

### **GPU** architecture

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- **GPU** architecture
  - Kepler
  - Maxwell
- Performance limiters
- Memory
  - Global memory
  - Shared memory
  - Read-only and constant cache
  - Registers and local memory
- Grid configuration
- Arithmetics
  - Branching
  - **Intrinsics**
- Asyncrony
- Case study

## Kepler SMX



## **Kepler GK110**

- 7.1 billion transistors
- \$\infty\$ 15 SMX
- 6 GB Memory
- S > 1 TFLOP FP64
- \$\infty\$ 1.5 MB L2 cache
- 384-bit GDDR5
- PCI Express 3

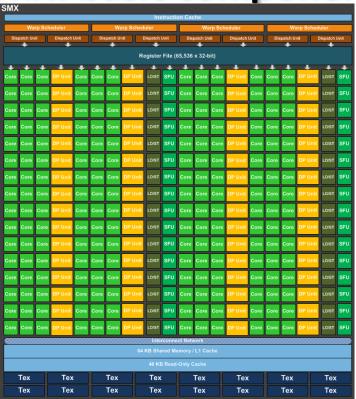




#### SMX:

- 192 CUDA cores
- 64 DP Unit
- 32 LD/ST Unit
- 32 SFU
- 64 KB L1 cache/ shared memory
- 16 texture processors
- 48 KB read-only cache
- 65536 32-bit register
- 4 warp schedulers
- 8 dispatch units
- Instruction cache

### **SMX:** basic components



## Global memory

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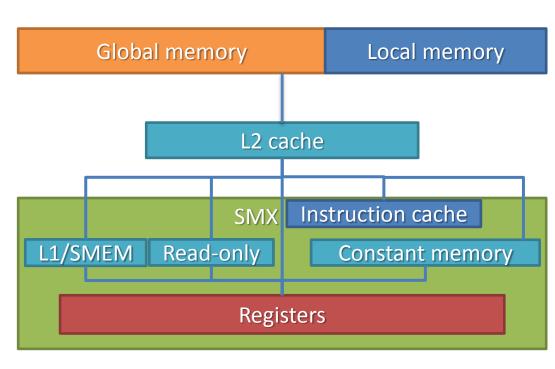
## Memory

#### What can I control:

- Global memory
- Shared memory

#### Partially:

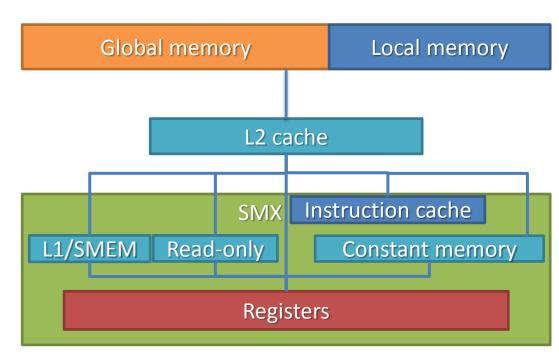
- Local memory
- Registers
- L1 cache
- Read-only cache





## **Global memory**

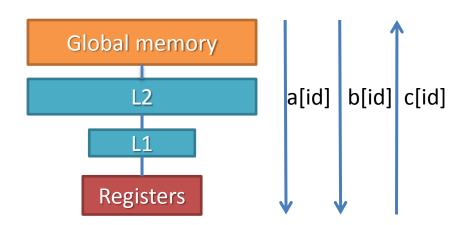
- S GDDR5
- Common for the entire device
- Accessible for host using cudaMemcpy() and UVA
- Accessible for threads via pointer
- Probable read/write conflicts
- Migh latency (200-400 cycles)





### Global memory requests

```
__global__ void vecAdd(double *a, double *b,
double *c, int n)
    int id = blockIdx.x*blockDim.x+threadIdx.x;
    if (id < n)
       c[id] = a[id] + b[id];
```



- Requests go trough L1 and L2 caches
  - L1 cache line 128 bytes
  - L2 cache line 32 bytes
- It is necessary to group requests high performance
- L1 cache switching on and off:
  - -Xptxas -dlcm=ca (on by default)
  - -Xptxas -dlcm=cg



