

Transforming Javascript Event-Loop Into a Pipeline

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Abstract—The development of a real-time web application often starts with a feature-oriented approach allowing to quickly react to users feedbacks. However, this approach poorly scales in performance. Yet, the audience can increase by an order of magnitude in a matter of hours. This first approach is unable to deal with the higher connections spikes. It leads the development team to adopt a scalable approach often linked to new development paradigm such as dataflow programming. This represents a disruptive and continuity-threatening shift of technology. To avoid this shift, we propose to abstract the feature-oriented development into a more scalable high-level language. Indeed, reasoning on this high-level language allows to dynamically cope with audience growth and decrease.

We propose a compilation approach that transforms a Javascript, single-threaded real-time web application into a network of small independent parts communicating by message streams. We named these parts *fluxions*, by contraction between a flow¹ and a function. The independence of these parts allows their execution to be parallel, and to organize an application on several processors to cope with its load, in a similar way network routers do with IP traffic. We test this approach by applying the compiler to a real web application. We transform this application to parallelize the execution of an independent part and present the result.

I. INTRODUCTION

The growth of real-time web services is partially due to Internet’s capacity to allow very quick releases of a minimal viable product (MVP). In a matter of hours, it is possible to release a prototype and start gathering a user community. “*Release early, release often*”, and “*Fail fast*” are the punchlines of the web entrepreneurial community. It is crucial for the prosperity of such project to quickly validate that the proposed solution meets the needs of its users. Indeed, misidentifying the market need is the first reason for startup failure². Hence the development team quickly concretizes an MVP using a feature-driven approach and iterates on it.

If the service successfully complies with users requirements, its community grows with its popularity. The service needs to be scalable to be able to respond to this growth. However, feature-based development best practices are hardly compatible with the required parallelism. The features are organized in modules which overlap and disturb the organization of a parallel execution [6, 13, 17]. Eventually this growth requires to discard the initial approach to adopt a more efficient

processing model. Many of the most efficient models decompose applications into execution units, disregarding the initial feature oriented approach. This decomposition may be spread over a cluster of commodity machines [10]. MapReduce [7] and the Staged Event-driven Architecture (SEDA) [21] are famous examples of that trend. Once split, the service parts are connected by an asynchronous messaging system. Many tools have been developed to express and manage these service parts and their communications. We can cite Spark [22], MillWheel [1], Naiad [16], Storm [20], and many others. However, these tools are in disruption from the initial approach. It requires the development team either to be trained or to hire experts, and more importantly, to start over the initial code base. This shift causes the development team to spend development resources in background without adding visible value for the users. It is a risk for the evolution of the project as the second and third reasons for startup failures are running out of cash, and missing the right competences².

The risks described above come from a disruption between the two levels of application expression, the feature level and the execution level. To lift these risks, we propose a tool to map one level onto the other, and allow an automatic transition from one to the other and back. This tool should allow a continuous development process.

We focus on web applications driven by users requests, developed in Javascript using the *Node.js*³ execution environment. Javascript is increasingly used to develop web applications. It is the most used language on Github⁴ and StackOverflow⁵. We think that it is possible to analyze this type of application as a stream of requests, passing through a pipeline of stages. Indeed, the event-loop used in *Node.js* is very similar to a pipeline architecture. We propose a compiler to transform a Javascript application into a network of independent parts communicating by message streams and executed in parallel. We named these parts *fluxions*, by contraction between a flux and a function. The contribution of this paper are twofold. First, the identification of the problems arising from the isolation of the global memory into fluxions. Second, a partial resolution of these problems, allowing an experimental compilation. We present a tool as a proof of concept for this

¹flux in french

²<https://www.cbinsights.com/blog/startup-failure-post-mortem/>

³<https://nodejs.org/>

⁴<http://github.info/>

⁵<http://stackoverflow.com/tags>

compilation approach. We start by describing in section II the execution environment targeted by this compiler. Then, we present the compiler in section III, and its evaluation in section IV. We compare our work with related works in section V. And finally, we conclude this paper.

II. FLUXIONAL EXECUTION MODEL

In this section, we present an execution model to provide scalability to web applications. To achieve this, the execution model provides a granularity of parallelism at the function level. Functions are encapsulated in autonomous execution containers with their state, so as to be reallocated and executed in parallel. This execution model is close to the actors model, as the execution containers are independent and communicate by messages. The communications are assimilated to stream of messages, similarly to the dataflow programming model. It allows to reason on the throughput of these streams, and to react to load increases [4].

