



LLVM Demystified

Building a Feedback Fuzzer

By
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Who Am I

- **Co-Founder @ ZINAD IT.**
- **Security Analyst at Hackerone.**
- **OSCP, OSCE.**
- **Love fuzzing and exploit writing.**

Agenda

- **What is LLVM ?**
- **Why should I care ?**
- **Is it hard to learn ?**
- **LLVM Passes.**
- **Building the taint analysis pass.**

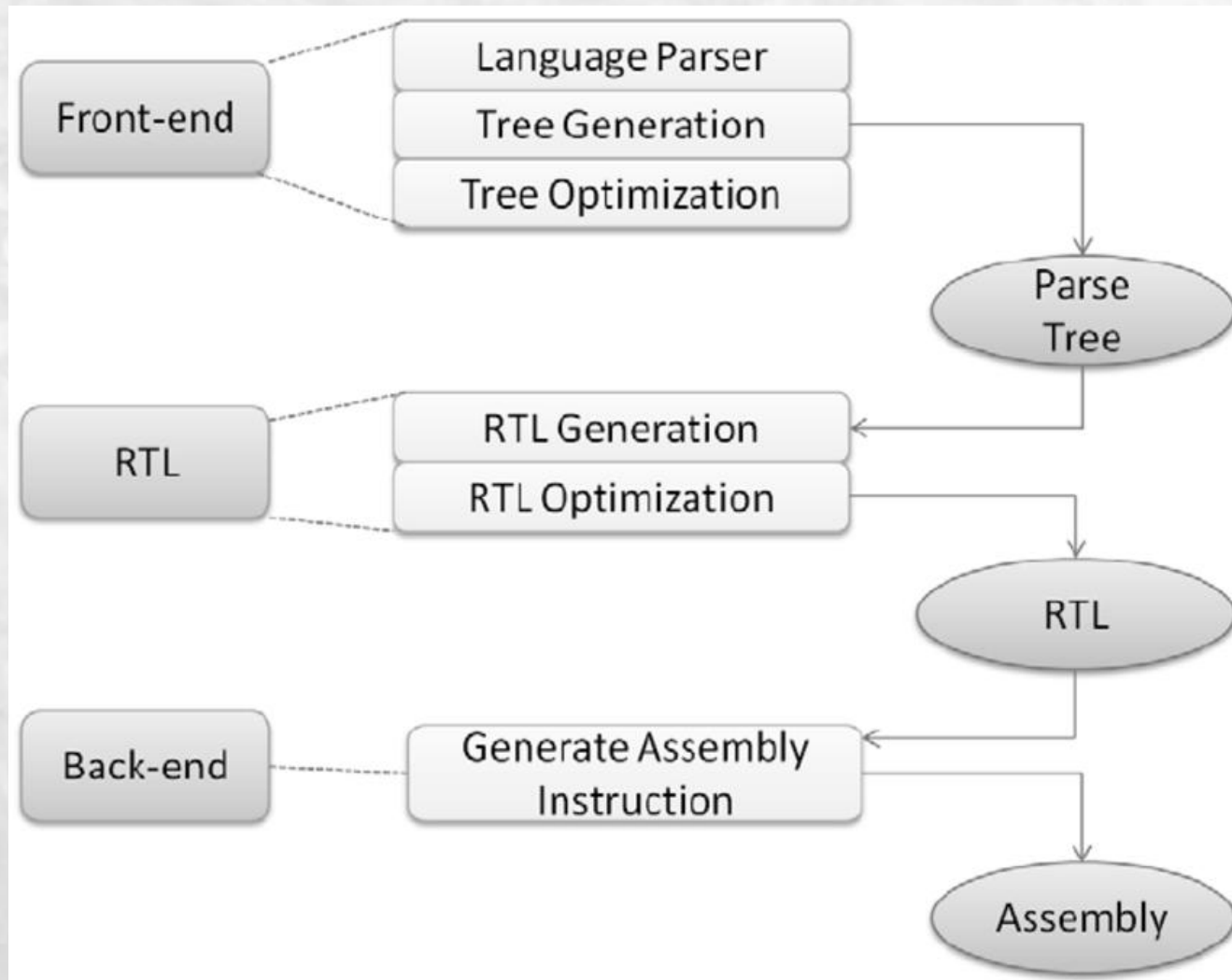
What is LLVM ?



