CPSC 424/524 Spring 2018 Assignment 3: Triangular Matrix Multiplication

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Building and files

- build-run-mpi.sh has instructions for running and building all non-serial versions of the program. It needs to be slightly modified to run different test cases.
- build-run.sh the original serial build and run script
- Makefile makefile for running on the cluster. module load Langs/Intel/15 MPI/OpenMPI/2.1.1-intel15 must be performed before running it.
- matmul.c barely modified serial version of the matrix multiplication program
- matmul parallel *.c contain the versions of matrix multiplication for tasks 2,3, and 4. Task 5 can use matmul parallel 4.c.
 - Define the preprocessor symbol #DEBUG to print out additional timing information both amount for communication and amount for calculation.
- parallel.c is the main function for tasks 2 through 5. It can also run with matmul.c in serial (if you run mpirun -c 1)
- readme.md this readme file
- serial.c barely modified version of the serial main program
- test_local.c main program for testing small cases of the matrix multiplication program without access to the large .dat files on the cluster.

All building and running is done in slightly modified versions of build-run-mpi.sh , which is then run with sbatch build-run-mpi.sh .

Task 1:

Here is the output of the serial program which is run with srun build-run.sh:

```
Matrix multiplication times:
  N TIME (secs) F-norm of Error
  1000
        0.1067 0.0000000000
                  0.0000000001
  2000
          0.8078
  4000
          16.7762
                    0.0000000003
  8000 141.0507
                    0.0000000014
real
    2m40.429s
      2m39.549s
user
sys 0m0.545s
```

Task 2:

Task 2a:

From the table for 2a, it is clear that there is a large imbalance in the amount of work processed by the lower numbered tasks compared to the higher numbered ones - this is almost certainly because while each task is assigned an equal number of rows, the first rows are much shorter and require less work to be done. Despite this, even task the last task (N=8000, #p=8, p=7) spends 10 seconds waiting for communication to complete, implying that we may be able to speed up the length of time it takes to finish the entire problem by streamlining the communication process.

The amount of time spent communicating does increase each time we double the number of processes, so eventually we will stop getting performance benefits from adding more tasks, but we the program would probably still be faster if we used p=16 compared to p=8.

Somehow, my code is faster than the serial program for large matrices even when p=1 and I have no idea why. That said, the times are roughly comparable, so I'm not too worried.

Here are the results on one node:

p=1

```
N TIME (secs) F-norm of Error
1000 0.1543 0.00000000000
2000 1.1806 0.00000000000
4000 14.0403 0.0000000000
8000 118.0142 0.00000000000
```

```
N TIME (secs) F-norm of Error
1000 0.1093 0.00000000000
2000 0.8251 0.00000000000
4000 10.1081 0.0000000000
8000 85.0111 0.00000000000
```

p=4

F-norm of Error 0.00000000000 0.00000000000
0.0000000000
0.0000000000
0.0000000000

p=8



Task 2a: Per process calculation vs communication tradeoff table

• p: total number of tasks.

