# $Tools \ for \ primal \ degenerate \ linear \ programs \\ + \\ Dual-guided \ pivot \ rules \ for \ LP \\ + \\ Vector \ Space \ Decomposition \ for \ LP$

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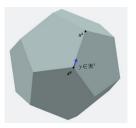
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What are the links between PS, DCA, IPS, and MMCC?

## Observation #1

## LP IN STANDARD FORM

$$\begin{array}{rcl} \min & \mathbf{c}^\mathsf{T} \mathbf{x} \\ & \mathbf{A} \mathbf{x} &= \mathbf{b} \\ & \mathbf{x} & \geq \mathbf{0} \end{array} \quad [\pi]$$



# From $x^0$ to $x^1$

- Find a *potential* improving direction  $\mathbf{y}^0 \in \mathbb{R}^n$ .
- **2** Determine step-size  $\rho^0 \in \mathbb{R}$ .
- **3** Compute  $x^1 := x^0 + \rho^0 y^0$ .

## Primal/dual formulations of the pricing in Primal Simplex

## $\overline{c_j} = 0, \forall j \in B$ : Pricing for $j \in N$ (non-basic variables)

Selection of an entering variable into basis  $\mathbf{A}_B$  relies on the minimum reduced cost of non-basic variables  $\boldsymbol{\pi}^\mathsf{T} = \mathbf{c}_B^\mathsf{T} \mathbf{A}_B^{-1}$   $\bar{c}_i = c_i - \boldsymbol{\pi}^\mathsf{T} \mathbf{a}_i, \ \forall j \in N.$ 

## PRIMAL/DUAL FORMULATIONS OF THE PRICING IN PRIMAL SIMPLEX

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## Find the minimum reduced cost value $\mu$ (Optimal if $\mu \geq 0$ )

$$\max \ \, \frac{\mu}{\mu} \quad \leq \quad c_j - \boldsymbol{\pi}^{\intercal} \mathbf{a}_j, \ \, \forall j \in N \qquad \quad [y_j]$$

\*\*\*  $\mu$  is the smallest reduced cost (given  $\pi$ ). \*\*\*

## PRIMAL/DUAL FORMULATIONS OF THE PRICING IN PRIMAL SIMPLEX

## $\overline{c_i} = 0, \forall j \in B$ : Pricing for $j \in N$ (non-basic variables)

Selection of an entering variable into basis  $A_B$  relies on the minimum reduced cost of non-basic variables  $\boldsymbol{\pi}^{\mathsf{T}} = \mathbf{c}_{\scriptscriptstyle R}^{\mathsf{T}} \mathbf{A}_{\scriptscriptstyle R}^{-1} \qquad \bar{c}_i = c_i - \boldsymbol{\pi}^{\mathsf{T}} \mathbf{a}_i, \ \forall i \in \mathcal{N}.$ 

## Find the minimum reduced cost value $\mu$ (Optimal if $\mu > 0$ )

 $\mu$  is the smallest reduced cost (given  $\pi$ ).

#### Equivalent to finding a convex combination of non-basic variables

TO FINDING A CONVEX COMBINATION OF NON-II
$$\mu = \min \sum_{j \in N} \bar{c}_j y_j$$

$$\sum_{j \in N} y_j = 1$$

$$y_j \geq 0, \ \forall j \in N$$

# Direction $\mathbf{y}^0 \in \mathbb{R}^n$

 $\mathbf{y}^{0}$ 

Entering variable only impacts the current basic variables, at most m variables. All but one non-basic variables remain at zero.

 $\mathbf{y}^0 \in \mathbb{R}^n$  makes changes on at most m+1 components.

# Direction $\mathbf{y}^0 \in \mathbb{R}^n$

 $y^0$ 

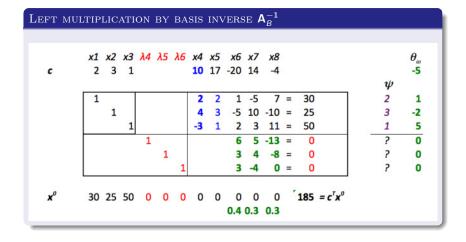
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Direction  $\mathbf{y}^0 \in \mathbb{R}^n$ : the selected entering variable, the non-selected non-basic variables, and the basic ones.

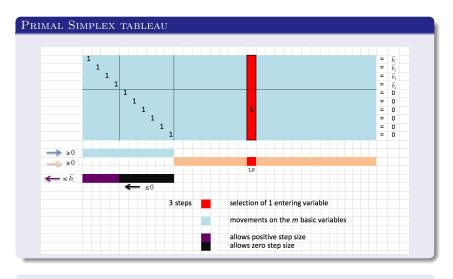
Step size computed such that  $\mathbf{x^1} := \mathbf{x^0} + \rho^0 \, \mathbf{y^0} \geq \mathbf{0}$ .

# Observation $\#\ 2$ : Degenerate solution on simplex-tableau

# Observation # 2: Degenerate solution on simplex-tableau

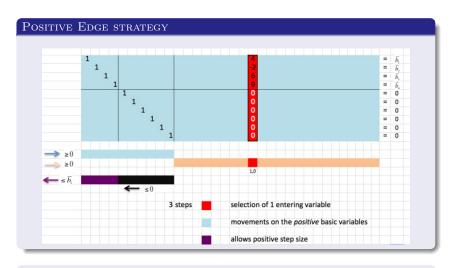


# Observation # 3: Structures



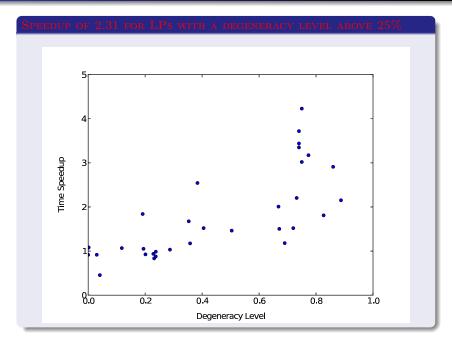
\*\*\* Changes on at most m+1 components. \*\*\*

# Observation # 3: Structures



\*\*\* Changes on at most p+1 components. \*\*\* Non-degenerate pivot.

# Positive Edge : identification in O(m)



## Positive Edge

The *Positive Edge* is a pricing rule for the Primal Simplex : it identifies variables  $x_j, j \notin B$  such that  $\bar{\mathbf{a}}_{Zj} = \mathbf{0}$  ... without computing  $\bar{\mathbf{a}}_{Zj}$ .

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If such a variable  $x_j$  has a negative reduced cost  $\bar{c}_j < 0$ , it strictly improves the objective function when entered into the basis.

# COMPUTATIONAL COMPLEXITY OF THE POSITIVE EDGE

Let  $\mathbf{v} \neq \mathbf{0}$  be a random vector of |Z| non-zero components.

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$$\mathbf{v}^{\mathsf{T}} \bar{\mathbf{a}}_{Zj} = \mathbf{v}^{\mathsf{T}} \ [\mathbf{A}_{B}^{-1}]_{Z} \ \mathbf{a}_{Zj} = \mathbf{v}^{\mathsf{T}} \ [\mathbf{A}_{B}^{-1}]_{Z} \ \mathbf{a}_{Zj}$$

Pre-compute once  $\mathbf{w}^{\mathsf{T}} := \mathbf{v}^{\mathsf{T}} [\mathbf{A}_{B}^{-1}]_{Z}$ , then compute  $\mathbf{w}^{\mathsf{T}} \mathbf{a}_{Zj}$ , on the **original data**.

