

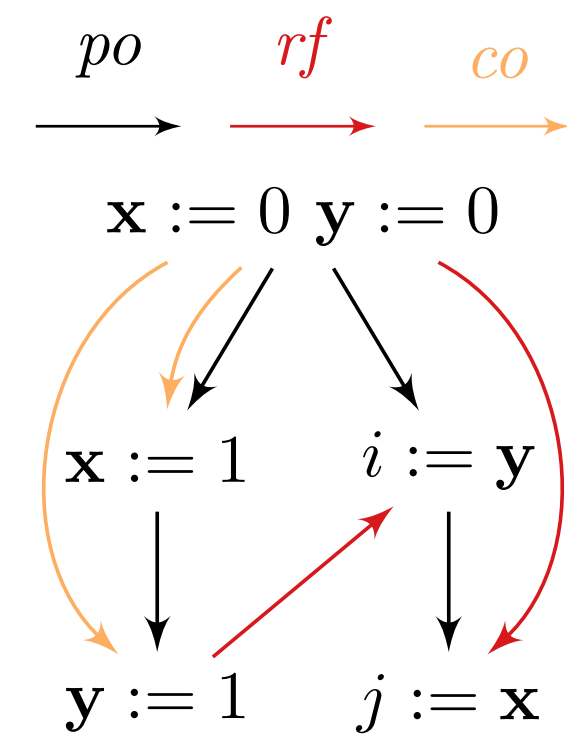
EXAMPLE

x := 0, **y** := 0

x := 1 | *i* := **y**
y := 1 | *j* := **x**

$\neg(i = 1 \wedge j = 0)$

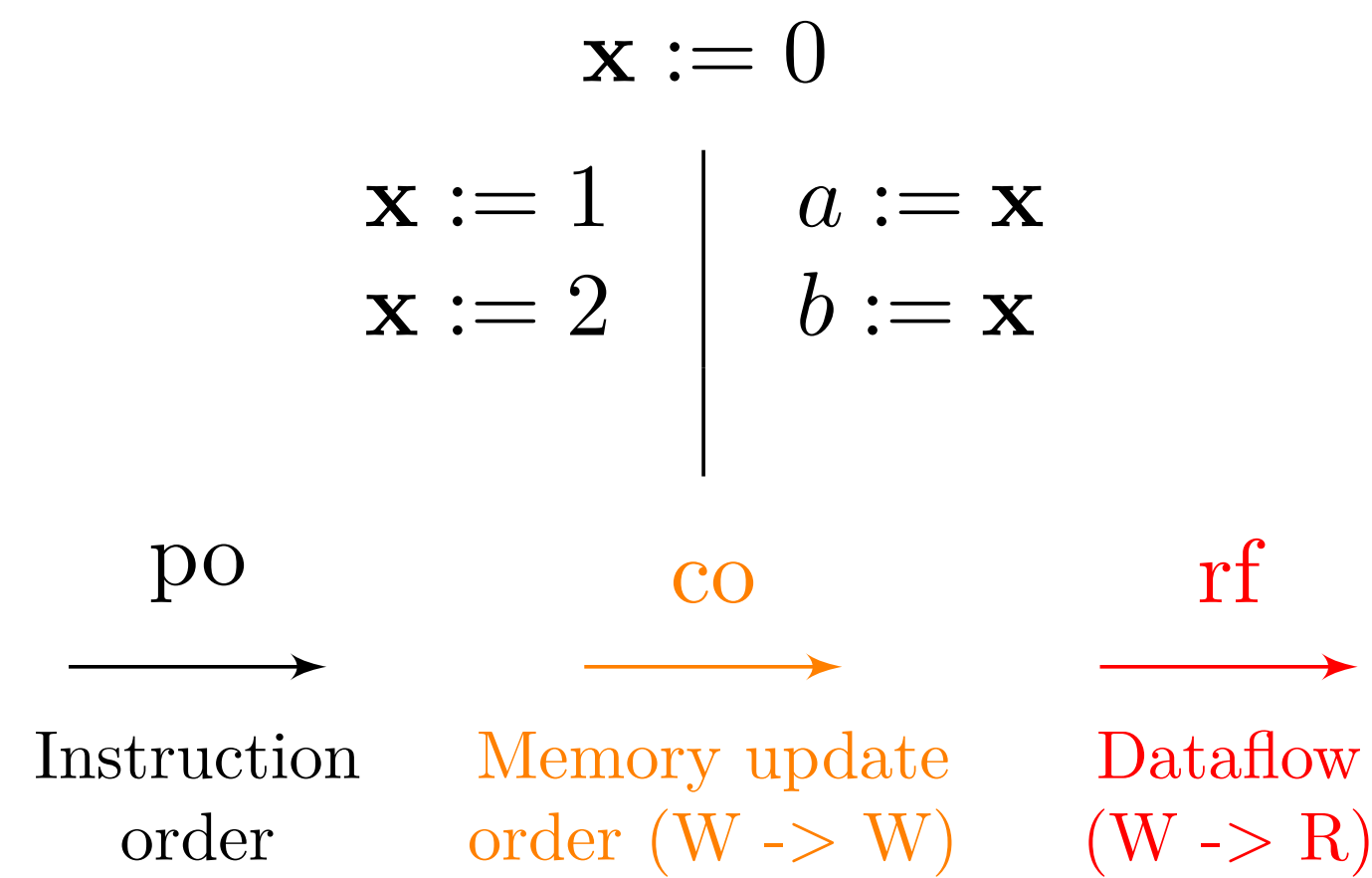
