

## 5 Computational Tests

In this section, we present the results obtained by the branch-and-cut algorithm described in Section 4. The branch-and-cut algorithm (called ‘BC’ from now on) was coded in C++ using the LEMON graph library [10]. All tests were performed on an OSX platform (iMac 2020), running on an Intel Core i9-10910 processor clocked at 3.6 GHz with 64 GB of RAM. The mathematical formulation was solved using the ILOG Concert Technology library and CPLEX 20.1 in single thread mode. A time limit of one hour and a memory limit of 5GB were imposed. All remaining CPLEX parameters were left to their default value.

The computational tests were carried out on the instances proposed in [2] and named *Set1*. This set of instances was obtained by adapting the instances of the Generalized Traveling Salesman Problem proposed in [13]. The number of vertices ranges from 52 to 1084 but we only consider the instances with up to 198 nodes for our computational tests as for higher dimensions the exact approach is not practicable. The number of clusters is equal to  $\sim 20\%$  the number of vertices.  $T_{max}$  is set to  $\omega \times GTSP^*$ , where  $GTSP^*$  is the best-known solution value of the GTSP (taken from [13]) and  $\omega$  is set to 0.4, 0.6, and 0.8. Finally, two rules,  $g_1$  and  $g_2$ , are used to assign the profit to the clusters. The first rule sets the profit of each cluster  $C_g$  equal to  $|C_g|$ . The second rule assigns to vertex  $j$  a profit equal to  $1 + (7141j + 73) \bmod(100)$  and the profit of a cluster is given by the sum of the values of the profit of all vertices belonging to it.

Another set of instances for SOP, named *Set2*, was introduced in [2]. The instances in this set are the same as those in *Set1*. More specifically, the instances contain the same vertices and the same number of clusters as the instances in *Set1*, but the vertices are assigned in a random way to the clusters. Both sets contain 27 instances each. The datasets and the tables of the paper are available here [https://github.com/fcarrabs/Set\\_Orienteering\\_Problem](https://github.com/fcarrabs/Set_Orienteering_Problem)

The section is organized as follows. We first analyze the effectiveness of the valid inequalities presented in Section 3. Specifically, in Section 5.1, we present the results of tests in which we run the branch-and-cut algorithm with the full set of valid inequalities and we compare it with the version of the algorithm in which we discard one inequality at a time. These tests enable us to determine the best version of the algorithm which is used to run the final tests. In Section 5.2 we provide the detailed results of the final

	C-BC Time	NoCond Time Gap%	NoCover Time Gap%	NoCluCover Time Gap%	NoPath Time Gap%
$\omega = 0.4$ and $g_1$					
AVG	622.50	610.86 0.80%	1229.98 -2.73%	615.23 0.00%	1361.12 -8.08%
#Worse		0	7	0	11
$\omega = 0.4$ and $g_2$					
AVG	808.27	696.90 0.31%	1230.15 -2.89%	751.66 0.00%	1479.63 -8.06%
#Worse		0	7	0	13
$\omega = 0.6$ and $g_1$					
AVG	934.94	1004.03 0.13%	1573.66 -1.83%	932.39 0.00%	1765.57 -2.43%
#Worse		1	5	0	6
$\omega = 0.6$ and $g_2$					
AVG	1208.51	1104.96 -0.15%	1653.26 -0.28%	1188.54 0.00%	1580.25 -0.66%
#Worse		1	2	0	2
$\omega = 0.8$ and $g_1$					
AVG	909.19	1222.59 -0.83%	2174.25 -4.99%	908.76 0.00%	2130.63 -4.90%
#Worse		2	10	0	10
$\omega = 0.8$ and $g_2$					
AVG	1768.56	1744.82 0.09%	2290.29 -2.98%	1845.61 -0.08%	2362.75 -3.30%
#Worse		1	8	1	8

Table 1: Summary of the computational results for the five versions of the branch-and-cut algorithm on all instances.

version of the algorithm and we provide some statistics on the performance of the algorithm. Finally, in Section 5.3 we compare the performance of our algorithm with two benchmark approaches available in the literature.

## 5.1 Valid Inequalities Effectiveness

As mentioned above, this section is dedicated to evaluate the effectiveness of the valid inequalities presented in Section 3. Specifically, we run the branch-and-cut algorithm with the full set of inequalities and we compare it with the versions in which we discard a single inequality at a time. Tests are made over all instances presented above.

The results of this comparison are summarized in Table 1 for the instances with  $\omega = 0.4, 0.6, 0.8$  and profits  $g_1$  and  $g_2$ . The detailed results are reported in Tables 12–15 of Appendix A. The first column of the table reports the computational time, in seconds, of the branch-and-cut version with the full set of inequalities (*C-BC*). Four groups of two columns follow, corresponding to the four versions of the branch-and-cut algorithm in which one inequality is excluded: no conditional cuts (*NoCond*), no cover

inequalities (*NoCover*), no cluster cover inequalities (*NoCluCover*) and no path inequalities (*NoPath*). Under the *Time* and *Gap%* headings, for each branch-and-cut version, we report the corresponding computational time and the percentage gap between the upper bound obtained by C-BC and the upper bound of the version considered, respectively. This gap is calculated as  $gap = \frac{UB_{C-BC} - UB_*}{UB_{C-BC}}$ , where  $UB_{C-BC}$  and  $UB_*$  are the upper bound by C-BC and by the version considered, respectively. Note that positive values of the gap mean that the version of the algorithm without the inequality gives a better result than C-BC. Results are aggregated by values of  $\omega$  and classes of profit. The rows of the table are organized in groups of two lines reporting, for each version of the algorithm in which an inequality is excluded: the average values of time and gap (*AVG*) and the number of times in which the upper bound provided by the version without an inequality is worse than the one provided by C-BC (*#Worse*).

We note that, when  $\omega = 0.4$ , cover and path inequalities are indeed effective as their exclusion cause both a remarkable increase in computing time and a deterioration in the value of the upper bound.

In more detail, for both  $g_1$  and  $g_2$  profits, we observe that the upper bound of NoCover worsens 7 times while the computational time increases by  $\sim 52\%$ , at least. For NoPath the upper bound worsens 11 and 13 times, for  $g_1$  and  $g_2$  respectively, and the computational time increases by  $\sim 96\%$ , at least. For  $\omega = 0.6$ , the contribution of these two inequalities is less impactful but still relevant. For  $\omega = 0.8$ , the *#Worse* value of the version without path inequalities is similar to the one observed for  $\omega = 0.4$  and profit  $g_1$  while it is lower for profit  $g_2$ . Instead, for the version without the cover inequalities, the time, *Gap%*, and *#Worse* value are worse than the ones observed for  $\omega = 0.4$  and  $\omega = 0.6$ . The effectiveness of conditional cuts is instead more debatable: indeed, they are not effective when  $\omega = 0.4$ , as their removal improve the upper bound and reduces the computational time. However, when  $\omega = 0.6$ , the computational time increases when they are removed for the case  $g_1$  and the upper bound slightly deteriorates for the case  $g_2$ . Finally, when  $\omega = 0.8$ , they are effective in  $g_1$  instances and not effective in  $g_2$  instances. Given that the overall difference between the two versions of the algorithm (with or without conditional cuts) is slightly to the advantage of the first version, we decided to retain them in the final version of the branch-and-cut algorithm. Finally, as for cluster cover inequalities, we notice that they are almost never effective: in fact, the impact on the value of the upper bound is null while the

computational time reduces when they are discarded (apart the case  $\omega = 0.8$  and  $g_2$ ). For this reason, we decided to remove these inequalities from the final version of the branch-and-cut algorithm.

## 5.2 Branch-and-cut performance

The results presented above show that the best version of the branch-and-cut algorithm is the one that does not include cluster cover inequalities. This is the version for which we present the results in the remaining of this section, which is from now on called *BC*.

This section is organized as follows. We first present some statistics about cuts and valid inequalities separation. Then, we analyse the performance of the algorithm over all instances of the testbed and compare it with a version of the algorithm where none of the inequalities presented in Section 3 is used. The aim of this comparison is to show that indeed the inequalities pay-off (apart the cluster cover inequalities as mentioned in the former section).

Specifically, Tables 2–5 present statistics about the separation of subtour elimination constraints and the remaining inequalities for the instances with  $\omega = 0.4, 0.6, 0.8$ , respectively. The first column reports the name of the instance which contains the reference to the number of customers in the instance (last part of the instance’ name). In the second we have the number of nodes of the branch-and-bound tree explored at termination (*Nnodes*). Then we report statistics related to the separation of subtour elimination constraints, conditional cuts, cover inequalities and path inequalities, respectively. Specifically, for each family of inequality, we report the number of inequalities separated (*Num*) and the total computational time for separation (*Time*). The rows of the table are divided in two groups associated with the datasets *Set1* and *Set2*, respectively. Finally, the last row of the table (*AVG*) reports the average values of *Nnodes*, *Num* and *Time*.

Focusing first on Table 2, we see that subtour elimination constraints are by far the most widely violated constraints, followed by path inequalities. The remaining inequalities are much less often violated. As for the time for separation, we notice that, while subtour elimination constraints are separated very efficiently, the computational time for separating path inequalities is, in comparison, much larger. This justifies our choice of separating path inequalities only when no other inequality is violated. As for the other inequalities, we recall that conditional cuts are separated during the separation of subtour elimination constraints, which explains the short computing time

$\omega = 0.4$ and $g_1$										
	Instance	Nnodes	Subtour		Conditional		Cover		Path	
			Num	Time	Num	Time	Num	Time	Num	Time
Set1	11berlin52	47	148	0.01	3	0.00	5	0.11	26	0.14
	11eil51	44	163	0.01	0	0.00	4	0.07	18	0.03
	14st70	60	217	0.03	0	0.00	3	0.11	29	0.10
	16eil76	144	709	0.17	2	0.00	12	0.43	182	0.78
	16pr76	126	705	0.20	0	0.00	6	1.04	51	0.26
	20kroA100	327	2754	1.49	0	0.00	11	4.32	137	0.62
	20kroB100	179	1125	0.50	2	0.00	17	1.68	94	0.90
	20kroC100	86	534	0.16	1	0.00	5	0.62	36	0.23
	20kroD100	127	670	0.18	1	0.00	15	0.61	125	0.82
	20kroE100	131	743	0.21	1	0.00	3	0.77	24	0.14
	20rat99	107	433	0.05	2	0.00	3	0.19	50	0.58
	20rd100	186	1082	0.31	0	0.00	6	0.78	67	0.28
	21eil101	150	1229	1.27	0	0.00	11	1.14	112	1.85
	21lin105	243	1212	0.29	3	0.00	6	1.16	61	0.32
	22pr107	3	9	0.00	0	0.00	0	0.01	0	0.00
	25pr124	775	4938	3.13	2	0.00	41	9.91	310	6.27
	26bier127	948	10409	25.12	1	0.00	13	21.28	184	3.50
	26ch130	812	9358	19.93	1	0.00	48	14.64	223	9.59
	28pr136	305	2423	2.31	0	0.00	19	2.81	291	4.34
	29pr144	1669	16353	22.97	3	0.00	24	21.38	462	5.76
	30ch150	1468	12015	19.23	17	0.00	66	16.30	528	13.48
	30kroA150	1118	12733	27.13	2	0.00	57	23.60	358	13.11
	30kroB150	692	7658	20.54	7	0.00	29	17.78	531	29.73
	31pr152	1094	11838	14.13	4	0.00	12	14.46	305	4.00
	32u159	969	10084	14.10	1	0.00	23	10.61	125	3.97
	39rat195	685	8612	13.78	1	0.00	17	6.61	344	10.00
	40d198	279	1914	1.23	17	0.00	24	1.66	230	6.65
Set2	11berlin52	53	283	0.05	4	0.00	4	0.19	37	0.35
	11eil51	116	466	0.07	1	0.00	18	0.29	151	1.57
	14st70	41	293	0.06	0	0.00	2	0.13	33	0.21
	16eil76	103	672	0.26	0	0.00	4	0.36	120	1.05
	16pr76	225	1765	0.80	3	0.00	3	2.38	44	0.64
	20kroA100	243	2564	2.30	3	0.00	17	5.51	166	10.56
	20kroB100	277	2645	2.08	1	0.00	15	2.86	47	4.32
	20kroC100	260	2051	1.48	0	0.00	23	2.61	157	8.39
	20kroD100	135	1042	0.59	0	0.00	8	1.43	80	3.87
	20kroE100	87	790	0.47	1	0.00	6	0.76	39	3.39
	20rat99	56	416	0.10	8	0.00	2	0.16	5	1.26
	20rd100	234	1993	1.30	1	0.00	6	1.68	209	11.84
	21eil101	231	2506	2.93	2	0.00	16	2.43	110	5.54
	21lin105	496	3880	2.22	42	0.00	24	3.06	430	6.83
	22pr107	281	2789	0.89	18	0.00	21	1.25	191	4.49
	25pr124	556	6744	7.66	0	0.00	12	13.41	185	46.82
	26bier127	837	16399	47.56	23	0.00	23	32.75	260	18.22
	26ch130	955	10080	26.69	3	0.00	60	19.73	329	268.79
	28pr136	186	2864	4.03	0	0.00	5	2.12	77	4.56
	29pr144	402	9044	19.02	0	0.00	10	13.22	336	7.16
	30ch150	544	8392	21.25	0	0.00	23	9.61	228	112.49
	30kroA150	375	9272	28.64	4	0.00	11	15.80	173	11.96
	30kroB150	646	12812	42.42	11	0.00	51	33.13	444	85.94
	31pr152	457	9460	18.32	0	0.00	18	11.15	328	19.43
	32u159	333	5439	12.63	0	0.00	19	4.82	154	9.81
	39rat195	746	14727	38.77	0	0.00	29	8.98	154	59.56
	40d198	478	7665	12.09	0	0.00	11	4.83	158	47.74
AVG		409.76	4761.50	8.95	3.63	0.00	17.06	6.83	176.81	16.19

Table 2: Branch-and-cut statistics on the instances with  $\omega = 0.4$  and profit  $g_1$ .

of separation for this class. As for the cover inequalities, the separation time is short so this justifies the choice of keeping these inequalities despite the fact that they are violated quite rarely.

