Auto Scaling Online Learning

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We propose a framework and algorithms for scaling online machine learning up or down, according to demand, and the priorities of the system. Different systems have different needs in terms of the cost assigned to machines, the cost of a bad user experience, etc. In this project, we focus most on the part that is general to auto scaling any application (and not only machine learning algorithms), but propose a framework that takes some ML particularities into account.

1 Introduction

Online machine learning algorithms operate on a single instance at a time. They have become particularly popular in natural language processing and applications with streaming data, including classification, ranking, etc [1, 2, 3]. Online learning is particularly interesting in the scenarios where data keeps streaming in, such as a web search engine doing advertisement placement. It is also interesting for scenarios where the whole dataset is too large to fit in main memory, as online learning only operates on a single example at a time.

A lot of tasks that use online learning have a particular structure that can be broken down into 2 major components: 1 - Learning a model from data, and 2 - making predictions according to the model. Going back to the web search engine scenario as an example: the system needs to make predictions for every user doing a query - and must also learn from the feedback given by those users. We call machines that make predictions **predictors**, and machines that learn from the feedback **learners**.

This structure comes with multiple challenges. First, the the amount of data is always growing, so archiving it comes at a cost, both because of storage constraints and computation constraints. A solution to this is to just keep the current model in memory, and archive the rest in the background. A second challenge is the variable speed at which

data streams in. Imagine a learning problem where in the learning dataset is a live twitter stream for a hashtag. In this scenario the rate at which the data comes in is a function of the popularity of the hashtag. Finally, different applications have different costs for learning, and different requirements for the latency of predictions.

The problem we tackled in this project is the problem of automatically handling the resources needed for online learning. The ideal system would allocate the resources necessary to keep the prediction latency acceptable, while at the same time learning appropriately. Finally, the system would be able to handle bursts (such as increase in demand) and different learning requirements for different systems. Since online learning in a distributed system is a research problem on its [4, 5], we abstracted this part from our work, and focused on some of the systems challenges.

Current works in auto-scaling[6, 8] address many of the issues we address in this report (such as load prediction, framing the problem as a cost minimization problem, etc). However, one aspect that we thought was lacking in current work (at least from the papers we read) is dealing with node failures and uncertainty. We reformulate the cost function in terms of expected cost, in order to account for uncertainty pertaining future load predictions and node failures. Finally, we evaluate some baselines and our proposed approach with simulated loads.

2 Architecture

We assume an architecture similar to the one presented in Figure 1. There is a set of nodes acting as a Load Balancer (whose addresses are known to the clients of the machine learning service). When a client needs a predictor or a learner, it first requests an a node from the Load Balancer, and then it proceeds to make the request. In Figure 1, the client got assigned the node in red.

The Load Balancer reads the state of the current system from some distributed, reliable storage. The state is comprised of which nodes are up and acting as learners or predictors, and a measure of the **power** of each node - that is, how many requests each node can effectively handle in a time interval. This allows for load balancing even when heterogeneous nodes are present.

The state of the current system is modified every so often by the set of nodes denoted **Autoscaler** in Figure 1. The Autoscaler periodically sends heartbeat messages to nodes and to the load balancer, and tests the prediction / learning time. It also gets the status - i.e., the number of requests served, eventual node failures, etc. Finally, it then updates