Program with property



CEx. candidate

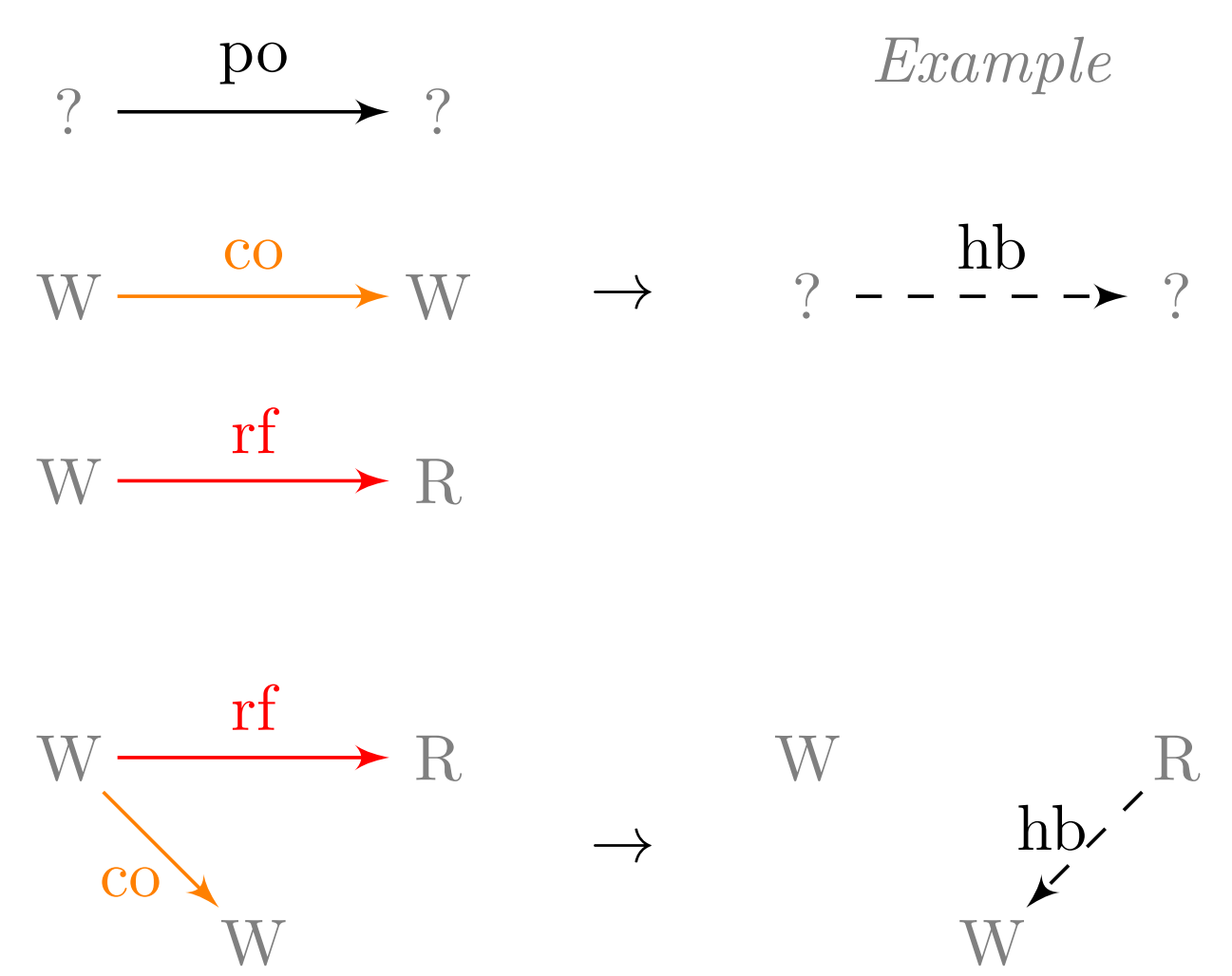
Bold variables are global, *italicized* variables are local.
The program is safe under SC and TSO but not PSO.