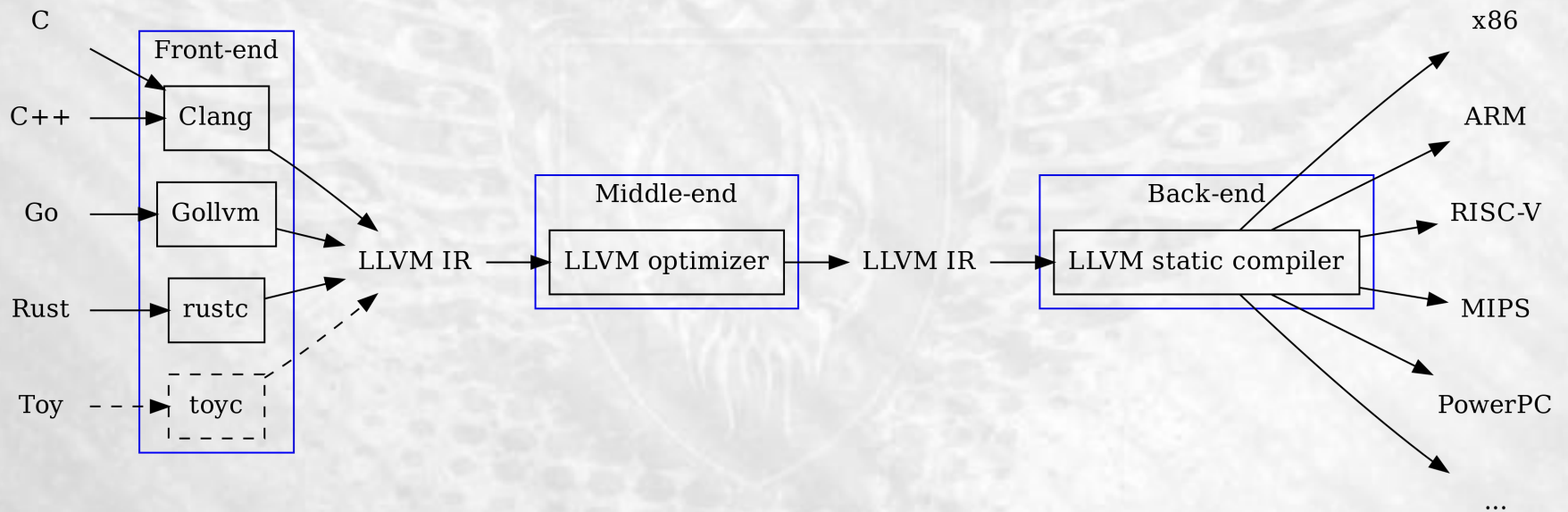
What is LLVM ?

- **A set of compiler and toolchain technologies, which can be used to develop a front end for any programming language and a back end for any instruction set architecture.**

Classic Compiler Architecture



LLVM Architecture



LLVM IR

- **LLVM uses IR to represent the code.**
- **Optimization passes can run on IR.**
- **You can write your own passes.**

IR Example

```
int main() {  
    return 12;  
}
```

Compile With:

```
clang -S -emit-llvm -O3 hello.c
```

IR Example : IR Code

```
define @main() {  
    ret i32 12  
}
```

IR Second Example

```
int main() {  
    int a = 12;  
    int b = 16;  
    return a+b;  
}
```

Compile With:

```
clang -S -emit-llvm -O0 hello2.c
```

IR Second Example : IR Code

```
int main() {  
    int a = 12;  
    int b = 16;  
    return a+b;  
}
```