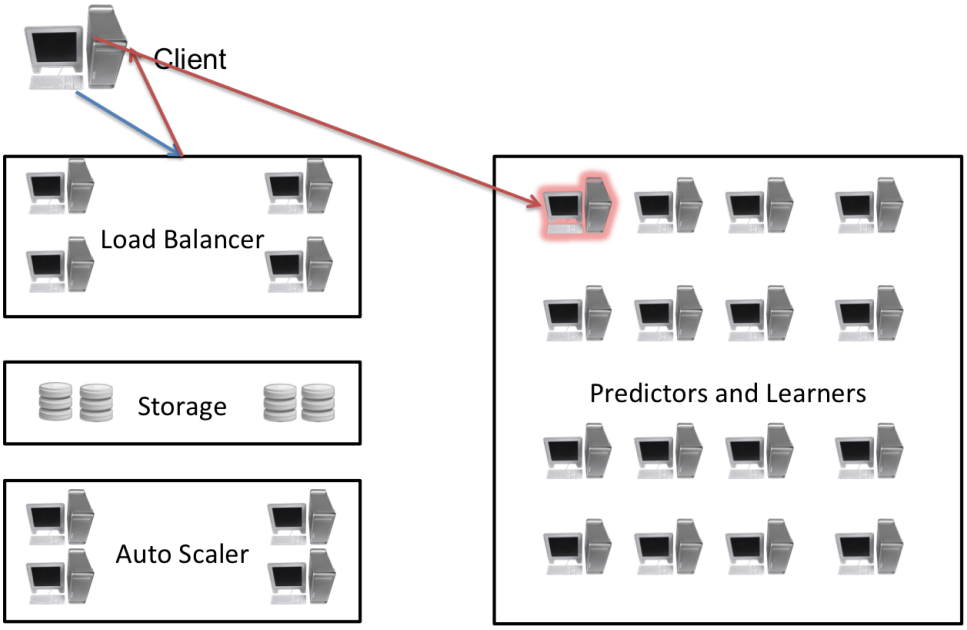


Figure 1: Architecture: Node asks the load balancer for a predictor, and then asks the predictor for a prediction.

the state in the distributed storage according to some policy. This architecture is very straightforward, and can be used with commodity machines for handling a lot of requests, while still being resilient to failures (depending of course on the Autoscaler policy).

3 Auto Scaling

3.1 Objective

The problem of auto scaling online learning can be formulated as minimizing a cost function over a sequence of time intervals. This function for each time interval is presented in Equation 1.

$$cost = \alpha SLA + \beta P + \gamma LR + \delta L \tag{1}$$

SLA is an indicator variable that is positive when a Service License Agreement (SLA) is violated in the time interval in question. This SLA is related to predictions, and can be defined as a threshold on the average time per prediction, on the maximum time per prediction, etc. There is a cost α associated with the SLA violation, that depends on the specific application that depends on the system. A bad user experience due to slow page loading on a search engine can be fatal to the success of the business, while a slower prediction for some internal testing service

can be less costly. **P** denotes the number of predictors active in the time interval, while β is a measure of the cost of each predictor. **LR** is a measure of the learning rate achieved by the machine learning algorithm. Some of the most popular online learning algorithms, such as Stochastic Gradient Descent, have a property that error ϵ decreases linearly with the number of training examples [7]: that is, the relation is $O(1/\epsilon)$, meaning that if training examples are available, adding more learners would offset the training error, and thus increase whatever measure of Learning Rate is being used - although probably not with linear speedup [4, 5]. **L** denotes the number of learners active in the time interval.

This objective formulation decouples learning from predicting, which brings advantages and disadvantages. Usually, the processor / memory requirements for learning and predicting are different - one would want a more powerful machine for learning than for predicting. Decoupling also ensures that failures in learners do not affect predictors, and vice versa. However, having learners and predictors together could be interesting in that learning involves predicting, and thus separating them creates the need to do predictions twice for the same examples: once in the predictor and once in the learner. In this project, we abstracted out the learning part, and focused on minimizing the first two terms in Equation 1. This choice was made assuming that a fixed, small number of learners would be sufficient for many applications, and that this number would not need to be changed very often - in contrast to predictors, who would need to be changed according to load.

3.2 Load Prediction

Ideally, we would employ a machine learning algorithm in order to predict the load in each future time interval. Features that we would consider include:

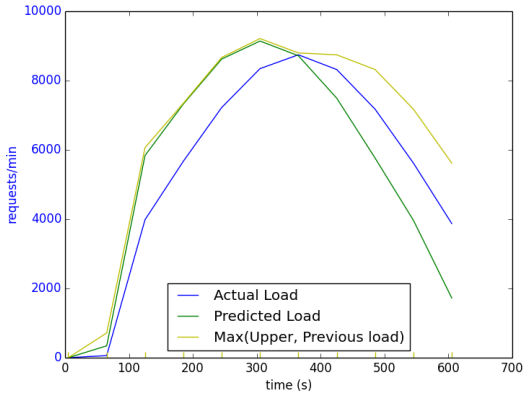
1. Time of day: traffic in most applications have load patterns that are very time-specific. Most people are sleeping at night, have a lunch break around the middle of the day, etc.