Similar considerations can be done for  $\omega$  equal to 0.4 and profit  $g_2$  (Table 3). However, an interesting observation here is that instances in  $g_2$  are more difficult to solve as witnessed by the larger number of nodes of the branch-and-bound tree. As a consequence, the number of inequalities separated and the computational time required to separate them almost always increases for all the types of inequalities.

The same trend is observed for values of  $\omega$  equal to 0.6 and 0.8 (reported in Tables 4 and 5, respectively) where again we notice that the instances with profit  $g_2$  are more difficult to solve with respect to the ones with profit  $g_1$ .

We now present the results of the comparison between BC and the branch-and-cut algorithm that solves formulation **clucut** only, i.e., with none of the valid inequalities proposed in Section 3. Formulation **clucut** is solved through branch-and-cut by separating the subtour elimination constraints (15). The separation algorithm is the same as the one used in BC.

The results of the comparison between **clucut** and BC are reported in Tables 6–8 for values of  $\omega$  equal to 0.4, 0.6 and 0.8, respectively.

Each table is organized as follows. The first column reports the name of the instance. Then, two groups of six columns report the results for profits  $g_1$  and  $g_2$ , respectively. Each group is composed by two subgroups of three columns, referring to the results by **clucut**, first, and BC, second. Specifically, we report the value of the best solution found at termination, the computational time (in seconds) and the optimality gap at termination (which is 0% in case the instance is solved to proven optimality). The rows of the table are divided in two groups associated with the datasets *Set1* and *Set2*, respectively. The last two rows of each table report the average computational time and the average optimality gap, and the number of instances solved to optimality, respectively.

The values on the AVG line in Table 6, related to instances with  $\omega = 0.4$ , show that BC is much faster than **clucut**. Indeed, the average computational time for BC is equal to 615 and 751 seconds for the instances with profits  $g_1$  and  $g_2$ , respectively. On the contrary, for **clucut** this time is equal to 1370 and 1360 seconds, respectively. Regarding the effectiveness, the average Gap% values are equal to 4.55% and 5.69% for BC and equal to 14.99% and 14.61% for **clucut**. Moreover, from line #Opt we observe that BC is able to solve 49 instances to optimality (out of 54) for  $g_1$  and 47 for  $g_2$ , while **clucut**

$\omega = 0.4$ and $g_2$										
	Instance	Nnodes	Subtour		Conditional		Cover		Path	
			Num	Time	Num	Time	Num	Time	Num	Time
Set1	11berlin52	60	182	0.02	1	0.00	6	0.14	38	0.12
	11eil51	159	332	0.03	0	0.00	21	0.24	152	0.50
	14st70	74	295	0.03	2	0.00	2	0.15	8	0.03
	16eil76	126	547	0.12	18	0.00	14	0.36	112	0.46
	16pr76	139	599	0.16	13	0.00	9	0.79	89	0.26
	20kroA100	402	2517	1.34	1	0.00	20	5.33	173	3.65
	20kroB100	296	1623	0.76	0	0.00	35	2.91	200	3.45
	20kroC100	194	796	0.24	3	0.00	14	1.11	162	1.53
	20kroD100	288	992	0.26	0	0.00	29	1.48	522	3.03
	20kroE100	168	944	0.28	1	0.00	11	1.15	103	0.98
	20rat99	76	278	0.03	2	0.00	4	0.12	34	0.15
	20rd100	217	1295	0.38	1	0.00	2	1.07	53	1.21
	21eil101	229	1816	1.83	2	0.00	6	1.60	72	0.98
	21lin105	229	944	0.23	13	0.00	11	1.01	128	1.14
	22pr107	7	29	0.00	0	0.00	1	0.02	17	0.03
	25pr124	1135	5835	3.89	4	0.00	85	12.45	629	17.94
	26bier127	2083	18913	48.53	1	0.00	57	37.82	504	32.86
	26ch130	1560	13028	27.41	1	0.00	116	21.11	513	20.82
	28pr136	865	7784	7.49	2	0.00	36	8.75	617	4.83
	29pr144	3682	12905	17.26	8	0.00	490	42.63	1889	98.98
	30ch150	1376	10970	17.88	1	0.00	40	14.12	548	33.02
	30kroA150	1473	12756	25.96	1	0.00	111	32.42	898	32.75
	30kroB150	2359	17929	47.76	3	0.00	258	54.98	1335	60.74
	31pr152	1099	11895	14.29	2	0.00	14	14.50	270	4.20
	32u159	884	9275	12.98	1	0.00	15	11.23	139	4.00
	39rat195	749	7499	12.40	4	0.00	81	8.67	920	30.59
	40d198	422	2967	1.84	2	0.00	10	2.45	142	2.17
Set2	11berlin52	36	211	0.03	3	0.00	4	0.14	35	0.47
	11eil51	71	308	0.05	2	0.00	5	0.14	24	0.23
	14st70	78	401	0.07	3	0.00	5	0.18	35	0.20
	16eil76	273	1224	0.47	1	0.00	20	0.95	274	7.02
	16pr76	214	1353	0.55	0	0.00	12	2.19	131	6.31
	20kroA100	342	2832	2.53	1	0.00	35	6.59	237	20.89
	20kroB100	396	2920	2.40	3	0.00	37	6.64	183	40.50
	20kroC100	808	3278	2.48	3	0.00	91	7.85	646	43.54
	20kroD100	342	2421	1.50	1	0.00	31	3.33	155	11.59
	20kroE100	154	1397	0.93	1	0.00	11	2.13	109	19.38
	20rat99	395	975	0.23	16	0.00	38	1.00	405	39.81
	20rd100	265	2561	1.62	0	0.00	9	1.86	128	5.44
	21eil101	356	3269	3.93	2	0.00	32	3.90	165	9.23
	21lin105	392	3677	2.01	4	0.00	9	2.48	177	2.47
	22pr107	253	2304	0.74	30	0.00	12	1.07	187	2.79
	25pr124	524	8959	9.81	4	0.00	13	15.41	131	11.97
	26bier127	530	9626	25.04	5	0.00	40	14.89	419	40.26
	26ch130	493	9396	24.93	1	0.00	23	14.11	282	79.01
	28pr136	491	4656	6.58	0	0.00	57	4.48	243	31.63
	29pr144	339	7470	15.52	0	0.00	9	9.13	319	7.80
	30ch150	657	10182	25.18	0	0.00	39	10.80	270	95.81
	30kroA150	436	10989	33.92	2	0.00	10	22.80	210	5.74
	30kroB150	395	6989	22.99	2	0.00	64	22.72	464	271.12
	31pr152	434	9557	18.94	1	0.00	9	10.56	251	8.01
	32u159	392	6500	14.98	1	0.00	11	6.06	143	7.72
	39rat195	655	9809	25.55	1	0.00	57	7.09	477	71.29
	40d198	744	6984	11.29	0	0.00	113	7.95	486	195.61
AVG		570.67	5096.17	9.22	3.22	0.00	42.48	8.61	312.09	25.86

Table 3: Branch-and-cut statistics on the instances with  $\omega = 0.4$  and profit  $g_2$ .

$\omega = 0.6$ and $g_1$										
	Instance	Nnodes	Subtour		Conditional		Cover		Path	
			Num	Time	Num	Time	Num	Time	Num	Time
Set1	11berlin52	214	813	0.12	0	0.00	14	0.72	100	0.80
	11eil51	105	509	0.09	0	0.00	12	0.33	65	0.50
	14st70	226	1316	0.44	2	0.00	11	0.76	109	2.90
	16eil76	167	1466	0.64	3	0.00	6	0.79	54	0.47
	16pr76	321	2092	1.24	3	0.00	16	2.50	158	2.23
	20kroA100	490	5432	5.55	2	0.00	7	8.74	80	1.81
	20kroB100	428	4196	4.48	0	0.00	11	5.95	144	1.03
	20kroC100	425	4048	4.02	1	0.00	24	5.86	261	2.23
	20kroD100	451	4369	3.64	1	0.00	9	5.67	105	3.26
	20kroE100	477	3959	2.43	0	0.00	4	4.29	46	0.49
	20rat99	313	2436	1.09	1	0.00	3	1.30	36	0.86
	20rd100	781	6982	7.07	0	0.00	23	6.55	176	5.94
	21eil101	450	4324	5.39	4	0.00	19	3.38	126	3.21
	21lin105	752	5506	3.84	7	0.00	9	6.23	90	1.18
	22pr107	700	7290	4.88	0	0.00	8	5.65	246	1.83
Set2	11berlin52	7	39	0.01	1	0.00	0	0.03	0	0.00
	11eil51	82	544	0.13	0	0.00	4	0.36	34	2.42
	14st70	336	2471	1.16	6	0.00	8	1.32	64	1.15
	16eil76	250	2542	1.67	3	0.00	6	1.58	5	4.12
	16pr76	316	2895	1.97	4	0.00	9	3.74	87	1.61
	20kroA100	356	5927	7.75	2	0.00	21	10.43	276	9.97
	20kroB100	336	5446	6.36	0	0.00	5	8.54	133	1.36
	20kroC100	328	4693	6.68	6	0.00	13	11.55	155	8.48
	20kroD100	420	5592	7.27	2	0.00	12	9.83	297	6.38
	20kroE100	314	3121	2.81	2	0.00	13	4.98	80	16.65
	20rat99	419	2929	2.01	2	0.00	3	1.99	21	53.31
	20rd100	313	4697	7.17	5	0.00	6	6.07	78	1.23
	21eil101	262	3696	5.72	0	0.00	3	2.71	60	2.99
	21lin105	515	7737	7.59	1	0.00	3	10.55	45	1.10
	22pr107	96	596	0.43	0	0.00	20	0.59	312	6.99
AVG		355.00	3588.77	3.45	1.93	0.00	10.07	4.43	114.77	4.88
$\omega = 0.6$ and $g_2$										
	Instance	Nnodes	Subtour		Conditional		Cover		Path	
			Num	Time	Num	Time	Num	Time	Num	Time
Set1	11berlin52	258	1063	0.17	5	0.00	10	1.00	67	0.37
	11eil51	73	414	0.08	0	0.00	4	0.21	53	0.41
	14st70	254	1075	0.34	1	0.00	25	0.82	292	6.31
	16eil76	224	1597	0.71	1	0.00	3	0.78	21	0.34
	16pr76	764	4298	2.59	5	0.00	54	5.80	210	6.45
	20kroA100	576	5376	5.60	0	0.00	37	9.47	273	7.17
	20kroB100	575	5233	5.57	1	0.00	16	8.84	130	1.67
	20kroC100	886	7800	7.97	5	0.00	37	12.03	200	7.86
	20kroD100	597	4214	3.43	1	0.00	39	6.09	396	9.01
	20kroE100	457	3425	2.21	0	0.00	7	3.43	147	1.53
	20rat99	311	2443	1.09	2	0.00	5	1.41	70	0.72
	20rd100	764	4193	4.34	0	0.00	71	6.76	609	20.49
	21eil101	966	6079	7.53	0	0.00	84	7.31	304	22.71
	21lin105	837	5216	3.69	0	0.00	14	7.10	160	3.75
	22pr107	734	7588	5.06	8	0.00	17	6.19	464	4.05
Set2	11berlin52	7	39	0.01	1	0.00	0	0.03	0	0.00
	11eil51	21	104	0.02	0	0.00	3	0.09	17	0.29
	14st70	420	3085	1.58	1	0.00	3	1.59	30	0.69
	16eil76	222	2109	1.34	4	0.00	8	1.25	47	1.24
	16pr76	260	2112	1.48	0	0.00	4	2.30	69	1.13
	20kroA100	335	5668	7.73	0	0.00	7	12.23	107	29.13
	20kroB100	453	6857	9.04	4	0.00	13	13.94	254	5.15
	20kroC100	348	5350	7.04	0	0.00	12	8.48	134	5.07
	20kroD100	372	5380	6.29	2	0.00	22	8.76	131	7.37
	20kroE100	408	5533	5.84	13	0.00	11	7.05	124	3.12
	20rat99	493	3264	2.18	1	0.00	3	2.36	96	78.91
	20rd100	310	5185	8.18	2	0.00	8	6.50	104	1.32
	21eil101	453	7298	11.22	4	0.00	7	5.98	67	22.19
	21lin105	631	9487	11.18	13	0.00	10	10.12	302	6.89
	22pr107	460	5835	5.73	44	0.00	53	4.85	495	28.38
AVG		448.97	4244.00	4.31	3.93	0.00	19.57	5.43	179.10	9.46

Table 4: Branch-and-cut statistics on the instances with  $\omega = 0.6$  and profit  $g_1$  and  $g_2$ .