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Determining if  $x_j$  satisfies the Positive Edge rule is done in O(m), as for  $\bar{c}_j = c_j - \pi^{\mathsf{T}} \mathbf{a}_j, j \notin B$ , where  $\pi = \mathbf{c}_B^{\mathsf{T}} A_B^{-1}$ .

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On large-scale instances ( $m \approx 100~000$ ,  $n \approx 450~000$ ), cpu time to identify all variables satisfying the PE rule : usual definition Positive Edge  $O(m^2n)$  O(mn)

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On large-scale instances ( $m \approx 100~000$ ,  $n \approx 450~000$ ), cpu time to identify all variables satisfying the PE rule : usual definition Positive Edge  $O(m^2n)$  O(mn) 2500 seconds

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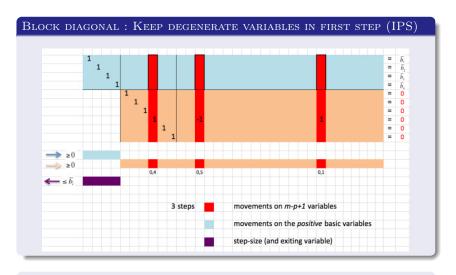
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On large-scale instances ( $m \approx 100~000$ ,  $n \approx 450~000$ ), cpu time to identify all variables satisfying the PE rule : usual definition Positive Edge  $O(m^2n)$  O(mn) 2500 seconds 0.5 seconds

# Observation #3: Structures



\*\*\* Changes on at most m+1 components. \*\*\*

## NOTATION

Vectors and matrices are written in **bold face**.

 $I_{\ell}$ : the  $\ell \times \ell$  identity matrix.

0 (1): a vector/matrix with all zeros (ones) entries of appropriate dimensions.

 $\mathbf{A}_{RC}$ : sub-matrix of  $\mathbf{A}$  containing the rows and columns indexed by R and C.

Standard notation : basis  $\mathbf{A}_B$ , inverse  $\mathbf{A}_B^{-1}$ ,  $\mathbf{c}_B^{\mathsf{T}}\mathbf{x}_B$ ,  $\mathbf{A}_B\mathbf{x}_B$ ,  $\boldsymbol{\pi} = \mathbf{c}_B\mathbf{A}_B^{-1}$  ...  $\mathbf{I}_F < \mathbf{x}_F < \mathbf{u}_F$ ,  $\mathbf{x}_L = \mathbf{I}_L$ ,  $\mathbf{x}_U = \mathbf{u}_U$ 

## Useful decomposition of $\mathbf{x} \in \mathbb{R}^n$ in $\mathbf{A}\mathbf{x} = \mathbf{b}, \ \ \mathbf{I} \leq \mathbf{x} \leq \mathbf{u}$

$$\mathbf{x} = \begin{bmatrix} \mathbf{x}_F \\ \mathbf{x}_L \\ \mathbf{x}_U \end{bmatrix} \qquad \mathbf{x} = \begin{bmatrix} \mathbf{x}_B \\ \mathbf{x}_N \end{bmatrix} = \begin{bmatrix} \mathbf{x}_F \\ \mathbf{x}_{B_L} \\ \mathbf{x}_{B_U} \\ \mathbf{x}_{N_L} \\ \mathbf{x}_{N_U} \end{bmatrix}$$

For 
$$\emptyset \subseteq S \subseteq B : \mathbf{x} = \begin{bmatrix} \mathbf{x}_S \\ \mathbf{x}_{\bar{S}} \end{bmatrix} = \begin{bmatrix} \mathbf{x}_{S_F} \\ \mathbf{x}_{S_L} \\ \mathbf{x}_{\bar{S}_U} \\ \mathbf{x}_{\bar{S}_U} \\ \mathbf{x}_{\bar{S}_U} \\ \mathbf{x}_{\bar{S}_U} \end{bmatrix} = \begin{bmatrix} \mathbf{x}_F \\ \mathbf{x}_L \\ \mathbf{x}_U \end{bmatrix}$$

# Dual Guided Pivot Rules for LPs

## LINEAR PROGRAM LP

$$z^* := \min c^T x$$
  
st.  $Ax = b$   
 $I \le x \le u$ 

# Dual Guided Pivot Rules for LPs

#### Linear program LP

$$\begin{split} z^{\star} := & & \text{min} & & c^{\mathsf{T}} x \\ \text{st.} & & & A x & = b \\ & & I \leq x & \leq u \end{split}$$

## Generic Algorithm with single parameter set S, $\emptyset \subseteq S \subseteq B$

- Let k = 0 and assume a feasible basic solution  $\mathbf{x}^k$  to LP.
- **2** For  $\emptyset \subseteq S \subseteq B$ , construct the residual problem  $LP_S(\mathbf{x}^k)$ .
- **3** Fix a subset of the dual variables in row-set  $R: \pi_R^{\mathsf{T}} = \mathbf{c}_S^{\mathsf{T}} \mathbf{A}_{RS}^{-1}$ .
- Determine the value of the smallest reduced cost  $\mu_s^k$ . If  $\mu_s^k \ge 0$ , STOP. Current solution  $\mathbf{x}^k$  is optimal for LP.
- **3** Retrieve direction  $\mathbf{y}_S^k \in \mathbb{R}^n$  and compute its maximum step-size  $\rho_S^k$ .
- Update  $x^{k+1} := x^k + \rho_S^k y_S^k;$   $z^{k+1} := z^k + \rho_S^k \mu_S^k;$ k := k + 1.
- Goto Step 2.

# Step 2 Construct residual problem $LP(\mathbf{x}^k)$

Assume a feasible solution 
$$\mathbf{x}^k = \begin{bmatrix} \mathbf{x}_F^k \\ \mathbf{x}_L^k \\ \mathbf{x}_U^k \end{bmatrix}$$

$$\mathbf{x} := \mathbf{x}^k + \mathbf{y}, \quad \mathbf{y} \in \mathbb{R}^n$$

# Step 2 Construct residual problem $LP(x^k)$

Assume a feasible solution 
$$\mathbf{x}^k = \begin{bmatrix} \mathbf{x}_F^2 \\ \mathbf{x}_L^k \\ \mathbf{x}_U^k \end{bmatrix}$$

$$\begin{split} \mathbf{x} &:= \mathbf{x}^k + \mathbf{y}, \quad \mathbf{y} \in \mathbb{R}^n \\ &= \mathbf{x}^k + (\vec{\mathbf{y}} - \mathbf{\ddot{y}}), & \quad \vec{\mathbf{y}}, \mathbf{\ddot{y}} \geq \mathbf{0}, \quad \vec{\mathbf{y}}^\mathsf{T} \mathbf{\ddot{y}} = \mathbf{0}, \quad \vec{\mathbf{y}} \leq \vec{\mathbf{r}}^k, \quad \mathbf{\ddot{y}} \leq \mathbf{\ddot{r}}^k \end{split}$$

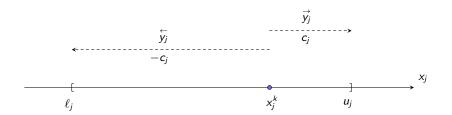
# STEP 2 CONSTRUCT RESIDUAL PROBLEM $LP(\mathbf{x}^k)$

Assume a feasible solution 
$$\mathbf{x}^k = \left[ egin{array}{c} \mathbf{x}_F^k \\ \mathbf{x}_L^k \\ \mathbf{x}_U^k \end{array} \right]$$