2. Trends in other services by the same company: if Facebook sees an increase in the number of users in the website at a particular time, it is more likely that internal services that use machine learning will also see an increase.

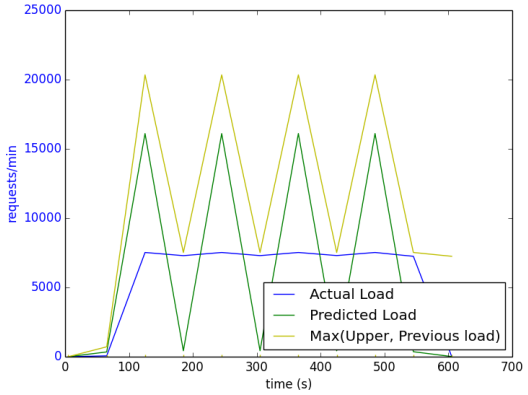
3. Usual load pattern: certain applications have very predictable load patterns. An internal service that always gets called in exponentially growing bursts would be an example.

4. Trend in the last time interval(s): by measuring the load at smaller intervals in the last time interval, one can build a "load trend" for the next time interval.

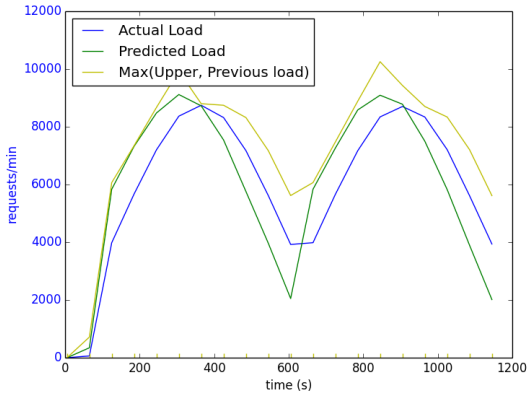
Since we used simulated data (we do not have access to such a rich dataset), we employed the trend in the last time interval in order to predict future load, by using traditional linear regression. In order to ensure a smoother decay (and thus prevent SLA violations due to wrong predictions), we took the max between the upper limit on a 95% confidence interval on the linear regression and the number of requests in the previous time interval as a prediction. Figure 2 illustrates the behaviour of our predictions on simulated data. We assume that we must predict the load at every minute, and collect the current load at every 5 seconds in order to build the trend line. The blue lines are the actual load, while the green and yellow lines indicate the linear regression prediction and the max we cited before, respectively. It is clear that the yellow line is more "conservative", and always stays on top of the actual load. This is what we used for predicting loads in our experiments.



(a) Load Pattern 1



(b) Load Pattern 2



(c) Load Pattern 3

Figure 2: Predicted vs Actual load, different load patterns.

3.3 Strategies

3.3.1 Naive Strategy

The Naive strategy, presented in Algorithm 1, is probably the first solution anyone would come up with when faced with the auto scaling problem. It is a reactive strategy that adds more nodes when the SLA is being violated, and removes nodes when the SLA is far from being violated. In our experiments, we assumed that the SLA was based on average prediction time being smaller then some constant T, and set $C = T/2$.

if *There is an SLA violation in the current time interval* **then**
| Add more nodes to the system
end
if *The mean prediction time is smaller than some constant C* **then**
| Remove nodes from the system
end

Algorithm 1: Naive Strategy

3.3.2 Power Strategy

The Power strategy, depicted in Algorithm 2, takes into account the power (number of requests the machine can handle) of each node, and compares the sum of the power of all the current nodes against the number of predicted requests for the next time interval. If the current power of the system is smaller than the predicted requests, it adds more nodes. If removing nodes does not cause the power of the system to be smaller than the number of predicted requests, it removes nodes.

while *Current Power* $<$ *predicted requests* **do**
| Add more nodes to the system
end
while *Current Power - power(*$node_i$*)* $>$ *predicted requests* **do**
| Remove $node_i$
end

Algorithm 2: Power Strategy

3.3.3 Smart Strategy

The two strategies presented before do not take into account the costs α and β. Furthermore, they do not take into account the probability of nodes failing, or the uncertainty associated with the load prediction. A reformulated objective function is presented in Equation 2. Here, we

are selecting the configuration of machines P and L that minimizes the expected cost in the next iteration.