$\omega = 0.8$ and $g_1$										
	Instance	Nnodes	Subtour		Conditional		Cover		Path	
			Num	Time	Num	Time	Num	Time	Num	Time
Set1	11berlin52	269	1286	0.22	22	0.00	10	1.03	81	0.45
	11eil51	101	472	0.10	1	0.00	5	0.23	31	0.15
	14st70	319	2313	0.97	2	0.00	4	1.29	40	0.37
	16eil76	294	2533	1.35	2	0.00	12	1.31	90	1.12
	16pr76	365	3602	2.33	2	0.00	12	3.83	116	5.05
	20kroA100	570	4960	5.36	0	0.00	54	7.68	293	17.85
	20kroB100	1125	12442	15.30	1	0.00	28	16.56	328	61.90
	20kroC100	482	4720	5.17	5	0.00	45	6.17	462	21.03
	20kroD100	473	5581	6.00	1	0.00	16	5.71	226	3.02
	20kroE100	788	5783	6.20	6	0.00	53	9.63	406	23.34
	20rat99	737	6604	5.94	3	0.00	27	4.97	155	11.53
	20rd100	418	4687	5.74	1	0.00	13	5.91	133	1.45
	21eil101	523	6211	8.46	3	0.00	7	4.37	48	2.18
	21lin105	748	7544	10.64	10	0.00	27	10.56	193	6.55
	22pr107	908	14048	16.36	9	0.00	30	13.24	612	13.21
Set2	11berlin52	0	4	0.00	0	0.00	0	0.01	0	0.00
	11eil51	69	363	0.07	8	0.00	2	0.20	9	0.07
	14st70	92	599	0.25	3	0.00	6	0.51	49	0.56
	16eil76	83	821	0.46	16	0.00	4	0.59	41	1.02
	16pr76	130	1214	0.67	20	0.00	9	1.56	95	1.46
	20kroA100	80	972	1.16	4	0.00	9	1.62	85	2.50
	20kroB100	420	6160	8.85	14	0.00	10	13.27	118	2.95
	20kroC100	417	6221	9.89	1	0.00	6	13.63	102	2.37
	20kroD100	669	10003	14.61	17	0.00	2	18.23	17	0.60
	20kroE100	456	5856	7.20	1	0.00	19	10.64	258	8.40
	20rat99	304	4146	3.49	1	0.00	2	3.14	16	1.45
	20rd100	300	4728	6.79	19	0.00	9	5.05	162	3.30
	21eil101	327	5445	8.45	6	0.00	3	4.32	33	10.76
	21lin105	36	329	0.39	1	0.00	4	0.52	48	0.87
	22pr107	45	408	0.52	3	0.00	9	0.60	95	2.91
AVG		384.93	4335.17	5.10	6.07	0.00	14.57	5.55	144.73	6.95
$\omega = 0.8$ and $g_2$										
	Instance	Nnodes	Subtour		Conditional		Cover		Path	
			Num	Time	Num	Time	Num	Time	Num	Time
Set1	11berlin52	384	1568	0.28	27	0.00	7	1.22	51	0.55
	11eil51	228	1091	0.24	1	0.00	9	0.55	27	0.50
	14st70	434	2729	1.23	14	0.00	5	1.41	41	0.32
	16eil76	300	2120	1.17	0	0.00	10	1.18	86	1.51
	16pr76	588	5379	3.92	8	0.00	14	6.25	89	2.00
	20kroA100	920	7834	8.65	0	0.00	74	12.71	486	27.40
	20kroB100	842	11643	13.59	1	0.00	25	14.68	181	2.77
	20kroC100	811	6911	8.04	3	0.00	43	8.73	394	23.33
	20kroD100	559	8551	9.20	1	0.00	12	8.89	103	2.97
	20kroE100	720	5812	6.43	0	0.00	31	8.23	718	25.94
	20rat99	1111	9701	8.69	6	0.00	31	7.22	287	10.00
	20rd100	530	8827	10.87	1	0.00	8	12.51	113	8.27
	21eil101	1048	9168	13.65	0	0.00	38	9.02	113	32.38
	21lin105	814	11973	15.85	2	0.00	12	14.93	77	3.44
	22pr107	750	10736	12.13	24	0.00	16	10.07	290	5.17
Set2	11berlin52	2	6	0.00	0	0.00	0	0.01	0	0.00
	11eil51	34	235	0.05	1	0.00	1	0.14	6	0.03
	14st70	220	2137	0.84	0	0.00	4	1.29	35	0.26
	16eil76	350	3974	2.79	8	0.00	3	2.57	32	4.96
	16pr76	510	5235	3.55	26	0.00	8	6.48	76	2.93
	20kroA100	498	6557	8.70	7	0.00	44	13.67	222	19.35
	20kroB100	561	6972	8.69	0	0.00	12	12.10	146	3.27
	20kroC100	397	6195	7.81	2	0.00	11	11.45	151	3.79
	20kroD100	734	10636	13.71	29	0.00	17	20.35	185	4.13
	20kroE100	581	8520	10.23	0	0.00	7	16.56	110	1.61
	20rat99	286	3537	3.92	4	0.00	10	3.00	117	29.47
	20rd100	219	2566	3.48	4	0.00	11	3.38	111	2.73
	21eil101	620	9369	15.33	10	0.00	7	8.14	87	53.05
	21lin105	761	9707	14.16	21	0.00	27	17.60	281	10.86
	22pr107	316	2880	3.74	1	0.00	26	3.77	383	8.51
AVG		537.60	6085.63	7.03	6.70	0.00	17.43	7.94	166.60	9.72

Table 5: Branch-and-cut statistics on the instances with  $\omega = 0.8$  and profit  $g_1$  and  $g_2$ .

$\omega = 0.4$													
Instance	$g_1$			BC			$g_2$			BC			
	Sol	Time	Gap%	Sol	Time	Gap%	Sol	Time	Gap%	Sol	Time	Gap%	
Set1	11berlin52	37	0.56	0.00%	37	0.41	0.00%	1829	0.64	0.00%	1829	0.46	0.00%
	11eil51	24	0.14	0.00%	24	0.21	0.00%	1279	0.26	0.00%	1279	1.03	0.00%
	14st70	33	0.45	0.00%	33	0.48	0.00%	1672	0.56	0.00%	1672	0.48	0.00%
	16eil76	40	3.38	0.00%	40	2.54	0.00%	2223	3.40	0.00%	2223	1.74	0.00%
	16pr76	47	4.57	0.00%	47	5.19	0.00%	2449	13.97	0.00%	2449	3.75	0.00%
	20kroA100	42	40.14	0.00%	42	27.14	0.00%	2151	69.19	0.00%	2151	30.50	0.00%
	20kroB100	49	13.84	0.00%	49	8.32	0.00%	2431	24.68	0.00%	2431	15.81	0.00%
	20kroC100	42	2.72	0.00%	42	2.07	0.00%	2174	2.62	0.00%	2174	4.64	0.00%
	20kroD100	39	3.23	0.00%	39	2.94	0.00%	1740	8.50	0.00%	1740	7.84	0.00%
	20kroE100	52	3.55	0.00%	52	3.68	0.00%	2415	2.49	0.00%	2415	5.18	0.00%
	20rat99	37	0.82	0.00%	37	1.61	0.00%	1905	0.75	0.00%	1905	0.59	0.00%
	20rd100	45	5.20	0.00%	45	5.62	0.00%	2228	13.91	0.00%	2228	11.38	0.00%
	21eil101	67	43.13	0.00%	67	12.20	0.00%	3365	56.74	0.00%	3365	15.78	0.00%
	21lin105	50	16.20	0.00%	50	31.96	0.00%	2489	12.82	0.00%	2489	12.89	0.00%
	22pr107	41	0.03	0.00%	41	0.04	0.00%	2123	0.05	0.00%	2123	0.10	0.00%
	25pr124	46	2130.68	0.00%	46	105.37	0.00%	2302	3516.36	0.00%	2302	175.17	0.00%
	26bier127	109	3759.00	8.49%	110	924.20	0.00%	5069	3685.02	15.49%	5420	2739.41	0.00%
	26ch130	70	3731.59	16.82%	70	499.02	0.00%	3423	3423.03	0.00%	3423	878.49	0.00%
	28pr136	53	281.85	0.00%	53	32.23	0.00%	2699	438.36	0.00%	2699	312.89	0.00%
	29pr144	6	3662.96	94.06%	60	1535.54	0.00%	3055	3767.94	39.15%	3055	1622.75	0.00%
	30ch150	61	3726.14	4.31%	61	570.16	0.00%	3131	1665.39	0.00%	3131	514.97	0.00%
	30kroA150	58	3742.55	28.62%	58	617.23	0.00%	3039	3732.85	12.73%	3039	742.87	0.00%
	30kroB150	66	3719.59	10.49%	66	343.15	0.00%	3172	3727.47	23.22%	3172	1933.77	0.00%
	31pr152	9	3651.60	91.43%	57	882.45	0.00%	2440	3649.65	54.71%	2915	1433.18	0.00%
	32u159	76	1710.73	0.00%	76	1296.21	0.00%	4002	2432.60	0.00%	4002	547.70	0.00%
	39rat195	71	1308.45	0.00%	71	287.79	0.00%	3656	975.52	0.00%	3656	251.21	0.00%
	40d198	70	501.47	0.00%	70	85.84	0.00%	3595	532.26	0.00%	3595	131.35	0.00%
Set2	11berlin52	50	0.85	0.00%	50	0.92	0.00%	2584	0.65	0.00%	2584	0.86	0.00%
	11eil51	37	0.32	0.00%	37	2.29	0.00%	1929	0.22	0.00%	1929	0.59	0.00%
	14st70	56	1.77	0.00%	56	0.76	0.00%	2736	1.67	0.00%	2736	0.84	0.00%
	16eil76	51	3.25	0.00%	51	2.70	0.00%	2518	5.58	0.00%	2518	10.63	0.00%
	16pr76	70	115.29	0.00%	70	129.47	0.00%	3550	107.33	0.00%	3550	30.61	0.00%
	20kroA100	80	1159.43	0.00%	80	40.97	0.00%	3894	681.68	0.00%	3894	54.43	0.00%
	20kroB100	86	533.20	0.00%	86	50.21	0.00%	4357	542.72	0.00%	4357	394.08	0.00%
	20kroC100	72	113.76	0.00%	72	27.22	0.00%	3586	175.16	0.00%	3586	95.66	0.00%
	20kroD100	78	24.67	0.00%	78	9.97	0.00%	3799	93.56	0.00%	3799	31.96	0.00%
	20kroE100	90	144.30	0.00%	90	7.54	0.00%	4614	21.82	0.00%	4614	27.59	0.00%
	20rat99	73	0.32	0.00%	73	1.69	0.00%	3624	1.06	0.00%	3624	42.50	0.00%
	20rd100	82	40.18	0.00%	82	27.52	0.00%	4181	39.30	0.00%	4181	28.32	0.00%
	21eil101	83	46.59	0.00%	83	29.75	0.00%	4264	71.45	0.00%	4264	45.93	0.00%
	21lin105	95	633.97	0.00%	95	329.39	0.00%	4814	748.22	0.00%	4814	339.21	0.00%
	22pr107	94	9.54	0.00%	94	13.10	0.00%	4740	69.03	0.00%	4740	18.69	0.00%
	25pr124	90	3624.48	25.62%	101	697.18	0.00%	4334	3622.27	28.41%	3859	3623.57	36.26%
	26bier127	11	3653.82	91.27%	124	3654.25	1.59%	6236	3673.75	1.53%	6004	3636.61	5.20%
	26ch130	6	3627.94	95.35%	111	2190.43	0.00%	250	3626.31	96.16%	5354	3635.65	16.25%
	28pr136	120	2239.39	0.00%	120	36.30	0.00%	6106	1609.05	0.00%	6106	146.88	0.00%
	29pr144	24	3636.55	83.22%	4	3629.55	97.20%	166	3629.58	97.71%	166	3626.60	97.71%
	30ch150	111	3627.12	25.50%	114	525.90	0.00%	124	3635.22	98.35%	6025	995.18	0.00%
	30kroA150	11	3624.33	92.62%	99	3631.89	33.56%	141	3628.84	98.13%	4478	3634.65	39.94%
	30kroB150	9	3629.74	93.92%	117	3640.74	19.16%	171	3627.88	97.73%	6190	3625.89	17.73%
	31pr152	89	3631.77	40.67%	9	3629.77	94.00%	431	3637.55	94.34%	431	3631.17	94.34%
	32u159	143	3627.89	7.14%	143	428.88	0.00%	7507	3620.71	4.37%	7507	919.53	0.00%
	39rat195	135	750.97	0.00%	135	472.05	0.00%	6813	1201.13	0.00%	6813	292.75	0.00%
	40d198	149	3432.08	0.00%	149	2728.54	0.00%	6700	3630.10	26.78%	7480	303.29	0.00%
AVG		1370.33	14.99%		615.23	4.55%		1360.35	14.61%		751.66	5.69%	
#Opt			38			49			39			47	

Table 6: Comparison between **clucut** and BC on the *Set1* and *Set2* instances with  $\omega = 0.4$ .

