Factorizing 15 Using Shor's Algorithm and My Investigation of Error

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Outline

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 - Difficult case
 - Easy case
- Results
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 - Measurement error mitigation
 - Depth vs. Error

Review

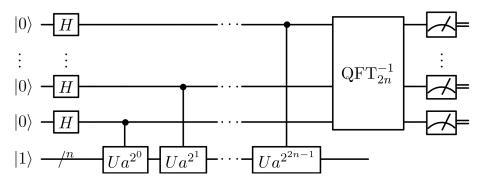


Figure: The circuit of Shor's algorithm.

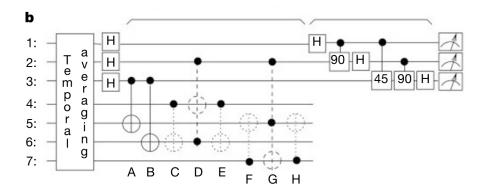
Method

Number of qubits: at least 6.

The options of a for N = 15:

- Easy case: 4, 11, 14 $(a^{2^k} \mod N = 1 \text{ for } k \ge 1)$
- Difficult case: 2, 7, 8, 13 ($a^{2^k} \mod N = 1 \text{ for } k \ge 2$)

Difficult case: Circuit (a = 7)

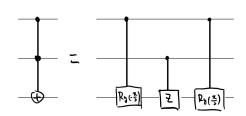


The $4n \mod 15$ gate

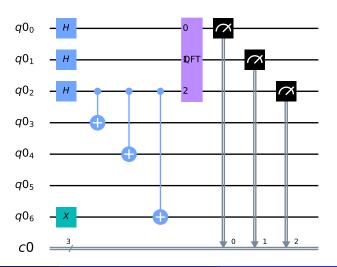
N		4N%15	
C	0000	0	0000
1	.0001	4	0100
2	0010	8	1000
3	0011	12	1100
4	0100	1	0001
5	0101	5	0101
6	0110	9	1001
7	0111	13	1101
8	1000	2	0010
9	1001	6	0110
10	1010	10	1010
11	1011	14	1110
12	1100	3	0011
13	1101	7	0111
14	1110	11	1011

Gate simplification

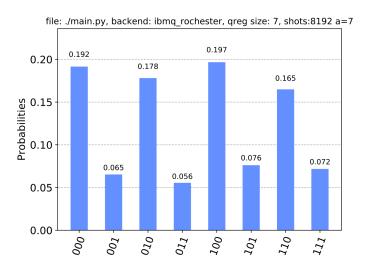
$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$



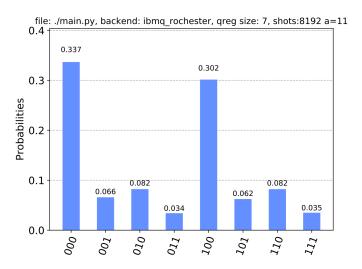
Easy case: Circuit



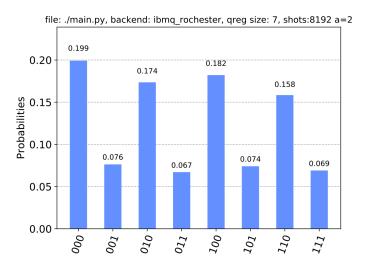
Result: a = 7



Result: a = 11



Result: a=2



Discussion

- How many qubits do we actually need in the first register?
- Do we need to know the answer in advance? Yes. While deciding the number of qubits of the first register.
- What do we need in order to factorize larger numbers? A powerful compiler is required (peephole compiler)

Can we factorize 21?

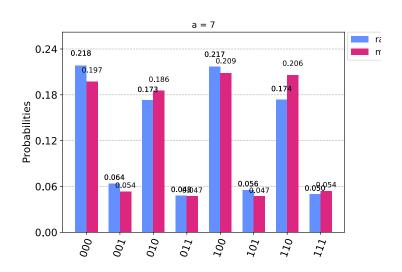
	_		31	'		3			
2	4	16	4	16	4	16	4	16	4
4	16	4	16	4	16	4	16	4	16
5	4	16	4	16	4	16	4	16	4
8	1	1	1	1	1	1	1	1	1
10	16	4	16	4	16	4	16	4	16
11	16	4	16	4	16	4	16	4	16
13	1	1	1	1	1	1	1	1	1
16	4	16	4	16	4	16	4	16	4
17	16	4	16	4	16	4	16	4	16
19	4	16	4	16	4	16	4	16	4
20	1	1	1	1	1	1	1	1	1