 $\mathbf{x} := \mathbf{x}^k + \mathbf{y}, \quad \mathbf{y} \in \mathbb{R}^n$ 

$$= \mathbf{x}^{k} + (\vec{\mathbf{y}} - \mathbf{\bar{y}}), \qquad \qquad \vec{\mathbf{y}}, \mathbf{\bar{y}} \ge \mathbf{0}, \quad \vec{\mathbf{y}}^{\mathsf{T}} \mathbf{\bar{y}} = 0, \quad \vec{\mathbf{y}} \le \mathbf{\bar{r}}^{k}, \quad \mathbf{\bar{y}} \le \mathbf{\bar{r}}^{k}$$

$$= \mathbf{x}^{k} + \left( \begin{bmatrix} \vec{\mathbf{y}}_{F} \\ \vec{\mathbf{y}}_{L} \\ \vec{\mathbf{y}}_{U} \end{bmatrix} - \begin{bmatrix} \mathbf{\bar{y}}_{F} \\ \mathbf{\bar{y}}_{L} \\ \mathbf{\bar{y}}_{U} \end{bmatrix} \right); \quad \vec{\mathbf{y}}, \mathbf{\bar{y}} \ge \mathbf{0}, \quad \vec{\mathbf{y}}^{\mathsf{T}} \mathbf{\bar{y}} = 0, \quad \vec{\mathbf{y}} \le \mathbf{\bar{r}}^{k}, \quad \mathbf{\bar{y}} \le \mathbf{\bar{r}}^{k}$$



# STEP 2 CONSTRUCT RESIDUAL PROBLEM $LP(\mathbf{x}^k)$

Residual problem 
$$LP(\mathbf{x}^k)$$
;  $\mathbf{x} := \mathbf{x}^k + (\vec{\mathbf{y}} - \mathbf{\ddot{y}})$  (change of variables)

$$z^* := \min \quad \mathbf{c}^{\mathsf{T}} \mathbf{x}^k \quad + \quad \mathbf{c}^{\mathsf{T}} (\vec{\mathbf{y}} - \mathbf{\bar{y}})$$

$$\text{st.} \quad \mathbf{A} \mathbf{x}^k \quad + \quad \mathbf{A} (\vec{\mathbf{y}} - \mathbf{\bar{y}}) \quad = \quad \mathbf{b}$$

$$\mathbf{0} \le \vec{\mathbf{y}} \le \vec{\mathbf{r}}^k$$

$$\mathbf{0} \le \mathbf{\bar{y}} \le \mathbf{\bar{r}}^k$$

# STEP 2 CONSTRUCT RESIDUAL PROBLEM $LP(\mathbf{x}^k)$

Residual problem 
$$LP(\mathbf{x}^k)$$
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$$\begin{split} z^{\star} &:= \min \quad \mathbf{c}^{\mathsf{T}} \mathbf{x}^{k} \quad + \quad \mathbf{c}^{\mathsf{T}} (\vec{\mathbf{y}} - \mathbf{\tilde{y}}) \\ &\text{st.} \quad \mathbf{A} \mathbf{x}^{k} \quad + \quad \mathbf{A} (\vec{\mathbf{y}} - \mathbf{\tilde{y}}) \quad = \quad \mathbf{b} \\ &\mathbf{0} \leq \vec{\mathbf{y}} \leq \mathbf{\tilde{r}}^{k} \\ &\mathbf{0} \leq \mathbf{\tilde{y}} \leq \mathbf{\tilde{r}}^{k} \end{split}$$

### RESIDUAL PROBLEM $LP(\mathbf{x}^k)$

$$z^\star := \mathbf{c}^\intercal \mathbf{x}^k + \min$$
  $\mathbf{c}^\intercal (\vec{\mathbf{y}} - \mathbf{\bar{y}})$  st.  $\mathbf{A} (\vec{\mathbf{y}} - \mathbf{\bar{y}}) = \mathbf{0}$   $\mathbf{0} \le \vec{\mathbf{y}} \le \vec{\mathbf{r}}^k$   $\mathbf{0} < \mathbf{\bar{y}} < \mathbf{\bar{r}}^k$ 

# STEP 2 CONSTRUCT RESIDUAL PROBLEM $LP(x^k)$

RESIDUAL PROBLEM WITH 
$$S = F$$
,  $\bar{S} = L \cup U$ 

$$z^{\star} := \mathbf{c}^{\mathsf{T}} \mathbf{x}^{k} +$$

$$\min \quad \mathbf{c}_{F}^{\mathsf{T}} (\vec{\mathbf{y}}_{F} - \mathbf{\bar{y}}_{F}) + \mathbf{c}_{L}^{\mathsf{T}} (\vec{\mathbf{y}}_{L}) - \mathbf{c}_{U}^{\mathsf{T}} (\mathbf{\bar{y}}_{U})$$

$$\text{st.} \quad \mathbf{A}_{F} (\vec{\mathbf{y}}_{F} - \mathbf{\bar{y}}_{F}) + \mathbf{A}_{L} (\vec{\mathbf{y}}_{L}) - \mathbf{A}_{U} (\mathbf{\bar{y}}_{U}) = \mathbf{0}$$

$$\vec{\mathbf{y}}_{F} \geq \mathbf{0}, \ \vec{\mathbf{y}}_{F} \geq \mathbf{0}, \qquad \vec{\mathbf{y}}_{L} \geq \mathbf{0}, \qquad \mathbf{\bar{y}}_{U} \geq \mathbf{0}$$

$$\vec{\mathbf{y}}_{F} \leq \vec{\mathbf{r}}_{F}, \ \vec{\mathbf{y}}_{F} \leq \mathbf{\bar{r}}_{F}, \qquad \vec{\mathbf{y}}_{L} \leq \vec{\mathbf{r}}_{L}, \qquad \mathbf{\bar{y}}_{U} \leq \mathbf{\bar{r}}_{U}$$

# STEP 2 CONSTRUCT RESIDUAL PROBLEM $LP(x^k)$

### Residual problem – General case : $\emptyset \subseteq S \subseteq B$

$$\begin{array}{lll} \mathbf{c}^\intercal \mathbf{x}^k + & & \min \ \mathbf{c}_S^\intercal (\vec{y}_S - \ddot{\mathbf{y}}_S) \ & & \\ & \text{st.} & \mathbf{A}_S (\vec{y}_S - \ddot{\mathbf{y}}_S) \ & & \\ & & \vec{y}_S \geq \mathbf{0}, \ \ddot{\mathbf{y}}_S \geq \mathbf{0} \\ & & & \ddot{\mathbf{y}}_S \geq \mathbf{0}, \ \ddot{\mathbf{y}}_S \geq \mathbf{0} \\ & & & \ddot{\mathbf{y}}_{\bar{S}_L} \leq \mathbf{0}, \ \ddot{\mathbf{y}}_{\bar{S}_U} \leq \mathbf{0} \\ & & & \ddot{\mathbf{y}}_{\bar{S}_L} \leq \mathbf{0}, \ \ddot{\mathbf{y}}_{\bar{S}_U} \leq \mathbf{0} \\ & & & \ddot{\mathbf{y}}_{\bar{S}_L} \leq \mathbf{0}, \ \ddot{\mathbf{y}}_{\bar{S}_U} \leq \mathbf{0} \\ & & \ddot{\mathbf{y}}_{\bar{S}_L} \leq \mathbf{0}, \ \ddot{\mathbf{y}}_{\bar{S}_U} \leq \mathbf{0} \\ & & \ddot{\mathbf{y}}_{\bar{S}_L} \leq \mathbf{0}, \ \ddot{\mathbf{y}}_{\bar{S}_U} \leq \mathbf{0} \\ & & \ddot{\mathbf{y}}_{\bar{S}_L} \leq \mathbf{0}, \ \ddot{\mathbf{y}}_{\bar{S}_U} \leq \mathbf{0} \\ & & \ddot{\mathbf{y}}_{\bar{S}_L} \leq \mathbf{0}, \ \ddot{\mathbf{y}}_{\bar{S}_U} \leq \mathbf{0} \\ \end{array}$$

$$S = S_F \cup S_L \cup S_U; \qquad \bar{S} = \bar{S}_F \cup \bar{S}_L \cup \bar{S}_U$$