$$min_{P,L} : E[cost] = E[\alpha SLA + \beta P + \gamma LR + \delta L] \tag{2}$$

Keeping the learners fixed, as we did before, the reformulated objective becomes:

$$\begin{aligned} min_P : E[cost] &= E[\alpha SLA + \beta P] \\ &= \alpha E[SLA] + \beta P \end{aligned}$$

One way of defining a SLA violation is when we have less power in a time interval than needed for the number of requests. Ideally, we would integrate out the uncertainty over our future prediction, and calculate $E[SLA]$ as follows, given that ρ is the probability that a node fails in the time interval, and each r denotes a possible value for the future load.

$$E[SLA] = P(SLA|P, \rho) = \int P(power < r|P, \rho, r)dr$$

This would require a more reliable estimator of future load, one that could output a probability distribution over the possible values of r. There are multiple machine learning algorithms that can do this, but since we are only using the trend in the last interval as a feature, we did not think it worthwhile to pursue this route for this project. Therefore, we assumed that the prediction for the future load was correct. $E[SLA]$ then becomes as following, noting that k is the minimum number of nodes required for the system's power to be greater than the number of requests:

$$\begin{aligned} P(SLA|P, \rho, requests) &= P(power < requests|P, \rho) \\ &= P(|\text{non failing nodes}| < k) \\ &\approx CDFBinomial(P, k-1, 1-\rho) \end{aligned}$$

Note that the last approximation depends on the nodes' power being close to uniform. This is a reasonable assumption on most cases. The smart strategy, therefore, amounts to picking P such that the expected cost function presented in Equation 2 is minimized.

4 Experiments

For our experiments, we simulated the architecture on Figure 1 on a single machine, by having multiple servers on different ports, and

multiple clients with different threads. We made the time per prediction be constant, by having it be a sleep with $t = 0.1$ seconds (so we are not making real predictions, just simulating them). Finally, we set the time interval on which the autoscaler works to be one minute - so that the state of the system can potentially change every minute. The Autoscaler sends heartbeats to the Load Balancer, and finds out how many requests have been served every 5 seconds, in order to build a trend for load prediction. The Autoscaler sends heartbeats to the predictors every 30 seconds, in order to check (1), that they are still up, and (2), that their power remains the same (if not, it includes the change in the next state for smart load balancing).

Figures 3, 4 and 5 present the behaviour of the different load balancing strategies when faced with the loads previously introduced in Figure 2. The loads are now depicted with a granularity of 1 second, as opposed to 1 minute in Figure 2. The blue lines indicate the number of requests, and the green line indicates either the number of active nodes or whether or not an SLA violation ocurred. The vertical dotted lines indicate the change points - that is, the moments when the autoscaler changes (or not) the system's configuration, which happens every 60 seconds.

In Figures 3(a) and 3(b), we see the behaviour of the Naive strategy. As soon as SLA violations start to occur, it adds a new node into the system. However, once the new node is added, the average prediction time goes down dramatically, and the strategy proceeds to remove a node - which causes a lot of SLA violations, and thus subsequent additions of nodes. The Power Strategy (Figures 3(c) and 3(d)) does better, as it takes into account the power of the nodes. It is able to react to the increase in load, although it still suffers two SLA violations in the beginning. The Smart strategy with high SLA cost α presented in Figures 3(e) and 3(f) anticipates the rise in load, and adds enough nodes to handle it before it happens, suffering no SLA violations. In Figures 3(g) and 3(h), we set the cost of each node to be the same as the cost of an SLA violation. When this happens, the system has no incentive to add new nodes, and thus suffers SLA violations throughout the experiment. This shows how the smart strategy adapts to the costs α and β accordingly.

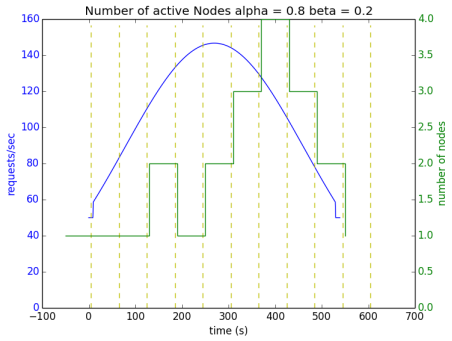
Figure 4 presentes a pathological case, where the load goes up and down abruptly right before or right after the change points. The Naive strategy is once again reactive. The Power strategy suffers the most from this, as it takes nodes up and down abruptly, and thus suffers the most SLA violations. The Smart strategy with high α is more conservative throughout, and suffers the least number of SLA violations.