The fluxional execution model executes programs written in our high-level fluxionnal language, whose grammar is presented in figure 1. An application $\langle \text{program} \rangle$ is partitioned into parts encapsulated in autonomous execution containers named *fluxions* $\langle \text{flx} \rangle$. In the following paragraphs, we present the *fluxions*. Then we present the messaging system to carry the communications between *fluxions*. Finally, we present an example application using this execution model.

A. Fluxions

A *fluxion* $\langle \text{flx} \rangle$ is named by a unique identifier $\langle \text{id} \rangle$ to receive messages, and might be part of one or more groups indicated by tags $\langle \text{tags} \rangle$. A *fluxion* is composed of a processing function $\langle \text{fn} \rangle$, and a local memory called a *context* $\langle \text{ctx} \rangle$. At a message reception, the *fluxion* modifies its *context*, and sends messages on its output streams $\langle \text{streams} \rangle$ to downstream *fluxions*. The *context* stores the state on which a *fluxion* relies between two message receptions. In addition to message passing, the execution model allows *fluxions* to communicate by sharing state between their *contexts*. The fluxions that need this synchronization are grouped with the same tag, and loose their independence.

There are two types of streams, *start* and *post*, which correspond to the nature of the rupture point yielding the stream. We differentiate the two types with two different arrows, double arrow (\Rightarrow or \gg) for *start* rupture points and simple arrow (\rightarrow or \rightarrow) for *post* rupture points. The two types of rupture points are further detailed in section III-A1.

B. Messaging system

The messaging system assures the stream communications between fluxions. It carries messages based on the names of the recipient fluxions. After the execution of a fluxion, it queues the resulting messages for the event loop to process.

The execution cycle of an example fluxional application is illustrated in figure 2. * Bigger circles represent registered fluxions, while smaller ones indicate the step in the execution. The source code for this application is in listing

$\langle \text{program} \rangle$	\models	$\langle \text{flx} \rangle \mid \langle \text{flx} \rangle \text{ eol } \langle \text{program} \rangle$
$\langle \text{flx} \rangle$	\models	$\text{flx } \langle \text{id} \rangle \langle \text{tags} \rangle \langle \text{ctx} \rangle \text{ eol } \langle \text{streams} \rangle \text{ eol } \langle \text{fn} \rangle$
$\langle \text{tags} \rangle$	\models	$\& \langle \text{list} \rangle \mid \text{empty string}$
$\langle \text{streams} \rangle$	\models	$\text{null} \mid \langle \text{stream} \rangle \mid \langle \text{stream} \rangle \text{ eol } \langle \text{streams} \rangle$
$\langle \text{stream} \rangle$	\models	$\langle \text{type} \rangle \langle \text{dest} \rangle [\langle \text{msg} \rangle]$
$\langle \text{dest} \rangle$	\models	$\langle \text{list} \rangle$
$\langle \text{ctx} \rangle$	\models	$\{ \langle \text{list} \rangle \}$
$\langle \text{msg} \rangle$	\models	$[\langle \text{list} \rangle]$
$\langle \text{list} \rangle$	\models	$\langle \text{id} \rangle \mid \langle \text{id} \rangle, \langle \text{list} \rangle$
$\langle \text{type} \rangle$	\models	$\gg \mid \rightarrow$
$\langle \text{id} \rangle$	\models	<i>Identifier</i>
$\langle \text{fn} \rangle$	\models	<i>Imperative language and stream indicators</i>

Figure 1. Syntax of a high-level language to represent a program in the fluxionnal form

1 and the fluxional code for this application is in listing 2. The fluxion *reply* has a context containing the variable count and template. The plain arrows represent the actual message paths in the messaging system, while the dashed arrows between fluxions represent the message streams as seen in the fluxionnal application.

The *main* fluxion is the first fluxion in the flow. When the application receives a request, this fluxion triggers the flow with a *start* message containing the request, ②. This first message is to be received by the next fluxion *handler*, ③ and ④. The fluxion *handler* sends back a message, ⑤, to be enqueued, ⑥. The system loops through steps ③ through ⑥ until the queue is empty. This cycle starts again for each new incoming request causing another *start* message.