# STEP 2 CONSTRUCT RESIDUAL PROBLEM $LP(\mathbf{x}^k)$

RESIDUAL PROBLEM WITH 
$$S = B$$
 (BASIC),  $\bar{S} = N$  (NON-BASIC)

$$c^{\mathsf{T}}x^k + \min c^{\mathsf{T}}_B(\vec{y}_B - \ddot{y}_B) + c^{\mathsf{T}}_N(\vec{y}_N - \ddot{y}_N) \\
\text{st.} \quad A_B(\vec{y}_B - \ddot{y}_B) + A_N(\vec{y}_N - \ddot{y}_N) = 0 \\
\vec{y}_B \ge 0, \ \ddot{y}_B \ge 0 \qquad \vec{y}_N \ge 0, \ \ddot{y}_N \ge 0 \\
\vec{y}_{N_L} \le 0, \ \vec{y}_{N_U} \le 0$$

$$\vec{y}_F \le \vec{r}_F, \ \ddot{y}_F \le \ddot{r}_F \\
\ddot{y}_{B_L} \le 0, \ \vec{y}_{B_U} \le 0 \qquad \vec{y}_{N_L} \le \vec{r}_{N_L}, \ \ddot{y}_{N_U} \le \ddot{r}_{N_U}$$

 $B = F \cup B_I \cup B_{II}$ ;  $N = N_I \cup N_{II}$ 

#### STEP 3 FIX SOME DUAL VARIABLES IN ROW-SET R

For 
$$\emptyset \subseteq S \subseteq B$$
, find  $\mathbf{y}^k = \begin{bmatrix} \mathbf{y}_S^k \\ \mathbf{y}_S^k \end{bmatrix} = \begin{bmatrix} (\mathbf{\vec{y}}_S^k - \mathbf{\bar{y}}_S^k) \\ (\mathbf{\vec{y}}_S^k - \mathbf{\bar{y}}_S^k) \end{bmatrix}$  of min reduced cost  $\mu^k$ .

For basic columns  $A_S$ , select a  $s \times s$  square matrix  $A_{RS}$ , a set of s independent rows.

$$\mathbf{T} = \begin{bmatrix} \mathbf{A}_{RS} & \mathbf{0} \\ \mathbf{A}_{ZS} & \mathbf{I}_{m-s} \end{bmatrix} \quad \mathbf{T}^{-1} = \begin{bmatrix} \mathbf{A}_{RS}^{-1} & \mathbf{0} \\ -\mathbf{A}_{ZS}\mathbf{A}_{RS}^{-1} & \mathbf{I}_{m-s} \end{bmatrix}$$

#### STEP 3 FIX SOME DUAL VARIABLES IN ROW-SET R

$$\emptyset \subset S \subset B$$

$$\begin{array}{lll} \text{Perform } \mathbf{T}^{-1} \mathbf{A}(\vec{\mathbf{y}} - \vec{\mathbf{y}}) = \mathbf{0}. & \text{General case} & \boldsymbol{\pi}^{\mathsf{T}} = \boldsymbol{\psi}^{\mathsf{T}} \mathbf{T}^{-1} \\ \mathbf{c}^{\mathsf{T}} \mathbf{x}^k + \min \ \mathbf{c}_{S}^{\mathsf{T}} (\vec{\mathbf{y}}_S - \vec{\mathbf{y}}_S) \ & \\ \text{st.} & (\vec{\mathbf{y}}_S - \vec{\mathbf{y}}_S) \ & \\ \mathbf{x}_S \cdot (\vec{\mathbf{y}}_S -$$

Observe  $\bar{c}_S = 0$  (basic variables).

#### STEP 3 FIX SOME DUAL VARIABLES IN ROW-SET R

Perform 
$$\mathbf{T}^{-1}\mathbf{A}(\vec{\mathbf{y}} - \mathbf{\ddot{\mathbf{y}}}) = \mathbf{0}$$
. General case : details  $\boldsymbol{\pi}^{\mathsf{T}} = \boldsymbol{\psi}^{\mathsf{T}}\mathbf{T}^{-1}$ 

$$\mathbf{c}^{\mathsf{T}}\mathbf{x}^k + \min \ \mathbf{c}^{\mathsf{T}}_{S}(\vec{\mathbf{y}}_{S} - \mathbf{\ddot{\mathbf{y}}}_{S}) + \mathbf{c}^{\mathsf{T}}_{S}(\vec{\mathbf{y}}_{\bar{S}} - \mathbf{\ddot{\mathbf{y}}}_{\bar{S}})$$
st.  $(\vec{\mathbf{y}}_{S} - \mathbf{\ddot{\mathbf{y}}}_{S}) + \mathbf{A}_{R\bar{S}}(\vec{\mathbf{y}}_{\bar{S}} - \mathbf{\ddot{\mathbf{y}}}_{\bar{S}}) = \mathbf{0} \ [\boldsymbol{\psi}_{R} = \mathbf{c}_{S}]$ 

$$\mathbf{\ddot{\mathbf{y}}}_{S} \geq \mathbf{0}, \ \mathbf{\ddot{\mathbf{y}}}_{S} \geq \mathbf{0}$$

$$\mathbf{\ddot{\mathbf{A}}}_{Z\bar{S}}(\vec{\mathbf{y}}_{\bar{S}} - \mathbf{\ddot{\mathbf{y}}}_{\bar{S}}) = \mathbf{0} \ [\boldsymbol{\psi}_{Z} \ unknown]$$

$$\mathbf{\ddot{\mathbf{y}}}_{\bar{S}} \geq \mathbf{0}, \ \mathbf{\ddot{\mathbf{y}}}_{\bar{S}} \geq \mathbf{0}$$

$$\mathbf{\ddot{\mathbf{y}}}_{\bar{S}_{L}} \leq \mathbf{0}, \ \mathbf{\ddot{\mathbf{y}}}_{\bar{S}_{U}} \leq \mathbf{0}$$

$$\mathbf{\ddot{\mathbf{y}}}_{S_{F}} \leq \mathbf{\ddot{\mathbf{r}}}_{S_{F}}, \ \mathbf{\ddot{\mathbf{y}}}_{S_{F}} \leq \mathbf{\ddot{\mathbf{r}}}_{S_{F}}$$

$$\mathbf{\ddot{\mathbf{y}}}_{S_{L}} \leq \mathbf{0}, \ \mathbf{\ddot{\mathbf{y}}}_{S_{U}} \leq \mathbf{0}$$

$$\mathbf{\ddot{\mathbf{y}}}_{\bar{S}_{L}} \leq \mathbf{\ddot{\mathbf{r}}}_{\bar{S}_{L}}, \ \mathbf{\ddot{\mathbf{y}}}_{\bar{S}_{U}} \leq \mathbf{\ddot{\mathbf{r}}}_{\bar{S}_{U}}$$

 $S = S_F \cup S_L \cup S_U$ :  $\bar{S} = \bar{S}_F \cup \bar{S}_U \cup \bar{S}_U$ 

# Step 4 Find minimum reduced cost $\mu_S^k$

#### Pricing of the variables

$$S \subseteq B$$
, hence  $\overline{c}_S = 0$ .