$\omega = 0.6$													
Instance	$g_1$						$g_2$						
	Sol	Time	Gap%	Sol	Time	Gap%	Sol	Time	Gap%	Sol	Time	Gap%	
Set1	11berlin52	43	3.41	0.00%	43	3.88	0.00%	2190	2.12	0.00%	2190	3.72	0.00%
	11eil51	39	0.53	0.00%	39	1.81	0.00%	1911	0.60	0.00%	1911	1.21	0.00%
	14st70	50	80.90	0.00%	50	19.87	0.00%	2589	33.47	0.00%	2589	18.62	0.00%
	16eil76	59	59.75	0.00%	59	8.02	0.00%	3119	63.80	0.00%	3119	19.60	0.00%
	16pr76	65	59.88	0.00%	65	124.85	0.00%	3275	1089.40	0.00%	3275	182.03	0.00%
	20kroA100	65	1573.50	0.00%	65	106.66	0.00%	3192	1524.53	0.00%	3192	135.17	0.00%
	20kroB100	59	3631.79	36.39%	66	97.08	0.00%	3203	1723.12	0.00%	3203	165.41	0.00%
	20kroC100	62	424.37	0.00%	62	72.43	0.00%	3110	1321.52	0.00%	3110	246.28	0.00%
	20kroD100	64	1901.65	0.00%	64	75.22	0.00%	3133	1740.32	0.00%	3133	81.68	0.00%
	20kroE100	63	84.31	0.00%	63	166.56	0.00%	2950	237.76	0.00%	2950	82.57	0.00%
	20rat99	52	102.58	0.00%	52	46.19	0.00%	2643	66.64	0.00%	2643	41.06	0.00%
	20rd100	72	321.63	0.00%	72	368.21	0.00%	3591	265.07	0.00%	3591	133.94	0.00%
	21eil101	82	762.19	0.00%	82	82.72	0.00%	4187	602.27	0.00%	4187	412.53	0.00%
	21lin105	78	472.41	0.00%	78	132.16	0.00%	3955	1087.02	0.00%	3955	162.10	0.00%
	22pr107	53	3622.72	36.14%	53	3625.48	31.17%	2697	3626.63	34.73%	2697	3626.40	30.44%
Set2	11berlin52	51	0.11	0.00%	51	0.14	0.00%	2608	0.11	0.00%	2608	0.14	0.00%
	11eil51	50	0.57	0.00%	50	3.76	0.00%	2575	0.51	0.00%	2575	0.56	0.00%
	14st70	64	1360.19	0.00%	64	300.55	0.00%	3218	2602.34	0.00%	3218	502.19	0.00%
	16eil76	74	353.70	0.00%	74	175.62	0.00%	3728	89.88	0.00%	3728	100.66	0.00%
	16pr76	74	3620.91	1.33%	74	1777.29	0.00%	3729	3334.55	0.00%	3729	481.00	0.00%
	20kroA100	91	3625.62	8.08%	95	3623.23	4.04%	3763	3628.82	24.86%	4554	3621.34	9.07%
	20kroB100	93	3630.43	6.06%	2	3621.31	97.98%	3578	3630.00	28.55%	4668	3623.16	6.79%
	20kroC100	5	3626.30	94.95%	90	3618.66	9.09%	3915	3622.41	21.83%	4534	3619.27	9.46%
	20kroD100	4	3624.16	95.96%	93	3618.94	6.06%	4394	3628.09	12.26%	4570	3618.92	8.75%
	20kroE100	97	3620.09	2.02%	97	2215.67	0.00%	4910	3621.33	1.96%	4910	3617.95	1.96%
	20rat99	87	140.58	0.00%	87	194.56	0.00%	4516	70.56	0.00%	4516	154.69	0.00%
	20rd100	97	3626.93	2.02%	99	2135.44	0.00%	5008	2876.95	0.00%	4957	3619.27	1.02%
	21eil101	95	3622.50	5.00%	97	962.97	0.00%	4933	3621.81	2.32%	4933	3622.82	2.32%
	21lin105	102	3641.03	1.92%	104	781.43	0.00%	5075	3629.93	2.93%	5103	3627.18	2.39%
	22pr107	106	225.08	0.00%	106	11.02	0.00%	5363	28.98	0.00%	5363	134.78	0.00%
AVG		1593.99	9.66%		932.39	4.94%		1592.35	4.31%		1188.54	2.41%	
#Opt			19			25			22			21	

Table 7: Comparison between **clucut** and BC on the *Set1* and *Set2* instances with  $\omega = 0.6$ .

solves 38 instances for  $g_1$  and 39 for  $g_2$ . It is worth noting that BC optimally solves all the instances of *Set1* while it does not find the optimal solution 12 times on the *Set2* dataset. This shows that, for our algorithm, the instances of *Set2* are more difficult to solve than the ones of *Set1*. We recall that the only difference is that customers are randomly assigned to clusters in *Set2*, while they are geographically clustered in *Set1*.

Similar considerations can be done for the results reported in Table 7 which are related to  $\omega = 0.6$ . Here we report the results for instances with up to 107 customers as no meaningful result was obtained for larger sizes. BC solves 25 instances to optimality (out of 30) for  $g_1$  and 21 for  $g_2$ , while **clucut** solves 19 and 22, respectively. Also, as before, BC performs much better in terms of both computational time and optimality gap. More in detail, the Gap% value for BC is equal to 4.94% and 2.41% for  $g_1$  and  $g_2$ , respectively,

$\omega = 0.8$													
Instance	$g_1$						$g_2$						
	Sol	Time	Gap%	Sol	Time	Gap%	Sol	Time	Gap%	Sol	Time	Gap%	
Set1	11berlin52	47	54.39	0.00%	47	6.36	0.00%	2384	11.80	0.00%	2384	12.23	0.00%
	11eil51	43	4.06	0.00%	43	2.10	0.00%	2114	6.50	0.00%	2114	6.94	0.00%
	14st70	65	773.70	0.00%	65	361.26	0.00%	3355	488.31	0.00%	3355	533.94	0.00%
	16eil76	69	476.84	0.00%	69	153.38	0.00%	3573	1334.32	0.00%	3573	85.04	0.00%
	16pr76	72	3618.59	1.37%	72	1653.24	0.00%	3611	3625.51	2.58%	3611	3619.83	2.25%
	20kroA100	73	3630.08	26.26%	79	224.73	0.00%	2713	3631.67	45.83%	4115	3034.79	0.00%
	20kroB100	77	3636.04	22.22%	86	2802.72	0.00%	4188	3627.16	16.37%	4117	3637.40	16.25%
	20kroC100	76	3630.97	23.23%	83	417.48	0.00%	3999	3624.35	20.15%	3999	287.95	0.00%
	20kroD100	77	3633.75	22.22%	85	430.87	0.00%	3854	3630.75	23.04%	4026	3625.22	19.61%
	20kroE100	78	3627.27	18.75%	80	347.91	0.00%	4002	2849.01	0.00%	4002	388.90	0.00%
	20rat99	69	3632.64	21.59%	79	1810.70	0.00%	3855	3625.82	13.06%	3992	2785.76	0.00%
	20rd100	90	3627.83	9.09%	91	813.84	0.00%	3892	3634.55	22.28%	4640	3626.99	7.35%
	21eil101	89	3630.17	11.00%	91	319.53	0.00%	4538	3635.25	10.14%	4717	1786.74	0.00%
	21lin105	87	3639.87	16.35%	90	289.91	0.00%	4245	3648.18	18.80%	4561	3638.53	10.43%
	22pr107	6	3637.27	94.34%	53	3645.25	50.00%	2156	3639.16	59.80%	2697	3636.09	49.71%
Set2	11berlin52	51	0.02	0.00%	51	0.03	0.00%	2608	0.06	0.00%	2608	0.07	0.00%
	11eil51	50	1.13	0.00%	50	0.79	0.00%	2575	0.45	0.00%	2575	0.51	0.00%
	14st70	69	7.23	0.00%	69	3.57	0.00%	3513	22.52	0.00%	3513	13.25	0.00%
	16eil76	75	12.27	0.00%	75	4.03	0.00%	3800	3.31	0.00%	3800	167.70	0.00%
	16pr76	75	1997.86	0.00%	75	8.15	0.00%	3800	2003.37	0.00%	3800	670.13	0.00%
	20kroA100	99	331.47	0.00%	99	9.73	0.00%	4086	3627.77	18.41%	4241	3623.81	15.32%
	20kroB100	69	3634.85	30.30%	99	1207.78	0.00%	83	3633.52	98.34%	4668	3624.05	6.79%
	20kroC100	4	3631.41	95.96%	94	3623.28	5.05%	249	3641.88	95.03%	3043	3622.69	39.24%
	20kroD100	5	3630.36	94.95%	95	3631.99	4.04%	3750	3624.62	25.12%	4776	3634.34	4.63%
	20kroE100	97	3624.73	2.02%	98	3620.10	1.01%	325	3631.13	93.51%	325	3627.43	93.51%
	20rat99	98	1335.88	0.00%	98	155.44	0.00%	5007	531.44	0.00%	5007	323.20	0.00%
	20rd100	99	134.14	0.00%	99	135.97	0.00%	5008	24.29	0.00%	5008	45.00	0.00%
	21eil101	99	3623.63	1.00%	100	1569.38	0.00%	4831	3628.04	4.34%	4933	3629.89	2.32%
	21lin105	104	3.76	0.00%	104	4.51	0.00%	5228	1953.40	0.00%	5228	1541.45	0.00%
	22pr107	106	11.65	0.00%	106	8.74	0.00%	5363	62.54	0.00%	5363	138.40	0.00%
AVG		2107.80	16.36%		908.76	2.00%		2246.69	18.89%		1845.61	8.91%	
#Opt			14			26			14			18	

Table 8: Comparison between **clucut** and BC on the *Set1* and *Set2* instances with  $\omega = 0.8$ .

while it is around double for **clucut**. Regarding the computational time, BC is nearly 41% faster than **clucut** for  $g_1$  and 25% faster for  $g_2$ . It is interesting to note that instances with  $\omega = 0.6$  are more difficult to solve than those with  $\omega = 0.4$ . Indeed, for both algorithms, the number of instances optimally solved decreases while the computational times increase. This was expected: the larger is the value of  $\omega$ , the larger is the set of feasible solutions so the more difficult is the problem to solve.

The results of Table 8, which are related to  $\omega = 0.8$ , show that the performance of **clucut** largely deteriorates with respect to the case  $\omega = 0.6$ . In fact, the number of instances solved to optimality reduces to 14 for both  $g_1$  and  $g_2$  while the average optimality gap increases to 16.36% for  $g_1$  and 18.89% for  $g_2$ . There is also an increase of the computational time of **clucut**, which now exceeds 2100 seconds. On the contrary, the results of BC are more stable. Indeed it solves 26 instances to optimality for  $g_1$  and 18 for  $g_2$ . Once again

BC largely outperforms **clucut** in both computational time and optimality gap at termination. It is worth noting that, when going from  $\omega = 0.6$  to  $\omega = 0.8$ , the average time of BC with profit  $g_1$  does not significantly change and the optimality gap reduces from 4.94% to 2.00%. On the contrary, for profit  $g_2$ , both the computational time and the optimality gap remarkably increase. Therefore, when going from  $\omega = 0.6$  to  $\omega = 0.8$ , there is a significant difference in the behaviour of the algorithm with respect to the type of profit considered.

Thus, by summarizing all results in Tables 6–8, except for cluster cover inequalities, we can conclude that the valid inequalities described in Section 3 greatly improves the efficacy of the exact algorithm.

### 5.3 Comparison with Benchmark Approaches

In this section, we compare BC with two exact approaches proposed in the literature, specifically, the one proposed in [2] (called **ACC** from now on) and the one proposed in [20] (called **PFS**). It is worth noting that BC and **ACC** were executed on the same machine and then their CPU times are directly comparable. For **PFS**, we scale the CPU time according to the processor performance.

In Tables 9 and 10 we compare the solutions found by BC with the **ACC** algorithm proposed in [2] on the dataset *Set1* and *Set2*, respectively. The formulation used in **ACC** is a polynomial formulation where subtour elimination constraints are modelled as Miller-Tucker-Zemlin (MTZ) constraints. The formulation uses three types of binary variables: arc variables, vertex visiting variables and cluster visiting variables. The results are restricted to the instances with 100 customers at most as **ACC** was not capable of providing any solution within the time limit for larger instances. The first three columns of the Tables report the name of the instance (*Instance*), the  $\omega$  value ( $\omega$ ) and the type of profit ( $p_g$ ). The next six columns are grouped in two parts referring to the results by **ACC**, first, and BC, second. Specifically, we report the value of the best solution found at termination (*Sol*), the computational time (*Time*), in seconds, and the optimality gap (*Gap%*) at termination (which is 0% in case the instance is solved to proven optimality). The last two rows of each table report the average computational time and the average optimality gap, and the number of instances solved to optimality, respectively.

The results of Table 9 highlight that BC is much better than **ACC** from

	Instance	$\omega$	$p_g$	ACC			BC		
				Sol	Time	Gap%	Sol	Time	Gap%
Set1	11berlin52	0.4	$g_1$	37	15.52	0.00%	37	0.41	0.00%
	11berlin52	0.4	$g_2$	1829	18.14	0.00%	1829	0.46	0.00%
	11berlin52	0.6	$g_1$	43	646.07	0.00%	43	3.88	0.00%
	11berlin52	0.6	$g_2$	2190	540.54	0.00%	2190	3.72	0.00%
	11berlin52	0.8	$g_1$	47	1877.97	0.00%	47	6.36	0.00%
	11berlin52	0.8	$g_2$	2384	1207.51	0.00%	2384	12.23	0.00%
	11eil51	0.4	$g_1$	24	12.51	0.00%	24	0.21	0.00%
	11eil51	0.4	$g_2$	1279	15.90	0.00%	1279	1.03	0.00%
	11eil51	0.6	$g_1$	39	9.38	0.00%	39	1.81	0.00%
	11eil51	0.6	$g_2$	1911	75.43	0.00%	1911	1.21	0.00%
	11eil51	0.8	$g_1$	43	1987.21	0.00%	43	2.10	0.00%
	11eil51	0.8	$g_2$	2114	1673.87	0.00%	2114	6.94	0.00%
	14st70	0.4	$g_1$	33	3610.11	18.98%	33	0.48	0.00%
	14st70	0.4	$g_2$	1672	3610.11	14.62%	1672	0.48	0.00%
	14st70	0.6	$g_1$	50	3610.11	14.35%	50	19.87	0.00%
	14st70	0.6	$g_2$	2589	3610.11	14.18%	2589	18.62	0.00%
	14st70	0.8	$g_1$	64	3610.10	5.88%	65	361.26	0.00%
	14st70	0.8	$g_2$	3229	3610.10	7.43%	3355	533.94	0.00%
	16eil76	0.4	$g_1$	40	2161.31	0.00%	40	2.54	0.00%
	16eil76	0.4	$g_2$	2223	2142.87	0.00%	2223	1.74	0.00%
	16eil76	0.6	$g_1$	59	3610.10	9.41%	59	8.02	0.00%
	16eil76	0.6	$g_2$	3119	3610.10	8.63%	3119	19.60	0.00%
	16eil76	0.8	$g_1$	67	3610.10	8.22%	69	153.38	0.00%
	16eil76	0.8	$g_2$	3525	3610.11	6.60%	3573	85.04	0.00%
	16pr76	0.4	$g_1$	47	3610.10	20.18%	47	5.19	0.00%
	16pr76	0.4	$g_2$	2449	3609.80	18.89%	2449	3.75	0.00%
	16pr76	0.6	$g_1$	65	3609.79	7.48%	65	124.85	0.00%
	16pr76	0.6	$g_2$	3275	3608.23	7.36%	3275	182.03	0.00%
	16pr76	0.8	$g_1$	71	3609.71	4.05%	72	1653.24	0.00%
	16pr76	0.8	$g_2$	3601	3608.42	3.43%	3611	3619.83	2.25%
	20kroA100	0.4	$g_1$	42	3609.53	26.60%	42	27.14	0.00%
	20kroA100	0.4	$g_2$	2151	3609.85	28.60%	2151	30.50	0.00%
	20kroA100	0.6	$g_1$	65	3609.89	18.68%	65	106.66	0.00%
	20kroA100	0.6	$g_2$	3164	3609.87	23.60%	3192	135.17	0.00%
	20kroA100	0.8	$g_1$	76	3610.03	23.23%	79	224.73	0.00%
	20kroA100	0.8	$g_2$	3919	3609.77	21.75%	4115	3034.79	0.00%
	AVG				2550.01	8.67%		288.70	0.06%
	#Opt					14			35