C. Service example

To illustrate the fluxional execution model, and the compiler, we present in listing 1 an example of a simple web application. This application reads a file, and sends it back along with a request counter.

```

1 var app = require('express')(),
2   fs = require('fs'),
3   count = 0;
4
5 app.get('/', function handler(req, res){
6   fs.readFile(__filename, function reply(err, data) {
7     count += 1;
8     res.send(err || template(count, data));
9   });
10 });
11
12 app.listen(8080);

```

Listing 1. Example web application

The handler function, line 5 to 10, receives the input stream of request. The count variable at line 3 counts the requests. This variable needs to be persisted in the fluxion *context*. The template function formats the output stream to be sent back to



the example in figure 2 is only partially explained in text, many parts (group, network, program, enqueue...) are not



I would suggest adding a brief paragraph to introduce Node.js and esprima and maybe move the url of Node.js stated in the footer of the last page to the first page where Node.js is already stated.



The discussion of the execution model is only discussed on a simple example (figure 2) and not discussed in general.

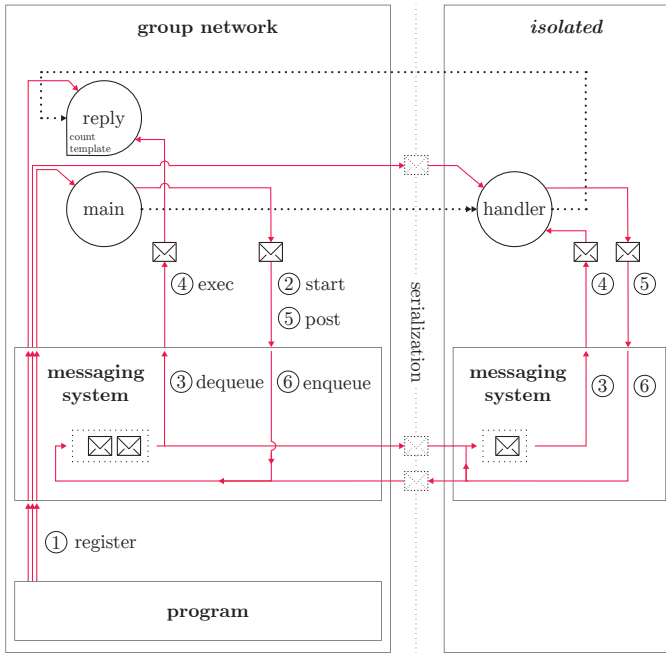


Figure 2. The fluxionnal execution model in details

the client. The `app.get` and `res.send` functions, respectively line 5 and 8, interface the application with the clients. And between these two interface functions is a chain of three functions to process the client requests : `app.get` \rightarrow `handler` \rightarrow `reply`. This application is transformed into the high-level fluxionnal language in listing 2 which is illustrated in Figure 2.



*
*

Transformation discussed in Sec 2.3 between Listing 1 and Listing 2 is not understandable many parts are not discussed in text. Even the parts discussed are difficult to follow (eg. "res" is created by the first fluxion main..."). how is this done? res is a shared variable, thus it requires to be protected how is this accomplished?

```

1 flx main & network
2 >> handler [res]
3   var app = require('express')(),
4     fs = require('fs'),
5     count = 0;
6
7   app.get('/', >> handler);
8   app.listen(8080);
9
10 flx handler
11 -> reply [res]
12   function handler(req, res) {
13     fs.readFile(__filename, -> reply);
14   }
15
16 flx reply & network {count, template}
17 -> null
18   function reply(error, data) {
19     count += 1;
20     res.send(err || template(count, data));
21   }

```

Listing 2. Example application expressed in the high-level fluxional language

The application is organized as follow. The flow of requests is received from the clients by the fluxion main, it continues in the fluxion handler, and finally goes through the fluxion reply to be sent back to the clients. The fluxions main and reply have the tag network. This tag indicates their dependency over the network interface, because they received the response from and send it back to the clients. The fluxion

handler doesn't have any dependencies, hence it can be executed in parallel, and replicated.

The last fluxion, reply, depends on the variable `res` created by the first fluxion, main. This variable is carried by the stream through the chain of fluxion to the fluxion reply that depends on it. The group network depends on this variable because it holds the references to the network sockets. This variable is created and consumed inside this group, therefore, it doesn't prevent the whole group from being replicated.

However, the last fluxion, reply, depends on its context to hold the variable `count` which prevents this replication. If it did not rely on this state, the group network would be stateless, and could be replicated to cope with the incoming traffic.

