



Flare-On 11 Challenge 7: fullspeed

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Overview

We can determine that the binary is a .NET native AOT binary since it has the .managed and .hydrated sections. For ease of reversing, we can take a memory dump after .hydrated is filled out with rehydrated data so we have access to strings.

By searching some of the strings used, we can figure out that this binary uses BouncyCastle for cryptography.

If we look at some of the more suspicious strings that are hex encoded, we see that they are passed to a function at 0x140107D80, which decrypts them with a simple XOR variant. The list of decrypted strings is below in Figure 1:

```
Unset
c90102faa48f18b5eac1f76bb40a1b9fb0d841712bbe3e5576a7a56976c2baeca47809765283aa078583e1e65172a
a 079 db 08 ea 2470350 c182487 b50 f7707 dd 46 a58 a1 d160 ff 79297 dcc 9b fad 6c fc 96 a81 c4 a 97564118 a4 0331 fe 0 fc 1326 fc 1426 fc 14
27f
9f939c02a7bd7fc263a4cce416f4c575f28d0c1315c4f0c282fca6709a5f9f7f9c251c9eede9eb1baa31602167fa5
087b5fe3ae6dcfb0e074b40f6208c8f6de4f4f0679d6933796d3b9bd659704fb85452f041fff14cf0e9aa7e45544f
127425c1d330ed537663e87459eaa1b1b53edfe305f6a79b184b3180033aab190eb9aa003e02e9dbf6d593c5e3b08
337
null
inf
inf
verify
verify failed
verify
null
too long
nul1
cd
ok
=== dirs ===
```





```
=== files ===
.
cat
exit
bad cmd
err
192.168.56.103;31337
;
err
```

Figure 1: decrypted strings

Looking at the xrefs to this function, we see a lot of interesting code. With extensive cross-referencing of strings and the <u>source of BouncyCastle</u> to figure out library functions, we can figure out that the binary connects to 192.168.56.103:31337 and performs an ECDH key exchange with a randomly generated 128-bit private key using the curve y2=x3+ax+b over $\mathbb{Z}/p\mathbb{Z}$ and the generator point Gx,Gy, where the following values are set from Figure 2:

```
Python

p =

0xc90102faa48f18b5eac1f76bb40a1b9fb0d841712bbe3e5576a7a56976c2baeca47809765283aa078583e1e6517
2a3fd

a =

0xa079db08ea2470350c182487b50f7707dd46a58a1d160ff79297dcc9bfad6cfc96a81c4a97564118a40331fe0fc
1327f

b =

0x9f939c02a7bd7fc263a4cce416f4c575f28d0c1315c4f0c282fca6709a5f9f7f9c251c9eede9eb1baa31602167f
a5380

gx =

0x087b5fe3ae6dcfb0e074b40f6208c8f6de4f4f0679d6933796d3b9bd659704fb85452f041fff14cf0e9aa7e4554
4f9d8

gy =

0x127425c1d330ed537663e87459eaa1b1b53edfe305f6a79b184b3180033aab190eb9aa003e02e9dbf6d593c5e3b
08182
```

Figure 2: Initial Generator values





Note that it's very unusual that the randomly generated private key is only 128 bits instead of the full 384 bits available! We'll see how this helps later.

Let's check what the order of the generator is using SageMath:

```
Python
sage: p =
0xc90102faa48f18b5eac1f76bb40a1b9fb0d841712bbe3e5576a7a56976c2baeca47809765283aa078583e1e6517
2a3fd
....: a =
0xa079db08ea2470350c182487b50f7707dd46a58a1d160ff79297dcc9bfad6cfc96a81c4a97564118a40331fe0fc
1327f
0x9f939c02a7bd7fc263a4cce416f4c575f28d0c1315c4f0c282fca6709a5f9f7f9c251c9eede9eb1baa31602167f
a5380
\dots: gx =
0x087b5fe3ae6dcfb0e074b40f6208c8f6de4f4f0679d6933796d3b9bd659704fb85452f041ffff14cf0e9aa7e4554
4f9d8
\dots: gy =
0x127425c1d330ed537663e87459eaa1b1b53edfe305f6a79b184b3180033aab190eb9aa003e02e9dbf6d593c5e3b
....: E = EllipticCurve(Zmod(p), [a,b])
\dots: gen = E(gx,gy)
....: factor(gen.order())
35809 * 46027 * 56369 * 57301 * 65063 * 111659 * 113111 *
```

The order is almost <u>smooth</u>! If it were smooth, then we could use <u>Pohlig-Hellman</u> to recover the private key, but unfortunately we still have a large prime factor in the order. However, we can still apply Pohlig-Hellman partially to gain information about the private key:

Recall that Pohlig-Hellman does the following (using multiplicative group notation):

- Given: generator g in a group of order n, element h = secret*g where secret < n is secret.
- For each prime (power) p in the factorization of the order:
 - o Compute g2 = (n/p)*g, h2 = (n/p)*g.
 - Solve the discrete logarithm problem using g2 as the generator and h2. This subgroup has order p.
 - The solution to the discrete logarithm problem is secret modulo p.
- Combine all the values for secret modulo p using the <u>Chinese remainder theorem</u> to recover secret.