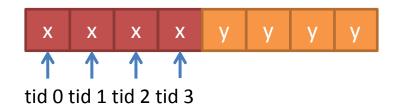
### **AoS vs SoA**

```
struct A{
    float x;
    float y;
};
struct A myArray[n]
```

```
x y x y x y x y

tid 0 tid 1 tid 2 tid 3
```

```
struct A{
    float x [n];
    float y [n];
};
struct A myArray
```



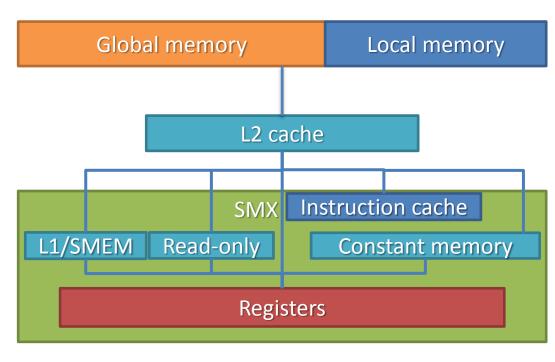
## Shared memory

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## **Shared memory**

- Allocated per block
- Accessible for all the threads of a block
- Possible read/write conflicts
- Solution
  Based on the same chip with L1 cache, combinations:
  - 16/48 KB
  - 32/32 KB
  - 48/16 KB
- Latency 20-60 cycles





## **Shared memory**

```
__global___ void staticReverse(int
*d, int n){
    __shared__ int s[64];
    int t = threadIdx.x;
    int tr = n-t-1;
    s[t] = d[t];
    __syncthreads();
    d[t] = s[tr];
}
```

- Static or dynamic allocation
- Secondary Barrier syncronization using \_\_syncthreads()

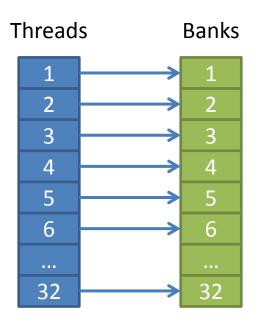


## Memory banks

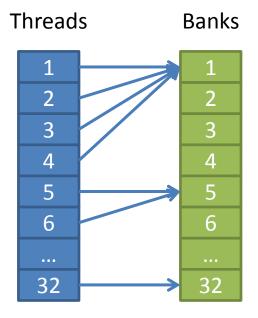
- 32 banks wide of 8 byte
  - Bank conflicts.
    - ✓ Program slows data requests are serialized
- Two configurations:
  - 4-byte access (by default)
  - 8-byte access
  - Use functions cudaDeviceSetSharedMemConfig() and cudaFuncSetSharedMemConfig() with the following arguments:
    - ✓ cudaSharedMemBankSizeFourByte
    - ✓ cudaSharedMemBankSizeEightByte