This execution model allows to parallelize the execution of an application. Some parts are arranged in pipeline, like the fluxion handler, some other parts are replicated, as could be the group network. This parallelization improves the scalability of the application. Indeed, as a fluxion contains its state and expresses its dependencies, it can be migrated. It allows to adapt the number of fluxions per core to adjust the resource usage in function of the desired throughput.

Our goal, as described in the introduction, is not to propose a new programming paradigm with this high-level language but to automate the architecture shift. We present the compiler to automate this architecture shift in the next section.

III. FLUXIONNAL COMPILER

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The source languages we focus on should present higher-order functions and be implemented as an event-loop with a global memory. Javascript is such a language : it doesn't require an event-loop, but it is often implemented on top of an event-loop. *Node.js* is an example of such an implementation. We developed a compiler that transforms a *Node.js* application into a fluxional application compliant with the execution model described in section II.



Section 3 describing the compiler is very high level and does not discuss in depth how the different steps are implemented in practice.

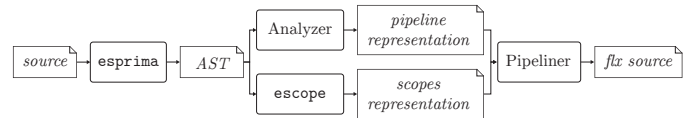


Figure 3. Compilation chain

The chain of compilation is described in figure 3. From the source of a *Node.js* application, the compiler extracts an Abstract Syntax Tree (AST) with *esprima*. From this AST, the *Analyzer* step identifies the limits of the different application parts and how they relate to form a pipeline. This first step outputs a pipeline representation of the application. Section III-A explains this first compilation step. In the pipeline representation, the stages are not yet independent and encapsulated into fluxions. From the AST, *escope* produces a representation of the memory scopes. The *Pipeliner* step analyzes the pipeline representation and the scopes representation to distribute the



I would also suggest adding labels to one of the fluxions in listing 2 to indicate its tags.

shared memory into independent groups of fluxions. Section III-B explains this second compilation step.

A. Analyzer step

The limit between two application parts is defined by a rupture point. The analyzer identifies these rupture points, and outputs a representation of the application in a pipeline form, with application parts as the stages, and rupture points as the message streams of this pipeline.

1) *Rupture points*: A rupture point is a call of a loosely coupled function. It is an asynchronous call without subsequent synchronization with the caller. In *Node.js*, I/O operations are asynchronous functions and indicate such rupture point between two application parts. Figure 4 shows a code example of a rupture point with the illustration of the execution of the two application parts isolated into fluxions. The two application parts are the caller of the asynchronous function call on one hand, and the callback provided to the asynchronous function call on the other hand.

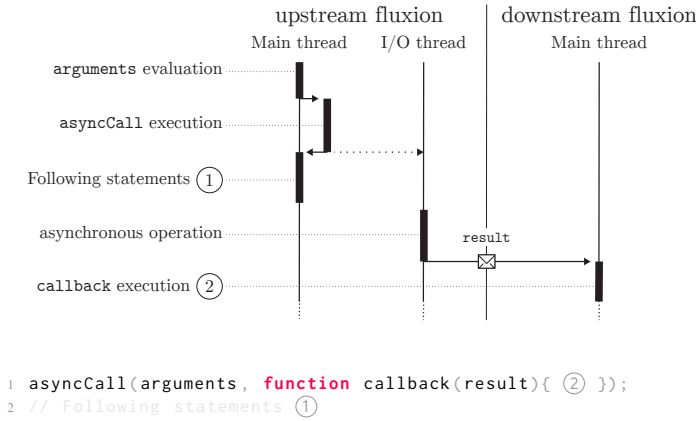


Figure 4. Rupture point interface

A callback is a function passed as a parameter to a function call. It is invoked by the callee to continue the execution with data not available in the caller context. We distinguish three kinds of callbacks, but only two are asynchronous : listeners and continuations. Similarly, there are two types of rupture points, respectively *start* and *post*.

Start rupture points are indicated by listeners. They are on the border between the application and the outside, continuously receiving incoming user requests. An example of a start rupture point is in listing 1, between the call to `app.get()`, and its listener handler. These rupture points indicate the input of a data stream in the program, and the beginning of a chain of fluxions to process this stream.

Post rupture points are indicated by continuations. They represent a continuity in the execution flow after an asynchronous operation yielding a unique result, such as reading a file, or querying a database. An example of a post rupture point is in listing 1, between the call to `fs.readFile()`, and its continuation reply.

