Algo 10 9 Novembre 2017

Multiplication de 2 polynomes de degre n

$$\begin{split} A(x) &= a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0 \\ B(x) &= b_n x^n + b_{n-1} x^{n-1} + \dots + b_1 x + b_0 \\ C(x) &= A(x) \times B(x) = c_{2n} x^{2n} + c_{2n-1} x^{2n-1} + \dots + c_0 \end{split}$$

Donc:

$$c_i = \sum_j a_j b_{i-j} \tag{1}$$

Il y a donc $(n+1)^2$ multiplications et n^2 addition

Ca c'est l'algo naif de multiplication de polynome.

Karatsuba algorithm

Voir le lien wikipedia. (Deja implem en TP de sup TinyBistro).

$$(P_1 \times P_2)(x) = Q_1(x) \times Q_2(x) \times x^n + (Q_1(x) \times R_2(x)) + Q_2(x) \times R_1(x))x^{n/2}$$

```
1 def Karatsuba(P1, P2, n):
  if n == 1:
2
     return [P1[0] + P2[0], 0]
4 Q1 = P1[n/2:]
5 	 R1 = P1[:n/2]
6 Q2 = P2[n/2:]
7 	 R2 = P1[:n/2]
8 A = Karatsuba(Q1, Q2, n/2)
9 B = Karatsuba(R1, R2, n/2)
10 C = Karatsuba(Q1 - R1, R2 - Q2, n/2)
for i in range(n):
     S[i] = B[i]
12
for i in range(n, 2 * n):
      S[i] = A[i]
14
     for i in range(n/2, 3 * n / 2):
15
16
       S[i] += A[i] + B[i] + C[i]
17
     return S
```

En complexite:

$$T(n) = 3T(\frac{n}{2} + \Theta(n))$$

$$T(n) = \Theta(1)$$

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