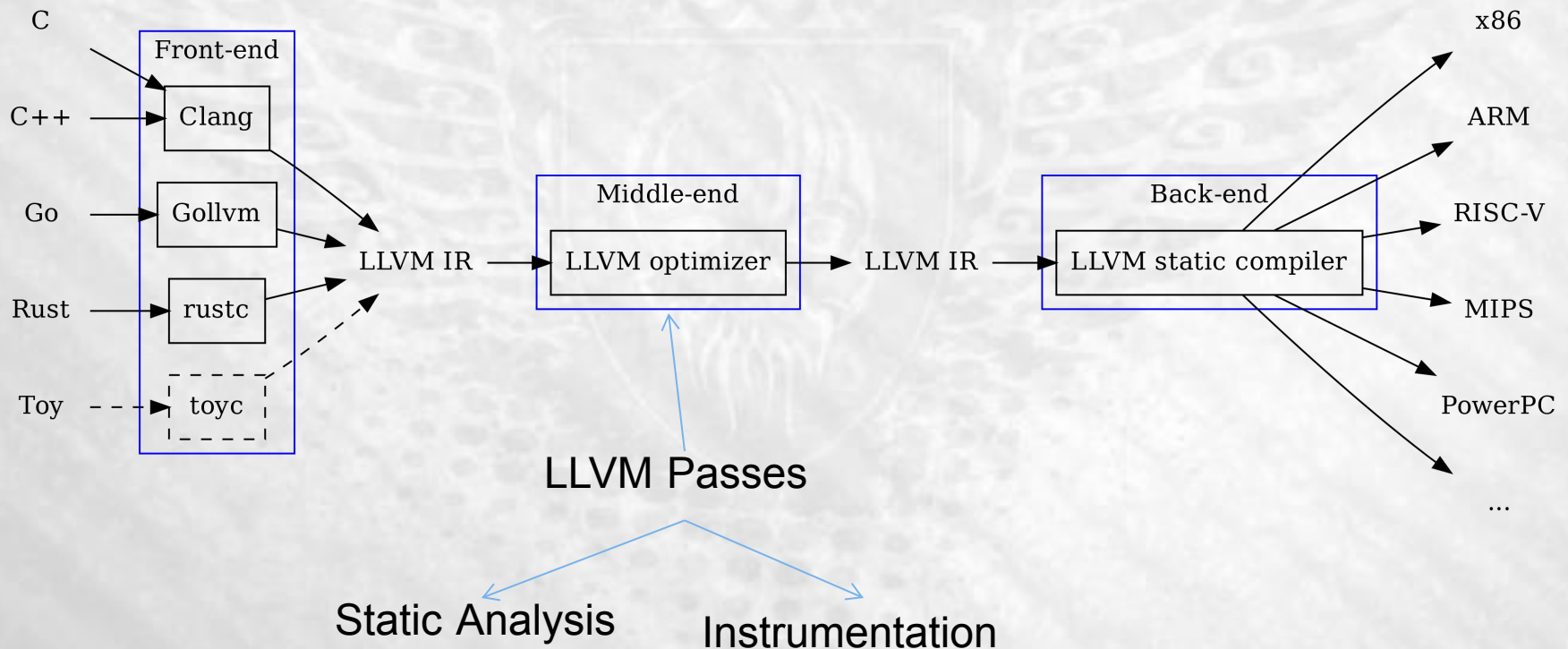
```
define i32 @main() {  
    %1 = alloca i32, align 4  
    %2 = alloca i32, align 4  
    %3 = alloca i32, align 4
```

```
    store i32 0, i32* %1, align 4  
    store i32 12, i32* %2, align 4  
    store i32 16, i32* %3, align 4
```