2) *Detection*: The compiler uses a list of common asynchronous callees, like the `express` and `file system` methods. This list can be augmented to match asynchronous callees individually for any application. To identify the callee, the analyzer walks the AST to find a call expression matching this list.

After the identification of the callee, the callback needs to be identified as well to be encapsulated in the downstream fluxion. For each asynchronous call detected, the compiler test if one of the arguments is of type function. Some callback functions are declared *in situ*, and are trivially detected. For variable identifier, and other expressions, the analyzer tries to detect their type. To do so, the analyzer walks back the AST to track their assignments and modifications, and to determine their last value.

B. Pipeliner step

A rupture point eventually breaks the chain of scopes between the upstream and downstream fluxion. The closure in the downstream fluxion cannot access the scope in the upstream fluxion as expected. The pipeliner step replaces the need for this closure, allowing application parts to rely only on independent memory stores and message passing. It determines the distribution using the scope representation, which represents the variables' dependencies between application parts. Depending on this representation, the compiler can replace the broken closures in three different ways. We present these three alternatives with the example figure 5.

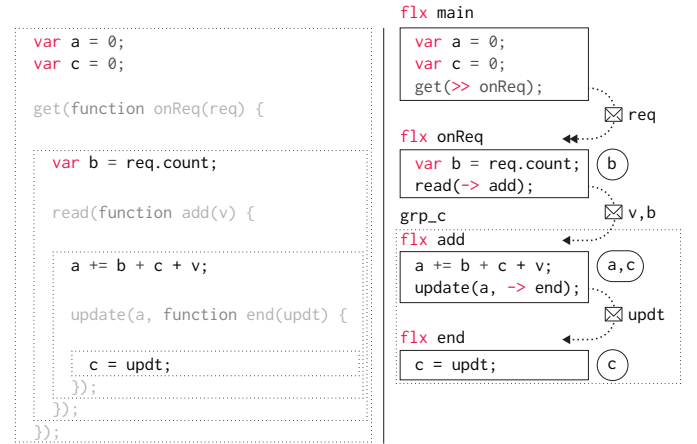


Figure 5. Variable management from Javascript to the high-level fluxionnal language

a) *Scope*: If a variable is modified inside only one application part in the current *post* chain, then the pipeliner adds it to the context of its fluxion.

In figure 5, the variable `a` is updated in the function `add`. The pipeliner step stores this variable in the context of the fluxion `add`.

b) *Stream*: If a variable is modified inside an application part, and read inside downstream application parts, then the pipeliner makes the upstream fluxion add this variable to the

message stream to be sent to the downstream fluxions. It is impossible to send variables to upstream fluxions, without race conditions. If the fluxion retro propagates the variable for an upstream fluxion to read, the upstream fluxion might use the old version while the new version is on its way.

In figure 5, the variable *b* is set in the function *onReq*, and read in the function *add*. The pipeliner step makes the fluxion *onReq* send the updated variable *b*, in addition to the variable *v*, in the message sent to the fluxion *add*.

Exceptionally, if a variable is defined inside a *post* chain, like *b*, then this variable can be streamed inside this *post* chain without restriction on the order of modification and read. Indeed, the execution of the upstream fluxion for the current *post* chain is assured to end before the execution of the downstream fluxion. Therefore, no reading of the variable by the upstream fluxion happens after the modification by the downstream fluxion.

c) Share: If a variable is needed for modification by several application parts, or is read by an upstream application part, then it needs to be synchronized between the fluxions. To respect the semantics of the source application, we cannot tolerate inconsistencies. Therefore, the pipeliner groups all the fluxions sharing this variable within the same tag. And it adds this variable to the contexts of each fluxions.

In figure 5, the variable *c* is set in the function *end*, and read in the function *add*. As the fluxion *add* is upstream of *end*, the pipeliner step groups the fluxion *add* and *end* with the tag *grp_c* to allow the two fluxions to share this variable.

IV. REAL CASE TEST

The goal of this test is to prove the possibility for an application to be compiled into a network of independent parts. We want to show the current limitations of this isolation and the modifications needed on the application to circumvent these limitations.

We present a test of our compiler on a real application, *gifsockets-server*⁶. This application was selected from the npm registry because it depends on *express*, it is tested, working, and simple enough to illustrate this evaluation. It is part of the selection from a previous work [5].