We can apply these steps to the small prime factors of the order to get the private key modulo the product of those factors! If the product was larger than the private key, then this would just be the private key and we would be done. Unfortunately, we still come up short: the product is only 112 bits, not the full 128 bits we need.





However, note that 112 is not much smaller than 128. We can do a \sim 2^16 brute force to recover the private key as follows: we know the secret modulo n, where n is approximately 112 bits. We also know the secret is up to 128 bits. Then we know secret = secret_modulo_n + k*n, where k is at most 128-112 = 16 bits. We can then bruteforce k to get candidates for secret, and test those candidates by deriving the ChaCha20 keys/nonces and decrypting the encrypted traffic with them.

Full solution script:

```
Python
from sage.all import EllipticCurve, Zmod, factor, Factorization
from scapy.all import rdpcap, Raw
from Crypto.Cipher import ChaCha20
import hashlib
import base64
pcap = rdpcap("capture.pcapng")
raw = b''
for pkt in pcap:
   if Raw in pkt:
       raw += pkt[Raw].load
0xc90102faa48f18b5eac1f76bb40a1b9fb0d841712bbe3e5576a7a56976c2baeca47809765283aa078583e1e6517
2a3fd
a =
0xa079db08ea2470350c182487b50f7707dd46a58a1d160ff79297dcc9bfad6cfc96a81c4a97564118a40331fe0fc
1327f
0x9f939c02a7bd7fc263a4cce416f4c575f28d0c1315c4f0c282fca6709a5f9f7f9c251c9eede9eb1baa31602167f
a5380
genx =
0x087b5fe3ae6dcfb0e074b40f6208c8f6de4f4f0679d6933796d3b9bd659704fb85452f041ffff14cf0e9aa7e4554
4f9d8
geny =
0x127425c1d330ed537663e87459eaa1b1b53edfe305f6a79b184b3180033aab190eb9aa003e02e9dbf6d593c5e3b
08182
xorkey =
71337
E = EllipticCurve(Zmod(p), [a,b])
gen = E(genx, geny)
def unpack(x):
   return int.from_bytes(x, byteorder='big', signed=False)
```

Google Cloud



```
ax = unpack(raw[48*0:48*1]) ^ xorkey
ay = unpack(raw[48*1:48*2]) ^ xorkey
bx = unpack(raw[48*2:48*3]) ^ xorkey
by = unpack(raw[48*3:48*4]) ^ xorkey
raw = raw[48*4:]
A = E(ax, ay)
B = E(bx, by)
print('calculating order...')
order = gen.order()
print(f'{order=}')
factors = factor(order, proof=False)
smooth_part = Factorization(factors[:-1]).value()
big_part = Factorization(factors[-1:]).value()
# apply pohlig-hellman partially by dividing out the big factor
secret_mod_smooth = (gen*big_part).discrete_log(A*big_part)
print(f'{secret_mod_smooth=}')
possible_secret = secret_mod_smooth
while True:
    AB = possible_secret * B
    h = hashlib.sha512(int(AB[0]).to_bytes(48, 'big')).digest()
   key = h[:32]
    iv = h[32:40]
    cipher = ChaCha20.new(key=key, nonce=iv)
    test = cipher.decrypt(raw[:6])
    if test == b'verify':
        print('secret found:', possible_secret)
        cipher.seek(∅)
        lines = cipher.decrypt(raw).decode().strip('\x00').split('\x00')
        print(lines)
        print(base64.b64decode(lines[-2]))
        break
    possible_secret += smooth_part
```

Solution script output:

```
Unset calculating order... order=309373396510199458922447942662567138904409224558720519847625055617635267803116168639895 11376879697740787911484829297
```





```
secret_mod_smooth=3914004671535485983675163411331184
secret found: 168606034648973740214207039875253762473
['verify', 'verify', 'ls', '=== dirs ===\r\nsecrets\r\n=== files ===\r\nfullspeed.exe\r\n',
'cd|secrets', 'ok', 'ls', '=== dirs ===\r\nsuper secrets\r\n=== files ===\r\n', 'cd|super
secrets', 'ok', 'ls', '=== dirs ===\r\n.hidden\r\n=== files ===\r\n', 'cd|.hidden', 'ok',
'ls', "=== dirs ===\r\nwait, dot folders aren't hidden on windows\r\n=== files ===\r\n",
"cd|wait, dot folders aren't hidden on windows", 'ok', 'ls', '=== dirs ===\r\n=== files
===\r\nflag.txt\r\n', 'cat|flag.txt', 'RDBudF9VNWVfeTB1cl9Pd25fQ3VSdjNzQGZsYXJlLW9uLmNvbQ==',
'exit']
b'D0nt_U5e_y0ur_Own_CuRv3s@flare-on.com'
```

Final Flag

Unset

D0nt_U5e_y0ur_Own_CuRv3s@flare-on.com