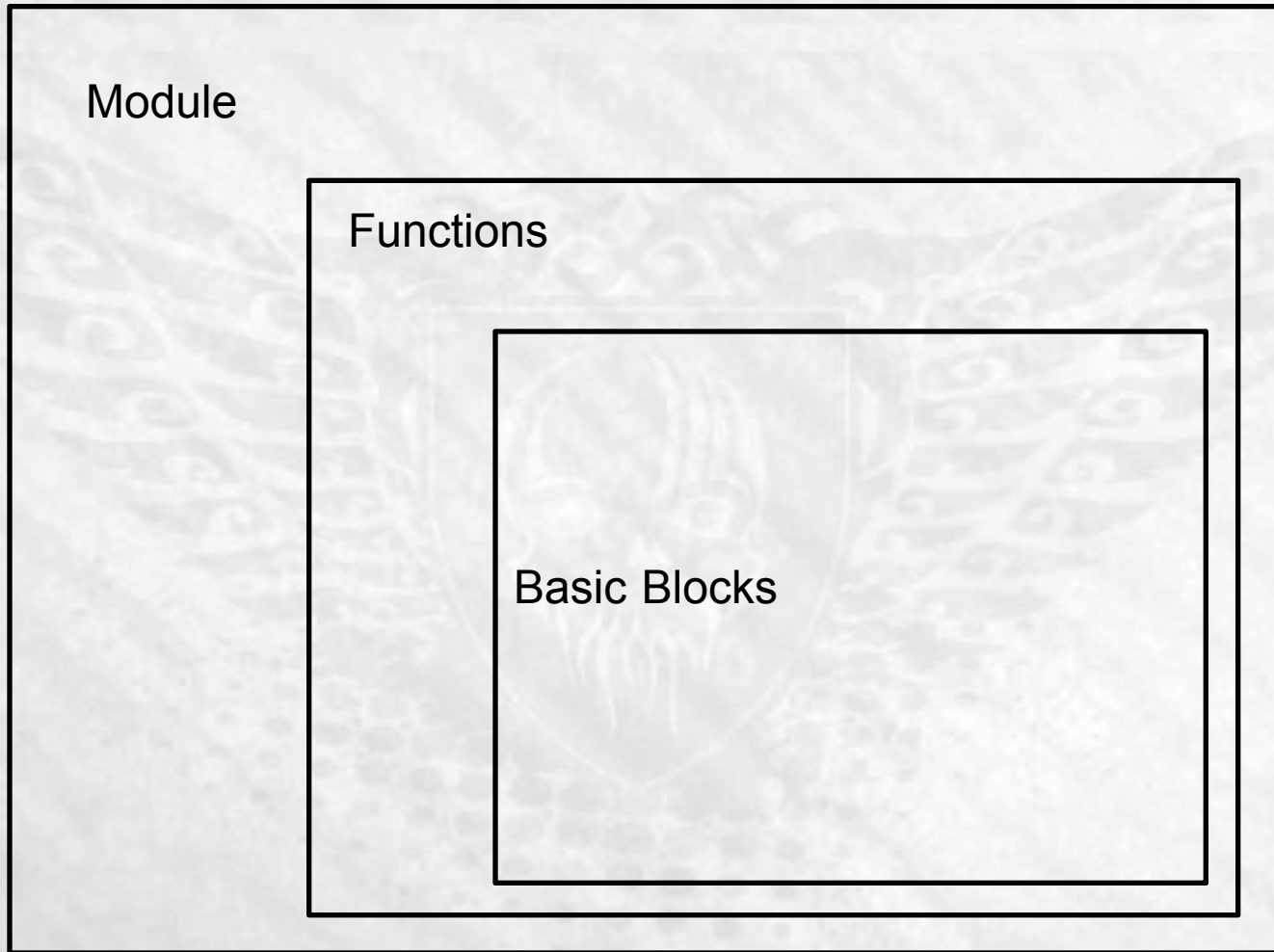
```
    %4 = load i32, i32* %2, align 4  
    %5 = load i32, i32* %3, align 4  
    %6 = add nsw i32 %4, %5  
    ret i32 %6
```

```
}
```

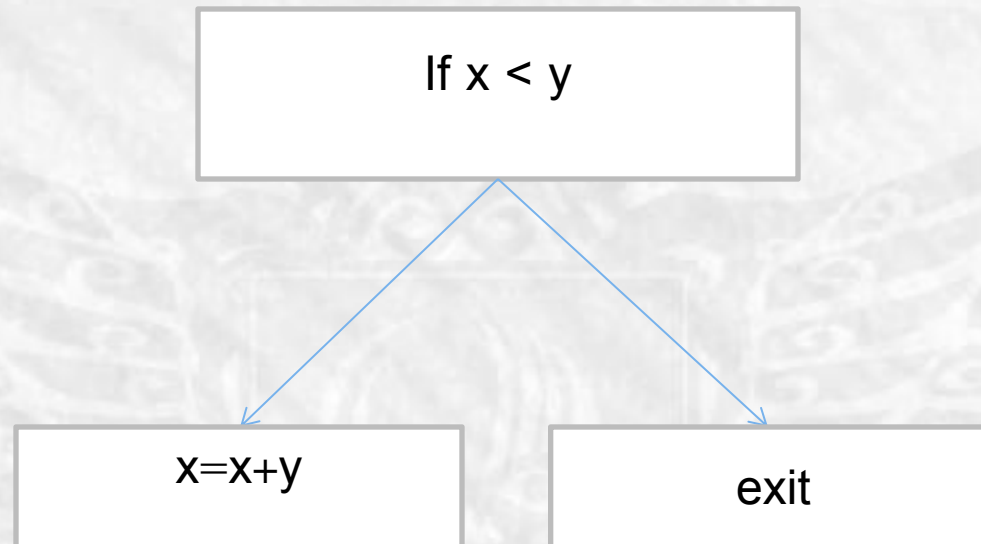
Why Should I Care ?