Table 9: Comparison between **ACC** and BC on the *Set1* instances.

the effectiveness and performance point of view. The AVG line shows that the average gap of BC is equal to 0.06% and 35 out of 36 instances are optimally solved. In the only case in which BC fails to find the optimal solution, the Gap% value is equal to 2.25%. Instead, the average gap of **ACC** goes up to 8.67% and the algorithm optimally solves only 14 instances within the time limit of one hour. It is worth noting that in 12 out of 22 instances not optimally solved by **ACC**, the Gap% value is greater than 14% with a peak

equal to 28.60%. Regarding the performance, with an average computational time equal to 288 seconds, BC is eight times faster than **ACC**. BC proves to be better than **ACC** even on dataset *Set2* but the gap is smaller here. Indeed, BC optimally solves 33 out of 36 instances while **ACC** solves 22 of them. The average gap is equal to 0.79% for BC and 3.49% for **ACC** and, about the performance, the average time of BC is equal to 426 seconds while it is 1644 seconds for **ACC**. Finally, in the worst case, the Gap% value is equal to 15.32% for BC and 24.50% for **ACC**.

The last comparison is carried out between BC and **PFS**. The formulation used in **PFS** involves only two types of binary variables, namely arc variables and vertex visiting variables. Subtour elimination constraints are dynamically separated. Moreover, a greedy construction procedure is used for creating an initial feasible solution which is used as a warm start for CPLEX. The results of the comparison are shown in Table 11 and are restricted to the instances tested in [20], i.e., 11berlin52, 11eil51, 14st70, 16eil76, for values of  $\omega$  equal to 0.4, 0.6 and 0.8. The headings of Table 11 are the same as in Table 10, the only difference being that the Gap% column is removed because all the instances are optimally solved by both algorithms. Indeed, in [20] only instances solved to optimality are reported. In order to have a fair comparative from the performance point of view, the CPU time reported in [20] has been scaled according to the scaling factor reported here: [https://www.cpubenchmark.net/CPU\\_mega\\_page.html](https://www.cpubenchmark.net/CPU_mega_page.html). From the results of Table 11, we observe that BC is around 17% faster than **PFS**. A detailed analysis reveals that in 14 out of 20 instances, BC is faster than **PFS**. Essentially, **PFS** wins on the smallest instances where probably the branch-and-cut pays on overhead associated with the separation of the valid inequalities. This is certified by the results of the two algorithms on the last 11 instances of the table, the largest ones, where BC is faster than **PFS** in 10 cases.

## 6 Conclusions

This paper deals with the Set Orienteering Problem (SOP), which is a variant of the Orienteering Problem recently introduced in the literature. Specifically, in the SOP, customers are grouped in clusters and the profit associated with a cluster is collected in case at least one customer from the cluster is visited. The SOP finds interesting applications in practice, especially related to mass distribution products. In general, routing problems with profits are

	Instance	$\omega$	$p_g$	ACC			BC		
				Sol	Time	Gap%	Sol	Time	Gap%
Set2	11berlin52	0.4	$g_1$	50	30.90	0.00%	50	0.92	0.00%
	11berlin52	0.4	$g_2$	2584	53.99	0.00%	2584	0.86	0.00%
	11berlin52	0.6	$g_1$	51	1.01	0.00%	51	0.14	0.00%
	11berlin52	0.6	$g_2$	2608	1.15	0.00%	2608	0.14	0.00%
	11berlin52	0.8	$g_1$	51	0.51	0.00%	51	0.03	0.00%
	11berlin52	0.8	$g_2$	2608	0.66	0.00%	2608	0.07	0.00%
	11eil51	0.4	$g_1$	37	3.98	0.00%	37	2.29	0.00%
	11eil51	0.4	$g_2$	1929	10.25	0.00%	1929	0.59	0.00%
	11eil51	0.6	$g_1$	50	17.71	0.00%	50	3.76	0.00%
	11eil51	0.6	$g_2$	2575	10.25	0.00%	2575	0.56	0.00%
	11eil51	0.8	$g_1$	50	4.72	0.00%	50	0.79	0.00%
	11eil51	0.8	$g_2$	2575	1.55	0.00%	2575	0.51	0.00%
	14st70	0.4	$g_1$	56	3609.95	12.50%	56	0.76	0.00%
	14st70	0.4	$g_2$	2736	3609.76	16.11%	2736	0.84	0.00%
	14st70	0.6	$g_1$	64	3609.40	7.25%	64	300.55	0.00%
	14st70	0.6	$g_2$	3218	3609.79	8.40%	3218	502.19	0.00%
	14st70	0.8	$g_1$	69	77.66	0.00%	69	3.57	0.00%
	14st70	0.8	$g_2$	3513	130.22	0.00%	3513	13.25	0.00%
	16eil76	0.4	$g_1$	51	1515.90	0.00%	51	2.70	0.00%
	16eil76	0.4	$g_2$	2518	3610.01	6.66%	2518	10.63	0.00%
	16eil76	0.6	$g_1$	74	1767.92	0.00%	74	175.62	0.00%
	16eil76	0.6	$g_2$	3728	1183.11	0.00%	3728	100.66	0.00%
	16eil76	0.8	$g_1$	75	249.92	0.00%	75	4.03	0.00%
	16eil76	0.8	$g_2$	3800	65.56	0.00%	3800	167.70	0.00%
	16pr76	0.4	$g_1$	70	3610.10	3.65%	70	129.47	0.00%
	16pr76	0.4	$g_2$	3402	3610.09	8.52%	3550	30.61	0.00%
	16pr76	0.6	$g_1$	74	3434.82	0.00%	74	1777.29	0.00%
	16pr76	0.6	$g_2$	3729	3610.10	1.66%	3729	481.00	0.00%
	16pr76	0.8	$g_1$	75	69.28	0.00%	75	8.15	0.00%
	16pr76	0.8	$g_2$	3800	21.90	0.00%	3800	670.13	0.00%
	20kroA100	0.4	$g_1$	75	3610.12	24.24%	80	40.97	0.00%
	20kroA100	0.4	$g_2$	3781	3610.11	24.50%	3894	54.43	0.00%
	20kroA100	0.6	$g_1$	95	3610.12	4.04%	95	3623.23	4.04%
	20kroA100	0.6	$g_2$	4741	3610.12	5.33%	4554	3621.34	9.07%
	20kroA100	0.8	$g_1$	98	3610.12	1.01%	99	9.73	0.00%
	20kroA100	0.8	$g_2$	4920	3610.12	1.76%	4241	3623.81	15.32%
	AVG				1644.25	3.49%		426.76	0.79%
	#Opt					22			33

Table 10: Comparison between **ACC** and BC on the *Set2* instances.

facing a new wave of interest from the research community thanks to their link with fast delivery services, where it often happens that not all customers requiring a service can be satisfied, thus a selection is needed (which is the main feature of routing problems with profits). Specifically, the SOP finds applications in delivery services where multiple options are associated with each customer regarding where a parcel can be delivered.

We focus on developing an exact algorithm for the SOP. Specifically, we



	Instance	$\omega$	$p_g$	PFS		BC	
				Sol	Time	Sol	Time
Set1	11berlin52	0.4	$g_1$	37	0.78	37	0.41
	11berlin52	0.4	$g_2$	1829	0.85	1829	0.46
	11berlin52	0.6	$g_1$	43	3.05	43	3.88
	11berlin52	0.6	$g_2$	2190	0.96	2190	3.72
	11berlin52	0.8	$g_1$	47	3.33	47	6.36
	11berlin52	0.8	$g_2$	2384	5.52	2384	12.23
	11eil51	0.4	$g_1$	24	1.83	24	0.21
	11eil51	0.4	$g_2$	1279	2.02	1279	1.03
	11eil51	0.6	$g_1$	39	1.20	39	1.81
	11eil51	0.6	$g_2$	1911	2.17	1911	1.21
	11eil51	0.8	$g_1$	43	11.88	43	2.10
	11eil51	0.8	$g_2$	2114	29.01	2114	6.94
	14st70	0.4	$g_1$	33	11.98	33	0.48
	14st70	0.4	$g_2$	1672	20.50	1672	0.48
	14st70	0.8	$g_1$	65	690.35	65	361.26
	14st70	0.8	$g_2$	3355	164.63	3355	533.94
	16eil76	0.4	$g_1$	40	62.00	40	2.54
	16eil76	0.4	$g_2$	2223	27.01	2223	1.74
	16eil76	0.6	$g_1$	59	46.27	59	8.02
	16eil76	0.6	$g_2$	3119	78.24	3119	19.60
AVG				58.18		48.42	

Table 11: Comparison between **PFS** and BC

propose a new formulation for the problem that has fewer variables than those proposed in the literature as it does not include vertex visiting variables. We show that the new formulation has a stronger relaxation than the formulation with vertex visiting variables. Then, we propose different classes of valid inequalities to strengthen the formulation. Exhaustive computational tests show that the resulting branch-and-cut algorithm is effective. The performance depends on the number of customers and on the value of  $\omega$ , which measures the maximum route length. Specifically, when  $\omega$  is small (and the vehicle route is short), the branch-and-cut algorithm is able to solve almost all instances with up to 200 vertices. For larger values of  $\omega$ , the algorithm fails to solve instances with more than 107 customers, while it solves the majority of those with fewer customers. Also, a comparison with two exact approaches proposed in the literature shows that our branch-and-cut algorithm scales better in terms of number of customers and can thus be considered as the new state-of-the art exact solution approach for the SOP.

As a future research direction, we plan to study the case in which multiple

vehicles are available and to adapt the branch-and-cut algorithm to this case. Also, a column generation approach might be suitable in the multiple vehicle case. Finally, it might be interesting to investigate whether similar formulations as the one proposed in this paper (which get rid of the vertex visiting variables) are effective in the solution of similar problems, as, for example, the Cluster Orienteering Problem.

## Appendix A Detailed Results

In this section we show the detailed results of the comparison among the branch-and-cut algorithm with the full set of inequalities and the versions in which we discard a single inequality at a time. These results are presented in Tables 12–15. The first column (*Instance*) shows the name of the instance. Then, five groups of columns follow, corresponding to the five versions of the branch-and-cut algorithm: the version with the full set of inequalities (*C-BC*), no conditional cuts (*NoCond*), no cover inequalities (*NoCover*), no cluster cover inequalities (*NoCluCover*) and no path inequalities (*NoPath*). For each version of the algorithm, we report the value of the upper bound (*UB*) at termination (which corresponds to the value of the optimal solution in case the computational time is lower than one hour) and the computational time (*Time*). Also, for each version of the algorithm in which one inequality is excluded, we report the percentage gap of the corresponding upper bound with respect to the upper bound obtained by C-BC, calculated as  $gap = \frac{UB_{C-BC} - UB_*}{UB_{C-BC}}$ , where  $UB_{C-BC}$  and  $UB_*$  are the upper bound of C-BC and of the version considered, respectively. Note that positive values of the gap mean that the version of the algorithm without the inequality gives a better result than C-BC. The second last row of each table reports the average values of time and gap while the last row report, for each version of the algorithm in which an inequality is excluded, the number of times in which the upper bound provided by that version is worse than the one provided by C-BC. These two rows are the same as the ones reported in Table 1.