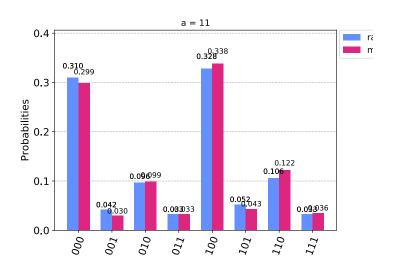
Measurement error mitigation

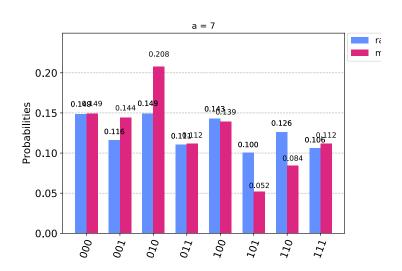
Assuming that there's some probability for each bitstring to flip to another while measured.

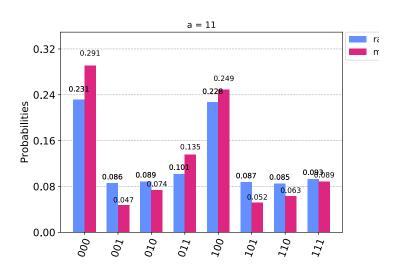
$$C_{noisy} = MC_{ideal}$$

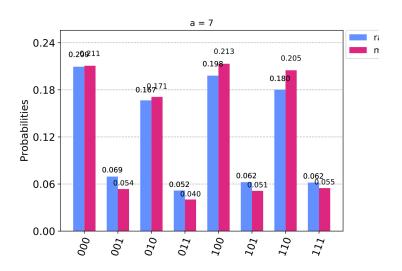
⇒Try each bitstring to construct the matrix!

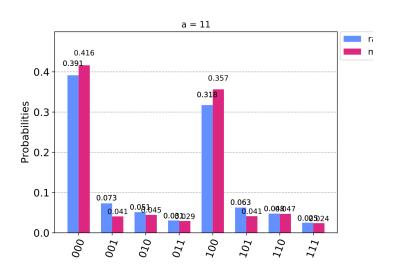












Discussion

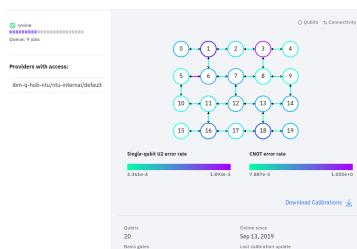
The result was not well. What was causing the error? Depth? What determines the depth?

- Hardware
- Transpiler/compiler

Hardware

ibmq_almaden v1.4.6





u1, u2, u3, cx, id

Maximum shots

8192

1.000e+0

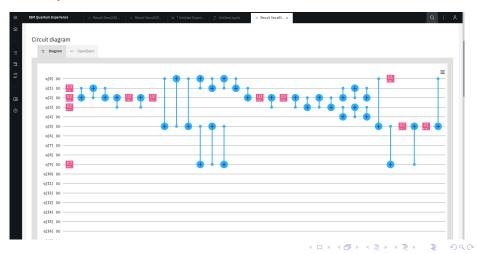
May 23, 2020 2:42 PM

Maximum experiments

900

Transpiling/compiling

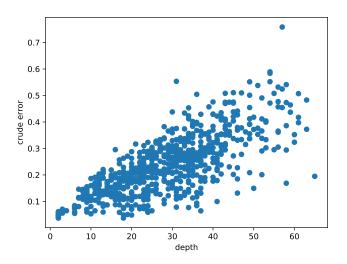
Use qk.compiler.transpile(qc, backend=be, initial_layout=layout, optimization_level=1) to assign manually.



Depth vs. error

- ullet Obviously depth $\uparrow\Rightarrow$ error \uparrow
 - But quantitatively?
- Create random circuits to test errors!.
- What to expect? Exponential decay of accuracy.

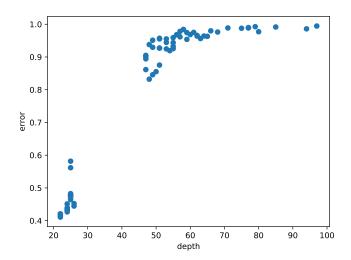
Depth vs. error: random circuits



Discussion

- Bad definition of error!
- Use a circuit with definite results.

Depth vs. error: adder



Depth vs. error: adder