N : Size of matrix

• #p : Task id

• calculating : Seconds that task spent calculating

• communicating : Seconds a task spent waiting for communication to complete

```
p N
       #p calculating communicating
1 1000 0 0.154257
                    0.000017
1
  2000
        0
           1.180626
                     0.000017
  4000
        0
           14.040238 0.000019
1 8000 0 118.014200 0.000019
2 1000 0 0.051110 0.058181
                    0.005574
2 1000 1 0.104436
  2000
        0
           0.383503
                     0.441595
2
  2000 1 0.800137
                     0.025683
2 4000 0 4.069784 6.038318
2
 4000 1 9.740579
                    0.368602
2 8000 0 39.813674 45.197437
2
  8000
        1 83.744027
                     1.269418
4
 1000
       0 0.015565
                     0.062778
4 1000 1 0.036512
                    0.040705
 1000 2 0.049549
                    0.028352
        3 0.055626
4 1000
                    0.023159
  2000
        0
           0.112637
                     0.454927
  2000
        1 0.272792
                     0.289623
 2000 2 0.376180
                    0.188934
4
 2000 3 0.427465
                    0.140764
                    4.953861
4
  4000 0 0.846208
  4000
        1
           2.255041
                     3.521447
        2 3.566092
4
  4000
                     2.222135
 4000 3 4.304769 1.496251
4
  8000 0 10.972409 46.989782
  8000 1 29.324692 28.540916
4
  8000
        2
           38.852703
                     19.060007
        3 44.075137 13.889009
4
  8000
 1000 0 0.004357
                    0.045225
8
 1000 1 0.011039
                    0.036936
8 1000 2 0.016431
                     0.031623
8
  1000
        3
           0.020570
                     0.028138
8
  1000
        4 0.023934
                     0.024864
 1000 5 0.026235
                    0.023224
8
8
 1000 6 0.027858
                    0.021690
 1000
       7 0.028416
8
                     0.021403
  2000
           0.032583
8
  2000 1 0.082960
                     0.240665
 2000 2 0.121056
                     0.203988
8
8
  2000 3 0.153598
                     0.172996
8
  2000 4 0.179858
                     0.147982
8
  2000
        5 0.198630
                     0.130860
8
  2000 6 0.211002
                    0.119658
 2000 7 0.216788 0.115520
8
  4000 0 0.238107 2.842225
  4000 1 0.624450
8
                     2.420885
8
  4000
           0.931620
  4000
        3 1.220106
                     1.836653
 4000 4 1.552247
8
                   1.510751
8
  4000 5 1.838238 1.230892
                    1.117819
  4000 6 1.956883
8
  4000
        7
           2.070927
                     1.010045
       0 2.507068 30.986574
8
  8000
 8000 1 8.270425 25.077264
8
8
  8000 2 12.808546 20.563267
8
  8000 3 17.372915 16.023761
  8000
           19.253519
                     14.167156
8
  8000
        5 21.606345 11.839339
 8000 6 22.385881 11.083529
8
8
 8000 7 23.038996 10.455570
```

Task 2b: (Round robin accross the sockets SLURM commands)

Spreading the processes across more nodes ends up creating some serious communication latency - adding about 10 seconds for changing from 2 to 4 nodes for 8 processes.

The added latency to each communication call does suggest that we might want to look at making our requests asynchronous to speed up total time, which will reduce apparent latency (but won't increase throughput or help us if we're waiting for another process to compute).

The load balancing strategy described in task 4 would be a good idea - there is some serious discrepancies between the calculation time of task 0 and of task 7 (or 3 when there are 4

total tasks), and load balancing would even that out.

Total time of additional runs

```
p, nodes TIME (secs) F-norm of Error

4,2 57.1226 0.00000000000

8,2 34.2154 0.00000000000

8,4 45.0222 0.0000000000
```

Per process calculation vs communication tradeoff table

- p : total number of tasks.
- node : Number of nodes the tasks are spread across
- #p : Task id
- calculating : Seconds that task spent calculating
- communicating : Seconds a task spent waiting for communication to complete

p node #p calculating communicating -- --- 4 1 0 10.972409 46.989782 4 1 1 29.324692 28.540916 4 1 2 38.852

Task 3:

Interesting aside:

I've been developing my program on my laptop and running on the server only to get final times. Interestingly, I get a segmentation fault when running MPI_Barrier on the server with an unfinished Isend but not on my laptop - it's an interesting difference in the implmentation of OpenMPI on the two computers. I think that MPI_Gatherv on my laptop functions more like a barrier than it does on the grace cluster.

Task 3a:

There wasn't as much of an improvement between task 2 and task 3 as I'd hoped. The program is perhaps more scalable because the time spent communicating for the high rank processes is significantly shorter than task 2, but the low rank processes show only a few percentage points of difference (probably because they are waiting for the poorly load balanced higher rank processes to finish calculating).