$\omega = 0.4$ and $g_1$															
Instance	C-BC		NoCond			NoCover			NoCluCover			NoPath			
	UB	Time	UB	Time	Gap%	UB	Time	Gap%	UB	Time	Gap%	UB	Time	Gap%	
Set1	11berlin52	37.0	0.39	37.0	0.52	0.00%	37.0	0.65	0.00%	37.0	0.41	0.00%	37.0	1.19	0.00%
	11eil51	24.0	0.34	24.0	0.34	0.00%	24.0	0.29	0.00%	24.0	0.21	0.00%	24.0	0.40	0.00%
	14st70	33.0	0.48	33.0	0.49	0.00%	33.0	0.58	0.00%	33.0	0.48	0.00%	33.0	0.70	0.00%
	16eil76	40.0	2.54	40.0	2.59	0.00%	40.0	1.52	0.00%	40.0	2.54	0.00%	40.0	4.24	0.00%
	16pr76	47.0	5.08	47.0	5.21	0.00%	47.0	8.04	0.00%	47.0	5.19	0.00%	47.0	7.46	0.00%
	20kroA100	42.0	27.09	42.0	27.97	0.00%	42.0	33.37	0.00%	42.0	27.14	0.00%	42.0	38.52	0.00%
	20kroB100	49.0	11.27	49.0	10.04	0.00%	49.0	15.13	0.00%	49.0	8.32	0.00%	49.0	22.92	0.00%
	20kroC100	42.0	2.05	42.0	2.04	0.00%	42.0	3.11	0.00%	42.0	2.07	0.00%	42.0	4.62	0.00%
	20kroD100	39.0	3.04	39.0	2.85	0.00%	39.0	4.26	0.00%	39.0	2.94	0.00%	39.0	7.11	0.00%
	20kroE100	52.0	5.07	52.0	4.39	0.00%	52.0	3.93	0.00%	52.0	3.68	0.00%	52.0	5.67	0.00%
	20rat99	37.0	1.60	37.0	0.78	0.00%	37.0	1.04	0.00%	37.0	1.61	0.00%	37.0	1.13	0.00%
	20rd100	45.0	5.99	45.0	5.76	0.00%	45.0	7.86	0.00%	45.0	5.62	0.00%	45.0	8.87	0.00%
	21eil101	67.0	23.12	67.0	30.11	0.00%	67.0	22.93	0.00%	67.0	12.20	0.00%	67.0	48.34	0.00%
	21lin105	50.0	33.20	50.0	6.06	0.00%	50.0	38.34	0.00%	50.0	31.96	0.00%	50.0	15.34	0.00%
	22pr107	41.0	0.04	41.0	0.04	0.00%	41.0	0.03	0.00%	41.0	0.04	0.00%	41.0	0.04	0.00%
	25pr124	46.0	136.08	46.0	114.73	0.00%	46.0	2773.53	0.00%	46.0	105.37	0.00%	46.0	2443.04	0.00%
	26bier127	110.0	1559.50	110.0	1019.16	0.00%	113.4	3730.33	-3.05%	110.0	924.20	0.00%	120.1	3663.76	-9.19%
	26ch130	70.0	158.72	70.0	159.48	0.00%	70.0	496.45	0.00%	70.0	499.02	0.00%	70.0	2568.43	0.00%
	28pr136	53.0	39.16	53.0	37.48	0.00%	53.0	324.79	0.00%	53.0	32.23	0.00%	53.0	324.88	0.00%
	29pr144	60.0	616.99	60.0	875.78	0.00%	60.0	2448.41	0.00%	60.0	1535.54	0.00%	100.0	3724.15	-66.71%
	30ch150	61.0	703.97	61.0	549.95	0.00%	63.0	3744.96	-3.28%	61.0	570.16	0.00%	71.5	3732.29	-17.16%
	30kroA150	58.0	1061.34	58.0	1058.20	0.00%	58.0	1722.95	0.00%	58.0	617.23	0.00%	105.9	3803.33	-82.52%
	30kroB150	66.0	362.01	66.0	371.70	0.00%	66.0	2126.92	0.00%	66.0	343.15	0.00%	132.0	3634.76	-100.00%
	31pr152	57.0	877.35	57.0	2122.48	0.00%	105.0	3671.32	-84.21%	57.0	882.45	0.00%	105.0	3678.55	-84.21%
	32u159	76.0	1373.83	76.0	1267.41	0.00%	76.0	1784.42	0.00%	76.0	1296.21	0.00%	76.0	2652.70	0.00%
	39rat195	71.0	291.73	71.0	839.51	0.00%	71.0	2104.62	0.00%	71.0	287.79	0.00%	71.0	2021.79	0.00%
	40d198	70.0	85.51	70.0	101.60	0.00%	70.0	346.87	0.00%	70.0	85.84	0.00%	70.0	902.71	0.00%
Set2	11berlin52	50.0	1.00	50.0	1.04	0.00%	50.0	0.59	0.00%	50.0	0.92	0.00%	50.0	1.07	0.00%
	11eil51	37.0	0.61	37.0	0.66	0.00%	37.0	0.56	0.00%	37.0	0.29	0.00%	37.0	0.72	0.00%
	14st70	56.0	0.78	56.0	0.77	0.00%	56.0	2.57	0.00%	56.0	0.76	0.00%	56.0	1.66	0.00%
	16eil76	51.0	7.95	51.0	8.11	0.00%	51.0	3.49	0.00%	51.0	2.70	0.00%	51.0	4.32	0.00%
	16pr76	70.0	135.46	70.0	146.88	0.00%	70.0	118.49	0.00%	70.0	129.47	0.00%	70.0	165.75	0.00%
	20kroA100	80.0	41.44	80.0	23.44	0.00%	80.0	1852.62	0.00%	80.0	40.97	0.00%	80.0	1038.06	0.00%
	20kroB100	86.0	50.14	86.0	71.47	0.00%	86.0	910.49	0.00%	86.0	50.21	0.00%	86.0	657.33	0.00%
	20kroC100	72.0	27.36	72.0	26.97	0.00%	72.0	137.27	0.00%	72.0	27.22	0.00%	72.0	116.40	0.00%
	20kroD100	78.0	10.11	78.0	10.13	0.00%	78.0	31.37	0.00%	78.0	9.97	0.00%	78.0	92.67	0.00%
	20kroE100	90.0	7.44	90.0	3.92	0.00%	90.0	255.75	0.00%	90.0	7.54	0.00%	90.0	68.45	0.00%
	20rat99	73.0	0.86	73.0	0.86	0.00%	73.0	6.50	0.00%	73.0	1.69	0.00%	73.0	0.67	0.00%
	20rd100	82.0	27.93	82.0	24.82	0.00%	82.0	138.71	0.00%	82.0	27.52	0.00%	82.0	57.48	0.00%
	21eil101	83.0	26.15	83.0	24.78	0.00%	83.0	33.38	0.00%	83.0	29.75	0.00%	83.0	49.03	0.00%
	21lin105	95.0	344.37	95.0	361.56	0.00%	95.0	648.17	0.00%	95.0	329.39	0.00%	95.0	1106.10	0.00%
	22pr107	94.0	13.22	94.0	13.92	0.00%	94.0	12.05	0.00%	94.0	13.10	0.00%	94.0	8.95	0.00%
	25pr124	101.0	698.51	101.0	696.02	0.00%	101.0	3249.68	0.00%	101.0	697.18	0.00%	121.0	3619.24	-19.80%
	26bier127	126.0	3653.99	126.0	3659.72	0.00%	126.0	3649.08	0.00%	126.0	3654.25	0.00%	126.0	3640.60	0.00%
	26ch130	111.0	2392.90	111.0	2345.38	0.00%	129.0	3628.53	-16.22%	111.0	2190.43	0.00%	129.0	3629.11	-16.22%
	28pr136	120.0	36.70	120.0	35.96	0.00%	120.0	2716.97	0.00%	120.0	36.30	0.00%	120.0	3097.60	0.00%
	29pr144	143.0	3629.13	143.0	3629.15	0.00%	143.0	3636.00	0.00%	143.0	3629.55	0.00%	143.0	3629.63	0.00%
	30ch150	114.0	575.59	114.0	578.36	0.00%	149.0	3631.51	-30.70%	114.0	525.90	0.00%	149.0	3625.85	-30.70%
	30kroA150	149.0	3632.46	110.0	2324.17	26.17%	149.0	3622.59	0.00%	149.0	3631.89	0.00%	149.0	3633.49	0.00%
	30kroB150	145.0	3640.89	120.0	3108.51	17.22%	148.0	3630.07	-2.10%	144.7	3640.74	0.16%	148.0	3629.60	-2.10%
	31pr152	150.0	3629.36	150.0	3629.78	0.00%	150.0	3630.57	0.00%	150.0	3629.77	0.00%	150.0	3627.45	0.00%
	32u159	143.0	426.69	143.0	428.30	0.00%	154.0	3629.74	-7.69%	143.0	428.88	0.00%	154.0	3621.58	-7.69%
	39rat195	135.0	504.49	135.0	501.76	0.00%	135.0	324.21	0.00%	135.0	472.05	0.00%	135.0	539.33	0.00%
	40d198	149.0	2711.98	149.0	2713.38	0.00%	149.0	1501.18	0.00%	149.0	2728.54	0.00%	149.0	521.36	0.00%
AVG	622.50		610.86		0.80%	1229.98		-2.73%	615.23		0.00%	1361.12		-8.08%	
#Worse					0			7			0			11	

Table 12: Computational results for the five versions of the branch-and-cut algorithm on the instances with  $\omega = 0.4$  and profit  $g_1$ .

$\omega = 0.4$ and $g_2$															
Instance	C-BC		NoCond			NoCover			NoCluCover			NoPath			
	UB	Time	UB	Time	Gap%	UB	Time	Gap%	UB	Time	Gap%	UB	Time	Gap%	
Set1	11berlin52	1829.0	0.59	1829.0	0.71	0.00%	1829.0	0.78	0.00%	1829.0	0.46	0.00%	1829.0	1.24	0.00%
	11eil51	1279.0	0.57	1279.0	0.64	0.00%	1279.0	0.52	0.00%	1279.0	1.03	0.00%	1279.0	0.77	0.00%
	14st70	1672.0	0.68	1672.0	0.66	0.00%	1672.0	0.81	0.00%	1672.0	0.48	0.00%	1672.0	0.88	0.00%
	16eil76	2223.0	5.77	2223.0	2.71	0.00%	2223.0	6.60	0.00%	2223.0	1.74	0.00%	2223.0	8.12	0.00%
	16pr76	2449.0	7.01	2449.0	8.01	0.00%	2449.0	11.52	0.00%	2449.0	3.75	0.00%	2449.0	7.97	0.00%
	20kroA100	2151.0	38.03	2151.0	53.20	0.00%	2151.0	42.72	0.00%	2151.0	30.50	0.00%	2151.0	111.96	0.00%
	20kroB100	2431.0	20.11	2431.0	19.48	0.00%	2431.0	11.65	0.00%	2431.0	15.81	0.00%	2431.0	43.78	0.00%
	20kroC100	2174.0	30.16	2174.0	29.84	0.00%	2174.0	11.85	0.00%	2174.0	4.64	0.00%	2174.0	4.54	0.00%
	20kroD100	1740.0	15.20	1740.0	15.15	0.00%	1740.0	21.36	0.00%	1740.0	7.84	0.00%	1740.0	13.49	0.00%
	20kroE100	2415.0	2.75	2415.0	2.76	0.00%	2415.0	10.36	0.00%	2415.0	5.18	0.00%	2415.0	6.87	0.00%
	20rat99	1905.0	1.24	1905.0	1.50	0.00%	1905.0	0.83	0.00%	1905.0	0.59	0.00%	1905.0	0.86	0.00%
	20rd100	2228.0	15.07	2228.0	10.21	0.00%	2228.0	7.85	0.00%	2228.0	11.38	0.00%	2228.0	11.71	0.00%
	21eil101	3365.0	34.06	3365.0	33.38	0.00%	3365.0	39.45	0.00%	3365.0	15.78	0.00%	3365.0	42.81	0.00%
	21lin105	2489.0	16.48	2489.0	26.24	0.00%	2489.0	14.38	0.00%	2489.0	12.89	0.00%	2489.0	21.65	0.00%
	22pr107	2123.0	0.11	2123.0	0.11	0.00%	2123.0	0.06	0.00%	2123.0	0.10	0.00%	2123.0	0.09	0.00%
	25pr124	2302.0	163.67	2302.0	167.95	0.00%	2302.0	2492.24	0.00%	2302.0	175.17	0.00%	4329.0	3628.61	-88.05%
	26bier127	5420.0	2264.48	5420.0	1362.98	0.00%	5420.0	2988.36	0.00%	5420.0	2739.41	0.00%	5973.6	3746.84	-10.21%
	26ch130	3423.0	1021.06	3423.0	495.95	0.00%	3423.0	2301.28	0.00%	3423.0	878.49	0.00%	3784.8	3729.98	-10.57%
	28spr136	2699.0	106.06	2699.0	106.76	0.00%	2699.0	160.09	0.00%	2699.0	312.89	0.00%	2699.0	450.63	0.00%
	29pr144	3055.0	2929.56	3055.0	2129.83	0.00%	3055.0	2643.36	0.00%	3055.0	1622.75	0.00%	5033.6	3753.32	-64.76%
	30ch150	3131.0	817.59	3131.0	954.14	0.00%	3131.0	2304.69	0.00%	3131.0	514.97	0.00%	3131.0	3606.56	0.00%
	30kroA150	3039.0	794.65	3039.0	840.56	0.00%	3039.0	1398.12	0.00%	3039.0	742.87	0.00%	4528.2	3741.84	-49.00%
	30kroB150	3172.0	1848.50	3172.0	1672.67	0.00%	3705.3	3744.28	-16.81%	3172.0	1933.77	0.00%	5722.4	3772.42	-80.40%
	31pr152	2915.0	1417.53	2915.0	1373.85	0.00%	5387.0	3649.19	-84.80%	2915.0	1433.18	0.00%	5387.0	3645.35	-84.80%
	32u159	4002.0	517.02	4002.0	420.29	0.00%	4002.0	280.93	0.00%	4002.0	547.70	0.00%	4002.0	994.63	0.00%
	39rat195	3656.0	454.65	3656.0	405.72	0.00%	3656.0	802.93	0.00%	3656.0	251.21	0.00%	3656.0	2780.70	0.00%
	40d198	3595.0	129.76	3595.0	135.22	0.00%	3595.0	608.51	0.00%	3595.0	131.35	0.00%	3595.0	437.66	0.00%
Set2	11berlin52	2584.0	0.85	2584.0	0.95	0.00%	2584.0	0.63	0.00%	2584.0	0.86	0.00%	2584.0	0.77	0.00%
	11eil51	1929.0	1.66	1929.0	1.60	0.00%	1929.0	2.39	0.00%	1929.0	0.59	0.00%	1929.0	0.65	0.00%
	14st70	2736.0	1.50	2736.0	1.39	0.00%	2736.0	1.74	0.00%	2736.0	0.84	0.00%	2736.0	1.99	0.00%
	16eil76	2518.0	32.22	2518.0	8.91	0.00%	2518.0	21.34	0.00%	2518.0	10.63	0.00%	2518.0	8.91	0.00%
	16pr76	3550.0	32.22	3550.0	31.50	0.00%	3550.0	117.36	0.00%	3550.0	30.61	0.00%	3550.0	301.43	0.00%
	20kroA100	3894.0	56.40	3894.0	63.12	0.00%	3894.0	641.03	0.00%	3894.0	54.43	0.00%	3894.0	1618.67	0.00%
	20kroB100	4357.0	458.83	4357.0	393.65	0.00%	4357.0	620.17	0.00%	4357.0	394.08	0.00%	4357.0	1675.12	0.00%
	20kroC100	3586.0	181.76	3586.0	104.91	0.00%	3586.0	366.21	0.00%	3586.0	95.66	0.00%	3586.0	222.14	0.00%
	20kroD100	3799.0	56.77	3799.0	62.48	0.00%	3799.0	149.19	0.00%	3799.0	31.96	0.00%	3799.0	273.25	0.00%
	20kroE100	4614.0	41.80	4614.0	10.00	0.00%	4614.0	38.30	0.00%	4614.0	27.59	0.00%	4614.0	43.07	0.00%
	20rat99	3624.0	12.99	3624.0	17.58	0.00%	3624.0	21.41	0.00%	3624.0	42.50	0.00%	3624.0	2.94	0.00%
	20rd100	4181.0	26.65	4181.0	32.46	0.00%	4181.0	79.52	0.00%	4181.0	28.32	0.00%	4181.0	42.35	0.00%
	21eil101	4264.0	15.50	4264.0	46.97	0.00%	4264.0	56.60	0.00%	4264.0	45.93	0.00%	4264.0	74.82	0.00%
	21lin105	4814.0	362.29	4814.0	368.71	0.00%	4814.0	1682.85	0.00%	4814.0	339.21	0.00%	4814.0	411.31	0.00%
	22pr107	4740.0	19.65	4740.0	16.97	0.00%	4740.0	11.05	0.00%	4740.0	18.69	0.00%	4740.0	52.44	0.00%
	25pr124	6054.0	3624.39	5035.0	746.37	16.83%	6054.0	3623.37	0.00%	6054.0	3623.57	0.00%	6054.0	3623.35	0.00%
	26bier127	6333.0	3635.77	6333.0	3633.48	0.00%	6333.0	3643.92	0.00%	6333.0	3636.61	0.00%	6333.0	3707.41	0.00%
	26ch130	6393.0	3634.77	6393.0	3633.04	0.00%	6503.0	3625.46	-1.72%	6393.0	3635.65	0.00%	6503.0	3627.20	-1.72%
	28spr136	6106.0	227.95	6106.0	228.75	0.00%	6106.0	1896.72	0.00%	6106.0	146.88	0.00%	6777.0	3624.38	-10.99%
	29pr144	7242.0	3626.17	7242.0	3626.51	0.00%	7242.0	3626.72	0.00%	7242.0	3626.60	0.00%	7242.0	3630.24	0.00%
	30ch150	6025.0	1209.80	6025.0	1214.47	0.00%	7533.0	3633.93	-25.03%	6025.0	995.18	0.00%	7533.0	3637.54	-25.03%
	30kroA150	7456.0	3634.86	7456.0	3633.38	0.00%	7533.0	3626.17	-1.03%	7456.0	3634.65	0.00%	7533.0	3624.25	-1.03%
	30kroB150	7524.0	3626.34	7524.0	3621.87	0.00%	7524.0	3627.64	0.00%	7524.0	3625.89	0.00%	7524.0	3628.93	0.00%
	31pr152	7613.0	3630.61	7613.0	3628.96	0.00%	7613.0	3634.35	0.00%	7613.0	3631.17	0.00%	7613.0	3628.25	0.00%
	32u159	7507.0	658.71	7507.0	435.99	0.00%	7850.0	3629.59	-4.57%	7507.0	919.53	0.00%	7860.0	3626.45	-4.70%
	39rat195	6813.0	1805.72	6813.0	1419.77	0.00%	6813.0	857.97	0.00%	6813.0	292.75	0.00%	6813.0	598.08	0.00%
	40d198	7480.0	378.10	7480.0	378.26	0.00%	9130.0	3637.92	-22.06%	7480.0	303.29	0.00%	7776.7	3648.54	-3.97%
AVG		808.27		696.90	0.31%		1230.15	-2.89%		751.66	0.00%		1479.63	-8.06%	
#Worse					0			7			0			13	