Finally, Figure 5 presents a load pattern similar to the one in Figure 3, but duplicated. All strategies do the same as they did in the first load pattern. However, this figure shows how effective auto-scaling can be at managing the cost of having multiple nodes: the Smart strategy is able to suffer no SLA violations, while only using two nodes when the load demands it, and staying with one node when it doesn't.

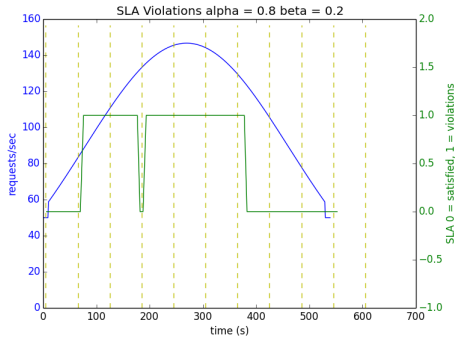
5 Conclusion

We have proposed an architecture and strategies for auto scaling online learning. In our methods and experiments, we abstracted out the online learning part, and presented and evaluated algorithms that are general to other auto scaling problems. We propose and evaluate a strategy that takes uncertainty and node failures into account, by modeling the cost function as an expectation.

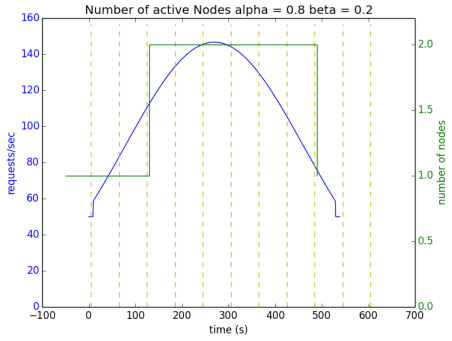
As future work, we would like to incorporate the machine learning aspect back into the system - possibly even coupling together learning and predicting (and thus modifying the cost function). This is an exciting area that is still new, and not a lot of research has been done on it.