Therefore pricing of  $\vec{y}_{\bar{S}}$  and  $\vec{y}_{\bar{S}}$  needed to get partial direction  $(\vec{y}_{\bar{S}}^k - \vec{y}_{\bar{S}}^k)$ 

followed by impact on  $(\vec{\mathbf{y}}_{S}^{k} - \vec{\mathbf{y}}_{S}^{k})$  to complete direction  $\mathbf{y}^{k} = \begin{bmatrix} \mathbf{y}_{S}^{k} \\ \mathbf{y}_{S}^{k} \end{bmatrix}$ .

#### PRIMAL/DUAL FORMULATIONS OF THE PRICING

$$\boldsymbol{\psi}_{\mathit{R}}^{\intercal} = \mathbf{c}_{\mathit{S}}^{\intercal} \mathbf{A}_{\mathit{RS}}^{-1}$$

$$\begin{array}{lll} \max \; \mu & \text{st. } \mu \mathbf{1}^\intercal & \leq & \mathbf{c}_{\bar{\mathsf{S}}}^\intercal - \boldsymbol{\psi}_R^\intercal \mathsf{A}_{R\bar{\mathsf{S}}} - \boldsymbol{\psi}_Z^\intercal \bar{\mathsf{A}}_{Z\bar{\mathsf{S}}} & \left[ \vec{\boldsymbol{\mathsf{y}}}_{\bar{\mathsf{S}}} \right] \\ & \mu \mathbf{1}^\intercal & \leq - (\mathbf{c}_{\bar{\mathsf{S}}}^\intercal - \boldsymbol{\psi}_R^\intercal \mathsf{A}_{R\bar{\mathsf{S}}} - \boldsymbol{\psi}_Z^\intercal \bar{\mathsf{A}}_{Z\bar{\mathsf{S}}}) & \left[ \mathbf{\bar{\mathsf{y}}}_{\bar{\mathsf{S}}} \right] \end{array}$$

### A convex combination of the variables $\vec{y}_{\bar{\varsigma}}$ and $\dot{\bar{y}}_{\bar{\varsigma}}$

\*\*\* Optimal solution :  $\mu_S^k$ ,  $\vec{\mathbf{y}}_{\bar{5}}^k$  and  $\dot{\mathbf{y}}_{\bar{5}}^k$ . \*\*\*

# Step 5 Retrieve direction $\mathbf{y}_{\mathsf{S}}^k$ of minimum reduced cost $\mu_{\mathsf{S}}^k$

### ... Complete direction $\mathbf{y}_{5}^{k}$

Given  $\mu^k < 0$  and  $(\vec{\mathbf{y}}_{\bar{S}}^k - \vec{\mathbf{y}}_{\bar{S}}^k)$ , find impacts on other variables  $(\vec{\mathbf{y}}_S^k - \vec{\mathbf{y}}_S^k)$ .

$$\begin{aligned} & (\vec{y}_S - \vec{y}_S) \quad + \quad \bar{\mathbf{A}}_{R\bar{S}} (\vec{\mathbf{y}_S^k} - \vec{\mathbf{y}_S^k}) \quad = \quad \mathbf{0} \\ \vec{\mathbf{y}}_S \geq \mathbf{0}, \ \, \mathbf{\tilde{y}}_S \geq \mathbf{0} \qquad \qquad \qquad \vec{y}_j \, \mathbf{\tilde{y}}_j = \mathbf{0}, \quad \forall j \in S_F \end{aligned}$$

## DIRECTION $\mathbf{y}_{s}^{k}$

$$\mathbf{y}_{S}^{k} = \begin{bmatrix} (\ddot{\mathbf{y}}_{S}^{k} - \ddot{\mathbf{y}}_{S}^{k}) \\ (\ddot{\mathbf{y}}_{S}^{k} - \ddot{\mathbf{y}}_{S}^{k}) \end{bmatrix} = \begin{bmatrix} -\bar{\mathbf{A}}_{R\bar{S}}(\ddot{\mathbf{y}}_{S}^{k} - \ddot{\mathbf{y}}_{S}^{k}) \\ (\ddot{\mathbf{y}}_{S}^{k} - \ddot{\mathbf{y}}_{S}^{k}) \end{bmatrix} = \begin{bmatrix} Impact \ level \ Pricing \ level \end{bmatrix}$$

Positive part; Negative part.

# Step 5 Compute maximum step-size $\rho_{\mathsf{S}}^k$

$$\rho \left[ \begin{array}{c} \vec{\mathbf{y}}_{S_F}^k \\ \vec{\mathbf{y}}_{S_L}^k \\ \vec{\mathbf{y}}_{S_U} \end{array} \right] \leq \left[ \begin{array}{c} \vec{\mathbf{r}}_{S_F}^k \\ \vec{\mathbf{r}}_{S_L}^k \\ \mathbf{0} \end{array} \right]; \quad \rho \left[ \begin{array}{c} \vec{\mathbf{y}}_{S_F}^k \\ \vec{\mathbf{y}}_{S_L}^k \\ \vec{\mathbf{y}}_{S_U}^k \end{array} \right] \leq \left[ \begin{array}{c} \vec{\mathbf{r}}_{S_F}^k \\ \mathbf{0} \\ \vec{\mathbf{r}}_{S_U}^k \end{array} \right]$$

$$\rho \left[ \begin{array}{c} \vec{\mathbf{y}}_{\bar{S}_F}^k \\ \vec{\mathbf{y}}_{\bar{S}_L}^k \\ \mathbf{0} \end{array} \right] \leq \left[ \begin{array}{c} \vec{\mathbf{r}}_{\bar{S}_F}^k \\ \vec{\mathbf{r}}_{\bar{S}_L}^k \\ \mathbf{0} \end{array} \right]; \quad \rho \left[ \begin{array}{c} \vec{\mathbf{y}}_{\bar{S}_F}^k \\ \mathbf{0} \\ \vec{\mathbf{y}}_{\bar{S}_U}^k \end{array} \right] \leq \left[ \begin{array}{c} \vec{\mathbf{r}}_{\bar{S}_F}^k \\ \mathbf{0} \\ \vec{\mathbf{r}}_{\bar{S}_U}^k \end{array} \right]$$

10 out of 12 types of residual upper bonds to verify.  $\emptyset \subseteq S \subseteq B$ 

# Step 5 Compute maximum step-size $\rho_{\mathsf{S}}^k$

$$\rho \left[ \begin{array}{c} \vec{\mathbf{y}}_{F}^{k} \\ \vec{\mathbf{y}}_{B_{L}}^{k} \\ \vec{\mathbf{y}}_{B_{U}} \end{array} \right] \leq \left[ \begin{array}{c} \vec{\mathbf{r}}_{F}^{k} \\ \vec{\mathbf{r}}_{B_{L}}^{k} \\ \mathbf{0} \end{array} \right]; \quad \rho \left[ \begin{array}{c} \vec{\mathbf{y}}_{F}^{k} \\ \vec{\mathbf{y}}_{B_{L}}^{k} \\ \vec{\mathbf{y}}_{B_{U}}^{k} \end{array} \right] \leq \left[ \begin{array}{c} \vec{\mathbf{r}}_{F}^{k} \\ \mathbf{0} \\ \vec{\mathbf{r}}_{B_{U}}^{k} \end{array} \right]$$

$$\rho \left[ \begin{array}{c} \vec{\mathbf{y}}_{N_L}^k \end{array} \right] \leq \left[ \begin{array}{c} \vec{\mathbf{r}}_{N_L}^k \end{array} \right]; \quad \rho \left[ \begin{array}{c} \overline{\mathbf{y}}_{N_U}^k \end{array} \right] \leq \left[ \begin{array}{c} \overline{\mathbf{r}}_{N_U}^k \end{array} \right]$$