MEMORY MODELS OVERVIEW



$$\forall r^x \exists w^x : w^x \xrightarrow{rf} r^x$$

$$\forall w_1^x, w_2^x : w_1^x \xrightarrow{co} w_2^x \vee w_2^x \xrightarrow{co} w_1^x$$



Find *co* and *rf* for a given *po* such that *hb* is acyclic.

TOOLS FOR WEAK MEMORY

Exhaustive Enumeration

Generate execution candidates, and check their consistency

Herd7 [2]

(memory model simulator)
Litmus tests CAT
memory model

Stateless Model Checking

Generate increasingly larger, always consistent executions (traces)

GenMC [5], **Nidhugg** [1],

...
(Subset of) C11
Custom library

Bounded Model Checking

Encode constraints of the memory model in the SMT query

Dartagnan [4]

(SV-COMP flavored) C
Subset of CAT

VIOLATION WITNESS EXAMPLE

Thread 0	waypoint type	value	line	column
	assume	$\backslash at(\mathbf{x}, 0) = 0$	0	<i>middle</i>
	assume	$\backslash at(\mathbf{y}, 0) = 0$	0	<i>end</i>
	thread_start	1, 2	1	0
Thread 1				
	assume	$\backslash at(\mathbf{x}, 1) = 1$	1	<i>end</i>
	assume	$\backslash at(\mathbf{y}, 1) = 1$	2	<i>end</i>
Thread 2				
	assume	$i = \backslash at(\mathbf{x}, 1)$	1	<i>end</i>
	assume	$j = \backslash at(\mathbf{y}, 1)$	2	<i>end</i>
	target	-	2	<i>end</i>

A violation witness, encoding a violation under PSO

CORRECTNESS WITNESS EXAMPLE

invariant type	value	line	column
location	$\backslash at(\mathbf{x}, 0) = 0$	0	<i>middle</i>
location	$\backslash at(\mathbf{y}, 0) = 0$	0	<i>end</i>
location	$\backslash at(\mathbf{x}, 1) = 1$	1 (<i>left</i>)	<i>end</i>
location	$\backslash at(\mathbf{y}, 1) = 1$	2 (<i>left</i>)	<i>end</i>
location	$\exists a : a \in \{0, 1\}$ $i = \backslash at(\mathbf{x}, a)$	1 (<i>right</i>)	<i>end</i>
location	$\exists a, b : a, b \in \{0, 1\}$ $j = \backslash at(\mathbf{y}, a)$ $i = \backslash at(\mathbf{x}, b)$ $b = 1 \implies a = 1$	2 (<i>right</i>)	<i>end</i>
location	$\neg(i = 1 \wedge j = 0)$	2 (<i>right</i>)	<i>end</i>

A correctness witness, encoding a proof over SC

MAPPING VERDICTS TO WITNESSES

$\backslash at(\mathbf{e}, \mathbf{id})$: Built-in ACSL construct (*abused a bit*)

- referring to the value of the expression **e** in the state at label **id** [3]
- Our *state labels* are integers, and denote ordering of memory events.
- Correctness: *state labels* are **symbolic** integers, and denote ordering of memory events.

FUTURE PLANS

- Implement witness serialization (THETA, CPACHECKER)
- Implement violation witness checking (THETA)
- Implement correctness witness checking (THETA)



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REPORT