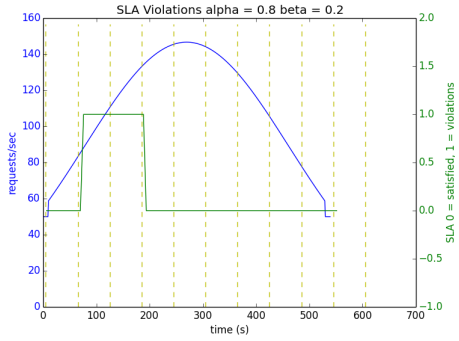
(a) Naive Strategy - active nodes



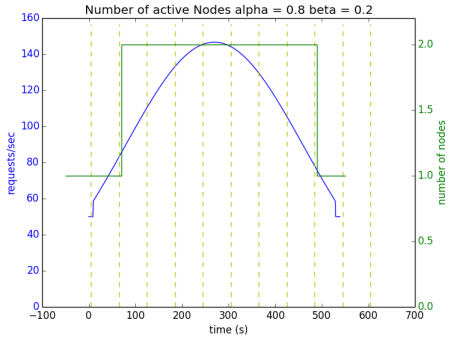
(b) Naive Strategy - SLA violations



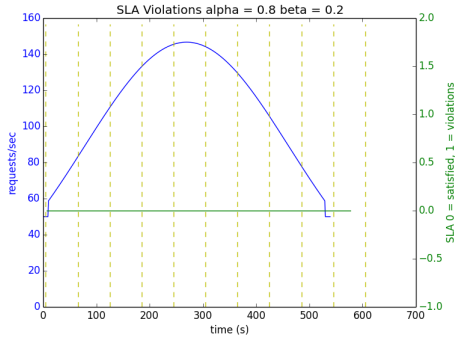
(c) Power Strategy - active nodes



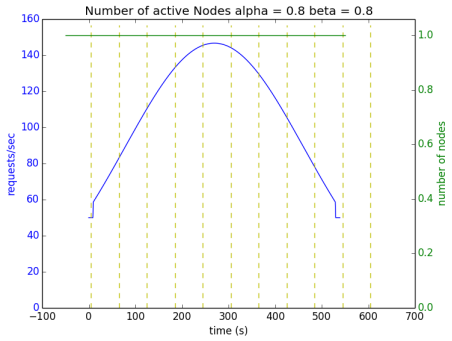
(d) Power Strategy - SLA violations



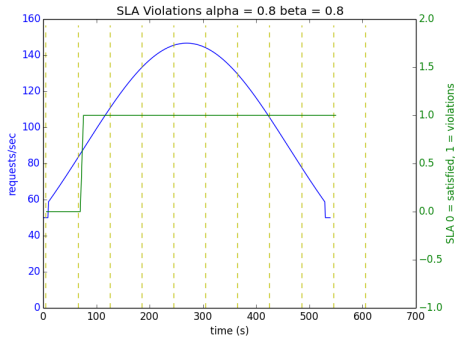
(e) Smart Strategy - active nodes



(f) Smart Strategy - SLA violations

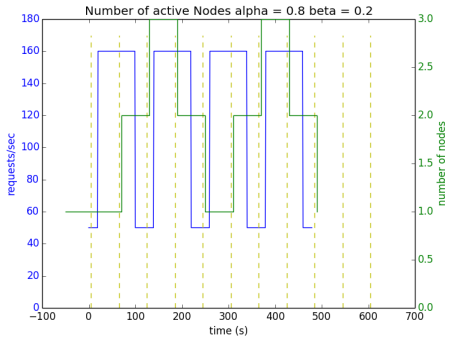


(g) Smart Strategy with high node cost - active nodes

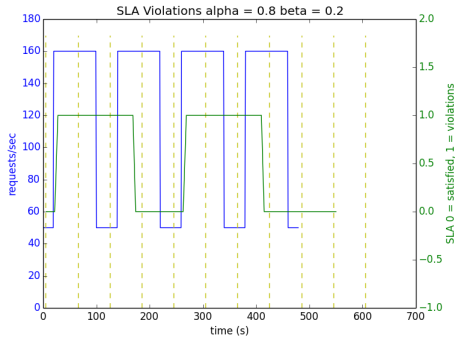


(h) Smart Strategy with high node cost - SLA violations

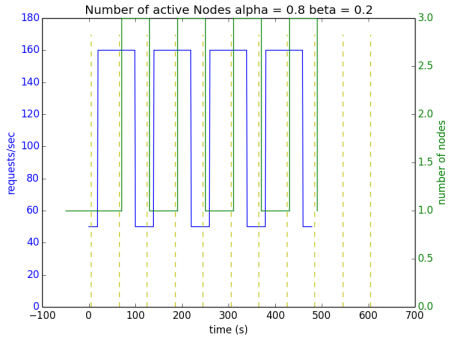
Figure 3: Load Pattern 1



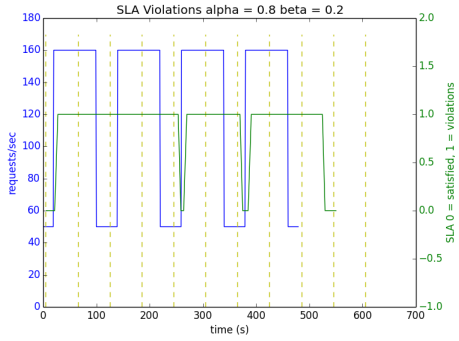
(a) Naive Strategy - active nodes



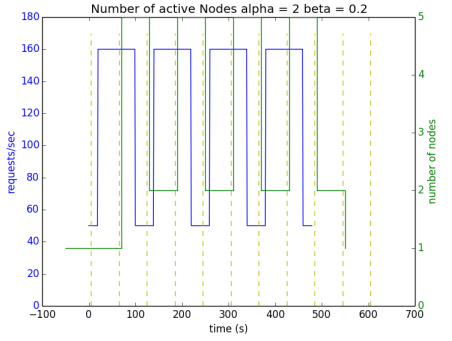