8 types of residual upper bonds to verify for Primal Simplex. S =

# Step 5 Compute maximum step-size $\rho_S^k$

$$\rho \left[ \begin{array}{c} \mathbf{\ddot{y}}_P^k \\ \mathbf{\ddot{y}}_Z^k \end{array} \right] \leq \left[ \begin{array}{c} \mathbf{\ddot{r}}_P^k \\ \mathbf{0} \end{array} \right]$$

Only 2 types of residual upper bonds to verify for PS in standard form.  $S = B = P \cup Z$  (positive and zero variables)

# Step 5 Compute maximum step-size $\rho_{\mathsf{S}}^k$

$$\rho \left[ \begin{array}{c} \vec{\mathbf{y}}_{F}^{k} \\ \vec{\mathbf{y}}_{L}^{k} \end{array} \right] \leq \left[ \begin{array}{c} \vec{\mathbf{r}}_{F}^{k} \\ \vec{\mathbf{r}}_{L}^{k} \end{array} \right]; \quad \rho \left[ \begin{array}{c} \mathbf{\tilde{y}}_{F}^{k} \\ \mathbf{\tilde{y}}_{U}^{k} \end{array} \right] \leq \left[ \begin{array}{c} \mathbf{\tilde{r}}_{F}^{k} \\ \mathbf{\tilde{r}}_{U}^{k} \end{array} \right]$$

4 types of strictly positive residual upper bonds to verify in MMCC.  $S = \emptyset$ 

# Step 5 Compute maximum step-size $\rho_{\mathsf{S}}^k$

$$\rho \left[ \begin{array}{c} \vec{\mathbf{y}}_{F}^{k} \end{array} \right] \leq \left[ \begin{array}{c} \vec{\mathbf{r}}_{F}^{k} \end{array} \right]; \quad \rho \left[ \begin{array}{c} \mathbf{\tilde{y}}_{F}^{k} \end{array} \right] \leq \left[ \begin{array}{c} \vec{\mathbf{r}}_{F}^{k} \end{array} \right]$$

$$\rho \left[\begin{array}{c} \vec{\mathbf{y}}_{L}^{k} \end{array}\right] \leq \left[\begin{array}{c} \vec{\mathbf{r}}_{L}^{k} \end{array}\right]; \quad \rho \left[\begin{array}{c} \overline{\mathbf{y}}_{U}^{k} \end{array}\right] \leq \left[\begin{array}{c} \overline{\mathbf{r}}_{U}^{k} \end{array}\right]$$

4 types of strictly positive residual upper bonds to verify in IPS. S = F

# Special case #1:S=B

$$T = [A_B \quad \emptyset], \quad T^{-1} = \begin{bmatrix} A_B^{-1} \\ \emptyset \end{bmatrix}$$
  
 $B = F \cup B_L \cup B_U; \quad N = N_L \cup N_U$ 

### Special case #1: S = B

$$\mathbf{T} = [\mathbf{A}_B \quad {}^{\bullet}], \quad \mathbf{T}^{-1} = \begin{bmatrix} \mathbf{A}_B^{-1} \\ \emptyset \end{bmatrix}$$
$$B = F \cup B_I \cup B_{II}: \quad N = N_I \cup N_{II}$$

#### Primal simplex method (Dantzig 1945)

$$\begin{array}{lll} \mathbf{c}^{\mathsf{T}}\mathbf{x}^k + \min \ \mathbf{c}_B^{\mathsf{T}}(\vec{\mathbf{y}}_B - \mathbf{\bar{y}}_B) & + & \mathbf{c}_N^{\mathsf{T}}(\vec{\mathbf{y}}_N - \mathbf{\bar{y}}_N) \\ & \mathrm{st.} & (\vec{\mathbf{y}}_B - \mathbf{\bar{y}}_B) & + & \bar{\mathbf{A}}_N(\vec{\mathbf{y}}_N - \mathbf{\bar{y}}_N) & = & \mathbf{0} \quad [\boldsymbol{\psi}^{\mathsf{T}} = \mathbf{c}_B] \\ & \vec{\mathbf{y}}_B \geq \mathbf{0}, \ \mathbf{\bar{y}}_B \geq \mathbf{0} & \\ & & \vec{\mathbf{y}}_N \geq \mathbf{0}, \ \mathbf{\bar{y}}_N \geq \mathbf{0} \\ & & \mathbf{\bar{y}}_N \geq \mathbf{0}, \ \mathbf{\bar{y}}_N \geq \mathbf{0} \\ & & \mathbf{\bar{y}}_N \geq \mathbf{0}, \ \mathbf{\bar{y}}_N \geq \mathbf{0} \\ & & \mathbf{\bar{y}}_F \leq \mathbf{\bar{r}}_F, \ \mathbf{\bar{y}}_F \leq \mathbf{\bar{r}}_F \\ & \mathbf{\bar{y}}_{B_I} \leq \mathbf{0}, \ \mathbf{\bar{y}}_{B_U} \leq \mathbf{0} & & \mathbf{\bar{y}}_{N_L} \leq \mathbf{\bar{r}}_{N_L}, \mathbf{\bar{y}}_{N_U} \leq \mathbf{\bar{r}}_{N_U} & *\rho_B \geq \mathbf{0}* \end{array}$$

#### Properties of PS

No equality constraints in the pricing problem.

Pricing contains • convex combination of the non-basic variables

• non-negativity restrictions (a cone).

Due to the step-size  $*\rho_B \ge 0*$ , possible degenerate pivots.

Oscillation of  $\mu_B$ ; it may even not converge towards 0.

# Special case $\#2: S = F \Rightarrow \bar{S} = L \cup U$

$$\mathbf{T} = \left[ egin{array}{ccc} \mathbf{A}_{RF} & \mathbf{0} & & \\ \mathbf{A}_{ZF} & \mathbf{I}_{m-r} & \\ \end{array} 
ight], \quad \mathbf{T}^{-1} = \left[ egin{array}{ccc} \mathbf{A}_{RF}^{-1} & \mathbf{0} & & \\ -\mathbf{A}_{ZS}\mathbf{A}_{RF}^{-1} & \mathbf{I}_{m-r} & \\ \end{array} 
ight].$$

# Special case $\#2: S = F \Rightarrow \bar{S} = L \cup U$

$$\mathbf{T} = \left[ \begin{array}{cc} \mathbf{A}_{RF} & \mathbf{0} \\ \mathbf{A}_{ZF} & \mathbf{I}_{m-r} \end{array} \right], \quad \mathbf{T}^{-1} = \left[ \begin{array}{cc} \mathbf{A}_{RF}^{-1} & \mathbf{0} \\ -\mathbf{A}_{ZS} \mathbf{A}_{RF}^{-1} & \mathbf{I}_{m-r} \end{array} \right].$$