## **Bank Conflicts**



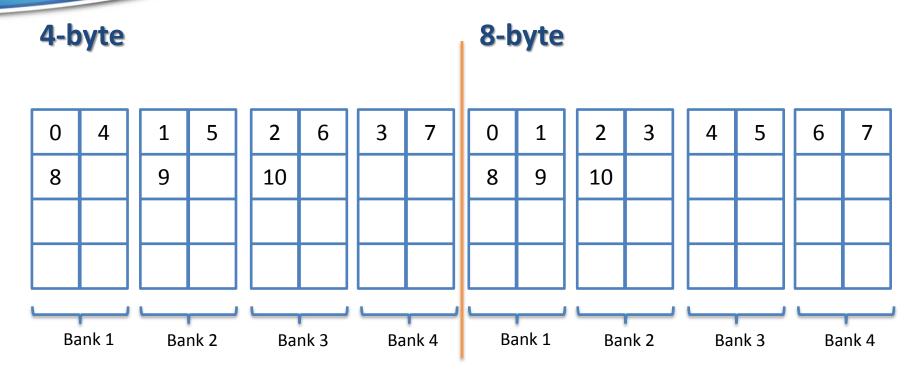
No conflicts



4 level conflict

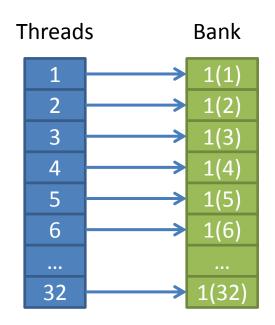


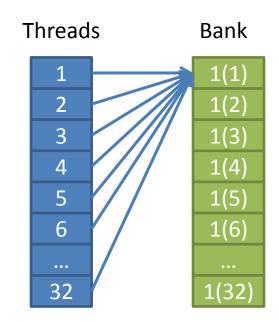
## Addressing types





### Conflicts



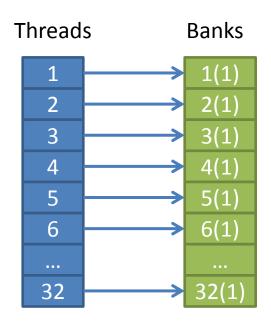


32 Threads request 32 elements from 1 bank.
32 level conflict

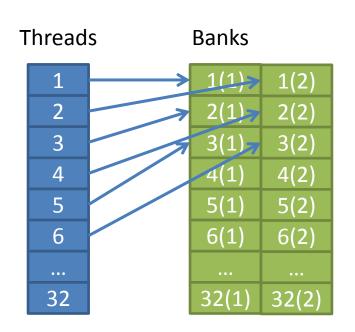
Threads request the same element Broadcast, no conflicts



### Conflicts



32 Threads request 32 8-byte elements from different Banks



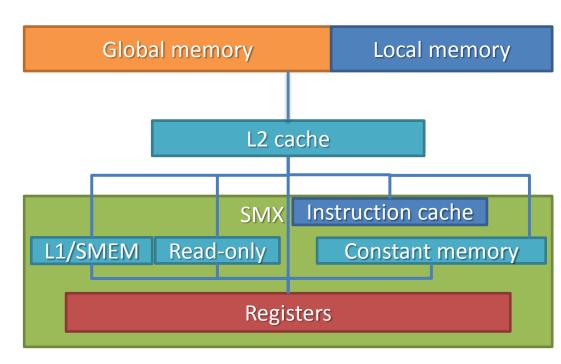
32 Threads request 32 4- byte elements

## Read-only and constant cache



## Read-only cache

- 48 KB on SMX separate of L1 cache (origins from texture cache)
- Latency is equal to L1 cache
- Caching 1D, 2D and 3D arrays
- Mardware data filtering, different patterns



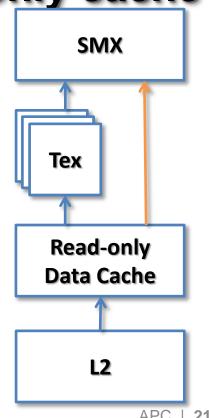


## Read-only cache

- const \_\_\_restrict
- Not using texture block
- Cached access to any global address
- Do not need textures preparations

#### What for?

- Separate of smem/L1 cache
- High bandwidth to L2
- Non-aligned access
- Caching 2-D arrays





## const \_\_restrict

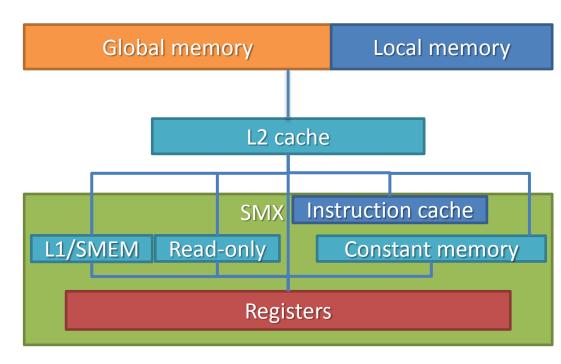
- Specify const \_\_restrict for parameters
- Compiler generates loading instruction using read-only cache

```
__global__
void saxpy (float x, float y,
            const float * __restrict input,
            float * output){
  size_t offset = threatIdx.x +
                  (blockIdx.x *
   blockDim.x);
  // Compiler will automatically use
   texture
  // for 'input'
  output[offset] = (input[offset] * x) + y;
```



### **Constant memory**

- Solution
  64 KB separate from L1 and read-only cache
- Access is faster than to global memory
- Accessible only for reading