How Passes See the Code



Basic Blocks



Familiar Right

```
enup= dword ptr 10h

push    ebp |
mov     ebp, esp
and     esp, 0FFFFFFF0h
sub     esp, 10h
call    _main
mov     [esp+10h+var_4], 0AAAh
nop
nop
nop
cmp     [esp+10h+var_4], 0
jz      short loc_40165E
```

```
add     eax, 1
add     eax, 2
add     eax, 3
add     eax, 1
add     eax, 2
add     eax, 3
add     eax, 1
add     eax, 2
add     eax, 3
add     eax, 1
add     eax, 2
add     eax, 3
add     eax, 1
add     eax, 2
add     eax, 3
add     eax, 1
add     eax, 2
add     eax, 3
jnp     short loc_40168B
```

```
loc_40165E:
sub     eax, 1
sub     eax, 2
sub     eax, 3
sub     eax, 1
sub     eax, 2
sub     eax, 3
sub     eax, 1
sub     eax, 2
sub     eax, 3
sub     eax, 1
sub     eax, 2
sub     eax, 3
```

```
loc_40168B:
mov     eax, 0
leave
retn
_main endp
```

LLVM Passes

- **Module Pass.**
- **Function Pass.**
- **Basic Block Pass.**



Pass Skeleton

```
1  bool runOnModule(Module &m) override {
2      for (Function &f: m) {
3          for (BasicBlock &bb : f) {
4              //Do something here.
5          }
6      }
7  }
```

Parsing Instructions

```
for (BasicBlock &bb : f) {  
    for (Instruction &i : bb) {  
        unsigned opcode = i.getOpcode();  
        switch (opcode) {  
            case llvm::Instruction::Store:  
            {  
                llvm::StoreInst *storeinst = llvm::dyn_cast<llvm::StoreInst>(&i);  
                //First operand of store.  
               StringRef operand1 = storeinst->getOperand(0)->getName();  
                //Second operand of store.  
                StringRef operand2 = storeinst->getOperand(1)->getName();  
                //Checking if the taintedVars vector contains this input.
```

Is it hard to learn ?

RTFM

Vs

RTFC



Building the Taint Analysis Pass

How it works ?

- We start by classify functions by thier type:

```
enum ApiType {  
    API_ENTRY,  
    API_SOURCE,  
    API_SINK,  
    API_SANITIZE,  
    API_COPY,  
    API_USERFUNC  
};
```


How it Works

- We loop through all the functions and mark the tainted ones as such

```
struct TaintNode {  
    std::string name;  
    int arg;  
    ApiType type;  
};
```

```
struct TaintEdge {  
    TaintNode src;  
    TaintNode dst;  
};
```

Taint Analysis

```
for (BasicBlock &bb : f) {
    for (Instruction &i : bb) {
        unsigned opcode = i.getOpcode();
        switch (opcode) {
            case llvm::Instruction::Store:
            {
                llvm::StoreInst *storeinst = llvm::dyn_cast<llvm::StoreInst>(&i);
                //First operand of store.
               StringRef operand1 = storeinst->getOperand(0)->getName();
                //Second operand of store.
               StringRef operand2 = storeinst->getOperand(1)->getName();
                //Checking if the taintedVars vector contains this input.
                if (std::find(taintedVars.begin(), taintedVars.end(), operand1) != taintedVars.end())
                {
                    // Element in vector.
                    // The instruction is loading operand1 into operand2.
                    // So let's add operand2 into our taint list.
                    taintedVars.push_back(operand2);
                }
                break;
            }
        }
    }
}
```

Printing the Stack Trace

- We keep track of the taints as edges.
- We follow the edges and see if a source is reaching a dangerous sink.

```
struct TaintEdge {  
    TaintNode src;  
    TaintNode dst;  
};
```

Running the Module

```
opt -load=libLLVM_TaintPass.so' -taintanalyzer -instnamer <  
'/media/fady/Data/Work/Defcon/LLVM/llvmPass/testcases/bof.ll'
```

```
=====
🔥 Vulnerability detected: dangerous use of 'strcpy' 🔥
🔍 Analysis Trace (argument index is zero based):
  ↓ Function : main argument 0
  ↺ Function : func1 argument 0
  ↺ Function : func2 argument 0
  ↺ Function : func3 argument 0
  ❌ Function : strcpy argument 0
=====
```



Q & A