### Improved Primal Simplex method (Elhallaoui et al. 2011)

$$\begin{array}{lllll} \mathbf{c}^{\intercal}\mathbf{x}^{k} + \min & \mathbf{c}_{F}^{\intercal}(\vec{\mathbf{y}}_{F} - \mathbf{\bar{y}}_{F}) & + & \mathbf{c}_{L}^{\intercal}\vec{\mathbf{y}}_{L} - \mathbf{c}_{U}^{\intercal}\mathbf{\bar{y}}_{U} \\ & \text{st.} & (\vec{\mathbf{y}}_{F} - \mathbf{\bar{y}}_{F}) & + & \mathbf{\bar{A}}_{RL}\vec{\mathbf{y}}_{L} - \mathbf{\bar{A}}_{RU}\mathbf{\bar{y}}_{U} & = & \mathbf{0} & [\boldsymbol{\psi}_{R} = \mathbf{c}_{F}] \\ & & \vec{\mathbf{y}}_{F} \geq \mathbf{0}, \ \mathbf{\bar{y}}_{F} \geq \mathbf{0} & \\ & & & \mathbf{\bar{A}}_{ZL}\vec{\mathbf{y}}_{L} - \mathbf{\bar{A}}_{ZU}\mathbf{\bar{y}}_{U} & = & \mathbf{0} & [\boldsymbol{\psi}_{Z}] \\ & & & \vec{\mathbf{y}}_{L} \geq \mathbf{0}, \ \mathbf{\bar{y}}_{U} \geq \mathbf{0} & \\ & & & \vec{\mathbf{y}}_{F} \leq \vec{\mathbf{r}}_{F}, \ \mathbf{\bar{y}}_{F} \leq \mathbf{\bar{r}}_{F}, & \vec{\mathbf{y}}_{L} \leq \vec{\mathbf{r}}_{L}, \mathbf{\bar{y}}_{U} \leq \mathbf{\bar{r}}_{U} & *\rho > 0* \end{array}$$

#### Properties of IPS

f equality constraints in the master problem, m - f in the pricing problem.

Non-degenerate pivots only ( $\rho_F > 0$ ).

 $z^0>z^1>z^2>\cdots=z^*$  cost strictly decreasing at each iteration  $(\rho_F>0)$ .

Oscillation of  $\mu_F$  but converging towards 0.

# Special case $\#3:S=\emptyset$

Select 
$$S = \emptyset$$
.  $T = \begin{bmatrix} \emptyset & I_m \end{bmatrix}$ ,  $T^{-1} = \begin{bmatrix} \emptyset & I_m \end{bmatrix}$ 

### Special case $\#3: S = \emptyset$

Select 
$$S = \emptyset$$
.  $\mathbf{T} = \begin{bmatrix} \emptyset & \mathbf{I_m} \end{bmatrix}, \mathbf{T}^{-1} = \begin{bmatrix} \emptyset \\ \mathbf{I_m} \end{bmatrix}$ 

#### MINIMUM MEAN CYCLE-CANCELING ALGORITHM ADAPTED FOR LP

$$\begin{array}{lll} \mathbf{c}^{\mathsf{T}}\mathbf{x}^k + & \min & \mathbf{c}^{\mathsf{T}}(\vec{\mathbf{y}} - \overline{\mathbf{y}}) & \\ & \mathrm{st.} & \mathbf{A}(\vec{\mathbf{y}} - \overline{\mathbf{y}}) &= \mathbf{0} & \textit{Directions} \\ & & \vec{\mathbf{y}}, \ \overline{\mathbf{y}} &\geq \mathbf{0} & \textit{in the cone} \\ & & & \overline{\mathbf{y}}_L, \ \overline{\mathbf{y}}_U &\leq \mathbf{0} & \textit{at vertex } \mathbf{x}^k \\ \\ & & & \vec{\mathbf{y}}_F, \overline{\mathbf{y}}_F, \ \overline{\mathbf{y}}_L, \ \overline{\mathbf{y}}_U &\leq \vec{\mathbf{r}}_F^k, \ \overline{\mathbf{r}}_F^k, \ \overline{\mathbf{r}}_L^k, \ \overline{\mathbf{r}}_U^k & *\textit{Step size } \rho_\emptyset > 0 * \end{array}$$

#### Properties of MMCC

All equality constraints in the pricing problem.

Upper bounds in the master problem.

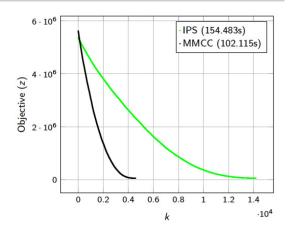
$$\begin{split} z^0 > z^1 > z^2 > \cdots = z^\star & \text{cost strictly decreasing at each iteration } (\rho_\emptyset > 0). \\ \mu^0 \le \mu^1 \le \mu^2 \le \cdots = 0 & \text{smallest reduced cost non decreasing}. \end{split}$$

MMCC is strongly polynomial for network flow problems in O(mn) phases.

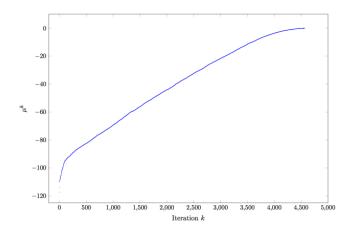
Goldberg and Tarjan (1989), Radzick and Goldberg (1994)

### Illustration of z (Network with n=1025, m=91,220)

#### Computational results

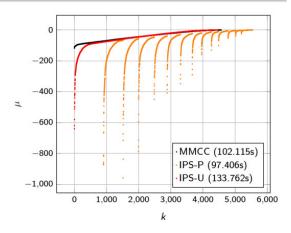


# Illustration of $\mu$ (Network with n=1025, m=91,220)



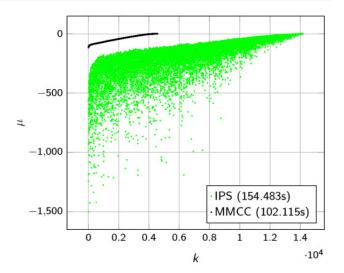
### Illustration of the phases on $\mu$ (n=1025, m=91,220)

#### COMPUTATIONAL RESULTS



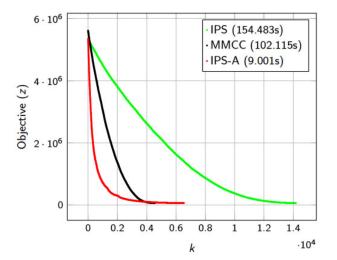
# $\mu$ for MMCC and IPS on a network (n=1025, m=91,220)

#### Computational results



## z for MMCC, IPS and IPS-Accelerated (n=1025, m=91,220)

### COMPUTATIONAL RESULTS



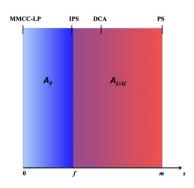
### Dual Guided Pivot Rules for LPs

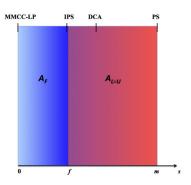
#### Linear program LP

$$\begin{aligned} z^{\star} &:= & & \text{min} & & c^{\mathsf{T}} x \\ \text{st.} & & & \mathbf{A} x &= \mathbf{b} \\ & & & \mathbf{I} \, \leq \, x & \leq \mathbf{u} \end{aligned}$$

### Generic Algorithm with single parameter set S, $\emptyset \subseteq S \subseteq B$

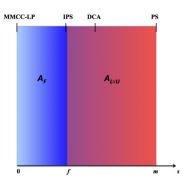
- **1** Let k = 0 and assume a feasible basic solution  $\mathbf{x}^k$  to LP.
- **2** For  $\emptyset \subseteq S \subseteq B$ , construct the residual problem  $LP_S(\mathbf{x}^k)$ .
- **3** Fix a subset of the dual variables in row-set  $R: \pi_R^T = \mathbf{c}_S^T \mathbf{A}_{RS}^{-1}$ .
- Determine the value of the smallest reduced cost  $\mu_s^k$ . If  $\mu_s^k \ge 0$ , STOP. Current solution  $\mathbf{x}^k$  is optimal for LP.
- **3** Retrieve direction  $\mathbf{y}_S^k \in \mathbb{R}^n$  and compute its maximum step-size  $\rho_S^k$ .
- Update  $x^{k+1} := x^k + \rho_S^k y_S^k;$   $z^{k+1} := z^k + \rho_S^k \mu_S^k;$ k := k + 1.
- Goto Step 2.