### **Constants**

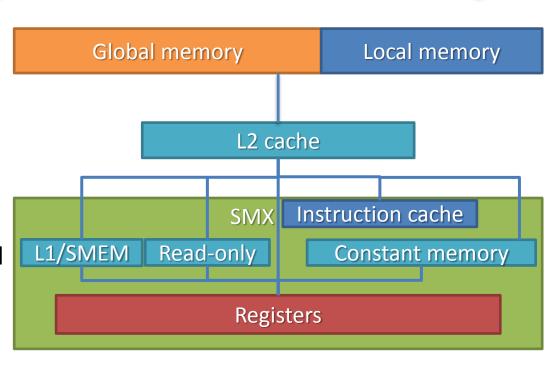
- \_\_\_constant\_
  - Cached access
  - Located in global memory
    - ✓ Max size 64 KB
    - ✓ Cached with each SMX separately
- Can't be changed from kernel
  - Doesn't support dynamic arrays
  - cudaMemcpyToSymbol()
  - cudaMemcpyFromSymbol()
- What for?
  - Separate of smem/L1 and read-only cache
  - Latency is lees than for L1

## Registers and local memory



## Registers and local memory

- SMX
  65536 32-bit registers for
- The fastest memory, used for computing
- Registers store:
  - Static arrays with predefined size
  - Variables
- Amount of registers allocated per thread defines the amount of blocks issued on SMX concurently





## Registers and local memory

- Compiler controls registers allocation
  - No more than 255 registers per thread
  - Could be artificially limited by –maxrregcount (compiler flag)
- If code requires more registers than hardware has there is a register spilling
  - A region of global memory (local memory) imitates registers
    - ✓ L1 and L2 cache requests
    - ✓ High latency

## Grid configuration



## **Grid configuration**

- Solution of Street S
  - Two main reasons:
    - ✓ Give enough parallel tasks to SMX
    - ✓ Share work between SMXs.
  - What to consider:
    - ✓ Amount of threads per block
    - ✓ Amount of blocks
    - ✓ Amount of work per block



## **Device occupancy**

- Device occupancy: amount of concurently issued threads per SMX
  - Could be measured as an amount of threads (warps)
  - Percent of maximum threads per SMX
- Different factors:
  - Amount of registers per thread
    - ✓ All SMX registers are divided between threads
  - Amount of shared memory per block
    - Shared memory is divided between blocks on SMX



# Solution Block size is multiple of warp size

- Even if you allocate less threads it will be rounded up by hardware
- Blocks could be too small
  - SMX can run up to 16 blocks concurrently
    - ✓ This can limit performance

#### **Block size**

SMX resources:
-Registers

-Shared memory





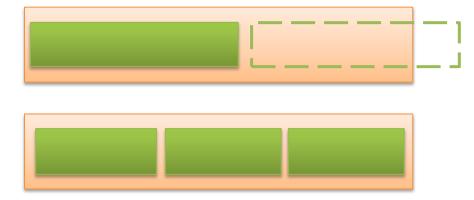


### Blocks could be too big:

- There is enough SMX resources to run additional threads but not for one more big block
- Block starts execution only when there is enough resources for all of it's threads

#### **Block size**

SMX resources:
-Registers
-Shared memory





## Occupancy calculator

- In SDK
  - CUDA Samples -> documentation -> CUDA Occupancy Calculator
- Melps to calculate the occupancy using the following parameters:
  - Architecture
  - Configuration of SMEM/L1
  - Threads per block
  - Registers per thread
  - Shared memory per block



