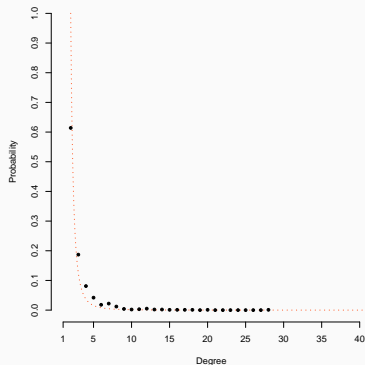
## ‘Small-world’ effect

*‘...Milgram’s letter-passing experiment in the 1960s, in which people were asked to get a letter from an initial holder to a distant target person by passing it from acquaintance to acquaintance through the social network. The letters that made it to the target did so in a remarkably small number of steps, around six on average.’*

*From Newman (2010)*

# Degree distribution

$\Pr[K = k] = \Pr[\text{'a randomly selected vertex } v_i \in \mathcal{V} \text{ has degree } k \geq 0']$





# Power laws

$$\ln y = -\alpha \ln x + c$$

$$y = Cx^{-\alpha}, \quad C = e^c$$

- Sizes of city populations and wars
- Frequency of use of words in human languages
- Occurrence of personal names in most cultures
- Number of papers scientists write
- Number of hits on Web pages
- Sales of almost every branded commodity
- Number of species in biological taxa

# Scale-free networks

- $\Pr[K = k] \propto k^{-\alpha}$  (at least asymptotically), usually  $2 \leq \alpha \leq 3$
- **Self-similar** (scale invariance)
- Presence of **hubs** (Milgram's 'sociometric superstars')
  - Fault tolerance
  - 'Small-world' effect

# Scale-free networks

## Examples

- Social and collaboration networks
- Many kinds of communication networks (e.g. the Internet)
- **Many biological networks**

# Models of network formation

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# What is complexity?

## **Complex**

Many interacting parts, not necessarily 'difficult'

## **Complicated**

Difficult, not necessarily 'complex'

# What is complexity?



*Via Wikimedia Commons*

## Spontaneous emergence of global structure out of local interactions

- Robust to damage and perturbations
  - Not controlled by (internal or external) agents
  - Collective, distributed process
- Outcome is not arbitrary, but 'prefers' certain situations  
→ Natural selection

# Self-organisation

## Examples

- Protein folding
- Pattern formation and morphogenesis
- Social structures and herd behaviour



Topology is crucial to understanding  
stochastic processes on networks

# Models of network formation

Random



Prefential attachment



# Erdős-Rényi model

## Algorithm

Construct a graph with  $n$  vertices and include each edge with probability  $p$  (independently from every other edge)

# Erdős-Rényi model

## Algorithm

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## Degree distribution

$$\Pr[K = k] = \binom{n-1}{k} p^k (1-p)^{n-1-k}$$

# Barabási-Albert model

## Algorithm

- Start with an initial connected graph with  $n$  vertices
- Add a new vertex and connect it to  $n^* \leq n$  existing vertices with probability proportional to their degree  
→ **Preferential attachment**

# Barabási-Albert model

## Algorithm

- Start with an initial connected graph with  $n$  vertices
- Add a new vertex and connect it to  $n^* \leq n$  existing vertices with probability proportional to their degree  
→ **Preferential attachment**

## Degree distribution

$$\Pr[K = k] \propto k^{-3}$$

# What about ‘-omics’?

- Very few longitudinal datasets
  - Difficult to understand evolution
- Few cross-platform studies
  - Can understand interactions between complexity levels
- Many single-platform studies
  - Can characterise “co-” networks