Honestly, load balancing is probably more important than making communication asynchronous for this problem with this many processes.

p=1

```
N TIME (secs) F-norm of Error
1000 0.1540 0.0000000000
2000 1.1786 0.00000000000
4000 14.0739 0.0000000000
8000 117.9903 0.00000000000
```

p=2

N	TIME (secs)	F-norm of Error
1000	0.1074	0.0000000000
2000	0.8238	0.0000000000
4000	9.4087	0.0000000000
8000	83.8489	0.0000000000

p=4

N	TIME (secs)	F-norm of Error
1000	0.0691	0.0000000000
2000	0.5033	0.0000000000
4000	5.2822	0.0000000000
8000	51.0482	0.0000000000

```
N TIME (secs) F-norm of Error
1000 0.0413 0.00000000000
2000 0.2825 0.0000000001
4000 2.7062 0.0000000002
8000 29.7321 0.00000000000
```

Task 3a: Per process calculation vs communication tradeoff table

• p : total number of tasks.

Size of matrix

• #p : Task id

• calculating : Seconds that task spent calculating

• communicating: Seconds a task spent waiting for communication to complete

ŗ	o 1	N	#p	calculating c	communicating
-					
1	1 1	1000	0	0.154022	0.000016
1	1 2	2000	0	1.178573	0.000009
1	1 4	4000	0	14.073851	0.000015
1	1 8	8000	0	117.990323	0.000014
2	2 1	1000	0	0.051080	0.056288
2	2 1	1000	1	0.104701	0.003424
2	2 2	2000	0	0.383458	0.440373
2	2 2	2000	1	0.808689	0.015969
2		4000		4.149853	5.258834
2	2 4	4000	1	9.345853	0.064124
		8000		39.853979	43.994954
		8000		83.595920	0.255661
2		1000		0.015381	0.053669
		1000		0.036668	0.011069
		1000		0.049572	0.019173
2		1000		0.056145	0.013480
_		2000		0.112228	0.391119
_		2000		0.272539	0.063865
_		2000		0.376729	0.124515
2		2000		0.426019	0.124313
2		2000 4000			
		4000 4000		0.847508	4.434644
4				2.407402	1.333655
4		4000 4000		3.453503	1.816430
4		4000		4.522799	0.760413
4		8000		10.978843	40.069344
4		8000		29.417425	5.842403
4		8000		39.835418	11.163131
4		8000		44.858624	6.191687
		1000		0.004699	0.036638
8		1000		0.011167	0.007971
8		1000		0.016560	0.007659
8		1000		0.020629	0.009523
8		1000		0.024134	0.012360
8	3 1	1000	5	0.027080	0.013844
8	3 1	1000	6	0.027915	0.013137
8	3 1	1000	7	0.028445	0.013179
8	3 2	2000	0	0.032533	0.250013
8	3 2	2000	1	0.081673	0.032853
8	3 2	2000	2	0.122260	0.028986
8		2000		0.153922	0.045857
8		2000		0.181195	0.063588
8		2000		0.202200	0.078424
		2000		0.212499	0.069333
		2000		0.218251	0.065063
8		4000		0.239247	2.466912
		4000		0.620603	0.550439
8		4000		0.930059	0.599245
		4000 4000		1.209316	0.599245
8		4000 4000		1.550454	0.813007
8		4000 4000		1.857511	0.837457
		4000		1.965945	0.735021
		4000		2.158771	0.548249
		8000		2.319516	27.412585
8		8000		8.217945	4.101042
8		8000		12.054720	3.617743
		8000		17.198141	3.838697
8		8000		19.167104	6.533131
		8000		22.405130	7.279827
8		8000		22.441349	7.266907
8	3 8	8000	7	23.820234	5.913296

Task 3b:

The table(s) below confirm my analysis in task 3a. The high amount of variance (and low minimum) in time spent waiting for communication to finish by each process indicates that almost all of the time spent waiting for communication could be fixed by load balancing, which would probably result in a dramatic improvement.

Total times of additional runs

Per process calculation vs communication tradeoff table

- p : total number of tasks.
- node : Number of nodes the tasks are spread across
- #p : Task id
- calculating : Seconds that task spent calculating
- communicating : Seconds a task spent waiting for communication to complete