Table 13: Computational results for the five versions of the branch-and-cut algorithm on the instances with  $\omega = 0.4$  and profit  $g_2$ .

$\omega = 0.6$ and $g_1$																
	Instance	C-BC			NoCond			NoCover			NoCluCover			NoPath		
		UB	Time		UB	Time	Gap%	UB	Time	Gap%	UB	Time	Gap%	UB	Time	Gap%
Set1	11berlin52	43.0	3.91		43.0	3.86	0.00%	43.0	3.46	0.00%	43.0	3.88	0.00%	43.0	3.18	0.00%
	11eil51	39.0	1.75		39.0	1.96	0.00%	39.0	0.56	0.00%	39.0	1.81	0.00%	39.0	0.85	0.00%
	14st70	50.0	25.46		50.0	17.32	0.00%	50.0	28.17	0.00%	50.0	19.87	0.00%	50.0	37.84	0.00%
	16eil76	59.0	3.26		59.0	5.47	0.00%	59.0	26.91	0.00%	59.0	8.02	0.00%	59.0	19.06	0.00%
	16pr76	65.0	227.57		65.0	207.38	0.00%	65.0	487.74	0.00%	65.0	124.85	0.00%	65.0	2670.91	0.00%
	20kroA100	65.0	107.03		65.0	99.77	0.00%	65.0	2310.71	0.00%	65.0	106.66	0.00%	65.0	1640.40	0.00%
	20kroB100	66.0	98.45		66.0	99.49	0.00%	92.8	3629.78	-40.67%	66.0	97.08	0.00%	98.0	3629.25	-48.48%
	20kroC100	62.0	74.74		62.0	95.02	0.00%	62.0	585.94	0.00%	62.0	72.43	0.00%	62.0	863.48	0.00%
	20kroD100	64.0	75.17		64.0	63.98	0.00%	64.0	2621.97	0.00%	64.0	75.22	0.00%	64.0	2162.33	0.00%
	20kroE100	63.0	174.86		63.0	175.40	0.00%	63.0	104.17	0.00%	63.0	166.56	0.00%	63.0	168.14	0.00%
	20rat99	52.0	49.10		52.0	58.21	0.00%	52.0	125.12	0.00%	52.0	46.19	0.00%	52.0	150.13	0.00%
20rd100	72.0	211.81		72.0	209.05	0.00%	72.0	265.78	0.00%	72.0	368.21	0.00%	72.0	384.49	0.00%	
21eil101	82.0	91.61		82.0	123.22	0.00%	82.0	259.97	0.00%	82.0	82.72	0.00%	82.0	201.01	0.00%	
21lin105	78.0	248.31		78.0	342.03	0.00%	78.0	339.01	0.00%	78.0	132.16	0.00%	86.0	3635.17	-10.26%	
22pr107	77.0	3623.71		77.0	3623.48	0.00%	83.0	3626.10	-7.79%	77.0	3625.48	0.00%	83.0	3649.39	-7.79%	
Set2	11berlin52	51.0	0.15		51.0	0.15	0.00%	51.0	0.12	0.00%	51.0	0.14	0.00%	51.0	0.14	0.00%
	11eil51	50.0	3.74		50.0	3.88	0.00%	50.0	0.73	0.00%	50.0	3.76	0.00%	50.0	0.96	0.00%
	14st70	64.0	334.63		64.0	538.55	0.00%	64.0	1208.40	0.00%	64.0	300.55	0.00%	64.0	834.69	0.00%
	16eil76	74.0	176.04		74.0	621.02	0.00%	74.0	274.06	0.00%	74.0	175.62	0.00%	74.0	340.50	0.00%
	16pr76	74.0	1779.79		74.0	839.91	0.00%	75.0	3620.74	-1.35%	74.0	1777.29	0.00%	75.0	3619.31	-1.35%
	20kroA100	99.0	3623.43		99.0	3619.14	0.00%	99.0	3624.66	0.00%	99.0	3623.23	0.00%	99.0	3622.15	0.00%
	20kroB100	99.0	3621.47		99.0	3621.26	0.00%	99.0	3626.66	0.00%	99.0	3621.31	0.00%	99.0	3622.42	0.00%
	20kroC100	99.0	3618.73		99.0	3619.69	0.00%	99.0	3625.41	0.00%	99.0	3618.66	0.00%	99.0	3627.18	0.00%
	20kroD100	99.0	3618.70		93.0	3039.27	6.06%	99.0	3621.29	0.00%	99.0	3618.94	0.00%	99.0	3624.66	0.00%
	20kroE100	97.0	2207.74		99.0	3616.08	-2.06%	99.0	3622.52	-2.06%	97.0	2215.67	0.00%	99.0	3618.93	-2.06%
	20rat99	87.0	188.38		87.0	223.63	0.00%	87.0	256.23	0.00%	87.0	194.56	0.00%	87.0	70.03	0.00%
20rd100	99.0	2115.45		99.0	1210.77	0.00%	99.0	1885.67	0.00%	99.0	2135.44	0.00%	99.0	2752.34	0.00%	
21eil101	97.0	958.46		97.0	960.57	0.00%	100.0	3630.72	-3.09%	97.0	962.97	0.00%	100.0	3622.44	-3.09%	
21lin105	104.0	773.88		104.0	3070.16	0.00%	104.0	3638.71	0.00%	104.0	781.43	0.00%	104.0	3647.31	0.00%	
22pr107	106.0	11.02		106.0	11.22	0.00%	106.0	158.53	0.00%	106.0	11.02	0.00%	106.0	748.37	0.00%	
AVG		934.94		1004.03	0.13%		1573.66	-1.83%		932.39	0.00%		1765.57	-2.43%		
#Worse						1		5			0			6		
$\omega = 0.6$ and $g_2$																
	Instance	C-BC			NoCond			NoCover			NoCluCover			NoPath		
		UB	Time		UB	Time	Gap%	UB	Time	Gap%	UB	Time	Gap%	UB	Time	Gap%
Set1	11berlin52	2190.0	2.98		2190.0	2.82	0.00%	2190.0	2.17	0.00%	2190.0	3.72	0.00%	2190.0	3.70	0.00%
	11eil51	1911.0	4.63		1911.0	4.79	0.00%	1911.0	1.87	0.00%	1911.0	1.21	0.00%	1911.0	0.89	0.00%
	14st70	2589.0	25.40		2589.0	15.89	0.00%	2589.0	45.11	0.00%	2589.0	18.62	0.00%	2589.0	16.48	0.00%
	16eil76	3119.0	23.29		3119.0	23.06	0.00%	3119.0	82.51	0.00%	3119.0	19.60	0.00%	3119.0	46.33	0.00%
	16pr76	3275.0	56.14		3275.0	55.32	0.00%	3275.0	41.63	0.00%	3275.0	182.03	0.00%	3275.0	520.02	0.00%
	20kroA100	3192.0	219.32		3192.0	219.64	0.00%	3192.0	1467.65	0.00%	3192.0	135.17	0.00%	3192.0	2890.51	0.00%
	20kroB100	3203.0	161.54		3203.0	159.52	0.00%	3203.0	3506.17	0.00%	3203.0	165.41	0.00%	4753.0	3630.43	-48.39%
	20kroC100	3110.0	503.42		3110.0	241.53	0.00%	3110.0	727.83	0.00%	3110.0	246.28	0.00%	3110.0	597.71	0.00%
	20kroD100	3133.0	80.11		3133.0	127.92	0.00%	3133.0	1584.18	0.00%	3133.0	81.68	0.00%	3133.0	2556.97	0.00%
	20kroE100	2950.0	136.87		2950.0	136.95	0.00%	2950.0	989.92	0.00%	2950.0	82.57	0.00%	2950.0	301.43	0.00%
	20rat99	2643.0	43.82		2643.0	50.08	0.00%	2643.0	67.76	0.00%	2643.0	41.06	0.00%	2643.0	131.01	0.00%
20rd100	3591.0	390.73		3591.0	287.11	0.00%	3591.0	318.44	0.00%	3591.0	133.94	0.00%	3591.0	247.61	0.00%	
21eil101	4187.0	416.81		4187.0	326.02	0.00%	4187.0	358.10	0.00%	4187.0	412.53	0.00%	4187.0	644.74	0.00%	
21lin105	3955.0	181.35		3955.0	182.64	0.00%	3955.0	1324.81	0.00%	3955.0	162.10	0.00%	3955.0	1058.50	0.00%	
22pr107	3877.0	3626.58		4132.0	3627.09	-6.58%	4132.0	3625.61	-6.58%	3877.0	3626.40	0.00%	2697.0	3045.43	30.44%	
Set2	11berlin52	2608.0	0.15		2608.0	0.15	0.00%	2608.0	0.12	0.00%	2608.0	0.14	0.00%	2608.0	0.15	0.00%
	11eil51	2575.0	0.57		2575.0	0.56	0.00%	2575.0	0.88	0.00%	2575.0	0.56	0.00%	2575.0	1.79	0.00%
	14st70	3218.0	528.44		3218.0	334.69	0.00%	3218.0	2829.94	0.00%	3218.0	502.19	0.00%	3218.0	795.15	0.00%
	16eil76	3728.0	101.79		3728.0	87.29	0.00%	3728.0	441.42	0.00%	3728.0	100.66	0.00%	3728.0	92.92	0.00%
	16pr76	3729.0	476.85		3729.0	473.60	0.00%	3800.0	3623.70	-1.90%	3729.0	481.00	0.00%	3800.0	3623.80	-1.90%
	20kroA100	5008.0	3621.09		5008.0	3621.25	0.00%	5008.0	3627.00	0.00%	5008.0	3621.34	0.00%	5008.0	3626.83	0.00%
	20kroB100	5008.0	3622.99		5008.0	3622.28	0.00%	5008.0	3631.70	0.00%	5008.0	3623.16	0.00%	5008.0	3629.08	0.00%
	20kroC100	5008.0	3618.73		5008.0	3618.58	0.00%	5008.0	3620.71	0.00%	5008.0	3619.27	0.00%	5008.0	3620.53	0.00%
	20kroD100	5008.0	3618.38		5008.0	3617.24	0.00%	5008.0	3622.52	0.00%	5008.0	3618.92	0.00%	5008.0	3630.76	0.00%
	20kroE100	5008.0	3617.79		4910.0	3423.90	1.96%	5008.0	3618.89	0.00%	5008.0	3617.95	0.00%	5008.0	3622.17	0.00%
	20rat99	4516.0	173.82		4516.0	220.51	0.00%	4516.0	123.91	0.00%	4516.0	154.69	0.00%	4516.0	72.77	0.00%
20rd100	5008.0	3619.60		5008.0	1374.48	0.00%	5008.0	3028.85	0.00%	5008.0	3619.27	0.00%	5008.0	3620.90	0.00%	
21eil101	5050.0	3622.75		5037.0	3617.47	0.26%	5050.0	3621.34	0.00%	5050.0	3622.82	0.00%	5050.0	3628.80	0.00%	
21lin105	5228.0	3627.07		5228.0	3641.74	0.00%	5228.0	3641.79	0.00%	5228.0	3627.18	0.00%	5228.0	1629.92	0.00%	
22pr107	5363.0	132.19		5363.0	34.54	0.00%	5363.0	21.29	0.00%	5363.0	134.78	0.00%	5363.0	120.14	0.00%	
AVG		1208.51		1104.96	-0.15%		1653.26	-0.28%		1188.54	0.00%		1580.25	-0.66%		
#Worse						1		2			0			2		