(b) Naive Strategy - SLA violations



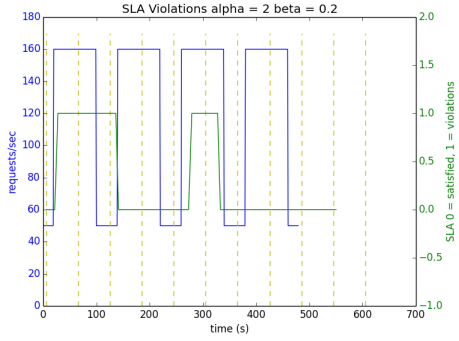
(c) Power Strategy - active nodes



(d) Power Strategy - SLA violations

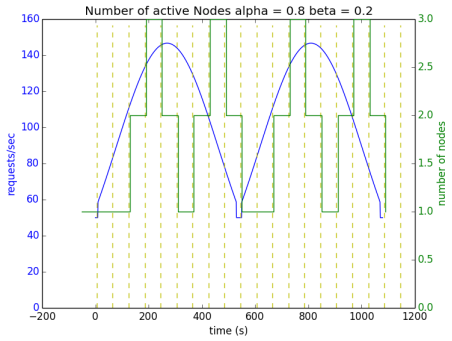


(e) Smart Strategy - active nodes

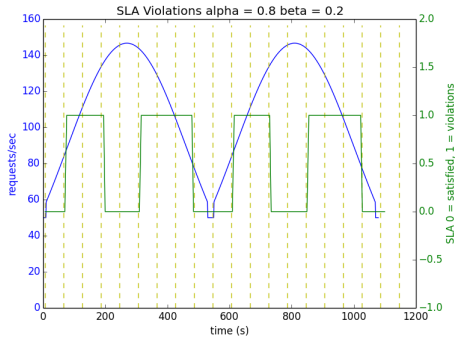


(f) Smart Strategy - SLA violations

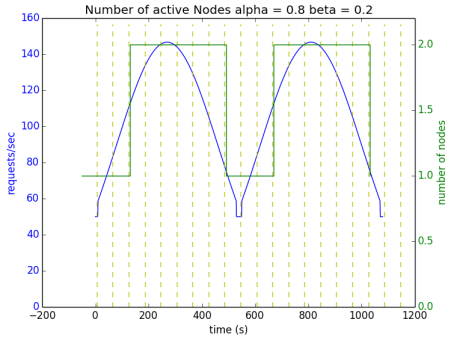
Figure 4: Load Pattern 2



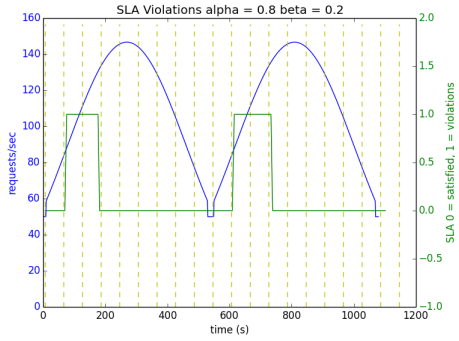
(a) Naive Strategy - active nodes



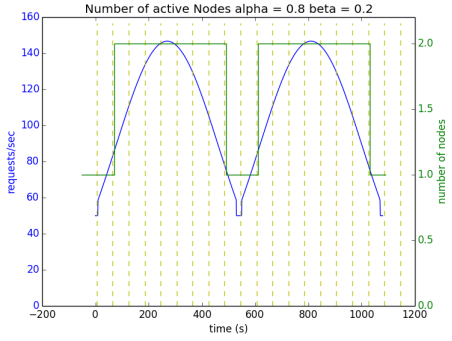
(b) Naive Strategy - SLA violations



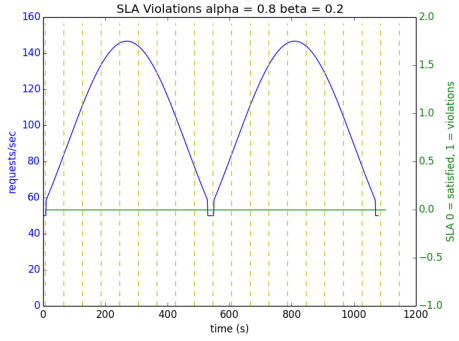
(c) Power Strategy - active nodes



(d) Power Strategy - SLA violations



(e) Smart Strategy - active nodes



(f) Smart Strategy - SLA violations

Figure 5: Load Pattern 3

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