- $S \cap \{L \cup U\} \neq \emptyset$ : It may come up with degenerate pivots and not converge.
  - Primal Simplex method (PS). S =

#### Properties

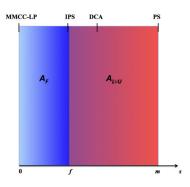


- $S \cap \{L \cup U\} \neq \emptyset$ : It may come up with degenerate pivots and not converge.
  - Primal Simplex method (PS).

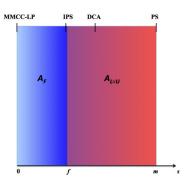
$$S = B$$
, \*  $ρ$ <sub>B</sub> ≥ 0 \*

- $\emptyset \subseteq S \subseteq F$ : It ensures a non-degenerate pivot at every iteration.
  - Improved Primal Simplex algorithm (IPS).  $S = F_1 * \rho_F > 0 *$

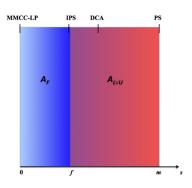
$$S = F$$
, \*  $\rho_F > 0$ 



- $S \cap \{L \cup U\} \neq \emptyset$ : It may come up with degenerate pivots and not converge.
  - Primal Simplex method (PS). S = B, \*  $\rho_B \ge 0$  \*
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  - Minimum mean cycle-canceling algorithm (MMCC)  $S = \emptyset$ , \*  $\rho_{\emptyset} > 0$  \* Strongly polynomial for network flow problems.
  - \*  $S \subset F$ : Optimal direction  $\mathbf{y}_S^k$  can be an <u>interior</u> ray.

#### VECTOR SPACE DECOMPOSITION

- \*\*\* Imagine the same transformation is kept for a while...
  - $\bullet$  T :=  $[\Lambda, \Lambda^{\perp}]$
  - $\Lambda = \begin{bmatrix} \Lambda_R \\ \Lambda_Z \end{bmatrix}$ , where  $\Lambda_R$  is a set of s independent rows.

$$\bullet \ \mathsf{T} = \left[ \begin{array}{cc} \mathsf{\Lambda}_R & \mathsf{0} \\ \mathsf{\Lambda}_Z & \mathsf{I}_{m-s} \end{array} \right]$$

$$\mathsf{T}^{-1} = \left[ \begin{array}{cc} \mathsf{\Lambda}_R^{-1} & \mathsf{0} \\ \mathsf{\Lambda}_Z \mathsf{\Lambda}_R^{-1} & \mathsf{I}_{m-s} \end{array} \right]$$

- $T^{-1}$  splits row-space  $\mathbb{R}^m$  of  $LP(\mathbf{x}^k)$  into two vector subspaces  $\mathbf{V}$  and  $\mathbf{V}^{\perp}$ .
- Vector subspace basis  $\Lambda$  spans V of dimension  $0 \le s \le m$ .
- Vector  $\mathbf{a} \in \mathbf{V}$  if and only if  $\bar{\mathbf{a}}_Z = \mathbf{0}$ , where  $\bar{\mathbf{a}} = \mathbf{T}^{-1}\mathbf{a} = \begin{bmatrix} \bar{\mathbf{a}}_R \\ \mathbf{0} \end{bmatrix}$ .
- $\emptyset \subseteq S \subseteq B$ : index subset of basic columns spanned by  $\Lambda$ .
- Algorithmic properties derived according to subset S.

#### At iteration k, a dynamic Dantzig-Wolfe decomposition

$$\begin{array}{lll} \mathbf{c}^{\intercal}\mathbf{x}^{k} + & \min \ \mathbf{c}_{S}^{\intercal}(\vec{y}_{S} - \mathbf{\bar{y}}_{S}) \ + & \mathbf{c}_{\bar{S}}^{\intercal}(\vec{y}_{\bar{S}} - \mathbf{\bar{y}}_{\bar{S}}) \\ & \text{st.} & \mathbf{A}_{RS}(\vec{y}_{S} - \mathbf{\bar{y}}_{S}) \ + & \mathbf{A}_{R\bar{S}}(\vec{y}_{\bar{S}} - \mathbf{\bar{y}}_{\bar{S}}) \ & = \mathbf{0} \ [\pi_{R}] \\ & & \bar{\mathbf{A}}_{Z\bar{S}}(\mathbf{\bar{y}}_{\bar{S}} - \mathbf{\bar{y}}_{\bar{S}}) \ & = \mathbf{0} \\ & & \bar{\mathbf{y}}_{S} \geq \mathbf{0}, \ \mathbf{\bar{y}}_{S} \geq \mathbf{0} \\ & & \mathbf{\bar{y}}_{\bar{S}} \geq \mathbf{0}, \ \mathbf{\bar{y}}_{\bar{S}} \geq \mathbf{0} \\ & & \bar{\mathbf{y}}_{\bar{S}_{L}} \leq \mathbf{0}, \ \mathbf{\bar{y}}_{\bar{S}_{U}} \leq \mathbf{0} \\ & & \bar{\mathbf{y}}_{S_{F}} \leq \mathbf{\bar{r}}_{S_{F}}, \ \mathbf{\bar{y}}_{S_{F}} \leq \mathbf{\bar{r}}_{\bar{S}_{F}} \\ & \bar{\mathbf{y}}_{S_{I}} \leq \mathbf{0}, \ \mathbf{\bar{y}}_{S_{U}} \leq \mathbf{0} \\ & & \bar{\mathbf{y}}_{\bar{S}_{I}} \leq \mathbf{\bar{r}}_{\bar{S}_{F}}, \ \mathbf{\bar{y}}_{\bar{S}_{F}} \leq \mathbf{\bar{r}}_{\bar{S}_{F}} \\ & \bar{\mathbf{y}}_{S_{I}} \leq \mathbf{0}, \ \mathbf{\bar{y}}_{S_{U}} \leq \mathbf{0} \\ & & \bar{\mathbf{y}}_{\bar{S}_{I}} \leq \mathbf{\bar{r}}_{\bar{S}_{I}}, \ \mathbf{\bar{y}}_{\bar{S}_{IJ}} \leq \mathbf{\bar{r}}_{\bar{S}_{IJ}} \end{array}$$

- Pricing domain keeps the m-s equality rows corresponding to  $\mathbf{V}^{\perp}$  and  $\mathbf{\bar{y}}_{\bar{5}} \geq \mathbf{0}, \ \mathbf{\bar{y}}_{\bar{5}} \geq \mathbf{0}, \ \mathbf{\bar{y}}_{\bar{5}_{ij}} = \mathbf{0}, \ \mathbf{\bar{y}}_{\bar{5}_{ij}} = \mathbf{0}.$
- Master domain keeps the s equality rows corresponding to V, and all remaining lower and upper bounds of variables.

### Special case $\#4: F \subseteq S \subseteq B$

# Dynamic Constraint Aggregation for Set Partitioning (Elhallaoui et al. 2005)

The partition of the row-set is derived from the f groups of identical rows of  $A_F$ .

#### Properties

 $F \subseteq S \subseteq B$  for fractional solutions but S = F (or  $S = \emptyset$ ) for binary solutions. Degenerate pivots may occur,  $\rho \ge 0$ , if r > f.

### $\mathsf{A}_{\mathsf{S}}$ : A SUBSET OF $\mathsf{s}$ COLUMNS OF THE SIMPLEX BASIS $\mathsf{A}_{\mathsf{B}}$