Table 14: Computational results for the five versions of the branch-and-cut algorithm on the instances with  $\omega = 0.6$  and profits  $g_1$  and  $g_2$ .

$\omega = 0.8$ and $g_1$															
	C-BC			NoCond			NoCover			NoCluCover			NoPath		
	Instance	UB	Time	UB	Time	Gap%	UB	Time	Gap%	UB	Time	Gap%	UB	Time	Gap%
Set1	11berlin52	47.0	6.35	47.0	8.06	0.00%	47.0	33.77	0.00%	47.0	6.36	0.00%	47.0	25.64	0.00%
	11eil51	43.0	2.21	43.0	5.67	0.00%	43.0	6.30	0.00%	43.0	2.10	0.00%	43.0	7.00	0.00%
	14st70	65.0	388.64	65.0	65.07	0.00%	65.0	2173.28	0.00%	65.0	361.26	0.00%	65.0	1216.32	0.00%
	16eil76	69.0	175.45	69.0	92.37	0.00%	69.0	668.67	0.00%	69.0	153.38	0.00%	69.0	741.48	0.00%
	16pr76	72.0	1645.28	72.0	1610.18	0.00%	75.0	3625.52	-4.17%	72.0	1653.24	0.00%	73.0	3619.85	-1.39%
	20kroA100	79.0	223.57	79.0	224.33	0.00%	99.0	3628.66	-25.32%	79.0	224.73	0.00%	99.0	3630.21	-25.32%
	20kroB100	86.0	2791.01	99.0	3642.47	-15.12%	99.0	3635.00	-15.12%	86.0	2802.72	0.00%	99.0	3637.24	-15.12%
	20kroC100	83.0	448.35	83.0	394.61	0.00%	99.0	3630.59	-19.28%	83.0	417.48	0.00%	99.0	3630.37	-19.28%
	20kroD100	85.0	465.21	85.0	262.62	0.00%	99.0	3631.93	-16.47%	85.0	430.87	0.00%	99.0	3632.48	-16.47%
	20kroE100	80.0	334.23	80.0	182.45	0.00%	99.0	3625.67	-23.75%	80.0	347.91	0.00%	99.0	3625.78	-23.75%
	20rat99	79.0	1795.76	79.0	1587.52	0.00%	88.0	3624.21	-11.39%	79.0	1810.70	0.00%	88.0	3622.83	-11.39%
	20rd100	91.0	807.77	91.0	849.85	0.00%	99.0	3627.06	-8.79%	91.0	813.84	0.00%	99.0	3630.76	-8.79%
	21eil101	91.0	312.88	100.0	3628.12	-9.89%	100.0	3639.89	-9.89%	91.0	319.53	0.00%	100.0	3637.50	-9.89%
	21lin105	90.0	270.19	90.0	1541.15	0.00%	104.0	3639.01	-15.56%	90.0	289.91	0.00%	104.0	3644.62	-15.56%
	22pr107	106.0	3645.23	106.0	3642.42	0.00%	106.0	3643.92	0.00%	106.0	3645.25	0.00%	106.0	3632.86	0.00%
Set2	11berlin52	51.0	0.03	51.0	0.03	0.00%	51.0	0.02	0.00%	51.0	0.03	0.00%	51.0	0.03	0.00%
	11eil51	50.0	0.84	50.0	0.51	0.00%	50.0	0.83	0.00%	50.0	0.79	0.00%	50.0	1.31	0.00%
	14st70	69.0	3.56	69.0	6.05	0.00%	69.0	134.31	0.00%	69.0	3.57	0.00%	69.0	17.66	0.00%
	16eil76	75.0	4.04	75.0	3.24	0.00%	75.0	13.35	0.00%	75.0	4.03	0.00%	75.0	2.65	0.00%
	16pr76	75.0	8.46	75.0	18.80	0.00%	75.0	3643.29	0.00%	75.0	8.15	0.00%	75.0	1583.03	0.00%
	20kroA100	99.0	10.19	99.0	27.53	0.00%	99.0	528.28	0.00%	99.0	9.73	0.00%	99.0	713.57	0.00%
	20kroB100	99.0	1202.32	99.0	3626.85	0.00%	99.0	3634.82	0.00%	99.0	1207.78	0.00%	99.0	3625.98	0.00%
	20kroC100	99.0	3623.05	99.0	3622.40	0.00%	99.0	3631.01	0.00%	99.0	3623.28	0.00%	99.0	3631.17	0.00%
	20kroD100	99.0	3632.61	99.0	3634.61	0.00%	99.0	3629.19	0.00%	99.0	3631.99	0.00%	99.0	3622.21	0.00%
	20kroE100	99.0	3619.87	99.0	3626.20	0.00%	99.0	3637.93	0.00%	99.0	3620.10	0.00%	99.0	3630.18	0.00%
	20rat99	98.0	154.46	98.0	632.05	0.00%	98.0	939.68	0.00%	98.0	155.44	0.00%	98.0	1272.54	0.00%
	20rd100	99.0	135.40	99.0	98.07	0.00%	99.0	441.48	0.00%	99.0	135.97	0.00%	99.0	234.98	0.00%
	21eil101	100.0	1555.04	100.0	3625.01	0.00%	100.0	2144.42	0.00%	100.0	1569.38	0.00%	100.0	3625.32	0.00%
	21lin105	104.0	4.65	104.0	10.31	0.00%	104.0	3.26	0.00%	104.0	4.51	0.00%	104.0	6.79	0.00%
	22pr107	106.0	8.94	106.0	9.16	0.00%	106.0	12.09	0.00%	106.0	8.74	0.00%	106.0	16.52	0.00%
AVG		909.19		1222.59	-0.83%		2174.25	-4.99%		908.76	0.00%		2130.63	-4.90%	
#Worse					2			10			0			10	
$\omega = 0.8$ and $g_2$															
	C-BC			NoCond			NoCover			NoCluCover			NoPath		
	Instance	UB	Time	UB	Time	Gap%	UB	Time	Gap%	UB	Time	Gap%	UB	Time	Gap%
Set1	11berlin52	2384.0	10.57	2384.0	12.43	0.00%	2384.0	18.16	0.00%	2384.0	12.23	0.00%	2384.0	19.31	0.00%
	11eil51	2114.0	9.92	2114.0	12.42	0.00%	2114.0	9.15	0.00%	2114.0	6.94	0.00%	2114.0	7.18	0.00%
	14st70	3355.0	566.19	3355.0	134.57	0.00%	3355.0	659.94	0.00%	3355.0	533.94	0.00%	3355.0	372.47	0.00%
	16eil76	3573.0	60.10	3573.0	59.11	0.00%	3573.0	1294.30	0.00%	3573.0	85.04	0.00%	3573.0	2992.97	0.00%
	16pr76	3611.0	857.82	3611.0	843.61	0.00%	3694.0	3622.85	-2.30%	3694.0	3619.83	-2.30%	3765.0	3626.08	-4.26%
	20kroA100	4115.0	3088.95	4115.0	3015.71	0.00%	5008.0	3635.46	-21.70%	4115.0	3034.79	0.00%	5008.0	3632.68	-21.70%
	20kroB100	4916.0	3637.48	4916.0	3640.34	0.00%	5008.0	3628.18	-1.87%	4916.0	3637.40	0.00%	5008.0	3630.01	-1.87%
	20kroC100	3999.0	304.44	3999.0	574.05	0.00%	5008.0	3626.03	-25.23%	3999.0	287.95	0.00%	5008.0	3624.21	-25.23%
	20kroD100	5008.0	3624.83	5008.0	3625.03	0.00%	5008.0	3621.51	0.00%	5008.0	3625.22	0.00%	5008.0	3629.87	0.00%
	20kroE100	4002.0	388.55	4002.0	384.94	0.00%	4703.0	3622.29	-17.52%	4002.0	388.90	0.00%	5008.0	3631.36	-25.14%
	20rat99	3992.0	3085.27	4297.0	3626.93	-7.64%	4434.0	3622.18	-11.07%	3992.0	2785.76	0.00%	4434.0	3625.47	-11.07%
	20rd100	5008.0	3627.15	5008.0	3629.51	0.00%	5008.0	3638.36	0.00%	5008.0	3626.99	0.00%	5008.0	3633.61	0.00%
	21eil101	4717.0	1873.07	4717.0	1878.69	0.00%	5050.0	3633.56	-7.06%	4717.0	1786.74	0.00%	5050.0	3626.43	-7.06%
	21lin105	5092.0	3638.94	4572.7	3664.68	10.20%	5228.0	3625.12	-2.67%	5092.0	3638.53	0.00%	5228.0	3646.06	-2.67%
	22pr107	5363.0	3636.41	5363.0	3630.58	0.00%	5363.0	3637.33	0.00%	5363.0	3636.09	0.00%	5363.0	3650.05	0.00%
Set2	11berlin52	2608.0	0.08	2608.0	0.07	0.00%	2608.0	0.06	0.00%	2608.0	0.07	0.00%	2608.0	0.07	0.00%
	11eil51	2575.0	0.51	2575.0	0.50	0.00%	2575.0	0.48	0.00%	2575.0	0.51	0.00%	2575.0	0.62	0.00%
	14st70	3513.0	13.74	3513.0	13.62	0.00%	3513.0	24.55	0.00%	3513.0	13.25	0.00%	3513.0	81.31	0.00%
	16eil76	3800.0	163.65	3800.0	33.08	0.00%	3800.0	2.22	0.00%	3800.0	167.70	0.00%	3800.0	117.77	0.00%
	16pr76	3800.0	665.83	3800.0	248.03	0.00%	3800.0	984.88	0.00%	3800.0	670.13	0.00%	3800.0	677.96	0.00%
	20kroA100	5008.0	3624.63	5008.0	3138.14	0.00%	5008.0	3626.01	0.00%	5008.0	3623.81	0.00%	5008.0	3632.02	0.00%
	20kroB100	5008.0	3623.69	5008.0	3623.81	0.00%	5008.0	3629.08	0.00%	5008.0	3624.05	0.00%	5008.0	3636.64	0.00%
	20kroC100	5008.0	3622.66	5008.0	3622.57	0.00%	5008.0	3643.09	0.00%	5008.0	3622.69	0.00%	5008.0	3636.16	0.00%
	20kroD100	5008.0	3635.24	5008.0	3633.28	0.00%	5008.0	3623.96	0.00%	5008.0	3634.34	0.00%	5008.0	3624.31	0.00%
	20kroE100	5008.0	3627.77	5008.0	3627.91	0.00%	5008.0	3630.97	0.00%	5008.0	3627.43	0.00%	5008.0	3639.56	0.00%
	20rat99	5007.0	324.77	5007.0	110.66	0.00%	5007.0	209.81	0.00%	5007.0	323.20	0.00%	5007.0	1748.72	0.00%
	20rd100	5008.0	43.69	5008.0	72.08	0.00%	5008.0	10.58	0.00%	5008.0	45.00	0.00%	5008.0	10.01	0.00%
	21eil101	5050.0	3631.53	5050.0	3626.64	0.00%	5050.0	3630.00	0.00%	5050.0	3629.89	0.00%	5050.0	3643.53	0.00%
	21lin105	5228.0	1532.00	5228.0	1827.56	0.00%	5228.0	3518.19	0.00%	5228.0	1541.45	0.00%	5228.0	2311.88	0.00%
	22pr107	5363.0	137.27	5363.0	33.96	0.00%	5363.0	280.41	0.00%	5363.0	138.40	0.00%	5363.0	774.23	0.00%
AVG		1768.56		1744.82	0.09%		2290.29	-2.98%		1845.61	-0.08%		2362.75	-3.30%	
#Worse					1			8			1			8	

Table 15: Computational results for the five versions of the branch-and-cut algorithm on the instances with  $\omega = 0.8$  and profits  $g_1$  and  $g_2$ .