#### **CUDA GPU Occupancy Calculator**

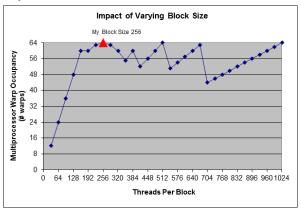
	e for help)	
1.) Select Compute Capability (click): 1.b) Select Shared Memory Size Config (bytes)	3,5 49152	<u>(He</u>
2.) Enter your resource usage:		
Threads Per Block	256	<u>(He</u>
Registers Per Thread	32	
Shared Memory Per Block (bytes)	4096	
(Don't edit anything below this line)		
3.) GPU Occupancy Data is displayed here and in the		
Active Threads per Multiprocessor	2048	<u>(He</u>
Active Warps per Multiprocessor	64	
Active Thread Blocks per Multiprocessor	8	
Occupancy of each Multiprocessor	100%	
Physical Limits for GPU Compute Capability:	3,5	
	<b>3,5</b>	
Threads per Warp		
Threads per Warp Warps per Multiprocessor	32	
Threads per Warp Warps per Multiprocessor Threads per Multiprocessor	32 64	
Physical Limits for GPU Compute Capability: Threads per Warp Warps per Multiprocessor Threads per Multiprocessor Thread Blocks per Multiprocessor Total # of 32-bit registers per Multiprocessor	32 64 2048	
Threads per Warp Warps per Multiprocessor Threads per Multiprocessor Thread Blocks per Multiprocessor	32 64 2048 16	
Threads per Warp Warps per Multiprocessor Threads per Multiprocessor Thread Blocks per Multiprocessor Total # of 32-bit registers per Multiprocessor	32 64 2048 16 65536	
Threads per Warp Warps per Multiprocessor Threads per Multiprocessor Thread Blocks per Multiprocessor Total # of 32-bit registers per Multiprocessor Register allocation unit size	32 64 2048 16 65536 256	
Threads per Warp Warps per Multiprocessor Threads per Multiprocessor Thread Blocks per Multiprocessor Total # of 32-bit registers per Multiprocessor Register allocation unit size Register allocation granularity	32 64 2048 16 65536 256 warp	
Threads per Warp Warps per Multiprocessor Threads per Multiprocessor Thread Blocks per Multiprocessor Total # of 32-bit registers per Multiprocessor Register allocation unit size Register allocation granularity Registers per Thread	32 64 2048 16 65536 256 warp 255	
Threads per Warp Warps per Multiprocessor Threads per Multiprocessor Thread Blocks per Multiprocessor Total # of 32-bit registers per Multiprocessor Register allocation unit size Register allocation granularity Registers per Thread Shared Memory per Multiprocessor (bytes)	32 64 2048 16 65536 256 warp 255 49152	

## Occupancy calculator

Click Here for detailed instructions on how to use this occupancy calculator.

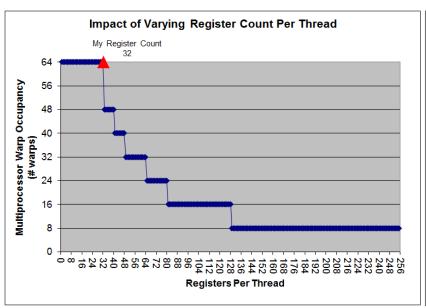
For more information on NVIDIA CUDA, visit http://developer.nvidia.com/cuda

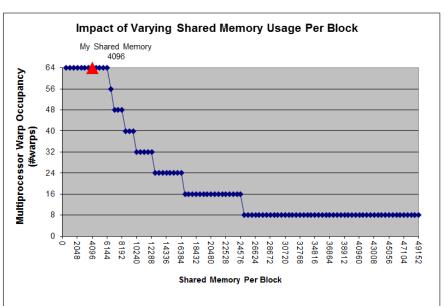
Your chosen resource usage is indicated by the red triangle on the graphs. The other data points represent the range of possible block sizes, register counts, and shared memory allocation.





## Occupancy calculator







### Guidelines

#### Selection Block size:

- Start with 128-256 treads per block
  - ✓ Increase or decrease according to your function parameters
- Multiple of warp size
- If the device occupancy is low:
  - ✓ Check the amount of resources registers and shared memory.

#### Grid size

- 1000 and more blocks
  - ✓ Uniform distribution of work over GPU
  - ✓ Code will be ready for execution on different architectures

# Arithmetic optimizations

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## **Control flow**

- Instructions are issued by warps
  - The entire warp issues the same instruction
- Scheduling
  - Warps are dynamically scheduled
  - Warp scheduler chooses a warp which has instructions ready to run
  - Warp switching helps to hide latency

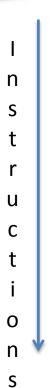


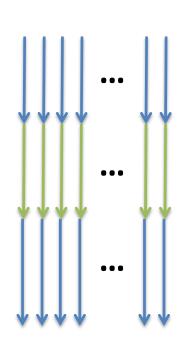
# **Control flow**

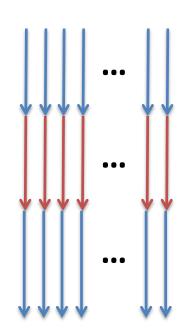
```
n
                          if ()
S
                            //then-clause
                          else
u
C
                            //else-clause
0
n
S
```