This application is a real-time chat using gif-based communication channels. The server transforms the received text into a gif frame, and pushes it back to a never-ending gif to be displayed on the client. Listing 3 is a simplified version of this application.

```
1 var express = require('express'),
2   app = express(),
3   routes = require('gifsockets-middleware'),
4   getRawBody = require('raw-body');
5
6 function bodyParser(limit) {
7   return function saveBody(req, res, next) {
8     getRawBody(req, {
9       expected: req.headers['content-length'],
10      limit: limit
11    }, function (err, buffer) {
12      req.body = buffer;
13      next();
14    });
15  };
16 }
```

```
14   });
15   };
16 }
17
18 app.post('/image/text', bodyParser(1 * 1024 * 1024), routes
19   .writeTextToImages);
20 app.listen(8000);
```

Listing 3. Simplified version of *gifsockets-server*

On line 18, the application registers two functions to process the requests received on the url */image/text*. The closure *saveBody*, line 7, returned by *bodyParser*, line 6, and the method *routes.writeTextToImages* from the external module *gifsockets-middleware*, line 3. The closure *saveBody* calls the asynchronous function *getRawBody* to get the request body. Its callback handles the errors, and calls *next* to continue processing the request with the next function, *routes.writeTextToImages*.

A. Compilation

We compile this application with the compiler detailed in section III. The function call *app.post*, line 18, is a rupture point. However, its callbacks, *bodyParser* and *routes.writeTextToImages* are evaluated as functions only at runtime. For this reason, the compiler ignores this rupture point, to avoid interfering with the evaluation.

The compilation result is in listing 4. The compiler detects a rupture point : the function *getRawBody* and its anonymous callback, line 11. It encapsulates this callback in a fluxion named *anonymous_1000*. The callback is replaced with a stream placeholder to send the message stream to this downstream fluxion. The variables *req*, and *next* are appended to this message stream, to propagate their value from the main fluxion to the *anonymous_1000* fluxion.

When *anonymous_1000* is not isolated from the main fluxion, as if they belong to the same group, the compilation result works as expected. The variables used in the fluxion, *req* and *next*, are still shared between the two fluxions. Our goal is to isolate the two fluxions, to be able to safely parallelize their executions.

```
1 flx main
2 >> anonymous_1000 [req, next]
3   var express = require('express'),
4     app = express(),
5     routes = require('gifsockets-middleware'),
6     getRawBody = require('raw-body');
7
8   function bodyParser(limit) {
9     return function saveBody(req, res, next) {
10       getRawBody(req, {
11         expected: req.headers['content-length'],
12         limit: limit
13       }, >> anonymous_1000);
14     };
15   }
16
17   app.post('/image/text', bodyParser(1 * 1024 * 1024),
18     routes.writeTextToImages);
19   app.listen(8000);
20
21 flx anonymous_1000
22 -> null
23   function (err, buffer) {
24     req.body = buffer;
25     next();
26   }
```

Listing 4. Compilation result of *gifsockets-server*

⁶<https://github.com/twolfson/gifsockets-server>

B. Isolation

In listing 4, the fluxion anonymous_1000 modifies the object req, line 23, to store the text of the received request, and it calls next to continue the execution, line 24. These operations produce side-effects that should propagate in the whole application, but the isolation prevents this propagation. Isolating the fluxion anonymous_1000 produces runtime exceptions. We detail in the next paragraph, how we handle this situation to allow the application to be parallelized. This test highlights the current limitations of the compiler, and presents future works to circumvent them.

1) *Variable req*: The variable req is read in fluxion main, lines 10 and 11. Then its property body is associated to buffer in fluxion anonymous_1000, line 23. The compiler is unable to identify further usages of this variable. However, the side effect resulting from this association impacts a variable in the scope of the next callback, routes.writeTextToImages. We modified the application to explicitly propagate this side-effect to the next callback through the function next. We explain further modification of this function in the next paragraph.