р	node	#p	calculating	communicating
4		0	10.978843	40.069344
4	1	1	29.417425	5.842403
4		2	39.835418	11.163131
4		3	44.858624	6.191687
4		0	9.159979	41.915732
4		1	26.346534	13.454560
4		2	39.721998	11.300924
4	2	3	44.854408	6.235647
8	1	0	2.319516	27.412585
8	1	1	8.217945	4.101042
8	1	2	12.054720	3.617743
8	1	3	17.198141	3.838697
8	1	4	19.167104	6.533131
8	1	5	22.405130	7.279827
8	1	6	22.441349	7.266907
8	1	7	23.820234	5.913296
8	2	0	1.765878	37.675139
8	2	1	4.785282	15.772475
8	2	2	10.535894	16.404557
8	2	3	14.108697	19.088410
8	2	4	33.422029	5.939334
8	2	5	37.347937	2.068089
8	2	6	22.300875	17.137172
8	2	7	22.539932	16.929151
8	4	0	1.765878	37.675139
8	4	1	4.785282	15.772475
8	4	2	10.535894	16.404557
8	4	3	14.108697	19.088410
8	4	4	33.422029	5.939334
8	4	5	37.347937	2.068089
8	4	6	22.300875	17.137172
8	4	7	22.539932	16.929151

Task 4:

Task 2 and task 3 clearly have load balancing issues. I decided that the simplist way to balance the load better is to ensure that each process has roughly the same number of elements of the matrix, which is easy enough to do by adding rows to the first process until it has (N*(N+1)/2) * num_processes) rows, and repeat that with the later processes, leaving any extra matrix rows to the last process.

As you can see, this decreased the running time dramatically (almost a factor of 2 compared to task 3). Intuitively, this process also seems more scalable - we waste far less running time waiting the last few processes to complete.

I'm not entirely sure why running 4 tasks on 1 node took longer than spreading it across 2 nodes.

Total times of matrix multiplications

```
p,nodes TIME (secs) F-norm of Error

4,1 48.3193 0.0000000007

4,2 38.2903 0.00000000007

8,1 21.3519 0.0000000009

8,2 26.4483 0.0000000009

8,4 26.4483 0.0000000009
```

Per process calculation vs communication tradeoff table

- p : total number of tasks.
- node : Number of nodes the tasks are spread across
- #p : Task id
- calculating : Seconds that task spent calculating
- communicating : Seconds a task spent waiting for communication to complete

```
p node #p calculating communicating
-- ---- --
4 1
       0 48.071516 0.247734
4 1
      1 40.917110 7.346713
4 1 2 35.727831 12.567271
4 1 3 31.566219 16.754695
4 2
          38.019593 0.270688
       0
4 2 1 33.210402 5.031797
4 2 2 28.178327 10.083532
4 2 3 24.931262 13.373344
      0 20.997349 0.354565
8 1
8 1
      1
          20.742037 0.522906
      2 18.897301 2.386293
8 1
8 1 3 17.617265 3.686062
8 1 4 16.078185 5.228097
8 1 5 15.317710 6.019837
8 1 6 14.404397 6.945837
      7 13.525022 7.841147
8 1
8 2 0 15.992409 10.455903
8 2 1 16.232806 3.881858
8 2
      2 16.223746 5.810235
8 2
           15.353427 8.564842
8 2 4 25.981504 0.429699
8 2 5 24.702734 1.732883
8 2 6 12.867661 13.578189
      7 12.160585 14.315729
0 15.992409 10.455903
8 2
8 4
8 4 1 16.232806 3.881858
8 4 2 16.223746 5.810235
8 4 3 15.353427 8.564842
      4 25.981504 0.429699
5 24.702734 1.732883
8 4
8 4
8 4 6 12.867661 13.578189
8 4
     7 12.160585 14.315729
```

Task 5:

I didn't need to change anything from my task 4 to get this to work as I had planned for this situation from the very beginning of the assignment.

```
N TIME (secs) F-norm of Error

7633 19.3182 0.0000000008

Process 0 spent 16.548526 time calculating and 2.769722 time communicating
Process 1 spent 16.662782 time calculating and 2.530805 time communicating
Process 2 spent 19.064767 time calculating and 0.183009 time communicating
Process 3 spent 16.645378 time calculating and 2.623345 time communicating
Process 4 spent 13.903794 time calculating and 5.382887 time communicating
Process 5 spent 13.010575 time calculating and 6.281878 time communicating
Process 6 spent 10.668788 time calculating and 8.649129 time communicating
```