# Warp control flow

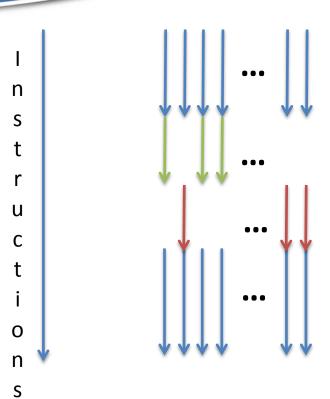


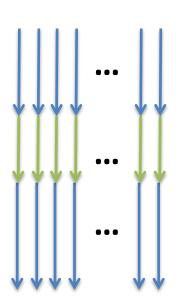






# Branching







- Sometiment of the second of
  - Maximum amount of instructions per cycle
  - Warp branching forces to issue extra instructions
- Serialization
  - Warp threads issue the same instruction serial
  - Most often cases are:
    - ✓ Shared memory bank conflicts
    - ✓ Uncoalesced memory access



- Compare resulting performance with arcitecture capabilities
  - Profiler measures performance as IPC (instructions per cycle)
  - Peak performance is listed in Programming Guide
- Serialization checking:
  - Amount of repeats: instructions\_issued instructions executed
  - Profiler reports
    - ✓ serial running percent executed instructions / issued instructions.



- S Branching
  - Profiler counters: divergent\_branch, branch
  - Allows to determine only branching serialization, all other reasons are not counted
- Make sure that you use 64-bit aritmetics only where you need
  - fp64 bandwidth is less than fp32
  - Values without"f" (12.2 instead of12.2f) interpreted asfp64

# Intrinsics



### Intrinsics

- Set and performance depends on the architecture
- Usually a sequence of software instructions can be replaced with one hardware instruction with less tolerance

- Use intrinsics if it's possible (\_\_sin(), \_\_sincos(), \_\_exp())
  - Available for many math functions
  - Tolerance is 2-3 bits less but performance is better.



### Performance test

#### Test structure

- Generating 2 sets of 10000000 random values
- Measuring time for 2 kernels:
  - ✓ Usual instruction
  - ✓ Intrinsic
- Output
  - ✓ Time
  - ✓ First 10 values to compare tolerance

## Example

```
__global__ void coss (float *x,float
                    *y,float *r,int n){
    int indexX = blockIdx.x *
         (blockDim.x) + threadIdx.x;
    if (indexX<n){</pre>
        float a,b,c;
        b=x[indexX];
        c=y[indexX];
        #pragma unroll
        for (int i=0; i<1000; i++)
            a = cosf(b) + cosf(c);
        r[indexX]=a;
```

```
_global___ void coss_i(float *x,float
                  *y,float *r,int n){
  int indexX = blockIdx.x *
       (blockDim.x) + threadIdx.x;
  if (indexX<n){</pre>
      float a,b,c;
      b=x[indexX];
      c=y[indexX];
      #pragma unroll
      for (int i=0; i<1000; i++)
      a+=__cosf(b)+__cosf(c);
          r[indexX]=a;
```



## Results

#### tolerance (for cos)

	Intrinsic GPU
GPU res	res
1697,11	1697,85
1485,11	1485,71
1683,31	1684,31
1346,15	1346,71
1453,87	1454,47
1833,71	1834,61
1383,51	1384,29
1923,46	1924,46
1836,79	1837,68
1842,56	1843,43

#### Time

	<b>GPU</b> time	Intrinsic GPU time
cos	1372,43	21,92
sin	1298,64	21,92
exp10	22,61	22,05
exp	22,65	22,05
fadd	21,85	21,84
fdiv	22,49	22,48
fmaf	21,84	21,84
powf	3,32	1,10
fsqrt	3,00	3,00



# Compiler flags

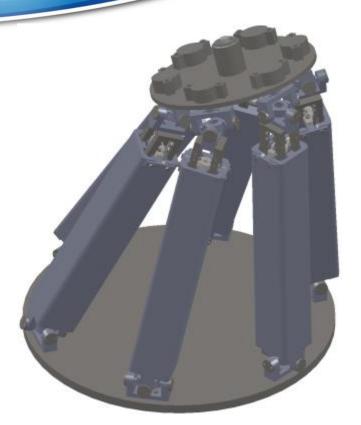
### Additional compiler flags:

- -ftz=true: replacing subnormal values with 0
- -prec-div=false : fast division instruction (some precision losses)
- -prec-sqrt=false : fast square root calculation (some precision losses)
- -use\_fast\_math : use intrinsics