2) *Closure next*: The function next is a closure provided by the express Router to continue the execution with the next function to handle the client request. Because it indirectly relies on network sockets, it is impossible to isolate its execution with the anonymous_1000 fluxion. Instead, we modify express, so as to be compatible with the fluxionnal execution model. We explain the modification below.

```
1 flx main & express
2 >> anonymous_1000 [req, next]
3   var express = require('express'),
4     app = express(),
5     routes = require('gifsockets-middleware'),
6     getRawBody = require('raw-body');
7
8   function bodyParser(limit) {
9     return function saveBody(req, res, next) {
10       getRawBody(req, {
11         expected: req.headers['content-length'],
12         limit: limit
13       }, >> anonymous_1000);
14     };
15   }
16
17   app.post('/image/text', bodyParser(1 * 1024 * 1024),
18     routes.writeTextToImages);
19   app.listen(8000);
20
21 flx anonymous_1000
22 -> express_dispatcher
23   function (err, buffer) {
24     req.body = buffer;
25     next_placeholder(req, -> express_dispatcher);
26   }
27
28 flx express_dispatcher & express
29 -> null
30   merge(req, msg.req);
31   next();
```

Listing 5. Simplified modification on the compiled result

Originally, the function next is the continuation to allow the anonymous callback on line 11, to continue the execution with the next function to handle the request. To isolate the anonymous callback, this function is replaced on both ends. The result of this replacement is illustrated in listing 5. The express Router registers a fluxion named

express_dispatcher, line 27, to continue the execution after the fluxion anonymous_1000. This fluxion is in the same group express as the main fluxion, hence it has access to network sockets, to the original variable req, and to the original function next. The call to the original next function in the anonymous callback is replaced by a placeholder to push the stream to the fluxion express_dispatcher, line 24. The fluxion express_dispatcher receives the stream from the upstream fluxion anonymous_1000, merges back the modification in the variable req to propagate the side effects, before calling the original function next to continue the execution, line 30.

After the modifications detailed above, the server works as expected for the subset of functionalities we modified. The isolated fluxion correctly receives, and returns its serialized messages. The client successfully receives a gif frame containing the text.

C. Future works

We intend to implement the compilation process presented into the runtime. A just-in-time compiler would allow to identify callbacks dynamically evaluated, and to analyze the memory to identify side-effects propagations instead of relying only on the source code. Moreover, this memory analysis would allow the closure serialization required to compile application using higher-order functions.

V. RELATED WORKS

The idea to split a task into independent parts goes back to the Actor's model [12], functional programming [13] and the following works on DataFlow leading up to Flow-based Programming (FBP) and Functional Reactive programming (FRP) [8]. Both FBP and FRP, recently got some attention in the Javascript community with the projects *NoFlo*⁷, *Bacon.js*⁸ and *react*⁹.

The execution model we presented in section II, is inspired by some works on scalability for very large systems, like the Staged Event-Driven Architecture (SEDA) by Matt Welsh [21] and later by the MapReduce architecture [7]. It also drew its inspiration from more recent work following SEDA. Among the best-known following works, we cited in the introduction Spark [22], MillWheel [1], Naiad [16] and Storm [20]. The first part of our work stands upon these thorough studies. However, we believe that it is difficult for most developers to distribute the state of an application. This belief motivated us to propose a compiler from an imperative programming model to these more scalable, distributed execution engines.

The transformation of an imperative programming model to be executed onto a parallel execution engine was recently addressed by Fernandez *et. al.* [9]. However, as in similar works [18], it requires annotations from developers, therefore partially conserves the disruption with the feature-based development. Our approach avoids the need for annotations,

⁷<http://noflojs.org/>

⁸<https://baconjs.github.io/>

⁹<https://facebook.github.io/react/>

thus targets a broader range of developers, and not only ones experienced with parallel development.

A great body of work focus on parallelizing sequential programs [3, 14, 15, 19]. Because of the synchronous execution of a sequential program, the speedup of parallelization is inherently limited. On the other hand, our approach is based on an asynchronous programming model. Hence the attainable speedup is not limited by the main synchronous thread of execution [2, 11].

Our compiler uses the *estools*¹⁰ suite to parse, manipulate and generate source code from Abstract Syntax Tree (AST). Our implementation is based on the work by Ryan Dahl : *Node.js*, as well as on one of the best-known *Node.js* web framework : *Express*¹¹. *

VI. CONCLUSION

In this paper, we presented our work on a high-level language allowing to represent a web application as a network of independent parts communicating by message streams. We presented a compiler to transform a *Node.js* web application into this high-level representation. To identify two independent parts, the compiler spots rupture points in the application, possibly leading to memory isolation and thus, parallelism. We presented an example of a compiled application to show the limits of this approach. The parallelism of this approach allows code-mobility which may lead to a better scalability. We believe it can enable the scalability required by highly concurrent web applications without discarding the familiar, feature-based programming models.

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¹⁰<https://github.com/estools>

¹¹<http://expressjs.com/>

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