# Example

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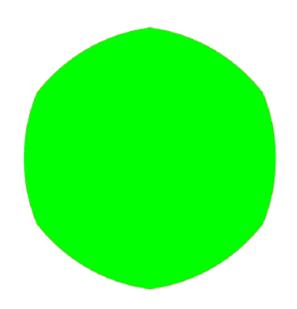


# Stewart platform

- Scanning Stewart platform workspace
- 350 slices by 1000\*1000 points
- Uses asinf()



# **Output data**



- MexapodSliceKernel generates bmp images of slices
- MexapodKernel calculates the circumcircle radius



#### 1: Block size

- Scanning workspace:
  - HexapodSliceKernel:
    - √ 27 registers per thread, 3977 bytes SMEM per warp
    - ✓ No more than 12 warps per SM
- Initial grid configuration: 16\*16 threads per block, 63 \*63 blocks, 350 kernel call
  - Blocks are too big 1 SM holds only 1 block
  - Device occupancy— 16%
  - Time: 5.6 s.



#### 2: Block size

### Optimized:

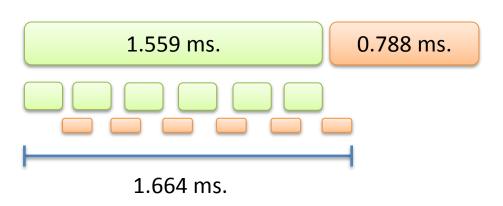
- 10\*10 threads per block, 100\*100\*350 blocks
- 12472 bytes SMEM per block
- Kernel and copy called only once
- Occupancy— 25%
- Time- 2,4 s.
- Not finished yet!



## 3: Asynchrony

### Asynchronous copy

- Split kernel into several iterations
- cudaMallocHost()
- cudaMemcpyAsync()
- cudaDeviceSynchronize()





#### 4: Block size

### Optimized:

- 16\*16 threads per block 63\*63\*50 blocks.
- 72 SMEM per block, 29 registers.
- Kernel and copy are called 7 times.
- Occupancy— 66,7%
- Time– 1.096 s.
- Speedup

   5.2 times.



#### Instruction Level Parallelism:

- Each thread calculates 2 output elements
- 8\*16 threads per block 63\*63\*50 blocks.
- 0 bytes SMEM per block, 47 registers.
- Kernel and copy called 7 times.
- Occupancy 41,7%(83,4%)
- Time– 1.090 s.

### 6:ILP

### ILP (adding #pragma unroll):

- Each thread calculates 2 output elements
- 8\*16 threads per block 63\*63\*50 blocks.
- 0 bytes SMEM per block, 47 registers.
- Kernel and copy are called 7 times.
- Occupancy 41,7%(83,4%)
- Time- 1,023 s.



#### 7:Fast math

#### Using intrinsics:

- 8\*16 per block
   63\*63\*50 blocks.
- 0 bytes SMEM per block, 40 registers.
- Kernel and copy are called 7 times.
- Occupancy –
   50%(100%)
- Time- **0,433** s.

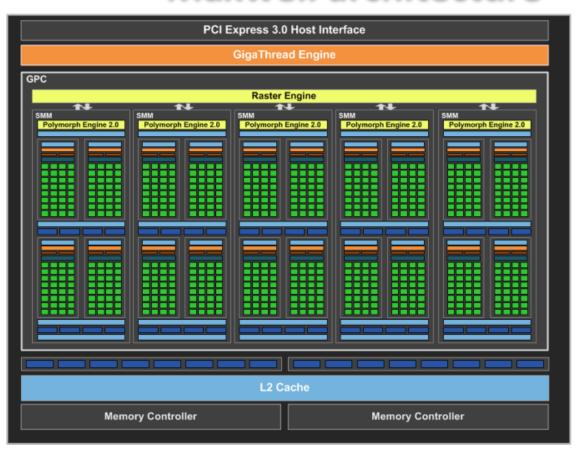


Time is 12,9 times less!!!



- S 5 SMM
- § 640 CUDA cores
- PCI-e 3.0
- 2 MB L2 cache

### Maxwell architecture





### **SMM** architecture

- 128 CUDA cores
- 8 texture processors
- 65536 32-byte registers
- 4 warp schedulers
- 64 KB L1/ texture cache
- 64 KB Shared memory
- 32 LD/ST unit
- 32 SF unit





### Results

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